$D \ R \ A \ F \ T$ Document for a Standard Message-Passing Interface

Message Passing Interface Forum

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Chapter 14

Tool Support

14.1 Introduction

This chapter discusses interfaces that allow debuggers, performance analyzers, and other tools to extract information about the operation of MPI processes. Specifically, this chapter defines both the MPI profiling interface (Section 14.2), which supports the transparent interception and inspection of MPI calls, and the MPI tool information interface (Section 14.3), which supports the inspection and manipulation of MPI control and performance variables. The interfaces described in this chapter are all defined in the context of an MPI process, i.e., are callable from the same code that invokes other MPI functions.

14.2 Profiling Interface

14.2.1 Requirements

To meet the requirements for the MPI profiling interface, an implementation of the MPI functions must

1. provide a mechanism through which all of the MPI defined functions, except those allowed as macros (See Section 2.6.4), may be accessed with a name shift. This requires, in C and Fortran, an alternate entry point name, with the prefix PMPI_ for each MPI function in each provided language binding and language support method. For routines implemented as macros, it is still required that the PMPI_ version be supplied and work as expected, but it is not possible to replace at link time the MPI_ version with a user-defined version.

For Fortran, the different support methods cause several linker names. Therefore, several profiling routines (with these linker names) are needed for each Fortran MPI routine, as described in Section 17.1.5 on page 605.

- 2. ensure that those MPI functions that are not replaced may still be linked into an executable image without causing name clashes.
- 3. document the implementation of different language bindings of the MPI interface if they are layered on top of each other, so that the profiler developer knows whether she must implement the profile interface for each binding, or can economize by implementing it only for the lowest level routines.

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4. where the implementation of different language bindings is done through a layered approach (e.g., the Fortran binding is a set of "wrapper" functions that call the C implementation), ensure that these wrapper functions are separable from the rest of the library.

This separability is necessary to allow a separate profiling library to be correctly implemented, since (at least with Unix linker semantics) the profiling library must contain these wrapper functions if it is to perform as expected. This requirement allows the person who builds the profiling library to extract these functions from the original MPI library and add them into the profiling library without bringing along any other unnecessary code.

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5. provide a no-op routine MPI_PCONTROL in the MPI library.

14.2.2 Discussion

15The objective of the MPI profiling interface is to ensure that it is relatively easy for authors 16 of profiling (and other similar) tools to interface their codes to MPI implementations on 17different machines. 18

Since MPI is a machine independent standard with many different implementations, 19 it is unreasonable to expect that the authors of profiling tools for MPI will have access to 20the source code that implements MPI on any particular machine. It is therefore necessary 21to provide a mechanism by which the implementors of such tools can collect whatever 22performance information they wish without access to the underlying implementation. 23

We believe that having such an interface is important if MPI is to be attractive to end 24users, since the availability of many different tools will be a significant factor in attracting 25users to the MPI standard. 26

The profiling interface is just that, an interface. It says *nothing* about the way in which 27 it is used. There is therefore no attempt to lay down what information is collected through 28the interface, or how the collected information is saved, filtered, or displayed. 29

While the initial impetus for the development of this interface arose from the desire to 30 permit the implementation of profiling tools, it is clear that an interface like that specified 31 may also prove useful for other purposes, such as "internetworking" multiple MPI imple-32mentations. Since all that is defined is an interface, there is no objection to its being used 33 wherever it is useful. 34

As the issues being addressed here are intimately tied up with the way in which ex-35 ecutable images are built, which may differ greatly on different machines, the examples 36 given below should be treated solely as one way of implementing the objective of the MPI 37 profiling interface. The actual requirements made of an implementation are those detailed 38 in the Requirements section above, the whole of the rest of this section is only present as 39 justification and discussion of the logic for those requirements. 40

The examples below show one way in which an implementation could be constructed to meet the requirements on a Unix system (there are doubtless others that would be equally 42valid). 43

Logic of the Design 14.2.3 45

46Provided that an MPI implementation meets the requirements above, it is possible for 47the implementor of the profiling system to intercept the MPI calls that are made by the 48

user program. She can then collect whatever information she requires before calling the underlying MPI implementation (through its name shifted entry points) to achieve the desired effects.

14.2.4 Miscellaneous Control of Profiling

There is a clear requirement for the user code to be able to control the profiler dynamically at run time. This capability is normally used for (at least) the purposes of

Profiling level (integer)

- Enabling and disabling profiling depending on the state of the calculation.
- Flushing trace buffers at non-critical points in the calculation.
- Adding user events to a trace file.

These requirements are met by use of MPI_PCONTROL.

MPI_PCONTROL(level, ...)

IN level

```
int MPI_Pcontrol(const int level, ...)
```

```
MPI_Pcontrol(level) BIND(C)
    INTEGER, INTENT(IN) :: level
```

MPI_PCONTROL(LEVEL) INTEGER LEVEL

MPI libraries themselves make no use of this routine, and simply return immediately to the user code. However the presence of calls to this routine allows a profiling package to be explicitly called by the user.

Since MPI has no control of the implementation of the profiling code, we are unable to specify precisely the semantics that will be provided by calls to MPI_PCONTROL. This vagueness extends to the number of arguments to the function, and their datatypes.

However to provide some level of portability of user codes to different profiling libraries, we request the following meanings for certain values of level.

- level==0 Profiling is disabled.
- level==1 Profiling is enabled at a normal default level of detail.
- level==2 Profile buffers are flushed, which may be a no-op in some profilers.
- All other values of level have profile library defined effects and additional arguments.

We also request that the default state after MPI_INIT has been called is for profiling to be enabled at the normal default level. (i.e., as if MPI_PCONTROL had just been called with the argument 1). This allows users to link with a profiling library and to obtain profile output without having to modify their source code at all.

The provision of MPI_PCONTROL as a no-op in the standard MPI library supports the collection of more detailed profiling information with source code that can still link against the standard MPI library. 48

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1 14.2.5 Profiler Implementation Example

A profiler can accumulate the total amount of data sent by the MPI_SEND function, along with the total elapsed time spent in the function as the following example shows:

```
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     Example 14.1
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     static int totalBytes = 0;
7
     static double totalTime = 0.0;
8
9
     int MPI_Send(const void* buffer, int count, MPI_Datatype datatype,
10
                    int dest, int tag, MPI_Comm comm)
11
12
     {
        double tstart = MPI_Wtime();
                                               /* Pass on all arguments */
13
        int extent;
14
        int result
                        = PMPI_Send(buffer,count,datatype,dest,tag,comm);
15
16
        totalTime += MPI_Wtime() - tstart;
                                                          /* and time
                                                                                 */
17
18
        MPI_Type_size(datatype, &extent); /* Compute size */
19
        totalBytes += count*extent;
20
21
        return result;
22
     }
23
24
            MPI Library Implementation Example
25
     14.2.6
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     If the MPI library is implemented in C on a Unix system, then there are various options,
27
     including the two presented here, for supporting the name-shift requirement. The choice
28
     between these two options depends partly on whether the linker and compiler support weak
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```

₃₀ symbols.

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32 Systems with Weak Symbols

If the compiler and linker support weak external symbols (e.g., Solaris 2.x, other System
 V.4 machines), then only a single library is required as the following example shows:

```
<sup>36</sup> Example 14.2
```

The effect of this **#pragma** is to define the external symbol MPI_Example as a weak definition. This means that the linker will not complain if there is another definition of the symbol (for instance in the profiling library); however if no other definition exists, then the linker will use the weak definition.

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Systems Without Weak Symbols

In the absence of weak symbols then one possible solution would be to use the C macro preprocessor as the following example shows:

Example 14.3

```
#ifdef PROFILELIB
# ifdef __STDC__
# define FUNCTION(name) P##name
# else
# define FUNCTION(name) P/**/name
# endif
#else
# define FUNCTION(name) name
#endif
```

Each of the user visible functions in the library would then be declared thus

```
int FUNCTION(MPI_Example)(/* appropriate args */)
{
```

/* Useful content */
}

The same source file can then be compiled to produce both versions of the library, depending on the state of the PROFILELIB macro symbol.

It is required that the standard MPI library be built in such a way that the inclusion of MPI functions can be achieved one at a time. This is a somewhat unpleasant requirement, since it may mean that each external function has to be compiled from a separate file. However this is necessary so that the author of the profiling library need only define those MPI functions that she wishes to intercept, references to any others being fulfilled by the normal MPI library. Therefore the link step can look something like this

```
% cc ... -lmyprof -lpmpi -lmpi
```

Here libmyprof.a contains the profiler functions that intercept some of the MPI functions, libpmpi.a contains the "name shifted" MPI functions, and libmpi.a contains the normal definitions of the MPI functions.

14.2.7 Complications

Multiple Counting

Since parts of the MPI library may themselves be implemented using more basic MPI func-tions (e.g., a portable implementation of the collective operations implemented using point to point communications), there is potential for profiling functions to be called from within an MPI function that was called from a profiling function. This could lead to "double counting" of the time spent in the inner routine. Since this effect could actually be useful under some circumstances (e.g., it might allow one to answer the question "How much time is spent in the point to point routines when they are called from collective functions?"), we have decided not to enforce any restrictions on the author of the MPI library that would

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overcome this. Therefore the author of the profiling library should be aware of this problem,
 and guard against it. In a single-threaded world this is easily achieved through use of a
 static variable in the profiling code that remembers if you are already inside a profiling
 routine. It becomes more complex in a multi-threaded environment (as does the meaning
 of the times recorded).

Linker Oddities

The Unix linker traditionally operates in one pass: the effect of this is that functions from libraries are only included in the image if they are needed at the time the library is scanned. When combined with weak symbols, or multiple definitions of the same function, this can cause odd (and unexpected) effects.

Consider, for instance, an implementation of MPI in which the Fortran binding is 13 achieved by using wrapper functions on top of the C implementation. The author of the 14 profile library then assumes that it is reasonable only to provide profile functions for the C 15binding, since Fortran will eventually call these, and the cost of the wrappers is assumed 16 to be small. However, if the wrapper functions are not in the profiling library, then none 17of the profiled entry points will be undefined when the profiling library is called. Therefore 18 none of the profiling code will be included in the image. When the standard MPI library 19 is scanned, the Fortran wrappers will be resolved, and will also pull in the base versions of 20the MPI functions. The overall effect is that the code will link successfully, but will not be 21profiled. 22

To overcome this we must ensure that the Fortran wrapper functions are included in the profiling version of the library. We ensure that this is possible by requiring that these be separable from the rest of the base MPI library. This allows them to be copied out of the base library and into the profiling one using a tool such as **ar**.

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28 Fortran Support Methods

The different Fortran support methods and possible options for the support of subarrays (depending on whether the compiler can support TYPE(*), DIMENSION(..) choice buffers) imply different linker names for the same Fortran MPI routine. The rules and implications for the profiling interface are described in Section 17.1.5 on page 605.

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³⁴₃₅ 14.2.8 Multiple Levels of Interception

The scheme given here does not directly support the nesting of profiling functions, since it provides only a single alternative name for each MPI function. Consideration was given to an implementation that would allow multiple levels of call interception, however we were unable to construct an implementation of this that did not have the following disadvantages

- 40 41
- assuming a particular implementation language,
- 42 43
- imposing a run time cost even when no profiling was taking place.

Since one of the objectives of MPI is to permit efficient, low latency implementations, and
 it is not the business of a standard to require a particular implementation language, we
 decided to accept the scheme outlined above.

⁴⁷ Note, however, that it is possible to use the scheme above to implement a multi-level
⁴⁸ system, since the function called by the user may call many different profiling functions

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before calling the underlying MPI function. This capability has been demonstrated in the P^N MPI tool infrastructure [?].

14.3 The MPI Tool Information Interface

MPI implementations often use internal variables to control their operation and performance. Understanding and manipulating these variables can provide a more efficient execution environment or improve performance for many applications. This section describes the MPI tool information interface, which provides a mechanism for MPI implementors to expose variables, each of which represents a particular property, setting, or performance measurement from within the MPI implementation. The interface is split into two parts: the first part provides information about and supports the setting of control variables through which the MPI implementation tunes its configuration. The second part provides access to performance variables that can provide insight into internal performance information of the MPI implementation.

To avoid restrictions on the MPI implementation, the MPI tool information interface allows the implementation to specify which control and performance variables exist. Additionally, the user of the MPI tool information interface can obtain metadata about each available variable, such as its datatype, and a textual description. The MPI tool information interface provides the necessary routines to find all variables that exist in a particular MPI implementation, to query their properties, to retrieve descriptions about their meaning, and to access and, if appropriate, to alter their values. Variables and categories across connected processes with the same name are required to have the same meaning. The intent of this requirement is to enforce consistency across connected processes. For example, variables describing the number of packets sent on heterogeneous network devices should have different names to reflect their potentially different meanings. Specification of what must be the same for variables and categories is stated in their respective sections.

The MPI tool information interface can be used independently from the MPI communication functionality. In particular, the routines of this interface can be called before MPI_INIT (or equivalent) and after MPI_FINALIZE. In order to support this behavior cleanly, the MPI tool information interface uses separate initialization and finalization routines. All identifiers used in the MPI tool information interface have the prefix MPI_T_.

On success, all MPI tool information interface routines return MPI_SUCCESS, otherwise they return an appropriate and unique return code indicating the reason why the call was not successfully completed. Details on return codes can be found in Section 14.3.9. However, unsuccessful calls to the MPI tool information interface are not fatal and do not impact the execution of subsequent MPI routines.

Since the MPI tool information interface primarily focuses on tools and support libraries, MPI implementations are only required to provide C bindings for functions and constants introduced in this section. Except where otherwise noted, all conventions and principles governing the C bindings of the MPI API also apply to the MPI tool information interface, which is available by including the mpi.h header file. All routines in this interface have local semantics.

Advice to users. The number and type of control variables and performance variables can vary between MPI implementations, platforms and different builds of the same implementation on the same platform as well as between runs. Hence, any application relying on a particular variable will not be portable. Further, there is no guaran-

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tee that the number of variables and variable indices are the same across connected processes.

This interface is primarily intended for performance monitoring tools, support tools, and libraries controlling the application's environment. When maximum portability is desired, application programmers should either avoid using the MPI tool information interface or avoid being dependent on the existence of a particular control or performance variable. (*End of advice to users.*)

14.3.1 Verbosity Levels

The MPI tool information interface provides access to internal configuration and perfor-mance information through a set of control and performance variables defined by the MPI implementation. Since some implementations may export a large number of variables, variables are classified by a verbosity level that categorizes both their intended audience (end users, performance tuners or MPI implementors) and a relative measure of level of detail (basic, detailed or all). These verbosity levels are described by a single integer. Table 14.1 lists the constants for all possible verbosity levels. The values of the con-stants are monotonic in the order listed in the table; i.e., MPI_T_VERBOSITY_USER_BASIC < MPI_T_VERBOSITY_USER_DETAIL < ... < MPI_T_VERBOSITY_MPIDEV_ALL.

20		
21	MPI_T_VERBOSITY_USER_BASIC	Basic information of interest to users
22	MPI_T_VERBOSITY_USER_DETAIL	Detailed information of interest to users
23	MPI_T_VERBOSITY_USER_ALL	All remaining information of interest to users
24	MPI_T_VERBOSITY_TUNER_BASIC	Basic information required for tuning
25	MPI_T_VERBOSITY_TUNER_DETAIL	Detailed information required for tuning
26	MPI_T_VERBOSITY_TUNER_ALL	All remaining information required for tuning
27	MPI_T_VERBOSITY_MPIDEV_BASIC	Basic information for MPI implementors
28	MPI_T_VERBOSITY_MPIDEV_DETAIL	Detailed information for MPI implementors
29	MPI_T_VERBOSITY_MPIDEV_ALL	All remaining information for MPI implementors

Table 14.1: MPI tool information interface verbosity levels

14.3.2 Binding MPI Tool Information Interface Variables to MPI Objects

Each MPI tool information interface variable provides access to a particular control setting or performance property of the MPI implementation. A variable may refer to a specific MPI object such as a communicator, datatype, or one-sided communication window, or the variable may refer more generally to the MPI environment of the process. Except for the last case, the variable must be bound to exactly one MPI object before it can be used. Table 14.2 lists all MPI object types to which an MPI tool information interface variable can be bound, together with the matching constant that MPI tool information interface routines return to identify the object type.

Rationale. Some variables have meanings tied to a specific MPI object. Examples
 include the number of send or receive operations that use a particular datatype, the
 number of times a particular error handler has been called, or the communication pro tocol and "eager limit" used for a particular communicator. Creating a new MPI tool

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Constant	MPI object
MPI_T_BIND_NO_OBJECT	N/A; applies globally to entire MPI process
MPI_T_BIND_MPI_COMM	MPI communicators
MPI_T_BIND_MPI_DATATYPE	MPI datatypes
MPI_T_BIND_MPI_ERRHANDLER	MPI error handlers
MPI_T_BIND_MPI_FILE	MPI file handles
MPI_T_BIND_MPI_GROUP	MPI groups
MPI_T_BIND_MPI_OP	MPI reduction operators
MPI_T_BIND_MPI_REQUEST	MPI requests
MPI_T_BIND_MPI_WIN	MPI windows for one-sided communication
MPI_T_BIND_MPI_MESSAGE	MPI message object
MPI_T_BIND_MPI_INFO	MPI info object

Table 14.2: Constants to identify associations of variables

information interface variable for each MPI object would cause the number of variables to grow without bound, since they cannot be reused to avoid naming conflicts. By associating MPI tool information interface variables with a specific MPI object, the MPI implementation only must specify and maintain a single variable, which can then be applied to as many MPI objects of the respective type as created during the program's execution. (*End of rationale.*)

14.3.3 Convention for Returning Strings

Several MPI tool information interface functions return one or more strings. These functions have two arguments for each string to be returned: an [OUT]OUT parameter that identifies a pointer to the buffer in which the string will be returned, and an [IN/OUT]IN/OUT parameter to pass the length of the buffer. The user is responsible for the memory allocation of the buffer and must pass the size of the buffer (n) as the length argument. Let n be the length value specified to the function. On return, the function writes at most n - 1 of the string's characters into the buffer, followed by a null terminator. If the returned string's length is greater than or equal to n, the string will be truncated to n - 1 characters. In this case, the length of the string plus one (for the terminating null character) is returned in the length argument. If the user passes the null pointer as the buffer argument or passes 0 as the length argument, the function does not return the string and only returns the length of the string plus one in the length argument. If the user passes the null pointer as the null pointer as the length of the string plus one in the length argument. If the user passes the null pointer as the null pointer as the length of the string plus one in the length argument. If the user passes the null pointer as the length of the string plus one in the length argument. If the user passes the null pointer as the length argument argument is ignored and nothing is returned.

14.3.4 Initialization and Finalization

The MPI tool information interface requires a separate set of initialization and finalization routines.

IN	required	desired level of thread support (integer)
OUT	provided	provided level of thread support (integer)
int MPI.	_T_init_thread(int	t required, int *provided)
MPI tool any other MPI_T_I izes the t this rout beyond in argument and thein listed in thread su informati The MPI_INIT the user is all process inherit th well as all Process	information interface r of its routines. A us NIT_THREAD, which thread environment for ine when the MPI to ncreasing the reference t required is used to sp r semantics are ident Section 12.4. The ca upport that will be ion interface routines MPI specification do T. If the MPI tool infi is responsible for ensu- sses it is used in. Pro- ne status of the MPI Il active sessions and resses created at runti	hat use the MPI tool information interface must initialize ce in the processes that will use the interface before cal ser can initialize the MPI tool information interface by cal h can be called multiple times. In addition, this routine initial for all routines in the MPI tool information interface. Cal cool information interface is already initialized has no efficient count of how often the interface has been initialized. If specify the desired level of thread support. The possible va- tical to the ones that can be used with MPI_INIT_THRE call returns in provided information about the actual level provided by the MPI implementation for calls to MPI s. It can be one of the four values listed in Section 12.4. oes not require all MPI processes to exist before the cal formation interface is used before MPI_INIT_has been cal suring that the MPI tool information interface is initialized tool information interface (whether it is initialized or no l handles) from the process from which they are created. time as a result of calls to MPI's dynamic process managem on before they can use the MPI tool information interface.
		MPI_T_INIT_THREAD is called before MPI_INIT_THRE
the	requested and grant	MPI_T_INIT_THREAD is called before MPI_INIT_THRE ted thread level for MPI_T_INIT_THREAD may influence lue of MPI_INIT_THREAD. The same is true for the rev

Advice to implementors. MPI implementations should strive to make as many control or performance variables available before MPI_INIT (instead of adding them within MPI_INIT) to allow tools the most flexibility. In particular, control variables should be available before MPI_INIT if their value cannot be changed after MPI_INIT. (*End* of advice to implementors.)

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- MPI_T_FINALIZE()
- int MPI_T_finalize(void)

This routine finalizes the use of the MPI tool information interface and may be called as often as the corresponding MPI_T_INIT_THREAD routine up to the current point of execution. Calling it more times returns a corresponding error code. As long as the number of calls to MPI_T_FINALIZE is smaller than the number of calls to MPI_T_INIT_THREAD up to the current point of execution, the MPI tool information interface remains initialized

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and calls to its routines are permissible. Further, additional calls to MPI_T_INIT_THREAD after one or more calls to MPI_T_FINALIZE are permissible.

Once MPI_T_FINALIZE is called the same number of times as the routine MPI_T_INIT_THREAD up to the current point of execution, the MPI tool information interface is no longer initialized. The interface can be reinitialized by subsequent calls to MPI_T_INIT_THREAD.

At the end of the program execution, unless MPI_ABORT is called, an application must have called MPI_T_INIT_THREAD and MPI_T_FINALIZE an equal number of times.

14.3.5 Datatype System

All variables managed through the MPI tool information interface represent their values through typed buffers of a given length and type using an MPI datatype (similar to regular send/receive buffers). Since the initialization of the MPI tool information interface is separate from the initialization of MPI, MPI tool information interface routines can be called before MPI_INIT. Consequently, these routines can also use MPI datatypes before MPI_INIT. Therefore, within the context of the MPI tool information interface, it is permissible to use a subset of MPI datatypes as specified below before a call to MPI_INIT (or equivalent).

MPI_INT
MPI_UNSIGNED
MPI_UNSIGNED_LONG
MPI_UNSIGNED_LONG_LONG
MPI_COUNT
MPI_CHAR
MPI_DOUBLE

Table 14.3: MPI datatypes that can be used by the MPI tool information interface

Rationale. The MPI tool information interface relies mainly on unsigned datatypes for integer values since most variables are expected to represent counters or resource sizes. MPI_INT is provided for additional flexibility and is expected to be used mainly for control variables and enumeration types (see below).

Providing all basic datatypes, in particular providing all signed and unsigned variants of integer types, would lead to a larger number of types, which tools need to interpret. This would cause unnecessary complexity in the implementation of tools based on the MPI tool information interface. (*End of rationale.*)

The MPI tool information interface only relies on a subset of the basic MPI datatypes and does not use any derived MPI datatypes. Table 14.3 lists all MPI datatypes that can be returned by the MPI tool information interface to represent its variables.

The use of the datatype MPI_CHAR in the MPI tool information interface implies a nullterminated character array, i.e., a string in the C language. If a variable has type MPI_CHAR, the value of the count parameter returned by MPI_T_CVAR_HANDLE_ALLOC and MPI_T_PVAR_HANDLE_ALLOC must be large enough to include any valid value, including its terminating null character. The contents of returned MPI_CHAR arrays are only defined from index 0 through the location of the first null character.

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1 2 3 4 5	<i>Rationale.</i> The MPI tool information interface requires a significantly simpler type system than MPI itself. Therefore, only its required subset must be present before MPI_INIT (or equivalent) and MPI implementations do not need to initialize the complete MPI datatype system. (<i>End of rationale.</i>)					
6 7 8 9 10 11 12 13 14 15	For variables of type MPI_INT, an MPI implementation can provide additional information by associating names with a fixed number of values. We refer to this information in the following as an enumeration. In this case, the respective calls that provide additional metadata for each control or performance variable, i.e., MPI_T_CVAR_GET_INFO (Section 14.3.6) and MPI_T_PVAR_GET_INFO (Section 14.3.7), return a handle of type MPI_T_enum that can be passed to the following functions to extract additional information. Thus, the MPI implementation can describe variables with a fixed set of values that each represents a particular state. Each enumeration type can have N different values, with a fixed N that can be queried using MPI_T_ENUM_GET_INFO.					
16 17	MPI_T_EN	IUM_GET_INFO(enumtype, nu	ım, name, name_len)			
18	IN	enumtype	enumeration to be queried (handle)			
19 20	OUT	num	number of discrete values represented by this enumer- ation (integer)			
21 22 23	OUT	name	buffer to return the string containing the name of the enumeration (string)			
24 25	INOUT	name_len	length of the string and/or buffer for $name\xspace$ (integer)			
25 26 27	int MPI_T	_enum_get_info(MPI_T_enum *name_len)	n enumtype, int *num, char *name, int			
28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48	<pre>*name_len) If enumtype is a valid enumeration, this routine returns the number of items represented by this enumeration type as well as its name. N must be greater than 0, i.e., the enumeration must represent at least one value. The arguments name and name_len are used to return the name of the enumeration as described in Section 14.3.3. The routine is required to return a name of at least length one. This name must be unique with respect to all other names for enumerations that the MPI implementation uses. Names associated with individual values in each enumeration enumtype can be queried using MPI_T_ENUM_GET_ITEM.</pre>					

1 0			icx, value, name, name_icn)	
	IN	enumtype	enumeration to be queried (handle)	2
	IN	index	number of the value to be queried in this enumeration (integer)	4
	OUT	value	variable value (integer)	e
	OUT	name	buffer to return the string containing the name of the enumeration item (string)	7
	INOUT	name_len	length of the string and/or buffer for $name\xspace$ (integer)	1
				1

MPI_T_ENUM_GET_ITEM(enumtype, index, value, name, name_len)

The arguments name and name_len are used to return the name of the enumeration item as described in Section 14.3.3.

If completed successfully, the routine returns the name/value pair that describes the enumeration at the specified index. The call is further required to return a name of at least length one. This name must be unique with respect to all other names of items for the same enumeration.

14.3.6 Control Variables

The routines described in this section of the MPI tool information interface specification focus on the ability to list, query, and possibly set control variables exposed by the MPI implementation. These variables can typically be used by the user to fine tune properties and configuration settings of the MPI implementation. On many systems, such variables can be set using environment variables, although other configuration mechanisms may be available, such as configuration files or central configuration registries. A typical example that is available in several existing MPI implementations is the ability to specify an "eager limit," i.e., an upper bound on the size of messages sent or received using an eager protocol.

Control Variable Query Functions

An MPI implementation exports a set of N control variables through the MPI tool information interface. If N is zero, then the MPI implementation does not export any control variables, otherwise the provided control variables are indexed from 0 to N-1. This index number is used in subsequent calls to identify the individual variables.

An MPI implementation is allowed to increase the number of control variables during the execution of an MPI application when new variables become available through dynamic loading. However, MPI implementations are not allowed to change the index of a control variable or to delete a variable once it has been added to the set. When a variable becomes inactive, e.g., through dynamic unloading, accessing its value should return a corresponding error code.

Advice to users. While the MPI tool information interface guarantees that indices or variable properties do not change during a particular run of an MPI program, it does not provide a similar guarantee between runs. (*End of advice to users.*)

	$\frac{1}{2}$	The fo	ollowing function ca	an be used to query the number of control variables, num_cvar :		
	3	MPI_T_C\	/AR_GET_NUM(nu	ım_cvar)		
	4 5	OUT	num_cvar	returns number of control variables (integer)		
	6 7	int MPI_7	[_cvar_get_num(i	nt *num_cvar)		
	8 9	The f each varia		$/AR_GET_INFO$ provides access to additional information for		
	10 11					
	12 13	MPI_T_C\	AR_GET_INFO(cv/ desc_len, bind	ar_index, name, name_len, verbosity, datatype, enumtype, desc, d, scope)		
	14 15 16	IN	cvar_index	index of the control variable to be queried, value be- tween 0 and $num_cvar - 1$ (integer)		
	17 18	OUT	name	buffer to return the string containing the name of the control variable (string)		
	19	INOUT	name_len	length of the string and/or buffer for name (integer)		
	20 21	OUT	verbosity	verbosity level of this variable (integer)		
	21 22 23	OUT	datatype	MPI datatype of the information stored in the control variable (handle)		
	24 25	OUT	enumtype	optional descriptor for enumeration information (han- dle)		
	26 27	OUT	desc	buffer to return the string containing a description of the control variable (string)		
	28 29	INOUT	desc_len	length of the string and/or buffer for desc (integer)		
	30 31	OUT	bind	type of MPI object to which this variable must be bound (integer)		
	32 33 34	OUT	scope	scope of when changes to this variable are possible (integer)		
	35 36 37 38	int MPI_7	<pre>*verbosity,</pre>	<pre>int cvar_index, char *name, int *name_len, int MPI_Datatype *datatype, MPI_T_enum *enumtype, char *desc_len, int *bind, int *scope)</pre>		
	39 40	After a successful call to MPI_T_CVAR_GET_INFO for a particular variable, subsequent calls to this routine that query information about the same variable must return the same				
ticket378.	41	informatio	n. An MPI implem	nentation is not allowed to alter any of the returned values.		
	42	If any OUT parameter to MPI_T_CVAR_GET_INFO is a NULL pointer, the implemen-				
	43		~ *	eter and not return a value for the parameter.		
	44 45		ed in Section 14.3.3	d name_len are used to return the name of the control variable		
	46			y, the routine is required to return a name of at least length		
	47			que with respect to all other names for control variables used		
	48	by the MPI implementation.				

The argument verbosity returns the verbosity level of the variable (see Section 14.3.1).

The argument datatype returns the MPI datatype that is used to represent the control variable.

If the variable is of type MPI_INT, MPI can optionally specify an enumeration for the values represented by this variable and return it in enumtype. In this case, MPI returns an enumeration identifier, which can then be used to gather more information as described in Section 14.3.5. Otherwise, enumtype is set to MPI_T_ENUM_NULL. If the datatype is not MPI_INT or the argument enumtype is the null pointer, no enumeration type is returned.

The arguments desc and desc_len are used to return a description of the control variable as described in Section 14.3.3.

Returning a description is optional. If an MPI implementation does not to return a description, the first character for desc must be set to the null character and desc_len must be set to one at the return of this call.

The parameter bind returns the type of the MPI object to which the variable must be bound or the value MPI_T_BIND_NO_OBJECT (see Section 14.3.2).

The scope of a variable determines whether changing a variable's value is either local to the process or must be done by the user across multiple processes. The latter is further split into variables that require changes in a group of processes and those that require collective changes among all connected processes. Both cases can require all processes either to be set to consistent (but potentially different) values or to equal values on every participating process. The description provided with the variable must contain an explanation about the requirements and/or restrictions for setting the particular variable.

On successful return from MPI_T_CVAR_GET_INFO, the argument scope will be set to one of the constants listed in Table 14.4.

If the name of a control variable is the same across connected processes, the following OUT parameters must be identical: verbosity, datatype, enumtype, desc, desc_len, bind, scope.

Scope Constant	Description
MPI_T_SCOPE_CONSTANT	read-only, value is constant
MPI_T_SCOPE_READONLY	read-only, cannot be written, but can change
MPI_T_SCOPE_LOCAL	may be writeable, writing is a local operation
MPI_T_SCOPE_GROUP	may be writeable, must be done to a group of processes,
	all processes in a group must be set to consistent values
MPI_T_SCOPE_GROUP_EQ	may be writeable, must be done to a group of processes,
	all processes in a group must be set to the same value
MPI_T_SCOPE_ALL	may be writeable, must be done to all processes,
	all connected processes must be set to consistent values
MPI_T_SCOPE_ALL_EQ	may be writeable, must be done to all processes,
	all connected processes must be set to the same value

Table 14.4: Scopes for control variables

Advice to users. The scope of a variable only indicates if a variable might be changeable; it is not a guarantee that it can be changed at any time. (*End of advice to users.*)

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²⁴ ticket383.

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1 2 3 MPI_T_CVAR_GET_INDEX(name, cvar_index) 4 IN name name of the control variable (string) 56 OUT cvar_index index of the control variable (integer) 7 8 int MPI_T_cvar_get_index(const char *name, int *cvar_index) 9 MPI_T_CVAR_GET_INDEX is a function for retrieving the index of a control variable 10 given a known variable name. The **name** parameter is provided by the caller, and **cvar_inde**x 11 is returned by the MPI implementation. The name parameter is a string terminated with a 12 null character. 13 This routine returns MPI_SUCCESS on success and returns MPI_T_ERR_INVALID_NAME 14 if name does not match the name of any control variable provided by the implementation 15 at the time of the call. 16 17This routine is provided to enable fast retrieval of control variables by Rationale. 18 a tool, assuming it knows the name of the variable for which it is looking. The 19 number of variables exposed by the implementation can change over time, so it is not 20possible for the tool to simply iterate over the list of variables once at initialization. 21Although using MPI implementation specific variable names is not portable across MPI 22 implementations, tool developers may choose to take this route for lower overhead at 23 runtime because the tool will not have to iterate over the entire set of variables to 24 find a specific one. (*End of rationale.*) 2526 27 Example: Printing All Control Variables 2829 30 Example 14.4 31 The following example shows how the MPI tool information interface can be used to 32query and to print the names of all available control variables. 33 34 #include <stdio.h> #include <stdlib.h> 35 #include <mpi.h> 36 37 int main(int argc, char *argv[]) { 38 int i, err, num, namelen, bind, verbose, scope; 39 int threadsupport; 40 char name[100]; 41 42 MPI_Datatype datatype; 43err=MPI_T_init_thread(MPI_THREAD_SINGLE,&threadsupport); 44 if (err!=MPI_SUCCESS) 45 return err; 46 47err=MPI_T_cvar_get_num(&num); 48

```
if (err!=MPI_SUCCESS)
    return err;
  for (i=0; i<num; i++) {</pre>
    namelen=100;
    err=MPI_T_cvar_get_info(i, name, &namelen,
            &verbose, &datatype, NULL,
            NULL, NULL, /*no description */
            &bind, &scope);
    if (err!=MPI_SUCCESS || err!=MPI_T_ERR_INVALID_INDEX) return err;
    printf("Var %i: %s\n", i, name);
  }
  err=MPI_T_finalize();
  if (err!=MPI_SUCCESS)
    return 1;
  else
    return 0;
}
```

Handle Allocation and Deallocation

Before reading or writing the value of a variable, a user must first allocate a handle of type MPI_T_cvar_handle for the variable by binding it to an MPI object (see also Section 14.3.2).

Rationale. Handles used in the MPI tool information interface are distinct from handles used in the remaining parts of the MPI standard because they must be usable before MPI_INIT and after MPI_FINALIZE. Further, accessing handles, in particular for performance variables, can be time critical and having a separate handle space enables optimizations. (*End of rationale.*)

MPI_T_CVAR_HANDLE_ALLOC(cvar_index, obj_handle, handle, count)

IN	cvar_index	index of control variable for which handle is to be al-	34
		located (index)	35
			36
IN	obj_handle	reference to a handle of the MPI object to which this	37
		variable is supposed to be bound (pointer)	38
OUT	handle	allocated handle (handle)	39
OUT	count	number of elements used to represent this variable (in-	40
001	count	- · · · · · · · · · · · · · · · · · · ·	41
		$\operatorname{teger})$	42
			42

This routine binds the control variable specified by the argument index to an MPI object. ⁴⁶ The object is passed in the argument obj_handle as an address to a local variable that stores ⁴⁷ the object's handle. The argument obj_handle is ignored if the MPI_T_CVAR_GET_INFO ⁴⁸

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1				PI_T_BIND_NO_OBJECT in the argument bind. The		
2				ble is returned in the argument		
3	handle. Upon successful return, count contains the number of elements (of the datatype					
4	returned by a previous $MPI_T_CVAR_GET_INFO$ call) used to represent this variable.					
5						
6	Advi	tce to users. T	he count can l	be different based on the MPI object to which the		
7	contr	rol variable was	bound. For e	example, variables bound to communicators could		
8	have	a count that m	atches the size	e of the communicator.		
9	It is	not portable to	pass reference	es to predefined MPI object handles, such as		
10		-	-	e, since their implementation depends on the MPI		
11				dles should be stored in a local variable and the		
12		e ,		d be passed into MPI_T_CVAR_HANDLE_ALLOC.		
13		l of advice to us		u be passed into with_t_cvan_trandel_alloc.		
14	(Enc	i of advice to us	ers.)			
15	The v	alue of over ind	ev should be i	n the range 0 to $num_cvar - 1$, where num_cvar		
16				bles as determined from a prior call to		
17				the MPI object it references must be consistent		
18			0 x	nent in a prior call to MPI_T_CVAR_GET_INFO.		
19			0	returned by MPI_T_CVAR_GET_INFO equals		
20			0	obj_handle is ignored.		
20		ID_NO_OBJECT,	the argument	obj_nancie is ignored.		
21						
23	MPI_T_C\	/AR_HANDLE_F	REE(handle)			
24	INOUT	handle	× ,	handle to be freed (handle)		
25		nanue		handle to be need (handle)		
26	int MDT 7	r arran handla	fmaa(MDT T	cvar_handle *handle)		
27	IIIC MFI_I	cvar_nanute_	_1166(WF1_1_0			
28	When	a handle is no l	onger needed,	a user of the MPI tool information interface should		
29	call MPI_T_CVAR_HANDLE_FREE to free the handle and the associated resources in the					
30	MPI imple	ementation. Or	n a successful	return, MPI sets the handle to		
31	MPI_T_CV	AR_HANDLE_NU	LL.			
32						
33	Control Va	riable Access Fu	nctions			
34						
35						
36		/AR_READ(hand	lle buf)			
37			ale, bul)			
38	IN	handle		handle to the control variable to be read (handle)		
39	OUT	buf		initial address of storage location for variable value		
40				(choice)		
41						
42	int MPI_7	C_cvar_read(MF	PI_T_cvar_ham	ndle handle, void* buf)		
43						
44		-		ntrol variable identified by the argument handle and		
45				y the parameter buf . The user must ensure that the		
46				ne entire value of the control variable (based on the		
47			-	corresponding calls to MPI_T_CVAR_GET_INFO		
48	and MPI_	T_CVAR_HAND	LE_ALLOC, re	spectively).		

MPI_T_0	CVAR_WRITE(handle, buf)		1
IN	handle	handle to the control variable to be written (handle)	2
IN	buf	initial address of storage location for variable value (choice)	3 4 5
int MPI	_T_cvar_write(MPI_T_cvar_h	andle handle, const void* buf)	6 7
the data buffer is returned and MPI If th MPI_T_C in all co scope, an the group MPI_T_C In b the scope variable If it returns e	stored in the buffer identified b of the appropriate size to hold t datatype and count from prior _T_CVAR_HANDLE_ALLOC, re- te variable has a global scope CVAR_GET_INFO call), any wr nnected (as defined in Section my write call to this variable m o, which must be described by CVAR_GET_INFO. oth cases, the user must ensur e is either MPI_T_SCOPE_ALL_E in all processes must be set to is not possible to change the ither MPI_T_ERR_CVAR_SET_N	(as returned by a prior corresponding ite call to this variable must be issued by the user 10.5.4) MPI processes. If the variable has group nust be issued by the user in all MPI processes in the MPI implementation in the description by the e that the writes in all processes are consistent. If Q or MPI_T_SCOPE_GROUP_EQ this means that the the same value. variable at the time the call is made, the function OT_NOW, if there may be a later time at which the AR_SET_NEVER, if the variable cannot be set for the	8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25
	Reading the Value of a Control		26 27
Exampl The control v	e 14.5 following example shows a ro	utine that can be used to query the value with a e example assumes that the variable is intended to	28 29 30 31 32 33
int err		MPI_Comm comm, int *val) {	34 35 36 37
	example assumes that the be bound to a communicator		38 39 40
	T_cvar_handle_alloc(index =MPI_SUCCESS) return err;	,&comm,&handle,&count);	41 42 43
	following assumes that the esented by a single intege		44 45 46
err=MPI	T_cvar_read(handle,val);		47 48

```
if (err!=MPI_SUCCESS) return err;
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     err=MPI_T_cvar_handle_free(&handle);
     return err;
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     }
\mathbf{5}
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14.3.7

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Performance Variables

The following section focuses on the ability to list and to query performance variables provided by the MPI implementation. Performance variables provide insight into MPI im-10 plementation specific internals and can represent information such as the state of the MPI 11 implementation (e.g., waiting blocked, receiving, not active), aggregated timing data for 12submodules, or queue sizes and lengths. 13

Rationale. The interface for performance variables is separate from the interface for control variables, since performance variables have different requirements and parameters. By keeping them separate, the interface provides cleaner semantics and allows for more performance optimization opportunities. (*End of rationale.*)

19 Performance Variable Classes 20

21Each performance variable is associated with a class that describes its basic semantics, 22possible datatypes, basic behavior, its starting value, whether it can overflow, and when 23and how an MPI implementation can change the variable's value. The starting value is the 24value that is assigned to the variable the first time that it is used or whenever it is reset. 25

If a performance variable belongs to a class that can overflow, Advice to users. it is up to the user to protect against this overflow, e.g., by frequently reading and resetting the variable value. (End of advice to users.)

Advice to implementors. MPI implementations should use large enough datatypes for each performance variable to avoid overflows under normal circumstances. (End of advice to implementors.)

The classes are defined by the following constants:

MPI_T_PVAR_CLASS_STATE

A performance variable in this class represents a set of discrete states. Variables of this class are represented by MPI_INT and can be set by the MPI implementation at any time. Variables of this type should be described further using an enumeration, as discussed in Section 14.3.5. The starting value is the current state of the implementation at the time that the starting value is set. MPI implementations must ensure that variables of this class cannot overflow.

MPI_T_PVAR_CLASS_LEVEL

A performance variable in this class represents a value that describes the utilization 44 level of a resource. The value of a variable of this class can change at any time to match 45 the current utilization level of the resource. Values returned from variables in this class 46 are non-negative and represented by one of the following datatypes: MPI_UNSIGNED, 47MPI_UNSIGNED_LONG, MPI_UNSIGNED_LONG_LONG, MPI_DOUBLE. The starting value 48

is the current utilization level of the resource at the time that the starting value is set. MPI implementations must ensure that variables of this class cannot overflow.

• MPI_T_PVAR_CLASS_SIZE

A performance variable in this class represents a value that is the [fixed]size of a resource. Values returned from variables in this class are non-negative and represented by one of the following datatypes: MPI_UNSIGNED, MPI_UNSIGNED_LONG, MPI_UNSIGNED_LONG, MPI_DOUBLE. The starting value is the current [utilization level]size of the resource at the time that the starting value is set. MPI implementations must ensure that variables of this class cannot overflow.

MPI_T_PVAR_CLASS_PERCENTAGE

The value of a performance variable in this class represents the percentage utilization of a finite resource. The value of a variable of this class can change at any time to match the current utilization level of the resource. It will be returned as an MPI_DOUBLE datatype. The value must always be between 0.0 (resource not used at all) and 1.0 (resource completely used). The starting value is the current percentage utilization level of the resource at the time that the starting value is set. MPI implementations must ensure that variables of this class cannot overflow.

MPI_T_PVAR_CLASS_HIGHWATERMARK

A performance variable in this class represents a value that describes the high watermark utilization of a resource. The value of a variable of this class is non-negative and grows monotonically from the initialization or reset of the variable. It can be represented by one of the following datatypes: MPI_UNSIGNED, MPI_UNSIGNED_LONG, MPI_UNSIGNED_LONG_LONG, MPI_DOUBLE. The starting value is the current utilization level of the resource at the time that the starting value is set. MPI implementations must ensure that variables of this class cannot overflow.

MPI_T_PVAR_CLASS_LOWWATERMARK

A performance variable in this class represents a value that describes the low watermark utilization of a resource. The value of a variable of this class is non-negative and decreases monotonically from the initialization or reset of the variable. It can be represented by one of the following datatypes: MPI_UNSIGNED, MPI_UNSIGNED_LONG, MPI_UNSIGNED_LONG_LONG, MPI_DOUBLE. The starting value is the current utilization level of the resource at the time that the starting value is set. MPI implementations must ensure that variables of this class cannot overflow.

MPI_T_PVAR_CLASS_COUNTER

A performance variable in this class counts the number of occurrences of a specific event (e.g., the number of memory allocations within an MPI library). The value of a variable of this class increases monotonically from the initialization or reset of the performance variable by one for each specific event that is observed. Values must be non-negative and represented by one of the following datatypes: MPI_UNSIGNED, MPI_UNSIGNED_LONG, MPI_UNSIGNED_LONG_LONG. The starting value for variables of this class is 0. Variables of this class can overflow.

MPI_T_PVAR_CLASS_AGGREGATE

The value of a performance variable in this class is an an aggregated value that ⁴⁷ represents a sum of arguments processed during a specific event (e.g., the amount ⁴⁸

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of memory allocated by all memory allocations). This class is similar to the counter class, but instead of counting individual events, the value can be incremented by arbitrary amounts. The value of a variable of this class increases monotonically from the initialization or reset of the performance variable. It must be non-negative and represented by one of the following datatypes: MPI_UNSIGNED, MPI_UNSIGNED_LONG, MPI_UNSIGNED_LONG_LONG, MPI_DOUBLE. The starting value for variables of this class is 0. Variables of this class can overflow.

MPI_T_PVAR_CLASS_TIMER

The value of a performance variable in this class represents the aggregated time that the MPI implementation spends executing a particular event, type of event, or section of the MPI library. This class has the same basic semantics as MPI_T_PVAR_CLASS_AGGREGATE, but explicitly records a timing value. The value of a variable of this class increases monotonically from the initialization or reset of the performance variable. It must be non-negative and represented by one of the following datatypes: MPI_UNSIGNED, MPI_UNSIGNED_LONG, MPI_UNSIGNED_LONG_LONG, MPI_DOUBLE. The starting value for variables of this class is 0. If the type MPI_DOUBLE is used, the units that represent time in this datatype must match the units used by MPI_WTIME. Otherwise, the time units should be documented, e.g., in the description returned by MPI_T_PVAR_GET_INFO. Variables of this class can overflow.

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MPI_T_PVAR_CLASS_GENERIC

This class can be used to describe a variable that does not fit into any of the other classes. For variables in this class, the starting value is variable-specific and implementation-defined.

27Performance Variable Query Functions 28

An MPI implementation exports a set of N performance variables through the MPI tool 29 information interface. If N is zero, then the MPI implementation does not export any 30 performance variables; otherwise the provided performance variables are indexed from 0 to 31 N-1. This index number is used in subsequent calls to identify the individual variables. 32

An MPI implementation is allowed to increase the number of performance variables 33 during the execution of an MPI application when new variables become available through 34 dynamic loading. However, MPI implementations are not allowed to change the index of 35 a performance variable or to delete a variable once it has been added to the set. When 36 a variable becomes inactive, e.g., through dynamic unloading, accessing its value should 37 return a corresponding error code. 38

The following function can be used to query the number of performance variables, N:

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MPI_T_PVAR_GET_NUM(num_pvar) num_pvar

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OUT

returns number of performance variables (integer)

45int MPI_T_pvar_get_num(int *num_pvar)

The function MPI_T_PVAR_GET_INFO provides access to additional information for 47each variable. 48

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MPI_T_PVAR_GET_INFO(pvar_index, name, name_len, verbosity, varclass, datatype	,
enumtype, desc, desc_len, bind, readonly, continuous, atomic)	

IN	pvar_index	index of the performance variable to be queried be-	3
		tween 0 and $num_pvar - 1$ (integer)	4
OUT	name	buffer to return the string containing the name of the	5 6
		performance variable (string)	7
INOUT	name_len	length of the string and/or buffer for $name\xspace$ (integer)	8
OUT	verbosity	verbosity level of this variable (integer)	9
OUT	var_class	class of performance variable (integer)	10 11
OUT	datatype	MPI datatype of the information stored in the perfor-	11
		mance variable (handle)	13
OUT	enumtype	optional descriptor for enumeration information (han-	14
	51	dle)	15
OUT	desc	buffer to return the string containing a description of	16 17
		the performance variable (string)	17
INOUT	desc_len	length of the string and/or buffer for desc (integer)	19
OUT	bind	type of MPI object to which this variable must be	20
		bound (integer)	21
OUT	readonly	flag indicating whether the variable can be	22
	, , , , , , , , , , , , , , , , , , ,	written/reset (integer)	23 24
OUT	continuous	flag indicating whether the variable can be started and	25
		stopped or is continuously active (integer)	26
OUT	atomic	flag indicating whether the variable can be atomically	27
		read and reset (integer)	28

After a successful call to MPI_T_PVAR_GET_INFO for a particular variable, subsequent calls to this routine that query information about the same variable must return the same information. An MPI implementation is not allowed to alter any of the returned values.

If any OUT parameter to MPI_T_PVAR_GET_INFO is a NULL pointer, the implementation will ignore the parameter and not return a value for the parameter.

The arguments name and name_len are used to return the name of the performance variable as described in Section 14.3.3. If completed successfully, the routine is required to return a name of at least length one.

The argument verbosity returns the verbosity level of the variable (see Section 14.3.1).

The class of the performance variable is returned in the parameter var_class. The class must be one of the constants defined in Section 14.3.7.

The combination of the name and the class of the performance variable must be unique with respect to all other names for performance variables used by the MPI implementation.

Unofficial Draft for Comment Only

³⁷ ticket378.

Advice to implementors. Groups of variables that belong closely together, but have 1 different classes, can have the same name. This choice is useful, e.g., to refer to 2 3 multiple variables that describe a single resource (like the level, the total size, as well as high and low watermarks). (End of advice to implementors.) 4 $\mathbf{5}$ The argument datatype returns the MPI datatype that is used to represent the perfor-6 mance variable. 7 If the variable is of type MPI_INT, MPI can optionally specify an enumeration for the 8 values represented by this variable and return it in enumtype. In this case, MPI returns an 9 enumeration identifier, which can then be used to gather more information as described in 10 Section 14.3.5. Otherwise, enumtype is set to MPI_T_ENUM_NULL. If the datatype is not 11 ticket0. 12 MPI_INT or the argument enumtype is the null pointer, no [enumeration] enumeration type is returned. 13 ticket0. 14 Returning a description is optional. If an MPI implementation does not to return a description, the first character for desc must be set to the null character and desc_len must 15 be set to one at the return from this function. 16 The parameter bind returns the type of the MPI object to which the variable must be 17bound or the value MPI_T_BIND_NO_OBJECT (see Section 14.3.2). 18 Upon return, the argument readonly is set to zero if the variable can be written or reset 19 by the user. It is set to one if the variable can only be read. 20Upon return, the argument continuous is set to zero if the variable can be started and 21stopped by the user, i.e., it is possible for the user to control if and when the value of a 22variable is updated. It is set to one if the variable is always active and cannot be controlled 23 by the user. 24 Upon return, the argument atomic is set to zero if the variable cannot be read and 25reset atomically. Only variables for which the call sets atomic to one can be used in a call 26ticket383. 27 to MPI_T_PVAR_READRESET. If the name and class of a performance variable are the same across connected processes, 28 the following OUT parameters must be identical: verbosity, varclass, datatype, enumtype, 29 ticket377. 30 desc, desc_len, bind, readonly, continuous, atomic. 31 32MPI_T_PVAR_GET_INDEX(name, var_class, pvar_index) 33 IN the name of the performance variable (string) name 34 35IN var_class the class of the performance variable (integer) 36 OUT pvar_index the index of the performance variable (integer) 37 38 int MPI_T_pvar_get_index(const char *name, int var_class, int *pvar_index) 39 40 MPI_T_PVAR_GET_INDEX is a function for retrieving the index of a performance 41variable given a known variable name and class. The name and var_class parameters are 42provided by the caller, and pvar_index is returned by the MPI implementation. The name 43parameter is string terminated with a null character. 44 This routine returns MPI_SUCCESS on success and returns MPI_T_ERR_INVALID_NAME if 45name does not match the name of any performance variable provided by the implementation 46 at the time of the call. 4748

Rationale. This routine is provided to enable fast retrieval of performance variables by a tool, assuming it knows the name of the variable for which it is looking. The number of variables exposed by the implementation can change over time, so it is not possible for the tool to simply iterate over the list of variables once at initialization. Although using MPI implementation specific variable names is not portable across MPI implementations, tool developers may choose to take this route for lower overhead at runtime because the tool will not have to iterate over the entire set of variables to find a specific one. (*End of rationale.*)

Performance Experiment Sessions

Within a single program, multiple components can use the MPI tool information interface. To avoid collisions with respect to accesses to performance variables, users of the MPI tool information interface must first create a session. Subsequent calls that access performance variables can then be made within the context of this session. Any call executed in a session must not influence the results in any other session.

MPI_T_F	VAR_SESSION_C	REATE(session)	19
OUT	session	identifier of performance session (handle)	20
		-	21
int MPI_	T_pvar_session_	_create(MPI_T_pvar_session *session)	22 23
Thia	call encoted a new	anging for according performance workships and returns a handle	24
		session for accessing performance variables and returns a handle nent session of type MPI_T_pvar_session.	25
101 01115 5	ession in the argui		26
			27
MPI_T_F	VAR_SESSION_FI	REE(session)	28
INOUT	session	identifier of performance experiment session (handle)	29
		, , , ,	30
int MPI_	T_pvar_session_	_free(MPI_T_pvar_session *session)	31
	-	-	32
		ing session. Calls to the MPI tool information interface can no	33
0		context of a session after it is freed. On a successful return, MPI	34 35
sets the s	session identifier to	MPI_T_PVAR_SESSION_NULL.	36
Handla A	llocation and Deall	acation	37
папше А	nocation and Deali		38
	· ·	e variable, a user must first allocate a handle of type	39
MPI_T_p	var_handle for the	variable by binding it to an MPI object (see also Section 14.3.2).	40
			41
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1	MPI_T_P\	/AR_HANDLE_ALLOC(session, pvar_index, obj_handle, handle, count)
2	IN	session	identifier of performance experiment session (handle)
3 4 5	IN	pvar_index	index of performance variable for which handle is to be allocated (integer)
6 7	IN	obj_handle	reference to a handle of the MPI object to which this variable is supposed to be bound (pointer)
8	001	handle	allocated handle (handle)
9	OUT	count	number of elements used to represent this variable (in-
11			teger)
12		·	MDT T
13 14		-	<pre>MPI_T_pvar_session session, int pvar_index, le, MPI_T_pvar_handle *handle, int *count)</pre>
15	1 nis 1	-	ormance variable specified by the argument index to an
16	. MPT objec		ed by the parameter session. The object is passed in the
18	argument	-	ss to a local variable that stores the object's handle. The f the MPI_T_PVAR_GET_INFO call for this performance
19		•	D_OBJECT in the argument bind. The handle allocated to
20	reference t		l in the argument handle. Upon successful return, count
21 22	contains th		o (of the datatype returned by a previous
23		/AR_GET_INFO call) us	sed to represent this variable.
24	Advi	ce to users. The cour	at can be different based on the MPI object to which the
25	-		bound. For example, variables bound to communicators
26		d have a count that ma	tches the size of the communicator.
27 28	It is		eferences to predefined MPI object handles, such as
29	MPI_	,	s routine, since their implementation depends on the MPI pject handle should be stored in a local variable and the
30		e ,	e should be passed into MPI_T_PVAR_HANDLE_ALLOC.
31	(End	l of advice to users.)	
32			
34	ine v		e in the range 0 to $num_pvar - 1$, where num_pvar is the variables as determined from a prior call to
35		-	ype of the MPI object it references must be consistent
ticket0. $\frac{36}{37}$	with the ty	ype returned in the bine	d argument in a prior call to MPI_T_PVAR_GET_INFO. ent equals MPI_T_BIND_NO_OBJECT, the argument
ticket386. ³⁸		is ignored.]	
39	FOT a		of this section that take both handle and session as IN
40			nt passed in is not associated with the session argument,
42		R_INVALID_HANDLE is 1	eturned.
43			
44		/AR_HANDLE_FREE(se	ssion, handle)
45 46	IIN	session	identifier of performance experiment session (handle)
40	INOUT	handle	handle to be freed (handle)
48	3		

When a handle is no longer needed, a user of the MPI tool information interface should call MPI_T_PVAR_HANDLE_FREE to free the handle in the session identified by the parameter session and the associated resources in the MPI implementation. On a successful return, MPI sets the handle to MPI_T_PVAR_HANDLE_NULL.

Starting and Stopping of Performance Variables

Performance variables that have the continuous flag set during the query operation are continuously operating once a handle has been allocated. Such variables may be queried at any time, but they cannot be started or stopped by the user. All other variables are in a stopped state after their handle has been allocated; their values are not updated until they have been started by the user.

MPI_T_PVAR_START(session, handle)

IN	session	identifier of performance experiment session (handle)
IN	handle	handle of a performance variable (handle)

int MPI_T_pvar_start(MPI_T_pvar_session session, MPI_T_pvar_handle handle)

This functions starts the performance variable with the handle identified by the parameter handle in the session identified by the parameter session.

If the constant MPI_T_PVAR_ALL_HANDLES is passed in handle, the MPI implementation attempts to start all variables within the session identified by the parameter session for which handles have been allocated. In this case, the routine returns MPI_SUCCESS if all variables are started successfully (even if there are no non-continuous variables to be started), otherwise MPI_T_ERR_PVAR_NO_STARTSTOP is returned. Continuous variables and variables that are already started are ignored when MPI_T_PVAR_ALL_HANDLES is specified.

MPI_T_PVAR_STOP(session, handle)

IN	session	identifier of performance experiment session (handle)
IN	handle	handle of a performance variable (handle)

int MPI_T_pvar_stop(MPI_T_pvar_session session, MPI_T_pvar_handle handle)

This functions stops the performance variable with the handle identified by the parameter handle in the session identified by the parameter session.

If the constant MPI_T_PVAR_ALL_HANDLES is passed in handle, the MPI implementation attempts to stop all variables within the session identified by the parameter session for which handles have been allocated. In this case, the routine returns MPI_SUCCESS if all variables are stopped successfully (even if there are no non-continuous variables to be started), otherwise MPI_T_ERR_PVAR_NO_STARTSTOP is returned. Continuous variables and variables that are already stopped are ignored when MPI_T_PVAR_ALL_HANDLES is specified.

ticket391.

²⁸ ticket391.

1 2	Performan	ce Variable Access	Functions	
3 4	MPI_T_P	VAR_READ(sessior	ı, handle, buf)	
5	IN	session	identifier of performance experiment session (handle)	
6 7	IN	handle	handle of a performance variable (handle)	
8 9 10	OUT	buf	initial address of storage location for variable value (choice)	
11 12 13	int MPI_	T_pvar_read(MPI void* buf)	_T_pvar_session session, MPI_T_pvar_handle handle,	
14 15 16 17 18 19 20 21 22 22 23	handle ha buffer ide is of the a the dataty MPI_T_P The o tion MPI_	ndle in the session ntified by the para appropriate size to ype and count ret VAR_GET_INFO at constant MPI_T_PV .T_PVAR_READ.	AD call queries the value of the performance variable with the identified by the parameter session and stores the result in the ameter buf . The user is responsible to ensure that the buffer behold the entire value of the performance variable (based on surned by the corresponding previous calls to and MPI_T_PVAR_HANDLE_ALLOC, respectively).	
24	MPI_T_P	VAR_WRITE(session	on,handle, buf)	
25	IN	session	identifier of performance experiment session (handle)	
26 27	IN	handle	handle of a performance variable (handle)	
28 29	IN	buf	initial address of storage location for variable value (choice)	
30 31 32	int MPI_	T_pvar_write(MP const void	I_T_pvar_session session, MPI_T_pvar_handle handle, * buf)	
333 34 35 36 37 38 338 40 41 41 42 43 44 45 46 44	The MPI_T_PVAR_WRITE call attempts to write the value of the performance variable with the handle identified by the parameter handle in the session identified by the parameter session. The value to be written is passed in the buffer identified by the parameter buf. The user must ensure that the buffer is of the appropriate size to hold the entire value of the per- formance variable (based on the datatype and count returned by the corresponding previous calls to MPI_T_PVAR_GET_INFO and MPI_T_PVAR_HANDLE_ALLOC, respectively). If it is not possible to change the variable, the function returns MPI_T_ERR_PVAR_NO_WRITE. The constant MPI_T_PVAR_ALL_HANDLES cannot be used as an argument for the func- tion MPI_T_PVAR_WRITE.			

	VAR_RESET(session		2
IN	session	identifier of performance experiment session (handle)	3
IN	handle	handle of a performance variable (handle)	4
			5
int MPI_	T_pvar_reset(MPI_	T_pvar_session session, MPI_T_pvar_handle handle)	6
The	MPI_T_PVAR_RESE	T call sets the performance variable with the handle identified	7
		starting value specified in Section 14.3.7. If it is not possible	8
o change	e the variable, the fu	nction returns MPI_T_ERR_PVAR_NO_WRITE.	9
		$AR_ALL_HANDLES$ is passed in handle, the MPI implementation	10
-		s within the session identified by the parameter session for	11 12
		cated. In this case, the routine returns MPI_SUCCESS if all	
		y (even if there are no valid handles or all are read-only),	$^{13}_{14}$ ticket 39
		NO_WRITE is returned. Read-only variables are ignored when	15
	VAR_ALL_HANDLES is	s specified.	16
			17
MPI_T_P	VAR_READRESET(s	ession, handle, buf)	18
IN	session	identifier of performance experiment session (handle)	19
IN	handle	handle of a performance variable (handle)	20
			21
OUT	buf	initial address of storage location for variable value (choice)	22
		(choice)	23
		(choice)	23 24
int MPT	T nuar readreset (23 24 25
int MPI_	-	MPI_T_pvar_session session, MPI_T_pvar_handle	24
	handle, void	MPI_T_pvar_session session, MPI_T_pvar_handle 1* buf)	24 25
This	handle, void call atomically com	MPI_T_pvar_session session, MPI_T_pvar_handle d* buf) bines the functionality of MPI_T_PVAR_READ and	24 25 26
This MPI_T_P	handle, void call atomically com VAR_RESET with th	MPI_T_pvar_session session, MPI_T_pvar_handle d* buf) bines the functionality of MPI_T_PVAR_READ and he same semantics as if these two calls were called separately.	24 25 26 27
This MPI_T_P f atomic	handle, void call atomically com VAR_RESET with the operations on this	MPI_T_pvar_session session, MPI_T_pvar_handle 1* buf) bines the functionality of MPI_T_PVAR_READ and he same semantics as if these two calls were called separately. variable are not supported, this routine returns	24 25 26 27 28 29 30
This MPI_T_P f atomic MPI_T_EF	handle, void call atomically com VAR_RESET with the operations on this RR_PVAR_NO_ATOMIC	MPI_T_pvar_session session, MPI_T_pvar_handle d* buf) bines the functionality of MPI_T_PVAR_READ and he same semantics as if these two calls were called separately. variable are not supported, this routine returns C.	24 25 26 27 28 29 30 31
This MPI_T_P f atomic MPI_T_EF The	handle, void call atomically com VAR_RESET with the operations on this RR_PVAR_NO_ATOMIC constant MPI_T_PVA	MPI_T_pvar_session session, MPI_T_pvar_handle d* buf) bines the functionality of MPI_T_PVAR_READ and he same semantics as if these two calls were called separately. variable are not supported, this routine returns C. R_ALL_HANDLES cannot be used as an argument for the func-	24 25 26 27 28 29 30 31 32
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of MPI_T_PVAR_ALL_HANDLES where appropriate. Further, this opens up additional performance optimizations for the implementation of handles. (*End of rationale.*)

Example: Tool to Detect Receives with Long Unexpected Message Queues

Example 14.6

The following example shows a sample tool to identify receive operations that occur during times with long message queues. This examples assumes that the MPI implementation exports a variable with the name "MPI_T_UMQ_LENGTH" to represent the current length of the unexpected message queue. The tool is implemented as a PMPI tool using the MPI profiling interface.

The tool consists of three parts: (1) the initialization (by intercepting the call to MPI_INIT), (2) the test for long unexpected message queues (by intercepting calls to MPI_RECV), and (3) the clean-up phase (by intercepting the call to MPI_FINALIZE). To capture all receives, the example would have to be extended to have similar wrappers for all receive operations.

Part 1— Initialization: During initialization, the tool searches for the variable and, once
 the right index is found, allocates a session and a handle for the variable with the found
 index, and starts the performance variable.

```
22
     #include <stdio.h>
23
     #include <stdlib.h>
24
     #include <string.h>
25
     #include <assert.h>
26
     #include <mpi.h>
27
28
     /* Global variables for the tool */
29
     static MPI_T_pvar_session session;
30
     static MPI_T_pvar_handle handle;
31
32
     int MPI_Init(int *argc, char ***argv ) {
33
     int err, num, i, index, namelen, verbosity;
34
              int var_class, bind, threadsup;
35
     int readonly, continuous, atomic, count;
36
     char name[18];
37
     MPI_Comm comm;
38
     MPI_Datatype datatype;
39
     MPI_T_enum enumtype;
40
41
     err=PMPI_Init(argc,argv);
42
     if (err!=MPI_SUCCESS) return err;
43
44
     err=PMPI_T_init_thread(MPI_THREAD_SINGLE,&threadsup);
45
     if (err!=MPI_SUCCESS) return err;
46
47
     err=PMPI_T_pvar_get_num(&num);
48
```

1

2 3 4

5 6

7

8

9

10

```
if (err!=MPI_SUCCESS) return err;
                                                                                     1
index=-1;
                                                                                     2
i=0;
                                                                                      3
while ((i<num) && (index<0) && (err==MPI_SUCCESS)) {</pre>
                                                                                      4
/* Pass a buffer that is at least one character longer than */
                                                                                     5
/* the name of the variable being searched for to avoid */
                                                                                     6
/* finding variables that have a name that has a prefix */
                                                                                     7
/* equal to the name of the variable being searched. */
                                                                                     8
namelen=18;
                                                                                      9
err=PMPI_T_pvar_get_info(i, name, &namelen, &verbosity,
                                                                                     10
&var_class, &datatype, &enumtype, NULL, NULL, &bind,
                                                                                     11
&readonly, &continuous, &atomic);
                                                                                     12
if (strcmp(name,"MPI_T_UMQ_LENGTH")==0) index=i;
                                                                                     13
i++; }
                                                                                     14
if (err!=MPI_SUCCESS) return err;
                                                                                     15
                                                                                     16
/* this could be handled in a more flexible way for a generic tool */
                                                                                     17
assert(index>=0);
                                                                                     18
assert(var_class==MPI_T_PVAR_CLASS_LEVEL);
                                                                                     19
assert(datatype==MPI_INT);
                                                                                     20
assert(bind==MPI_T_BIND_MPI_COMM);
                                                                                     21
                                                                                     22
/* Create a session */
                                                                                     23
err=PMPI_T_pvar_session_create(&session);
                                                                                     24
if (err!=MPI_SUCCESS) return err;
                                                                                     25
                                                                                     26
/* Get a handle and bind to MPI_COMM_WORLD */
                                                                                     27
comm=MPI_COMM_WORLD;
                                                                                     28
err=PMPI_T_pvar_handle_alloc(session, index, &comm, &handle, &count);
                                                                                     29
if (err!=MPI_SUCCESS) return err;
                                                                                     30
                                                                                     31
/* this could be handled in a more flexible way for a generic tool */
                                                                                     32
assert(count==1);
                                                                                     33
                                                                                     34
/* Start variable */
                                                                                     35
err=PMPI_T_pvar_start(session, handle);
                                                                                     36
if (err!=MPI_SUCCESS) return err;
                                                                                     37
                                                                                     38
return MPI_SUCCESS;
                                                                                     39
}
                                                                                     40
                                                                                     41
Part 2 — Testing the Queue Lengths During Receives: During every receive operation, the
                                                                                     42
tool reads the unexpected queue length through the matching performance variable and
                                                                                     43
compares it against a predefined threshold.
                                                                                     44
                                                                                     45
#define THRESHOLD 5
                                                                                     46
                                                                                     47
```

int MPI_Recv(void *buf, int count, MPI_Datatype datatype, int source, int tag, 48

```
MPI_Comm comm, MPI_Status *status)
1
     {
\mathbf{2}
3
     int value, err;
4
     if (comm==MPI_COMM_WORLD) {
5
     err=PMPI_T_pvar_read(session, handle, &value);
6
     if ((err==MPI_SUCCESS) && (value>THRESHOLD))
7
     {
8
                                /* tool identified receive called with long UMQ */
9
     /* execute tool functionality, */
10
     /* e.g., gather and print call stack */
11
     }
12
     }
13
14
     return PMPI_Recv(buf, count, datatype, source, tag, comm, status);
15
     }
16
17
     Part 3 — Termination: In the wrapper for MPI_FINALIZE, the MPI tool information inter-
18
19
     face is finalized.
20
     int MPI_Finalize()
21
     {
22
     int err;
23
     err=PMPI_T_pvar_handle_free(session, &handle);
24
     err=PMPI_T_pvar_session_free(&session);
25
     err=PMPI_T_finalize();
26
     return PMPI_Finalize();
27
     }
28
29
            Variable Categorization
30
     14.3.8
31
```

MPI implementations can optionally group performance and control variables into categories to express logical relationships between various variables. For example, an MPI implementation could group all control and performance variables that refer to message transfers in the MPI implementation and thereby distinguish them from variables that refer to local resources such as memory allocations or other interactions with the operating system.

Categories can also contain other categories to form a hierarchical grouping. Categories can never include themselves, either directly or transitively within other included categories. Expanding on the example above, this allows MPI to refine the grouping of variables referring to message transfers into variables to control and to monitor message queues, message matching activities and communication protocols. Each of these groups of variables would be represented by a separate category and these categories would then be listed in a single category representing variables for message transfers.

The category information may be queried in a fashion similar to the mechanism for querying variable information. The MPI implementation exports a set of N categories via the MPI tool information interface. If N = 0, then the MPI implementation does not export any categories, otherwise the provided categories are indexed from 0 to N - 1. This index

32

number is used in subsequent calls to functions of the MPI tool information interface to 1 identify the individual categories. 2 An MPI implementation is permitted to increase the number of categories during the 3 execution of an MPI program when new categories become available through dynamic load-4 ing. However, MPI implementations are not allowed to change the index of a category or 5delete it once it has been added to the set. 6 Similarly, MPI implementations are allowed to add variables to categories, but they 7 are not allowed to remove variables from categories or change the order in which they are 8 returned. 9 The following function can be used to query the number of [control variables] categories, 10 ticket 387. N.11 12 13 MPI_T_CATEGORY_GET_NUM(num_cat) 14 OUT current number of categories (integer) num_cat 1516 17int MPI_T_category_get_num(int *num_cat) 18 Individual category information can then be queried by calling the following function: 19 2021MPI_T_CATEGORY_GET_INFO(cat_index, name, name_len, desc, desc_len, num_cvars, 22num_pvars, num_categories) 23 IN cat_index index of the category to be queried (integer) 24OUT 25name buffer to return the string containing the name of the 26category (string) 27INOUT name_len length of the string and/or buffer for name (integer) 28OUT desc buffer to return the string containing the description 29 of the category (string) 30 31 INOUT desc_len length of the string and/or buffer for desc (integer) 32 OUT number of control variables in the category (integer) num_cvars 33 OUT number of performance variables in the category (innum_pvars 34 teger) 35OUT 36 num_categories number of categories contained in the category (integer) 37 38 39 int MPI_T_category_get_info(int cat_index, char *name, int *name_len, 40 char *desc, int *desc_len, int *num_cvars, int *num_pvars, 41int *num_categories) 42 The arguments name and name_len are used to return the name of the category as 43described in Section 14.3.3. 44 The routine is required to return a name of at least length one. This name must be 45 unique with respect to all other names for categories used by the MPI implementation. $_{46}$ ticket 378. If any OUT parameter to MPI_T_CATEGORY_GET_INFO is a NULL pointer, the im-47plementation will ignore the parameter and not return a value for the parameter. 48

Unofficial Draft for Comment Only

1		arguments desc an in Section 14.3.3.	d $desc_len$ are used to return the description of the category as			
3	Retu	Returning a description is optional. If an MPI implementation decides not to return a description, the first character for desc must be set to the null character and desc_len must				
4	-	one at the return				
6			ne number of control variables, performance variables and other			
7			queried category in the arguments num_cvars, num_pvars, and			
ticket383. 8		gories, respectively				
9	If th	e name of a catego	ory is the same across connected processes, the following OUT			
ticket377. 10						
11	11					
12		ATECORY CET I	NDEX(name, cat_index)			
13						
14		name	the name of the category (string)			
15 16	001	cat_index	the index of the category (integer)			
17 18	int MPI	_T_category_get_	index(const char *name, int *cat_index)			
19	MPI_	_T_CATEGORY_G	ET_INDEX is a function for retrieving the index of a category			
20	given a ki	nown category nan	ne. The name parameter is provided by the caller, and cat_index			
21		is returned by the MPI implementation. The name parameter is a string terminated with a				
22		null character.				
23			$PI_SUCCESS$ on success and returns $MPI_T_ERR_INVALID_NAME$			
24			e name of any category provided by the implementation at the			
25		he call.				
26 27	Dat	Rationale. This routine is provided to enable fast retrieval of a category in				
28		by a tool, assuming it knows the name of the category for which it is looking. The				
29			exposed by the implementation can change over time, so it is not			
30	\mathbf{pos}	sible for the tool to	o simply iterate over the list of categories once at initialization.			
31		0 0	implementation specific category names is not portable across			
32		MPI implementations, tool developers may choose to take this route for lower overhead				
33			ne tool will not have to iterate over the entire set of categories			
34	to f	ind a specific one.	(End of rationale.)			
35						
36						
37		CATEGORY_GET_	CVARS(cat_index, len, indices)			
38 39	IN	cat_index	index of the category to be queried, in the range $[0, N-$			
39 40			1] (integer)			
40	INI	len	the length of the indices array (integer)			
42		indices	an integer array of size len, indicating control variable			
43		mulces	indices (array of integers)			
44			indices (array of integers)			
45	int MDT	T category got	cvars(int cat_index, int len, int indices[])			
46						
47			ET_CVARS can be used to query which control variables are			
48	contained	l in a particular ca	tegory. A category contains zero or more control variables.			

MPI_T_C	ATEGORY_GET_PVA	RS(cat_index,len,indices)	1
IN	cat_index	index of the category to be queried, in the range $[0, N-$	2
		1] (integer)	3
	lon		4
IN	len	the length of the indices array (integer)	5
OUT	indices	an integer array of size len, indicating performance	6
		variable indices (array of integers)	7
			8

MPL T CATEGORY GET PVARS(cat index len indices)

int MPI_T_category_get_pvars(int cat_index, int len, int indices[])

MPI_T_CATEGORY_GET_PVARS can be used to query which performance variables are contained in a particular category. A category contains zero or more performance variables.

MPI_T_CA	ATEGORY_GET_CATEGORIES	(cat_index,len,indices)	15
IN	cat_index	index of the category to be queried, in the range $[0, N-1]$ (integer)	16 17
IN	len	the length of the indices array (integer)	18 19
OUT	indices	an integer array of size len, indicating category indices	20
		(array of integers)	21 22

int MPI_T_category_get_categories(int cat_index, int len, int indices[])

MPI_T_CATEGORY_GET_CATEGORIES can be used to query which other categories are contained in a particular category. A category contains zero or more other categories.

As mentioned above, MPI implementations can grow the number of categories as well as the number of variables or other categories within a category. In order to allow users of the MPI tool information interface to check quickly whether new categories have been added or new variables or categories have been added to a category, MPI maintains a virtual timestamp. This timestamp is monotonically increasing during the execution and is returned by the following function:

			34
MPI_T_C	ATEGORY_CHANGED(stamp)		35
OUT	stamp	a virtual time stamp to indicate the last change to the	36
		categories (integer)	37
			38

int MPI_T_category_changed(int *stamp)

If two subsequent calls to this routine return the same timestamp, it is guaranteed that the category information has not changed between the two calls. If the timestamp retrieved from the second call is higher, then some categories have been added or expanded.

Advice to users. The timestamp value is purely virtual and only intended to check for changes in the category information. It should not be used for any other purpose. (End of advice to users.)

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4344

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The index values returned in indices by MPI_T_CATEGORY_GET_CVARS, 1 MPI_T_CATEGORY_GET_PVARS and MPI_T_CATEGORY_GET_CATEGORIES can be used 2 as input to MPI_T_CVAR_GET_INFO, MPI_T_PVAR_GET_INFO and 3 MPI_T_CATEGORY_GET_INFO, respectively. 4 The user is responsible for allocating the arrays passed into the functions $\mathbf{5}$ MPI_T_CATEGORY_GET_CVARS, MPI_T_CATEGORY_GET_PVARS and 6 MPI_T_CATEGORY_GET_CATEGORIES. Starting from array index 0, each function writes 7 up to len elements into the array. If the category contains more than len elements, the 8 function returns an arbitrary subset of size len. Otherwise, the entire set of elements is 9 returned in the beginning entries of the array, and any remaining array entries are not 10 modified. 1112Return Codes for the MPI Tool Information Interface 13 14.3.914 All functions defined as part of the MPI tool information interface return an integer error 15code (see Table 14.5) to indicate whether the function was completed successfully or was 16 aborted. In the latter case the error code indicates the reason for not completing the routine. 17Such errors neither impact the execution of the MPI process nor invoke MPI error handlers. 18 The MPI process continues executing regardless of the return code from the call. The MPI 19 implementation is not required to check all user-provided parameters; if a user passes invalid 20parameter values to any routine the behavior of the implementation is undefined. 21All error codes with the prefix MPI_T_ must be unique values and cannot overlap with 22 any other error codes or error classes returned by the MPI implementation. Further, they 23 shall be treated as MPI error classes as defined in Section 8.4 on page 347 and follow the 24same rules and restrictions. In particular, they must satisfy: 25 $0 = MPI_SUCCESS < MPI_T_ERR_... \le MPI_ERR_LASTCODE.$ 2728All MPI tool information interface functions must return error classes, Rationale. 29 because applications cannot portably call MPI_ERROR_CLASS before 30 31

26

MPI_INIT or MPI_INIT_THREAD to map an arbitrary error code to an error class. (End of rationale.)

32 33 34

35 36

37 38

Profiling Interface 14.3.10

All requirements for the profiling interfaces, as described in Section 14.2, also apply to the MPI tool information interface. All rules, guidelines, and recommendations from Section 14.2 apply equally to calls defined as part of the MPI tool information interface.

- 39 40
- 41 42

- 44
- 45
- 46
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		1
Return Code	Description	2 3
Return Codes for All Functions in the N		4
MPI_SUCCESS	Call completed successfully	5
[ticket400.] MPI_T_ERR_INVALID	[ticket400.]Invalid use of the interface or bad parame	etær values
MPI_T_ERR_MEMORY	Out of memory	7
MPI_T_ERR_NOT_INITIALIZED	Interface not initialized	8
MPI_T_ERR_CANNOT_INIT	Interface not in the state to be initialized	9
Return Codes for Datatype Functions:		10
MPI_T_ERR_INVALID_INDEX	The enumeration index is invalid	11
	[ticket406.][or has been deleted.]	12
MPI_T_ERR_INVALID_ITEM	The item index queried is out of range	13
	(for MPI_T_ENUM_GET_ITEM only)	14
Return Codes for variable and category	r query functions: [ticket377.][MPI_T_*_GET_INFO]M	PT * GET *
MPI_T_ERR_INVALID_INDEX	The variable or category index is invalid	16
[ticket377.]MPI_T_ERR_INVALID_NAME	[ticket377.]The variable or category name is invalid	17
Return Codes for Handle Functions: MI		
MPI_T_ERR_INVALID_INDEX	The variable index is invalid [ticket406.] [or has been	<u>19</u> deleted]
MPI_T_ERR_INVALID_HANDLE	The handle is invalid	20
MPI_T_ERR_OUT_OF_HANDLES	No more handles available	21
Return Codes for Session Functions: M		22
		23
MPI_T_ERR_OUT_OF_SESSIONS	No more sessions available	24
MPI_T_ERR_INVALID_SESSION	Session argument is not a valid session	25
Return Codes for Control Variable Acce	ess Functions:	26
MPI_T_CVAR_READ, WRITE		27
MPI_T_ERR_CVAR_SET_NOT_NOW	Variable cannot be set at this moment	28
MPI_T_ERR_CVAR_SET_NEVER	Variable cannot be set until end of execution	29
MPI_T_ERR_INVALID_HANDLE	The handle is invalid	30
Return Codes for Performance Variable		31
MPI_T_PVAR_{START STOP READ W		32
MPI_T_ERR_INVALID_HANDLE	The handle is invalid	33
MPI_T_ERR_INVALID_SESSION	Session argument is not a valid session	34
MPI_T_ERR_PVAR_NO_STARTSTOP	Variable cannot be started or stopped	35
	(for MPI_T_PVAR_START and	36
	MPI_T_PVAR_STOP)	37
MPI_T_ERR_PVAR_NO_WRITE	Variable cannot be written or reset	38
	(for MPI_T_PVAR_WRITE and	39
	MPI_T_PVAR_RESET)	40
MPI_T_ERR_PVAR_NO_ATOMIC	Variable cannot be read and written atomically	41
	(for MPI_T_PVAR_READRESET)	42
Return Codes for Category Functions:		-43
MPI_T_ERR_INVALID_INDEX	The category index is invalid	-44
	The category much is myand	-45
Table 14.5: Return codes used in fu	nctions of the MPI tool information interface	46

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