

# MPI: A Message-Passing Interface Standard

Version 3.0

(Draft)

Unofficial, for comment only

Message Passing Interface Forum

July 5, 2012

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ticket0. 1 This document describes the Message-Passing Interface (MPI) standard, version [2.2]3.0.  
2 The MPI standard includes point-to-point message-passing, collective communications, group  
3 and communicator concepts, process topologies, environmental management, process cre-  
4 ation and management, one-sided communications, extended collective operations, external  
5 interfaces, I/O, some miscellaneous topics, and a profiling interface. Language bindings for  
ticket0. 6 C, C++ and Fortran are defined.

7 [Technically, this version of the standard is based on “MPI: A Message-Passing Interface  
8 Standard, version 2.1, June 23, 2008. The MPI Forum added seven new routines and a  
9 number of enhancements and clarifications to the standard.]

10 Historically, the evolution of the standards is from MPI-1.0 (June 1994) to MPI-1.1  
11 (June 12, 1995) to MPI-1.2 (July 18, 1997), with several clarifications and additions and  
12 published as part of the MPI-2 document, to MPI-2.0 (July 18, 1997), with new functionality,  
13 to MPI-1.3 (May 30, 2008), combining for historical reasons the documents 1.1 and 1.2  
14 and some errata documents to one combined document, and to MPI-2.1 (June 23, 2008),  
ticket0. 15 combining the previous documents. [This version, MPI-2.2, is based on MPI-2.1 and provides  
16 additional clarifications and errata corrections as well as a few enhancements.]Version MPI-  
17 2.2 (September 2009) added additional clarifications and seven new routines. This version,  
18 MPI-3.0, is an extension of MPI-2.2.  
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Version 3.0: xx, x, 2011. Coincident with the development of MPI-2.2, the MPI Forum began discussions of a major extension to MPI. This document contains the MPI-3 Standard. This draft version of the MPI-3 standard extends the collective operations by including nonblocking versions. Unlike MPI-2.2, this standard is considered a major update to the MPI standard. As with previous versions, new features have been adopted only when there were compelling needs for the users. Some features, however, may have more than a minor impact on existing MPI implementations.

Version 2.2: September 4, 2009. This document contains mostly corrections and clarifications to the [MPI 2.1]MPI-2.1 document. A few extensions have been added; however all correct [MPI 2.1]MPI-2.1 programs are correct [MPI 2.2]MPI-2.2 programs. New features were adopted only when there were compelling needs for users, open source implementations, and minor impact on existing MPI implementations.

Version 2.1: June 23, 2008. This document combines the previous documents MPI-1.3 (May 30, 2008) and MPI-2.0 (July 18, 1997). Certain parts of MPI-2.0, such as some sections of Chapter 4, Miscellany, and Chapter 7, Extended Collective Operations have been merged into the Chapters of MPI-1.3. Additional errata and clarifications collected by the MPI Forum are also included in this document.

Version 1.3: May 30, 2008. This document combines the previous documents MPI-1.1 (June 12, 1995) and the MPI-1.2 Chapter in MPI-2 (July 18, 1997). Additional errata collected by the MPI Forum referring to MPI-1.1 and MPI-1.2 are also included in this document.

Version 2.0: July 18, 1997. Beginning after the release of MPI-1.1, the MPI Forum began meeting to consider corrections and extensions. MPI-2 has been focused on process creation and management, one-sided communications, extended collective communications, external interfaces and parallel I/O. A miscellany chapter discusses items that [don't]do not fit elsewhere, in particular language interoperability.

Version 1.2: July 18, 1997. The MPI-2 Forum introduced MPI-1.2 as Chapter 3 in the standard ["]“MPI-2: Extensions to the Message-Passing Interface”, July 18, 1997. This section contains clarifications and minor corrections to Version 1.1 of the MPI Standard. The only new function in MPI-1.2 is one for identifying to which version of the MPI Standard the implementation conforms. There are small differences between MPI-1 and MPI-1.1. There are very few differences between MPI-1.1 and MPI-1.2, but large differences between MPI-1.2 and MPI-2.

Version 1.1: June, 1995. Beginning in March, 1995, the Message-Passing Interface Forum reconvened to correct errors and to make clarifications in the MPI document of May 5, 1994, referred to below as Version 1.0. These discussions resulted in Version 1.1[, which is this document]. The changes from Version 1.0 are minor. A version of this document with all changes marked is available. [This paragraph is an example of a change.]

1 Version 1.0: May, 1994. The Message-Passing Interface Forum (MPIF), with participation  
ticket0. 2 from over 40 organizations, has been meeting since January 1993 to discuss and to define a  
3 set of library interface standards for message passing. MPIF is not sanctioned or supported  
4 by any official standards organization.

5 The goal of the Message-Passing Interface, simply stated, is to develop a widely used  
6 standard for writing message-passing programs. As such the interface should establish a  
ticket0. 7 practical, portable, efficient, and flexible standard for message-passing.

8 [This is the final report, Version 1.0, of the Message-Passing Interface Forum. ]This  
9 document contains all the technical features proposed for the interface. This copy of the  
10 draft was processed by L<sup>A</sup>T<sub>E</sub>X on May 5, 1994.

11 Please send comments on MPI to [mpi-comments@mpi-forum.org](mailto:mpi-comments@mpi-forum.org). Your comment will  
12 be forwarded to MPI Forum committee members who will attempt to respond.

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All chapters have been revisited to achieve a consistent MPI-2.2 text. Those who served as authors for the necessary modifications are:

- William Gropp, Frontmatter, Introduction, and Bibliography; MPI 2.2 chair.
- Richard Graham, Point-to-Point Communication and Datatypes
- Adam Moody, Collective Communication
- Torsten Hoefler, Collective Communication and Process Topologies
- Richard Treumann, Groups, Contexts, and Communicators
- Jesper Larsson Träff, Process Topologies, Info-Object and One-Sided Communications
- George Bosilca, Datatypes and Environmental Management
- David Solt, Process Creation and Management
- Bronis R. de Supinski, External Interfaces, and Profiling
- Rajeev Thakur, I/O
- Jeffrey M. Squyres, Language Bindings and MPI 2.2 Secretary
- Rolf Rabenseifner, Deprecated Functions, Annex Change-Log, and Annex Language Bindings
- Alexander Supalov, Annex Language Bindings

The following list includes some of the active participants who attended MPI-2 Forum meetings and in the e-mail discussions of the errata items and are not mentioned above.

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### MPI-3:

MPI-3 is a significant effort to extend and modernize the MPI Standard.

The editors and organizers of the MPI-3 have been: *Taken from MPI-2.2 with minor corrections. Need to separate the working groups list (which is currently reviewers) from the primary authors . Also, did I miss active steering committee members?*

- William Gropp, Steering committee, Frontmatter, Introduction, Groups, Contexts, and Communicators, One-Sided Communications, and Bibliography
- Richard Graham, Steering committee, Point-to-Point Communication; Meeting Convener, and MPI-3 chair
- Adam Moody, Collective Communication
- Torsten Hoefler, Collective Communication and Process Topologies
- George Bosilca, Datatypes and Environmental Management
- David Solt, Process Creation and Management
- Bronis R. de Supinski, External Interfaces, and Profiling
- Rajeev Thakur, I/O and One-Sided Communications
- Darius Buntinas, Info Object
- Jeffrey M. Squyres, Language Bindings and MPI 3.0 Secretary



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# Chapter 1

## Introduction to MPI

### 1.1 Overview and Goals

MPI (Message-Passing Interface) is a *message-passing library interface specification*. All parts of this definition are significant. MPI addresses primarily the message-passing parallel programming model, in which data is moved from the address space of one process to that of another process through cooperative operations on each process. [Extensions to the “classical” message-passing model are provided in collective operations, remote-memory access operations, dynamic process creation, and parallel I/O.] MPI is a *specification*, not an implementation; there are multiple implementations of MPI. This specification is for a *library interface*; MPI is not a language, and all MPI operations are expressed as functions, subroutines, or methods, according to the appropriate language bindings, which for C, C++, [Fortran-77, and Fortran-95]and Fortran, are part of the MPI standard. The standard has been defined through an open process by a community of parallel computing vendors, computer scientists, and application developers. The next few sections provide an overview of the history of MPI’s development.

The main advantages of establishing a message-passing standard are portability and ease of use. In a distributed memory communication environment in which the higher level routines and/or abstractions are built upon lower level message-passing routines the benefits of standardization are particularly apparent. Furthermore, the definition of a message-passing standard, such as that proposed here, provides vendors with a clearly defined base set of routines that they can implement efficiently, or in some cases [provide hardware support for]for which they can provide hardware support, thereby enhancing scalability.

The goal of the Message-Passing Interface simply stated is to develop a widely used standard for writing message-passing programs. As such the interface should establish a practical, portable, efficient, and flexible standard for message passing.

A complete list of goals follows.

- Design an application programming interface (not necessarily for compilers or a system implementation library).
- Allow efficient communication: Avoid memory-to-memory copying, allow overlap of computation and communication, and offload to communication co-processor, where available.
- Allow for implementations that can be used in a heterogeneous environment.

- Allow convenient C, C++, [Fortran-77, and Fortran-95]and Fortran bindings for the interface.
- Assume a reliable communication interface: the user need not cope with communication failures. Such failures are dealt with by the underlying communication subsystem.
- Define an interface that can be implemented on many vendor's platforms, with no significant changes in the underlying communication and system software.
- Semantics of the interface should be language independent.
- The interface should be designed to allow for thread safety.

## 1.2 Background of MPI-1.0

MPI sought to make use of the most attractive features of a number of existing message-passing systems, rather than selecting one of them and adopting it as the standard. Thus, MPI was strongly influenced by work at the IBM T. J. Watson Research Center [1, 2], Intel's NX/2 [50], Express [13], nCUBE's Vertex [46], p4 [8, 9], and PARMACS [5, 10]. Other important contributions have come from Zipcode [53, 54], Chimp [17, 18], PVM [4, 15], Chameleon [27], and PICL [25].

The MPI standardization effort involved about 60 people from 40 organizations mainly from the United States and Europe. Most of the major vendors of concurrent computers were involved in MPI, along with researchers from universities, government laboratories, and industry. The standardization process began with the Workshop on Standards for Message-Passing in a Distributed Memory Environment, sponsored by the Center for Research on Parallel Computing, held April 29-30, 1992, in Williamsburg, Virginia [61]. At this workshop the basic features essential to a standard message-passing interface were discussed, and a working group established to continue the standardization process.

A preliminary draft proposal, known as MPI1, was put forward by Dongarra, Hempel, Hey, and Walker in November 1992, and a revised version was completed in February 1993 [16]. MPI1 embodied the main features that were identified at the Williamsburg workshop as being necessary in a message passing standard. Since MPI1 was primarily intended to promote discussion and "get the ball rolling," it focused mainly on point-to-point communications. MPI1 brought to the forefront a number of important standardization issues, but did not include any collective communication routines and was not thread-safe.

In November 1992, a meeting of the MPI working group was held in Minneapolis, at which it was decided to place the standardization process on a more formal footing, and to generally adopt the procedures and organization of the High Performance Fortran Forum. Subcommittees were formed for the major component areas of the standard, and an email discussion service established for each. In addition, the goal of producing a draft MPI standard by the Fall of 1993 was set. To achieve this goal the MPI working group met every 6 weeks for two days throughout the first 9 months of 1993, and presented the draft MPI standard at the Supercomputing 93 conference in November 1993. These meetings and the email discussion together constituted the MPI Forum, membership of which has been open to all members of the high performance computing community.

### 1.3 Background of MPI-1.1, MPI-1.2, and MPI-2.0

Beginning in March 1995, the MPI Forum began meeting to consider corrections and extensions to the original MPI Standard document [22]. The first product of these deliberations was Version 1.1 of the MPI specification, released in June of 1995 [23] (see <http://www.mpi-forum.org> for official MPI document releases). At that time, effort focused in five areas.

1. Further corrections and clarifications for the MPI-1.1 document.
2. Additions to MPI-1.1 that do not significantly change its types of functionality (new datatype constructors, language interoperability, etc.).
3. Completely new types of functionality (dynamic processes, one-sided communication, parallel I/O, etc.) that are what everyone thinks of as “MPI-2 functionality.”
4. Bindings for Fortran 90 and C++. MPI-2 specifies C++ bindings for both MPI-1 and MPI-2 functions, and extensions to the Fortran 77 binding of MPI-1 and MPI-2 to handle Fortran 90 issues.
5. Discussions of areas in which the MPI process and framework seem likely to be useful, but where more discussion and experience are needed before standardization (e.g. zero-copy semantics on shared-memory machines, real-time specifications).

Corrections and clarifications (items of type 1 in the above list) were collected in Chapter 3 of the MPI-2 document: “Version 1.2 of MPI.” That chapter also contains the function for identifying the version number. Additions to MPI-1.1 (items of types 2, 3, and 4 in the above list) are in the remaining chapters of the MPI-2 document, and constitute the specification for MPI-2. Items of type 5 in the above list have been moved to a separate document, the “MPI Journal of Development” (JOD), and are not part of the MPI-2 Standard.

This structure makes it easy for users and implementors to understand what level of MPI compliance a given implementation has:

- MPI-1 compliance will mean compliance with MPI-1.3. This is a useful level of compliance. It means that the implementation conforms to the clarifications of MPI-1.1 function behavior given in Chapter 3 of the MPI-2 document. Some implementations may require changes to be MPI-1 compliant.
- MPI-2 compliance will mean compliance with all of MPI-2.1.
- The MPI Journal of Development is not part of the MPI Standard.

It is to be emphasized that forward compatibility is preserved. That is, a valid MPI-1.1 program is both a valid MPI-1.3 program and a valid MPI-2.1 program, and a valid MPI-1.3 program is a valid MPI-2.1 program.

### 1.4 Background of MPI-1.3 and MPI-2.1

After the release of MPI-2.0, the MPI Forum kept working on errata and clarifications for both standard documents (MPI-1.1 and MPI-2.0). The short document “Errata for MPI-1.1” was released October 12, 1998. On July 5, 2001, a first ballot of errata and clarifications for

MPI-2.0 was released, and a second ballot was voted on May 22, 2002. Both votes were done electronically. Both ballots were combined into one document: “Errata for MPI-2”, May 15, 2002. This errata process was then interrupted, but the Forum and its e-mail reflectors kept working on new requests for clarification.

Restarting regular work of the MPI Forum was initiated in three meetings, at EuroPVM/MPI’06 in Bonn, at EuroPVM/MPI’07 in Paris, and at SC’07 in Reno. In December 2007, a steering committee started the organization of new MPI Forum meetings at regular 8-weeks intervals. At the January 14-16, 2008 meeting in Chicago, the MPI Forum decided to combine the existing and future MPI documents to one [single] document for each version of the MPI standard. For technical and historical reasons, this series was started with MPI-1.3. Additional Ballots 3 and 4 solved old questions from the errata list started in 1995 up to new questions from the last years. After all documents (MPI-1.1, MPI-2, Errata for MPI-1.1 (Oct. 12, 1998), and MPI-2.1 Ballots 1-4) were combined into one draft document, for each chapter, a chapter author and review team were defined. They cleaned up the document to achieve a consistent MPI-2.1 document. The final MPI-2.1 standard document was finished in June 2008, and finally released with a second vote in September 2008 in the meeting at Dublin, just before EuroPVM/MPI’08. The major work of the current MPI Forum is the preparation of MPI-3.

## 1.5 Background of MPI-2.2

MPI-2.2 is a minor update to the MPI-2.1 standard. This version addresses additional errors and ambiguities that were not corrected in the MPI-2.1 standard as well as a small number of extensions to MPI-2.1 that met the following criteria:

- Any correct MPI-2.1 program is a correct MPI-2.2 program.
- Any extension must have significant benefit for users.
- Any extension must not require significant implementation effort. To that end, all such changes are accompanied by an open source implementation.

The discussions of MPI-2.2 proceeded concurrently with the MPI-3 discussions; in some cases, extensions were proposed for MPI-2.2 but were later moved to MPI-3.

## 1.6 Background of MPI-3.0

MPI-3.0 is a major update to the MPI standard. Areas of particular interest are the extension of collective operations to include nonblocking, with other areas under consideration. This *draft* contains the MPI Forum’s current draft of nonblocking collective routines.

A new Fortran `mpi_f08` module is introduced to provide extended compile-time argument checking and buffer handling in nonblocking routines. This new Fortran support method provides protection against the optimization problems[ with] with asynchronous accesses to the buffers of nonblocking calls. The existing `mpi` module is enhanced to provide basic compile-time argument checking for MPI calls. The use of `mpif.h` is strongly discouraged.

## 1.7 Who Should Use This Standard?

This standard is intended for use by all those who want to write portable message-passing programs in Fortran, C and C++. This includes individual application programmers, developers of software designed to run on parallel machines, and creators of environments and tools. In order to be attractive to this wide audience, the standard must provide a simple, easy-to-use interface for the basic user while not semantically precluding the high-performance message-passing operations available on advanced machines.

## 1.8 What Platforms Are Targets For Implementation?

The attractiveness of the message-passing paradigm at least partially stems from its wide portability. Programs expressed this way may run on distributed-memory multiprocessors, networks of workstations, and combinations of all of these. In addition, shared-memory implementations, including those for multi-core processors and hybrid architectures, are possible. The paradigm will not be made obsolete by architectures combining the shared- and distributed-memory views, or by increases in network speeds. It thus should be both possible and useful to implement this standard on a great variety of machines, including those “machines” consisting of collections of other machines, parallel or not, connected by a communication network.

The interface is suitable for use by fully general MIMD programs, as well as those written in the more restricted style of SPMD. MPI provides many features intended to improve performance on scalable parallel computers with specialized interprocessor communication hardware. Thus, we expect that native, high-performance implementations of MPI will be provided on such machines. At the same time, implementations of MPI on top of standard Unix interprocessor communication protocols will provide portability to workstation clusters and heterogenous networks of workstations.

## 1.9 What Is Included In The Standard?

The standard includes:

- Point-to-point communication,
- Datatypes,
- Collective operations,
- Process groups,
- Communication contexts,
- Process topologies,
- Environmental [M]management and inquiry,
- The [i]Info object,
- Process creation and management,
- One-sided communication,

- External interfaces,
- Parallel file I/O,
- Language [B]bindings for Fortran, C and C++,
- Profiling interface.

ticket0.

## 1.10 What Is Not Included In The Standard?

The standard does not specify:

- Operations that require more operating system support than is currently standard; for example, interrupt-driven receives, remote execution, or active messages,
- Program construction tools,
- Debugging facilities.

There are many features that have been considered and not included in this standard. This happened for a number of reasons, one of which is the time constraint that was self-imposed in finishing the standard. Features that are not included can always be offered as extensions by specific implementations. Perhaps future versions of MPI will address some of these issues.

## 1.11 Organization of this Document

The following is a list of the remaining chapters in this document, along with a brief description of each.

- Chapter 2, MPI Terms and Conventions, explains notational terms and conventions used throughout the MPI document.
- Chapter 3, Point to Point Communication, defines the basic, pairwise communication subset of MPI. *Send* and *receive* are found here, along with many associated functions designed to make basic communication powerful and efficient.
- Chapter 4, Datatypes, defines a method to describe any data layout, e.g., an array of structures in the memory, which can be used as message send or receive buffer.
- Chapter 5, Collective Communications, defines process-group collective communication operations. Well known examples of this are barrier and broadcast over a group of processes (not necessarily all the processes). With MPI-2, the semantics of collective communication was extended to include intercommunicators. It also adds two new collective operations. **MPI-3 adds nonblocking collective operations.**
- Chapter 6, Groups, Contexts, Communicators, and Caching, shows how groups of processes are formed and manipulated, how unique communication contexts are obtained, and how the two are bound together into a *communicator*.

- Chapter 7, Process Topologies, explains a set of utility functions meant to assist in the mapping of process groups (a linearly ordered set) to richer topological structures such as multi-dimensional grids.
- Chapter 8, MPI Environmental Management, explains how the programmer can manage and make inquiries of the current MPI environment. These functions are needed for the writing of correct, robust programs, and are especially important for the construction of highly-portable message-passing programs.
- Chapter 9, The Info Object, defines an opaque object, that is used as input `[of]in` several MPI routines.
- Chapter 10, Process Creation and Management, defines routines that allow for creation of processes.
- Chapter 11, One-Sided Communications, defines communication routines that can be completed by a single process. These include shared-memory operations (put/get) and remote accumulate operations.
- Chapter 12, External Interfaces, defines routines designed to allow developers to layer on top of MPI. This includes generalized requests, routines that decode MPI opaque objects, and threads.
- Chapter 13, I/O, defines MPI support for parallel I/O.
- Chapter 14.2, Profiling Interface, explains a simple name-shifting convention that any MPI implementation must support. One motivation for this is the ability to put performance profiling calls into MPI without the need for access to the MPI source code. The name shift is merely an interface, it says nothing about how the actual profiling should be done and in fact, the name shift can be useful for other purposes.
- Chapter 15, Deprecated Functions, describes routines that are kept for reference. However usage of these functions is discouraged, as they may be deleted in future versions of the standard.
- Chapter 16, Language Bindings, describes the C++ binding, discusses Fortran issues, and describes language interoperability aspects between C, C++, and Fortran.

The Appendices are:

- Annex A, Language Bindings Summary, gives specific syntax in C, C++, and Fortran, for all MPI functions, constants, and types.
- Annex B, Change-Log, summarizes major changes since the previous version of the standard.
- Several Index pages `[are showing]show` the locations of examples, constants and pre-defined handles, callback routine`[s]` prototypes, and all MPI functions.

MPI provides various interfaces to facilitate interoperability of distinct MPI implementations. Among these are the canonical data representation for MPI I/O and for

MPI\_PACK\_EXTERNAL and MPI\_UNPACK\_EXTERNAL. The definition of an actual binding of these interfaces that will enable interoperability is outside the scope of this document.

A separate document consists of ideas that were discussed in the MPI Forum and deemed to have value, but are not included in the MPI Standard. They are part of the “Journal of Development” (JOD), lest good ideas be lost and in order to provide a starting point for further work. The chapters in the JOD are

- Chapter 2, **Spawning Independent Processes**, includes some elements of dynamic process management, in particular management of processes with which the spawning processes do not intend to communicate, that the Forum discussed at length but ultimately decided not to include in the MPI Standard.
- Chapter 3, **Threads and MPI**, describes some of the expected interaction between an MPI implementation and a thread library in a multi-threaded environment.
- Chapter 4, **Communicator ID**, describes an approach to providing identifiers for communicators.
- Chapter 5, **Miscellany**, discusses Miscellaneous topics in the MPI JOD, in particular single-copy routines for use in shared-memory environments and new datatype constructors.
- Chapter 6, **Toward a Full Fortran 90 Interface**, describes an approach to providing a more elaborate Fortran 90 interface.
- Chapter 7, **Split Collective Communication**, describes a specification for certain non-blocking collective operations.
- Chapter 8, **Real-Time MPI**, discusses MPI support for real time processing.



## Chapter 2

# MPI Terms and Conventions

This chapter explains notational terms and conventions used throughout the MPI document, some of the choices that have been made, and the rationale behind those choices. It is similar to the MPI-1 Terms and Conventions chapter but differs in some major and minor ways. Some of the major areas of difference are the naming conventions, some semantic definitions, file objects, Fortran 90 *vs* Fortran 77, C++, processes, and interaction with signals.

### 2.1 Document Notation

*Rationale.* Throughout this document, the rationale for the design choices made in the interface specification is set off in this format. Some readers may wish to skip these sections, while readers interested in interface design may want to read them carefully. (*End of rationale.*)

*Advice to users.* Throughout this document, material aimed at users and that illustrates usage is set off in this format. Some readers may wish to skip these sections, while readers interested in programming in MPI may want to read them carefully. (*End of advice to users.*)

*Advice to implementors.* Throughout this document, material that is primarily commentary to implementors is set off in this format. Some readers may wish to skip these sections, while readers interested in MPI implementations may want to read them carefully. (*End of advice to implementors.*)

### 2.2 Naming Conventions

In many cases MPI names for C functions are of the form `MPI_Class_action_subset`. This convention originated with MPI-1. Since MPI-2 an attempt has been made to standardize the names of MPI functions according to the following rules. The C++ bindings in particular follow these rules (see Section 2.6.4 on page 20).

1. In C, all routines associated with a particular type of MPI object should be of the form `MPI_Class_action_subset` or, if no subset exists, of the form `MPI_Class_action`. In Fortran, all routines associated with a particular type of MPI object should be of the form `MPI_CLASS_ACTION_SUBSET` or, if no subset exists, of the form

MPI\_CLASS\_ACTION. For C and Fortran we use the C++ terminology to define the Class. In C++, the routine is a method on **Class** and is named MPI::Class::Action\_subset. If the routine is associated with a certain class, but does not make sense as an object method, it is a static member function of the class.

2. If the routine is not associated with a class, the name should be of the form MPI\_Action\_subset in C and MPI\_ACTION\_SUBSET in Fortran, and in C++ should be scoped in the MPI namespace, MPI::Action\_subset.
3. The names of certain actions have been standardized. In particular, **Create** creates a new object, **Get** retrieves information about an object, **Set** sets this information, **Delete** deletes information, **Is** asks whether or not an object has a certain property.

C and Fortran names for some MPI functions (that were defined during the MPI-1 process) violate these rules in several cases. The most common exceptions are the omission of the **Class** name from the routine and the omission of the **Action** where one can be inferred.

MPI identifiers are limited to 30 characters (31 with the profiling interface). This is done to avoid exceeding the limit on some compilation systems.

## 2.3 Procedure Specification

MPI procedures are specified using a language-independent notation. The arguments of procedure calls are marked as IN, OUT or INOUT. The meanings of these are:

- IN: the call may use the input value but does not update the argument **from the perspective of the caller at any time during the call's execution**,
- OUT: the call may update the argument but does not use its input value,
- INOUT: the call may both use and update the argument.

There is one special case — if an argument is a handle to an opaque object (these terms are defined in Section 2.5.1), and the object is updated by the procedure call, then the argument is marked INOUT or OUT. It is marked this way even though the handle itself is not modified — we use the INOUT or OUT attribute to denote that what the handle *references* is updated. Thus, in C++, IN arguments are usually either references or pointers to **const** objects.

*Rationale.* The definition of MPI tries to avoid, to the largest possible extent, the use of INOUT arguments, because such use is error-prone, especially for scalar arguments. (*End of rationale.*)

MPI's use of IN, OUT and INOUT is intended to indicate to the user how an argument is to be used, but does not provide a rigorous classification that can be translated directly into all language bindings (e.g., **INTENT** in Fortran 90 bindings or **const** in C bindings). For instance, the “constant” MPI\_BOTTOM can usually be passed to OUT buffer arguments. Similarly, MPI\_STATUS\_IGNORE can be passed as the OUT status argument.

A common occurrence for MPI functions is an argument that is used as IN by some processes and OUT by other processes. Such an argument is, syntactically, an INOUT argument

and is marked as such, although, semantically, it is not used in one call both for input and for output on a single process.

Another frequent situation arises when an argument value is needed only by a subset of the processes. When an argument is not significant at a process then an arbitrary value can be passed as an argument.

Unless specified otherwise, an argument of type OUT or type INOUT cannot be aliased with any other argument passed to an MPI procedure. An example of argument aliasing in C appears below. If we define a C procedure like this,

```
void copyIntBuffer( int *pin, int *pout, int len )
{
    int i;
    for (i=0; i<len; ++i) *pout++ = *pin++;
}
```

then a call to it in the following code fragment has aliased arguments.

```
int a[10];
copyIntBuffer( a, a+3, 7);
```

Although the C language allows this, such usage of MPI procedures is forbidden unless otherwise specified. Note that Fortran prohibits aliasing of arguments.

All MPI functions are first specified in the language-independent notation. Immediately below this, [the ISO C version of the function is shown followed by a version of the same function in Fortran and then the C++ binding.]language dependent bindings follow:

- The ISO C version of the function.
- The Fortran version used with `USE mpi_f08`.
- The Fortran version of the same function used with `USE mpi` or `INCLUDE 'mpif.h'`
- The C++ binding (which is deprecated).

“Fortran” in this document refers to Fortran 90 and higher; see Section 2.6.

## 2.4 Semantic Terms

When discussing MPI procedures the following semantic terms are used.

**nonblocking** A procedure is nonblocking if the procedure may return before the operation completes, and before the user is allowed to reuse resources (such as buffers) specified in the call. A nonblocking request is **started** by the call that initiates it, e.g., `MPI_ISEND`. The word complete is used with respect to operations, requests, and communications. An **operation completes** when the user is allowed to reuse resources, and any output buffers have been updated; i.e. a call to `MPI_TEST` will return `flag = true`. A **request is completed** by a call to wait, which returns, or a test or get status call which returns `flag = true`. This completing call has two effects: the status is extracted from the request; in the case of test and wait, if the request was nonpersistent, it is **freed**, and becomes **inactive** if it was persistent. A **communication completes** when all participating operations complete.

**blocking** A procedure is blocking if return from the procedure indicates the user is allowed to reuse resources specified in the call.

**local** A procedure is local if completion of the procedure depends only on the local executing process.

**non-local** A procedure is non-local if completion of the operation may require the execution of some MPI procedure on another process. Such an operation may require communication occurring with another user process.

**collective** A procedure is collective if all processes in a process group need to invoke the procedure. A collective call may or may not be synchronizing. Collective calls over the same communicator must be executed in the same order by all members of the process group.

**predefined** A predefined datatype is a datatype with a predefined (constant) name (such as `MPI_INT`, `MPI_FLOAT_INT`, or `MPI_UB`) or a datatype constructed with `MPI_TYPE_CREATE_F90_INTEGER`, `MPI_TYPE_CREATE_F90_REAL`, or `MPI_TYPE_CREATE_F90_COMPLEX`. The former are **named** whereas the latter are **unnamed**.

**derived** A derived datatype is any datatype that is not predefined.

**portable** A datatype is portable, if it is a predefined datatype, or it is derived from a portable datatype using only the type constructors `MPI_TYPE_CONTIGUOUS`, `MPI_TYPE_VECTOR`, `MPI_TYPE_INDEXED`, `MPI_TYPE_CREATE_INDEXED_BLOCK`, `MPI_TYPE_CREATE_SUBARRAY`, `MPI_TYPE_DUP`, and `MPI_TYPE_CREATE_DARRAY`. Such a datatype is portable because all displacements in the datatype are in terms of extents of one predefined datatype. Therefore, if such a datatype fits a data layout in one memory, it will fit the corresponding data layout in another memory, if the same declarations were used, even if the two systems have different architectures. On the other hand, if a datatype was constructed using `MPI_TYPE_CREATE_HINDEXED`, `MPI_TYPE_CREATE_HINDEXED_BLOCK`, `MPI_TYPE_CREATE_HVECTOR` or `MPI_TYPE_CREATE_STRUCT`, then the datatype contains explicit byte displacements (e.g., providing padding to meet alignment restrictions). These displacements are unlikely to be chosen correctly if they fit data layout on one memory, but are used for data layouts on another process, running on a processor with a different architecture.

**equivalent** Two datatypes are equivalent if they appear to have been created with the same sequence of calls (and arguments) and thus have the same typemap. Two equivalent datatypes do not necessarily have the same cached attributes or the same names.

## 2.5 Data Types

### 2.5.1 Opaque Objects

MPI manages **system memory** that is used for buffering messages and for storing internal representations of various MPI objects such as groups, communicators, datatypes, etc. This memory is not directly accessible to the user, and objects stored there are **opaque**: their

size and shape is not visible to the user. Opaque objects are accessed via **handles**, which exist in user space. MPI procedures that operate on opaque objects are passed handle arguments to access these objects. In addition to their use by MPI calls for object access, handles can participate in assignments and comparisons.

In Fortran with `USE mpi` or `INCLUDE 'mpif.h'`, all handles have type `INTEGER`. In Fortran with `USE mpi_f08`, and in C and C++, a different handle type is defined for each category of objects. With Fortran `USE mpi_f08`, the handles are defined as Fortran `BIND(C)` derived types that consist of only one element `INTEGER :: MPI_VAL`. The internal handle value is identical to the Fortran `INTEGER` value used in the `mpi` module and `mpif.h`. The type names are identical to the names in C, except that they are not case sensitive. For example:

```
TYPE, BIND(C) :: MPI_Comm
  INTEGER      :: MPI_VAL
END TYPE MPI_Comm
```

In addition, handles themselves are distinct objects in C++. The C and C++ types must support the use of the assignment and equality operators.

*Advice to implementors.* In Fortran, the handle can be an index into a table of opaque objects in a system table; in C it can be such an index or a pointer to the object. C++ handles can simply “wrap up” a table index or pointer. (*End of advice to implementors.*)

*Rationale.* Since the Fortran integer values are equivalent, applications can easily convert MPI handles between all three supported Fortran methods. For example, an integer communicator handle `COMM` can be converted directly into an exactly equivalent `mpi_f08` communicator handle named `comm_f08` by `comm_f08%MPI_VAL=COMM`, and vice versa. The use of the `INTEGER` defined handles and the `BIND(C)` derived type handles is different: Fortran 2003 (and later) define that `BIND(C)` derived types can be used within user defined common blocks, but it is up to the rules of the companion C compiler how many numerical storage units are used for these `BIND(C)` derived type handles. Most compilers use one unit for both, the `INTEGER` handles and the handles defined as `BIND(C)` derived types. (*End of rationale.*)

*Advice to users.* If a user wants to substitute `mpif.h` or the `mpi` module by the `mpi_f08` module and the application program stores a handle in a Fortran common block then it is necessary to change the Fortran support method in all application routines that use this common block, because the number of numerical storage units of such a handle can be different in the two modules. (*End of advice to users.*)

Opaque objects are allocated and deallocated by calls that are specific to each object type. These are listed in the sections where the objects are described. The calls accept a handle argument of matching type. In an allocate call this is an `OUT` argument that returns a valid reference to the object. In a call to deallocate this is an `INOUT` argument which returns with an “invalid handle” value. MPI provides an “invalid handle” constant for each object type. Comparisons to this constant are used to test for validity of the handle.

A call to a deallocate routine invalidates the handle and marks the object for deallocation. The object is not accessible to the user after the call. However, MPI need not

1 deallocate the object immediately. Any operation pending (at the time of the deallocate)  
 2 that involves this object will complete normally; the object will be deallocated afterwards.

3 An opaque object and its handle are significant only at the process where the object  
 4 was created and cannot be transferred to another process.

5 MPI provides certain predefined opaque objects and predefined, static handles to these  
 6 objects. The user must not free such objects. In C++, this is enforced by declaring the  
 7 handles to these predefined objects to be `static const`.

8  
 9 *Rationale.* This design hides the internal representation used for MPI data struc-  
 10 tures, thus allowing similar calls in C, C++, and Fortran. It also avoids conflicts with  
 11 the typing rules in these languages, and easily allows future extensions of functional-  
 12 ity. The mechanism for opaque objects used here loosely follows the POSIX Fortran  
 13 binding standard.

14 The explicit separation of handles in user space and objects in system space allows  
 15 space-reclaiming and deallocation calls to be made at appropriate points in the user  
 16 program. If the opaque objects were in user space, one would have to be very careful  
 17 not to go out of scope before any pending operation requiring that object completed.  
 18 The specified design allows an object to be marked for deallocation, the user program  
 19 can then go out of scope, and the object itself still persists until any pending operations  
 20 are complete.

21 The requirement that handles support assignment/comparison is made since such  
 22 operations are common. This restricts the domain of possible implementations. The  
 23 alternative would have been to allow handles to have been an arbitrary, opaque type.  
 24 This would force the introduction of routines to do assignment and comparison, adding  
 25 complexity, and was therefore ruled out. (*End of rationale.*)

26  
 27 *Advice to users.* A user may accidentally create a dangling reference by assigning to a  
 28 handle the value of another handle, and then deallocating the object associated with  
 29 these handles. Conversely, if a handle variable is deallocated before the associated  
 30 object is freed, then the object becomes inaccessible (this may occur, for example, if  
 31 the handle is a local variable within a subroutine, and the subroutine is exited before  
 32 the associated object is deallocated). It is the user's responsibility to avoid adding or  
 33 deleting references to opaque objects, except as a result of MPI calls that allocate or  
 34 deallocate such objects. (*End of advice to users.*)

35  
 36 *Advice to implementors.* The intended semantics of opaque objects is that opaque  
 37 objects are separate from one another; each call to allocate such an object copies  
 38 all the information required for the object. Implementations may avoid excessive  
 39 copying by substituting referencing for copying. For example, a derived datatype  
 40 may contain references to its components, rather than copies of its components; a  
 41 call to `MPI_COMM_GROUP` may return a reference to the group associated with the  
 42 communicator, rather than a copy of this group. In such cases, the implementation  
 43 must maintain reference counts, and allocate and deallocate objects in such a way that  
 44 the visible effect is as if the objects were copied. (*End of advice to implementors.*)

## 45 2.5.2 Array Arguments

46  
 47 An MPI call may need an argument that is an array of opaque objects, or an array of  
 48 handles. The array-of-handles is a regular array with entries that are handles to objects

of the same type in consecutive locations in the array. Whenever such an array is used, an additional `len` argument is required to indicate the number of valid entries (unless this number can be derived otherwise). The valid entries are at the beginning of the array; `len` indicates how many of them there are, and need not be the size of the entire array. The same approach is followed for other array arguments. In some cases `NULL` handles are considered valid entries. When a `NULL` argument is desired for an array of statuses, one uses `MPI_STATUSES_IGNORE`.

### 2.5.3 State

MPI procedures use at various places arguments with *state* types. The values of such a data type are all identified by names, and no operation is defined on them. For example, the `MPI_TYPE_CREATE_SUBARRAY` routine has a state argument `order` with values `MPI_ORDER_C` and `MPI_ORDER_FORTRAN`.

### 2.5.4 Named Constants

MPI procedures sometimes assign a special meaning to a special value of a basic type argument; e.g., `tag` is an integer-valued argument of point-to-point communication operations, with a special wild-card value, `MPI_ANY_TAG`. Such arguments will have a range of regular values, which is a proper subrange of the range of values of the corresponding basic type; special values (such as `MPI_ANY_TAG`) will be outside the regular range. The range of regular values, such as `tag`, can be queried using environmental inquiry functions (Chapter 7 of the MPI-1 document). The range of other values, such as `source`, depends on values given by other MPI routines (in the case of `source` it is the communicator size).

MPI also provides predefined named constant handles, such as `MPI_COMM_WORLD`.

All named constants, with the exceptions noted below for Fortran, can be used in initialization expressions or assignments, but not necessarily in array declarations or as labels in C/C++ `switch` or Fortran `select/case` statements. This implies named constants to be link-time but not necessarily compile-time constants. The named constants listed below are required to be compile-time constants in both C/C++ and Fortran. These constants do not change values during execution. Opaque objects accessed by constant handles are defined and do not change value between MPI initialization (`MPI_INIT`) and MPI completion (`MPI_FINALIZE`). The handles themselves are constants and can be also used in initialization expressions or assignments.

The constants that are required to be compile-time constants (and can thus be used for array length declarations and labels in C/C++ `switch` and Fortran `case/select` statements) are:

```
MPI_MAX_PROCESSOR_NAME
MPI_MAX_LIBRARY_VERSION_STRING
MPI_MAX_ERROR_STRING
MPI_MAX_DATAREP_STRING
MPI_MAX_INFO_KEY
MPI_MAX_INFO_VAL
MPI_MAX_OBJECT_NAME
MPI_MAX_PORT_NAME
MPI_STATUS_SIZE (Fortran only)
MPI_ADDRESS_KIND (Fortran only)
```

ticket204.

ticket265.1.



1 **MPI\_COUNT\_KIND** (Fortran only)  
 2 **MPI\_INTEGER\_KIND** (Fortran only)  
 3 **MPI\_OFFSET\_KIND** (Fortran only)  
 4 **MPI\_SUBARRAYS\_SUPPORTED** (Fortran only)  
 5 **MPI\_ASYNC\_PROTECTS\_NONBLOCKING** (Fortran only)

6 and their C++ counterparts where appropriate.

7 The constants that cannot be used in initialization expressions or assignments in For-  
 8 tran are:

9 **MPI\_BOTTOM**  
 10 **MPI\_STATUS\_IGNORE**  
 11 **MPI\_STATUSES\_IGNORE**  
 12 **MPI\_ERRCODES\_IGNORE**  
 13 **MPI\_IN\_PLACE**  
 14 **MPI\_ARGV\_NULL**  
 15 **MPI\_ARGVS\_NULL**  
 16 **MPI\_UNWEIGHTED**

17  
 18 *Advice to implementors.* In Fortran the implementation of these special constants  
 19 may require the use of language constructs that are outside the Fortran standard.  
 20 Using special values for the constants (e.g., by defining them through **PARAMETER**  
 21 statements) is not possible because an implementation cannot distinguish these val-  
 22 ues from [legal]valid data. Typically, these constants are implemented as predefined  
 23 static variables (e.g., a variable in an MPI-declared **COMMON** block), relying on the fact  
 24 that the target compiler passes data by address. Inside the subroutine, this address  
 25 can be extracted by some mechanism outside the Fortran standard (e.g., by Fortran  
 26 extensions or by implementing the function in C). (*End of advice to implementors.*)

## 28 2.5.5 Choice

29 MPI functions sometimes use arguments with a *choice* (or union) data type. Distinct calls to  
 30 the same routine may pass by reference actual arguments of different types. The mechanism  
 31 for providing such arguments will differ from language to language. For Fortran with the  
 32 include file **mpif.h** or the **mpi** module, the document uses **<type>** to represent a choice  
 33 variable; with the Fortran **mpi\_f08** module, such arguments are declared with the Fortran  
 34 2008 + TR 29113 syntax **TYPE(\*), DIMENSION(..)**; for C and C++, we use **void \***.  
 35

36  
 37 *Advice to implementors.* Implementors can freely choose how to implement choice  
 38 arguments in the **mpi** module, e.g., with a non-standard compiler-dependent method  
 39 that has the quality of the call mechanism in the implicit Fortran interfaces, or with  
 40 the method defined for the **mpi\_f08** module. See details in Section 16.2.1 on page 640.  
 41 (*End of advice to implementors.*)

## 43 2.5.6 Addresses

44  
 45 Some MPI procedures use *address* arguments that represent an absolute address in the  
 46 calling program. The datatype of such an argument is **MPI\_Aint** in C, **MPI::Aint** in C++  
 47 and **INTEGER (KIND=MPI\_ADDRESS\_KIND)** in Fortran. These types must have the same  
 48 width and encode address values in the same manner such that address values in one



language may be passed directly to another language without conversion. There is the MPI constant `MPI_BOTTOM` to indicate the start of the address range.

### 2.5.7 File Offsets

For I/O there is a need to give the size, displacement, and offset into a file. These quantities can easily be larger than 32 bits which can be the default size of a Fortran integer. To overcome this, these quantities are declared to be `INTEGER (KIND=MPI_OFFSET_KIND)` in Fortran. In C one uses `MPI_Offset` whereas in C++ one uses `MPI::Offset`. These types must have the same width and encode address values in the same manner such that offset values in one language may be passed directly to another language without conversion.

### 2.5.8 Counts

As described above, MPI defines types (e.g., `MPI_Aint`) to address locations within memory and other types (e.g., `MPI_Offset`) to address locations within files. In addition, some MPI procedures use *count* arguments that represent a number of MPI datatypes on which to operate. At times, one needs a single type that can be used to address locations within either memory or files as well as express *count* values, and that type is `MPI_Count` in C and `INTEGER (KIND=MPI_COUNT_KIND)` in Fortran. These types must have the same width and encode values in the same manner such that count values in one language may be passed directly to another language without conversion. The size of the `MPI_Count` type is determined by the MPI implementation with the restriction that it must be minimally capable of encoding any value that may be stored in a variable of type `int`, `MPI_Aint`, or `MPI_Offset` in C and of type `INTEGER`, `INTEGER (KIND=MPI_ADDRESS_KIND)`, or `INTEGER (KIND=MPI_OFFSET_KIND)` in Fortran.

*Rationale.* Count values logically need to be large enough to encode any value used for expressing element counts, type maps in memory, type maps in file views, etc. For backward compatibility reasons, many MPI routines still use `int` in C and `INTEGER` in Fortran as the type of count arguments. (*End of rationale.*)

## 2.6 Language Binding

This section defines the rules for MPI language binding in general and for Fortran, ISO C, and C++, in particular. (Note that ANSI C has been replaced by ISO C.) The C++ language bindings have been deprecated. Defined here are various object representations, as well as the naming conventions used for expressing this standard. The actual calling sequences are defined elsewhere.

MPI bindings are for Fortran 90 or later, though they [are] were originally designed to be usable in Fortran 77 environments. With the `mpi_f08` module, two new Fortran features, *assumed type* and *assumed rank*, are also required, see Section 2.5.5 on page 16.

Since the word `PARAMETER` is a keyword in the Fortran language, we use the word “argument” to denote the arguments to a subroutine. These are normally referred to as parameters in C and C++, however, we expect that C and C++ programmers will understand the word “argument” (which has no specific meaning in C/C++), thus allowing us to avoid unnecessary confusion for Fortran programmers.

Since Fortran is case insensitive, linkers may use either lower case or upper case when resolving Fortran names. Users of case sensitive languages should avoid the “mpi\_” and “pmpi\_” prefixes.

## 2.6.1 Deprecated Names and Functions

A number of chapters refer to deprecated or replaced MPI-1 constructs. These are constructs that continue to be part of the MPI standard, as documented in Chapter 15, but that users are recommended not to continue using, since better solutions were provided **beginning** with MPI-2. For example, the Fortran binding for MPI-1 functions that have address arguments uses `INTEGER`. This is not consistent with the C binding, and causes problems on machines with 32 bit `INTEGER`s and 64 bit addresses. In MPI-2, these functions were given new names with new bindings for the address arguments. The use of the old functions is deprecated. For consistency, here and in a few other cases, new C functions are also provided, even though the new functions are equivalent to the old functions. The old names are deprecated. Another example is provided by the MPI-1 predefined datatypes `MPI_UB` and `MPI_LB`. They are deprecated, since their use is awkward and error-prone. The MPI-2 function `MPI_TYPE_CREATE_RESIZED` provides a more convenient mechanism to achieve the same effect.

Table 2.1 shows a list of all of the deprecated constructs. Note that the constants `MPI_LB` and `MPI_UB` are replaced by the function `MPI_TYPE_CREATE_RESIZED`; this is because their principal use was as input datatypes to `MPI_TYPE_STRUCT` to create resized datatypes. Also note that some C typedefs and Fortran subroutine names are included in this list; they are the types of callback functions.

## 2.6.2 Fortran Binding Issues

Originally, MPI-1.1 provided bindings for Fortran 77. These bindings are retained, but they are now interpreted in the context of the Fortran 90 standard. MPI can still be used with most Fortran 77 compilers, as noted below. When the term “Fortran” is used it means Fortran 90 or later; it means Fortran 2008 + TR 29113 and later if the `mpi_f08` module is used.

All MPI names have an `MPI_` prefix, and all characters are capitals. Programs must not declare variables, parameters, or functions with names beginning with the prefix `MPI_`. To avoid conflicting with the profiling interface, programs should also avoid functions with the prefix `PMPI_`. This is mandated to avoid possible name collisions.

All MPI Fortran subroutines have a return code in the last argument. **With `USE mpi_f08`, this last argument is declared as `OPTIONAL`, except for user-defined callback functions (e.g., `COMM_COPY_ATTR_FUNCTION`) and their predefined callbacks (e.g., `MPI_NULL_COPY_FN`).** A few MPI operations which are functions do not have the return code argument. The return code value for successful completion is `MPI_SUCCESS`. Other error codes are implementation dependent; see the error codes in Chapter 8 and Annex A.

Constants representing the maximum length of a string are one smaller in Fortran than in C and C++ as discussed in Section 16.3.9.

Handles are represented in Fortran as `INTEGER`s, or as a `BIND(C)` derived type with the `mpi_f08` module; see Section 2.5.1 on page 12. Binary-valued variables are of type `LOGICAL`.

Array arguments are indexed from one.

The older MPI Fortran binding [is]s (`mpif.h` and `use mpi`) are inconsistent with the

Deprecated	MPI-2 Replacement
MPI_ADDRESS	MPI_GET_ADDRESS
MPI_TYPE_HINDEXED	MPI_TYPE_CREATE_HINDEXED
MPI_TYPE_HVECTOR	MPI_TYPE_CREATE_HVECTOR
MPI_TYPE_STRUCT	MPI_TYPE_CREATE_STRUCT
MPI_TYPE_EXTENT	MPI_TYPE_GET_EXTENT
MPI_TYPE_UB	MPI_TYPE_GET_EXTENT
MPI_TYPE_LB	MPI_TYPE_GET_EXTENT
MPI_LB	MPI_TYPE_CREATE_RESIZED
MPI_UB	MPI_TYPE_CREATE_RESIZED
MPI_ERRHANDLER_CREATE	MPI_COMM_CREATE_ERRHANDLER
MPI_ERRHANDLER_GET	MPI_COMM_GET_ERRHANDLER
MPI_ERRHANDLER_SET	MPI_COMM_SET_ERRHANDLER
MPI_Handler_function	MPI_Comm_errhandler_function
MPI_KEYVAL_CREATE	MPI_COMM_CREATE_KEYVAL
MPI_KEYVAL_FREE	MPI_COMM_FREE_KEYVAL
MPI_DUP_FN	MPI_COMM_DUP_FN
MPI_NULL_COPY_FN	MPI_COMM_NULL_COPY_FN
MPI_NULL_DELETE_FN	MPI_COMM_NULL_DELETE_FN
MPI_Copy_function	MPI_Comm_copy_attr_function
COPY_FUNCTION	COMM_COPY_ATTR_[ticket250-V.][FN]FUNCTION
MPI_Delete_function	MPI_Comm_delete_attr_function
DELETE_FUNCTION	COMM_DELETE_ATTR_[ticket250-V.][FN]FUNCTION
MPI_ATTR_DELETE	MPI_COMM_DELETE_ATTR
MPI_ATTR_GET	MPI_COMM_GET_ATTR
MPI_ATTR_PUT	MPI_COMM_SET_ATTR

Table 2.1: Deprecated constructs

Fortran[ 90] standard in several respects. These inconsistencies, such as register optimization problems, have implications for user codes that are discussed in detail in Section 16.2.16. [They are also inconsistent with Fortran 77.]

### 2.6.3 C Binding Issues

We use the ISO C declaration format. All MPI names have an `MPI_` prefix, defined constants are in all capital letters, and defined types and functions have one capital letter after the prefix. Programs must not declare variables or functions with names beginning with the prefix `MPI_`. To support the profiling interface, programs should not declare functions with names beginning with the prefix `PMPI_`.

The definition of named constants, function prototypes, and type definitions must be supplied in an include file `mpi.h`.

Almost all C functions return an error code. The successful return code will be `MPI_SUCCESS`, but failure return codes are implementation dependent.

Type declarations are provided for handles to each category of opaque objects.

Array arguments are indexed from zero.

Logical flags are integers with value 0 meaning “false” and a non-zero value meaning

“true.”

Choice arguments are pointers of type `void *`.

Address arguments are of MPI defined type `MPI_Aint`. File displacements are of type `MPI_Offset`. `MPI_Aint` is defined to be an integer of the size needed to hold any valid address on the target architecture. `MPI_Offset` is defined to be an integer of the size needed to hold any valid file size on the target architecture.

## 2.6.4 C++ Binding Issues

The C++ language bindings have been deprecated. There are places in the standard that give rules for C and not for C++. In these cases, the C rule should be applied to the C++ case, as appropriate. In particular, the values of constants given in the text are the ones for C and Fortran. A cross index of these with the C++ names is given in Annex A.

We use the ISO C++ declaration format. All MPI names are declared within the scope of a namespace called `MPI` and therefore are referenced with an `MPI::` prefix. Defined constants are in all capital letters, and class names, defined types, and functions have only their first letter capitalized. Programs must not declare variables or functions in the `MPI` namespace. This is mandated to avoid possible name collisions.

The definition of named constants, function prototypes, and type definitions must be supplied in an include file `mpi.h`.

*Advice to implementors.* The file `mpi.h` may contain both the C and C++ definitions. Usually one can simply use the defined value (generally `__cplusplus`, but not required) to see if one is using C++ to protect the C++ definitions. It is possible that a C compiler will require that the source protected this way be `[legal]valid` C code. In this case, all the C++ definitions can be placed in a different include file and the `#include` directive can be used to include the necessary C++ definitions in the `mpi.h` file. (*End of advice to implementors.*)

C++ functions that create objects or return information usually place the object or information in the return value. Since the language neutral prototypes of MPI functions include the C++ return value as an OUT parameter, semantic descriptions of MPI functions refer to the C++ return value by that parameter name. The remaining C++ functions return `void`.

In some circumstances, MPI permits users to indicate that they do not want a return value. For example, the user may indicate that the status is not filled in. Unlike C and Fortran where this is achieved through a special input value, in C++ this is done by having two bindings where one has the optional argument and one does not.

C++ functions do not return error codes. If the default error handler has been set to `MPI::ERRORS_THROW_EXCEPTIONS`, the C++ exception mechanism is used to signal an error by throwing an `MPI::Exception` object.

It should be noted that the default error handler (i.e., `MPI::ERRORS_ARE_FATAL`) on a given type has not changed. User error handlers are also permitted. `MPI::ERRORS_RETURN` simply returns control to the calling function; there is no provision for the user to retrieve the error code.

User callback functions that return integer error codes should not throw exceptions; the returned error will be handled by the MPI implementation by invoking the appropriate error handler.

*Advice to users.* C++ programmers that want to handle MPI errors on their own should use the `MPI::ERRORS_THROW_EXCEPTIONS` error handler, rather than `MPI::ERRORS_RETURN`, that is used for that purpose in C. Care should be taken using exceptions in mixed language situations. (*End of advice to users.*)

Opaque object handles must be objects in themselves, and have the assignment and equality operators overridden to perform semantically like their C and Fortran counterparts.

Array arguments are indexed from zero.

Logical flags are of type `bool`.

Choice arguments are pointers of type `void *`.

Address arguments are of MPI-defined integer type `MPI::Aint`, defined to be an integer of the size needed to hold any valid address on the target architecture. Analogously, `MPI::Offset` is an integer to hold file offsets.

Most MPI functions are methods of MPI C++ classes. MPI class names are generated from the language neutral MPI types by dropping the `MPI_` prefix and scoping the type within the MPI namespace. For example, `MPI_DATATYPE` becomes `MPI::Datatype`.

The names of MPI functions generally follow the naming rules given. In some circumstances, the MPI function is related to a function defined already for MPI-1 with a name that does not follow the naming conventions. In this circumstance, the language neutral name is in analogy to the MPI name even though this gives an MPI-2 name that violates the naming conventions. The C and Fortran names are the same as the language neutral name in this case. However, the C++ names do reflect the naming rules and can differ from the C and Fortran names. Thus, the analogous name in C++ to the MPI name may be different than the language neutral name. This results in the C++ name differing from the language neutral name. An example of this is the language neutral name of `MPI_FINALIZED` and a C++ name of `MPI::Is_finalized`.

In C++, function `typedefs` are made publicly within appropriate classes. However, these declarations then become somewhat cumbersome, as with the following:

```
{typedef MPI::Grequest::Query_function(); (binding deprecated, see Section 15.2)}
```

would look like the following:

```
namespace MPI {
    class Request {
        // ...
    };

    class Grequest : public MPI::Request {
        // ...
        typedef Query_function(void* extra_state, MPI::Status& status);
    };
};
```

Rather than including this scaffolding when declaring C++ `typedefs`, we use an abbreviated form. In particular, we explicitly indicate the class and namespace scope for the `typedef` of the function. Thus, the example above is shown in the text as follows:

```
typedef int MPI::Grequest::Query_function(void* extra_state,
                                           MPI::Status& status)
```

The C++ bindings presented in Annex A.5 and throughout this document were generated by applying a simple set of name generation rules to the MPI function specifications. While these guidelines may be sufficient in most cases, they may not be suitable for all situations. In cases of ambiguity or where a specific semantic statement is desired, these guidelines may be superseded as the situation dictates.

1. All functions, types, and constants are declared within the scope of a `namespace` called `MPI`.
2. Arrays of MPI handles are always left in the argument list (whether they are IN or OUT arguments).
3. If the argument list of an MPI function contains a scalar IN handle, and it makes sense to define the function as a method of the object corresponding to that handle, the function is made a member function of the corresponding MPI class. The member functions are named according to the corresponding MPI function name, but without the “MPI\_” prefix and without the object name prefix (if applicable). In addition:
  - (a) The scalar IN handle is dropped from the argument list, and `this` corresponds to the dropped argument.
  - (b) The function is declared `const`.
4. MPI functions are made into class functions (static) when they belong on a class but do not have a unique scalar IN or INOUT parameter of that class.
5. If the argument list contains a single OUT argument that is not of type `MPI_STATUS` (or an array), that argument is dropped from the list and the function returns that value.

**Example 2.1** The C++ binding for `MPI_COMM_SIZE` is  
`int MPI::Comm::Get_size(void) const.`

6. If there are multiple OUT arguments in the argument list, one is chosen as the return value and is removed from the list.
7. If the argument list does not contain any OUT arguments, the function returns `void`.

**Example 2.2** The C++ binding for `MPI_REQUEST_FREE` is  
`void MPI::Request::Free(void)`

8. MPI functions to which the above rules do not apply are not members of any class, but are defined in the `MPI` namespace.

**Example 2.3** The C++ binding for `MPI_BUFFER_ATTACH` is  
`void MPI::Attach_buffer(void* buffer, int size).`

9. All class names, defined types, and function names have only their first letter capitalized. Defined constants are in all capital letters.
10. Any IN pointer, reference, or array argument must be declared `const`.

11. Handles are passed by reference.
12. Array arguments are denoted with square brackets (`[]`), not pointers, as this is more semantically precise.

### 2.6.5 Functions and Macros

An implementation is allowed to implement `MPI_WTIME`, `MPI_WTICK`, `PMPI_WTIME`, `PMPI_WTICK`, and the handle-conversion functions (`MPI_Group_f2c`, etc.) in Section 16.3.4, and no others, as macros in C.

*Advice to implementors.* Implementors should document which routines are implemented as macros. (*End of advice to implementors.*)

*Advice to users.* If these routines are implemented as macros, they will not work with the MPI profiling interface. (*End of advice to users.*)

## 2.7 Processes

An MPI program consists of autonomous processes, executing their own code, in an MIMD style. The codes executed by each process need not be identical. The processes communicate via calls to MPI communication primitives. Typically, each process executes in its own address space, although shared-memory implementations of MPI are possible.

This document specifies the behavior of a parallel program assuming that only MPI calls are used. The interaction of an MPI program with other possible means of communication, I/O, and process management is not specified. Unless otherwise stated in the specification of the standard, MPI places no requirements on the result of its interaction with external mechanisms that provide similar or equivalent functionality. This includes, but is not limited to, interactions with external mechanisms for process control, shared and remote memory access, file system access and control, interprocess communication, process signaling, and terminal I/O. High quality implementations should strive to make the results of such interactions intuitive to users, and attempt to document restrictions where deemed necessary.

*Advice to implementors.* Implementations that support such additional mechanisms for functionality supported within MPI are expected to document how these interact with MPI. (*End of advice to implementors.*)

The interaction of MPI and threads is defined in Section 12.4.

## 2.8 Error Handling

MPI provides the user with reliable message transmission. A message sent is always received correctly, and the user does not need to check for transmission errors, time-outs, or other error conditions. In other words, MPI does not provide mechanisms for dealing with failures in the communication system. If the MPI implementation is built on an unreliable underlying mechanism, then it is the job of the implementor of the MPI subsystem to insulate the user from this unreliability, or to reflect unrecoverable errors as failures. Whenever possible,



such failures will be reflected as errors in the relevant communication call. Similarly, MPI itself provides no mechanisms for handling processor failures.

Of course, MPI programs may still be erroneous. A **program error** can occur when an MPI call is made with an incorrect argument (non-existing destination in a send operation, buffer too small in a receive operation, etc.). This type of error would occur in any implementation. In addition, a **resource error** may occur when a program exceeds the amount of available system resources (number of pending messages, system buffers, etc.). The occurrence of this type of error depends on the amount of available resources in the system and the resource allocation mechanism used; this may differ from system to system. A high-quality implementation will provide generous limits on the important resources so as to alleviate the portability problem this represents.

In C and Fortran, almost all MPI calls return a code that indicates successful completion of the operation. Whenever possible, MPI calls return an error code if an error occurred during the call. By default, an error detected during the execution of the MPI library causes the parallel computation to abort, except for file operations. However, MPI provides mechanisms for users to change this default and to handle recoverable errors. The user may specify that no error is fatal, and handle error codes returned by MPI calls by himself or herself. Also, the user may provide his or her own error-handling routines, which will be invoked whenever an MPI call returns abnormally. The MPI error handling facilities are described in Section 8.3. The return values of C++ functions are not error codes. If the default error handler has been set to `MPI::ERRORS_THROW_EXCEPTIONS`, the C++ exception mechanism is used to signal an error by throwing an `MPI::Exception` object. See also Section 16.1.8 on page 636.

Several factors limit the ability of MPI calls to return with meaningful error codes when an error occurs. MPI may not be able to detect some errors; other errors may be too expensive to detect in normal execution mode; finally some errors may be “catastrophic” and may prevent MPI from returning control to the caller in a consistent state.

Another subtle issue arises because of the nature of asynchronous communications: MPI calls may initiate operations that continue asynchronously after the call returned. Thus, the operation may return with a code indicating successful completion, yet later cause an error exception to be raised. If there is a subsequent call that relates to the same operation (e.g., a call that verifies that an asynchronous operation has completed) then the error argument associated with this call will be used to indicate the nature of the error. In a few cases, the error may occur after all calls that relate to the operation have completed, so that no error value can be used to indicate the nature of the error (e.g., an error on the receiver in a send with the ready mode). Such an error must be treated as fatal, since information cannot be returned for the user to recover from it.

This document does not specify the state of a computation after an erroneous MPI call has occurred. The desired behavior is that a relevant error code be returned, and the effect of the error be localized to the greatest possible extent. E.g., it is highly desirable that an erroneous receive call will not cause any part of the receiver’s memory to be overwritten, beyond the area specified for receiving the message.

Implementations may go beyond this document in supporting in a meaningful manner MPI calls that are defined here to be erroneous. For example, MPI specifies strict type matching rules between matching send and receive operations: it is erroneous to send a floating point variable and receive an integer. Implementations may go beyond these type matching rules, and provide automatic type conversion in such situations. It will be helpful to generate warnings for such non-conforming behavior.



MPI defines a way for users to create new error codes as defined in Section 8.5.

## 2.9 Implementation Issues

There are a number of areas where an MPI implementation may interact with the operating environment and system. While MPI does not mandate that any services (such as signal handling) be provided, it does strongly suggest the behavior to be provided if those services are available. This is an important point in achieving portability across platforms that provide the same set of services.

### 2.9.1 Independence of Basic Runtime Routines

MPI programs require that library routines that are part of the basic language environment (such as `write` in Fortran and `printf` and `malloc` in ISO C) and are executed after `MPI_INIT` and before `MPI_FINALIZE` operate independently and that their *completion* is independent of the action of other processes in an MPI program.

Note that this in no way prevents the creation of library routines that provide parallel services whose operation is collective. However, the following program is expected to complete in an ISO C environment regardless of the size of `MPI_COMM_WORLD` (assuming that `printf` is available at the executing nodes).

```
int rank;
MPI_Init((void *)0, (void *)0);
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
if (rank == 0) printf("Starting program\n");
MPI_Finalize();
```

The corresponding Fortran and C++ programs are also expected to complete.

An example of what is *not* required is any particular ordering of the action of these routines when called by several tasks. For example, MPI makes neither requirements nor recommendations for the output from the following program (again assuming that I/O is available at the executing nodes).

```
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
printf("Output from task rank %d\n", rank);
```

In addition, calls that fail because of resource exhaustion or other error are not considered a violation of the requirements here (however, they are required to complete, just not to complete successfully).

### 2.9.2 Interaction with Signals

MPI does not specify the interaction of processes with signals and does not require that MPI be signal safe. The implementation may reserve some signals for its own use. It is required that the implementation document which signals it uses, and it is strongly recommended that it not use `SIGALRM`, `SIGFPE`, or `SIGIO`. Implementations may also prohibit the use of MPI calls from within signal handlers.

In multithreaded environments, users can avoid conflicts between signals and the MPI library by catching signals only on threads that do not execute MPI calls. High quality single-threaded implementations will be signal safe: an MPI call suspended by a signal will resume and complete normally after the signal is handled.

## 2.10 Examples

The examples in this document are for illustration purposes only. They are not intended to specify the standard. Furthermore, the examples have not been carefully checked or verified.

## Chapter 3

# Point-to-Point Communication

### 3.1 Introduction

Sending and receiving of messages by processes is the basic MPI communication mechanism. The basic point-to-point communication operations are **send** and **receive**. Their use is illustrated in the example below.

```
#include "mpi.h"
int main( int argc, char **argv )
{
    char message[20];
    int myrank;
    MPI_Status status;
    MPI_Init( &argc, &argv );
    MPI_Comm_rank( MPI_COMM_WORLD, &myrank );
    if (myrank == 0) /* code for process zero */
    {
        strcpy(message,"Hello, there");
        MPI_Send(message, strlen(message)+1, MPI_CHAR, 1, 99, MPI_COMM_WORLD);
    }
    else if (myrank == 1) /* code for process one */
    {
        MPI_Recv(message, 20, MPI_CHAR, 0, 99, MPI_COMM_WORLD, &status);
        printf("received :%s:\n", message);
    }
    MPI_Finalize();
    return 0;
}
```

In this example, process zero (`myrank = 0`) sends a message to process one using the **send** operation `MPI_SEND`. The operation specifies a **send buffer** in the sender memory from which the message data is taken. In the example above, the send buffer consists of the storage containing the variable `message` in the memory of process zero. The location, size and type of the send buffer are specified by the first three parameters of the send operation. The message sent will contain the 13 characters of this variable. In addition, the send operation associates an **envelope** with the message. This envelope specifies the

message destination and contains distinguishing information that can be used by the **receive** operation to select a particular message. The last three parameters of the send operation, along with the rank of the sender, specify the envelope for the message sent. Process one (**myrank** = 1) receives this message with the **receive** operation **MPI\_RECV**. The message to be received is selected according to the value of its envelope, and the message data is stored into the **receive buffer**. In the example above, the receive buffer consists of the storage containing the string **message** in the memory of process one. The first three parameters of the receive operation specify the location, size and type of the receive buffer. The next three parameters are used for selecting the incoming message. The last parameter is used to return information on the message just received.

The next sections describe the blocking send and receive operations. We discuss send, receive, blocking communication semantics, type matching requirements, type conversion in heterogeneous environments, and more general communication modes. Nonblocking communication is addressed next, followed by **probing and canceling a message**, channel-like constructs and send-receive operations, [ **Nonblocking communication is addressed next, followed by channel-like constructs and send-receive operations,** ] ending with a description of the “dummy” process, **MPI\_PROC\_NULL**.

## 3.2 Blocking Send and Receive Operations

### 3.2.1 Blocking Send

The syntax of the blocking send operation is given below.

**MPI\_SEND**(buf, count, datatype, dest, tag, comm)

IN	buf	initial address of send buffer (choice)
IN	count	number of elements in send buffer (non-negative integer)
IN	datatype	datatype of each send buffer element (handle)
IN	dest	rank of destination (integer)
IN	tag	message tag (integer)
IN	comm	communicator (handle)

```
int MPI_Send(const void* buf, int count, MPI_Datatype datatype, int dest,
             int tag, MPI_Comm comm)
```

```
MPI_Send(buf, count, datatype, dest, tag, comm, ierror) BIND(C)
  TYPE(*), DIMENSION(..), INTENT(IN) :: buf
  INTEGER, INTENT(IN) :: count, dest, tag
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  TYPE(MPI_Comm), INTENT(IN) :: comm
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_SEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, IERROR)
  <type> BUF(*)
  INTEGER COUNT, DATATYPE, DEST, TAG, COMM, IERROR
```

```
{void MPI::Comm::Send(const void* buf, int count, const
    MPI::Datatype& datatype, int dest, int tag) const(binding
    deprecated, see Section 15.2) }
```

The blocking semantics of this call are described in Section 3.4.

### 3.2.2 Message Data

The send buffer specified by the `MPI_SEND` operation consists of `count` successive entries of the type indicated by `datatype`, starting with the entry at address `buf`. Note that we specify the message length in terms of number of *elements*, not number of *bytes*. The former is machine independent and closer to the application level.

The data part of the message consists of a sequence of `count` values, each of the type indicated by `datatype`. `count` may be zero, in which case the data part of the message is empty. The basic datatypes that can be specified for message data values correspond to the basic datatypes of the host language. Possible values of this argument for Fortran and the corresponding Fortran types are listed in Table 3.1.

MPI datatype	Fortran datatype
<code>MPI_INTEGER</code>	<code>INTEGER</code>
<code>MPI_REAL</code>	<code>REAL</code>
<code>MPI_DOUBLE_PRECISION</code>	<code>DOUBLE PRECISION</code>
<code>MPI_COMPLEX</code>	<code>COMPLEX</code>
<code>MPI_LOGICAL</code>	<code>LOGICAL</code>
<code>MPI_CHARACTER</code>	<code>CHARACTER(1)</code>
<code>MPI_BYTE</code>	
<code>MPI_PACKED</code>	

Table 3.1: Predefined MPI datatypes corresponding to Fortran datatypes

Possible values for this argument for C and the corresponding C types are listed in Table 3.2.

The datatypes `MPI_BYTE` and `MPI_PACKED` do not correspond to a Fortran or C datatype. A value of type `MPI_BYTE` consists of a byte (8 binary digits). A byte is uninterpreted and is different from a character. Different machines may have different representations for characters, or may use more than one byte to represent characters. On the other hand, a byte has the same binary value on all machines. The use of the type `MPI_PACKED` is explained in Section 4.2.

MPI requires support of these datatypes, which match the basic datatypes of Fortran and ISO C. Additional MPI datatypes should be provided if the host language has additional data types: `MPI_DOUBLE_COMPLEX` for double precision complex in Fortran declared to be of type `DOUBLE COMPLEX`; `MPI_REAL2`, `MPI_REAL4` and `MPI_REAL8` for Fortran reals, declared to be of type `REAL*2`, `REAL*4` and `REAL*8`, respectively; `MPI_INTEGER1`, `MPI_INTEGER2` and `MPI_INTEGER4` for Fortran integers, declared to be of type `INTEGER*1`, `INTEGER*2` and `INTEGER*4`, respectively; etc.

*Rationale.* One goal of the design is to allow for MPI to be implemented as a library, with no need for additional preprocessing or compilation. Thus, one cannot assume that a communication call has information on the datatype of variables in the

MPI datatype	C datatype
MPI_CHAR	char (treated as printable character)
MPI_SHORT	signed short int
MPI_INT	signed int
MPI_LONG	signed long int
MPI_LONG_LONG_INT	signed long long int
MPI_LONG_LONG (as a synonym)	signed long long int
MPI_SIGNED_CHAR	signed char (treated as integral value)
MPI_UNSIGNED_CHAR	unsigned char (treated as integral value)
MPI_UNSIGNED_SHORT	unsigned short int
MPI_UNSIGNED	unsigned int
MPI_UNSIGNED_LONG	unsigned long int
MPI_UNSIGNED_LONG_LONG	unsigned long long int
MPI_FLOAT	float
MPI_DOUBLE	double
MPI_LONG_DOUBLE	long double
MPI_WCHAR	wchar_t (defined in <stddef.h>) (treated as printable character)
MPI_C_BOOL	_Bool
MPI_INT8_T	int8_t
MPI_INT16_T	int16_t
MPI_INT32_T	int32_t
MPI_INT64_T	int64_t
MPI_UINT8_T	uint8_t
MPI_UINT16_T	uint16_t
MPI_UINT32_T	uint32_t
MPI_UINT64_T	uint64_t
MPI_C_COMPLEX	float _Complex
MPI_C_FLOAT_COMPLEX (as a synonym)	float _Complex
MPI_C_DOUBLE_COMPLEX	double _Complex
MPI_C_LONG_DOUBLE_COMPLEX	long double _Complex
MPI_BYTE	
MPI_PACKED	

Table 3.2: Predefined MPI datatypes corresponding to C datatypes

communication buffer; this information must be supplied by an explicit argument. The need for such datatype information will become clear in Section 3.3.2. (*End of rationale.*)

*Rationale.* The datatypes MPI\_C\_BOOL, MPI\_INT8\_T, MPI\_INT16\_T, MPI\_INT32\_T, MPI\_UINT8\_T, MPI\_UINT16\_T, MPI\_UINT32\_T, MPI\_C\_COMPLEX, MPI\_C\_FLOAT\_COMPLEX, MPI\_C\_DOUBLE\_COMPLEX, and

MPI datatype	C datatype	Fortran datatype	
MPI_AINT	MPI_Aint	INTEGER (KIND=MPI_ADDRESS_KIND)	1
MPI_OFFSET	MPI_Offset	INTEGER (KIND=MPI_OFFSET_KIND)	2
[ticket265.]MPI_COUNT	[ticket265.]MPI_Count	[ticket265.]INTEGER (KIND=MPI_COUNT_KIND)	3

Table 3.3: Predefined MPI datatypes corresponding to both C and Fortran datatypes

MPI\_C\_LONG\_DOUBLE\_COMPLEX have no corresponding C++ bindings. This was intentionally done to avoid potential collisions with the C preprocessor and namespaced C++ names. C++ applications can use the C bindings with no loss of functionality. (*End of rationale.*)

The datatypes MPI\_AINT[ and], MPI\_OFFSET [ ], and MPI\_COUNT correspond to the MPI-defined C types MPI\_Aint[ and], MPI\_Offset[ ], and MPI\_Count and their Fortran equivalents INTEGER (KIND=MPI\_ADDRESS\_KIND)[ and], INTEGER (KIND=MPI\_OFFSET\_KIND), and INTEGER (KIND=MPI\_COUNT\_KIND). This is described in Table 3.3. See Section 16.3.10 for information on interlanguage communication with these types.

### 3.2.3 Message Envelope

In addition to the data part, messages carry information that can be used to distinguish messages and selectively receive them. This information consists of a fixed number of fields, which we collectively call the **message envelope**. These fields are

source  
destination  
tag  
communicator

The message source is implicitly determined by the identity of the message sender. The other fields are specified by arguments in the send operation.

The message destination is specified by the **dest** argument.

The integer-valued message tag is specified by the **tag** argument. This integer can be used by the program to distinguish different types of messages. The range of valid tag values is 0,...,UB, where the value of UB is implementation dependent. It can be found by querying the value of the attribute MPI\_TAG\_UB, as described in Chapter 8. MPI requires that UB be no less than 32767.

The **comm** argument specifies the **communicator** that is used for the send operation. Communicators are explained in Chapter 6; below is a brief summary of their usage.

A communicator specifies the communication context for a communication operation. Each communication context provides a separate “communication universe[:]:” messages are always received within the context they were sent, and messages sent in different contexts do not interfere.

The communicator also specifies the set of processes that share this communication context. This **process group** is ordered and processes are identified by their rank within this group. Thus, the range of valid values for **dest** is 0, ... , n-1, where n is the number of

processes in the group. (If the communicator is an inter-communicator, then destinations are identified by their rank in the remote group. See Chapter 6.)

A predefined communicator `MPI_COMM_WORLD` is provided by MPI. It allows communication with all processes that are accessible after MPI initialization and processes are identified by their rank in the group of `MPI_COMM_WORLD`.

*Advice to users.* Users that are comfortable with the notion of a flat name space for processes, and a single communication context, as offered by most existing communication libraries, need only use the predefined variable `MPI_COMM_WORLD` as the `comm` argument. This will allow communication with all the processes available at initialization time.

Users may define new communicators, as explained in Chapter 6. Communicators provide an important encapsulation mechanism for libraries and modules. They allow modules to have their own disjoint communication universe and their own process numbering scheme. (*End of advice to users.*)

*Advice to implementors.* The message envelope would normally be encoded by a fixed-length message header. However, the actual encoding is implementation dependent. Some of the information (e.g., source or destination) may be implicit, and need not be explicitly carried by messages. Also, processes may be identified by relative ranks, or absolute ids, etc. (*End of advice to implementors.*)

### 3.2.4 Blocking Receive

The syntax of the blocking receive operation is given below.

`MPI_RECV (buf, count, datatype, source, tag, comm, status)`

OUT	buf	initial address of receive buffer (choice)
IN	count	number of elements in receive buffer (non-negative integer)
IN	datatype	datatype of each receive buffer element (handle)
IN	source	rank of source or <code>MPI_ANY_SOURCE</code> (integer)
IN	tag	message tag or <code>MPI_ANY_TAG</code> (integer)
IN	comm	communicator (handle)
OUT	status	status object (Status)

```
int MPI_Recv(void* buf, int count, MPI_Datatype datatype, int source,
            int tag, MPI_Comm comm, MPI_Status *status)
```

```
MPI_Recv(buf, count, datatype, source, tag, comm, status, ierror) BIND(C)
    TYPE(*), DIMENSION(..) :: buf
    INTEGER, INTENT(IN) :: count, source, tag
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Status) :: status
```



```

    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_RECV(BUF, COUNT, DATATYPE, SOURCE, TAG, COMM, STATUS, IERROR)
    <type> BUF(*)
    INTEGER COUNT, DATATYPE, SOURCE, TAG, COMM, STATUS(MPI_STATUS_SIZE),
    IERROR
{void MPI::Comm::Recv(void* buf, int count, const MPI::Datatype& datatype,
    int source, int tag, MPI::Status& status) const(binding
    deprecated, see Section 15.2) }
{void MPI::Comm::Recv(void* buf, int count, const MPI::Datatype& datatype,
    int source, int tag) const(binding deprecated, see Section 15.2) }

```

The blocking semantics of this call are described in Section 3.4.

The receive buffer consists of the storage containing `count` consecutive elements of the type specified by `datatype`, starting at address `buf`. The length of the received message must be less than or equal to the length of the receive buffer. An overflow error occurs if all incoming data does not fit, without truncation, into the receive buffer.

If a message that is shorter than the receive buffer arrives, then only those locations corresponding to the (shorter) message are modified.

*Advice to users.* The `MPI_PROBE` function described in Section 3.8 can be used to receive messages of unknown length. (*End of advice to users.*)

*Advice to implementors.* Even though no specific behavior is mandated by MPI for erroneous programs, the recommended handling of overflow situations is to return in `status` information about the source and tag of the incoming message. The receive operation will return an error code. A quality implementation will also ensure that no memory that is outside the receive buffer will ever be overwritten.

In the case of a message shorter than the receive buffer, MPI is quite strict in that it allows no modification of the other locations. A more lenient statement would allow for some optimizations but this is not allowed. The implementation must be ready to end a copy into the receiver memory exactly at the end of the receive buffer, even if it is an odd address. (*End of advice to implementors.*)

The selection of a message by a receive operation is governed by the value of the message envelope. A message can be received by a receive operation if its envelope matches the `source`, `tag` and `comm` values specified by the receive operation. The receiver may specify a wildcard `MPI_ANY_SOURCE` value for `source`, and/or a wildcard `MPI_ANY_TAG` value for `tag`, indicating that any source and/or tag are acceptable. It cannot specify a wildcard value for `comm`. Thus, a message can be received by a receive operation only if it is addressed to the receiving process, has a matching communicator, has matching source unless `source=MPI_ANY_SOURCE` in the pattern, and has a matching tag unless `tag=MPI_ANY_TAG` in the pattern.

The message tag is specified by the `tag` argument of the receive operation. The argument `source`, if different from `MPI_ANY_SOURCE`, is specified as a rank within the process group associated with that same communicator (remote process group, for intercommunicators). Thus, the range of valid values for the `source` argument is  $\{0, \dots, n-1\} \cup \{\text{MPI\_ANY\_SOURCE}\}$ , where  $n$  is the number of processes in this group.

Note the asymmetry between send and receive operations: A receive operation may accept messages from an arbitrary sender, on the other hand, a send operation must specify a unique receiver. This matches a “push” communication mechanism, where data transfer is effected by the sender (rather than a “pull” mechanism, where data transfer is effected by the receiver).

Source = destination is allowed, that is, a process can send a message to itself. (However, it is unsafe to do so with the blocking send and receive operations described above, since this may lead to deadlock. See Section 3.5.)

*Advice to implementors.* Message context and other communicator information can be implemented as an additional tag field. It differs from the regular message tag in that wild card matching is not allowed on this field, and that value setting for this field is controlled by communicator manipulation functions. (*End of advice to implementors.*)

### 3.2.5 Return Status

The source or tag of a received message may not be known if wildcard values were used in the receive operation. Also, if multiple requests are completed by a single MPI function (see Section 3.7.5), a distinct error code may need to be returned for each request. The information is returned by the `status` argument of `MPI_RECV`. The type of `status` is MPI-defined. Status variables need to be explicitly allocated by the user, that is, they are not system objects.

In C, `status` is a structure that contains three fields named `MPI_SOURCE`, `MPI_TAG`, and `MPI_ERROR`; the structure may contain additional fields. Thus, `status.MPI_SOURCE`, `status.MPI_TAG` and `status.MPI_ERROR` contain the source, tag, and error code, respectively, of the received message.

In Fortran with `USE mpi` or `INCLUDE 'mpif.h'`, `status` is an array of `INTEGER`s of size `MPI_STATUS_SIZE`. The constants `MPI_SOURCE`, `MPI_TAG` and `MPI_ERROR` are the indices of the entries that store the source, tag and error fields. Thus, `status(MPI_SOURCE)`, `status(MPI_TAG)` and `status(MPI_ERROR)` contain, respectively, the source, tag and error code of the received message.

With Fortran `USE mpi_f08`, `status` is defined as the Fortran `BIND(C)` derived type `TYPE(MPI_Status)` containing three public fields named `MPI_SOURCE`, `MPI_TAG`, and `MPI_ERROR`. `TYPE(MPI_Status)` may contain additional, implementation-specific fields. Thus, `status%MPI_SOURCE`, `status%MPI_TAG` and `status%MPI_ERROR` contain the source, tag, and error code of a received message respectively. Additionally, within both the `mpi` and the `mpi_f08` modules, the constants `MPI_STATUS_SIZE`, `MPI_SOURCE`, `MPI_TAG`, `MPI_ERROR`, and `TYPE(MPI_Status)` are defined to allow conversion between both `status` representations. Conversion routines are provided in Section 16.3.5 on page 696.

*Rationale.* The Fortran `TYPE(MPI_Status)` is defined as a `BIND(C)` derived type so that it can be used at any location where the status integer array representation can be used, e.g., in user defined common blocks. (*End of rationale.*)

*Rationale.* It is allowed to have the same name (e.g., `MPI_SOURCE`) defined as a constant (e.g., Fortran parameter) and as a field of a derived type. (*End of rationale.*)

In C++, the `status` object is handled through the following methods:

```

{int MPI::Status::Get_source() const(binding deprecated, see Section 15.2) }
{void MPI::Status::Set_source(int source)(binding deprecated, see Section 15.2) }
{int MPI::Status::Get_tag() const(binding deprecated, see Section 15.2) }
{void MPI::Status::Set_tag(int tag)(binding deprecated, see Section 15.2) }
{int MPI::Status::Get_error() const(binding deprecated, see Section 15.2) }
{void MPI::Status::Set_error(int error)(binding deprecated, see Section 15.2) }

```

In general, message-passing calls do not modify the value of the error code field of status variables. This field may be updated only by the functions in Section 3.7.5 which return multiple statuses. The field is updated if and only if such function returns with an error code of MPI\_ERR\_IN\_STATUS.

*Rationale.* The error field in status is not needed for calls that return only one status, such as MPI\_WAIT, since that would only duplicate the information returned by the function itself. The current design avoids the additional overhead of setting it, in such cases. The field is needed for calls that return multiple statuses, since each request may have had a different failure. (*End of rationale.*)

The status argument also returns information on the length of the message received. However, this information is not directly available as a field of the status variable and a call to MPI\_GET\_COUNT is required to “decode” this information.

MPI\_GET\_COUNT(status, datatype, count)

IN	status	return status of receive operation (Status)
IN	datatype	datatype of each receive buffer entry (handle)
OUT	count	number of received entries (integer)

```

int MPI_Get_count(const MPI_Status *status, MPI_Datatype datatype,
                  int *count)

```

```

MPI_Get_count(status, datatype, count, ierror) BIND(C)
  TYPE(MPI_Status), INTENT(IN) :: status
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  INTEGER, INTENT(OUT) :: count
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_GET_COUNT(STATUS, DATATYPE, COUNT, IERROR)
  INTEGER STATUS(MPI_STATUS_SIZE), DATATYPE, COUNT, IERROR

```

```

{int MPI::Status::Get_count(const MPI::Datatype& datatype) const(binding
  deprecated, see Section 15.2) }

```

Returns the number of entries received. (Again, we count *entries*, each of type *datatype*, not *bytes*.) The *datatype* argument should match the argument provided by the receive call that set the status variable. If the number of entries received exceeds the limits of the count parameter, then MPI\_GET\_COUNT sets the value of count to MPI\_UNDEFINED. [(We

shall later see, in Section 4.1.11, that `MPI_GET_COUNT` may return, in certain situations, the value `MPI_UNDEFINED`.)] There are other situations where the value of `count` can be set to `MPI_UNDEFINED`; see Section 4.1.11.

*Rationale.* Some message-passing libraries use `INOUT count`, `tag` and `source` arguments, thus using them both to specify the selection criteria for incoming messages and return the actual envelope values of the received message. The use of a separate status argument prevents errors that are often attached with `INOUT` argument (e.g., using the `MPI_ANY_TAG` constant as the tag in a receive). Some libraries use calls that refer implicitly to the “last message received.” This is not thread safe.

The `datatype` argument is passed to `MPI_GET_COUNT` so as to improve performance. A message might be received without counting the number of elements it contains, and the count value is often not needed. Also, this allows the same function to be used after a call to `MPI_PROBE` or `MPI_IProbe`. With a status from `MPI_PROBE` or `MPI_IProbe`, the same datatypes are allowed as in a call to `MPI_RECV` to receive this message. (*End of rationale.*)

The value returned as the `count` argument of `MPI_GET_COUNT` for a datatype of length zero where zero bytes have been transferred is zero. If the number of bytes transferred is greater than zero, `MPI_UNDEFINED` is returned.

*Rationale.* Zero-length datatypes may be created in a number of cases. An important case is `MPI_TYPE_CREATE_DARRAY`, where the definition of the particular darray results in an empty block on some MPI process. Programs written in an SPMD style will not check for this special case and may want to use `MPI_GET_COUNT` to check the status. (*End of rationale.*)

*Advice to users.* The buffer size required for the receive can be affected by data conversions and by the stride of the receive datatype. In most cases, the safest approach is to use the same datatype with `MPI_GET_COUNT` and the receive. (*End of advice to users.*)

All send and receive operations use the `buf`, `count`, `datatype`, `source`, `dest`, `tag`, `comm` and `status` arguments in the same way as the blocking `MPI_SEND` and `MPI_RECV` operations described in this section.

### 3.2.6 Passing `MPI_STATUS_IGNORE` for Status

Every call to `MPI_RECV` includes a `status` argument, wherein the system can return details about the message received. There are also a number of other MPI calls where `status` is returned. An object of type `MPI_STATUS` is not an MPI opaque object; its structure is declared in `mpi.h` and `mpif.h`, and it exists in the user’s program. In many cases, application programs are constructed so that it is unnecessary for them to examine the `status` fields. In these cases, it is a waste for the user to allocate a status object, and it is particularly wasteful for the MPI implementation to fill in fields in this object.

To cope with this problem, there are two predefined constants, `MPI_STATUS_IGNORE` and `MPI_STATUSES_IGNORE`, which when passed to a receive, `probe`, wait, or test function, inform the implementation that the status fields are not to be filled in. Note that `MPI_STATUS_IGNORE` is not a special type of `MPI_STATUS` object; rather, it is a special

value for the argument. In C one would expect it to be NULL, not the address of a special MPI\_STATUS.

MPI\_STATUS\_IGNORE, and the array version MPI\_STATUSES\_IGNORE, can be used everywhere a status argument is passed to a receive, wait, or test function. MPI\_STATUS\_IGNORE cannot be used when status is an IN argument. Note that in Fortran MPI\_STATUS\_IGNORE and MPI\_STATUSES\_IGNORE are objects like MPI\_BOTTOM (not usable for initialization or assignment). See Section 2.5.4.

In general, this optimization can apply to all functions for which status or an array of statuses is an OUT argument. Note that this converts status into an INOUT argument. The functions that can be passed MPI\_STATUS\_IGNORE are all the various forms of MPI\_RECV, MPI\_PROBE, MPI\_TEST, and MPI\_WAIT, as well as MPI\_REQUEST\_GET\_STATUS. When an array is passed, as in the MPI\_{TEST|WAIT}{ALL|SOME} functions, a separate constant, MPI\_STATUSES\_IGNORE, is passed for the array argument. It is possible for an MPI function to return MPI\_ERR\_IN\_STATUS even when MPI\_STATUS\_IGNORE or MPI\_STATUSES\_IGNORE has been passed to that function.

MPI\_STATUS\_IGNORE and MPI\_STATUSES\_IGNORE are not required to have the same values in C and Fortran.

It is not allowed to have some of the statuses in an array of statuses for MPI\_{TEST|WAIT}{ALL|SOME} functions set to MPI\_STATUS\_IGNORE; one either specifies ignoring *all* of the statuses in such a call with MPI\_STATUSES\_IGNORE, or *none* of them by passing normal statuses in all positions in the array of statuses.

[In the deprecated C++ bindings, there] There are no C++ bindings for MPI\_STATUS\_IGNORE or MPI\_STATUSES\_IGNORE. To allow an OUT or INOUT MPI::Status argument to be ignored, all MPI C++ bindings that have OUT or INOUT MPI::Status parameters are overloaded with a second version that omits the OUT or INOUT MPI::Status parameter.

**Example 3.1** The [deprecated C++ bindings for MPI\_PROBE are:

```
void MPI::Comm::Probe(int source, int tag, MPI::Status& status) const
void MPI::Comm::Probe(int source, int tag) const
```

## 3.3 Data Type Matching and Data Conversion

### 3.3.1 Type Matching Rules

One can think of message transfer as consisting of the following three phases.

1. Data is pulled out of the send buffer and a message is assembled.
2. A message is transferred from sender to receiver.
3. Data is pulled from the incoming message and disassembled into the receive buffer.

Type matching has to be observed at each of these three phases: The type of each variable in the sender buffer has to match the type specified for that entry by the send operation; the type specified by the send operation has to match the type specified by the receive operation; and the type of each variable in the receive buffer has to match the type specified for that entry by the receive operation. A program that fails to observe these three rules is erroneous.

To define type matching more precisely, we need to deal with two issues: matching of types of the host language with types specified in communication operations; and matching of types at sender and receiver.

The types of a send and receive match (phase two) if both operations use identical names. That is, `MPI_INTEGER` matches `MPI_INTEGER`, `MPI_REAL` matches `MPI_REAL`, and so on. There is one exception to this rule, discussed in Section 4.2, the type `MPI_PACKED` can match any other type.

The type of a variable in a host program matches the type specified in the communication operation if the datatype name used by that operation corresponds to the basic type of the host program variable. For example, an entry with type name `MPI_INTEGER` matches a Fortran variable of type `INTEGER`. A table giving this correspondence for Fortran and C appears in Section 3.2.2. There are two exceptions to this last rule: an entry with type name `MPI_BYTE` or `MPI_PACKED` can be used to match any byte of storage (on a byte-addressable machine), irrespective of the datatype of the variable that contains this byte. The type `MPI_PACKED` is used to send data that has been explicitly packed, or receive data that will be explicitly unpacked, see Section 4.2. The type `MPI_BYTE` allows one to transfer the binary value of a byte in memory unchanged.

To summarize, the type matching rules fall into the three categories below.

- Communication of typed values (e.g., with datatype different from `MPI_BYTE`), where the datatypes of the corresponding entries in the sender program, in the send call, in the receive call and in the receiver program must all match.
- Communication of untyped values (e.g., of datatype `MPI_BYTE`), where both sender and receiver use the datatype `MPI_BYTE`. In this case, there are no requirements on the types of the corresponding entries in the sender and the receiver programs, nor is it required that they be the same.
- Communication involving packed data, where `MPI_PACKED` is used.

The following examples illustrate the first two cases.

**Example 3.2** Sender and receiver specify matching types.

```
CALL MPI_COMM_RANK(comm, rank, ierr)
IF (rank.EQ.0) THEN
  CALL MPI_SEND(a(1), 10, MPI_REAL, 1, tag, comm, ierr)
ELSE IF (rank.EQ.1) THEN
  CALL MPI_RECV(b(1), 15, MPI_REAL, 0, tag, comm, status, ierr)
END IF
```

This code is correct if both `a` and `b` are real arrays of size  $\geq 10$ . (In Fortran, it might be correct to use this code even if `a` or `b` have size  $< 10$ : e.g., when `a(1)` can be equivalenced to an array with ten reals.)

**Example 3.3** Sender and receiver do not specify matching types.

```

CALL MPI_COMM_RANK(comm, rank, ierr)
IF (rank.EQ.0) THEN
    CALL MPI_SEND(a(1), 10, MPI_REAL, 1, tag, comm, ierr)
ELSE IF (rank.EQ.1) THEN
    CALL MPI_RECV(b(1), 40, MPI_BYTE, 0, tag, comm, status, ierr)
END IF

```

This code is erroneous, since sender and receiver do not provide matching datatype arguments.

**Example 3.4** Sender and receiver specify communication of untyped values.

```

CALL MPI_COMM_RANK(comm, rank, ierr)
IF (rank.EQ.0) THEN
    CALL MPI_SEND(a(1), 40, MPI_BYTE, 1, tag, comm, ierr)
ELSE IF (rank.EQ.1) THEN
    CALL MPI_RECV(b(1), 60, MPI_BYTE, 0, tag, comm, status, ierr)
END IF

```

This code is correct, irrespective of the type and size of **a** and **b** (unless this results in an out of bound memory access).

*Advice to users.* If a buffer of type **MPI\_BYTE** is passed as an argument to **MPI\_SEND**, then MPI will send the data stored at contiguous locations, starting from the address indicated by the **buf** argument. This may have unexpected results when the data layout is not as a casual user would expect it to be. For example, some Fortran compilers implement variables of type **CHARACTER** as a structure that contains the character length and a pointer to the actual string. In such an environment, sending and receiving a Fortran **CHARACTER** variable using the **MPI\_BYTE** type will not have the anticipated result of transferring the character string. For this reason, the user is advised to use typed communications whenever possible. (*End of advice to users.*)

#### Type **MPI\_CHARACTER**

The type **MPI\_CHARACTER** matches one character of a Fortran variable of type **CHARACTER**, rather than the entire character string stored in the variable. Fortran variables of type **CHARACTER** or substrings are transferred as if they were arrays of characters. This is illustrated in the example below.

#### Example 3.5

Transfer of Fortran **CHARACTER**s.

```

CHARACTER*10 a
CHARACTER*10 b

CALL MPI_COMM_RANK(comm, rank, ierr)
IF (rank.EQ.0) THEN
    CALL MPI_SEND(a, 5, MPI_CHARACTER, 1, tag, comm, ierr)
ELSE IF (rank.EQ.1) THEN
    CALL MPI_RECV(b(6:10), 5, MPI_CHARACTER, 0, tag, comm, status, ierr)
END IF

```



The last five characters of string `b` at process 1 are replaced by the first five characters of string `a` at process 0.

*Rationale.* The alternative choice would be for `MPI_CHARACTER` to match a character of arbitrary length. This runs into problems.

A Fortran character variable is a constant length string, with no special termination symbol. There is no fixed convention on how to represent characters, and how to store their length. Some compilers pass a character argument to a routine as a pair of arguments, one holding the address of the string and the other holding the length of string. Consider the case of an MPI communication call that is passed a communication buffer with type defined by a derived datatype (Section 4.1). If this communicator buffer contains variables of type `CHARACTER` then the information on their length will not be passed to the MPI routine.

This problem forces us to provide explicit information on character length with the MPI call. One could add a length parameter to the type `MPI_CHARACTER`, but this does not add much convenience and the same functionality can be achieved by defining a suitable derived datatype. (*End of rationale.*)

*Advice to implementors.* Some compilers pass Fortran `CHARACTER` arguments as a structure with a length and a pointer to the actual string. In such an environment, the MPI call needs to dereference the pointer in order to reach the string. (*End of advice to implementors.*)

### 3.3.2 Data Conversion

One of the goals of MPI is to support parallel computations across heterogeneous environments. Communication in a heterogeneous environment may require data conversions. We use the following terminology.

**type conversion** changes the datatype of a value, e.g., by rounding a `REAL` to an `INTEGER`.

**representation conversion** changes the binary representation of a value, e.g., from Hex floating point to IEEE floating point.

The type matching rules imply that MPI communication never entails type conversion. On the other hand, MPI requires that a representation conversion be performed when a typed value is transferred across environments that use different representations for the datatype of this value. MPI does not specify rules for representation conversion. Such conversion is expected to preserve integer, logical or character values, and to convert a floating point value to the nearest value that can be represented on the target system.

Overflow and underflow exceptions may occur during floating point conversions. Conversion of integers or characters may also lead to exceptions when a value that can be represented in one system cannot be represented in the other system. An exception occurring during representation conversion results in a failure of the communication. An error occurs either in the send operation, or the receive operation, or both.

If a value sent in a message is untyped (i.e., of type `MPI_BYTE`), then the binary representation of the byte stored at the receiver is identical to the binary representation of the byte loaded at the sender. This holds true, whether sender and receiver run in the same or in distinct environments. No representation conversion is required. (Note that



representation conversion may occur when values of type `MPI_CHARACTER` or `MPI_CHAR` are transferred, for example, from an EBCDIC encoding to an ASCII encoding.)

No conversion need occur when an MPI program executes in a homogeneous system, where all processes run in the same environment.

Consider the three examples, 3.2–3.4. The first program is correct, assuming that `a` and `b` are `REAL` arrays of size  $\geq 10$ . If the sender and receiver execute in different environments, then the ten real values that are fetched from the send buffer will be converted to the representation for reals on the receiver site before they are stored in the receive buffer. While the number of real elements fetched from the send buffer equal the number of real elements stored in the receive buffer, the number of bytes stored need not equal the number of bytes loaded. For example, the sender may use a four byte representation and the receiver an eight byte representation for reals.

The second program is erroneous, and its behavior is undefined.

The third program is correct. The exact same sequence of forty bytes that were loaded from the send buffer will be stored in the receive buffer, even if sender and receiver run in a different environment. The message sent has exactly the same length (in bytes) and the same binary representation as the message received. If `a` and `b` are of different types, or if they are of the same type but different data representations are used, then the bits stored in the receive buffer may encode values that are different from the values they encoded in the send buffer.

Data representation conversion also applies to the envelope of a message: source, destination and tag are all integers that may need to be converted.

*Advice to implementors.* The current definition does not require messages to carry data type information. Both sender and receiver provide complete data type information. In a heterogeneous environment, one can either use a machine independent encoding such as XDR, or have the receiver convert from the sender representation to its own, or even have the sender do the conversion.

Additional type information might be added to messages in order to allow the system to detect mismatches between datatype at sender and receiver. This might be particularly useful in a slower but safer debug mode. (*End of advice to implementors.*)

MPI requires support for inter-language communication, i.e., if messages are sent by a C or C++ process and received by a Fortran process, or vice-versa. The behavior is defined in Section 16.3 on page 692.

## 3.4 Communication Modes

The send call described in Section 3.2.1 is **blocking**: it does not return until the message data and envelope have been safely stored away so that the sender is free to modify the send buffer. The message might be copied directly into the matching receive buffer, or it might be copied into a temporary system buffer.

Message buffering decouples the send and receive operations. A blocking send can complete as soon as the message was buffered, even if no matching receive has been executed by the receiver. On the other hand, message buffering can be expensive, as it entails additional memory-to-memory copying, and it requires the allocation of memory for buffering. MPI offers the choice of several communication modes that allow one to control the choice of the communication protocol.

The send call described in Section 3.2.1 uses the **standard** communication mode. In this mode, it is up to MPI to decide whether outgoing messages will be buffered. MPI may buffer outgoing messages. In such a case, the send call may complete before a matching receive is invoked. On the other hand, buffer space may be unavailable, or MPI may choose not to buffer outgoing messages, for performance reasons. In this case, the send call will not complete until a matching receive has been posted, and the data has been moved to the receiver.

Thus, a send in standard mode can be started whether or not a matching receive has been posted. It may complete before a matching receive is posted. The standard mode send is **non-local**: successful completion of the send operation may depend on the occurrence of a matching receive.

*Rationale.* The reluctance of MPI to mandate whether standard sends are buffering or not stems from the desire to achieve portable programs. Since any system will run out of buffer resources as message sizes are increased, and some implementations may want to provide little buffering, MPI takes the position that correct (and therefore, portable) programs do not rely on system buffering in standard mode. Buffering may improve the performance of a correct program, but it doesn't affect the result of the program. If the user wishes to guarantee a certain amount of buffering, the user-provided buffer system of Section 3.6 should be used, along with the buffered-mode send. (*End of rationale.*)

There are three additional communication modes.

A **buffered** mode send operation can be started whether or not a matching receive has been posted. It may complete before a matching receive is posted. However, unlike the standard send, this operation is **local**, and its completion does not depend on the occurrence of a matching receive. Thus, if a send is executed and no matching receive is posted, then MPI must buffer the outgoing message, so as to allow the send call to complete. An error will occur if there is insufficient buffer space. The amount of available buffer space is controlled by the user — see Section 3.6. Buffer allocation by the user may be required for the buffered mode to be effective.

A send that uses the **synchronous** mode can be started whether or not a matching receive was posted. However, the send will complete successfully only if a matching receive is posted, and the receive operation has started to receive the message sent by the synchronous send. Thus, the completion of a synchronous send not only indicates that the send buffer can be reused, but it also indicates that the receiver has reached a certain point in its execution, namely that it has started executing the matching receive. If both sends and receives are blocking operations then the use of the synchronous mode provides synchronous communication semantics: a communication does not complete at either end before both processes rendezvous at the communication. A send executed in this mode is **non-local**.

A send that uses the **ready** communication mode may be started *only* if the matching receive is already posted. Otherwise, the operation is erroneous and its outcome is undefined. On some systems, this allows the removal of a hand-shake operation that is otherwise required and results in improved performance. The completion of the send operation does not depend on the status of a matching receive, and merely indicates that the send buffer can be reused. A send operation that uses the ready mode has the same semantics as a standard send operation, or a synchronous send operation; it is merely that the sender provides additional information to the system (namely that a matching receive is already

posted), that can save some overhead. In a correct program, therefore, a ready send could be replaced by a standard send with no effect on the behavior of the program other than performance.

Three additional send functions are provided for the three additional communication modes. The communication mode is indicated by a one letter prefix: B for buffered, S for synchronous, and R for ready.

**MPI\_BSEND** (buf, count, datatype, dest, tag, comm)

IN	buf	initial address of send buffer (choice)
IN	count	number of elements in send buffer (non-negative integer)
IN	datatype	datatype of each send buffer element (handle)
IN	dest	rank of destination (integer)
IN	tag	message tag (integer)
IN	comm	communicator (handle)

```
int MPI_Bsend(const void* buf, int count, MPI_Datatype datatype, int dest,
             int tag, MPI_Comm comm)
```

```
MPI_Bsend(buf, count, datatype, dest, tag, comm, ierror) BIND(C)
  TYPE(*), DIMENSION(..), INTENT(IN) :: buf
  INTEGER, INTENT(IN) :: count, dest, tag
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  TYPE(MPI_Comm), INTENT(IN) :: comm
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_BSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, IERROR)
  <type> BUF(*)
  INTEGER COUNT, DATATYPE, DEST, TAG, COMM, IERROR
```

```
{void MPI::Comm::Bsend(const void* buf, int count, const
  MPI::Datatype& datatype, int dest, int tag) const(binding
  deprecated, see Section 15.2) }
```

Send in buffered mode.

1 MPI\_SSEND (buf, count, datatype, dest, tag, comm)

2	IN	buf	initial address of send buffer (choice)
3			
4	IN	count	number of elements in send buffer (non-negative integer)
5			
6	IN	datatype	datatype of each send buffer element (handle)
7	IN	dest	rank of destination (integer)
8			
9	IN	tag	message tag (integer)
10	IN	comm	communicator (handle)
11			

ticket140. 12 int MPI\_Ssend(const void\* buf, int count, MPI\_Datatype datatype, int dest,  
13 int tag, MPI\_Comm comm)

ticket-248T. 14 MPI\_Ssend(buf, count, datatype, dest, tag, comm, ierror) BIND(C)  
15 TYPE(\*), DIMENSION(..), INTENT(IN) :: buf  
16 INTEGER, INTENT(IN) :: count, dest, tag  
17 TYPE(MPI\_Datatype), INTENT(IN) :: datatype  
18 TYPE(MPI\_Comm), INTENT(IN) :: comm  
19 INTEGER, OPTIONAL, INTENT(OUT) :: ierror

21 MPI\_SSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, IERROR)  
22 <type> BUF(\*)  
23 INTEGER COUNT, DATATYPE, DEST, TAG, COMM, IERROR

25 {void MPI::Comm::Ssend(const void\* buf, int count, const  
26 MPI::Datatype& datatype, int dest, int tag) const(*binding*  
27 *deprecated, see Section 15.2*) }

28 Send in synchronous mode.  
29

31 MPI\_RSEND (buf, count, datatype, dest, tag, comm)

32	IN	buf	initial address of send buffer (choice)
33			
34	IN	count	number of elements in send buffer (non-negative integer)
35			
36	IN	datatype	datatype of each send buffer element (handle)
37	IN	dest	rank of destination (integer)
38			
39	IN	tag	message tag (integer)
40	IN	comm	communicator (handle)
41			

ticket140. 42 int MPI\_Rsend(const void\* buf, int count, MPI\_Datatype datatype, int dest,  
43 int tag, MPI\_Comm comm)

ticket-248T. 44 MPI\_Rsend(buf, count, datatype, dest, tag, comm, ierror) BIND(C)  
45 TYPE(\*), DIMENSION(..), INTENT(IN) :: buf  
46 INTEGER, INTENT(IN) :: count, dest, tag  
47 TYPE(MPI\_Datatype), INTENT(IN) :: datatype  
48

```

TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_RSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, IERROR)
  <type> BUF(*)
  INTEGER COUNT, DATATYPE, DEST, TAG, COMM, IERROR

{void MPI::Comm::Rsend(const void* buf, int count, const
    MPI::Datatype& datatype, int dest, int tag) const(binding
    deprecated, see Section 15.2) }
```

Send in ready mode.

There is only one receive operation, but it matches any of the send modes. The receive operation described in the last section is **blocking**: it returns only after the receive buffer contains the newly received message. A receive can complete before the matching send has completed (of course, it can complete only after the matching send has started).

In a multi-threaded implementation of MPI, the system may de-schedule a thread that is blocked on a send or receive operation, and schedule another thread for execution in the same address space. In such a case it is the user's responsibility not to modify a communication buffer until the communication completes. Otherwise, the outcome of the computation is undefined.

*Advice to implementors.* Since a synchronous send cannot complete before a matching receive is posted, one will not normally buffer messages sent by such an operation.

It is recommended to choose buffering over blocking the sender, whenever possible, for standard sends. The programmer can signal his or her preference for blocking the sender until a matching receive occurs by using the synchronous send mode.

A possible communication protocol for the various communication modes is outlined below.

**ready send:** The message is sent as soon as possible.

**synchronous send:** The sender sends a request-to-send message. The receiver stores this request. When a matching receive is posted, the receiver sends back a permission-to-send message, and the sender now sends the message.

**standard send:** First protocol may be used for short messages, and second protocol for long messages.

**buffered send:** The sender copies the message into a buffer and then sends it with a nonblocking send (using the same protocol as for standard send).

Additional control messages might be needed for flow control and error recovery. Of course, there are many other possible protocols.

Ready send can be implemented as a standard send. In this case there will be no performance advantage (or disadvantage) for the use of ready send.

A standard send can be implemented as a synchronous send. In such a case, no data buffering is needed. However, users may expect some buffering.

In a multi-threaded environment, the execution of a blocking communication should block only the executing thread, allowing the thread scheduler to de-schedule this thread and schedule another thread for execution. (*End of advice to implementors.*)

### 3.5 Semantics of Point-to-Point Communication

A valid MPI implementation guarantees certain general properties of point-to-point communication, which are described in this section.

**Order** Messages are *non-overtaking*: If a sender sends two messages in succession to the same destination, and both match the same receive, then this operation cannot receive the second message if the first one is still pending. If a receiver posts two receives in succession, and both match the same message, then the second receive operation cannot be satisfied by this message, if the first one is still pending. This requirement facilitates matching of sends to receives. It guarantees that message-passing code is deterministic, if processes are single-threaded and the wildcard `MPI_ANY_SOURCE` is not used in receives. (Some of the calls described later, such as `MPI_CANCEL` or `MPI_WAITANY`, are additional sources of nondeterminism.)

If a process has a single thread of execution, then any two communications executed by this process are ordered. On the other hand, if the process is multi-threaded, then the semantics of thread execution may not define a relative order between two send operations executed by two distinct threads. The operations are logically concurrent, even if one physically precedes the other. In such a case, the two messages sent can be received in any order. Similarly, if two receive operations that are logically concurrent receive two successively sent messages, then the two messages can match the two receives in either order.

**Example 3.6** An example of non-overtaking messages.

```
CALL MPI_COMM_RANK(comm, rank, ierr)
IF (rank.EQ.0) THEN
    CALL MPI_BSEND(buf1, count, MPI_REAL, 1, tag, comm, ierr)
    CALL MPI_BSEND(buf2, count, MPI_REAL, 1, tag, comm, ierr)
ELSE IF (rank.EQ.1) THEN
    CALL MPI_RECV(buf1, count, MPI_REAL, 0, MPI_ANY_TAG, comm, status, ierr)
    CALL MPI_RECV(buf2, count, MPI_REAL, 0, tag, comm, status, ierr)
END IF
```

The message sent by the first send must be received by the first receive, and the message sent by the second send must be received by the second receive.

**Progress** If a pair of matching send and receives have been initiated on two processes, then at least one of these two operations will complete, independently of other actions in the system: the send operation will complete, unless the receive is satisfied by another message, and completes; the receive operation will complete, unless the message sent is consumed by another matching receive that was posted at the same destination process.

**Example 3.7** An example of two, intertwined matching pairs.

```

CALL MPI_COMM_RANK(comm, rank, ierr)
IF (rank.EQ.0) THEN
    CALL MPI_BSEND(buf1, count, MPI_REAL, 1, tag1, comm, ierr)
    CALL MPI_SSEND(buf2, count, MPI_REAL, 1, tag2, comm, ierr)
ELSE IF (rank.EQ.1) THEN
    CALL MPI_RECV(buf1, count, MPI_REAL, 0, tag2, comm, status, ierr)
    CALL MPI_RECV(buf2, count, MPI_REAL, 0, tag1, comm, status, ierr)
END IF

```

Both processes invoke their first communication call. Since the first send of process zero uses the buffered mode, it must complete, irrespective of the state of process one. Since no matching receive is posted, the message will be copied into buffer space. (If insufficient buffer space is available, then the program will fail.) The second send is then invoked. At that point, a matching pair of send and receive operation is enabled, and both operations must complete. Process one next invokes its second receive call, which will be satisfied by the buffered message. Note that process one received the messages in the reverse order they were sent.

**Fairness** MPI makes no guarantee of *fairness* in the handling of communication. Suppose that a send is posted. Then it is possible that the destination process repeatedly posts a receive that matches this send, yet the message is never received, because it is each time overtaken by another message, sent from another source. Similarly, suppose that a receive was posted by a multi-threaded process. Then it is possible that messages that match this receive are repeatedly received, yet the receive is never satisfied, because it is overtaken by other receives posted at this node (by other executing threads). It is the programmer's responsibility to prevent starvation in such situations.

**Resource limitations** Any pending communication operation consumes system resources that are limited. Errors may occur when lack of resources prevent the execution of an MPI call. A quality implementation will use a (small) fixed amount of resources for each pending send in the ready or synchronous mode and for each pending receive. However, buffer space may be consumed to store messages sent in standard mode, and must be consumed to store messages sent in buffered mode, when no matching receive is available. The amount of space available for buffering will be much smaller than program data memory on many systems. Then, it will be easy to write programs that overrun available buffer space.

MPI allows the user to provide buffer memory for messages sent in the buffered mode. Furthermore, MPI specifies a detailed operational model for the use of this buffer. An MPI implementation is required to do no worse than implied by this model. This allows users to avoid buffer overflows when they use buffered sends. Buffer allocation and use is described in Section 3.6.

A buffered send operation that cannot complete because of a lack of buffer space is erroneous. When such a situation is detected, an error is signalled that may cause the program to terminate abnormally. On the other hand, a standard send operation that cannot complete because of lack of buffer space will merely block, waiting for buffer space to become available or for a matching receive to be posted. This behavior is preferable in many situations. Consider a situation where a producer repeatedly produces new values and sends them to a consumer. Assume that the producer produces new values faster than the consumer can consume them. If buffered sends are used, then a buffer overflow

will result. Additional synchronization has to be added to the program so as to prevent this from occurring. If standard sends are used, then the producer will be automatically throttled, as its send operations will block when buffer space is unavailable.

In some situations, a lack of buffer space leads to deadlock situations. This is illustrated by the examples below.

**Example 3.8** An exchange of messages.

```
CALL MPI_COMM_RANK(comm, rank, ierr)
IF (rank.EQ.0) THEN
    CALL MPI_SEND(sendbuf, count, MPI_REAL, 1, tag, comm, ierr)
    CALL MPI_RECV(recvbuf, count, MPI_REAL, 1, tag, comm, status, ierr)
ELSE IF (rank.EQ.1) THEN
    CALL MPI_RECV(recvbuf, count, MPI_REAL, 0, tag, comm, status, ierr)
    CALL MPI_SEND(sendbuf, count, MPI_REAL, 0, tag, comm, ierr)
END IF
```

This program will succeed even if no buffer space for data is available. The standard send operation can be replaced, in this example, with a synchronous send.

**Example 3.9** An errant attempt to exchange messages.

```
CALL MPI_COMM_RANK(comm, rank, ierr)
IF (rank.EQ.0) THEN
    CALL MPI_RECV(recvbuf, count, MPI_REAL, 1, tag, comm, status, ierr)
    CALL MPI_SEND(sendbuf, count, MPI_REAL, 1, tag, comm, ierr)
ELSE IF (rank.EQ.1) THEN
    CALL MPI_RECV(recvbuf, count, MPI_REAL, 0, tag, comm, status, ierr)
    CALL MPI_SEND(sendbuf, count, MPI_REAL, 0, tag, comm, ierr)
END IF
```

The receive operation of the first process must complete before its send, and can complete only if the matching send of the second processor is executed. The receive operation of the second process must complete before its send and can complete only if the matching send of the first process is executed. This program will always deadlock. The same holds for any other send mode.

**Example 3.10** An exchange that relies on buffering.

```
CALL MPI_COMM_RANK(comm, rank, ierr)
IF (rank.EQ.0) THEN
    CALL MPI_SEND(sendbuf, count, MPI_REAL, 1, tag, comm, ierr)
    CALL MPI_RECV(recvbuf, count, MPI_REAL, 1, tag, comm, status, ierr)
ELSE IF (rank.EQ.1) THEN
    CALL MPI_SEND(sendbuf, count, MPI_REAL, 0, tag, comm, ierr)
    CALL MPI_RECV(recvbuf, count, MPI_REAL, 0, tag, comm, status, ierr)
END IF
```

The message sent by each process has to be copied out before the send operation returns and the receive operation starts. For the program to complete, it is necessary that at least one of the two messages sent be buffered. Thus, this program can succeed only if the communication system can buffer at least count words of data.



*Advice to users.* When standard send operations are used, then a deadlock situation may occur where both processes are blocked because buffer space is not available. The same will certainly happen, if the synchronous mode is used. If the buffered mode is used, and not enough buffer space is available, then the program will not complete either. However, rather than a deadlock situation, we shall have a buffer overflow error.

A program is “safe” if no message buffering is required for the program to complete. One can replace all sends in such program with synchronous sends, and the program will still run correctly. This conservative programming style provides the best portability, since program completion does not depend on the amount of buffer space available or on the communication protocol used.

Many programmers prefer to have more leeway and opt to use the “unsafe” programming style shown in Example 3.10. In such cases, the use of standard sends is likely to provide the best compromise between performance and robustness: quality implementations will provide sufficient buffering so that “common practice” programs will not deadlock. The buffered send mode can be used for programs that require more buffering, or in situations where the programmer wants more control. This mode might also be used for debugging purposes, as buffer overflow conditions are easier to diagnose than deadlock conditions.

Nonblocking message-passing operations, as described in Section 3.7, can be used to avoid the need for buffering outgoing messages. This prevents deadlocks due to lack of buffer space, and improves performance, by allowing overlap of computation and communication, and avoiding the overheads of allocating buffers and copying messages into buffers. (*End of advice to users.*)

## 3.6 Buffer Allocation and Usage

A user may specify a buffer to be used for buffering messages sent in buffered mode. Buffering is done by the sender.

`MPI_BUFFER_ATTACH(buffer, size)`

IN	buffer	initial buffer address (choice)
IN	size	buffer size, in bytes (non-negative integer)

`int MPI_Buffer_attach(void* buffer, int size)`

```
MPI_Buffer_attach(buffer, size, ierror) BIND(C)
  TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buffer
  INTEGER, INTENT(IN) :: size
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

`MPI_BUFFER_ATTACH(BUFFER, SIZE, IERROR)`

```
<type> BUFFER(*)
INTEGER SIZE, IERROR
```

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```

1 {void MPI::Attach_buffer(void* buffer, int size) (binding deprecated, see
2   Section 15.2) }
```

Provides to MPI a buffer in the user's memory to be used for buffering outgoing messages. The buffer is used only by messages sent in buffered mode. Only one buffer can be attached to a process at a time. In C and C++, *buffer* is the starting address of a memory region. In Fortran, one can pass the first element of a memory region or a whole array, which must be 'simply contiguous' (for 'simply contiguous', see also Section 16.2.12 on page 671).

```

11 MPI_BUFFER_DETACH(buffer_addr, size)
```

12	OUT	buffer_addr	initial buffer address (choice)
13			
14	OUT	size	buffer size, in bytes (non-negative integer)

```

15
16 int MPI_Buffer_detach(void* buffer_addr, int* size)
```

```

17 MPI_Buffer_detach(buffer_addr, size, ierror) BIND(C)
18   USE, INTRINSIC :: ISO_C_BINDING, ONLY : C_PTR
19   TYPE(C_PTR), INTENT(OUT) :: buffer_addr
20   INTEGER, INTENT(OUT) :: size
21   INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```

22
23 MPI_BUFFER_DETACH(BUFFER_ADDR, SIZE, IERROR)
24   <type> BUFFER_ADDR(*)
25   INTEGER SIZE, IERROR
```

```

26
27 {int MPI::Detach_buffer(void*& buffer) (binding deprecated, see Section 15.2) }
```

Detach the buffer currently associated with MPI. The call returns the address and the size of the detached buffer. This operation will block until all messages currently in the buffer have been transmitted. Upon return of this function, the user may reuse or deallocate the space taken by the buffer.

**Example 3.11** Calls to attach and detach buffers.

```

33
34
35 #define BUFFSIZE 10000
36 int size;
37 char *buff;
38 MPI_Buffer_attach( malloc(BUFFSIZE), BUFFSIZE);
39 /* a buffer of 10000 bytes can now be used by MPI_Bsend */
40 MPI_Buffer_detach( &buff, &size);
41 /* Buffer size reduced to zero */
42 MPI_Buffer_attach( buff, size);
43 /* Buffer of 10000 bytes available again */
```

*Advice to users.* Even though the C functions `MPI_Buffer_attach` and `MPI_Buffer_detach` both have a first argument of type `void*`, these arguments are used differently: A pointer to the buffer is passed to `MPI_Buffer_attach`; the address of the pointer is passed to `MPI_Buffer_detach`, so that this call can return the pointer value.

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In Fortran with the `mpi` module or `mpif.h`, the type of the `buffer_addr` argument is wrongly defined and the argument is therefore unused. In Fortran with the `mpi_f08` module, the address of the buffer is returned as `TYPE(C_PTR)`, see also Example 8.1 on page 354 about the use of `C_PTR` pointers. (*End of advice to users.*)

*Rationale.* Both arguments are defined to be of type `void*` (rather than `void*` and `void**`, respectively), so as to avoid complex type casts. E.g., in the last example, `&buff`, which is of type `char**`, can be passed as argument to `MPI_Buffer_detach` without type casting. If the formal parameter had type `void**` then we would need a type cast before and after the call. (*End of rationale.*)

The statements made in this section describe the behavior of MPI for buffered-mode sends. When no buffer is currently associated, MPI behaves as if a zero-sized buffer is associated with the process.

MPI must provide as much buffering for outgoing messages *as if* outgoing message data were buffered by the sending process, in the specified buffer space, using a circular, contiguous-space allocation policy. We outline below a model implementation that defines this policy. MPI may provide more buffering, and may use a better buffer allocation algorithm than described below. On the other hand, MPI may signal an error whenever the simple buffering allocator described below would run out of space. In particular, if no buffer is explicitly associated with the process, then any buffered send may cause an error.

MPI does not provide mechanisms for querying or controlling buffering done by standard mode sends. It is expected that vendors will provide such information for their implementations.

*Rationale.* There is a wide spectrum of possible implementations of buffered communication: buffering can be done at sender, at receiver, or both; buffers can be dedicated to one sender-receiver pair, or be shared by all communications; buffering can be done in real or in virtual memory; it can use dedicated memory, or memory shared by other processes; buffer space may be allocated statically or be changed dynamically; etc. It does not seem feasible to provide a portable mechanism for querying or controlling buffering that would be compatible with all these choices, yet provide meaningful information. (*End of rationale.*)

### 3.6.1 Model Implementation of Buffered Mode

The model implementation uses the packing and unpacking functions described in Section 4.2 and the nonblocking communication functions described in Section 3.7.

We assume that a circular queue of pending message entries (PME) is maintained. Each entry contains a communication request handle that identifies a pending nonblocking send, a pointer to the next entry and the packed message data. The entries are stored in successive locations in the buffer. Free space is available between the queue tail and the queue head.

A buffered send call results in the execution of the following code.

- Traverse sequentially the PME queue from head towards the tail, deleting all entries for communications that have completed, up to the first entry with an uncompleted request; update queue head to point to that entry.

- Compute the number,  $n$ , of bytes needed to store an entry for the new message. An upper bound on  $n$  can be computed as follows: A call to the function `MPI_PACK_SIZE(count, datatype, comm, size)`, with the `count`, `datatype` and `comm` arguments used in the `MPI_BSEND` call, returns an upper bound on the amount of space needed to buffer the message data (see Section 4.2). The MPI constant `MPI_BSEND_OVERHEAD` provides an upper bound on the additional space consumed by the entry (e.g., for pointers or envelope information).
- Find the next contiguous empty space of  $n$  bytes in buffer (space following queue tail, or space at start of buffer if queue tail is too close to end of buffer). If space is not found then raise buffer overflow error.
- Append to end of PME queue in contiguous space the new entry that contains request handle, next pointer and packed message data; `MPI_PACK` is used to pack data.
- Post nonblocking send (standard mode) for packed data.
- Return

### 3.7 Nonblocking Communication

One can improve performance on many systems by overlapping communication and computation. This is especially true on systems where communication can be executed autonomously by an intelligent communication controller. Light-weight threads are one mechanism for achieving such overlap. An alternative mechanism that often leads to better performance is to use **nonblocking communication**. A nonblocking **send start** call initiates the send operation, but does not complete it. The send start call can return before the message was copied out of the send buffer. A separate **send complete** call is needed to complete the communication, i.e., to verify that the data has been copied out of the send buffer. With suitable hardware, the transfer of data out of the sender memory may proceed concurrently with computations done at the sender after the send was initiated and before it completed. Similarly, a nonblocking **receive start call** initiates the receive operation, but does not complete it. The call can return before a message is stored into the receive buffer. A separate **receive complete** call is needed to complete the receive operation and verify that the data has been received into the receive buffer. With suitable hardware, the transfer of data into the receiver memory may proceed concurrently with computations done after the receive was initiated and before it completed. The use of nonblocking receives may also avoid system buffering and memory-to-memory copying, as information is provided early on the location of the receive buffer.

Nonblocking send start calls can use the same four modes as blocking sends: **standard**, **buffered**, **synchronous** and **ready**. These carry the same meaning. Sends of all modes, **ready** excepted, can be started whether a matching receive has been posted or not; a nonblocking **ready** send can be started only if a matching receive is posted. In all cases, the send start call is local: it returns immediately, irrespective of the status of other processes. If the call causes some system resource to be exhausted, then it will fail and return an error code. Quality implementations of MPI should ensure that this happens only in “pathological” cases. That is, an MPI implementation should be able to support a large number of pending nonblocking operations.

The send-complete call returns when data has been copied out of the send buffer. It may carry additional meaning, depending on the send mode.

If the send mode is **synchronous**, then the send can complete only if a matching receive has started. That is, a receive has been posted, and has been matched with the send. In this case, the send-complete call is non-local. Note that a synchronous, nonblocking send may complete, if matched by a nonblocking receive, before the receive complete call occurs. (It can complete as soon as the sender “knows” the transfer will complete, but before the receiver “knows” the transfer will complete.)

If the send mode is **buffered** then the message must be buffered if there is no pending receive. In this case, the send-complete call is local, and must succeed irrespective of the status of a matching receive.

If the send mode is **standard** then the send-complete call may return before a matching receive is posted, if the message is buffered. On the other hand, the send-complete may not complete until a matching receive is posted, and the message was copied into the receive buffer.

Nonblocking sends can be matched with blocking receives, and vice-versa.

*Advice to users.* The completion of a send operation may be delayed, for standard mode, and must be delayed, for synchronous mode, until a matching receive is posted. The use of nonblocking sends in these two cases allows the sender to proceed ahead of the receiver, so that the computation is more tolerant of fluctuations in the speeds of the two processes.

Nonblocking sends in the buffered and ready modes have a more limited impact, e.g., the blocking version of buffered send is capable of completing regardless of when a matching receive call is made. However, separating the start from the completion of these sends still gives some opportunity for optimization within the MPI library. For example, starting a buffered send gives an implementation more flexibility in determining if and how the message is buffered. There are also advantages for both nonblocking buffered and ready modes when data copying can be done concurrently with computation.

The message-passing model implies that communication is initiated by the sender. The communication will generally have lower overhead if a receive is already posted when the sender initiates the communication (data can be moved directly to the receive buffer, and there is no need to queue a pending send request). However, a receive operation can complete only after the matching send has occurred. The use of nonblocking receives allows one to achieve lower communication overheads without blocking the receiver while it waits for the send. (*End of advice to users.*)

### 3.7.1 Communication Request Objects

Nonblocking communications use opaque **request** objects to identify communication operations and match the operation that initiates the communication with the operation that terminates it. These are system objects that are accessed via a handle. A request object identifies various properties of a communication operation, such as the send mode, the communication buffer that is associated with it, its context, the tag and destination arguments to be used for a send, or the tag and source arguments to be used for a receive. In addition, this object stores information about the status of the pending communication operation.

### 3.7.2 Communication Initiation

We use the same naming conventions as for blocking communication: a prefix of B, S, or R is used for buffered, synchronous or ready mode. In addition a prefix of I (for immediate) indicates that the call is nonblocking.

`MPI_ISEND(buf, count, datatype, dest, tag, comm, request)`

IN	buf	initial address of send buffer (choice)
IN	count	number of elements in send buffer (non-negative integer)
IN	datatype	datatype of each send buffer element (handle)
IN	dest	rank of destination (integer)
IN	tag	message tag (integer)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```
int MPI_Isend(const void* buf, int count, MPI_Datatype datatype, int dest,
              int tag, MPI_Comm comm, MPI_Request *request)
```

```
MPI_Isend(buf, count, datatype, dest, tag, comm, request, ierror) BIND(C)
  TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
  INTEGER, INTENT(IN) :: count, dest, tag
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  TYPE(MPI_Comm), INTENT(IN) :: comm
  TYPE(MPI_Request), INTENT(OUT) :: request
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_ISEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)
  <type> BUF(*)
  INTEGER COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR
```

```
{MPI::Request MPI::Comm::Isend(const void* buf, int count, const
                               MPI::Datatype& datatype, int dest, int tag) const(binding
                               deprecated, see Section 15.2) }
```

Start a standard mode, nonblocking send.

```

MPI_IBSEND(buf, count, datatype, dest, tag, comm, request)
    IN      buf          initial address of send buffer (choice)
    IN      count        number of elements in send buffer (non-negative integer)
    IN      datatype     datatype of each send buffer element (handle)
    IN      dest         rank of destination (integer)
    IN      tag          message tag (integer)
    IN      comm         communicator (handle)
    OUT     request      communication request (handle)

int MPI_Ibsend(const void* buf, int count, MPI_Datatype datatype, int dest,
              int tag, MPI_Comm comm, MPI_Request *request)

MPI_Ibsend(buf, count, datatype, dest, tag, comm, request, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count, dest, tag
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_IBSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)
    <type> BUF(*)
    INTEGER COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR

{MPI::Request MPI::Comm::Ibsend(const void* buf, int count, const
    MPI::Datatype& datatype, int dest, int tag) const(binding
    deprecated, see Section 15.2) }

    Start a buffered mode, nonblocking send.

MPI_ISSEND(buf, count, datatype, dest, tag, comm, request)
    IN      buf          initial address of send buffer (choice)
    IN      count        number of elements in send buffer (non-negative integer)
    IN      datatype     datatype of each send buffer element (handle)
    IN      dest         rank of destination (integer)
    IN      tag          message tag (integer)
    IN      comm         communicator (handle)
    OUT     request      communication request (handle)

int MPI_Issend(const void* buf, int count, MPI_Datatype datatype, int dest,
              int tag, MPI_Comm comm, MPI_Request *request)

```

```

1 MPI_Issend(buf, count, datatype, dest, tag, comm, request, ierror) BIND(C)
2   TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
3   INTEGER, INTENT(IN) :: count, dest, tag
4   TYPE(MPI_Datatype), INTENT(IN) :: datatype
5   TYPE(MPI_Comm), INTENT(IN) :: comm
6   TYPE(MPI_Request), INTENT(OUT) :: request
7   INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

8 MPI_ISSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)
9   <type> BUF(*)
10  INTEGER COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR

```

```

12 {MPI::Request MPI::Comm::Issend(const void* buf, int count, const
13     MPI::Datatype& datatype, int dest, int tag) const(binding
14     deprecated, see Section 15.2) }

```

Start a synchronous mode, nonblocking send.

```

18 MPI_IRSEND(buf, count, datatype, dest, tag, comm, request)

```

IN	buf	initial address of send buffer (choice)
IN	count	number of elements in send buffer (non-negative integer)
IN	datatype	datatype of each send buffer element (handle)
IN	dest	rank of destination (integer)
IN	tag	message tag (integer)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```

31 int MPI_Irsend(const void* buf, int count, MPI_Datatype datatype, int dest,
32     int tag, MPI_Comm comm, MPI_Request *request)

```

```

34 MPI_Irsend(buf, count, datatype, dest, tag, comm, request, ierror) BIND(C)
35   TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
36   INTEGER, INTENT(IN) :: count, dest, tag
37   TYPE(MPI_Datatype), INTENT(IN) :: datatype
38   TYPE(MPI_Comm), INTENT(IN) :: comm
39   TYPE(MPI_Request), INTENT(OUT) :: request
40   INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

41 MPI_IRSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)
42   <type> BUF(*)
43  INTEGER COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR

```

```

44 {MPI::Request MPI::Comm::Irsend(const void* buf, int count, const
45     MPI::Datatype& datatype, int dest, int tag) const(binding
46     deprecated, see Section 15.2) }

```

Start a ready mode nonblocking send.



MPI\_Irecv(buf, count, datatype, source, tag, comm, request)

OUT	buf	initial address of receive buffer (choice)
IN	count	number of elements in receive buffer (non-negative integer)
IN	datatype	datatype of each receive buffer element (handle)
IN	source	rank of source or MPI_ANY_SOURCE (integer)
IN	tag	message tag or MPI_ANY_TAG (integer)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```
int MPI_Irecv(void* buf, int count, MPI_Datatype datatype, int source,
              int tag, MPI_Comm comm, MPI_Request *request)
```

```
MPI_Irecv(buf, count, datatype, source, tag, comm, request, ierror) BIND(C)
  TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
  INTEGER, INTENT(IN) :: count, source, tag
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  TYPE(MPI_Comm), INTENT(IN) :: comm
  TYPE(MPI_Request), INTENT(OUT) :: request
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_RECV(buf, count, datatype, source, tag, comm, request, ierror)
  <type> BUF(*)
  INTEGER COUNT, DATATYPE, SOURCE, TAG, COMM, REQUEST, IERROR
```

```
{MPI::Request MPI::Comm::Irecv(void* buf, int count, const
                               MPI::Datatype& datatype, int source, int tag) const(binding
                               deprecated, see Section 15.2) }
```

Start a nonblocking receive.

These calls allocate a communication request object and associate it with the request handle (the argument `request`). The request can be used later to query the status of the communication or wait for its completion.

A nonblocking send call indicates that the system may start copying data out of the send buffer. The sender should not modify any part of the send buffer after a nonblocking send operation is called, until the send completes.

A nonblocking receive call indicates that the system may start writing data into the receive buffer. The receiver should not access any part of the receive buffer after a nonblocking receive operation is called, until the receive completes.

*Advice to users.* To prevent problems with the argument copying and register optimization done by Fortran compilers, please note the hints in [subsections “Problems Due to Data Copying and Sequence Association,” and “A Problem with Register Optimization” in Section 16.2.10 on pages 671 and 677. ]Sections 16.2.10-16.2.20, especially in Sections 16.2.12 and 16.2.13 on pages 671-674 about “Problems Due to Data Copying and Sequence Association with Subscript Triplets” and “Vector Subscripts”, and in Sections 16.2.16 to 16.2.19 on pages 677 to 686 about “Optimization Problems”,

“Code Movements and Register Optimization”, “Temporary Data Movements” and “Permanent Data Movements”. (*End of advice to users.*)

### 3.7.3 Communication Completion

The functions `MPI_WAIT` and `MPI_TEST` are used to complete a nonblocking communication. The completion of a send operation indicates that the sender is now free to update the locations in the send buffer (the send operation itself leaves the content of the send buffer unchanged). It does not indicate that the message has been received, rather, it may have been buffered by the communication subsystem. However, if a `synchronous` mode send was used, the completion of the send operation indicates that a matching receive was initiated, and that the message will eventually be received by this matching receive.

The completion of a receive operation indicates that the receive buffer contains the received message, the receiver is now free to access it, and that the status object is set. It does not indicate that the matching send operation has completed (but indicates, of course, that the send was initiated).

We shall use the following terminology: A **null** handle is a handle with value `MPI_REQUEST_NULL`. A persistent request and the handle to it are **inactive** if the request is not associated with any ongoing communication (see Section 3.9). A handle is **active** if it is neither null nor inactive. An **empty** status is a status which is set to return `tag = MPI_ANY_TAG`, `source = MPI_ANY_SOURCE`, `error = MPI_SUCCESS`, and is also internally configured so that calls to `MPI_GET_COUNT` [and], `MPI_GET_ELEMENTS`, and `MPI_GET_ELEMENTS_X` return `count = 0` and `MPI_TEST_CANCELLED` returns false. We set a status variable to empty when the value returned by it is not significant. Status is set in this way so as to prevent errors due to accesses of stale information.

The fields in a **status** object returned by a call to `MPI_WAIT`, `MPI_TEST`, or any of the other derived functions (`MPI_{TEST|WAIT}{ALL|SOME|ANY}`), where the **request** corresponds to a send call, are undefined, with two exceptions: The error status field will contain valid information if the wait or test call returned with `MPI_ERR_IN_STATUS`; and the returned status can be queried by the call `MPI_TEST_CANCELLED`.

Error codes belonging to the error class `MPI_ERR_IN_STATUS` should be returned only by the MPI completion functions that take arrays of `MPI_STATUS`. For the functions `MPI_TEST`, `MPI_TESTANY`, `MPI_WAIT`, and `MPI_WAITANY`, which return a single `MPI_STATUS` value, the normal MPI error return process should be used (not the `MPI_ERROR` field in the `MPI_STATUS` argument).

`MPI_WAIT(request, status)`

INOUT	request	request (handle)
OUT	status	status object (Status)

`int MPI_Wait(MPI_Request *request, MPI_Status *status)`

```

MPI_Wait(request, status, ierror) BIND(C)
    TYPE(MPI_Request), INTENT(INOUT) :: request
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_WAIT(REQUEST, STATUS, IERROR)
    INTEGER REQUEST, STATUS(MPI_STATUS_SIZE), IERROR

{void MPI::Request::Wait(MPI::Status& status) (binding deprecated, see
    Section 15.2) }

{void MPI::Request::Wait() (binding deprecated, see Section 15.2) }

```

A call to `MPI_WAIT` returns when the operation identified by `request` is complete. If the [communication object associated with this request was created by a nonblocking send or receive call, then the object]request is an active persistent request, it is marked inactive. Any other type of request is [ deallocated by the call to `MPI_WAIT` ] and the request handle is set to `MPI_REQUEST_NULL`. `MPI_WAIT` is a non-local operation.

The call returns, in `status`, information on the completed operation. The content of the status object for a receive operation can be accessed as described in Section 3.2.5. The status object for a send operation may be queried by a call to `MPI_TEST_CANCELLED` (see Section 3.8).

One is allowed to call `MPI_WAIT` with a null or inactive `request` argument. In this case the operation returns immediately with empty `status`.

*Advice to users.* Successful return of `MPI_WAIT` after a `MPI_IBSEND` implies that the user send buffer can be reused — i.e., data has been sent out or copied into a buffer attached with `MPI_BUFFER_ATTACH`. Note that, at this point, we can no longer cancel the send (see Section 3.8). If a matching receive is never posted, then the buffer cannot be freed. This runs somewhat counter to the stated goal of `MPI_CANCEL` (always being able to free program space that was committed to the communication subsystem). (*End of advice to users.*)

*Advice to implementors.* In a multi-threaded environment, a call to `MPI_WAIT` should block only the calling thread, allowing the thread scheduler to schedule another thread for execution. (*End of advice to implementors.*)

```

MPI_TEST(request, flag, status)

```

INOUT	request	communication request (handle)
OUT	flag	true if operation completed (logical)
OUT	status	status object (Status)

```

int MPI_Test(MPI_Request *request, int *flag, MPI_Status *status)

```

```

MPI_Test(request, flag, status, ierror) BIND(C)
    TYPE(MPI_Request), INTENT(INOUT) :: request
    LOGICAL, INTENT(OUT) :: flag
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_TEST(REQUEST, FLAG, STATUS, IERROR)
    LOGICAL FLAG
    INTEGER REQUEST, STATUS(MPI_STATUS_SIZE), IERROR

```

```

1  {bool MPI::Request::Test(MPI::Status& status) (binding deprecated, see
2      Section 15.2) }

```

```

3
4  {bool MPI::Request::Test() (binding deprecated, see Section 15.2) }

```

A call to MPI\_TEST returns flag = true if the operation identified by request is complete. In such a case, the status object is set to contain information on the completed operation[; if the communication object was created by a nonblocking send or receive, then it]. If the request is an active persistent request, it is marked as inactive. Any other type of request is deallocated and the request handle is set to MPI\_REQUEST\_NULL. The call returns flag = false[, otherwise.] if the operation identified by request is not complete. In this case, the value of the status object is undefined. MPI\_TEST is a local operation.

The return status object for a receive operation carries information that can be accessed as described in Section 3.2.5. The status object for a send operation carries information that can be accessed by a call to MPI\_TEST\_CANCELLED (see Section 3.8).

One is allowed to call MPI\_TEST with a null or inactive request argument. In such a case the operation returns with flag = true and empty status.

The functions MPI\_WAIT and MPI\_TEST can be used to complete both sends and receives.

*Advice to users.* The use of the nonblocking MPI\_TEST call allows the user to schedule alternative activities within a single thread of execution. An event-driven thread scheduler can be emulated with periodic calls to MPI\_TEST. (*End of advice to users.*)

**Example 3.12** Simple usage of nonblocking operations and MPI\_WAIT.

```

27 CALL MPI_COMM_RANK(comm, rank, ierr)
28 IF (rank.EQ.0) THEN
29     CALL MPI_ISEND(a(1), 10, MPI_REAL, 1, tag, comm, request, ierr)
30     **** do some computation to mask latency ****
31     CALL MPI_WAIT(request, status, ierr)
32 ELSE IF (rank.EQ.1) THEN
33     CALL MPI_IRECV(a(1), 15, MPI_REAL, 0, tag, comm, request, ierr)
34     **** do some computation to mask latency ****
35     CALL MPI_WAIT(request, status, ierr)
36 END IF

```

A request object can be deallocated without waiting for the associated communication to complete, by using the following operation.

```

41 MPI_REQUEST_FREE(request)

```

```

43     INOUT      request                communication request (handle)

```

```

45 int MPI_Request_free(MPI_Request *request)

```

```

47 MPI_Request_free(request, ierror) BIND(C)
48     TYPE(MPI_Request), INTENT(INOUT) :: request

```

```

    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_REQUEST_FREE(REQUEST, IERROR)
    INTEGER REQUEST, IERROR
{void MPI::Request::Free() (binding deprecated, see Section 15.2) }

```

Mark the request object for deallocation and set `request` to `MPI_REQUEST_NULL`. An ongoing communication that is associated with the request will be allowed to complete. The request will be deallocated only after its completion.

*Rationale.* The `MPI_REQUEST_FREE` mechanism is provided for reasons of performance and convenience on the sending side. (*End of rationale.*)

*Advice to users.* Once a request is freed by a call to `MPI_REQUEST_FREE`, it is not possible to check for the successful completion of the associated communication with calls to `MPI_WAIT` or `MPI_TEST`. Also, if an error occurs subsequently during the communication, an error code cannot be returned to the user — such an error must be treated as fatal. An active receive request should never be freed as the receiver will have no way to verify that the receive has completed and the receive buffer can be reused. (*End of advice to users.*)

**Example 3.13** An example using `MPI_REQUEST_FREE`.

```

CALL MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
IF (rank.EQ.0) THEN
    DO i=1, n
        CALL MPI_ISEND(outval, 1, MPI_REAL, 1, 0, MPI_COMM_WORLD, req, ierr)
        CALL MPI_REQUEST_FREE(req, ierr)
        CALL MPI_Irecv(inval, 1, MPI_REAL, 1, 0, MPI_COMM_WORLD, req, ierr)
        CALL MPI_WAIT(req, status, ierr)
    END DO
ELSE IF (rank.EQ.1) THEN
    CALL MPI_Irecv(inval, 1, MPI_REAL, 0, 0, MPI_COMM_WORLD, req, ierr)
    CALL MPI_WAIT(req, status, ierr)
    DO I=1, n-1
        CALL MPI_ISEND(outval, 1, MPI_REAL, 0, 0, MPI_COMM_WORLD, req, ierr)
        CALL MPI_REQUEST_FREE(req, ierr)
        CALL MPI_Irecv(inval, 1, MPI_REAL, 0, 0, MPI_COMM_WORLD, req, ierr)
        CALL MPI_WAIT(req, status, ierr)
    END DO
    CALL MPI_ISEND(outval, 1, MPI_REAL, 0, 0, MPI_COMM_WORLD, req, ierr)
    CALL MPI_WAIT(req, status, ierr)
END IF

```

### 3.7.4 Semantics of Nonblocking Communications

The semantics of nonblocking communication is defined by suitably extending the definitions in Section 3.5.

**Order** Nonblocking communication operations are ordered according to the execution order of the calls that initiate the communication. The non-overtaking requirement of Section 3.5 is extended to nonblocking communication, with this definition of order being used.

**Example 3.14** Message ordering for nonblocking operations.

```

CALL MPI_COMM_RANK(comm, rank, ierr)
IF (RANK.EQ.0) THEN
    CALL MPI_ISEND(a, 1, MPI_REAL, 1, 0, comm, r1, ierr)
    CALL MPI_ISEND(b, 1, MPI_REAL, 1, 0, comm, r2, ierr)
ELSE IF (rank.EQ.1) THEN
    CALL MPI_Irecv(a, 1, MPI_REAL, 0, MPI_ANY_TAG, comm, r1, ierr)
    CALL MPI_Irecv(b, 1, MPI_REAL, 0, 0, comm, r2, ierr)
END IF
CALL MPI_WAIT(r1, status, ierr)
CALL MPI_WAIT(r2, status, ierr)

```

The first send of process zero will match the first receive of process one, even if both messages are sent before process one executes either receive.

**Progress** A call to MPI\_WAIT that completes a receive will eventually terminate and return if a matching send has been started, unless the send is satisfied by another receive. In particular, if the matching send is nonblocking, then the receive should complete even if no call is executed by the sender to complete the send. Similarly, a call to MPI\_WAIT that completes a send will eventually return if a matching receive has been started, unless the receive is satisfied by another send, and even if no call is executed to complete the receive.

**Example 3.15** An illustration of progress semantics.

```

CALL MPI_COMM_RANK(comm, rank, ierr)
IF (RANK.EQ.0) THEN
    CALL MPI_SSEND(a, 1, MPI_REAL, 1, 0, comm, ierr)
    CALL MPI_SEND(b, 1, MPI_REAL, 1, 1, comm, ierr)
ELSE IF (rank.EQ.1) THEN
    CALL MPI_Irecv(a, 1, MPI_REAL, 0, 0, comm, r, ierr)
    CALL MPI_RECV(b, 1, MPI_REAL, 0, 1, comm, status, ierr)
    CALL MPI_WAIT(r, status, ierr)
END IF

```

This code should not deadlock in a correct MPI implementation. The first synchronous send of process zero must complete after process one posts the matching (nonblocking) receive even if process one has not yet reached the completing wait call. Thus, process zero will continue and execute the second send, allowing process one to complete execution.

If an MPI\_TEST that completes a receive is repeatedly called with the same arguments, and a matching send has been started, then the call will eventually return `flag = true`, unless the send is satisfied by another receive. If an MPI\_TEST that completes a send is repeatedly called with the same arguments, and a matching receive has been started, then the call will eventually return `flag = true`, unless the receive is satisfied by another send.

## 3.7.5 Multiple Completions

It is convenient to be able to wait for the completion of any, some, or all the operations in a list, rather than having to wait for a specific message. A call to `MPI_WAITANY` or `MPI_TESTANY` can be used to wait for the completion of one out of several operations. A call to `MPI_WAITALL` or `MPI_TESTALL` can be used to wait for all pending operations in a list. A call to `MPI_WAITSOME` or `MPI_TESTSOME` can be used to complete all enabled operations in a list.

`MPI_WAITANY` (count, array\_of\_requests, index, status)

IN	count	list length (non-negative integer)
INOUT	array_of_requests	array of requests (array of handles)
OUT	index	index of handle for operation that completed (integer)
OUT	status	status object (Status)

```
int MPI_Waitany(int count, MPI_Request [*]array_of_requests[], int *index,
               MPI_Status *status)
```

```
MPI_Waitany(count, array_of_requests, index, status, ierror) BIND(C)
  INTEGER, INTENT(IN) :: count
  TYPE(MPI_Request), INTENT(INOUT) :: array_of_requests(count)
  INTEGER, INTENT(OUT) :: index
  TYPE(MPI_Status) :: status
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_WAITANY(COUNT, ARRAY_OF_REQUESTS, INDEX, STATUS, IERROR)
  INTEGER COUNT, ARRAY_OF_REQUESTS(*), INDEX, STATUS(MPI_STATUS_SIZE),
  IERROR
```

```
{static int MPI::Request::Waitany(int count,
    MPI::Request array_of_requests[], MPI::Status& status) (binding
    deprecated, see Section 15.2) }
```

```
{static int MPI::Request::Waitany(int count,
    MPI::Request array_of_requests[]) (binding deprecated, see
    Section 15.2) }
```

Blocks until one of the operations associated with the active requests in the array has completed. If more than one operation is enabled and can terminate, one is arbitrarily chosen. Returns in `index` the index of that request in the array and returns in `status` the status of the completing [\[communication\]operation](#). (The array is indexed from zero in C, and from one in Fortran.) If the request [\[was allocated by a nonblocking communication operation, then it\]](#)is an active persistent request, it is marked inactive. Any other type of request is deallocated and the request handle is set to `MPI_REQUEST_NULL`.

The `array_of_requests` list may contain null or inactive handles. If the list contains no active handles (list has length zero or all entries are null or inactive), then the call returns immediately with `index = MPI_UNDEFINED`, and a empty status.



The execution of `MPI_WAITANY(count, array_of_requests, index, status)` has the same effect as the execution of `MPI_WAIT(&array_of_requests[i], status)`, where `i` is the value returned by `index` (unless the value of `index` is `MPI_UNDEFINED`). `MPI_WAITANY` with an array containing one active entry is equivalent to `MPI_WAIT`.

`MPI_TESTANY(count, array_of_requests, index, flag, status)`

IN	count	list length (non-negative integer)
INOUT	array_of_requests	array of requests (array of handles)
OUT	index	index of operation that completed, or <code>MPI_UNDEFINED</code> if none completed (integer)
OUT	flag	true if one of the operations is complete (logical)
OUT	status	status object (Status)

```
int MPI_Testany(int count, MPI_Request [*]array_of_requests[], int *index,
               int *flag, MPI_Status *status)
```

```
MPI_Testany(count, array_of_requests, index, flag, status, ierror) BIND(C)
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Request), INTENT(INOUT) :: array_of_requests(count)
    INTEGER, INTENT(OUT) :: index
    LOGICAL, INTENT(OUT) :: flag
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TESTANY(COUNT, ARRAY_OF_REQUESTS, INDEX, FLAG, STATUS, IERROR)
    LOGICAL FLAG
    INTEGER COUNT, ARRAY_OF_REQUESTS(*), INDEX, STATUS(MPI_STATUS_SIZE),
    IERROR
```

```
{static bool MPI::Request::Testany(int count,
    MPI::Request array_of_requests[], int& index,
    MPI::Status& status) (binding deprecated, see Section 15.2) }

{static bool MPI::Request::Testany(int count,
    MPI::Request array_of_requests[], int& index) (binding deprecated,
    see Section 15.2) }
```

Tests for completion of either one or none of the operations associated with active handles. In the former case, it returns `flag = true`, returns in `index` the index of this request in the array, and returns in `status` the status of that operation; if the request was allocated by a nonblocking communication call then the request. If the request is an active persistent request, it is marked as inactive. Any other type of request is deallocated and the handle is set to `MPI_REQUEST_NULL`. (The array is indexed from zero in C, and from one in Fortran.) In the latter case (no operation completed), it returns `flag = false`, returns a value of `MPI_UNDEFINED` in `index` and `status` is undefined.

The array may contain null or inactive handles. If the array contains no active handles then the call returns immediately with `flag = true`, `index = MPI_UNDEFINED`, and an empty `status`.



If the array of requests contains active handles then the execution of `MPI_TESTANY(count, array_of_requests, index, status)` has the same effect as the execution of `MPI_TEST(&array_of_requests[i], flag, status)`, for  $i=0, 1, \dots, \text{count}-1$ , in some arbitrary order, until one call returns `flag = true`, or all fail. In the former case, `index` is set to the last value of  $i$ , and in the latter case, it is set to `MPI_UNDEFINED`. `MPI_TESTANY` with an array containing one active entry is equivalent to `MPI_TEST`.

`MPI_WAITALL( count, array_of_requests, array_of_statuses)`

IN	count	lists length (non-negative integer)
INOUT	array_of_requests	array of requests (array of handles)
OUT	array_of_statuses	array of status objects (array of Status)

```
int MPI_Waitall(int count, MPI_Request [*]array_of_requests[],
                MPI_Status [*]array_of_statuses())
```

```
MPI_Waitall(count, array_of_requests, array_of_statuses, ierror) BIND(C)
  INTEGER, INTENT(IN) :: count
  TYPE(MPI_Request), INTENT(INOUT) :: array_of_requests(count)
  TYPE(MPI_Status) :: array_of_statuses(*)
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_WAITALL(COUNT, ARRAY_OF_REQUESTS, ARRAY_OF_STATUSES, IERROR)
  INTEGER COUNT, ARRAY_OF_REQUESTS(*)
  INTEGER ARRAY_OF_STATUSES(MPI_STATUS_SIZE,*), IERROR
```

```
{static void MPI::Request::Waitall(int count,
  MPI::Request array_of_requests[],
  MPI::Status array_of_statuses[]) (binding deprecated, see
  Section 15.2) }
```

```
{static void MPI::Request::Waitall(int count,
  MPI::Request array_of_requests[]) (binding deprecated, see
  Section 15.2) }
```

Blocks until all communication operations associated with active handles in the list complete, and return the status of all these operations (this includes the case where no handle in the list is active). Both arrays have the same number of valid entries. The  $i$ -th entry in `array_of_statuses` is set to the return status of the  $i$ -th operation. **Requests that were created by nonblocking communication operations** Active persistent requests are marked inactive. Requests of any other type are deallocated and the corresponding handles in the array are set to `MPI_REQUEST_NULL`. The list may contain null or inactive handles. The call sets to empty the status of each such entry.

The error-free execution of `MPI_WAITALL(count, array_of_requests, array_of_statuses)` has the same effect as the execution of `MPI_WAIT(&array_of_request[i], &array_of_statuses[i])`, for  $i=0, \dots, \text{count}-1$ , in some arbitrary order. `MPI_WAITALL` with an array of length one is equivalent to `MPI_WAIT`.

When one or more of the communications completed by a call to `MPI_WAITALL` fail, it is desirable to return specific information on each communication. The function

MPI\_WAITALL will return in such case the error code MPI\_ERR\_IN\_STATUS and will set the error field of each status to a specific error code. This code will be MPI\_SUCCESS, if the specific communication completed; it will be another specific error code, if it failed; or it can be MPI\_ERR\_PENDING if it has neither failed nor completed. The function MPI\_WAITALL will return MPI\_SUCCESS if no request had an error, or will return another error code if it failed for other reasons (such as invalid arguments). In such cases, it will not update the error fields of the statuses.

*Rationale.* This design streamlines error handling in the application. The application code need only test the (single) function result to determine if an error has occurred. It needs to check each individual status only when an error occurred. (*End of rationale.*)

MPI\_TESTALL(count, array\_of\_requests, flag, array\_of\_statuses)

IN	count	lists length (non-negative integer)
INOUT	array_of_requests	array of requests (array of handles)
OUT	flag	(logical)
OUT	array_of_statuses	array of status objects (array of Status)

```
int MPI_Testall(int count, MPI_Request [*]array_of_requests[], int *flag,
                MPI_Status [*]array_of_statuses[])
```

```
MPI_Testall(count, array_of_requests, flag, array_of_statuses, ierror)
    BIND(C)
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Request), INTENT(INOUT) :: array_of_requests(count)
    LOGICAL, INTENT(OUT) :: flag
    TYPE(MPI_Status) :: array_of_statuses(*)
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TESTALL(COUNT, ARRAY_OF_REQUESTS, FLAG, ARRAY_OF_STATUSES, IERROR)
    LOGICAL FLAG
    INTEGER COUNT, ARRAY_OF_REQUESTS(*),
    ARRAY_OF_STATUSES(MPI_STATUS_SIZE,*), IERROR
```

```
{static bool MPI::Request::Testall(int count,
    MPI::Request array_of_requests[],
    MPI::Status array_of_statuses[]) (binding deprecated, see
    Section 15.2) }
```

```
{static bool MPI::Request::Testall(int count,
    MPI::Request array_of_requests[]) (binding deprecated, see
    Section 15.2) }
```

Returns `flag = true` if all communications associated with active handles in the array have completed (this includes the case where no handle in the list is active). In this case, each status entry that corresponds to an active `[handle]request` is set to the status of the corresponding `[communication; if the request was allocated by a nonblocking communication`

call then it is deallocated, and the handle is]operation. Active persistent requests are marked inactive. Requests of any other type are deallocated and the corresponding handles in the array are set to MPI\_REQUEST\_NULL. Each status entry that corresponds to a null or inactive handle is set to empty.

Otherwise, flag = false is returned, no request is modified and the values of the status entries are undefined. This is a local operation.

Errors that occurred during the execution of MPI\_TESTALL are handled as errors in MPI\_WAITALL.

MPI\_WAITSSOME(incount, array\_of\_requests, outcount, array\_of\_indices, array\_of\_statuses)

IN	incount	length of array_of_requests (non-negative integer)
INOUT	array_of_requests	array of requests (array of handles)
OUT	outcount	number of completed requests (integer)
OUT	array_of_indices	array of indices of operations that completed (array of integers)
OUT	array_of_statuses	array of status objects for operations that completed (array of Status)

```
int MPI_Waitssome(int incount, MPI_Request [*]array_of_requests[],
                  int *outcount, int [*]array_of_indices[],
                  MPI_Status [*]array_of_statuses[])
```

```
MPI_Waitssome(incount, array_of_requests, outcount, array_of_indices,
               array_of_statuses, ierror) BIND(C)
```

```
INTEGER, INTENT(IN) :: incount
TYPE(MPI_Request), INTENT(INOUT) :: array_of_requests(incount)
INTEGER, INTENT(OUT) :: outcount, array_of_indices(*)
TYPE(MPI_Status) :: array_of_statuses(*)
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_WAITSSOME(INCOUNT, ARRAY_OF_REQUESTS, OUTCOUNT, ARRAY_OF_INDICES,
               ARRAY_OF_STATUSES, IERROR)
INTEGER INCOUNT, ARRAY_OF_REQUESTS(*), OUTCOUNT, ARRAY_OF_INDICES(*),
ARRAY_OF_STATUSES(MPI_STATUS_SIZE,*), IERROR
```

```
{static int MPI::Request::Waitssome(int incount,
MPI::Request array_of_requests[], int array_of_indices[],
MPI::Status array_of_statuses[]) (binding deprecated, see
Section 15.2) }
```

```
{static int MPI::Request::Waitssome(int incount,
MPI::Request array_of_requests[],
int array_of_indices[]) (binding deprecated, see Section 15.2) }
```

Waits until at least one of the operations associated with active handles in the list have completed. Returns in outcount the number of requests from the list array\_of\_requests that

have completed. Returns in the first `outcount` locations of the array `array_of_indices` the indices of these operations (index within the array `array_of_requests`; the array is indexed from zero in C and from one in Fortran). Returns in the first `outcount` locations of the array `array_of_status` the status for these completed operations. [If a request that completed was allocated by a nonblocking communication call, then it] Completed active persistent requests are marked as inactive. Any other type or request that completed is deallocated, and the associated handle is set to `MPI_REQUEST_NULL`.

If the list contains no active handles, then the call returns immediately with `outcount = MPI_UNDEFINED`.

When one or more of the communications completed by `MPI_WAITSSOME` fails, then it is desirable to return specific information on each communication. The arguments `outcount`, `array_of_indices` and `array_of_statuses` will be adjusted to indicate completion of all communications that have succeeded or failed. The call will return the error code `MPI_ERR_IN_STATUS` and the error field of each status returned will be set to indicate success or to indicate the specific error that occurred. The call will return `MPI_SUCCESS` if no request resulted in an error, and will return another error code if it failed for other reasons (such as invalid arguments). In such cases, it will not update the error fields of the statuses.

`MPI_TESTSSOME(incount, array_of_requests, outcount, array_of_indices, array_of_statuses)`

IN	<code>incount</code>	length of <code>array_of_requests</code> (non-negative integer)
INOUT	<code>array_of_requests</code>	array of requests (array of handles)
OUT	<code>outcount</code>	number of completed requests (integer)
OUT	<code>array_of_indices</code>	array of indices of operations that completed (array of integers)
OUT	<code>array_of_statuses</code>	array of status objects for operations that completed (array of Status)

```
int MPI_Testsome(int incount, MPI_Request [*]array_of_requests[],
                 int *outcount, int [*]array_of_indices[],
                 MPI_Status [*]array_of_statuses[])
```

```
MPI_Testsome(incount, array_of_requests, outcount, array_of_indices,
              array_of_statuses, ierror) BIND(C)
    INTEGER, INTENT(IN) :: incount
    TYPE(MPI_Request), INTENT(INOUT) :: array_of_requests(incount)
    INTEGER, INTENT(OUT) :: outcount, array_of_indices(*)
    TYPE(MPI_Status) :: array_of_statuses(*)
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TESTSSOME(INCOUNT, ARRAY_OF_REQUESTS, OUTCOUNT, ARRAY_OF_INDICES,
               ARRAY_OF_STATUSES, IERROR)
    INTEGER INCOUNT, ARRAY_OF_REQUESTS(*), OUTCOUNT, ARRAY_OF_INDICES(*),
    ARRAY_OF_STATUSES(MPI_STATUS_SIZE,*), IERROR
```

```

{static int MPI::Request::Testsome(int incout,
    MPI::Request array_of_requests[], int array_of_indices[],
    MPI::Status array_of_statuses[]) (binding deprecated, see
    Section 15.2) }
{static int MPI::Request::Testsome(int incout,
    MPI::Request array_of_requests[],
    int array_of_indices[]) (binding deprecated, see Section 15.2) }

```

Behaves like MPI\_WAITSSOME, except that it returns immediately. If no operation has completed it returns `outcount = 0`. If there is no active handle in the list it returns `outcount = MPI_UNDEFINED`.

MPI\_TESTSSOME is a local operation, which returns immediately, whereas MPI\_WAITSSOME will block until a communication completes, if it was passed a list that contains at least one active handle. Both calls fulfill a fairness requirement: If a request for a receive repeatedly appears in a list of requests passed to MPI\_WAITSSOME or MPI\_TESTSSOME, and a matching send has been posted, then the receive will eventually succeed, unless the send is satisfied by another receive; and similarly for send requests.

Errors that occur during the execution of MPI\_TESTSSOME are handled as for MPI\_WAITSSOME.

*Advice to users.* The use of MPI\_TESTSSOME is likely to be more efficient than the use of MPI\_TESTANY. The former returns information on all completed communications, with the latter, a new call is required for each communication that completes.

A server with multiple clients can use MPI\_WAITSSOME so as not to starve any client. Clients send messages to the server with service requests. The server calls MPI\_WAITSSOME with one receive request for each client, and then handles all receives that completed. If a call to MPI\_WAITANY is used instead, then one client could starve while requests from another client always sneak in first. (*End of advice to users.*)

*Advice to implementors.* MPI\_TESTSSOME should complete as many pending communications as possible. (*End of advice to implementors.*)

**Example 3.16** Client-server code (starvation can occur).

```

CALL MPI_COMM_SIZE(comm, size, ierr)
CALL MPI_COMM_RANK(comm, rank, ierr)
IF(rank .GT. 0) THEN      ! client code
  DO WHILE(.TRUE.)
    CALL MPI_ISEND(a, n, MPI_REAL, 0, tag, comm, request, ierr)
    CALL MPI_WAIT(request, status, ierr)
  END DO
ELSE      ! rank=0 -- server code
  DO i=1, size-1
    CALL MPI_Irecv(a(1,i), n, MPI_REAL, i, tag,
      comm, request_list(i), ierr)
  END DO
  DO WHILE(.TRUE.)

```

```

1      CALL MPI_WAITANY(size-1, request_list, index, status, ierr)
2      CALL DO_SERVICE(a(1,index)) ! handle one message
3      CALL MPI_Irecv(a(1, index), n, MPI_REAL, index, tag,
4                    comm, request_list(index), ierr)
5  END DO
6 END IF

```

**Example 3.17** Same code, using MPI\_WAITSSOME.

```

11 CALL MPI_COMM_SIZE(comm, size, ierr)
12 CALL MPI_COMM_RANK(comm, rank, ierr)
13 IF(rank .GT. 0) THEN ! client code
14   DO WHILE(.TRUE.)
15     CALL MPI_ISEND(a, n, MPI_REAL, 0, tag, comm, request, ierr)
16     CALL MPI_WAIT(request, status, ierr)
17   END DO
18 ELSE ! rank=0 -- server code
19   DO i=1, size-1
20     CALL MPI_Irecv(a(1,i), n, MPI_REAL, i, tag,
21                   comm, request_list(i), ierr)
22   END DO
23   DO WHILE(.TRUE.)
24     CALL MPI_WAITSSOME(size, request_list, numdone,
25                       indices, statuses, ierr)
26     DO i=1, numdone
27       CALL DO_SERVICE(a(1, indices(i)))
28       CALL MPI_Irecv(a(1, indices(i)), n, MPI_REAL, 0, tag,
29                     comm, request_list(indices(i)), ierr)
30     END DO
31   END DO
32 END IF

```

### 3.7.6 Non-destructive Test of status

This call is useful for accessing the information associated with a request, without freeing the request (in case the user is expected to access it later). It allows one to layer libraries more conveniently, since multiple layers of software may access the same completed request and extract from it the status information.

**MPI\_REQUEST\_GET\_STATUS( request, flag, status )**

IN	request	request (handle)
OUT	flag	boolean flag, same as from MPI_TEST (logical)
OUT	status	MPI_STATUS object if flag is true (Status)

```

int MPI_Request_get_status(MPI_Request request, int *flag,
                           MPI_Status *status)
MPI_Request_get_status(request, flag, status, ierror) BIND(C)
    TYPE(MPI_Request), INTENT(IN) :: request
    LOGICAL, INTENT(OUT) :: flag
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_REQUEST_GET_STATUS( REQUEST, FLAG, STATUS, IERROR)
    INTEGER REQUEST, STATUS(MPI_STATUS_SIZE), IERROR
    LOGICAL FLAG
{bool MPI::Request::Get_status(MPI::Status& status) const(binding deprecated,
    see Section 15.2) }
{bool MPI::Request::Get_status() const(binding deprecated, see Section 15.2) }

```

Sets `flag=true` if the operation is complete, and, if so, returns in `status` the request status. However, unlike `test` or `wait`, it does not deallocate or inactivate the request; a subsequent call to `test`, `wait` or `free` should be executed with that request. It sets `flag=false` if the operation is not complete.

One is allowed to call `MPI_REQUEST_GET_STATUS` with a null or inactive request argument. In such a case the operation returns with `flag=true` and empty `status`.

## 3.8 Probe and Cancel

The `MPI_PROBE` [ and ], `MPI_IPROBE`[], `MPI_MPROBE`, and `MPI_IMPROBE` operations allow incoming messages to be checked for, without actually receiving them. The user can then decide how to receive them, based on the information returned by the probe (basically, the information returned by `status`). In particular, the user may allocate memory for the receive buffer, according to the length of the probed message.

The `MPI_CANCEL` operation allows pending communications to be canceled. This is required for cleanup. Posting a send or a receive ties up user resources (send or receive buffers), and a cancel may be needed to free these resources gracefully.

### 3.8.1 Probe

```

MPI_IPROBE(source, tag, comm, flag, status)
    IN      source      rank of source or MPI_ANY_SOURCE (integer)
    IN      tag          message tag or MPI_ANY_TAG (integer)
    IN      comm         communicator (handle)
    OUT     flag         (logical)
    OUT     status       status object (Status)

int MPI_Iprobe(int source, int tag, MPI_Comm comm, int *flag,
               MPI_Status *status)

```

```

1
2 MPI_Iprobe(source, tag, comm, flag, status, ierror) BIND(C)
3     INTEGER, INTENT(IN) :: source, tag
4     TYPE(MPI_Comm), INTENT(IN) :: comm
5     LOGICAL, INTENT(OUT) :: flag
6     TYPE(MPI_Status) :: status
7     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
8
9 MPI_IPROBE(SOURCE, TAG, COMM, FLAG, STATUS, IERROR)
10     LOGICAL FLAG
11     INTEGER SOURCE, TAG, COMM, STATUS(MPI_STATUS_SIZE), IERROR
12
13 {bool MPI::Comm::Iprobe(int source, int tag, MPI::Status& status)
14     const(binding deprecated, see Section 15.2) }
15
16 {bool MPI::Comm::Iprobe(int source, int tag) const(binding deprecated, see
17     Section 15.2) }

```

MPI\_IPROBE(source, tag, comm, flag, status) returns flag = true if there is a message that can be received and that matches the pattern specified by the arguments source, tag, and comm. The call matches the same message that would have been received by a call to MPI\_RECV(..., source, tag, comm, status) executed at the same point in the program, and returns in status the same value that would have been returned by MPI\_RECV(). Otherwise, the call returns flag = false, and leaves status undefined.

If MPI\_IPROBE returns flag = true, then the content of the status object can be subsequently accessed as described in Section 3.2.5 to find the source, tag and length of the probed message.

A subsequent receive executed with the same communicator, and the source and tag returned in status by MPI\_IPROBE will receive the message that was matched by the probe, if no other intervening receive occurs after the probe, and the send is not successfully cancelled before the receive. If the receiving process is multi-threaded, it is the user's responsibility to ensure that the last condition holds.

The source argument of MPI\_PROBE can be MPI\_ANY\_SOURCE, and the tag argument can be MPI\_ANY\_TAG, so that one can probe for messages from an arbitrary source and/or with an arbitrary tag. However, a specific communication context must be provided with the comm argument.

It is not necessary to receive a message immediately after it has been probed for, and the same message may be probed for several times before it is received.

```

38
39 MPI_PROBE(source, tag, comm, status)
40     IN          source          rank of source or MPI_ANY_SOURCE (integer)
41     IN          tag             message tag or MPI_ANY_TAG (integer)
42     IN          comm            communicator (handle)
43     OUT         status          status object (Status)
44
45

```

```

46 int MPI_Probe(int source, int tag, MPI_Comm comm, MPI_Status *status)

```

```

47 MPI_Probe(source, tag, comm, status, ierror) BIND(C)
48

```



```

    INTEGER, INTENT(IN) :: source, tag
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_PROBE(SOURCE, TAG, COMM, STATUS, IERROR)
    INTEGER SOURCE, TAG, COMM, STATUS(MPI_STATUS_SIZE), IERROR

{void MPI::Comm::Probe(int source, int tag, MPI::Status& status)
    const(binding deprecated, see Section 15.2) }

{void MPI::Comm::Probe(int source, int tag) const(binding deprecated, see
    Section 15.2) }
```

MPI\_PROBE behaves like MPI\_IPROBE except that it is a blocking call that returns only after a matching message has been found.

The MPI implementation of MPI\_PROBE and MPI\_IPROBE needs to guarantee progress: if a call to MPI\_PROBE has been issued by a process, and a send that matches the probe has been initiated by some process, then the call to MPI\_PROBE will return, unless the message is received by another concurrent receive operation (that is executed by another thread at the probing process). Similarly, if a process busy waits with MPI\_IPROBE and a matching message has been issued, then the call to MPI\_IPROBE will eventually return flag = true unless the message is received by another concurrent receive operation or matched by a concurrent matched probe.

ticket38.

### Example 3.18

Use blocking probe to wait for an incoming message.

```

CALL MPI_COMM_RANK(comm, rank, ierr)
IF (rank.EQ.0) THEN
    CALL MPI_SEND(i, 1, MPI_INTEGER, 2, 0, comm, ierr)
ELSE IF (rank.EQ.1) THEN
    CALL MPI_SEND(x, 1, MPI_REAL, 2, 0, comm, ierr)
ELSE IF (rank.EQ.2) THEN
    DO i=1, 2
        CALL MPI_PROBE(MPI_ANY_SOURCE, 0,
                        comm, status, ierr)
        IF (status(MPI_SOURCE) .EQ. 0) THEN
100          CALL MPI_RECV(i, 1, MPI_INTEGER, 0, 0, comm, status, ierr)
        ELSE
200          CALL MPI_RECV(x, 1, MPI_REAL, 1, 0, comm, status, ierr)
        END IF
    END DO
END IF
```

Each message is received with the right type.

**Example 3.19** A similar program to the previous example, but now it has a problem.

```

1      CALL MPI_COMM_RANK(comm, rank, ierr)
2      IF (rank.EQ.0) THEN
3          CALL MPI_SEND(i, 1, MPI_INTEGER, 2, 0, comm, ierr)
4      ELSE IF (rank.EQ.1) THEN
5          CALL MPI_SEND(x, 1, MPI_REAL, 2, 0, comm, ierr)
6      ELSE IF (rank.EQ.2) THEN
7          DO i=1, 2
8              CALL MPI_PROBE(MPI_ANY_SOURCE, 0,
9                          comm, status, ierr)
10             IF (status(MPI_SOURCE) .EQ. 0) THEN
11 100          CALL MPI_RECV(i, 1, MPI_INTEGER, MPI_ANY_SOURCE,
12                      0, comm, status, ierr)
13             ELSE
14 200          CALL MPI_RECV(x, 1, MPI_REAL, MPI_ANY_SOURCE,
15                      0, comm, status, ierr)
16             END IF
17         END DO
18     END IF

```

[ We slightly modified Example 3.18, using `MPI_ANY_SOURCE` as the source argument in the two receive calls in statements labeled 100 and 200. The program is now incorrect: the receive operation may receive a message that is distinct from the message probed by the preceding call to `MPI_PROBE`.] In Example 3.19, the two receive calls in statements labeled 100 and 200 in Example 3.18 slightly modified, using `MPI_ANY_SOURCE` as the source argument. The program is now incorrect: the receive operation may receive a message that is distinct from the message probed by the preceding call to `MPI_PROBE`.

*Advice to users.* In a multithreaded MPI program, `MPI_PROBE` and `MPI_Iprobe` might need special care. If a thread probes for a message and then immediately posts a matching receive, the receive may match a message other than that found by the probe since another thread could concurrently receive that original message [29]. `MPI_Mprobe` and `MPI_improbe` solve this problem by matching the incoming message so that it may only be received with `MPI_Mrecv` or `MPI_imrecv` on the corresponding message handle. (*End of advice to users.*)

*Advice to implementors.* A call to `MPI_PROBE(source, tag, comm, status)` will match the message that would have been received by a call to `MPI_RECV(..., source, tag, comm, status)` executed at the same point. Suppose that this message has source `s`, tag `t` and communicator `c`. If the tag argument in the probe call has value `MPI_ANY_TAG` then the message probed will be the earliest pending message from source `s` with communicator `c` and any tag; in any case, the message probed will be the earliest pending message from source `s` with tag `t` and communicator `c` (this is the message that would have been received, so as to preserve message order). This message continues as the earliest pending message from source `s` with tag `t` and communicator `c`, until it is received. A receive operation subsequent to the probe that uses the same communicator as the probe and uses the tag and source values returned by the probe, must receive this message, unless it has already been received by another receive operation. (*End of advice to implementors.*)

### 3.8.2 Matching Probe

The function `MPI_PROBE` checks for incoming messages without receiving them. Since the list of incoming messages is global among the threads of each MPI process, it can be hard to use this functionality in threaded environments [29, 26].

Like `MPI_PROBE` and `MPI_IPROBE`, the `MPI_MPROBE` and `MPI_IMPROBE` operations allow incoming messages to be queried without actually receiving them, except that `MPI_MPROBE` and `MPI_IMPROBE` provide a mechanism to receive the specific message that was matched regardless of other intervening probe or receive operations. This gives the application an opportunity to decide how to receive the message, based on the information returned by the probe. In particular, the user may allocate memory for the receive buffer, according to the length of the probed message.

`MPI_IMPROBE(source, tag, comm, flag, message, status)`

IN	source	rank of source or <code>MPI_ANY_SOURCE</code> (integer)
IN	tag	message tag or <code>MPI_ANY_TAG</code> (integer)
IN	comm	communicator (handle)
OUT	flag	flag (logical)
OUT	message	returned message (handle)
OUT	status	status object (Status)

```
int MPI_Improbe(int source, int tag, MPI_Comm comm, int *flag,
                MPI_Message *message, MPI_Status *status)
```

```
MPI_Improbe(source, tag, comm, flag, message, status, ierror) BIND(C)
    INTEGER, INTENT(IN) :: source, tag
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, INTENT(OUT) :: flag
    TYPE(MPI_Message), INTENT(OUT) :: message
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_IMPROBE(SOURCE, TAG, COMM, FLAG, MESSAGE, STATUS, IERROR)
    INTEGER SOURCE, TAG, COMM, FLAG, MESSAGE, STATUS(MPI_STATUS_SIZE),
    IERROR
```

`MPI_IMPROBE(source, tag, comm, flag, message, status)` returns `flag = true` if there is a message that can be received and that matches the pattern specified by the arguments `source`, `tag`, and `comm`. The call matches the same message that would have been received by a call to `MPI_RECV(..., source, tag, comm, status)` executed at the same point in the program and returns in `status` the same value that would have been returned by `MPI_RECV`. In addition, it returns in `message` a handle to the matched message. Otherwise, the call returns `flag = false`, and leaves `status` and `message` undefined.

A matched receive (`MPI_MRECV` or `MPI_IMRECV`) executed with the message handle will receive the message that was matched by the probe. Unlike `MPI_IPROBE`, no other probe or receive operation may match the message returned by `MPI_IMPROBE`.

Each message returned by `MPI_IMPROBE` must be received with either `MPI_MRECV` or `MPI_IMRECV`.

The source argument of `MPI_IMPROBE` can be `MPI_ANY_SOURCE`, and the tag argument can be `MPI_ANY_TAG`, so that one can probe for messages from an arbitrary source and/or with an arbitrary tag. However, a specific communication context must be provided with the `comm` argument.

A synchronous send operation that is matched with `MPI_IMPROBE` or `MPI_MPROBE` will complete successfully only if both a matching receive is posted with `MPI_MRECV` or `MPI_IMRECV`, and the receive operation has started to receive the message sent by the synchronous send.

A matching probe with `MPI_PROC_NULL` as source returns `flag = true`, `message = MPI_MESSAGE_NULL`, and the status object returns `source = MPI_PROC_NULL`, `tag = MPI_ANY_TAG`, and `count = 0`; see Section 3.11.

`MPI_MPROBE(source, tag, comm, message, status)`

IN	source	rank of source or <code>MPI_ANY_SOURCE</code> (integer)
IN	tag	message tag or <code>MPI_ANY_TAG</code> (integer)
IN	comm	communicator (handle)
OUT	message	returned message (handle)
OUT	status	status object (Status)

`int MPI_Mprobe(int source, int tag, MPI_Comm comm, MPI_Message *message, MPI_Status *status)`

`MPI_Mprobe(source, tag, comm, message, status, ierror) BIND(C)`

```

INTEGER, INTENT(IN) :: source, tag
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Message), INTENT(OUT) :: message
TYPE(MPI_Status) :: status
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

`MPI_MPROBE(SOURCE, TAG, COMM, MESSAGE, STATUS, IERROR)`

`INTEGER SOURCE, TAG, COMM, MESSAGE, STATUS(MPI_STATUS_SIZE), IERROR`

`MPI_MPROBE` behaves like `MPI_IMPROBE` except that it is a blocking call that returns only after a matching message has been found.

The implementation of `MPI_MPROBE` and `MPI_IMPROBE` needs to guarantee progress in the same way as in the case of `MPI_PROBE` and `MPI_IPROBE`.

### 3.8.3 Matched Receives

The functions `MPI_MRECV` and `MPI_IMRECV` receive messages that have been previously matched by a matching probe (Section 3.8.2).

`MPI_MRECV(buf, count, datatype, message, status)`

OUT	buf	initial address of receive buffer (choice)
IN	count	number of elements in receive buffer (non-negative integer)
IN	datatype	datatype of each receive buffer element (handle)
INOUT	message	message (handle)
OUT	status	status object (Status)

```
int MPI_Mrecv(void* buf, int count, MPI_Datatype datatype,
              MPI_Message *message, MPI_Status *status)
```

```
MPI_Mrecv(buf, count, datatype, message, status, ierror) BIND(C)
```

```
TYPE(*), DIMENSION(..) :: buf
INTEGER, INTENT(IN) :: count
TYPE(MPI_Datatype), INTENT(IN) :: datatype
TYPE(MPI_Message), INTENT(INOUT) :: message
TYPE(MPI_Status) :: status
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_MRECV(BUF, COUNT, DATATYPE, MESSAGE, STATUS, IERROR)
```

```
<type> BUF(*)
INTEGER COUNT, DATATYPE, MESSAGE, STATUS(MPI_STATUS_SIZE), IERROR
```

This call receives a message matched by a matching probe operation (Section 3.8.2).

The receive buffer consists of the storage containing `count` consecutive elements of the type specified by `datatype`, starting at address `buf`. The length of the received message must be less than or equal to the length of the receive buffer. An overflow error occurs if all incoming data does not fit, without truncation, into the receive buffer.

If the message is shorter than the receive buffer, then only those locations corresponding to the (shorter) message are modified.

On return from this function, the message handle is set to `MPI_MESSAGE_NULL`. All errors that occur during the execution of this operation are handled according to the error handler set for the communicator used in the matching probe call that produced the message handle.

If `MPI_MRECV` is called with `MPI_MESSAGE_NULL` as the message argument, the call returns immediately with the status object set to `source = MPI_PROC_NULL`, `tag = MPI_ANY_TAG`, and `count = 0`, as if a receive from `MPI_PROC_NULL` was issued, see Section 3.11.

```
1 MPI_IMRECV(buf, count, datatype, message, request)
```

2	OUT	buf	initial address of receive buffer (choice)
3			
4	IN	count	number of elements in receive buffer (non-negative integer)
5			
6	IN	datatype	datatype of each receive buffer element (handle)
7	INOUT	message	message (handle)
8			
9	OUT	request	communication request (handle)

```
10
11 int MPI_Imrecv(void* buf, int count, MPI_Datatype datatype,
12               MPI_Message *message, MPI_Request *request)
```

```
13 MPI_Imrecv(buf, count, datatype, message, request, ierror) BIND(C)
14   TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
15   INTEGER, INTENT(IN) :: count
16   TYPE(MPI_Datatype), INTENT(IN) :: datatype
17   TYPE(MPI_Message), INTENT(INOUT) :: message
18   TYPE(MPI_Request), INTENT(OUT) :: request
19   INTEGER, OPTIONAL, INTENT(OUT) :: ierror
20
```

```
21 MPI_IMRECV(BUF, COUNT, DATATYPE, MESSAGE, REQUEST, IERROR)
22   <type> BUF(*)
23   INTEGER COUNT, DATATYPE, MESSAGE, REQUEST, IERROR
24
```

25 MPI\_IMRECV is the nonblocking variant of MPI\_MRECV and starts a nonblocking  
 26 receive of a matched message. Completion semantics are similar to MPI\_IRECV as described  
 27 in Section 3.7.2. On return from this function, the message handle is set to  
 28 MPI\_MESSAGE\_NULL.

29 *Advice to implementors.* If reception of a matched message is started with  
 30 MPI\_IMRECV, then it is possible to cancel the returned request with MPI\_CANCEL. If  
 31 MPI\_CANCEL succeeds, the matched message must be found by a subsequent message  
 32 probe (MPI\_PROBE, MPI\_IPROBE, MPI\_MPROBE, or MPI\_ImPROBE), received by  
 33 a subsequent receive operation or canceled by the sender. See Section 3.8.4 for details  
 34 about MPI\_CANCEL. The cancellation of operations initiated with MPI\_IMRECV may  
 35 fail. (*End of advice to implementors.*)

### 37 3.8.4 Cancel

```
40 MPI_CANCEL(request)
```

41	IN	request	communication request (handle)
----	----	---------	--------------------------------

```
43
44 int MPI_Cancel(MPI_Request *request)
```

```
45 MPI_Cancel(request, ierror) BIND(C)
46   TYPE(MPI_Request), INTENT(IN) :: request
47   INTEGER, OPTIONAL, INTENT(OUT) :: ierror
48
```

```
MPI_CANCEL(REQUEST, IERROR)
    INTEGER REQUEST, IERROR
```

```
{void MPI::Request::Cancel() const(binding deprecated, see Section 15.2) }
```

A call to `MPI_CANCEL` marks for cancellation a pending, nonblocking communication operation (send or receive). The cancel call is local. It returns immediately, possibly before the communication is actually canceled. It is still necessary to complete a communication that has been marked for cancellation, using a call to `MPI_REQUEST_FREE`, `MPI_WAIT` or `MPI_TEST` (or any of the derived operations).

If a communication is marked for cancellation, then a `MPI_WAIT` call for that communication is guaranteed to return, irrespective of the activities of other processes (i.e., `MPI_WAIT` behaves as a local function); similarly if `MPI_TEST` is repeatedly called in a busy wait loop for a canceled communication, then `MPI_TEST` will eventually be successful.

`MPI_CANCEL` can be used to cancel a communication that uses a persistent request (see Section 3.9), in the same way it is used for nonpersistent requests. A successful cancellation cancels the active communication, but not the request itself. After the call to `MPI_CANCEL` and the subsequent call to `MPI_WAIT` or `MPI_TEST`, the request becomes inactive and can be activated for a new communication.

The successful cancellation of a buffered send frees the buffer space occupied by the pending message.

Either the cancellation succeeds, or the communication succeeds, but not both. If a send is marked for cancellation, then it must be the case that either the send completes normally, in which case the message sent was received at the destination process, or that the send is successfully canceled, in which case no part of the message was received at the destination. Then, any matching receive has to be satisfied by another send. If a receive is marked for cancellation, then it must be the case that either the receive completes normally, or that the receive is successfully canceled, in which case no part of the receive buffer is altered. Then, any matching send has to be satisfied by another receive.

If the operation has been canceled, then information to that effect will be returned in the status argument of the operation that completes the communication.

*Rationale.* Although the IN request handle parameter should not need to be passed by reference, the C binding has listed the argument type as `MPI_Request*` since MPI-1.0. This function signature therefore cannot be changed without breaking existing MPI applications. (*End of rationale.*)

```
MPI_TEST_CANCELLED(status, flag)
```

IN	status	status object (Status)
OUT	flag	(logical)

```
int MPI_Test_cancelled(const MPI_Status *status, int *flag)
```

```
MPI_Test_cancelled(status, flag, ierror) BIND(C)
    TYPE(MPI_Status), INTENT(IN) :: status
    LOGICAL, INTENT(OUT) :: flag
```

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```

1      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
2
3      MPI_TEST_CANCELLED(STATUS, FLAG, IERROR)
4          LOGICAL FLAG
5          INTEGER STATUS(MPI_STATUS_SIZE), IERROR
6
7      {bool MPI::Status::Is_cancelled() const(binding deprecated, see Section 15.2) }

```

Returns `flag = true` if the communication associated with the status object was canceled successfully. In such a case, all other fields of `status` (such as `count` or `tag`) are undefined. Returns `flag = false`, otherwise. If a receive operation might be canceled then one should call `MPI_TEST_CANCELLED` first, to check whether the operation was canceled, before checking on the other fields of the return status.

*Advice to users.* Cancel can be an expensive operation that should be used only exceptionally. (*End of advice to users.*)

*Advice to implementors.* If a send operation uses an “eager” protocol (data is transferred to the receiver before a matching receive is posted), then the cancellation of this send may require communication with the intended receiver in order to free allocated buffers. On some systems this may require an interrupt to the intended receiver. Note that, while communication may be needed to implement `MPI_CANCEL`, this is still a local operation, since its completion does not depend on the code executed by other processes. If processing is required on another process, this should be transparent to the application (hence the need for an interrupt and an interrupt handler). (*End of advice to implementors.*)

### 3.9 Persistent Communication Requests

Often a communication with the same argument list is repeatedly executed within the inner loop of a parallel computation. In such a situation, it may be possible to optimize the communication by binding the list of communication arguments to a **persistent** communication request once and, then, repeatedly using the request to initiate and complete messages. The persistent request thus created can be thought of as a communication port or a “half-channel.” It does not provide the full functionality of a conventional channel, since there is no binding of the send port to the receive port. This construct allows reduction of the overhead for communication between the process and communication controller, but not of the overhead for communication between one communication controller and another. It is not necessary that messages sent with a persistent request be received by a receive operation using a persistent request, or vice versa.

A persistent communication request is created using one of the five following calls. These calls involve no communication.



MPI\_SEND\_INIT(buf, count, datatype, dest, tag, comm, request)

IN	buf	initial address of send buffer (choice)
IN	count	number of elements sent (non-negative integer)
IN	datatype	type of each element (handle)
IN	dest	rank of destination (integer)
IN	tag	message tag (integer)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```
int MPI_Send_init(const void* buf, int count, MPI_Datatype datatype,
                 int dest, int tag, MPI_Comm comm, MPI_Request *request)
```

```
MPI_Send_init(buf, count, datatype, dest, tag, comm, request, ierror)
```

```
    BIND(C)
```

```
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count, dest, tag
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_SEND_INIT(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)
```

```
    <type> BUF(*)
```

```
    INTEGER [REQUEST, ]COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR
```

```
{MPI::Prequest MPI::Comm::Send_init(const void* buf, int count, const
    MPI::Datatype& datatype, int dest, int tag) const(binding
    deprecated, see Section 15.2) }
```

Creates a persistent communication request for a standard mode send operation, and binds to it all the arguments of a send operation.

MPI\_BSEND\_INIT(buf, count, datatype, dest, tag, comm, request)

IN	buf	initial address of send buffer (choice)
IN	count	number of elements sent (non-negative integer)
IN	datatype	type of each element (handle)
IN	dest	rank of destination (integer)
IN	tag	message tag (integer)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```
int MPI_Bsend_init(const void* buf, int count, MPI_Datatype datatype,
                  int dest, int tag, MPI_Comm comm, MPI_Request *request)
```

```

1  MPI_Bsend_init(buf, count, datatype, dest, tag, comm, request, ierror)
2      BIND(C)
3      TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
4      INTEGER, INTENT(IN) :: count, dest, tag
5      TYPE(MPI_Datatype), INTENT(IN) :: datatype
6      TYPE(MPI_Comm), INTENT(IN) :: comm
7      TYPE(MPI_Request), INTENT(OUT) :: request
8      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
9
10 MPI_BSEND_INIT(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)
11 <type> BUF(*)
ticket250-V. 12  INTEGER [REQUEST, ]COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR
13 {MPI::Prequest MPI::Comm::Bsend_init(const void* buf, int count, const
14     MPI::Datatype& datatype, int dest, int tag) const(binding
15     deprecated, see Section 15.2) }
16
17     Creates a persistent communication request for a buffered mode send.
18
19
20 MPI_SSEND_INIT(buf, count, datatype, dest, tag, comm, request)
21     IN        buf                initial address of send buffer (choice)
22     IN        count              number of elements sent (non-negative integer)
23     IN        datatype           type of each element (handle)
24     IN        dest               rank of destination (integer)
25     IN        tag               message tag (integer)
26     IN        comm              communicator (handle)
27     IN        comm
28     OUT       request            communication request (handle)
29
30
ticket140. 31 int MPI_Ssend_init(const void* buf, int count, MPI_Datatype datatype,
32     int dest, int tag, MPI_Comm comm, MPI_Request *request)
ticket229.2. 33
ticket-248T. 34 MPI_Ssend_init(buf, count, datatype, dest, tag, comm, request, ierror)
35     BIND(C)
36     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
37     INTEGER, INTENT(IN) :: count, dest, tag
38     TYPE(MPI_Datatype), INTENT(IN) :: datatype
39     TYPE(MPI_Comm), INTENT(IN) :: comm
40     TYPE(MPI_Request), INTENT(OUT) :: request
41     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
42
43 MPI_SSEND_INIT(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)
44 <type> BUF(*)
45     INTEGER COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR
46
47 {MPI::Prequest MPI::Comm::Ssend_init(const void* buf, int count, const
48     MPI::Datatype& datatype, int dest, int tag) const(binding
    deprecated, see Section 15.2) }

```

Creates a persistent communication object for a synchronous mode send operation.

`MPI_RSEND_INIT(buf, count, datatype, dest, tag, comm, request)`

IN	buf	initial address of send buffer (choice)
IN	count	number of elements sent (non-negative integer)
IN	datatype	type of each element (handle)
IN	dest	rank of destination (integer)
IN	tag	message tag (integer)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```
int MPI_Rsend_init(const void* buf, int count, MPI_Datatype datatype,
                  int dest, int tag, MPI_Comm comm, MPI_Request *request)
```

```
MPI_Rsend_init(buf, count, datatype, dest, tag, comm, request, ierror)
```

```
    BIND(C)
```

```
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count, dest, tag
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_RSEND_INIT(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)
```

```
    <type> BUF(*)
```

```
    INTEGER COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR
```

```
{MPI::Prequest MPI::Comm::Rsend_init(const void* buf, int count, const
    MPI::Datatype& datatype, int dest, int tag) const(binding
    deprecated, see Section 15.2) }
```

Creates a persistent communication object for a ready mode send operation.

`MPI_RECV_INIT(buf, count, datatype, source, tag, comm, request)`

OUT	buf	initial address of receive buffer (choice)
IN	count	number of elements received (non-negative integer)
IN	datatype	type of each element (handle)
IN	source	rank of source or MPI_ANY_SOURCE (integer)
IN	tag	message tag or MPI_ANY_TAG (integer)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```

1  int MPI_Recv_init(void* buf, int count, MPI_Datatype datatype, int source,
2                      int tag, MPI_Comm comm, MPI_Request *request)

```

```

3  MPI_Recv_init(buf, count, datatype, source, tag, comm, request, ierror)
4  BIND(C)

```

```

5  TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
6  INTEGER, INTENT(IN) :: count, source, tag
7  TYPE(MPI_Datatype), INTENT(IN) :: datatype
8  TYPE(MPI_Comm), INTENT(IN) :: comm
9  TYPE(MPI_Request), INTENT(OUT) :: request
10  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

11
12 MPI_RECV_INIT(BUF, COUNT, DATATYPE, SOURCE, TAG, COMM, REQUEST, IERROR)
13 <type> BUF(*)
14 INTEGER COUNT, DATATYPE, SOURCE, TAG, COMM, REQUEST, IERROR

```

```

15 {MPI::Prequest MPI::Comm::Recv_init(void* buf, int count, const
16     MPI::Datatype& datatype, int source, int tag) const(binding
17     deprecated, see Section 15.2) }

```

Creates a persistent communication request for a receive operation. The argument `buf` is marked as OUT because the user gives permission to write on the receive buffer by passing the argument to `MPI_RECV_INIT`.

A persistent communication request is inactive after it was created — no active communication is attached to the request.

A communication (send or receive) that uses a persistent request is initiated by the function `MPI_START`.

```

27
28 MPI_START(request)

```

```

29     INOUT      request                      communication request (handle)

```

```

30
31 int MPI_Start(MPI_Request *request)

```

```

32
33 MPI_Start(request, ierror) BIND(C)
34     TYPE(MPI_Request), INTENT(INOUT) :: request
35     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

36
37 MPI_START(REQUEST, IERROR)
38     INTEGER REQUEST, IERROR

```

```

39 {void MPI::Prequest::Start()(binding deprecated, see Section 15.2) }

```

The argument, `request`, is a handle returned by one of the previous five calls. The associated request should be inactive. The request becomes active once the call is made.

If the request is for a send with ready mode, then a matching receive should be posted before the call is made. The communication buffer should not be modified after the call, and until the operation completes.

The call is local, with similar semantics to the nonblocking communication operations described in Section 3.7. That is, a call to `MPI_START` with a request created by `MPI_SEND_INIT` starts a communication in the same manner as a call to `MPI_ISEND`; a

call to `MPI_START` with a request created by `MPI_BSEND_INIT` starts a communication in the same manner as a call to `MPI_IBSEND`; and so on.

`MPI_STARTALL(count, array_of_requests)`

IN	count	list length (non-negative integer)
INOUT	array_of_requests	array of requests (array of handle)

```
int MPI_Startall(int count, MPI_Request [*]array_of_requests[])
```

```
MPI_Startall(count, array_of_requests, ierror) BIND(C)
  INTEGER, INTENT(IN) :: count
  TYPE(MPI_Request), INTENT(INOUT) :: array_of_requests(count)
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_STARTALL(COUNT, ARRAY_OF_REQUESTS, IERROR)
  INTEGER COUNT, ARRAY_OF_REQUESTS(*), IERROR
```

```
{static void MPI::Prequest::Startall(int count,
    MPI::Prequest array_of_requests[]) (binding deprecated, see
    Section 15.2) }
```

Start all communications associated with requests in `array_of_requests`. A call to `MPI_STARTALL(count, array_of_requests)` has the same effect as calls to `MPI_START (&array_of_requests[i])`, executed for `i=0, ..., count-1`, in some arbitrary order.

A communication started with a call to `MPI_START` or `MPI_STARTALL` is completed by a call to `MPI_WAIT`, `MPI_TEST`, or one of the derived functions described in Section 3.7.5. The request becomes inactive after successful completion of such call. The request is not deallocated and it can be activated anew by an `MPI_START` or `MPI_STARTALL` call.

A persistent request is deallocated by a call to `MPI_REQUEST_FREE` (Section 3.7.3).

The call to `MPI_REQUEST_FREE` can occur at any point in the program after the persistent request was created. However, the request will be deallocated only after it becomes inactive. Active receive requests should not be freed. Otherwise, it will not be possible to check that the receive has completed. It is preferable, in general, to free requests when they are inactive. If this rule is followed, then the functions described in this section will be invoked in a sequence of the form,

**Create (Start Complete)\* Free**

where `*` indicates zero or more repetitions. If the same communication object is used in several concurrent threads, it is the user's responsibility to coordinate calls so that the correct sequence is obeyed.

A send operation initiated with `MPI_START` can be matched with any receive operation and, likewise, a receive operation initiated with `MPI_START` can receive messages generated by any send operation.

*Advice to users.* To prevent problems with the argument copying and register optimization done by Fortran compilers, please note the hints in [subsections "Problems

Due to Data Copying and Sequence Association,” and “A Problem with Register Optimization and Temporary Memory Modifications” in Section 16.2.10 on pages 671 and 677. ]Sections 16.2.10-16.2.20, especially in Sections 16.2.12 and 16.2.13 on pages 671-674 about “Problems Due to Data Copying and Sequence Association with Subscript Triplets” and “Vector Subscripts”, and in Sections 16.2.16 to 16.2.19 on pages 677 to 686 about “Optimization Problems”, “Code Movements and Register Optimization”, “Temporary Data Movements” and “Permanent Data Movements”. (*End of advice to users.*)

### 3.10 Send-Receive

The **send-receive** operations combine in one call the sending of a message to one destination and the receiving of another message, from another process. The two (source and destination) are possibly the same. A send-receive operation is very useful for executing a shift operation across a chain of processes. If blocking sends and receives are used for such a shift, then one needs to order the sends and receives correctly (for example, even processes send, then receive, odd processes receive first, then send) so as to prevent cyclic dependencies that may lead to deadlock. When a send-receive operation is used, the communication subsystem takes care of these issues. The send-receive operation can be used in conjunction with the functions described in Chapter 7 in order to perform shifts on various logical topologies. Also, a send-receive operation is useful for implementing remote procedure calls.

A message sent by a send-receive operation can be received by a regular receive operation or probed by a probe operation; a send-receive operation can receive a message sent by a regular send operation.

**MPI\_SENDRECV**(sendbuf, sendcount, sendtype, dest, sendtag, recvbuf, recvcount, recvtype, source, recvtag, comm, status)

IN	sendbuf	initial address of send buffer (choice)
IN	sendcount	number of elements in send buffer (non-negative integer)
IN	sendtype	type of elements in send buffer (handle)
IN	dest	rank of destination (integer)
IN	sendtag	send tag (integer)
OUT	recvbuf	initial address of receive buffer (choice)
IN	recvcount	number of elements in receive buffer (non-negative integer)
IN	recvtype	type of elements in receive buffer (handle)
IN	source	rank of source or MPI_ANY_SOURCE (integer)
IN	recvtag	receive tag or MPI_ANY_TAG (integer)
IN	comm	communicator (handle)
OUT	status	status object (Status)

```

int MPI_Sendrecv(const void *sendbuf, int sendcount, MPI_Datatype sendtype,
                 int dest, int sendtag, void *recvbuf, int recvcount,
                 MPI_Datatype recvtype, int source, int recvtag, MPI_Comm comm,
                 MPI_Status *status)
MPI_Sendrecv(sendbuf, sendcount, sendtype, dest, sendtag, recvbuf,
              recvcount, recvtype, source, recvtag, comm, status, ierror)
              BIND(C)
              TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
              TYPE(*), DIMENSION(..) :: recvbuf
              INTEGER, INTENT(IN) :: sendcount, dest, sendtag, recvcount, source,
              recvtag
              TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
              TYPE(MPI_Comm), INTENT(IN) :: comm
              TYPE(MPI_Status) :: status
              INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_SENDRCV(SENDBUF, SENDCOUNT, SENDTYPE, DEST, SENDTAG, RECVBUF,
             RECVCOUNT, RECVTYPE, SOURCE, RECVTAG, COMM, STATUS, IERROR)
<type> SENDBUF(*), RECVBUF(*)
INTEGER SENDCOUNT, SENDTYPE, DEST, SENDTAG, RECVCOUNT, RECVTYPE,
SOURCE, RECVTAG, COMM, STATUS(MPI_STATUS_SIZE), IERROR

{void MPI::Comm::Sendrecv(const void *sendbuf, int sendcount, const
                          MPI::Datatype& sendtype, int dest, int sendtag, void *recvbuf,
                          int recvcount, const MPI::Datatype& recvtype, int source,
                          int recvtag, MPI::Status& status) const(binding deprecated, see
                          Section 15.2) }

{void MPI::Comm::Sendrecv(const void *sendbuf, int sendcount, const
                          MPI::Datatype& sendtype, int dest, int sendtag, void *recvbuf,
                          int recvcount, const MPI::Datatype& recvtype, int source,
                          int recvtag) const(binding deprecated, see Section 15.2) }

```

Execute a blocking send and receive operation. Both send and receive use the same communicator, but possibly different tags. The send buffer and receive buffers must be disjoint, and may have different lengths and datatypes.

The semantics of a send-receive operation is what would be obtained if the caller forked two concurrent threads, one to execute the send, and one to execute the receive, followed by a join of these two threads.

```

1 MPI_SENDRECV_REPLACE(buf, count, datatype, dest, sendtag, source, recvtag, comm, sta-
2 tus)

```

3	INOUT	buf	initial address of send and receive buffer (choice)
4	IN	count	number of elements in send and receive buffer (non-
5			negative integer)
6			
7	IN	datatype	type of elements in send and receive buffer (handle)
8	IN	dest	rank of destination (integer)
9			
10	IN	sendtag	send message tag (integer)
11	IN	source	rank of source or MPI_ANY_SOURCE (integer)
12	IN	recvtag	receive message tag or MPI_ANY_TAG (integer)
13			
14	IN	comm	communicator (handle)
15	OUT	status	status object (Status)
16			

```

17 int MPI_Sendrecv_replace(void* buf, int count, MPI_Datatype datatype,
18 int dest, int sendtag, int source, int recvtag, MPI_Comm comm,
19 MPI_Status *status)

```

```

20
21 MPI_Sendrecv_replace(buf, count, datatype, dest, sendtag, source, recvtag,
22 comm, status, ierror) BIND(C)
23 TYPE(*), DIMENSION(..) :: buf
24 INTEGER, INTENT(IN) :: count, dest, sendtag, source, recvtag
25 TYPE(MPI_Datatype), INTENT(IN) :: datatype
26 TYPE(MPI_Comm), INTENT(IN) :: comm
27 TYPE(MPI_Status) :: status
28 INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

29 MPI_SENDRECV_REPLACE(BUF, COUNT, DATATYPE, DEST, SENDTAG, SOURCE, RECVTAG,
30 COMM, STATUS, IERROR)
31 <type> BUF(*)
32 INTEGER COUNT, DATATYPE, DEST, SENDTAG, SOURCE, RECVTAG, COMM,
33 STATUS(MPI_STATUS_SIZE), IERROR
34

```

```

35 {void MPI::Comm::Sendrecv_replace(void* buf, int count, const
36 MPI::Datatype& datatype, int dest, int sendtag, int source,
37 int recvtag, MPI::Status& status) const(binding deprecated, see
38 Section 15.2) }

```

```

39 {void MPI::Comm::Sendrecv_replace(void* buf, int count, const
40 MPI::Datatype& datatype, int dest, int sendtag, int source,
41 int recvtag) const(binding deprecated, see Section 15.2) }
42

```

Execute a blocking send and receive. The same buffer is used both for the send and for the receive, so that the message sent is replaced by the message received.

*Advice to implementors.* Additional intermediate buffering is needed for the “replace” variant. (*End of advice to implementors.*)



### 3.11 Null Processes

In many instances, it is convenient to specify a “dummy” source or destination for communication. This simplifies the code that is needed for dealing with boundaries, for example, in the case of a non-circular shift done with calls to send-receive.

The special value `MPI_PROC_NULL` can be used instead of a rank wherever a source or a destination argument is required in a call. A communication with process `MPI_PROC_NULL` has no effect. A send to `MPI_PROC_NULL` succeeds and returns as soon as possible. A receive from `MPI_PROC_NULL` succeeds and returns as soon as possible with no modifications to the receive buffer. When a receive with `source = MPI_PROC_NULL` is executed then the status object returns `source = MPI_PROC_NULL`, `tag = MPI_ANY_TAG` and `count = 0`. A matching probe (cf. Section 3.8.2) with `MPI_PROC_NULL` as source returns `flag = true`, `message = MPI_MESSAGE_NULL`, and the status object returns `source = MPI_PROC_NULL`, `tag = MPI_ANY_TAG`, and `count = 0`.

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# Chapter 4

## Datatypes

Basic datatypes were introduced in Section 3.2.2 Message Data on page 29 and in Section 3.3 Data Type Matching and Data Conversion on page 37. In this chapter, this model is extended to describe any data layout. We consider general datatypes that allow one to transfer efficiently heterogeneous and noncontiguous data. We conclude with the description of calls for explicit packing and unpacking of messages.

### 4.1 Derived Datatypes

Up to here, all point to point communication have involved only buffers containing a sequence of identical basic datatypes. This is too constraining on two accounts. One often wants to pass messages that contain values with different datatypes (e.g., an integer count, followed by a sequence of real numbers); and one often wants to send noncontiguous data (e.g., a sub-block of a matrix). One solution is to pack noncontiguous data into a contiguous buffer at the sender site and unpack it at the receiver site. This has the disadvantage of requiring additional memory-to-memory copy operations at both sites, even when the communication subsystem has scatter-gather capabilities. Instead, MPI provides mechanisms to specify more general, mixed, and noncontiguous communication buffers. It is up to the implementation to decide whether data should be first packed in a contiguous buffer before being transmitted, or whether it can be collected directly from where it resides.

The general mechanisms provided here allow one to transfer directly, without copying, objects of various shape and size. It is not assumed that the MPI library is cognizant of the objects declared in the host language. Thus, if one wants to transfer a structure, or an array section, it will be necessary to provide in MPI a definition of a communication buffer that mimics the definition of the structure or array section in question. These facilities can be used by library designers to define communication functions that can transfer objects defined in the host language — by decoding their definitions as available in a symbol table or a dope vector. Such higher-level communication functions are not part of MPI.

More general communication buffers are specified by replacing the basic datatypes that have been used so far with derived datatypes that are constructed from basic datatypes using the constructors described in this section. These methods of constructing derived datatypes can be applied recursively.

A **general datatype** is an opaque object that specifies two things:

- A sequence of basic datatypes

- A sequence of integer (byte) displacements

The displacements are not required to be positive, distinct, or in increasing order. Therefore, the order of items need not coincide with their order in store, and an item may appear more than once. We call such a pair of sequences (or sequence of pairs) a **type map**. The sequence of basic datatypes (displacements ignored) is the **type signature** of the datatype.

Let

$$Typemap = \{(type_0, disp_0), \dots, (type_{n-1}, disp_{n-1})\},$$

be such a type map, where  $type_i$  are basic types, and  $disp_i$  are displacements. Let

$$Typesig = \{type_0, \dots, type_{n-1}\}$$

be the associated type signature. This type map, together with a base address  $buf$ , specifies a communication buffer: the communication buffer that consists of  $n$  entries, where the  $i$ -th entry is at address  $buf + disp_i$  and has type  $type_i$ . A message assembled from such a communication buffer will consist of  $n$  values, of the types defined by  $Typesig$ .

Most datatype constructors have replication count or block length arguments. Allowed values are non-negative integers. If the value is zero, no elements are generated in the type map and there is no effect on datatype bounds or extent.

We can use a handle to a general datatype as an argument in a send or receive operation, instead of a basic datatype argument. The operation `MPI_SEND(buf, 1, datatype,...)` will use the send buffer defined by the base address `buf` and the general datatype associated with `datatype`; it will generate a message with the type signature determined by the `datatype` argument. `MPI_RECV(buf, 1, datatype,...)` will use the receive buffer defined by the base address `buf` and the general datatype associated with `datatype`.

General datatypes can be used in all send and receive operations. We discuss, in Section 4.1.11, the case where the second argument `count` has value  $> 1$ .

The basic datatypes presented in Section 3.2.2 are particular cases of a general datatype, and are predefined. Thus, `MPI_INT` is a predefined handle to a datatype with type map  $\{(int, 0)\}$ , with one entry of type `int` and displacement zero. The other basic datatypes are similar.

The **extent** of a datatype is defined to be the span from the first byte to the last byte occupied by entries in this datatype, rounded up to satisfy alignment requirements. That is, if

$$Typemap = \{(type_0, disp_0), \dots, (type_{n-1}, disp_{n-1})\},$$

then

$$\begin{aligned} lb(Typemap) &= \min_j disp_j, \\ ub(Typemap) &= \max_j (disp_j + sizeof(type_j)) + \epsilon, \text{ and} \\ extent(Typemap) &= ub(Typemap) - lb(Typemap). \end{aligned} \tag{4.1}$$

If  $type_i$  requires alignment to a byte address that is a multiple of  $k_i$ , then  $\epsilon$  is the least non-negative increment needed to round  $extent(Typemap)$  to the next multiple of  $\max_i k_i$ .

**In Fortran, whether the alignments  $k_i$  are computed according to the alignments used by the compiler in common blocks, SEQUENCE derived types, BIND(C) derived types, or derived**

types that are neither **SEQUENCE** nor **BIND(C)**, is implementation-dependent. The complete definition of **extent** is given [on page 113]in Section 4.1.6 on page 112.

**Example 4.1** Assume that  $Type = \{(\text{double}, 0), (\text{char}, 8)\}$  (a **double** at displacement zero, followed by a **char** at displacement eight). Assume, furthermore, that doubles have to be strictly aligned at addresses that are multiples of eight. Then, the extent of this datatype is 16 (9 rounded to the next multiple of 8). A datatype that consists of a character immediately followed by a double will also have an extent of 16.

*Rationale.* The definition of extent is motivated by the assumption that the amount of padding added at the end of each structure in an array of structures is the least needed to fulfill alignment constraints. More explicit control of the extent is provided in Section 4.1.6. Such explicit control is needed in cases where the assumption does not hold, for example, where union types are used. In Fortran, structures can be expressed with several language features, e.g., common blocks, **SEQUENCE** derived types, or **BIND(C)** derived types. The compiler may use different alignments, and therefore, it is recommended to use **MPI\_TYPE\_CREATE\_RESIZED** for arrays of structures if an alignment may cause an alignment-gap at the end of a structure as described in Section 4.1.6 on page 112 and in Section 16.2.15 on page 675. (*End of rationale.*)

#### 4.1.1 Type Constructors with Explicit Addresses

In Fortran, the functions **MPI\_TYPE\_CREATE\_HVECTOR**, **MPI\_TYPE\_CREATE\_HINDEXED**, **MPI\_TYPE\_CREATE\_STRUCT**, and **MPI\_GET\_ADDRESS** accept arguments of type **INTEGER(KIND=MPI\_ADDRESS\_KIND)**, wherever arguments of type **MPI\_Aint** and **MPI::Aint** are used in C and C++. On Fortran 77 systems that do not support the Fortran 90 **KIND** notation, and where addresses are 64 bits whereas default **INTEGER**s are 32 bits, these arguments will be of type **INTEGER\*8**.

#### 4.1.2 Datatype Constructors

**Contiguous** The simplest datatype constructor is **MPI\_TYPE\_CONTIGUOUS** which allows replication of a datatype into contiguous locations.

**MPI\_TYPE\_CONTIGUOUS**(count, oldtype, newtype)

IN	count	replication count (non-negative integer)
IN	oldtype	old datatype (handle)
OUT	newtype	new datatype (handle)

```
int MPI_Type_contiguous(int count, MPI_Datatype oldtype,
                        MPI_Datatype *newtype)
```

```
MPI_Type_contiguous(count, oldtype, newtype, ierror) BIND(C)
  INTEGER, INTENT(IN) :: count
  TYPE(MPI_Datatype), INTENT(IN) :: oldtype
  TYPE(MPI_Datatype), INTENT(OUT) :: newtype
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
1 MPI_TYPE_CONTIGUOUS(COUNT, OLDTYPE, NEWTYPE, IERROR)
```

```
2     INTEGER COUNT, OLDTYPE, NEWTYPE, IERROR
```

```
3 {MPI::Datatype MPI::Datatype::Create_contiguous(int count) const(binding  
4     deprecated, see Section 15.2) }
```

6 `newtype` is the datatype obtained by concatenating `count` copies of  
7 `oldtype`. Concatenation is defined using *extent* as the size of the concatenated copies.

9 **Example 4.2** Let `oldtype` have type map  $\{(\text{double}, 0), (\text{char}, 8)\}$ , with extent 16, and let  
10 `count` = 3. The type map of the datatype returned by `newtype` is

```
11     {(double, 0), (char, 8), (double, 16), (char, 24), (double, 32), (char, 40)};
```

13 i.e., alternating `double` and `char` elements, with displacements 0, 8, 16, 24, 32, 40.

15 In general, assume that the type map of `oldtype` is

```
16     {(type0, disp0), ..., (typen-1, dispn-1)},
```

18 with extent *ex*. Then `newtype` has a type map with `count · n` entries defined by:

```
19     {(type0, disp0), ..., (typen-1, dispn-1), (type0, disp0 + ex), ..., (typen-1, dispn-1 + ex),  
20     ..., (type0, disp0 + ex · (count - 1)), ..., (typen-1, dispn-1 + ex · (count - 1))}.
```

24 **Vector** The function `MPI_TYPE_VECTOR` is a more general constructor that allows repli-  
25 cation of a datatype into locations that consist of equally spaced blocks. Each block is  
26 obtained by concatenating the same number of copies of the old datatype. The spacing  
27 between blocks is a multiple of the extent of the old datatype.

```
31 MPI_TYPE_VECTOR( count, blocklength, stride, oldtype, newtype)
```

34	IN	count	number of blocks (non-negative integer)
35	IN	blocklength	number of elements in each block (non-negative integer)
36			
37	IN	stride	number of elements between start of each block (integer)
38			
39	IN	oldtype	old datatype (handle)
40	OUT	newtype	new datatype (handle)

```
41 int MPI_Type_vector(int count, int blocklength, int stride,  
42     MPI_Datatype oldtype, MPI_Datatype *newtype)
```

```
43 MPI_Type_vector(count, blocklength, stride, oldtype, newtype, ierror)
```

```
44     BIND(C)
```

```
45     INTEGER, INTENT(IN) :: count, blocklength, stride
```

```

TYPE(MPI_Datatype), INTENT(IN) :: oldtype
TYPE(MPI_Datatype), INTENT(OUT) :: newtype
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_TYPE_VECTOR(COUNT, BLOCKLENGTH, STRIDE, OLDTYPE, NEWTYPE, IERROR)
    INTEGER COUNT, BLOCKLENGTH, STRIDE, OLDTYPE, NEWTYPE, IERROR

{MPI::Datatype MPI::Datatype::Create_vector(int count, int blocklength,
    int stride) const(binding deprecated, see Section 15.2) }
```

**Example 4.3** Assume, again, that `oldtype` has type map  $\{(\text{double}, 0), (\text{char}, 8)\}$ , with extent 16. A call to `MPI_TYPE_VECTOR( 2, 3, 4, oldtype, newtype)` will create the datatype with type map,

$$\{(\text{double}, 0), (\text{char}, 8), (\text{double}, 16), (\text{char}, 24), (\text{double}, 32), (\text{char}, 40),$$

$$(\text{double}, 64), (\text{char}, 72), (\text{double}, 80), (\text{char}, 88), (\text{double}, 96), (\text{char}, 104)\}.$$

That is, two blocks with three copies each of the old type, with a stride of 4 elements ( $4 \cdot 16$  bytes) between the blocks.

**Example 4.4** A call to `MPI_TYPE_VECTOR(3, 1, -2, oldtype, newtype)` will create the datatype,

$$\{(\text{double}, 0), (\text{char}, 8), (\text{double}, -32), (\text{char}, -24), (\text{double}, -64), (\text{char}, -56)\}.$$

In general, assume that `oldtype` has type map,

$$\{(type_0, disp_0), \dots, (type_{n-1}, disp_{n-1})\},$$

with extent  $ex$ . Let  $bl$  be the blocklength. The newly created datatype has a type map with  $count \cdot bl \cdot n$  entries:

$$\{(type_0, disp_0), \dots, (type_{n-1}, disp_{n-1}),$$

$$(type_0, disp_0 + ex), \dots, (type_{n-1}, disp_{n-1} + ex), \dots,$$

$$(type_0, disp_0 + (bl - 1) \cdot ex), \dots, (type_{n-1}, disp_{n-1} + (bl - 1) \cdot ex),$$

$$(type_0, disp_0 + stride \cdot ex), \dots, (type_{n-1}, disp_{n-1} + stride \cdot ex), \dots,$$

$$(type_0, disp_0 + (stride + bl - 1) \cdot ex), \dots, (type_{n-1}, disp_{n-1} + (stride + bl - 1) \cdot ex), \dots,$$

$$(type_0, disp_0 + stride \cdot (count - 1) \cdot ex), \dots,$$

$$(type_{n-1}, disp_{n-1} + stride \cdot (count - 1) \cdot ex), \dots,$$

$$(type_0, disp_0 + (stride \cdot (count - 1) + bl - 1) \cdot ex), \dots,$$

$$(type_{n-1}, disp_{n-1} + (stride \cdot (count - 1) + bl - 1) \cdot ex)\}.$$

A call to `MPI_TYPE_CONTIGUOUS(count, oldtype, newtype)` is equivalent to a call to `MPI_TYPE_VECTOR(count, 1, 1, oldtype, newtype)`, or to a call to `MPI_TYPE_VECTOR(1, count, n, oldtype, newtype)`,  $n$  arbitrary.

Hvector The function `MPI_TYPE_CREATE_HVECTOR` is identical to `MPI_TYPE_VECTOR`, except that `stride` is given in bytes, rather than in elements. The use for both types of vector constructors is illustrated in Section 4.1.14. (H stands for “heterogeneous”).

`MPI_TYPE_CREATE_HVECTOR( count, blocklength, stride, oldtype, newtype)`

IN	count	number of blocks (non-negative integer)
IN	blocklength	number of elements in each block (non-negative integer)
IN	stride	number of bytes between start of each block (integer)
IN	oldtype	old datatype (handle)
OUT	newtype	new datatype (handle)

```
int MPI_Type_create_hvector(int count, int blocklength, MPI_Aint stride,
    MPI_Datatype oldtype, MPI_Datatype *newtype)
```

```
MPI_Type_create_hvector(count, blocklength, stride, oldtype, newtype,
    ierror) BIND(C)
    INTEGER, INTENT(IN) :: count, blocklength
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: stride
    TYPE(MPI_Datatype), INTENT(IN) :: oldtype
    TYPE(MPI_Datatype), INTENT(OUT) :: newtype
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TYPE_CREATE_HVECTOR(COUNT, BLOCKLENGTH, STRIDE, OLDTYPE, NEWTYPE,
    IERROR)
    INTEGER COUNT, BLOCKLENGTH, OLDTYPE, NEWTYPE, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) STRIDE
```

```
{MPI::Datatype MPI::Datatype::Create_hvector(int count, int blocklength,
    MPI::Aint stride) const(binding deprecated, see Section 15.2) }
```

This function replaces `MPI_TYPE_HVECTOR`, whose use is deprecated. See also Chapter 15.

Assume that `oldtype` has type map,

$$\{(type_0, disp_0), \dots, (type_{n-1}, disp_{n-1})\},$$

with extent  $ex$ . Let  $bl$  be the `blocklength`. The newly created datatype has a type map with  $count \cdot bl \cdot n$  entries:

$$\begin{aligned} &\{(type_0, disp_0), \dots, (type_{n-1}, disp_{n-1}), \\ &(type_0, disp_0 + ex), \dots, (type_{n-1}, disp_{n-1} + ex), \dots, \\ &(type_0, disp_0 + (bl - 1) \cdot ex), \dots, (type_{n-1}, disp_{n-1} + (bl - 1) \cdot ex), \\ &(type_0, disp_0 + stride), \dots, (type_{n-1}, disp_{n-1} + stride), \dots, \end{aligned}$$



$$\begin{aligned}
& (type_0, disp_0 + stride + (bl - 1) \cdot ex), \dots, \\
& (type_{n-1}, disp_{n-1} + stride + (bl - 1) \cdot ex), \dots, \\
& (type_0, disp_0 + stride \cdot (count - 1)), \dots, (type_{n-1}, disp_{n-1} + stride \cdot (count - 1)), \dots, \\
& (type_0, disp_0 + stride \cdot (count - 1) + (bl - 1) \cdot ex), \dots, \\
& (type_{n-1}, disp_{n-1} + stride \cdot (count - 1) + (bl - 1) \cdot ex)\}.
\end{aligned}$$

**Indexed** The function `MPI_TYPE_INDEXED` allows replication of an old datatype into a sequence of blocks (each block is a concatenation of the old datatype), where each block can contain a different number of copies and have a different displacement. All block displacements are multiples of the old type extent.

`MPI_TYPE_INDEXED( count, array_of_blocklengths, array_of_displacements, oldtype, newtype)`

IN	count	number of blocks – also number of entries in array_of_displacements and array_of_blocklengths (non-negative integer)
IN	array_of_blocklengths	number of elements per block (array of non-negative integers)
IN	array_of_displacements	displacement for each block, in multiples of oldtype extent (array of integer)
IN	oldtype	old datatype (handle)
OUT	newtype	new datatype (handle)

```

int MPI_Type_indexed(int count, int *array_of_blocklengths,
                    int *array_of_displacements, MPI_Datatype oldtype,
                    MPI_Datatype *newtype)

```

```

MPI_Type_indexed(count, array_of_blocklengths, array_of_displacements,
                oldtype, newtype, ierror) BIND(C)
    INTEGER, INTENT(IN) :: count, array_of_blocklengths(count),
    array_of_displacements(count)
    TYPE(MPI_Datatype), INTENT(IN) :: oldtype
    TYPE(MPI_Datatype), INTENT(OUT) :: newtype
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_TYPE_INDEXED(COUNT, ARRAY_OF_BLOCKLENGTHS, ARRAY_OF_DISPLACEMENTS,
                OLDTYPE, NEWTYPE, IERROR)
    INTEGER COUNT, ARRAY_OF_BLOCKLENGTHS(*), ARRAY_OF_DISPLACEMENTS(*),
    OLDTYPE, NEWTYPE, IERROR

```

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```

1 {MPI::Datatype MPI::Datatype::Create_indexed(int count,
2         const int array_of_blocklengths[],
3         const int array_of_displacements[]) const(binding deprecated, see
4         Section 15.2) }
5
6

```

### Example 4.5

Let `oldtype` have type map  $\{(\text{double}, 0), (\text{char}, 8)\}$ , with extent 16. Let  $B = (3, 1)$  and let  $D = (4, 0)$ . A call to `MPI_TYPE_INDEXED(2, B, D, oldtype, newtype)` returns a datatype with type map,

$$\{(\text{double}, 64), (\text{char}, 72), (\text{double}, 80), (\text{char}, 88), (\text{double}, 96), (\text{char}, 104),$$

$$(\text{double}, 0), (\text{char}, 8)\}.$$

That is, three copies of the old type starting at displacement 64, and one copy starting at displacement 0.

In general, assume that `oldtype` has type map,

$$\{(type_0, disp_0), \dots, (type_{n-1}, disp_{n-1})\},$$

with extent  $ex$ . Let  $B$  be the `array_of_blocklength` argument and  $D$  be the `array_of_displacements` argument. The newly created datatype has  $n \cdot \sum_{i=0}^{\text{count}-1} B[i]$  entries:

$$\{(type_0, disp_0 + D[0] \cdot ex), \dots, (type_{n-1}, disp_{n-1} + D[0] \cdot ex), \dots,$$

$$(type_0, disp_0 + (D[0] + B[0] - 1) \cdot ex), \dots, (type_{n-1}, disp_{n-1} + (D[0] + B[0] - 1) \cdot ex), \dots,$$

$$(type_0, disp_0 + D[\text{count}-1] \cdot ex), \dots, (type_{n-1}, disp_{n-1} + D[\text{count}-1] \cdot ex), \dots,$$

$$(type_0, disp_0 + (D[\text{count}-1] + B[\text{count}-1] - 1) \cdot ex), \dots,$$

$$(type_{n-1}, disp_{n-1} + (D[\text{count}-1] + B[\text{count}-1] - 1) \cdot ex)\}.$$

A call to `MPI_TYPE_VECTOR(count, blocklength, stride, oldtype, newtype)` is equivalent to a call to `MPI_TYPE_INDEXED(count, B, D, oldtype, newtype)` where

$$D[j] = j \cdot \text{stride}, \quad j = 0, \dots, \text{count} - 1,$$

and

$$B[j] = \text{blocklength}, \quad j = 0, \dots, \text{count} - 1.$$

Hindexed The function `MPI_TYPE_CREATE_HINDEXED` is identical to `MPI_TYPE_INDEXED`, except that block displacements in `array_of_displacements` are specified in bytes, rather than in multiples of the `oldtype` extent.

`MPI_TYPE_CREATE_HINDEXED( count, array_of_blocklengths, array_of_displacements, oldtype, newtype)`

IN	count	number of blocks — also number of entries in <code>array_of_displacements</code> and <code>array_of_blocklengths</code> (non-negative integer)
IN	array_of_blocklengths	number of elements in each block (array of non-negative integers)
IN	array_of_displacements	byte displacement of each block (array of integer)
IN	oldtype	old datatype (handle)
OUT	newtype	new datatype (handle)

```
int MPI_Type_create_hindexed(int count, int array_of_blocklengths[],
                             MPI_Aint array_of_displacements[], MPI_Datatype oldtype,
                             MPI_Datatype *newtype)
```

```
MPI_Type_create_hindexed(count, array_of_blocklengths,
                          array_of_displacements, oldtype, newtype, ierror) BIND(C)
INTEGER, INTENT(IN) :: count, array_of_blocklengths(count)
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) ::
array_of_displacements(count)
TYPE(MPI_Datatype), INTENT(IN) :: oldtype
TYPE(MPI_Datatype), INTENT(OUT) :: newtype
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TYPE_CREATE_HINDEXED(COUNT, ARRAY_OF_BLOCKLENGTHS,
                          ARRAY_OF_DISPLACEMENTS, OLDTYPE, NEWTYPE, IERROR)
INTEGER COUNT, ARRAY_OF_BLOCKLENGTHS(*), OLDTYPE, NEWTYPE, IERROR
INTEGER(KIND=MPI_ADDRESS_KIND) ARRAY_OF_DISPLACEMENTS(*)
```

```
{MPI::Datatype MPI::Datatype::Create_hindexed(int count,
const int array_of_blocklengths[],
const MPI::Aint array_of_displacements[]) const(binding
deprecated, see Section 15.2) }
```

This function replaces `MPI_TYPE_HINDEXED`, whose use is deprecated. See also Chapter 15.

Assume that `oldtype` has type map,

$$\{(type_0, disp_0), \dots, (type_{n-1}, disp_{n-1})\},$$

with extent *ex*. Let `B` be the `array_of_blocklength` argument and `D` be the `array_of_displacements` argument. The newly created datatype has a type map with  $n \cdot$

$\sum_{i=0}^{\text{count}-1} B[i]$  entries:

$$\{(type_0, disp_0 + D[0]), \dots, (type_{n-1}, disp_{n-1} + D[0]), \dots, \\ (type_0, disp_0 + D[0] + (B[0] - 1) \cdot ex), \dots, \\ (type_{n-1}, disp_{n-1} + D[0] + (B[0] - 1) \cdot ex), \dots, \\ (type_0, disp_0 + D[\text{count}-1]), \dots, (type_{n-1}, disp_{n-1} + D[\text{count}-1]), \dots, \\ (type_0, disp_0 + D[\text{count}-1] + (B[\text{count}-1] - 1) \cdot ex), \dots, \\ (type_{n-1}, disp_{n-1} + D[\text{count}-1] + (B[\text{count}-1] - 1) \cdot ex)\}.$$

**Indexed\_block** This function is the same as `MPI_TYPE_INDEXED` except that the block-length is the same for all blocks. There are many codes using indirect addressing arising from unstructured grids where the blocksize is always 1 (gather/scatter). The following convenience function allows for constant blocksize and arbitrary displacements.

`MPI_TYPE_CREATE_INDEXED_BLOCK(count, blocklength, array_of_displacements, oldtype, newtype)`

IN	count	length of array of displacements (non-negative integer)
IN	blocklength	size of block (non-negative integer)
IN	array_of_displacements	array of displacements (array of integer)
IN	oldtype	old datatype (handle)
OUT	newtype	new datatype (handle)

```
int MPI_Type_create_indexed_block(int count, int blocklength,
    int array_of_displacements[], MPI_Datatype oldtype,
    MPI_Datatype *newtype)
```

```
MPI_Type_create_indexed_block(count, blocklength, array_of_displacements,
    oldtype, newtype, ierror) BIND(C)
    INTEGER, INTENT(IN) :: count, blocklength,
    array_of_displacements(count)
    TYPE(MPI_Datatype), INTENT(IN) :: oldtype
    TYPE(MPI_Datatype), INTENT(OUT) :: newtype
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TYPE_CREATE_INDEXED_BLOCK(COUNT, BLOCKLENGTH, ARRAY_OF_DISPLACEMENTS,
    OLDTYPE, NEWTYPE, IERROR)
    INTEGER COUNT, BLOCKLENGTH, ARRAY_OF_DISPLACEMENTS(*), OLDTYPE,
    NEWTYPE, IERROR
```

```
{MPI::Datatype MPI::Datatype::Create_indexed_block(int count,
    int blocklength,
```

```
const int array_of_displacements[]) const(binding deprecated, see
Section 15.2) }
```

Struct `MPI_TYPE_STRUCT` is the most general type constructor. It further generalizes `MPI_TYPE_CREATE_HINDEXED` in that it allows each block to consist of replications of different datatypes.

```
MPI_TYPE_CREATE_STRUCT(count, array_of_blocklengths, array_of_displacements,
array_of_types, newtype)
```

IN	count	number of blocks (non-negative integer) — also number of entries in arrays <code>array_of_types</code> , <code>array_of_displacements</code> and <code>array_of_blocklengths</code>
IN	array_of_blocklength	number of elements in each block (array of non-negative integer)
IN	array_of_displacements	byte displacement of each block (array of integer)
IN	array_of_types	type of elements in each block (array of handles to datatype objects)
OUT	newtype	new datatype (handle)

```
int MPI_Type_create_struct(int count, int array_of_blocklengths[],
MPI_Aint array_of_displacements[],
MPI_Datatype array_of_types[], MPI_Datatype *newtype)
```

```
MPI_Type_create_struct(count, array_of_blocklengths,
array_of_displacements, array_of_types, newtype, ierror)
BIND(C)
INTEGER, INTENT(IN) :: count, array_of_blocklengths(count)
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) ::
array_of_displacements(count)
TYPE(MPI_Datatype), INTENT(IN) :: array_of_types(count)
TYPE(MPI_Datatype), INTENT(OUT) :: newtype
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TYPE_CREATE_STRUCT(COUNT, ARRAY_OF_BLOCKLENGTHS,
ARRAY_OF_DISPLACEMENTS, ARRAY_OF_TYPES, NEWTYPE, IERROR)
INTEGER COUNT, ARRAY_OF_BLOCKLENGTHS(*), ARRAY_OF_TYPES(*), NEWTYPE,
IERROR
INTEGER(KIND=MPI_ADDRESS_KIND) ARRAY_OF_DISPLACEMENTS(*)
```

```
{static MPI::Datatype MPI::Datatype::Create_struct(int count,
const int array_of_blocklengths[], const MPI::Aint
array_of_displacements[],
const MPI::Datatype array_of_types[]) (binding deprecated, see
Section 15.2) }
```

This function replaces `MPI_TYPE_STRUCT`, whose use is deprecated. See also Chapter 15.

**Example 4.6** Let `type1` have type map,

$$\{(\text{double}, 0), (\text{char}, 8)\},$$

with extent 16. Let  $B = (2, 1, 3)$ ,  $D = (0, 16, 26)$ , and  $T = (\text{MPI\_FLOAT}, \text{type1}, \text{MPI\_CHAR})$ . Then a call to `MPI_TYPE_STRUCT(3, B, D, T, newtype)` returns a datatype with type map,

$$\{(\text{float}, 0), (\text{float}, 4), (\text{double}, 16), (\text{char}, 24), (\text{char}, 26), (\text{char}, 27), (\text{char}, 28)\}.$$

That is, two copies of `MPI_FLOAT` starting at 0, followed by one copy of `type1` starting at 16, followed by three copies of `MPI_CHAR`, starting at 26. (We assume that a float occupies four bytes.)

In general, let  $T$  be the `array_of_types` argument, where  $T[i]$  is a handle to,

$$\text{typemap}_i = \{(type_0^i, disp_0^i), \dots, (type_{n_i-1}^i, disp_{n_i-1}^i)\},$$

with extent  $ex_i$ . Let  $B$  be the `array_of_blocklength` argument and  $D$  be the `array_of_displacements` argument. Let  $c$  be the count argument. Then the newly created datatype has a type map with  $\sum_{i=0}^{c-1} B[i] \cdot n_i$  entries:

$$\begin{aligned} &\{(type_0^0, disp_0^0 + D[0]), \dots, (type_{n_0}^0, disp_{n_0}^0 + D[0]), \dots, \\ &(type_0^0, disp_0^0 + D[0] + (B[0] - 1) \cdot ex_0), \dots, (type_{n_0}^0, disp_{n_0}^0 + D[0] + (B[0]-1) \cdot ex_0), \dots, \\ &(type_0^{c-1}, disp_0^{c-1} + D[c-1]), \dots, (type_{n_{c-1}-1}^{c-1}, disp_{n_{c-1}-1}^{c-1} + D[c-1]), \dots, \\ &(type_0^{c-1}, disp_0^{c-1} + D[c-1] + (B[c-1] - 1) \cdot ex_{c-1}), \dots, \\ &(type_{n_{c-1}-1}^{c-1}, disp_{n_{c-1}-1}^{c-1} + D[c-1] + (B[c-1]-1) \cdot ex_{c-1})\}. \end{aligned}$$

A call to `MPI_TYPE_CREATE_HINDEXED(count, B, D, oldtype, newtype)` is equivalent to a call to `MPI_TYPE_CREATE_STRUCT(count, B, D, T, newtype)`, where each entry of  $T$  is equal to `oldtype`.

## 4.1.3 Subarray Datatype Constructor

```
MPI_TYPE_CREATE_SUBARRAY(ndims, array_of_sizes, array_of_subsizes, array_of_starts,
                          order, oldtype, newtype)
```

IN	ndims	number of array dimensions (positive integer)
IN	array_of_sizes	number of elements of type <code>oldtype</code> in each dimension of the full array (array of positive integers)
IN	array_of_subsizes	number of elements of type <code>oldtype</code> in each dimension of the subarray (array of positive integers)
IN	array_of_starts	starting coordinates of the subarray in each dimension (array of non-negative integers)
IN	order	array storage order flag (state)
IN	oldtype	array element datatype (handle)
OUT	newtype	new datatype (handle)

```
int MPI_Type_create_subarray(int ndims, int array_of_sizes[],
                           int array_of_subsizes[], int array_of_starts[], int order,
                           MPI_Datatype oldtype, MPI_Datatype *newtype)
```

```
MPI_Type_create_subarray(ndims, array_of_sizes, array_of_subsizes,
                        array_of_starts, order, oldtype, newtype, ierror) BIND(C)
INTEGER, INTENT(IN) :: ndims, array_of_sizes(ndims),
array_of_subsizes(ndims), array_of_starts(ndims), order
TYPE(MPI_Datatype), INTENT(IN) :: oldtype
TYPE(MPI_Datatype), INTENT(OUT) :: newtype
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TYPE_CREATE_SUBARRAY(NDIMS, ARRAY_OF_SIZES, ARRAY_OF_SUBSIZES,
                        ARRAY_OF_STARTS, ORDER, OLDTYPE, NEWTYPE, IERROR)
INTEGER NDIMS, ARRAY_OF_SIZES(*), ARRAY_OF_SUBSIZES(*),
ARRAY_OF_STARTS(*), ORDER, OLDTYPE, NEWTYPE, IERROR
```

```
{MPI::Datatype MPI::Datatype::Create_subarray(int ndims,
        const int array_of_sizes[], const int array_of_subsizes[],
        const int array_of_starts[], int order) const(binding deprecated,
        see Section 15.2) }
```

The subarray type constructor creates an MPI datatype describing an  $n$ -dimensional subarray of an  $n$ -dimensional array. The subarray may be situated anywhere within the full array, and may be of any nonzero size up to the size of the larger array as long as it is confined within this array. This type constructor facilitates creating filetypes to access arrays distributed in blocks among processes to a single file that contains the global array, see MPI I/O, especially Section 13.1.1 on page 509.

This type constructor can handle arrays with an arbitrary number of dimensions and works for both C and Fortran ordered matrices (i.e., row-major or column-major). Note that a C program may use Fortran order and a Fortran program may use C order.

The `ndims` parameter specifies the number of dimensions in the full data array and gives the number of elements in `array_of_sizes`, `array_of_subsizes`, and `array_of_starts`.

The number of elements of type `oldtype` in each dimension of the  $n$ -dimensional array and the requested subarray are specified by `array_of_sizes` and `array_of_subsizes`, respectively. For any dimension  $i$ , it is erroneous to specify `array_of_subsizes[i] < 1` or `array_of_subsizes[i] > array_of_sizes[i]`.

The `array_of_starts` contains the starting coordinates of each dimension of the subarray. Arrays are assumed to be indexed starting from zero. For any dimension  $i$ , it is erroneous to specify `array_of_starts[i] < 0` or `array_of_starts[i] > (array_of_sizes[i] - array_of_subsizes[i])`.

*Advice to users.* In a Fortran program with arrays indexed starting from 1, if the starting coordinate of a particular dimension of the subarray is  $n$ , then the entry in `array_of_starts` for that dimension is  $n-1$ . (*End of advice to users.*)

The `order` argument specifies the storage order for the subarray as well as the full array. It must be set to one of the following:

**MPI\_ORDER\_C** The ordering used by C arrays, (i.e., row-major order)

**MPI\_ORDER\_FORTRAN** The ordering used by Fortran arrays, (i.e., column-major order)

A `ndims`-dimensional subarray (`newtype`) with no extra padding can be defined by the function `Subarray()` as follows:

```
newtype = Subarray(ndims, {size0, size1, ..., sizendims-1},
                    {subsize0, subsize1, ..., subsizendims-1},
                    {start0, start1, ..., startndims-1}, oldtype)
```

Let the typemap of `oldtype` have the form:

$$\{(type_0, disp_0), (type_1, disp_1), \dots, (type_{n-1}, disp_{n-1})\}$$

where  $type_i$  is a predefined MPI datatype, and let  $ex$  be the extent of `oldtype`. Then we define the `Subarray()` function recursively using the following three equations. Equation 4.2 defines the base step. Equation 4.3 defines the recursion step when `order = MPI_ORDER_FORTRAN`, and Equation 4.4 defines the recursion step when `order = MPI_ORDER_C`.

$$\begin{aligned} & \text{Subarray}(1, \{size_0\}, \{subsize_0\}, \{start_0\}, \\ & \quad \{(type_0, disp_0), (type_1, disp_1), \dots, (type_{n-1}, disp_{n-1})\}) \\ &= \{(MPI\_LB, 0), \\ & \quad (type_0, disp_0 + start_0 \times ex), \dots, (type_{n-1}, disp_{n-1} + start_0 \times ex), \\ & \quad (type_0, disp_0 + (start_0 + 1) \times ex), \dots, (type_{n-1}, \\ & \quad \quad disp_{n-1} + (start_0 + 1) \times ex), \dots \\ & \quad (type_0, disp_0 + (start_0 + subsize_0 - 1) \times ex), \dots, \\ & \quad \quad (type_{n-1}, disp_{n-1} + (start_0 + subsize_0 - 1) \times ex), \\ & \quad (MPI\_UB, size_0 \times ex)\} \end{aligned} \tag{4.2}$$



$$\begin{aligned}
& \text{Subarray}(\text{ndims}, \{size_0, size_1, \dots, size_{\text{ndims}-1}\}, \\
& \quad \{subsize_0, subsize_1, \dots, subsize_{\text{ndims}-1}\}, \\
& \quad \{start_0, start_1, \dots, start_{\text{ndims}-1}\}, \text{oldtype}) \\
= & \text{Subarray}(\text{ndims} - 1, \{size_1, size_2, \dots, size_{\text{ndims}-1}\}, \\
& \quad \{subsize_1, subsize_2, \dots, subsize_{\text{ndims}-1}\}, \\
& \quad \{start_1, start_2, \dots, start_{\text{ndims}-1}\}, \\
& \quad \text{Subarray}(1, \{size_0\}, \{subsize_0\}, \{start_0\}, \text{oldtype}))
\end{aligned}
\tag{4.3}$$

$$\begin{aligned}
& \text{Subarray}(\text{ndims}, \{size_0, size_1, \dots, size_{\text{ndims}-1}\}, \\
& \quad \{subsize_0, subsize_1, \dots, subsize_{\text{ndims}-1}\}, \\
& \quad \{start_0, start_1, \dots, start_{\text{ndims}-1}\}, \text{oldtype}) \\
= & \text{Subarray}(\text{ndims} - 1, \{size_0, size_1, \dots, size_{\text{ndims}-2}\}, \\
& \quad \{subsize_0, subsize_1, \dots, subsize_{\text{ndims}-2}\}, \\
& \quad \{start_0, start_1, \dots, start_{\text{ndims}-2}\}, \\
& \quad \text{Subarray}(1, \{size_{\text{ndims}-1}\}, \{subsize_{\text{ndims}-1}\}, \{start_{\text{ndims}-1}\}, \text{oldtype}))
\end{aligned}
\tag{4.4}$$

For an example use of `MPI_TYPE_CREATE_SUBARRAY` in the context of I/O see Section 13.9.2.

#### 4.1.4 Distributed Array Datatype Constructor

The distributed array type constructor supports HPF-like [42] data distributions. However, unlike in HPF, the storage order may be specified for C arrays as well as for Fortran arrays.

*Advice to users.* One can create an HPF-like file view using this type constructor as follows. Complementary filetypes are created by having every process of a group call this constructor with identical arguments (with the exception of `rank` which should be set appropriately). These filetypes (along with identical `disp` and `etype`) are then used to define the view (via `MPI_FILE_SET_VIEW`), see MPI I/O, especially Section 13.1.1 on page 509 and Section 13.3 on page 522. Using this view, a collective data access operation (with identical offsets) will yield an HPF-like distribution pattern. (*End of advice to users.*)

```

1 MPI_TYPE_CREATE_DARRAY(size, rank, ndims, array_of_gsizes, array_of_distrib,
2   array_of_dargs, array_of_psize, order, oldtype, newtype)
3
4   IN      size                size of process group (positive integer)
5   IN      rank                rank in process group (non-negative integer)
6   IN      ndims               number of array dimensions as well as process grid
7   dimensions (positive integer)
8   IN      array_of_gsizes      number of elements of type oldtype in each dimension
9   of global array (array of positive integers)
10
11  IN      array_of_distrib      distribution of array in each dimension (array of state)
12  IN      array_of_dargs        distribution argument in each dimension (array of positive
13  integers)
14  IN      array_of_psize        size of process grid in each dimension (array of positive
15  integers)
16
17  IN      order                 array storage order flag (state)
18  IN      oldtype               old datatype (handle)
19  OUT     newtype               new datatype (handle)
20

```

```

21
22 int MPI_Type_create_darray(int size, int rank, int ndims,
23   int array_of_gsizes[], int array_of_distrib[], int
24   array_of_dargs[], int array_of_psize[], int order,
25   MPI_Datatype oldtype, MPI_Datatype *newtype)

```

```

26 MPI_Type_create_darray(size, rank, ndims, array_of_gsizes,
27   array_of_distrib, array_of_dargs, array_of_psize, order,
28   oldtype, newtype, ierror) BIND(C)
29   INTEGER, INTENT(IN) :: size, rank, ndims, array_of_gsizes(ndims),
30   array_of_distrib(ndims), array_of_dargs(ndims),
31   array_of_psize(ndims), order
32   TYPE(MPI_Datatype), INTENT(IN) :: oldtype
33   TYPE(MPI_Datatype), INTENT(OUT) :: newtype
34   INTEGER, OPTIONAL, INTENT(OUT) :: ierror
35
36 MPI_TYPE_CREATE_DARRAY(SIZE, RANK, NDIMS, ARRAY_OF_GSIZES,
37   ARRAY_OF_DISTRIBS, ARRAY_OF_DARGS, ARRAY_OF_PSIZE, ORDER,
38   OLDTYPE, NEWTYPE, IERROR)
39   INTEGER SIZE, RANK, NDIMS, ARRAY_OF_GSIZES(*), ARRAY_OF_DISTRIBS(*),
40   ARRAY_OF_DARGS(*), ARRAY_OF_PSIZE(*), ORDER, OLDTYPE, NEWTYPE, IERROR
41 {MPI::Datatype MPI::Datatype::Create_darray(int size, int rank, int ndims,
42   const int array_of_gsizes[], const int array_of_distrib[],
43   const int array_of_dargs[], const int array_of_psize[],
44   int order) const(binding deprecated, see Section 15.2) }
45

```

MPI\_TYPE\_CREATE\_DARRAY can be used to generate the datatypes corresponding to the distribution of an `ndims`-dimensional array of `oldtype` elements onto an `ndims`-dimensional grid of logical processes. Unused dimensions of `array_of_psize` should be set to 1. (See

Example 4.7, page 109.) For a call to `MPI_TYPE_CREATE_DARRAY` to be correct, the equation  $\prod_{i=0}^{ndims-1} array\_of\_psizes[i] = size$  must be satisfied. The ordering of processes in the process grid is assumed to be row-major, as in the case of virtual Cartesian process topologies .

*Advice to users.* For both Fortran and C arrays, the ordering of processes in the process grid is assumed to be row-major. This is consistent with the ordering used in virtual Cartesian process topologies in MPI. To create such virtual process topologies, or to find the coordinates of a process in the process grid, etc., users may use the corresponding process topology functions, see Chapter 7 on page 301. (*End of advice to users.*)

Each dimension of the array can be distributed in one of three ways:

- `MPI_DISTRIBUTE_BLOCK` - Block distribution
- `MPI_DISTRIBUTE_CYCLIC` - Cyclic distribution
- `MPI_DISTRIBUTE_NONE` - Dimension not distributed.

The constant `MPI_DISTRIBUTE_DFLT_DARG` specifies a default distribution argument. The distribution argument for a dimension that is not distributed is ignored. For any dimension  $i$  in which the distribution is `MPI_DISTRIBUTE_BLOCK`, it is erroneous to specify  $array\_of\_dargs[i] * array\_of\_psizes[i] < array\_of\_gsizes[i]$ .

For example, the HPF layout `ARRAY(CYCLIC(15))` corresponds to `MPI_DISTRIBUTE_CYCLIC` with a distribution argument of 15, and the HPF layout `ARRAY(BLOCK)` corresponds to `MPI_DISTRIBUTE_BLOCK` with a distribution argument of `MPI_DISTRIBUTE_DFLT_DARG`.

The `order` argument is used as in `MPI_TYPE_CREATE_SUBARRAY` to specify the storage order. Therefore, arrays described by this type constructor may be stored in Fortran (column-major) or C (row-major) order. Valid values for `order` are `MPI_ORDER_FORTRAN` and `MPI_ORDER_C`.

This routine creates a new MPI datatype with a typemap defined in terms of a function called “cyclic()” (see below).

Without loss of generality, it suffices to define the typemap for the `MPI_DISTRIBUTE_CYCLIC` case where `MPI_DISTRIBUTE_DFLT_DARG` is not used.

`MPI_DISTRIBUTE_BLOCK` and `MPI_DISTRIBUTE_NONE` can be reduced to the `MPI_DISTRIBUTE_CYCLIC` case for dimension  $i$  as follows.

`MPI_DISTRIBUTE_BLOCK` with  $array\_of\_dargs[i]$  equal to `MPI_DISTRIBUTE_DFLT_DARG` is equivalent to `MPI_DISTRIBUTE_CYCLIC` with  $array\_of\_dargs[i]$  set to

$$(array\_of\_gsizes[i] + array\_of\_psizes[i] - 1) / array\_of\_psizes[i].$$

If  $array\_of\_dargs[i]$  is not `MPI_DISTRIBUTE_DFLT_DARG`, then `MPI_DISTRIBUTE_BLOCK` and `MPI_DISTRIBUTE_CYCLIC` are equivalent.

`MPI_DISTRIBUTE_NONE` is equivalent to `MPI_DISTRIBUTE_CYCLIC` with  $array\_of\_dargs[i]$  set to  $array\_of\_gsizes[i]$ .

Finally, `MPI_DISTRIBUTE_CYCLIC` with  $array\_of\_dargs[i]$  equal to `MPI_DISTRIBUTE_DFLT_DARG` is equivalent to `MPI_DISTRIBUTE_CYCLIC` with  $array\_of\_dargs[i]$  set to 1.

For `MPI_ORDER_FORTRAN`, an  $ndims$ -dimensional distributed array (`newtype`) is defined by the following code fragment:

```

1  oldtype[0] = oldtype;
2  for ( i = 0; i < ndims; i++ ) {
3      oldtype[i+1] = cyclic(array_of_dargs[i],
4                          array_of_gsizes[i],
5                          r[i],
6                          array_of_psize[i],
7                          oldtype[i]);
8  }
9  newtype = oldtype[ndims];

```

For MPI\_ORDER\_C, the code is:

```

12 oldtype[0] = oldtype;
13 for ( i = 0; i < ndims; i++ ) {
14     oldtype[i + 1] = cyclic(array_of_dargs[ndims - i - 1],
15                           array_of_gsizes[ndims - i - 1],
16                           r[ndims - i - 1],
17                           array_of_psize[ndims - i - 1],
18                           oldtype[i]);
19 }
20 newtype = oldtype[ndims];

```

where  $r[i]$  is the position of the process (with rank rank) in the process grid at dimension  $i$ . The values of  $r[i]$  are given by the following code fragment:

```

26 t_rank = rank;
27 t_size = 1;
28 for (i = 0; i < ndims; i++)
29     t_size *= array_of_psize[i];
30 for (i = 0; i < ndims; i++) {
31     t_size = t_size / array_of_psize[i];
32     r[i] = t_rank / t_size;
33     t_rank = t_rank % t_size;
34 }

```

Let the typemap of oldtype have the form:

$$\{(type_0, disp_0), (type_1, disp_1), \dots, (type_{n-1}, disp_{n-1})\}$$

where  $type_i$  is a predefined MPI datatype, and let  $ex$  be the extent of oldtype.

Given the above, the function cyclic() is defined as follows:

```

42 cyclic(darg, gsize, r, psize, oldtype)
43 =  {(MPI_LB, 0),
44     (type0, disp0 + r × darg × ex), ...,
45     (typen-1, dispn-1 + r × darg × ex),
46     (type0, disp0 + (r × darg + 1) × ex), ...,
47     (typen-1, dispn-1 + (r × darg + 1) × ex),

```

```

...
(type0, disp0 + ((r + 1) × darg - 1) × ex), ...,
      (typen-1, dispn-1 + ((r + 1) × darg - 1) × ex),

(type0, disp0 + r × darg × ex + psize × darg × ex), ...,
      (typen-1, dispn-1 + r × darg × ex + psize × darg × ex),
(type0, disp0 + (r × darg + 1) × ex + psize × darg × ex), ...,
      (typen-1, dispn-1 + (r × darg + 1) × ex + psize × darg × ex),
...
(type0, disp0 + ((r + 1) × darg - 1) × ex + psize × darg × ex), ...,
      (typen-1, dispn-1 + ((r + 1) × darg - 1) × ex + psize × darg × ex),
      ⋮
(type0, disp0 + r × darg × ex + psize × darg × ex × (count - 1)), ...,
      (typen-1, dispn-1 + r × darg × ex + psize × darg × ex × (count - 1)),
(type0, disp0 + (r × darg + 1) × ex + psize × darg × ex × (count - 1)), ...,
      (typen-1, dispn-1 + (r × darg + 1) × ex
        + psize × darg × ex × (count - 1)),
...
(type0, disp0 + (r × darg + darglast - 1) × ex
      + psize × darg × ex × (count - 1)), ...,
      (typen-1, dispn-1 + (r × darg + darglast - 1) × ex
        + psize × darg × ex × (count - 1)),
(MPI_UB, gsize × ex)

```

where *count* is defined by this code fragment:

```

nblocks = (gsize + (darg - 1)) / darg;
count = nblocks / psize;
left_over = nblocks - count * psize;
if (r < left_over)
    count = count + 1;

```

Here, *nblocks* is the number of blocks that must be distributed among the processors. Finally, *darg<sub>last</sub>* is defined by this code fragment:

```

if ((num_in_last_cyclic = gsize % (psize * darg)) == 0)
    darg_last = darg;
else
    darg_last = num_in_last_cyclic - darg * r;
    if (darg_last > darg)
        darg_last = darg;
    if (darg_last <= 0)
        darg_last = darg;

```

**Example 4.7** Consider generating the filetypes corresponding to the HPF distribution:

```

1      <oldtype> FILEARRAY(100, 200, 300)
2      !HPF$ PROCESSORS PROCESSES(2, 3)
3      !HPF$ DISTRIBUTE FILEARRAY(CYCLIC(10), *, BLOCK) ONTO PROCESSES

```

This can be achieved by the following Fortran code, assuming there will be six processes attached to the run:

```

7      ndims = 3
8      array_of_gsizes(1) = 100
9      array_of_distribs(1) = MPI_DISTRIBUTE_CYCLIC
10     array_of_dargs(1) = 10
11     array_of_gsizes(2) = 200
12     array_of_distribs(2) = MPI_DISTRIBUTE_NONE
13     array_of_dargs(2) = 0
14     array_of_gsizes(3) = 300
15     array_of_distribs(3) = MPI_DISTRIBUTE_BLOCK
16     array_of_dargs(3) = MPI_DISTRIBUTE_DFLT_DARG
17     array_of_psizes(1) = 2
18     array_of_psizes(2) = 1
19     array_of_psizes(3) = 3
20     call MPI_COMM_SIZE(MPI_COMM_WORLD, size, ierr)
21     call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
22     call MPI_TYPE_CREATE_DARRAY(size, rank, ndims, array_of_gsizes, &
23                                array_of_distribs, array_of_dargs, array_of_psizes, &
24                                MPI_ORDER_FORTRAN, oldtype, newtype, ierr)

```

#### 4.1.5 Address and Size Functions

The displacements in a general datatype are relative to some initial buffer address. **Absolute addresses** can be substituted for these displacements: we treat them as displacements relative to “address zero,” the start of the address space. This initial address zero is indicated by the constant `MPI_BOTTOM`. Thus, a datatype can specify the absolute address of the entries in the communication buffer, in which case the `buf` argument is passed the value `MPI_BOTTOM`.

The address of a location in memory can be found by invoking the function `MPI_GET_ADDRESS`.

```

38 MPI_GET_ADDRESS(location, address)

```

IN	location	location in caller memory (choice)
OUT	address	address of location (integer)

```

42 int MPI_Get_address(void *location, MPI_Aint *address)

```

```

43 MPI_Get_address(location, address, ierror) BIND(C)
44     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: location
45     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: address
46     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```
MPI_GET_ADDRESS(LOCATION, ADDRESS, IERROR)
```

```
<type> LOCATION(*)
```

```
INTEGER IERROR
```

```
INTEGER(KIND=MPI_ADDRESS_KIND) ADDRESS
```

```
{MPI::Aint MPI::Get_address(void* location) (binding deprecated, see Section 15.2)
}
```

This function replaces MPI\_ADDRESS, whose use is deprecated. See also Chapter 15. Returns the (byte) address of location.

*Advice to users.* Current Fortran MPI codes will run unmodified, and will port to any system. However, they may fail if addresses larger than  $2^{32} - 1$  are used in the program. New codes should be written so that they use the new functions. This provides compatibility with C/C++ and avoids errors on 64 bit architectures. However, such newly written codes may need to be (slightly) rewritten to port to old Fortran 77 environments that do not support KIND declarations. (*End of advice to users.*)

*Rationale.* In the `mpi_f08` module, the `location` argument is not defined with `INTENT(IN)` because existing applications may use `MPI_GET_ADDRESS` (and the deprecated `MPI_ADDRESS`) as a substitute for `MPI_F_SYNC_REG` that was not defined before MPI-3.0. (*End of rationale.*)

**Example 4.8** Using MPI\_GET\_ADDRESS for an array.

```
REAL A(100,100)
INTEGER(KIND=MPI_ADDRESS_KIND) I1, I2, DIFF
CALL MPI_GET_ADDRESS(A(1,1), I1, IERROR)
CALL MPI_GET_ADDRESS(A(10,10), I2, IERROR)
DIFF = I2 - I1
! The value of DIFF is 909*sizeofreal; the values of I1 and I2 are
! implementation dependent.
```

*Advice to users.* C users may be tempted to avoid the usage of MPI\_GET\_ADDRESS and rely on the availability of the address operator `&`. Note, however, that `& cast-expression` is a pointer, not an address. ISO C does not require that the value of a pointer (or the pointer cast to int) be the absolute address of the object pointed at — although this is commonly the case. Furthermore, referencing may not have a unique definition on machines with a segmented address space. The use of MPI\_GET\_ADDRESS to “reference” C variables guarantees portability to such machines as well. (*End of advice to users.*)

*Advice to users.* To prevent problems with the argument copying and register optimization done by Fortran compilers, please note the hints in [subsections “Problems Due to Data Copying and Sequence Association,” and “A Problem with Register Optimization” in Section 16.2.10 on pages 671 and 677. ]Sections 16.2.10-16.2.20, especially in Sections 16.2.12 and 16.2.13 on pages 671-674 about “Problems Due to

Data Copying and Sequence Association with Subscript Triplets” and “Vector Subscripts”, and in Sections 16.2.16 to 16.2.19 on pages 677 to 686 about “Optimization Problems”, “Code Movements and Register Optimization”, “Temporary Data Movements” and “Permanent Data Movements”. (*End of advice to users.*)

ticket238-J.

The following auxiliary function provides useful information on derived datatypes.

MPI\_TYPE\_SIZE(datatype, size)

IN	datatype	datatype (handle)
OUT	size	datatype size (integer)

int MPI\_Type\_size(MPI\_Datatype datatype, int \*size)

```
MPI_Type_size(datatype, size, ierror) BIND(C)
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  INTEGER, INTENT(OUT) :: size
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TYPE_SIZE(DATATYPE, SIZE, IERROR)
  INTEGER DATATYPE, SIZE, IERROR
```

```
{int MPI::Datatype::Get_size() const(binding deprecated, see Section 15.2) }
```

MPI\_TYPE\_SIZE returns the total size, in bytes, of the entries in the type signature associated with `datatype`; i.e., the total size of the data in a message that would be created with this datatype. Entries that occur multiple times in the datatype are counted with their multiplicity.

#### 4.1.6 Lower-Bound and Upper-Bound Markers

It is often convenient to define explicitly the lower bound and upper bound of a type map, and override the definition given on page 113. This allows one to define a datatype that has “holes” at its beginning or its end, or a datatype with entries that extend above the upper bound or below the lower bound. Examples of such usage are provided in Section 4.1.14. Also, the user may want to override the alignment rules that are used to compute upper bounds and extents. E.g., a C compiler may allow the user to override default alignment rules for some of the structures within a program. The user has to specify explicitly the bounds of the datatypes that match these structures.

To achieve this, we add two additional “pseudo-datatypes,” `MPI_LB` and `MPI_UB`, that can be used, respectively, to mark the lower bound or the upper bound of a datatype. These pseudo-datatypes occupy no space ( $\text{extent}(\text{MPI\_LB}) = \text{extent}(\text{MPI\_UB}) = 0$ ). They do not affect the size or count of a datatype, and do not affect the content of a message created with this datatype. However, they do affect the definition of the extent of a datatype and, therefore, affect the outcome of a replication of this datatype by a datatype constructor.

**Example 4.9** Let  $D = (-3, 0, 6)$ ;  $T = (\text{MPI\_LB}, \text{MPI\_INT}, \text{MPI\_UB})$ , and  $B = (1, 1, 1)$ . Then a call to `MPI_TYPE_STRUCT(3, B, D, T, type1)` creates a new datatype that has an extent of 9 (from -3 to 5, 5 included), and contains an integer at displacement 0. This is the datatype defined by the sequence  $\{(\text{lb}, -3), (\text{int}, 0), (\text{ub}, 6)\}$ . If this type is replicated



twice by a call to `MPI_TYPE_CONTIGUOUS(2, type1, type2)` then the newly created type can be described by the sequence  $\{(\text{lb}, -3), (\text{int}, 0), (\text{int}, 9), (\text{ub}, 15)\}$ . (An entry of type `ub` can be deleted if there is another entry of type `ub` with a higher displacement; an entry of type `lb` can be deleted if there is another entry of type `lb` with a lower displacement.)

In general, if

$$\text{Typemap} = \{(type_0, disp_0), \dots, (type_{n-1}, disp_{n-1})\},$$

then the **lower bound** of *Typemap* is defined to be

$$lb(\text{Typemap}) = \begin{cases} \min_j disp_j & \text{if no entry has basic type lb} \\ \min_j \{disp_j \text{ such that } type_j = \text{lb}\} & \text{otherwise} \end{cases}$$

Similarly, the **upper bound** of *Typemap* is defined to be

$$ub(\text{Typemap}) = \begin{cases} \max_j disp_j + \text{sizeof}(type_j) + \epsilon & \text{if no entry has basic type ub} \\ \max_j \{disp_j \text{ such that } type_j = \text{ub}\} & \text{otherwise} \end{cases}$$

Then

$$\text{extent}(\text{Typemap}) = ub(\text{Typemap}) - lb(\text{Typemap})$$

If  $type_i$  requires alignment to a byte address that is a multiple of  $k_i$ , then  $\epsilon$  is the least non-negative increment needed to round  $\text{extent}(\text{Typemap})$  to the next multiple of  $\max_i k_i$ .

In Fortran, whether the alignments  $k_i$  are computed according to the alignments used by the compiler in common blocks, **SEQUENCE** derived types, **BIND(C)** derived types, or derived types that are neither **SEQUENCE** nor **BIND(C)**, is implementation-dependent.

The formal definitions given for the various datatype constructors apply now, with the amended definition of **extent**.

*Rationale.* Before Fortran 2003, `MPI_TYPE_CREATE_STRUCT` could be applied to Fortran common blocks and **SEQUENCE** derived types. With Fortran 2003, this list was extended by **BIND(C)** derived types and MPI implementors have implemented the alignments  $k_i$  differently, i.e., some based on the alignments used in **SEQUENCE** derived types, and others according to **BIND(C)** derived types. (*End of rationale.*)

*Advice to implementors.* In Fortran, it is generally recommended to use **BIND(C)** derived types instead of common blocks or **SEQUENCE** derived types. Therefore it is recommended to calculate the alignments  $k_i$  based on **BIND(C)** derived types. (*End of advice to implementors.*)

*Advice to users.* Structures combining different basic datatypes should be defined so that there will be no gaps based on alignment rules, and if used as an array of structures, then also without such an alignment-gap at the end of the structure. In MPI communication, the content of such gaps would not be communicated into the receiver's buffer. For example, such an alignment-gap may occur between an odd number of `floats` or `REALs` before a `double` or `DOUBLE PRECISION` data. Such gaps may be added explicitly to both the structure and the MPI derived datatype handle because the communication of a contiguous derived datatype may be significantly faster than the communication of one that is non-contiguous due to such alignment-gaps.

Example: Instead of

```

1      TYPE, BIND(C) :: my_data
2          REAL, DIMENSION(3) :: x
3          ! there may be a gap of the size of one REAL
4          ! if the alignment of a DOUBLE PRECISION is
5          ! two times the size of a REAL
6          DOUBLE PRECISION: p
7      END TYPE

```

one should define

```

11     TYPE, BIND(C) :: my_data
12         REAL, DIMENSION(3) :: x
13         REAL :: gap1
14         DOUBLE PRECISION: p
15     END TYPE

```

and also including `gap1` in the matching MPI derived datatype. It is required that all processes in a communication add the same gaps, i.e., defined with the same basic datatype. Both, the original and the modified structures are portable, but may have different performance implications for the communication and memory accesses during computation on systems with different alignment values.

In principle, a compiler may define an additional alignment rule for structures, e.g., to use at least 4 or 8 byte alignment although the content may have a  $max_i k_i$  alignment less than this structure alignment. To keep an application portable, it is therefore recommended to always resize structure derived datatype handles if used in an array of structures, see the Example in Section 16.2.15 on page 675. (*End of advice to users.*)

#### 4.1.7 Extent and Bounds of Datatypes

The following function replaces the three functions `MPI_TYPE_UB`, `MPI_TYPE_LB` and `MPI_TYPE_EXTENT`. It also returns address sized integers, in the Fortran binding. The use of `MPI_TYPE_UB`, `MPI_TYPE_LB` and `MPI_TYPE_EXTENT` is deprecated.

```

37 MPI_TYPE_GET_EXTENT(datatype, lb, extent)

```

IN	datatype	datatype to get information on (handle)
OUT	lb	lower bound of datatype (integer)
OUT	extent	extent of datatype (integer)

```

43 int MPI_Type_get_extent(MPI_Datatype datatype, MPI_Aint *lb,
44                        MPI_Aint *extent)

```

```

46 MPI_Type_get_extent(datatype, lb, extent, ierror) BIND(C)
47     TYPE(MPI_Datatype), INTENT(IN) :: datatype
48     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: lb, extent

```

```

    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_TYPE_GET_EXTENT(DATATYPE, LB, EXTENT, IERROR)
    INTEGER DATATYPE, IERROR
    INTEGER(KIND = MPI_ADDRESS_KIND) LB, EXTENT
{void MPI::Datatype::Get_extent(MPI::Aint& lb, MPI::Aint& extent)
    const(binding deprecated, see Section 15.2) }
```

Returns the lower bound and the extent of `datatype` (as defined in Section 4.1.6 on page 112).

MPI allows one to change the extent of a datatype, using lower bound and upper bound markers (`MPI_LB` and `MPI_UB`). This is useful, as it allows to control the stride of successive datatypes that are replicated by datatype constructors, or are replicated by the `count` argument in a send or receive call. However, the current mechanism for achieving it is painful; also it is restrictive. `MPI_LB` and `MPI_UB` are “sticky”: once present in a datatype, they cannot be overridden (e.g., the upper bound can be moved up, by adding a new `MPI_UB` marker, but cannot be moved down below an existing `MPI_UB` marker). A new type constructor is provided to facilitate these changes. The use of `MPI_LB` and `MPI_UB` is deprecated.

```

MPI_TYPE_CREATE_RESIZED(oldtype, lb, extent, newtype)

IN      oldtype      input datatype (handle)
IN      lb           new lower bound of datatype (integer)
IN      extent       new extent of datatype (integer)
OUT     newtype      output datatype (handle)

int MPI_Type_create_resized(MPI_Datatype oldtype, MPI_Aint lb, MPI_Aint
    extent, MPI_Datatype *newtype)
```

```

MPI_Type_create_resized(oldtype, lb, extent, newtype, ierror) BIND(C)
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: lb, extent
    TYPE(MPI_Datatype), INTENT(IN) :: oldtype
    TYPE(MPI_Datatype), INTENT(OUT) :: newtype
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```

MPI_TYPE_CREATE_RESIZED(OLDTYPE, LB, EXTENT, NEWTYPE, IERROR)
    INTEGER OLDTYPE, NEWTYPE, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) LB, EXTENT
{MPI::Datatype MPI::Datatype::Create_resized(const MPI::Aint lb,
    const MPI::Aint extent) const(binding deprecated, see Section 15.2) }
```

Returns in `newtype` a handle to a new datatype that is identical to `oldtype`, except that the lower bound of this new datatype is set to be `lb`, and its upper bound is set to be `lb + extent`. Any previous `lb` and `ub` markers are erased, and a new pair of lower bound and upper bound markers are put in the positions indicated by the `lb` and `extent` arguments. This affects the behavior of the datatype when used in communication operations, with `count > 1`, and when used in the construction of new derived datatypes.

*Advice to users.* It is strongly recommended that users use these two new functions, rather than the old MPI-1 functions to set and access lower bound, upper bound and extent of datatypes. (*End of advice to users.*)

#### 4.1.8 True Extent of Datatypes

Suppose we implement gather (see also Section 5.5 on page 160) as a spanning tree implemented on top of point-to-point routines. Since the receive buffer is only valid on the root process, one will need to allocate some temporary space for receiving data on intermediate nodes. However, the datatype extent cannot be used as an estimate of the amount of space that needs to be allocated, if the user has modified the extent using the MPI\_UB and MPI\_LB values. A function is provided which returns the true extent of the datatype.

MPI\_TYPE\_GET\_TRUE\_EXTENT(datatype, true\_lb, true\_extent)

IN	datatype	datatype to get information on (handle)
OUT	true_lb	true lower bound of datatype (integer)
OUT	true_extent	true size of datatype (integer)

```
int MPI_Type_get_true_extent(MPI_Datatype datatype, MPI_Aint *true_lb,
                             MPI_Aint *true_extent)
```

```
MPI_Type_get_true_extent(datatype, true_lb, true_extent, ierror) BIND(C)
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: true_lb, true_extent
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TYPE_GET_TRUE_EXTENT(DATATYPE, TRUE_LB, TRUE_EXTENT, IERROR)
  INTEGER DATATYPE, IERROR
  INTEGER(KIND = MPI_ADDRESS_KIND) TRUE_LB, TRUE_EXTENT
```

```
{void MPI::Datatype::Get_true_extent(MPI::Aint& true_lb,
                                     MPI::Aint& true_extent) const(binding deprecated, see Section 15.2) }
```

true\_lb returns the offset of the lowest unit of store which is addressed by the datatype, i.e., the lower bound of the corresponding typemap, ignoring MPI\_LB markers. true\_extent returns the true size of the datatype, i.e., the extent of the corresponding typemap, ignoring MPI\_LB and MPI\_UB markers, and performing no rounding for alignment. If the typemap associated with datatype is

$$Typemap = \{(type_0, disp_0), \dots, (type_{n-1}, disp_{n-1})\}$$

Then

$$true\_lb(Typemap) = \min_j \{disp_j : type_j \neq lb, ub\},$$

$$true\_ub(Typemap) = \max_j \{disp_j + sizeof(type_j) : type_j \neq lb, ub\},$$

and

$$true\_extent(Typemap) = true\_ub(Typemap) - true\_lb(Typemap).$$

(Readers should compare this with the definitions in Section 4.1.6 on page 112 and Section 4.1.7 on page 114, which describe the function `MPI_TYPE_GET_EXTENT`.)

The `true_extent` is the minimum number of bytes of memory necessary to hold a datatype, uncompressed.

#### 4.1.9 Commit and Free

A datatype object has to be **committed** before it can be used in a communication. As an argument in datatype constructors, uncommitted and also committed datatypes can be used. There is no need to commit basic datatypes. They are “pre-committed.”

`MPI_TYPE_COMMIT(datatype)`

INOUT      datatype                      datatype that is committed (handle)

```
int MPI_Type_commit(MPI_Datatype *datatype)
```

```
MPI_Type_commit(datatype, ierror) BIND(C)
    TYPE(MPI_Datatype), INTENT(INOUT) :: datatype
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TYPE_COMMIT(DATATYPE, IERROR)
    INTEGER DATATYPE, IERROR
```

```
{void MPI::Datatype::Commit() (binding deprecated, see Section 15.2) }
```

The commit operation commits the datatype, that is, the formal description of a communication buffer, not the content of that buffer. Thus, after a datatype has been committed, it can be repeatedly reused to communicate the changing content of a buffer or, indeed, the content of different buffers, with different starting addresses.

*Advice to implementors.* The system may “compile” at commit time an internal representation for the datatype that facilitates communication, e.g. change from a compacted representation to a flat representation of the datatype, and select the most convenient transfer mechanism. (*End of advice to implementors.*)

`MPI_TYPE_COMMIT` will accept a committed datatype; in this case, it is equivalent to a no-op.

**Example 4.10** The following code fragment gives examples of using `MPI_TYPE_COMMIT`.

```
INTEGER type1, type2
CALL MPI_TYPE_CONTIGUOUS(5, MPI_REAL, type1, ierr)
    ! new type object created
CALL MPI_TYPE_COMMIT(type1, ierr)
    ! now type1 can be used for communication
type2 = type1
    ! type2 can be used for communication
    ! (it is a handle to same object as type1)
CALL MPI_TYPE_VECTOR(3, 5, 4, MPI_REAL, type1, ierr)
    ! new uncommitted type object created
```

```

1 CALL MPI_TYPE_COMMIT(type1, ierr)
2     ! now type1 can be used anew for communication
3
4

```

```

5 MPI_TYPE_FREE(datatype)
6
7     INOUT    datatype                datatype that is freed (handle)
8

```

```

9 int MPI_Type_free(MPI_Datatype *datatype)

```

```

10 MPI_Type_free(datatype, ierror) BIND(C)
11     TYPE(MPI_Datatype), INTENT(INOUT) :: datatype
12     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
13

```

```

14 MPI_TYPE_FREE(DATATYPE, IERROR)
15     INTEGER DATATYPE, IERROR
16

```

```

17 {void MPI::Datatype::Free() (binding deprecated, see Section 15.2) }

```

Marks the datatype object associated with **datatype** for deallocation and sets **datatype** to `MPI_DATATYPE_NULL`. Any communication that is currently using this datatype will complete normally. Freeing a datatype does not affect any other datatype that was built from the freed datatype. The system behaves as if input datatype arguments to derived datatype constructors are passed by value.

*Advice to implementors.* The implementation may keep a reference count of active communications that use the datatype, in order to decide when to free it. Also, one may implement constructors of derived datatypes so that they keep pointers to their datatype arguments, rather than copying them. In this case, one needs to keep track of active datatype definition references in order to know when a datatype object can be freed. (*End of advice to implementors.*)

#### 4.1.10 Duplicating a Datatype

```

34 MPI_TYPE_DUP(oldtype, newtype)

```

```

35     IN      [ticket252-W.]oldtype      datatype (handle)
36
37     OUT     newtype                    copy of oldtype (handle)
38

```

```

39 int MPI_Type_dup(MPI_Datatype oldtype, MPI_Datatype *newtype)

```

```

40 MPI_Type_dup(oldtype, newtype, ierror) BIND(C)
41     TYPE(MPI_Datatype), INTENT(IN) :: oldtype
42     TYPE(MPI_Datatype), INTENT(OUT) :: newtype
43     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
44

```

```

45 MPI_TYPE_DUP(OLDTYPE, NEWTYPE, IERROR)
46     INTEGER OLDTYPE, NEWTYPE, IERROR
47

```

```

48 {MPI::Datatype MPI::Datatype::Dup() const (binding deprecated, see Section 15.2) }

```

`MPI_TYPE_DUP` is a type constructor which duplicates the existing type with associated key values. For each key value, the respective copy callback function determines the attribute value associated with this key in the new communicator; one particular action that a copy callback may take is to delete the attribute from the new datatype. Returns in `newtype` a new datatype with exactly the same properties as `oldtype` and any copied cached information, see Section 6.7.4 on page 288. The new datatype has identical upper bound and lower bound and yields the same net result when fully decoded with the functions in Section 4.1.13. The `newtype` has the same committed state as the old `oldtype`.

#### 4.1.11 Use of General Datatypes in Communication

Handles to derived datatypes can be passed to a communication call wherever a datatype argument is required. A call of the form `MPI_SEND(buf, count, datatype, ...)`, where `count > 1`, is interpreted as if the call was passed a new datatype which is the concatenation of `count` copies of `datatype`. Thus, `MPI_SEND(buf, count, datatype, dest, tag, comm)` is equivalent to,

```
MPI_TYPE_CONTIGUOUS(count, datatype, newtype)
MPI_TYPE_COMMIT(newtype)
MPI_SEND(buf, 1, newtype, dest, tag, comm).
```

Similar statements apply to all other communication functions that have a `count` and `datatype` argument.

Suppose that a send operation `MPI_SEND(buf, count, datatype, dest, tag, comm)` is executed, where `datatype` has type map,

$$\{(type_0, disp_0), \dots, (type_{n-1}, disp_{n-1})\},$$

and extent `extent`. (Empty entries of “pseudo-type” `MPI_UB` and `MPI_LB` are not listed in the type map, but they affect the value of `extent`.) The send operation sends  $n \cdot \text{count}$  entries, where entry  $i \cdot n + j$  is at location  $addr_{i,j} = \text{buf} + \text{extent} \cdot i + disp_j$  and has type `type_j`, for  $i = 0, \dots, \text{count} - 1$  and  $j = 0, \dots, n - 1$ . These entries need not be contiguous, nor distinct; their order can be arbitrary.

The variable stored at address  $addr_{i,j}$  in the calling program should be of a type that matches `type_j`, where type matching is defined as in Section 3.3.1. The message sent contains  $n \cdot \text{count}$  entries, where entry  $i \cdot n + j$  has type `type_j`.

Similarly, suppose that a receive operation `MPI_RECV(buf, count, datatype, source, tag, comm, status)` is executed, where `datatype` has type map,

$$\{(type_0, disp_0), \dots, (type_{n-1}, disp_{n-1})\},$$

with extent `extent`. (Again, empty entries of “pseudo-type” `MPI_UB` and `MPI_LB` are not listed in the type map, but they affect the value of `extent`.) This receive operation receives  $n \cdot \text{count}$  entries, where entry  $i \cdot n + j$  is at location  $\text{buf} + \text{extent} \cdot i + disp_j$  and has type `type_j`. If the incoming message consists of  $k$  elements, then we must have  $k \leq n \cdot \text{count}$ ; the  $i \cdot n + j$ -th element of the message should have a type that matches `type_j`.

Type matching is defined according to the type signature of the corresponding datatypes, that is, the sequence of basic type components. Type matching does not depend on some aspects of the datatype definition, such as the displacements (layout in memory) or the intermediate types used.

**Example 4.11** This example shows that type matching is defined in terms of the basic types that a derived type consists of.

```

...
CALL MPI_TYPE_CONTIGUOUS( 2, MPI_REAL, type2, ...)
CALL MPI_TYPE_CONTIGUOUS( 4, MPI_REAL, type4, ...)
CALL MPI_TYPE_CONTIGUOUS( 2, type2, type22, ...)
...
CALL MPI_SEND( a, 4, MPI_REAL, ...)
CALL MPI_SEND( a, 2, type2, ...)
CALL MPI_SEND( a, 1, type22, ...)
CALL MPI_SEND( a, 1, type4, ...)
...
CALL MPI_RECV( a, 4, MPI_REAL, ...)
CALL MPI_RECV( a, 2, type2, ...)
CALL MPI_RECV( a, 1, type22, ...)
CALL MPI_RECV( a, 1, type4, ...)

```

Each of the sends matches any of the receives.

A datatype may specify overlapping entries. The use of such a datatype in a receive operation is erroneous. (This is erroneous even if the actual message received is short enough not to write any entry more than once.)

Suppose that `MPI_RECV(buf, count, datatype, dest, tag, comm, status)` is executed, where `datatype` has type map,

$$\{(type_0, disp_0), \dots, (type_{n-1}, disp_{n-1})\}.$$

The received message need not fill all the receive buffer, nor does it need to fill a number of locations which is a multiple of  $n$ . Any number,  $k$ , of basic elements can be received, where  $0 \leq k \leq \text{count} \cdot n$ . The number of basic elements received can be retrieved from `status` using the query function `MPI_GET_ELEMENTS`.

`MPI_GET_ELEMENTS( status, datatype, count)`

IN	status	return status of receive operation (Status)
IN	datatype	datatype used by receive operation (handle)
OUT	count	number of received basic elements (integer)

`int MPI_Get_elements(MPI_Status *status, MPI_Datatype datatype, int *count)`

```

MPI_Get_elements(status, datatype, count, ierror) BIND(C)
  TYPE(MPI_Status), INTENT(IN) :: status
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  INTEGER, INTENT(OUT) :: count
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_GET_ELEMENTS(STATUS, DATATYPE, COUNT, IERROR)
  INTEGER STATUS(MPI_STATUS_SIZE), DATATYPE, COUNT, IERROR

```



```
{int MPI::Status::Get_elements(const MPI::Datatype& datatype) const(binding
    deprecated, see Section 15.2) }
```

The previously defined function, `MPI_GET_COUNT` (Section 3.2.5), has a different behavior. It returns the number of “top-level entries” received, i.e. the number of “copies” of type `datatype`. In the previous example, `MPI_GET_COUNT` may return any integer value  $k$ , where  $0 \leq k \leq \text{count}$ . If `MPI_GET_COUNT` returns  $k$ , then the number of basic elements received (and the value returned by `MPI_GET_ELEMENTS`) is  $n \cdot k$ . If the number of basic elements received is not a multiple of  $n$ , that is, if the receive operation has not received an integral number of `datatype` “copies,” then `MPI_GET_COUNT` returns the value `MPI_UNDEFINED`. The `datatype` argument should match the argument provided by the receive call that set the status variable.

**Example 4.12** Usage of `MPI_GET_COUNT` and `MPI_GET_ELEMENTS`.

```
...
CALL MPI_TYPE_CONTIGUOUS(2, MPI_REAL, Type2, ierr)
CALL MPI_TYPE_COMMIT(Type2, ierr)
...
CALL MPI_COMM_RANK(comm, rank, ierr)
IF (rank.EQ.0) THEN
    CALL MPI_SEND(a, 2, MPI_REAL, 1, 0, comm, ierr)
    CALL MPI_SEND(a, 3, MPI_REAL, 1, 0, comm, ierr)
ELSE IF (rank.EQ.1) THEN
    CALL MPI_RECV(a, 2, Type2, 0, 0, comm, stat, ierr)
    CALL MPI_GET_COUNT(stat, Type2, i, ierr)      ! returns i=1
    CALL MPI_GET_ELEMENTS(stat, Type2, i, ierr)  ! returns i=2
    CALL MPI_RECV(a, 2, Type2, 0, 0, comm, stat, ierr)
    CALL MPI_GET_COUNT(stat, Type2, i, ierr)      ! returns i=MPI_UNDEFINED
    CALL MPI_GET_ELEMENTS(stat, Type2, i, ierr)  ! returns i=3
END IF
```

The function `MPI_GET_ELEMENTS` can also be used after a probe to find the number of elements in the probed message. Note that the two functions `MPI_GET_COUNT` and `MPI_GET_ELEMENTS` return the same values when they are used with basic datatypes.

*Rationale.* The extension given to the definition of `MPI_GET_COUNT` seems natural: one would expect this function to return the value of the `count` argument, when the receive buffer is filled. Sometimes `datatype` represents a basic unit of data one wants to transfer, for example, a record in an array of records (structures). One should be able to find out how many components were received without bothering to divide by the number of elements in each component. However, on other occasions, `datatype` is used to define a complex layout of data in the receiver memory, and does not represent a basic unit of data for transfers. In such cases, one needs to use the function `MPI_GET_ELEMENTS`. (*End of rationale.*)

*Advice to implementors.* The definition implies that a receive cannot change the value of storage outside the entries defined to compose the communication buffer. In particular, the definition implies that padding space in a structure should not be modified when such a structure is copied from one process to another. This would prevent

the obvious optimization of copying the structure, together with the padding, as one contiguous block. The implementation is free to do this optimization when it does not impact the outcome of the computation. The user can “force” this optimization by explicitly including padding as part of the message. (*End of advice to implementors.*)

#### 4.1.12 Correct Use of Addresses

Successively declared variables in C or Fortran are not necessarily stored at contiguous locations. Thus, care must be exercised that displacements do not cross from one variable to another. Also, in machines with a segmented address space, addresses are not unique and address arithmetic has some peculiar properties. Thus, the use of **addresses**, that is, displacements relative to the start address `MPI_BOTTOM`, has to be restricted.

Variables belong to the same **sequential storage** if they belong to the same array, to the same **COMMON** block in Fortran, or to the same structure in C. Valid addresses are defined recursively as follows:

1. The function `MPI_GET_ADDRESS` returns a valid address, when passed as argument a variable of the calling program.
2. The `buf` argument of a communication function evaluates to a valid address, when passed as argument a variable of the calling program.
3. If `v` is a valid address, and `i` is an integer, then `v+i` is a valid address, provided `v` and `v+i` are in the same sequential storage.
4. If `v` is a valid address then `MPI_BOTTOM + v` is a valid address.

A correct program uses only valid addresses to identify the locations of entries in communication buffers. Furthermore, if `u` and `v` are two valid addresses, then the (integer) difference `u - v` can be computed only if both `u` and `v` are in the same sequential storage. No other arithmetic operations can be meaningfully executed on addresses.

The rules above impose no constraints on the use of derived datatypes, as long as they are used to define a communication buffer that is wholly contained within the same sequential storage. However, the construction of a communication buffer that contains variables that are not within the same sequential storage must obey certain restrictions. Basically, a communication buffer with variables that are not within the same sequential storage can be used only by specifying in the communication call `buf = MPI_BOTTOM`, `count = 1`, and using a **datatype** argument where all displacements are valid (absolute) addresses.

*Advice to users.* It is not expected that MPI implementations will be able to detect erroneous, “out of bound” displacements — unless those overflow the user address space — since the MPI call may not know the extent of the arrays and records in the host program. (*End of advice to users.*)

*Advice to implementors.* There is no need to distinguish (absolute) addresses and (relative) displacements on a machine with contiguous address space: `MPI_BOTTOM` is zero, and both addresses and displacements are integers. On machines where the distinction is required, addresses are recognized as expressions that involve `MPI_BOTTOM`. (*End of advice to implementors.*)

## 4.1.13 Decoding a Datatype

MPI datatype objects allow users to specify an arbitrary layout of data in memory. There are several cases where accessing the layout information in opaque datatype objects would be useful. The opaque datatype object has found a number of uses outside MPI. Furthermore, a number of tools wish to display internal information about a datatype. To achieve this, datatype decoding functions are provided. The two functions in this section are used together to decode datatypes to recreate the calling sequence used in their initial definition. These can be used to allow a user to determine the type map and type signature of a datatype.

`MPI_TYPE_GET_ENVELOPE(datatype, num_integers, num_addresses, num_datatypes, combiner)`

IN	<code>datatype</code>	datatype to access (handle)
OUT	<code>num_integers</code>	number of input integers used in the call constructing combiner (non-negative integer)
OUT	<code>num_addresses</code>	number of input addresses used in the call constructing combiner (non-negative integer)
OUT	<code>num_datatypes</code>	number of input datatypes used in the call constructing combiner (non-negative integer)
OUT	<code>combiner</code>	combiner (state)

```
int MPI_Type_get_envelope(MPI_Datatype datatype, int *num_integers,
                          int *num_addresses, int *num_datatypes, int *combiner)
```

```
MPI_Type_get_envelope(datatype, num_integers, num_addresses, num_datatypes,
                      combiner, ierror) BIND(C)
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    INTEGER, INTENT(OUT) :: num_integers, num_addresses, num_datatypes,
    combiner
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TYPE_GET_ENVELOPE(DATATYPE, NUM_INTEGERS, NUM_ADDRESSES, NUM_DATATYPES,
                      COMBINER, IERROR)
    INTEGER DATATYPE, NUM_INTEGERS, NUM_ADDRESSES, NUM_DATATYPES, COMBINER,
    IERROR
```

```
{void MPI::Datatype::Get_envelope(int& num_integers, int& num_addresses,
    int& num_datatypes, int& combiner) const(binding deprecated, see
    Section 15.2) }
```

For the given datatype, `MPI_TYPE_GET_ENVELOPE` returns information on the number and type of input arguments used in the call that created the `datatype`. The number-of-arguments values returned can be used to provide sufficiently large arrays in the decoding routine `MPI_TYPE_GET_CONTENTS`. This call and the meaning of the returned values is described below. The combiner reflects the MPI datatype constructor call that was used in creating datatype.

*Rationale.* By requiring that the combiner reflect the constructor used in the creation of the datatype, the decoded information can be used to effectively recreate the calling sequence used in the original creation. One call is effectively the same as another when the information obtained from `MPI_TYPE_GET_CONTENTS` may be used with either to produce the same outcome. C calls `MPI_Type_hindexed` and `MPI_Type_create_hindexed` are always effectively the same while the Fortran call `MPI_TYPE_HINDEXED` will be different than either of these in some MPI implementations. This is the most useful information and was felt to be reasonable even though it constrains implementations to remember the original constructor sequence even if the internal representation is different.

The decoded information keeps track of datatype duplications. This is important as one needs to distinguish between a predefined datatype and a dup of a predefined datatype. The former is a constant object that cannot be freed, while the latter is a derived datatype that can be freed. (*End of rationale.*)

The list below has the values that can be returned in combiner on the left and the call associated with them on the right.

<code>MPI_COMBINER_NAMED</code>	a named predefined datatype
<code>MPI_COMBINER_DUP</code>	<code>MPI_TYPE_DUP</code>
<code>MPI_COMBINER_CONTIGUOUS</code>	<code>MPI_TYPE_CONTIGUOUS</code>
<code>MPI_COMBINER_VECTOR</code>	<code>MPI_TYPE_VECTOR</code>
<code>MPI_COMBINER_HVECTOR_INTEGER</code>	<code>MPI_TYPE_HVECTOR</code> from Fortran
<code>MPI_COMBINER_HVECTOR</code>	<code>MPI_TYPE_HVECTOR</code> from C or C++ and in some case Fortran or <code>MPI_TYPE_CREATE_HVECTOR</code>
<code>MPI_COMBINER_INDEXED</code>	<code>MPI_TYPE_INDEXED</code>
<code>MPI_COMBINER_HINDEXED_INTEGER</code>	<code>MPI_TYPE_HINDEXED</code> from Fortran
<code>MPI_COMBINER_HINDEXED</code>	<code>MPI_TYPE_HINDEXED</code> from C or C++ and in some case Fortran or <code>MPI_TYPE_CREATE_HINDEXED</code>
<code>MPI_COMBINER_INDEXED_BLOCK</code>	<code>MPI_TYPE_CREATE_INDEXED_BLOCK</code>
<code>MPI_COMBINER_STRUCT_INTEGER</code>	<code>MPI_TYPE_STRUCT</code> from Fortran
<code>MPI_COMBINER_STRUCT</code>	<code>MPI_TYPE_STRUCT</code> from C or C++ and in some case Fortran or <code>MPI_TYPE_CREATE_STRUCT</code>
<code>MPI_COMBINER_SUBARRAY</code>	<code>MPI_TYPE_CREATE_SUBARRAY</code>
<code>MPI_COMBINER_DARRAY</code>	<code>MPI_TYPE_CREATE_DARRAY</code>
<code>MPI_COMBINER_F90_REAL</code>	<code>MPI_TYPE_CREATE_F90_REAL</code>
<code>MPI_COMBINER_F90_COMPLEX</code>	<code>MPI_TYPE_CREATE_F90_COMPLEX</code>
<code>MPI_COMBINER_F90_INTEGER</code>	<code>MPI_TYPE_CREATE_F90_INTEGER</code>
<code>MPI_COMBINER_RESIZED</code>	<code>MPI_TYPE_CREATE_RESIZED</code>

Table 4.1: combiner values returned from `MPI_TYPE_GET_ENVELOPE`

If combiner is `MPI_COMBINER_NAMED` then datatype is a named predefined datatype.

For deprecated calls with address arguments, we sometimes need to differentiate whether the call used an integer or an address size argument. For example, there are two combin-  
ers for hvector: `MPI_COMBINER_HVECTOR_INTEGER` and `MPI_COMBINER_HVECTOR`. The  
former is used if it was the MPI-1 call from Fortran, and the latter is used if it was the  
MPI-1 call from C or C++. However, on systems where `MPI_ADDRESS_KIND =`  
`MPI_INTEGER_KIND` (i.e., where integer arguments and address size arguments are the same),  
the combiner `MPI_COMBINER_HVECTOR` may be returned for a datatype constructed by a  
call to `MPI_TYPE_HVECTOR` from Fortran. Similarly, `MPI_COMBINER_HINDEXED` may  
be returned for a datatype constructed by a call to `MPI_TYPE_HINDEXED` from Fortran,  
and `MPI_COMBINER_STRUCT` may be returned for a datatype constructed by a call to  
`MPI_TYPE_STRUCT` from Fortran. On such systems, one need not differentiate construc-  
tors that take address size arguments from constructors that take integer arguments, since  
these are the same. The preferred calls all use address sized arguments so two combin-  
ers are not required for them.

*Rationale.* For recreating the original call, it is important to know if address informa-  
tion may have been truncated. The deprecated calls from Fortran for a few routines  
could be subject to truncation in the case where the default `INTEGER` size is smaller  
than the size of an address. (*End of rationale.*)

The actual arguments used in the creation call for a datatype can be obtained from the  
call:

`MPI_TYPE_GET_CONTENTS(datatype, max_integers, max_addresses, max_datatypes, ar-  
ray_of_integers, array_of_addresses, array_of_datatypes)`

IN	datatype	datatype to access (handle)
IN	max_integers	number of elements in array_of_integers (non-negative integer)
IN	max_addresses	number of elements in array_of_addresses (non-negative integer)
IN	max_datatypes	number of elements in array_of_datatypes (non-negative integer)
OUT	array_of_integers	contains integer arguments used in constructing datatype (array of integers)
OUT	array_of_addresses	contains address arguments used in constructing datatype (array of integers)
OUT	array_of_datatypes	contains datatype arguments used in constructing datatype (array of handles)

```
int MPI_Type_get_contents(MPI_Datatype datatype, int max_integers,
    int max_addresses, int max_datatypes, int array_of_integers[],
    MPI_Aint array_of_addresses[],
    MPI_Datatype array_of_datatypes[])
```

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```

1 MPI_Type_get_contents(datatype, max_integers, max_addresses, max_datatypes,
2     array_of_integers, array_of_addresses, array_of_datatypes,
3     ierror) BIND(C)
4     TYPE(MPI_Datatype), INTENT(IN) :: datatype
5     INTEGER, INTENT(IN) :: max_integers, max_addresses, max_datatypes
6     INTEGER, INTENT(OUT) :: array_of_integers(max_integers)
7     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) ::
8     array_of_addresses(max_addresses)
9     TYPE(MPI_Datatype), INTENT(OUT) :: array_of_datatypes(max_datatypes)
10    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
11
12 MPI_TYPE_GET_CONTENTS(DATATYPE, MAX_INTEGERS, MAX_ADDRESSES, MAX_DATATYPES,
13     ARRAY_OF_INTEGERS, ARRAY_OF_ADDRESSES, ARRAY_OF_DATATYPES,
14     IERROR)
15     INTEGER DATATYPE, MAX_INTEGERS, MAX_ADDRESSES, MAX_DATATYPES,
16     ARRAY_OF_INTEGERS(*), ARRAY_OF_DATATYPES(*), IERROR
17     INTEGER(KIND=MPI_ADDRESS_KIND) ARRAY_OF_ADDRESSES(*)
18 {void MPI::Datatype::Get_contents(int max_integers, int max_addresses,
19     int max_datatypes, int array_of_integers[],
20     MPI::Aint array_of_addresses[],
21     MPI::Datatype array_of_datatypes[]) const(binding deprecated, see
22     Section 15.2) }
23

```

`datatype` must be a predefined unnamed or a derived datatype; the call is erroneous if `datatype` is a predefined named datatype.

The values given for `max_integers`, `max_addresses`, and `max_datatypes` must be at least as large as the value returned in `num_integers`, `num_addresses`, and `num_datatypes`, respectively, in the call `MPI_TYPE_GET_ENVELOPE` for the same `datatype` argument.

*Rationale.* The arguments `max_integers`, `max_addresses`, and `max_datatypes` allow for error checking in the call. (*End of rationale.*)

The datatypes returned in `array_of_datatypes` are handles to datatype objects that are equivalent to the datatypes used in the original construction call. If these were derived datatypes, then the returned datatypes are new datatype objects, and the user is responsible for freeing these datatypes with `MPI_TYPE_FREE`. If these were predefined datatypes, then the returned datatype is equal to that (constant) predefined datatype and cannot be freed.

The committed state of returned derived datatypes is undefined, i.e., the datatypes may or may not be committed. Furthermore, the content of attributes of returned datatypes is undefined.

Note that `MPI_TYPE_GET_CONTENTS` can be invoked with a `datatype` argument that was constructed using `MPI_TYPE_CREATE_F90_REAL`, `MPI_TYPE_CREATE_F90_INTEGER`, or `MPI_TYPE_CREATE_F90_COMPLEX` (an unnamed predefined datatype). In such a case, an empty `array_of_datatypes` is returned.

*Rationale.* The definition of datatype equivalence implies that equivalent predefined datatypes are equal. By requiring the same handle for named predefined datatypes, it is possible to use the `==` or `.EQ.` comparison operator to determine the datatype involved. (*End of rationale.*)

*Advice to implementors.* The datatypes returned in `array_of_datatypes` must appear to the user as if each is an equivalent copy of the datatype used in the type constructor call. Whether this is done by creating a new datatype or via another mechanism such as a reference count mechanism is up to the implementation as long as the semantics are preserved. (*End of advice to implementors.*)

*Rationale.* The committed state and attributes of the returned datatype is deliberately left vague. The datatype used in the original construction may have been modified since its use in the constructor call. Attributes can be added, removed, or modified as well as having the datatype committed. The semantics given allow for a reference count implementation without having to track these changes. (*End of rationale.*)

In the deprecated datatype constructor calls, the address arguments in Fortran are of type `INTEGER`. In the preferred calls, the address arguments are of type `INTEGER(KIND=MPI_ADDRESS_KIND)`. The call `MPI_TYPE_GET_CONTENTS` returns all addresses in an argument of type `INTEGER(KIND=MPI_ADDRESS_KIND)`. This is true even if the deprecated calls were used. Thus, the location of values returned can be thought of as being returned by the C bindings. It can also be determined by examining the preferred calls for datatype constructors for the deprecated calls that involve addresses.

*Rationale.* By having all address arguments returned in the `array_of_addresses` argument, the result from a C and Fortran decoding of a `datatype` gives the result in the same argument. It is assumed that an integer of type `INTEGER(KIND=MPI_ADDRESS_KIND)` will be at least as large as the `INTEGER` argument used in datatype construction with the old MPI-1 calls so no loss of information will occur. (*End of rationale.*)

The following defines what values are placed in each entry of the returned arrays depending on the datatype constructor used for `datatype`. It also specifies the size of the arrays needed which is the values returned by `MPI_TYPE_GET_ENVELOPE`. In Fortran, the following calls were made:

```

PARAMETER (LARGE = 1000)
INTEGER TYPE, NI, NA, ND, COMBINER, I(LARGE), D(LARGE), IERROR
INTEGER(KIND=MPI_ADDRESS_KIND) A(LARGE)
! CONSTRUCT DATATYPE TYPE (NOT SHOWN)
CALL MPI_TYPE_GET_ENVELOPE(TYPE, NI, NA, ND, COMBINER, IERROR)
IF ((NI .GT. LARGE) .OR. (NA .GT. LARGE) .OR. (ND .GT. LARGE)) THEN
    WRITE (*, *) "NI, NA, OR ND = ", NI, NA, ND, &
    " RETURNED BY MPI_TYPE_GET_ENVELOPE IS LARGER THAN LARGE = ", LARGE
    CALL MPI_ABORT(MPI_COMM_WORLD, 99, IERROR)
ENDIF
CALL MPI_TYPE_GET_CONTENTS(TYPE, NI, NA, ND, I, A, D, IERROR)

```

or in C the analogous calls of:

```

#define LARGE 1000
int ni, na, nd, combiner, i[LARGE];

```

```

1 MPI_Aint a[LARGE];
2 MPI_Datatype type, d[LARGE];
3 /* construct datatype type (not shown) */
4 MPI_Type_get_envelope(type, &ni, &na, &nd, &combiner);
5 if ((ni > LARGE) || (na > LARGE) || (nd > LARGE)) {
6     fprintf(stderr, "ni, na, or nd = %d %d %d returned by ", ni, na, nd);
7     fprintf(stderr, "MPI_Type_get_envelope is larger than LARGE = %d\n",
8             LARGE);
9     MPI_Abort(MPI_COMM_WORLD, 99);
10 };
11 MPI_Type_get_contents(type, ni, na, nd, i, a, d);
12

```

The C++ code is in analogy to the C code above with the same values returned.

In the descriptions that follow, the lower case name of arguments is used.

If combiner is MPI\_COMBINER\_NAMED then it is erroneous to call

**MPI\_TYPE\_GET\_CONTENTS.**

If combiner is MPI\_COMBINER\_DUP then

Constructor argument	C & C++ location	Fortran location
oldtype	d[0]	D(1)

and ni = 0, na = 0, nd = 1.

If combiner is MPI\_COMBINER\_CONTIGUOUS then

Constructor argument	C & C++ location	Fortran location
count	i[0]	I(1)
oldtype	d[0]	D(1)

and ni = 1, na = 0, nd = 1.

If combiner is MPI\_COMBINER\_VECTOR then

Constructor argument	C & C++ location	Fortran location
count	i[0]	I(1)
blocklength	i[1]	I(2)
stride	i[2]	I(3)
oldtype	d[0]	D(1)

and ni = 3, na = 0, nd = 1.

If combiner is MPI\_COMBINER\_HVECTOR\_INTEGER or MPI\_COMBINER\_HVECTOR then

Constructor argument	C & C++ location	Fortran location
count	i[0]	I(1)
blocklength	i[1]	I(2)
stride	a[0]	A(1)
oldtype	d[0]	D(1)

and ni = 2, na = 1, nd = 1.

If combiner is MPI\_COMBINER\_INDEXED then

Constructor argument	C & C++ location	Fortran location
count	i[0]	I(1)
array_of_blocklengths	i[1] to i[i[0]]	I(2) to I(I(1)+1)
array_of_displacements	i[i[0]+1] to i[2*i[0]]	I(I(1)+2) to I(2*I(1)+1)
oldtype	d[0]	D(1)



and  $ni = 2*count+1$ ,  $na = 0$ ,  $nd = 1$ .

If combiner is `MPI_COMBINER_HINDEXED_INTEGER` or `MPI_COMBINER_HINDEXED` then

Constructor argument	C & C++ location	Fortran location
count	$i[0]$	$I(1)$
array_of_blocklengths	$i[1]$ to $i[i[0]]$	$I(2)$ to $I(I(1)+1)$
array_of_displacements	$a[0]$ to $a[i[0]-1]$	$A(1)$ to $A(I(1))$
oldtype	$d[0]$	$D(1)$

and  $ni = count+1$ ,  $na = count$ ,  $nd = 1$ .

If combiner is `MPI_COMBINER_INDEXED_BLOCK` then

Constructor argument	C & C++ location	Fortran location
count	$i[0]$	$I(1)$
blocklength	$i[1]$	$I(2)$
array_of_displacements	$i[2]$ to $i[i[0]+1]$	$I(3)$ to $I(I(1)+2)$
oldtype	$d[0]$	$D(1)$

and  $ni = count+2$ ,  $na = 0$ ,  $nd = 1$ .

If combiner is `MPI_COMBINER_STRUCT_INTEGER` or `MPI_COMBINER_STRUCT` then

Constructor argument	C & C++ location	Fortran location
count	$i[0]$	$I(1)$
array_of_blocklengths	$i[1]$ to $i[i[0]]$	$I(2)$ to $I(I(1)+1)$
array_of_displacements	$a[0]$ to $a[i[0]-1]$	$A(1)$ to $A(I(1))$
array_of_types	$d[0]$ to $d[i[0]-1]$	$D(1)$ to $D(I(1))$

and  $ni = count+1$ ,  $na = count$ ,  $nd = count$ .

If combiner is `MPI_COMBINER_SUBARRAY` then

Constructor argument	C & C++ location	Fortran location
ndims	$i[0]$	$I(1)$
array_of_sizes	$i[1]$ to $i[i[0]]$	$I(2)$ to $I(I(1)+1)$
array_of_subsizes	$i[i[0]+1]$ to $i[2*i[0]]$	$I(I(1)+2)$ to $I(2*I(1)+1)$
array_of_starts	$i[2*i[0]+1]$ to $i[3*i[0]]$	$I(2*I(1)+2)$ to $I(3*I(1)+1)$
order	$i[3*i[0]+1]$	$I(3*I(1)+2)$
oldtype	$d[0]$	$D(1)$

and  $ni = 3*ndims+2$ ,  $na = 0$ ,  $nd = 1$ .

If combiner is `MPI_COMBINER_DARRAY` then

Constructor argument	C & C++ location	Fortran location
size	$i[0]$	$I(1)$
rank	$i[1]$	$I(2)$
ndims	$i[2]$	$I(3)$
array_of_gsizes	$i[3]$ to $i[i[2]+2]$	$I(4)$ to $I(I(3)+3)$
array_of_distribs	$i[i[2]+3]$ to $i[2*i[2]+2]$	$I(I(3)+4)$ to $I(2*I(3)+3)$
array_of_dargs	$i[2*i[2]+3]$ to $i[3*i[2]+2]$	$I(2*I(3)+4)$ to $I(3*I(3)+3)$
array_of_psize	$i[3*i[2]+3]$ to $i[4*i[2]+2]$	$I(3*I(3)+4)$ to $I(4*I(3)+3)$
order	$i[4*i[2]+3]$	$I(4*I(3)+4)$
oldtype	$d[0]$	$D(1)$

and  $ni = 4*ndims+4$ ,  $na = 0$ ,  $nd = 1$ .

If combiner is `MPI_COMBINER_F90_REAL` then

Constructor argument	C & C++ location	Fortran location
p	i[0]	I(1)
r	i[1]	I(2)

and ni = 2, na = 0, nd = 0.

If combiner is MPI\_COMBINER\_F90\_COMPLEX then

Constructor argument	C & C++ location	Fortran location
p	i[0]	I(1)
r	i[1]	I(2)

and ni = 2, na = 0, nd = 0.

If combiner is MPI\_COMBINER\_F90\_INTEGER then

Constructor argument	C & C++ location	Fortran location
r	i[0]	I(1)

and ni = 1, na = 0, nd = 0.

If combiner is MPI\_COMBINER\_RESIZED then

Constructor argument	C & C++ location	Fortran location
lb	a[0]	A(1)
extent	a[1]	A(2)
oldtype	d[0]	D(1)

and ni = 0, na = 2, nd = 1.

#### 4.1.14 Examples

The following examples illustrate the use of derived datatypes.

**Example 4.13** Send and receive a section of a 3D array.

```

REAL a(100,100,100), e(9,9,9)
INTEGER oneslice, twoslice, threeslice, sizeofreal, myrank, ierr
INTEGER status(MPI_STATUS_SIZE)

C    extract the section a(1:17:2, 3:11, 2:10)
C    and store it in e(:, :, :).

CALL MPI_COMM_RANK(MPI_COMM_WORLD, myrank, ierr)

CALL MPI_TYPE_EXTENT( MPI_REAL, sizeofreal, ierr)

C    create datatype for a 1D section
CALL MPI_TYPE_VECTOR( 9, 1, 2, MPI_REAL, oneslice, ierr)

C    create datatype for a 2D section
CALL MPI_TYPE_HVECTOR(9, 1, 100*sizeofreal, oneslice, twoslice, ierr)

C    create datatype for the entire section
CALL MPI_TYPE_HVECTOR( 9, 1, 100*100*sizeofreal, twoslice,
                      threeslice, ierr)

```

```

CALL MPI_TYPE_COMMIT( threeslice, ierr)
CALL MPI_SENDRECV(a(1,3,2), 1, threeslice, myrank, 0, e, 9*9*9,
                  MPI_REAL, myrank, 0, MPI_COMM_WORLD, status, ierr)

```

**Example 4.14** Copy the (strictly) lower triangular part of a matrix.

```

REAL a(100,100), b(100,100)
INTEGER disp(100), blocklen(100), ltype, myrank, ierr
INTEGER status(MPI_STATUS_SIZE)

C    copy lower triangular part of array a
C    onto lower triangular part of array b

CALL MPI_COMM_RANK(MPI_COMM_WORLD, myrank, ierr)

C    compute start and size of each column
DO i=1, 100
    disp(i) = 100*(i-1) + i
    blocklen(i) = 100-i
END DO

C    create datatype for lower triangular part
CALL MPI_TYPE_INDEXED( 100, blocklen, disp, MPI_REAL, ltype, ierr)

CALL MPI_TYPE_COMMIT(ltype, ierr)
CALL MPI_SENDRECV( a, 1, ltype, myrank, 0, b, 1,
                  ltype, myrank, 0, MPI_COMM_WORLD, status, ierr)

```

**Example 4.15** Transpose a matrix.

```

REAL a(100,100), b(100,100)
INTEGER row, xpose, sizeofreal, myrank, ierr
INTEGER status(MPI_STATUS_SIZE)

C    transpose matrix a onto b

CALL MPI_COMM_RANK(MPI_COMM_WORLD, myrank, ierr)

CALL MPI_TYPE_EXTENT( MPI_REAL, sizeofreal, ierr)

C    create datatype for one row
CALL MPI_TYPE_VECTOR( 100, 1, 100, MPI_REAL, row, ierr)

C    create datatype for matrix in row-major order
CALL MPI_TYPE_HVECTOR( 100, 1, sizeofreal, row, xpose, ierr)

```

```

1      CALL MPI_TYPE_COMMIT( xpose, ierr)
2
3  C    send matrix in row-major order and receive in column major order
4      CALL MPI_SENDRECV( a, 1, xpose, myrank, 0, b, 100*100,
5                          MPI_REAL, myrank, 0, MPI_COMM_WORLD, status, ierr)
6

```

**Example 4.16** Another approach to the transpose problem:

```

9      REAL a(100,100), b(100,100)
10     INTEGER disp(2), blocklen(2), type(2), row, row1, sizeofreal
11     INTEGER myrank, ierr
12     INTEGER status(MPI_STATUS_SIZE)
13
14     CALL MPI_COMM_RANK(MPI_COMM_WORLD, myrank, ierr)
15
16  C    transpose matrix a onto b
17
18     CALL MPI_TYPE_EXTENT( MPI_REAL, sizeofreal, ierr)
19
20  C    create datatype for one row
21     CALL MPI_TYPE_VECTOR( 100, 1, 100, MPI_REAL, row, ierr)
22
23  C    create datatype for one row, with the extent of one real number
24     disp(1) = 0
25     disp(2) = sizeofreal
26     type(1)  = row
27     type(2)  = MPI_UB
28     blocklen(1) = 1
29     blocklen(2) = 1
30     CALL MPI_TYPE_STRUCT( 2, blocklen, disp, type, row1, ierr)
31
32     CALL MPI_TYPE_COMMIT( row1, ierr)
33
34  C    send 100 rows and receive in column major order
35     CALL MPI_SENDRECV( a, 100, row1, myrank, 0, b, 100*100,
36                         MPI_REAL, myrank, 0, MPI_COMM_WORLD, status, ierr)
37
38

```

**Example 4.17** We manipulate an array of structures.

```

41  struct Partstruct
42  {
43      int    class; /* particle class */
44      double d[6];  /* particle coordinates */
45      char   b[7];  /* some additional information */
46  };
47
48  struct Partstruct  particle[1000];

```

```

int          i, dest, rank, tag;
MPI_Comm     comm;

/* build datatype describing structure */

MPI_Datatype Particletype;
MPI_Datatype type[3] = {MPI_INT, MPI_DOUBLE, MPI_CHAR};
int          blocklen[3] = {1, 6, 7};
MPI_Aint     disp[3];
MPI_Aint     base;

/* compute displacements of structure components */

MPI_Address( particle, disp);
MPI_Address( particle[0].d, disp+1);
MPI_Address( particle[0].b, disp+2);
base = disp[0];
for (i=0; i < 3; i++) disp[i] -= base;

MPI_Type_struct( 3, blocklen, disp, type, &Particletype);

/* If compiler does padding in mysterious ways,
   the following may be safer */

MPI_Datatype type1[4] = {MPI_INT, MPI_DOUBLE, MPI_CHAR, MPI_UB};
int          blocklen1[4] = {1, 6, 7, 1};
MPI_Aint     disp1[4];

/* compute displacements of structure components */

MPI_Address( particle, disp1);
MPI_Address( particle[0].d, disp1+1);
MPI_Address( particle[0].b, disp1+2);
MPI_Address( particle+1, disp1+3);
base = disp1[0];
for (i=0; i < 4; i++) disp1[i] -= base;

/* build datatype describing structure */

MPI_Type_struct( 4, blocklen1, disp1, type1, &Particletype);

/* 4.1:
   send the entire array */

```

```

1  MPI_Type_commit( &Particletype);
2  MPI_Send( particle, 1000, Particletype, dest, tag, comm);
3
4
5      /* 4.2:
6      send only the entries of class zero particles,
7      preceded by the number of such entries */
8
9  MPI_Datatype Zparticles; /* datatype describing all particles
10                          with class zero (needs to be recomputed
11                          if classes change) */
12  MPI_Datatype Ztype;
13
14  MPI_Aint      zdisp[1000];
15  int           zblock[1000], j, k;
16  int           zzblock[2] = {1,1};
17  MPI_Aint      zzdisp[2];
18  MPI_Datatype  zztype[2];
19
20  /* compute displacements of class zero particles */
21  j = 0;
22  for(i=0; i < 1000; i++)
23      if (particle[i].class == 0)
24          {
25              zdisp[j] = i;
26              zblock[j] = 1;
27              j++;
28          }
29
30  /* create datatype for class zero particles */
31  MPI_Type_indexed( j, zblock, zdisp, Particletype, &Zparticles);
32
33  /* prepend particle count */
34  MPI_Address(&j, zzdisp);
35  MPI_Address(particle, zzdisp+1);
36  zztype[0] = MPI_INT;
37  zztype[1] = Zparticles;
38  MPI_Type_struct(2, zzblock, zzdisp, zztype, &Ztype);
39
40  MPI_Type_commit( &Ztype);
41  MPI_Send( MPI_BOTTOM, 1, Ztype, dest, tag, comm);
42
43
44      /* A probably more efficient way of defining Zparticles */
45
46  /* consecutive particles with index zero are handled as one block */
47  j=0;
48  for (i=0; i < 1000; i++)

```

```

if (particle[i].index == 0)
{
    for (k=i+1; (k < 1000)&&(particle[k].index == 0) ; k++);
    zdisp[j] = i;
    zblock[j] = k-i;
    j++;
    i = k;
}
MPI_Type_indexed( j, zblock, zdisp, Particletype, &Zparticles);

/* 4.3:
   send the first two coordinates of all entries */

MPI_Datatype Allpairs;      /* datatype for all pairs of coordinates */

MPI_Aint sizeofentry;

MPI_Type_extent( Particletype, &sizeofentry);

/* sizeofentry can also be computed by subtracting the address
   of particle[0] from the address of particle[1] */

MPI_Type_hvector( 1000, 2, sizeofentry, MPI_DOUBLE, &Allpairs);
MPI_Type_commit( &Allpairs);
MPI_Send( particle[0].d, 1, Allpairs, dest, tag, comm);

/* an alternative solution to 4.3 */

MPI_Datatype Onepair;      /* datatype for one pair of coordinates, with
                           the extent of one particle entry */

MPI_Aint disp2[3];
MPI_Datatype type2[3] = {MPI_LB, MPI_DOUBLE, MPI_UB};
int blocklen2[3] = {1, 2, 1};

MPI_Address( particle, disp2);
MPI_Address( particle[0].d, disp2+1);
MPI_Address( particle+1, disp2+2);
base = disp2[0];
for (i=0; i<2; i++) disp2[i] -= base;

MPI_Type_struct( 3, blocklen2, disp2, type2, &Onepair);
MPI_Type_commit( &Onepair);
MPI_Send( particle[0].d, 1000, Onepair, dest, tag, comm);

```

**Example 4.18** The same manipulations as in the previous example, but use absolute

```

1  addresses in datatypes.
2
3  struct Partstruct
4  {
5      int class;
6      double d[6];
7      char b[7];
8  };
9
10 struct Partstruct particle[1000];
11
12      /* build datatype describing first array entry */
13
14 MPI_Datatype Particletype;
15 MPI_Datatype type[3] = {MPI_INT, MPI_DOUBLE, MPI_CHAR};
16 int          block[3] = {1, 6, 7};
17 MPI_Aint      disp[3];
18
19 MPI_Address( particle, disp);
20 MPI_Address( particle[0].d, disp+1);
21 MPI_Address( particle[0].b, disp+2);
22 MPI_Type_struct( 3, block, disp, type, &Particletype);
23
24 /* Particletype describes first array entry -- using absolute
25    addresses */
26
27      /* 5.1:
28         send the entire array */
29
30 MPI_Type_commit( &Particletype);
31 MPI_Send( MPI_BOTTOM, 1000, Particletype, dest, tag, comm);
32
33
34      /* 5.2:
35         send the entries of class zero,
36         preceded by the number of such entries */
37
38 MPI_Datatype Zparticles, Ztype;
39
40 MPI_Aint      zdisp[1000];
41 int           zblock[1000], i, j, k;
42 int           zzblock[2] = {1,1};
43 MPI_Datatype  zztype[2];
44 MPI_Aint      zzdisp[2];
45
46 j=0;
47 for (i=0; i < 1000; i++)
48     if (particle[i].index == 0)

```



```

    {
        for (k=i+1; (k < 1000)&&(particle[k].index == 0) ; k++);
        zdisp[j] = i;
        zblock[j] = k-i;
        j++;
        i = k;
    }
MPI_Type_indexed( j, zblock, zdisp, Particletype, &Zparticles);
/* Zparticles describe particles with class zero, using
   their absolute addresses*/

/* prepend particle count */
MPI_Address(&j, zzdisp);
zzdisp[1] = MPI_BOTTOM;
zztype[0] = MPI_INT;
zztype[1] = Zparticles;
MPI_Type_struct(2, zzbblock, zzdisp, zztype, &Ztype);

MPI_Type_commit( &Ztype);
MPI_Send( MPI_BOTTOM, 1, Ztype, dest, tag, comm);

```

**Example 4.19** Handling of unions.

```

union {
    int    ival;
    float  fval;
} u[1000];

int    utype;

/* All entries of u have identical type; variable
   utype keeps track of their current type */

MPI_Datatype  type[2];
int           blocklen[2] = {1,1};
MPI_Aint      disp[2];
MPI_Datatype  mpi_utype[2];
MPI_Aint      i,j;

/* compute an MPI datatype for each possible union type;
   assume values are left-aligned in union storage. */

MPI_Address( u, &i);
MPI_Address( u+1, &j);
disp[0] = 0; disp[1] = j-i;
type[1] = MPI_UB;

```

```

1
2  type[0] = MPI_INT;
3  MPI_Type_struct(2, blocklen, disp, type, &mpi_utype[0]);
4
5  type[0] = MPI_FLOAT;
6  MPI_Type_struct(2, blocklen, disp, type, &mpi_utype[1]);
7
8  for(i=0; i<2; i++) MPI_Type_commit(&mpi_utype[i]);
9
10 /* actual communication */
11
12 MPI_Send(u, 1000, mpi_utype[utype], dest, tag, comm);
13

```

**Example 4.20** This example shows how a datatype can be decoded. The routine `printdatatype` prints out the elements of the datatype. Note the use of `MPI_Type_free` for datatypes that are not predefined.

```

18 /*
19  Example of decoding a datatype.
20
21  Returns 0 if the datatype is predefined, 1 otherwise
22  */
23 #include <stdio.h>
24 #include <stdlib.h>
25 #include "mpi.h"
26 int printdatatype( MPI_Datatype datatype )
27 {
28     int *array_of_ints;
29     MPI_Aint *array_of_adds;
30     MPI_Datatype *array_of_dtypes;
31     int num_ints, num_adds, num_dtypes, combiner;
32     int i;
33
34     MPI_Type_get_envelope( datatype,
35                           &num_ints, &num_adds, &num_dtypes, &combiner );
36     switch (combiner) {
37     case MPI_COMBINER_NAMED:
38         printf( "Datatype is named:" );
39         /* To print the specific type, we can match against the
40          predefined forms. We can NOT use a switch statement here
41          We could also use MPI_TYPE_GET_NAME if we preferred to use
42          names that the user may have changed.
43          */
44         if (datatype == MPI_INT)    printf( "MPI_INT\n" );
45         else if (datatype == MPI_DOUBLE) printf( "MPI_DOUBLE\n" );
46         ... else test for other types ...
47         return 0;
48         break;

```

```

case MPI_COMBINER_STRUCT:
case MPI_COMBINER_STRUCT_INTEGER:
    printf( "Datatype is struct containing" );
    array_of_ints = (int *)malloc( num_ints * sizeof(int) );
    array_of_adds =
        (MPI_Aint *) malloc( num_adds * sizeof(MPI_Aint) );
    array_of_dtypes = (MPI_Datatype *)
        malloc( num_dtypes * sizeof(MPI_Datatype) );
    MPI_Type_get_contents( datatype, num_ints, num_adds, num_dtypes,
        array_of_ints, array_of_adds, array_of_dtypes );
    printf( " %d datatypes:\n", array_of_ints[0] );
    for (i=0; i<array_of_ints[0]; i++) {
        printf( "blocklength %d, displacement %ld, type:\n",
            array_of_ints[i+1], array_of_adds[i] );
        if (printdatatype( array_of_dtypes[i] )) {
            /* Note that we free the type ONLY if it
               is not predefined */
            MPI_Type_free( &array_of_dtypes[i] );
        }
    }
    free( array_of_ints );
    free( array_of_adds );
    free( array_of_dtypes );
    break;
    ... other combiner values ...
default:
    printf( "Unrecognized combiner type\n" );
}
return 1;
}

```

## 4.2 Pack and Unpack

Some existing communication libraries provide pack/unpack functions for sending noncontiguous data. In these, the user explicitly packs data into a contiguous buffer before sending it, and unpacks it from a contiguous buffer after receiving it. Derived datatypes, which are described in Section 4.1, allow one, in most cases, to avoid explicit packing and unpacking. The user specifies the layout of the data to be sent or received, and the communication library directly accesses a noncontiguous buffer. The pack/unpack routines are provided for compatibility with previous libraries. Also, they provide some functionality that is not otherwise available in MPI. For instance, a message can be received in several parts, where the receive operation done on a later part may depend on the content of a former part. Another use is that outgoing messages may be explicitly buffered in user supplied space, thus overriding the system buffering policy. Finally, the availability of pack and unpack operations facilitates the development of additional communication libraries layered on top of MPI.

```

1 MPI_PACK(inbuf, incount, datatype, outbuf, outsize, position, comm)
2     IN      inbuf          input buffer start (choice)
3
4     IN      incount        number of input data items (non-negative integer)
5
6     IN      datatype       datatype of each input data item (handle)
7
8     OUT     outbuf         output buffer start (choice)
9
10    IN      outsize        output buffer size, in bytes (non-negative integer)
11
12    INOUT   position       current position in buffer, in bytes (integer)
13
14    IN      comm           communicator for packed message (handle)
15
16 int MPI_Pack(void* inbuf, int incount, MPI_Datatype datatype, void *outbuf,
17             int outsize, int *position, MPI_Comm comm)
18
19 MPI_Pack(inbuf, incount, datatype, outbuf, outsize, position, comm, ierror)
20
21     BIND(C)
22
23     TYPE(*), DIMENSION(..), INTENT(IN) :: inbuf
24     TYPE(*), DIMENSION(..) :: outbuf
25     INTEGER, INTENT(IN) :: incount, outsize
26     TYPE(MPI_Datatype), INTENT(IN) :: datatype
27     INTEGER, INTENT(INOUT) :: position
28     TYPE(MPI_Comm), INTENT(IN) :: comm
29     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
30
31 MPI_PACK(INBUF, INCOUNT, DATATYPE, OUTBUF, OUTSIZE, POSITION, COMM, IERROR)
32
33     <type> INBUF(*), OUTBUF(*)
34
35     INTEGER INCOUNT, DATATYPE, OUTSIZE, POSITION, COMM, IERROR
36
37 {void MPI::Datatype::Pack(const void* inbuf, int incount, void *outbuf,
38     int outsize, int& position, const MPI::Comm &comm)
39     const(binding deprecated, see Section 15.2) }
40
41
42
43
44
45
46
47
48

```

Packs the message in the send buffer specified by `inbuf`, `incount`, `datatype` into the buffer space specified by `outbuf` and `outsize`. The input buffer can be any communication buffer allowed in `MPI_SEND`. The output buffer is a contiguous storage area containing `outsize` bytes, starting at the address `outbuf` (length is counted in bytes, not elements, as if it were a communication buffer for a message of type `MPI_PACKED`).

The input value of `position` is the first location in the output buffer to be used for packing. `position` is incremented by the size of the packed message, and the output value of `position` is the first location in the output buffer following the locations occupied by the packed message. The `comm` argument is the communicator that will be subsequently used for sending the packed message.

```
MPI_UNPACK(inbuf, insize, position, outbuf, outcount, datatype, comm)
```

IN	inbuf	input buffer start (choice)
IN	insize	size of input buffer, in bytes (non-negative integer)
INOUT	position	current position in bytes (integer)
OUT	outbuf	output buffer start (choice)
IN	outcount	number of items to be unpacked (integer)
IN	datatype	datatype of each output data item (handle)
IN	comm	communicator for packed message (handle)

```
int MPI_Unpack(void* inbuf, int insize, int *position, void *outbuf,
               int outcount, MPI_Datatype datatype, MPI_Comm comm)
```

```
MPI_Unpack(inbuf, insize, position, outbuf, outcount, datatype, comm,
           ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN) :: inbuf
    TYPE(*), DIMENSION(..) :: outbuf
    INTEGER, INTENT(IN) :: insize, outcount
    INTEGER, INTENT(INOUT) :: position
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_UNPACK(INBUF, INSIZE, POSITION, OUTBUF, OUTCOUNT, DATATYPE, COMM,
           IERROR)
    <type> INBUF(*), OUTBUF(*)
    INTEGER INSIZE, POSITION, OUTCOUNT, DATATYPE, COMM, IERROR
```

```
{void MPI::Datatype::Unpack(const void* inbuf, int insize, void *outbuf,
                           int outcount, int& position, const MPI::Comm& comm)
  const(binding deprecated, see Section 15.2) }
```

Unpacks a message into the receive buffer specified by `outbuf`, `outcount`, `datatype` from the buffer space specified by `inbuf` and `insize`. The output buffer can be any communication buffer allowed in `MPI_RECV`. The input buffer is a contiguous storage area containing `insize` bytes, starting at address `inbuf`. The input value of `position` is the first location in the input buffer occupied by the packed message. `position` is incremented by the size of the packed message, so that the output value of `position` is the first location in the input buffer after the locations occupied by the message that was unpacked. `comm` is the communicator used to receive the packed message.

*Advice to users.* Note the difference between `MPI_RECV` and `MPI_UNPACK`: in `MPI_RECV`, the count argument specifies the maximum number of items that can be received. The actual number of items received is determined by the length of the incoming message. In `MPI_UNPACK`, the count argument specifies the actual number of items that are unpacked; the “size” of the corresponding message is the increment in `position`. The reason for this change is that the “incoming message size” is not predetermined since the user decides how much to unpack; nor is it easy to

determine the “message size” from the number of items to be unpacked. In fact, in a heterogeneous system, this number may not be determined *a priori*. (*End of advice to users.*)

To understand the behavior of pack and unpack, it is convenient to think of the data part of a message as being the sequence obtained by concatenating the successive values sent in that message. The pack operation stores this sequence in the buffer space, as if sending the message to that buffer. The unpack operation retrieves this sequence from buffer space, as if receiving a message from that buffer. (It is helpful to think of internal Fortran files or `sscanf` in C, for a similar function.)

Several messages can be successively packed into one **packing unit**. This is effected by several successive **related** calls to `MPI_PACK`, where the first call provides `position = 0`, and each successive call inputs the value of `position` that was output by the previous call, and the same values for `outbuf`, `outcount` and `comm`. This packing unit now contains the equivalent information that would have been stored in a message by one send call with a send buffer that is the “concatenation” of the individual send buffers.

A packing unit can be sent using type `MPI_PACKED`. Any point to point or collective communication function can be used to move the sequence of bytes that forms the packing unit from one process to another. This packing unit can now be received using any receive operation, with any datatype: the type matching rules are relaxed for messages sent with type `MPI_PACKED`.

A message sent with any type (including `MPI_PACKED`) can be received using the type `MPI_PACKED`. Such a message can then be unpacked by calls to `MPI_UNPACK`.

A packing unit (or a message created by a regular, “typed” send) can be unpacked into several successive messages. This is effected by several successive related calls to `MPI_UNPACK`, where the first call provides `position = 0`, and each successive call inputs the value of `position` that was output by the previous call, and the same values for `inbuf`, `insize` and `comm`.

The concatenation of two packing units is not necessarily a packing unit; nor is a substring of a packing unit necessarily a packing unit. Thus, one cannot concatenate two packing units and then unpack the result as one packing unit; nor can one unpack a substring of a packing unit as a separate packing unit. Each packing unit, that was created by a related sequence of pack calls, or by a regular send, must be unpacked as a unit, by a sequence of related unpack calls.

*Rationale.* The restriction on “atomic” packing and unpacking of packing units allows the implementation to add at the head of packing units additional information, such as a description of the sender architecture (to be used for type conversion, in a heterogeneous environment) (*End of rationale.*)

The following call allows the user to find out how much space is needed to pack a message and, thus, manage space allocation for buffers.

MPI\_PACK\_SIZE(incount, datatype, comm, size)

IN	incount	count argument to packing call (non-negative integer)
IN	datatype	datatype argument to packing call (handle)
IN	comm	communicator argument to packing call (handle)
OUT	size	upper bound on size of packed message, in bytes (non-negative integer)

```
int MPI_Pack_size(int incount, MPI_Datatype datatype, MPI_Comm comm,
                 int *size)
```

```
MPI_Pack_size(incount, datatype, comm, size, ierror) BIND(C)
```

```
    INTEGER, INTENT(IN) :: incount
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, INTENT(OUT) :: size
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_PACK_SIZE(INCOUNT, DATATYPE, COMM, SIZE, IERROR)
```

```
    INTEGER INCOUNT, DATATYPE, COMM, SIZE, IERROR
```

```
{int MPI::Datatype::Pack_size(int incount, const MPI::Comm& comm)
    const(binding deprecated, see Section 15.2) }
```

A call to MPI\_PACK\_SIZE(incount, datatype, comm, size) returns in size an upper bound on the increment in position that is effected by a call to MPI\_PACK(inbuf, incount, datatype, outbuf, outcount, position, comm).

*Rationale.* The call returns an upper bound, rather than an exact bound, since the exact amount of space needed to pack the message may depend on the context (e.g., first message packed in a packing unit may take more space). (*End of rationale.*)

**Example 4.21** An example using MPI\_PACK.

```
int      position, i, j, a[2];
char     buff[1000];

MPI_Comm_rank(MPI_COMM_WORLD, &myrank);
if (myrank == 0)
{
    /* SENDER CODE */

    position = 0;
    MPI_Pack(&i, 1, MPI_INT, buff, 1000, &position, MPI_COMM_WORLD);
    MPI_Pack(&j, 1, MPI_INT, buff, 1000, &position, MPI_COMM_WORLD);
    MPI_Send( buff, position, MPI_PACKED, 1, 0, MPI_COMM_WORLD);
}
else /* RECEIVER CODE */
    MPI_Recv( a, 2, MPI_INT, 0, 0, MPI_COMM_WORLD);
```

**Example 4.22** An elaborate example.

```

1  int    position, i;
2
3  float a[1000];
4
5  char  buff[1000];
6
7  MPI_Comm_rank(MPI_Comm_world, &myrank);
8  if (myrank == 0)
9  {
10     /* SENDER CODE */
11
12     int len[2];
13     MPI_Aint disp[2];
14     MPI_Datatype type[2], newtype;
15
16     /* build datatype for i followed by a[0]...a[i-1] */
17
18     len[0] = 1;
19     len[1] = i;
20     MPI_Address( &i, disp);
21     MPI_Address( a, disp+1);
22     type[0] = MPI_INT;
23     type[1] = MPI_FLOAT;
24     MPI_Type_struct( 2, len, disp, type, &newtype);
25     MPI_Type_commit( &newtype);
26
27     /* Pack i followed by a[0]...a[i-1]*/
28
29     position = 0;
30     MPI_Pack( MPI_BOTTOM, 1, newtype, buff, 1000, &position, MPI_COMM_WORLD);
31
32     /* Send */
33
34     MPI_Send( buff, position, MPI_PACKED, 1, 0,
35              MPI_COMM_WORLD);
36
37     /* *****
38      One can replace the last three lines with
39      MPI_Send( MPI_BOTTOM, 1, newtype, 1, 0, MPI_COMM_WORLD);
40      ***** */
41 }
42 else if (myrank == 1)
43 {
44     /* RECEIVER CODE */
45
46     MPI_Status status;
47
48     /* Receive */

```



```

MPI_Recv( buff, 1000, MPI_PACKED, 0, 0, MPI_COMM_WORLD, &status);

/* Unpack i */

position = 0;
MPI_Unpack(buff, 1000, &position, &i, 1, MPI_INT, MPI_COMM_WORLD);

/* Unpack a[0]...a[i-1] */
MPI_Unpack(buff, 1000, &position, a, i, MPI_FLOAT, MPI_COMM_WORLD);
}

```

**Example 4.23** Each process sends a count, followed by count characters to the root; the root concatenates all characters into one string.

```

int  count, gsize, counts[64], totalcount, k1, k2, k,
     displs[64], position, concat_pos;
char chr[100], *lbuf, *rbuf, *cbuf;

MPI_Comm_size(comm, &gsize);
MPI_Comm_rank(comm, &myrank);

/* allocate local pack buffer */
MPI_Pack_size(1, MPI_INT, comm, &k1);
MPI_Pack_size(count, MPI_CHAR, comm, &k2);
k = k1+k2;
lbuf = (char *)malloc(k);

/* pack count, followed by count characters */
position = 0;
MPI_Pack(&count, 1, MPI_INT, lbuf, k, &position, comm);
MPI_Pack(chr, count, MPI_CHAR, lbuf, k, &position, comm);

if (myrank != root) {
    /* gather at root sizes of all packed messages */
    MPI_Gather( &position, 1, MPI_INT, NULL, 0,
               MPI_DATATYPE_NULL, root, comm);

    /* gather at root packed messages */
    MPI_Gatherv( lbuf, position, MPI_PACKED, NULL,
                 NULL, NULL, NULL, root, comm);
} else { /* root code */
    /* gather sizes of all packed messages */
    MPI_Gather( &position, 1, MPI_INT, counts, 1,
               MPI_INT, root, comm);

    /* gather all packed messages */

```

```

1     displs[0] = 0;
2     for (i=1; i < gsize; i++)
3         displs[i] = displs[i-1] + counts[i-1];
4     totalcount = displs[gsz-1] + counts[gsz-1];
5     rbuf = (char *)malloc(totalcount);
6     cbuf = (char *)malloc(totalcount);
7     MPI_Gatherv( lbuf, position, MPI_PACKED, rbuf,
8                 counts, displs, MPI_PACKED, root, comm);
9
10    /* unpack all messages and concatenate strings */
11    concat_pos = 0;
12    for (i=0; i < gsize; i++) {
13        position = 0;
14        MPI_Unpack( rbuf+displs[i], totalcount-displs[i],
15                  &position, &count, 1, MPI_INT, comm);
16        MPI_Unpack( rbuf+displs[i], totalcount-displs[i],
17                  &position, cbuf+concat_pos, count, MPI_CHAR, comm);
18        concat_pos += count;
19    }
20    cbuf[concat_pos] = '\0';
21 }

```

### 4.3 Canonical MPI\_PACK and MPI\_UNPACK

These functions read/write data to/from the buffer in the “external32” data format specified in Section 13.5.2, and calculate the size needed for packing. Their first arguments specify the data format, for future extensibility, but currently the only valid value of the `datarep` argument is “external32.”

*Advice to users.* These functions could be used, for example, to send typed data in a portable format from one MPI implementation to another. (*End of advice to users.*)

The buffer will contain exactly the packed data, without headers. `MPI_BYTE` should be used to send and receive data that is packed using `MPI_PACK_EXTERNAL`.

*Rationale.* `MPI_PACK_EXTERNAL` specifies that there is no header on the message and further specifies the exact format of the data. Since `MPI_PACK` may (and is allowed to) use a header, the datatype `MPI_PACKED` cannot be used for data packed with `MPI_PACK_EXTERNAL`. (*End of rationale.*)

MPI\_PACK\_EXTERNAL(datarep, inbuf, incout, datatype, outbuf, outsize, position )

IN	datarep	data representation (string)
IN	inbuf	input buffer start (choice)
IN	incout	number of input data items (integer)
IN	datatype	datatype of each input data item (handle)
OUT	outbuf	output buffer start (choice)
IN	outsize	output buffer size, in bytes (integer)
INOUT	position	current position in buffer, in bytes (integer)

```
int MPI_Pack_external(char *datarep, void *inbuf, int incout,
    MPI_Datatype datatype, void *outbuf, MPI_Aint outsize,
    MPI_Aint *position)
```

```
MPI_Pack_external(datarep, inbuf, incout, datatype, outbuf, outsize,
    position, ierror) BIND(C)
    CHARACTER(LEN=*) INTENT(IN) :: datarep
    TYPE(*), DIMENSION(..) INTENT(IN) :: inbuf
    TYPE(*), DIMENSION(..) :: outbuf
    INTEGER, INTENT(IN) :: incout
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: outsize
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(INOUT) :: position
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_PACK_EXTERNAL(DATAREP, INBUF, INCOUNT, DATATYPE, OUTBUF, OUTSIZE,
    POSITION, IERROR)
    INTEGER INCOUNT, DATATYPE, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) OUTSIZE, POSITION
    CHARACTER*(*) DATAREP
    <type> INBUF(*), OUTBUF(*)
```

```
{void MPI::Datatype::Pack_external(const char* datarep, const void* inbuf,
    int incout, void* outbuf, MPI::Aint outsize,
    MPI::Aint& position) const(binding deprecated, see Section 15.2) }
```

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```

1  MPI_UNPACK_EXTERNAL(datarep, inbuf, insize, position, outbuf, outsize, position )
2
3  IN      datarep      data representation (string)
4  IN      inbuf        input buffer start (choice)
5  IN      insize       input buffer size, in bytes (integer)
6  INOUT   position     current position in buffer, in bytes (integer)
7  OUT     outbuf       output buffer start (choice)
8  IN      outcount     number of output data items (integer)
9  IN      datatype     datatype of output data item (handle)
10
11
12
13  int MPI_Unpack_external(char *datarep, void *inbuf, MPI_Aint insize,
14                        MPI_Aint *position, void *outbuf, int outcount,
15                        MPI_Datatype datatype)
16
17  MPI_Unpack_external(datarep, inbuf, insize, position, outbuf, outcount,
18                      datatype, ierror) BIND(C)
19  CHARACTER(LEN=*) , INTENT(IN) :: datarep
20  TYPE(*), DIMENSION(..), INTENT(IN) :: inbuf
21  TYPE(*), DIMENSION(..) :: outbuf
22  INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: insize
23  INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(INOUT) :: position
24  INTEGER, INTENT(IN) :: outcount
25  TYPE(MPI_Datatype), INTENT(IN) :: datatype
26  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
27
28  MPI_UNPACK_EXTERNAL(DATAREP, INBUF, INSIZE, POSITION, OUTBUF, OUTCOUNT,
29                      DATATYPE, IERROR)
30  INTEGER OUTCOUNT, DATATYPE, IERROR
31  INTEGER(KIND=MPI_ADDRESS_KIND) INSIZE, POSITION
32  CHARACTER*(*) DATAREP
33  <type> INBUF(*), OUTBUF(*)
34
35  {void MPI::Datatype::Unpack_external(const char* datarep,
36                                     const void* inbuf, MPI::Aint insize, MPI::Aint& position,
37                                     void* outbuf, int outcount) const (binding deprecated, see
38                                     Section 15.2) }
39
40  MPI_PACK_EXTERNAL_SIZE( datarep, incount, datatype, size )
41
42  IN      datarep      data representation (string)
43  IN      incount       number of input data items (integer)
44  IN      datatype     datatype of each input data item (handle)
45  OUT     size         output buffer size, in bytes (integer)
46
47  int MPI_Pack_external_size(char *datarep, int incount,
48                          MPI_Datatype datatype, MPI_Aint *size)

```

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```

MPI_Pack_external_size(datarep, incount, datatype, size, ierror) BIND(C)
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    INTEGER, INTENT(IN) :: incount
    CHARACTER(LEN=*), INTENT(IN) :: datarep
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: size
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_PACK_EXTERNAL_SIZE(DATAREP, INCOUNT, DATATYPE, SIZE, IERROR)
    INTEGER INCOUNT, DATATYPE, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) SIZE
    CHARACTER*(*) DATAREP

{MPI::Aint MPI::Datatype::Pack_external_size(const char* datarep,
    int incount) const(binding deprecated, see Section 15.2) }
```

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# Chapter 5

## Collective Communication

### 5.1 Introduction and Overview

Collective communication is defined as communication that involves a group or groups of processes. The functions of this type provided by MPI are the following:

- **MPI\_BARRIER**, **MPI\_IBARRIER**: Barrier synchronization across all members of a group (Section 5.3 and Section 5.12.1).
- **MPI\_BCAST**, **MPI\_IBCAST**: Broadcast from one member to all members of a group (Section 5.4 and Section 5.12.2). This is shown as “broadcast” in Figure 5.1.
- **MPI\_GATHER**, **MPI\_IGATHER**, **MPI\_GATHERV**, **MPI\_IGATHERV**: Gather data from all members of a group to one member (Section 5.5 and Section 5.12.3). This is shown as “gather” in Figure 5.1.
- **MPI\_SCATTER**, **MPI\_ISCATTER**, **MPI\_SCATTERV**, **MPI\_ISCATTERV**: Scatter data from one member to all members of a group (Section 5.6 and Section 5.12.4). This is shown as “scatter” in Figure 5.1.
- **MPI\_ALLGATHER**, **MPI\_IALLGATHER**, **MPI\_ALLGATHERV**, **MPI\_IALLGATHERV**: A variation on Gather where all members of a group receive the result (Section 5.7 and Section 5.12.5). This is shown as “allgather” in Figure 5.1.
- **MPI\_ALLTOALL**, **MPI\_IALLTOALL**, **MPI\_ALLTOALLV**, **MPI\_IALLTOALLV**, **MPI\_ALLTOALLW**, **MPI\_IALLTOALLW**: Scatter/Gather data from all members to all members of a group (also called complete exchange) (Section 5.8 and Section 5.12.6). This is shown as “complete exchange” in Figure 5.1.
- **MPI\_ALLREDUCE**, **MPI\_IALLREDUCE**, **MPI\_REDUCE**, **MPI\_IREDUCE**: Global reduction operations such as sum, max, min, or user-defined functions, where the result is returned to all members of a group (Section 5.9.6 and Section 5.12.8) and a variation where the result is returned to only one member (Section 5.9 and Section 5.12.7).
- **MPI\_REDUCE\_SCATTER\_BLOCK**, **MPI\_IREDUCE\_SCATTER\_BLOCK**, **MPI\_REDUCE\_SCATTER**, **MPI\_IREDUCE\_SCATTER**: A combined reduction and scatter operation (Section 5.10, Section 5.12.9, and Section 5.12.10).

- `MPI_SCAN`, `MPI_ISCAN`, `MPI_EXSCAN`, `MPI_IEXSCAN`: Scan across all members of a group (also called prefix) (Section 5.11, Section 5.11.2, Section 5.12.11, and Section 5.12.12).

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One of the key arguments in a call to a collective routine is a communicator that defines the group or groups of participating processes and provides a context for the operation. This is discussed further in Section 5.2. The syntax and semantics of the collective operations are defined to be consistent with the syntax and semantics of the point-to-point operations. Thus, general datatypes are allowed and must match between sending and receiving processes as specified in Chapter 4. Several collective routines such as broadcast and gather have a single originating or receiving process. Such a process is called the *root*. Some arguments in the collective functions are specified as “significant only at root,” and are ignored for all participants except the root. The reader is referred to Chapter 4 for information concerning communication buffers, general datatypes and type matching rules, and to Chapter 6 for information on how to define groups and create communicators.

The type-matching conditions for the collective operations are more strict than the corresponding conditions between sender and receiver in point-to-point. Namely, for collective operations, the amount of data sent must exactly match the amount of data specified by the receiver. Different type maps (the layout in memory, see Section 4.1) between sender and receiver are still allowed.

Collective [routine calls]operations can (but are not required to) [return]complete as soon as [their]the caller’s participation in the collective communication is [complete]finished. A blocking operation is complete as soon as the call returns. A nonblocking (immediate) call requires a separate completion call (cf. Section 3.7). The completion of a [call]collective operation indicates that the caller is [now] free to modify locations in the communication buffer. It does not indicate that other processes in the group have completed or even started the operation (unless otherwise implied by the description of the operation). [Thus, a collective communication call may, or may not, have the effect of synchronizing all calling processes. This statement excludes, of course, the barrier function]Thus, a collective communication operation may, or may not, have the effect of synchronizing all calling processes. This statement excludes, of course, the barrier operation.

Collective communication calls may use the same communicators as point-to-point communication; MPI guarantees that messages generated on behalf of collective communication calls will not be confused with messages generated by point-to-point communication. The collective operations do not have a message tag argument. A more detailed discussion of correct use of collective routines is found in Section 5.13.

*Rationale.* The equal-data restriction (on type matching) was made so as to avoid the complexity of providing a facility analogous to the status argument of `MPI_RECV` for discovering the amount of data sent. Some of the collective routines would require an array of status values.

The statements about synchronization are made so as to allow a variety of implementations of the collective functions.

[The collective operations do not accept a message tag argument. If future revisions of MPI define nonblocking collective functions, then tags (or a similar mechanism) might need to be added so as to allow the dis-ambiguation of multiple, pending, collective operations.] (*End of rationale.*)



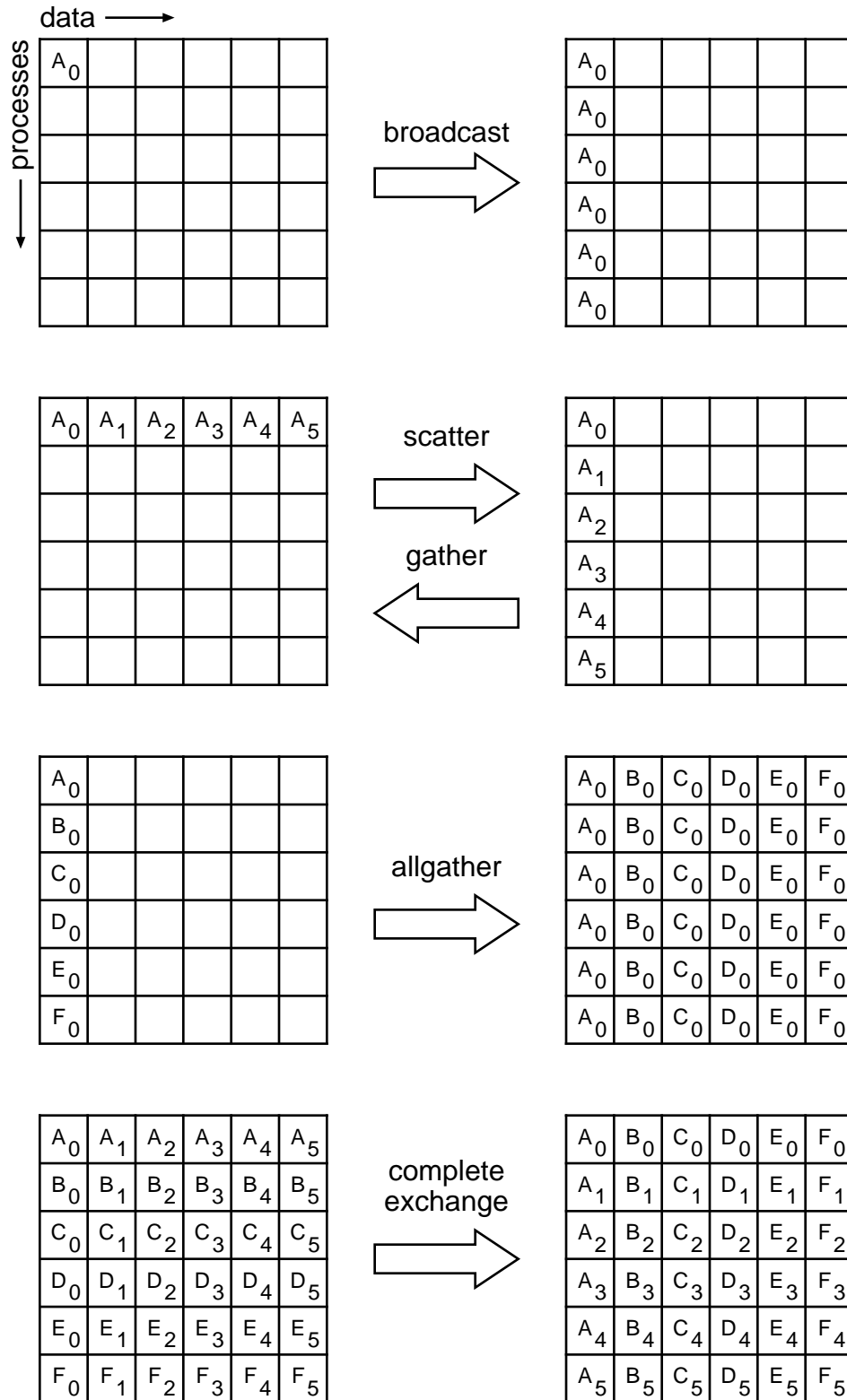


Figure 5.1: Collective move functions illustrated for a group of six processes. In each case, each row of boxes represents data locations in one process. Thus, in the broadcast, initially just the first process contains the data  $A_0$ , but after the broadcast all processes contain it.

*Advice to users.* It is dangerous to rely on synchronization side-effects of the collective operations for program correctness. For example, even though a particular implementation may provide a broadcast routine with a side-effect of synchronization, the standard does not require this, and a program that relies on this will not be portable.

On the other hand, a correct, portable program must allow for the fact that a collective call *may* be synchronizing. Though one cannot rely on any synchronization side-effect, one must program so as to allow it. These issues are discussed further in Section 5.13. (*End of advice to users.*)

*Advice to implementors.* While vendors may write optimized collective routines matched to their architectures, a complete library of the collective communication routines can be written entirely using the MPI point-to-point communication functions and a few auxiliary functions. If implementing on top of point-to-point, a hidden, special communicator might be created for the collective operation so as to avoid interference with any on-going point-to-point communication at the time of the collective call. This is discussed further in Section 5.13. (*End of advice to implementors.*)

Many of the descriptions of the collective routines provide illustrations in terms of blocking MPI point-to-point routines. These are intended solely to indicate what data is sent or received by what process. Many of these examples are *not* correct MPI programs; for purposes of simplicity, they often assume infinite buffering.

## 5.2 Communicator Argument

The key concept of the collective functions is to have a group or groups of participating processes. The routines do not have group identifiers as explicit arguments. Instead, there is a communicator argument. Groups and communicators are discussed in full detail in Chapter 6. For the purposes of this chapter, it is sufficient to know that there are two types of communicators: *intra-communicators* and *inter-communicators*. An intracommunicator can be thought of as an identifier for a single group of processes linked with a context. An intercommunicator identifies two distinct groups of processes linked with a context.

### 5.2.1 Specifics for Intracommunicator Collective Operations

All processes in the group identified by the intracommunicator must call the collective routine.

In many cases, collective communication can occur “in place” for intracommunicators, with the output buffer being identical to the input buffer. This is specified by providing a special argument value, `MPI_IN_PLACE`, instead of the send buffer or the receive buffer argument, depending on the operation performed.

*Rationale.* The “in place” operations are provided to reduce unnecessary memory motion by both the MPI implementation and by the user. Note that while the simple check of testing whether the send and receive buffers have the same address will work for some cases (e.g., `MPI_ALLREDUCE`), they are inadequate in others (e.g., `MPI_GATHER`, with root not equal to zero). Further, Fortran explicitly prohibits aliasing of arguments; the approach of using a special value to denote “in place” operation eliminates that difficulty. (*End of rationale.*)

*Advice to users.* By allowing the “in place” option, the receive buffer in many of the collective calls becomes a send-and-receive buffer. For this reason, a Fortran binding that includes INTENT must mark these as INOUT, not OUT.

Note that MPI\_IN\_PLACE is a special kind of value; it has the same restrictions on its use that MPI\_BOTTOM has. [ Some intracommunicator collective operations do not support the “in place” option (e.g., MPI\_ALLTOALLV).] (*End of advice to users.*)

### 5.2.2 Applying Collective Operations to Intercommunicators

To understand how collective operations apply to intercommunicators, we can view most MPI intracommunicator collective operations as fitting one of the following categories (see, for instance, [56]):

**All-To-All** All processes contribute to the result. All processes receive the result.

- MPI\_ALLGATHER, MPI\_IALLGATHER, MPI\_ALLGATHERV, MPI\_IALLGATHERV
- MPI\_ALLTOALL, MPI\_IALLTOALL, MPI\_ALLTOALLV, MPI\_IALLTOALLV, MPI\_ALLTOALLW, MPI\_IALLTOALLW
- MPI\_ALLREDUCE, MPI\_IALLREDUCE, MPI\_REDUCE\_SCATTER\_BLOCK, MPI\_IREDUCE\_SCATTER\_BLOCK, MPI\_REDUCE\_SCATTER, MPI\_IREDUCE\_SCATTER
- MPI\_BARRIER, MPI\_IBARRIER

**All-To-One** All processes contribute to the result. One process receives the result.

- MPI\_GATHER, MPI\_IGATHER, MPI\_GATHERV, MPI\_IGATHERV
- MPI\_REDUCE, MPI\_IREDUCE

**One-To-All** One process contributes to the result. All processes receive the result.

- MPI\_BCAST, MPI\_IBCAST
- MPI\_SCATTER, MPI\_ISCATTER, MPI\_SCATTERV, MPI\_ISCATTERV

**Other** Collective operations that do not fit into one of the above categories.

- MPI\_SCAN, MPI\_ISCAN, MPI\_EXSCAN, MPI\_IEXSCAN

The data movement patterns of MPI\_SCAN, MPI\_ISCAN [and], MPI\_EXSCAN, and MPI\_IEXSCAN do not fit this taxonomy.

The application of collective communication to intercommunicators is best described in terms of two groups. For example, an all-to-all MPI\_ALLGATHER operation can be described as collecting data from all members of one group with the result appearing in all members of the other group (see Figure 5.2). As another example, a one-to-all MPI\_BCAST operation sends data from one member of one group to all members of the other group. Collective computation operations such as MPI\_REDUCE\_SCATTER have a similar interpretation (see Figure 5.3). For intracommunicators, these two groups are the same. For intercommunicators, these two groups are distinct. For the all-to-all operations, each such operation is described in two phases, so that it has a symmetric, full-duplex behavior.

The following collective operations also apply to intercommunicators:

- MPI\_BARRIER, MPI\_IBARRIER
- MPI\_BCAST, MPI\_IBCAST
- MPI\_GATHER, MPI\_IGATHER, MPI\_GATHERV, MPI\_IGATHERV,
- MPI\_SCATTER, MPI\_ISCATTER, MPI\_SCATTERV, MPI\_ISCATTERV,
- MPI\_ALLGATHER, MPI\_IALLGATHER, MPI\_ALLGATHERV, MPI\_IALLGATHERV,
- MPI\_ALLTOALL, MPI\_IALLTOALL, MPI\_ALLTOALLV, MPI\_IALLTOALLV,  
MPI\_ALLTOALLW, MPI\_IALLTOALLW,
- MPI\_ALLREDUCE, MPI\_IALLREDUCE, MPI\_REDUCE, MPI\_IREDUCE,
- MPI\_REDUCE\_SCATTER\_BLOCK, MPI\_IREDUCE\_SCATTER\_BLOCK,  
MPI\_REDUCE\_SCATTER, MPI\_IREDUCE\_SCATTER.

In C++, the bindings for these functions are in the `MPI::Comm` class. However, since the collective operations do not make sense on a C++ `MPI::Comm` (as it is neither an intercommunicator nor an intracommunicator), the functions are all pure virtual.

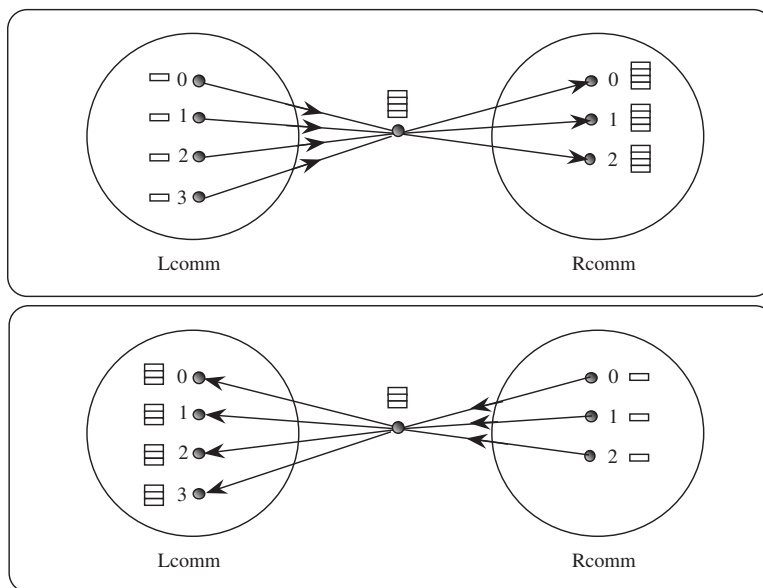


Figure 5.2: Intercommunicator allgather. The focus of data to one process is represented, not mandated by the semantics. The two phases do allgathers in both directions.

### 5.2.3 Specifics for Intercommunicator Collective Operations

All processes in both groups identified by the intercommunicator must call the collective routine.

Note that the “in place” option for intracommunicators does not apply to intercommunicators since in the intercommunicator case there is no communication from a process to itself.

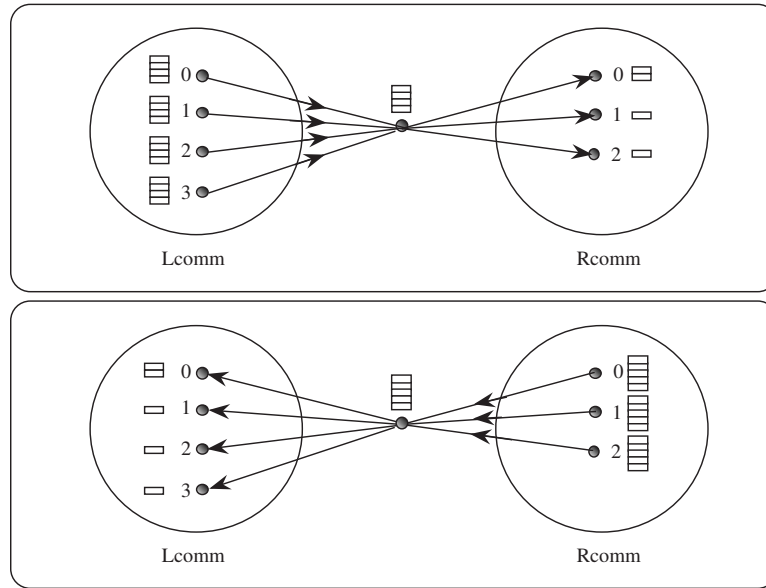


Figure 5.3: Intercommunicator reduce-scatter. The focus of data to one process is represented, not mandated by the semantics. The two phases do reduce-scatters in both directions.

For intercommunicator collective communication, if the operation is in the All-To-One or One-To-All categories, then the transfer is unidirectional. The direction of the transfer is indicated by a special value of the root argument. In this case, for the group containing the root process, all processes in the group must call the routine using a special argument for the root. For this, the root process uses the special root value `MPI_ROOT`; all other processes in the same group as the root use `MPI_PROC_NULL`. All processes in the other group (the group that is the remote group relative to the root process) must call the collective routine and provide the rank of the root. If the operation is in the All-To-All category, then the transfer is bidirectional.

*Rationale.* Operations in the All-To-One and One-To-All categories are unidirectional by nature, and there is a clear way of specifying direction. Operations in the All-To-All category will often occur as part of an exchange, where it makes sense to communicate in both directions at once. (*End of rationale.*)

### 5.3 Barrier Synchronization

`MPI_BARRIER(comm)`

IN            `comm`                            communicator (handle)

`int MPI_Barrier(MPI_Comm comm)`

`MPI_Barrier(comm, ierror) BIND(C)`  
`TYPE(MPI_Comm), INTENT(IN) :: comm`

```
1      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
2      MPI_BARRIER(COMM, IERROR)
```

```
3      INTEGER COMM, IERROR
```

```
4      {void MPI::Comm::Barrier() const = 0(binding deprecated, see Section 15.2) }
```

6 If `comm` is an intracommunicator, `MPI_BARRIER` blocks the caller until all group members have called it. The call returns at any process only after all group members have entered the call.

7 If `comm` is an intercommunicator, `MPI_BARRIER` involves two groups. The call returns at processes in one group (group A) of the intercommunicator only after all members of the other group (group B) have entered the call (and vice versa). A process may return from the call before all processes in its own group have entered the call.

## 15 5.4 Broadcast

```
18      MPI_BCAST(buffer, count, datatype, root, comm)
```

```
19      INOUT      buffer                starting address of buffer (choice)
```

```
20      IN         count                number of entries in buffer (non-negative integer)
```

```
21      IN         datatype             data type of buffer (handle)
```

```
22      IN         root                rank of broadcast root (integer)
```

```
23      IN         comm                communicator (handle)
```

```
24      int MPI_Bcast(void* buffer, int count, MPI_Datatype datatype, int root,
```

```
25      MPI_Comm comm)
```

```
26      MPI_Bcast(buffer, count, datatype, root, comm, ierror) BIND(C)
```

```
27      TYPE(*), DIMENSION(..) :: buffer
```

```
28      INTEGER, INTENT(IN) :: count, root
```

```
29      TYPE(MPI_Datatype), INTENT(IN) :: datatype
```

```
30      TYPE(MPI_Comm), INTENT(IN) :: comm
```

```
31      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
32      MPI_BCAST(BUFFER, COUNT, DATATYPE, ROOT, COMM, IERROR)
```

```
33      <type> BUFFER(*)
```

```
34      INTEGER COUNT, DATATYPE, ROOT, COMM, IERROR
```

```
35      {void MPI::Comm::Bcast(void* buffer, int count,
```

```
36      const MPI::Datatype& datatype, int root) const = 0(binding
```

```
37      deprecated, see Section 15.2) }
```

38 If `comm` is an intracommunicator, `MPI_BCAST` broadcasts a message from the process with rank `root` to all processes of the group, itself included. It is called by all members of the group using the same arguments for `comm` and `root`. On return, the content of `root`'s buffer is copied to all other processes.

General, derived datatypes are allowed for `datatype`. The type signature of `count`, `datatype` on any process must be equal to the type signature of `count`, `datatype` at the root. This implies that the amount of data sent must be equal to the amount received, pairwise between each process and the root. `MPI_BCAST` and all other data-movement collective routines make this restriction. Distinct type maps between sender and receiver are still allowed.

The “in place” option is not meaningful here.

If `comm` is an intercommunicator, then the call involves all processes in the intercommunicator, but with one group (group A) defining the root process. All processes in the other group (group B) pass the same value in argument `root`, which is the rank of the root in group A. The root passes the value `MPI_ROOT` in `root`. All other processes in group A pass the value `MPI_PROC_NULL` in `root`. Data is broadcast from the root to all processes in group B. The buffer arguments of the processes in group B must be consistent with the buffer argument of the root.

#### 5.4.1 Example using `MPI_BCAST`

The examples in this section use intracommunicators.

##### Example 5.1

Broadcast 100 ints from process 0 to every process in the group.

```
MPI_Comm comm;
int array[100];
int root=0;
...
MPI_Bcast(array, 100, MPI_INT, root, comm);
```

As in many of our example code fragments, we assume that some of the variables (such as `comm` in the above) have been assigned appropriate values.

## 5.5 Gather

`MPI_GATHER(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype, root, comm)`

IN	sendbuf	starting address of send buffer (choice)
IN	sendcount	number of elements in send buffer (non-negative integer)
IN	sendtype	data type of send buffer elements (handle)
OUT	recvbuf	address of receive buffer (choice, significant only at root)
IN	recvcount	number of elements for any single receive (non-negative integer, significant only at root)
IN	recvtype	data type of recv buffer elements (significant only at root) (handle)
IN	root	rank of receiving process (integer)
IN	comm	communicator (handle)

```
int MPI_Gather(const void* sendbuf, int sendcount, MPI_Datatype sendtype,
              void* recvbuf, int recvcount, MPI_Datatype recvtype, int root,
              MPI_Comm comm)
```

```
MPI_Gather(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype,
           root, comm, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
    TYPE(*), DIMENSION(..) :: recvbuf
    INTEGER, INTENT(IN) :: sendcount, recvcount, root
    TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_GATHER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT, RECVTYPE,
          ROOT, COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, ROOT, COMM, IERROR
```

```
{void MPI::Comm::Gather(const void* sendbuf, int sendcount, const
    MPI::Datatype& sendtype, void* recvbuf, int recvcount,
    const MPI::Datatype& recvtype, int root) const = 0(binding
    deprecated, see Section 15.2) }
```

If `comm` is an intracommunicator, each process (root process included) sends the contents of its send buffer to the root process. The root process receives the messages and stores them in rank order. The outcome is *as if* each of the `n` processes in the group (including the root process) had executed a call to

```
MPI_Send(sendbuf, sendcount, sendtype, root, ...),
```



and the root had executed `n` calls to

```
MPI_Recv(recvbuf + i · recvcnt · extent(recvtype), recvcnt, recvtype, i, ...),
```

where `extent(recvtype)` is the type extent obtained from a call to `MPI_Type_get_extent()`.

An alternative description is that the `n` messages sent by the processes in the group are concatenated in rank order, and the resulting message is received by the root as if by a call to `MPI_RECV(recvbuf, recvcnt·n, recvtype, ...)`.

The receive buffer is ignored for all non-root processes.

General, derived datatypes are allowed for both `sendtype` and `recvtype`. The type signature of `sendcount`, `sendtype` on each process must be equal to the type signature of `recvcnt`, `recvtype` at the root. This implies that the amount of data sent must be equal to the amount of data received, pairwise between each process and the root. Distinct type maps between sender and receiver are still allowed.

All arguments to the function are significant on process `root`, while on other processes, only arguments `sendbuf`, `sendcount`, `sendtype`, `root`, and `comm` are significant. The arguments `root` and `comm` must have identical values on all processes.

The specification of counts and types should not cause any location on the root to be written more than once. Such a call is erroneous.

Note that the `recvcnt` argument at the root indicates the number of items it receives from *each* process, not the total number of items it receives.

The “in place” option for intracommunicators is specified by passing `MPI_IN_PLACE` as the value of `sendbuf` at the root. In such a case, `sendcount` and `sendtype` are ignored, and the contribution of the root to the gathered vector is assumed to be already in the correct place in the receive buffer.

If `comm` is an intercommunicator, then the call involves all processes in the intercommunicator, but with one group (group A) defining the root process. All processes in the other group (group B) pass the same value in argument `root`, which is the rank of the root in group A. The root passes the value `MPI_ROOT` in `root`. All other processes in group A pass the value `MPI_PROC_NULL` in `root`. Data is gathered from all processes in group B to the root. The send buffer arguments of the processes in group B must be consistent with the receive buffer argument of the root.

1	MPI_GATHERV(sendbuf, sendcount, sendtype, recvbuf, recvcounts, displs, recvtype, root,		
2	comm)		
3	IN	sendbuf	starting address of send buffer (choice)
4	IN	sendcount	number of elements in send buffer (non-negative integer)
5			
6			
7	IN	sendtype	data type of send buffer elements (handle)
8	OUT	recvbuf	address of receive buffer (choice, significant only at root)
9			
10			
11	IN	recvcounts	non-negative integer array (of length group size) containing the number of elements that are received from each process (significant only at root)
12			
13			
14	IN	displs	integer array (of length group size). Entry i specifies the displacement relative to <code>recvbuf</code> at which to place the incoming data from process i (significant only at root)
15			
16			
17			
18	IN	recvtype	data type of recv buffer elements (significant only at root) (handle)
19			
20			
21	IN	root	rank of receiving process (integer)
22	IN	comm	communicator (handle)

```

int MPI_Gatherv(const void* sendbuf, int sendcount, MPI_Datatype sendtype,
               void* recvbuf, const int recvcounts[], const int displs[],
               MPI_Datatype recvtype, int root, MPI_Comm comm)

```

```

MPI_Gatherv(sendbuf, sendcount, sendtype, recvbuf, recvcounts, displs,
            recvtype, root, comm, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
    TYPE(*), DIMENSION(..) :: recvbuf
    INTEGER, INTENT(IN) :: sendcount, recvcounts(*), displs(*), root
    TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_GATHERV(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNTS, DISPLS,
            RECVTYPE, ROOT, COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNT, SENDTYPE, RECVCOUNTS(*), DISPLS(*), RECVTYPE, ROOT,
    COMM, IERROR

```

```

{void MPI::Comm::Gatherv(const void* sendbuf, int sendcount, const
    MPI::Datatype& sendtype, void* recvbuf,
    const int recvcounts[], const int displs[],
    const MPI::Datatype& recvtype, int root) const = 0(binding
    deprecated, see Section 15.2) }

```

MPI\_GATHERV extends the functionality of MPI\_GATHER by allowing a varying count of data from each process, since `recvcounts` is now an array. It also allows more flexibility

as to where the data is placed on the root, by providing the new argument, `displs`.

If `comm` is an intracommunicator, the outcome is *as if* each process, including the root process, sends a message to the root,

```
MPI_Send(sendbuf, sendcount, sendtype, root, ...),
```

and the root executes `n` receives,

```
MPI_Recv(recvbuf + displs[j] * extent(recvtype), recvcnt[j], recvtype, i, ...).
```

The data received from process `j` is placed into `recvbuf` of the root process beginning at offset `displs[j]` elements (in terms of the `recvtype`).

The receive buffer is ignored for all non-root processes.

The type signature implied by `sendcount`, `sendtype` on process `i` must be equal to the type signature implied by `recvcnt[i]`, `recvtype` at the root. This implies that the amount of data sent must be equal to the amount of data received, pairwise between each process and the root. Distinct type maps between sender and receiver are still allowed, as illustrated in Example 5.6.

All arguments to the function are significant on process `root`, while on other processes, only arguments `sendbuf`, `sendcount`, `sendtype`, `root`, and `comm` are significant. The arguments `root` and `comm` must have identical values on all processes.

The specification of counts, types, and displacements should not cause any location on the root to be written more than once. Such a call is erroneous.

The “in place” option for intracommunicators is specified by passing `MPI_IN_PLACE` as the value of `sendbuf` at the root. In such a case, `sendcount` and `sendtype` are ignored, and the contribution of the root to the gathered vector is assumed to be already in the correct place in the receive buffer.

If `comm` is an intercommunicator, then the call involves all processes in the intercommunicator, but with one group (group A) defining the root process. All processes in the other group (group B) pass the same value in argument `root`, which is the rank of the root in group A. The root passes the value `MPI_ROOT` in `root`. All other processes in group A pass the value `MPI_PROC_NULL` in `root`. Data is gathered from all processes in group B to the root. The send buffer arguments of the processes in group B must be consistent with the receive buffer argument of the root.

### 5.5.1 Examples using MPI\_GATHER, MPI\_GATHERV

The examples in this section use intracommunicators.

#### Example 5.2

Gather 100 ints from every process in group to root. See Figure 5.4.

```
MPI_Comm comm;
int gsize, sendarray[100];
int root, *rbuf;
...
MPI_Comm_size(comm, &gsize);
rbuf = (int *)malloc(gsize*100*sizeof(int));
MPI_Gather(sendarray, 100, MPI_INT, rbuf, 100, MPI_INT, root, comm);
```

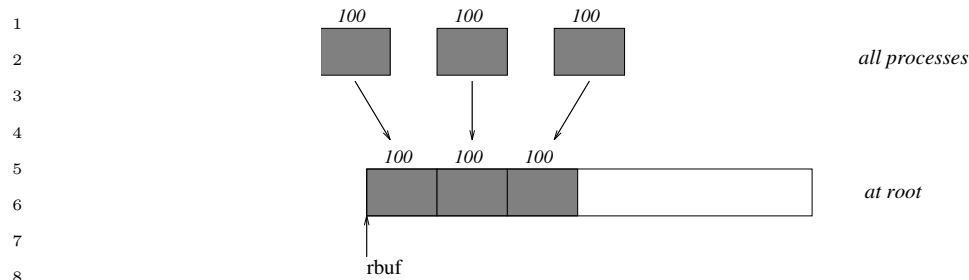


Figure 5.4: The root process gathers 100 ints from each process in the group.

### Example 5.3

Previous example modified – only the root allocates memory for the receive buffer.

```

MPI_Comm comm;
int gsize, sendarray[100];
int root, myrank, *rbuf;
...
MPI_Comm_rank(comm, &myrank);
if (myrank == root) {
    MPI_Comm_size(comm, &gsize);
    rbuf = (int *)malloc(gsize*100*sizeof(int));
}
MPI_Gather(sendarray, 100, MPI_INT, rbuf, 100, MPI_INT, root, comm);

```

### Example 5.4

Do the same as the previous example, but use a derived datatype. Note that the type cannot be the entire set of `gsize*100` ints since type matching is defined pairwise between the root and each process in the gather.

```

MPI_Comm comm;
int gsize, sendarray[100];
int root, *rbuf;
MPI_Datatype rtype;
...
MPI_Comm_size(comm, &gsize);
MPI_Type_contiguous(100, MPI_INT, &rtype);
MPI_Type_commit(&rtype);
rbuf = (int *)malloc(gsize*100*sizeof(int));
MPI_Gather(sendarray, 100, MPI_INT, rbuf, 1, rtype, root, comm);

```

### Example 5.5

Now have each process send 100 ints to root, but place each set (of 100) `stride` ints apart at receiving end. Use `MPI_GATHERV` and the `displs` argument to achieve this effect. Assume `stride`  $\geq 100$ . See Figure 5.5.

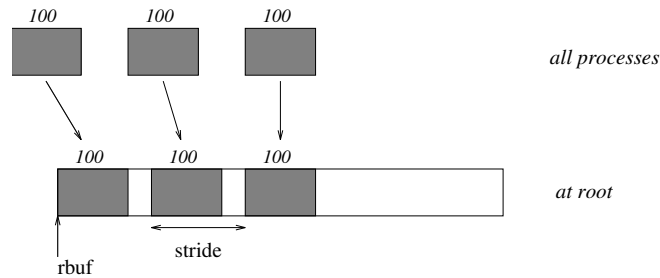


Figure 5.5: The root process gathers 100 ints from each process in the group, each set is placed `stride` ints apart.

```

MPI_Comm comm;
int gsize, sendarray[100];
int root, *rbuf, stride;
int *displs, i, *rcounts;

...

MPI_Comm_size(comm, &gsize);
rbuf = (int *)malloc(gsize*stride*sizeof(int));
displs = (int *)malloc(gsize*sizeof(int));
rcounts = (int *)malloc(gsize*sizeof(int));
for (i=0; i<gsize; ++i) {
    displs[i] = i*stride;
    rcounts[i] = 100;
}
MPI_Gatherv(sendarray, 100, MPI_INT, rbuf, rcounts, displs, MPI_INT,
            root, comm);

```

Note that the program is erroneous if `stride < 100`.

### Example 5.6

Same as Example 5.5 on the receiving side, but send the 100 ints from the 0th column of a 100×150 int array, in C. See Figure 5.6.

```

MPI_Comm comm;
int gsize, sendarray[100][150];
int root, *rbuf, stride;
MPI_Datatype stype;
int *displs, i, *rcounts;

...

MPI_Comm_size(comm, &gsize);
rbuf = (int *)malloc(gsize*stride*sizeof(int));
displs = (int *)malloc(gsize*sizeof(int));
rcounts = (int *)malloc(gsize*sizeof(int));
for (i=0; i<gsize; ++i) {

```

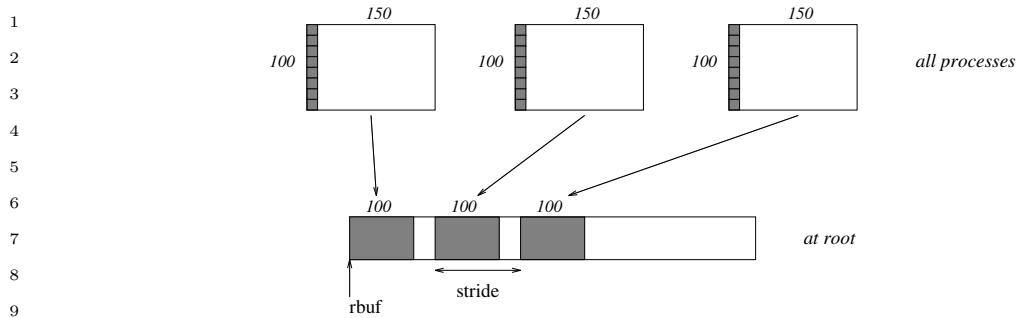


Figure 5.6: The root process gathers column 0 of a 100×150 C array, and each set is placed `stride` ints apart.

```

    displs[i] = i*stride;
    rcounts[i] = 100;
}
/* Create datatype for 1 column of array
 */
MPI_Type_vector(100, 1, 150, MPI_INT, &stype);
MPI_Type_commit(&stype);
MPI_Gatherv(sendarray, 1, stype, rbuf, rcounts, displs, MPI_INT,
            root, comm);

```

### Example 5.7

Process  $i$  sends  $(100-i)$  ints from the  $i$ -th column of a  $100 \times 150$  int array, in C. It is received into a buffer with stride, as in the previous two examples. See Figure 5.7.

```

MPI_Comm comm;
int gsize, sendarray[100][150], *sptr;
int root, *rbuf, stride, myrank;
MPI_Datatype stype;
int *displs, i, *rcounts;

...

MPI_Comm_size(comm, &gsize);
MPI_Comm_rank(comm, &myrank);
rbuf = (int *)malloc(gsize*stride*sizeof(int));
displs = (int *)malloc(gsize*sizeof(int));
rcounts = (int *)malloc(gsize*sizeof(int));
for (i=0; i<gsize; ++i) {
    displs[i] = i*stride;
    rcounts[i] = 100-i;    /* note change from previous example */
}
/* Create datatype for the column we are sending
 */
MPI_Type_vector(100-myrank, 1, 150, MPI_INT, &stype);
MPI_Type_commit(&stype);

```

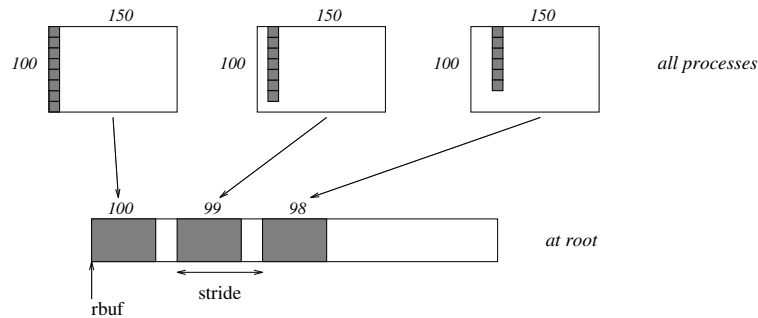


Figure 5.7: The root process gathers  $100-i$  ints from column  $i$  of a  $100 \times 150$  C array, and each set is placed `stride` ints apart.

```
/* sptr is the address of start of "myrank" column
 */
sptr = &sendarray[0][myrank];
MPI_Gatherv(sptr, 1, stype, rbuf, rcounts, displs, MPI_INT,
            root, comm);
```

Note that a different amount of data is received from each process.

### Example 5.8

Same as Example 5.7, but done in a different way at the sending end. We create a datatype that causes the correct striding at the sending end so that we read a column of a C array. A similar thing was done in Example 4.16, Section 4.1.14.

```
MPI_Comm comm;
int gsize, sendarray[100][150], *sptr;
int root, *rbuf, stride, myrank, blocklen[2];
MPI_Aint disp[2];
MPI_Datatype stype, type[2];
int *displs, i, *rcounts;

...

MPI_Comm_size(comm, &gsize);
MPI_Comm_rank(comm, &myrank);
rbuf = (int *)malloc(gsize*stride*sizeof(int));
displs = (int *)malloc(gsize*sizeof(int));
rcounts = (int *)malloc(gsize*sizeof(int));
for (i=0; i<gsize; ++i) {
    displs[i] = i*stride;
    rcounts[i] = 100-i;
}
/* Create datatype for one int, with extent of entire row
 */
disp[0] = 0;      disp[1] = 150*sizeof(int);
type[0] = MPI_INT; type[1] = MPI_UB;
blocklen[0] = 1;  blocklen[1] = 1;
```

```

1 MPI_Type_create_struct(2, blocklen, disp, type, &stype);
2 MPI_Type_commit(&stype);
3 sptr = &sendarray[0][myrank];
4 MPI_Gatherv(sptr, 100-myrank, stype, rbuf, rcounts, displs, MPI_INT,
5                                                     root, comm);
6
7

```

### Example 5.9

Same as Example 5.7 at sending side, but at receiving side we make the stride between received blocks vary from block to block. See Figure 5.8.

```

11 MPI_Comm comm;
12 int gsize, sendarray[100][150], *sptr;
13 int root, *rbuf, *stride, myrank, bufsize;
14 MPI_Datatype stype;
15 int *displs, i, *rcounts, offset;
16
17 ...
18
19 MPI_Comm_size(comm, &gsize);
20 MPI_Comm_rank(comm, &myrank);
21
22 stride = (int *)malloc(gsize*sizeof(int));
23
24 ...
25 /* stride[i] for i = 0 to gsize-1 is set somehow
26    */
27
28 /* set up displs and rcounts vectors first
29    */
30 displs = (int *)malloc(gsize*sizeof(int));
31 rcounts = (int *)malloc(gsize*sizeof(int));
32 offset = 0;
33 for (i=0; i<gsize; ++i) {
34     displs[i] = offset;
35     offset += stride[i];
36     rcounts[i] = 100-i;
37 }
38 /* the required buffer size for rbuf is now easily obtained
39    */
40 bufsize = displs[gsize-1]+rcounts[gsize-1];
41 rbuf = (int *)malloc(bufsize*sizeof(int));
42 /* Create datatype for the column we are sending
43    */
44 MPI_Type_vector(100-myrank, 1, 150, MPI_INT, &stype);
45 MPI_Type_commit(&stype);
46 sptr = &sendarray[0][myrank];
47 MPI_Gatherv(sptr, 1, stype, rbuf, rcounts, displs, MPI_INT,
48                                                     root, comm);

```



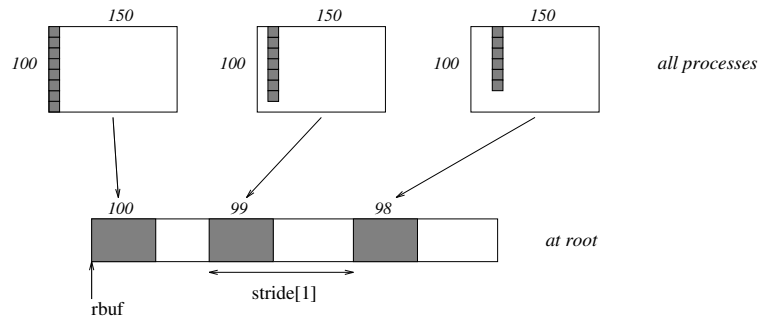


Figure 5.8: The root process gathers  $100-i$  ints from column  $i$  of a  $100 \times 150$  C array, and each set is placed  $\text{stride}[i]$  ints apart (a varying stride).

### Example 5.10

Process  $i$  sends  $\text{num}$  ints from the  $i$ -th column of a  $100 \times 150$  int array, in C. The complicating factor is that the various values of  $\text{num}$  are not known to *root*, so a separate gather must first be run to find these out. The data is placed contiguously at the receiving end.

```

MPI_Comm comm;
int gsize, sendarray[100][150], *sptr;
int root, *rbuf, myrank, blocklen[2];
MPI_Aint disp[2];
MPI_Datatype stype, type[2];
int *displs, i, *rcounts, num;

...

MPI_Comm_size(comm, &gsize);
MPI_Comm_rank(comm, &myrank);

/* First, gather nums to root
 */
rcounts = (int *)malloc(gsize*sizeof(int));
MPI_Gather(&num, 1, MPI_INT, rcounts, 1, MPI_INT, root, comm);
/* root now has correct rcounts, using these we set displs[] so
 * that data is placed contiguously (or concatenated) at receive end
 */
displs = (int *)malloc(gsize*sizeof(int));
displs[0] = 0;
for (i=1; i<gsize; ++i) {
    displs[i] = displs[i-1]+rcounts[i-1];
}
/* And, create receive buffer
 */
rbuf = (int *)malloc(gsize*(displs[gsize-1]+rcounts[gsize-1])
                    *sizeof(int));

/* Create datatype for one int, with extent of entire row

```

```

1      */
2      disp[0] = 0;          disp[1] = 150*sizeof(int);
3      type[0] = MPI_INT; type[1] = MPI_UB;
4      blocklen[0] = 1;   blocklen[1] = 1;
5      MPI_Type_create_struct( 2, blocklen, disp, type, &stype );
6      MPI_Type_commit(&stype);
7      sptr = &sendarray[0][myrank];
8      MPI_Gatherv(sptr, num, stype, rbuf, rcounts, displs, MPI_INT,
9                  root, comm);
10
11
12
13
14

```

## 5.6 Scatter

MPI\_SCATTER(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype, root, comm)

IN	sendbuf	address of send buffer (choice, significant only at root)
IN	sendcount	number of elements sent to each process (non-negative integer, significant only at root)
IN	sendtype	data type of send buffer elements (significant only at root) (handle)
OUT	recvbuf	address of receive buffer (choice)
IN	recvcount	number of elements in receive buffer (non-negative integer)
IN	recvtype	data type of receive buffer elements (handle)
IN	root	rank of sending process (integer)
IN	comm	communicator (handle)

```

30 int MPI_Scatter(const void* sendbuf, int sendcount, MPI_Datatype sendtype,
31               void* recvbuf, int recvcount, MPI_Datatype recvtype, int root,
32               MPI_Comm comm)
33

```

```

34 MPI_Scatter(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype,
35            root, comm, ierror) BIND(C)
36     TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
37     TYPE(*), DIMENSION(..)  :: recvbuf
38     INTEGER, INTENT(IN) :: sendcount, recvcount, root
39     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
40     TYPE(MPI_Comm), INTENT(IN) :: comm
41     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
42

```

```

43 MPI_SCATTER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT, RECVTYPE,
44            ROOT, COMM, IERROR)
45     <type> SENDBUF(*), RECVBUF(*)
46     INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, ROOT, COMM, IERROR
47

```

```

48 {void MPI::Comm::Scatter(const void* sendbuf, int sendcount, const
49     MPI::Datatype& sendtype, void* recvbuf, int recvcount,

```

```
const MPI::Datatype& recvtype, int root) const = 0(binding
deprecated, see Section 15.2) }
```

MPI\_SCATTER is the inverse operation to MPI\_GATHER.

If `comm` is an intracommunicator, the outcome is *as if* the root executed `n` send operations,

```
MPI_Send(sendbuf + i · sendcount · extent(sendtype), sendcount, sendtype, i, ...),
```

and each process executed a receive,

```
MPI_Recv(recvbuf, recvcount, recvtype, i, ...).
```

An alternative description is that the root sends a message with `MPI_Send(sendbuf, sendcount·n, sendtype, ...)`. This message is split into `n` equal segments, the  $i$ -th segment is sent to the  $i$ -th process in the group, and each process receives this message as above.

The send buffer is ignored for all non-root processes.

The type signature associated with `sendcount`, `sendtype` at the root must be equal to the type signature associated with `recvcount`, `recvtype` at all processes (however, the type maps may be different). This implies that the amount of data sent must be equal to the amount of data received, pairwise between each process and the root. Distinct type maps between sender and receiver are still allowed.

All arguments to the function are significant on process `root`, while on other processes, only arguments `recvbuf`, `recvcount`, `recvtype`, `root`, and `comm` are significant. The arguments `root` and `comm` must have identical values on all processes.

The specification of counts and types should not cause any location on the root to be read more than once.

*Rationale.* Though not needed, the last restriction is imposed so as to achieve symmetry with MPI\_GATHER, where the corresponding restriction (a multiple-write restriction) is necessary. (*End of rationale.*)

The “in place” option for intracommunicators is specified by passing `MPI_IN_PLACE` as the value of `recvbuf` at the root. In such a case, `recvcount` and `recvtype` are ignored, and root “sends” no data to itself. The scattered vector is still assumed to contain  $n$  segments, where  $n$  is the group size; the  $root$ -th segment, which root should “send to itself,” is not moved.

If `comm` is an intercommunicator, then the call involves all processes in the intercommunicator, but with one group (group A) defining the root process. All processes in the other group (group B) pass the same value in argument `root`, which is the rank of the root in group A. The root passes the value `MPI_ROOT` in `root`. All other processes in group A pass the value `MPI_PROC_NULL` in `root`. Data is scattered from the root to all processes in group B. The receive buffer arguments of the processes in group B must be consistent with the send buffer argument of the root.

1	MPI_SCATTERV(sendbuf, sendcounts, displs, sendtype, recvbuf, recvcount, recvtype, root,		
2	comm)		
3	IN	sendbuf	address of send buffer (choice, significant only at root)
4	IN	sendcounts	non-negative integer array (of length group size) specifying the number of elements to send to each <a href="#">[processor]</a> <a href="#">rank</a> <a href="#">ticket0</a> .
5			
6			
7	IN	displs	integer array (of length group size). Entry <i>i</i> specifies the displacement (relative to <code>sendbuf</code> ) from which to take the outgoing data to process <i>i</i>
8			
9			
10	IN	sendtype	data type of send buffer elements (handle)
11	OUT	recvbuf	address of receive buffer (choice)
12	IN	recvcount	number of elements in receive buffer (non-negative integer)
13			
14			
15	IN	recvtype	data type of receive buffer elements (handle)
16			
17	IN	root	rank of sending process (integer)
18	IN	comm	communicator (handle)
19			

```

20 int MPI_Scatterv(const void* sendbuf, const int sendcounts[], const
21                 int displs[], MPI_Datatype sendtype, void* recvbuf,
22                 int recvcount, MPI_Datatype recvtype, int root, MPI_Comm comm)

```

```

23 MPI_Scatterv(sendbuf, sendcounts, displs, sendtype, recvbuf, recvcount,
24             recvtype, root, comm, ierror) BIND(C)
25     TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
26     TYPE(*), DIMENSION(..) :: recvbuf
27     INTEGER, INTENT(IN) :: sendcounts(*), displs(*), recvcount, root
28     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
29     TYPE(MPI_Comm), INTENT(IN) :: comm
30     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
31
32 MPI_SCATTERV(SENDBUF, SENDCOUNTS, DISPLS, SENDTYPE, RECVBUF, REVCOUNT,
33             RECVTYPE, ROOT, COMM, IERROR)
34 <type> SENDBUF(*), RECVBUF(*)
35 INTEGER SENDCOUNTS(*), DISPLS(*), SENDTYPE, REVCOUNT, RECVTYPE, ROOT,
36 COMM, IERROR

```

```

37
38 {void MPI::Comm::Scatterv(const void* sendbuf, const int sendcounts[],
39                          const int displs[], const MPI::Datatype& sendtype,
40                          void* recvbuf, int recvcount, const MPI::Datatype& recvtype,
41                          int root) const = 0(binding deprecated, see Section 15.2) }

```

MPI\_SCATTERV is the inverse operation to MPI\_GATHERV.

MPI\_SCATTERV extends the functionality of MPI\_SCATTER by allowing a varying count of data to be sent to each process, since `sendcounts` is now an array. It also allows more flexibility as to where the data is taken from on the root, by providing an additional argument, `displs`.

If `comm` is an intracommunicator, the outcome is as if the root executed `n` send operations,

```
MPI_Send(sendbuf + displs[i] * extent(sendtype), sendcounts[i], sendtype, i, ...),
```

and each process executed a receive,

```
MPI_Recv(recvbuf, recvcount, recvtype, i, ...).
```

The send buffer is ignored for all non-root processes.

The type signature implied by `sendcount[i]`, `sendtype` at the root must be equal to the type signature implied by `recvcount`, `recvtype` at process `i` (however, the type maps may be different). This implies that the amount of data sent must be equal to the amount of data received, pairwise between each process and the root. Distinct type maps between sender and receiver are still allowed.

All arguments to the function are significant on process `root`, while on other processes, only arguments `recvbuf`, `recvcount`, `recvtype`, `root`, and `comm` are significant. The arguments `root` and `comm` must have identical values on all processes.

The specification of counts, types, and displacements should not cause any location on the root to be read more than once.

The “in place” option for intracommunicators is specified by passing `MPI_IN_PLACE` as the value of `recvbuf` at the root. In such a case, `recvcount` and `recvtype` are ignored, and root “sends” no data to itself. The scattered vector is still assumed to contain  $n$  segments, where  $n$  is the group size; the *root*-th segment, which root should “send to itself,” is not moved.

If `comm` is an intercommunicator, then the call involves all processes in the intercommunicator, but with one group (group A) defining the root process. All processes in the other group (group B) pass the same value in argument `root`, which is the rank of the root in group A. The root passes the value `MPI_ROOT` in `root`. All other processes in group A pass the value `MPI_PROC_NULL` in `root`. Data is scattered from the root to all processes in group B. The receive buffer arguments of the processes in group B must be consistent with the send buffer argument of the root.

### 5.6.1 Examples using MPI\_SCATTER, MPI\_SCATTERV

The examples in this section use intracommunicators.

#### Example 5.11

The reverse of Example 5.2. Scatter sets of 100 ints from the root to each process in the group. See Figure 5.9.

```
MPI_Comm comm;
int gsize,*sendbuf;
int root, rbuf[100];
...
MPI_Comm_size(comm, &gsize);
sendbuf = (int *)malloc(gsize*100*sizeof(int));
...
MPI_Scatter(sendbuf, 100, MPI_INT, rbuf, 100, MPI_INT, root, comm);
```

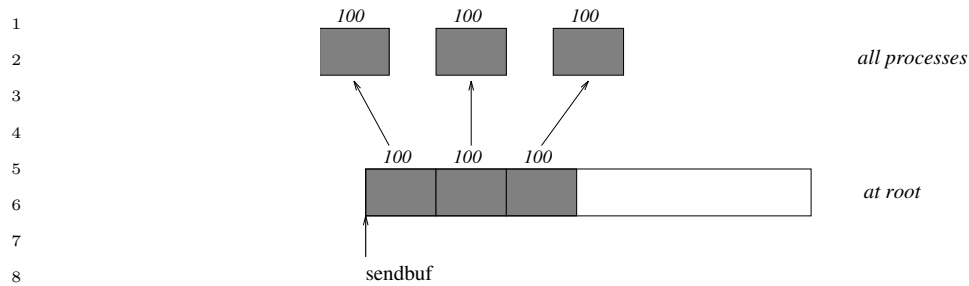


Figure 5.9: The root process scatters sets of 100 ints to each process in the group.

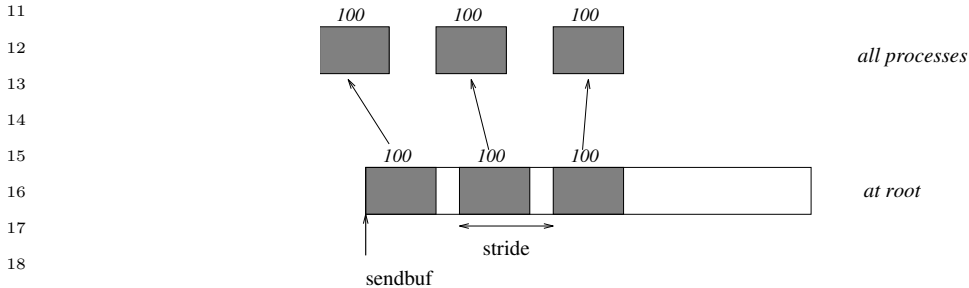


Figure 5.10: The root process scatters sets of 100 ints, moving by `stride` ints from send to send in the scatter.

### Example 5.12

The reverse of Example 5.5. The root process scatters sets of 100 ints to the other processes, but the sets of 100 are *stride* ints apart in the sending buffer. Requires use of `MPI_SCATTERV`. Assume *stride*  $\geq$  100. See Figure 5.10.

```

MPI_Comm comm;
int gsize,*sendbuf;
int root, rbuf[100], i, *displs, *scounts;

...

MPI_Comm_size(comm, &gsize);
sendbuf = (int *)malloc(gsize*stride*sizeof(int));
...
displs = (int *)malloc(gsize*sizeof(int));
scount = (int *)malloc(gsize*sizeof(int));
for (i=0; i<gsize; ++i) {
    displs[i] = i*stride;
    scounts[i] = 100;
}
MPI_Scatterv(sendbuf, scounts, displs, MPI_INT, rbuf, 100, MPI_INT,
             root, comm);

```

### Example 5.13

The reverse of Example 5.9. We have a varying stride between blocks at sending (root) side, at the receiving side we receive into the  $i$ -th column of a  $100 \times 150$  C array. See Figure 5.11.

```

MPI_Comm comm;
int gsize, recvarray[100][150], *rptr;
int root, *sendbuf, myrank, *stride;
MPI_Datatype rtype;
int i, *displs, *counts, offset;
...
MPI_Comm_size(comm, &gsize);
MPI_Comm_rank(comm, &myrank);

stride = (int *)malloc(gsize*sizeof(int));
...
/* stride[i] for i = 0 to gsize-1 is set somehow
 * sendbuf comes from elsewhere
 */
...
displs = (int *)malloc(gsize*sizeof(int));
counts = (int *)malloc(gsize*sizeof(int));
offset = 0;
for (i=0; i<gsize; ++i) {
    displs[i] = offset;
    offset += stride[i];
    counts[i] = 100 - i;
}
/* Create datatype for the column we are receiving
 */
MPI_Type_vector(100-myrank, 1, 150, MPI_INT, &rtype);
MPI_Type_commit(&rtype);
rptr = &recvarray[0][myrank];
MPI_Scatterv(sendbuf, counts, displs, MPI_INT, rptr, 1, rtype,
            root, comm);

```

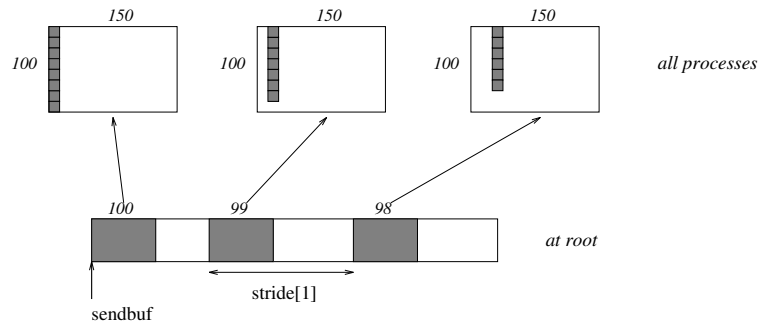


Figure 5.11: The root scatters blocks of  $100-i$  ints into column  $i$  of a  $100 \times 150$  C array. At the sending side, the blocks are `stride[i]` ints apart.

## 5.7 Gather-to-all

`MPI_ALLGATHER(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype, comm)`

IN	<code>sendbuf</code>	starting address of send buffer (choice)
IN	<code>sendcount</code>	number of elements in send buffer (non-negative integer)
IN	<code>sendtype</code>	data type of send buffer elements (handle)
OUT	<code>recvbuf</code>	address of receive buffer (choice)
IN	<code>recvcount</code>	number of elements received from any process (non-negative integer)
IN	<code>recvtype</code>	data type of receive buffer elements (handle)
IN	<code>comm</code>	communicator (handle)

```

int MPI_Allgather(const void* sendbuf, int sendcount,
                  MPI_Datatype sendtype, void* recvbuf, int recvcount,
                  MPI_Datatype recvtype, MPI_Comm comm)
MPI_Allgather(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype,
              comm, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
    TYPE(*), DIMENSION(..) :: recvbuf
    INTEGER, INTENT(IN) :: sendcount, recvcount
    TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_ALLGATHER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT, RECVTYPE,
              COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, COMM, IERROR

```



```
{void MPI::Comm::Allgather(const void* sendbuf, int sendcount, const
    MPI::Datatype& sendtype, void* recvbuf, int recvcount,
    const MPI::Datatype& recvtype) const = 0(binding deprecated, see
    Section 15.2) }
```

MPI\_ALLGATHER can be thought of as MPI\_GATHER, but where all processes receive the result, instead of just the root. The block of data sent from the  $j$ -th process is received by every process and placed in the  $j$ -th block of the buffer `recvbuf`.

The type signature associated with `sendcount`, `sendtype`, at a process must be equal to the type signature associated with `recvcount`, `recvtype` at any other process.

If `comm` is an intracommunicator, the outcome of a call to `MPI_ALLGATHER(...)` is as if all processes executed `n` calls to

```
MPI_Gather(sendbuf, sendcount, sendtype, recvbuf, recvcount,
           recvtype, root, comm)
```

for `root = 0, ..., n-1`. The rules for correct usage of `MPI_ALLGATHER` are easily found from the corresponding rules for `MPI_GATHER`.

The “in place” option for intracommunicators is specified by passing the value `MPI_IN_PLACE` to the argument `sendbuf` at all processes. `sendcount` and `sendtype` are ignored. Then the input data of each process is assumed to be in the area where that process would receive its own contribution to the receive buffer.

If `comm` is an intercommunicator, then each process of one group (group A) contributes `sendcount` data items; these data are concatenated and the result is stored at each process in the other group (group B). Conversely the concatenation of the contributions of the processes in group B is stored at each process in group A. The send buffer arguments in group A must be consistent with the receive buffer arguments in group B, and vice versa.

*Advice to users.* The communication pattern of `MPI_ALLGATHER` executed on an intercommunication domain need not be symmetric. The number of items sent by processes in group A (as specified by the arguments `sendcount`, `sendtype` in group A and the arguments `recvcount`, `recvtype` in group B), need not equal the number of items sent by processes in group B (as specified by the arguments `sendcount`, `sendtype` in group B and the arguments `recvcount`, `recvtype` in group A). In particular, one can move data in only one direction by specifying `sendcount = 0` for the communication in the reverse direction.

*(End of advice to users.)*

1 MPI\_ALLGATHERV(sendbuf, sendcount, sendtype, recvbuf, recvcunts, displs, recvtype, comm)

3	IN	sendbuf	starting address of send buffer (choice)
4	IN	sendcount	number of elements in send buffer (non-negative integer)
5	IN	sendtype	data type of send buffer elements (handle)
6	OUT	recvbuf	address of receive buffer (choice)
7	IN	recvcunts	non-negative integer array (of length group size) containing the number of elements that are received from each process
8	IN	displs	integer array (of length group size). Entry <i>i</i> specifies the displacement (relative to <i>recvbuf</i> ) at which to place the incoming data from process <i>i</i>
9	IN	recvtype	data type of receive buffer elements (handle)
10	IN	comm	communicator (handle)

ticket140. 20 int MPI\_Allgatherv(const void\* sendbuf, int sendcount,  
ticket140. 21 MPI\_Datatype sendtype, void\* recvbuf, const int recvcunts[],  
ticket140. 22 const int displs[], MPI\_Datatype recvtype, MPI\_Comm comm)

ticket229.1. 23 MPI\_Allgatherv(sendbuf, sendcount, sendtype, recvbuf, recvcunts, displs,  
ticket-248T. 24 recvtype, comm, ierror) BIND(C)  
25 TYPE(\*), DIMENSION(..), INTENT(IN) :: sendbuf  
26 TYPE(\*), DIMENSION(..) :: recvbuf  
27 INTEGER, INTENT(IN) :: sendcount, recvcunts(\*), displs(\*)  
28 TYPE(MPI\_Datatype), INTENT(IN) :: sendtype, recvtype  
29 TYPE(MPI\_Comm), INTENT(IN) :: comm  
30 INTEGER, OPTIONAL, INTENT(OUT) :: ierror

32 MPI\_ALLGATHERV(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNTS, DISPLS,  
33 RECVTYPE, COMM, IERROR)  
34 <type> SENDBUF(\*), RECVBUF(\*)  
35 INTEGER SENDCOUNT, SENDTYPE, RECVCOUNTS(\*), DISPLS(\*), RECVTYPE, COMM,  
36 IERROR

37 {void MPI::Comm::Allgatherv(const void\* sendbuf, int sendcount, const  
38 MPI::Datatype& sendtype, void\* recvbuf,  
39 const int recvcunts[], const int displs[],  
40 const MPI::Datatype& recvtype) const = 0(*binding deprecated, see*  
41 *Section 15.2*) }

43 MPI\_ALLGATHERV can be thought of as MPI\_GATHERV, but where all processes re-  
44 ceive the result, instead of just the root. The block of data sent from the *j*-th process is  
45 received by every process and placed in the *j*-th block of the buffer *recvbuf*. These blocks  
46 need not all be the same size.

47 The type signature associated with *sendcount*, *sendtype*, at process *j* must be equal to  
48 the type signature associated with *recvcunts[j]*, *recvtype* at any other process.

If `comm` is an intracommunicator, the outcome is as if all processes executed calls to

```
MPI_GATHERV(sendbuf, sendcount, sendtype, recvbuf, recvcounts, displs,
            recvtype, root, comm),
```

for `root = 0, ..., n-1`. The rules for correct usage of `MPI_ALLGATHERV` are easily found from the corresponding rules for `MPI_GATHERV`.

The “in place” option for intracommunicators is specified by passing the value `MPI_IN_PLACE` to the argument `sendbuf` at all processes. In such a case, `sendcount` and `sendtype` are ignored, and the input data of each process is assumed to be in the area where that process would receive its own contribution to the receive buffer.

If `comm` is an intercommunicator, then each process of one group (group A) contributes `sendcount` data items; these data are concatenated and the result is stored at each process in the other group (group B). Conversely the concatenation of the contributions of the processes in group B is stored at each process in group A. The send buffer arguments in group A must be consistent with the receive buffer arguments in group B, and vice versa.

### 5.7.1 Example using `MPI_ALLGATHER`

The example in this section uses intracommunicators.

#### Example 5.14

The all-gather version of Example 5.2. Using `MPI_ALLGATHER`, we will gather 100 ints from every process in the group to every process.

```
MPI_Comm comm;
int gsize, sendarray[100];
int *rbuf;
...
MPI_Comm_size(comm, &gsize);
rbuf = (int *)malloc(gsize*100*sizeof(int));
MPI_Allgather(sendarray, 100, MPI_INT, rbuf, 100, MPI_INT, comm);
```

After the call, every process has the group-wide concatenation of the sets of data.

## 5.8 All-to-All Scatter/Gather

```

MPI_ALLTOALL(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype, comm)

    IN      sendbuf      starting address of send buffer (choice)
    IN      sendcount    number of elements sent to each process (non-negative
                        integer)
    IN      sendtype     data type of send buffer elements (handle)
    OUT     recvbuf      address of receive buffer (choice)
    IN      recvcount    number of elements received from any process (non-
                        negative integer)
    IN      recvtype     data type of receive buffer elements (handle)
    IN      comm         communicator (handle)

```

```

int MPI_Alltoall(const void* sendbuf, int sendcount, MPI_Datatype sendtype,
                void* recvbuf, int recvcount, MPI_Datatype recvtype,
                MPI_Comm comm)

```

```

MPI_Alltoall(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype,
             comm, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
    TYPE(*), DIMENSION(..) :: recvbuf
    INTEGER, INTENT(IN) :: sendcount, recvcount
    TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_ALLTOALL(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT, RECVTYPE,
             COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, COMM, IERROR

```

```

{void MPI::Comm::Alltoall(const void* sendbuf, int sendcount, const
                        MPI::Datatype& sendtype, void* recvbuf, int recvcount,
                        const MPI::Datatype& recvtype) const = 0 (binding deprecated, see
                        Section 15.2) }

```

MPI\_ALLTOALL is an extension of MPI\_ALLGATHER to the case where each process sends distinct data to each of the receivers. The  $j$ -th block sent from process  $i$  is received by process  $j$  and is placed in the  $i$ -th block of `recvbuf`.

The type signature associated with `sendcount`, `sendtype`, at a process must be equal to the type signature associated with `recvcount`, `recvtype` at any other process. This implies that the amount of data sent must be equal to the amount of data received, pairwise between every pair of processes. As usual, however, the type maps may be different.

If `comm` is an intracommunicator, the outcome is as if each process executed a send to each process (itself included) with a call to,

```

MPI_Send(sendbuf + i · sendcount · extent(sendtype), sendcount, sendtype, i, ...),

```

and a receive from every other process with a call to,

```
MPI_Recv(recvbuf + i · recvcnt · extent(recvtype), recvcnt, recvtype, i, ...).
```

All arguments on all processes are significant. The argument `comm` must have identical values on all processes.

The “in place” option for intracommunicators is specified by passing `MPI_IN_PLACE` to the argument `sendbuf` at *all* processes. In such a case, `sendcount` and `sendtype` are ignored. The data to be sent is taken from the `recvbuf` and replaced by the received data. Data sent and received must have the same type map as specified by `recvcnt` and `recvtype`.

*Rationale.* For large `MPI_ALLTOALL` instances, allocating both send and receive buffers may consume too much memory. The “in place” option effectively halves the application memory consumption and is useful in situations where the data to be sent will not be used by the sending process after the `MPI_ALLTOALL` exchange (e.g., in parallel Fast Fourier Transforms). (*End of rationale.*)

*Advice to implementors.* Users may opt to use the “in place” option in order to conserve memory. Quality MPI implementations should thus strive to minimize system buffering. (*End of advice to implementors.*)

If `comm` is an intercommunicator, then the outcome is as if each process in group A sends a message to each process in group B, and vice versa. The *j*-th send buffer of process *i* in group A should be consistent with the *i*-th receive buffer of process *j* in group B, and vice versa.

*Advice to users.* When a complete exchange is executed on an intercommunication domain, then the number of data items sent from processes in group A to processes in group B need not equal the number of items sent in the reverse direction. In particular, one can have unidirectional communication by specifying `sendcount = 0` in the reverse direction.

(*End of advice to users.*)

1	MPI_ALLTOALLV(sendbuf, sendcounts, sdispls, sendtype, recvbuf, recvcounts, rdispls,			
2	recvtype, comm)			
3	IN	sendbuf	starting address of send buffer (choice)	
4	IN	sendcounts	non-negative integer array (of length group size) specifying the number of elements to send to each [processor]rank	ticket0.
5				
6				
7	IN	sdispls	integer array (of length group size). Entry j specifies the displacement (relative to sendbuf) from which to take the outgoing data destined for process j	ticket109.
8				
9				
10	IN	sendtype	data type of send buffer elements (handle)	
11	OUT	recvbuf	address of receive buffer (choice)	
12	IN	recvcounts	non-negative integer array (of length group size) specifying the number of elements that can be received from each [processor]rank	ticket0.
13				
14				
15	IN	rdispls	integer array (of length group size). Entry i specifies the displacement (relative to recvbuf) at which to place the incoming data from process i	ticket109.
16				
17				
18				
19	IN	recvtype	data type of receive buffer elements (handle)	
20	IN	comm	communicator (handle)	
21				
22				

```

23 int MPI_Alltoallv(const void* sendbuf, const int sendcounts[], const
24                 int sdispls[], MPI_Datatype sendtype, void* recvbuf, const
25                 int recvcounts[], const int rdispls[], MPI_Datatype recvtype,
26                 MPI_Comm comm)

```

```

27 MPI_Alltoallv(sendbuf, sendcounts, sdispls, sendtype, recvbuf, recvcounts,
28               rdispls, recvtype, comm, ierror) BIND(C)
29
30 TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
31 TYPE(*), DIMENSION(..) :: recvbuf
32 INTEGER, INTENT(IN) :: sendcounts(*), sdispls(*), recvcounts(*),
33 rdispls(*)
34 TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
35 TYPE(MPI_Comm), INTENT(IN) :: comm
36 INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

37 MPI_ALLTOALLV(SENDBUF, SENDCOUNTS, SDISPLS, SENDTYPE, RECVBUF, RECVCOUNTS,
38               RDISPLS, RECVTYPE, COMM, IERROR)
39
40 <type> SENDBUF(*), RECVBUF(*)
41 INTEGER SENDCOUNTS(*), SDISPLS(*), SENDTYPE, RECVCOUNTS(*), RDISPLS(*),
42 RECVTYPE, COMM, IERROR

```

```

43 {void MPI::Comm::Alltoallv(const void* sendbuf, const int sendcounts[],
44                           const int sdispls[], const MPI::Datatype& sendtype,
45                           void* recvbuf, const int recvcounts[], const int rdispls[],
46                           const MPI::Datatype& recvtype) const = 0(binding deprecated, see
47                           Section 15.2) }
48

```

MPI\_ALLTOALLV adds flexibility to MPI\_ALLTOALL in that the location of data for the send is specified by `sdispls` and the location of the placement of the data on the receive side is specified by `rdispls`.

If `comm` is an intracommunicator, then the  $j$ -th block sent from process  $i$  is received by process  $j$  and is placed in the  $i$ -th block of `recvbuf`. These blocks need not all have the same size.

The type signature associated with `sendcounts[j]`, `sendtype` at process  $i$  must be equal to the type signature associated with `recvcounts[i]`, `recvtype` at process  $j$ . This implies that the amount of data sent must be equal to the amount of data received, pairwise between every pair of processes. Distinct type maps between sender and receiver are still allowed.

The outcome is as if each process sent a message to every other process with,

```
MPI_Send(sendbuf + sdispls[i] · extent(sendtype), sendcounts[i], sendtype, i, ...),
```

and received a message from every other process with a call to

```
MPI_Recv(recvbuf + rdispls[i] · extent(recvtype), recvcounts[i], recvtype, i, ...).
```

All arguments on all processes are significant. The argument `comm` must have identical values on all processes.

The “in place” option for intracommunicators is specified by passing `MPI_IN_PLACE` to the argument `sendbuf` at *all* processes. In such a case, `sendcounts`, `sdispls` and `sendtype` are ignored. The data to be sent is taken from the `recvbuf` and replaced by the received data. Data sent and received must have the same type map as specified by the `recvcounts` array and the `recvtype`, and is taken from the locations of the receive buffer specified by `rdispls`.

*Advice to users.* Specifying the “in place” option (which must be given on all processes) implies that the same amount and type of data is sent and received between any two processes in the group of the communicator. Different pairs of processes can exchange different amounts of data. Users must ensure that `recvcounts[j]` and `recvtype` on process  $i$  match `recvcounts[i]` and `recvtype` on process  $j$ . This symmetric exchange can be useful in applications where the data to be sent will not be used by the sending process after the MPI\_ALLTOALLV exchange. (*End of advice to users.*)

If `comm` is an intercommunicator, then the outcome is as if each process in group A sends a message to each process in group B, and vice versa. The  $j$ -th send buffer of process  $i$  in group A should be consistent with the  $i$ -th receive buffer of process  $j$  in group B, and vice versa.

*Rationale.* The definitions of MPI\_ALLTOALL and MPI\_ALLTOALLV give as much flexibility as one would achieve by specifying  $n$  independent, point-to-point communications, with two exceptions: all messages use the same datatype, and messages are scattered from (or gathered to) sequential storage. (*End of rationale.*)

*Advice to implementors.* Although the discussion of collective communication in terms of point-to-point operation implies that each message is transferred directly from sender to receiver, implementations may use a tree communication pattern. Messages can be forwarded by intermediate nodes where they are split (for scatter) or concatenated (for gather), if this is more efficient. (*End of advice to implementors.*)

	1	MPI_ALLTOALLW(sendbuf, sendcounts, sdispls, sendtypes, recvbuf, recvcounts, rdispls,	
	2	recvtypes, comm)	
	3		
	4	IN	sendbuf starting address of send buffer (choice)
	5	IN	sendcounts non-negative integer array (of length group size) speci-
ticket0.	6		fying the number of elements to send to each [processor]rank
	7	IN	sdispls integer array (of length group size). Entry j specifies
	8		the displacement in bytes (relative to sendbuf) from
	9		which to take the outgoing data destined for process
	10		j (array of integers)
	11	IN	sendtypes array of datatypes (of length group size). Entry j
	12		specifies the type of data to send to process j (array
	13		of handles)
	14		
	15	OUT	recvbuf address of receive buffer (choice)
	16	IN	recvcounts non-negative integer array (of length group size) spec-
	17		ifying the number of elements that can be received
ticket0.	18		from each [processor]rank
	19	IN	rdispls integer array (of length group size). Entry i specifies
	20		the displacement in bytes (relative to recvbuf) at which
	21		to place the incoming data from process i (array of
	22		integers)
	23	IN	recvtypes array of datatypes (of length group size). Entry i
	24		specifies the type of data received from process i (ar-
	25		ray of handles)
	26		
	27	IN	comm communicator (handle)
	28		
ticket140.	29	int MPI_Alltoallw(const void* sendbuf, const int sendcounts[], const	
ticket140.	30	int sdispls[], const MPI_Datatype sendtypes[], void* recvbuf,	
ticket140.	31	const int recvcounts[], const int rdispls[], const	
ticket140.	32	MPI_Datatype recvtypes[], MPI_Comm comm)	
ticket140.	33		
ticket140.	34	MPI_Alltoallw(sendbuf, sendcounts, sdispls, sendtypes, recvbuf, recvcounts,	
ticket140.	35	rdispls, recvtypes, comm, ierror) BIND(C)	
ticket-248T.	36	TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf	
	37	TYPE(*), DIMENSION(..) :: recvbuf	
	38	INTEGER, INTENT(IN) :: sendcounts(*), sdispls(*), recvcounts(*),	
	39	rdispls(*)	
	40	TYPE(MPI_Datatype), INTENT(IN) :: sendtypes(*)	
	41	TYPE(MPI_Datatype), INTENT(IN) :: recvtypes(*)	
	42	TYPE(MPI_Comm), INTENT(IN) :: comm	
	43	INTEGER, OPTIONAL, INTENT(OUT) :: ierror	
	44	MPI_ALLTOALLW(SENDBUF, SENDCOUNTS, SDISPLS, SENDTYPES, RECVBUF, RECVCOUNTS,	
	45	RDISPLS, RECVTYPES, COMM, IERROR)	
	46	<type> SENDBUF(*), RECVBUF(*)	
	47	INTEGER SENDCOUNTS(*), SDISPLS(*), SENDTYPES(*), RECVCOUNTS(*),	
	48	RDISPLS(*), RECVTYPES(*), COMM, IERROR	



```
{void MPI::Comm::Alltoallw(const void* sendbuf, const int sendcounts[],
    const int sdispls[], const MPI::Datatype sendtypes[], void*
    recvbuf, const int recvcounst[], const int rdispls[], const
    MPI::Datatype recvtypes[]) const = 0(binding deprecated, see
    Section 15.2) }
```

MPI\_ALLTOALLW is the most general form of complete exchange. Like MPI\_TYPE\_CREATE\_STRUCT, the most general type constructor, MPI\_ALLTOALLW allows separate specification of count, displacement and datatype. In addition, to allow maximum flexibility, the displacement of blocks within the send and receive buffers is specified in bytes.

If `comm` is an intracommunicator, then the  $j$ -th block sent from process  $i$  is received by process  $j$  and is placed in the  $i$ -th block of `recvbuf`. These blocks need not all have the same size.

The type signature associated with `sendcounts[j]`, `sendtypes[j]` at process  $i$  must be equal to the type signature associated with `recvcounst[i]`, `recvtypes[i]` at process  $j$ . This implies that the amount of data sent must be equal to the amount of data received, pairwise between every pair of processes. Distinct type maps between sender and receiver are still allowed.

The outcome is as if each process sent a message to every other process with

```
MPI_Send(sendbuf + sdispls[i], sendcounts[i], sendtypes[i], i, ...),
```

and received a message from every other process with a call to

```
MPI_Recv(recvbuf + rdispls[i], recvcounst[i], recvtypes[i], i, ...).
```

All arguments on all processes are significant. The argument `comm` must describe the same communicator on all processes.

Like for MPI\_ALLTOALLV, the “in place” option for intracommunicators is specified by passing MPI\_IN\_PLACE to the argument `sendbuf` at *all* processes. In such a case, `sendcounts`, `sdispls` and `sendtypes` are ignored. The data to be sent is taken from the `recvbuf` and replaced by the received data. Data sent and received must have the same type map as specified by the `recvcounst` and `recvtypes` arrays, and is taken from the locations of the receive buffer specified by `rdispls`.

If `comm` is an intercommunicator, then the outcome is as if each process in group A sends a message to each process in group B, and vice versa. The  $j$ -th send buffer of process  $i$  in group A should be consistent with the  $i$ -th receive buffer of process  $j$  in group B, and vice versa.

*Rationale.* The MPI\_ALLTOALLW function generalizes several MPI functions by carefully selecting the input arguments. For example, by making all but one process have `sendcounts[i] = 0`, this achieves an MPI\_SCATTERW function. (*End of rationale.*)

## 5.9 Global Reduction Operations

The functions in this section perform a global reduce operation (for example sum, maximum, and logical and) across all members of a group. The reduction operation can be either one of a predefined list of operations, or a user-defined operation. The global reduction functions come in several flavors: a reduce that returns the result of the reduction to one member of a group, an all-reduce that returns this result to all members of a group, and two scan (parallel

prefix) operations. In addition, a reduce-scatter operation combines the functionality of a reduce and of a scatter operation.

### 5.9.1 Reduce

**MPI\_REDUCE**(sendbuf, recvbuf, count, datatype, op, root, comm)

IN	sendbuf	address of send buffer (choice)
OUT	recvbuf	address of receive buffer (choice, significant only at root)
IN	count	number of elements in send buffer (non-negative integer)
IN	datatype	data type of elements of send buffer (handle)
IN	op	reduce operation (handle)
IN	root	rank of root process (integer)
IN	comm	communicator (handle)

```
int MPI_Reduce(const void* sendbuf, void* recvbuf, int count,
               MPI_Datatype datatype, MPI_Op op, int root, MPI_Comm comm)
```

```
MPI_Reduce(sendbuf, recvbuf, count, datatype, op, root, comm, ierror)
           BIND(C)
           TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
           TYPE(*), DIMENSION(..) :: recvbuf
           INTEGER, INTENT(IN) :: count, root
           TYPE(MPI_Datatype), INTENT(IN) :: datatype
           TYPE(MPI_Op), INTENT(IN) :: op
           TYPE(MPI_Comm), INTENT(IN) :: comm
           INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_REDUCE(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, ROOT, COMM, IERROR)
<type> SENDBUF(*), RECVBUF(*)
INTEGER COUNT, DATATYPE, OP, ROOT, COMM, IERROR
```

```
{void MPI::Comm::Reduce(const void* sendbuf, void* recvbuf, int count,
                        const MPI::Datatype& datatype, const MPI::Op& op, int root)
  const = 0(binding deprecated, see Section 15.2) }
```

If `comm` is an intracommunicator, **MPI\_REDUCE** combines the elements provided in the input buffer of each process in the group, using the operation `op`, and returns the combined value in the output buffer of the process with rank `root`. The input buffer is defined by the arguments `sendbuf`, `count` and `datatype`; the output buffer is defined by the arguments `recvbuf`, `count` and `datatype`; both have the same number of elements, with the same type. The routine is called by all group members using the same arguments for `count`, `datatype`, `op`, `root` and `comm`. Thus, all processes provide input buffers [and output buffers] of the same length, with elements of the same type[.] as the output buffer at the root. Each process can provide one element, or a sequence of elements, in which case the combine operation

is executed element-wise on each entry of the sequence. For example, if the operation is `MPI_MAX` and the send buffer contains two elements that are floating point numbers (`count = 2` and `datatype = MPI_FLOAT`), then `recvbuf(1) = global max(sendbuf(1))` and `recvbuf(2) = global max(sendbuf(2))`.

Section 5.9.2, lists the set of predefined operations provided by MPI. That section also enumerates the datatypes to which each operation can be applied.

In addition, users may define their own operations that can be overloaded to operate on several datatypes, either basic or derived. This is further explained in Section 5.9.5.

The operation `op` is always assumed to be associative. All predefined operations are also assumed to be commutative. Users may define operations that are assumed to be associative, but not commutative. The “canonical” evaluation order of a reduction is determined by the ranks of the processes in the group. However, the implementation can take advantage of associativity, or associativity and commutativity in order to change the order of evaluation. This may change the result of the reduction for operations that are not strictly associative and commutative, such as floating point addition.

*Advice to implementors.* It is strongly recommended that `MPI_REDUCE` be implemented so that the same result be obtained whenever the function is applied on the same arguments, appearing in the same order. Note that this may prevent optimizations that take advantage of the physical location of `[processors]ranks`. (*End of advice to implementors.*)

*Advice to users.* Some applications may not be able to ignore the non-associative nature of floating-point operations or may use user-defined operations (see Section 5.9.5) that require a special reduction order and cannot be treated as associative. Such applications should enforce the order of evaluation explicitly. For example, in the case of operations that require a strict left-to-right (or right-to-left) evaluation order, this could be done by gathering all operands at a single process (e.g., with `MPI_GATHER`), applying the reduction operation in the desired order (e.g., with `MPI_REDUCE_LOCAL`), and if needed, broadcast or scatter the result to the other processes (e.g., with `MPI_BCAST`). (*End of advice to users.*)

The `datatype` argument of `MPI_REDUCE` must be compatible with `op`. Predefined operators work only with the MPI types listed in Section 5.9.2 and Section 5.9.4. Furthermore, the `datatype` and `op` given for predefined operators must be the same on all processes.

Note that it is possible for users to supply different user-defined operations to `MPI_REDUCE` in each process. MPI does not define which operations are used on which operands in this case. User-defined operators may operate on general, derived datatypes. In this case, each argument that the reduce operation is applied to is one element described by such a datatype, which may contain several basic values. This is further explained in Section 5.9.5.

*Advice to users.* Users should make no assumptions about how `MPI_REDUCE` is implemented. It is safest to ensure that the same function is passed to `MPI_REDUCE` by each process. (*End of advice to users.*)

Overlapping datatypes are permitted in “send” buffers. Overlapping datatypes in “receive” buffers are erroneous and may give unpredictable results.

The “in place” option for intracommunicators is specified by passing the value `MPI_IN_PLACE` to the argument `sendbuf` at the root. In such a case, the input data is taken at the root from the receive buffer, where it will be replaced by the output data.

If `comm` is an intercommunicator, then the call involves all processes in the intercommunicator, but with one group (group A) defining the root process. All processes in the other group (group B) pass the same value in argument `root`, which is the rank of the root in group A. The root passes the value `MPI_ROOT` in `root`. All other processes in group A pass the value `MPI_PROC_NULL` in `root`. Only send buffer arguments are significant in group B and only receive buffer arguments are significant at the root.

## 5.9.2 Predefined Reduction Operations

The following predefined operations are supplied for `MPI_REDUCE` and related functions `MPI_ALLREDUCE`, `MPI_REDUCE_SCATTER_BLOCK`, `MPI_REDUCE_SCATTER`, `MPI_SCAN`, [and] `MPI_EXSCAN`, all nonblocking variants of those (see Section 5.12), and `MPI_REDUCE_LOCAL`. These operations are invoked by placing the following in `op`.

Name	Meaning
<code>MPI_MAX</code>	maximum
<code>MPI_MIN</code>	minimum
<code>MPI_SUM</code>	sum
<code>MPI_PROD</code>	product
<code>MPI_LAND</code>	logical and
<code>MPI_BAND</code>	bit-wise and
<code>MPI_LOR</code>	logical or
<code>MPI BOR</code>	bit-wise or
<code>MPI_LXOR</code>	logical exclusive or (xor)
<code>MPI_BXOR</code>	bit-wise exclusive or (xor)
<code>MPI_MAXLOC</code>	max value and location
<code>MPI_MINLOC</code>	min value and location

The two operations `MPI_MINLOC` and `MPI_MAXLOC` are discussed separately in Section 5.9.4. For the other predefined operations, we enumerate below the allowed combinations of `op` and `datatype` arguments. First, define groups of MPI basic datatypes in the following way.

C integer:	<code>MPI_INT</code> , <code>MPI_LONG</code> , <code>MPI_SHORT</code> , <code>MPI_UNSIGNED_SHORT</code> , <code>MPI_UNSIGNED</code> , <code>MPI_UNSIGNED_LONG</code> , <code>MPI_LONG_LONG_INT</code> , <code>MPI_LONG_LONG</code> (as synonym), <code>MPI_UNSIGNED_LONG_LONG</code> , <code>MPI_SIGNED_CHAR</code> , <code>MPI_UNSIGNED_CHAR</code> , <code>MPI_INT8_T</code> , <code>MPI_INT16_T</code> , <code>MPI_INT32_T</code> , <code>MPI_INT64_T</code> , <code>MPI_UINT8_T</code> , <code>MPI_UINT16_T</code> ,
------------	--------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------

ticket265. Fortran integer: MPI\_UINT32\_T, MPI\_UINT64\_T  
MPI\_INTEGER, MPI\_AINT, MPI\_COUNT,  
MPI\_OFFSET,  
and handles returned from  
MPI\_TYPE\_CREATE\_F90\_INTEGER,  
and if available: MPI\_INTEGER1,  
MPI\_INTEGER2, MPI\_INTEGER4,  
MPI\_INTEGER8, MPI\_INTEGER16  
Floating point: MPI\_FLOAT, MPI\_DOUBLE, MPI\_REAL,  
MPI\_DOUBLE\_PRECISION  
MPI\_LONG\_DOUBLE  
and handles returned from  
MPI\_TYPE\_CREATE\_F90\_REAL,  
and if available: MPI\_REAL2,  
MPI\_REAL4, MPI\_REAL8, MPI\_REAL16  
Logical: MPI\_LOGICAL, MPI\_C\_BOOL  
Complex: MPI\_COMPLEX,  
MPI\_C\_FLOAT\_COMPLEX,  
MPI\_C\_DOUBLE\_COMPLEX,  
MPI\_C\_LONG\_DOUBLE\_COMPLEX,  
and handles returned from  
MPI\_TYPE\_CREATE\_F90\_COMPLEX,  
and if available: MPI\_DOUBLE\_COMPLEX,  
MPI\_COMPLEX4, MPI\_COMPLEX8,  
MPI\_COMPLEX16, MPI\_COMPLEX32  
Byte: MPI\_BYTE  
Now, the valid datatypes for each option [is]are specified below.

Op	Allowed Types
MPI_MAX, MPI_MIN	C integer, Fortran integer, Floating point
MPI_SUM, MPI_PROD	C integer, Fortran integer, Floating point, Complex
MPI_LAND, MPI_LOR, MPI_LXOR	C integer, Logical
MPI_BAND, MPI_BOR, MPI_BXOR	C integer, Fortran integer, Byte

The following examples use intracommunicators.

### Example 5.15

A routine that computes the dot product of two vectors that are distributed across a group of processes and returns the answer at node zero.

```

1  SUBROUTINE PAR_BLAS1(m, a, b, c, comm)
2  REAL a(m), b(m)          ! local slice of array
3  REAL c                   ! result (at node zero)
4  REAL sum
5  INTEGER m, comm, i, ierr
6
7  ! local sum
8  sum = 0.0
9  DO i = 1, m
10     sum = sum + a(i)*b(i)
11 END DO
12
13 ! global sum
14 CALL MPI_REDUCE(sum, c, 1, MPI_REAL, MPI_SUM, 0, comm, ierr)
15 RETURN
16 END

```

### Example 5.16

A routine that computes the product of a vector and an array that are distributed across a group of processes and returns the answer at node zero.

```

22 SUBROUTINE PAR_BLAS2(m, n, a, b, c, comm)
23 REAL a(m), b(m,n)      ! local slice of array
24 REAL c(n)              ! result
25 REAL sum(n)
26 INTEGER n, comm, i, j, ierr
27
28 ! local sum
29 DO j= 1, n
30     sum(j) = 0.0
31     DO i = 1, m
32         sum(j) = sum(j) + a(i)*b(i,j)
33     END DO
34 END DO
35
36 ! global sum
37 CALL MPI_REDUCE(sum, c, n, MPI_REAL, MPI_SUM, 0, comm, ierr)
38
39 ! return result at node zero (and garbage at the other nodes)
40 RETURN
41 END

```

### 5.9.3 Signed Characters and Reductions

The types `MPI_SIGNED_CHAR` and `MPI_UNSIGNED_CHAR` can be used in reduction operations. `MPI_CHAR`, `MPI_WCHAR`, and `MPI_CHARACTER` (which represent printable characters) cannot be used in reduction operations. In a heterogeneous environment, `MPI_CHAR`, `MPI_WCHAR`, and `MPI_CHARACTER` will be translated so as to preserve the printable

character, whereas `MPI_SIGNED_CHAR` and `MPI_UNSIGNED_CHAR` will be translated so as to preserve the integer value.

*Advice to users.* The types `MPI_CHAR`, `MPI_WCHAR`, and `MPI_CHARACTER` are intended for characters, and so will be translated to preserve the printable representation, rather than the integer value, if sent between machines with different character codes. The types `MPI_SIGNED_CHAR` and `MPI_UNSIGNED_CHAR` should be used in C if the integer value should be preserved. (*End of advice to users.*)

#### 5.9.4 MINLOC and MAXLOC

The operator `MPI_MINLOC` is used to compute a global minimum and also an index attached to the minimum value. `MPI_MAXLOC` similarly computes a global maximum and index. One application of these is to compute a global minimum (maximum) and the rank of the process containing this value.

The operation that defines `MPI_MAXLOC` is:

$$\begin{pmatrix} u \\ i \end{pmatrix} \circ \begin{pmatrix} v \\ j \end{pmatrix} = \begin{pmatrix} w \\ k \end{pmatrix}$$

where

$$w = \max(u, v)$$

and

$$k = \begin{cases} i & \text{if } u > v \\ \min(i, j) & \text{if } u = v \\ j & \text{if } u < v \end{cases}$$

`MPI_MINLOC` is defined similarly:

$$\begin{pmatrix} u \\ i \end{pmatrix} \circ \begin{pmatrix} v \\ j \end{pmatrix} = \begin{pmatrix} w \\ k \end{pmatrix}$$

where

$$w = \min(u, v)$$

and

$$k = \begin{cases} i & \text{if } u < v \\ \min(i, j) & \text{if } u = v \\ j & \text{if } u > v \end{cases}$$

Both operations are associative and commutative. Note that if `MPI_MAXLOC` is applied to reduce a sequence of pairs  $(u_0, 0), (u_1, 1), \dots, (u_{n-1}, n-1)$ , then the value returned is  $(u, r)$ , where  $u = \max_i u_i$  and  $r$  is the index of the first global maximum in the sequence. Thus, if each process supplies a value and its rank within the group, then a reduce operation with `op = MPI_MAXLOC` will return the maximum value and the rank of the first process with that value. Similarly, `MPI_MINLOC` can be used to return a minimum and its index. More generally, `MPI_MINLOC` computes a *lexicographic minimum*, where elements are ordered

according to the first component of each pair, and ties are resolved according to the second component.

The reduce operation is defined to operate on arguments that consist of a pair: value and index. For both Fortran and C, types are provided to describe the pair. The potentially mixed-type nature of such arguments is a problem in Fortran. The problem is circumvented, for Fortran, by having the MPI-provided type consist of a pair of the same type as value, and coercing the index to this type also. In C, the MPI-provided pair type has distinct types and the index is an `int`.

In order to use `MPI_MINLOC` and `MPI_MAXLOC` in a reduce operation, one must provide a `datatype` argument that represents a pair (value and index). MPI provides nine such predefined datatypes. The operations `MPI_MAXLOC` and `MPI_MINLOC` can be used with each of the following datatypes.

Fortran:

Name	Description
<code>MPI_2REAL</code>	pair of <code>REAL</code> s
<code>MPI_2DOUBLE_PRECISION</code>	pair of <code>DOUBLE PRECISION</code> variables
<code>MPI_2INTEGER</code>	pair of <code>INTEGER</code> s

C:

Name	Description
<code>MPI_FLOAT_INT</code>	<code>float</code> and <code>int</code>
<code>MPI_DOUBLE_INT</code>	<code>double</code> and <code>int</code>
<code>MPI_LONG_INT</code>	<code>long</code> and <code>int</code>
<code>MPI_2INT</code>	pair of <code>int</code>
<code>MPI_SHORT_INT</code>	<code>short</code> and <code>int</code>
<code>MPI_LONG_DOUBLE_INT</code>	<code>long double</code> and <code>int</code>

The datatype `MPI_2REAL` is *as if* defined by the following (see Section 4.1).

```
MPI_TYPE_CONTIGUOUS(2, MPI_REAL, MPI_2REAL)
```

Similar statements apply for `MPI_2INTEGER`, `MPI_2DOUBLE_PRECISION`, and `MPI_2INT`.

The datatype `MPI_FLOAT_INT` is *as if* defined by the following sequence of instructions.

```
type[0] = MPI_FLOAT
type[1] = MPI_INT
disp[0] = 0
disp[1] = sizeof(float)
block[0] = 1
block[1] = 1
MPI_TYPE_CREATE_STRUCT(2, block, disp, type, MPI_FLOAT_INT)
```

Similar statements apply for `MPI_LONG_INT` and `MPI_DOUBLE_INT`.

The following examples use intracommunicators.

### Example 5.17

Each process has an array of 30 `doubles`, in C. For each of the 30 locations, compute the value and rank of the process containing the largest value.



```

...
/* each process has an array of 30 double: ain[30]
*/
double ain[30], aout[30];
int ind[30];
struct {
    double val;
    int rank;
} in[30], out[30];
int i, myrank, root;

MPI_Comm_rank(comm, &myrank);
for (i=0; i<30; ++i) {
    in[i].val = ain[i];
    in[i].rank = myrank;
}
MPI_Reduce(in, out, 30, MPI_DOUBLE_INT, MPI_MAXLOC, root, comm);
/* At this point, the answer resides on process root
*/
if (myrank == root) {
    /* read ranks out
    */
    for (i=0; i<30; ++i) {
        aout[i] = out[i].val;
        ind[i] = out[i].rank;
    }
}

```

**Example 5.18**

Same example, in Fortran.

```

...
! each process has an array of 30 double: ain(30)

DOUBLE PRECISION ain(30), aout(30)
INTEGER ind(30)
DOUBLE PRECISION in(2,30), out(2,30)
INTEGER i, myrank, root, ierr

CALL MPI_COMM_RANK(comm, myrank, ierr)
DO I=1, 30
    in(1,i) = ain(i)
    in(2,i) = myrank    ! myrank is coerced to a double
END DO

CALL MPI_REDUCE(in, out, 30, MPI_2DOUBLE_PRECISION, MPI_MAXLOC, root,
               comm, ierr)

```

```

1      ! At this point, the answer resides on process root
2
3      IF (myrank .EQ. root) THEN
4          ! read ranks out
5          DO I= 1, 30
6              aout(i) = out(1,i)
7              ind(i) = out(2,i) ! rank is coerced back to an integer
8          END DO
9      END IF

```

### Example 5.19

Each process has a non-empty array of values. Find the minimum global value, the rank of the process that holds it and its index on this process.

```

15 #define LEN 1000
16
17 float val[LEN];          /* local array of values */
18 int count;               /* local number of values */
19 int myrank, minrank, minindex;
20 float minval;
21
22 struct {
23     float value;
24     int index;
25 } in, out;
26
27     /* local minloc */
28 in.value = val[0];
29 in.index = 0;
30 for (i=1; i < count; i++)
31     if (in.value > val[i]) {
32         in.value = val[i];
33         in.index = i;
34     }
35
36     /* global minloc */
37 MPI_Comm_rank(comm, &myrank);
38 in.index = myrank*LEN + in.index;
39 MPI_Reduce( &in, &out, 1, MPI_FLOAT_INT, MPI_MINLOC, root, comm );
40     /* At this point, the answer resides on process root
41     */
42 if (myrank == root) {
43     /* read answer out
44     */
45     minval = out.value;
46     minrank = out.index / LEN;
47     minindex = out.index % LEN;
48 }

```

*Rationale.* The definition of MPI\_MINLOC and MPI\_MAXLOC given here has the advantage that it does not require any special-case handling of these two operations: they are handled like any other reduce operation. A programmer can provide his or her own definition of MPI\_MAXLOC and MPI\_MINLOC, if so desired. The disadvantage is that values and indices have to be first interleaved, and that indices and values have to be coerced to the same type, in Fortran. (*End of rationale.*)

### 5.9.5 User-Defined Reduction Operations

MPI\_OP\_CREATE([function]user\_fn, commute, op)

IN [ticket252-W.] [function]user\_fn user defined function (function)  
 IN commute true if commutative; false otherwise.  
 OUT op operation (handle)

int MPI\_Op\_create(MPI\_User\_function\* [function]user\_fn, int commute,  
 MPI\_Op\* op)

MPI\_Op\_create(user\_fn, commute, op, ierror) BIND(C)  
 PROCEDURE(MPI\_User\_function) :: user\_fn  
 LOGICAL, INTENT(IN) :: commute  
 TYPE(MPI\_Op), INTENT(OUT) :: op  
 INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI\_OP\_CREATE( [FUNCTION]USER\_FN, COMMUTE, OP, IERROR)  
 EXTERNAL [FUNCTION]USER\_FN  
 LOGICAL COMMUTE  
 INTEGER OP, IERROR

{void MPI::Op::Init(MPI::User\_function\* [function]user\_fn,  
 bool commute) (*binding deprecated, see Section 15.2*) }

MPI\_OP\_CREATE binds a user-defined reduction operation to an op handle that can subsequently be used in MPI\_REDUCE, MPI\_ALLREDUCE, MPI\_REDUCE\_SCATTER\_BLOCK, MPI\_REDUCE\_SCATTER, MPI\_SCAN, [MPI\_EXSCAN, all nonblocking variants of those (see Section 5.12), and MPI\_REDUCE\_LOCAL. The user-defined operation is assumed to be associative. If commute = true, then the operation should be both commutative and associative. If commute = false, then the order of operands is fixed and is defined to be in ascending, process rank order, beginning with process zero. The order of evaluation can be changed, taking advantage of the associativity of the operation. If commute = true then the order of evaluation can be changed, taking advantage of commutativity and associativity.

The argument [function]user\_fn is the user-defined function, which must have the following four arguments: invec, inoutvec, len and datatype.

The ISO C prototype for the function is the following.

typedef void MPI\_User\_function(void\* invec, void\* inoutvec, int \*len,  
 MPI\_Datatype \*datatype);

The Fortran declarations of the user-defined function `user_fn` appear[s] below.

**ABSTRACT INTERFACE**

```
SUBROUTINE MPI_User_function(invec, inoutvec, len, datatype) BIND(C)
    USE, INTRINSIC :: ISO_C_BINDING, ONLY : C_PTR
    TYPE(C_PTR), VALUE :: invec, inoutvec
    INTEGER :: len
    TYPE(MPI_Datatype) :: datatype
```

```
SUBROUTINE USER_FUNCTION(INVEC, INOUTVEC, LEN, [TYPE]DATATYPE)
    <type> INVEC(LEN), INOUTVEC(LEN)
    INTEGER LEN, [TYPE]DATATYPE
```

The C++ declaration of the user-defined function appears below.

```
{typedef void MPI::User_function(const void* invec, void* inoutvec, int
    len, const Datatype& datatype); (binding deprecated, see
    Section 15.2)}
```

The `datatype` argument is a handle to the data type that was passed into the call to `MPI_REDUCE`. The user reduce function should be written such that the following holds: Let `u[0], ... , u[len-1]` be the `len` elements in the communication buffer described by the arguments `invec`, `len` and `datatype` when the function is invoked; let `v[0], ... , v[len-1]` be `len` elements in the communication buffer described by the arguments `inoutvec`, `len` and `datatype` when the function is invoked; let `w[0], ... , w[len-1]` be `len` elements in the communication buffer described by the arguments `inoutvec`, `len` and `datatype` when the function returns; then `w[i] = u[i] ◦ v[i]`, for `i=0, ... , len-1`, where `◦` is the reduce operation that the function computes.

Informally, we can think of `invec` and `inoutvec` as arrays of `len` elements that `[function]user_fn` is combining. The result of the reduction over-writes values in `inoutvec`, hence the name. Each invocation of the function results in the pointwise evaluation of the reduce operator on `len` elements: i.e., the function returns in `inoutvec[i]` the value `invec[i] ◦ inoutvec[i]`, for `i = 0, ... , count - 1`, where `◦` is the combining operation computed by the function.

*Rationale.* The `len` argument allows `MPI_REDUCE` to avoid calling the function for each element in the input buffer. Rather, the system can choose to apply the function to chunks of input. In C, it is passed in as a reference for reasons of compatibility with Fortran.

By internally comparing the value of the `datatype` argument to known, global handles, it is possible to overload the use of a single user-defined function for several, different data types. (*End of rationale.*)

General datatypes may be passed to the user function. However, use of datatypes that are not contiguous is likely to lead to inefficiencies.

No MPI communication function may be called inside the user function. `MPI_ABORT` may be called inside the function in case of an error.

*Advice to users.* Suppose one defines a library of user-defined reduce functions that are overloaded: the `datatype` argument is used to select the right execution path at each invocation, according to the types of the operands. The user-defined reduce function

cannot “decode” the `datatype` argument that it is passed, and cannot identify, by itself, the correspondence between the datatype handles and the datatype they represent. This correspondence was established when the datatypes were created. Before the library is used, a library initialization preamble must be executed. This preamble code will define the datatypes that are used by the library, and store handles to these datatypes in global, static variables that are shared by the user code and the library code.

The Fortran version of `MPI_REDUCE` will invoke a user-defined reduce function using the Fortran calling conventions and will pass a Fortran-type datatype argument; the C version will use C calling convention and the C representation of a datatype handle. Users who plan to mix languages should define their reduction functions accordingly. (*End of advice to users.*)

*Advice to implementors.* We outline below a naive and inefficient implementation of `MPI_REDUCE` not supporting the “in place” option.

```

MPI_Comm_size(comm, &groupsize);
MPI_Comm_rank(comm, &rank);
if (rank > 0) {
    MPI_Recv(tempbuf, count, datatype, rank-1,...);
    User_reduce(tempbuf, sendbuf, count, datatype);
}
if (rank < groupsize-1) {
    MPI_Send(sendbuf, count, datatype, rank+1, ...);
}
/* answer now resides in process groupsize-1 ... now send to root
*/
if (rank == root) {
    MPI_Irecv(recvbuf, count, datatype, groupsize-1,..., &req);
}
if (rank == groupsize-1) {
    MPI_Send(sendbuf, count, datatype, root, ...);
}
if (rank == root) {
    MPI_Wait(&req, &status);
}

```

The reduction computation proceeds, sequentially, from process 0 to process `groupsize-1`. This order is chosen so as to respect the order of a possibly non-commutative operator defined by the function `User_reduce()`. A more efficient implementation is achieved by taking advantage of associativity and using a logarithmic tree reduction. Commutativity can be used to advantage, for those cases in which the `commute` argument to `MPI_OP_CREATE` is true. Also, the amount of temporary buffer required can be reduced, and communication can be pipelined with computation, by transferring and reducing the elements in chunks of size `len < count`.

The predefined reduce operations can be implemented as a library of user-defined operations. However, better performance might be achieved if `MPI_REDUCE` handles these functions as a special case. (*End of advice to implementors.*)

```

1 MPI_OP_FREE(op)
2     INOUT    op                operation (handle)
3

```

```

4
5 int MPI_Op_free(MPI_Op *op)
6 MPI_Op_free(op, ierror) BIND(C)
7     TYPE(MPI_Op), INTENT(INOUT) :: op
8     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
9
10 MPI_OP_FREE(OP, IERROR)
11     INTEGER OP, IERROR

```

```

12 {void MPI::Op::Free() (binding deprecated, see Section 15.2) }
13

```

Marks a user-defined reduction operation for deallocation and sets `op` to `MPI_OP_NULL`.

#### Example of User-defined Reduce

It is time for an example of user-defined reduction. The example in this section uses an intracommunicator.

**Example 5.20** Compute the product of an array of complex numbers, in C.

```

21
22 typedef struct {
23     double real,imag;
24 } Complex;
25
26 /* the user-defined function
27 */
28 void myProd(void *inP, void *inoutP, int *len, MPI_Datatype *dptr)
29 {
30     int i;
31     Complex c;
32     Complex *in = (Complex *)inP, *inout = (Complex *)inoutP;
33
34     for (i=0; i< *len; ++i) {
35         c.real = inout->real*in->real -
36             inout->imag*in->imag;
37         c.imag = inout->real*in->imag +
38             inout->imag*in->real;
39         *inout = c;
40         in++; inout++;
41     }
42 }
43
44 /* and, to call it...
45 */
46 ...
47
48 /* each process has an array of 100 Complexes

```

```

    */
    Complex a[100], answer[100];
    MPI_Op myOp;
    MPI_Datatype ctype;

    /* explain to MPI how type Complex is defined
    */
    MPI_Type_contiguous(2, MPI_DOUBLE, &ctype);
    MPI_Type_commit(&ctype);
    /* create the complex-product user-op
    */
    MPI_Op_create( myProd, 1, &myOp );

    MPI_Reduce(a, answer, 100, ctype, myOp, root, comm);

    /* At this point, the answer, which consists of 100 Complexes,
    * resides on process root
    */

```

**Example 5.21** How to use the `mpi_f08` interface of the Fortran `MPI_User_function`.

```

subroutine my_user_function( invec, inoutvec, len, type )    bind(c)
  use, intrinsic :: iso_c_binding, only : c_ptr, c_f_pointer
  type(c_ptr), value :: invec, inoutvec
  integer :: len
  type(MPI_Datatype) :: type
  real, pointer :: invec_r(:), inoutvec_r(:)
  if (type%MPI_VAL == MPI_REAL%MPI_VAL) then
    call c_f_pointer(invec, invec_r, (/ len /) )
    call c_f_pointer(inoutvec, inoutvec_r, (/ len /) )
    inoutvec_r = invec_r + inoutvec_r
  end if
end subroutine

```

### 5.9.6 All-Reduce

MPI includes a variant of the reduce operations where the result is returned to all processes in a group. MPI requires that all processes from the same group participating in these operations receive identical results.

1	MPI_ALLREDUCE(sendbuf, recvbuf, count, datatype, op, comm)		
2	IN	sendbuf	starting address of send buffer (choice)
3	OUT	recvbuf	starting address of receive buffer (choice)
4	IN	count	number of elements in send buffer (non-negative integer)
5	IN	datatype	data type of elements of send buffer (handle)
6	IN	op	operation (handle)
7	IN	comm	communicator (handle)

```

12 int MPI_Allreduce(const void* sendbuf, void* recvbuf, int count,
13                 MPI_Datatype datatype, MPI_Op op, MPI_Comm comm)
14
15 MPI_Allreduce(sendbuf, recvbuf, count, datatype, op, comm, ierror) BIND(C)
16     TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
17     TYPE(*), DIMENSION(..) :: recvbuf
18     INTEGER, INTENT(IN) :: count
19     TYPE(MPI_Datatype), INTENT(IN) :: datatype
20     TYPE(MPI_Op), INTENT(IN) :: op
21     TYPE(MPI_Comm), INTENT(IN) :: comm
22     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
23
24 MPI_ALLREDUCE(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, COMM, IERROR)
25     <type> SENDBUF(*), RECVBUF(*)
26     INTEGER COUNT, DATATYPE, OP, COMM, IERROR
27
28 {void MPI::Comm::Allreduce(const void* sendbuf, void* recvbuf, int count,
29                           const MPI::Datatype& datatype, const MPI::Op& op)
30                           const = 0(binding deprecated, see Section 15.2) }

```

If comm is an intracommunicator, MPI\_ALLREDUCE behaves the same as MPI\_REDUCE except that the result appears in the receive buffer of all the group members.

*Advice to implementors.* The all-reduce operations can be implemented as a reduce, followed by a broadcast. However, a direct implementation can lead to better performance. (*End of advice to implementors.*)

The “in place” option for intracommunicators is specified by passing the value MPI\_IN\_PLACE to the argument sendbuf at all processes. In this case, the input data is taken at each process from the receive buffer, where it will be replaced by the output data.

If comm is an intercommunicator, then the result of the reduction of the data provided by processes in group A is stored at each process in group B, and vice versa. Both groups should provide count and datatype arguments that specify the same type signature.

The following example uses an intracommunicator.

### Example 5.22

A routine that computes the product of a vector and an array that are distributed across a group of processes and returns the answer at all nodes (see also Example 5.16).



```

SUBROUTINE PAR_BLAS2(m, n, a, b, c, comm)
REAL a(m), b(m,n)      ! local slice of array
REAL c(n)              ! result
REAL sum(n)
INTEGER n, comm, i, j, ierr

! local sum
DO j= 1, n
  sum(j) = 0.0
  DO i = 1, m
    sum(j) = sum(j) + a(i)*b(i,j)
  END DO
END DO

! global sum
CALL MPI_ALLREDUCE(sum, c, n, MPI_REAL, MPI_SUM, comm, ierr)

! return result at all nodes
RETURN
END

```

### 5.9.7 Process-[l]Local [r]Reduction

The functions in this section are of importance to library implementors who may want to implement special reduction patterns that are otherwise not easily covered by the standard MPI operations.

The following function applies a reduction operator to local arguments.

**MPI\_REDUCE\_LOCAL**( inbuf, inoutbuf, count, datatype, op)

IN	inbuf	input buffer (choice)
INOUT	inoutbuf	combined input and output buffer (choice)
IN	count	number of elements in inbuf and inoutbuf buffers (non-negative integer)
IN	datatype	data type of elements of inbuf and inoutbuf buffers (handle)
IN	op	operation (handle)

```

int MPI_Reduce_local(const void* inbuf, void* inoutbuf, int count,
                    MPI_Datatype datatype, MPI_Op op)

```

```

MPI_Reduce_local(inbuf, inoutbuf, count, datatype, op, ierror) BIND(C)
  TYPE(*), DIMENSION(..), INTENT(IN) :: inbuf
  TYPE(*), DIMENSION(..) :: inoutbuf
  INTEGER, INTENT(IN) :: count
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  TYPE(MPI_Op), INTENT(IN) :: op

```

```

1      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
2
3 MPI_REDUCE_LOCAL(INBUF, INOUTBUF, COUNT, DATATYPE, OP, IERROR)
4     <type> INBUF(*), INOUTBUF(*)
5     INTEGER COUNT, DATATYPE, OP, IERROR
6
7 {void MPI::Op::Reduce_local(const void* inbuf, void* inoutbuf, int count,
8     const MPI::Datatype& datatype) const(binding deprecated, see
9     Section 15.2) }

```

The function applies the operation given by `op` element-wise to the elements of `inbuf` and `inoutbuf` with the result stored element-wise in `inoutbuf`, as explained for user-defined operations in Section 5.9.5. Both `inbuf` and `inoutbuf` (input as well as result) have the same number of elements given by `count` and the same datatype given by `datatype`. The `MPI_IN_PLACE` option is not allowed.

Reduction operations can be queried for their commutativity.

```

17 MPI_OP_COMMUTATIVE( op, commute)
18
19 IN      op      operation (handle)
20 OUT    commute  true if op is commutative, false otherwise (logical)
21
22
23 int MPI_Op_commutative(MPI_Op op, int *commute)
24
25 MPI_Op_commutative(op, commute, ierror) BIND(C)
26     TYPE(MPI_Op), INTENT(IN) :: op
27     LOGICAL, INTENT(OUT) :: commute
28     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
29
30 MPI_OP_COMMUTATIVE(OP, COMMUTE, IERROR)
31     LOGICAL COMMUTE
32     INTEGER OP, IERROR
33
34 {bool MPI::Op::Is_commutative() const(binding deprecated, see Section 15.2) }

```

## 5.10 Reduce-Scatter

MPI includes variants of the reduce operations where the result is scattered to all processes in a group on return. One variant scatters equal-sized blocks to all processes, while another variant scatters blocks that may vary in size for each process.

## 5.10.1 MPI\_REDUCE\_SCATTER\_BLOCK

MPI_REDUCE_SCATTER_BLOCK( sendbuf, recvbuf, recvcnt, datatype, op, comm)		
IN	sendbuf	starting address of send buffer (choice)
OUT	recvbuf	starting address of receive buffer (choice)
IN	recvcnt	element count per block (non-negative integer)
IN	datatype	data type of elements of send and receive buffers (handle)
IN	op	operation (handle)
IN	comm	communicator (handle)

```
int MPI_Reduce_scatter_block(const void* sendbuf, void* recvbuf,
                           int recvcnt, MPI_Datatype datatype, MPI_Op op,
                           MPI_Comm comm)
```

```
MPI_Reduce_scatter_block(sendbuf, recvbuf, recvcnt, datatype, op, comm,
                          ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
    TYPE(*), DIMENSION(..) :: recvbuf
    INTEGER, INTENT(IN) :: recvcnt
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Op), INTENT(IN) :: op
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_REDUCE_SCATTER_BLOCK(SENDBUF, RECVBUF, REVCOUNT, DATATYPE, OP, COMM,
                          IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER REVCOUNT, DATATYPE, OP, COMM, IERROR
```

```
{void MPI::Comm::Reduce_scatter_block(const void* sendbuf, void* recvbuf,
                                     int recvcnt, const MPI::Datatype& datatype,
                                     const MPI::Op& op) const = 0(binding deprecated, see Section 15.2) }
```

If `comm` is an intracommunicator, `MPI_REDUCE_SCATTER_BLOCK` first performs a global, element-wise reduction on vectors of `count = n*recvcnt` elements in the send buffers defined by `sendbuf`, `count` and `datatype`, using the operation `op`, where `n` is the number of processes in the group of `comm`. The routine is called by all group members using the same arguments for `recvcnt`, `datatype`, `op` and `comm`. The resulting vector is treated as `n` consecutive blocks of `recvcnt` elements that are scattered to the processes of the group. The `i`-th block is sent to process `i` and stored in the receive buffer defined by `recvbuf`, `recvcnt`, and `datatype`.

*Advice to implementors.* The `MPI_REDUCE_SCATTER_BLOCK` routine is functionally equivalent to: an `MPI_REDUCE` collective operation with `count` equal to `recvcnt*n`, followed by an `MPI_SCATTER` with `sendcount` equal to `recvcnt`. However, a direct implementation may run faster. (*End of advice to implementors.*)

ticket0. 1 The “in place” option for intracommunicators is specified by passing `MPI_IN_PLACE`  
 2 in the `sendbuf` argument on *all* processes. In this case, the input data is taken from the  
 3 receive buffer.

4 If `comm` is an intercommunicator, then the result of the reduction of the data provided  
 5 by processes in one group (group A) is scattered among processes in the other group (group  
 6 B) and vice versa. Within each group, all processes provide the same value for the `recvcount`  
 7 argument, and provide input vectors of `count = n*recvcount` elements stored in the send  
 8 buffers, where `n` is the size of the group. The number of elements `count` must be the same  
 9 for the two groups. The resulting vector from the other group is scattered in blocks of  
 10 `recvcount` elements among the processes in the group.

11  
 12 *Rationale.* The last restriction is needed so that the length of the send buffer of  
 13 one group can be determined by the local `recvcount` argument of the other group.  
 14 Otherwise, a communication is needed to figure out how many elements are reduced.  
 15 (*End of rationale.*)

### 17 5.10.2 MPI\_REDUCE\_SCATTER

18 `MPI_REDUCE_SCATTER` extends the functionality of `MPI_REDUCE_SCATTER_BLOCK`  
 19 such that the scattered blocks can vary in size. Block sizes are determined by the `recvcounts`  
 20 array, such that the *i*-th block contains `recvcounts[i]` elements.

21  
 22  
 23 `MPI_REDUCE_SCATTER( sendbuf, recvbuf, recvcounts, datatype, op, comm)`

24	IN	<code>sendbuf</code>	starting address of send buffer (choice)
25	OUT	<code>recvbuf</code>	starting address of receive buffer (choice)
26	IN	<code>recvcounts</code>	non-negative integer array (of length group size) specifying the number of elements of the result distributed to each process.
27	IN	<code>datatype</code>	data type of elements of send and receive buffers (handle)
28	IN	<code>op</code>	operation (handle)
29	IN	<code>comm</code>	communicator (handle)

30  
 31  
 32  
 33  
 34  
 35  
 36  
 37  
 38  
 39  
 40  
 41  
 42  
 43  
 44  
 45  
 46  
 47  
 48

```

int MPI_Reduce_scatter(const void* sendbuf, void* recvbuf, const
    int recvcounts[], MPI_Datatype datatype, MPI_Op op,
    MPI_Comm comm)
MPI_Reduce_scatter(sendbuf, recvbuf, recvcounts, datatype, op, comm,
    ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
TYPE(*), DIMENSION(..) :: recvbuf
INTEGER, INTENT(IN) :: recvcounts(*)
TYPE(MPI_Datatype), INTENT(IN) :: datatype
TYPE(MPI_Op), INTENT(IN) :: op
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_REDUCE_SCATTER(SENDBUF, RECVBUFF, RECVCOUNTS, DATATYPE, OP, COMM,
                    IERROR)
<type> SENDBUF(*), RECVBUFF(*)
INTEGER RECVCOUNTS(*), DATATYPE, OP, COMM, IERROR
{void MPI::Comm::Reduce_scatter(const void* sendbuf, void* recvbuf,
    int recvcounts[], const MPI::Datatype& datatype,
    const MPI::Op& op) const = 0(binding deprecated, see Section 15.2) }
```

If `comm` is an intracommunicator, `MPI_REDUCE_SCATTER` first performs a global, element-wise reduction on vectors of `count =  $\sum_{i=0}^{n-1} \text{recvcounts}[i]$`  elements in the send buffers defined by `sendbuf`, `count` and `datatype`, using the operation `op`, where `n` is the number of processes in the group of `comm`. The routine is called by all group members using the same arguments for `recvcounts`, `datatype`, `op` and `comm`. The resulting vector is treated as `n` consecutive blocks where the number of elements of the `i`-th block is `recvcounts[i]`. The blocks are scattered to the processes of the group. The `i`-th block is sent to process `i` and stored in the receive buffer defined by `recvbuf`, `recvcounts[i]` and `datatype`.

*Advice to implementors.* The `MPI_REDUCE_SCATTER` routine is functionally equivalent to: an `MPI_REDUCE` collective operation with `count` equal to the sum of `recvcounts[i]` followed by `MPI_SCATTERV` with `sendcounts` equal to `recvcounts`. However, a direct implementation may run faster. (*End of advice to implementors.*)

The “in place” option for intracommunicators is specified by passing `MPI_IN_PLACE` in the `sendbuf` argument. In this case, the input data is taken from the receive buffer. It is not required to specify the “in place” option on all processes, since the processes for which `recvcounts[i]==0` may not have allocated a receive buffer.

If `comm` is an intercommunicator, then the result of the reduction of the data provided by processes in one group (group A) is scattered among processes in the other group (group B), and vice versa. Within each group, all processes provide the same `recvcounts` argument, and provide input vectors of `count =  $\sum_{i=0}^{n-1} \text{recvcounts}[i]$`  elements stored in the send buffers, where `n` is the size of the group. The resulting vector from the other group is scattered in blocks of `recvcounts[i]` elements among the processes in the group. The number of elements `count` must be the same for the two groups.

*Rationale.* The last restriction is needed so that the length of the send buffer can be determined by the sum of the local `recvcounts` entries. Otherwise, a communication is needed to figure out how many elements are reduced. (*End of rationale.*)

## 5.11 Scan

### 5.11.1 Inclusive Scan

`MPI_SCAN(sendbuf, recvbuf, count, datatype, op, comm)`

IN	<code>sendbuf</code>	starting address of send buffer (choice)
OUT	<code>recvbuf</code>	starting address of receive buffer (choice)
IN	<code>count</code>	number of elements in input buffer (non-negative integer)
IN	<code>datatype</code>	data type of elements of input buffer (handle)
IN	<code>op</code>	operation (handle)
IN	<code>comm</code>	communicator (handle)

```
int MPI_Scan(const void* sendbuf, void* recvbuf, int count,
             MPI_Datatype datatype, MPI_Op op, MPI_Comm comm)
```

```
MPI_Scan(sendbuf, recvbuf, count, datatype, op, comm, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
    TYPE(*), DIMENSION(..) :: recvbuf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Op), INTENT(IN) :: op
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_SCAN(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER COUNT, DATATYPE, OP, COMM, IERROR
```

```
{void MPI::Intracomm::Scan(const void* sendbuf, void* recvbuf, int count,
                           const MPI::Datatype& datatype, const MPI::Op& op) const(binding
                           deprecated, see Section 15.2) }
```

If `comm` is an intracommunicator, `MPI_SCAN` is used to perform a prefix reduction on data distributed across the group. The operation returns, in the receive buffer of the process with rank `i`, the reduction of the values in the send buffers of processes with ranks `0, ..., i` (inclusive). The type of operations supported, their semantics, and the constraints on send and receive buffers are as for `MPI_REDUCE`.

The “in place” option for intracommunicators is specified by passing `MPI_IN_PLACE` in the `sendbuf` argument. In this case, the input data is taken from the receive buffer, and replaced by the output data.

This operation is invalid for intercommunicators.

## 5.11.2 Exclusive Scan

MPI\_EXSCAN(sendbuf, recvbuf, count, datatype, op, comm)

IN	sendbuf	starting address of send buffer (choice)
OUT	recvbuf	starting address of receive buffer (choice)
IN	count	number of elements in input buffer (non-negative integer)
IN	datatype	data type of elements of input buffer (handle)
IN	op	operation (handle)
IN	comm	intracommunicator (handle)

```
int MPI_Exscan(const void* sendbuf, void* recvbuf, int count,
               MPI_Datatype datatype, MPI_Op op, MPI_Comm comm)
```

```
MPI_Exscan(sendbuf, recvbuf, count, datatype, op, comm, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
    TYPE(*), DIMENSION(..) :: recvbuf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Op), INTENT(IN) :: op
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_EXSCAN(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER COUNT, DATATYPE, OP, COMM, IERROR
```

```
{void MPI::Intracomm::Exscan(const void* sendbuf, void* recvbuf, int count,
                             const MPI::Datatype& datatype, const MPI::Op& op) const(binding
                             deprecated, see Section 15.2) }
```

If `comm` is an intracommunicator, `MPI_EXSCAN` is used to perform a prefix reduction on data distributed across the group. The value in `recvbuf` on the process with rank 0 is undefined, and `recvbuf` is not significant on process 0. The value in `recvbuf` on the process with rank 1 is defined as the value in `sendbuf` on the process with rank 0. For processes with rank  $i > 1$ , the operation returns, in the receive buffer of the process with rank  $i$ , the reduction of the values in the send buffers of processes with ranks  $0, \dots, i - 1$  (inclusive). The type of operations supported, their semantics, and the constraints on send and receive buffers, are as for `MPI_REDUCE`.

The “in place” option for intracommunicators is specified by passing `MPI_IN_PLACE` in the `sendbuf` argument. In this case, the input data is taken from the receive buffer, and replaced by the output data. The receive buffer on rank 0 is not changed by this operation.

This operation is invalid for intercommunicators.

*Rationale.* The exclusive scan is more general than the inclusive scan. Any inclusive scan operation can be achieved by using the exclusive scan and then locally combining

the local contribution. Note that for non-invertable operations such as MPI\_MAX, the exclusive scan cannot be computed with the inclusive scan. (*End of rationale.*)

### 5.11.3 Example using MPI\_SCAN

The example in this section uses an intracommunicator.

#### Example 5.23

This example uses a user-defined operation to produce a *segmented scan*. A segmented scan takes, as input, a set of values and a set of logicals, and the logicals delineate the various segments of the scan. For example:

<i>values</i>	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$	$v_7$	$v_8$
<i>logicals</i>	0	0	1	1	1	0	0	1
<i>result</i>	$v_1$	$v_1 + v_2$	$v_3$	$v_3 + v_4$	$v_3 + v_4 + v_5$	$v_6$	$v_6 + v_7$	$v_8$

The operator that produces this effect is,

$$\begin{pmatrix} u \\ i \end{pmatrix} \circ \begin{pmatrix} v \\ j \end{pmatrix} = \begin{pmatrix} w \\ j \end{pmatrix},$$

where,

$$w = \begin{cases} u + v & \text{if } i = j \\ v & \text{if } i \neq j \end{cases}.$$

Note that this is a non-commutative operator. C code that implements it is given below.

```

typedef struct {
    double val;
    int log;
} SegScanPair;

/* the user-defined function
*/
void segScan(SegScanPair *in, SegScanPair *inout, int *len,
             MPI_Datatype *dptr)
{
    int i;
    SegScanPair c;

    for (i=0; i< *len; ++i) {
        if (in->log == inout->log)
            c.val = in->val + inout->val;
        else
            c.val = inout->val;
        c.log = inout->log;
        *inout = c;
        in++; inout++;
    }
}

```



Note that the `inout` argument to the user-defined function corresponds to the right-hand operand of the operator. When using this operator, we must be careful to specify that it is non-commutative, as in the following.

```

int i,base;
SegScanPair  a, answer;
MPI_Op      myOp;
MPI_Datatype type[2] = {MPI_DOUBLE, MPI_INT};
MPI_Aint     disp[2];
int          blocklen[2] = { 1, 1};
MPI_Datatype sspair;

/* explain to MPI how type SegScanPair is defined
 */
MPI_Get_address( &a, disp);
MPI_Get_address( &a.log, disp+1);
base = disp[0];
for (i=0; i<2; ++i) disp[i] -= base;
MPI_Type_create_struct( 2, blocklen, disp, type, &sspair );
MPI_Type_commit( &sspair );
/* create the segmented-scan user-op
 */
MPI_Op_create(segScan, 0, &myOp);
...
MPI_Scan( &a, &answer, 1, sspair, myOp, comm );

```

ticket109.

## 5.12 Nonblocking Collective Operations

As described in Section 3.7, performance of many applications can be improved by overlapping communication and computation, and many systems enable this. Nonblocking collective operations combine the potential benefits of nonblocking point-to-point operations, to exploit overlap and to avoid synchronization, with the optimized implementation and message scheduling provided by collective operations [30, 34]. One way of doing this would be to perform a blocking collective operation in a separate thread. An alternative mechanism that often leads to better performance (e.g., avoids context switching, scheduler overheads, and thread management) is to use nonblocking collective communication [32].

The nonblocking collective communication model is similar to the model used for nonblocking point-to-point communication. A nonblocking call initiates a collective operation, which must be completed in a separate completion call. Once initiated, the operation may progress independently of any computation or other communication at participating processes. In this manner, nonblocking collective operations can mitigate possible synchronizing effects of collective operations by running them in the “background.” In addition to enabling communication-computation overlap, nonblocking collective operations can perform collective operations on overlapping communicators, which would lead to deadlocks with blocking operations. Their semantic advantages can also be useful in combination with point-to-point communication.

As in the nonblocking point-to-point case, all calls are local and return immediately, irrespective of the status of other processes. The call initiates the operation, which indicates

that the system may start to copy data out of the send buffer and into the receive buffer. Once initiated, all associated send buffers and buffers associated with input arguments (such as arrays of counts, displacements, or datatypes in the vector versions of the collectives) should not be modified, and all associated receive buffers should not be accessed, until the collective operation completes. The call returns a request handle, which must be passed to a completion call.

All completion calls (e.g., `MPI_WAIT`) described in Section 3.7.3 are supported for nonblocking collective operations. Similarly to the blocking case, nonblocking collective operations are considered to be complete when the local part of the operation is finished, i.e., for the caller, the semantics of the operation are guaranteed and all buffers can be safely accessed and modified. Completion does not indicate that other processes have completed or even started the operation (unless otherwise implied by the description of the operation). Completion of a particular nonblocking collective operation also does not indicate completion of any other posted nonblocking collective (or send-receive) operations, whether they are posted before or after the completed operation.

*Advice to users.* Users should be aware that implementations are allowed, but not required (with exception of `MPI_IBARRIER`), to synchronize processes during the completion of a nonblocking collective operation. (*End of advice to users.*)

Upon returning from a completion call in which a nonblocking collective operation completes, the `MPI_ERROR` field in the associated status object is set appropriately, see Section 3.2.5 on page 34. The values of the `MPI_SOURCE` and `MPI_TAG` fields are undefined. It is valid to mix different request types (i.e., any combination of collective requests, I/O requests, generalized requests, or point-to-point requests) in functions that enable multiple completions (e.g., `MPI_WAITALL`). It is erroneous to call `MPI_REQUEST_FREE` or `MPI_CANCEL` for a request associated with a nonblocking collective operation. Nonblocking collective requests are not persistent.

*Rationale.* Freeing an active nonblocking collective request could cause similar problems as discussed for point-to-point requests (see Section 3.7.3). Cancelling a request is not supported because the semantics of this operation are not well-defined. (*End of rationale.*)

Multiple nonblocking collective operations can be outstanding on a single communicator. If the nonblocking call causes some system resource to be exhausted, then it will fail and generate an MPI exception. Quality implementations of MPI should ensure that this happens only in pathological cases. That is, an MPI implementation should be able to support a large number of pending nonblocking operations.

Unlike point-to-point operations, nonblocking collective operations do not match with blocking collective operations, and collective operations do not have a tag argument. All processes must call collective operations (blocking and nonblocking) in the same order per communicator. In particular, once a process calls a collective operation, all other processes in the communicator must eventually call the same collective operation, and no other collective operation with the same communicator in between. This is consistent with the ordering rules for blocking collective operations in threaded environments.

*Rationale.* Matching blocking and nonblocking collective operations is not allowed because the implementation might use different communication algorithms for the two

cases. Blocking collective operations may be optimized for minimal time to completion, while nonblocking collective operations may balance time to completion with CPU overhead and asynchronous progression.

The use of tags for collective operations can prevent certain hardware optimizations. (*End of rationale.*)

*Advice to users.* If program semantics require matching blocking and nonblocking collective operations, then a nonblocking collective operation can be initiated and immediately completed with a blocking wait to emulate blocking behavior. (*End of advice to users.*)

In terms of data movements, each nonblocking collective operation has the same effect as its blocking counterpart for intracommunicators and intercommunicators after completion. Likewise, upon completion, nonblocking collective reduction operations have the same effect as their blocking counterparts, and the same restrictions and recommendations on reduction orders apply.

The use of the “in place” option is allowed exactly as described for the corresponding blocking collective operations. When using the “in place” option, message buffers function as both send and receive buffers. Such buffers should not be modified or accessed until the operation completes.

Progression rules for nonblocking collective operations are similar to progression of nonblocking point-to-point operations, refer to Section 3.7.4.

*Advice to implementors.* Nonblocking collective operations can be implemented with local execution schedules [33] using nonblocking point-to-point communication and a reserved tag-space. (*End of advice to implementors.*)

### 5.12.1 Nonblocking Barrier Synchronization

`MPI_IBARRIER(comm , request)`

IN	<code>comm</code>	communicator (handle)
OUT	<code>request</code>	communication request (handle)

`int MPI_Ibarrier(MPI_Comm comm, MPI_Request *request)`

`MPI_Ibarrier(comm, request, ierror) BIND(C)`

`TYPE(MPI_Comm), INTENT(IN) :: comm`

`TYPE(MPI_Request), INTENT(OUT) :: request`

`INTEGER, OPTIONAL, INTENT(OUT) :: ierror`

`MPI_IBARRIER(COMM, REQUEST, IERROR)`

`INTEGER COMM, REQUEST, IERROR`

[ {`MPI::Request MPI::Comm::Ibarrier()` `const = 0` (*binding deprecated, see Section 15.2*) } ]

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1       ]

2       MPI\_IBARRIER is a nonblocking version of MPI\_BARRIER. By calling MPI\_IBARRIER,

3       a process notifies that it has reached the barrier. The call returns immediately, indepen-

4       dent of whether other processes have called MPI\_IBARRIER. The usual barrier semantics

5       are enforced at the corresponding completion operation (test or wait), which in the intra-

6       communicator case will complete only after all other processes in the communicator have

7       called MPI\_IBARRIER. In the intercommunicator case, it will complete when all processes

8       in the remote group have called MPI\_IBARRIER.

9

10       *Advice to users.* A nonblocking barrier can be used to hide latency. Moving indepen-

11       dent computations between the MPI\_IBARRIER and the subsequent completion call

12       can overlap the barrier latency and therefore shorten possible waiting times. The se-

13       mantic properties are also useful when mixing collective operations and point-to-point

14       messages. (*End of advice to users.*)

### 15       5.12.2 Nonblocking Broadcast

16

17

18

19       MPI\_IBCAST(buffer, count, datatype, root, comm, request)

20	INOUT	buffer	starting address of buffer (choice)
21	IN	count	number of entries in buffer (non-negative integer)
22	IN	datatype	data type of buffer (handle)
23	IN	root	rank of broadcast root (integer)
24	IN	comm	communicator (handle)
25	OUT	request	communication request (handle)

26

27

28

29       int MPI\_Ibcast(const void\* buffer, int count, MPI\_Datatype datatype,

30                   int root, MPI\_Comm comm, MPI\_Request \*request)

31

32       MPI\_Ibcast(buffer, count, datatype, root, comm, request, ierror) BIND(C)

33       TYPE(\*), DIMENSION(..), ASYNCHRONOUS :: buffer

34       INTEGER, INTENT(IN) :: count, root

35       TYPE(MPI\_Datatype), INTENT(IN) :: datatype

36       TYPE(MPI\_Comm), INTENT(IN) :: comm

37       TYPE(MPI\_Request), INTENT(OUT) :: request

38       INTEGER, OPTIONAL, INTENT(OUT) :: ierror

39

40       MPI\_IBCAST(BUFFER, COUNT, DATATYPE, ROOT, COMM, REQUEST, IERROR)

41       <type> BUFFER(\*)

42       INTEGER COUNT, DATATYPE, ROOT, COMM, REQUEST, IERROR

43       [ {MPI::Request MPI::Comm::Ibcast(void\* buffer, int count,

44           const MPI::Datatype& datatype, int root) const = 0(*binding*

45           *deprecated, see Section 15.2*) }

46

47       ]

48       This call starts a nonblocking variant of MPI\_BCAST (see Section 5.4).

## Example using MPI\_IBCAST

The example in this section uses an intracommunicator.

**Example 5.24**

Start a broadcast of 100 ints from process 0 to every process in the group, perform some computation on independent data, and then complete the outstanding broadcast operation.

```
MPI_Comm comm;
int array1[100], array2[100];
int root=0;
MPI_Request req;
...
MPI_Ibcast(array1, 100, MPI_INT, root, comm, &req);
compute(array2, 100);
MPI_Wait(&req, MPI_STATUS_IGNORE);
```

## 5.12.3 Nonblocking Gather

`MPI_IGATHER(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype, root, comm, request)`

IN	sendbuf	starting address of send buffer (choice)
IN	sendcount	number of elements in send buffer (non-negative integer)
IN	sendtype	data type of send buffer elements (handle)
OUT	recvbuf	address of receive buffer (choice, significant only at root)
IN	recvcount	number of elements for any single receive (non-negative integer, significant only at root)
IN	recvtype	data type of recv buffer elements (significant only at root) (handle)
IN	root	rank of receiving process (integer)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```
int MPI_Igather(const void* sendbuf, int sendcount, MPI_Datatype sendtype,
               void* recvbuf, int recvcount, MPI_Datatype recvtype, int root,
               MPI_Comm comm, MPI_Request *request)
```

```
MPI_Igather(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype,
            root, comm, request, ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
INTEGER, INTENT(IN) :: sendcount, recvcount, root
TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
```

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```

1      TYPE(MPI_Comm), INTENT(IN) :: comm
2      TYPE(MPI_Request), INTENT(OUT) :: request
3      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
4
5      MPI_IGATHER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT, RECVTYPE,
6                  ROOT, COMM, REQUEST, IERROR)
7      <type> SENDBUF(*), RECVBUF(*)
8      INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, ROOT, COMM, REQUEST,
9      IERROR

```

```

10     [ {MPI::Request MPI::Comm::Igather(const void* sendbuf, int sendcount,
11                                         const MPI::Datatype& sendtype, void* recvbuf, int recvcount,
12                                         const MPI::Datatype& recvtype, int root) const = 0(binding
13                                             deprecated, see Section 15.2) }
14
15 ]

```

This call starts a nonblocking variant of MPI\_GATHER (see Section 5.5).

```

18      MPI_IGATHERV(sendbuf, sendcount, sendtype, recvbuf, recvcounts, displs, recvtype, root,
19                  comm, request)

```

IN	sendbuf	starting address of send buffer (choice)
IN	sendcount	number of elements in send buffer (non-negative integer)
IN	sendtype	data type of send buffer elements (handle)
OUT	recvbuf	address of receive buffer (choice, significant only at root)
IN	recvcounts	non-negative integer array (of length group size) containing the number of elements that are received from each process (significant only at root)
IN	displs	integer array (of length group size). Entry <i>i</i> specifies the displacement relative to <i>recvbuf</i> at which to place the incoming data from process <i>i</i> (significant only at root)
IN	recvtype	data type of recv buffer elements (significant only at root) (handle)
IN	root	rank of receiving process (integer)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```

42      int MPI_Igatherv(const void* sendbuf, int sendcount, MPI_Datatype sendtype,
43                      void* recvbuf, const int recvcounts[], const int displs[],
44                      MPI_Datatype recvtype, int root, MPI_Comm comm,
45                      MPI_Request *request)

```

```

46      MPI_Igatherv(sendbuf, sendcount, sendtype, recvbuf, recvcounts, displs,
47                  recvtype, root, comm, request, ierror) BIND(C)

```

```

TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
INTEGER, INTENT(IN) :: sendcount, root
INTEGER, INTENT(IN), ASYNCHRONOUS :: recvcounts(*), displs(*)
TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Request), INTENT(OUT) :: request
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_IGATHERV(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNTS, DISPLS,
             RECVTYPE, ROOT, COMM, REQUEST, IERROR)
<type> SENDBUF(*), RECVBUF(*)
INTEGER SENDCOUNT, SENDTYPE, RECVCOUNTS(*), DISPLS(*), RECVTYPE, ROOT,
COMM, REQUEST, IERROR

[ {MPI::Request MPI::Comm::Igatherv(const void* sendbuf, int sendcount,
    const MPI::Datatype& sendtype, void* recvbuf,
    const int recvcounts[], const int displs[],
    const MPI::Datatype& recvtype, int root) const = 0(binding
    deprecated, see Section 15.2) }
]

```

This call starts a nonblocking variant of MPI\_GATHERV (see Section 5.5).

#### 5.12.4 Nonblocking Scatter

```

MPI_ISCATTER(sendbuf, sendcount, sendtype, recvbuf, recvcnt, recvtype, root, comm,
             request)

```

IN	sendbuf	address of send buffer (choice, significant only at root)
IN	sendcount	number of elements sent to each process (non-negative integer, significant only at root)
IN	sendtype	data type of send buffer elements (significant only at root) (handle)
OUT	recvbuf	address of receive buffer (choice)
IN	recvcnt	number of elements in receive buffer (non-negative integer)
IN	recvtype	data type of receive buffer elements (handle)
IN	root	rank of sending process (integer)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```

int MPI_Iscatter(const void* sendbuf, int sendcount, MPI_Datatype sendtype,
                void* recvbuf, int recvcnt, MPI_Datatype recvtype, int root,
                MPI_Comm comm, MPI_Request *request)

```

```

1 MPI_Iscatter(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype,
2             root, comm, request, ierror) BIND(C)
3     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
4     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
5     INTEGER, INTENT(IN) :: sendcount, recvcount, root
6     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
7     TYPE(MPI_Comm), INTENT(IN) :: comm
8     TYPE(MPI_Request), INTENT(OUT) :: request
9     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

10 MPI_ISCATTER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT, RECVTYPE,
11             ROOT, COMM, REQUEST, IERROR)
12     <type> SENDBUF(*), RECVBUF(*)
13     INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, ROOT, COMM, REQUEST,
14     IERROR

```

```

15     [ {MPI::Request MPI::Comm::Iscatter(const void* sendbuf, int sendcount,
16         const MPI::Datatype& sendtype, void* recvbuf, int recvcount,
17         const MPI::Datatype& recvtype, int root) const = 0(binding
18         deprecated, see Section 15.2) }
19
20 ]

```

This call starts a nonblocking variant of MPI\_SCATTER (see Section 5.6).

```

21
22
23
24 MPI_ISCATTERV(sendbuf, sendcounts, displs, sendtype, recvbuf, recvcount, recvtype, root,
25             comm, request)

```

IN	sendbuf	address of send buffer (choice, significant only at root)
IN	sendcounts	non-negative integer array (of length group size) specifying the number of elements to send to each <a href="#">processor</a> rank
IN	displs	integer array (of length group size). Entry <i>i</i> specifies the displacement (relative to <code>sendbuf</code> ) from which to take the outgoing data to process <i>i</i>
IN	sendtype	data type of send buffer elements (handle)
OUT	recvbuf	address of receive buffer (choice)
IN	recvcount	number of elements in receive buffer (non-negative integer)
IN	recvtype	data type of receive buffer elements (handle)
IN	root	rank of sending process (integer)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```

44
45 int MPI_Iscatterv(const void* sendbuf, const int sendcounts[], const
46                 int displs[], MPI_Datatype sendtype, void* recvbuf,
47                 int recvcount, MPI_Datatype recvtype, int root, MPI_Comm comm,
48                 MPI_Request *request)

```



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```

MPI_Iscatterv(sendbuf, sendcounts, displs, sendtype, recvbuf, recvcount,
              recvtype, root, comm, request, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN), ASYNCHRONOUS :: sendcounts(*), displs(*)
    INTEGER, INTENT(IN) :: recvcount, root
    TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_ISCATTERV(SENDBUF, SENDCOUNTS, DISPLS, SENDTYPE, RECVBUF, RECVCOUNT,
              RECVTYPE, ROOT, COMM, REQUEST, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNTS(*), DISPLS(*), SENDTYPE, RECVCOUNT, RECVTYPE, ROOT,
    COMM, REQUEST, IERROR

[ {MPI::Request MPI::Comm::Iscatterv(const void* sendbuf,
    const int sendcounts[], const int displs[],
    const MPI::Datatype& sendtype, void* recvbuf, int recvcount,
    const MPI::Datatype& recvtype, int root) const = 0(binding
    deprecated, see Section 15.2) }
]

```

This call starts a nonblocking variant of MPI\_SCATTERV (see Section 5.6).

### 5.12.5 Nonblocking Gather-to-all

```

MPI_IALLGATHER(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype, comm,
               request)

```

IN	sendbuf	starting address of send buffer (choice)
IN	sendcount	number of elements in send buffer (non-negative integer)
IN	sendtype	data type of send buffer elements (handle)
OUT	recvbuf	address of receive buffer (choice)
IN	recvcount	number of elements received from any process (non-negative integer)
IN	recvtype	data type of receive buffer elements (handle)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```

int MPI_Iallgather(const void* sendbuf, int sendcount,
                  MPI_Datatype sendtype, void* recvbuf, int recvcount,
                  MPI_Datatype recvtype, MPI_Comm comm, MPI_Request *request)

```

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```

1 MPI_Iallgather(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype,
2               comm, request, ierror) BIND(C)
3
4     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
5     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
6     INTEGER, INTENT(IN) :: sendcount, recvcount
7     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
8     TYPE(MPI_Comm), INTENT(IN) :: comm
9     TYPE(MPI_Request), INTENT(OUT) :: request
10    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

11 MPI_IALLGATHER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT, RECVTYPE,
12               COMM, REQUEST, IERROR)
13
14 <type> SENDBUF(*), RECVBUF(*)
15 INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, COMM, REQUEST, IERROR

```

```

16 [ {MPI::Request MPI::Comm::Iallgather(const void* sendbuf,
17   int sendcount, const MPI::Datatype& sendtype, void* recvbuf,
18   int recvcount, const MPI::Datatype& recvtype) const = 0(binding
19   deprecated, see Section 15.2) }

```

```

20 ]

```

This call starts a nonblocking variant of MPI\_ALLGATHER (see Section 5.7).

```

23
24 MPI_IALLGATHERV(sendbuf, sendcount, sendtype, recvbuf, recvcounts, displs, recvtype, comm,
25                request)

```

26	IN	sendbuf	starting address of send buffer (choice)
27	IN	sendcount	number of elements in send buffer (non-negative integer)
28	IN	sendtype	data type of send buffer elements (handle)
29	OUT	recvbuf	address of receive buffer (choice)
30	IN	recvcounts	non-negative integer array (of length group size) containing the number of elements that are received from each process
31	IN	displs	integer array (of length group size). Entry <i>i</i> specifies the displacement (relative to <i>recvbuf</i> ) at which to place the incoming data from process <i>i</i>
32	IN	recvtype	data type of receive buffer elements (handle)
33	IN	comm	communicator (handle)
34	OUT	request	communication request (handle)

```

35
36
37
38
39
40
41
42
43
44 int MPI_Iallgatherv(const void* sendbuf, int sendcount,
45                    MPI_Datatype sendtype, void* recvbuf, const int recvcounts[],
46                    const int displs[], MPI_Datatype recvtype, MPI_Comm comm,
47                    MPI_Request* request)

```

```

MPI_Iallgatherv(sendbuf, sendcount, sendtype, recvbuf, recvcoun- 1
ts, displs,                                           2
                recvtype, comm, request, ierror) BIND(C) 3
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf 4
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf        5
    INTEGER, INTENT(IN) :: sendcount                      6
    INTEGER, INTENT(IN), ASYNCHRONOUS :: recvcoun-         7
ts(*), displs(*)
    TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype 8
    TYPE(MPI_Comm), INTENT(IN) :: comm                    9
    TYPE(MPI_Request), INTENT(OUT) :: request             10
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror              11

```

```

MPI_IALLGATHERV(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECV- 12
COUNTS, DISPLS,
                RECVTYPE, COMM, REQUEST, IERROR)          13
<type> SENDBUF(*), RECVBUF(*)                             14
INTEGER SENDCOUNT, SENDTYPE, REVCOUNTS(*), DISPLS(*), RECVTYPE, COMM, 15
REQUEST, IERROR

```

```

[ {MPI::Request MPI::Comm::Iallgatherv(const void* sendbuf, 16
    int sendcount, const MPI::Datatype& sendtype, void* recvbuf, 17
    const int recvcoun-
ts[], const int displs[], 18
    const MPI::Datatype& recvtype) const = 0(binding deprecated, see 19
    Section 15.2) }

```

```

]

```

This call starts a nonblocking variant of MPI\_ALLGATHERV (see Section 5.7).

### 5.12.6 Nonblocking All-to-All Scatter/Gather

```

MPI_IALLTOALL(sendbuf, sendcount, sendtype, recvbuf, recvcoun- 29
t, recvtype, comm, request)

```

IN	sendbuf	starting address of send buffer (choice)
IN	sendcount	number of elements sent to each process (non-negative integer)
IN	sendtype	data type of send buffer elements (handle)
OUT	recvbuf	address of receive buffer (choice)
IN	recvcoun-	number of elements received from any process (non-
	t	negative integer)
IN	recvtype	data type of receive buffer elements (handle)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```

int MPI_Ialltoall(const void* sendbuf, int sendcount, 45
    MPI_Datatype sendtype, void* recvbuf, int recvcoun- 46
t, MPI_Datatype recvtype, MPI_Comm comm, MPI_Request *request) 47

```

```

1 MPI_Ialltoall(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype,
2               comm, request, ierror) BIND(C)
3     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
4     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
5     INTEGER, INTENT(IN) :: sendcount, recvcount
6     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
7     TYPE(MPI_Comm), INTENT(IN) :: comm
8     TYPE(MPI_Request), INTENT(OUT) :: request
9     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

10 MPI_IALLTOALL(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT, RECVTYPE,
11              COMM, REQUEST, IERROR)
12     <type> SENDBUF(*), RECVBUF(*)
13     INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, COMM, REQUEST, IERROR
14     [ {MPI::Request MPI::Comm::Ialltoall(const void* sendbuf, int sendcount,
15                                         const MPI::Datatype& sendtype, void* recvbuf, int recvcount,
16                                         const MPI::Datatype& recvtype) const = 0 (binding deprecated, see
17                                         Section 15.2) }
18
19 ]

```

This call starts a nonblocking variant of MPI\_ALLTOALL (see Section 5.8).

```

23 MPI_IALLTOALLV(sendbuf, sendcounts, sdispls, sendtype, recvbuf, recvcounts, rdispls,
24               recvtype, comm, request)

```

IN	sendbuf	starting address of send buffer (choice)
IN	sendcounts	non-negative integer array (of length group size) specifying the number of elements to send to each [processor]rank
IN	sdispls	integer array (of length group size). Entry j specifies the displacement (relative to sendbuf) from which to take the outgoing data destined for process j
IN	sendtype	data type of send buffer elements (handle)
OUT	recvbuf	address of receive buffer (choice)
IN	recvcounts	non-negative integer array (of length group size) specifying the number of elements that can be received from each [processor]rank
IN	rdispls	integer array (of length group size). Entry i specifies the displacement (relative to recvbuf) at which to place the incoming data from process i
IN	recvtype	data type of receive buffer elements (handle)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```

47 int MPI_Ialltoallv(const void* sendbuf, const int sendcounts[], const
48                   int sdispls[], MPI_Datatype sendtype, void* recvbuf, const

```

```

        int recvcnts[], const int rdispls[], MPI_Datatype recvtype,
        MPI_Comm comm, MPI_Request *request)
MPI_Ialltoallv(sendbuf, sendcounts, sdispls, sendtype, recvbuf, recvcnts,
        rdispls, recvtype, comm, request, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN), ASYNCHRONOUS :: sendcounts(*), sdispls(*),
    recvcnts(*), rdispls(*)
    TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_IALLTOALLV(SENDBUF, SENDCOUNTS, SDISPLS, SENDTYPE, RECVBUF, RECVCOUNTS,
        RDISPLS, RECVTYPE, COMM, REQUEST, IERROR)
<type> SENDBUF(*), RECVBUF(*)
INTEGER SENDCOUNTS(*), SDISPLS(*), SENDTYPE, RECVCOUNTS(*), RDISPLS(*),
RECVTYPE, COMM, REQUEST, IERROR

[ {MPI::Request MPI::Comm::Ialltoallv(const void* sendbuf,
        const int sendcounts[], const int sdispls[],
        const MPI::Datatype& sendtype, void* recvbuf,
        const int recvcnts[], const int rdispls[],
        const MPI::Datatype& recvtype) const = 0(binding deprecated, see
        Section 15.2) }

```

This call starts a nonblocking variant of MPI\_ALLTOALLV (see Section 5.8).

```

1  MPI_IALLTOALLW(sendbuf, sendcounts, sdispls, sendtypes, recvbuf, recvcounts, rdispls,
2                  recvtypes, comm, request)
3
4  IN      sendbuf      starting address of send buffer (choice)
5
6  IN      sendcounts   integer array (of length group size) specifying the num-
7                      ber of elements to send to each [processor]rank (array
8                      of non-negative integers)
9
10     IN      sdispls   integer array (of length group size). Entry j specifies
11                     the displacement in bytes (relative to sendbuf) from
12                     which to take the outgoing data destined for process
13                     j (array of integers)
14
15     IN      sendtypes array of datatypes (of length group size). Entry j
16                     specifies the type of data to send to process j (array
17                     of handles)
18
19     OUT     recvbuf   address of receive buffer (choice)
20
21     IN      recvcounts integer array (of length group size) specifying the num-
22                     ber of elements that can be received from each [processor]rank
23                     (array of non-negative integers)
24
25     IN      rdispls   integer array (of length group size). Entry i specifies
26                     the displacement in bytes (relative to recvbuf) at which
27                     to place the incoming data from process i (array of
28                     integers)
29
30     IN      recvtypes array of datatypes (of length group size). Entry i
31                     specifies the type of data received from process i (ar-
32                     ray of handles)
33
34     IN      comm      communicator (handle)
35
36     OUT     request   communication request (handle)

```

```

ticket140. 31  int MPI_Ialltoallw(const void* sendbuf, const int sendcounts[], const
ticket140. 32                      int sdispls[], const MPI_Datatype sendtypes[], void* recvbuf,
ticket140. 33                      const int recvcounts[], const int rdispls[], const
ticket140. 34                      MPI_Datatype recvtypes[], MPI_Comm comm, MPI_Request *request)
ticket140. 35
ticket140. 36  MPI_Ialltoallw(sendbuf, sendcounts, sdispls, sendtypes, recvbuf,
ticket140. 37                      recvcounts, rdispls, recvtypes, comm, request, ierror) BIND(C)
ticket229.1. 38  TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
ticket-248T. 39  TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
40  INTEGER, INTENT(IN), ASYNCHRONOUS :: sendcounts(*), sdispls(*),
41  recvcounts(*), rdispls(*)
42  TYPE(MPI_Datatype), INTENT(IN), ASYNCHRONOUS :: sendtypes(*),
43  recvtypes(*)
44  TYPE(MPI_Comm), INTENT(IN) :: comm
45  TYPE(MPI_Request), INTENT(OUT) :: request
46  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
47
48

```

```

MPI_IALLTOALLW(SENDBUF, SENDCOUNTS, SDISPLS, SENDTYPES, RECVBUF,
               RECVCOUNTS, RDISPLS, RECVTYPES, COMM, REQUEST, IERROR)
<type> SENDBUF(*), RECVBUF(*)
INTEGER SENDCOUNTS(*), SDISPLS(*), SENDTYPES(*), RECVCOUNTS(*),
RDISPLS(*), RECVTYPES(*), COMM, REQUEST, IERROR

```

```

[ {MPI::Request MPI::Comm::Ialltoallw(const void* sendbuf, const int
    sendcounts[], const int sdispls[], const MPI::Datatype
    sendtypes[], void* recvbuf, const int recvcounst[], const int
    rdispls[], const MPI::Datatype recvtypes[]) const = 0(binding
    deprecated, see Section 15.2) }

```

```

]

```

This call starts a nonblocking variant of MPI\_ALLTOALLW (see Section 5.8).

### 5.12.7 Nonblocking Reduce

```

MPI_IREDUCE(sendbuf, recvbuf, count, datatype, op, root, comm, request)

```

IN	sendbuf	address of send buffer (choice)
OUT	recvbuf	address of receive buffer (choice, significant only at root)
IN	count	number of elements in send buffer (non-negative integer)
IN	datatype	data type of elements of send buffer (handle)
IN	op	reduce operation (handle)
IN	root	rank of root process (integer)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```

int MPI_Ireduce(const void* sendbuf, void* recvbuf, int count,
               MPI_Datatype datatype, MPI_Op op, int root, MPI_Comm comm,
               MPI_Request *request)

```

```

MPI_Ireduce(sendbuf, recvbuf, count, datatype, op, root, comm, request,
            ierror) BIND(C)

```

```

TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
INTEGER, INTENT(IN) :: count, root
TYPE(MPI_Datatype), INTENT(IN) :: datatype
TYPE(MPI_Op), INTENT(IN) :: op
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Request), INTENT(OUT) :: request
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_IREDUCE(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, ROOT, COMM, REQUEST,
            IERROR)

```

```

1      <type> SENDBUF(*), RECVBUF(*)
2      INTEGER COUNT, DATATYPE, OP, ROOT, COMM, REQUEST, IERROR
3
4      [ {MPI::Request MPI::Comm::Ireduce(const void* sendbuf, void* recvbuf,
5          int count, const MPI::Datatype& datatype, const MPI::Op& op,
6          int root) const = 0(binding deprecated, see Section 15.2) }
7
8      ]

```

This call starts a nonblocking variant of `MPI_REDUCE` (see Section 5.9.1).

*Advice to implementors.* The implementation is explicitly allowed to use different algorithms for blocking and nonblocking reduction operations that might change the order of evaluation of the operations. However, as for `MPI_REDUCE`, it is strongly recommended that `MPI_IREDUCE` be implemented so that the same result be obtained whenever the function is applied on the same arguments, appearing in the same order. Note that this may prevent optimizations that take advantage of the physical location of processes. (*End of advice to implementors.*)

*Advice to users.* For operations which are not truly associative, the result delivered upon completion of the nonblocking reduction may not exactly equal the result delivered by the blocking reduction, even when specifying the same arguments in the same order. (*End of advice to users.*)

### 5.12.8 Nonblocking All-Reduce

`MPI_IALLREDUCE(sendbuf, recvbuf, count, datatype, op, comm, request)`

IN	sendbuf	starting address of send buffer (choice)
OUT	recvbuf	starting address of receive buffer (choice)
IN	count	number of elements in send buffer (non-negative integer)
IN	datatype	data type of elements of send buffer (handle)
IN	op	operation (handle)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```

38      int MPI_Iallreduce(const void* sendbuf, void* recvbuf, int count,
39          MPI_Datatype datatype, MPI_Op op, MPI_Comm comm,
40          MPI_Request *request)

```

```

41      MPI_Iallreduce(sendbuf, recvbuf, count, datatype, op, comm, request,
42          ierror) BIND(C)
43      TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
44      TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
45      INTEGER, INTENT(IN) :: count
46      TYPE(MPI_Datatype), INTENT(IN) :: datatype
47      TYPE(MPI_Op), INTENT(IN) :: op

```



```

TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Request), INTENT(OUT) :: request
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_IALLREDUCE(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, COMM, REQUEST,
               IERROR)
<type> SENDBUF(*), RECVBUF(*)
INTEGER COUNT, DATATYPE, OP, COMM, REQUEST, IERROR

[ {MPI::Request MPI::Comm::Iallreduce(const void* sendbuf,
    void* recvbuf, int count, const MPI::Datatype& datatype,
    const MPI::Op& op) const = 0(binding deprecated, see Section 15.2) }
]

```

This call starts a nonblocking variant of MPI\_ALLREDUCE (see Section 5.9.6).

ticket272.

### 5.12.9 Nonblocking Reduce-Scatter with Equal Blocks

```
MPI_IREDUCE_SCATTER_BLOCK(sendbuf, recvbuf, recvcnt, datatype, op, comm, request)
```

IN	sendbuf	starting address of send buffer (choice)
OUT	recvbuf	starting address of receive buffer (choice)
IN	recvcnt	element count per block (non-negative integer)
IN	datatype	data type of elements of send and receive buffers (handle)
IN	op	operation (handle)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```

int MPI_Ireduce_scatter_block(const void* sendbuf, void* recvbuf,
    int recvcnt, MPI_Datatype datatype, MPI_Op op,
    MPI_Comm comm, MPI_Request *request)

```

ticket140.

```

MPI_Ireduce_scatter_block(sendbuf, recvbuf, recvcnt, datatype, op, comm,
    request, ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
INTEGER, INTENT(IN) :: recvcnt
TYPE(MPI_Datatype), INTENT(IN) :: datatype
TYPE(MPI_Op), INTENT(IN) :: op
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Request), INTENT(OUT) :: request
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_IREDUCE_SCATTER_BLOCK(SENDBUF, RECVBUF, RECVCOUNT, DATATYPE, OP, COMM,
    REQUEST, IERROR)

```

ticket-248T.

```

1      <type> SENDBUF(*), RECVBUF(*)
2      INTEGER RECVCOUNT, DATATYPE, OP, COMM, REQUEST, IERROR
3
4      [ {MPI::Request MPI::Comm::Ireduce_scatter_block(const void* sendbuf,
5              void* recvbuf, int recvcnt, const MPI::Datatype& datatype,
6              const MPI::Op& op) const = 0(binding deprecated, see Section 15.2) }
7
8      ]

```

This call starts a nonblocking variant of MPI\_REDUCE\_SCATTER\_BLOCK (see Section 5.10.1).

### 5.12.10 Nonblocking Reduce-Scatter

```

14 MPI_IREDUCE_SCATTER(sendbuf, recvbuf, recvcnt, datatype, op, comm, request)
15
16     IN      sendbuf      starting address of send buffer (choice)
17     OUT     recvbuf      starting address of receive buffer (choice)
18     IN      recvcnt      non-negative integer array specifying the number of
19                        elements in result distributed to each process. Array
20                        must be identical on all calling processes.
21
22     IN      datatype     data type of elements of input buffer (handle)
23     IN      op           operation (handle)
24     IN      comm         communicator (handle)
25     OUT     request      communication request (handle)
26
27

```

```

28 int MPI_Ireduce_scatter(const void* sendbuf, void* recvbuf, const
29     int recvcnt[], MPI_Datatype datatype, MPI_Op op,
30     MPI_Comm comm, MPI_Request *request)
31
32

```

```

33 MPI_Ireduce_scatter(sendbuf, recvbuf, recvcnt, datatype, op, comm,
34     request, ierror) BIND(C)
35     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
36     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
37     INTEGER, INTENT(IN), ASYNCHRONOUS :: recvcnt(*)
38     TYPE(MPI_Datatype), INTENT(IN) :: datatype
39     TYPE(MPI_Op), INTENT(IN) :: op
40     TYPE(MPI_Comm), INTENT(IN) :: comm
41     TYPE(MPI_Request), INTENT(OUT) :: request
42     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
43
44

```

```

45 MPI_IREDUCE_SCATTER(SENDBUF, RECVBUF, RECVCOUNTS, DATATYPE, OP, COMM,
46     REQUEST, IERROR)
47     <type> SENDBUF(*), RECVBUF(*)
48     INTEGER RECVCOUNTS(*), DATATYPE, OP, COMM, REQUEST, IERROR
49
50     [ {MPI::Request MPI::Comm::Ireduce_scatter(const void* sendbuf,
51         void* recvbuf, int recvcnt[],

```

```

        const MPI::Datatype& datatype, const MPI::Op& op)
        const = 0(binding deprecated, see Section 15.2) }
    ]

```

This call starts a nonblocking variant of MPI\_REDUCE\_SCATTER (see Section 5.10.2).

### 5.12.11 Nonblocking Inclusive Scan

MPI\_ISCAN(sendbuf, recvbuf, count, datatype, op, comm, request)

IN	sendbuf	starting address of send buffer (choice)
OUT	recvbuf	starting address of receive buffer (choice)
IN	count	number of elements in input buffer (non-negative integer)
IN	datatype	data type of elements of input buffer (handle)
IN	op	operation (handle)
IN	comm	communicator (handle)
OUT	request	communication request (handle)

```

int MPI_Iscan(const void* sendbuf, void* recvbuf, int count,
              MPI_Datatype datatype, MPI_Op op, MPI_Comm comm,
              MPI_Request *request)

```

```

MPI_Iscan(sendbuf, recvbuf, count, datatype, op, comm, request, ierror)
    BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Op), INTENT(IN) :: op
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_ISCAN(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, COMM, REQUEST, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER COUNT, DATATYPE, OP, COMM, REQUEST, IERROR

```

```

    [ {MPI::Request MPI::Intracomm::Iscale(const void* sendbuf,
        void* recvbuf, int count, const MPI::Datatype& datatype,
        const MPI::Op& op) const(binding deprecated, see Section 15.2) }
    ]

```

This call starts a nonblocking variant of MPI\_SCAN (see Section 5.11).

### 5.12.12 Nonblocking Exclusive Scan

**MPI\_IEXSCAN**(sendbuf, recvbuf, count, datatype, op, comm, request)

IN	sendbuf	starting address of send buffer (choice)
OUT	recvbuf	starting address of receive buffer (choice)
IN	count	number of elements in input buffer (non-negative integer)
IN	datatype	data type of elements of input buffer (handle)
IN	op	operation (handle)
IN	comm	intracommunicator (handle)
OUT	request	communication request (handle)

```
int MPI_Iexscan(const void* sendbuf, void* recvbuf, int count,
               MPI_Datatype datatype, MPI_Op op, MPI_Comm comm,
               MPI_Request *request)
```

```
MPI_Iexscan(sendbuf, recvbuf, count, datatype, op, comm, request, ierror)
    BIND(C)
```

```
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Op), INTENT(IN) :: op
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_IEXSCAN(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, COMM, REQUEST, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER COUNT, DATATYPE, OP, COMM, REQUEST, IERROR
```

```
[ {MPI::Request MPI::Intracomm::Iexscan(const void* sendbuf, void*
    recvbuf, int count, const MPI::Datatype& datatype, const
    MPI::Op& op) const (binding deprecated, see Section 15.2) }
```

```
]
```

This call starts a nonblocking variant of **MPI\_EXSCAN** (see Section 5.11.2).

## 5.13 Correctness

A correct, portable program must invoke collective communications so that deadlock will not occur, whether collective communications are synchronizing or not. The following examples illustrate dangerous use of collective routines on intracommunicators.

### Example 5.25

The following is erroneous.

```

switch(rank) {
    case 0:
        MPI_Bcast(buf1, count, type, 0, comm);
        MPI_Bcast(buf2, count, type, 1, comm);
        break;
    case 1:
        MPI_Bcast(buf2, count, type, 1, comm);
        MPI_Bcast(buf1, count, type, 0, comm);
        break;
}

```

We assume that the group of `comm` is  $\{0,1\}$ . Two processes execute two broadcast operations in reverse order. If the operation is synchronizing then a deadlock will occur.

Collective operations must be executed in the same order at all members of the communication group.

### Example 5.26

The following is erroneous.

```

switch(rank) {
    case 0:
        MPI_Bcast(buf1, count, type, 0, comm0);
        MPI_Bcast(buf2, count, type, 2, comm2);
        break;
    case 1:
        MPI_Bcast(buf1, count, type, 1, comm1);
        MPI_Bcast(buf2, count, type, 0, comm0);
        break;
    case 2:
        MPI_Bcast(buf1, count, type, 2, comm2);
        MPI_Bcast(buf2, count, type, 1, comm1);
        break;
}

```

Assume that the group of `comm0` is  $\{0,1\}$ , of `comm1` is  $\{1, 2\}$  and of `comm2` is  $\{2,0\}$ . If the broadcast is a synchronizing operation, then there is a cyclic dependency: the broadcast in `comm2` completes only after the broadcast in `comm0`; the broadcast in `comm0` completes only after the broadcast in `comm1`; and the broadcast in `comm1` completes only after the broadcast in `comm2`. Thus, the code will deadlock.

Collective operations must be executed in an order so that no cyclic dependencies occur. Nonblocking collective operations can alleviate this issue.

### Example 5.27

The following is erroneous.

```

1  switch(rank) {
2      case 0:
3          MPI_Bcast(buf1, count, type, 0, comm);
4          MPI_Send(buf2, count, type, 1, tag, comm);
5          break;
6      case 1:
7          MPI_Recv(buf2, count, type, 0, tag, comm, status);
8          MPI_Bcast(buf1, count, type, 0, comm);
9          break;
10 }

```

Process zero executes a broadcast, followed by a blocking send operation. Process one first executes a blocking receive that matches the send, followed by broadcast call that matches the broadcast of process zero. This program may deadlock. The broadcast call on process zero *may* block until process one executes the matching broadcast call, so that the send is not executed. Process one will definitely block on the receive and so, in this case, never executes the broadcast.

The relative order of execution of collective operations and point-to-point operations should be such, so that even if the collective operations and the point-to-point operations are synchronizing, no deadlock will occur.

### Example 5.28

An unsafe, non-deterministic program.

```

24 switch(rank) {
25     case 0:
26         MPI_Bcast(buf1, count, type, 0, comm);
27         MPI_Send(buf2, count, type, 1, tag, comm);
28         break;
29     case 1:
30         MPI_Recv(buf2, count, type, MPI_ANY_SOURCE, tag, comm, status);
31         MPI_Bcast(buf1, count, type, 0, comm);
32         MPI_Recv(buf2, count, type, MPI_ANY_SOURCE, tag, comm, status);
33         break;
34     case 2:
35         MPI_Send(buf2, count, type, 1, tag, comm);
36         MPI_Bcast(buf1, count, type, 0, comm);
37         break;
38 }

```

All three processes participate in a broadcast. Process 0 sends a message to process 1 after the broadcast, and process 2 sends a message to process 1 before the broadcast. Process 1 receives before and after the broadcast, with a wildcard source argument.

Two possible executions of this program, with different matchings of sends and receives, are illustrated in Figure 5.12. Note that the second execution has the peculiar effect that a send executed after the broadcast is received at another node before the broadcast. This example illustrates the fact that one should not rely on collective communication functions to have particular synchronization effects. A program that works correctly only when the first execution occurs (only when broadcast is synchronizing) is erroneous.

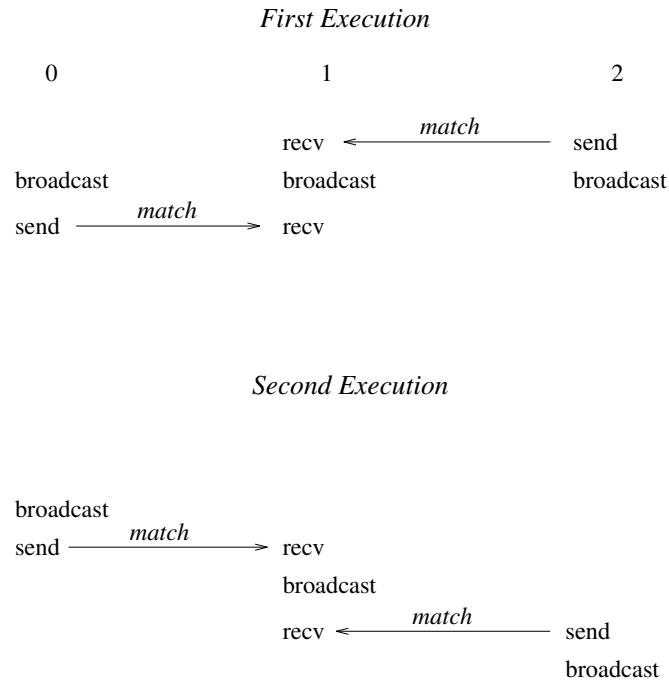


Figure 5.12: A race condition causes non-deterministic matching of sends and receives. One cannot rely on synchronization from a broadcast to make the program deterministic.

Finally, in multithreaded implementations, one can have more than one, concurrently executing, collective communication call at a process. In these situations, it is the user’s responsibility to ensure that the same communicator is not used concurrently by two different collective communication calls at the same process.

*Advice to implementors.* Assume that broadcast is implemented using point-to-point MPI communication. Suppose the following two rules are followed.

1. All receives specify their source explicitly (no wildcards).
2. Each process sends all messages that pertain to one collective call before sending any message that pertain to a subsequent collective call.

Then, messages belonging to successive broadcasts cannot be confused, as the order of point-to-point messages is preserved.

It is the implementor’s responsibility to ensure that point-to-point messages are not confused with collective messages. One way to accomplish this is, whenever a communicator is created, to also create a “hidden communicator” for collective communication. One could achieve a similar effect more cheaply, for example, by using a hidden tag or context bit to indicate whether the communicator is used for point-to-point or collective communication. (*End of advice to implementors.*)

### Example 5.29

Blocking and nonblocking collective operations can be interleaved, i.e., a blocking collective operation can be posted even if there is a nonblocking collective operation outstanding.

```

1  MPI_Request req;
2
3  MPI_Ibarrier(comm, &req);
4  MPI_Bcast(buf1, count, type, 0, comm);
5  MPI_Wait(&req, MPI_STATUS_IGNORE);
6

```

Each process starts a nonblocking barrier operation, participates in a blocking broadcast and then waits until every other process started the barrier operation. This effectively turns the broadcast into a synchronizing broadcast with possible communication/communication overlap (MPI\_Bcast is allowed, but not required to synchronize).

### Example 5.30

The starting order of collective operations on a particular communicator defines their matching. The following example shows an erroneous matching of different collective operations on the same communicator.

```

16 MPI_Request req;
17 switch(rank) {
18     case 0:
19         /* erroneous matching */
20         MPI_Ibarrier(comm, &req);
21         MPI_Bcast(buf1, count, type, 0, comm);
22         MPI_Wait(&req, MPI_STATUS_IGNORE);
23         break;
24     case 1:
25         /* erroneous matching */
26         MPI_Bcast(buf1, count, type, 0, comm);
27         MPI_Ibarrier(comm, &req);
28         MPI_Wait(&req, MPI_STATUS_IGNORE);
29         break;
30 }
31

```

This ordering would match MPI\_Ibarrier on rank 0 with MPI\_Bcast on rank 1 which is erroneous and the program behavior is undefined. However, if such an order is required, the user must create different duplicate communicators and perform the operations on them. If started with two processes, the following program would be correct:

```

37 MPI_Request req;
38 MPI_Comm dupcomm;
39 MPI_Comm_dup(comm, &dupcomm);
40 switch(rank) {
41     case 0:
42         MPI_Ibarrier(comm, &req);
43         MPI_Bcast(buf1, count, type, 0, dupcomm);
44         MPI_Wait(&req, MPI_STATUS_IGNORE);
45         break;
46     case 1:
47         MPI_Bcast(buf1, count, type, 0, dupcomm);
48         MPI_Ibarrier(comm, &req);

```



```

        MPI_Wait(&req, MPI_STATUS_IGNORE);
        break;
}

```

*Advice to users.* The use of different communicators offers some flexibility regarding the matching of nonblocking collective operations. In this sense, communicators could be used as an equivalent to tags. However, communicator construction might induce overheads so that this should be used carefully. (*End of advice to users.*)

### Example 5.31

Nonblocking collective operations can rely on the same progression rules as nonblocking point-to-point messages. Thus, if started with two processes, the following program is a valid MPI program and is guaranteed to terminate:

```

MPI_Request req;

switch(rank) {
    case 0:
        MPI_Ibarrier(comm, &req);
        MPI_Wait(&req, MPI_STATUS_IGNORE);
        MPI_Send(buf, count, dtype, 1, tag, comm);
        break;
    case 1:
        MPI_Ibarrier(comm, &req);
        MPI_Recv(buf, count, dtype, 0, tag, comm, MPI_STATUS_IGNORE);
        MPI_Wait(&req, MPI_STATUS_IGNORE);
        break;
}

```

The MPI library must progress the barrier in the MPI\_Recv call. Thus, the MPI\_Wait call in rank 0 will eventually complete, which enables the matching MPI\_Send so all calls eventually return.

### Example 5.32

Blocking and nonblocking collective operations do not match. The following example is erroneous.

```

MPI_Request req;

switch(rank) {
    case 0:
        /* erroneous false matching of Alltoall and Ialltoall */
        MPI_Ialltoall(sbuf, scnt, stype, rbuf, rcnt, rtype, comm, &req);
        MPI_Wait(&req, MPI_STATUS_IGNORE);
        break;
    case 1:
        /* erroneous false matching of Alltoall and Ialltoall */
        MPI_Alltoall(sbuf, scnt, stype, rbuf, rcnt, rtype, comm);
        break;
}

```

**Example 5.33**

Collective and point-to-point requests can be mixed in functions that enable multiple completions. If started with two processes, the following program is valid.

```

MPI_Request reqs[2];

switch(rank) {
    case 0:
        MPI_Ibarrier(comm, &reqs[0]);
        MPI_Send(buf, count, dtype, 1, tag, comm);
        MPI_Wait(&reqs[0], MPI_STATUS_IGNORE);
        break;
    case 1:
        MPI_Irecv(buf, count, dtype, 0, tag, comm, &reqs[0]);
        MPI_Ibarrier(comm, &reqs[1]);
        MPI_Waitall(2, reqs, MPI_STATUSES_IGNORE);
        break;
}

```

The Waitall call returns only after the barrier and the receive completed.

**Example 5.34**

Multiple nonblocking collective operations can be outstanding on a single communicator and match in order.

```

MPI_Request reqs[3];

compute(buf1);
MPI_Ibcast(buf1, count, type, 0, comm, &reqs[0]);
compute(buf2);
MPI_Ibcast(buf2, count, type, 0, comm, &reqs[1]);
compute(buf3);
MPI_Ibcast(buf3, count, type, 0, comm, &reqs[2]);
MPI_Waitall(3, reqs, MPI_STATUSES_IGNORE);

```

*Advice to users.* Pipelining and double-buffering techniques can efficiently be used to overlap computation and communication. However, having too many outstanding requests might have a negative impact on performance. (*End of advice to users.*)

*Advice to implementors.* The use of pipelining may generate many outstanding requests. A high-quality hardware-supported implementation with limited resources should be able to fall back to a software implementation if its resources are exhausted. In this way, the implementation could limit the number of outstanding requests only by the available memory. (*End of advice to implementors.*)

**Example 5.35**

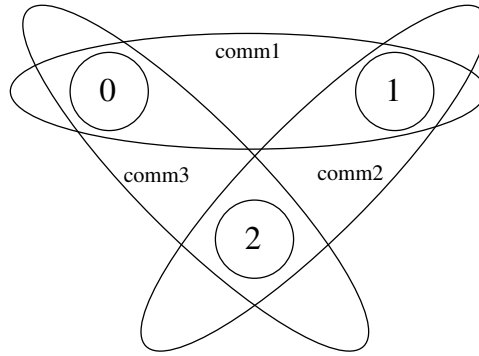


Figure 5.13: Example with overlapping communicators.

Nonblocking collective operations can also be used to enable simultaneous collective operations on multiple overlapping communicators (see Figure 5.13). The following example is started with three processes and three communicators. The first communicator `comm1` includes ranks 0 and 1, `comm2` includes ranks 1 and 2 and `comm3` spans ranks 0 and 2. It is not possible to perform a blocking collective operation on all communicators because there exists no deadlock-free order to invoke them. However, nonblocking collective operations can easily be used to achieve this task.

```
MPI_Request reqs[2];

switch(rank) {
    case 0:
        MPI_Iallreduce(sbuf1, rbuf1, count, dtype, MPI_SUM, comm1, &reqs[0]);
        MPI_Iallreduce(sbuf3, rbuf3, count, dtype, MPI_SUM, comm3, &reqs[1]);
        break;
    case 1:
        MPI_Iallreduce(sbuf1, rbuf1, count, dtype, MPI_SUM, comm1, &reqs[0]);
        MPI_Iallreduce(sbuf2, rbuf2, count, dtype, MPI_SUM, comm2, &reqs[1]);
        break;
    case 2:
        MPI_Iallreduce(sbuf2, rbuf2, count, dtype, MPI_SUM, comm2, &reqs[0]);
        MPI_Iallreduce(sbuf3, rbuf3, count, dtype, MPI_SUM, comm3, &reqs[1]);
        break;
}
MPI_Waitall(2, reqs, MPI_STATUSES_IGNORE);
```

*Advice to users.* This method can be useful if overlapping neighboring regions (halo or ghost zones) are used in collective operations. The sequence of the two calls in each process is irrelevant because the two nonblocking operations are performed on different communicators. (*End of advice to users.*)

### Example 5.36

The progress of multiple outstanding nonblocking collective operations is completely independent.

```
1  MPI_Request reqs[2];
2
3  compute(buf1);
4  MPI_Ibcast(buf1, count, type, 0, comm, &reqs[0]);
5  compute(buf2);
6  MPI_Ibcast(buf2, count, type, 0, comm, &reqs[1]);
7  MPI_Wait(&reqs[1], MPI_STATUS_IGNORE);
8  /* nothing is known about the status of the first bcast here */
9  MPI_Wait(&reqs[0], MPI_STATUS_IGNORE);
```

Finishing the second MPI\_IBCAST is completely independent of the first one. This means that it is not guaranteed that the first broadcast operation is finished or even started after the second one is completed via reqs[1].

## Chapter 6

# Groups, Contexts, Communicators, and Caching

### 6.1 Introduction

This chapter introduces MPI features that support the development of parallel libraries. Parallel libraries are needed to encapsulate the distracting complications inherent in parallel implementations of key algorithms. They help to ensure consistent correctness of such procedures, and provide a “higher level” of portability than MPI itself can provide. As such, libraries prevent each programmer from repeating the work of defining consistent data structures, data layouts, and methods that implement key algorithms (such as matrix operations). Since the best libraries come with several variations on parallel systems (different data layouts, different strategies depending on the size of the system or problem, or type of floating point), this too needs to be hidden from the user.

We refer the reader to [55] and [3] for further information on writing libraries in MPI, using the features described in this chapter.

#### 6.1.1 Features Needed to Support Libraries

The key features needed to support the creation of robust parallel libraries are as follows:

- Safe communication space, that guarantees that libraries can communicate as they need to, without conflicting with communication extraneous to the library,
- Group scope for collective operations, that allow libraries to avoid unnecessarily synchronizing uninvolved processes (potentially running unrelated code),
- Abstract process naming to allow libraries to describe their communication in terms suitable to their own data structures and algorithms,
- The ability to “adorn” a set of communicating processes with additional user-defined attributes, such as extra collective operations. This mechanism should provide a means for the user or library writer effectively to extend a message-passing notation.

In addition, a unified mechanism or object is needed for conveniently denoting communication context, the group of communicating processes, to house abstract process naming, and to store adornments.

### 6.1.2 MPI's Support for Libraries

The corresponding concepts that MPI provides, specifically to support robust libraries, are as follows:

- **Contexts** of communication,
- **Groups** of processes,
- **Virtual topologies**,
- **Attribute caching**,
- **Communicators**.

**Communicators** (see [20, 53, 58]) encapsulate all of these ideas in order to provide the appropriate scope for all communication operations in MPI. Communicators are divided into two kinds: intra-communicators for operations within a single group of processes and inter-communicators for operations between two groups of processes.

**Caching.** Communicators (see below) provide a “caching” mechanism that allows one to associate new attributes with communicators, on a par with MPI built-in features. This can be used by advanced users to adorn communicators further, and by MPI to implement some communicator functions. For example, the virtual-topology functions described in Chapter 7 are likely to be supported this way.

**Groups.** Groups define an ordered collection of processes, each with a rank, and it is this group that defines the low-level names for inter-process communication (ranks are used for sending and receiving). Thus, groups define a scope for process names in point-to-point communication. In addition, groups define the scope of collective operations. Groups may be manipulated separately from communicators in MPI, but only communicators can be used in communication operations.

**Intra-communicators.** The most commonly used means for message passing in MPI is via intra-communicators. Intra-communicators contain an instance of a group, contexts of communication for both point-to-point and collective communication, and the ability to include virtual topology and other attributes. These features work as follows:

- **Contexts** provide the ability to have separate safe “universes” of message-passing in MPI. A context is akin to an additional tag that differentiates messages. The system manages this differentiation process. The use of separate communication contexts by distinct libraries (or distinct library invocations) insulates communication internal to the library execution from external communication. This allows the invocation of the library even if there are pending communications on “other” communicators, and avoids the need to synchronize entry or exit into library code. Pending point-to-point communications are also guaranteed not to interfere with collective communications within a single communicator.
- **Groups** define the participants in the communication (see above) of a communicator.

- A **virtual topology** defines a special mapping of the ranks in a group to and from a topology. Special constructors for communicators are defined in Chapter 7 to provide this feature. Intra-communicators as described in this chapter do not have topologies.
- **Attributes** define the local information that the user or library has added to a communicator for later reference.

*Advice to users.* The practice in many communication libraries is that there is a unique, predefined communication universe that includes all processes available when the parallel program is initiated; the processes are assigned consecutive ranks. Participants in a point-to-point communication are identified by their rank; a collective communication (such as broadcast) always involves all processes. This practice can be followed in MPI by using the predefined communicator `MPI_COMM_WORLD`. *Users who are satisfied with this practice can plug in `MPI_COMM_WORLD` wherever a communicator argument is required, and can consequently disregard the rest of this chapter. (End of advice to users.)*

**Inter-communicators.** The discussion has dealt so far with **intra-communication**: communication within a group. MPI also supports **inter-communication**: communication between two non-overlapping groups. When an application is built by composing several parallel modules, it is convenient to allow one module to communicate with another using local ranks for addressing within the second module. This is especially convenient in a client-server computing paradigm, where either client or server are parallel. The support of inter-communication also provides a mechanism for the extension of MPI to a dynamic model where not all processes are preallocated at initialization time. In such a situation, it becomes necessary to support communication across “universes.” Inter-communication is supported by objects called **inter-communicators**. These objects bind two groups together with communication contexts shared by both groups. For inter-communicators, these features work as follows:

- **Contexts** provide the ability to have a separate safe “universe” of message-passing between the two groups. A send in the local group is always a receive in the remote group, and vice versa. The system manages this differentiation process. The use of separate communication contexts by distinct libraries (or distinct library invocations) insulates communication internal to the library execution from external communication. This allows the invocation of the library even if there are pending communications on “other” communicators, and avoids the need to synchronize entry or exit into library code.
- A local and remote group specify the recipients and destinations for an inter-communicator.
- Virtual topology is undefined for an inter-communicator.
- As before, attributes cache defines the local information that the user or library has added to a communicator for later reference.

MPI provides mechanisms for creating and manipulating inter-communicators. They are used for point-to-point and collective communication in an related manner to intra-communicators. Users who do not need inter-communication in their applications can safely

ignore this extension. Users who require inter-communication between overlapping groups must layer this capability on top of MPI.

## 6.2 Basic Concepts

In this section, we turn to a more formal definition of the concepts introduced above.

### 6.2.1 Groups

A **group** is an ordered set of process identifiers (henceforth processes); processes are implementation-dependent objects. Each process in a group is associated with an integer **rank**. Ranks are contiguous and start from zero. Groups are represented by opaque **group objects**, and hence cannot be directly transferred from one process to another. A group is used within a communicator to describe the participants in a communication “universe” and to rank such participants (thus giving them unique names within that “universe” of communication).

There is a special pre-defined group: `MPI_GROUP_EMPTY`, which is a group with no members. The predefined constant `MPI_GROUP_NULL` is the value used for invalid group handles.

*Advice to users.* `MPI_GROUP_EMPTY`, which is a valid handle to an empty group, should not be confused with `MPI_GROUP_NULL`, which in turn is an invalid handle. The former may be used as an argument to group operations; the latter, which is returned when a group is freed, is not a valid argument. (*End of advice to users.*)

*Advice to implementors.* A group may be represented by a virtual-to-real process-address-translation table. Each communicator object (see below) would have a pointer to such a table.

Simple implementations of MPI will enumerate groups, such as in a table. However, more advanced data structures make sense in order to improve scalability and memory usage with large numbers of processes. Such implementations are possible with MPI. (*End of advice to implementors.*)

### 6.2.2 Contexts

A **context** is a property of communicators (defined next) that allows partitioning of the communication space. A message sent in one context cannot be received in another context. Furthermore, where permitted, collective operations are independent of pending point-to-point operations. Contexts are not explicit MPI objects; they appear only as part of the realization of communicators (below).

*Advice to implementors.* Distinct communicators in the same process have distinct contexts. A context is essentially a system-managed tag (or tags) needed to make a communicator safe for point-to-point and MPI-defined collective communication. Safety means that collective and point-to-point communication within one communicator do not interfere, and that communication over distinct communicators don’t interfere.



A possible implementation for a context is as a supplemental tag attached to messages on send and matched on receive. Each intra-communicator stores the value of its two tags (one for point-to-point and one for collective communication). Communicator-generating functions use a collective communication to agree on a new group-wide unique context.

Analogously, in inter-communication, two context tags are stored per communicator, one used by group A to send and group B to receive, and a second used by group B to send and for group A to receive.

Since contexts are not explicit objects, other implementations are also possible. (*End of advice to implementors.*)

### 6.2.3 Intra-Communicators

Intra-communicators bring together the concepts of group and context. To support implementation-specific optimizations, and application topologies (defined in the next chapter, Chapter 7), communicators may also “cache” additional information (see Section 6.7). MPI communication operations reference communicators to determine the scope and the “communication universe” in which a point-to-point or collective operation is to operate.

Each communicator contains a group of valid participants; this group always includes the local process. The source and destination of a message is identified by process rank within that group.

For collective communication, the intra-communicator specifies the set of processes that participate in the collective operation (and their order, when significant). Thus, the communicator restricts the “spatial” scope of communication, and provides machine-independent process addressing through ranks.

Intra-communicators are represented by opaque **intra-communicator objects**, and hence cannot be directly transferred from one process to another.

### 6.2.4 Predefined Intra-Communicators

An initial intra-communicator `MPI_COMM_WORLD` of all processes the local process can communicate with after initialization (itself included) is defined once `MPI_INIT` or `MPI_INIT_THREAD` has been called. In addition, the communicator `MPI_COMM_SELF` is provided, which includes only the process itself.

The predefined constant `MPI_COMM_NULL` is the value used for invalid communicator handles.

In a static-process-model implementation of MPI, all processes that participate in the computation are available after MPI is initialized. For this case, `MPI_COMM_WORLD` is a communicator of all processes available for the computation; this communicator has the same value in all processes. In an implementation of MPI where processes can dynamically join an MPI execution, it may be the case that a process starts an MPI computation without having access to all other processes. In such situations, `MPI_COMM_WORLD` is a communicator incorporating all processes with which the joining process can immediately communicate. Therefore, `MPI_COMM_WORLD` may simultaneously represent disjoint groups in different processes.

All MPI implementations are required to provide the `MPI_COMM_WORLD` communicator. It cannot be deallocated during the life of a process. The group corresponding to this communicator does not appear as a pre-defined constant, but it may be accessed using

MPI\_COMM\_GROUP (see below). MPI does not specify the correspondence between the process rank in MPI\_COMM\_WORLD and its (machine-dependent) absolute address. Neither does MPI specify the function of the host process, if any. Other implementation-dependent, predefined communicators may also be provided.

## 6.3 Group Management

This section describes the manipulation of process groups in MPI. These operations are local and their execution does not require interprocess communication.

### 6.3.1 Group Accessors

MPI\_GROUP\_SIZE(group, size)

IN	group	group (handle)
OUT	size	number of processes in the group (integer)

```
int MPI_Group_size(MPI_Group group, int *size)
```

```
MPI_Group_size(group, size, ierror) BIND(C)
  TYPE(MPI_Group), INTENT(IN) :: group
  INTEGER, INTENT(OUT) :: size
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_GROUP_SIZE(GROUP, SIZE, IERROR)
  INTEGER GROUP, SIZE, IERROR
```

```
{int MPI::Group::Get_size() const(binding deprecated, see Section 15.2) }
```

MPI\_GROUP\_RANK(group, rank)

IN	group	group (handle)
OUT	rank	rank of the calling process in group, or MPI_UNDEFINED if the process is not a member (integer)

```
int MPI_Group_rank(MPI_Group group, int *rank)
```

```
MPI_Group_rank(group, rank, ierror) BIND(C)
  TYPE(MPI_Group), INTENT(IN) :: group
  INTEGER, INTENT(OUT) :: rank
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_GROUP_RANK(GROUP, RANK, IERROR)
  INTEGER GROUP, RANK, IERROR
```

```
{int MPI::Group::Get_rank() const(binding deprecated, see Section 15.2) }
```

```

MPI_GROUP_TRANSLATE_RANKS (group1, n, ranks1, group2, ranks2)
    IN      group1      group1 (handle)
    IN      n           number of ranks in ranks1 and ranks2 arrays (integer)
    IN      ranks1      array of zero or more valid ranks in group1
    IN      group2      group2 (handle)
    OUT     ranks2      array of corresponding ranks in group2,
                        MPI_UNDEFINED when no correspondence exists.

int MPI_Group_translate_ranks (MPI_Group group1, int n, const int *ranks1,
                             MPI_Group group2, int *ranks2)

MPI_Group_translate_ranks(group1, n, ranks1, group2, ranks2, ierror)
    BIND(C)
    TYPE(MPI_Group), INTENT(IN) :: group1, group2
    INTEGER, INTENT(IN) :: n, ranks1(n)
    INTEGER, INTENT(OUT) :: ranks2(n)
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_GROUP_TRANSLATE_RANKS(GROUP1, N, RANKS1, GROUP2, RANKS2, IERROR)
    INTEGER GROUP1, N, RANKS1(*), GROUP2, RANKS2(*), IERROR

{static void MPI::Group::Translate_ranks (const MPI::Group& group1, int n,
    const int ranks1[], const MPI::Group& group2,
    int ranks2[]) (binding deprecated, see Section 15.2) }

    This function is important for determining the relative numbering of the same processes
    in two different groups. For instance, if one knows the ranks of certain processes in the group
    of MPI_COMM_WORLD, one might want to know their ranks in a subset of that group.

    MPI_PROC_NULL is a valid rank for input to MPI_GROUP_TRANSLATE_RANKS, which
    returns MPI_PROC_NULL as the translated rank.

MPI_GROUP_COMPARE(group1, group2, result)
    IN      group1      first group (handle)
    IN      group2      second group (handle)
    OUT     result      result (integer)

int MPI_Group_compare(MPI_Group group1, MPI_Group group2, int *result)

MPI_Group_compare(group1, group2, result, ierror) BIND(C)
    TYPE(MPI_Group), INTENT(IN) :: group1, group2
    INTEGER, INTENT(OUT) :: result
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_GROUP_COMPARE(GROUP1, GROUP2, RESULT, IERROR)
    INTEGER GROUP1, GROUP2, RESULT, IERROR

```

```

1 {static int MPI::Group::Compare(const MPI::Group& group1,
2     const MPI::Group& group2) (binding deprecated, see Section 15.2) }
3

```

MPI\_IDENT results if the group members and group order is exactly the same in both groups. This happens for instance if group1 and group2 are the same handle. MPI\_SIMILAR results if the group members are the same but the order is different. MPI\_UNEQUAL results otherwise.

### 6.3.2 Group Constructors

Group constructors are used to subset and superset existing groups. These constructors construct new groups from existing groups. These are local operations, and distinct groups may be defined on different processes; a process may also define a group that does not include itself. Consistent definitions are required when groups are used as arguments in communicator-building functions. MPI does not provide a mechanism to build a group from scratch, but only from other, previously defined groups. The base group, upon which all other groups are defined, is the group associated with the initial communicator MPI\_COMM\_WORLD (accessible through the function MPI\_COMM\_GROUP).

*Rationale.* In what follows, there is no group duplication function analogous to MPI\_COMM\_DUP, defined later in this chapter. There is no need for a group duplicator. A group, once created, can have several references to it by making copies of the handle. The following constructors address the need for subsets and supersets of existing groups. (*End of rationale.*)

*Advice to implementors.* Each group constructor behaves as if it returned a new group object. When this new group is a copy of an existing group, then one can avoid creating such new objects, using a reference-count mechanism. (*End of advice to implementors.*)

```

30 MPI_COMM_GROUP(comm, group)
31

```

IN	comm	communicator (handle)
OUT	group	group corresponding to comm (handle)

```

35 int MPI_Comm_group(MPI_Comm comm, MPI_Group *group)
36

```

```

37 MPI_Comm_group(comm, group, ierror) BIND(C)
38     TYPE(MPI_Comm), INTENT(IN) :: comm
39     TYPE(MPI_Group), INTENT(OUT) :: group
40     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
41

```

```

42 MPI_COMM_GROUP(COMM, GROUP, IERROR)
43     INTEGER COMM, GROUP, IERROR
44

```

```

45 {MPI::Group MPI::Comm::Get_group() const (binding deprecated, see Section 15.2) }
46

```

MPI\_COMM\_GROUP returns in group a handle to the group of comm.

```

MPI_GROUP_UNION(group1, group2, newgroup)
    IN      group1      first group (handle)
    IN      group2      second group (handle)
    OUT     newgroup     union group (handle)

int MPI_Group_union(MPI_Group group1, MPI_Group group2,
    MPI_Group *newgroup)

MPI_Group_union(group1, group2, newgroup, ierror) BIND(C)
    TYPE(MPI_Group), INTENT(IN) :: group1, group2
    TYPE(MPI_Group), INTENT(OUT) :: newgroup
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_GROUP_UNION(GROUP1, GROUP2, NEWGROUP, IERROR)
    INTEGER GROUP1, GROUP2, NEWGROUP, IERROR

{static MPI::Group MPI::Group::Union(const MPI::Group& group1,
    const MPI::Group& group2) (binding deprecated, see Section 15.2) }

MPI_GROUP_INTERSECTION(group1, group2, newgroup)
    IN      group1      first group (handle)
    IN      group2      second group (handle)
    OUT     newgroup     intersection group (handle)

int MPI_Group_intersection(MPI_Group group1, MPI_Group group2,
    MPI_Group *newgroup)

MPI_Group_intersection(group1, group2, newgroup, ierror) BIND(C)
    TYPE(MPI_Group), INTENT(IN) :: group1, group2
    TYPE(MPI_Group), INTENT(OUT) :: newgroup
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_GROUP_INTERSECTION(GROUP1, GROUP2, NEWGROUP, IERROR)
    INTEGER GROUP1, GROUP2, NEWGROUP, IERROR

{static MPI::Group MPI::Group::Intersect(const MPI::Group& group1,
    const MPI::Group& group2) (binding deprecated, see Section 15.2) }

MPI_GROUP_DIFFERENCE(group1, group2, newgroup)
    IN      group1      first group (handle)
    IN      group2      second group (handle)
    OUT     newgroup     difference group (handle)

int MPI_Group_difference(MPI_Group group1, MPI_Group group2,
    MPI_Group *newgroup)

```

```

1
ticket-248T. 2 MPI_Group_difference(group1, group2, newgroup, ierror) BIND(C)
3 TYPE(MPI_Group), INTENT(IN) :: group1, group2
4 TYPE(MPI_Group), INTENT(OUT) :: newgroup
5 INTEGER, OPTIONAL, INTENT(OUT) :: ierror
6
7 MPI_GROUP_DIFFERENCE(GROUP1, GROUP2, NEWGROUP, IERROR)
8 INTEGER GROUP1, GROUP2, NEWGROUP, IERROR
9 {static MPI::Group MPI::Group::Difference(const MPI::Group& group1,
10 const MPI::Group& group2) (binding deprecated, see Section 15.2) }
11

```

The set-like operations are defined as follows:

**union** All elements of the first group (group1), followed by all elements of second group (group2) not in first.

**intersect** all elements of the first group that are also in the second group, ordered as in first group.

**difference** all elements of the first group that are not in the second group, ordered as in the first group.

Note that for these operations the order of processes in the output group is determined primarily by order in the first group (if possible) and then, if necessary, by order in the second group. Neither union nor intersection are commutative, but both are associative.

The new group can be empty, that is, equal to MPI\_GROUP\_EMPTY.

```

25
26 MPI_GROUP_INCL(group, n, ranks, newgroup)

```

27	IN	group	group (handle)
28	IN	n	number of elements in array ranks (and size of
29			newgroup) (integer)
30			
31	IN	ranks	ranks of processes in group to appear in
32			newgroup (array of integers)
33	OUT	newgroup	new group derived from above, in the order defined by
34			ranks (handle)
35			

```

36
ticket140. 37 int MPI_Group_incl(MPI_Group group, int n, const int *ranks,
38 MPI_Group *newgroup)

```

```

ticket229.1. 39 MPI_Group_incl(group, n, ranks, newgroup, ierror) BIND(C)
ticket-248T. 40 TYPE(MPI_Group), INTENT(IN) :: group
41 INTEGER, INTENT(IN) :: n, ranks(n)
42 TYPE(MPI_Group), INTENT(OUT) :: newgroup
43 INTEGER, OPTIONAL, INTENT(OUT) :: ierror
44
45 MPI_GROUP_INCL(GROUP, N, RANKS, NEWGROUP, IERROR)
46 INTEGER GROUP, N, RANKS(*), NEWGROUP, IERROR
47 {MPI::Group MPI::Group::Incl(int n, const int ranks[]) const (binding
48 deprecated, see Section 15.2) }

```

The function `MPI_GROUP_INCL` creates a group `newgroup` that consists of the `n` processes in `group` with ranks `ranks[0], ..., ranks[n-1]`; the process with rank `i` in `newgroup` is the process with rank `ranks[i]` in `group`. Each of the `n` elements of `ranks` must be a valid rank in `group` and all elements must be distinct, or else the program is erroneous. If `n = 0`, then `newgroup` is `MPI_GROUP_EMPTY`. This function can, for instance, be used to reorder the elements of a group. See also `MPI_GROUP_COMPARE`.

`MPI_GROUP_EXCL(group, n, ranks, newgroup)`

IN	<code>group</code>	group (handle)
IN	<code>n</code>	number of elements in array <code>ranks</code> (integer)
IN	<code>ranks</code>	array of integer ranks in <code>group</code> not to appear in <code>newgroup</code>
OUT	<code>newgroup</code>	new group derived from above, preserving the order defined by <code>group</code> (handle)

```
int MPI_Group_excl(MPI_Group group, int n, const int *ranks,
                  MPI_Group *newgroup)
```

```
MPI_Group_excl(group, n, ranks, newgroup, ierror) BIND(C)
    TYPE(MPI_Group), INTENT(IN) :: group
    INTEGER, INTENT(IN) :: n, ranks(n)
    TYPE(MPI_Group), INTENT(OUT) :: newgroup
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_GROUP_EXCL(GROUP, N, RANKS, NEWGROUP, IERROR)
    INTEGER GROUP, N, RANKS(*), NEWGROUP, IERROR
```

```
{MPI::Group MPI::Group::Excl(int n, const int ranks[]) const(binding
    deprecated, see Section 15.2) }
```

The function `MPI_GROUP_EXCL` creates a group of processes `newgroup` that is obtained by deleting from `group` those processes with ranks `ranks[0] ... ranks[n-1]`. The ordering of processes in `newgroup` is identical to the ordering in `group`. Each of the `n` elements of `ranks` must be a valid rank in `group` and all elements must be distinct; otherwise, the program is erroneous. If `n = 0`, then `newgroup` is identical to `group`.

`MPI_GROUP_RANGE_INCL(group, n, ranges, newgroup)`

IN	<code>group</code>	group (handle)
IN	<code>n</code>	number of triplets in array <code>ranges</code> (integer)
IN	<code>ranges</code>	a one-dimensional array of integer triplets, of the form (first rank, last rank, stride) indicating ranks in <code>group</code> of processes to be included in <code>newgroup</code>
OUT	<code>newgroup</code>	new group derived from above, in the order defined by <code>ranges</code> (handle)

```

1 int MPI_Group_range_incl(MPI_Group group, int n, int ranges[][3],
2     MPI_Group *newgroup)

```

```

3 MPI_Group_range_incl(group, n, ranges, newgroup, ierror) BIND(C)
4     TYPE(MPI_Group), INTENT(IN) :: group
5     INTEGER, INTENT(IN) :: n, ranges(3,n)
6     TYPE(MPI_Group), INTENT(OUT) :: newgroup
7     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

9 MPI_GROUP_RANGE_INCL(GROUP, N, RANGES, NEWGROUP, IERROR)
10     INTEGER GROUP, N, RANGES(3,*), NEWGROUP, IERROR

```

```

11 {MPI::Group MPI::Group::Range_incl(int n, const int ranges[][3])
12     const(binding deprecated, see Section 15.2) }

```

If `ranges` consist of the triplets

$$(first_1, last_1, stride_1), \dots, (first_n, last_n, stride_n)$$

then `newgroup` consists of the sequence of processes in `group` with ranks

$$first_1, first_1 + stride_1, \dots, first_1 + \left\lfloor \frac{last_1 - first_1}{stride_1} \right\rfloor stride_1, \dots$$

$$first_n, first_n + stride_n, \dots, first_n + \left\lfloor \frac{last_n - first_n}{stride_n} \right\rfloor stride_n.$$

Each computed rank must be a valid rank in `group` and all computed ranks must be distinct, or else the program is erroneous. Note that we may have  $first_i > last_i$ , and  $stride_i$  may be negative, but cannot be zero.

The functionality of this routine is specified to be equivalent to expanding the array of ranges to an array of the included ranks and passing the resulting array of ranks and other arguments to `MPI_GROUP_INCL`. A call to `MPI_GROUP_INCL` is equivalent to a call to `MPI_GROUP_RANGE_INCL` with each rank `i` in `ranks` replaced by the triplet `(i,i,1)` in the argument `ranges`.

```

34 MPI_GROUP_RANGE_EXCL(group, n, ranges, newgroup)

```

IN	group	group (handle)
IN	n	number of elements in array ranges (integer)
IN	ranges	a one-dimensional array of integer triplets of the form (first rank, last rank, stride), indicating the ranks in group of processes to be excluded from the output group <code>newgroup</code> .
OUT	newgroup	new group derived from above, preserving the order in group (handle)

```

45 int MPI_Group_range_excl(MPI_Group group, int n, int ranges[][3],
46     MPI_Group *newgroup)

```

```

47 MPI_Group_range_excl(group, n, ranges, newgroup, ierror) BIND(C)

```



```

TYPE(MPI_Group), INTENT(IN) :: group
INTEGER, INTENT(IN) :: n, ranges(3,n)
TYPE(MPI_Group), INTENT(OUT) :: newgroup
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_GROUP_RANGE_EXCL(GROUP, N, RANGES, NEWGROUP, IERROR)
  INTEGER GROUP, N, RANGES(3,*), NEWGROUP, IERROR

{MPI::Group MPI::Group::Range_excl(int n, const int ranges[] [3])
  const(binding deprecated, see Section 15.2) }
```

Each computed rank must be a valid rank in group and all computed ranks must be distinct, or else the program is erroneous.

The functionality of this routine is specified to be equivalent to expanding the array of ranges to an array of the excluded ranks and passing the resulting array of ranks and other arguments to MPI\_GROUP\_EXCL. A call to MPI\_GROUP\_EXCL is equivalent to a call to MPI\_GROUP\_RANGE\_EXCL with each rank *i* in *ranks* replaced by the triplet (*i*,*i*,1) in the argument *ranges*.

*Advice to users.* The range operations do not explicitly enumerate ranks, and therefore are more scalable if implemented efficiently. Hence, we recommend MPI programmers to use them whenever possible, as high-quality implementations will take advantage of this fact. (*End of advice to users.*)

*Advice to implementors.* The range operations should be implemented, if possible, without enumerating the group members, in order to obtain better scalability (time and space). (*End of advice to implementors.*)

### 6.3.3 Group Destructors

```

MPI_GROUP_FREE(group)
  INOUT    group                group (handle)

int MPI_Group_free(MPI_Group *group)

MPI_Group_free(group, ierror) BIND(C)
  TYPE(MPI_Group), INTENT(INOUT) :: group
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_GROUP_FREE(GROUP, IERROR)
  INTEGER GROUP, IERROR

{void MPI::Group::Free() (binding deprecated, see Section 15.2) }
```

This operation marks a group object for deallocation. The handle *group* is set to MPI\_GROUP\_NULL by the call. Any on-going operation using this group will complete normally.

*Advice to implementors.* One can keep a reference count that is incremented for each call to MPI\_COMM\_GROUP, MPI\_COMM\_CREATE and MPI\_COMM\_DUP, and

decremented for each call to `MPI_GROUP_FREE` or `MPI_COMM_FREE`; the group object is ultimately deallocated when the reference count drops to zero. (*End of advice to implementors.*)

## 6.4 Communicator Management

This section describes the manipulation of communicators in MPI. Operations that access communicators are local and their execution does not require interprocess communication. Operations that create communicators are collective and may require interprocess communication.

*Advice to implementors.* High-quality implementations should amortize the overheads associated with the creation of communicators (for the same group, or subsets thereof) over several calls, by allocating multiple contexts with one collective communication. (*End of advice to implementors.*)

### 6.4.1 Communicator Accessors

The following are all local operations.

`MPI_COMM_SIZE(comm, size)`

IN	comm	communicator (handle)
OUT	size	number of processes in the group of comm (integer)

`int MPI_Comm_size(MPI_Comm comm, int *size)`

```
MPI_Comm_size(comm, size, ierror) BIND(C)
  TYPE(MPI_Comm), INTENT(IN) :: comm
  INTEGER, INTENT(OUT) :: size
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_COMM_SIZE(COMM, SIZE, IERROR)
  INTEGER COMM, SIZE, IERROR
```

```
{int MPI::Comm::Get_size() const(binding deprecated, see Section 15.2) }
```

*Rationale.* This function is equivalent to accessing the communicator's group with `MPI_COMM_GROUP` (see above), computing the size using `MPI_GROUP_SIZE`, and then freeing the temporary group via `MPI_GROUP_FREE`. However, this function is so commonly used, that this shortcut was introduced. (*End of rationale.*)

*Advice to users.* This function indicates the number of processes involved in a communicator. For `MPI_COMM_WORLD`, it indicates the total number of processes available (for this version of MPI, there is no standard way to change the number of processes once initialization has taken place).

This call is often used with the next call to determine the amount of concurrency available for a specific library or program. The following call, `MPI_COMM_RANK`

indicates the rank of the process that calls it in the range from 0...size-1, where size is the return value of MPI\_COMM\_SIZE. (*End of advice to users.*)

MPI\_COMM\_RANK(comm, rank)

IN	comm	communicator (handle)
OUT	rank	rank of the calling process in group of comm (integer)

```
int MPI_Comm_rank(MPI_Comm comm, int *rank)
```

```
MPI_Comm_rank(comm, rank, ierror) BIND(C)
  TYPE(MPI_Comm), INTENT(IN) :: comm
  INTEGER, INTENT(OUT) :: rank
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_COMM_RANK(COMM, RANK, IERROR)
  INTEGER COMM, RANK, IERROR
```

```
{int MPI::Comm::Get_rank() const(binding deprecated, see Section 15.2) }
```

*Rationale.* This function is equivalent to accessing the communicator's group with MPI\_COMM\_GROUP (see above), computing the rank using MPI\_GROUP\_RANK, and then freeing the temporary group via MPI\_GROUP\_FREE. However, this function is so commonly used, that this shortcut was introduced. (*End of rationale.*)

*Advice to users.* This function gives the rank of the process in the particular communicator's group. It is useful, as noted above, in conjunction with MPI\_COMM\_SIZE.

Many programs will be written with the master-slave model, where one process (such as the rank-zero process) will play a supervisory role, and the other processes will serve as compute nodes. In this framework, the two preceding calls are useful for determining the roles of the various processes of a communicator. (*End of advice to users.*)

MPI\_COMM\_COMPARE(comm1, comm2, result)

IN	comm1	first communicator (handle)
IN	comm2	second communicator (handle)
OUT	result	result (integer)

```
int MPI_Comm_compare(MPI_Comm comm1, MPI_Comm comm2, int *result)
```

```
MPI_Comm_compare(comm1, comm2, result, ierror) BIND(C)
  TYPE(MPI_Comm), INTENT(IN) :: comm1, comm2
  INTEGER, INTENT(OUT) :: result
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
1 MPI_COMM_COMPARE(COMM1, COMM2, RESULT, IERROR)
```

```
2     INTEGER COMM1, COMM2, RESULT, IERROR
```

```
3 {static int MPI::Comm::Compare(const MPI::Comm& comm1,
4     const MPI::Comm& comm2) (binding deprecated, see Section 15.2) }
```

6 MPI\_IDENT results if and only if comm1 and comm2 are handles for the same object (identical groups and same contexts). MPI\_CONGRUENT results if the underlying groups are identical in constituents and rank order; these communicators differ only by context. MPI\_SIMILAR results if the group members of both communicators are the same but the rank order differs. MPI\_UNEQUAL results otherwise.

## 12 6.4.2 Communicator Constructors

14 The following are collective functions that are invoked by all processes in the group or groups associated with comm.

17 *Rationale.* Note that there is a chicken-and-egg aspect to MPI in that a communicator is needed to create a new communicator. The base communicator for all MPI communicators is predefined outside of MPI, and is MPI\_COMM\_WORLD. This model was arrived at after considerable debate, and was chosen to increase “safety” of programs written in MPI. (*End of rationale.*)

22 The MPI interface provides four communicator construction routines that apply to both intracommunicators and intercommunicators. The construction routine MPI\_INTERCOMM\_CREATE (discussed later) applies only to intercommunicators.

26 An intracommunicator involves a single group while an intercommunicator involves two groups. Where the following discussions address intercommunicator semantics, the two groups in an intercommunicator are called the *left* and *right* groups. A process in an intercommunicator is a member of either the left or the right group. From the point of view of that process, the group that the process is a member of is called the *local* group; the other group (relative to that process) is the *remote* group. The left and right group labels give us a way to describe the two groups in an intercommunicator that is not relative to any particular process (as the local and remote groups are).

```
35 MPI_COMM_DUP(comm, newcomm)
```

```
36     IN      comm      communicator (handle)
```

```
38     OUT    newcomm    copy of comm (handle)
```

```
40 int MPI_Comm_dup(MPI_Comm comm, MPI_Comm *newcomm)
```

```
41 MPI_Comm_dup(comm, newcomm, ierror) BIND(C)
```

```
42     TYPE(MPI_Comm), INTENT(IN) :: comm
```

```
44     TYPE(MPI_Comm), INTENT(OUT) :: newcomm
```

```
45     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
46 MPI_COMM_DUP(COMM, NEWCOMM, IERROR)
```

```
47     INTEGER COMM, NEWCOMM, IERROR
```

```

{MPI::Intracomm MPI::Intracomm::Dup() const(binding deprecated, see Section 15.2)
    }
{MPI::Intercomm MPI::Intercomm::Dup() const(binding deprecated, see Section 15.2)
    }
{MPI::Cartcomm MPI::Cartcomm::Dup() const(binding deprecated, see Section 15.2) }
{MPI::Graphcomm MPI::Graphcomm::Dup() const(binding deprecated, see Section 15.2)
    }
{MPI::Distgraphcomm MPI::Distgraphcomm::Dup() const(binding deprecated, see
    Section 15.2) }
{MPI::Comm& MPI::Comm::Clone() const = 0(binding deprecated, see Section 15.2) }
{MPI::Intracomm& MPI::Intracomm::Clone() const(binding deprecated, see
    Section 15.2) }
{MPI::Intercomm& MPI::Intercomm::Clone() const(binding deprecated, see
    Section 15.2) }
{MPI::Cartcomm& MPI::Cartcomm::Clone() const(binding deprecated, see Section 15.2)
    }
{MPI::Graphcomm& MPI::Graphcomm::Clone() const(binding deprecated, see
    Section 15.2) }
{MPI::Distgraphcomm& MPI::Distgraphcomm::Clone() const(binding deprecated, see
    Section 15.2) }

```

**MPI\_COMM\_DUP** Duplicates the existing communicator `comm` with associated key values. For each key value, the respective copy callback function determines the attribute value associated with this key in the new communicator; one particular action that a copy callback may take is to delete the attribute from the new communicator. Returns in `newcomm` a new communicator with the same group or groups, any copied cached information, but a new context (see Section 6.7.1). Please see Section 16.1.7 on page 634 for further discussion about the C++ bindings for `Dup()` and `Clone()`.

*Advice to users.* This operation is used to provide a parallel library call with a duplicate communication space that has the same properties as the original communicator. This includes any attributes (see below), and topologies (see Chapter 7). This call is valid even if there are pending point-to-point communications involving the communicator `comm`. A typical call might involve a `MPI_COMM_DUP` at the beginning of the parallel call, and an `MPI_COMM_FREE` of that duplicated communicator at the end of the call. Other models of communicator management are also possible.

This call applies to both intra- and inter-communicators. (*End of advice to users.*)

*Advice to implementors.* One need not actually copy the group information, but only add a new reference and increment the reference count. Copy on write can be used for the cached information. (*End of advice to implementors.*)

```

1 MPI_COMM_CREATE(comm, group, newcomm)
2     IN      comm      communicator (handle)
3     IN      group      Group, which is a subset of the group of
4                      comm (handle)
5     OUT     newcomm    new communicator (handle)
6
7
8 int MPI_Comm_create(MPI_Comm comm, MPI_Group group, MPI_Comm *newcomm)
9
10 MPI_Comm_create(comm, group, newcomm, ierror) BIND(C)
11     TYPE(MPI_Comm), INTENT(IN) :: comm
12     TYPE(MPI_Group), INTENT(IN) :: group
13     TYPE(MPI_Comm), INTENT(OUT) :: newcomm
14     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
15
16 MPI_COMM_CREATE(COMM, GROUP, NEWCOMM, IERROR)
17     INTEGER COMM, GROUP, NEWCOMM, IERROR
18 {MPI::Intercomm MPI::Intercomm::Create(const MPI::Group& group)
19     const(binding deprecated, see Section 15.2) }
20
21 {MPI::Intracomm MPI::Intracomm::Create(const MPI::Group& group)
22     const(binding deprecated, see Section 15.2) }

```

If `comm` is an intracommunicator, this function returns a new communicator `newcomm` with communication group defined by the `group` argument. No cached information propagates from `comm` to `newcomm`. Each process must call with a `group` argument that is a subgroup of the `group` associated with `comm`; this could be `MPI_GROUP_EMPTY`. The processes may specify different values for the `group` argument. If a process calls with a non-empty `group` then all processes in that `group` must call the function with the same `group` as argument, that is the same processes in the same order. Otherwise the call is erroneous. This implies that the set of groups specified across the processes must be disjoint. If the calling process is a member of the group given as `group` argument, then `newcomm` is a communicator with `group` as its associated group. In the case that a process calls with a `group` to which it does not belong, e.g., `MPI_GROUP_EMPTY`, then `MPI_COMM_NULL` is returned as `newcomm`. The function is collective and must be called by all processes in the group of `comm`.

*Rationale.* The interface supports the original mechanism from MPI-1.1, which required the same `group` in all processes of `comm`. It was extended in MPI-2.2 to allow the use of disjoint subgroups in order to allow implementations to eliminate unnecessary communication that `MPI_COMM_SPLIT` would incur when the user already knows the membership of the disjoint subgroups. (*End of rationale.*)

*Rationale.* The requirement that the entire group of `comm` participate in the call stems from the following considerations:

- It allows the implementation to layer `MPI_COMM_CREATE` on top of regular collective communications.
- It provides additional safety, in particular in the case where partially overlapping groups are used to create new communicators.

- It permits implementations sometimes to avoid communication related to context creation.

(*End of rationale.*)

*Advice to users.* MPI\_COMM\_CREATE provides a means to subset a group of processes for the purpose of separate MIMD computation, with separate communication space. newcomm, which emerges from MPI\_COMM\_CREATE can be used in subsequent calls to MPI\_COMM\_CREATE (or other communicator constructors) further to subdivide a computation into parallel sub-computations. A more general service is provided by MPI\_COMM\_SPLIT, below. (*End of advice to users.*)

*Advice to implementors.* When calling MPI\_COMM\_DUP, all processes call with the same group (the group associated with the communicator). When calling MPI\_COMM\_CREATE, the processes provide the same group or disjoint subgroups. For both calls, it is theoretically possible to agree on a group-wide unique context with no communication. However, local execution of these functions requires use of a larger context name space and reduces error checking. Implementations may strike various compromises between these conflicting goals, such as bulk allocation of multiple contexts in one collective operation.

Important: If new communicators are created without synchronizing the processes involved then the communication system should be able to cope with messages arriving in a context that has not yet been allocated at the receiving process. (*End of advice to implementors.*)

If comm is an intercommunicator, then the output communicator is also an intercommunicator where the local group consists only of those processes contained in group (see Figure 6.1). The group argument should only contain those processes in the local group of the input intercommunicator that are to be a part of newcomm. All processes in the same local group of comm must specify the same value for group, i.e., the same members in the same order. If either group does not specify at least one process in the local group of the intercommunicator, or if the calling process is not included in the group, MPI\_COMM\_NULL is returned.

*Rationale.* In the case where either the left or right group is empty, a null communicator is returned instead of an intercommunicator with MPI\_GROUP\_EMPTY because the side with the empty group must return MPI\_COMM\_NULL. (*End of rationale.*)

**Example 6.1** The following example illustrates how the first node in the left side of an intercommunicator could be joined with all members on the right side of an intercommunicator to form a new intercommunicator.

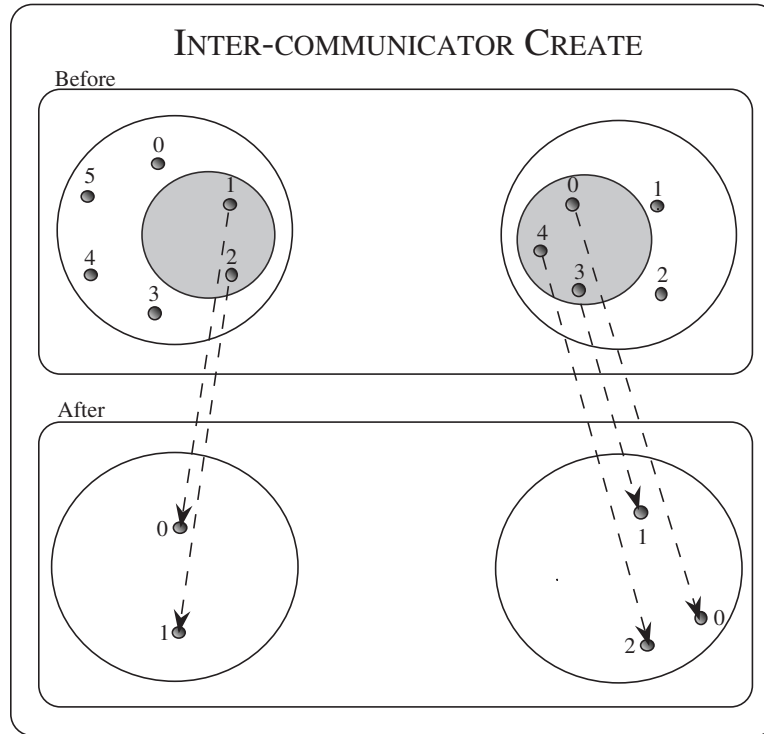


Figure 6.1: Intercommunicator create using `MPI_COMM_CREATE` extended to intercommunicators. The input groups are those in the grey circle.

```

MPI_Comm inter_comm, new_inter_comm;
MPI_Group local_group, group;
int      rank = 0; /* rank on left side to include in
                    new inter-comm */

/* Construct the original intercommunicator: "inter_comm" */
...

/* Construct the group of processes to be in new
intercommunicator */
if (/* I'm on the left side of the intercommunicator */) {
    MPI_Comm_group ( inter_comm, &local_group );
    MPI_Group_incl ( local_group, 1, &rank, &group );
    MPI_Group_free ( &local_group );
}
else
    MPI_Comm_group ( inter_comm, &group );

MPI_Comm_create ( inter_comm, group, &new_inter_comm );
MPI_Group_free( &group );

```



MPI\_COMM\_SPLIT(comm, color, key, newcomm)

IN	comm	communicator (handle)
IN	color	control of subset assignment (integer)
IN	key	control of rank assignment (integer)
OUT	newcomm	new communicator (handle)

```
int MPI_Comm_split(MPI_Comm comm, int color, int key, MPI_Comm *newcomm)
```

```
MPI_Comm_split(comm, color, key, newcomm, ierror) BIND(C)
```

```
    TYPE(MPI_Comm), INTENT(IN) :: comm
```

```
    INTEGER, INTENT(IN) :: color, key
```

```
    TYPE(MPI_Comm), INTENT(OUT) :: newcomm
```

```
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_COMM_SPLIT(COMM, COLOR, KEY, NEWCOMM, IERROR)
```

```
    INTEGER COMM, COLOR, KEY, NEWCOMM, IERROR
```

```
{MPI::Intercomm MPI::Intercomm::Split(int color, int key) const(binding
    deprecated, see Section 15.2) }
```

```
{MPI::Intracomm MPI::Intracomm::Split(int color, int key) const(binding
    deprecated, see Section 15.2) }
```

This function partitions the group associated with `comm` into disjoint subgroups, one for each value of `color`. Each subgroup contains all processes of the same color. Within each subgroup, the processes are ranked in the order defined by the value of the argument `key`, with ties broken according to their rank in the old group. A new communicator is created for each subgroup and returned in `newcomm`. A process may supply the color value `MPI_UNDEFINED`, in which case `newcomm` returns `MPI_COMM_NULL`. This is a collective call, but each process is permitted to provide different values for `color` and `key`.

With an intracommunicator `comm`, a call to `MPI_COMM_CREATE(comm, group, newcomm)` is equivalent to a call to `MPI_COMM_SPLIT(comm, color, key, newcomm)`, where processes that are members of their `group` argument provide `color = number of the group` (based on a unique numbering of all disjoint groups) and `key = rank in group`, and all processes that are not members of their `group` argument provide `color = MPI_UNDEFINED`.

The value of `color` must be non-negative.

*Advice to users.* This is an extremely powerful mechanism for dividing a single communicating group of processes into  $k$  subgroups, with  $k$  chosen implicitly by the user (by the number of colors asserted over all the processes). Each resulting communicator will be non-overlapping. Such a division could be useful for defining a hierarchy of computations, such as for multigrid, or linear algebra. For intracommunicators, `MPI_COMM_SPLIT` provides similar capability as `MPI_COMM_CREATE` to split a communicating group into disjoint subgroups. `MPI_COMM_SPLIT` is useful when some processes do not have complete information of the other members in their group, but all processes know (the color of) the group to which they belong. In this case, the MPI implementation discovers the other group members via communication. `MPI_COMM_CREATE` is useful when all processes have complete information

of the members of their group. In this case, MPI can avoid the extra communication required to discover group membership.

Multiple calls to `MPI_COMM_SPLIT` can be used to overcome the requirement that any call have no overlap of the resulting communicators (each process is of only one color per call). In this way, multiple overlapping communication structures can be created. Creative use of the `color` and `key` in such splitting operations is encouraged.

Note that, for a fixed color, the keys need not be unique. It is `MPI_COMM_SPLIT`'s responsibility to sort processes in ascending order according to this key, and to break ties in a consistent way. If all the keys are specified in the same way, then all the processes in a given color will have the relative rank order as they did in their parent group.

Essentially, making the key value zero for all processes of a given color means that one doesn't really care about the rank-order of the processes in the new communicator. (*End of advice to users.*)

*Rationale.* `color` is restricted to be non-negative, so as not to conflict with the value assigned to `MPI_UNDEFINED`. (*End of rationale.*)

The result of `MPI_COMM_SPLIT` on an intercommunicator is that those processes on the left with the same color as those processes on the right combine to create a new intercommunicator. The `key` argument describes the relative rank of processes on each side of the intercommunicator (see Figure 6.2). For those colors that are specified only on one side of the intercommunicator, `MPI_COMM_NULL` is returned. `MPI_COMM_NULL` is also returned to those processes that specify `MPI_UNDEFINED` as the color.

*Advice to users.* For intercommunicators, `MPI_COMM_SPLIT` is more general than `MPI_COMM_CREATE`. A single call to `MPI_COMM_SPLIT` can create a set of disjoint intercommunicators, while a call to `MPI_COMM_CREATE` creates only one. (*End of advice to users.*)

**Example 6.2** (Parallel client-server model). The following client code illustrates how clients on the left side of an intercommunicator could be assigned to a single server from a pool of servers on the right side of an intercommunicator.

```
/* Client code */
MPI_Comm multiple_server_comm;
MPI_Comm single_server_comm;
int      color, rank, num_servers;

/* Create intercommunicator with clients and servers:
   multiple_server_comm */
...

/* Find out the number of servers available */
MPI_Comm_remote_size ( multiple_server_comm, &num_servers );

/* Determine my color */
MPI_Comm_rank ( multiple_server_comm, &rank );
```

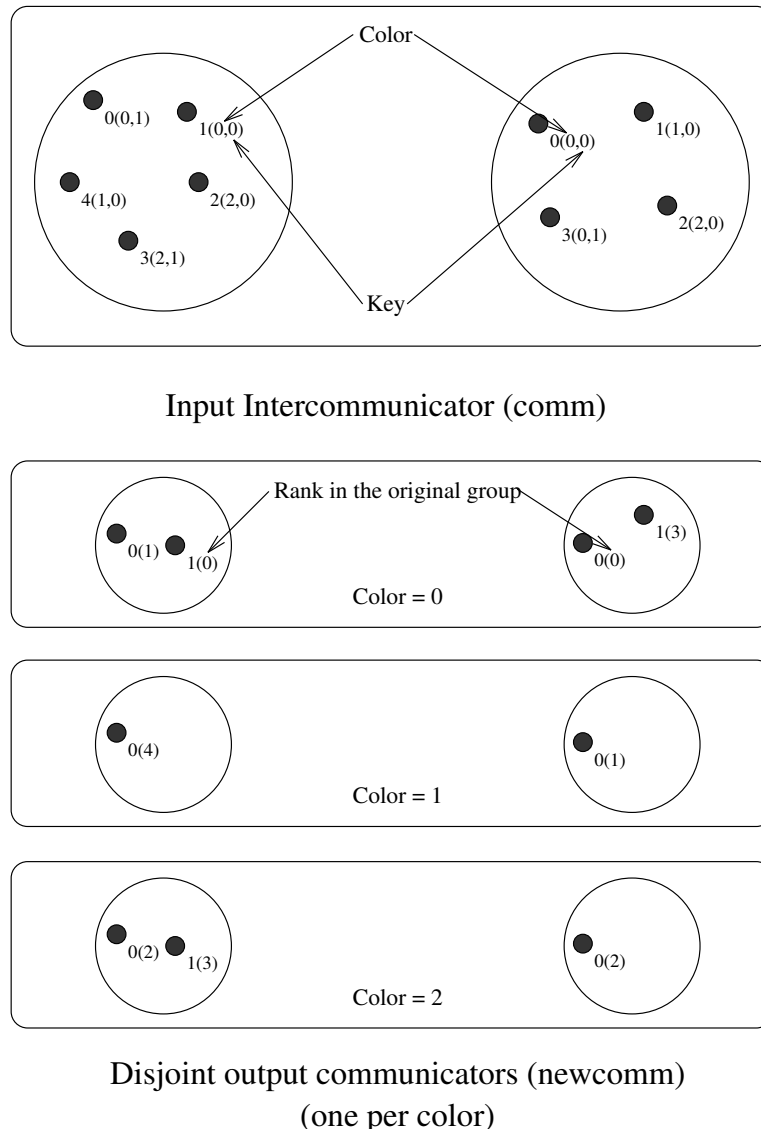


Figure 6.2: Intercommunicator construction achieved by splitting an existing intercommunicator with `MPI_COMM_SPLIT` extended to intercommunicators.

```

1      color = rank % num_servers;
2
3      /* Split the intercommunicator */
4      MPI_Comm_split ( multiple_server_comm, color, rank,
5                      &single_server_comm );
6

```

The following is the corresponding server code:

```

8      /* Server code */
9      MPI_Comm multiple_client_comm;
10     MPI_Comm single_server_comm;
11     int rank;
12
13     /* Create intercommunicator with clients and servers:
14        multiple_client_comm */
15     ...
16
17     /* Split the intercommunicator for a single server per group
18        of clients */
19     MPI_Comm_rank ( multiple_client_comm, &rank );
20     MPI_Comm_split ( multiple_client_comm, rank, 0,
21                     &single_server_comm );
22

```

### 6.4.3 Communicator Destructors

**MPI\_COMM\_FREE(comm)**

INOUT comm                                      communicator to be destroyed (handle)

```

30 int MPI_Comm_free(MPI_Comm *comm)

```

```

31 MPI_Comm_free(comm, ierror) BIND(C)
32     TYPE(MPI_Comm), INTENT(INOUT) :: comm
33     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
34

```

```

35 MPI_COMM_FREE(COMM, IERROR)
36     INTEGER COMM, IERROR
37

```

```

38 {void MPI::Comm::Free() (binding deprecated, see Section 15.2) }
39

```

This collective operation marks the communication object for deallocation. The handle is set to `MPI_COMM_NULL`. Any pending operations that use this communicator will complete normally; the object is actually deallocated only if there are no other active references to it. This call applies to intra- and inter-communicators. The delete callback functions for all cached attributes (see Section 6.7) are called in arbitrary order.

*Advice to implementors.* A reference-count mechanism may be used: the reference count is incremented by each call to `MPI_COMM_DUP`, and decremented by each call to `MPI_COMM_FREE`. The object is ultimately deallocated when the count reaches zero.

Though collective, it is anticipated that this operation will normally be implemented to be local, though a debugging version of an MPI library might choose to synchronize. (*End of advice to implementors.*)

## 6.5 Motivating Examples

### 6.5.1 Current Practice #1

Example #1a:

```
int main(int argc, char **argv)
{
    int me, size;
    ...
    MPI_Init ( &argc, &argv );
    MPI_Comm_rank (MPI_COMM_WORLD, &me);
    MPI_Comm_size (MPI_COMM_WORLD, &size);

    (void)printf ("Process %d size %d\n", me, size);
    ...
    MPI_Finalize();
}
```

Example #1a is a do-nothing program that initializes itself [legally](#), and refers to the “all” communicator, and prints a message. It terminates itself [legally](#) too. This example does not imply that MPI supports `printf`-like communication itself.

Example #1b (supposing that `size` is even):

```
int main(int argc, char **argv)
{
    int me, size;
    int SOME_TAG = 0;
    ...
    MPI_Init(&argc, &argv);

    MPI_Comm_rank(MPI_COMM_WORLD, &me); /* local */
    MPI_Comm_size(MPI_COMM_WORLD, &size); /* local */

    if((me % 2) == 0)
    {
        /* send unless highest-numbered process */
        if((me + 1) < size)
            MPI_Send(..., me + 1, SOME_TAG, MPI_COMM_WORLD);
    }
    else
        MPI_Recv(..., me - 1, SOME_TAG, MPI_COMM_WORLD, &status);

    ...
    MPI_Finalize();
}
```

Example #1b schematically illustrates message exchanges between “even” and “odd” processes in the “all” communicator.

### 6.5.2 Current Practice #2

```

1  int main(int argc, char **argv)
2  {
3      int me, count;
4      void *data;
5      ...
6
7      MPI_Init(&argc, &argv);
8      MPI_Comm_rank(MPI_COMM_WORLD, &me);
9
10     if(me == 0)
11     {
12         /* get input, create buffer ‘data’ */
13         ...
14     }
15
16     MPI_Bcast(data, count, MPI_BYTE, 0, MPI_COMM_WORLD);
17
18     ...
19
20     MPI_Finalize();
21 }

```

This example illustrates the use of a collective communication.

### 6.5.3 (Approximate) Current Practice #3

```

31  int main(int argc, char **argv)
32  {
33      int me, count, count2;
34      void *send_buf, *recv_buf, *send_buf2, *recv_buf2;
35      MPI_Group MPI_GROUP_WORLD, grprem;
36      MPI_Comm commslave;
37      static int ranks[] = {0};
38      ...
39      MPI_Init(&argc, &argv);
40      MPI_Comm_group(MPI_COMM_WORLD, &MPI_GROUP_WORLD);
41      MPI_Comm_rank(MPI_COMM_WORLD, &me); /* local */
42
43      MPI_Group_excl(MPI_GROUP_WORLD, 1, ranks, &grprem); /* local */
44      MPI_Comm_create(MPI_COMM_WORLD, grprem, &commslave);
45
46      if(me != 0)
47      {
48          /* compute on slave */

```

```

...
MPI_Reduce(send_buf,recv_buf,count, MPI_INT, MPI_SUM, 1, commslave);
...
MPI_Comm_free(&commslave);
}
/* zero falls through immediately to this reduce, others do later... */
MPI_Reduce(send_buf2, recv_buf2, count2,
           MPI_INT, MPI_SUM, 0, MPI_COMM_WORLD);

MPI_Group_free(&MPI_GROUP_WORLD);
MPI_Group_free(&grpem);
MPI_Finalize();
}

```

This example illustrates how a group consisting of all but the zeroth process of the “all” group is created, and then how a communicator is formed (`commslave`) for that new group. The new communicator is used in a collective call, and all processes execute a collective call in the `MPI_COMM_WORLD` context. This example illustrates how the two communicators (that inherently possess distinct contexts) protect communication. That is, communication in `MPI_COMM_WORLD` is insulated from communication in `commslave`, and vice versa.

In summary, “group safety” is achieved via communicators because distinct contexts within communicators are enforced to be unique on any process.

#### 6.5.4 Example #4

The following example is meant to illustrate “safety” between point-to-point and collective communication. MPI guarantees that a single communicator can do safe point-to-point and collective communication.

```

#define TAG_ARBITRARY 12345
#define SOME_COUNT    50

int main(int argc, char **argv)
{
    int me;
    MPI_Request request[2];
    MPI_Status status[2];
    MPI_Group MPI_GROUP_WORLD, subgroup;
    int ranks[] = {2, 4, 6, 8};
    MPI_Comm the_comm;
    ...
    MPI_Init(&argc, &argv);
    MPI_Comm_group(MPI_COMM_WORLD, &MPI_GROUP_WORLD);

    MPI_Group_incl(MPI_GROUP_WORLD, 4, ranks, &subgroup); /* local */
    MPI_Group_rank(subgroup, &me); /* local */

    MPI_Comm_create(MPI_COMM_WORLD, subgroup, &the_comm);
}

```

```

1      if(me != MPI_UNDEFINED)
2      {
3          MPI_Irecv(buff1, count, MPI_DOUBLE, MPI_ANY_SOURCE, TAG_ARBITRARY,
4                  the_comm, request);
5          MPI_Isend(buff2, count, MPI_DOUBLE, (me+1)%4, TAG_ARBITRARY,
6                  the_comm, request+1);
7          for(i = 0; i < SOME_COUNT; i++)
8              MPI_Reduce(..., the_comm);
9          MPI_Waitall(2, request, status);
10
11         MPI_Comm_free(&the_comm);
12     }
13
14     MPI_Group_free(&MPI_GROUP_WORLD);
15     MPI_Group_free(&subgroup);
16     MPI_Finalize();
17 }

```

### 6.5.5 Library Example #1

The main program:

```

22     int main(int argc, char **argv)
23     {
24         int done = 0;
25         user_lib_t *libh_a, *libh_b;
26         void *dataset1, *dataset2;
27         ...
28         MPI_Init(&argc, &argv);
29         ...
30         init_user_lib(MPI_COMM_WORLD, &libh_a);
31         init_user_lib(MPI_COMM_WORLD, &libh_b);
32         ...
33         user_start_op(libh_a, dataset1);
34         user_start_op(libh_b, dataset2);
35         ...
36         while(!done)
37         {
38             /* work */
39             ...
40             MPI_Reduce(..., MPI_COMM_WORLD);
41             ...
42             /* see if done */
43             ...
44         }
45         user_end_op(libh_a);
46         user_end_op(libh_b);
47
48

```



```

    uninit_user_lib(libh_a);
    uninit_user_lib(libh_b);
    MPI_Finalize();
}

```

The user library initialization code:

```

void init_user_lib(MPI_Comm comm, user_lib_t **handle)
{
    user_lib_t *save;

    user_lib_initsave(&save); /* local */
    MPI_Comm_dup(comm, &(save -> comm));

    /* other inits */
    ...

    *handle = save;
}

```

User start-up code:

```

void user_start_op(user_lib_t *handle, void *data)
{
    MPI_Irecv( ..., handle->comm, &(handle -> irecv_handle) );
    MPI_Isend( ..., handle->comm, &(handle -> isend_handle) );
}

```

User communication clean-up code:

```

void user_end_op(user_lib_t *handle)
{
    MPI_Status status;
    MPI_Wait(& handle -> isend_handle, &status);
    MPI_Wait(& handle -> irecv_handle, &status);
}

```

User object clean-up code:

```

void uninit_user_lib(user_lib_t *handle)
{
    MPI_Comm_free(&(handle -> comm));
    free(handle);
}

```

### 6.5.6 Library Example #2

The main program:

```

int main(int argc, char **argv)
{

```

```

1      int ma, mb;
2      MPI_Group MPI_GROUP_WORLD, group_a, group_b;
3      MPI_Comm comm_a, comm_b;
4
5      static int list_a[] = {0, 1};
6  #if defined(EXAMPLE_2B) | defined(EXAMPLE_2C)
7      static int list_b[] = {0, 2, 3};
8  #else/* EXAMPLE_2A */
9      static int list_b[] = {0, 2};
10 #endif
11     int size_list_a = sizeof(list_a)/sizeof(int);
12     int size_list_b = sizeof(list_b)/sizeof(int);
13
14     ...
15     MPI_Init(&argc, &argv);
16     MPI_Comm_group(MPI_COMM_WORLD, &MPI_GROUP_WORLD);
17
18     MPI_Group_incl(MPI_GROUP_WORLD, size_list_a, list_a, &group_a);
19     MPI_Group_incl(MPI_GROUP_WORLD, size_list_b, list_b, &group_b);
20
21     MPI_Comm_create(MPI_COMM_WORLD, group_a, &comm_a);
22     MPI_Comm_create(MPI_COMM_WORLD, group_b, &comm_b);
23
24     if(comm_a != MPI_COMM_NULL)
25         MPI_Comm_rank(comm_a, &ma);
26     if(comm_b != MPI_COMM_NULL)
27         MPI_Comm_rank(comm_b, &mb);
28
29     if(comm_a != MPI_COMM_NULL)
30         lib_call(comm_a);
31
32     if(comm_b != MPI_COMM_NULL)
33     {
34         lib_call(comm_b);
35         lib_call(comm_b);
36     }
37
38     if(comm_a != MPI_COMM_NULL)
39         MPI_Comm_free(&comm_a);
40     if(comm_b != MPI_COMM_NULL)
41         MPI_Comm_free(&comm_b);
42     MPI_Group_free(&group_a);
43     MPI_Group_free(&group_b);
44     MPI_Group_free(&MPI_GROUP_WORLD);
45     MPI_Finalize();
46 }

```

The library:

```

void lib_call(MPI_Comm comm)
{
    int me, done = 0;
    MPI_Status status;
    MPI_Comm_rank(comm, &me);
    if(me == 0)
        while(!done)
        {
            MPI_Recv(..., MPI_ANY_SOURCE, MPI_ANY_TAG, comm, &status);
            ...
        }
    else
    {
        /* work */
        MPI_Send(..., 0, ARBITRARY_TAG, comm);
        ....
    }
#ifdef EXAMPLE_2C
    /* include (resp, exclude) for safety (resp, no safety): */
    MPI_Barrier(comm);
#endif
}

```

The above example is really three examples, depending on whether or not one includes rank 3 in `list_b`, and whether or not a synchronize is included in `lib_call`. This example illustrates that, despite contexts, subsequent calls to `lib_call` with the same context need not be safe from one another (colloquially, “back-masking”). Safety is realized if the `MPI_Barrier` is added. What this demonstrates is that libraries have to be written carefully, even with contexts. When rank 3 is excluded, then the synchronize is not needed to get safety from back masking.

Algorithms like “reduce” and “allreduce” have strong enough source selectivity properties so that they are inherently okay (no backmasking), provided that MPI provides basic guarantees. So are multiple calls to a typical tree-broadcast algorithm with the same root or different roots (see [58]). Here we rely on two guarantees of MPI: pairwise ordering of messages between processes in the same context, and source selectivity — deleting either feature removes the guarantee that backmasking cannot be required.

Algorithms that try to do non-deterministic broadcasts or other calls that include wildcard operations will not generally have the good properties of the deterministic implementations of “reduce,” “allreduce,” and “broadcast.” Such algorithms would have to utilize the monotonically increasing tags (within a communicator scope) to keep things straight.

All of the foregoing is a supposition of “collective calls” implemented with point-to-point operations. MPI implementations may or may not implement collective calls using point-to-point operations. These algorithms are used to illustrate the issues of correctness and safety, independent of how MPI implements its collective calls. See also Section 6.9.

## 6.6 Inter-Communication

This section introduces the concept of inter-communication and describes the portions of MPI that support it. It describes support for writing programs that contain user-level servers.

All communication described thus far has involved communication between processes that are members of the same group. This type of communication is called “intra-communication” and the communicator used is called an “intra-communicator,” as we have noted earlier in the chapter.

In modular and multi-disciplinary applications, different process groups execute distinct modules and processes within different modules communicate with one another in a pipeline or a more general module graph. In these applications, the most natural way for a process to specify a target process is by the rank of the target process within the target group. In applications that contain internal user-level servers, each server may be a process group that provides services to one or more clients, and each client may be a process group that uses the services of one or more servers. It is again most natural to specify the target process by rank within the target group in these applications. This type of communication is called “inter-communication” and the communicator used is called an “inter-communicator,” as introduced earlier.

An inter-communication is a point-to-point communication between processes in different groups. The group containing a process that initiates an inter-communication operation is called the “local group,” that is, the sender in a send and the receiver in a receive. The group containing the target process is called the “remote group,” that is, the receiver in a send and the sender in a receive. As in intra-communication, the target process is specified using a (communicator, rank) pair. Unlike intra-communication, the rank is relative to a second, remote group.

All inter-communicator constructors are blocking and require that the local and remote groups be disjoint.

*Advice to users.* The groups must be disjoint for several reasons. Primarily, this is the intent of the intercommunicators — to provide a communicator for communication between disjoint groups. This is reflected in the definition of `MPI_INTERCOMM_MERGE`, which allows the user to control the ranking of the processes in the created intracommunicator; this ranking makes little sense if the groups are not disjoint. In addition, the natural extension of collective operations to inter-communicators makes the most sense when the groups are disjoint. (*End of advice to users.*)

Here is a summary of the properties of inter-communication and inter-communicators:

- The syntax of point-to-point and collective communication is the same for both inter- and intra-communication. The same communicator can be used both for send and for receive operations.
- A target process is addressed by its rank in the remote group, both for sends and for receives.
- Communications using an inter-communicator are guaranteed not to conflict with any communications that use a different communicator.
- A communicator will provide either intra- or inter-communication, never both.

The routine `MPI_COMM_TEST_INTER` may be used to determine if a communicator is an inter- or intra-communicator. Inter-communicators can be used as arguments to some of the other communicator access routines. Inter-communicators cannot be used as input to some of the constructor routines for intra-communicators (for instance, `MPI_CART_CREATE`).

*Advice to implementors.* For the purpose of point-to-point communication, communicators can be represented in each process by a tuple consisting of:

**group**  
**send\_context**  
**receive\_context**  
**source**

For inter-communicators, **group** describes the remote group, and **source** is the rank of the process in the local group. For intra-communicators, **group** is the communicator group (remote=local), **source** is the rank of the process in this group, and **send context** and **receive context** are identical. A group can be represented by a rank-to-absolute-address translation table.

The inter-communicator cannot be discussed sensibly without considering processes in both the local and remote groups. Imagine a process **P** in group  $\mathcal{P}$ , which has an inter-communicator  $C_{\mathcal{P}}$ , and a process **Q** in group  $\mathcal{Q}$ , which has an inter-communicator  $C_{\mathcal{Q}}$ . Then

- $C_{\mathcal{P}}.\text{group}$  describes the group  $\mathcal{Q}$  and  $C_{\mathcal{Q}}.\text{group}$  describes the group  $\mathcal{P}$ .
- $C_{\mathcal{P}}.\text{send\_context} = C_{\mathcal{Q}}.\text{receive\_context}$  and the context is unique in  $\mathcal{Q}$ ;  
 $C_{\mathcal{P}}.\text{receive\_context} = C_{\mathcal{Q}}.\text{send\_context}$  and this context is unique in  $\mathcal{P}$ .
- $C_{\mathcal{P}}.\text{source}$  is rank of **P** in  $\mathcal{P}$  and  $C_{\mathcal{Q}}.\text{source}$  is rank of **Q** in  $\mathcal{Q}$ .

Assume that **P** sends a message to **Q** using the inter-communicator. Then **P** uses the **group** table to find the absolute address of **Q**; **source** and **send\_context** are appended to the message.

Assume that **Q** posts a receive with an explicit source argument using the inter-communicator. Then **Q** matches **receive\_context** to the message context and source argument to the message source.

The same algorithm is appropriate for intra-communicators as well.

In order to support inter-communicator accessors and constructors, it is necessary to supplement this model with additional structures, that store information about the local communication group, and additional safe contexts. (*End of advice to implementors.*)

### 6.6.1 Inter-communicator Accessors

`MPI_COMM_TEST_INTER(comm, flag)`

IN	comm	communicator (handle)
OUT	flag	(logical)

```
1 int MPI_Comm_test_inter(MPI_Comm comm, int *flag)
```

```
2 MPI_Comm_test_inter(comm, flag, ierror) BIND(C)
3   TYPE(MPI_Comm), INTENT(IN) :: comm
4   LOGICAL, INTENT(OUT) :: flag
5   INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
7 MPI_COMM_TEST_INTER(COMM, FLAG, IERROR)
8   INTEGER COMM, IERROR
9   LOGICAL FLAG
```

```
11 {bool MPI::Comm::Is_inter() const(binding deprecated, see Section 15.2) }
```

This local routine allows the calling process to determine if a communicator is an inter-communicator or an intra-communicator. It returns true if it is an inter-communicator, otherwise false.

When an inter-communicator is used as an input argument to the communicator accessors described above under intra-communication, the following table describes behavior.

MPI_COMM_SIZE	returns the size of the local group.
MPI_COMM_GROUP	returns the local group.
MPI_COMM_RANK	returns the rank in the local group

Table 6.1: MPI\_COMM\_\* Function Behavior (in Inter-Communication Mode)

Furthermore, the operation MPI\_COMM\_COMPARE is valid for inter-communicators. Both communicators must be either intra- or inter-communicators, or else MPI\_UNEQUAL results. Both corresponding local and remote groups must compare correctly to get the results MPI\_CONGRUENT and MPI\_SIMILAR. In particular, it is possible for MPI\_SIMILAR to result because either the local or remote groups were similar but not identical.

The following accessors provide consistent access to the remote group of an inter-communicator:

The following are all local operations.

```
35 MPI_COMM_REMOTE_SIZE(comm, size)
```

IN	comm	inter-communicator (handle)
OUT	size	number of processes in the remote group of comm (integer)

```
41 int MPI_Comm_remote_size(MPI_Comm comm, int *size)
```

```
42 MPI_Comm_remote_size(comm, size, ierror) BIND(C)
43   TYPE(MPI_Comm), INTENT(IN) :: comm
44   INTEGER, INTENT(OUT) :: size
45   INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
47 MPI_COMM_REMOTE_SIZE(COMM, SIZE, IERROR)
48   INTEGER COMM, SIZE, IERROR
```

```
{int MPI::Intercomm::Get_remote_size() const(binding deprecated, see Section 15.2)
    }
```

```
MPI_COMM_REMOTE_GROUP(comm, group)
```

```
IN      comm      inter-communicator (handle)
OUT     group      remote group corresponding to comm (handle)
```

```
int MPI_Comm_remote_group(MPI_Comm comm, MPI_Group *group)
```

```
MPI_Comm_remote_group(comm, group, ierror) BIND(C)
```

```
    TYPE(MPI_Comm), INTENT(IN) :: comm
```

```
    TYPE(MPI_Group), INTENT(OUT) :: group
```

```
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_COMM_REMOTE_GROUP(COMM, GROUP, IERROR)
```

```
    INTEGER COMM, GROUP, IERROR
```

```
{MPI::Group MPI::Intercomm::Get_remote_group() const(binding deprecated, see
    Section 15.2) }
```

*Rationale.* Symmetric access to both the local and remote groups of an inter-communicator is important, so this function, as well as `MPI_COMM_REMOTE_SIZE` have been provided. (*End of rationale.*)

### 6.6.2 Inter-communicator Operations

This section introduces four blocking inter-communicator operations.

`MPI_INTERCOMM_CREATE` is used to bind two intra-communicators into an inter-communicator; the function `MPI_INTERCOMM_MERGE` creates an intra-communicator by merging the local and remote groups of an inter-communicator. The functions `MPI_COMM_DUP` and `MPI_COMM_FREE`, introduced previously, duplicate and free an inter-communicator, respectively.

Overlap of local and remote groups that are bound into an inter-communicator is prohibited. If there is overlap, then the program is erroneous and is likely to deadlock. (If a process is multithreaded, and MPI calls block only a thread, rather than a process, then “dual membership” can be supported. It is then the user’s responsibility to make sure that calls on behalf of the two “roles” of a process are executed by two independent threads.)

The function `MPI_INTERCOMM_CREATE` can be used to create an inter-communicator from two existing intra-communicators, in the following situation: At least one selected member from each group (the “group leader”) has the ability to communicate with the selected member from the other group; that is, a “peer” communicator exists to which both leaders belong, and each leader knows the rank of the other leader in this peer communicator. Furthermore, members of each group know the rank of their leader.

Construction of an inter-communicator from two intra-communicators requires separate collective operations in the local group and in the remote group, as well as a point-to-point communication between a process in the local group and a process in the remote group.

In standard MPI implementations (with static process allocation at initialization), the `MPI_COMM_WORLD` communicator (or preferably a dedicated duplicate thereof) can be this peer communicator. For applications that have used `spawn` or `join`, it may be necessary to first create an intracommunicator to be used as peer.

The application topology functions described in Chapter 7 do not apply to inter-communicators. Users that require this capability should utilize `MPI_INTERCOMM_MERGE` to build an intra-communicator, then apply the graph or cartesian topology capabilities to that intra-communicator, creating an appropriate topology-oriented intra-communicator. Alternatively, it may be reasonable to devise one’s own application topology mechanisms for this case, without loss of generality.

```
MPI_INTERCOMM_CREATE(local_comm, local_leader, peer_comm, remote_leader, tag,
                      newintercomm)
```

IN	<code>local_comm</code>	local intra-communicator (handle)
IN	<code>local_leader</code>	rank of local group leader in <code>local_comm</code> (integer)
IN	<code>peer_comm</code>	“peer” communicator; significant only at the <code>local_leader</code> (handle)
IN	<code>remote_leader</code>	rank of remote group leader in <code>peer_comm</code> ; significant only at the <code>local_leader</code> (integer)
IN	<code>tag</code>	“safe” tag (integer)
OUT	<code>newintercomm</code>	new inter-communicator (handle)

```
int MPI_Intercomm_create(MPI_Comm local_comm, int local_leader,
                        MPI_Comm peer_comm, int remote_leader, int tag,
                        MPI_Comm *newintercomm)
```

```
MPI_Intercomm_create(local_comm, local_leader, peer_comm, remote_leader,
                    tag, newintercomm, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: local_comm, peer_comm
    INTEGER, INTENT(IN) :: local_leader, remote_leader, tag
    TYPE(MPI_Comm), INTENT(OUT) :: newintercomm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_INTERCOMM_CREATE(LOCAL_COMM, LOCAL_LEADER, PEER_COMM, REMOTE_LEADER,
                    TAG, NEWINTERCOMM, IERROR)
    INTEGER LOCAL_COMM, LOCAL_LEADER, PEER_COMM, REMOTE_LEADER, TAG,
    NEWINTERCOMM, IERROR
```

```
{MPI::Intercomm MPI::Intracomm::Create_intercomm(int local_leader, const
    MPI::Comm& peer_comm, int remote_leader, int tag) const(binding
    deprecated, see Section 15.2) }
```

This call creates an inter-communicator. It is collective over the union of the local and remote groups. Processes should provide identical `local_comm` and `local_leader` arguments within each group. Wildcards are not permitted for `remote_leader`, `local_leader`, and `tag`.

This call uses point-to-point communication with communicator `peer_comm`, and with tag `tag` between the leaders. Thus, care must be taken that there be



no pending communication on `peer_comm` that could interfere with this communication.

*Advice to users.* We recommend using a dedicated peer communicator, such as a duplicate of `MPI_COMM_WORLD`, to avoid trouble with peer communicators. (*End of advice to users.*)

`MPI_INTERCOMM_MERGE(intercomm, high, newintracomm)`

IN	intercomm	Inter-Communicator (handle)
IN	high	(logical)
OUT	newintracomm	new intra-communicator (handle)

```
int MPI_Intercomm_merge(MPI_Comm intercomm, int high,
                        MPI_Comm *newintracomm)
```

```
MPI_Intercomm_merge(intercomm, high, newintracomm, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: intercomm
    LOGICAL, INTENT(IN) :: high
    TYPE(MPI_Comm), INTENT(OUT) :: newintracomm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_INTERCOMM_MERGE(INTERCOMM, HIGH, NEWINTRACOMM, IERROR)
    INTEGER INTERCOMM, NEWINTRACOMM, IERROR
    LOGICAL HIGH
```

```
{MPI::Intracomm MPI::Intercomm::Merge(bool high) const(binding deprecated, see
    Section 15.2) }
```

This function creates an intra-communicator from the union of the two groups that are associated with `intercomm`. All processes should provide the same `high` value within each of the two groups. If processes in one group provided the value `high = false` and processes in the other group provided the value `high = true` then the union orders the “low” group before the “high” group. If all processes provided the same `high` argument then the order of the union is arbitrary. This call is blocking and collective within the union of the two groups.

The error handler on the new intercommunicator in each process is inherited from the communicator that contributes the local group. Note that this can result in different processes in the same communicator having different error handlers.

*Advice to implementors.* The implementation of `MPI_INTERCOMM_MERGE`, `MPI_COMM_FREE` and `MPI_COMM_DUP` are similar to the implementation of `MPI_INTERCOMM_CREATE`, except that contexts private to the input inter-communicator are used for communication between group leaders rather than contexts inside a bridge communicator. (*End of advice to implementors.*)

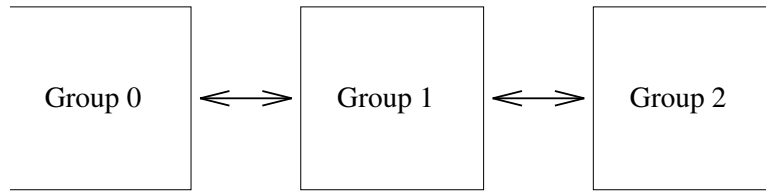


Figure 6.3: Three-group pipeline[ticket0.][.]

### 6.6.3 Inter-Communication Examples

#### Example 1: Three-Group “Pipeline”

Groups 0 and 1 communicate. Groups 1 and 2 communicate. Therefore, group 0 requires one inter-communicator, group 1 requires two inter-communicators, and group 2 requires 1 inter-communicator.

```

int main(int argc, char **argv)
{
    MPI_Comm    myComm;          /* intra-communicator of local sub-group */
    MPI_Comm    myFirstComm;     /* inter-communicator */
    MPI_Comm    mySecondComm;   /* second inter-communicator (group 1 only) */
    int membershipKey;
    int rank;

    MPI_Init(&argc, &argv);
    MPI_Comm_rank(MPI_COMM_WORLD, &rank);

    /* User code must generate membershipKey in the range [0, 1, 2] */
    membershipKey = rank % 3;

    /* Build intra-communicator for local sub-group */
    MPI_Comm_split(MPI_COMM_WORLD, membershipKey, rank, &myComm);

    /* Build inter-communicators. Tags are hard-coded. */
    if (membershipKey == 0)
    {
        /* Group 0 communicates with group 1. */
        MPI_Intercomm_create( myComm, 0, MPI_COMM_WORLD, 1,
                             1, &myFirstComm);
    }
    else if (membershipKey == 1)
    {
        /* Group 1 communicates with groups 0 and 2. */
        MPI_Intercomm_create( myComm, 0, MPI_COMM_WORLD, 0,
                             1, &myFirstComm);
        MPI_Intercomm_create( myComm, 0, MPI_COMM_WORLD, 2,
                             12, &mySecondComm);
    }
    else if (membershipKey == 2)
    {
        /* Group 2 communicates with group 1. */

```

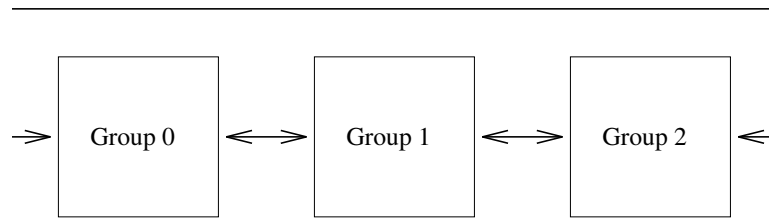


Figure 6.4: Three-group ring[ticket0.][.]

```

    MPI_Intercomm_create( myComm, 0, MPI_COMM_WORLD, 1,
                          12, &myFirstComm);
}

/* Do work ... */

switch(membershipKey) /* free communicators appropriately */
{
case 1:
    MPI_Comm_free(&mySecondComm);
case 0:
case 2:
    MPI_Comm_free(&myFirstComm);
    break;
}

MPI_Finalize();
}

```

## Example 2: Three-Group “Ring”

Groups 0 and 1 communicate. Groups 1 and 2 communicate. Groups 0 and 2 communicate. Therefore, each requires two inter-communicators.

```

int main(int argc, char **argv)
{
    MPI_Comm  myComm;      /* intra-communicator of local sub-group */
    MPI_Comm  myFirstComm; /* inter-communicators */
    MPI_Comm  mySecondComm;
    MPI_Status status;
    int membershipKey;
    int rank;

    MPI_Init(&argc, &argv);
    MPI_Comm_rank(MPI_COMM_WORLD, &rank);
    ...

    /* User code must generate membershipKey in the range [0, 1, 2] */
    membershipKey = rank % 3;
}

```

```

1
2      /* Build intra-communicator for local sub-group */
3      MPI_Comm_split(MPI_COMM_WORLD, membershipKey, rank, &myComm);
4
5      /* Build inter-communicators.  Tags are hard-coded. */
6      if (membershipKey == 0)
7      {
8          /* Group 0 communicates with groups 1 and 2. */
9          MPI_Intercomm_create( myComm, 0, MPI_COMM_WORLD, 1,
10                             1, &myFirstComm);
11          MPI_Intercomm_create( myComm, 0, MPI_COMM_WORLD, 2,
12                             2, &mySecondComm);
13      }
14      else if (membershipKey == 1)
15      {
16          /* Group 1 communicates with groups 0 and 2. */
17          MPI_Intercomm_create( myComm, 0, MPI_COMM_WORLD, 0,
18                             1, &myFirstComm);
19          MPI_Intercomm_create( myComm, 0, MPI_COMM_WORLD, 2,
20                             12, &mySecondComm);
21      }
22      else if (membershipKey == 2)
23      {
24          /* Group 2 communicates with groups 0 and 1. */
25          MPI_Intercomm_create( myComm, 0, MPI_COMM_WORLD, 0,
26                             2, &myFirstComm);
27          MPI_Intercomm_create( myComm, 0, MPI_COMM_WORLD, 1,
28                             12, &mySecondComm);
29      }
30
31      /* Do some work ... */
32
33      /* Then free communicators before terminating... */
34      MPI_Comm_free(&myFirstComm);
35      MPI_Comm_free(&mySecondComm);
36      MPI_Comm_free(&myComm);
37      MPI_Finalize();
38  }

```

## 6.7 Caching

MPI provides a “caching” facility that allows an application to attach arbitrary pieces of information, called **attributes**, to three kinds of MPI objects, communicators, windows and datatypes. More precisely, the caching facility allows a portable library to do the following:

- pass information between calls by associating it with an MPI intra- or inter-communicator, window or datatype,
- quickly retrieve that information, and
- be guaranteed that out-of-date information is never retrieved, even if the object is freed and its handle subsequently reused by MPI.

The caching capabilities, in some form, are required by built-in MPI routines such as collective communication and application topology. Defining an interface to these capabilities as part of the MPI standard is valuable because it permits routines like collective communication and application topologies to be implemented as portable code, and also because it makes MPI more extensible by allowing user-written routines to use standard MPI calling sequences.

*Advice to users.* The communicator `MPI_COMM_SELF` is a suitable choice for posting process-local attributes, via this attributing-caching mechanism. (*End of advice to users.*)

*Rationale.* In one extreme one can allow caching on all opaque handles. The other extreme is to only allow it on communicators. Caching has a cost associated with it and should only be allowed when it is clearly needed and the increased cost is modest. This is the reason that windows and datatypes were added but not other handles. (*End of rationale.*)

One difficulty is the potential for size differences between Fortran integers and C pointers. To overcome this problem with attribute caching on communicators, functions are also given for this case. The functions to cache on datatypes and windows also address this issue. For a general discussion of the address size problem, see Section 16.3.6.

*Advice to implementors.* High-quality implementations should raise an error when a keyval that was created by a call to `MPI_XXX_CREATE_KEYVAL` is used with an object of the wrong type with a call to `MPI_YYY_GET_ATTR`, `MPI_YYY_SET_ATTR`, `MPI_YYY_DELETE_ATTR`, or `MPI_YYY_FREE_KEYVAL`. To do so, it is necessary to maintain, with each keyval, information on the type of the associated user function. (*End of advice to implementors.*)

### 6.7.1 Functionality

Attributes can be attached to communicators, windows, and datatypes. Attributes are local to the process and specific to the communicator to which they are attached. Attributes are not propagated by MPI from one communicator to another except when the communicator is duplicated using `MPI_COMM_DUP` (and even then the application must give specific permission through callback functions for the attribute to be copied).

*Advice to users.* Attributes in C are of type `void *`. Typically, such an attribute will be a pointer to a structure that contains further information, or a handle to an MPI object. In Fortran, attributes are of type `INTEGER`. Such attribute can be a handle to an MPI object, or just an integer-valued attribute. (*End of advice to users.*)

*Advice to implementors.* Attributes are scalar values, equal in size to, or larger than a C-language pointer. Attributes can always hold an MPI handle. (*End of advice to implementors.*)

The caching interface defined here requires that attributes be stored by MPI opaquely within a communicator, window, and datatype. Accessor functions include the following:

- obtain a key value (used to identify an attribute); the user specifies “callback” functions by which MPI informs the application when the communicator is destroyed or copied.
- store and retrieve the value of an attribute;

*Advice to implementors.* Caching and callback functions are only called synchronously, in response to explicit application requests. This avoids problems that result from repeated crossings between user and system space. (This synchronous calling rule is a general property of MPI.)

The choice of key values is under control of MPI. This allows MPI to optimize its implementation of attribute sets. It also avoids conflict between independent modules caching information on the same communicators.

A much smaller interface, consisting of just a callback facility, would allow the entire caching facility to be implemented by portable code. However, with the minimal callback interface, some form of table searching is implied by the need to handle arbitrary communicators. In contrast, the more complete interface defined here permits rapid access to attributes through the use of pointers in communicators (to find the attribute table) and cleverly chosen key values (to retrieve individual attributes). In light of the efficiency “hit” inherent in the minimal interface, the more complete interface defined here is seen to be superior. (*End of advice to implementors.*)

MPI provides the following services related to caching. They are all process local.

## 6.7.2 Communicators

Functions for caching on communicators are:

`MPI_COMM_CREATE_KEYVAL(comm_copy_attr_fn, comm_delete_attr_fn, comm_keyval, extra_state)`

IN	<code>comm_copy_attr_fn</code>	copy callback function for <code>comm_keyval</code> (function)
IN	<code>comm_delete_attr_fn</code>	delete callback function for <code>comm_keyval</code> (function)
OUT	<code>comm_keyval</code>	key value for future access (integer)
IN	<code>extra_state</code>	extra state for callback functions

```
int MPI_Comm_create_keyval(MPI_Comm_copy_attr_function *comm_copy_attr_fn,
                           MPI_Comm_delete_attr_function *comm_delete_attr_fn,
                           int *comm_keyval, void *extra_state)
```

```
MPI_Comm_create_keyval(comm_copy_attr_fn, comm_delete_attr_fn, comm_keyval,
                        extra_state, ierror) BIND(C)
PROCEDURE(MPI_Comm_copy_attr_function) :: comm_copy_attr_fn
PROCEDURE(MPI_Comm_delete_attr_function) :: comm_delete_attr_fn
INTEGER, INTENT(OUT) :: comm_keyval
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```

MPI_COMM_CREATE_KEYVAL(COMM_COPY_ATTR_FN, COMM_DELETE_ATTR_FN, COMM_KEYVAL,
                        EXTRA_STATE, IERROR)
EXTERNAL COMM_COPY_ATTR_FN, COMM_DELETE_ATTR_FN
INTEGER COMM_KEYVAL, IERROR
INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE

{static int MPI::Comm::Create_keyval(MPI::Comm::Copy_attr_function*
                                     comm_copy_attr_fn,
                                     MPI::Comm::Delete_attr_function* comm_delete_attr_fn,
                                     void* extra_state) (binding deprecated, see Section 15.2) }
```

Generates a new attribute key. Keys are locally unique in a process, and opaque to user, though they are explicitly stored in integers. Once allocated, the key value can be used to associate attributes and access them on any locally defined communicator.

This function replaces `MPI_KEYVAL_CREATE`, whose use is deprecated. The C binding is identical. The Fortran binding differs in that `extra_state` is an address-sized integer. Also, the copy and delete callback functions have Fortran bindings that are consistent with address-sized attributes.

The C callback functions are:

```

typedef int MPI_Comm_copy_attr_function(MPI_Comm oldcomm, int comm_keyval,
                                         void *extra_state, void *attribute_val_in,
                                         void *attribute_val_out, int *flag);
```

and

```

typedef int MPI_Comm_delete_attr_function(MPI_Comm comm, int comm_keyval,
                                         void *attribute_val, void *extra_state);
```

which are the same as the MPI-1.1 calls but with a new name. The old names are deprecated.

With the `mpi_f08` module, the Fortran callback functions are:

**ABSTRACT INTERFACE**

```

SUBROUTINE MPI_Comm_copy_attr_function(oldcomm, comm_keyval, extra_state,
attribute_val_in, attribute_val_out, flag, ierror) BIND(C)
  TYPE(MPI_Comm) :: oldcomm
  INTEGER :: comm_keyval, ierror
  INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state, attribute_val_in,
attribute_val_out
  LOGICAL :: flag
```

and

**ABSTRACT INTERFACE**

```

SUBROUTINE MPI_Comm_delete_attr_function(comm, comm_keyval,
attribute_val, extra_state, ierror) BIND(C)
  TYPE(MPI_Comm) :: comm
  INTEGER :: comm_keyval, ierror
  INTEGER(KIND=MPI_ADDRESS_KIND) :: attribute_val, extra_state
```

[The] With the `mpi` module and `mpif.h`, the Fortran callback functions are:

```

SUBROUTINE COMM_COPY_ATTR_[FN]FUNCTION(OLDCOMM, COMM_KEYVAL, EXTRA_STATE,
```

```

1      ATTRIBUTE_VAL_IN, ATTRIBUTE_VAL_OUT, FLAG, IERROR)
2      INTEGER OLDCOMM, COMM_KEYVAL, IERROR
3      INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE, ATTRIBUTE_VAL_IN,
4          ATTRIBUTE_VAL_OUT
5      LOGICAL FLAG
6
7  and
ticket250-V. 8  SUBROUTINE COMM_DELETE_ATTR_ [FN] FUNCTION(COMM, COMM_KEYVAL, ATTRIBUTE_VAL,
9          EXTRA_STATE, IERROR)
10     INTEGER COMM, COMM_KEYVAL, IERROR
11     INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL, EXTRA_STATE

```

The C++ callbacks are:

```

14 {typedef int MPI::Comm::Copy_attr_function(const MPI::Comm& oldcomm,
15     int comm_keyval, void* extra_state, void* attribute_val_in,
16     void* attribute_val_out, bool& flag); (binding deprecated, see
17     Section 15.2)}
18
19 and
20 {typedef int MPI::Comm::Delete_attr_function(MPI::Comm& comm,
21     int comm_keyval, void* attribute_val, void* extra_state);
22     (binding deprecated, see Section 15.2)}

```

The `comm_copy_attr_fn` function is invoked when a communicator is duplicated by `MPI_COMM_DUP`. `comm_copy_attr_fn` should be of type `MPI_Comm_copy_attr_function`. The copy callback function is invoked for each key value in `oldcomm` in arbitrary order. Each call to the copy callback is made with a key value and its corresponding attribute. If it returns `flag = 0` or `.FALSE.`, then the attribute is deleted in the duplicated communicator. Otherwise (`flag = 1` or `.TRUE.`), the new attribute value is set to the value returned in `attribute_val_out`. The function returns `MPI_SUCCESS` on success and an error code on failure (in which case `MPI_COMM_DUP` will fail).

The argument `comm_copy_attr_fn` may be specified as `MPI_COMM_NULL_COPY_FN` or `MPI_COMM_DUP_FN` from either C, C++, or Fortran. `MPI_COMM_NULL_COPY_FN` is a function that does nothing other than returning `flag = 0` or `.FALSE.` (depending on whether the keyval was created with a C/C++ or Fortran binding to `MPI_COMM_CREATE_KEYVAL`) and `MPI_SUCCESS`. `MPI_COMM_DUP_FN` is a simple-minded copy function that sets `flag = 1` or `.TRUE.`, returns the value of `attribute_val_in` in `attribute_val_out`, and returns `MPI_SUCCESS`. These replace the MPI-1 predefined callbacks `MPI_NULL_COPY_FN` and `MPI_DUP_FN`, whose use is deprecated.

*Advice to users.* Even though both formal arguments `attribute_val_in` and `attribute_val_out` are of type `void *`, their usage differs. The C copy function is passed by MPI in `attribute_val_in` the *value* of the attribute, and in `attribute_val_out` the *address* of the attribute, so as to allow the function to return the (new) attribute value. The use of type `void *` for both is to avoid messy type casts.

A valid copy function is one that completely duplicates the information by making a full duplicate copy of the data structures implied by an attribute; another might just make another reference to that data structure, while using a reference-count



mechanism. Other types of attributes might not copy at all (they might be specific to oldcomm only). (*End of advice to users.*)

*Advice to implementors.* A C interface should be assumed for copy and delete functions associated with key values created in C; a Fortran calling interface should be assumed for key values created in Fortran. (*End of advice to implementors.*)

Analogous to `comm_copy_attr_fn` is a callback deletion function, defined as follows. The `comm_delete_attr_fn` function is invoked when a communicator is deleted by `MPI_COMM_FREE` or when a call is made explicitly to `MPI_COMM_DELETE_ATTR`. `comm_delete_attr_fn` should be of type `MPI_Comm_delete_attr_function`.

This function is called by `MPI_COMM_FREE`, `MPI_COMM_DELETE_ATTR`, and `MPI_COMM_SET_ATTR` to do whatever is needed to remove an attribute. The function returns `MPI_SUCCESS` on success and an error code on failure (in which case `MPI_COMM_FREE` will fail).

The argument `comm_delete_attr_fn` may be specified as `MPI_COMM_NULL_DELETE_FN` from either C, C++, or Fortran. `MPI_COMM_NULL_DELETE_FN` is a function that does nothing, other than returning `MPI_SUCCESS`. `MPI_COMM_NULL_DELETE_FN` replaces `MPI_NULL_DELETE_FN`, whose use is deprecated.

If an attribute copy function or attribute delete function returns other than `MPI_SUCCESS`, then the call that caused it to be invoked (for example, `MPI_COMM_FREE`), is erroneous.

The special key value `MPI_KEYVAL_INVALID` is never returned by `MPI_KEYVAL_CREATE`. Therefore, it can be used for static initialization of key values.

*Advice to implementors.* To be able to use the predefined C functions `MPI_COMM_NULL_COPY_FN` or `MPI_COMM_DUP_FN` as `comm_copy_attr_fn` argument and/or `MPI_COMM_NULL_DELETE_FN` as the `comm_delete_attr_fn` argument in a call to the C++ routine `MPI::Comm::Create_keyval`, this routine may be overloaded with 3 additional routines that accept the C functions as the first, the second, or both input arguments (instead of an argument that matches the C++ prototype). (*End of advice to implementors.*)

*Advice to users.* If a user wants to write a “wrapper” routine that internally calls `MPI::Comm::Create_keyval` and `comm_copy_attr_fn` and/or `comm_delete_attr_fn` are arguments of this wrapper routine, and if this wrapper routine should be callable with both user-defined C++ copy and delete functions and with the predefined C functions, then the same overloading as described above in the advice to implementors may be necessary. (*End of advice to users.*)

*Advice to implementors.* The predefined Fortran functions `MPI_COMM_NULL_COPY_FN`, `MPI_COMM_DUP_FN`, and `MPI_COMM_NULL_DELETE_FN` are defined in the `mpi` module (and `mpif.h`) and the `mpi_f08` module with the same name, but with different interfaces. Each function can coexist twice with the same name in the same MPI library, one routine as an implicit interface outside of the `mpi` module, i.e., declared as `EXTERNAL`, and the other routine within `mpi_f08` declared with `CONTAINS`. These routines have different link names, which are also different to the link names used for the routines used in C and C++. (*End of advice to implementors.*)

*Advice to users.* Callbacks, including the predefined Fortran functions `MPI_COMM_NULL_COPY_FN`, `MPI_COMM_DUP_FN`, and `MPI_COMM_NULL_DELETE_FN` should not be passed from one application routine that uses the `mpi_f08` module to another application routine that uses the `mpi` module or `mpif.h`, and vice versa, see also the advice to users on page 701. (*End of advice to users.*)

MPI\_COMM\_FREE\_KEYVAL(comm\_keyval)

INOUT	comm_keyval	key value (integer)
-------	-------------	---------------------

```
int MPI_Comm_free_keyval(int *comm_keyval)
```

```
MPI_Comm_free_keyval(comm_keyval, ierror) BIND(C)
```

```
INTEGER, INTENT(INOUT) :: comm_keyval
```

```
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

MPI\_COMM\_FREE\_KEYVAL(COMM\_KEYVAL, IERROR)

```
INTEGER COMM_KEYVAL, IERROR
```

```
{static void MPI::Comm::Free_keyval(int& comm_keyval) (binding deprecated, see  
Section 15.2) }
```

Frees an extant attribute key. This function sets the value of `keyval` to `MPI_KEYVAL_INVALID`. Note that it is not erroneous to free an attribute key that is in use, because the actual free does not transpire until after all references (in other communicators on the process) to the key have been freed. These references need to be explicitly freed by the program, either via calls to `MPI_COMM_DELETE_ATTR` that free one attribute instance, or by calls to `MPI_COMM_FREE` that free all attribute instances associated with the freed communicator.

This call is identical to the MPI-1 call `MPI_KEYVAL_FREE` but is needed to match the new communicator-specific creation function. The use of `MPI_KEYVAL_FREE` is deprecated.

MPI\_COMM\_SET\_ATTR(comm, comm\_keyval, attribute\_val)

INOUT	comm	communicator from which attribute will be attached (handle)
-------	------	----------------------------------------------------------------

IN	comm_keyval	key value (integer)
----	-------------	---------------------

IN	attribute_val	attribute value
----	---------------	-----------------

```
int MPI_Comm_set_attr(MPI_Comm comm, int comm_keyval, void *attribute_val)
```

```
MPI_Comm_set_attr(comm, comm_keyval, attribute_val, ierror) BIND(C)
```

```
TYPE(MPI_Comm), INTENT(IN) :: comm
```

```
INTEGER, INTENT(IN) :: comm_keyval
```

```
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: attribute_val
```

```
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

MPI\_COMM\_SET\_ATTR(COMM, COMM\_KEYVAL, ATTRIBUTE\_VAL, IERROR)

```

INTEGER COMM, COMM_KEYVAL, IERROR
INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL
{void MPI::Comm::Set_attr(int comm_keyval, const void* attribute_val)
    const(binding deprecated, see Section 15.2) }

```

This function stores the stipulated attribute value `attribute_val` for subsequent retrieval by `MPI_COMM_GET_ATTR`. If the value is already present, then the outcome is as if `MPI_COMM_DELETE_ATTR` was first called to delete the previous value (and the callback function `comm_delete_attr_fn` was executed), and a new value was next stored. The call is erroneous if there is no key with value `keyval`; in particular `MPI_KEYVAL_INVALID` is an erroneous key value. The call will fail if the `comm_delete_attr_fn` function returned an error code other than `MPI_SUCCESS`.

This function replaces `MPI_ATTR_PUT`, whose use is deprecated. The C binding is identical. The Fortran binding differs in that `attribute_val` is an address-sized integer.

```

MPI_COMM_GET_ATTR(comm, comm_keyval, attribute_val, flag)

```

IN	comm	communicator to which the attribute is attached (handle)
IN	comm_keyval	key value (integer)
OUT	attribute_val	attribute value, unless <code>flag = false</code>
OUT	flag	false if no attribute is associated with the key (logical)

```

int MPI_Comm_get_attr(MPI_Comm comm, int comm_keyval, void *attribute_val,
    int *flag)

```

```

MPI_Comm_get_attr(comm, comm_keyval, attribute_val, flag, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, INTENT(IN) :: comm_keyval
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val
    LOGICAL, INTENT(OUT) :: flag
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_COMM_GET_ATTR(COMM, COMM_KEYVAL, ATTRIBUTE_VAL, FLAG, IERROR)
    INTEGER COMM, COMM_KEYVAL, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL
    LOGICAL FLAG

```

```

{bool MPI::Comm::Get_attr(int comm_keyval, void* attribute_val)
    const(binding deprecated, see Section 15.2) }

```

Retrieves attribute value by key. The call is erroneous if there is no key with value `keyval`. On the other hand, the call is correct if the key value exists, but no attribute is attached on `comm` for that key; in such case, the call returns `flag = false`. In particular `MPI_KEYVAL_INVALID` is an erroneous key value.

*Advice to users.* The call to `MPI_Comm_set_attr` passes in `attribute_val` the *value* of the attribute; the call to `MPI_Comm_get_attr` passes in `attribute_val` the *address* of the location where the attribute value is to be returned. Thus, if the attribute value itself is



```
MPI_WIN_CREATE_KEYVAL(win_copy_attr_fn, win_delete_attr_fn, win_keyval, extra_state)
```

IN	win_copy_attr_fn	copy callback function for win_keyval (function)
IN	win_delete_attr_fn	delete callback function for win_keyval (function)
OUT	win_keyval	key value for future access (integer)
IN	extra_state	extra state for callback functions

```
int MPI_Win_create_keyval(MPI_Win_copy_attr_function *win_copy_attr_fn,
    MPI_Win_delete_attr_function *win_delete_attr_fn,
    int *win_keyval, void *extra_state)
```

```
MPI_Win_create_keyval(win_copy_attr_fn, win_delete_attr_fn, win_keyval,
    extra_state, ierror) BIND(C)
    PROCEDURE(MPI_Win_copy_attr_function) :: win_copy_attr_fn
    PROCEDURE(MPI_Win_delete_attr_function) :: win_delete_attr_fn
    INTEGER, INTENT(OUT) :: win_keyval
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_WIN_CREATE_KEYVAL(WIN_COPY_ATTR_FN, WIN_DELETE_ATTR_FN, WIN_KEYVAL,
    EXTRA_STATE, IERROR)
    EXTERNAL WIN_COPY_ATTR_FN, WIN_DELETE_ATTR_FN
    INTEGER WIN_KEYVAL, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE
```

```
{static int MPI::Win::Create_keyval(MPI::Win::Copy_attr_function*
    win_copy_attr_fn,
    MPI::Win::Delete_attr_function* win_delete_attr_fn,
    void* extra_state) (binding deprecated, see Section 15.2) }
```

The argument `win_copy_attr_fn` may be specified as `MPI_WIN_NULL_COPY_FN` or `MPI_WIN_DUP_FN` from either C, C++, or Fortran. `MPI_WIN_NULL_COPY_FN` is a function that does nothing other than returning `flag = 0` and `MPI_SUCCESS`. `MPI_WIN_DUP_FN` is a simple-minded copy function that sets `flag = 1`, returns the value of `attribute_val_in` in `attribute_val_out`, and returns `MPI_SUCCESS`.

The argument `win_delete_attr_fn` may be specified as `MPI_WIN_NULL_DELETE_FN` from either C, C++, or Fortran. `MPI_WIN_NULL_DELETE_FN` is a function that does nothing, other than returning `MPI_SUCCESS`.

The C callback functions are:

```
typedef int MPI_Win_copy_attr_function(MPI_Win oldwin, int win_keyval,
    void *extra_state, void *attribute_val_in,
    void *attribute_val_out, int *flag);
```

and

```
typedef int MPI_Win_delete_attr_function(MPI_Win win, int win_keyval,
    void *attribute_val, void *extra_state);
```

With the `mpi_f08` module, the Fortran callback functions are:

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```

1  ABSTRACT INTERFACE
2
3      SUBROUTINE MPI_Win_copy_attr_function(oldwin, win_keyval, extra_state,
4      attribute_val_in, attribute_val_out, flag, ierror) BIND(C)
5          TYPE(MPI_Win) :: oldwin
6          INTEGER :: win_keyval, ierror
7          INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state, attribute_val_in,
8      attribute_val_out
9          LOGICAL :: flag

```

and

```

11 ABSTRACT INTERFACE
12
13     SUBROUTINE MPI_Win_delete_attr_function(win, win_keyval, attribute_val,
14     extra_state, ierror) BIND(C)
15         TYPE(MPI_Win) :: win
16         INTEGER :: win_keyval, ierror
17         INTEGER(KIND=MPI_ADDRESS_KIND) :: attribute_val, extra_state

```

[The] With the `mpi` module and `mpif.h`, the Fortran callback functions are:

```

20 SUBROUTINE WIN_COPY_ATTR_[FN]FUNCTION(OLDWIN, WIN_KEYVAL, EXTRA_STATE,
21     ATTRIBUTE_VAL_IN, ATTRIBUTE_VAL_OUT, FLAG, IERROR)
22     INTEGER OLDWIN, WIN_KEYVAL, IERROR
23     INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE, ATTRIBUTE_VAL_IN,
24     ATTRIBUTE_VAL_OUT
25     LOGICAL FLAG

```

and

```

27 SUBROUTINE WIN_DELETE_ATTR_[FN]FUNCTION(WIN, WIN_KEYVAL, ATTRIBUTE_VAL,
28     EXTRA_STATE, IERROR)
29     INTEGER WIN, WIN_KEYVAL, IERROR
30     INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL, EXTRA_STATE

```

The C++ callbacks are:

```

34 {typedef int MPI::Win::Copy_attr_function(const MPI::Win& oldwin,
35     int win_keyval, void* extra_state, void* attribute_val_in,
36     void* attribute_val_out, bool& flag); (binding deprecated, see
37     Section 15.2)}
```

and

```

39 {typedef int MPI::Win::Delete_attr_function(MPI::Win& win, int win_keyval,
40     void* attribute_val, void* extra_state); (binding deprecated, see
41     Section 15.2)}
```

If an attribute copy function or attribute delete function returns other than `MPI_SUCCESS`, then the call that caused it to be invoked (for example, `MPI_WIN_FREE`), is erroneous.

MPI\_WIN\_FREE\_KEYVAL(win\_keyval)

INOUT win\_keyval key value (integer)

int MPI\_Win\_free\_keyval(int \*win\_keyval)

MPI\_Win\_free\_keyval(win\_keyval, ierror) BIND(C)

INTEGER, INTENT(INOUT) :: win\_keyval

INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI\_WIN\_FREE\_KEYVAL(WIN\_KEYVAL, IERROR)

INTEGER WIN\_KEYVAL, IERROR

{static void MPI::Win::Free\_keyval(int& win\_keyval) (*binding deprecated, see Section 15.2*) }

MPI\_WIN\_SET\_ATTR(win, win\_keyval, attribute\_val)

INOUT win window to which attribute will be attached (handle)

IN win\_keyval key value (integer)

IN attribute\_val attribute value

int MPI\_Win\_set\_attr(MPI\_Win win, int win\_keyval, void \*attribute\_val)

MPI\_Win\_set\_attr(win, win\_keyval, attribute\_val, ierror) BIND(C)

TYPE(MPI\_Win), INTENT(IN) :: win

INTEGER, INTENT(IN) :: win\_keyval

INTEGER(KIND=MPI\_ADDRESS\_KIND), INTENT(IN) :: attribute\_val

INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI\_WIN\_SET\_ATTR(WIN, WIN\_KEYVAL, ATTRIBUTE\_VAL, IERROR)

INTEGER WIN, WIN\_KEYVAL, IERROR

INTEGER(KIND=MPI\_ADDRESS\_KIND) ATTRIBUTE\_VAL

{void MPI::Win::Set\_attr(int win\_keyval, const void\* attribute\_val) (*binding deprecated, see Section 15.2*) }

MPI\_WIN\_GET\_ATTR(win, win\_keyval, attribute\_val, flag)

IN win window to which the attribute is attached (handle)

IN win\_keyval key value (integer)

OUT attribute\_val attribute value, unless flag = false

OUT flag false if no attribute is associated with the key (logical)

int MPI\_Win\_get\_attr(MPI\_Win win, int win\_keyval, void \*attribute\_val, int \*flag)

MPI\_Win\_get\_attr(win, win\_keyval, attribute\_val, flag, ierror) BIND(C)

```

1      TYPE(MPI_Win), INTENT(IN) :: win
2      INTEGER, INTENT(IN) :: win_keyval
3      INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val
4      LOGICAL, INTENT(OUT) :: flag
5      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
6
7      MPI_WIN_GET_ATTR(WIN, WIN_KEYVAL, ATTRIBUTE_VAL, FLAG, IERROR)
8          INTEGER WIN, WIN_KEYVAL, IERROR
9          INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL
10         LOGICAL FLAG
11
12     {bool MPI::Win::Get_attr(int win_keyval, void* attribute_val) const(binding
13         deprecated, see Section 15.2) }
14
15     MPI_WIN_DELETE_ATTR(win, win_keyval)
16
17         INOUT    win                window from which the attribute is deleted (handle)
18         IN       win_keyval         key value (integer)
19
20     int MPI_Win_delete_attr(MPI_Win win, int win_keyval)
21
22     MPI_Win_delete_attr(win, win_keyval, ierror) BIND(C)
23         TYPE(MPI_Win), INTENT(IN) :: win
24         INTEGER, INTENT(IN) :: win_keyval
25         INTEGER, OPTIONAL, INTENT(OUT) :: ierror
26
27     MPI_WIN_DELETE_ATTR(WIN, WIN_KEYVAL, IERROR)
28         INTEGER WIN, WIN_KEYVAL, IERROR
29
30     {void MPI::Win::Delete_attr(int win_keyval)(binding deprecated, see Section 15.2)
31         }
32
33
34
35
36
37
38
39
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41
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45
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47
48

```

#### 6.7.4 Datatypes

The new functions for caching on datatypes are:

```

37     MPI_TYPE_CREATE_KEYVAL(type_copy_attr_fn, type_delete_attr_fn, type_keyval, extra_state)
38
39         IN       type_copy_attr_fn    copy callback function for type_keyval (function)
40         IN       type_delete_attr_fn  delete callback function for type_keyval (function)
41         OUT      type_keyval          key value for future access (integer)
42         IN       extra_state          extra state for callback functions
43
44
45     int MPI_Type_create_keyval(MPI_Type_copy_attr_function *type_copy_attr_fn,
46         MPI_Type_delete_attr_function *type_delete_attr_fn,
47         int *type_keyval, void *extra_state)
48

```



```

MPI_Type_create_keyval(type_copy_attr_fn, type_delete_attr_fn, type_keyval,
    extra_state, ierror) BIND(C)
    PROCEDURE(MPI_Type_copy_attr_function) :: type_copy_attr_fn
    PROCEDURE(MPI_Type_delete_attr_function) :: type_delete_attr_fn
    INTEGER, INTENT(OUT) :: type_keyval
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_TYPE_CREATE_KEYVAL(TYPE_COPY_ATTR_FN, TYPE_DELETE_ATTR_FN, TYPE_KEYVAL,
    EXTRA_STATE, IERROR)
    EXTERNAL TYPE_COPY_ATTR_FN, TYPE_DELETE_ATTR_FN
    INTEGER TYPE_KEYVAL, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE

{static int MPI::Datatype::Create_keyval(MPI::Datatype::Copy_attr_function*
    type_copy_attr_fn, MPI::Datatype::Delete_attr_function*
    type_delete_attr_fn, void* extra_state) (binding deprecated, see
    Section 15.2) }

```

The argument `type_copy_attr_fn` may be specified as `MPI_TYPE_NULL_COPY_FN` or `MPI_TYPE_DUP_FN` from either C, C++, or Fortran. `MPI_TYPE_NULL_COPY_FN` is a function that does nothing other than returning `flag = 0` and `MPI_SUCCESS`. `MPI_TYPE_DUP_FN` is a simple-minded copy function that sets `flag = 1`, returns the value of `attribute_val_in` in `attribute_val_out`, and returns `MPI_SUCCESS`.

The argument `type_delete_attr_fn` may be specified as `MPI_TYPE_NULL_DELETE_FN` from either C, C++, or Fortran. `MPI_TYPE_NULL_DELETE_FN` is a function that does nothing, other than returning `MPI_SUCCESS`.

The C callback functions are:

```

typedef int MPI_Type_copy_attr_function(MPI_Datatype oldtype,
    int type_keyval, void *extra_state, void *attribute_val_in,
    void *attribute_val_out, int *flag);

```

and

```

typedef int MPI_Type_delete_attr_function(MPI_Datatype datatype,
    int type_keyval, void *attribute_val, void *extra_state);

```

With the `mpi_f08` module, the Fortran callback functions are:

```

ABSTRACT INTERFACE
    SUBROUTINE MPI_Type_copy_attr_function(oldtype, type_keyval, extra_state,
        attribute_val_in, attribute_val_out, flag, ierror) BIND(C)
        TYPE(MPI_Datatype) :: oldtype
        INTEGER :: type_keyval, ierror
        INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state, attribute_val_in,
            attribute_val_out
        LOGICAL :: flag

```

and

```

ABSTRACT INTERFACE

```

```

1      SUBROUTINE MPI_Type_delete_attr_function(datatype, type_keyval,
2      attribute_val, extra_state, ierror) BIND(C)
3          TYPE(MPI_Datatype) :: datatype
4          INTEGER :: type_keyval, ierror
5          INTEGER(KIND=MPI_ADDRESS_KIND) :: attribute_val, extra_state
6

```

ticket230-B. [The]With the `mpi` module and `mpif.h`, the Fortran callback functions are:

```

ticket250-V. 8      SUBROUTINE TYPE_COPY_ATTR_[FN]FUNCTION(OLDTYPE, TYPE_KEYVAL, EXTRA_STATE,
9      ATTRIBUTE_VAL_IN, ATTRIBUTE_VAL_OUT, FLAG, IERROR)
10
11      INTEGER OLDTYPE, TYPE_KEYVAL, IERROR
12      INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE,
13      ATTRIBUTE_VAL_IN, ATTRIBUTE_VAL_OUT
14      LOGICAL FLAG
15

```

and

```

ticket250-V. 16     SUBROUTINE TYPE_DELETE_ATTR_[FN]FUNCTION(DATATYPE, TYPE_KEYVAL,
ticket252-W. 17     ATTRIBUTE_VAL, EXTRA_STATE, IERROR)
18
19     INTEGER DATATYPE, TYPE_KEYVAL, IERROR
20     INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL, EXTRA_STATE
21

```

The C++ callbacks are:

```

22     {typedef int
23
24         MPI::Datatype::Copy_attr_function(const MPI::Datatype& oldtype,
25         int type_keyval, void* extra_state,
26         const void* attribute_val_in, void* attribute_val_out,
27         bool& flag); (binding deprecated, see Section 15.2)}
28

```

and

```

ticket252-W. 29     {typedef int MPI::Datatype::Delete_attr_function(MPI::Datatype& datatype,
30     int type_keyval, void* attribute_val, void* extra_state);
31     (binding deprecated, see Section 15.2)}
32

```

If an attribute copy function or attribute delete function returns other than `MPI_SUCCESS`, then the call that caused it to be invoked (for example, `MPI_TYPE_FREE`), is erroneous.

`MPI_TYPE_FREE_KEYVAL`(type\_keyval)

INOUT    type\_keyval                      key value (integer)

```

40
41     int MPI_Type_free_keyval(int *type_keyval)
42

```

ticket-248T.

```

43     MPI_Type_free_keyval(type_keyval, ierror) BIND(C)
44     INTEGER, INTENT(INOUT) :: type_keyval
45     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
46
47     MPI_TYPE_FREE_KEYVAL(TYPE_KEYVAL, IERROR)
48     INTEGER TYPE_KEYVAL, IERROR
49

```

```
{static void MPI::Datatype::Free_keyval(int& type_keyval) (binding deprecated,  
                                         see Section 15.2) }
```

```
MPI_TYPE_SET_ATTR(datatype, type_keyval, attribute_val)
```

INOUT	[ticket252-W.] <b>datatype</b>	datatype to which attribute will be attached (handle)
IN	type_keyval	key value (integer)
IN	attribute_val	attribute value

```
int MPI_Type_set_attr(MPI_Datatype datatype, int type_keyval,  
                     void *attribute_val)
```

```
MPI_Type_set_attr(datatype, type_keyval, attribute_val, ierror) BIND(C)  
  TYPE(MPI_Datatype), INTENT(IN) :: datatype  
  INTEGER, INTENT(IN) :: type_keyval  
  INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: attribute_val  
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TYPE_SET_ATTR(DATATYPE, TYPE_KEYVAL, ATTRIBUTE_VAL, IERROR)  
  INTEGER DATATYPE, TYPE_KEYVAL, IERROR  
  INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL
```

```
{void MPI::Datatype::Set_attr(int type_keyval, const void*  
                             attribute_val) (binding deprecated, see Section 15.2) }
```

```
MPI_TYPE_GET_ATTR(datatype, type_keyval, attribute_val, flag)
```

IN	[ticket252-W.] <b>datatype</b>	datatype to which the attribute is attached (handle)
IN	type_keyval	key value (integer)
OUT	attribute_val	attribute value, unless flag = false
OUT	flag	false if no attribute is associated with the key (logical)

```
int MPI_Type_get_attr(MPI_Datatype datatype, int type_keyval, void  
                     *attribute_val, int *flag)
```

```
MPI_Type_get_attr(datatype, type_keyval, attribute_val, flag, ierror)  
  BIND(C)  
  TYPE(MPI_Datatype), INTENT(IN) :: datatype  
  INTEGER, INTENT(IN) :: type_keyval  
  INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val  
  LOGICAL, INTENT(OUT) :: flag  
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TYPE_GET_ATTR(DATATYPE, TYPE_KEYVAL, ATTRIBUTE_VAL, FLAG, IERROR)  
  INTEGER DATATYPE, TYPE_KEYVAL, IERROR  
  INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL  
  LOGICAL FLAG
```

```

1  {bool MPI::Datatype::Get_attr(int type_keyval, void* attribute_val)
2      const(binding deprecated, see Section 15.2) }
3
4
5  MPI_TYPE_DELETE_ATTR(datatype, type_keyval)
6
7      INOUT    [ticket252-W.]datatype      datatype from which the attribute is deleted (handle)
8
9      IN       type_keyval                key value (integer)
10
11  int MPI_Type_delete_attr(MPI_Datatype datatype, int type_keyval)
12
13  MPI_Type_delete_attr(datatype, type_keyval, ierror) BIND(C)
14      TYPE(MPI_Datatype), INTENT(IN) :: datatype
15      INTEGER, INTENT(IN) :: type_keyval
16      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
17
18  MPI_TYPE_DELETE_ATTR(DATATYPE, TYPE_KEYVAL, IERROR)
19
20  {void MPI::Datatype::Delete_attr(int type_keyval)(binding deprecated, see
21      Section 15.2) }
22
23
24
25
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27
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```

### 6.7.5 Error Class for Invalid Keyval

Key values for attributes are system-allocated, by MPI\_{TYPE,COMM,WIN}\_CREATE\_KEYVAL. Only such values can be passed to the functions that use key values as input arguments. In order to signal that an erroneous key value has been passed to one of these functions, there is a new MPI error class: MPI\_ERR\_KEYVAL. It can be returned by MPI\_ATTR\_PUT, MPI\_ATTR\_GET, MPI\_ATTR\_DELETE, MPI\_KEYVAL\_FREE, MPI\_{TYPE,COMM,WIN}\_DELETE\_ATTR, MPI\_{TYPE,COMM,WIN}\_SET\_ATTR, MPI\_{TYPE,COMM,WIN}\_GET\_ATTR, MPI\_{TYPE,COMM,WIN}\_FREE\_KEYVAL, MPI\_COMM\_DUP, MPI\_COMM\_DISCONNECT, and MPI\_COMM\_FREE. The last three are included because keyval is an argument to the copy and delete functions for attributes.

### 6.7.6 Attributes Example

*Advice to users.* This example shows how to write a collective communication operation that uses caching to be more efficient after the first call. The coding style assumes that MPI function results return only error statuses. (*End of advice to users.*)

```

/* key for this module's stuff: */
static int gop_key = MPI_KEYVAL_INVALID;

typedef struct
{
    int ref_count;          /* reference count */
    /* other stuff, whatever else we want */
} gop_stuff_type;

```

```

void Efficient_Collective_Op (MPI_Comm comm, ...)
{
    gop_stuff_type *gop_stuff;
    MPI_Group      group;
    int            foundflag;

    MPI_Comm_group(comm, &group);

    if (gop_key == MPI_KEYVAL_INVALID) /* get a key on first call ever */
    {
        if ( ! MPI_Comm_create_keyval( gop_stuff_copier,
                                       gop_stuff_destructor,
                                       &gop_key, (void *)0));
        /* get the key while assigning its copy and delete callback
           behavior. */

        MPI_Abort (comm, 99);
    }

    MPI_Comm_get_attr (comm, gop_key, &gop_stuff, &foundflag);
    if (foundflag)
    { /* This module has executed in this group before.
       We will use the cached information */
    }
    else
    { /* This is a group that we have not yet cached anything in.
       We will now do so.
       */

        /* First, allocate storage for the stuff we want,
           and initialize the reference count */

        gop_stuff = (gop_stuff_type *) malloc (sizeof(gop_stuff_type));
        if (gop_stuff == NULL) { /* abort on out-of-memory error */ }

        gop_stuff -> ref_count = 1;

        /* Second, fill in *gop_stuff with whatever we want.
           This part isn't shown here */

        /* Third, store gop_stuff as the attribute value */
        MPI_Comm_set_attr ( comm, gop_key, gop_stuff);
    }
    /* Then, in any case, use contents of *gop_stuff
       to do the global op ... */
}

/* The following routine is called by MPI when a group is freed */

```

```

1
2  int gop_stuff_destructor (MPI_Comm comm, int keyval, void *gop_stuffP,
3                          void *extra)
4  {
5      gop_stuff_type *gop_stuff = (gop_stuff_type *)gop_stuffP;
6      if (keyval != gop_key) { /* abort -- programming error */ }
7
8      /* The group's being freed removes one reference to gop_stuff */
9      gop_stuff -> ref_count -= 1;
10
11     /* If no references remain, then free the storage */
12     if (gop_stuff -> ref_count == 0) {
13         free((void *)gop_stuff);
14     }
15     return MPI_SUCCESS;
16 }
17
18 /* The following routine is called by MPI when a group is copied */
19 int gop_stuff_copier (MPI_Comm comm, int keyval, void *extra,
20                     void *gop_stuff_inP, void *gop_stuff_outP, int *flag)
21 {
22     gop_stuff_type *gop_stuff_in = (gop_stuff_type *)gop_stuff_inP;
23     gop_stuff_type **gop_stuff_out = (gop_stuff_type **)gop_stuff_outP;
24     if (keyval != gop_key) { /* abort -- programming error */ }
25
26     /* The new group adds one reference to this gop_stuff */
27     gop_stuff_in -> ref_count += 1;
28     *gop_stuff_out = gop_stuff_in;
29     return MPI_SUCCESS;
30 }
31
32
33
34
35
36
37
38
39
40
41
42
43
44
45
46

```

## 6.8 Naming Objects

There are many occasions on which it would be useful to allow a user to associate a printable identifier with an MPI communicator, window, or datatype, for instance error reporting, debugging, and profiling. The names attached to opaque objects do not propagate when the object is duplicated or copied by MPI routines. For communicators this can be achieved using the following two functions.

**MPI\_COMM\_SET\_NAME** (comm, comm\_name)

INOUT	comm	communicator whose identifier is to be set (handle)
IN	comm_name	the character string which is remembered as the name (string)

```
int MPI_Comm_set_name(MPI_Comm comm, const char *comm_name)
```

```

MPI_Comm_set_name(comm, comm_name, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    CHARACTER(LEN=*), INTENT(IN) :: comm_name
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_COMM_SET_NAME(COMM, COMM_NAME, IERROR)
    INTEGER COMM, IERROR
    CHARACTER*(*) COMM_NAME

```

```

{void MPI::Comm::Set_name(const char* comm_name) (binding deprecated, see
    Section 15.2) }

```

MPI\_COMM\_SET\_NAME allows a user to associate a name string with a communicator. The character string which is passed to MPI\_COMM\_SET\_NAME will be saved inside the MPI library (so it can be freed by the caller immediately after the call, or allocated on the stack). Leading spaces in name are significant but trailing ones are not.

MPI\_COMM\_SET\_NAME is a local (non-collective) operation, which only affects the name of the communicator as seen in the process which made the MPI\_COMM\_SET\_NAME call. There is no requirement that the same (or any) name be assigned to a communicator in every process where it exists.

*Advice to users.* Since MPI\_COMM\_SET\_NAME is provided to help debug code, it is sensible to give the same name to a communicator in all of the processes where it exists, to avoid confusion. (*End of advice to users.*)

The length of the name which can be stored is limited to the value of MPI\_MAX\_OBJECT\_NAME in Fortran and MPI\_MAX\_OBJECT\_NAME-1 in C and C++ to allow for the null terminator. Attempts to put names longer than this will result in truncation of the name. MPI\_MAX\_OBJECT\_NAME must have a value of at least 64.

*Advice to users.* Under circumstances of store exhaustion an attempt to put a name of any length could fail, therefore the value of MPI\_MAX\_OBJECT\_NAME should be viewed only as a strict upper bound on the name length, not a guarantee that setting names of less than this length will always succeed. (*End of advice to users.*)

*Advice to implementors.* Implementations which pre-allocate a fixed size space for a name should use the length of that allocation as the value of MPI\_MAX\_OBJECT\_NAME. Implementations which allocate space for the name from the heap should still define MPI\_MAX\_OBJECT\_NAME to be a relatively small value, since the user has to allocate space for a string of up to this size when calling MPI\_COMM\_GET\_NAME. (*End of advice to implementors.*)

```

MPI_COMM_GET_NAME (comm, comm_name, resultlen)

```

IN	comm	communicator whose name is to be returned (handle)
OUT	comm_name	the name previously stored on the communicator, or an empty string if no such name exists (string)
OUT	resultlen	length of returned name (integer)

```

int MPI_Comm_get_name(MPI_Comm comm, char *comm_name, int *resultlen)

```

```

1
ticket-248T. 2 MPI_Comm_get_name(comm, comm_name, resultlen, ierror) BIND(C)
3     TYPE(MPI_Comm), INTENT(IN) :: comm
4     CHARACTER(LEN=MPI_MAX_OBJECT_NAME), INTENT(OUT) :: comm_name
5     INTEGER, INTENT(OUT) :: resultlen
6     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
7
8 MPI_COMM_GET_NAME(COMM, COMM_NAME, RESULTLEN, IERROR)
9     INTEGER COMM, RESULTLEN, IERROR
10    CHARACTER*(*) COMM_NAME
11
12 {void MPI::Comm::Get_name(char* comm_name, int& resultlen) const(binding
13     deprecated, see Section 15.2) }

```

MPI\_COMM\_GET\_NAME returns the last name which has previously been associated with the given communicator. The name may be set and got from any language. The same name will be returned independent of the language used. `name` should be allocated so that it can hold a resulting string of length MPI\_MAX\_OBJECT\_NAME characters.

MPI\_COMM\_GET\_NAME returns a copy of the set name in `name`.

In C, a null character is additionally stored at `name[resultlen]`. The value of `resultlen` cannot be larger [then]than MPI\_MAX\_OBJECT\_NAME-1. In Fortran, `name` is padded on the right with blank characters. The value of `resultlen` cannot be larger [then]than MPI\_MAX\_OBJECT\_NAME.

If the user has not associated a name with a communicator, or an error occurs, MPI\_COMM\_GET\_NAME will return an empty string (all spaces in Fortran, "" in C and C++). The three predefined communicators will have predefined names associated with them. Thus, the names of MPI\_COMM\_WORLD, MPI\_COMM\_SELF, and the communicator returned by MPI\_COMM\_GET\_PARENT (if not MPI\_COMM\_NULL) will have the default of MPI\_COMM\_WORLD, MPI\_COMM\_SELF, and MPI\_COMM\_PARENT. The fact that the system may have chosen to give a default name to a communicator does not prevent the user from setting a name on the same communicator; doing this removes the old name and assigns the new one.

*Rationale.* We provide separate functions for setting and getting the name of a communicator, rather than simply providing a predefined attribute key for the following reasons:

- It is not, in general, possible to store a string as an attribute from Fortran.
- It is not easy to set up the delete function for a string attribute unless it is known to have been allocated from the heap.
- To make the attribute key useful additional code to call `strdup` is necessary. If this is not standardized then users have to write it. This is extra unneeded work which we can easily eliminate.
- The Fortran binding is not trivial to write (it will depend on details of the Fortran compilation system), and will not be portable. Therefore it should be in the library rather than in user code.

(*End of rationale.*)



*Advice to users.* The above definition means that it is safe simply to print the string returned by `MPI_COMM_GET_NAME`, as it is always a valid string even if there was no name.

Note that associating a name with a communicator has no effect on the semantics of an MPI program, and will (necessarily) increase the store requirement of the program, since the names must be saved. Therefore there is no requirement that users use these functions to associate names with communicators. However debugging and profiling MPI applications may be made easier if names are associated with communicators, since the debugger or profiler should then be able to present information in a less cryptic manner. (*End of advice to users.*)

The following functions are used for setting and getting names of datatypes. The constant `MPI_MAX_OBJECT_NAME` also applies to these names.

`MPI_TYPE_SET_NAME` (`datatype`, `type_name`)

INOUT	[ticket252-W.] <code>datatype</code>	datatype whose identifier is to be set (handle)
IN	<code>type_name</code>	the character string which is remembered as the name (string)

`int MPI_Type_set_name(MPI_Datatype datatype, const char *type_name)`

`MPI_Type_set_name(datatype, type_name, ierror) BIND(C)`  
`TYPE(MPI_Datatype), INTENT(IN) :: datatype`  
`CHARACTER(LEN=*) , INTENT(IN) :: type_name`  
`INTEGER, OPTIONAL, INTENT(OUT) :: ierror`

`MPI_TYPE_SET_NAME(DATATYPE, TYPE_NAME, IERROR)`

`INTEGER DATATYPE, IERROR`  
`CHARACTER*(*) TYPE_NAME`

`{void MPI::Datatype::Set_name(const char* type_name) (binding deprecated, see Section 15.2) }`

`MPI_TYPE_GET_NAME` (`datatype`, `type_name`, `resultlen`)

IN	[ticket252-W.] <code>datatype</code>	datatype whose name is to be returned (handle)
OUT	<code>type_name</code>	the name previously stored on the datatype, or a empty string if no such name exists (string)
OUT	<code>resultlen</code>	length of returned name (integer)

`int MPI_Type_get_name(MPI_Datatype datatype, char *type_name, int *resultlen)`

`MPI_Type_get_name(datatype, type_name, resultlen, ierror) BIND(C)`  
`TYPE(MPI_Datatype), INTENT(IN) :: datatype`  
`CHARACTER(LEN=MPI_MAX_OBJECT_NAME), INTENT(OUT) :: type_name`

```

1      INTEGER, INTENT(OUT) :: resultlen
2      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
3
4  MPI_TYPE_GET_NAME(DATATYPE, TYPE_NAME, RESULTLEN, IERROR)
5      INTEGER DATATYPE, RESULTLEN, IERROR
6      CHARACTER*(*) TYPE_NAME
7
8  {void MPI::Datatype::Get_name(char* type_name, int& resultlen) const(binding
9      deprecated, see Section 15.2) }

```

Named predefined datatypes have the default names of the datatype name. For example, MPI\_WCHAR has the default name of MPI\_WCHAR.

The following functions are used for setting and getting names of windows. The constant MPI\_MAX\_OBJECT\_NAME also applies to these names.

```

15  MPI_WIN_SET_NAME (win, win_name)
16
17      INOUT    win                window whose identifier is to be set (handle)
18      IN       win_name          the character string which is remembered as the name
19                                  (string)

```

```

21  int MPI_Win_set_name(MPI_Win win, const char *win_name)
22
23  MPI_Win_set_name(win, win_name, ierror) BIND(C)
24      TYPE(MPI_Win), INTENT(IN) :: win
25      CHARACTER(LEN=*), INTENT(IN) :: win_name
26      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
27
28  MPI_WIN_SET_NAME(WIN, WIN_NAME, IERROR)
29      INTEGER WIN, IERROR
30      CHARACTER*(*) WIN_NAME
31
32  {void MPI::Win::Set_name(const char* win_name)(binding deprecated, see
33      Section 15.2) }

```

```

35  MPI_WIN_GET_NAME (win, win_name, resultlen)
36
37      IN       win                window whose name is to be returned (handle)
38      OUT      win_name          the name previously stored on the window, or a empty
39                                  string if no such name exists (string)
40      OUT      resultlen         length of returned name (integer)

```

```

42  int MPI_Win_get_name(MPI_Win win, char *win_name, int *resultlen)
43
44  MPI_Win_get_name(win, win_name, resultlen, ierror) BIND(C)
45      TYPE(MPI_Win), INTENT(IN) :: win
46      CHARACTER(LEN=MPI_MAX_OBJECT_NAME), INTENT(OUT) :: win_name
47      INTEGER, INTENT(OUT) :: resultlen
48      INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_WIN_GET_NAME(WIN, WIN_NAME, RESULTLEN, IERROR)
    INTEGER WIN, RESULTLEN, IERROR
    CHARACTER*(*) WIN_NAME

{void MPI::Win::Get_name(char* win_name, int& resultlen) const(binding
    deprecated, see Section 15.2) }
```

## 6.9 Formalizing the Loosely Synchronous Model

In this section, we make further statements about the loosely synchronous model, with particular attention to intra-communication.

### 6.9.1 Basic Statements

When a caller passes a communicator (that contains a context and group) to a callee, that communicator must be free of side effects throughout execution of the subprogram: there should be no active operations on that communicator that might involve the process. This provides one model in which libraries can be written, and work “safely.” For libraries so designated, the callee has permission to do whatever communication it likes with the communicator, and under the above guarantee knows that no other communications will interfere. Since we permit good implementations to create new communicators without synchronization (such as by preallocated contexts on communicators), this does not impose a significant overhead.

This form of safety is analogous to other common computer-science usages, such as passing a descriptor of an array to a library routine. The library routine has every right to expect such a descriptor to be valid and modifiable.

### 6.9.2 Models of Execution

In the loosely synchronous model, transfer of control to a **parallel procedure** is effected by having each executing process invoke the procedure. The invocation is a collective operation: it is executed by all processes in the execution group, and invocations are similarly ordered at all processes. However, the invocation need not be synchronized.

We say that a parallel procedure is *active* in a process if the process belongs to a group that may collectively execute the procedure, and some member of that group is currently executing the procedure code. If a parallel procedure is active in a process, then this process may be receiving messages pertaining to this procedure, even if it does not currently execute the code of this procedure.

#### Static communicator allocation

This covers the case where, at any point in time, at most one invocation of a parallel procedure can be active at any process, and the group of executing processes is fixed. For example, all invocations of parallel procedures involve all processes, processes are single-threaded, and there are no recursive invocations.

In such a case, a communicator can be statically allocated to each procedure. The static allocation can be done in a preamble, as part of initialization code. If the parallel procedures can be organized into libraries, so that only one procedure of each library can

be concurrently active in each processor, then it is sufficient to allocate one communicator per library.

#### Dynamic communicator allocation

Calls of parallel procedures are well-nested if a new parallel procedure is always invoked in a subset of a group executing the same parallel procedure. Thus, processes that execute the same parallel procedure have the same execution stack.

In such a case, a new communicator needs to be dynamically allocated for each new invocation of a parallel procedure. The allocation is done by the caller. A new communicator can be generated by a call to `MPI_COMM_DUP`, if the callee execution group is identical to the caller execution group, or by a call to `MPI_COMM_SPLIT` if the caller execution group is split into several subgroups executing distinct parallel routines. The new communicator is passed as an argument to the invoked routine.

The need for generating a new communicator at each invocation can be alleviated or avoided altogether in some cases: If the execution group is not split, then one can allocate a stack of communicators in a preamble, and next manage the stack in a way that mimics the stack of recursive calls.

One can also take advantage of the well-ordering property of communication to avoid confusing caller and callee communication, even if both use the same communicator. To do so, one needs to abide by the following two rules:

- messages sent before a procedure call (or before a return from the procedure) are also received before the matching call (or return) at the receiving end;
- messages are always selected by source (no use is made of `MPI_ANY_SOURCE`).

#### ticket0. The General [c]Case

In the general case, there may be multiple concurrently active invocations of the same parallel procedure within the same group; invocations may not be well-nested. A new communicator needs to be created for each invocation. It is the user's responsibility to make sure that, should two distinct parallel procedures be invoked concurrently on overlapping sets of processes, then communicator creation be properly coordinated.

# Chapter 7

## Process Topologies

### 7.1 Introduction

This chapter discusses the MPI topology mechanism. A topology is an extra, optional attribute that one can give to an intra-communicator; topologies cannot be added to inter-communicators. A topology can provide a convenient naming mechanism for the processes of a group (within a communicator), and additionally, may assist the runtime system in mapping the processes onto hardware.

As stated in Chapter 6, a process group in MPI is a collection of  $n$  processes. Each process in the group is assigned a rank between 0 and  $n-1$ . In many parallel applications a linear ranking of processes does not adequately reflect the logical communication pattern of the processes (which is usually determined by the underlying problem geometry and the numerical algorithm used). Often the processes are arranged in topological patterns such as two- or three-dimensional grids. More generally, the logical process arrangement is described by a graph. In this chapter we will refer to this logical process arrangement as the “virtual topology.”

A clear distinction must be made between the virtual process topology and the topology of the underlying, physical hardware. The virtual topology can be exploited by the system in the assignment of processes to physical processors, if this helps to improve the communication performance on a given machine. How this mapping is done, however, is outside the scope of MPI. The description of the virtual topology, on the other hand, depends only on the application, and is machine-independent. The functions that are described in this chapter deal [only ]with machine-independent mapping and communication on virtual process topologies.

*Rationale.* Though physical mapping is not discussed, the existence of the virtual topology information may be used as advice by the runtime system. There are well-known techniques for mapping grid/torus structures to hardware topologies such as hypercubes or grids. For more complicated graph structures good heuristics often yield nearly optimal results [44]. On the other hand, if there is no way for the user to specify the logical process arrangement as a “virtual topology,” a random mapping is most likely to result. On some machines, this will lead to unnecessary contention in the interconnection network. Some details about predicted and measured performance improvements that result from good process-to-processor mapping on modern wormhole-routing architectures can be found in [11, 12].

Besides possible performance benefits, the virtual topology can function as a convenient, process-naming structure, with significant benefits for program readability and notational power in message-passing programming. (*End of rationale.*)

## 7.2 Virtual Topologies

The communication pattern of a set of processes can be represented by a graph. The nodes represent processes, and the edges connect processes that communicate with each other. MPI provides message-passing between any pair of processes in a group. There is no requirement for opening a channel explicitly. Therefore, a “missing link” in the user-defined process graph does not prevent the corresponding processes from exchanging messages. It means rather that this connection is neglected in the virtual topology. This strategy implies that the topology gives no convenient way of naming this pathway of communication. Another possible consequence is that an automatic mapping tool (if one exists for the runtime environment) will not take account of this edge when mapping.

Specifying the virtual topology in terms of a graph is sufficient for all applications. However, in many applications the graph structure is regular, and the detailed set-up of the graph would be inconvenient for the user and might be less efficient at run time. A large fraction of all parallel applications use process topologies like rings, two- or higher-dimensional grids, or tori. These structures are completely defined by the number of dimensions and the numbers of processes in each coordinate direction. Also, the mapping of grids and tori is generally an easier problem [then]than that of general graphs. Thus, it is desirable to address these cases explicitly.

Process coordinates in a Cartesian structure begin their numbering at 0. Row-major numbering is always used for the processes in a Cartesian structure. This means that, for example, the relation between group rank and coordinates for four processes in a  $(2 \times 2)$  grid is as follows.

```
coord (0,0):  rank 0
coord (0,1):  rank 1
coord (1,0):  rank 2
coord (1,1):  rank 3
```

## 7.3 Embedding in MPI

The support for virtual topologies as defined in this chapter is consistent with other parts of MPI, and, whenever possible, makes use of functions that are defined elsewhere. Topology information is associated with communicators. It is added to communicators using the caching mechanism described in Chapter 6.

## 7.4 Overview of the Functions

[The functions `MPI_GRAPH_CREATE`, `MPI_DIST_GRAPH_CREATE_ADJACENT`, `MPI_DIST_GRAPH_CREATE` and `MPI_CART_CREATE` are used to create general (graph) virtual topologies and Cartesian topologies, respectively. These topology creation functions are collective. As with other collective calls, the program must be written to work correctly, whether the call synchronizes or not.]MPI supports three topology types: Cartesian,

graph, and distributed graph. The function `MPI_CART_CREATE` is used to create Cartesian topologies, the function `MPI_GRAPH_CREATE` is used to create graph topologies, and the functions `MPI_DIST_GRAPH_CREATE_ADJACENT` and `MPI_DIST_GRAPH_CREATE` are used to create distributed graph topologies.

The topology creation functions take as input an existing communicator `comm_old`, which defines the set of processes on which the topology is to be mapped. For `MPI_GRAPH_CREATE` and `MPI_CART_CREATE`, all input arguments must have identical values on all processes of the group of `comm_old`. [For `MPI_DIST_GRAPH_CREATE_ADJACENT` and `MPI_DIST_GRAPH_CREATE` the input communication graph is distributed across the calling processes.] When calling `MPI_GRAPH_CREATE`, each process specifies all nodes and edges in the graph. In contrast, the functions `MPI_DIST_GRAPH_CREATE_ADJACENT` or `MPI_DIST_GRAPH_CREATE` are used to specify the graph in a distributed fashion, whereby each process only specifies a subset of the edges in the graph such that the entire graph structure is defined collectively across the set of processes. Therefore the processes provide different values for the arguments specifying the graph. However, all processes must give the same value for `reorder` and the `info` argument. In all cases, a new communicator `comm_topol` is created that carries the topological structure as cached information (see Chapter 6). In analogy to function `MPI_COMM_CREATE`, no cached information propagates from `comm_old` to `comm_topol`.

`MPI_CART_CREATE` can be used to describe Cartesian structures of arbitrary dimension. For each coordinate direction one specifies whether the process structure is periodic or not. Note that an  $n$ -dimensional hypercube is an  $n$ -dimensional torus with 2 processes per coordinate direction. Thus, special support for hypercube structures is not necessary. The local auxiliary function `MPI_DIMS_CREATE` can be used to compute a balanced distribution of processes among a given number of dimensions.

*Rationale.* Similar functions are contained in EXPRESS [13] and PARMACS. (*End of rationale.*)

[The function `MPI_TOPO_TEST` can be used to inquire about the topology associated with a communicator. The topological information can be extracted from the communicator using the functions `MPI_GRAPHDIMS_GET` and `MPI_GRAPH_GET`, for general graphs, and `MPI_CARTDIM_GET` and `MPI_CART_GET`, for Cartesian topologies. Several additional functions are provided to manipulate Cartesian topologies: the functions `MPI_CART_RANK` and `MPI_CART_COORDS` translate Cartesian coordinates into a group rank, and vice-versa; the function `MPI_CART_SUB` can be used to extract a Cartesian subspace (analogous to `MPI_COMM_SPLIT`). The function `MPI_CART_SHIFT` provides the information needed to communicate with neighbors in a Cartesian dimension. The two functions `MPI_GRAPH_NEIGHBORS_COUNT` and `MPI_GRAPH_NEIGHBORS` can be used to extract the neighbors of a node in a graph. For distributed graphs, the functions `MPI_DIST_NEIGHBORS_COUNT` and `MPI_DIST_NEIGHBORS` can be used to extract the neighbors of the calling node. The function `MPI_CART_SUB` is collective over the input communicator's group; all other functions are local.] MPI defines functions to query a communicator for topology information. The function `MPI_TOPO_TEST` is used to query for the type of topology associated with a communicator. Depending on the topology type, different information can be extracted. For a graph topology, the functions `MPI_GRAPHDIMS_GET` and `MPI_GRAPH_GET` return the values that were specified in the call to `MPI_GRAPH_CREATE`. Additionally, the functions `MPI_GRAPH_NEIGHBORS_COUNT`



and `MPI_GRAPH_NEIGHBORS` can be used to obtain the neighbors of an arbitrary node in the graph. For a distributed graph topology, the functions `MPI_DIST_GRAPH_NEIGHBORS_COUNT` and `MPI_DIST_GRAPH_NEIGHBORS` can be used to obtain the neighbors of the calling process. For a Cartesian topology, the functions `MPI_CARTDIM_GET` and `MPI_CART_GET` return the values that were specified in the call to `MPI_CART_CREATE`. Additionally, the functions `MPI_CART_RANK` and `MPI_CART_COORDS` translate Cartesian coordinates into a group rank, and vice-versa. The function `MPI_CART_SHIFT` provides the information needed to communicate with neighbors along a Cartesian dimension. All of these query functions are local.

For Cartesian topologies, the function `MPI_CART_SUB` can be used to extract a Cartesian subspace (analogous to `MPI_COMM_SPLIT`). This function is collective over the input communicator's group.

Two additional functions, `MPI_GRAPH_MAP` and `MPI_CART_MAP` are presented in the last section. In general these functions are not called by the user directly. However, together with the communicator manipulation functions presented in Chapter 6, they are sufficient to implement all other topology functions. Section 7.5.8 outlines such an implementation.

The neighborhood collective communication routines `MPI_NEIGHBOR_ALLGATHER`, `MPI_NEIGHBOR_ALLGATHERV`, `MPI_NEIGHBOR_ALLTOALL`, `MPI_NEIGHBOR_ALLTOALLV`, and `MPI_NEIGHBOR_ALLTOALLW` communicate with the nearest neighbors on the topology associated with the communicator. The nonblocking variants are `MPI_INEIGHBOR_ALLGATHER`, `MPI_INEIGHBOR_ALLGATHERV`, `MPI_INEIGHBOR_ALLTOALL`, `MPI_INEIGHBOR_ALLTOALLV`, and `MPI_INEIGHBOR_ALLTOALLW`.

## 7.5 Topology Constructors

### 7.5.1 Cartesian Constructor

`MPI_CART_CREATE(comm_old, ndims, dims, periods, reorder, comm_cart)`

IN	<code>comm_old</code>	input communicator (handle)
IN	<code>ndims</code>	number of dimensions of Cartesian grid (integer)
IN	<code>dims</code>	integer array of size <code>ndims</code> specifying the number of processes in each dimension
IN	<code>periods</code>	logical array of size <code>ndims</code> specifying whether the grid is periodic ( <code>true</code> ) or not ( <code>false</code> ) in each dimension
IN	<code>reorder</code>	ranking may be reordered ( <code>true</code> ) or not ( <code>false</code> ) (logical)
OUT	<code>comm_cart</code>	communicator with new Cartesian topology (handle)

```
int MPI_Cart_create(MPI_Comm comm_old, int ndims, const int [*dims]dims[],
                  const int [*periods]periods[], int reorder,
                  MPI_Comm *comm_cart)
```

```
MPI_Cart_create(comm_old, ndims, dims, periods, reorder, comm_cart, ierror)
BIND(C)
```



```

TYPE(MPI_Comm), INTENT(IN) :: comm_old
INTEGER, INTENT(IN) :: ndims, dims(ndims)
LOGICAL, INTENT(IN) :: periods(ndims), reorder
TYPE(MPI_Comm), INTENT(OUT) :: comm_cart
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_CART_CREATE(COMM_OLD, NDIMS, DIMS, PERIODS, REORDER, COMM_CART, IERROR)
  INTEGER COMM_OLD, NDIMS, DIMS(*), COMM_CART, IERROR
  LOGICAL PERIODS(*), REORDER

{MPI::Cartcomm MPI::Intracomm::Create_cart(int ndims, const int dims[],
      const bool periods[], bool reorder) const (binding deprecated, see
      Section 15.2) }

```

MPI\_CART\_CREATE returns a handle to a new communicator to which the Cartesian topology information is attached. If `reorder = false` then the rank of each process in the new group is identical to its rank in the old group. Otherwise, the function may reorder the processes (possibly so as to choose a good embedding of the virtual topology onto the physical machine). If the total size of the Cartesian grid is smaller than the size of the group of `[comm]comm_old`, then some processes are returned MPI\_COMM\_NULL, in analogy to MPI\_COMM\_SPLIT. If `ndims` is zero then a zero-dimensional Cartesian topology is created. The call is erroneous if it specifies a grid that is larger than the group size or if `ndims` is negative.

### 7.5.2 Cartesian Convenience Function: MPI\_DIMS\_CREATE

For Cartesian topologies, the function MPI\_DIMS\_CREATE helps the user select a balanced distribution of processes per coordinate direction, depending on the number of processes in the group to be balanced and optional constraints that can be specified by the user. One use is to partition all the processes (the size of MPI\_COMM\_WORLD's group) into an  $n$ -dimensional topology.

```

MPI_DIMS_CREATE(nnodes, ndims, dims)

IN      nnodes      number of nodes in a grid (integer)
IN      ndims       number of Cartesian dimensions (integer)
INOUT   dims        integer array of size ndims specifying the number of
                    nodes in each dimension

int MPI_Dims_create(int nnodes, int ndims, int *dims)

MPI_Dims_create(nnodes, ndims, dims, ierror) BIND(C)
  INTEGER, INTENT(IN) :: nnodes, ndims
  INTEGER, INTENT(INOUT) :: dims(ndims)
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_DIMS_CREATE(NNODES, NDIMS, DIMS, IERROR)
  INTEGER NNODES, NDIMS, DIMS(*), IERROR

```

```
1 {void MPI::Compute_dims(int nnodes, int ndims, int dims[]) (binding deprecated,
2 see Section 15.2) }
```

3  
4 The entries in the array `dims` are set to describe a Cartesian grid with `ndims` dimensions and a total of `nnodes` nodes. The dimensions are set to be as close to each other as possible, using an appropriate divisibility algorithm. The caller may further constrain the operation of this routine by specifying elements of array `dims`. If `dims[i]` is set to a positive number, the routine will not modify the number of nodes in dimension `i`; only those entries where `dims[i] = 0` are modified by the call.

5  
6  
7  
8  
9 Negative input values of `dims[i]` are erroneous. An error will occur if `nnodes` is not a multiple of  $\prod_{i, \text{dims}[i] \neq 0} \text{dims}[i]$ .

10  
11  
12 For `dims[i]` set by the call, `dims[i]` will be ordered in non-increasing order. Array `dims` is suitable for use as input to routine `MPI_CART_CREATE`. `MPI_DIMS_CREATE` is local.

### 16 Example 7.1

17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48	dims before call	function call	dims on return
	(0,0)	MPI_DIMS_CREATE(6, 2, dims)	(3,2)
	(0,0)	MPI_DIMS_CREATE(7, 2, dims)	(7,1)
	(0,3,0)	MPI_DIMS_CREATE(6, 3, dims)	(2,3,1)
	(0,3,0)	MPI_DIMS_CREATE(7, 3, dims)	erroneous call

### 7.5.3 [General (Graph)]Graph Constructor

MPI\_GRAPH\_CREATE(comm\_old, nnodes, index, edges, reorder, comm\_graph)

IN	comm_old	input communicator (handle)
IN	nnodes	number of nodes in graph (integer)
IN	index	array of integers describing node degrees (see below)
IN	edges	array of integers describing graph edges (see below)
IN	reorder	ranking may be reordered (true) or not (false) (logical)
OUT	comm_graph	communicator with graph topology added (handle)

```
int MPI_Graph_create(MPI_Comm comm_old, int nnodes, const
    int [*index]index[], const int [*edges]edges[], int reorder,
    MPI_Comm *comm_graph)
```

```
MPI_Graph_create(comm_old, nnodes, index, edges, reorder, comm_graph,
    ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm_old
    INTEGER, INTENT(IN) :: nnodes, index(nnodes), edges(*)
    LOGICAL, INTENT(IN) :: reorder
    TYPE(MPI_Comm), INTENT(OUT) :: comm_graph
```

```

    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_GRAPH_CREATE(COMM_OLD, NNODES, INDEX, EDGES, REORDER, COMM_GRAPH,
    IERROR)
    INTEGER COMM_OLD, NNODES, INDEX(*), EDGES(*), COMM_GRAPH, IERROR
    LOGICAL REORDER
{MPI::Graphcomm MPI::Intracomm::Create_graph(int nnodes, const int index[],
    const int edges[], bool reorder) const(binding deprecated, see
    Section 15.2) }
```

MPI\_GRAPH\_CREATE returns a handle to a new communicator to which the graph topology information is attached. If `reorder = false` then the rank of each process in the new group is identical to its rank in the old group. Otherwise, the function may reorder the processes. If the size, `nnodes`, of the graph is smaller than the size of the group of `[comm]comm_old`, then some processes are returned MPI\_COMM\_NULL, in analogy to MPI\_CART\_CREATE and MPI\_COMM\_SPLIT. If the graph is empty, i.e., `nnodes == 0`, then MPI\_COMM\_NULL is returned in all processes. The call is erroneous if it specifies a graph that is larger than the group size of the input communicator.

The three parameters `nnodes`, `index` and `edges` define the graph structure. `nnodes` is the number of nodes of the graph. The nodes are numbered from 0 to `nnodes-1`. The *i*-th entry of array `index` stores the total number of neighbors of the first *i* graph nodes. The lists of neighbors of nodes 0, 1, ..., `nnodes-1` are stored in consecutive locations in array `edges`. The array `edges` is a flattened representation of the edge lists. The total number of entries in `index` is `nnodes` and the total number of entries in `edges` is equal to the number of graph edges.

The definitions of the arguments `nnodes`, `index`, and `edges` are illustrated with the following simple example.

### Example 7.2

Assume there are four processes 0, 1, 2, 3 with the following adjacency matrix:

process	neighbors
0	1, 3
1	0
2	3
3	0, 2

Then, the input arguments are:

```

nnodes = 4
index = 2, 3, 4, 6
edges = 1, 3, 0, 3, 0, 2
```

Thus, in C, `index[0]` is the degree of node zero, and `index[i] - index[i-1]` is the degree of node *i*, *i*=1, ..., `nnodes-1`; the list of neighbors of node zero is stored in `edges[j]`, for  $0 \leq j \leq \text{index}[0] - 1$  and the list of neighbors of node *i*, *i* > 0, is stored in `edges[j]`,  $\text{index}[i-1] \leq j \leq \text{index}[i] - 1$ .

In Fortran, `index(1)` is the degree of node zero, and `index(i+1) - index(i)` is the degree of node *i*, *i*=1, ..., `nnodes-1`; the list of neighbors of node zero is stored in `edges(j)`, for  $1 \leq j \leq \text{index}(1)$  and the list of neighbors of node *i*, *i* > 0, is stored in

`edges(j), index(i) + 1 ≤ j ≤ index(i + 1).`

A single process is allowed to be defined multiple times in the list of neighbors of a process (i.e., there may be multiple edges between two processes). A process is also allowed to be a neighbor to itself (i.e., a self loop in the graph). The adjacency matrix is allowed to be non-symmetric.

*Advice to users.* Performance implications of using multiple edges or a non-symmetric adjacency matrix are not defined. The definition of a node-neighbor edge does not imply a direction of the communication. (*End of advice to users.*)

*Advice to implementors.* The following topology information is likely to be stored with a communicator:

- Type of topology (Cartesian/graph),
- For a Cartesian topology:
  1. `ndims` (number of dimensions),
  2. `dims` (numbers of processes per coordinate direction),
  3. `periods` (periodicity information),
  4. `own_position` (own position in grid, could also be computed from rank and `dims`)

- For a graph topology:

1. `index`,
2. `edges`,

which are the vectors defining the graph structure.

For a graph structure the number of nodes is equal to the number of processes in the group. Therefore, the number of nodes does not have to be stored explicitly. An additional zero entry at the start of array `index` simplifies access to the topology information. (*End of advice to implementors.*)

#### 7.5.4 [Distributed (Graph)]Distributed Graph Constructor

[The general graph constructor assumes] `MPI_GRAPH_CREATE` requires that each process passes the full (global) communication graph to the call. This limits the scalability of this constructor. With the distributed graph interface, the communication graph is specified in a fully distributed fashion. Each process specifies only the part of the communication graph of which it is aware. Typically, this could be the set of processes from which the process will eventually receive or get data, or the set of processes to which the process will send or put data, or some combination of such edges. Two different interfaces can be used to create a distributed graph topology. `MPI_DIST_GRAPH_CREATE_ADJACENT` creates a distributed graph communicator with each process specifying [all]each of its incoming and outgoing (adjacent) edges in the logical communication graph and thus requires minimal communication during creation. [MPI\_DIST\_GRAPH\_CREATE provides full flexibility, and processes can indicate that communication will occur between other pairs of processes.] `MPI_DIST_GRAPH_CREATE` provides full flexibility such that any process can indicate that communication will occur between any pair of processes in the graph.

To provide better possibilities for optimization by the MPI library, the distributed graph constructors permit weighted communication edges and take an `info` argument that

can further influence process reordering or other optimizations performed by the MPI library. For example, hints can be provided on how edge weights are to be interpreted, the quality of the reordering, and/or the time permitted for the MPI library to process the graph.

```
MPI_DIST_GRAPH_CREATE_ADJACENT(comm_old, indegree, sources, sourceweights, out-
                                degree, destinations, destweights, info, reorder, comm_dist_graph)
```

IN	comm_old	input communicator (handle)
IN	indegree	size of <code>sources</code> and <code>sourceweights</code> arrays (non-negative integer)
IN	sources	ranks of processes for which the calling process is a destination (array of non-negative integers)
IN	sourceweights	weights of the edges into the calling process (array of non-negative integers)
IN	outdegree	size of <code>destinations</code> and <code>destweights</code> arrays (non-negative integer)
IN	destinations	ranks of processes for which the calling process is a source (array of non-negative integers)
IN	destweights	weights of the edges out of the calling process (array of non-negative integers)
IN	info	hints on optimization and interpretation of weights (handle)
IN	reorder	the ranks may be reordered ( <b>true</b> ) or not ( <b>false</b> ) (logical)
OUT	comm_dist_graph	communicator with distributed graph topology (handle)

```
int MPI_Dist_graph_create_adjacent(MPI_Comm comm_old, int indegree, const
    int sources[], const int sourceweights[], int outdegree, const
    int destinations[], const int destweights[], MPI_Info info,
    int reorder, MPI_Comm *comm_dist_graph)
```

```

MPI_Dist_graph_create_adjacent(comm_old, indegree, sources, sourceweights,
                               outdegree, destinations, destweights, info, reorder,
                               comm_dist_graph, ierror) BIND(C)
TYPE(MPI_Comm), INTENT(IN) :: comm_old
INTEGER, INTENT(IN) :: indegree, sources(indegree), outdegree,
destinations(outdegree)
INTEGER, INTENT(IN) :: sourceweights(*), destweights(*)
TYPE(MPI_Info), INTENT(IN) :: info
LOGICAL, INTENT(IN) :: reorder
TYPE(MPI_Comm), INTENT(OUT) :: comm_dist_graph
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

0	
1	
2	
3	
4	
5	
6	
7	
8	
9	
0	
1	
2	
3	
4	
5	
6	
7	
8	
9	
0	
1	ticket140.
2	ticket140.
3	ticket140.
4	ticket140.
5	ticket229.2.
6	ticket-248T.
7	
8	
9	
0	
1	
2	
3	
4	
5	
6	
7	
8	

```

1 MPI_DIST_GRAPH_CREATE_ADJACENT(COMM_OLD, INDEGREE, SOURCES, SOURCEWEIGHTS,
2     OUTDEGREE, DESTINATIONS, DESTWEIGHTS, INFO, REORDER,
3     COMM_DIST_GRAPH, IERROR)
4     INTEGER COMM_OLD, INDEGREE, SOURCES(*), SOURCEWEIGHTS(*), OUTDEGREE,
5     DESTINATIONS(*), DESTWEIGHTS(*), INFO, COMM_DIST_GRAPH, IERROR
6     LOGICAL REORDER
7
8 {MPI::Distgraphcomm MPI::Intracomm::Dist_graph_create_adjacent(int
9     indegree, const int sources[], const int sourceweights[],
10    int outdegree, const int destinations[],
11    const int destweights[], const MPI::Info& info, bool reorder)
12    const(binding deprecated, see Section 15.2) }
13
14 {MPI::Distgraphcomm
15     MPI::Intracomm::Dist_graph_create_adjacent(int indegree,
16     const int sources[], int outdegree, const int destinations[],
17     const MPI::Info& info, bool reorder) const(binding deprecated, see
18     Section 15.2) }

```

MPI\_DIST\_GRAPH\_CREATE\_ADJACENT returns a handle to a new communicator to which the distributed graph topology information is attached. *[Each process passes all information about the edges to its neighbors]* Each process passes all information about its incoming and outgoing edges in the virtual distributed graph topology. The calling processes must ensure that each edge of the graph is described in the source and in the destination process with the same weights. If there are multiple edges for a given (source,dest) pair, then the sequence of the weights of these edges does not matter. The complete communication topology is the combination of all edges shown in the sources arrays of all processes in comm\_old, which must be identical to the combination of all edges shown in the destinations arrays. Source and destination ranks must be process ranks of comm\_old. This allows a fully distributed specification of the communication graph. Isolated processes (i.e., processes with no outgoing or incoming edges, that is, processes that have specified indegree and outdegree as zero and *[that ]* thus do not occur as source or destination rank in the graph specification) are allowed.

The call creates a new communicator comm\_dist\_graph of distributed graph topology type to which topology information has been attached. The number of processes in comm\_dist\_graph is identical to the number of processes in comm\_old. The call to MPI\_DIST\_GRAPH\_CREATE\_ADJACENT is collective.

Weights are specified as non-negative integers and can be used to influence the process remapping strategy and other internal MPI optimizations. For instance, approximate count arguments of later communication calls along specific edges could be used as their edge weights. Multiplicity of edges can likewise indicate more intense communication between pairs of processes. However, the exact meaning of edge weights is not specified by the MPI standard and is left to the implementation. In C or Fortran, an application can supply the special value MPI\_UNWEIGHTED for the weight array to indicate that all edges have the same (effectively no) weight. In C++, this constant does not exist and the weight arguments may be omitted from the argument list. It is erroneous to supply MPI\_UNWEIGHTED, or in C++ omit the weight arrays, for some but not all processes of comm\_old. Note that MPI\_UNWEIGHTED is not a special weight value; rather it is a special value for the total array argument. In C, one would expect it to be NULL. In Fortran, MPI\_UNWEIGHTED is an

object like MPI\_BOTTOM (not usable for initialization or assignment). See Section 2.5.4.

The meaning of the `info` and `reorder` arguments is defined in the description of the following routine.

`MPI_DIST_GRAPH_CREATE(comm_old, n, sources, degrees, destinations, weights, info, reorder, comm_dist_graph)`

IN	<code>comm_old</code>	input communicator (handle)
IN	<code>n</code>	number of source nodes for which this process specifies edges (non-negative integer)
IN	<code>sources</code>	array containing the <code>n</code> source nodes for which this process specifies edges (array of non-negative integers)
IN	<code>degrees</code>	array specifying the number of destinations for each source node in the source node array (array of non-negative integers)
IN	<code>destinations</code>	destination nodes for the source nodes in the source node array (array of non-negative integers)
IN	<code>weights</code>	weights for source to destination edges (array of non-negative integers)
IN	<code>info</code>	hints on optimization and interpretation of weights (handle)
IN	<code>reorder</code>	the process may be reordered (true) or not (false) (logical)
OUT	<code>comm_dist_graph</code>	communicator with distributed graph topology added (handle)

```

int MPI_Dist_graph_create(MPI_Comm comm_old, int n, const int sources[],
    const int degrees[], const int destinations[], const
    int weights[], MPI_Info info, int reorder,
    MPI_Comm *comm_dist_graph)
MPI_Dist_graph_create(comm_old, n, sources, degrees, destinations, weights,
    info, reorder, comm_dist_graph, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm_old
    INTEGER, INTENT(IN) :: n, sources(n), degrees(n), destinations(*)
    INTEGER, INTENT(IN) :: weights(*)
    TYPE(MPI_Info), INTENT(IN) :: info
    LOGICAL, INTENT(IN) :: reorder
    TYPE(MPI_Comm), INTENT(OUT) :: comm_dist_graph
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_DIST_GRAPH_CREATE(COMM_OLD, N, SOURCES, DEGREES, DESTINATIONS, WEIGHTS,
    INFO, REORDER, COMM_DIST_GRAPH, IERROR)
    INTEGER COMM_OLD, N, SOURCES(*), DEGREES(*), DESTINATIONS(*),
    WEIGHTS(*), INFO, COMM_DIST_GRAPH, IERROR
    LOGICAL REORDER

```



```

1 {MPI::Distgraphcomm MPI::Intracomm::Dist_graph_create(int n,
2               const int sources[], const int degrees[], const int
3               destinations[], const int weights[], const MPI::Info& info,
4               bool reorder) const(binding deprecated, see Section 15.2) }
5
6 {MPI::Distgraphcomm MPI::Intracomm::Dist_graph_create(int n,
7               const int sources[], const int degrees[],
8               const int destinations[], const MPI::Info& info, bool reorder)
9               const(binding deprecated, see Section 15.2) }

```

MPI\_DIST\_GRAPH\_CREATE returns a handle to a new communicator to which the distributed graph topology information is attached. Concretely, each process calls the constructor with a set of directed (source,destination) communication edges as described below. Every process passes an array of  $n$  source nodes in the `sources` array. For each source node, a non-negative number of destination nodes is specified in the `degrees` array. The destination nodes are stored in the corresponding consecutive segment of the `destinations` array. More precisely, if the  $i$ -th node in `sources` is  $s$ , this specifies `degrees[i]` edges  $(s,d)$  with  $d$  of the  $j$ -th such edge stored in `destinations[degrees[0]+...+degrees[i-1]+j]`. The weight of this edge is stored in `weights[degrees[0]+...+degrees[i-1]+j]`. Both the `sources` and the `destinations` arrays may contain the same node more than once, and the order in which nodes are listed as destinations or sources is not significant. Similarly, different processes may specify edges with the same source and destination nodes. Source and destination nodes must be process ranks of `comm_old`. Different processes may specify different numbers of source and destination nodes, as well as different source to destination edges. This allows a fully distributed specification of the communication graph. Isolated processes (i.e., processes with no outgoing or incoming edges, that is, processes that do not occur as source or destination node in the graph specification) are allowed.

The call creates a new communicator `comm_dist_graph` of distributed graph topology type to which topology information has been attached. The number of processes in `comm_dist_graph` is identical to the number of processes in `comm_old`. The call to [ `MPI_Dist_graph_create`] `MPI_DIST_GRAPH_CREATE` is collective.

If `reorder = false`, all processes will have the same rank in `comm_dist_graph` as in `comm_old`. If `reorder = true` then the MPI library is free to remap to other processes (of `comm_old`) in order to improve communication on the edges of the communication graph. The weight associated with each edge is a hint to the MPI library about the amount or intensity of communication on that edge, and may be used to compute a “best” reordering.

Weights are specified as non-negative integers and can be used to influence the process remapping strategy and other internal MPI optimizations. For instance, approximate count arguments of later communication calls along specific edges could be used as their edge weights. Multiplicity of edges can likewise indicate more intense communication between pairs of processes. However, the exact meaning of edge weights is not specified by the MPI standard and is left to the implementation. In C or Fortran, an application can supply the special value `MPI_UNWEIGHTED` for the weight array to indicate that all edges have the same (effectively no) weight. In C++, this constant does not exist and the weights argument may be omitted from the argument list. It is erroneous to supply `MPI_UNWEIGHTED`, or in C++ omit the weight arrays, for some but not all processes of `comm_old`. Note that `MPI_UNWEIGHTED` is not a special weight value; rather it is a special value for the total array argument. In C, one would expect it to be `NULL`. In Fortran, `MPI_UNWEIGHTED` is



an object like `MPI_BOTTOM` (not usable for initialization or assignment). See Section 2.5.4

The meaning of the `weights` argument can be influenced by the `info` argument. Info arguments can be used to guide the mapping; possible options include minimizing the maximum number of edges between processes on different SMP nodes, or minimizing the sum of all such edges. An MPI implementation is not obliged to follow specific hints, and it is valid for an MPI implementation not to do any reordering. An MPI implementation may specify more info key-value pairs. All processes must specify the same set of key-value info pairs.

*Advice to implementors.* MPI implementations must document any additionally supported key-value info pairs. `MPI_INFO_NULL` is always valid, and may indicate the default creation of the distributed graph topology to the MPI library.

An implementation does not explicitly need to construct the topology from its distributed parts. However, all processes can construct the full topology from the distributed specification and use this in a call to `MPI_GRAPH_CREATE` to create the topology. This may serve as a reference implementation of the functionality, and may be acceptable for small communicators. However, a scalable high-quality implementation would save the topology graph in a distributed way. (*End of advice to implementors.*)

**Example 7.3** As for Example 7.2, assume there are four processes 0, 1, 2, 3 with the following adjacency matrix and unit edge weights:

process	neighbors
0	1, 3
1	0
2	3
3	0, 2

With `MPI_DIST_GRAPH_CREATE`, this graph could be constructed in many different ways. One way would be that each process specifies its outgoing edges. The arguments per process would be:

process	n	sources	degrees	destinations	weights
0	1	0	2	1,3	1,1
1	1	1	1	0	1
2	1	2	1	3	1
3	1	3	2	0,2	1,1

Another way would be to pass the whole graph on process 0, which could be done with the following arguments per process:

process	n	sources	degrees	destinations	weights
0	4	0,1,2,3	2,1,1,2	1,3,0,3,0,2	1,1,1,1,1,1
1	0	-	-	-	-
2	0	-	-	-	-
3	0	-	-	-	-

In both cases above, the application could supply `MPI_UNWEIGHTED` instead of explic-

itly providing identical weights.

MPI\_DIST\_GRAPH\_CREATE\_ADJACENT could be used to specify this graph using the following arguments:

process	indegree	sources	sourceweights	outdegree	destinations	destweights
0	2	1,3	1,1	2	1,3	1,1
1	1	0	1	1	0	1
2	1	3	1	1	3	1
3	2	0,2	1,1	2	0,2	1,1

**Example 7.4** A two-dimensional  $P \times Q$  torus where all processes communicate along the dimensions and along the diagonal edges. This cannot be [modelled]modeled with Cartesian topologies, but can easily be captured with MPI\_DIST\_GRAPH\_CREATE as shown in the following code. In this example, the communication along the dimensions is twice as heavy as the communication along the diagonals:

```

/*
Input:      dimensions P, Q
Condition:  number of processes equal to P*Q; otherwise only
            ranks smaller than P*Q participate
*/
int rank, x, y;
int sources[1], degrees[1];
int destinations[8], weights[8];
MPI_Comm comm_dist_graph;

MPI_Comm_rank(MPI_COMM_WORLD, &rank);

/* get x and y dimension */
y=rank/P; x=rank%P;

/* get my communication partners along x dimension */
destinations[0] = P*y+(x+1)%P; weights[0] = 2;
destinations[1] = P*y+(P+x-1)%P; weights[1] = 2;

/* get my communication partners along y dimension */
destinations[2] = P*((y+1)%Q)+x; weights[2] = 2;
destinations[3] = P*((Q+y-1)%Q)+x; weights[3] = 2;

/* get my communication partners along diagonals */
destinations[4] = P*((y+1)%Q)+(x+1)%P; weights[4] = 1;
destinations[5] = P*((Q+y-1)%Q)+(x+1)%P; weights[5] = 1;
destinations[6] = P*((y+1)%Q)+(P+x-1)%P; weights[6] = 1;
destinations[7] = P*((Q+y-1)%Q)+(P+x-1)%P; weights[7] = 1;

sources[0] = rank;
degrees[0] = 8;
MPI_Dist_graph_create(MPI_COMM_WORLD, 1, sources, degrees, destinations,

```

```
weights, MPI_INFO_NULL, 1, &comm_dist_graph);
```

### 7.5.5 Topology Inquiry Functions

If a topology has been defined with one of the above functions, then the topology information can be looked up using inquiry functions. They all are local calls.

**MPI\_TOPO\_TEST(comm, status)**

IN	comm	communicator (handle)
OUT	status	topology type of communicator comm (state)

```
int MPI_Topo_test(MPI_Comm comm, int *status)
```

```
MPI_Topo_test(comm, status, ierror) BIND(C)
  TYPE(MPI_Comm), INTENT(IN) :: comm
  INTEGER, INTENT(OUT) :: status
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TOPO_TEST(COMM, STATUS, IERROR)
  INTEGER COMM, STATUS, IERROR
```

```
{int MPI::Comm::Get_topology() const(binding deprecated, see Section 15.2) }
```

The function **MPI\_TOPO\_TEST** returns the type of topology that is assigned to a communicator.

The output value **status** is one of the following:

<b>MPI_GRAPH</b>	graph topology
<b>MPI_CART</b>	Cartesian topology
<b>MPI_DIST_GRAPH</b>	distributed graph topology
<b>MPI_UNDEFINED</b>	no topology

**MPI\_GRAPHDIMS\_GET(comm, nnodes, nedges)**

IN	comm	communicator for group with graph structure (handle)
OUT	nnodes	number of nodes in graph (integer) (same as number of processes in the group)
OUT	nedges	number of edges in graph (integer)

```
int MPI_Graphdims_get(MPI_Comm comm, int *nnodes, int *nedges)
```

```
MPI_Graphdims_get(comm, nnodes, nedges, ierror) BIND(C)
  TYPE(MPI_Comm), INTENT(IN) :: comm
  INTEGER, INTENT(OUT) :: nnodes, nedges
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_GRAPHDIMS_GET(COMM, NNODES, NEDGES, IERROR)
  INTEGER COMM, NNODES, NEDGES, IERROR
```

```
1 {void MPI::Graphcomm::Get_dims(int nnodes[], int nedges[]) const(binding
2 deprecated, see Section 15.2) }
```

Functions MPI\_GRAPHDIMS\_GET and MPI\_GRAPH\_GET retrieve the graph-topology information that was associated with a communicator by MPI\_GRAPH\_CREATE.

The information provided by MPI\_GRAPHDIMS\_GET can be used to dimension the vectors `index` and `edges` correctly for the following call to MPI\_GRAPH\_GET.

```
9 MPI_GRAPH_GET(comm, maxindex, maxedges, index, edges)
```

11	IN	comm	communicator with graph structure (handle)
12	IN	maxindex	length of vector <code>index</code> in the calling program
13			(integer)
14	IN	maxedges	length of vector <code>edges</code> in the calling program
15			(integer)
16	OUT	index	array of integers containing the graph structure (for
17			details see the definition of MPI_GRAPH_CREATE)
18			
19	OUT	edges	array of integers containing the graph structure

```
21 int MPI_Graph_get(MPI_Comm comm, int maxindex, int maxedges,
22 int [*index]index[], int [*edges]edges[])
```

```
ticket125. 24 MPI_Graph_get(comm, maxindex, maxedges, index, edges, ierror) BIND(C)
ticket126. 25 TYPE(MPI_Comm), INTENT(IN) :: comm
26 INTEGER, INTENT(IN) :: maxindex, maxedges
27 INTEGER, INTENT(OUT) :: index(maxindex), edges(maxedges)
28 INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
29 MPI_GRAPH_GET(COMM, MAXINDEX, MAXEDGES, INDEX, EDGES, IERROR)
30 INTEGER COMM, MAXINDEX, MAXEDGES, INDEX(*), EDGES(*), IERROR
```

```
32 {void MPI::Graphcomm::Get_topo(int maxindex, int maxedges, int index[],
33 int edges[]) const(binding deprecated, see Section 15.2) }
```

```
36 MPI_CARTDIM_GET(comm, ndims)
```

37	IN	comm	communicator with Cartesian structure (handle)
38			
39	OUT	ndims	number of dimensions of the Cartesian structure (in-
40			teger)

```
42 int MPI_Cartdim_get(MPI_Comm comm, int *ndims)
```

```
ticket-248T. 43 MPI_Cartdim_get(comm, ndims, ierror) BIND(C)
44 TYPE(MPI_Comm), INTENT(IN) :: comm
45 INTEGER, INTENT(OUT) :: ndims
46 INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
48 MPI_CARTDIM_GET(COMM, NDIMS, IERROR)
```

```

    INTEGER COMM, NDIMS, IERROR
    {int MPI::Cartcomm::Get_dim() const(binding deprecated, see Section 15.2) }

```

The functions `MPI_CARTDIM_GET` and `MPI_CART_GET` return the Cartesian topology information that was associated with a communicator by `MPI_CART_CREATE`. If `comm` is associated with a zero-dimensional Cartesian topology, `MPI_CARTDIM_GET` returns `ndims=0` and `MPI_CART_GET` will keep all output arguments unchanged.

```

MPI_CART_GET(comm, maxdims, dims, periods, coords)

```

IN	comm	communicator with Cartesian structure (handle)
IN	maxdims	length of vectors <code>dims</code> , <code>periods</code> , and <code>coords</code> in the calling program (integer)
OUT	dims	number of processes for each Cartesian dimension (array of integer)
OUT	periods	periodicity (true/false) for each Cartesian dimension (array of logical)
OUT	coords	coordinates of calling process in Cartesian structure (array of integer)

```

int MPI_Cart_get(MPI_Comm comm, int maxdims, int [*dims]dims[],
                int [*periods]periods[], int [*coords]coords[])

```

```

MPI_Cart_get(comm, maxdims, dims, periods, coords, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, INTENT(IN) :: maxdims
    INTEGER, INTENT(OUT) :: dims(maxdims), coords(maxdims)
    LOGICAL, INTENT(OUT) :: periods(maxdims)
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_CART_GET(COMM, MAXDIMS, DIMS, PERIODS, COORDS, IERROR)
    INTEGER COMM, MAXDIMS, DIMS(*), COORDS(*), IERROR
    LOGICAL PERIODS(*)

```

```

{void MPI::Cartcomm::Get_topo(int maxdims, int dims[], bool periods[],
    int coords[]) const(binding deprecated, see Section 15.2) }

```

```

MPI_CART_RANK(comm, coords, rank)

```

IN	comm	communicator with Cartesian structure (handle)
IN	coords	integer array (of size <code>ndims</code> ) specifying the Cartesian coordinates of a process
OUT	rank	rank of specified process (integer)

```

int MPI_Cart_rank(MPI_Comm comm, const int [*coords]coords[], int *rank)

```

```

MPI_Cart_rank(comm, coords, rank, ierror) BIND(C)

```

```

1      TYPE(MPI_Comm), INTENT(IN) :: comm
2      INTEGER, INTENT(IN) :: coords(*)
3      INTEGER, INTENT(OUT) :: rank
4      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
5
6 MPI_CART_RANK(COMM, COORDS, RANK, IERROR)
7     INTEGER COMM, COORDS(*), RANK, IERROR
8
9 {int MPI::Cartcomm::Get_cart_rank(const int coords[]) const (binding
10    deprecated, see Section 15.2) }

```

For a process group with Cartesian structure, the function `MPI_CART_RANK` translates the logical process coordinates to process ranks as they are used by the point-to-point routines.

For dimension `i` with `periods(i) = true`, if the coordinate, `coords(i)`, is out of range, that is, `coords(i) < 0` or `coords(i) ≥ dims(i)`, it is shifted back to the interval  $0 \leq \text{coords}(i) < \text{dims}(i)$  automatically. Out-of-range coordinates are erroneous for non-periodic dimensions.

If `comm` is associated with a zero-dimensional Cartesian topology, `coords` is not significant and 0 is returned in `rank`.

```

21 MPI_CART_COORDS(comm, rank, maxdims, coords)
22
23     IN      comm      communicator with Cartesian structure (handle)
24     IN      rank      rank of a process within group of comm (integer)
25     IN      maxdims   length of vector coords in the calling program (inte-
26                       ger)
27
28     OUT     coords    integer array (of size ndims) containing the Cartesian
29                       coordinates of specified process (array of integers)

```

```

31 int MPI_Cart_coords(MPI_Comm comm, int rank, int maxdims,
32    int [*coords] coords[])

```

```

33 MPI_Cart_coords(comm, rank, maxdims, coords, ierror) BIND(C)
34     TYPE(MPI_Comm), INTENT(IN) :: comm
35     INTEGER, INTENT(IN) :: rank, maxdims
36     INTEGER, INTENT(OUT) :: coords(maxdims)
37     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
38

```

```

39 MPI_CART_COORDS(COMM, RANK, MAXDIMS, COORDS, IERROR)
40     INTEGER COMM, RANK, MAXDIMS, COORDS(*), IERROR
41
42 {void MPI::Cartcomm::Get_coords(int rank, int maxdims, int coords[])
43    const (binding deprecated, see Section 15.2) }

```

The inverse mapping, rank-to-coordinates translation is provided by `MPI_CART_COORDS`.

If `comm` is associated with a zero-dimensional Cartesian topology, `coords` will be unchanged.

```

MPI_GRAPH_NEIGHBORS_COUNT(comm, rank, nneighbors)
    IN      comm      communicator with graph topology (handle)
    IN      rank      rank of process in group of comm (integer)
    OUT     nneighbors number of neighbors of specified process (integer)

int MPI_Graph_neighbors_count(MPI_Comm comm, int rank, int *nneighbors)

MPI_Graph_neighbors_count(comm, rank, nneighbors, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, INTENT(IN) :: rank
    INTEGER, INTENT(OUT) :: nneighbors
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_GRAPH_NEIGHBORS_COUNT(COMM, RANK, NNEIGHBORS, IERROR)
    INTEGER COMM, RANK, NNEIGHBORS, IERROR

{int MPI::Graphcomm::Get_neighbors_count(int rank) const(binding deprecated,
    see Section 15.2) }

MPI_GRAPH_NEIGHBORS(comm, rank, maxneighbors, neighbors)
    IN      comm      communicator with graph topology (handle)
    IN      rank      rank of process in group of comm (integer)
    IN      maxneighbors size of array neighbors (integer)
    OUT     neighbors  ranks of processes that are neighbors to specified process (array of integer)

int MPI_Graph_neighbors(MPI_Comm comm, int rank, int maxneighbors,
    int [*neighbors]neighbors[])

MPI_Graph_neighbors(comm, rank, maxneighbors, neighbors, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, INTENT(IN) :: rank, maxneighbors
    INTEGER, INTENT(OUT) :: neighbors(maxneighbors)
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_GRAPH_NEIGHBORS(COMM, RANK, MAXNEIGHBORS, NEIGHBORS, IERROR)
    INTEGER COMM, RANK, MAXNEIGHBORS, NEIGHBORS(*), IERROR

{void MPI::Graphcomm::Get_neighbors(int rank, int maxneighbors, int
    neighbors[]) const(binding deprecated, see Section 15.2) }

    MPI_GRAPH_NEIGHBORS_COUNT and MPI_GRAPH_NEIGHBORS provide adjacency
    information for a [general] graph topology. The returned count and array of neighbors for
    the queried rank will both include all neighbors and reflect the same edge ordering as
    was specified by the original call to MPI_GRAPH_CREATE. Specifically,
    MPI_GRAPH_NEIGHBORS_COUNT and MPI_GRAPH_NEIGHBORS will return values based

```

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on the original `index` and `edges` array passed to `MPI_GRAPH_CREATE` (assuming that `index[-1]` effectively equals zero):

- The `[count]`number of neighbors (`nneighbors`) returned from `MPI_GRAPH_NEIGHBORS_COUNT` will be `(index[rank] - index[rank-1])`.
- The `neighbors` array returned from `MPI_GRAPH_NEIGHBORS` will be `edges[index[rank-1]]` through `edges[index[rank]-1]`.

### Example 7.5

Assume there are four processes 0, 1, 2, 3 with the following adjacency matrix (note that some neighbors are listed multiple times):

process	neighbors
0	1, 1, 3
1	0, 0
2	3
3	0, 2, 2

Thus, the input arguments to `MPI_GRAPH_CREATE` are:

```
nnodes = 4
index = 3, 5, 6, 9
edges = 1, 1, 3, 0, 0, 3, 0, 2, 2
```

Therefore, calling `MPI_GRAPH_NEIGHBORS_COUNT` and `MPI_GRAPH_NEIGHBORS` for each of the 4 processes will return:

Input rank	Count	Neighbors
0	3	1, 1, 3
1	2	0, 0
2	1	3
3	3	0, 2, 2

### Example 7.6

Suppose that `comm` is a communicator with a shuffle-exchange topology. The group has  $2^n$  members. Each process is labeled by  $a_1, \dots, a_n$  with  $a_i \in \{0, 1\}$ , and has three neighbors:  $\text{exchange}(a_1, \dots, a_n) = a_1, \dots, a_{n-1}, \bar{a}_n$  ( $\bar{a} = 1 - a$ ),  $\text{shuffle}(a_1, \dots, a_n) = a_2, \dots, a_n, a_1$ , and  $\text{unshuffle}(a_1, \dots, a_n) = a_n, a_1, \dots, a_{n-1}$ . The graph adjacency list is illustrated below for  $n = 3$ .



node	exchange neighbors(1)	shuffle neighbors(2)	unshuffle neighbors(3)
0 (000)	1	0	0
1 (001)	0	2	4
2 (010)	3	4	1
3 (011)	2	6	5
4 (100)	5	1	2
5 (101)	4	3	6
6 (110)	7	5	3
7 (111)	6	7	7

Suppose that the communicator `comm` has this topology associated with it. The following code fragment cycles through the three types of neighbors and performs an appropriate permutation for each.

```

C  assume: each process has stored a real number A.
C  extract neighborhood information
    CALL MPI_COMM_RANK(comm, myrank, ierr)
    CALL MPI_GRAPH_NEIGHBORS(comm, myrank, 3, neighbors, ierr)
C  perform exchange permutation
    CALL MPI_SENDRECV_REPLACE(A, 1, MPI_REAL, neighbors(1), 0,
+    neighbors(1), 0, comm, status, ierr)
C  perform shuffle permutation
    CALL MPI_SENDRECV_REPLACE(A, 1, MPI_REAL, neighbors(2), 0,
+    neighbors(3), 0, comm, status, ierr)
C  perform unshuffle permutation
    CALL MPI_SENDRECV_REPLACE(A, 1, MPI_REAL, neighbors(3), 0,
+    neighbors(2), 0, comm, status, ierr)

```

`MPI_DIST_GRAPH_NEIGHBORS_COUNT` and `MPI_DIST_GRAPH_NEIGHBORS` provide adjacency information for a distributed graph topology.

`MPI_DIST_GRAPH_NEIGHBORS_COUNT(comm, indegree, outdegree, weighted)`

IN	comm	communicator with distributed graph topology (handle)
OUT	indegree	number of edges into this process (non-negative integer)
OUT	outdegree	number of edges out of this process (non-negative integer)
OUT	weighted	false if <code>MPI_UNWEIGHTED</code> was supplied during creation, true otherwise (logical)

```

int MPI_Dist_graph_neighbors_count(MPI_Comm comm, int *indegree,
    int *outdegree, int *weighted)

```

```

MPI_Dist_graph_neighbors_count(comm, indegree, outdegree, weighted, ierror)
    BIND(C)

```

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```

1      TYPE(MPI_Comm), INTENT(IN) :: comm
2      INTEGER, INTENT(OUT) :: indegree, outdegree
3      LOGICAL, INTENT(OUT) :: weighted
4      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
5
6 MPI_DIST_GRAPH_NEIGHBORS_COUNT(COMM, INDEGREE, OUTDEGREE, WEIGHTED, IERROR)
7     INTEGER COMM, INDEGREE, OUTDEGREE, IERROR
8     LOGICAL WEIGHTED
9
10 {void MPI::Distgraphcomm::Get_dist_neighbors_count(int rank,
11     int indegree[], int outdegree[], bool& weighted) const(binding
12     deprecated, see Section 15.2) }
13
14 MPI_DIST_GRAPH_NEIGHBORS(comm, maxindegree, sources, sourceweights, maxoutdegree,
15     destinations, destweights)
16
17 IN      comm      communicator with distributed graph topology (han-
18     dle)
19 IN      maxindegree      size of sources and sourceweights arrays (non-negative
20     integer)
21 OUT     sources      processes for which the calling process is a destination
22     (array of non-negative integers)
23 OUT     sourceweights      weights of the edges into the calling process (array of
24     non-negative integers)
25
26 IN      maxoutdegree      size of destinations and destweights arrays (non-negative
27     integer)
28 OUT     destinations      processes for which the calling process is a source (ar-
29     ray of non-negative integers)
30
31 OUT     destweights      weights of the edges out of the calling process (array
32     of non-negative integers)
33
34 int MPI_Dist_graph_neighbors(MPI_Comm comm, int maxindegree, int sources[],
35     int sourceweights[], int maxoutdegree, int destinations[],
36     int destweights[])
37
38 MPI_Dist_graph_neighbors(comm, maxindegree, sources, sourceweights,
39     maxoutdegree, destinations, destweights, ierror) BIND(C)
40     TYPE(MPI_Comm), INTENT(IN) :: comm
41     INTEGER, INTENT(IN) :: maxindegree, maxoutdegree
42     INTEGER, INTENT(OUT) :: sources(maxindegree),
43     destinations(maxoutdegree)
44     INTEGER :: sourceweights(*), destweights(*)
45     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
46
47 MPI_DIST_GRAPH_NEIGHBORS(COMM, MAXINDEGREE, SOURCES, SOURCEWEIGHTS,
48     MAXOUTDEGREE, DESTINATIONS, DESTWEIGHTS, IERROR)
49     INTEGER COMM, MAXINDEGREE, SOURCES(*), SOURCEWEIGHTS(*), MAXOUTDEGREE,

```

```

DESTINATIONS(*), DESTWEIGHTS(*), IERROR
{void MPI::Distgraphcomm::Get_dist_neighbors(int maxindegree,
      int sources[], int sourceweights[], int maxoutdegree,
      int destinations[], int destweights[]) (binding deprecated, see
      Section 15.2) }
```

These calls are local. The number of edges into and out of the process returned by MPI\_DIST\_GRAPH\_NEIGHBORS\_COUNT are the total number of such edges given in the call to MPI\_DIST\_GRAPH\_CREATE\_ADJACENT or MPI\_DIST\_GRAPH\_CREATE (potentially by processes other than the calling process in the case of MPI\_DIST\_GRAPH\_CREATE). Multiply defined edges are all counted and returned by MPI\_DIST\_GRAPH\_NEIGHBORS in some order. If MPI\_UNWEIGHTED is supplied for sourceweights or destweights or both, or if MPI\_UNWEIGHTED was supplied during the construction of the graph then no weight information is returned in that array or those arrays. [The ]If the communicator was created with MPI\_DIST\_GRAPH\_CREATE\_ADJACENT then for each rank in comm, the order of the values in sources and destinations is identical to the input that was used by the process with the same rank in comm\_old in the creation call. If the communicator was created with MPI\_DIST\_GRAPH\_CREATE then the only requirement on the order of values in sources and destinations is that two calls to the routine with same input argument comm will return the same sequence of edges. If maxindegree or maxoutdegree is smaller than the numbers returned by MPI\_DIST\_GRAPH\_NEIGHBOR\_COUNT, then only the first part of the full list is returned. [Note, that the order of returned edges does need not to be identical to the order that was provided in the creation of comm for the case that MPI\_DIST\_GRAPH\_CREATE\_ADJACENT was used. ]

*Advice to implementors.* Since the query calls are defined to be local, each process needs to store the list of its neighbors with incoming and outgoing edges. Communication is required at the collective MPI\_DIST\_GRAPH\_CREATE call in order to compute the neighbor lists for each process from the distributed graph specification. (*End of advice to implementors.*)

### 7.5.6 Cartesian Shift Coordinates

If the process topology is a Cartesian structure, an MPI\_SENDRECV operation is likely to be used along a coordinate direction to perform a shift of data. As input, MPI\_SENDRECV takes the rank of a source process for the receive, and the rank of a destination process for the send. If the function MPI\_CART\_SHIFT is called for a Cartesian process group, it provides the calling process with the above identifiers, which then can be passed to MPI\_SENDRECV. The user specifies the coordinate direction and the size of the step (positive or negative). The function is local.

1	MPI_CART_SHIFT(comm, direction, disp, rank_source, rank_dest)		
2	IN	comm	communicator with Cartesian structure (handle)
3	IN	direction	coordinate dimension of shift (integer)
4	IN	disp	displacement (> 0: upwards shift, < 0: downwards shift) (integer)
5	OUT	rank_source	rank of source process (integer)
6	OUT	rank_dest	rank of destination process (integer)

```

10
11 int MPI_Cart_shift(MPI_Comm comm, int direction, int disp,
12                   int *rank_source, int *rank_dest)
13
14 MPI_Cart_shift(comm, direction, disp, rank_source, rank_dest, ierror)
15     BIND(C)
16     TYPE(MPI_Comm), INTENT(IN) :: comm
17     INTEGER, INTENT(IN) :: direction, disp
18     INTEGER, INTENT(OUT) :: rank_source, rank_dest
19     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
20
21 MPI_CART_SHIFT(COMM, DIRECTION, DISP, RANK_SOURCE, RANK_DEST, IERROR)
22     INTEGER COMM, DIRECTION, DISP, RANK_SOURCE, RANK_DEST, IERROR
23
24 {void MPI::Cartcomm::Shift(int direction, int disp, int& rank_source,
25                             int& rank_dest) const(binding deprecated, see Section 15.2) }
```

The `[direction]direction` argument indicates the coordinate dimension to be traversed by the shift. The dimensions are numbered from 0 to `ndims-1`, where `ndims` is the number of dimensions.

Depending on the periodicity of the Cartesian group in the specified coordinate direction, `MPI_CART_SHIFT` provides the identifiers for a circular or an end-off shift. In the case of an end-off shift, the value `MPI_PROC_NULL` may be returned in `rank_source` or `rank_dest`, indicating that the source or the destination for the shift is out of range.

It is erroneous to call `MPI_CART_SHIFT` with a direction that is either negative or greater than or equal to the number of dimensions in the Cartesian communicator. This implies that it is erroneous to call `MPI_CART_SHIFT` with a `comm` that is associated with a zero-dimensional Cartesian topology.

### Example 7.7

The communicator, `comm`, has a two-dimensional, periodic, Cartesian topology associated with it. A two-dimensional array of `REALs` is stored one element per process, in variable `A`. One wishes to skew this array, by shifting column `i` (vertically, i.e., along the column) by `i` steps.

```

42 ....
43
44 C find process rank
45     CALL MPI_COMM_RANK(comm, rank, ierr)
46 C find Cartesian coordinates
47     CALL MPI_CART_COORDS(comm, rank, maxdims, coords, ierr)
48 C compute shift source and destination
```

```

CALL MPI_CART_SHIFT(comm, 0, coords(2), source, dest, ierr)
C skew array
CALL MPI_SENDRECV_REPLACE(A, 1, MPI_REAL, dest, 0, source, 0, comm,
+                          status, ierr)

```

*Advice to users.* In Fortran, the dimension indicated by `DIRECTION = i` has `DIMS(i+1)` nodes, where `DIMS` is the array that was used to create the grid. In C, the dimension indicated by `direction = i` is the dimension specified by `dims[i]`. (*End of advice to users.*)

### 7.5.7 Partitioning of Cartesian [s]Structures

`MPI_CART_SUB(comm, remain_dims, newcomm)`

IN	comm	communicator with Cartesian structure (handle)
IN	remain_dims	the i-th entry of <code>remain_dims</code> specifies whether the i-th dimension is kept in the subgrid ( <code>true</code> ) or is dropped ( <code>false</code> ) (logical vector)
OUT	newcomm	communicator containing the subgrid that includes the calling process (handle)

```

int MPI_Cart_sub(MPI_Comm comm, const int [*remain_dims]remain_dims[],
MPI_Comm *newcomm)

```

```

MPI_Cart_sub(comm, remain_dims, newcomm, ierror) BIND(C)
  TYPE(MPI_Comm), INTENT(IN) :: comm
  LOGICAL, INTENT(IN) :: remain_dims(*)
  TYPE(MPI_Comm), INTENT(OUT) :: newcomm
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_CART_SUB(COMM, REMAIN_DIMS, NEWCOMM, IERROR)
  INTEGER COMM, NEWCOMM, IERROR
  LOGICAL REMAIN_DIMS(*)

```

```

{MPI::Cartcomm MPI::Cartcomm::Sub(const bool remain_dims[]) const(binding
  deprecated, see Section 15.2) }

```

If a Cartesian topology has been created with `MPI_CART_CREATE`, the function `MPI_CART_SUB` can be used to partition the communicator group into subgroups that form lower-dimensional Cartesian subgrids, and to build for each subgroup a communicator with the associated subgrid Cartesian topology. If all entries in `remain_dims` are false or `comm` is already associated with a zero-dimensional Cartesian topology then `newcomm` is associated with a zero-dimensional Cartesian topology. (This function is closely related to `MPI_COMM_SPLIT`.)

#### Example 7.8

Assume that `MPI_CART_CREATE(..., comm)` has defined a  $(2 \times 3 \times 4)$  grid. Let `remain_dims = (true, false, true)`. Then a call to,

```

MPI_CART_SUB(comm, remain_dims, comm_new),

```

will create three communicators each with eight processes in a  $2 \times 4$  Cartesian topology. If `remain_dims = (false, false, true)` then the call to `MPI_CART_SUB(comm, remain_dims, comm_new)` will create six non-overlapping communicators, each with four processes, in a one-dimensional Cartesian topology.

### 7.5.8 Low-Level Topology Functions

The two additional functions introduced in this section can be used to implement all other topology functions. In general they will not be called by the user directly, unless he or she is creating additional virtual topology capability other than that provided by MPI. **The two calls are both local.**

`MPI_CART_MAP(comm, ndims, dims, periods, newrank)`

IN	<code>comm</code>	input communicator (handle)
IN	<code>ndims</code>	number of dimensions of Cartesian structure (integer)
IN	<code>dims</code>	integer array of size <code>ndims</code> specifying the number of processes in each coordinate direction
IN	<code>periods</code>	logical array of size <code>ndims</code> specifying the periodicity specification in each coordinate direction
OUT	<code>newrank</code>	reordered rank of the calling process; MPI_UNDEFINED if calling process does not belong to grid (integer)

```
int MPI_Cart_map(MPI_Comm comm, int ndims, const int [*dims]dims[], const
                int *periods, int *newrank)
```

```
MPI_Cart_map(comm, ndims, dims, periods, newrank, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, INTENT(IN) :: ndims, dims(ndims)
    LOGICAL, INTENT(IN) :: periods(ndims)
    INTEGER, INTENT(OUT) :: newrank
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_CART_MAP(COMM, NDIMS, DIMS, PERIODS, NEWRANK, IERROR)
    INTEGER COMM, NDIMS, DIMS(*), NEWRANK, IERROR
    LOGICAL PERIODS(*)
```

```
{int MPI::Cartcomm::Map(int ndims, const int dims[], const bool periods[])
    const(binding deprecated, see Section 15.2) }
```

`MPI_CART_MAP` computes an “optimal” placement for the calling process on the physical machine. A possible implementation of this function is to always return the rank of the calling process, that is, not to perform any reordering.

*Advice to implementors.* The function `MPI_CART_CREATE(comm, ndims, dims, periods, reorder, comm_cart)`, with `reorder = true` can be implemented by calling `MPI_CART_MAP(comm, ndims, dims, periods, newrank)`, then calling `MPI_COMM_SPLIT(comm, color, key, comm_cart)`, with `color = 0` if `newrank`  $\neq$

MPI\_UNDEFINED, color = MPI\_UNDEFINED otherwise, and key = newrank. If ndims is zero then a zero-dimensional Cartesian topology is created.

The function MPI\_CART\_SUB(comm, remain\_dims, comm\_new) can be implemented by a call to MPI\_COMM\_SPLIT(comm, color, key, comm\_new), using a single number encoding of the lost dimensions as color and a single number encoding of the preserved dimensions as key.

All other Cartesian topology functions can be implemented locally, using the topology information that is cached with the communicator. (*End of advice to implementors.*)

The corresponding [new] function for [general] graph structures is as follows.

MPI\_GRAPH\_MAP(comm, nnodes, index, edges, newrank)

IN	comm	input communicator (handle)
IN	nnodes	number of graph nodes (integer)
IN	index	integer array specifying the graph structure, see MPI_GRAPH_CREATE
IN	edges	integer array specifying the graph structure
OUT	newrank	reordered rank of the calling process; MPI_UNDEFINED if the calling process does not belong to graph (integer)

```
int MPI_Graph_map(MPI_Comm comm, int nnodes, const int [*index]index[],
                  const int [*edges]edges[], int *newrank)
```

```
MPI_Graph_map(comm, nnodes, index, edges, newrank, ierror) BIND(C)
  TYPE(MPI_Comm), INTENT(IN) :: comm
  INTEGER, INTENT(IN) :: nnodes, index(nnodes), edges(*)
  INTEGER, INTENT(OUT) :: newrank
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_GRAPH_MAP(COMM, NNODES, INDEX, EDGES, NEWRANK, IERROR)
  INTEGER COMM, NNODES, INDEX(*), EDGES(*), NEWRANK, IERROR
```

```
{int MPI::Graphcomm::Map(int nnodes, const int index[], const int edges[])
    const(binding deprecated, see Section 15.2) }
```

*Advice to implementors.* The function MPI\_GRAPH\_CREATE(comm, nnodes, index, edges, reorder, comm\_graph), with reorder = true can be implemented by calling MPI\_GRAPH\_MAP(comm, nnodes, index, edges, newrank), then calling MPI\_COMM\_SPLIT(comm, color, key, comm\_graph), with color = 0 if newrank ≠ MPI\_UNDEFINED, color = MPI\_UNDEFINED otherwise, and key = newrank.

All other graph topology functions can be implemented locally, using the topology information that is cached with the communicator. (*End of advice to implementors.*)



## 7.6 Neighborhood Collective Communication on Process Topologies

MPI process topologies specify a communication graph, but they implement no communication function themselves. Many applications require sparse nearest neighbor communications that can be expressed as graph topologies. We now describe several collective operations that perform communication along the edges of a process topology. All these functions are collective; i.e., they must be called by all processes in the specified communicator. See Section 5 on page 151 for an overview of other dense (global) collective communication operations and the semantics of collective operations.

If the graph was created with `MPI_DIST_GRAPH_CREATE_ADJACENT` with `sources` and `destinations` containing 0, ..., `n-1`, where `n` is the number of processes in the group of `comm_old` (i.e., the graph is fully connected and includes also an edge from each node to itself), then the sparse neighborhood communication routine performs the same data exchange as the corresponding dense (fully-connected) collective operation. In the case of a Cartesian communicator, only nearest neighbor communication is provided, corresponding to `rank_source` and `rank_dest` in `MPI_CART_SHIFT` with input `disp=1`.

*Rationale.* Neighborhood collective communications enable communication on a process topology. This high-level specification of data exchange among neighboring processes enables optimizations in the MPI library because the communication pattern is known statically (the topology). Thus, the implementation can compute optimized message schedules during creation of the topology [35]. This functionality can significantly simplify the implementation of neighbor exchanges [31]. (*End of rationale.*)

For a distributed graph topology, created with `MPI_DIST_GRAPH_CREATE`, the sequence of neighbors in the send and receive buffers at each process is defined as the sequence returned by `MPI_DIST_GRAPH_NEIGHBORS` for destinations and sources, respectively. For a general graph topology, created with `MPI_GRAPH_CREATE`, the order of neighbors in the send and receive buffers is defined as the sequence of neighbors as returned by `MPI_GRAPH_NEIGHBORS`. Note that general graph topologies should generally be replaced by the distributed graph topologies.

For a Cartesian topology, created with `MPI_CART_CREATE`, the sequence of neighbors in the send and receive buffers at each process is defined by order of the dimensions, first the neighbor in the negative direction and then in the positive direction with displacement 1. The numbers of sources and destinations in the communication routines are `2*ndims` with `ndims` defined in `MPI_CART_CREATE`. If a neighbor does not exist, i.e., at the border of a Cartesian topology in the case of a non-periodic virtual grid dimension (i.e., `periods[...] == false`), then this neighbor is defined to be `MPI_PROC_NULL`.

If a neighbor in any of the functions is `MPI_PROC_NULL`, then the neighborhood collective communication behaves like a point-to-point communication with `MPI_PROC_NULL` in this direction. That is, the buffer is still part of the sequence of neighbors but it is neither communicated nor updated.

### 7.6.1 Neighborhood Gather

In this function, each process  $i$  gathers data items from each process  $j$  if an edge  $(j, i)$  exists in the topology graph, and each process  $i$  sends the same data items to all processes  $j$  where an edge  $(i, j)$  exists. The send buffer is sent to each neighboring process and the  $l$ -th block in the receive buffer is received from the  $l$ -th neighbor.



```

MPI_NEIGHBOR_ALLGATHER(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype,
                        comm)
    IN      sendbuf      starting address of send buffer (choice)
    IN      sendcount    number of elements sent to each neighbor (non-negative
                        integer)
    IN      sendtype     data type of send buffer elements (handle)
    OUT     recvbuf      starting address of receive buffer (choice)
    IN      recvcount    number of elements received from each neighbor (non-
                        negative integer)
    IN      recvtype     data type of receive buffer elements (handle)
    IN      comm         communicator with topology structure (handle)

int MPI_Neighbor_allgather(const void* sendbuf, int sendcount, MPI_Datatype
                        sendtype, void* recvbuf, int recvcount, MPI_Datatype recvtype,
                        MPI_Comm comm)
MPI_Neighbor_allgather(sendbuf, sendcount, sendtype, recvbuf, recvcount,
                        recvtype, comm, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
    TYPE(*), DIMENSION(..) :: recvbuf
    INTEGER, INTENT(IN) :: sendcount, recvcount
    TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_NEIGHBOR_ALLGATHER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT,
                        RECVTYPE, COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, COMM, IERROR

    This function supports Cartesian communicators, graph communicators, and distributed
    graph communicators as described in Section 7.6 on page 328. If comm is a distributed graph
    communicator, the outcome is as if each process executed sends to each of its outgoing
    neighbors and receives from each of its incoming neighbors:

MPI_Dist_graph_neighbors_count(comm,&indegree,&outdegree,&weighted);
int *srcs=(int*)malloc(indegree*sizeof(int));
int *dsts=(int*)malloc(outdegree*sizeof(int));
MPI_Dist_graph_neighbors(comm,indegree,srcs,MPI_UNWEIGHTED,
                        outdegree,dsts,MPI_UNWEIGHTED);

int k,l;

for(k=0; k<outdegree; ++k)
    MPI_Isend(sendbuf,sendcount,sendtype,dsts[k],...);

for(l=0; l<indegree; ++l)
    MPI_Irecv(recvbuf+l*recvcount*extent(recvtype),recvcount,recvtype,

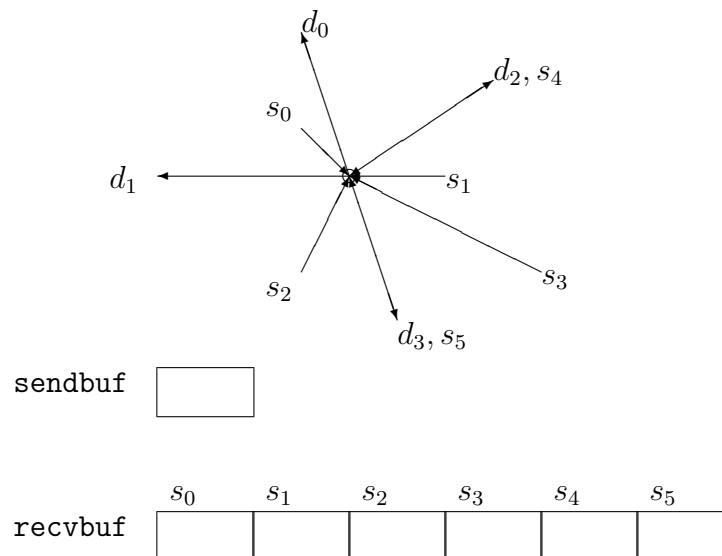
```

```

1      srcs[1],...);
2
3  MPI_Waitall(...);
4

```

Figure 7.6.1 shows the neighborhood gather communication of one process with outgoing neighbors  $d_0 \dots d_3$  and incoming neighbors  $s_0 \dots s_5$ . The process will send its `sendbuf` to all four destinations (outgoing neighbors) and it will receive the contribution from all six sources (incoming neighbors) into separate locations of its receive buffer.



All arguments are significant on all processes and the argument `comm` must have identical values on all processes.

The type signature associated with `sendcount`, `sendtype`, at a process must be equal to the type signature associated with `recvcount`, `recvtype` at all other processes. This implies that the amount of data sent must be equal to the amount of data received, pairwise between every pair of communicating processes. Distinct type maps between sender and receiver are still allowed.

*Rationale.* For optimization reasons, the same type signature is required independently of whether the topology graph is connected or not. (*End of rationale.*)

The “in place” option is not meaningful for this operation.

The vector variant of `MPI_NEIGHBOR_ALLGATHER` allows one to gather different numbers of elements from each neighbor.

```

MPI_NEIGHBOR_ALLGATHERV(sendbuf, sendcount, sendtype, recvbuf, recvcoun- 1
                        ts, displs, recvtype, comm)                                2

IN      sendbuf      starting address of send buffer (choice)                    3
IN      sendcount    number of elements sent to each neighbor (non-negative 4
                        integer)                                                5
IN      sendtype     data type of send buffer elements (handle)                6
OUT     recvbuf      starting address of receive buffer (choice)              7
IN      recvcoun-    non-negative integer array (of length indegree) con- 8
ts      ts           taining the number of elements that are received from 9
IN      displs       integer array (of length indegree). Entry i specifies 10
                        the displacement (relative to recvbuf) at which to place 11
                        the incoming data from neighbor i                      12
IN      recvtype     data type of receive buffer elements (handle)            13
IN      comm         communicator with topology structure (handle)            14
                                                    15
                                                    16
                                                    17
                                                    18
                                                    19

int MPI_Neighbor_allgatherv(const void* sendbuf, int sendcount,            20
                        MPI_Datatype sendtype, void* recvbuf, const int recvcoun- 21
                        ts[], MPI_Datatype recvtype, MPI_Comm comm)          22
                                                    23
MPI_Neighbor_allgatherv(sendbuf, sendcount, sendtype, recvbuf, recvcoun-    ticket140.
ts, displs, recvtype, comm, ierror) BIND(C)                                ticket140.
                                                    ticket140.
                                                    ticket-248T.
TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
TYPE(*), DIMENSION(..) :: recvbuf
INTEGER, INTENT(IN) :: sendcount, recvcoun-
ts(*), displs(*)
TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_NEIGHBOR_ALLGATHERV(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNTS,
                        DISPLS, RECVTYPE, COMM, IERROR)
<type> SENDBUF(*), RECVBUF(*)
INTEGER SENDCOUNT, SENDTYPE, REVCOUNTS(*), DISPLS(*), RECVTYPE, COMM,
IERROR

```

This function supports Cartesian communicators, graph communicators, and distributed graph communicators as described in Section 7.6 on page 328. If `comm` is a distributed graph communicator, the outcome is as if each process executed sends to each of its outgoing neighbors and receives from each of its incoming neighbors:

```

MPI_Dist_graph_neighbors_count(comm,&indegree,&outdegree,&weighted);
int *srcs=(int*)malloc(indegree*sizeof(int));
int *dsts=(int*)malloc(outdegree*sizeof(int));
MPI_Dist_graph_neighbors(comm,indegree,srcs,MPI_UNWEIGHTED,
                        outdegree,dsts,MPI_UNWEIGHTED);
int k,l;

```

```

1
2  for(k=0; k<outdegree; ++k)
3      MPI_Isend(sendbuf, sendcount, sendtype, dsts[k], ...);
4
5  for(l=0; l<indegree; ++l)
6      MPI_Irecv(recvbuf+displs[l]*extent(recvtype), recvcunts[l], recvtype,
7              srcs[l], ...);
8
9  MPI_Waitall(...);
10

```

The type signature associated with `sendcount`, `sendtype`, at process  $j$  must be equal to the type signature associated with `recvcunts[l]`, `recvtype` at any other process with `srcs[l]==j`. This implies that the amount of data sent must be equal to the amount of data received, pairwise between every pair of communicating processes. Distinct type maps between sender and receiver are still allowed. The data received from the  $l$ -th neighbor is placed into `recvbuf` beginning at offset `displs[l]` elements (in terms of the `recvtype`).

The “in place” option is not meaningful for this operation.

All arguments are significant on all processes and the argument `comm` must have identical values on all processes.

## 7.6.2 Neighbor Alltoall

In this function, each process  $i$  receives data items from each process  $j$  if an edge  $(j, i)$  exists in the topology graph or Cartesian topology. Similarly, each process  $i$  sends data items to all processes  $j$  where an edge  $(i, j)$  exists. This call is more general than `MPI_NEIGHBOR_ALLGATHER` in that different data items can be sent to each neighbor. The  $k$ -th block in send buffer is sent to the  $k$ -th neighboring process and the  $l$ -th block in the receive buffer is received from the  $l$ -th neighbor.

```

MPI_NEIGHBOR_ALLTOALL(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype, comm)

```

IN	<code>sendbuf</code>	starting address of send buffer (choice)
IN	<code>sendcount</code>	number of elements sent to each neighbor (non-negative integer)
IN	<code>sendtype</code>	data type of send buffer elements (handle)
OUT	<code>recvbuf</code>	starting address of receive buffer (choice)
IN	<code>recvcount</code>	number of elements received from each neighbor (non-negative integer)
IN	<code>recvtype</code>	data type of receive buffer elements (handle)
IN	<code>comm</code>	communicator with topology structure (handle)

```

int MPI_Neighbor_alltoall(const void* sendbuf, int sendcount, MPI_Datatype
    sendtype, void* recvbuf, int recvcount, MPI_Datatype recvtype,
    MPI_Comm comm)

```

```

MPI_Neighbor_alltoall(sendbuf, sendcount, sendtype, recvbuf, recvcount,
                      recvtype, comm, ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
TYPE(*), DIMENSION(..) :: recvbuf
INTEGER, INTENT(IN) :: sendcount, recvcount
TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_NEIGHBOR_ALLTOALL(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT,
                      RECVTYPE, COMM, IERROR)
<type> SENDBUF(*), RECVBUF(*)
INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, COMM, IERROR

```

This function supports Cartesian communicators, graph communicators, and distributed graph communicators as described in Section 7.6 on page 328. If `comm` is a distributed graph communicator, the outcome is as if each process executed sends to each of its outgoing neighbors and receives from each of its incoming neighbors:

```

MPI_Dist_graph_neighbors_count(comm,&indegree,&outdegree,&weighted);
int *srcs=(int*)malloc(indegree*sizeof(int));
int *dsts=(int*)malloc(outdegree*sizeof(int));
MPI_Dist_graph_neighbors(comm,indegree,srcs,MPI_UNWEIGHTED,
                        outdegree,dsts,MPI_UNWEIGHTED);

int k,l;

for(k=0; k<outdegree; ++k)
    MPI_Isend(sendbuf+k*sendcount*extent(sendtype),sendcount,sendtype,
              dsts[k],...);

for(l=0; l<indegree; ++l)
    MPI_Irecv(recvbuf+l*recvcount*extent(recvtype),recvcount,recvtype,
              srcs[l],...);

MPI_Waitall(...);

```

The type signature associated with `sendcount`, `sendtype`, at a process must be equal to the type signature associated with `recvcount`, `recvtype` at any other process. This implies that the amount of data sent must be equal to the amount of data received, pairwise between every pair of communicating processes. Distinct type maps between sender and receiver are still allowed.

The “in place” option is not meaningful for this operation.

All arguments are significant on all processes and the argument `comm` must have identical values on all processes.

The vector variant of `MPI_NEIGHBOR_ALLTOALL` allows sending/receiving different numbers of elements to and from each neighbor.

```

1  MPI_NEIGHBOR_ALLTOALLV(sendbuf, sendcounts, sdispls, sendtype, recvbuf, recvcoun-
2      rdispls, recvtype, comm)
3
4      IN      sendbuf      starting address of send buffer (choice)
5
6      IN      sendcounts   non-negative integer array (of length outdegree) speci-
7                          fying the number of elements to send to each neighbor
8
9      IN      sdispls      integer array (of length outdegree). Entry j specifies
10                      the displacement (relative to sendbuf) from which to
11                      send the outgoing data to neighbor j
12
13      IN      sendtype     data type of send buffer elements (handle)
14
15      OUT     recvbuf      starting address of receive buffer (choice)
16
17      IN      recvcoun-    non-negative integer array (of length indegree) speci-
18                      fying the number of elements that can be received
19                      from each neighbor
20
21      IN      rdispls      integer array (of length indegree). Entry i specifies
22                      the displacement (relative to recvbuf) at which to place
23                      the incoming data from neighbor i
24
25      IN      recvtype     data type of receive buffer elements (handle)
26
27      IN      comm         communicator with topology structure (handle)

```

```

ticket140. 23  int MPI_Neighbor_alltoallv(const void* sendbuf, const int sendcounts[],
ticket140. 24      const int sdispls[], MPI_Datatype sendtype, void* recvbuf,
ticket140. 25      const int recvcoun-    const int rdispls[], MPI_Datatype
ticket140. 26      recvtype, MPI_Comm comm)

```

```

ticket-248T. 27  MPI_Neighbor_alltoallv(sendbuf, sendcounts, sdispls, sendtype, recvbuf,
28      recvcoun-    recvcoun-    rdispls, recvtype, comm, ierror) BIND(C)
29
30      TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
31      TYPE(*), DIMENSION(..) :: recvbuf
32      INTEGER, INTENT(IN) :: sendcounts(*), sdispls(*), recvcoun-    (*),
33      rdispls(*)
34      TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
35      TYPE(MPI_Comm), INTENT(IN) :: comm
36      INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

37  MPI_NEIGHBOR_ALLTOALLV(SENDBUF, SENDCOUNTS, SDISPLS, SENDTYPE, RECVBUF,
38      RECVCOUNTS, RDISPLS, RECVTYPE, COMM, IERROR)
39
40      <type> SENDBUF(*), RECVBUF(*)
41      INTEGER SENDCOUNTS(*), SDISPLS(*), SENDTYPE, RECVCOUNTS(*), RDISPLS(*),
42      RECVTYPE, COMM, IERROR

```

This function supports Cartesian communicators, graph communicators, and distributed graph communicators as described in Section 7.6 on page 328. If `comm` is a distributed graph communicator, the outcome is as if each process executed sends to each of its outgoing neighbors and receives from each of its incoming neighbors:

```

47  MPI_Dist_graph_neighbors_count(comm,&indegree,&outdegree,&weighted);
48

```

```

int *srcs=(int*)malloc(indegree*sizeof(int));
int *dsts=(int*)malloc(outdegree*sizeof(int));
MPI_Dist_graph_neighbors(comm,indegree,srcs,MPI_UNWEIGHTED,
                        outdegree,dsts,MPI_UNWEIGHTED);
int k,l;

for(k=0; k<outdegree; ++k)
    MPI_Isend(sendbuf+sdispls[k]*extent(sendtype),sendcounts[k],sendtype,
              dsts[k],...);

for(l=0; l<indegree; ++l)
    MPI_Irecv(recvbuf+rdispls[l]*extent(recvtype),recvcounts[l],recvtype,
              srcs[l],...);

MPI_Waitall(...);

```

The type signature associated with `sendcounts[k]`, `sendtype` with `dsts[k]==j` at process `i` must be equal to the type signature associated with `recvcounts[l]`, `recvtype` with `srcs[l]==i` at process `j`. This implies that the amount of data sent must be equal to the amount of data received, pairwise between every pair of communicating processes. Distinct type maps between sender and receiver are still allowed. The data in the `sendbuf` beginning at offset `sdispls[k]` elements (in terms of the `sendtype`) is sent to the `k`-th outgoing neighbor. The data received from the `l`-th incoming neighbor is placed into `recvbuf` beginning at offset `rdispls[l]` elements (in terms of the `recvtype`).

The “in place” option is not meaningful for this operation.

All arguments are significant on all processes and the argument `comm` must have identical values on all processes.

`MPI_NEIGHBOR_ALLTOALLW` allows one to send and receive with different datatypes to and from each neighbor.

```

1  MPI_NEIGHBOR_ALLTOALLW(sendbuf, sendcounts, sdispls, sendtypes, recvbuf, recvcounts,
2      rdispls, recvtypes, comm)
3
4  IN      sendbuf      starting address of send buffer (choice)
5
6  IN      sendcounts    non-negative integer array (of length outdegree) speci-
7      fying the number of elements to send to each neighbor
8
9  IN      sdispls       integer array (of length outdegree). Entry j specifies
10     the displacement in bytes (relative to sendbuf) from
11     which to take the outgoing data destined for neighbor
12     j (array of integers)
13
14  IN      sendtypes     array of datatypes (of length outdegree). Entry j spec-
15     ifies the type of data to send to neighbor j (array of
16     handles)
17
18  OUT     recvbuf       starting address of receive buffer (choice)
19
20  IN      recvcounts     non-negative integer array (of length indegree) spec-
21     ifying the number of elements that can be received
22     from each neighbor
23
24  IN      rdispls       integer array (of length indegree). Entry i specifies
25     the displacement in bytes (relative to recvbuf) at which
26     to place the incoming data from neighbor i (array of
27     integers)
28
29  IN      recvtypes     array of datatypes (of length indegree). Entry i spec-
30     ifies the type of data received from neighbor i (array
31     of handles)
32
33  IN      comm          communicator with topology structure (handle)

```

```

ticket140. 29  int MPI_Neighbor_alltoallw(const void* sendbuf, const int sendcounts[],
ticket140. 30      const [int]MPI_Aint sdispls[], const MPI_Datatype sendtypes[],
ticket140. 31      void* recvbuf, const int recvcounts[], const [int]MPI_Aint
ticket299. 32      rdispls[], const MPI_Datatype recvtypes[], MPI_Comm comm)

```

```

ticket140. 33  MPI_Neighbor_alltoallw(sendbuf, sendcounts, sdispls, sendtypes, recvbuf,
ticket140. 34      recvcounts, rdispls, recvtypes, comm, ierror) BIND(C)
ticket140. 35      TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
ticket299. 36      TYPE(*), DIMENSION(..) :: recvbuf
ticket140. 37      INTEGER, INTENT(IN) :: sendcounts(*), recvcounts(*)
ticket-248T. 38      INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: sdispls(*), rdispls(*)
39      TYPE(MPI_Datatype), INTENT(IN) :: sendtypes(*), recvtypes(*)
40      TYPE(MPI_Comm), INTENT(IN) :: comm
41      INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

43  MPI_NEIGHBOR_ALLTOALLW(SENDBUF, SENDCOUNTS, SDISPLS, SENDTYPES, RECVBUF,
44      RECVCOUNTS, RDISPLS, RECVTYPES, COMM, IERROR)
45      <type> SENDBUF(*), RECVBUF(*)
46      INTEGER(KIND=MPI_ADDRESS_KIND) SDISPLS(*), RDISPLS(*)
47      INTEGER SENDCOUNTS(*), SENDTYPES(*), RECVCOUNTS(*), RECVTYPES(*), COMM,
48      IERROR

```



This function supports Cartesian communicators, graph communicators, and distributed graph communicators as described in Section 7.6 on page 328. If `comm` is a distributed graph communicator, the outcome is as if each process executed sends to each of its outgoing neighbors and receives from each of its incoming neighbors:

```
MPI_Dist_graph_neighbors_count(comm,&indegree,&outdegree,&weighted);
int *srcs=(int*)malloc(indegree*sizeof(int));
int *dsts=(int*)malloc(outdegree*sizeof(int));
MPI_Dist_graph_neighbors(comm,indegree,srcs,MPI_UNWEIGHTED,
                        outdegree,dsts,MPI_UNWEIGHTED);

int k,l;

for(k=0; k<outdegree; ++k)
    MPI_Isend(sendbuf+sdispls[k],sendcounts[k], sendtypes[k],dsts[k],...);

for(l=0; l<indegree; ++l)
    MPI_Irecv(recvbuf+rdispls[l],recvcounts[l], recvtypes[l],srcs[l],...);

MPI_Waitall(...);
```

The type signature associated with `sendcounts[k]`, `sendtypes[k]` with `dsts[k]==j` at process `i` must be equal to the type signature associated with `recvcounts[l]`, `recvtypes[l]` with `srcs[l]==i` at process `j`. This implies that the amount of data sent must be equal to the amount of data received, pairwise between every pair of communicating processes. Distinct type maps between sender and receiver are still allowed.

The “in place” option is not meaningful for this operation.

All arguments are significant on all processes and the argument `comm` must have identical values on all processes.

## 7.7 Nonblocking Neighborhood Communication on Process Topologies

Nonblocking variants of the neighborhood collective operations allow relaxed synchronization and overlapping of computation and communication. The semantics are similar to nonblocking collective operations as described in Section 5.12.

### 7.7.1 Nonblocking Neighborhood Gather

**MPI\_INEIGHBOR\_ALLGATHER**(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype, comm, request)

IN	sendbuf	starting address of send buffer (choice)
IN	sendcount	number of elements sent to each neighbor (non-negative integer)
IN	sendtype	data type of send buffer elements (handle)
OUT	recvbuf	starting address of receive buffer (choice)
IN	recvcount	number of elements received from each neighbor (non-negative integer)
IN	recvtype	data type of receive buffer elements (handle)
IN	comm	communicator with topology structure (handle)
OUT	request	communication request (handle)

```

int MPI_Ineighbor_allgather(const void* sendbuf, int sendcount,
    MPI_Datatype sendtype, void* recvbuf, int recvcount,
    MPI_Datatype recvtype, MPI_Comm comm, MPI_Request *request)
MPI_Ineighbor_allgather(sendbuf, sendcount, sendtype, recvbuf, recvcount,
    recvtype, comm, request, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN) :: sendcount, recvcount
    TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_INEIGHBOR_ALLGATHER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT,
    RECVTYPE, COMM, REQUEST, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, COMM, REQUEST, IERROR

This call starts a nonblocking variant of MPI_NEIGHBOR_ALLGATHER.

```

MPI_INEIGHBOR_ALLGATHERV(sendbuf, sendcount, sendtype, recvbuf, recvcunts, displs,			1
recvtype, comm, request)			2
IN	sendbuf	starting address of send buffer (choice)	3
IN	sendcount	number of elements sent to each neighbor (non-negative integer)	4
IN	sendtype	data type of send buffer elements (handle)	5
OUT	recvbuf	starting address of receive buffer (choice)	6
IN	recvcunts	non-negative integer array (of length indegree) containing the number of elements that are received from each neighbor	7
IN	displs	integer array (of length indegree). Entry i specifies the displacement (relative to recvbuf) at which to place the incoming data from neighbor i	8
IN	recvtype	data type of receive buffer elements (handle)	9
IN	comm	communicator with topology structure (handle)	10
OUT	request	communication request (handle)	11
			12
int MPI_Ineighbor_allgatherv(const void* sendbuf, int sendcount,			13
MPI_Datatype sendtype, void* recvbuf, const int recvcunts[],			14
const int displs[], MPI_Datatype recvtype, MPI_Comm comm,			15
MPI_Request *request)			16
			17
MPI_Ineighbor_allgatherv(sendbuf, sendcount, sendtype, recvbuf, recvcunts,			18
displs, recvtype, comm, request, ierror) BIND(C)			19
TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf			20
TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf			21
INTEGER, INTENT(IN) :: sendcount			22
INTEGER, INTENT(IN), ASYNCHRONOUS :: recvcunts(*), displs(*)			23
TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype			24
TYPE(MPI_Comm), INTENT(IN) :: comm			25
TYPE(MPI_Request), INTENT(OUT) :: request			26
INTEGER, OPTIONAL, INTENT(OUT) :: ierror			27
			28
MPI_INEIGHBOR_ALLGATHERV(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNTS,			29
DISPLS, RECVTYPE, COMM, REQUEST, IERROR)			30
<type> SENDBUF(*), RECVBUF(*)			31
INTEGER SENDCOUNT, SENDTYPE, RECVCOUNTS(*), DISPLS(*), RECVTYPE, COMM,			32
REQUEST, IERROR			33
			34
This call starts a nonblocking variant of MPI_NEIGHBOR_ALLGATHERV.			35
			36
			37
			38
			39
			40
			41
			42
			43
			44
			45
			46
			47
			48

## 7.7.2 Nonblocking Neighborhood Alltoall

```
MPI_INEIGHBOR_ALLTOALL(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype, comm,
                        request)
```

IN	sendbuf	starting address of send buffer (choice)
IN	sendcount	number of elements sent to each neighbor (non-negative integer)
IN	sendtype	data type of send buffer elements (handle)
OUT	recvbuf	starting address of receive buffer (choice)
IN	recvcount	number of elements received from each neighbor (non-negative integer)
IN	recvtype	data type of receive buffer elements (handle)
IN	comm	communicator with topology structure (handle)
OUT	request	communication request (handle)

```
int MPI_Ineighbor_alltoall(const void* sendbuf, int sendcount, MPI_Datatype
                          sendtype, void* recvbuf, int recvcount, MPI_Datatype recvtype,
                          MPI_Comm comm, MPI_Request *request)
```

```
MPI_Ineighbor_alltoall(sendbuf, sendcount, sendtype, recvbuf, recvcount,
                      recvtype, comm, request, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN) :: sendcount, recvcount
    TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_INEIGHBOR_ALLTOALL(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT,
                      RECVTYPE, COMM, REQUEST, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, COMM, REQUEST, IERROR
```

This call starts a nonblocking variant of MPI\_NEIGHBOR\_ALLTOALL.

MPI_INEIGHBOR_ALLTOALLV(sendbuf, sendcounts, sdispls, sendtype, recvbuf, recvcoun-			1
tdispls, recvtype, comm, request)			2
IN	sendbuf	starting address of send buffer (choice)	3
IN	sendcounts	non-negative integer array (of length outdegree) speci-	4
		fying the number of elements to send to each neighbor	5
IN	sdispls	integer array (of length outdegree). Entry j specifies	6
		the displacement (relative to sendbuf) from which send	7
		the outgoing data to neighbor j	8
IN	sendtype	data type of send buffer elements (handle)	9
OUT	recvbuf	starting address of receive buffer (choice)	10
IN	recvcoun-	non-negative integer array (of length indegree) spec-	11
		ifying the number of elements that can be received	12
		from each neighbor	13
IN	rdispls	integer array (of length indegree). Entry i specifies	14
		the displacement (relative to recvbuf) at which to place	15
		the incoming data from neighbor i	16
IN	recvtype	data type of receive buffer elements (handle)	17
IN	comm	communicator with topology structure (handle)	18
OUT	request	communication request (handle)	19
			20
int MPI_Ineighbor_alltoallv(const void* sendbuf, const int sendcounts[],			21
const int sdispls[], MPI_Datatype sendtype, void* recvbuf,			22
const int recvcoun-			23
tdispls[], MPI_Datatype			24
recvtype, MPI_Comm comm, MPI_Request *request)			25
			26
MPI_Ineighbor_alltoallv(sendbuf, sendcounts, sdispls, sendtype, recvbuf,			27
recvcoun-			28
tdispls, recvtype, comm, request, ierror) BIND(C)			29
TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf			30
TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf			31
INTEGER, INTENT(IN), ASYNCHRONOUS :: sendcounts(*), sdispls(*),			32
recvcoun-			33
tdispls(*)			34
TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype			35
TYPE(MPI_Comm), INTENT(IN) :: comm			36
TYPE(MPI_Request), INTENT(OUT) :: request			37
INTEGER, OPTIONAL, INTENT(OUT) :: ierror			38
			39
MPI_INEIGHBOR_ALLTOALLV(SENDBUF, SENDCOUNTS, SDISPLS, SENDTYPE, RECVBUF,			40
RECVCOUNTS, RDISPLS, RECVTYPE, COMM, REQUEST, IERROR)			41
<type> SENDBUF(*), RECVBUF(*)			42
INTEGER SENDCOUNTS(*), SDISPLS(*), SENDTYPE, RECVCOUNTS(*), RDISPLS(*),			43
RECVTYPE, COMM, REQUEST, IERROR			44
			45
This call starts a nonblocking variant of MPI_NEIGHBOR_ALLTOALLV.			46
			47
			48

```

1  MPI_INEIGHBOR_ALLTOALLW(sendbuf, sendcounts, sdispls, sendtypes, recvbuf, recvcounts,
2      rdispls, recvtypes, comm, request)
3
4  IN      sendbuf      starting address of send buffer (choice)
5
6  IN      sendcounts    non-negative integer array (of length outdegree) speci-
7      fying the number of elements to send to each neighbor
8
9  IN      sdispls       integer array (of length outdegree). Entry j specifies
10     the displacement in bytes (relative to sendbuf) from
11     which to take the outgoing data destined for neighbor
12     j (array of integers)
13
14  IN      sendtypes     array of datatypes (of length outdegree). Entry j spec-
15     ifies the type of data to send to neighbor j (array of
16     handles)
17
18  OUT     recvbuf       starting address of receive buffer (choice)
19
20  IN      recvcounts    non-negative integer array (of length indegree) spec-
21     ifying the number of elements that can are received
22     from each neighbor
23
24  IN      rdispls       integer array (of length indegree). Entry i specifies
25     the displacement in bytes (relative to recvbuf) at which
26     to place the incoming data from neighbor i (array of
27     integers)
28
29  IN      recvtypes     array of datatypes (of length indegree). Entry i spec-
30     ifies the type of data received from neighbor i (array
31     of handles)
32
33  IN      comm          communicator with topology structure (handle)
34
35  OUT     request       communication request (handle)

```

```

36  int MPI_Ineighbor_alltoallw(const void* sendbuf, const int sendcounts[],
37      const [int]MPI_Aint sdispls[], const MPI_Datatype sendtypes[],
38      void* recvbuf, const int recvcounts[], const [int]MPI_Aint
39      rdispls[], const MPI_Datatype recvtypes[], MPI_Comm comm,
40      MPI_Request *request)
41
42  MPI_Ineighbor_alltoallw(sendbuf, sendcounts, sdispls, sendtypes, recvbuf,
43      recvcounts, rdispls, recvtypes, comm, request, ierror) BIND(C)
44  TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
45  TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
46  INTEGER, INTENT(IN), ASYNCHRONOUS :: sendcounts(*), recvcounts(*)
47  INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN), ASYNCHRONOUS ::
48  sdispls(*), rdispls(*)
49  TYPE(MPI_Datatype), INTENT(IN), ASYNCHRONOUS :: sendtypes(*),
50  recvtypes(*)
51  TYPE(MPI_Comm), INTENT(IN) :: comm
52  TYPE(MPI_Request), INTENT(OUT) :: request
53  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_INEIGHBOR_ALLTOALLW(SENDBUF, SENDCOUNTS, SDISPLS, SENDTYPES, RECVBUF,
    RECVCOUNTS, RDISPLS, RECVTYPES, COMM, REQUEST, IERROR)
<type> SENDBUF(*), RECVBUF(*)
INTEGER(KIND=MPI_ADDRESS_KIND) SDISPLS(*), RDISPLS(*)
INTEGER SENDCOUNTS(*), SENDTYPES(*), RECVCOUNTS(*), RECVTYPES(*), COMM,
REQUEST, IERROR

```

This call starts a nonblocking variant of MPI\_NEIGHBOR\_ALLTOALLW.

## 7.8 An Application Example

**Example 7.9** [] The example in [Figure 7.1]Figures 7.2-7.4 shows how the grid definition and inquiry functions can be used in an application program. A partial differential equation, for instance the Poisson equation, is to be solved on a rectangular domain. First, the processes organize themselves in a two-dimensional structure. Each process then inquires about the ranks of its neighbors in the four directions (up, down, right, left). The numerical problem is solved by an iterative method, the details of which are hidden in the subroutine `relax`.

In each relaxation step each process computes new values for the solution grid function at [all]the points `u(1:100,1:100)` owned by the process. Then the values at inter-process boundaries have to be exchanged with neighboring processes. For example, the [exchange subroutine might contain a call like `MPI_SEND(...,neigh_rank(1),...)` to send updated values to the left-hand neighbor (`i-1,j`). ]newly calculated values in `u(1,1:100)` must be sent into the halo cells `u(101,1:100)` of the left-hand neighbor with coordinates `(own_coord(1)-1,own_coord(2))`

[]

```

1
2
3   integer ndims, num_neigh
4   logical reorder
5   parameter (ndims=2, num_neigh=4, reorder=.true.)
6   integer comm, comm_cart, dims(ndims), neigh_def(ndims), ierr
7   integer neigh_rank(num_neigh), own_position(ndims), i, j
8   logical periods(ndims)
9   real*8 u(0:101,0:101), f(0:101,0:101)
10  data dims / ndims * 0 /
11  comm = MPI_COMM_WORLD
12  C   Set process grid size and periodicity
13  call MPI_DIMS_CREATE(comm, ndims, dims,ierr)
14  periods(1) = .TRUE.
15  periods(2) = .TRUE.
16  C   Create a grid structure in WORLD group and inquire about own position
17  call MPI_CART_CREATE (comm, ndims, dims, periods, reorder, comm_cart,ierr)
18  call MPI_CART_GET (comm_cart, ndims, dims, periods, own_position,ierr)
19  C   Look up the ranks for the neighbors. Own process coordinates are (i,j).
20  C   Neighbors are (i-1,j), (i+1,j), (i,j-1), (i,j+1)
21  i = own_position(1)
22  j = own_position(2)
23  neigh_def(1) = i-1
24  neigh_def(2) = j
25  call MPI_CART_RANK (comm_cart, neigh_def, neigh_rank(1),ierr)
26  neigh_def(1) = i+1
27  neigh_def(2) = j
28  call MPI_CART_RANK (comm_cart, neigh_def, neigh_rank(2),ierr)
29  neigh_def(1) = i
30  neigh_def(2) = j-1
31  call MPI_CART_RANK (comm_cart, neigh_def, neigh_rank(3),ierr)
32  neigh_def(1) = i
33  neigh_def(2) = j+1
34  call MPI_CART_RANK (comm_cart, neigh_def, neigh_rank(4),ierr)
35  C   Initialize the grid functions and start the iteration
36  call init (u, f)
37  do 10 it=1,100
38      call relax (u, f)
39  C   Exchange data with neighbor processes
40      call exchange (u, comm_cart, neigh_rank, num_neigh)
41  10 continue
42  call output (u)
43  end
44
45
46
47
48

```

Figure 7.1: Set-up of process structure for two-dimensional parallel Poisson solver.



```

1
2
3
4
5
6
7
8 INTEGER ndims, num_neigh
9 LOGICAL reorder
10 PARAMETER (ndims=2, num_neigh=4, reorder=.true.)
11 INTEGER comm, comm_cart, dims(ndims), ierr
12 INTEGER neigh_rank(num_neigh), own_coords(ndims), i, j, it
13 LOGICAL periods(ndims)
14 REAL u(0:101,0:101), f(0:101,0:101)
15 DATA dims / ndims * 0 /
16 comm = MPI_COMM_WORLD
17 ! Set process grid size and periodicity
18 CALL MPI_DIMS_CREATE(comm, ndims, dims,ierr)
19 periods(1) = .TRUE.
20 periods(2) = .TRUE.
21 ! Create a grid structure in WORLD group and inquire about own position
22 CALL MPI_CART_CREATE (comm, ndims, dims, periods, reorder, &
23                      comm_cart,ierr)
24 CALL MPI_CART_GET (comm_cart, ndims, dims, periods, own_coords,ierr)
25 i = own_coords(1)
26 j = own_coords(2)
27 ! Look up the ranks for the neighbors. Own process coordinates are (i,j).
28 ! Neighbors are (i-1,j), (i+1,j), (i,j-1), (i,j+1) modulo (dims(1),dims(2))
29 CALL MPI_CART_SHIFT (comm_cart, 0,1, neigh_rank(1),neigh_rank(2), ierr)
30 CALL MPI_CART_SHIFT (comm_cart, 1,1, neigh_rank(3),neigh_rank(4), ierr)
31 ! Initialize the grid functions and start the iteration
32 CALL init (u, f)
33 DO it=1,100
34     CALL relax (u, f)
35     ! Exchange data with neighbor processes
36     CALL exchange (u, comm_cart, neigh_rank, num_neigh)
37 END DO
38 CALL output (u)
39
40
41
42
43
44
45
46
47
48

```

Figure 7.2: Set-up of process structure for two-dimensional parallel Poisson solver.

```

1
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7
8
9
10
11
12 SUBROUTINE exchange (u, comm_cart, neigh_rank, num_neigh)
13 REAL u(0:101,0:101)
14 INTEGER comm_cart, num_neigh, neigh_rank(num_neigh)
15 REAL sndbuf(100,num_neigh), rcvbuf(100,num_neigh)
16 INTEGER ierr
17 sndbuf(1:100,1) = u( 1,1:100)
18 sndbuf(1:100,2) = u(100,1:100)
19 sndbuf(1:100,3) = u(1:100, 1)
20 sndbuf(1:100,4) = u(1:100,100)
21 CALL MPI_NEIGHBOR_ALLTOALL (sndbuf, 100, MPI_REAL, rcvbuf, 100, MPI_REAL, &
22                               comm_cart, ierr)
23 ! instead of
24 ! DO i=1,num_neigh
25 !   CALL MPI_Irecv(rcvbuf(1,i),100,MPI_REAL,neigh_rank(i),...,rq(2*i-1),ierr)
26 !   CALL MPI_Isend(sndbuf(1,i),100,MPI_REAL,neigh_rank(i),...,rq(2*i ),ierr)
27 ! END DO
28 ! CALL MPI_WAITALL (2*num_neigh, rq, statuses, ierr)
29
30 u( 0,1:100) = rcvbuf(1:100,1)
31 u(101,1:100) = rcvbuf(1:100,2)
32 u(1:100, 0) = rcvbuf(1:100,3)
33 u(1:100,101) = rcvbuf(1:100,4)
34 END
35
36
37
38
39
40
41
42
43
44
45
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```

Figure 7.3: Communication routine with local data copying and sparse neighborhood all-to-all.

```

SUBROUTINE exchange (u, comm_cart, neigh_rank, num_neigh)
USE MPI
REAL u(0:101,0:101)
INTEGER comm_cart, num_neigh, neigh_rank(num_neigh)
INTEGER sndcounts(num_neigh), sdispls(num_neigh), sndtypes(num_neigh)
INTEGER rcvcounts(num_neigh), rdispls(num_neigh), rcvtypes(num_neigh)
INTEGER (KIND=MPI_ADDRESS_KIND) lb, sizeofreal
INTEGER type_vec, i, ierr
!   The following initialization need to be done only once
!   before the first call of exchange.
CALL MPI_TYPE_GET_EXTENT(MPI_REAL, lb, sizeofreal, ierr)
CALL MPI_TYPE_VECTOR (100, 1, 102, MPI_REAL, type_vec, ierr)
CALL MPI_TYPE_COMMIT (type_vec, ierr)
sndtypes(1) = type_vec
sndtypes(2) = type_vec
sndtypes(3) = MPI_REAL
sndtypes(4) = MPI_REAL
DO i=1,num_neigh
    sndcounts(i) = 100
    rcvcounts(i) = 100
    rcvtypes(i) = sndtypes(i)
END DO
sdispls(1) = ( 1 + 1*102) * sizeofreal ! first element of u( 1,1:100)
sdispls(2) = (100 + 1*102) * sizeofreal ! first element of u(100,1:100)
sdispls(3) = ( 1 + 1*102) * sizeofreal ! first element of u(1:100, 1)
sdispls(4) = ( 1 + 100*102) * sizeofreal ! first element of u(1:100,100)
rdispls(1) = ( 0 + 1*102) * sizeofreal ! first element of u( 0,1:100)
rdispls(2) = (101 + 1*102) * sizeofreal ! first element of u(101,1:100)
rdispls(3) = ( 1 + 0*102) * sizeofreal ! first element of u(1:100, 0)
rdispls(4) = ( 1 + 101*102) * sizeofreal ! first element of u(1:100,101)

! the following communication has to be done in each call of exchange
CALL MPI_NEIGHBOR_ALLTOALLW (u, sndcounts, sdispls, sndtypes, &
                             u, rcvcounts, rdispls, rcvtypes, comm_cart, ierr)

!   The following finalizing need to be done only once
!   after the last call of exchange.
CALL MPI_TYPE_FREE (type_vec, ierr)
END

```

Figure 7.4: Communication routine with sparse neighborhood all-to-all-w and without local data copying.

1  
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3  
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## Chapter 8

# MPI Environmental Management

This chapter discusses routines for getting and, where appropriate, setting various parameters that relate to the MPI implementation and the execution environment (such as error handling). The procedures for entering and leaving the MPI execution environment are also described here.

### 8.1 Implementation Information

#### 8.1.1 Version Inquiries

In order to cope with changes to the MPI Standard, there are both compile-time and run-time ways to determine which version of the standard is in use in the environment one is using.

The “version” will be represented by two separate integers, for the version and subversion: In C and C++,

```
#define MPI_VERSION    2
#define MPI_SUBVERSION 2
```

in Fortran,

```
INTEGER [ticket240-L.] :: MPI_VERSION, MPI_SUBVERSION
PARAMETER (MPI_VERSION    = 2)
PARAMETER (MPI_SUBVERSION = 2)
```

For runtime determination,

`MPI_GET_VERSION( version, subversion )`

OUT	version	version number (integer)
OUT	subversion	subversion number (integer)

```
int MPI_Get_version(int *version, int *subversion)
```

```
MPI_Get_version(version, subversion, ierror) BIND(C)
```

```
    INTEGER, INTENT(OUT) :: version, subversion
```

```
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```

1 MPI_GET_VERSION(VERSION, SUBVERSION, IERROR)
2     INTEGER VERSION, SUBVERSION, IERROR
3
4 {void MPI::Get_version(int& version, int& subversion) (binding deprecated, see
5     Section 15.2) }
```

MPI\_GET\_VERSION is one of the few functions that can be called before MPI\_INIT and after MPI\_FINALIZE. Valid (MPI\_VERSION, MPI\_SUBVERSION) pairs in this and previous versions of the MPI standard are (2,2), (2,1), (2,0), and (1,2).

### 8.1.2 Environmental Inquiries

A set of attributes that describe the execution environment are attached to the communicator MPI\_COMM\_WORLD when MPI is initialized. The value of these attributes can be inquired by using the function MPI\_COMM\_GET\_ATTR described in Chapter 6. It is erroneous to delete these attributes, free their keys, or change their values.

The list of predefined attribute keys include

**MPI\_TAG\_UB** Upper bound for tag value.

**MPI\_HOST** Host process rank, if such exists, MPI\_PROC\_NULL, otherwise.

**MPI\_IO** rank of a node that has regular I/O facilities (possibly myrank). Nodes in the same communicator may return different values for this parameter.

**MPI\_WTIME\_IS\_GLOBAL** Boolean variable that indicates whether clocks are synchronized.

Vendors may add implementation specific parameters (such as node number, real memory size, virtual memory size, etc.)

These predefined attributes do not change value between MPI initialization (MPI\_INIT and MPI completion (MPI\_FINALIZE), and cannot be updated or deleted by users.

*Advice to users.* Note that in the C binding, the value returned by these attributes is a *pointer* to an `int` containing the requested value. (*End of advice to users.*)

The required parameter values are discussed in more detail below:

#### Tag Values

Tag values range from 0 to the value returned for MPI\_TAG\_UB inclusive. These values are guaranteed to be unchanging during the execution of an MPI program. In addition, the tag upper bound value must be *at least* 32767. An MPI implementation is free to make the value of MPI\_TAG\_UB larger than this; for example, the value  $2^{30} - 1$  is also a legal value for MPI\_TAG\_UB.

The attribute MPI\_TAG\_UB has the same value on all processes of MPI\_COMM\_WORLD.

#### Host Rank

The value returned for MPI\_HOST gets the rank of the HOST process in the group associated with communicator MPI\_COMM\_WORLD, if there is such. MPI\_PROC\_NULL is returned if there is no host. MPI does not specify what it means for a process to be a HOST, nor does it requires that a HOST exists.

The attribute MPI\_HOST has the same value on all processes of MPI\_COMM\_WORLD.

## IO Rank

The value returned for `MPI_IO` is the rank of a processor that can provide language-standard I/O facilities. For Fortran, this means that all of the Fortran I/O operations are supported (e.g., `OPEN`, `REWIND`, `WRITE`). For C and C++, this means that all of the ISO C and C++, I/O operations are supported (e.g., `fopen`, `fprintf`, `lseek`).

If every process can provide language-standard I/O, then the value `MPI_ANY_SOURCE` will be returned. Otherwise, if the calling process can provide language-standard I/O, then its rank will be returned. Otherwise, if some process can provide language-standard I/O then the rank of one such process will be returned. The same value need not be returned by all processes. If no process can provide language-standard I/O, then the value `MPI_PROC_NULL` will be returned.

*Advice to users.* Note that input is not collective, and this attribute does *not* indicate which process can or does provide input. (*End of advice to users.*)

## Clock Synchronization

The value returned for `MPI_WTIME_IS_GLOBAL` is 1 if clocks at all processes in `MPI_COMM_WORLD` are synchronized, 0 otherwise. A collection of clocks is considered synchronized if explicit effort has been taken to synchronize them. The expectation is that the variation in time, as measured by calls to `MPI_WTIME`, will be less than one half the round-trip time for an MPI message of length zero. If time is measured at a process just before a send and at another process just after a matching receive, the second time should be always higher than the first one.

The attribute `MPI_WTIME_IS_GLOBAL` need not be present when the clocks are not synchronized (however, the attribute key `MPI_WTIME_IS_GLOBAL` is always valid). This attribute may be associated with communicators other than `MPI_COMM_WORLD`.

The attribute `MPI_WTIME_IS_GLOBAL` has the same value on all processes of `MPI_COMM_WORLD`.

`MPI_GET_PROCESSOR_NAME( name, resultlen )`

OUT	name	A unique specifier for the actual (as opposed to virtual) node.
OUT	resultlen	Length (in printable characters) of the result returned in name

```
int MPI_Get_processor_name(char *name, int *resultlen)
```

```
MPI_Get_processor_name(name, resultlen, ierror) BIND(C)
  CHARACTER(LEN=MPI_MAX_PROCESSOR_NAME), INTENT(OUT) :: name
  INTEGER, INTENT(OUT) :: resultlen
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_GET_PROCESSOR_NAME( NAME, RESULTLEN, IERROR)
  CHARACTER*(*) NAME
  INTEGER RESULTLEN, IERROR
```

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```

1 {void MPI::Get_processor_name(char* name, int& resultlen) (binding deprecated,
2 see Section 15.2) }
```

This routine returns the name of the processor on which it was called at the moment of the call. The name is a character string for maximum flexibility. From this value it must be possible to identify a specific piece of hardware; possible values include “processor 9 in rack 4 of mpp.cs.org” and “231” (where 231 is the actual processor number in the running homogeneous system). The argument `name` must represent storage that is at least `MPI_MAX_PROCESSOR_NAME` characters long. `MPI_GET_PROCESSOR_NAME` may write up to this many characters into `name`.

The number of characters actually written is returned in the output argument, `resultlen`. In C, a null character is additionally stored at `name[resultlen]`. The `resultlen` cannot be larger than `MPI_MAX_PROCESSOR_NAME-1`. In Fortran, `name` is padded on the right with blank characters. The `resultlen` cannot be larger than `MPI_MAX_PROCESSOR_NAME`.

*Rationale.* This function allows MPI implementations that do process migration to return the current processor. Note that nothing in MPI *requires* or defines process migration; this definition of `MPI_GET_PROCESSOR_NAME` simply allows such an implementation. (*End of rationale.*)

*Advice to users.* The user must provide at least `MPI_MAX_PROCESSOR_NAME` space to write the processor name — processor names can be this long. The user should examine the output argument, `resultlen`, to determine the actual length of the name. (*End of advice to users.*)

□ The constant `MPI_BSEND_OVERHEAD` provides an upper bound on the fixed overhead per message buffered by a call to `MPI_BSEND` (see Section 3.6.1).

## 8.2 Memory Allocation

In some systems, message-passing and remote-memory-access (RMA) operations run faster when accessing specially allocated memory (e.g., memory that is shared by the other processes in the communicating group on an SMP). MPI provides a mechanism for allocating and freeing such special memory. The use of such memory for message-passing or RMA is not mandatory, and this memory can be used without restrictions as any other dynamically allocated memory. However, implementations may restrict the use of the `MPI_WIN_LOCK` and `MPI_WIN_UNLOCK` functions to windows allocated in such memory (see Section 11.5.3.)

```

39 MPI_ALLOC_MEM(size, info, baseptr)
```

IN	size	size of memory segment in bytes (non-negative integer)
IN	info	info argument (handle)
OUT	baseptr	pointer to beginning of memory segment allocated

```

46 int MPI_Alloc_mem(MPI_Aint size, MPI_Info info, void *baseptr)
```

```

47 MPI_Alloc_mem(size, info, baseptr, ierror) BIND(C)
```



```

USE, INTRINSIC :: ISO_C_BINDING, ONLY : C_PTR
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: size
TYPE(MPI_Info), INTENT(IN) :: info
TYPE(C_PTR), INTENT(OUT) :: baseptr
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_ALLOC_MEM(SIZE, INFO, BASEPTR, IERROR)
  INTEGER INFO, IERROR
  INTEGER(KIND=MPI_ADDRESS_KIND) SIZE, BASEPTR

{void* MPI::Alloc_mem(MPI::Aint size, const MPI::Info& info) (binding
  deprecated, see Section 15.2) }

```

If the Fortran compiler provides `TYPE(C_PTR)`, then the following interface must be provided in the `mpi` module and should be provided in `mpif.h` through overloading, i.e., with the same routine name as the routine with `INTEGER(KIND=MPI_ADDRESS_KIND) BASEPTR`, but with a different linker name:

```

INTERFACE MPI_ALLOC_MEM
  SUBROUTINE MPI_ALLOC_MEM_CPTR(SIZE, INFO, BASEPTR, IERROR)
    USE, INTRINSIC :: ISO_C_BINDING, ONLY : C_PTR
    INTEGER :: INFO, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) :: SIZE
    TYPE(C_PTR) :: BASEPTR
  END SUBROUTINE
END INTERFACE

```

The linker name base of this overloaded function is `MPI_ALLOC_MEM_CPTR`. The implied linker names are described in Section 16.2.5 on page 649.

The `info` argument can be used to provide directives that control the desired location of the allocated memory. Such a directive does not affect the semantics of the call. Valid `info` values are implementation-dependent; a null directive value of `info = MPI_INFO_NULL` is always valid.

The function `MPI_ALLOC_MEM` may return an error code of class `MPI_ERR_NO_MEM` to indicate it failed because memory is exhausted.

```

MPI_FREE_MEM(base)
  IN      base
          initial address of memory segment allocated by
          MPI_ALLOC_MEM (choice)

```

```

int MPI_Free_mem(void *base)

MPI_Free_mem(base, ierror) BIND(C)
  TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: base
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_FREE_MEM(BASE, IERROR)
  <type> BASE(*)
  INTEGER IERROR

```

```
1 {void MPI::Free_mem(void *base) (binding deprecated, see Section 15.2) }
```

2 The function `MPI_FREE_MEM` may return an error code of class `MPI_ERR_BASE` to  
3 indicate an invalid base argument.  
4

5 *Rationale.* The C and C++ bindings of `MPI_ALLOC_MEM` and `MPI_FREE_MEM`  
6 are similar to the bindings for the `malloc` and `free` C library calls: a call to  
7 `MPI_Alloc_mem(..., &base)` should be paired with a call to `MPI_Free_mem(base)` (one  
8 less level of indirection). Both arguments are declared to be of same type `void*` so  
9 as to facilitate type casting. The Fortran binding is consistent with the C and C++  
10 bindings: the Fortran `MPI_ALLOC_MEM` call returns in `baseptr` the `TYPE(C_PTR)`  
11 `pointer` or the (integer valued) address of the allocated memory. The `base` argument  
12 of `MPI_FREE_MEM` is a choice argument, which passes (a reference to) the variable  
13 stored at that location. (*End of rationale.*)  
14

15 *Advice to implementors.* If `MPI_ALLOC_MEM` allocates special memory, then a  
16 design similar to the design of C `malloc` and `free` functions has to be used, in order  
17 to find out the size of a memory segment, when the segment is freed. If no special  
18 memory is used, `MPI_ALLOC_MEM` simply invokes `malloc`, and `MPI_FREE_MEM`  
19 invokes `free`.  
20

21 A call to `MPI_ALLOC_MEM` can be used in shared memory systems to allocate mem-  
22 ory in a shared memory segment. (*End of advice to implementors.*)  
23

24 **Example 8.1** Example of use of `MPI_ALLOC_MEM`, in Fortran with  
25 `TYPE(C_PTR)` pointers. We assume 4-byte REALs.  
26

```
27 USE mpi_f08 ! or USE mpi (not guaranteed with INCLUDE 'mpif.h')
28 USE, INTRINSIC :: ISO_C_BINDING
29 TYPE(C_PTR) :: p
30 REAL, DIMENSION(:,:), POINTER :: a ! no memory is allocated
31 INTEGER, DIMENSION(2) :: shape
32 INTEGER(KIND=MPI_ADDRESS_KIND) :: size
33 shape = (/100,100/)
34 size = 4 * shape(1) * shape(2) ! assuming 4 bytes per REAL
35 CALL MPI_Alloc_mem(size,MPI_INFO_NULL,p,ierr) ! memory is allocated and
36 CALL C_F_POINTER(p, a, shape) ! intrinsic ! now accessible via a(i,j)
37 ... ! in ISO_C_BINDING
38 a(3,5) = 2.71;
39 ...
40 CALL MPI_Free_mem(a, ierr) ! memory is freed
41
```

42 **Example 8.2** [Example of use of `MPI_ALLOC_MEM`, in Fortran with [pointer sup-  
43 port]non-standard *Cray-pointer*. We assume 4-byte REALs, and assume that these pointers  
44 are address-sized.  
45  
46  
47  
48

```

REAL A
POINTER (P, A(100,100)) ! no memory is allocated
[ticket245-Q.]INTEGER(KIND=MPI_ADDRESS_KIND) SIZE
[ticket245-Q.]SIZE = 4*100*100
CALL MPI_ALLOC_MEM([ticket245-Q.][4*100*100]SIZE, MPI_INFO_NULL, P, IERR)
! memory is allocated
...
A(3,5) = 2.71;
...
CALL MPI_FREE_MEM(A, IERR) ! memory is freed

```

[Since standard Fortran does not support (C-like) pointers, this] This code is not Fortran 77 or Fortran 90 code. Some compilers [(in particular, at the time of writing, g77 and Fortran compilers for Intel) do] may not support this code or need a special option, e.g., the GNU gFortran compiler needs `-fcray-pointer`.

*Advice to implementors.* Some compilers map Cray-pointers to address-sized integers, some to `TYPE(C_PTR)` pointers (e.g., Cray Fortran, version 7.3.3). From the user's viewpoint, this mapping is irrelevant because Examples 8.2 should work correctly with an MPI-3.0 (or later) library if Cray-pointers are available. (*End of advice to implementors.*)

**Example 8.3** Same example, in C

```

float (*f)[100][100] ;
/* no memory is allocated */
MPI_Alloc_mem(sizeof(float)*100*100, MPI_INFO_NULL, &f);
/* memory allocated */
...
(*f)[5][3] = 2.71;
...
MPI_Free_mem(f);

```

## 8.3 Error Handling

An MPI implementation cannot or may choose not to handle some errors that occur during MPI calls. These can include errors that generate exceptions or traps, such as floating point errors or access violations. The set of errors that are handled by MPI is implementation-dependent. Each such error generates an **MPI exception**.

The above text takes precedence over any text on error handling within this document. Specifically, text that states that errors *will* be handled should be read as *may* be handled.

A user can associate error handlers to three types of objects: communicators, windows, and files. The specified error handling routine will be used for any MPI exception that occurs during a call to MPI for the respective object. MPI calls that are not related to any objects are considered to be attached to the communicator `MPI_COMM_WORLD`. The attachment of error handlers to objects is purely local: different processes may attach different error handlers to corresponding objects.

Several predefined error handlers are available in MPI:

**MPI\_ERRORS\_ARE\_FATAL** The handler, when called, causes the program to abort on all executing processes. This has the same effect as if **MPI\_ABORT** was called by the process that invoked the handler.

**MPI\_ERRORS\_RETURN** The handler has no effect other than returning the error code to the user.

Implementations may provide additional predefined error handlers and programmers can code their own error handlers.

The error handler **MPI\_ERRORS\_ARE\_FATAL** is associated by default with **MPI\_COMM\_WORLD** after initialization. Thus, if the user chooses not to control error handling, every error that MPI handles is treated as fatal. Since (almost) all MPI calls return an error code, a user may choose to handle errors in its main code, by testing the return code of MPI calls and executing a suitable recovery code when the call was not successful. In this case, the error handler **MPI\_ERRORS\_RETURN** will be used. Usually it is more convenient and more efficient not to test for errors after each MPI call, and have such error handled by a non trivial MPI error handler.

After an error is detected, the state of MPI is undefined. That is, using a user-defined error handler, or **MPI\_ERRORS\_RETURN**, does *not* necessarily allow the user to continue to use MPI after an error is detected. The purpose of these error handlers is to allow a user to issue user-defined error messages and to take actions unrelated to MPI (such as flushing I/O buffers) before a program exits. An MPI implementation is free to allow MPI to continue after an error but is not required to do so.

*Advice to implementors.* A good quality implementation will, to the greatest possible extent, circumscribe the impact of an error, so that normal processing can continue after an error handler was invoked. The implementation documentation will provide information on the possible effect of each class of errors. (*End of advice to implementors.*)

An MPI error handler is an opaque object, which is accessed by a handle. MPI calls are provided to create new error handlers, to associate error handlers with objects, and to test which error handler is associated with an object. C and C++ have distinct typedefs for user defined error handling callback functions that accept communicator, file, and window arguments. In Fortran there are three user routines.

An error handler object is created by a call to **MPI\_XXX\_CREATE\_ERRHANDLER**(function, errhandler), where XXX is, respectively, **COMM**, **WIN**, or **FILE**.

An error handler is attached to a communicator, window, or file by a call to **MPI\_XXX\_SET\_ERRHANDLER**. The error handler must be either a predefined error handler, or an error handler that was created by a call to **MPI\_XXX\_CREATE\_ERRHANDLER**, with matching XXX. The predefined error handlers **MPI\_ERRORS\_RETURN** and **MPI\_ERRORS\_ARE\_FATAL** can be attached to communicators, windows, and files. In C++, the predefined error handler **MPI::ERRORS\_THROW\_EXCEPTIONS** can also be attached to communicators, windows, and files.

The error handler currently associated with a communicator, window, or file can be retrieved by a call to **MPI\_XXX\_GET\_ERRHANDLER**.

The MPI function **MPI\_ERRHANDLER\_FREE** can be used to free an error handler that was created by a call to **MPI\_XXX\_CREATE\_ERRHANDLER**.

**MPI\_{COMM,WIN,FILE}\_GET\_ERRHANDLER** behave as if a new error handler object is created. That is, once the error handler is no longer needed, **MPI\_ERRHANDLER\_FREE**

should be called with the error handler returned from `MPI_ERRHANDLER_GET` or `MPI_{COMM,WIN,FILE}_GET_ERRHANDLER` to mark the error handler for deallocation. This provides behavior similar to that of `MPI_COMM_GROUP` and `MPI_GROUP_FREE`.

*Advice to implementors.* High-quality implementation should raise an error when an error handler that was created by a call to `MPI_XXX_CREATE_ERRHANDLER` is attached to an object of the wrong type with a call to `MPI_YYY_SET_ERRHANDLER`. To do so, it is necessary to maintain, with each error handler, information on the typedef of the associated user function. (*End of advice to implementors.*)

The syntax for these calls is given below.

### 8.3.1 Error Handlers for Communicators

MPI_COMM_CREATE_ERRHANDLER([function]comm_errhandler_fn, errhandler)	
IN	[ticket252-W.][function]comm_errhandler_fn user defined error handling procedure (function)
OUT	errhandler MPI error handler (handle)

```

int MPI_Comm_create_errhandler(MPI_Comm_errhandler_function
                               *\[function\]comm_errhandler_fn, MPI_Errhandler *errhandler)

MPI_Comm_create_errhandler(comm_errhandler_fn, errhandler, ierror) BIND(C)
  PROCEDURE(MPI_Comm_errhandler_function) :: comm_errhandler_fn
  TYPE(MPI_Errhandler), INTENT(OUT) :: errhandler
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_COMM_CREATE_ERRHANDLER(\[FUNCTION\]COMM_ERRHANDLER_FN, ERRHANDLER,
                           IERROR)
  EXTERNAL \[FUNCTION\]COMM_ERRHANDLER_FN
  INTEGER ERRHANDLER, IERROR

{static MPI::Errhandler
    MPI::Comm::Create_errhandler(MPI::Comm::Errhandler_function*
                                  \[function\]comm_errhandler_fn)(binding deprecated, see Section 15.2) }

```

Creates an error handler that can be attached to communicators. This function is identical to `MPI_ERRHANDLER_CREATE`, whose use is deprecated.

The user routine should be, in C, a function of type `MPI_Comm_errhandler_function`, which is defined as

```
typedef void MPI_Comm_errhandler_function(MPI_Comm *, int *, ...);
```

The first argument is the communicator in use. The second is the error code to be returned by the MPI routine that raised the error. If the routine would have returned `MPI_ERR_IN_STATUS`, it is the error code returned in the status for the request that caused the error handler to be invoked. The remaining arguments are “`stdargs`” arguments whose number and meaning is implementation-dependent. An implementation should clearly document these arguments. Addresses are used so that the handler may be written in Fortran. This typedef replaces `MPI_Handler_function`, whose use is deprecated.

ticket230-B. 1 With the Fortran `mpi_f08` module, the user routine `comm_errhandler_fn` should be of the  
 ticket-248T. 2 form:

```
3 ABSTRACT INTERFACE
4   SUBROUTINE MPI_Comm_errhandler_function(comm, error_code) BIND(C)
5       TYPE(MPI_Comm) :: comm
6       INTEGER :: error_code
```

ticket230-B. 8 [In Fortran] With the Fortran `mpi` module and `mpif.h`, the user routine  
 ticket230-B. 9 `COMM_ERRHANDLER_FN` should be of the form:

```
10
11 SUBROUTINE COMM_ERRHANDLER_FUNCTION(COMM, ERROR_CODE)
12     INTEGER COMM, ERROR_CODE
```

13  
 14 In C++, the user routine should be of the form:

```
15 {typedef void MPI::Comm::Errhandler_function(MPI::Comm &, int *, ...);
16     (binding deprecated, see Section 15.2)}
```

17  
 18  
 19 *Rationale.* The variable argument list is provided because it provides an ISO-  
 20 standard hook for providing additional information to the error handler; without this  
 21 hook, ISO C prohibits additional arguments. (*End of rationale.*)

22  
 23 *Advice to users.* A newly created communicator inherits the error handler that  
 24 is associated with the “parent” communicator. In particular, the user can specify  
 25 a “global” error handler for all communicators by associating this handler with the  
 26 communicator `MPI_COMM_WORLD` immediately after initialization. (*End of advice to*  
 27 *users.*)

28  
 29  
 30 `MPI_COMM_SET_ERRHANDLER(comm, errhandler)`

31 `INOUT comm` communicator (handle)

32 `IN errhandler` new error handler for communicator (handle)

33  
 34  
 35 `int MPI_Comm_set_errhandler(MPI_Comm comm, MPI_Errhandler errhandler)`

ticket-248T. 36  
 37 `MPI_Comm_set_errhandler(comm, errhandler, ierror) BIND(C)`  
 38 `TYPE(MPI_Comm), INTENT(IN) :: comm`  
 39 `TYPE(MPI_Errhandler), INTENT(IN) :: errhandler`  
 40 `INTEGER, OPTIONAL, INTENT(OUT) :: ierror`

41 `MPI_COMM_SET_ERRHANDLER(COMM, ERRHANDLER, IERROR)`  
 42 `INTEGER COMM, ERRHANDLER, IERROR`

43  
 44 `{void MPI::Comm::Set_errhandler(const MPI::Errhandler& errhandler) (binding`  
 45 `deprecated, see Section 15.2) }`

46  
 47 Attaches a new error handler to a communicator. The error handler must be either  
 48 a predefined error handler, or an error handler created by a call to

MPI\_COMM\_CREATE\_ERRHANDLER. This call is identical to MPI\_ERRHANDLER\_SET, whose use is deprecated.

MPI\_COMM\_GET\_ERRHANDLER(comm, errhandler)

IN	comm	communicator (handle)
OUT	errhandler	error handler currently associated with communicator (handle)

```
int MPI_Comm_get_errhandler(MPI_Comm comm, MPI_Errhandler *errhandler)
```

```
MPI_Comm_get_errhandler(comm, errhandler, ierror) BIND(C)
```

```
    TYPE(MPI_Comm), INTENT(IN) :: comm
```

```
    TYPE(MPI_Errhandler), INTENT(OUT) :: errhandler
```

```
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_COMM_GET_ERRHANDLER(COMM, ERRHANDLER, IERROR)
```

```
    INTEGER COMM, ERRHANDLER, IERROR
```

```
{MPI::Errhandler MPI::Comm::Get_errhandler() const(binding deprecated, see  
    Section 15.2) }
```

Retrieves the error handler currently associated with a communicator. This call is identical to MPI\_ERRHANDLER\_GET, whose use is deprecated.

Example: A library function may register at its entry point the current error handler for a communicator, set its own private error handler for this communicator, and restore before exiting the previous error handler.

### 8.3.2 Error Handlers for Windows

MPI\_WIN\_CREATE\_ERRHANDLER([function]win\_errhandler\_fn, errhandler)

IN	[ticket252-W.][function]win_errhandler_fn	user defined error handling procedure (function)
----	-------------------------------------------	--------------------------------------------------

OUT	errhandler	MPI error handler (handle)
-----	------------	----------------------------

```
int MPI_Win_create_errhandler(MPI_Win_errhandler_function  
    *[function]win_errhandler_fn, MPI_Errhandler *errhandler)
```

```
MPI_Win_create_errhandler(win_errhandler_fn, errhandler, ierror) BIND(C)
```

```
    PROCEDURE(MPI_Win_errhandler_function) :: win_errhandler_fn
```

```
    TYPE(MPI_Errhandler), INTENT(OUT) :: errhandler
```

```
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_WIN_CREATE_ERRHANDLER( [FUNCTION]WIN_ERRHANDLER_FN, ERRHANDLER, IERROR)
```

```
    EXTERNAL [FUNCTION]WIN_ERRHANDLER_FN
```

```
    INTEGER ERRHANDLER, IERROR
```

```

1  {static MPI::Errhandler
2      MPI::Win::Create_errhandler(MPI::Win::Errhandler_function*
3      [function]win_errhandler_fn) (binding deprecated, see Section 15.2) }

```

Creates an error handler that can be attached to a window object. The user routine should be, in C, a function of type `MPI_Win_errhandler_function` which is defined as

```
typedef void MPI_Win_errhandler_function(MPI_Win *, int *, ...);
```

The first argument is the window in use, the second is the error code to be returned.

With the Fortran `mpi_f08` module, the user routine `win_errhandler_fn` should be of the form:

```

11 ABSTRACT INTERFACE
12     SUBROUTINE MPI_Win_errhandler_function(win, error_code) BIND(C)
13         TYPE(MPI_Win) :: win
14         INTEGER :: error_code

```

[In Fortran] With the Fortran `mpi` module and `mpif.h`, the user routine `WIN_ERRHANDLER_FN` should be of the form:

```

18 SUBROUTINE WIN_ERRHANDLER_FUNCTION(WIN, ERROR_CODE)
19     INTEGER WIN, ERROR_CODE

```

In C++, the user routine should be of the form:

```

23 {typedef void MPI::Win::Errhandler_function(MPI::Win &, int *, ...);
24     (binding deprecated, see Section 15.2)}

```

```
27 MPI_WIN_SET_ERRHANDLER(win, errhandler)
```

INOUT	win	window (handle)
IN	errhandler	new error handler for window (handle)

```
32 int MPI_Win_set_errhandler(MPI_Win win, MPI_Errhandler errhandler)
```

```

34 MPI_Win_set_errhandler(win, errhandler, ierror) BIND(C)
35     TYPE(MPI_Win), INTENT(IN) :: win
36     TYPE(MPI_Errhandler), INTENT(IN) :: errhandler
37     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

38 MPI_WIN_SET_ERRHANDLER(WIN, ERRHANDLER, IERROR)
39     INTEGER WIN, ERRHANDLER, IERROR

```

```

41 {void MPI::Win::Set_errhandler(const MPI::Errhandler& errhandler) (binding
42     deprecated, see Section 15.2) }

```

Attaches a new error handler to a window. The error handler must be either a pre-defined error handler, or an error handler created by a call to `MPI_WIN_CREATE_ERRHANDLER`.



MPI\_WIN\_GET\_ERRHANDLER(win, errhandler)

IN	win	window (handle)
OUT	errhandler	error handler currently associated with window (handle)

int MPI\_Win\_get\_errhandler(MPI\_Win win, MPI\_Errhandler \*errhandler)

```

MPI_Win_get_errhandler(win, errhandler, ierror) BIND(C)
  TYPE(MPI_Win), INTENT(IN) :: win
  TYPE(MPI_Errhandler), INTENT(OUT) :: errhandler
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_WIN_GET_ERRHANDLER(WIN, ERRHANDLER, IERROR)
  INTEGER WIN, ERRHANDLER, IERROR

```

```

{MPI::Errhandler MPI::Win::Get_errhandler() const(binding deprecated, see
  Section 15.2) }

```

Retrieves the error handler currently associated with a window.

### 8.3.3 Error Handlers for Files

MPI\_FILE\_CREATE\_ERRHANDLER([function]file\_errhandler\_fn, errhandler)

IN	[ticket252-W.][function]file_errhandler_fn	user defined error handling procedure (function)
OUT	errhandler	MPI error handler (handle)

```

int MPI_File_create_errhandler(MPI_File_errhandler_function
  *[function]file_errhandler_fn, MPI_Errhandler *errhandler)

```

```

MPI_File_create_errhandler(file_errhandler_fn, errhandler, ierror) BIND(C)
  PROCEDURE(MPI_File_errhandler_function) :: file_errhandler_fn
  TYPE(MPI_Errhandler), INTENT(OUT) :: errhandler
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_FILE_CREATE_ERRHANDLER([FUNCTION]FILE_ERRHANDLER_FN, ERRHANDLER,
  IERROR)
  EXTERNAL [FUNCTION]FILE_ERRHANDLER_FN
  INTEGER ERRHANDLER, IERROR

```

```

{static MPI::Errhandler
  MPI::File::Create_errhandler(MPI::File::Errhandler_function*
    [function]file_errhandler_fn)(binding deprecated, see Section 15.2) }

```

Creates an error handler that can be attached to a file object. The user routine should be, in C, a function of type MPI\_File\_errhandler\_function, which is defined as

```
typedef void MPI_File_errhandler_function(MPI_File *, int *, ...);
```

The first argument is the file in use, the second is the error code to be returned.

ticket230-B. 1 With the Fortran `mpi_f08` module, the user routine `file_errhandler_fn` should be of the form:

ticket248T. 2 **ABSTRACT INTERFACE**  
 3     **SUBROUTINE** `MPI_File_errhandler_function`(`file`, `error_code`) **BIND(C)**  
 4         **TYPE**(`MPI_File`) :: `file`  
 5         **INTEGER** :: `error_code`  
 6

ticket230-B. 7 [In Fortran] With the Fortran `mpi` module and `mpif.h`, the user routine  
 ticket230-B. 8 **FILE\_ERRHANDLER\_FN** should be of the form:  
 9

10 **SUBROUTINE** `FILE_ERRHANDLER_FUNCTION`(`FILE`, `ERROR_CODE`)  
 11     **INTEGER** `FILE`, `ERROR_CODE`  
 12

13 In C++, the user routine should be of the form:

14 {**typedef** `void MPI::File::Errhandler_function`(`MPI::File` &, `int` \*, ...);  
 15     *(binding deprecated, see Section 15.2)*}  
 16

17  
 18 **MPI\_FILE\_SET\_ERRHANDLER**(`file`, `errhandler`)  
 19

20     **INOUT**     `file`                             file (handle)  
 21     **IN**         `errhandler`                     new error handler for file (handle)  
 22

23 **int** `MPI_File_set_errhandler`(`MPI_File` `file`, `MPI_Errhandler` `errhandler`)  
 24

25 **MPI\_File\_set\_errhandler**(`file`, `errhandler`, `ierror`) **BIND(C)**  
 26     **TYPE**(`MPI_File`), **INTENT(IN)** :: `file`  
 27     **TYPE**(`MPI_Errhandler`), **INTENT(IN)** :: `errhandler`  
 28     **INTEGER**, **OPTIONAL**, **INTENT(OUT)** :: `ierror`  
 29

30 **MPI\_FILE\_SET\_ERRHANDLER**(`FILE`, `ERRHANDLER`, `IERROR`)  
 31     **INTEGER** `FILE`, `ERRHANDLER`, `IERROR`  
 32

32 {**void** `MPI::File::Set_errhandler`(`const MPI::Errhandler`& `errhandler`) *(binding*  
 33     *deprecated, see Section 15.2)* }  
 34

35     Attaches a new error handler to a file. The error handler must be either a predefined  
 36     error handler, or an error handler created by a call to `MPI_FILE_CREATE_ERRHANDLER`.  
 37

38 **MPI\_FILE\_GET\_ERRHANDLER**(`file`, `errhandler`)  
 39

40     **IN**         `file`                             file (handle)  
 41     **OUT**        `errhandler`                     error handler currently associated with file (handle)  
 42

43 **int** `MPI_File_get_errhandler`(`MPI_File` `file`, `MPI_Errhandler` \*`errhandler`)  
 44

45 **MPI\_File\_get\_errhandler**(`file`, `errhandler`, `ierror`) **BIND(C)**  
 46     **TYPE**(`MPI_File`), **INTENT(IN)** :: `file`  
 47     **TYPE**(`MPI_Errhandler`), **INTENT(OUT)** :: `errhandler`  
 48     **INTEGER**, **OPTIONAL**, **INTENT(OUT)** :: `ierror`  
 49

```
MPI_FILE_GET_ERRHANDLER(FILE, ERRHANDLER, IERROR)
```

```
INTEGER FILE, ERRHANDLER, IERROR
```

```
{MPI::Errhandler MPI::File::Get_errhandler() const(binding deprecated, see  

Section 15.2) }
```

Retrieves the error handler currently associated with a file.

#### 8.3.4 Freeing Errorhandlers and Retrieving Error Strings

```
MPI_ERRHANDLER_FREE( errhandler )
```

```
INOUT    errhandler          MPI error handler (handle)
```

```
int MPI_Errhandler_free(MPI_Errhandler *errhandler)
```

```
MPI_Errhandler_free(errhandler, ierror) BIND(C)  

  TYPE(MPI_Errhandler), INTENT(INOUT) :: errhandler  

  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_ERRHANDLER_FREE(ERRHANDLER, IERROR)
```

```
INTEGER ERRHANDLER, IERROR
```

```
{void MPI::Errhandler::Free()(binding deprecated, see Section 15.2) }
```

Marks the error handler associated with `errhandler` for deallocation and sets `errhandler` to `MPI_ERRHANDLER_NULL`. The error handler will be deallocated after all the objects associated with it (communicator, window, or file) have been deallocated.

```
MPI_ERROR_STRING( errorcode, string, resultlen )
```

```
IN          errorcode          Error code returned by an MPI routine  

OUT         string              Text that corresponds to the errorcode  

OUT         resultlen           Length (in printable characters) of the result returned  

                                in string
```

```
int MPI_Error_string(int errorcode, char *string, int *resultlen)
```

```
MPI_Error_string(errorcode, string, resultlen, ierror) BIND(C)  

  INTEGER, INTENT(IN) :: errorcode  

  CHARACTER(LEN=MPI_MAX_ERROR_STRING), INTENT(OUT) :: string  

  INTEGER, INTENT(OUT) :: resultlen  

  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_ERROR_STRING(ERRORCODE, STRING, RESULTLEN, IERROR)
```

```
INTEGER ERRORCODE, RESULTLEN, IERROR
```

```
CHARACTER*(*) STRING
```

```
{void MPI::Get_error_string(int errorcode, char* name,  

  int& resultlen)(binding deprecated, see Section 15.2) }
```

Returns the error string associated with an error code or class. The argument `string` must represent storage that is at least `MPI_MAX_ERROR_STRING` characters long. The number of characters actually written is returned in the output argument, `resultlen`.

*Rationale.* The form of this function was chosen to make the Fortran and C bindings similar. A version that returns a pointer to a string has two difficulties. First, the return string must be statically allocated and different for each error message (allowing the pointers returned by successive calls to `MPI_ERROR_STRING` to point to the correct message). Second, in Fortran, a function declared as returning `CHARACTER*(*)` can not be referenced in, for example, a `PRINT` statement. (*End of rationale.*)

## 8.4 Error Codes and Classes

The error codes returned by MPI are left entirely to the implementation (with the exception of `MPI_SUCCESS`). This is done to allow an implementation to provide as much information as possible in the error code (for use with `MPI_ERROR_STRING`).

To make it possible for an application to interpret an error code, the routine `MPI_ERROR_CLASS` converts any error code into one of a small set of standard error codes, called *error classes*. Valid error classes are shown in Table 8.1 and Table 8.2.

The error classes are a subset of the error codes: an MPI function may return an error class number; and the function `MPI_ERROR_STRING` can be used to compute the error string associated with an error class. An MPI error class is a valid MPI error code. Specifically, the values defined for MPI error classes are valid MPI error codes.

The error codes satisfy,

$$0 = \text{MPI\_SUCCESS} < \text{MPI\_ERR\_...} \leq \text{MPI\_ERR\_LASTCODE}.$$

*Rationale.* The difference between `MPI_ERR_UNKNOWN` and `MPI_ERR_OTHER` is that `MPI_ERROR_STRING` can return useful information about `MPI_ERR_OTHER`.

Note that `MPI_SUCCESS = 0` is necessary to be consistent with C practice; the separation of error classes and error codes allows us to define the error classes this way. Having a known `LASTCODE` is often a nice sanity check as well. (*End of rationale.*)

`MPI_ERROR_CLASS( errorcode, errorclass )`

IN	errorcode	Error code returned by an MPI routine
OUT	errorclass	Error class associated with errorcode

`int MPI_Error_class(int errorcode, int *errorclass)`

```
MPI_Error_class(errorcode, errorclass, ierror) BIND(C)
  INTEGER, INTENT(IN) :: errorcode
  INTEGER, INTENT(OUT) :: errorclass
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_ERROR_CLASS(ERRORCODE, ERRORCLASS, IERROR)
  INTEGER ERRORCODE, ERRORCLASS, IERROR
```

MPI_SUCCESS	No error
MPI_ERR_BUFFER	Invalid buffer pointer
MPI_ERR_COUNT	Invalid count argument
MPI_ERR_TYPE	Invalid datatype argument
MPI_ERR_TAG	Invalid tag argument
MPI_ERR_COMM	Invalid communicator
MPI_ERR_RANK	Invalid rank
MPI_ERR_REQUEST	Invalid request (handle)
MPI_ERR_ROOT	Invalid root
MPI_ERR_GROUP	Invalid group
MPI_ERR_OP	Invalid operation
MPI_ERR_TOPOLOGY	Invalid topology
MPI_ERR_DIMS	Invalid dimension argument
MPI_ERR_ARG	Invalid argument of some other kind
MPI_ERR_UNKNOWN	Unknown error
MPI_ERR_TRUNCATE	Message truncated on receive
MPI_ERR_OTHER	Known error not in this list
MPI_ERR_INTERN	Internal MPI (implementation) error
MPI_ERR_IN_STATUS	Error code is in status
MPI_ERR_PENDING	Pending request
MPI_ERR_KEYVAL	Invalid keyval has been passed
MPI_ERR_NO_MEM	MPI_ALLOC_MEM failed because memory is exhausted
MPI_ERR_BASE	Invalid base passed to MPI_FREE_MEM
MPI_ERR_INFO_KEY	Key longer than MPI_MAX_INFO_KEY
MPI_ERR_INFO_VALUE	Value longer than MPI_MAX_INFO_VAL
MPI_ERR_INFO_NOKEY	Invalid key passed to MPI_INFO_DELETE
MPI_ERR_SPAWN	Error in spawning processes
MPI_ERR_PORT	Invalid port name passed to MPI_COMM_CONNECT
MPI_ERR_SERVICE	Invalid service name passed to MPI_UNPUBLISH_NAME
MPI_ERR_NAME	Invalid service name passed to MPI_LOOKUP_NAME
MPI_ERR_WIN	Invalid win argument
MPI_ERR_SIZE	Invalid size argument
MPI_ERR_DISP	Invalid disp argument
MPI_ERR_INFO	Invalid info argument
MPI_ERR_LOCKTYPE	Invalid locktype argument
MPI_ERR_ASSERT	Invalid assert argument
MPI_ERR_RMA_CONFLICT	Conflicting accesses to window
MPI_ERR_RMA_SYNC	Wrong synchronization of RMA calls

Table 8.1: Error classes (Part 1)

1	MPI_ERR_FILE	Invalid file handle
2	MPI_ERR_NOT_SAME	Collective argument not identical on all
3		processes, or collective routines called in
4		a different order by different processes
5	MPI_ERR_AMODE	Error related to the <code>amode</code> passed to
6		<code>MPI_FILE_OPEN</code>
7	MPI_ERR_UNSUPPORTED_DATAREP	Unsupported <code>datarep</code> passed to
8		<code>MPI_FILE_SET_VIEW</code>
9	MPI_ERR_UNSUPPORTED_OPERATION	Unsupported operation, such as seeking on
10		a file which supports sequential access only
11	MPI_ERR_NO_SUCH_FILE	File does not exist
12	MPI_ERR_FILE_EXISTS	File exists
13	MPI_ERR_BAD_FILE	Invalid file name (e.g., path name too long)
14	MPI_ERR_ACCESS	Permission denied
15	MPI_ERR_NO_SPACE	Not enough space
16	MPI_ERR_QUOTA	Quota exceeded
17	MPI_ERR_READ_ONLY	Read-only file or file system
18	MPI_ERR_FILE_IN_USE	File operation could not be completed, as
19		the file is currently open by some process
20	MPI_ERR_DUP_DATAREP	Conversion functions could not be regis-
21		tered because a data representation identi-
22		fier that was already defined was passed to
23		<code>MPI_REGISTER_DATAREP</code>
24	MPI_ERR_CONVERSION	An error occurred in a user supplied data
25		conversion function.
26	MPI_ERR_IO	Other I/O error
27	MPI_ERR_LASTCODE	Last error code

Table 8.2: Error classes (Part 2)

```
{int MPI::Get_error_class(int errorcode) (binding deprecated, see Section 15.2) }
```

The function `MPI_ERROR_CLASS` maps each standard error code (error class) onto itself.

## 8.5 Error Classes, Error Codes, and Error Handlers

Users may want to write a layered library on top of an existing MPI implementation, and this library may have its own set of error codes and classes. An example of such a library is an I/O library based on MPI, see Chapter 13 on page 509. For this purpose, functions are needed to:

1. add a new error class to the ones an MPI implementation already knows.
2. associate error codes with this error class, so that `MPI_ERROR_CLASS` works.
3. associate strings with these error codes, so that `MPI_ERROR_STRING` works.
4. invoke the error handler associated with a communicator, window, or object.

Several functions are provided to do this. They are all local. No functions are provided to free error classes or codes: it is not expected that an application will generate them in significant numbers.

`MPI_ADD_ERROR_CLASS(errorclass)`

OUT      errorclass      value for the new error class (integer)

`int MPI_Add_error_class(int *errorclass)`

`MPI_Add_error_class(errorclass, ierror) BIND(C)`

`INTEGER, INTENT(OUT) :: errorclass`

`INTEGER, OPTIONAL, INTENT(OUT) :: ierror`

`MPI_ADD_ERROR_CLASS(ERRORCLASS, IERROR)`

`INTEGER ERRORCLASS, IERROR`

`{int MPI::Add_error_class() (binding deprecated, see Section 15.2) }`

Creates a new error class and returns the value for it.

*Rationale.* To avoid conflicts with existing error codes and classes, the value is set by the implementation and not by the user. (*End of rationale.*)

*Advice to implementors.* A high-quality implementation will return the value for a new `errorclass` in the same deterministic way on all processes. (*End of advice to implementors.*)

*Advice to users.* Since a call to `MPI_ADD_ERROR_CLASS` is local, the same `errorclass` may not be returned on all processes that make this call. Thus, it is not safe to assume that registering a new error on a set of processes at the same time will yield the same `errorclass` on all of the processes. However, if an implementation returns the new `errorclass` in a deterministic way, and they are always generated in the same order on the same set of processes (for example, all processes), then the value will be the same. However, even if a deterministic algorithm is used, the value can vary across processes. This can happen, for example, if different but overlapping groups of processes make a series of calls. As a result of these issues, getting the “same” error on multiple processes may not cause the same value of error code to be generated. (*End of advice to users.*)

The value of `MPI_ERR_LASTCODE` is a constant value and is not affected by new user-defined error codes and classes. Instead, a predefined attribute key `MPI_LASTUSED` is associated with `MPI_COMM_WORLD`. The attribute value corresponding to this key is the current maximum error class including the user-defined ones. This is a local value and may be different on different processes. The value returned by this key is always greater than or equal to `MPI_ERR_LASTCODE`.

*Advice to users.* The value returned by the key `MPI_LASTUSED` will not change unless the user calls a function to explicitly add an error class/code. In a multi-threaded environment, the user must take extra care in assuming this value has not

changed. Note that error codes and error classes are not necessarily dense. A user may not assume that each error class below MPI\_LASTUSEDCLASS is valid. (*End of advice to users.*)

MPI\_ADD\_ERROR\_CODE(errorclass, errorcode)

IN	errorclass	error class (integer)
OUT	errorcode	new error code to associated with errorclass (integer)

int MPI\_Add\_error\_code(int errorclass, int \*errorcode)

```
MPI_Add_error_code(errorclass, errorcode, ierror) BIND(C)
  INTEGER, INTENT(IN) :: errorclass
  INTEGER, INTENT(OUT) :: errorcode
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_ADD_ERROR_CODE(ERRORCLASS, ERRORCODE, IERROR)
  INTEGER ERRORCLASS, ERRORCODE, IERROR
```

{int MPI::Add\_error\_code(int errorclass) (*binding deprecated, see Section 15.2*) }

Creates new error code associated with errorclass and returns its value in errorcode.

*Rationale.* To avoid conflicts with existing error codes and classes, the value of the new error code is set by the implementation and not by the user. (*End of rationale.*)

*Advice to implementors.* A high-quality implementation will return the value for a new errorcode in the same deterministic way on all processes. (*End of advice to implementors.*)

MPI\_ADD\_ERROR\_STRING(errorcode, string)

IN	errorcode	error code or class (integer)
IN	string	text corresponding to errorcode (string)

int MPI\_Add\_error\_string(int errorcode, char \*string)

```
MPI_Add_error_string(errorcode, string, ierror) BIND(C)
  INTEGER, INTENT(IN) :: errorcode
  CHARACTER(LEN=*), INTENT(IN) :: string
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_ADD_ERROR_STRING(ERRORCODE, STRING, IERROR)
  INTEGER ERRORCODE, IERROR
  CHARACTER*(*) STRING
```

{void MPI::Add\_error\_string(int errorcode, const char\* string) (*binding deprecated, see Section 15.2*) }



Associates an error string with an error code or class. The string must be no more than `MPI_MAX_ERROR_STRING` characters long. The length of the string is as defined in the calling language. The length of the string does not include the null terminator in C or C++. Trailing blanks will be stripped in Fortran. Calling `MPI_ADD_ERROR_STRING` for an errorcode that already has a string will replace the old string with the new string. It is erroneous to call `MPI_ADD_ERROR_STRING` for an error code or class with a value  $\leq$  `MPI_ERR_LASTCODE`.

If `MPI_ERROR_STRING` is called when no string has been set, it will return a empty string (all spaces in Fortran, "" in C and C++).

Section 8.3 on page 355 describes the methods for creating and associating error handlers with communicators, files, and windows.

`MPI_COMM_CALL_ERRHANDLER (comm, errorcode)`

IN	comm	communicator with error handler (handle)
IN	errorcode	error code (integer)

`int MPI_Comm_call_errhandler(MPI_Comm comm, int errorcode)`

`MPI_Comm_call_errhandler(comm, errorcode, ierror) BIND(C)`

`TYPE(MPI_Comm), INTENT(IN) :: comm`

`INTEGER, INTENT(IN) :: errorcode`

`INTEGER, OPTIONAL, INTENT(OUT) :: ierror`

`MPI_COMM_CALL_ERRHANDLER(COMM, ERRORCODE, IERROR)`

`INTEGER COMM, ERRORCODE, IERROR`

`{void MPI::Comm::Call_errhandler(int errorcode) const(binding deprecated, see Section 15.2) }`

This function invokes the error handler assigned to the communicator with the error code supplied. This function returns `MPI_SUCCESS` in C and C++ and the same value in `IERROR` if the error handler was successfully called (assuming the process is not aborted and the error handler returns).

*Advice to users.* Users should note that the default error handler is `MPI_ERRORS_ARE_FATAL`. Thus, calling `MPI_COMM_CALL_ERRHANDLER` will abort the `comm` processes if the default error handler has not been changed for this communicator or on the parent before the communicator was created. (*End of advice to users.*)

`MPI_WIN_CALL_ERRHANDLER (win, errorcode)`

IN	win	window with error handler (handle)
IN	errorcode	error code (integer)

`int MPI_Win_call_errhandler(MPI_Win win, int errorcode)`

```

1  MPI_Win_call_errhandler(win, errorcode, ierror) BIND(C)
2      TYPE(MPI_Win), INTENT(IN) :: win
3      INTEGER, INTENT(IN) :: errorcode
4      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
5
6  MPI_WIN_CALL_ERRHANDLER(WIN, ERRORCODE, IERROR)
7      INTEGER WIN, ERRORCODE, IERROR
8
9  {void MPI::Win::Call_errhandler(int errorcode) const(binding deprecated, see
10                                     Section 15.2) }

```

This function invokes the error handler assigned to the window with the error code supplied. This function returns MPI\_SUCCESS in C and C++ and the same value in IERROR if the error handler was successfully called (assuming the process is not aborted and the error handler returns).

*Advice to users.* As with communicators, the default error handler for windows is MPI\_ERRORS\_ARE\_FATAL. (*End of advice to users.*)

```

19  MPI_FILE_CALL_ERRHANDLER(fh, errorcode)
20
21      IN          fh                      file with error handler (handle)
22      IN          errorcode                error code (integer)

```

```

24  int MPI_File_call_errhandler(MPI_File fh, int errorcode)
25
26  MPI_File_call_errhandler(fh, errorcode, ierror) BIND(C)
27      TYPE(MPI_File), INTENT(IN) :: fh
28      INTEGER, INTENT(IN) :: errorcode
29      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
30
31  MPI_FILE_CALL_ERRHANDLER(FH, ERRORCODE, IERROR)
32      INTEGER FH, ERRORCODE, IERROR
33
34  {void MPI::File::Call_errhandler(int errorcode) const(binding deprecated, see
35                                     Section 15.2) }

```

This function invokes the error handler assigned to the file with the error code supplied. This function returns MPI\_SUCCESS in C and C++ and the same value in IERROR if the error handler was successfully called (assuming the process is not aborted and the error handler returns).

*Advice to users.* Unlike errors on communicators and windows, the default behavior for files is to have MPI\_ERRORS\_RETURN. (*End of advice to users.*)

*Advice to users.* Users are warned that handlers should not be called recursively with MPI\_COMM\_CALL\_ERRHANDLER, MPI\_FILE\_CALL\_ERRHANDLER, or MPI\_WIN\_CALL\_ERRHANDLER. Doing this can create a situation where an infinite recursion is created. This can occur if MPI\_COMM\_CALL\_ERRHANDLER, MPI\_FILE\_CALL\_ERRHANDLER, or MPI\_WIN\_CALL\_ERRHANDLER is called inside an error handler.

Error codes and classes are associated with a process. As a result, they may be used in any error handler. Error handlers should be prepared to deal with any error code they are given. Furthermore, it is good practice to only call an error handler with the appropriate error codes. For example, file errors would normally be sent to the file error handler. (*End of advice to users.*)

## 8.6 Timers and Synchronization

MPI defines a timer. A timer is specified even though it is not “message-passing,” because timing parallel programs is important in “performance debugging” and because existing timers (both in POSIX 1003.1-1988 and 1003.4D 14.1 and in Fortran 90) are either inconvenient or do not provide adequate access to high-resolution timers. See also Section 2.6.5 on page 23.

`MPI_WTIME()`

`double MPI_Wtime(void)`

`DOUBLE PRECISION MPI_Wtime() BIND(C)`

`DOUBLE PRECISION MPI_WTIME()`

`{double MPI::Wtime() (binding deprecated, see Section 15.2) }`

`MPI_WTIME` returns a floating-point number of seconds, representing elapsed wall-clock time since some time in the past.

The “time in the past” is guaranteed not to change during the life of the process. The user is responsible for converting large numbers of seconds to other units if they are preferred.

This function is portable (it returns seconds, not “ticks”), it allows high-resolution, and carries no unnecessary baggage. One would use it like this:

```
{
    double starttime, endtime;
    starttime = MPI_Wtime();
    .... stuff to be timed ...
    endtime   = MPI_Wtime();
    printf("That took %f seconds\n",endtime-starttime);
}
```

The times returned are local to the node that called them. There is no requirement that different nodes return “the same time.” (But see also the discussion of `MPI_WTIME_IS_GLOBAL`).

`MPI_WTICK()`

`double MPI_Wtick(void)`

`DOUBLE PRECISION MPI_Wtick() BIND(C)`

```

1  DOUBLE PRECISION MPI_WTICK()
2
3  {double MPI::Wtick() (binding deprecated, see Section 15.2) }

```

MPI\_WTICK returns the resolution of MPI\_WTIME in seconds. That is, it returns, as a double precision value, the number of seconds between successive clock ticks. For example, if the clock is implemented by the hardware as a counter that is incremented every millisecond, the value returned by MPI\_WTICK should be  $10^{-3}$ .

## 8.7 Startup

One goal of MPI is to achieve *source code portability*. By this we mean that a program written using MPI and complying with the relevant language standards is portable as written, and must not require any source code changes when moved from one system to another. This explicitly does *not* say anything about how an MPI program is started or launched from the command line, nor what the user must do to set up the environment in which an MPI program will run. However, an implementation may require some setup to be performed before other MPI routines may be called. To provide for this, MPI includes an initialization routine MPI\_INIT.

```

21  MPI_INIT()

```

```

23  int MPI_Init(int *argc, char ***argv)

```

```

24  MPI_Init(ierr) BIND(C)
25      INTEGER, OPTIONAL, INTENT(OUT) ::  ierr

```

```

27  MPI_INIT(IERROR)

```

```

28      INTEGER IERROR

```

```

30  {void MPI::Init(int& argc, char**& argv) (binding deprecated, see Section 15.2) }

```

```

31  {void MPI::Init() (binding deprecated, see Section 15.2) }

```

All MPI programs must contain exactly one call to an MPI initialization routine: MPI\_INIT or MPI\_INIT\_THREAD. Subsequent calls to any initialization routines are erroneous. The only MPI functions that may be invoked before the MPI initialization routines are called are MPI\_GET\_VERSION, MPI\_INITIALIZED, and MPI\_FINALIZED. The version for ISO C accepts the argc and argv that are provided by the arguments to main or NULL:

```

38  int main(int argc, char **argv)
39  {
40      MPI_Init(&argc, &argv);
41
42      /* parse arguments */
43      /* main program      */
44
45      MPI_Finalize();      /* see below */
46  }

```

The Fortran version takes only IERROR.

Conforming implementations of MPI are required to allow applications to pass `NULL` for both the `argc` and `argv` arguments of `main` in C and C++. In C++, there is an alternative binding for `MPI::Init` that does not have these arguments at all.

*Rationale.* In some applications, libraries may be making the call to `MPI_Init`, and may not have access to `argc` and `argv` from `main`. It is anticipated that applications requiring special information about the environment or information supplied by `mpiexec` can get that information from environment variables. (*End of rationale.*)

## `MPI_FINALIZE()`

```
int MPI_Finalize(void)
```

```
MPI_Finalize(ierr) BIND(C)
    INTEGER, OPTIONAL, INTENT(OUT) :: ierr
```

```
MPI_FINALIZE(IERROR)
    INTEGER IERROR
```

```
{void MPI::Finalize() (binding deprecated, see Section 15.2) }
```

This routine cleans up all MPI state. Each process must call `MPI_FINALIZE` before it exits. Unless there has been a call to `MPI_ABORT`, each process must ensure that all pending nonblocking communications are (locally) complete before calling `MPI_FINALIZE`. Further, at the instant at which the last process calls `MPI_FINALIZE`, all pending sends must be matched by a receive, and all pending receives must be matched by a send.

For example, the following program is correct:

Process 0	Process 1
-----	-----
<code>MPI_Init();</code>	<code>MPI_Init();</code>
<code>MPI_Send(dest=1);</code>	<code>MPI_Recv(src=0);</code>
<code>MPI_Finalize();</code>	<code>MPI_Finalize();</code>

Without the matching receive, the program is erroneous:

Process 0	Process 1
-----	-----
<code>MPI_Init();</code>	<code>MPI_Init();</code>
<code>MPI_Send (dest=1);</code>	
<code>MPI_Finalize();</code>	<code>MPI_Finalize();</code>

A successful return from a blocking communication operation or from `MPI_WAIT` or `MPI_TEST` tells the user that the buffer can be reused and means that the communication is completed by the user, but does not guarantee that the local process has no more work to do. A successful return from `MPI_REQUEST_FREE` with a request handle generated by an `MPI_ISEND` nullifies the handle but provides no assurance of operation completion. The `MPI_ISEND` is complete only when it is known by some means that a matching receive has

completed. `MPI_FINALIZE` guarantees that all local actions required by communications the user has completed will, in fact, occur before it returns.

`MPI_FINALIZE` guarantees nothing about pending communications that have not been completed (completion is assured only by `MPI_WAIT`, `MPI_TEST`, or `MPI_REQUEST_FREE` combined with some other verification of completion).

**Example 8.4** This program is correct:

```

rank 0                                rank 1
=====
...
MPI_Isend();                          MPI_Recv();
MPI_Request_free();                   MPI_Barrier();
MPI_Barrier();                       MPI_Finalize();
MPI_Finalize();                      exit();
exit();

```

**Example 8.5** This program is erroneous and its behavior is undefined:

```

rank 0                                rank 1
=====
...
MPI_Isend();                          MPI_Recv();
MPI_Request_free();                   MPI_Finalize();
MPI_Finalize();                      exit();
exit();

```

If no `MPI_BUFFER_DETACH` occurs between an `MPI_BSEND` (or other buffered send) and `MPI_FINALIZE`, the `MPI_FINALIZE` implicitly supplies the `MPI_BUFFER_DETACH`.

**Example 8.6** This program is correct, and after the `MPI_Finalize`, it is as if the buffer had been detached.

```

rank 0                                rank 1
=====
...
buffer = malloc(1000000);             MPI_Recv();
MPI_Buffer_attach();                  MPI_Finalize();
MPI_Bsend();                          exit();
MPI_Finalize();
free(buffer);
exit();

```

**Example 8.7** In this example, `MPI_lprobe()` must return a `FALSE` flag. `MPI_Test_cancelled()` must return a `TRUE` flag, independent of the relative order of execution of `MPI_Cancel()` in process 0 and `MPI_Finalize()` in process 1.

The `MPI_lprobe()` call is there to make sure the implementation knows that the “tag1” message exists at the destination, without being able to claim that the user knows about it.

```

rank 0                                     rank 1
=====
MPI_Init();                               MPI_Init();
MPI_Isend(tag1);                           MPI_Barrier();
MPI_Barrier();                             MPI_Iprobe(tag2);
                                           MPI_Barrier();
MPI_Barrier();                             MPI_Finalize();
                                           exit();

MPI_Cancel();
MPI_Wait();
MPI_Test_cancelled();
MPI_Finalize();
exit();

```

*Advice to implementors.* An implementation may need to delay the return from MPI\_FINALIZE until all potential future message cancellations have been processed. One possible solution is to place a barrier inside MPI\_FINALIZE (*End of advice to implementors.*)

Once MPI\_FINALIZE returns, no MPI routine (not even MPI\_INIT) may be called, except for MPI\_GET\_VERSION, MPI\_INITIALIZED, and MPI\_FINALIZED. Each process must complete any pending communication it initiated before it calls MPI\_FINALIZE. If the call returns, each process may continue local computations, or exit, without participating in further MPI communication with other processes. MPI\_FINALIZE is collective over all connected processes. If no processes were spawned, accepted or connected then this means over MPI\_COMM\_WORLD; otherwise it is collective over the union of all processes that have been and continue to be connected, as explained in Section 10.5.4 on page 414.

*Advice to implementors.* Even though a process has completed all the communication it initiated, such communication may not yet be completed from the viewpoint of the underlying MPI system. E.g., a blocking send may have completed, even though the data is still buffered at the sender. The MPI implementation must ensure that a process has completed any involvement in MPI communication before MPI\_FINALIZE returns. Thus, if a process exits after the call to MPI\_FINALIZE, this will not cause an ongoing communication to fail. (*End of advice to implementors.*)

Although it is not required that all processes return from MPI\_FINALIZE, it is required that at least process 0 in MPI\_COMM\_WORLD return, so that users can know that the MPI portion of the computation is over. In addition, in a POSIX environment, they may desire to supply an exit code for each process that returns from MPI\_FINALIZE.

**Example 8.8** The following illustrates the use of requiring that at least one process return and that it be known that process 0 is one of the processes that return. One wants code like the following to work no matter how many processes return.

```

1      ...
2      MPI_Comm_rank(MPI_COMM_WORLD, &myrank);
3      ...
4      MPI_Finalize();
5      if (myrank == 0) {
6          resultfile = fopen("outfile","w");
7          dump_results(resultfile);
8          fclose(resultfile);
9      }
10     exit(0);

```

```

13 MPI_INITIALIZED( flag )

```

```

15     OUT      flag                      Flag is true if MPI_INIT has been called and false
16                                           otherwise.

```

```

18 int MPI_Initialized(int *flag)

```

```

20 MPI_Initialized(flag, ierror) BIND(C)
21     LOGICAL, INTENT(OUT) :: flag
22     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

23 MPI_INITIALIZED(FLAG, IERROR)
24     LOGICAL FLAG
25     INTEGER IERROR

```

```

27 {bool MPI::Is_initialized() (binding deprecated, see Section 15.2) }

```

This routine may be used to determine whether MPI\_INIT has been called. MPI\_INITIALIZED returns true if the calling process has called MPI\_INIT. Whether MPI\_FINALIZE has been called does not affect the behavior of MPI\_INITIALIZED. It is one of the few routines that may be called before MPI\_INIT is called.

```

34 MPI_ABORT( comm, errorcode )

```

```

35     IN      comm                      communicator of tasks to abort
36
37     IN      errorcode                 error code to return to invoking environment

```

```

39 int MPI_Abort(MPI_Comm comm, int errorcode)

```

```

41 MPI_Abort(comm, errorcode, ierror) BIND(C)
42     TYPE(MPI_Comm), INTENT(IN) :: comm
43     INTEGER, INTENT(IN) :: errorcode
44     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

45 MPI_ABORT(COMM, ERRORCODE, IERROR)
46     INTEGER COMM, ERRORCODE, IERROR

```

```

47 {void MPI::Comm::Abort(int errorcode) (binding deprecated, see Section 15.2) }

```



This routine makes a “best attempt” to abort all tasks in the group of `comm`. This function does not require that the invoking environment take any action with the error code. However, a Unix or POSIX environment should handle this as a **return errorcode** from the main program.

It may not be possible for an MPI implementation to abort only the processes represented by `comm` if this is a subset of the processes. In this case, the MPI implementation should attempt to abort all the connected processes but should not abort any unconnected processes. If no processes were spawned, accepted or connected then this has the effect of aborting all the processes associated with `MPI_COMM_WORLD`.

*Rationale.* The communicator argument is provided to allow for future extensions of MPI to environments with, for example, dynamic process management. In particular, it allows but does not require an MPI implementation to abort a subset of `MPI_COMM_WORLD`. (*End of rationale.*)

*Advice to users.* Whether the errorcode is returned from the executable or from the MPI process startup mechanism (e.g., `mpiexec`), is an aspect of quality of the MPI library but not mandatory. (*End of advice to users.*)

*Advice to implementors.* Where possible, a high-quality implementation will try to return the errorcode from the MPI process startup mechanism (e.g. `mpiexec` or singleton init). (*End of advice to implementors.*)

### 8.7.1 Allowing User Functions at Process Termination

There are times in which it would be convenient to have actions happen when an MPI process finishes. For example, a routine may do initializations that are useful until the MPI job (or that part of the job that being terminated in the case of dynamically created processes) is finished. This can be accomplished in MPI by attaching an attribute to `MPI_COMM_SELF` with a callback function. When `MPI_FINALIZE` is called, it will first execute the equivalent of an `MPI_COMM_FREE` on `MPI_COMM_SELF`. This will cause the delete callback function to be executed on all keys associated with `MPI_COMM_SELF`, in the reverse order that they were set on `MPI_COMM_SELF`. If no key has been attached to `MPI_COMM_SELF`, then no callback is invoked. The “freeing” of `MPI_COMM_SELF` occurs before any other parts of MPI are affected. Thus, for example, calling `MPI_FINALIZED` will return `false` in any of these callback functions. Once done with `MPI_COMM_SELF`, the order and rest of the actions taken by `MPI_FINALIZE` is not specified.

*Advice to implementors.* Since attributes can be added from any supported language, the MPI implementation needs to remember the creating language so the correct callback is made. Implementations that use the attribute delete callback on `MPI_COMM_SELF` internally should register their internal callbacks before returning from `MPI_INIT` / `MPI_INIT_THREAD`, so that libraries or applications will not have portions of the MPI implementation shut down before the application-level callbacks are made. (*End of advice to implementors.*)

## 8.7.2 Determining Whether MPI Has Finished

One of the goals of MPI was to allow for layered libraries. In order for a library to do this cleanly, it needs to know if MPI is active. In MPI the function `MPI_INITIALIZED` was provided to tell if MPI had been initialized. The problem arises in knowing if MPI has been finalized. Once MPI has been finalized it is no longer active and cannot be restarted. A library needs to be able to determine this to act accordingly. To achieve this the following function is needed:

`MPI_FINALIZED(flag)`

OUT      flag                                      true if MPI was finalized (logical)

`int MPI_Finalized(int *flag)`

`MPI_Finalized(flag, ierror) BIND(C)`

`LOGICAL, INTENT(OUT) :: flag`

`INTEGER, OPTIONAL, INTENT(OUT) :: ierror`

`MPI_FINALIZED(FLAG, IERROR)`

`LOGICAL FLAG`

`INTEGER IERROR`

`{bool MPI::Is_finalized() (binding deprecated, see Section 15.2) }`

This routine returns true if `MPI_FINALIZE` has completed. It is legal to call `MPI_FINALIZED` before `MPI_INIT` and after `MPI_FINALIZE`.

*Advice to users.* MPI is “active” and it is thus safe to call MPI functions if `MPI_INIT` has completed and `MPI_FINALIZE` has not completed. If a library has no other way of knowing whether MPI is active or not, then it can use `MPI_INITIALIZED` and `MPI_FINALIZED` to determine this. For example, MPI is “active” in callback functions that are invoked during `MPI_FINALIZE`. (*End of advice to users.*)

## 8.8 Portable MPI Process Startup

A number of implementations of MPI provide a startup command for MPI programs that is of the form

`mpirun <mpirun arguments> <program> <program arguments>`

Separating the command to start the program from the program itself provides flexibility, particularly for network and heterogeneous implementations. For example, the startup script need not run on one of the machines that will be executing the MPI program itself.

Having a standard startup mechanism also extends the portability of MPI programs one step further, to the command lines and scripts that manage them. For example, a validation suite script that runs hundreds of programs can be a portable script if it is written using such a standard startup mechanism. In order that the “standard” command not be confused with existing practice, which is not standard and not portable among implementations, instead of `mpirun` MPI specifies `mpiexec`.

While a standardized startup mechanism improves the usability of MPI, the range of environments is so diverse (e.g., there may not even be a command line interface) that MPI cannot mandate such a mechanism. Instead, MPI specifies an `mpiexec` startup command and recommends but does not require it, as advice to implementors. However, if an implementation does provide a command called `mpiexec`, it must be of the form described below.

It is suggested that

```
mpiexec -n <numprocs> <program>
```

be at least one way to start `<program>` with an initial `MPI_COMM_WORLD` whose group contains `<numprocs>` processes. Other arguments to `mpiexec` may be implementation-dependent.

*Advice to implementors.* Implementors, if they do provide a special startup command for MPI programs, are advised to give it the following form. The syntax is chosen in order that `mpiexec` be able to be viewed as a command-line version of `MPI_COMM_SPAWN` (See Section 10.3.4).

Analogous to `MPI_COMM_SPAWN`, we have

```
mpiexec -n      <maxprocs>
           -soft  <      >
           -host  <      >
           -arch  <      >
           -wdir  <      >
           -path  <      >
           -file  <      >
           ...
           <command line>
```

for the case where a single command line for the application program and its arguments will suffice. See Section 10.3.4 for the meanings of these arguments. For the case corresponding to `MPI_COMM_SPAWN_MULTIPLE` there are two possible formats:

Form A:

```
mpiexec { <above arguments> } : { ... } : { ... } : ... : { ... }
```

As with `MPI_COMM_SPAWN`, all the arguments are optional. (Even the `-n x` argument is optional; the default is implementation dependent. It might be 1, it might be taken from an environment variable, or it might be specified at compile time.) The names and meanings of the arguments are taken from the keys in the `info` argument to `MPI_COMM_SPAWN`. There may be other, implementation-dependent arguments as well.

Note that Form A, though convenient to type, prevents colons from being program arguments. Therefore an alternate, file-based form is allowed:

Form B:

```
1      mpiexec -configfile <filename>
```

```
2
3      where the lines of <filename> are of the form separated by the colons in Form A.
4      Lines beginning with '#' are comments, and lines may be continued by terminating
5      the partial line with '\'.
6
```

```
7
8      Example 8.9 Start 16 instances of myprog on the current or default machine:
```

```
9
10     mpiexec -n 16 myprog
```

```
11
12     Example 8.10 Start 10 processes on the machine called ferrari:
```

```
13
14     mpiexec -n 10 -host ferrari myprog
```

```
15
16
17     Example 8.11 Start three copies of the same program with different command-line
18     arguments:
```

```
19
20     mpiexec myprog infile1 : myprog infile2 : myprog infile3
```

```
21
22
23     Example 8.12 Start the ocean program on five Suns and the atmos program on 10
24     RS/6000's:
```

```
25
26     mpiexec -n 5 -arch sun ocean : -n 10 -arch rs6000 atmos
```

```
27
28     It is assumed that the implementation in this case has a method for choosing hosts of
29     the appropriate type. Their ranks are in the order specified.
```

```
30
31     Example 8.13 Start the ocean program on five Suns and the atmos program on 10
32     RS/6000's (Form B):
```

```
33
34     mpiexec -configfile myfile
```

```
35
36     where myfile contains
```

```
37
38         -n 5  -arch sun    ocean
39         -n 10 -arch rs6000 atmos
```

```
40
41     (End of advice to implementors.)
```

# Chapter 9

## The Info Object

Many of the routines in MPI take an argument `info`. `info` is an opaque object with a handle of type `MPI_Info` in C and Fortran with the `mpi_f08` module, `MPI::Info` in C++, and `INTEGER` in Fortran with the `mpi` module or the include file `mpif.h`. It stores an unordered set of (key,value) pairs (both key and value are strings). A key can have only one value. MPI reserves several keys and requires that if an implementation uses a reserved key, it must provide the specified functionality. An implementation is not required to support these keys and may support any others not reserved by MPI.

An implementation must support info objects as caches for arbitrary (key, value) pairs, regardless of whether it recognizes the key. Each function that takes hints in the form of an `MPI_Info` must be prepared to ignore any key it does not recognize. This description of info objects does not attempt to define how a particular function should react if it recognizes a key but not the associated value. `MPI_INFO_GET_NKEYS`, `MPI_INFO_GET_NTHKEY`, `MPI_INFO_GET_VALUELEN`, and `MPI_INFO_GET` must retain all (key,value) pairs so that layered functionality can also use the `Info` object.

Keys have an implementation-defined maximum length of `MPI_MAX_INFO_KEY`, which is at least 32 and at most 255. Values have an implementation-defined maximum length of `MPI_MAX_INFO_VAL`. In Fortran, leading and trailing spaces are stripped from both. Returned values will never be larger than these maximum lengths. Both key and value are case sensitive.

*Rationale.* Keys have a maximum length because the set of known keys will always be finite and known to the implementation and because there is no reason for keys to be complex. The small maximum size allows applications to declare keys of size `MPI_MAX_INFO_KEY`. The limitation on value sizes is so that an implementation is not forced to deal with arbitrarily long strings. (*End of rationale.*)

*Advice to users.* `MPI_MAX_INFO_VAL` might be very large, so it might not be wise to declare a string of that size. (*End of advice to users.*)

When it is an argument to a nonblocking routine, `info` is parsed before that routine returns, so that it may be modified or freed immediately after return.

When the descriptions refer to a key or value as being a boolean, an integer, or a list, they mean the string representation of these types. An implementation may define its own rules for how info value strings are converted to other types, but to ensure portability, every implementation must support the following representations. Legal values for a boolean must

include the strings “true” and “false” (all lowercase). For integers, legal values must include string representations of decimal values of integers that are within the range of a standard integer type in the program. (However it is possible that not every legal integer is a legal value for a given key.) On positive numbers, + signs are optional. No space may appear between a + or – sign and the leading digit of a number. For comma separated lists, the string must contain legal elements separated by commas. Leading and trailing spaces are stripped automatically from the types of info values described above and for each element of a comma separated list. These rules apply to all info values of these types. Implementations are free to specify a different interpretation for values of other info keys.

**MPI\_INFO\_CREATE**(info)

OUT      info                                  info object created (handle)

**int MPI\_Info\_create**(MPI\_Info \*info)

**MPI\_Info\_create**(info, ierror) BIND(C)  
 TYPE(MPI\_Info), INTENT(OUT) :: info  
 INTEGER, OPTIONAL, INTENT(OUT) :: ierror

**MPI\_INFO\_CREATE**(INFO, IERROR)

INTEGER INFO, IERROR

{static MPI::Info MPI::Info::Create() (*binding deprecated, see Section 15.2*) }

**MPI\_INFO\_CREATE** creates a new info object. The newly created object contains no key/value pairs.

**MPI\_INFO\_SET**(info, key, value)

INOUT    info                                  info object (handle)

IN        key                                  key (string)

IN        value                                value (string)

**int MPI\_Info\_set**(MPI\_Info info, const char \*key, const char \*value)

**MPI\_Info\_set**(info, key, value, ierror) BIND(C)  
 TYPE(MPI\_Info), INTENT(IN) :: info  
 CHARACTER(LEN=\*), INTENT(IN) :: key, value  
 INTEGER, OPTIONAL, INTENT(OUT) :: ierror

**MPI\_INFO\_SET**(INFO, KEY, VALUE, IERROR)

INTEGER INFO, IERROR

CHARACTER\*(\*) KEY, VALUE

{void MPI::Info::Set(const char\* key, const char\* value) (*binding deprecated, see Section 15.2*) }

**MPI\_INFO\_SET** adds the (key,value) pair to info, and overrides the value if a value for the same key was previously set. key and value are null-terminated strings in C. In Fortran,

leading and trailing spaces in key and value are stripped. If either key or value are larger than the allowed maximums, the errors MPI\_ERR\_INFO\_KEY or MPI\_ERR\_INFO\_VALUE are raised, respectively.

MPI\_INFO\_DELETE(info, key)

INOUT	info	info object (handle)
IN	key	key (string)

```
int MPI_Info_delete(MPI_Info info, const char *key)
```

```
MPI_Info_delete(info, key, ierror) BIND(C)
    TYPE(MPI_Info), INTENT(IN) :: info
    CHARACTER(LEN=*), INTENT(IN) :: key
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_INFO_DELETE(INFO, KEY, IERROR)
    INTEGER INFO, IERROR
    CHARACTER*(*) KEY
```

```
{void MPI::Info::Delete(const char* key) (binding deprecated, see Section 15.2) }
```

MPI\_INFO\_DELETE deletes a (key,value) pair from info. If key is not defined in info, the call raises an error of class MPI\_ERR\_INFO\_NOKEY.

MPI\_INFO\_GET(info, key, valuelen, value, flag)

IN	info	info object (handle)
IN	key	key (string)
IN	valuelen	length of value arg (integer)
OUT	value	value (string)
OUT	flag	true if key defined, false if not (boolean)

```
int MPI_Info_get(MPI_Info info, const char *key, int valuelen, char *value,
    int *flag)
```

```
MPI_Info_get(info, key, valuelen, value, flag, ierror) BIND(C)
    TYPE(MPI_Info), INTENT(IN) :: info
    CHARACTER(LEN=*), INTENT(IN) :: key
    INTEGER, INTENT(IN) :: valuelen
    CHARACTER(LEN=valuelen), INTENT(OUT) :: value
    LOGICAL, INTENT(OUT) :: flag
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_INFO_GET(INFO, KEY, VALUELEN, VALUE, FLAG, IERROR)
    INTEGER INFO, VALUELEN, IERROR
    CHARACTER*(*) KEY, VALUE
    LOGICAL FLAG
```

```

1  {bool MPI::Info::Get(const char* key, int valuelen, char* value)
2      const(binding deprecated, see Section 15.2) }
3

```

This function retrieves the value associated with key in a previous call to MPI\_INFO\_SET. If such a key exists, it sets flag to true and returns the value in value, otherwise it sets flag to false and leaves value unchanged. valuelen is the number of characters available in value. If it is less than the actual size of the value, the value is truncated. In C, valuelen should be one less than the amount of allocated space to allow for the null terminator.

If key is larger than MPI\_MAX\_INFO\_KEY, the call is erroneous.

```

12 MPI_INFO_GET_VALUELEN(info, key, valuelen, flag)
13

```

14	IN	info	info object (handle)
15	IN	key	key (string)
16	OUT	valuelen	length of value arg (integer)
17	OUT	flag	true if key defined, false if not (boolean)

```

19
ticket140. 20 int MPI_Info_get_valuelen(MPI_Info info, const char *key, int *valuelen,
21                      int *flag)
ticket-248T. 22

```

```

23 MPI_Info_get_valuelen(info, key, valuelen, flag, ierror) BIND(C)
24     TYPE(MPI_Info), INTENT(IN) :: info
25     CHARACTER(LEN=*), INTENT(IN) :: key
26     INTEGER, INTENT(OUT) :: valuelen
27     LOGICAL, INTENT(OUT) :: flag
28     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
29

```

```

29 MPI_INFO_GET_VALUELEN(INFO, KEY, VALUELEN, FLAG, IERROR)
30     INTEGER INFO, VALUELEN, IERROR
31     LOGICAL FLAG
32     CHARACTER*(*) KEY
33

```

```

34 {bool MPI::Info::Get_valuelen(const char* key, int& valuelen) const(binding
35     deprecated, see Section 15.2) }
36

```

Retrieves the length of the value associated with key. If key is defined, valuelen is set to the length of its associated value and flag is set to true. If key is not defined, valuelen is not touched and flag is set to false. The length returned in C or C++ does not include the end-of-string character.

If key is larger than MPI\_MAX\_INFO\_KEY, the call is erroneous.

```

42 MPI_INFO_GET_NKEYS(info, nkeys)
43

```

44	IN	info	info object (handle)
45	OUT	nkeys	number of defined keys (integer)

```

47 int MPI_Info_get_nkeys(MPI_Info info, int *nkeys)
48

```



ticket-248T.

```

MPI_Info_get_nkeys(info, nkeys, ierror) BIND(C)
    TYPE(MPI_Info), INTENT(IN) :: info
    INTEGER, INTENT(OUT) :: nkeys
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_INFO_GET_NKEYS(INFO, NKEYS, IERROR)
    INTEGER INFO, NKEYS, IERROR

{int MPI::Info::Get_nkeys() const(binding deprecated, see Section 15.2) }

MPI_INFO_GET_NKEYS returns the number of currently defined keys in info.

```

```

MPI_INFO_GET_NTHKEY(info, n, key)

```

IN	info	info object (handle)
IN	n	key number (integer)
OUT	key	key (string)

```

int MPI_Info_get_nthkey(MPI_Info info, int n, char *key)

```

```

MPI_Info_get_nthkey(info, n, key, ierror) BIND(C)
    TYPE(MPI_Info), INTENT(IN) :: info
    INTEGER, INTENT(IN) :: n
    CHARACTER(LEN=*), INTENT(OUT) :: key
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_INFO_GET_NTHKEY(INFO, N, KEY, IERROR)
    INTEGER INFO, N, IERROR
    CHARACTER*(*) KEY

```

```

{void MPI::Info::Get_nthkey(int n, char* key) const(binding deprecated, see
    Section 15.2) }

```

This function returns the  $n$ th defined key in `info`. Keys are numbered  $0 \dots N - 1$  where  $N$  is the value returned by `MPI_INFO_GET_NKEYS`. All keys between 0 and  $N - 1$  are guaranteed to be defined. The number of a given key does not change as long as `info` is not modified with `MPI_INFO_SET` or `MPI_INFO_DELETE`.

```

MPI_INFO_DUP(info, newinfo)

```

IN	info	info object (handle)
OUT	newinfo	info object (handle)

```

int MPI_Info_dup(MPI_Info info, MPI_Info *newinfo)

```

```

MPI_Info_dup(info, newinfo, ierror) BIND(C)
    TYPE(MPI_Info), INTENT(IN) :: info
    TYPE(MPI_Info), INTENT(OUT) :: newinfo
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```
1 MPI_INFO_DUP(INFO, NEWINFO, IERROR)
```

```
2     INTEGER INFO, NEWINFO, IERROR
```

```
3 {MPI::Info MPI::Info::Dup() const(binding deprecated, see Section 15.2) }
```

5 MPI\_INFO\_DUP duplicates an existing info object, creating a new object, with the  
6 same (key,value) pairs and the same ordering of keys.

```
8 MPI_INFO_FREE(info)
```

```
10     INOUT    info                      info object (handle)
```

```
12 int MPI_Info_free(MPI_Info *info)
```

```
14 MPI_Info_free(info, ierror) BIND(C)
```

```
15     TYPE(MPI_Info), INTENT(INOUT) :: info
```

```
16     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
17 MPI_INFO_FREE(INFO, IERROR)
```

```
18     INTEGER INFO, IERROR
```

```
20 {void MPI::Info::Free() (binding deprecated, see Section 15.2) }
```

21 This function frees info and sets it to MPI\_INFO\_NULL. The value of an info argument is  
22 interpreted each time the info is passed to a routine. Changes to an info after return from  
23 a routine do not affect that interpretation.

# Chapter 10

## Process Creation and Management

### 10.1 Introduction

MPI is primarily concerned with communication rather than process or resource management. However, it is necessary to address these issues to some degree in order to define a useful framework for communication. This chapter presents a set of MPI interfaces that allow for a variety of approaches to process management while placing minimal restrictions on the execution environment.

The MPI model for process creation allows both the creation of an initial set of processes related by their membership in a common MPI\_COMM\_WORLD and the creation and management of processes after an MPI application has been started. A major impetus for the later form of process creation comes from the PVM [24] research effort. This work has provided a wealth of experience with process management and resource control that illustrates their benefits and potential pitfalls.

The MPI Forum decided not to address resource control because it was not able to design a portable interface that would be appropriate for the broad spectrum of existing and potential resource and process controllers. Resource control can encompass a wide range of abilities, including adding and deleting nodes from a virtual parallel machine, reserving and scheduling resources, managing compute partitions of an MPP, and returning information about available resources. assumes that resource control is provided externally — probably by computer vendors, in the case of tightly coupled systems, or by a third party software package when the environment is a cluster of workstations.

The reasons for including process management in MPI are both technical and practical. Important classes of message-passing applications require process control. These include task farms, serial applications with parallel modules, and problems that require a run-time assessment of the number and type of processes that should be started. On the practical side, users of workstation clusters who are migrating from PVM to MPI may be accustomed to using PVM's capabilities for process and resource management. The lack of these features would be a practical stumbling block to migration.

The following goals are central to the design of MPI process management:

- The MPI process model must apply to the vast majority of current parallel environments. These include everything from tightly integrated MPPs to heterogeneous networks of workstations.
- MPI must not take over operating system responsibilities. It should instead provide a

clean interface between an application and system software.

- MPI must guarantee communication determinism in the presense of dynamic processes, i.e., dynamic process management must not introduce unavoidable race conditions.
- MPI must not contain features that compromise performance.

The process management model addresses these issues in two ways. First, MPI remains primarily a communication library. It does not manage the parallel environment in which a parallel program executes, though it provides a minimal interface between an application and external resource and process managers.

Second, MPI maintains a consistent concept of a communicator, regardless of how its members came into existence. A communicator is never changed once created, and it is always created using deterministic collective operations.

## 10.2 The Dynamic Process Model

The dynamic process model allows for the creation and cooperative termination of processes after an MPI application has started. It provides a mechanism to establish communication between the newly created processes and the existing MPI application. It also provides a mechanism to establish communication between two existing MPI applications, even when one did not “start” the other.

### 10.2.1 Starting Processes

MPI applications may start new processes through an interface to an external process manager.

MPI\_COMM\_SPAWN starts MPI processes and establishes communication with them, returning an intercommunicator. MPI\_COMM\_SPAWN\_MULTIPLE starts several different binaries (or the same binary with different arguments), placing them in the same MPI\_COMM\_WORLD and returning an intercommunicator.

MPI uses the existing group abstraction to represent processes. A process is identified by a (group, rank) pair.

### 10.2.2 The Runtime Environment

The MPI\_COMM\_SPAWN and MPI\_COMM\_SPAWN\_MULTIPLE routines provide an interface between MPI and the *runtime environment* of an MPI application. The difficulty is that there is an enormous range of runtime environments and application requirements, and MPI must not be tailored to any particular one. Examples of such environments are:

- **MPP managed by a batch queueing system.** Batch queueing systems generally allocate resources before an application begins, enforce limits on resource use (CPU time, memory use, etc.), and do not allow a change in resource allocation after a job begins. Moreover, many MPPs have special limitations or extensions, such as a limit on the number of processes that may run on one processor, or the ability to gang-schedule processes of a parallel application.

- **Network of workstations with PVM.** PVM (Parallel Virtual Machine) allows a user to create a “virtual machine” out of a network of workstations. An application may extend the virtual machine or manage processes (create, kill, redirect output, etc.) through the PVM library. Requests to manage the machine or processes may be intercepted and handled by an external resource manager.
- **Network of workstations managed by a load balancing system.** A load balancing system may choose the location of spawned processes based on dynamic quantities, such as load average. It may transparently migrate processes from one machine to another when a resource becomes unavailable.
- **Large SMP with Unix.** Applications are run directly by the user. They are scheduled at a low level by the operating system. Processes may have special scheduling characteristics (gang-scheduling, processor affinity, deadline scheduling, processor locking, etc.) and be subject to OS resource limits (number of processes, amount of memory, etc.).

MPI assumes, implicitly, the existence of an environment in which an application runs. It does not provide “operating system” services, such as a general ability to query what processes are running, to kill arbitrary processes, to find out properties of the runtime environment (how many processors, how much memory, etc.).

Complex interaction of an MPI application with its runtime environment should be done through an environment-specific API. An example of such an API would be the PVM task and machine management routines — `pvm_addhosts`, `pvm_config`, `pvm_tasks`, etc., possibly modified to return an MPI (group,rank) when possible. A Condor or PBS API would be another possibility.

At some low level, obviously, MPI must be able to interact with the runtime system, but the interaction is not visible at the application level and the details of the interaction are not specified by the MPI standard.

In many cases, it is impossible to keep environment-specific information out of the MPI interface without seriously compromising MPI functionality. To permit applications to take advantage of environment-specific functionality, many MPI routines take an `info` argument that allows an application to specify environment-specific information. There is a tradeoff between functionality and portability: applications that make use of `info` are not portable.

MPI does not require the existence of an underlying “virtual machine” model, in which there is a consistent global view of an MPI application and an implicit “operating system” managing resources and processes. For instance, processes spawned by one task may not be visible to another; additional hosts added to the runtime environment by one process may not be visible in another process; tasks spawned by different processes may not be automatically distributed over available resources.

Interaction between MPI and the runtime environment is limited to the following areas:

- A process may start new processes with `MPI_COMM_SPAWN` and `MPI_COMM_SPAWN_MULTIPLE`.
- When a process spawns a child process, it may optionally use an `info` argument to tell the runtime environment where or how to start the process. This extra information may be opaque to MPI.

- An attribute `MPI_UNIVERSE_SIZE` on `MPI_COMM_WORLD` tells a program how “large” the initial runtime environment is, namely how many processes can usefully be started in all. One can subtract the size of `MPI_COMM_WORLD` from this value to find out how many processes might usefully be started in addition to those already running.

## 10.3 Process Manager Interface

### 10.3.1 Processes in MPI

A process is represented in MPI by a (group, rank) pair. A (group, rank) pair specifies a unique process but a process does not determine a unique (group, rank) pair, since a process may belong to several groups.

### 10.3.2 Starting Processes and Establishing Communication

The following routine starts a number of MPI processes and establishes communication with them, returning an intercommunicator.

*Advice to users.* It is possible in MPI to start a static SPMD or MPMD application by starting first one process and having that process start its siblings with `MPI_COMM_SPAWN`. This practice is discouraged primarily for reasons of performance. If possible, it is preferable to start all processes at once, as a single MPI application. (*End of advice to users.*)

`MPI_COMM_SPAWN(command, argv, maxprocs, info, root, comm, intercomm,  
array_of_errcodes)`

IN	command	name of program to be spawned (string, significant only at root)
IN	argv	arguments to <code>command</code> (array of strings, significant only at root)
IN	maxprocs	maximum number of processes to start (integer, significant only at root)
IN	info	a set of key-value pairs telling the runtime system where and how to start the processes (handle, significant only at root)
IN	root	rank of process in which previous arguments are examined (integer)
IN	comm	intracommunicator containing group of spawning processes (handle)
OUT	intercomm	intercommunicator between original group and the newly spawned group (handle)
OUT	array_of_errcodes	one code per process (array of integer)

```

int MPI_Comm_spawn(char *command, char *argv[], int maxprocs, MPI_Info
                    info, int root, MPI_Comm comm, MPI_Comm *intercomm,
                    int array_of_errcodes[])
MPI_Comm_spawn(command, argv, maxprocs, info, root, comm, intercomm,
                array_of_errcodes, ierror) BIND(C)
CHARACTER(LEN=*) INTENT(IN) :: command, argv(*)
INTEGER, INTENT(IN) :: maxprocs, root
TYPE(MPI_Info), INTENT(IN) :: info
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Comm), INTENT(OUT) :: intercomm
INTEGER :: array_of_errcodes(*)
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_COMM_SPAWN(COMMAND, ARGV, MAXPROCS, INFO, ROOT, COMM, INTERCOMM,
                ARRAY_OF_ERRCODES, IERROR)
CHARACTER*(*) COMMAND, ARGV(*)
INTEGER INFO, MAXPROCS, ROOT, COMM, INTERCOMM, ARRAY_OF_ERRCODES(*),
IERROR

{MPI::Intercomm MPI::Intracomm::Spawn(const char* command,
    const char* argv[], int maxprocs, const MPI::Info& info,
    int root, int array_of_errcodes[]) const(binding deprecated, see
    Section 15.2) }

{MPI::Intercomm MPI::Intracomm::Spawn(const char* command,
    const char* argv[], int maxprocs, const MPI::Info& info,
    int root) const(binding deprecated, see Section 15.2) }

```

MPI\_COMM\_SPAWN tries to start maxprocs identical copies of the MPI program specified by `command`, establishing communication with them and returning an intercommunicator. The spawned processes are referred to as children. The children have their own MPI\_COMM\_WORLD, which is separate from that of the parents. MPI\_COMM\_SPAWN is collective over `comm`, and also may not return until MPI\_INIT has been called in the children. Similarly, MPI\_INIT in the children may not return until all parents have called MPI\_COMM\_SPAWN. In this sense, MPI\_COMM\_SPAWN in the parents and MPI\_INIT in the children form a collective operation over the union of parent and child processes. The intercommunicator returned by MPI\_COMM\_SPAWN contains the parent processes in the local group and the child processes in the remote group. The ordering of processes in the local and remote groups is the same as the ordering of the group of the `comm` in the parents and of MPI\_COMM\_WORLD of the children, respectively. This intercommunicator can be obtained in the children through the function MPI\_COMM\_GET\_PARENT.

*Advice to users.* An implementation may automatically establish communication before MPI\_INIT is called by the children. Thus, completion of MPI\_COMM\_SPAWN in the parent does not necessarily mean that MPI\_INIT has been called in the children (although the returned intercommunicator can be used immediately). (*End of advice to users.*)

The **command** argument The **command** argument is a string containing the name of a program to be spawned. The string is null-terminated in C. In Fortran, leading and trailing spaces are stripped. MPI does not specify how to find the executable or how the working directory is determined. These rules are implementation-dependent and should be appropriate for the runtime environment.

*Advice to implementors.* The implementation should use a natural rule for finding executables and determining working directories. For instance, a homogeneous system with a global file system might look first in the working directory of the spawning process, or might search the directories in a **PATH** environment variable as do Unix shells. An implementation on top of PVM would use PVM's rules for finding executables (usually in `$HOME/pvm3/bin/$PVM_ARCH`). An MPI implementation running under POE on an IBM SP would use POE's method of finding executables. An implementation should document its rules for finding executables and determining working directories, and a high-quality implementation should give the user some control over these rules. (*End of advice to implementors.*)

If the program named in **command** does not call **MPI\_INIT**, but instead forks a process that calls **MPI\_INIT**, the results are undefined. Implementations may allow this case to work but are not required to.

*Advice to users.* MPI does not say what happens if the program you start is a shell script and that shell script starts a program that calls **MPI\_INIT**. Though some implementations may allow you to do this, they may also have restrictions, such as requiring that arguments supplied to the shell script be supplied to the program, or requiring that certain parts of the environment not be changed. (*End of advice to users.*)

The **argv** argument **argv** is an array of strings containing arguments that are passed to the program. The first element of **argv** is the first argument passed to **command**, not, as is conventional in some contexts, the command itself. The argument list is terminated by **NULL** in C and C++ and an empty string in Fortran. In Fortran, leading and trailing spaces are always stripped, so that a string consisting of all spaces is considered an empty string. The constant **MPI\_ARGV\_NULL** may be used in C, C++ and Fortran to indicate an empty argument list. In C and C++, this constant is the same as **NULL**.

#### **Example 10.1** Examples of **argv** in C and Fortran

To run the program "ocean" with arguments "-gridfile" and "ocean1.grd" in C:

```
char command[] = "ocean";
char *argv[] = {"-gridfile", "ocean1.grd", NULL};
MPI_Comm_spawn(command, argv, ...);
```

or, if not everything is known at compile time:



```

char *command;
char **argv;
command = "ocean";
argv=(char **)malloc(3 * sizeof(char *));
argv[0] = "-gridfile";
argv[1] = "ocean1.grd";
argv[2] = NULL;
MPI_Comm_spawn(command, argv, ...);

```

In Fortran:

```

CHARACTER*25 command, argv(3)
command = ' ocean '
argv(1) = ' -gridfile '
argv(2) = ' ocean1.grd'
argv(3) = ' '
call MPI_COMM_SPAWN(command, argv, ...)

```

Arguments are supplied to the program if this is allowed by the operating system. In C, the `MPI_COMM_SPAWN` argument `argv` differs from the `argv` argument of `main` in two respects. First, it is shifted by one element. Specifically, `argv[0]` of `main` is provided by the implementation and conventionally contains the name of the program (given by `command`). `argv[1]` of `main` corresponds to `argv[0]` in `MPI_COMM_SPAWN`, `argv[2]` of `main` to `argv[1]` of `MPI_COMM_SPAWN`, etc. Second, `argv` of `MPI_COMM_SPAWN` must be null-terminated, so that its length can be determined. Passing an `argv` of `MPI_ARGV_NULL` to `MPI_COMM_SPAWN` results in `main` receiving `argc` of 1 and an `argv` whose element 0 is (conventionally) the name of the program.

If a Fortran implementation supplies routines that allow a program to obtain its arguments, the arguments may be available through that mechanism. In C, if the operating system does not support arguments appearing in `argv` of `main()`, the MPI implementation may add the arguments to the `argv` that is passed to `MPI_INIT`.

The `maxprocs` argument MPI tries to spawn `maxprocs` processes. If it is unable to spawn `maxprocs` processes, it raises an error of class `MPI_ERR_SPAWN`.

An implementation may allow the `info` argument to change the default behavior, such that if the implementation is unable to spawn all `maxprocs` processes, it may spawn a smaller number of processes instead of raising an error. In principle, the `info` argument may specify an arbitrary set  $\{m_i : 0 \leq m_i \leq \text{maxprocs}\}$  of allowed values for the number of processes spawned. The set  $\{m_i\}$  does not necessarily include the value `maxprocs`. If an implementation is able to spawn one of these allowed numbers of processes, `MPI_COMM_SPAWN` returns successfully and the number of spawned processes,  $m$ , is given by the size of the remote group of `intercomm`. If  $m$  is less than `maxproc`, reasons why the other processes were not spawned are given in `array_of_errcodes` as described below. If it is not possible to spawn one of the allowed numbers of processes, `MPI_COMM_SPAWN` raises an error of class `MPI_ERR_SPAWN`.

A spawn call with the default behavior is called *hard*. A spawn call for which fewer than `maxprocs` processes may be returned is called *soft*. See Section 10.3.4 on page 398 for more information on the `soft` key for `info`.

*Advice to users.* By default, requests are hard and MPI errors are fatal. This means that by default there will be a fatal error if MPI cannot spawn all the requested processes. If you want the behavior “spawn as many processes as possible, up to  $N$ ,” you should do a soft spawn, where the set of allowed values  $\{m_i\}$  is  $\{0 \dots N\}$ . However, this is not completely portable, as implementations are not required to support soft spawning. (*End of advice to users.*)

**The info argument** The info argument to all of the routines in this chapter is an opaque handle of type `MPI_Info` in C and Fortran with the `mpi_f08` module, `MPI::Info` in C++ and `INTEGER` in Fortran with the `mpi` module or the include file `mpif.h`. It is a container for a number of user-specified (key,value) pairs. key and value are strings (null-terminated `char*` in C, `character(*)` in Fortran). Routines to create and manipulate the info argument are described in Section 9 on page 381.

For the `SPAWN` calls, info provides additional (and possibly implementation-dependent) instructions to MPI and the runtime system on how to start processes. An application may pass `MPI_INFO_NULL` in C or Fortran. Portable programs not requiring detailed control over process locations should use `MPI_INFO_NULL`.

MPI does not specify the content of the info argument, except to reserve a number of special key values (see Section 10.3.4 on page 398). The info argument is quite flexible and could even be used, for example, to specify the executable and its command-line arguments. In this case the `command` argument to `MPI_COMM_SPAWN` could be empty. The ability to do this follows from the fact that MPI does not specify how an executable is found, and the info argument can tell the runtime system where to “find” the executable “” (empty string). Of course a program that does this will not be portable across MPI implementations.

**The root argument** All arguments before the root argument are examined only on the process whose rank in `comm` is equal to `root`. The value of these arguments on other processes is ignored.

**The array\_of\_errcodes argument** The `array_of_errcodes` is an array of length `maxprocs` in which MPI reports the status of each process that MPI was requested to start. If all `maxprocs` processes were spawned, `array_of_errcodes` is filled in with the value `MPI_SUCCESS`. If only  $m$  ( $0 \leq m < \text{maxprocs}$ ) processes are spawned,  $m$  of the entries will contain `MPI_SUCCESS` and the rest will contain an implementation-specific error code indicating the reason MPI could not start the process. MPI does not specify which entries correspond to failed processes. An implementation may, for instance, fill in error codes in one-to-one correspondence with a detailed specification in the info argument. These error codes all belong to the error class `MPI_ERR_SPAWN` if there was no error in the argument list. In C or Fortran, an application may pass `MPI_ERRCODES_IGNORE` if it is not interested in the error codes. In C++ this constant does not exist, and the `array_of_errcodes` argument may be omitted from the argument list.

*Advice to implementors.* `MPI_ERRCODES_IGNORE` in Fortran is a special type of constant, like `MPI_BOTTOM`. See the discussion in Section 2.5.4 on page 15. (*End of advice to implementors.*)

MPI\_COMM\_GET\_PARENT(parent)

OUT      parent      the parent communicator (handle)

```
int MPI_Comm_get_parent(MPI_Comm *parent)
```

```
MPI_Comm_get_parent(parent, ierror) BIND(C)
  TYPE(MPI_Comm), INTENT(OUT) :: parent
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_COMM_GET_PARENT(PARENT, IERROR)
  INTEGER PARENT, IERROR
```

```
{static MPI::Intercomm MPI::Comm::Get_parent() (binding deprecated, see
  Section 15.2) }
```

If a process was started with MPI\_COMM\_SPAWN or MPI\_COMM\_SPAWN\_MULTIPLE, MPI\_COMM\_GET\_PARENT returns the “parent” intercommunicator of the current process. This parent intercommunicator is created implicitly inside of MPI\_INIT and is the same intercommunicator returned by SPAWN in the parents.

If the process was not spawned, MPI\_COMM\_GET\_PARENT returns MPI\_COMM\_NULL.

After the parent communicator is freed or disconnected, MPI\_COMM\_GET\_PARENT returns MPI\_COMM\_NULL.

*Advice to users.* MPI\_COMM\_GET\_PARENT returns a handle to a single intercommunicator. Calling MPI\_COMM\_GET\_PARENT a second time returns a handle to the same intercommunicator. Freeing the handle with MPI\_COMM\_DISCONNECT or MPI\_COMM\_FREE will cause other references to the intercommunicator to become invalid (dangling). Note that calling MPI\_COMM\_FREE on the parent communicator is not useful. (*End of advice to users.*)

*Rationale.* The desire of the Forum was to create a constant MPI\_COMM\_PARENT similar to MPI\_COMM\_WORLD. Unfortunately such a constant cannot be used (syntactically) as an argument to MPI\_COMM\_DISCONNECT, which is explicitly allowed. (*End of rationale.*)

### 10.3.3 Starting Multiple Executables and Establishing Communication

While MPI\_COMM\_SPAWN is sufficient for most cases, it does not allow the spawning of multiple binaries, or of the same binary with multiple sets of arguments. The following routine spawns multiple binaries or the same binary with multiple sets of arguments, establishing communication with them and placing them in the same MPI\_COMM\_WORLD.

```

1 MPI_COMM_SPAWN_MULTIPLE(count, array_of_commands, array_of_argv, array_of_maxprocs,
2   array_of_info, root, comm, intercomm, array_of_errcodes)

```

3	IN	count	number of commands (positive integer, significant to MPI only at root — see advice to users)
4			
5			
6	IN	array_of_commands	programs to be executed (array of strings, significant only at root)
7			
8	IN	array_of_argv	arguments for commands (array of array of strings, significant only at root)
9			
10	IN	array_of_maxprocs	maximum number of processes to start for each command (array of integer, significant only at root)
11			
12			
13	IN	array_of_info	info objects telling the runtime system where and how to start processes (array of handles, significant only at root)
14			
15			
16	IN	root	rank of process in which previous arguments are examined (integer)
17			
18	IN	comm	intracommunicator containing group of spawning processes (handle)
19			
20			
21	OUT	intercomm	intercommunicator between original group and newly spawned group (handle)
22			
23	OUT	array_of_errcodes	one error code per process (array of integer)
24			

```

25 int MPI_Comm_spawn_multiple(int count, char *array_of_commands[],
26   char **array_of_argv[], int array_of_maxprocs[],
27   MPI_Info array_of_info[], int root, MPI_Comm comm,
28   MPI_Comm *intercomm, int array_of_errcodes[])

```

```

29 MPI_Comm_spawn_multiple(count, array_of_commands, array_of_argv,
30   array_of_maxprocs, array_of_info, root, comm, intercomm,
31   array_of_errcodes, ierror) BIND(C)
32   INTEGER, INTENT(IN) :: count, array_of_maxprocs(*), root
33   CHARACTER(LEN=*), INTENT(IN) :: array_of_commands(*),
34   array_of_argv(count, *)
35   TYPE(MPI_Info), INTENT(IN) :: array_of_info(*)
36   TYPE(MPI_Comm), INTENT(IN) :: comm
37   TYPE(MPI_Comm), INTENT(OUT) :: intercomm
38   INTEGER :: array_of_errcodes(*)
39   INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

41 MPI_COMM_SPAWN_MULTIPLE(COUNT, ARRAY_OF_COMMANDS, ARRAY_OF_ARGV,
42   ARRAY_OF_MAXPROCS, ARRAY_OF_INFO, ROOT, COMM, INTERCOMM,
43   ARRAY_OF_ERRCODES, IERROR)
44   INTEGER COUNT, ARRAY_OF_INFO(*), ARRAY_OF_MAXPROCS(*), ROOT, COMM,
45   INTERCOMM, ARRAY_OF_ERRCODES(*), IERROR
46   CHARACTER*(*) ARRAY_OF_COMMANDS(*), ARRAY_OF_ARGV(COUNT, *)

```

```

{MPI::Intercomm MPI::Intracomm::Spawn_multiple(int count,
    const char* array_of_commands[], const char** array_of_argv[],
    const int array_of_maxprocs[],
    const MPI::Info array_of_info[], int root,
    int array_of_errcodes[]) (binding deprecated, see Section 15.2) }

{MPI::Intercomm MPI::Intracomm::Spawn_multiple(int count,
    const char* array_of_commands[], const char** array_of_argv[],
    const int array_of_maxprocs[],
    const MPI::Info array_of_info[], int root) (binding deprecated, see
    Section 15.2) }

```

MPI\_COMM\_SPAWN\_MULTIPLE is identical to MPI\_COMM\_SPAWN except that there are multiple executable specifications. The first argument, `count`, gives the number of specifications. Each of the next four arguments are simply arrays of the corresponding arguments in MPI\_COMM\_SPAWN. For the Fortran version of `array_of_argv`, the element `array_of_argv(i,j)` is the  $j$ -th argument to command number  $i$ .

*Rationale.* This may seem backwards to Fortran programmers who are familiar with Fortran's column-major ordering. However, it is necessary to do it this way to allow MPI\_COMM\_SPAWN to sort out arguments. Note that the leading dimension of `array_of_argv` *must* be the same as `count`. (*End of rationale.*)

*Advice to users.* The argument `count` is interpreted by MPI only at the root, as is `array_of_argv`. Since the leading dimension of `array_of_argv` is `count`, a non-positive value of `count` at a non-root node could theoretically cause a runtime bounds check error, even though `array_of_argv` should be ignored by the subroutine. If this happens, you should explicitly supply a reasonable value of `count` on the non-root nodes. (*End of advice to users.*)

In any language, an application may use the constant MPI\_ARGVS\_NULL (which is likely to be `(char ***)0` in C) to specify that no arguments should be passed to any commands. The effect of setting individual elements of `array_of_argv` to MPI\_ARGV\_NULL is not defined. To specify arguments for some commands but not others, the commands without arguments should have a corresponding `argv` whose first element is null (`(char *)0` in C and empty string in Fortran). **In Fortran at non-root processes, the `count` argument must be set to a value that is consistent with the provided `array_of_argv` although the content of these arguments has no meaning for this operation.**

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All of the spawned processes have the same MPI\_COMM\_WORLD. Their ranks in MPI\_COMM\_WORLD correspond directly to the order in which the commands are specified in MPI\_COMM\_SPAWN\_MULTIPLE. Assume that  $m_1$  processes are generated by the first command,  $m_2$  by the second, etc. The processes corresponding to the first command have ranks  $0, 1, \dots, m_1 - 1$ . The processes in the second command have ranks  $m_1, m_1 + 1, \dots, m_1 + m_2 - 1$ . The processes in the third have ranks  $m_1 + m_2, m_1 + m_2 + 1, \dots, m_1 + m_2 + m_3 - 1$ , etc.

*Advice to users.* Calling MPI\_COMM\_SPAWN multiple times would create many sets of children with different MPI\_COMM\_WORLDS whereas MPI\_COMM\_SPAWN\_MULTIPLE creates children with a single MPI\_COMM\_WORLD,

so the two methods are not completely equivalent. There are also two performance-related reasons why, if you need to spawn multiple executables, you may want to use `MPI_COMM_SPAWN_MULTIPLE` instead of calling `MPI_COMM_SPAWN` several times. First, spawning several things at once may be faster than spawning them sequentially. Second, in some implementations, communication between processes spawned at the same time may be faster than communication between processes spawned separately. (*End of advice to users.*)

The `array_of_errcodes` argument is a 1-dimensional array of size  $\sum_{i=1}^{count} n_i$ , where  $n_i$  is the  $i$ -th element of `array_of_maxprocs`. Command number  $i$  corresponds to the  $n_i$  contiguous slots in this array from element  $\sum_{j=1}^{i-1} n_j$  to  $\left[\sum_{j=1}^i n_j\right] - 1$ . Error codes are treated as for `MPI_COMM_SPAWN`.

#### Example 10.2 Examples of `array_of_argv` in C and Fortran

To run the program “ocean” with arguments “-gridfile” and “ocean1.grd” and the program “atmos” with argument “atmos.grd” in C:

```
char *array_of_commands[2] = {"ocean", "atmos"};
char **array_of_argv[2];
char *argv0[] = {"-gridfile", "ocean1.grd", (char *)0};
char *argv1[] = {"atmos.grd", (char *)0};
array_of_argv[0] = argv0;
array_of_argv[1] = argv1;
MPI_Comm_spawn_multiple(2, array_of_commands, array_of_argv, ...);
```

Here’s how you do it in Fortran:

```
CHARACTER*25 commands(2), array_of_argv(2, 3)
commands(1) = ' ocean '
array_of_argv(1, 1) = ' -gridfile '
array_of_argv(1, 2) = ' ocean1.grd'
array_of_argv(1, 3) = ' '

commands(2) = ' atmos '
array_of_argv(2, 1) = ' atmos.grd '
array_of_argv(2, 2) = ' '

call MPI_COMM_SPAWN_MULTIPLE(2, commands, array_of_argv, ...)
```

#### 10.3.4 Reserved Keys

The following keys are reserved. An implementation is not required to interpret these keys, but if it does interpret the key, it must provide the functionality described.

**host** Value is a hostname. The format of the hostname is determined by the implementation.

**arch** Value is an architecture name. Valid architecture names and what they mean are determined by the implementation.

**wdir** Value is the name of a directory on a machine on which the spawned process(es) execute(s). This directory is made the working directory of the executing process(es). The format of the directory name is determined by the implementation.

**path** Value is a directory or set of directories where the implementation should look for the executable. The format of path is determined by the implementation.

**file** Value is the name of a file in which additional information is specified. The format of the filename and internal format of the file are determined by the implementation.

**soft** Value specifies a set of numbers which are allowed values for the number of processes that `MPI_COMM_SPAWN` (et al.) may create. The format of the value is a comma-separated list of Fortran-90 triplets each of which specifies a set of integers and which together specify the set formed by the union of these sets. Negative values in this set and values greater than `maxprocs` are ignored. MPI will spawn the largest number of processes it can, consistent with some number in the set. The order in which triplets are given is not significant.

By Fortran-90 triplets, we mean:

1. `a` means  $a$
2. `a:b` means  $a, a + 1, a + 2, \dots, b$
3. `a:b:c` means  $a, a + c, a + 2c, \dots, a + ck$ , where for  $c > 0$ ,  $k$  is the largest integer for which  $a + ck \leq b$  and for  $c < 0$ ,  $k$  is the largest integer for which  $a + ck \geq b$ . If  $b > a$  then  $c$  must be positive. If  $b < a$  then  $c$  must be negative.

Examples:

1. `a:b` gives a range between  $a$  and  $b$
2. `0:N` gives full “soft” functionality
3. `1,2,4,8,16,32,64,128,256,512,1024,2048,4096` allows power-of-two number of processes.
4. `2:10000:2` allows even number of processes.
5. `2:10:2,7` allows 2, 4, 6, 7, 8, or 10 processes.

### 10.3.5 Spawn Example

Manager-worker Example [.] Using `MPI_COMM_SPAWN`.

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```
/* manager */
#include "mpi.h"
int main(int argc, char *argv[])
{
    int world_size, universe_size, *universe_sizep, flag;
    MPI_Comm everyone;          /* intercommunicator */
    char worker_program[100];

    MPI_Init(&argc, &argv);
    MPI_Comm_size(MPI_COMM_WORLD, &world_size);
```

```

1      if (world_size != 1)    error("Top heavy with management");
2
3      MPI_Comm_get_attr(MPI_COMM_WORLD, MPI_UNIVERSE_SIZE,
4                        &universe_sizep, &flag);
5      if (!flag) {
6          printf("This MPI does not support UNIVERSE_SIZE. How many\n\
7 processes total?");
8          scanf("%d", &universe_size);
9      } else universe_size = *universe_sizep;
10     if (universe_size == 1) error("No room to start workers");
11
12     /*
13      * Now spawn the workers. Note that there is a run-time determination
14      * of what type of worker to spawn, and presumably this calculation must
15      * be done at run time and cannot be calculated before starting
16      * the program. If everything is known when the application is
17      * first started, it is generally better to start them all at once
18      * in a single MPI_COMM_WORLD.
19      */
20
21     choose_worker_program(worker_program);
22     MPI_Comm_spawn(worker_program, MPI_ARGV_NULL, universe_size-1,
23                   MPI_INFO_NULL, 0, MPI_COMM_SELF, &everyone,
24                   MPI_ERRCODES_IGNORE);
25     /*
26      * Parallel code here. The communicator "everyone" can be used
27      * to communicate with the spawned processes, which have ranks 0,..
28      * MPI_UNIVERSE_SIZE-1 in the remote group of the intercommunicator
29      * "everyone".
30      */
31
32     MPI_Finalize();
33     return 0;
34 }
35
36 /* worker */
37
38 #include "mpi.h"
39 int main(int argc, char *argv[])
40 {
41     int size;
42     MPI_Comm parent;
43     MPI_Init(&argc, &argv);
44     MPI_Comm_get_parent(&parent);
45     if (parent == MPI_COMM_NULL) error("No parent!");
46     MPI_Comm_remote_size(parent, &size);
47     if (size != 1) error("Something's wrong with the parent");
48

```



```

/*
 * Parallel code here.
 * The manager is represented as the process with rank 0 in (the remote
 * group of) the parent communicator. If the workers need to communicate
 * among themselves, they can use MPI_COMM_WORLD.
 */

MPI_Finalize();
return 0;
}

```

## 10.4 Establishing Communication

This section provides functions that establish communication between two sets of MPI processes that do not share a communicator.

Some situations in which these functions are useful are:

1. Two parts of an application that are started independently need to communicate.
2. A visualization tool wants to attach to a running process.
3. A server wants to accept connections from multiple clients. Both clients and server may be parallel programs.

In each of these situations, MPI must establish communication channels where none existed before, and there is no parent/child relationship. The routines described in this section establish communication between the two sets of processes by creating an MPI intercommunicator, where the two groups of the intercommunicator are the original sets of processes.

Establishing contact between two groups of processes that do not share an existing communicator is a collective but asymmetric process. One group of processes indicates its willingness to accept connections from other groups of processes. We will call this group the (parallel) *server*, even if this is not a client/server type of application. The other group connects to the server; we will call it the *client*.

*Advice to users.* While the names *client* and *server* are used throughout this section, MPI does not guarantee the traditional robustness of client server systems. The functionality described in this section is intended to allow two cooperating parts of the same application to communicate with one another. For instance, a client that gets a segmentation fault and dies, or one that doesn't participate in a collective operation may cause a server to crash or hang. (*End of advice to users.*)

### 10.4.1 Names, Addresses, Ports, and All That

Almost all of the complexity in MPI client/server routines addresses the question “how does the client find out how to contact the server?” The difficulty, of course, is that there is no existing communication channel between them, yet they must somehow agree on a rendezvous point where they will establish communication.

1       Agreeing on a rendezvous point always involves a third party. The third party may  
2       itself provide the rendezvous point or may communicate rendezvous information from server  
3       to client. Complicating matters might be the fact that a client doesn't really care what  
4       server it contacts, only that it be able to get in touch with one that can handle its request.

5       Ideally, MPI can accommodate a wide variety of run-time systems while retaining the  
6       ability to write simple portable code. The following should be compatible with MPI:

- 7       • The server resides at a well-known internet address host:port.
- 8       • The server prints out an address to the terminal, the user gives this address to the  
9       client program.
- 10      • The server places the address information on a nameserver, where it can be retrieved  
11      with an agreed-upon name.
- 12      • The server to which the client connects is actually a broker, acting as a middleman  
13      between the client and the real server.

14       MPI does not require a nameserver, so not all implementations will be able to support  
15       all of the above scenarios. However, MPI provides an optional nameserver interface, and is  
16       compatible with external name servers.

17       A `port_name` is a *system-supplied* string that encodes a low-level network address at  
18       which a server can be contacted. Typically this is an IP address and a port number, but  
19       an implementation is free to use any protocol. The server establishes a `port_name` with  
20       the `MPI_OPEN_PORT` routine. It accepts a connection to a given port with  
21       `MPI_COMM_ACCEPT`. A client uses `port_name` to connect to the server.

22       By itself, the `port_name` mechanism is completely portable, but it may be clumsy  
23       to use because of the necessity to communicate `port_name` to the client. It would be more  
24       convenient if a server could specify that it be known by an *application-supplied* `service_name`  
25       so that the client could connect to that `service_name` without knowing the `port_name`.

26       An MPI implementation may allow the server to publish a (`port_name`, `service_name`)  
27       pair with `MPI_PUBLISH_NAME` and the client to retrieve the port name from the service  
28       name with `MPI_LOOKUP_NAME`. This allows three levels of portability, with increasing  
29       levels of functionality.

- 30      1. Applications that do not rely on the ability to publish names are the most portable.  
31      Typically the `port_name` must be transferred “by hand” from server to client.
- 32      2. Applications that use the `MPI_PUBLISH_NAME` mechanism are completely portable  
33      among implementations that provide this service. To be portable among all imple-  
34      mentations, these applications should have a fall-back mechanism that can be used  
35      when names are not published.
- 36      3. Applications may ignore MPI's name publishing functionality and use their own mech-  
37      anism (possibly system-supplied) to publish names. This allows arbitrary flexibility  
38      but is not portable.

## 10.4.2 Server Routines

A server makes itself available with two routines. First it must call `MPI_OPEN_PORT` to establish a port at which it may be contacted. Secondly it must call `MPI_COMM_ACCEPT` to accept connections from clients.

`MPI_OPEN_PORT(info, port_name)`

IN	info	implementation-specific information on how to establish an address (handle)
OUT	port_name	newly established port (string)

`int MPI_Open_port(MPI_Info info, char *port_name)`

```
MPI_Open_port(info, port_name, ierror) BIND(C)
  TYPE(MPI_Info), INTENT(IN) :: info
  CHARACTER(LEN=MPI_MAX_PORT_NAME), INTENT(OUT) :: port_name
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_OPEN_PORT(INFO, PORT_NAME, IERROR)
  CHARACTER*(*) PORT_NAME
  INTEGER INFO, IERROR
```

```
{void MPI::Open_port(const MPI::Info& info, char* port_name) (binding
  deprecated, see Section 15.2) }
```

This function establishes a network address, encoded in the `port_name` string, at which the server will be able to accept connections from clients. `port_name` is supplied by the system, possibly using information in the `info` argument.

MPI copies a system-supplied port name into `port_name`. `port_name` identifies the newly opened port and can be used by a client to contact the server. The maximum size string that may be supplied by the system is `MPI_MAX_PORT_NAME`.

*Advice to users.* The system copies the port name into `port_name`. The application must pass a buffer of sufficient size to hold this value. (*End of advice to users.*)

`port_name` is essentially a network address. It is unique within the communication universe to which it belongs (determined by the implementation), and may be used by any client within that communication universe. For instance, if it is an internet (host:port) address, it will be unique on the internet. If it is a low level switch address on an IBM SP, it will be unique to that SP.

*Advice to implementors.* These examples are not meant to constrain implementations. A `port_name` could, for instance, contain a user name or the name of a batch job, as long as it is unique within some well-defined communication domain. The larger the communication domain, the more useful MPI's client/server functionality will be. (*End of advice to implementors.*)

The precise form of the address is implementation-defined. For instance, an internet address may be a host name or IP address, or anything that the implementation can decode into an IP address. A port name may be reused after it is freed with `MPI_CLOSE_PORT` and released by the system.

*Advice to implementors.* Since the user may type in `port_name` by hand, it is useful to choose a form that is easily readable and does not have embedded spaces. (*End of advice to implementors.*)

`info` may be used to tell the implementation how to establish the address. It may, and usually will, be `MPI_INFO_NULL` in order to get the implementation defaults.

`MPI_CLOSE_PORT(port_name)`

IN            `port_name`                            a port (string)

int `MPI_Close_port(char *port_name)`

`MPI_Close_port(port_name, ierror) BIND(C)`  
     `CHARACTER(LEN=*)`, INTENT(IN) :: `port_name`  
     INTEGER, OPTIONAL, INTENT(OUT) :: `ierror`

`MPI_CLOSE_PORT(PORT_NAME, IERROR)`  
     `CHARACTER*(*)` `PORT_NAME`  
     INTEGER `IERROR`

{void `MPI::Close_port(const char* port_name)` (*binding deprecated, see Section 15.2*)  
     }

This function releases the network address represented by `port_name`.

`MPI_COMM_ACCEPT(port_name, info, root, comm, newcomm)`

IN            `port_name`                            port name (string, used only on root)  
 IN            `info`                                implementation-dependent information (handle, used only on root)  
 IN            `root`                                rank in `comm` of root node (integer)  
 IN            `comm`                                intracommunicator over which call is collective (handle)  
 OUT          `newcomm`                            intercommunicator with client as remote group (handle)

int `MPI_Comm_accept(char *port_name, MPI_Info info, int root,`  
     `MPI_Comm comm, MPI_Comm *newcomm)`

`MPI_Comm_accept(port_name, info, root, comm, newcomm, ierror) BIND(C)`  
     `CHARACTER(LEN=*)`, INTENT(IN) :: `port_name`  
     `TYPE(MPI_Info)`, INTENT(IN) :: `info`  
     INTEGER, INTENT(IN) :: `root`  
     `TYPE(MPI_Comm)`, INTENT(IN) :: `comm`  
     `TYPE(MPI_Comm)`, INTENT(OUT) :: `newcomm`  
     INTEGER, OPTIONAL, INTENT(OUT) :: `ierror`

`MPI_COMM_ACCEPT(PORT_NAME, INFO, ROOT, COMM, NEWCOMM, IERROR)`

```

    CHARACTER*(*) PORT_NAME
    INTEGER INFO, ROOT, COMM, NEWCOMM, IERROR
{MPI::Intercomm MPI::Intracomm::Accept(const char* port_name,
    const MPI::Info& info, int root) const(binding deprecated, see
    Section 15.2) }
```

MPI\_COMM\_ACCEPT establishes communication with a client. It is collective over the calling communicator. It returns an intercommunicator that allows communication with the client.

The `port_name` must have been established through a call to `MPI_OPEN_PORT`.

`info` is a implementation-defined string that may allow fine control over the `ACCEPT` call.

### 10.4.3 Client Routines

There is only one routine on the client side.

`MPI_COMM_CONNECT(port_name, info, root, comm, newcomm)`

IN	<code>port_name</code>	network address (string, used only on root)
IN	<code>info</code>	implementation-dependent information (handle, used only on root)
IN	<code>root</code>	rank in <code>comm</code> of root node (integer)
IN	<code>comm</code>	intracommunicator over which call is collective (handle)
OUT	<code>newcomm</code>	intercommunicator with server as remote group (handle)

```

int MPI_Comm_connect(char *port_name, MPI_Info info, int root,
    MPI_Comm comm, MPI_Comm *newcomm)
```

`MPI_Comm_connect(port_name, info, root, comm, newcomm, ierror)` **BIND(C)**

```

    CHARACTER(LEN=*), INTENT(IN) :: port_name
    TYPE(MPI_Info), INTENT(IN) :: info
    INTEGER, INTENT(IN) :: root
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Comm), INTENT(OUT) :: newcomm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

`MPI_COMM_CONNECT(PORT_NAME, INFO, ROOT, COMM, NEWCOMM, IERROR)`

```

    CHARACTER*(*) PORT_NAME
    INTEGER INFO, ROOT, COMM, NEWCOMM, IERROR
```

```

{MPI::Intercomm MPI::Intracomm::Connect(const char* port_name,
    const MPI::Info& info, int root) const(binding deprecated, see
    Section 15.2) }
```

This routine establishes communication with a server specified by `port_name`. It is collective over the calling communicator and returns an intercommunicator in which the remote group participated in an `MPI_COMM_ACCEPT`.

If the named port does not exist (or has been closed), `MPI_COMM_CONNECT` raises an error of class `MPI_ERR_PORT`.

If the port exists, but does not have a pending `MPI_COMM_ACCEPT`, the connection attempt will eventually time out after an implementation-defined time, or succeed when the server calls `MPI_COMM_ACCEPT`. In the case of a time out, `MPI_COMM_CONNECT` raises an error of class `MPI_ERR_PORT`.

*Advice to implementors.* The time out period may be arbitrarily short or long. However, a high quality implementation will try to queue connection attempts so that a server can handle simultaneous requests from several clients. A high quality implementation may also provide a mechanism, through the `info` arguments to `MPI_OPEN_PORT`, `MPI_COMM_ACCEPT` and/or `MPI_COMM_CONNECT`, for the user to control timeout and queuing behavior. (*End of advice to implementors.*)

MPI provides no guarantee of fairness in servicing connection attempts. That is, connection attempts are not necessarily satisfied in the order they were initiated and competition from other connection attempts may prevent a particular connection attempt from being satisfied.

`port_name` is the address of the server. It must be the same as the name returned by `MPI_OPEN_PORT` on the server. Some freedom is allowed here. If there are equivalent forms of `port_name`, an implementation may accept them as well. For instance, if `port_name` is (hostname:port), an implementation may accept (ip\_address:port) as well.

#### 10.4.4 Name Publishing

The routines in this section provide a mechanism for publishing names. A (`service_name`, `port_name`) pair is published by the server, and may be retrieved by a client using the `service_name` only. An MPI implementation defines the *scope* of the `service_name`, that is, the domain over which the `service_name` can be retrieved. If the domain is the empty set, that is, if no client can retrieve the information, then we say that name publishing is not supported. Implementations should document how the scope is determined. High-quality implementations will give some control to users through the `info` arguments to name publishing functions. Examples are given in the descriptions of individual functions.

`MPI_PUBLISH_NAME(service_name, info, port_name)`

IN	<code>service_name</code>	a service name to associate with the port (string)
IN	<code>info</code>	implementation-specific information (handle)
IN	<code>port_name</code>	a port name (string)

`int MPI_Publish_name(char *service_name, MPI_Info info, char *port_name)`

```

MPI_Publish_name(service_name, info, port_name, ierror) BIND(C)
    TYPE(MPI_Info), INTENT(IN) :: info
    CHARACTER(LEN=*), INTENT(IN) :: service_name, port_name

```

```

    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_PUBLISH_NAME(SERVICE_NAME, INFO, PORT_NAME, IERROR)
    INTEGER INFO, IERROR
    CHARACTER*(*) SERVICE_NAME, PORT_NAME

{void MPI::Publish_name(const char* service_name, const MPI::Info& info,
    const char* port_name) (binding deprecated, see Section 15.2) }
```

This routine publishes the pair (port\_name, service\_name) so that an application may retrieve a system-supplied port\_name using a well-known service\_name.

The implementation must define the *scope* of a published service name, that is, the domain over which the service name is unique, and conversely, the domain over which the (port name, service name) pair may be retrieved. For instance, a service name may be unique to a job (where job is defined by a distributed operating system or batch scheduler), unique to a machine, or unique to a Kerberos realm. The scope may depend on the info argument to MPI\_PUBLISH\_NAME.

MPI permits publishing more than one service\_name for a single port\_name. On the other hand, if service\_name has already been published within the scope determined by info, the behavior of MPI\_PUBLISH\_NAME is undefined. An MPI implementation may, through a mechanism in the info argument to MPI\_PUBLISH\_NAME, provide a way to allow multiple servers with the same service in the same scope. In this case, an implementation-defined policy will determine which of several port names is returned by MPI\_LOOKUP\_NAME.

Note that while service\_name has a limited scope, determined by the implementation, port\_name always has global scope within the communication universe used by the implementation (i.e., it is globally unique).

port\_name should be the name of a port established by MPI\_OPEN\_PORT and not yet deleted by MPI\_CLOSE\_PORT. If it is not, the result is undefined.

*Advice to implementors.* In some cases, an MPI implementation may use a name service that a user can also access directly. In this case, a name published by MPI could easily conflict with a name published by a user. In order to avoid such conflicts, MPI implementations should mangle service names so that they are unlikely to conflict with user code that makes use of the same service. Such name mangling will of course be completely transparent to the user.

The following situation is problematic but unavoidable, if we want to allow implementations to use nameservers. Suppose there are multiple instances of “ocean” running on a machine. If the scope of a service name is confined to a job, then multiple oceans can coexist. If an implementation provides site-wide scope, however, multiple instances are not possible as all calls to MPI\_PUBLISH\_NAME after the first may fail. There is no universal solution to this.

To handle these situations, a high-quality implementation should make it possible to limit the domain over which names are published. (*End of advice to implementors.*)

```

1 MPI_UNPUBLISH_NAME(service_name, info, port_name)
2     IN      service_name      a service name (string)
3
4     IN      info              implementation-specific information (handle)
5
6     IN      port_name         a port name (string)

```

```

7 int MPI_Unpublish_name(char *service_name, MPI_Info info, char *port_name)

```

```

8 MPI_Unpublish_name(service_name, info, port_name, ierror) BIND(C)
9     CHARACTER(LEN=*), INTENT(IN) :: service_name, port_name
10    TYPE(MPI_Info), INTENT(IN) :: info
11    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

12 MPI_UNPUBLISH_NAME(SERVICE_NAME, INFO, PORT_NAME, IERROR)
13     INTEGER INFO, IERROR
14     CHARACTER*(*) SERVICE_NAME, PORT_NAME

```

```

15 {void MPI::Unpublish_name(const char* service_name, const MPI::Info& info,
16     const char* port_name) (binding deprecated, see Section 15.2) }

```

This routine unpublishes a service name that has been previously published. Attempting to unpublish a name that has not been published or has already been unpublished is erroneous and is indicated by the error class MPI\_ERR\_SERVICE.

All published names must be unpublished before the corresponding port is closed and before the publishing process exits. The behavior of MPI\_UNPUBLISH\_NAME is implementation dependent when a process tries to unpublish a name that it did not publish.

If the info argument was used with MPI\_PUBLISH\_NAME to tell the implementation how to publish names, the implementation may require that info passed to MPI\_UNPUBLISH\_NAME contain information to tell the implementation how to unpublish a name.

```

31 MPI_LOOKUP_NAME(service_name, info, port_name)

```

```

32     IN      service_name      a service name (string)
33
34     IN      info              implementation-specific information (handle)
35
36     OUT     port_name         a port name (string)

```

```

37 int MPI_Lookup_name(char *service_name, MPI_Info info, char *port_name)

```

```

38 MPI_Lookup_name(service_name, info, port_name, ierror) BIND(C)
39     CHARACTER(LEN=*), INTENT(IN) :: service_name
40     TYPE(MPI_Info), INTENT(IN) :: info
41     CHARACTER(LEN=MPI_MAX_PORT_NAME), INTENT(OUT) :: port_name
42     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

43 MPI_LOOKUP_NAME(SERVICE_NAME, INFO, PORT_NAME, IERROR)
44     CHARACTER*(*) SERVICE_NAME, PORT_NAME
45     INTEGER INFO, IERROR

```



```
{void MPI::Lookup_name(const char* service_name, const MPI::Info& info,
                        char* port_name) (binding deprecated, see Section 15.2) }
```

This function retrieves a `port_name` published by `MPI_PUBLISH_NAME` with `service_name`. If `service_name` has not been published, it raises an error in the error class `MPI_ERR_NAME`. The application must supply a `port_name` buffer large enough to hold the largest possible port name (see discussion above under `MPI_OPEN_PORT`).

If an implementation allows multiple entries with the same `service_name` within the same scope, a particular `port_name` is chosen in a way determined by the implementation.

If the `info` argument was used with `MPI_PUBLISH_NAME` to tell the implementation how to publish names, a similar `info` argument may be required for `MPI_LOOKUP_NAME`.

#### 10.4.5 Reserved Key Values

The following key values are reserved. An implementation is not required to interpret these key values, but if it does interpret the key value, it must provide the functionality described.

`ip_port` Value contains IP port number at which to establish a port. (Reserved for `MPI_OPEN_PORT` only).

`ip_address` Value contains IP address at which to establish a port. If the address is not a valid IP address of the host on which the `MPI_OPEN_PORT` call is made, the results are undefined. (Reserved for `MPI_OPEN_PORT` only).

#### 10.4.6 Client/Server Examples

Simplest Example — Completely Portable.

The following example shows the simplest way to use the client/server interface. It does not use service names at all.

On the server side:

```
char myport[MPI_MAX_PORT_NAME];
MPI_Comm intercomm;
/* ... */
MPI_Open_port(MPI_INFO_NULL, myport);
printf("port name is: %s\n", myport);

MPI_Comm_accept(myport, MPI_INFO_NULL, 0, MPI_COMM_SELF, &intercomm);
/* do something with intercomm */
```

The server prints out the port name to the terminal and the user must type it in when starting up the client (assuming the MPI implementation supports stdin such that this works). On the client side:

```
MPI_Comm intercomm;
char name[MPI_MAX_PORT_NAME];
printf("enter port name: ");
gets(name);
MPI_Comm_connect(name, MPI_INFO_NULL, 0, MPI_COMM_SELF, &intercomm);
```

### 1 Ocean/Atmosphere - Relies on Name Publishing

2 In this example, the “ocean” application is the “server” side of a coupled ocean-atmosphere  
3 climate model. It assumes that the MPI implementation publishes names.  
4

```
5
6     MPI_Open_port(MPI_INFO_NULL, port_name);
7     MPI_Publish_name("ocean", MPI_INFO_NULL, port_name);
8
9     MPI_Comm_accept(port_name, MPI_INFO_NULL, 0, MPI_COMM_SELF, &intercomm);
10    /* do something with intercomm */
11    MPI_Unpublish_name("ocean", MPI_INFO_NULL, port_name);
12
```

13  
14 On the client side:

```
15
16     MPI_Lookup_name("ocean", MPI_INFO_NULL, port_name);
17     MPI_Comm_connect( port_name, MPI_INFO_NULL, 0, MPI_COMM_SELF,
18                      &intercomm);
19
```

### 20 Simple Client-Server Example.

21 This is a simple example; the server accepts only a single connection at a time and serves  
22 that connection until the client requests to be disconnected. The server is a single process.  
23

24 Here is the server. It accepts a single connection and then processes data until it  
25 receives a message with tag 1. A message with tag 0 tells the server to exit.

```
26 #include "mpi.h"
27 int main( int argc, char **argv )
28 {
29     MPI_Comm client;
30     MPI_Status status;
31     char port_name[MPI_MAX_PORT_NAME];
32     double buf[MAX_DATA];
33     int size, again;
34
35     MPI_Init( &argc, &argv );
36     MPI_Comm_size(MPI_COMM_WORLD, &size);
37     if (size != 1) error(FATAL, "Server too big");
38     MPI_Open_port(MPI_INFO_NULL, port_name);
39     printf("server available at %s\n",port_name);
40     while (1) {
41         MPI_Comm_accept( port_name, MPI_INFO_NULL, 0, MPI_COMM_WORLD,
42                         &client );
43         again = 1;
44         while (again) {
45             MPI_Recv( buf, MAX_DATA, MPI_DOUBLE,
46                     MPI_ANY_SOURCE, MPI_ANY_TAG, client, &status );
47             switch (status.MPI_TAG) {
48                 case 0: MPI_Comm_free( &client );
```

```

        MPI_Close_port(port_name);
        MPI_Finalize();
        return 0;
    case 1: MPI_Comm_disconnect( &client );
        again = 0;
        break;
    case 2: /* do something */
        ...
    default:
        /* Unexpected message type */
        MPI_Abort( MPI_COMM_WORLD, 1 );
    }
}
}
}

```

Here is the client.

```

#include "mpi.h"
int main( int argc, char **argv )
{
    MPI_Comm server;
    double buf[MAX_DATA];
    char port_name[MPI_MAX_PORT_NAME];

    MPI_Init( &argc, &argv );
    strcpy(port_name, argv[1] ); /* assume server's name is cmd-line arg */

    MPI_Comm_connect( port_name, MPI_INFO_NULL, 0, MPI_COMM_WORLD,
                      &server );

    while (!done) {
        tag = 2; /* Action to perform */
        MPI_Send( buf, n, MPI_DOUBLE, 0, tag, server );
        /* etc */
    }
    MPI_Send( buf, 0, MPI_DOUBLE, 0, 1, server );
    MPI_Comm_disconnect( &server );
    MPI_Finalize();
    return 0;
}

```

## 10.5 Other Functionality

### 10.5.1 Universe Size

Many “dynamic” MPI applications are expected to exist in a static runtime environment, in which resources have been allocated before the application is run. When a user (or

possibly a batch system) runs one of these quasi-static applications, she will usually specify a number of processes to start and a total number of processes that are expected. An application simply needs to know how many slots there are, i.e., how many processes it should spawn.

MPI provides an attribute on `MPI_COMM_WORLD`, `MPI_UNIVERSE_SIZE`, that allows the application to obtain this information in a portable manner. This attribute indicates the total number of processes that are expected. In Fortran, the attribute is the integer value. In C, the attribute is a pointer to the integer value. An application typically subtracts the size of `MPI_COMM_WORLD` from `MPI_UNIVERSE_SIZE` to find out how many processes it should spawn. `MPI_UNIVERSE_SIZE` is initialized in `MPI_INIT` and is not changed by MPI. If defined, it has the same value on all processes of `MPI_COMM_WORLD`. `MPI_UNIVERSE_SIZE` is determined by the application startup mechanism in a way not specified by MPI. (The size of `MPI_COMM_WORLD` is another example of such a parameter.)

Possibilities for how `MPI_UNIVERSE_SIZE` might be set include

- A `-universe_size` argument to a program that starts MPI processes.
- Automatic interaction with a batch scheduler to figure out how many processors have been allocated to an application.
- An environment variable set by the user.
- Extra information passed to `MPI_COMM_SPAWN` through the `info` argument.

An implementation must document how `MPI_UNIVERSE_SIZE` is set. An implementation may not support the ability to set `MPI_UNIVERSE_SIZE`, in which case the attribute `MPI_UNIVERSE_SIZE` is not set.

`MPI_UNIVERSE_SIZE` is a recommendation, not necessarily a hard limit. For instance, some implementations may allow an application to spawn 50 processes per processor, if they are requested. However, it is likely that the user only wants to spawn one process per processor.

`MPI_UNIVERSE_SIZE` is assumed to have been specified when an application was started, and is in essence a portable mechanism to allow the user to pass to the application (through the MPI process startup mechanism, such as `mpiexec`) a piece of critical runtime information. Note that no interaction with the runtime environment is required. If the runtime environment changes size while an application is running, `MPI_UNIVERSE_SIZE` is not updated, and the application must find out about the change through direct communication with the runtime system.

### 10.5.2 Singleton `MPI_INIT`

A high-quality implementation will allow any process (including those not started with a “parallel application” mechanism) to become an MPI process by calling `MPI_INIT`. Such a process can then connect to other MPI processes using the `MPI_COMM_ACCEPT` and `MPI_COMM_CONNECT` routines, or spawn other MPI processes. MPI does not mandate this behavior, but strongly encourages it where technically feasible.

*Advice to implementors.* To start MPI processes belonging to the same `MPI_COMM_WORLD` requires some special coordination. The processes must be started at the “same” time, they must have a mechanism to establish communication, etc.

Either the user or the operating system must take special steps beyond simply starting processes.

When an application enters `MPI_INIT`, clearly it must be able to determine if these special steps were taken. If a process enters `MPI_INIT` and determines that no special steps were taken (i.e., it has not been given the information to form an `MPI_COMM_WORLD` with other processes) it succeeds and forms a singleton MPI program, that is, one in which `MPI_COMM_WORLD` has size 1.

In some implementations, MPI may not be able to function without an “MPI environment.” For example, MPI may require that daemons be running or MPI may not be able to work at all on the front-end of an MPP. In this case, an MPI implementation may either

1. Create the environment (e.g., start a daemon) or
2. Raise an error if it cannot create the environment and the environment has not been started independently.

A high-quality implementation will try to create a singleton MPI process and not raise an error.

*(End of advice to implementors.)*

### 10.5.3 MPI\_APPNUM

There is a predefined attribute `MPI_APPNUM` of `MPI_COMM_WORLD`. In Fortran, the attribute is an integer value. In C, the attribute is a pointer to an integer value. If a process was spawned with `MPI_COMM_SPAWN_MULTIPLE`, `MPI_APPNUM` is the command number that generated the current process. Numbering starts from zero. If a process was spawned with `MPI_COMM_SPAWN`, it will have `MPI_APPNUM` equal to zero.

Additionally, if the process was not started by a spawn call, but by an implementation-specific startup mechanism that can handle multiple process specifications, `MPI_APPNUM` should be set to the number of the corresponding process specification. In particular, if it is started with

```
mpiexec spec0 [: spec1 : spec2 : ...]
```

`MPI_APPNUM` should be set to the number of the corresponding specification.

If an application was not spawned with `MPI_COMM_SPAWN` or `MPI_COMM_SPAWN_MULTIPLE`, and `MPI_APPNUM` doesn’t make sense in the context of the implementation-specific startup mechanism, `MPI_APPNUM` is not set.

MPI implementations may optionally provide a mechanism to override the value of `MPI_APPNUM` through the info argument. MPI reserves the following key for all `SPAWN` calls.

**appnum** Value contains an integer that overrides the default value for `MPI_APPNUM` in the child.

*Rationale.* When a single application is started, it is able to figure out how many processes there are by looking at the size of `MPI_COMM_WORLD`. An application consisting of multiple SPMD sub-applications has no way to find out how many sub-applications there are and to which sub-application the process belongs. While there are ways to

figure it out in special cases, there is no general mechanism. MPI\_APPNUM provides such a general mechanism. (*End of rationale.*)

#### 10.5.4 Releasing Connections

Before a client and server connect, they are independent MPI applications. An error in one does not affect the other. After establishing a connection with MPI\_COMM\_CONNECT and MPI\_COMM\_ACCEPT, an error in one may affect the other. It is desirable for a client and server to be able to disconnect, so that an error in one will not affect the other. Similarly, it might be desirable for a parent and child to disconnect, so that errors in the child do not affect the parent, or vice-versa.

- Two processes are **connected** if there is a communication path (direct or indirect) between them. More precisely:
  1. Two processes are connected if
    - (a) they both belong to the same communicator (inter- or intra-, including MPI\_COMM\_WORLD) *or*
    - (b) they have previously belonged to a communicator that was freed with MPI\_COMM\_FREE instead of MPI\_COMM\_DISCONNECT *or*
    - (c) they both belong to the group of the same window or filehandle.
  2. If A is connected to B and B to C, then A is connected to C.
- Two processes are **disconnected** (also **independent**) if they are not connected.
- By the above definitions, connectivity is a transitive property, and divides the universe of MPI processes into disconnected (independent) sets (equivalence classes) of processes.
- Processes which are connected, but don't share the same MPI\_COMM\_WORLD may become disconnected (independent) if the communication path between them is broken by using MPI\_COMM\_DISCONNECT.

The following additional rules apply to MPI routines in other chapters:

- MPI\_FINALIZE is collective over a set of connected processes.
- MPI\_ABORT does not abort independent processes. It may abort all processes in the caller's MPI\_COMM\_WORLD (ignoring its **comm** argument). Additionally, it may abort connected processes as well, though it makes a "best attempt" to abort only the processes in **comm**.
- If a process terminates without calling MPI\_FINALIZE, independent processes are not affected but the effect on connected processes is not defined.

MPI\_COMM\_DISCONNECT(comm)

INOUT     comm                             communicator (handle)

int MPI\_Comm\_disconnect(MPI\_Comm \*comm)

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```

MPI_Comm_disconnect(comm, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(INOUT) :: comm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_COMM_DISCONNECT(COMM, IERROR)
    INTEGER COMM, IERROR

```

```

{void MPI::Comm::Disconnect() (binding deprecated, see Section 15.2) }

```

This function waits for all pending communication on `comm` to complete internally, deallocates the communicator object, and sets the handle to `MPI_COMM_NULL`. It is a collective operation.

It may not be called with the communicator `MPI_COMM_WORLD` or `MPI_COMM_SELF`.

`MPI_COMM_DISCONNECT` may be called only if all communication is complete and matched, so that buffered data can be delivered to its destination. This requirement is the same as for `MPI_FINALIZE`.

`MPI_COMM_DISCONNECT` has the same action as `MPI_COMM_FREE`, except that it waits for pending communication to finish internally and enables the guarantee about the behavior of disconnected processes.

*Advice to users.* To disconnect two processes you may need to call `MPI_COMM_DISCONNECT`, `MPI_WIN_FREE` and `MPI_FILE_CLOSE` to remove all communication paths between the two processes. Notes that it may be necessary to disconnect several communicators (or to free several windows or files) before two processes are completely independent. (*End of advice to users.*)

*Rationale.* It would be nice to be able to use `MPI_COMM_FREE` instead, but that function explicitly does not wait for pending communication to complete. (*End of rationale.*)

### 10.5.5 Another Way to Establish MPI Communication

```

MPI_COMM_JOIN(fd, intercomm)

```

IN	fd	socket file descriptor
OUT	intercomm	new intercommunicator (handle)

```

int MPI_Comm_join(int fd, MPI_Comm *intercomm)

```

```

MPI_Comm_join(fd, intercomm, ierror) BIND(C)
    INTEGER, INTENT(IN) :: fd
    TYPE(MPI_Comm), INTENT(OUT) :: intercomm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_COMM_JOIN(FD, INTERCOMM, IERROR)
    INTEGER FD, INTERCOMM, IERROR

```

```

{static MPI::Intercomm MPI::Comm::Join(const int fd) (binding deprecated, see
    Section 15.2) }

```

MPI\_COMM\_JOIN is intended for MPI implementations that exist in an environment supporting the Berkeley Socket interface [45, 49]. Implementations that exist in an environment not supporting Berkeley Sockets should provide the entry point for MPI\_COMM\_JOIN and should return MPI\_COMM\_NULL.

This call creates an intercommunicator from the union of two MPI processes which are connected by a socket. MPI\_COMM\_JOIN should normally succeed if the local and remote processes have access to the same implementation-defined MPI communication universe.

*Advice to users.* An MPI implementation may require a specific communication medium for MPI communication, such as a shared memory segment or a special switch. In this case, it may not be possible for two processes to successfully join even if there is a socket connecting them and they are using the same MPI implementation. (*End of advice to users.*)

*Advice to implementors.* A high-quality implementation will attempt to establish communication over a slow medium if its preferred one is not available. If implementations do not do this, they must document why they cannot do MPI communication over the medium used by the socket (especially if the socket is a TCP connection). (*End of advice to implementors.*)

fd is a file descriptor representing a socket of type SOCK\_STREAM (a two-way reliable byte-stream connection). Nonblocking I/O and asynchronous notification via SIGIO must not be enabled for the socket. The socket must be in a connected state. The socket must be quiescent when MPI\_COMM\_JOIN is called (see below). It is the responsibility of the application to create the socket using standard socket API calls.

MPI\_COMM\_JOIN must be called by the process at each end of the socket. It does not return until both processes have called MPI\_COMM\_JOIN. The two processes are referred to as the local and remote processes.

MPI uses the socket to bootstrap creation of the intercommunicator, and for nothing else. Upon return from MPI\_COMM\_JOIN, the file descriptor will be open and quiescent (see below).

If MPI is unable to create an intercommunicator, but is able to leave the socket in its original state, with no pending communication, it succeeds and sets intercomm to MPI\_COMM\_NULL.

The socket must be quiescent before MPI\_COMM\_JOIN is called and after MPI\_COMM\_JOIN returns. More specifically, on entry to MPI\_COMM\_JOIN, a read on the socket will not read any data that was written to the socket before the remote process called MPI\_COMM\_JOIN. On exit from MPI\_COMM\_JOIN, a read will not read any data that was written to the socket before the remote process returned from MPI\_COMM\_JOIN. It is the responsibility of the application to ensure the first condition, and the responsibility of the MPI implementation to ensure the second. In a multithreaded application, the application must ensure that one thread does not access the socket while another is calling MPI\_COMM\_JOIN, or call MPI\_COMM\_JOIN concurrently.

*Advice to implementors.* MPI is free to use any available communication path(s) for MPI messages in the new communicator; the socket is only used for the initial handshaking. (*End of advice to implementors.*)



MPI\_COMM\_JOIN uses non-MPI communication to do its work. The interaction of non-MPI communication with pending MPI communication is not defined. Therefore, the result of calling MPI\_COMM\_JOIN on two connected processes (see Section 10.5.4 on page 414 for the definition of connected) is undefined.

The returned communicator may be used to establish MPI communication with additional processes, through the usual MPI communicator creation mechanisms.

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# Chapter 11

## One-Sided Communications

### 11.1 Introduction

Remote Memory Access (RMA) extends the communication mechanisms of MPI by allowing one process to specify all communication parameters, both for the sending side and for the receiving side. This mode of communication facilitates the coding of some applications with dynamically changing data access patterns where the data distribution is fixed or slowly changing. In such a case, each process can compute what data it needs to access or to update at other processes. [However, processes may not know which data in their own memory need to be accessed or to be updated by remote processes, and may not even know the identity of these processes.] However, the programmer may not be able to easily determine which data in a process may need to be accessed or to be updated by operations executed by a different process, and may not even know which processes may perform such updates. Thus, the transfer parameters are all available only on one side. Regular send/receive communication requires matching operations by sender and receiver. In order to issue the matching operations, an application needs to distribute the transfer parameters. This distribution may require all processes to participate in a time-consuming global computation, or to [periodically poll for potential communication requests to receive and act upon]poll for potential communication requests to receive and upon which to act periodically. The use of RMA communication mechanisms avoids the need for global computations or explicit polling. A generic example of this nature is the execution of an assignment of the form  $A = B(\text{map})$ , where  $\text{map}$  is a permutation vector, and  $A$ ,  $B$  and  $\text{map}$  are distributed in the same manner.

Message-passing communication achieves two effects: *communication* of data from sender to receiver; and *synchronization* of sender with receiver. The RMA design separates these two functions. [Three communication calls are provided: `MPI_PUT` (remote write), `MPI_GET` (remote read) and `MPI_ACCUMULATE` (remote update). A larger number of synchronization calls are provided that support different synchronization styles. The design is similar to that of weakly coherent memory systems: correct ordering of memory accesses has to be imposed by the user, using synchronization calls; the implementation can delay communication operations until the synchronization calls occur, for efficiency.] The following communication calls are provided:

- Remote write: `MPI_PUT`, `MPI_RPUT`
- Remote read: `MPI_GET`, `MPI_RGET`

- Remote update: `MPI_ACCUMULATE`, `MPI_RACCUMULATE`
- Remote read and update: `MPI_GET_ACCUMULATE`, `MPI_RGET_ACCUMULATE`, and `MPI_FETCH_AND_OP`
- Remote atomic swap operations: `MPI_COMPARE_AND_SWAP`

This chapter refers to an operations set that includes all remote update, remote read and update, and remote atomic swap operations as “accumulate” operations.

MPI supports two fundamentally different memory models: separate and unified. The first model makes no assumption about memory consistency and is highly portable. This model is similar to that of weakly coherent memory systems: the user must impose correct ordering of memory accesses through synchronization calls[; for efficiency, the implementation can delay communication operations until the synchronization calls occur]. The second model can exploit cache-coherent hardware and hardware-accelerated one-sided operations that are commonly available in high-performance systems. [In this model, communication can be independent of synchronization calls.] The two different models are discussed in detail in Section 11.4. Both models support a large number of synchronization calls to support different synchronization styles.

The design of the RMA functions allows implementors to take advantage [ , in many cases,] of fast or asynchronous communication mechanisms provided by various platforms, such as coherent or noncoherent shared memory, DMA engines, hardware-supported put/get operations, and communication coprocessors[ , etc]. The most frequently used RMA communication mechanisms can be layered on top of message-passing. [However, support for asynchronous communication agents in software (handlers, threads, etc.) is needed, for certain RMA functions, in a distributed memory environment.] However, certain RMA functions might need support for asynchronous communication agents in software (handlers, threads, etc.) in a distributed memory environment.

We shall denote by **origin** the process that performs the call, and by **target** the process in which the memory is accessed. Thus, in a put operation, source=origin and destination=target; in a get operation, source=target and destination=origin.

## 11.2 Initialization

[The initialization operation]MPI provides [three]the following window initialization functions, `MPI_WIN_CREATE`, `MPI_WIN_ALLOCATE`, `MPI_WIN_ALLOCATE_SHARED` and `MPI_WIN_CREATE_DYNAMIC` that are collective on an intracommunicator.

`MPI_WIN_CREATE` allows each process [in an intracommunicator group] to specify [ , in a collective operation,] a “window” in its memory that is made accessible to accesses by remote processes. The call returns an opaque object that represents the group of processes that own and access the set of windows, and the attributes of each window, as specified by the initialization call. `MPI_WIN_ALLOCATE` differs from `MPI_WIN_CREATE` in that the user does not pass allocated memory; `MPI_WIN_ALLOCATE` returns a pointer to memory allocated by the MPI implementation. `MPI_WIN_ALLOCATE_SHARED` differs from `MPI_WIN_ALLOCATE` in that the allocated memory can be accessed from all processes in the window’s group with direct load/store instructions. Some restrictions apply to the specified communicator. `MPI_WIN_CREATE_DYNAMIC` creates a window that allows the user to dynamically control which memory is exposed by the window.

## 11.2.1 Window Creation

MPI\_WIN\_CREATE(base, size, disp\_unit, info, comm, win)

IN	base	initial address of window (choice)
IN	size	size of window in bytes (non-negative integer)
IN	disp_unit	local unit size for displacements, in bytes (positive integer)
IN	info	info argument (handle)
IN	comm	intra-communicator (handle)
OUT	win	window object returned by the call (handle)

```
int MPI_Win_create(void *base, MPI_Aint size, int disp_unit, MPI_Info info,
                  MPI_Comm comm, MPI_Win *win)
```

```
MPI_Win_create(base, size, disp_unit, info, comm, win, ierror) BIND(C)
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: base
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: size
    INTEGER, INTENT(IN) :: disp_unit
    TYPE(MPI_Info), INTENT(IN) :: info
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Win), INTENT(OUT) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_WIN_CREATE(BASE, SIZE, DISP_UNIT, INFO, COMM, WIN, IERROR)
    <type> BASE(*)
    INTEGER(KIND=MPI_ADDRESS_KIND) SIZE
    INTEGER DISP_UNIT, INFO, COMM, WIN, IERROR
```

```
{static MPI::Win MPI::Win::Create(const void* base, MPI::Aint size, int
    disp_unit, const MPI::Info& info, const MPI::Intracomm&
    comm) (binding deprecated, see Section 15.2) }
```

This is a collective call executed by all processes in the group of `comm`. It returns a window object that can be used by these processes to perform RMA operations. Each process specifies a window of existing memory that it exposes to RMA accesses by the processes in the group of `comm`. The window consists of `size` bytes, starting at address `base`. In C and C++, `base` is the starting address of a memory region. In Fortran, one can pass the first element of a memory region or a whole array, which must be ‘simply contiguous’ (for ‘simply contiguous’, see also Section 16.2.12 on page 671). A process may elect to expose no memory by specifying `size = 0`.

The displacement unit argument is provided to facilitate address arithmetic in RMA operations: the target displacement argument of an RMA operation is scaled by the factor `disp_unit` specified by the target process, at window creation.

*Rationale.* The window size is specified using an address sized integer, so as to allow windows that span more than 4 GB of address space. (Even if the physical memory

size is less than 4 GB, the address range may be larger than 4 GB, if addresses are not contiguous.) (*End of rationale.*)

*Advice to users.* Common choices for `disp_unit` are 1 (no scaling), and (in C syntax) `sizeof(type)`, for a window that consists of an array of elements of type `type`. The later choice will allow one to use array indices in RMA calls, and have those scaled correctly to byte displacements, even in a heterogeneous environment. (*End of advice to users.*)

The `info` argument provides optimization hints to the runtime about the expected usage pattern of the window. The following info key[is]s are predefined:

`no_locks` — if set to true, then the implementation may assume that ~~the local window is never locked (by a call to `MPI_WIN_LOCK` or `MPI_WIN_LOCK_ALL`).~~ This implies that this window is not used for 3-party communication, and RMA can be implemented with no (less) asynchronous agent activity at this process.

`accumulate_ordering` — controls the ordering of accumulate operations at the target. See Section 11.7.2 for details.

`accumulate_ops` — if set to `same_op`, the implementation will assume that all concurrent accumulate calls to the same target address will use the same operation. If set to `same_op_no_op`, then the implementation will assume that all concurrent accumulate calls to the same target address will use the same operation or `MPI_NO_OP`. This can eliminate the need to protect access for certain operation types where the hardware can guarantee atomicity. The default is `same_op_no_op`.

*Advice to users.* If windows are passed to libraries, the user ~~needs to~~ ensure that the info keys specified at window creation are communicated to the called library, which might need to constrain the operations on the passed window. (*End of advice to users.*)

The various processes in the group of `comm` may specify completely different target windows, in location, size, displacement units and info arguments. As long as all the get, put and accumulate accesses to a particular process fit their specific target window this should pose no problem. The same area in memory may appear in multiple windows, each associated with a different window object. However, concurrent communications to distinct, overlapping windows may lead to [erroneous]undefined results.

*Rationale.* The reason for specifying the memory that may be accessed from another process in an RMA operation is to permit the programmer to specify what memory can be a target of RMA operations and for the implementation to enforce that specification. For example, with this definition, a server process can safely allow a client process to use RMA operations, knowing that (under the assumption that the MPI implementation does enforce the specified limits on the exposed memory) an error in the client cannot affect any memory other than what was explicitly exposed. (*End of rationale.*)

*Advice to users.* A window can be created in any part of the process memory. However, on some systems, the performance of windows in memory allocated by MPI\_ALLOC\_MEM (Section 8.2, page 352) will be better. Also, on some systems, performance is improved when window boundaries are aligned at “natural” boundaries (word, double-word, cache line, page frame, etc.). (*End of advice to users.*)

*Advice to implementors.* In cases where RMA operations use different mechanisms in different memory areas (e.g., load/store in a shared memory segment, and an asynchronous handler in private memory), the MPI\_WIN\_CREATE call needs to figure out which type of memory is used for the window. To do so, MPI maintains, internally, the list of memory segments allocated by MPI\_ALLOC\_MEM, or by other, implementation specific, mechanisms, together with information on the type of memory segment allocated. When a call to MPI\_WIN\_CREATE occurs, then MPI checks which segment contains each window, and decides, accordingly, which mechanism to use for RMA operations.

Vendors may provide additional, implementation-specific mechanisms to allocate or to specify memory regions that are preferable for use in one-sided communication. In particular, such mechanisms can be used to place static variables into such preferred regions.

Implementors should document any performance impact of window alignment. (*End of advice to implementors.*)

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### 11.2.2 Window That Allocates Memory

**MPI\_WIN\_ALLOCATE(size, disp\_unit, info, comm, baseptr, win)**

IN	size	size of window in bytes (non-negative integer)
IN	disp_unit	local unit size for displacements, in bytes (positive integer)
IN	info	info argument (handle)
IN	comm	intra-communicator (handle)
OUT	baseptr	initial address of window (choice)
OUT	win	window object returned by the call (handle)

```
int MPI_Win_allocate(MPI_Aint size, int disp_unit, MPI_Info info,
                    MPI_Comm comm, void *baseptr, MPI_Win *win)
```

```
MPI_Win_allocate(size, disp_unit, info, comm, baseptr, win, ierror) BIND(C)
USE, INTRINSIC :: ISO_C_BINDING, ONLY : C_PTR
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: size
INTEGER, INTENT(IN) :: disp_unit
TYPE(MPI_Info), INTENT(IN) :: info
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(C_PTR), INTENT(OUT) :: baseptr
```

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```

1      TYPE(MPI_Win), INTENT(OUT) :: win
2      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
3
4      MPI_WIN_ALLOCATE(SIZE, DISP_UNIT, INFO, COMM, BASEPTR, WIN, IERROR)
5      INTEGER DISP_UNIT, INFO, COMM, WIN, IERROR
6      INTEGER(KIND=MPI_ADDRESS_KIND) SIZE, BASEPTR

```

This is a collective call executed by all processes in the group of `comm`. On each process, it allocates memory of at least `size` bytes, returns a pointer to it, and returns a window object that can be used by all processes in `comm` to perform RMA operations. The returned memory consists of `size` bytes local to each process, starting at address `baseptr` and is associated with the window as if the user called `MPI_WIN_CREATE` on existing memory. The size argument may be different at each process and `size = 0` is valid; however, a library might allocate and expose more memory in order to create a fast, globally symmetric allocation. The discussion of and rationales for `MPI_ALLOC_MEM` and `MPI_FREE_MEM` in Section 8.2 also apply to `MPI_WIN_ALLOCATE`; in particular, see the rationale in Section 8.2 for an explanation of the type used for `baseptr`.

*Rationale.* By allocating (potentially aligned) memory instead of allowing the user to pass in an arbitrary buffer, this call can improve the performance for systems with remote direct memory access significantly. This also permits the collective allocation of memory and supports what is sometimes called the “symmetric allocation” model that can be more scalable (for example, the implementation can arrange to return an address for the allocated memory that is the same on all processes). (*End of rationale.*)

The `info` argument can be used to specify hints similar to the `info` argument for `MPI_WIN_CREATE` and `MPI_ALLOC_MEM`. The following `info` key is predefined:

`same_size` — if set to true, then the implementation may assume that the argument `size` is identical on all processes.

### 11.2.3 Window That Allocates Shared Memory

```

36      MPI_WIN_ALLOCATE_SHARED(size, info, comm, baseptr, win)

```

IN	size	size of local window in bytes (non-negative integer)
IN	info	info argument (handle)
IN	comm	intra-communicator (handle)
OUT	baseptr	address of local allocated window segment (choice)
OUT	win	window object returned by the call (handle)

```

44      int MPI_Win_allocate_shared(MPI_Aint size, MPI_Info info, MPI_Comm comm,
45                                void *baseptr, MPI_Win *win)

```

```

47      MPI_WIN_ALLOCATE_SHARED(SIZE, INFO, COMM, BASEPTR, WIN, IERROR)
48      INTEGER SIZE, INFO, COMM, WIN, IERROR

```



INTEGER(KIND=MPI\_ADDRESS\_KIND) SIZE, BASEPTR

This is a collective call executed by all processes in the group of `comm`. On each process  $i$ , it allocates memory of at least `size` bytes that is shared among all processes in `comm`, and returns a pointer to the locally allocated segment in `baseptr` that can be used for load/store accesses on the calling process. The locally allocated memory can be the target of load/store accesses by remote processes; the base pointers for other processes can be queried using the function `MPI_WIN_SHARED_QUERY`. The call also returns a window object that can be used by all processes in `comm` to perform RMA operations. The size argument may be different at each process and `size = 0` is valid; ~~however, a library might allocate and expose more memory in order to create a fast, globally symmetric allocation.~~ It is the user's responsibility to ensure that the communicator `comm` represents a group of processes that can create a shared memory segment that can be accessed by all processes in the group. The discussions of rationales for `MPI_ALLOC_MEM` and `MPI_FREE_MEM` in Section 8.2 also apply to `MPI_WIN_ALLOCATE_SHARED`; in particular, see the rationale in Section 8.2 for an explanation of the type used for `baseptr`. The allocated memory is contiguous across process ranks unless the info key `alloc_shared_noncontig` is specified. Contiguous across process ranks means that the first address in the memory segment of process  $i$  is consecutive with the last address in the memory segment of process  $i - 1$ . This may enable the user to calculate remote address offsets with local information only.

The info argument can be used to specify hints similar to the info argument for `MPI_WIN_CREATE`, `MPI_WIN_ALLOC`, and `MPI_ALLOC_MEM`. The additional info key `alloc_shared_noncontig` allows the library to optimize the layout of the shared memory segments in memory.

*Advice to users.* If the info key `alloc_shared_noncontig` is not set to true, the allocation strategy is to allocate contiguous memory across process ranks. This may limit the performance on some architectures because it does not allow the implementation to modify the data layout (e.g., padding to reduce access latency). (*End of advice to users.*)

*Advice to implementors.* If the user sets the info key `alloc_shared_noncontig` to true, the implementation can allocate the memory requested by each process in a location that is close to this process. This can be achieved by padding or allocating memory in special memory segments. Both techniques may make the address space across consecutive ranks noncontiguous. (*End of advice to implementors.*)

The consistency of load/store accesses from/to the shared memory as observed by the user program depends on the architecture. A consistent view can be created in the unified memory model (see Section 11.4) by utilizing the window synchronization functions (see Section 11.5) or explicitly completing outstanding store accesses (e.g., by calling `MPI_WIN_FLUSH`). MPI does not define semantics for accessing shared memory windows in the separate memory model.

**MPI\_WIN\_SHARED\_QUERY**(win, rank, size, baseptr)

IN	win	shared memory window object (handle)
IN	rank	rank in the group of window win (non-negative integer)
OUT	size	size of the window segment (non-negative integer)
OUT	baseptr	address for load/store access to window segment (choice)

```
int MPI_Win_shared_query(MPI_Win win, int rank, MPI_Aint *size,
                        void *baseptr)
```

```
MPI_WIN_SHARED_QUERY(WIN, RANK, SIZE, BASEPTR, IERROR)
    INTEGER WIN, RANK, IERROR
    INTEGER (KIND=MPI_ADDRESS_KIND) SIZE, BASEPTR
```

This function queries the process-local address for remote memory segments created with **MPI\_WIN\_ALLOCATE\_SHARED**. This function can return different process-local addresses for the same physical memory on different processes. The returned memory can be used for load/store accesses subject to the constraints defined in Section 11.7. This function can only be called with windows of type **MPI\_WIN\_FLAVOR\_SHARED**. If the passed window is not of flavor **MPI\_WIN\_FLAVOR\_SHARED**, the error **MPI\_ERR\_RMA\_WRONG\_FLAVOR** is raised. When rank is **MPI\_PROC\_NULL**, the pointer and size returned are the pointer and size of the memory segment belonging the lowest rank that specified `size > 0`. If all processes in the group attached to the window specified `size = 0`, then the call returns `size = 0` and a `baseptr` as if **MPI\_ALLOC\_MEM** was called with `size = 0`.

## 11.2.4 Window of Dynamically Attached Memory

The **MPI-2** RMA model requires the user to identify the local memory that may be a target of RMA calls at the time the window is created. This has advantages for both the programmer (only this memory can be updated by one-sided operations and provides greater safety) and the MPI implementation (special steps may be taken to make one-sided access to such memory more efficient). However, consider implementing a modifiable linked list using RMA operations; as new items are added to the list, memory must be allocated. In a C or C++ program, this memory is typically allocated using `malloc` or `new` respectively. In **MPI-2 RMA**, the programmer must create a window with a predefined amount of memory and then implement routines for allocating memory from within that memory. In addition, there is no easy way to handle the situation where the predefined amount of memory turns out to be inadequate. To support this model, the routine **MPI\_WIN\_CREATE\_DYNAMIC** creates a window that makes it possible to expose memory without remote synchronization. It must be used in combination with the local routines **MPI\_WIN\_ATTACH** and **MPI\_WIN\_DETACH**.

**MPI\_WIN\_CREATE\_DYNAMIC**(info, comm, win)

IN	info	info argument (handle)
IN	comm	intra-communicator (handle)
OUT	win	window object returned by the call (handle)

```

int MPI_Win_create_dynamic(MPI_Info info, MPI_Comm comm, MPI_Win *win)
MPI_Win_create_dynamic(info, comm, win, ierror) BIND(C)
    TYPE(MPI_Info), INTENT(IN) :: info
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Win), INTENT(OUT) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_WIN_CREATE_DYNAMIC(INFO, COMM, WIN, IERROR)
    INTEGER INFO, COMM, WIN, IERROR

```

This is a collective call executed by all processes in the group of `comm`. It returns a window `win` without memory attached. Existing process memory can be attached as described below. This routine returns a window object that can be used by these processes to perform RMA operations on attached memory. Because this window has special properties, it will sometimes be referred to as a *dynamic* window.

The `info` argument can be used to specify hints similar to the `info` argument for `MPI_WIN_CREATE`.

In the case of a window created with `MPI_WIN_CREATE_DYNAMIC`, the `target_disp` for all RMA functions is the address at the target; i.e., the effective `window_base` is `MPI_BOTTOM` and the `disp_unit` is one. Users should use `MPI_GET_ADDRESS` at the target process to determine the address of a target memory location and communicate this address to the origin process.

*Advice to implementors.* In environments with heterogeneous data representations, care must be exercised in communicating addresses between processes. For example, it is possible that an address valid at the target process (for example, a 64-bit pointer) cannot be expressed as an address at the origin (for example, the origin uses 32-bit pointers). For this reason, a portable MPI implementation should ensure that the type `MPI_AINT` (cf. Table 3.3 on Page 31) is able to store addresses from any process. (*End of advice to implementors.*)

Memory in this window may not be used as the target of one-sided accesses in this window until it is attached using the function `MPI_WIN_ATTACH`. That is, in addition to using `MPI_WIN_CREATE_DYNAMIC` to create an MPI window, the user must use `MPI_WIN_ATTACH` before any local memory may be the target of an MPI RMA operation. Only memory that is currently accessible may be attached.

```

MPI_WIN_ATTACH(win, base, size)

```

IN	win	window object (handle)
IN	base	initial address of memory to be attached
IN	size	size of memory to be attached in bytes

```

int MPI_Win_attach(MPI_Win win, void *base, MPI_Aint size)
MPI_Win_attach(win, base, size, ierror) BIND(C)
    TYPE(MPI_Win), INTENT(IN) :: win
    TYPE(*), ASYNCHRONOUS :: base

```

```

1      INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: size
2      INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

3
4  MPI_WIN_ATTACH(WIN, BASE, SIZE, IERROR)
5      INTEGER WIN, IERROR
6      <type> [base]BASE(*)
7      INTEGER (KIND=MPI_ADDRESS_[SIZE]KIND) [size]SIZE

```

Attaches a local memory region beginning at `base` for remote access within the given window. The memory region specified must not contain any part that is already attached to the window `win`, that is, attaching overlapping memory concurrently within the same window is erroneous. The argument `win` must be a window that was created with `MPI_WIN_CREATE_DYNAMIC`. Multiple (but non-overlapping) memory regions may be attached to the same window.

*Rationale.* Requiring that memory be explicitly attached before it is exposed to one-sided access by other processes can significantly simplify implementations and improve performance. The ability to make memory available for RMA operations without requiring a collective `MPI_WIN_CREATE` call is needed for some one-sided programming models. (*End of rationale.*)

*Advice to users.* [Memory registration] Attaching memory to a window may require the use of scarce resources; thus, attaching large regions of memory is not recommended in portable programs. [Memory registration] Attaching memory to a window may fail if sufficient resources are not available; this is similar to the behavior of `MPI_ALLOC_MEM`.

The user is also responsible for ensuring that [memory registration] `MPI_WIN_ATTACH` at the target has [completed] returned before a process attempts to target that memory with an MPI RMA call.

Performing an RMA operation to memory that has not been attached [from] to a window created with `MPI_WIN_CREATE_DYNAMIC` is erroneous. (*End of advice to users.*)

*Advice to implementors.* A high-quality implementation will attempt to make as much memory available for [registration] attaching as possible. Any limitations should be documented by the implementor. (*End of advice to implementors.*)

[Memory registration] Attaching memory is a local operation as defined by MPI, which means that the call is not collective and completes without requiring any MPI routine to be called in any other process. Memory may be detached with the routine `MPI_WIN_DETACH`. After memory has been detached, it may not be the target of an MPI RMA operation on that window (unless the memory is re-attached with `MPI_WIN_ATTACH`).

```

43  MPI_WIN_DETACH(win, base)

```

IN	win	window object (handle)
IN	base	initial address of memory to be detached

```

48  int MPI_Win_detach(MPI_Win win, const void *base)

```

ticket-248T.

```

1 MPI_Win_detach(win, base, ierror) BIND(C)
2
3     TYPE(MPI_Win), INTENT(IN) :: win
4     TYPE(*), ASYNCHRONOUS :: base
5     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
6
7 MPI_WIN_DETACH(WIN, BASE, IERROR)
8     INTEGER WIN, IERROR
9     <type> [base]BASE(*)

```

Detaches a previously attached memory region beginning at **base**. The arguments **base** and **win** must match the arguments passed to a previous call to **MPI\_WIN\_ATTACH**.

*Advice to users.* Detaching memory may permit the implementation to make more efficient use of special memory or provide memory that may be needed by a subsequent **MPI\_WIN\_ATTACH**. Users are encouraged to detach memory that is no longer needed. Memory should be detached before it is freed by the user. (*End of advice to users.*)

Memory becomes detached when the associated dynamic memory window is freed, see Section 11.2.5.

### 11.2.5 Window Destruction

```

20 MPI_WIN_FREE(win)
21
22
23
24 MPI_Win_free(win)
25     INOUT    win
26             window object (handle)

```

```

27 int MPI_Win_free(MPI_Win *win)

```

```

28 MPI_Win_free(win, ierror) BIND(C)
29     TYPE(MPI_Win), INTENT(INOUT) :: win
30     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

31 MPI_WIN_FREE(WIN, IERROR)
32     INTEGER WIN, IERROR

```

```

33 {void MPI::Win::Free() (binding deprecated, see Section 15.2) }

```

Frees the window object **win** and returns a null handle (equal to **MPI\_WIN\_NULL**). This is a collective call executed by all processes in the group associated with **win**. **MPI\_WIN\_FREE(win)** can be invoked by a process only after it has completed its involvement in RMA communications on window **win**: [i.e.]e.g., the process has called **MPI\_WIN\_FENCE**, or called **MPI\_WIN\_WAIT** to match a previous call to **MPI\_WIN\_POST** or called **MPI\_WIN\_COMPLETE** to match a previous call to **MPI\_WIN\_START** or called **MPI\_WIN\_UNLOCK** to match a previous call to **MPI\_WIN\_LOCK**. [When the call returns, the window memory can be freed.]The memory associated with windows created by a call to **MPI\_WIN\_CREATE** may be freed after the call returns. If the window was created with **MPI\_WIN\_ALLOCATE**, **MPI\_WIN\_FREE** will free the window memory that was allocated in **MPI\_WIN\_ALLOCATE**. Freeing a window that was created with a call to **MPI\_WIN\_CREATE\_DYNAMIC** detaches all associated memory; i.e., it has the same effect as if all attached memory was detached by a call to **MPI\_WIN\_DETACH**.

*Advice to implementors.* MPI\_WIN\_FREE requires a barrier synchronization: no process can return from free until all processes in the group of win called free. This[,] is to ensure that no process will attempt to access a remote window (e.g., with lock/unlock) after it was freed. The only exception to this rule is when the user sets the no\_locks info [argument]key to true when creating the window. In that case, the local window can be freed without barrier synchronization. (End of advice to implementors.)

### 11.2.6 Window Attributes

The following [three] attributes are cached with a window[,] when the window is created.

MPI_WIN_BASE	window base address.
MPI_WIN_SIZE	[ ]window size, in bytes.
MPI_WIN_DISP_UNIT	displacement unit associated with the window.
[ticket270.]MPI_WIN_CREATE_FLAVOR	how the window was created.
[ticket270.]MPI_WIN_MODEL	memory model for window.

In C, calls to MPI\_Win\_get\_attr(win, MPI\_WIN\_BASE, &base, &flag), MPI\_Win\_get\_attr(win, MPI\_WIN\_SIZE, &size, &flag)[ and], MPI\_Win\_get\_attr(win, MPI\_WIN\_DISP\_UNIT, &disp\_unit, &flag), MPI\_Win\_get\_attr(win, MPI\_WIN\_CREATE\_FLAVOR, &create\_kind, &flag), and MPI\_Win\_get\_attr(win, MPI\_WIN\_MODEL, &memory\_model, &flag) will return in base a pointer to the start of the window win, and will return in size[ and], disp\_unit, create\_kind, and memory\_model pointers to the size[ and], displacement unit of the window, the kind of routine used to create the window, and the memory model, respectively. [And similarly, in C++.]And similarly, in C++ (binding deprecated, see Section 15.2). A detailed listing of the type of the pointer in the attribute value argument to MPI\_Win\_get\_attr and MPI\_Win\_set\_attr is shown in Table 11.1.

Attribute	C Type
MPI_WIN_BASE	void *
MPI_WIN_SIZE	MPI_Aint *
MPI_WIN_DISP_UNIT	int *
MPI_WIN_CREATE_FLAVOR	int *
MPI_WIN_MODEL	int *

Table 11.1: C types of attribute value argument to MPI\_Win\_get\_attr and MPI\_Win\_set\_attr.

In Fortran, calls to MPI\_WIN\_GET\_ATTR(win, MPI\_WIN\_BASE, base, flag, ierror), MPI\_WIN\_GET\_ATTR(win, MPI\_WIN\_SIZE, size, flag, ierror)[ and], MPI\_WIN\_GET\_ATTR(win, MPI\_WIN\_DISP\_UNIT, disp\_unit, flag, ierror), MPI\_WIN\_GET\_ATTR(win, MPI\_WIN\_CREATE\_FLAVOR, create\_kind, flag, ierror), and MPI\_WIN\_GET\_ATTR(win, MPI\_WIN\_MODEL, memory\_model, flag, ierror) will return in base, size[ and], disp\_unit create\_kind and memory\_model the (integer representation of) the base address, the size[ and], the displacement unit of the window win, the kind of routine used to create the window, and the memory model, respectively.

The values of `create_kind` are

<code>MPI_WIN_FLAVOR_CREATE</code>	Window was created with <code>MPI_WIN_CREATE</code> .	
<code>MPI_WIN_FLAVOR_ALLOCATE</code>	Window was created with <code>MPI_WIN_ALLOCATE</code> .	
<code>MPI_WIN_FLAVOR_DYNAMIC</code>	Window was created with <code>MPI_WIN_CREATE_DYNAMIC</code> .	ticket284.
<code>MPI_WIN_FLAVOR_SHARED</code>	Window was created with <code>MPI_WIN_ALLOCATE_SHARED</code> .	

The values of `memory_model` are `MPI_WIN_SEPARATE` and `MPI_WIN_UNIFIED`. The meaning of these is described in Section 11.4.

In the case of windows created with `MPI_WIN_CREATE_DYNAMIC`, the base address is `MPI_BOTTOM` and the size is 0. In C, pointers are returned and in Fortran, the values are returned, for the respective attributes. (The window attribute access functions are defined in Section 6.7.3, page 284.) The value returned for an attribute on a window is constant over the lifetime of the window. ticket270.

The other “window attribute,” namely the group of processes attached to the window, can be retrieved using the call below.

`MPI_WIN_GET_GROUP(win, group)`

IN	<code>win</code>	window object (handle)
OUT	<code>group</code>	group of processes which share access to the window (handle)

`int MPI_Win_get_group(MPI_Win win, MPI_Group *group)`

`MPI_Win_get_group(win, group, ierror) BIND(C)`  
`TYPE(MPI_Win), INTENT(IN) :: win`  
`TYPE(MPI_Group), INTENT(OUT) :: group`  
`INTEGER, OPTIONAL, INTENT(OUT) :: ierror`

`MPI_WIN_GET_GROUP(WIN, GROUP, IERROR)`  
`INTEGER WIN, GROUP, IERROR`

`{MPI::Group MPI::Win::Get_group() const(binding deprecated, see Section 15.2)}`

`MPI_WIN_GET_GROUP` returns a duplicate of the group of the communicator used to create the window`[.]` associated with `win`. The group is returned in `group`. ticket0.

## 11.3 Communication Calls

MPI supports [three]the following RMA communication calls: `MPI_PUT` [transfers]and `MPI_RPUT` transfer data from the caller memory (origin) to the target memory; `MPI_GET` [transfers]and `MPI_RGET` transfer data from the target memory to the caller memory; [and] `MPI_ACCUMULATE` [updates]and `MPI_RACCUMULATE` update locations in the target memory, e.g., by adding to these locations values sent from the caller memory`[.]`; `MPI_GET_ACCUMULATE`, `MPI_RGET_ACCUMULATE` and `MPI_FETCH_AND_OP` atom- ically return the data before the accumulate operation; and `MPI_COMPARE_AND_SWAP` ticket270. ticket270. ticket270. ticket270. ticket270. ticket270. ticket270.



performs a remote compare and swap operation. These operations are *nonblocking*: the call initiates the transfer, but the transfer may continue after the call returns. The transfer is completed, both at the origin and at the target, when a subsequent *synchronization* call is issued by the caller on the involved window object. These synchronization calls are described in Section 11.5, page 451. Transfers can also be completed with calls to flush routines; see Section 11.5.4, page 464 for details. For the MPI\_RPUT, MPI\_RGET, MPI\_RACCUMULATE, and MPI\_RGET\_ACCUMULATE calls, the transfer can be locally completed by using the MPI test or wait operations described in Section 3.7.3, page 58.

The local communication buffer of an RMA call should not be updated, and the local communication buffer of a get call should not be accessed after the RMA call[,] until the [subsequent synchronization call completes.]operation completes at the origin.

[It is erroneous to have concurrent conflicting accesses to the same memory location in a window ]The outcome of concurrent conflicting accesses to the same memory locations is undefined; if a location is updated by a put or accumulate operation, then [this location cannot be accessed by a load or another RMA operation ]the outcome of [local] loads or other RMA operations is undefined until the updating operation has completed at the target. There is one exception to this rule; namely, the same location can be updated by several concurrent accumulate calls, the outcome being as if these updates occurred in some order. In addition, [if a window cannot concurrently be updated by a put or accumulate operation and by a local store operation. This, even if these two updates access different locations in the window. The last restriction enables more efficient implementations of RMA operations on many systems. ]the outcome of concurrent [local]load/store and RMA updates to the same memory location is undefined. These restrictions are described in more detail in Section 11.7, page 468.

The calls use general datatype arguments to specify communication buffers at the origin and at the target. Thus, a transfer operation may also gather data at the source and scatter it at the destination. However, all arguments specifying both communication buffers are provided by the caller.

For all [three]RMA calls, the target process may be identical with the origin process; i.e., a process may use an RMA operation to move data in its memory.

*Rationale.* The choice of supporting “self-communication” is the same as for message-passing. It simplifies some coding, and is very useful with accumulate operations, to allow atomic updates of local variables. (*End of rationale.*)

MPI\_PROC\_NULL is a valid target rank in [the MPI RMA calls MPI\_ACCUMULATE, MPI\_GET, and MPI\_PUT]all MPI RMA communication calls. The effect is the same as for MPI\_PROC\_NULL in MPI point-to-point communication. After any RMA operation with rank MPI\_PROC\_NULL, it is still necessary to finish the RMA epoch with the synchronization method that started the epoch.

### 11.3.1 Put

The execution of a put operation is similar to the execution of a send by the origin process and a matching receive by the target process. The obvious difference is that all arguments are provided by one call — the call executed by the origin process.



```
MPI_PUT(origin_addr, origin_count, origin_datatype, target_rank, target_disp, target_count,
        target_datatype, win)
```

IN	origin_addr	initial address of origin buffer (choice)
IN	origin_count	number of entries in origin buffer (non-negative integer)
IN	origin_datatype	datatype of each entry in origin buffer (handle)
IN	target_rank	rank of target (non-negative integer)
IN	target_disp	displacement from start of window to target buffer (non-negative integer)
IN	target_count	number of entries in target buffer (non-negative integer)
IN	target_datatype	datatype of each entry in target buffer (handle)
IN	win	window object used for communication (handle)

```
int MPI_Put(const void *origin_addr, int origin_count, MPI_Datatype
            origin_datatype, int target_rank, MPI_Aint target_disp, int
            target_count, MPI_Datatype target_datatype, MPI_Win win)
```

```
MPI_Put(origin_addr, origin_count, origin_datatype, target_rank,
        target_disp, target_count, target_datatype, win, ierror)
BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr
INTEGER, INTENT(IN) :: origin_count, target_rank, target_count
TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
TYPE(MPI_Win), INTENT(IN) :: win
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_PUT(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK,
        TARGET_DISP, TARGET_COUNT, TARGET_DATATYPE, WIN, IERROR)
<type> ORIGIN_ADDR(*)
INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK, TARGET_COUNT,
TARGET_DATATYPE, WIN, IERROR
```

```
{void MPI::Win::Put(const void* origin_addr, int origin_count,
                    const MPI::Datatype& origin_datatype, int target_rank,
                    MPI::Aint target_disp, int target_count,
                    const MPI::Datatype& target_datatype) const(binding deprecated,
                    see Section 15.2) }
```

Transfers `origin_count` successive entries of the type specified by the `origin_datatype`, starting at address `origin_addr` on the origin node to the target node specified by the `win`, `target_rank` pair. The data are written in the target buffer at address `target_addr = window_base + target_disp × disp_unit`, where `window_base` and `disp_unit` are the base address and window displacement unit specified at window initialization, by the target process.

The target buffer is specified by the arguments `target_count` and `target_datatype`.

The data transfer is the same as that which would occur if the origin process executed a send operation with arguments `origin_addr`, `origin_count`, `origin_datatype`, `target_rank`, `tag`, `comm`, and the target process executed a receive operation with arguments `target_addr`, `target_count`, `target_datatype`, `source`, `tag`, `comm`, where `target_addr` is the target buffer address computed as explained above, **the values of tag are arbitrary valid matching tag values**, and `comm` is a communicator for the group of `win`.

The communication must satisfy the same constraints as for a similar message-passing communication. The `target_datatype` may not specify overlapping entries in the target buffer. The message sent must fit, without truncation, in the target buffer. Furthermore, the target buffer must fit in the target window **or in attached memory in a dynamic window**.

The `target_datatype` argument is a handle to a datatype object defined at the origin process. However, this object is interpreted at the target process: the outcome is as if the target datatype object was defined at the target process[,] by the same sequence of calls used to define it at the origin process. The target datatype must contain only relative displacements, not absolute addresses. The same holds for `get` and `accumulate`. **In the case of windows created with `MPI_WIN_CREATE_DYNAMIC`, displacements in the target datatype must be relative to `MPI_BOTTOM`.**

*Advice to users.* The `target_datatype` argument is a handle to a datatype object that is defined at the origin process, even though it defines a data layout in the target process memory. This causes no problems in a homogeneous environment, or in a heterogeneous environment[,] if only portable datatypes are used (portable datatypes are defined in Section 2.4, page 11).

The performance of a `put` transfer can be significantly affected, on some systems, [from]by the choice of window location and the shape and location of the origin and target buffer: transfers to a target window in memory allocated by `MPI_ALLOC_MEM` or `MPI_WIN_ALLOCATE` may be much faster on shared memory systems; transfers from contiguous buffers will be faster on most, if not all, systems; the alignment of the communication buffers may also impact performance. (*End of advice to users.*)

*Advice to implementors.* A high-quality implementation will attempt to prevent remote accesses to memory outside the window that was exposed by the process. This, both for debugging purposes, and for protection with client-server codes that use RMA. I.e., a high-quality implementation will check, if possible, window bounds on each RMA call, and raise an MPI exception at the origin call if an out-of-bound situation occurred. Note that the condition can be checked at the origin. Of course, the added safety achieved by such checks has to be weighed against the added cost of such checks. (*End of advice to implementors.*)

## 11.3.2 Get

```
MPI_GET(origin_addr, origin_count, origin_datatype, target_rank, target_disp, target_count,
        target_datatype, win)
```

OUT	origin_addr	initial address of origin buffer (choice)
IN	origin_count	number of entries in origin buffer (non-negative integer)
IN	origin_datatype	datatype of each entry in origin buffer (handle)
IN	target_rank	rank of target (non-negative integer)
IN	target_disp	displacement from window start to the beginning of the target buffer (non-negative integer)
IN	target_count	number of entries in target buffer (non-negative integer)
IN	target_datatype	datatype of each entry in target buffer (handle)
IN	win	window object used for communication (handle)

```
int MPI_Get(void *origin_addr, int origin_count,
            MPI_Datatype origin_datatype, int target_rank,
            MPI_Aint target_disp, int target_count,
            MPI_Datatype target_datatype, MPI_Win win)
```

```
MPI_Get(origin_addr, origin_count, origin_datatype, target_rank,
        target_disp, target_count, target_datatype, win, ierror)
    BIND(C)
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: origin_addr
    INTEGER, INTENT(IN) :: origin_count, target_rank, target_count
    TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
    TYPE(MPI_Win), INTENT(IN) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_GET(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK,
        TARGET_DISP, TARGET_COUNT, TARGET_DATATYPE, WIN, IERROR)
    <type> ORIGIN_ADDR(*)
    INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
    INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK, TARGET_COUNT,
    TARGET_DATATYPE, WIN, IERROR
```

```
{void MPI::Win::Get(void *origin_addr, int origin_count,
                    const MPI::Datatype& origin_datatype, int target_rank,
                    MPI::Aint target_disp, int target_count,
                    const MPI::Datatype& target_datatype) const(binding deprecated,
                    see Section 15.2) }
```

Similar to MPI\_PUT, except that the direction of data transfer is reversed. Data are copied from the target memory to the origin. The origin\_datatype may not specify

overlapping entries in the origin buffer. The target buffer must be contained within the target window **or within attached memory in a dynamic window**, and the copied data must fit, without truncation, in the origin buffer.

### 11.3.3 Examples for Communication Calls

These examples show the use of the **MPI\_GET** function. As all **MPI RMA** communication functions are nonblocking, they must be completed. In the following, this is accomplished with the routine **MPI\_WIN\_FENCE**, introduced in Section 11.5.

**Example 11.1** We show how to implement the generic indirect assignment  $A = B(\text{map})$ , where  $A$ ,  $B$  and  $\text{map}$  have the same distribution, and  $\text{map}$  is a permutation. To simplify, we assume a block distribution with equal size blocks.

```

SUBROUTINE MAPVALS(A, B, map, m, comm, p)
  USE MPI
  INTEGER m, map(m), comm, p
  REAL A(m), B(m)

  INTEGER otype(p), oindex(m), & ! used to construct origin datatypes
    ttype(p), tindex(m), & ! used to construct target datatypes
    count(p), total(p), &
    win, ierr
  INTEGER (KIND=MPI_ADDRESS_KIND) lowerbound, sizeofreal

  ! This part does the work that depends on the locations of B.
  ! Can be reused while this does not change

  CALL MPI_TYPE_GET_EXTENT(MPI_REAL, lowerbound, sizeofreal, ierr)
  CALL MPI_WIN_CREATE(B, m*sizeofreal, sizeofreal, MPI_INFO_NULL, &
    comm, win, ierr)

  ! This part does the work that depends on the value of map and
  ! the locations of the arrays.
  ! Can be reused while these do not change

  ! Compute number of entries to be received from each process

  DO i=1,p
    count(i) = 0
  END DO
  DO i=1,m
    j = map(i)/m+1
    count(j) = count(j)+1
  END DO

  total(1) = 0
  DO i=2,p

```

```

    total(i) = total(i-1) + count(i-1)
END DO

DO i=1,p
    count(i) = 0
END DO

! compute origin and target indices of entries.
! entry i at current process is received from location
! k at process (j-1), where map(i) = (j-1)*m + (k-1),
! j = 1..p and k = 1..m

DO i=1,m
    j = map(i)/m+1
    k = MOD(map(i),m)+1
    count(j) = count(j)+1
    oindex(total(j) + count(j)) = i
    tindex(total(j) + count(j)) = k
END DO

! create origin and target datatypes for each get operation
DO i=1,p
    CALL MPI_TYPE_CREATE_INDEXED_BLOCK(count(i), 1, oindex(total(i)+1), &
                                     MPI_REAL, otype(i), ierr)
    CALL MPI_TYPE_COMMIT(otype(i), ierr)
    CALL MPI_TYPE_CREATE_INDEXED_BLOCK(count(i), 1, tindex(total(i)+1), &
                                     MPI_REAL, ttype(i), ierr)
    CALL MPI_TYPE_COMMIT(ttype(i), ierr)
END DO

! this part does the assignment itself
CALL MPI_WIN_FENCE(0, win, ierr)
DO i=1,p
    CALL MPI_GET(A, 1, otype(i), i-1, 0, 1, ttype(i), win, ierr)
END DO
CALL MPI_WIN_FENCE(0, win, ierr)

CALL MPI_WIN_FREE(win, ierr)
DO i=1,p
    CALL MPI_TYPE_FREE(otype(i), ierr)
    CALL MPI_TYPE_FREE(ttype(i), ierr)
END DO
RETURN
END

```

### Example 11.2

A simpler version can be written that does not require that a datatype be built for the target buffer. But, one then needs a separate get call for each entry, as illustrated below. This code is much simpler, but usually much less efficient, for large arrays.

```

SUBROUTINE MAPVALS(A, B, map, m, comm, p)
USE MPI
INTEGER m, map(m), comm, p
REAL A(m), B(m)
INTEGER win, ierr
INTEGER (KIND=MPI_ADDRESS_KIND) lowerbound, sizeofreal

CALL MPI_TYPE_GET_EXTENT(MPI_REAL, lowerbound, sizeofreal, ierr)
CALL MPI_WIN_CREATE(B, m*sizeofreal, sizeofreal, MPI_INFO_NULL, &
                    comm, win, ierr)

CALL MPI_WIN_FENCE(0, win, ierr)
DO i=1,m
  j = map(i)/m
  k = MOD(map(i),m)
  CALL MPI_GET(A(i), 1, MPI_REAL, j, k, 1, MPI_REAL, win, ierr)
END DO
CALL MPI_WIN_FENCE(0, win, ierr)
CALL MPI_WIN_FREE(win, ierr)
RETURN
END

```

#### 11.3.4 Accumulate Functions

It is often useful in a put operation to combine the data moved to the target process with the data that resides at that process, rather than replacing the data there. This will allow, for example, the accumulation of a sum by having all involved processes add their contribution to the sum variable in the memory of one process. The accumulate functions have slightly different semantics than the put and get functions; see Section 11.7 for details.

#### Accumulate Function

```
MPI_ACCUMULATE(origin_addr, origin_count, origin_datatype, target_rank, target_disp,
                target_count, target_datatype, op, win)
```

IN	origin_addr	initial address of buffer (choice)
IN	origin_count	number of entries in buffer (non-negative integer)
IN	origin_datatype	datatype of each entry (handle)
IN	target_rank	rank of target (non-negative integer)
IN	target_disp	displacement from start of window to beginning of target buffer (non-negative integer)
IN	target_count	number of entries in target buffer (non-negative integer)
IN	target_datatype	datatype of each entry in target buffer (handle)
IN	op	reduce operation (handle)
IN	win	window object (handle)

```
int MPI_Accumulate(const void *origin_addr, int origin_count,
                  MPI_Datatype origin_datatype, int target_rank,
                  MPI_Aint target_disp, int target_count,
                  MPI_Datatype target_datatype, MPI_Op op, MPI_Win win)
```

```
MPI_Accumulate(origin_addr, origin_count, origin_datatype, target_rank,
               target_disp, target_count, target_datatype, op, win, ierror)
BIND(C)
```

```
TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr
INTEGER, INTENT(IN) :: origin_count, target_rank, target_count
TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
TYPE(MPI_Op), INTENT(IN) :: op
TYPE(MPI_Win), INTENT(IN) :: win
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_ACCUMULATE(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK,
               TARGET_DISP, TARGET_COUNT, TARGET_DATATYPE, OP, WIN, IERROR)
<type> ORIGIN_ADDR(*)
INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK, TARGET_COUNT,
TARGET_DATATYPE, OP, WIN, IERROR
```

```
{void MPI::Win::Accumulate(const void* origin_addr, int origin_count, const
                          MPI::Datatype& origin_datatype, int target_rank, MPI::Aint
                          target_disp, int target_count, const MPI::Datatype&
                          target_datatype, const MPI::Op& op) const (binding deprecated, see
                          Section 15.2) }
```

Accumulate the contents of the origin buffer (as defined by `origin_addr`, `origin_count` and `origin_datatype`) to the buffer specified by arguments `target_count` and `target_datatype`, at offset `target_disp`, in the target window specified by `target_rank` and `win`, using the operation

op. This is like MPI\_PUT except that data is combined into the target area instead of overwriting it.

Any of the predefined operations for MPI\_REDUCE can be used. User-defined functions cannot be used. For example, if op is MPI\_SUM, each element of the origin buffer is added to the corresponding element in the target, replacing the former value in the target.

Each datatype argument must be a predefined datatype or a derived datatype, where all basic components are of the same predefined datatype. Both datatype arguments must be constructed from the same predefined datatype. The operation op applies to elements of that predefined type. The parameter target\_datatype must not specify overlapping entries, and the target buffer must fit in the target window.

A new predefined operation, MPI\_REPLACE, is defined. It corresponds to the associative function  $f(a, b) = b$ ; i.e., the current value in the target memory is replaced by the value supplied by the origin.

MPI\_REPLACE can be used only in MPI\_ACCUMULATE, MPI\_RACCUMULATE, MPI\_GET\_ACCUMULATE, MPI\_FETCH\_AND\_OP, and MPI\_RGET\_ACCUMULATE, but not in collective reduction operations[,] such as MPI\_REDUCE.

*Advice to users.* MPI\_PUT is a special case of MPI\_ACCUMULATE, with the operation MPI\_REPLACE. Note, however, that MPI\_PUT and MPI\_ACCUMULATE have different constraints on concurrent updates. (*End of advice to users.*)

**Example 11.3** We want to compute  $B(j) = \sum_{\text{map}(i)=j} A(i)$ . The arrays A, B and map are distributed in the same manner. We write the simple version.

```

SUBROUTINE SUM(A, B, map, m, comm, p)
USE MPI
INTEGER m, map(m), comm, p, win, ierr
REAL A(m), B(m)
INTEGER (KIND=MPI_ADDRESS_KIND) lowerbound, sizeofreal

CALL MPI_TYPE_GET_EXTENT(MPI_REAL, lowerbound, sizeofreal, ierr)
CALL MPI_WIN_CREATE(B, m*sizeofreal, sizeofreal, MPI_INFO_NULL, &
                    comm, win, ierr)

CALL MPI_WIN_FENCE(0, win, ierr)
DO i=1,m
  j = map(i)/m
  k = MOD(map(i),m)
  CALL MPI_ACCUMULATE(A(i), 1, MPI_REAL, j, k, 1, MPI_REAL, &
                    MPI_SUM, win, ierr)
END DO
CALL MPI_WIN_FENCE(0, win, ierr)

CALL MPI_WIN_FREE(win, ierr)
RETURN
END

```

This code is identical to the code in Example 11.2, page 437, except that a call to get has been replaced by a call to accumulate. (Note that, if map is one-to-one, then the



code computes  $B = A(\text{map}^{-1})$ , which is the reverse assignment to the one computed in that previous example.) In a similar manner, we can replace in Example 11.1, page 436, the call to get by a call to accumulate, thus performing the computation with only one communication between any two processes.

### Get Accumulate Function

It is often useful to have fetch-and-accumulate semantics such that the remote data is returned to the caller before the sent data is accumulated into the remote data. The get and accumulate steps are executed atomically for each basic element in the datatype (see Section 11.7 for details). The predefined operation `MPI_REPLACE` provides fetch-and-set behavior.

```
MPI_GET_ACCUMULATE(origin_addr,    origin_count,    origin_datatype,    result_addr,
                    result_count, result_datatype, target_rank, target_disp, target_count, target_datatype,
                    op, win)
```

IN	origin_addr	initial address of buffer (choice)
IN	origin_count	number of entries in origin buffer (non-negative integer)
IN	origin_datatype	datatype of each entry in origin buffer (handle)
OUT	result_addr	initial address of result buffer (choice)
IN	result_count	number of entries in result buffer (non-negative integer)
IN	result_datatype	datatype of each entry in result buffer (handle)
IN	target_rank	rank of target (non-negative integer)
IN	target_disp	displacement from start of window to beginning of target buffer (non-negative integer)
IN	target_count	number of entries in target buffer (non-negative integer)
IN	target_datatype	datatype of each entry in target buffer (handle)
IN	op	reduce operation (handle)
IN	win	window object (handle)

```
int MPI_Get_accumulate(const void *origin_addr, int origin_count,
                      MPI_Datatype origin_datatype, void *result_addr,
                      int result_count, MPI_Datatype result_datatype,
                      int target_rank, MPI_Aint target_disp, int target_count,
                      MPI_Datatype target_datatype, MPI_Op op, MPI_Win win)
```

```
MPI_Get_accumulate(origin_addr, origin_count, origin_datatype, result_addr,
                  result_count, result_datatype, target_rank, target_disp,
                  target_count, target_datatype, op, win, ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr
TYPE(*), DIMENSION(..), ASYNCHRONOUS :: result_addr
```

```

1      INTEGER, INTENT(IN) :: origin_count, result_count, target_rank,
2      target_count
3      TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype,
4      result_datatype
5      INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
6      TYPE(MPI_Op), INTENT(IN) :: op
7      TYPE(MPI_Win), INTENT(IN) :: win
8      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
9
10     MPI_GET_ACCUMULATE(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE, RESULT_ADDR,
11                        RESULT_COUNT, RESULT_DATATYPE, TARGET_RANK, TARGET_DISP,
12                        TARGET_COUNT, TARGET_DATATYPE, OP, WIN, IERROR)
13     <type> ORIGIN_ADDR(*), RESULT_ADDR(*)
14     INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
15     INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, RESULT_COUNT, RESULT_DATATYPE,
16     TARGET_RANK, TARGET_COUNT, TARGET_DATATYPE, OP, WIN, IERROR

```

Accumulate `origin_count` elements of type `origin_datatype` from the origin buffer (`origin_addr`) to the buffer at offset `target_disp`, in the target window specified by `target_rank` and `win`, using the operation `op` and return in the result buffer `result_addr` the content of the target buffer before the accumulation.

The origin and result buffers (`origin_addr` and `result_addr`) must be disjoint. Each datatype argument must be a predefined datatype or a derived datatype where all basic components are of the same predefined datatype. All datatype arguments must be constructed from the same predefined datatype. The operation `op` applies to elements of that predefined type. `target_datatype` must not specify overlapping entries, and the target buffer must fit in the target window or in attached memory in a dynamic window. The operation is executed atomically for each basic datatype; see Section 11.7 for details.

Any of the predefined operations for `MPI_REDUCE`, and `MPI_NO_OP` or `MPI_REPLACE` can be specified as `op`. User-defined functions cannot be used. A new predefined operation, `MPI_NO_OP`, is defined. It corresponds to the associative function  $f(a,b) = a$ ; i.e., the current value in the target memory is returned in the result buffer at the origin and no operation is performed on the target buffer. `MPI_NO_OP` can be used only in `MPI_GET_ACCUMULATE`, `MPI_RGET_ACCUMULATE`, and `MPI_FETCH_AND_OP`. `MPI_NO_OP` cannot be used in `MPI_ACCUMULATE`, `MPI_RACCUMULATE`, or collective reduction operations, such as `MPI_REDUCE` and others.

*Advice to users.* `MPI_GET` is similar to `MPI_GET_ACCUMULATE`, with the operation `MPI_NO_OP`. Note, however, that `MPI_GET` and `MPI_GET_ACCUMULATE` have different constraints on concurrent updates. (*End of advice to users.*)

## Fetch and Op Function

The generic functionality of `MPI_GET_ACCUMULATE` might ~~significantly~~ limit the performance of fetch-and-increment or fetch-and-add calls that might be supported by special hardware operations. `MPI_FETCH_AND_OP` thus allows for a fast implementation of a commonly used subset of the functionality of `MPI_GET_ACCUMULATE`.

`MPI_FETCH_AND_OP(origin_addr, result_addr, datatype, target_rank, target_disp, op, win)`

IN	<code>origin_addr</code>	initial address of buffer (choice)
OUT	<code>result_addr</code>	initial address of result buffer (choice)
IN	<code>datatype</code>	datatype of the entry in origin, result, and target buffers (handle)
IN	<code>target_rank</code>	rank of target (non-negative integer)
IN	<code>target_disp</code>	displacement from start of window to beginning of target buffer (non-negative integer)
IN	<code>op</code>	reduce operation (handle)
IN	<code>win</code>	window object (handle)

```
int MPI_Fetch_and_op(const void *origin_addr, void *result_addr,
                    MPI_Datatype datatype, int target_rank, MPI_Aint target_disp,
                    MPI_Op op, MPI_Win win)
```

```
MPI_Fetch_and_op(origin_addr, result_addr, datatype, target_rank,
                 target_disp, op, win, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: result_addr
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    INTEGER, INTENT(IN) :: target_rank
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
    TYPE(MPI_Op), INTENT(IN) :: op
    TYPE(MPI_Win), INTENT(IN) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FETCH_AND_OP(ORIGIN_ADDR, RESULT_ADDR, DATATYPE, TARGET_RANK,
                 TARGET_DISP, OP, WIN, IERROR)
    <type> ORIGIN_ADDR(*), RESULT_ADDR(*)
    INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
    INTEGER DATATYPE, TARGET_RANK, OP, WIN, IERROR
```

Accumulate one element of type `datatype` from the origin buffer (`origin_addr`) to the buffer at offset `target_disp`, in the target window specified by `target_rank` and `win`, using the operation `op` and return in the result buffer `result_addr` the content of the target buffer before the accumulation.

The origin and result buffers (`origin_addr` and `result_addr`) must be disjoint. Any of the predefined operations for `MPI_REDUCE`, as well as `MPI_NO_OP` or `MPI_REPLACE`, can be specified as `op`. User-defined functions cannot be used. The `datatype` argument must be a predefined datatype. The operation is executed atomically.

### Compare and Swap Function

Another useful operation is an atomic compare and swap where the value at the origin is compared to the value at the target, which is atomically replaced by a third value only if origin and target are equal.

```

1 MPI_COMPARE_AND_SWAP(origin_addr, compare_addr, result_addr, datatype, target_rank,
2 target_disp, win)

```

3	IN	origin_addr	initial address of buffer (choice)
4	IN	compare_addr	initial address of compare buffer (choice)
5	OUT	result_addr	initial address of result buffer (choice)
6	IN	datatype	datatype of the element in all buffers (handle)
7	IN	target_rank	rank of target (non-negative integer)
8	IN	target_disp	displacement from start of window to beginning of target buffer (non-negative integer)
9	IN	win	window object (handle)

```

14 int MPI_Compare_and_swap(const void *origin_addr, const void *compare_addr,
15 void *result_addr, MPI_Datatype datatype, int target_rank,
16 MPI_Aint target_disp, MPI_Win win)

```

```

17 MPI_Compare_and_swap(origin_addr, compare_addr, result_addr, datatype,
18 target_rank, target_disp, win, ierror) BIND(C)
19 TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr,
20 compare_addr
21 TYPE(*), DIMENSION(..), ASYNCHRONOUS :: result_addr
22 TYPE(MPI_Datatype), INTENT(IN) :: datatype
23 INTEGER, INTENT(IN) :: target_rank
24 INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
25 TYPE(MPI_Win), INTENT(IN) :: win
26 INTEGER, OPTIONAL, INTENT(OUT) :: ierror
27
28 MPI_COMPARE_AND_SWAP(ORIGIN_ADDR, COMPARE_ADDR, RESULT_ADDR, DATATYPE,
29 TARGET_RANK, TARGET_DISP, WIN, IERROR)
30 <type> ORIGIN_ADDR(*), COMPARE_ADDR(*), RESULT_ADDR(*)
31 INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
32 INTEGER DATATYPE, TARGET_RANK, WIN, IERROR
33

```

This function compares one element of type `datatype` in the compare buffer `compare_addr` with the buffer at offset `target_disp` in the target window specified by `target_rank` and `win` and replaces the value at the target with the value in the origin buffer `origin_addr` if the compare buffer and the target buffer are identical. The original value at the target is returned in the buffer `result_addr`. The parameter `datatype` must belong to one of the following categories of predefined datatypes: C integer, Fortran integer, Logical, [Complex, ]or Byte as specified in Section 5.9.2 on page 188[, or can be of type `MPI_AINT` or `MPI_OFFSET`]. The origin and result buffers (`origin_addr` and `result_addr`) must be disjoint. [ Any of the predefined operations for `MPI_REDUCE`, and `MPI_NO_OP` or `MPI_REPLACE` can be specified as `op`. User-defined functions cannot be used. The outcome of accumulate operations with overlapping types of different sizes or target displacements is undefined, see Section 11.7.1. ]

### 11.3.5 Request-based RMA Communication Operations

Request-based RMA communication operations allow the user to associate a request handle with the RMA operations and test or wait for the completion of these requests using the functions described in Section 3.7.3, page 58. Request-based RMA operations are only valid within a passive-target epoch.

Upon returning from a completion call in which an RMA operation completes, the `MPI_ERROR` field in the associated status object is set appropriately (see Section 3.2.5 on page 34). [The values of the `MPI_SOURCE` and `MPI_TAG` fields are undefined.] All other fields of status and the results of status query functions (e.g., `MPI_GET_COUNT`) are undefined. It is valid to mix different request types ([i.e.] e.g., any combination of RMA requests, collective requests, I/O requests, generalized requests, or point-to-point requests) in functions that enable multiple completions (e.g., `MPI_WAITALL`). It is erroneous to call `MPI_REQUEST_FREE` or `MPI_CANCEL` for a request associated with an RMA operation. RMA requests are not persistent.

The end of the epoch, or explicit bulk synchronization using `MPI_WIN_FLUSH`, `MPI_WIN_FLUSH_ALL`, `MPI_WIN_FLUSH_LOCAL` or `MPI_WIN_FLUSH_LOCAL_ALL`, also indicates completion of the RMA operations. However, users must still wait or test on the request handle to allow the MPI implementation to clean up any resources associated with these requests; in such cases the wait operation will complete locally.

`MPI_RPUT(origin_addr, origin_count, origin_datatype, target_rank, target_disp, target_count, target_datatype, win, request)`

IN	<code>origin_addr</code>	initial address of origin buffer (choice)
IN	<code>origin_count</code>	number of entries in origin buffer (non-negative integer)
IN	<code>origin_datatype</code>	datatype of each entry in origin buffer (handle)
IN	<code>target_rank</code>	rank of target (non-negative integer)
IN	<code>target_disp</code>	displacement from start of window to target buffer (non-negative integer)
IN	<code>target_count</code>	number of entries in target buffer (non-negative integer)
IN	<code>target_datatype</code>	datatype of each entry in target buffer (handle)
IN	<code>win</code>	window object used for communication (handle)
OUT	<code>request</code>	RMA request (handle)

```
int MPI_Rput(const void *origin_addr, int origin_count,
             MPI_Datatype origin_datatype, int target_rank,
             MPI_Aint target_disp, int target_count,
             MPI_Datatype target_datatype, MPI_Win win,
             MPI_Request *request)
```

```
MPI_Rput(origin_addr, origin_count, origin_datatype, target_rank,
         target_disp, target_count, target_datatype, win, request,
         ierror) BIND(C)
```

```

1  TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr
2  INTEGER, INTENT(IN) :: origin_count, target_rank, target_count
3  TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype
4  INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
5  TYPE(MPI_Win), INTENT(IN) :: win
6  TYPE(MPI_Request), INTENT(OUT) :: request
7  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
8
9  MPI_RPUT(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK,
10         TARGET_DISP, TARGET_COUNT, TARGET_DATATYPE, WIN, REQUEST,
11         IERROR)
12  <type> ORIGIN_ADDR(*)
13  INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
14  INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK, TARGET_COUNT,
15  TARGET_DATATYPE, WIN, REQUEST, IERROR

```

MPI\_RPUT is similar to MPI\_PUT (Section 11.3.1), except that it allocates a communication request object and associates it with the request handle (the argument `request`). The completion of an MPI\_RPUT operation (i.e., after the corresponding test or wait) indicates that the sender is now free to update the locations in the origin buffer. It does not indicate that the data is available at the target window. If remote completion is required, MPI\_WIN\_FLUSH, MPI\_WIN\_FLUSH\_ALL, MPI\_WIN\_UNLOCK or MPI\_WIN\_UNLOCK\_ALL can be used.

```

25  MPI_RGET(origin_addr, origin_count, origin_datatype, target_rank, target_disp, target_count,
26         target_datatype, win, request)

```

27	OUT	origin_addr	initial address of origin buffer (choice)
28	IN	origin_count	number of entries in origin buffer (non-negative integer)
29			
30			
31	IN	origin_datatype	datatype of each entry in origin buffer (handle)
32	IN	target_rank	rank of target (non-negative integer)
33	IN	target_disp	displacement from window start to the beginning of the target buffer (non-negative integer)
34			
35			
36	IN	target_count	number of entries in target buffer (non-negative integer)
37			
38	IN	target_datatype	datatype of each entry in target buffer (handle)
39	IN	win	window object used for communication (handle)
40			
41	OUT	request	RMA request (handle)

```

42
43  int MPI_Rget(void *origin_addr, int origin_count,
44             MPI_Datatype origin_datatype, int target_rank,
45             MPI_Aint target_disp, int target_count,
46             MPI_Datatype target_datatype, MPI_Win win,
47             MPI_Request *request)
48

```

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```

MPI_Rget(origin_addr, origin_count, origin_datatype, target_rank,
        target_disp, target_count, target_datatype, win, request,
        ierror) BIND(C)
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: origin_addr
    INTEGER, INTENT(IN) :: origin_count, target_rank, target_count
    TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
    TYPE(MPI_Win), INTENT(IN) :: win
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_RGET(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK,
        TARGET_DISP, TARGET_COUNT, TARGET_DATATYPE, WIN, REQUEST,
        IERROR)
    <type> ORIGIN_ADDR(*)
    INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
    INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK, TARGET_COUNT,
    TARGET_DATATYPE, WIN, REQUEST, IERROR

```

MPI\_RGET is similar to MPI\_GET (Section 11.3.2), except that it allocates a communication request object and associates it with the request handle (the argument `request`) that can be used to wait or test for completion. The completion of an MPI\_RGET operation indicates that the data is available in the origin buffer. If `origin_addr` points to memory attached to a window, then the data becomes available in the private copy of this window.

```

MPI_RACCUMULATE(origin_addr, origin_count, origin_datatype, target_rank, target_disp,
        target_count, target_datatype, op, win, request)

```

IN	<code>origin_addr</code>	initial address of buffer (choice)
IN	<code>origin_count</code>	number of entries in buffer (non-negative integer)
IN	<code>origin_datatype</code>	datatype of each buffer entry (handle)
IN	<code>target_rank</code>	rank of target (non-negative integer)
IN	<code>target_disp</code>	displacement from start of window to beginning of target buffer (non-negative integer)
IN	<code>target_count</code>	number of entries in target buffer (non-negative integer)
IN	<code>target_datatype</code>	datatype of each entry in target buffer (handle)
IN	<code>op</code>	reduce operation (handle)
IN	<code>win</code>	window object (handle)
OUT	<code>request</code>	RMA request (handle)

```

int MPI_Raccumulate(const void *origin_addr, int origin_count,
        MPI_Datatype origin_datatype, int target_rank,
        MPI_Aint target_disp, int target_count,
        MPI_Datatype target_datatype, MPI_Op op, MPI_Win win,
        MPI_Request *request)

```

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```

1
ticket-248T. 2 MPI_Raccumulate(origin_addr, origin_count, origin_datatype, target_rank,
3             target_disp, target_count, target_datatype, op, win, request,
4             ierror) BIND(C)
5             TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr
6             INTEGER, INTENT(IN) :: origin_count, target_rank, target_count
7             TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype
8             INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
9             TYPE(MPI_Op), INTENT(IN) :: op
10            TYPE(MPI_Win), INTENT(IN) :: win
11            TYPE(MPI_Request), INTENT(OUT) :: request
12            INTEGER, OPTIONAL, INTENT(OUT) :: ierror
13
14 MPI_RACCUMULATE(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK,
15                TARGET_DISP, TARGET_COUNT, TARGET_DATATYPE, OP, WIN, REQUEST,
16                IERROR)
17    <type> ORIGIN_ADDR(*)
18    INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
19    INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK, TARGET_COUNT,
20    TARGET_DATATYPE, OP, WIN, REQUEST, IERROR

```

MPI\_RACCUMULATE is similar to MPI\_ACCUMULATE (Section 11.3.4), except that it allocates a communication request object and associates it with the request handle (the argument request) that can be used to wait or test for completion. The completion of an MPI\_RACCUMULATE operation indicates that the origin buffer is free to be updated. It does not indicate that the operation has completed at the target window.



```

MPI_RGET_ACCUMULATE(origin_addr,  origin_count,  origin_datatype,  result_addr,
                    result_count, result_datatype, target_rank, target_disp, target_count, target_datatype,
                    op, win, request)

```

IN	origin_addr	initial address of buffer (choice)
IN	origin_count	number of entries in origin buffer (non-negative integer)
IN	origin_datatype	datatype of each buffer entry (handle)
OUT	result_addr	initial address of result buffer (choice)
IN	result_count	number of entries in result buffer (non-negative integer)
IN	result_datatype	datatype of each buffer entry (handle)
IN	target_rank	rank of target (non-negative integer)
IN	target_disp	displacement from start of window to beginning of target buffer (non-negative integer)
IN	target_count	number of entries in target buffer (non-negative integer)
IN	target_datatype	datatype of each buffer entry (handle)
IN	op	reduce operation (handle)
IN	win	window object (handle)
OUT	request	RMA request (handle)

```

int MPI_Rget_accumulate(const void *origin_addr, int origin_count,
    MPI_Datatype origin_datatype, void *result_addr,
    int result_count, MPI_Datatype result_datatype,
    int target_rank, MPI_Aint target_disp, int target_count,
    MPI_Datatype target_datatype, MPI_Op op, MPI_Win win,
    MPI_Request *request)

```

```

MPI_Rget_accumulate(origin_addr, origin_count, origin_datatype,
    result_addr, result_count, result_datatype, target_rank,
    target_disp, target_count, target_datatype, op, win, request,
    ierror) BIND(C)

```

```

TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr
TYPE(*), DIMENSION(..), ASYNCHRONOUS :: result_addr
INTEGER, INTENT(IN) :: origin_count, result_count, target_rank,
target_count
TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype,
result_datatype
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
TYPE(MPI_Op), INTENT(IN) :: op
TYPE(MPI_Win), INTENT(IN) :: win
TYPE(MPI_Request), INTENT(OUT) :: request
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

1 MPI_RGET_ACCUMULATE(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE,
2     RESULT_ADDR, RESULT_COUNT, RESULT_DATATYPE, TARGET_RANK,
3     TARGET_DISP, TARGET_COUNT, TARGET_DATATYPE, OP, WIN, REQUEST,
4     IERROR)
5     <type> ORIGIN_ADDR(*), RESULT_ADDR(*)
6     INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
7     INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, RESULT_COUNT, RESULT_DATATYPE,
8     TARGET_RANK, TARGET_COUNT, TARGET_DATATYPE, OP, WIN, REQUEST, IERROR
9

```

MPI\_RGET\_ACCUMULATE is similar to MPI\_GET\_ACCUMULATE (Section 11.3.4), except that it allocates a communication request object and associates it with the request handle (the argument `request`) that can be used to wait or test for completion. The completion of an MPI\_RGET\_ACCUMULATE operation indicates that the data is available in the result buffer and the origin buffer is free to be updated. It does not indicate that the operation has been completed at the target window.

## 11.4 Memory Model

The memory semantics of RMA are best understood by using the concept of public and private window copies. We assume that systems have a public memory region that is addressable by all processes (e.g., the shared memory in shared memory machines or the exposed main memory in distributed memory machines). In addition, most machines have fast private buffers (e.g., transparent caches or explicit communication buffers) local to each process where copies of data elements from the main memory can be stored for faster access. Such buffers are either coherent, i.e., all updates to main memory are reflected in all private copies consistently, or non-coherent, i.e., conflicting accesses to main memory need to be synchronized and updated in all private copies explicitly. Coherent systems allow direct updates to remote memory without any participation of the remote side. Non-coherent systems, however, need to call RMA functions in order to reflect updates to the public window in their private memory. Thus, in coherent memory, the public and the private window are identical while they remain logically separate in the non-coherent case. MPI thus differentiates between two memory models called *RMA unified*, if public and private window are logically identical, and *RMA separate*, otherwise.

In the RMA separate model, there is only one instance of each variable in process memory, but a distinct *public* copy of the variable for each window that contains it. A load accesses the instance in process memory (this includes MPI sends). A local store accesses and updates the instance in process memory (this includes MPI receives), but the update may affect other public copies of the same locations. A get on a window accesses the public copy of that window. A put or accumulate on a window accesses and updates the public copy of that window, but the update may affect the private copy of the same locations in process memory, and public copies of other overlapping windows. This is illustrated in Figure 11.1.

In the RMA unified model, public and private copies are identical and updates via put or accumulate calls are observed by load operations without additional RMA calls. A store access to a window is visible to remote get or accumulate calls without additional RMA calls. These stronger semantics of the RMA unified model allow the user to omit some synchronization calls and potentially improve performance.

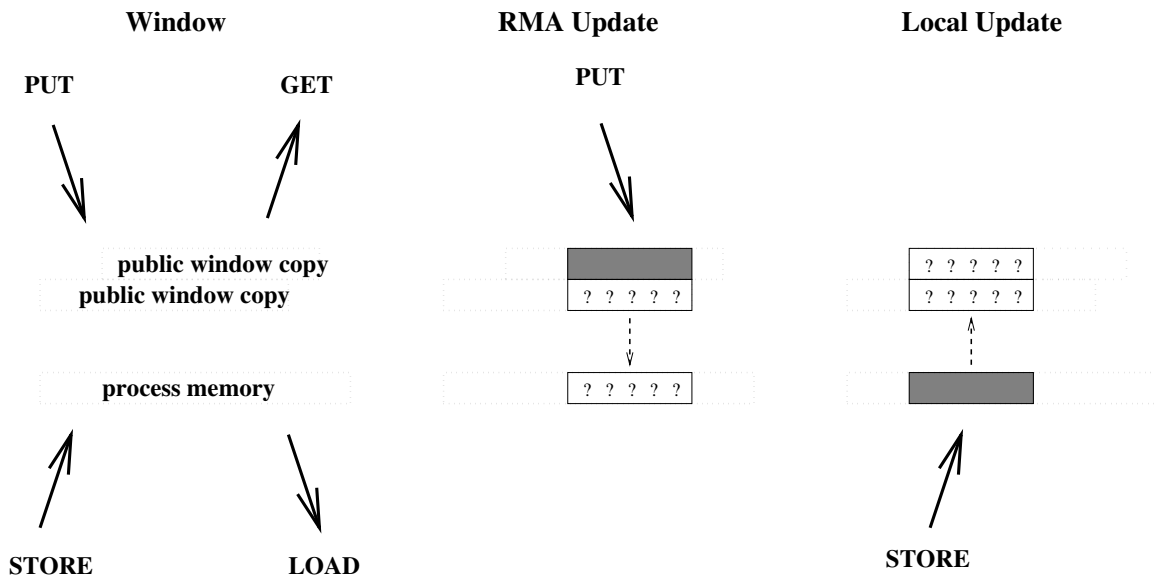


Figure 11.1: Schematic description of [ticket270.]the public/private window operations in the MPI\_WIN\_SEPARATE memory model for two overlapping windows.

*Advice to users.* If accesses in the RMA unified model are not synchronized (with locks or flushes, see Section 11.5.3), load and store operations might observe changes to the memory while they are in progress. The order in which data is written is not specified unless further synchronization is used. This might lead to inconsistent views on memory and programs that assume that a transfer is complete by only checking parts of the message are erroneous. (*End of advice to users.*)

The memory model for a particular RMA window can be determined by accessing the attribute MPI\_WIN\_MODEL. If the memory model is the unified model, the value of this attribute is MPI\_WIN\_UNIFIED; otherwise, the value is MPI\_WIN\_SEPARATE.

## 11.5 Synchronization Calls

RMA communications fall in two categories:

- **active target** communication, where data is moved from the memory of one process to the memory of another, and both are explicitly involved in the communication. This communication pattern is similar to message passing, except that all the data transfer arguments are provided by one process, and the second process only participates in the synchronization.
- **passive target** communication, where data is moved from the memory of one process to the memory of another, and only the origin process is explicitly involved in the transfer. Thus, two origin processes may communicate by accessing the same location in a target window. The process that owns the target window may be distinct from the two communicating processes, in which case it does not participate explicitly in the communication. This communication paradigm is closest to a shared memory model, where shared data can be accessed by all processes, irrespective of location.

RMA communication calls with argument `win` must occur at a process only within an **access epoch** for `win`. Such an epoch starts with an RMA synchronization call on `win`; it proceeds with zero or more RMA communication calls (e.g., `MPI_PUT`, `MPI_GET` or `MPI_ACCUMULATE`) on `win`; it completes with another synchronization call on `win`. This allows users to amortize one synchronization with multiple data transfers and provide implementors more flexibility in the implementation of RMA operations.

Distinct access epochs for `win` at the same process must be disjoint. On the other hand, epochs pertaining to different `win` arguments may overlap. Local operations or other MPI calls may also occur during an epoch.

In active target communication, a target window can be accessed by RMA operations only within an **exposure epoch**. Such an epoch is started and completed by RMA synchronization calls executed by the target process. Distinct exposure epochs at a process on the same window must be disjoint, but such an exposure epoch may overlap with exposure epochs on other windows or with access epochs for the same or other `win` arguments. There is a one-to-one matching between access epochs at origin processes and exposure epochs on target processes: RMA operations issued by an origin process for a target window will access that target window during the same exposure epoch if and only if they were issued during the same access epoch.

In passive target communication the target process does not execute RMA synchronization calls, and there is no concept of an exposure epoch.

MPI provides three synchronization mechanisms:

1. The `MPI_WIN_FENCE` collective synchronization call supports a simple synchronization pattern that is often used in parallel computations: namely a loosely-synchronous model, where global computation phases alternate with global communication phases. This mechanism is most useful for loosely synchronous algorithms where the graph of communicating processes changes very frequently, or where each process communicates with many others.

This call is used for active target communication. An access epoch at an origin process or an exposure epoch at a target process are started and completed by calls to `MPI_WIN_FENCE`. A process can access windows at all processes in the group of `win` during such an access epoch, and the local window can be accessed by all processes in the group of `win` during such an exposure epoch.

2. The four functions `MPI_WIN_START`, `MPI_WIN_COMPLETE`, `MPI_WIN_POST` and `MPI_WIN_WAIT` can be used to restrict synchronization to the minimum: only pairs of communicating processes synchronize, and they do so only when a synchronization is needed to order correctly RMA accesses to a window with respect to local accesses to that same window. This mechanism may be more efficient when each process communicates with few (logical) neighbors, and the communication graph is fixed or changes infrequently.

These calls are used for active target communication. An access epoch is started at the origin process by a call to `MPI_WIN_START` and is terminated by a call to `MPI_WIN_COMPLETE`. The start call has a group argument that specifies the group of target processes for that epoch. An exposure epoch is started at the target process by a call to `MPI_WIN_POST` and is completed by a call to `MPI_WIN_WAIT`. The post call has a group argument that specifies the set of origin processes for that epoch.

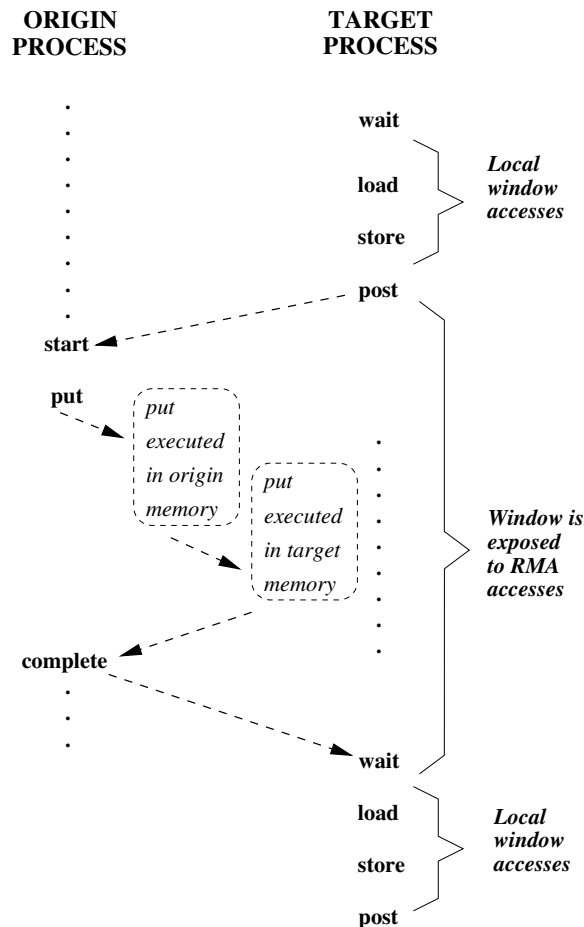


Figure 11.2: Active target communication. Dashed arrows represent synchronizations (ordering of events).

3. [Finally, shared and exclusive locks are provided by the two functions `MPI_WIN_LOCK` and `MPI_WIN_UNLOCK`.] Finally, shared lock access is provided by the functions `MPI_WIN_LOCK`, `MPI_WIN_LOCK_ALL`, `MPI_WIN_UNLOCK`, and `MPI_WIN_UNLOCK_ALL`. `MPI_WIN_LOCK` and `MPI_WIN_UNLOCK` also provide exclusive lock capability. Lock synchronization is useful for MPI applications that emulate a shared memory model via MPI calls; e.g., in a “billboard” model, where processes can, at random times, access or update different parts of the billboard.

These [two]four calls provide passive target communication. An access epoch is started by a call to `MPI_WIN_LOCK` or `MPI_WIN_LOCK_ALL` and terminated by a call to `MPI_WIN_UNLOCK` or `MPI_WIN_UNLOCK_ALL`, respectively. [Only one target window can be accessed during that epoch with `win`. ]

Figure 11.2 illustrates the general synchronization pattern for active target communication. The synchronization between `post` and `start` ensures that the `put` call of the origin process does not start until the target process exposes the window (with the `post` call); the target process will expose the window only after preceding local accesses to the window have completed. The synchronization between `complete` and `wait` ensures that the `put` call of the origin process completes before the window is unexposed (with the `wait` call). The

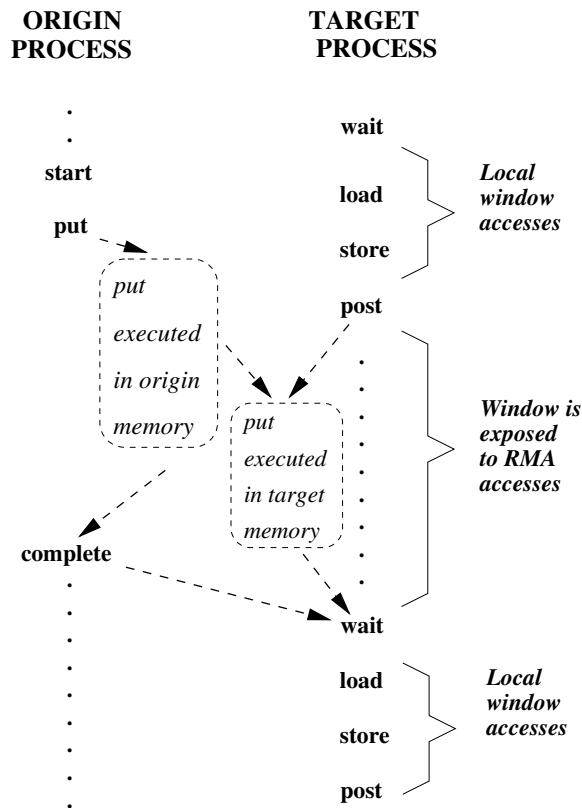


Figure 11.3: Active target communication, with weak synchronization. Dashed arrows represent synchronizations (ordering of events)

target process will execute following local accesses to the target window only after the **wait** returned.

Figure 11.2 shows operations occurring in the natural temporal order implied by the synchronizations: the **post** occurs before the matching **start**, and **complete** occurs before the matching **wait**. However, such **strong** synchronization is more than needed for correct ordering of window accesses. The semantics of MPI calls allow **weak** synchronization, as illustrated in Figure 11.3. The access to the target window is delayed until the window is exposed, after the **post**. However the **start** may complete earlier; the **put** and **complete** may also terminate earlier, if put data is buffered by the implementation. The synchronization calls order correctly window accesses, but do not necessarily synchronize other operations. This weaker synchronization semantic allows for more efficient implementations.

Figure 11.4 illustrates the general synchronization pattern for passive target communication. The first origin process communicates data to the second origin process, through the memory of the target process; the target process is not explicitly involved in the communication. The **lock** and **unlock** calls ensure that the two RMA accesses do not occur concurrently. However, they do *not* ensure that the **put** by origin 1 will precede the **get** by origin 2.

*Rationale.* RMA does not define fine-grained mutexes in memory (only logical coarse-grained process locks). MPI provides the primitives (compare and swap, accumulate[s], send/[recv]receive, etc.) needed to implement high-level synchronization operations. (End of rationale.)

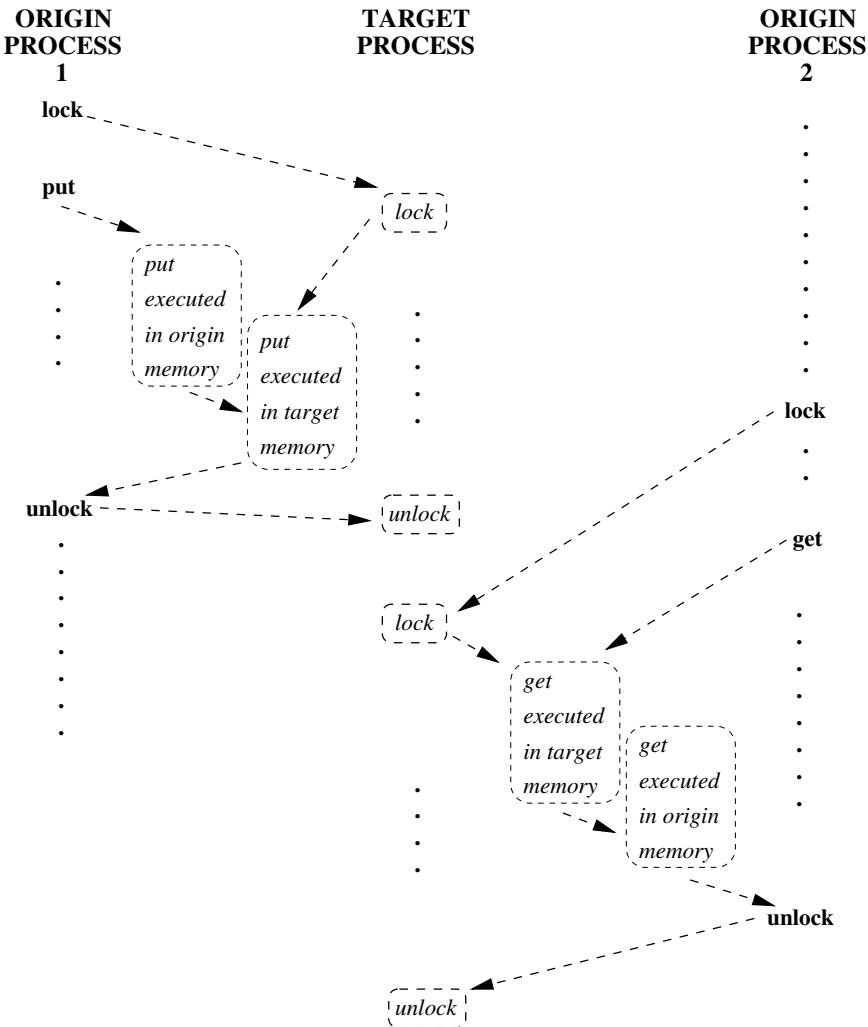


Figure 11.4: Passive target communication. Dashed arrows represent synchronizations (ordering of events).

11.5.1 Fence

MPI\_WIN\_FENCE(assert, win)

IN	assert	program assertion (integer)
IN	win	window object (handle)

int MPI\_Win\_fence(int assert, MPI\_Win win)

```
MPI_Win_fence(assert, win, ierror) BIND(C)
  INTEGER, INTENT(IN) :: assert
  TYPE(MPI_Win), INTENT(IN) :: win
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

MPI\_WIN\_FENCE(ASSERT, WIN, IERROR)

```

1      INTEGER ASSERT, WIN, IERROR
2
3      {void MPI::Win::Fence(int assert) const(binding deprecated, see Section 15.2) }
```

The MPI call MPI\_WIN\_FENCE(assert, win) synchronizes RMA calls on win. The call is collective on the group of win. All RMA operations on win originating at a given process and started before the fence call will complete at that process before the fence call returns. They will be completed at their target before the fence call returns at the target. RMA operations on win started by a process after the fence call returns will access their target window only after MPI\_WIN\_FENCE has been called by the target process.

The call completes an RMA access epoch if it was preceded by another fence call and the local process issued RMA communication calls on win between these two calls. The call completes an RMA exposure epoch if it was preceded by another fence call and the local window was the target of RMA accesses between these two calls. The call starts an RMA access epoch if it is followed by another fence call and by RMA communication calls issued between these two fence calls. The call starts an exposure epoch if it is followed by another fence call and the local window is the target of RMA accesses between these two fence calls. Thus, the fence call is equivalent to calls to a subset of **post**, **start**, **complete**, **wait**.

A fence call usually entails a barrier synchronization: a process completes a call to MPI\_WIN\_FENCE only after all other processes in the group entered their matching call. However, a call to MPI\_WIN\_FENCE that is known not to end any epoch (in particular, a call with `assert = MPI_MODE_NOPRECEDE`) does not necessarily act as a barrier.

The `assert` argument is used to provide assertions on the context of the call that may be used for various optimizations. This is described in Section 11.5.5. A value of `assert = 0` is always valid.

*Advice to users.* Calls to MPI\_WIN\_FENCE should both precede and follow calls to **[put, get or accumulate]RMA communication functions** that are synchronized with fence calls. (*End of advice to users.*)

## 11.5.2 General Active Target Synchronization

```

33 MPI_WIN_START(group, assert, win)
```

IN	group	group of target processes (handle)
IN	assert	program assertion (integer)
IN	win	window object (handle)

```

39 int MPI_Win_start(MPI_Group group, int assert, MPI_Win win)
```

```

41 MPI_Win_start(group, assert, win, ierror) BIND(C)
42     TYPE(MPI_Group), INTENT(IN) :: group
43     INTEGER, INTENT(IN) :: assert
44     TYPE(MPI_Win), INTENT(IN) :: win
45     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```

47 MPI_WIN_START(GROUP, ASSERT, WIN, IERROR)
48     INTEGER GROUP, ASSERT, WIN, IERROR
```



```
{void MPI::Win::Start(const MPI::Group& group, int assert) const(binding
    deprecated, see Section 15.2) }
```

Starts an RMA access epoch for `win`. RMA calls issued on `win` during this epoch must access only windows at processes in `group`. Each process in `group` must issue a matching call to `MPI_WIN_POST`. RMA accesses to each target window will be delayed, if necessary, until the target process executed the matching call to `MPI_WIN_POST`. `MPI_WIN_START` is allowed to block until the corresponding `MPI_WIN_POST` calls are executed, but is not required to.

The `assert` argument is used to provide assertions on the context of the call that may be used for various optimizations. This is described in Section 11.5.5. A value of `assert = 0` is always valid.

`MPI_WIN_COMPLETE(win)`

IN            `win`                            window object (handle)

```
int MPI_Win_complete(MPI_Win win)
```

```
MPI_Win_complete(win, ierror) BIND(C)
    TYPE(MPI_Win), INTENT(IN) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_WIN_COMPLETE(WIN, IERROR)
    INTEGER WIN, IERROR
```

```
{void MPI::Win::Complete() const(binding deprecated, see Section 15.2) }
```

Completes an RMA access epoch on `win` started by a call to `MPI_WIN_START`. All RMA communication calls issued on `win` during this epoch will have completed at the origin when the call returns.

`MPI_WIN_COMPLETE` enforces completion of preceding RMA calls at the origin, but not at the target. A put or accumulate call may not have completed at the target when it has completed at the origin.

Consider the sequence of calls in the example below.

#### Example 11.4

```
MPI_Win_start(group, flag, win);
MPI_Put(...,win);
MPI_Win_complete(win);
```

The call to `MPI_WIN_COMPLETE` does not return until the put call has completed at the origin; and the target window will be accessed by the put operation only after the call to `MPI_WIN_START` has matched a call to `MPI_WIN_POST` by the target process. This still leaves much choice to implementors. The call to `MPI_WIN_START` can block until the matching call to `MPI_WIN_POST` occurs at all target processes. One can also have implementations where the call to `MPI_WIN_START` is nonblocking, but the call to `MPI_PUT` blocks until the matching call to `MPI_WIN_POST` occurred; or implementations where the first two calls are nonblocking, but the call to `MPI_WIN_COMPLETE` blocks until the call to `MPI_WIN_POST` occurred; or even implementations where all three calls

can complete before any target process called `MPI_WIN_POST` — the data put must be buffered, in this last case, so as to allow the put to complete at the origin ahead of its completion at the target. However, once the call to `MPI_WIN_POST` is issued, the sequence above must complete, without further dependencies.

`MPI_WIN_POST(group, assert, win)`

IN	group	group of origin processes (handle)
IN	assert	program assertion (integer)
IN	win	window object (handle)

`int MPI_Win_post(MPI_Group group, int assert, MPI_Win win)`

```

MPI_Win_post(group, assert, win, ierror) BIND(C)
    TYPE(MPI_Group), INTENT(IN) :: group
    INTEGER, INTENT(IN) :: assert
    TYPE(MPI_Win), INTENT(IN) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

`MPI_WIN_POST(GROUP, ASSERT, WIN, IERROR)`  
`INTEGER GROUP, ASSERT, WIN, IERROR`

`{void MPI::Win::Post(const MPI::Group& group, int assert) const` (*binding deprecated, see Section 15.2*) `}`

Starts an RMA exposure epoch for the local window associated with `win`. Only processes in `group` should access the window with RMA calls on `win` during this epoch. Each process in `group` must issue a matching call to `MPI_WIN_START`. `MPI_WIN_POST` does not block.

`MPI_WIN_WAIT(win)`

IN	win	window object (handle)
----	-----	------------------------

`int MPI_Win_wait(MPI_Win win)`

```

MPI_Win_wait(win, ierror) BIND(C)
    TYPE(MPI_Win), INTENT(IN) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

`MPI_WIN_WAIT(WIN, IERROR)`  
`INTEGER WIN, IERROR`

`{void MPI::Win::Wait() const` (*binding deprecated, see Section 15.2*) `}`

Completes an RMA exposure epoch started by a call to `MPI_WIN_POST` on `win`. This call matches calls to `MPI_WIN_COMPLETE(win)` issued by each of the origin processes that were granted access to the window during this epoch. The call to `MPI_WIN_WAIT` will block until all matching calls to `MPI_WIN_COMPLETE` have occurred. This guarantees that all these origin processes have completed their RMA accesses to the local window. When the call returns, all these RMA accesses will have completed at the target window.

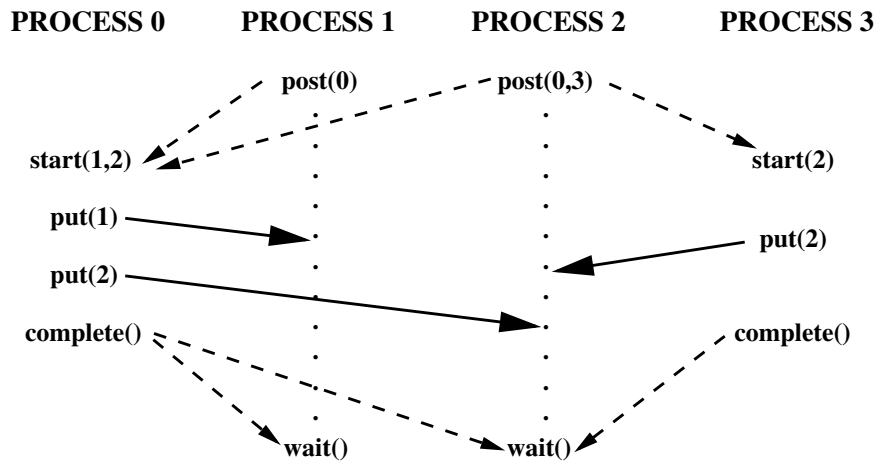


Figure 11.5: Active target communication. Dashed arrows represent synchronizations and solid arrows represent data transfer.

Figure 11.5 illustrates the use of these four functions. Process 0 puts data in the windows of processes 1 and 2 and process 3 puts data in the window of process 2. Each start call lists the ranks of the processes whose windows will be accessed; each post call lists the ranks of the processes that access the local window. The figure illustrates a possible timing for the events, assuming strong synchronization; in a weak synchronization, the start, put or complete calls may occur ahead of the matching post calls.

**MPI\_WIN\_TEST(win, flag)**

IN	win	window object (handle)
OUT	flag	success flag (logical)

```
int MPI_Win_test(MPI_Win win, int *flag)
```

```
MPI_Win_test(win, flag, ierror) BIND(C)
  TYPE(MPI_Win), INTENT(IN) :: win
  LOGICAL, INTENT(OUT) :: flag
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_WIN_TEST(WIN, FLAG, IERROR)
  INTEGER WIN, IERROR
  LOGICAL FLAG
```

```
{bool MPI::Win::Test() const(binding deprecated, see Section 15.2) }
```

This is the nonblocking version of MPI\_WIN\_WAIT. It returns `flag = true` if all accesses to the local window by the group to which it was exposed by the corresponding MPI\_WIN\_POST call have been completed as signalled by matching MPI\_WIN\_COMPLETE calls, and `flag = false` otherwise. In the former case MPI\_WIN\_WAIT would have returned immediately. The effect of return of MPI\_WIN\_TEST with `flag = true` is the same as the effect of a return of MPI\_WIN\_WAIT. If `flag = false` is returned, then the call has no visible

effect.

`MPI_WIN_TEST` should be invoked only where `MPI_WIN_WAIT` can be invoked. Once the call has returned `flag = true`, it must not be invoked anew, until the window is posted anew.

Assume that window `win` is associated with a “hidden” communicator `wincomm`, used for communication by the processes of `win`. The rules for matching of post and start calls and for matching complete and wait call can be derived from the rules for matching sends and receives, by considering the following (partial) model implementation.

**`MPI_WIN_POST(group,0,win)`** initiate a nonblocking send with tag `tag0` to each process in `group`, using `wincomm`. No need to wait for the completion of these sends.

**`MPI_WIN_START(group,0,win)`** initiate a nonblocking receive with tag `tag0` from each process in `group`, using `wincomm`. An RMA access to a window in target process `i` is delayed until the receive from `i` is completed.

**`MPI_WIN_COMPLETE(win)`** initiate a nonblocking send with tag `tag1` to each process in the group of the preceding start call. No need to wait for the completion of these sends.

**`MPI_WIN_WAIT(win)`** initiate a nonblocking receive with tag `tag1` from each process in the group of the preceding post call. Wait for the completion of all receives.

No races can occur in a correct program: each of the sends matches a unique receive, and vice versa.

*Rationale.* The design for general active target synchronization requires the user to provide complete information on the communication pattern, at each end of a communication link: each origin specifies a list of targets, and each target specifies a list of origins. This provides maximum flexibility (hence, efficiency) for the implementor: each synchronization can be initiated by either side, since each “knows” the identity of the other. This also provides maximum protection from possible races. On the other hand, the design requires more information than RMA needs, in general: in general, it is sufficient for the origin to know the rank of the target, but not vice versa. Users that want more “anonymous” communication will be required to use the fence or lock mechanisms. (*End of rationale.*)

*Advice to users.* Assume a communication pattern that is represented by a directed graph  $G = \langle V, E \rangle$ , where  $V = \{0, \dots, n-1\}$  and  $ij \in E$  if origin process  $i$  accesses the window at target process  $j$ . Then each process  $i$  issues a call to `MPI_WIN_POST(ingroupi, ...)`, followed by a call to `MPI_WIN_START(outgroupi, ...)`, where  $outgroup_i = \{j : ij \in E\}$  and  $ingroup_i = \{j : ji \in E\}$ . A call is a noop, and can be skipped, if the group argument is empty. After the communications calls, each process that issued a start will issue a complete. Finally, each process that issued a post will issue a wait.

Note that each process may call with a group argument that has different members. (*End of advice to users.*)

## 11.5.3 Lock

MPI\_WIN\_LOCK(lock\_type, rank, assert, win)

IN	lock_type	either MPI_LOCK_EXCLUSIVE or MPI_LOCK_SHARED (state)
IN	rank	rank of locked window (non-negative integer)
IN	assert	program assertion (integer)
IN	win	window object (handle)

int MPI\_Win\_lock(int lock\_type, int rank, int assert, MPI\_Win win)

MPI\_Win\_lock(lock\_type, rank, assert, win, ierror) BIND(C)

INTEGER, INTENT(IN) :: lock\_type, rank, assert

TYPE(MPI\_Win), INTENT(IN) :: win

INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI\_WIN\_LOCK(LOCK\_TYPE, RANK, ASSERT, WIN, IERROR)

INTEGER LOCK\_TYPE, RANK, ASSERT, WIN, IERROR

{void MPI::Win::Lock(int lock\_type, int rank, int assert) const(*binding  
deprecated, see Section 15.2*) }

Starts an RMA access epoch. Only the window at the process with rank `rank` can be accessed by RMA operations on `win` during that epoch.

MPI\_WIN\_LOCK\_ALL(assert, win)

IN	assert	program assertion (integer)
----	--------	-----------------------------

IN	win	window object (handle)
----	-----	------------------------

int MPI\_Win\_lock\_all(int assert, MPI\_Win win)

MPI\_Win\_lock\_all(assert, win, ierror) BIND(C)

INTEGER, INTENT(IN) :: assert

TYPE(MPI\_Win), INTENT(IN) :: win

INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI\_WIN\_LOCK\_ALL(ASSERT, WIN, IERROR)

INTEGER ASSERT, WIN, IERROR

Starts an RMA access epoch to all processes in `win`, with a lock type of `MPI_LOCK_SHARED`. During the epoch, the calling process can access the window memory on all processes in `win` by using RMA operations. A window locked with `MPI_WIN_LOCK_ALL` must be unlocked with `MPI_WIN_UNLOCK_ALL`. This routine is not collective — the `ALL` refers to a lock on all members of the group of the window.

*Advice to users.* There may be additional overheads associated with using `MPI_WIN_LOCK` and `MPI_WIN_LOCK_ALL` concurrently on the same window. These

overheads could be avoided by specifying the assertion `MPI_MODE_NOCHECK` when possible (see Section 11.5.5). (*End of advice to users.*)

`MPI_WIN_UNLOCK(rank, win)`

IN rank rank of window (non-negative integer)

IN win window object (handle)

`int MPI_Win_unlock(int rank, MPI_Win win)`

`MPI_Win_unlock(rank, win, ierror) BIND(C)`

INTEGER, INTENT(IN) :: rank

TYPE(MPI\_Win), INTENT(IN) :: win

INTEGER, OPTIONAL, INTENT(OUT) :: ierror

`MPI_WIN_UNLOCK(RANK, WIN, IERROR)`

INTEGER RANK, WIN, IERROR

{void MPI::Win::Unlock(int rank) const(*binding deprecated, see Section 15.2*) }

Completes an RMA access epoch started by a call to `MPI_WIN_LOCK(...,win)`. RMA operations issued during this period will have completed both at the origin and at the target when the call returns.

`MPI_WIN_UNLOCK_ALL(win)`

IN win window object (handle)

`int MPI_Win_unlock_all(MPI_Win win)`

`MPI_Win_unlock_all(win, ierror) BIND(C)`

TYPE(MPI\_Win), INTENT(IN) :: win

INTEGER, OPTIONAL, INTENT(OUT) :: ierror

`MPI_WIN_UNLOCK_ALL(WIN, IERROR)`

INTEGER WIN, IERROR

Completes a shared RMA access epoch started by a call to `MPI_WIN_LOCK_ALL(assert, win)`. RMA operations issued during this epoch will have completed both at the origin and at the target when the call returns.

Locks are used to protect accesses to the locked target window effected by RMA calls issued between the lock and unlock calls, and to protect [local] load/store accesses to a locked local or shared memory window executed between the lock and unlock call. Accesses that are protected by an exclusive lock will not be concurrent at the window site with other accesses to the same window that are lock protected. Accesses that are protected by a shared lock will not be concurrent at the window site with accesses protected by an exclusive lock to the same window.

It is erroneous to have a window locked and exposed (in an exposure epoch) concurrently. [I.e.]For example, a process may not call `MPI_WIN_LOCK` to lock a target window

if the target process has called `MPI_WIN_POST` and has not yet called `MPI_WIN_WAIT`; it is erroneous to call `MPI_WIN_POST` while the local window is locked.

*Rationale.* An alternative is to require MPI to enforce mutual exclusion between exposure epochs and locking periods. But this would entail additional overheads when locks or active target synchronization do not interact in support of those rare interactions between the two mechanisms. The programming style that we encourage here is that a set of windows is used with only one synchronization mechanism at a time, with shifts from one mechanism to another being rare and involving global synchronization. (*End of rationale.*)

*Advice to users.* Users need to use explicit synchronization code in order to enforce mutual exclusion between locking periods and exposure epochs on a window. (*End of advice to users.*)

Implementors may restrict the use of RMA communication that is synchronized by lock calls to windows in memory allocated by `MPI_ALLOC_MEM` (Section 8.2, page 352), `MPI_WIN_ALLOCATE` (Section 11.2.2, page 423), or attached with `MPI_WIN_ATTACH` (Section 11.2.4, page 426). Locks can be used portably only in such memory.

*Rationale.* The implementation of passive target communication when memory is not shared [requires] may require an asynchronous software agent. Such an agent can be implemented more easily, and can achieve better performance, if restricted to specially allocated memory. It can be avoided altogether if shared memory is used. It seems natural to impose restrictions that allows one to use shared memory for [3-rd] third party communication in shared memory machines.

The downside of this decision is that passive target communication cannot be used without taking advantage of nonstandard Fortran features: namely, the availability of C-like pointers; these are not supported by some Fortran compilers[(g77 and Windows/NT compilers, at the time of writing)]. [Also, passive target communication cannot be portably targeted to COMMON blocks, or other statically declared Fortran arrays.] (*End of rationale.*)

Consider the sequence of calls in the example below.

### Example 11.5

```
MPI_Win_lock(MPI_LOCK_EXCLUSIVE, rank, assert, win);
MPI_Put(..., rank, ..., win);
MPI_Win_unlock(rank, win);
```

The call to `MPI_WIN_UNLOCK` will not return until the put transfer has completed at the origin and at the target. This still leaves much freedom to implementors. The call to `MPI_WIN_LOCK` may block until an exclusive lock on the window is acquired; or, the call to `MPI_WIN_LOCK` may not block, while the call to `MPI_PUT` blocks until a lock is acquired; or, the first two calls may not block, while `MPI_WIN_UNLOCK` blocks until a lock is acquired — the update of the target window is then postponed until the call to `MPI_WIN_UNLOCK` occurs. However, if the call to `MPI_WIN_LOCK` is used to lock a local window, then the call must block until the lock is acquired, since the lock may protect local load/store accesses to the window issued after the lock call returns.

## 11.5.4 Flush and Sync

All flush and sync functions can be called only within ~~lock-unlock or lockall-unlockall~~ epochs.

### MPI\_WIN\_FLUSH(rank, win)

IN rank rank of target window (non-negative integer)  
IN win window object (handle)

```
int MPI_Win_flush(int rank, MPI_Win win)
```

```
MPI_Win_flush(rank, win, ierror) BIND(C)
  INTEGER, INTENT(IN) :: rank
  TYPE(MPI_Win), INTENT(IN) :: win
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_WIN_FLUSH(RANK, WIN, IERROR)
  INTEGER RANK, WIN, IERROR
```

MPI\_WIN\_FLUSH completes all outstanding RMA operations initiated by the calling process to the target rank on the specified window. The operations are completed both at the origin and at the target. ~~Flush completes locally in the sense used in this document, meaning that the call must return without requiring the target process to call any MPI routine.~~

### MPI\_WIN\_FLUSH\_ALL(win)

IN win window object (handle)

```
int MPI_Win_flush_all(MPI_Win win)
```

```
MPI_Win_flush_all(win, ierror) BIND(C)
  TYPE(MPI_Win), INTENT(IN) :: win
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_WIN_FLUSH_ALL(WIN, IERROR)
  INTEGER WIN, IERROR
```

All RMA operations issued by the calling process to any target on the specified window prior to this call and in the specified window will have completed both at the origin and at the target when this call returns. ~~MPI\_WIN\_FLUSH\_ALL completes locally in the sense used in this document, meaning that the call must return without requiring the target processes to call any MPI routine.~~

### MPI\_WIN\_FLUSH\_LOCAL(rank, win)

IN rank rank of target window (non-negative integer)  
IN win window object (handle)



```
int MPI_Win_flush_local(int rank, MPI_Win win)
```

```
MPI_Win_flush_local(rank, win, ierror) BIND(C)
```

```
    INTEGER, INTENT(IN) :: rank
```

```
    TYPE(MPI_Win), INTENT(IN) :: win
```

```
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_WIN_FLUSH_LOCAL(RANK, WIN, IERROR)
```

```
    INTEGER RANK, WIN, IERROR
```

Locally completes at the origin all outstanding RMA operations initiated by the calling process to the target process specified by rank on the specified window. For example, after this routine completes, the user may reuse any buffers provided to put, get, or accumulate operations. ~~MPI\_WIN\_FLUSH\_LOCAL completes locally in the sense used in this document, meaning that the call must return without requiring the target processes to call any MPI routine.~~

```
MPI_WIN_FLUSH_LOCAL_ALL(win)
```

```
    IN          win                      window object (handle)
```

```
int MPI_Win_flush_local_all(MPI_Win win)
```

```
MPI_Win_flush_local_all(win, ierror) BIND(C)
```

```
    TYPE(MPI_Win), INTENT(IN) :: win
```

```
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_WIN_FLUSH_LOCAL_ALL(WIN, IERROR)
```

```
    INTEGER WIN, IERROR
```

All RMA operations issued to any target prior to this call in this window will have completed at the origin when MPI\_WIN\_FLUSH\_LOCAL\_ALL returns.

~~MPI\_WIN\_FLUSH\_LOCAL\_ALL completes locally in the sense used in this document, meaning that the call must return without requiring the target processes to call any MPI routine.~~

```
MPI_WIN_SYNC(win)
```

```
    IN          win                      window object (handle)
```

```
int MPI_Win_sync(MPI_Win win)
```

```
MPI_Win_sync(win, ierror) BIND(C)
```

```
    TYPE(MPI_Win), INTENT(IN) :: win
```

```
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_WIN_SYNC(WIN, IERROR)
```

```
    INTEGER WIN, IERROR
```

The call MPI\_WIN\_SYNC synchronizes the private and public window copy of win. For the purposes of synchronizing the private and public window, MPI\_WIN\_SYNC has the effect of ending and reopening an access and exposure epoch on the window (note that it

ticket0. 1 does not actually end an epoch or complete any pending [MPI]MPI RMA operations).

### 11.5.5 Assertions

ticket270. 5 The `assert` argument in the calls `MPI_WIN_POST`, `MPI_WIN_START`,  
 ticket270. 6 `MPI_WIN_FENCE` [and], `MPI_WIN_LOCK`, and `MPI_WIN_LOCK_ALL` is used to provide as-  
 ticket270. 7 sertions on the context of the call that may be used to optimize performance. The `assert`  
 ticket270. 8 argument does not change program semantics if it provides correct information on the pro-  
 ticket270. 9 gram — it is erroneous to provide [s] incorrect information. Users may always provide `assert`  
 ticket270. 10 `= 0` to indicate a general case [,] where no guarantees are made.

ticket270. 12 *Advice to users.* Many implementations may not take advantage of the information in  
 ticket270. 13 `assert`; some of the information is relevant only for noncoherent [,] shared memory ma-  
 ticket270. 14 chines. Users should consult their implementation manual to find which information  
 ticket270. 15 is useful on each system. On the other hand, applications that provide correct asser-  
 ticket270. 16 tions whenever applicable are portable and will take advantage of assertion specific  
 ticket270. 17 optimizations [,] whenever available. (*End of advice to users.*)

ticket270. 18 *Advice to implementors.* Implementations can always ignore the  
 ticket270. 19 `assert` argument. Implementors should document which `assert` values are significant  
 ticket270. 20 on their implementation. (*End of advice to implementors.*)

ticket270. 22 `assert` is the bit-vector OR of zero or more of the following integer constants:  
 ticket270. 23 `MPI_MODE_NOCHECK`, `MPI_MODE_NOSTORE`, `MPI_MODE_NOPUT`,  
 ticket270. 24 `MPI_MODE_NOPRECEDE` and `MPI_MODE_NOSUCCEED`. The significant options are listed  
 ticket270. 25 below [,] for each call.

ticket270. 27 *Advice to users.* C/C++ users can use bit vector or (|) to combine these constants;  
 ticket270. 28 Fortran 90 users can use the bit-vector `IOR` intrinsic. Fortran 77 users can use (non-  
 ticket270. 29 portably) bit vector `IOR` on systems that support it. Alternatively, Fortran users can  
 ticket270. 30 portably use integer addition to OR the constants (each constant should appear at  
 ticket270. 31 most once in the addition!). (*End of advice to users.*)

#### **MPI\_WIN\_START:**

ticket270. 35 `MPI_MODE_NOCHECK` — the matching calls to `MPI_WIN_POST` have already com-  
 ticket270. 36 pleted on all target processes when the call to `MPI_WIN_START` is made. The  
 ticket270. 37 `nocheck` option can be specified in a start call if and only if it is specified in  
 ticket270. 38 each matching post call. This is similar to the optimization of “ready-send” that  
 ticket270. 39 may save a handshake when the handshake is implicit in the code. (However,  
 ticket270. 40 ready-send is matched by a regular receive, whereas both start and post must  
 ticket270. 41 specify the `nocheck` option.)

#### **MPI\_WIN\_POST:**

ticket270. 44 `MPI_MODE_NOCHECK` — the matching calls to `MPI_WIN_START` have not yet oc-  
 ticket270. 45 curred on any origin processes when the call to `MPI_WIN_POST` is made. The  
 ticket270. 46 `nocheck` option can be specified by a post call if and only if it is specified by each  
 ticket270. 47 matching start call.

ticket284.

`MPI_MODE_NOSTORE` — the local window was not updated by [\[local\]](#) stores (or local get or receive calls) since last synchronization. This may avoid the need for cache synchronization at the post call.

`MPI_MODE_NOPUT` — the local window will not be updated by put or accumulate calls after the post call, until the ensuing (wait) synchronization. This may avoid the need for cache synchronization at the wait call.

#### `MPI_WIN_FENCE:`

`MPI_MODE_NOSTORE` — the local window was not updated by [\[local\]](#) stores (or local get or receive calls) since last synchronization.

`MPI_MODE_NOPUT` — the local window will not be updated by put or accumulate calls after the fence call, until the ensuing (fence) synchronization.

`MPI_MODE_NOPRECEDE` — the fence does not complete any sequence of locally issued RMA calls. If this assertion is given by any process in the window group, then it must be given by all processes in the group.

[\[MPI\\_MODE\\_NOSUCCEED\]](#)`MPI_MODE_NOSUCCEED` — the fence does not start any sequence of locally issued RMA calls. If the assertion is given by any process in the window group, then it must be given by all processes in the group.

#### `MPI_WIN_LOCK`, `MPI_WIN_LOCK_ALL:`

`MPI_MODE_NOCHECK` — no other process holds, or will attempt to acquire a conflicting lock, while the caller holds the window lock. This is useful when mutual exclusion is achieved by other means, but the coherence operations that may be attached to the lock and unlock calls are still required.

*Advice to users.* Note that the nostore and noprecede flags provide information on what happened *before* the call; the noput and nosucceed flags provide information on what will happen *after* the call. (*End of advice to users.*)

### 11.5.6 Miscellaneous Clarifications

Once an RMA routine completes, it is safe to free any opaque objects passed as argument to that routine. For example, the `datatype` argument of a `MPI_PUT` call can be freed as soon as the call returns, even though the communication may not be complete.

As in message-passing, datatypes must be committed before they can be used in RMA communication.

[\[\[Moved: Section on Examples\]\]](#)

## 11.6 Error Handling

### 11.6.1 Error Handlers

Errors occurring during calls to [\[MPI\\_WIN\\_CREATE\(...,comm,...\)\]](#) routines that create MPI windows (e.g., `MPI_WIN_CREATE(...,comm,...)`) cause the error handler currently associated with `comm` to be invoked. All other RMA calls have an input `win` argument. When an error occurs during such a call, the error handler currently associated with `win` is invoked.

The default error handler associated with `win` is `MPI_ERRORS_ARE_FATAL`. Users may change this default by explicitly associating a new error handler with `win` (see Section 8.3, page 355).

### 11.6.2 Error Classes

The [following] error classes for one-sided communication are defined in Table 11.2. RMA routines may (and almost certainly will) use other MPI error classes, such as `MPI_ERR_OP` or `MPI_ERR_RANK`.

<code>MPI_ERR_WIN</code>	invalid <code>win</code> argument
<code>MPI_ERR_BASE</code>	invalid <code>base</code> argument
<code>MPI_ERR_SIZE</code>	invalid <code>size</code> argument
<code>MPI_ERR_DISP</code>	invalid <code>disp</code> argument
<code>MPI_ERR_LOCKTYPE</code>	invalid <code>locktype</code> argument
<code>MPI_ERR_ASSERT</code>	invalid <code>assert</code> argument
<code>MPI_ERR_RMA_CONFLICT</code>	conflicting accesses to window
<code>MPI_ERR_RMA_SYNC</code>	[ticket270.][wrong]invalid synchronization of RMA calls
[ticket270.] <code>MPI_ERR_RMA_RANGE</code>	[ticket270.]target memory is not part of the window (in the case of a window created with <code>MPI_WIN_CREATE_DYNAMIC</code> , target memory is not attached)
[ticket270.] <code>MPI_ERR_RMA_ATTACH</code>	[ticket270.]memory cannot be attached (e.g., because of resource exhaustion)
[ticket284.] <code>MPI_ERR_RMA_SHARED</code>	[ticket284.]memory cannot be shared (e.g., some process in the group of the specified communicator cannot expose shared memory)
[ticket284.] <code>MPI_ERR_RMA_WRONG_FLAVOR</code>	[ticket284.]passed window has the wrong flavor for the called function

Table 11.2: Error classes in one-sided communication routines

### 11.7 Semantics and Correctness

[The semantics of RMA operations is best understood by assuming that the system maintains a separate *public* copy of each window, in addition to the original location in process memory (the *private* window copy). There is only one instance of each variable in process memory, but a distinct *public* copy of the variable for each window that contains it. A load accesses the instance in process memory (this includes MPI sends). A store accesses and updates the instance in process memory (this includes MPI receives), but the update may affect other public copies of the same locations. A get on a window accesses the public copy of that window. A put or accumulate on a window accesses and updates the public copy of that window, but the update may affect the private copy of the same locations in process memory, and public copies of other overlapping windows. This is illustrated in Figure 11.1.]

The following rules specify the latest time at which an operation must complete at the origin or the target. The update performed by a get call in the origin process memory is visible when the get operation is complete at the origin (or earlier); the update performed by a put or accumulate call in the public copy of the target window is visible when the put or accumulate has completed at the target (or earlier). The rules also specify the latest time at which an update of one window copy becomes visible in another overlapping copy.

1. An RMA operation is completed at the origin by the ensuing call to MPI\_WIN\_COMPLETE, MPI\_WIN\_FENCE[ or MPI\_WIN\_UNLOCK], MPI\_WIN\_FLUSH, MPI\_WIN\_FLUSH\_ALL, MPI\_WIN\_FLUSH\_LOCAL, MPI\_WIN\_FLUSH\_LOCAL\_ALL, MPI\_WIN\_UNLOCK, or MPI\_WIN\_UNLOCK\_ALL that synchronizes this access at the origin.
2. If an RMA operation is completed at the origin by a call to MPI\_WIN\_FENCE then the operation is completed at the target by the matching call to MPI\_WIN\_FENCE by the target process.
3. If an RMA operation is completed at the origin by a call to MPI\_WIN\_COMPLETE then the operation is completed at the target by the matching call to MPI\_WIN\_WAIT by the target process.
4. If an RMA operation is completed at the origin by a call to MPI\_WIN\_UNLOCK, MPI\_WIN\_UNLOCK\_ALL, MPI\_WIN\_FLUSH(rank=target), or MPI\_WIN\_FLUSH\_ALL, then the operation is completed at the target by that same call[ to MPI\_WIN\_UNLOCK].
5. An update of a location in a private window copy in process memory becomes visible in the public window copy at latest when an ensuing call to MPI\_WIN\_POST, MPI\_WIN\_FENCE, [or MPI\_WIN\_UNLOCK]MPI\_WIN\_UNLOCK, MPI\_WIN\_UNLOCK\_ALL, or MPI\_WIN\_SYNC is executed on that window by the window owner. In the RMA unified memory model, an update of a location in a private window in process memory becomes visible without additional RMA calls.
6. An update by a put or accumulate call to a public window copy becomes visible in the private copy in process memory at latest when an ensuing call to MPI\_WIN\_WAIT, MPI\_WIN\_FENCE,[ or MPI\_WIN\_LOCK]MPI\_WIN\_LOCK, MPI\_WIN\_LOCK\_ALL, or MPI\_WIN\_SYNC is executed on that window by the window owner. In the RMA unified memory model, an update by a put or accumulate call to a public window copy becomes visible in the private copy in process memory without additional RMA calls.

The MPI\_WIN\_FENCE or MPI\_WIN\_WAIT call that completes the transfer from public copy to private copy (6) is the same call that completes the put or accumulate operation in the window copy (2, 3). If a put or accumulate access was synchronized with a lock, then the update of the public window copy is complete as soon as the updating process executed MPI\_WIN\_UNLOCK or MPI\_WIN\_UNLOCK\_ALL. [On the other hand]In the RMA separate memory model, the update of private copy in the process memory may be delayed until the target process executes a synchronization call on that window (6). Thus, updates to process memory can always be delayed in the RMA separate memory model until the process executes a suitable synchronization call, while they have to complete in the RMA unified

model without additional synchronization calls. [Updates to a public window copy can also be delayed until the window owner executes a synchronization call, if fences or post-start-complete-wait synchronization is used.] If fence or post-start-complete-wait synchronization is used, updates to a public window copy can be delayed in both memory models until the window owner executes a synchronization call. [Only when lock synchronization is used does it become necessary to update the public window copy, even if the window owner does not execute any related synchronization call.] When passive-target synchronization (lock/unlock or even flush) is used, it is necessary to update the public window copy in the RMA separate model, or the private window copy in the RMA unified model, even if the window owner does not execute any related synchronization call.

The rules above also define, by implication, when an update to a public window copy becomes visible in another overlapping public window copy. Consider, for example, two overlapping windows, win1 and win2. A call to `MPI_WIN_FENCE(0, win1)` by the window owner makes visible in the process memory previous updates to window win1 by remote processes. A subsequent call to `MPI_WIN_FENCE(0, win2)` makes these updates visible in the public copy of win2.

The behavior of some MPI RMA operations may be *undefined* in some situations. For example, the result of several origin processes performing concurrent `MPI_PUT` operations to the same target location is undefined. In addition, the result of a single origin process performing multiple `MPI_PUT` operation to the same target location within the same access epoch is also undefined. The result at the target may have all of the data from one of the `MPI_PUT` operations (the “last” one, in some sense), or bytes from some of each of the operations, or something else. In MPI-2, such operations were *erroneous*. That meant that an MPI implementation was permitted to signal an MPI exception. Thus, user programs or tools that used MPI RMA could not portably permit such operations, even if the application code could function correctly with such an undefined result. In MPI-3, these operations are not erroneous, but do not have a defined behavior.

*Rationale.* As discussed in [6], requiring operations such as overlapping puts to be erroneous makes it [very] difficult to use MPI RMA to implement programming models—such as Unified Parallel C (UPC) or SHMEM—that permit these operations. Further, while MPI-2 defined these operations as erroneous, the MPI Forum is unaware of any implementation that enforces this rule, as it would require significant overhead. Thus, relaxing this condition does not impact existing implementations or applications. (*End of rationale.*)

*Advice to implementors.* Overlapping accesses are undefined. However, to assist users in debugging code, implementations may wish to provide a mode in which such operations are detected and reported to the user. Note, however, that in MPI-3, such operations must not generate an MPI exception. (*End of advice to implementors.*)

A [correct program] program with a well-defined outcome in the `MPI_WIN_SEPARATE` memory model must obey the following rules.

1. A location in a window must not be accessed [locally] with load/store operations once an update to that location has started, until the update becomes visible in the private window copy in process memory.
2. A location in a window must not be accessed as a target of an RMA operation once an update to that location has started, until the update becomes visible in the public

window copy. There is one exception to this rule, in the case where the same variable is updated by two concurrent accumulates [that use the same operation, ] with the same predefined datatype, on the same window. Additional restrictions on the operation apply, see the info key `accumulate_ops` in Section 11.2.1.

3. A put or accumulate must not access a target window once a [local]load/store update or a put or accumulate update to another (overlapping) target window [have]has started on a location in the target window, until the update becomes visible in the public copy of the window. Conversely, a [local update in]store to process memory to a location in a window must not start once a put or accumulate update to that target window has started, until the put or accumulate update becomes visible in process memory. In both cases, the restriction applies to operations even if they access disjoint locations in the window.

[A program is erroneous if it violates these rules.]

*Rationale.* The last constraint on correct RMA accesses may seem unduly restrictive, as it forbids concurrent accesses to nonoverlapping locations in a window. The reason for this constraint is that, on some architectures, explicit coherence restoring operations may be needed at synchronization points. A different operation may be needed for locations that were [locally] updated by stores and for locations that were remotely updated by put or accumulate operations. Without this constraint, the MPI library will have to track precisely which locations in a window were updated by a put or accumulate call. The additional overhead of maintaining such information is considered prohibitive. (*End of rationale.*)

Note that `MPI_WIN_SYNC` may be used within a passive target epoch to synchronize the private and public window copies (that is, updates to one are made visible to the other).

In the `MPI_WIN_UNIFIED` memory model, the rules are much simpler because the public and private windows are the same. However, there are restrictions to avoid concurrent access to the same memory locations by different processes. The rules that a program with a well-defined outcome must obey in this case are:

1. A location in a window must not be accessed [locally]with load/store operations once an update to that location has started, until the update is complete, subject to the following special case.
2. [Locally accessing (but not updating)]Accessing a location in the window [with a load operation ]that is also the target of a remote update is valid (not erroneous) but the precise result will depend on the behavior of the implementation. Updates from a remote process will appear in the memory of the target, but there are no atomicity or ordering guarantees if more than one byte is updated. Updates are stable in the sense that once data appears in memory of the target, the data remains until replaced by another update. This permits polling on a location for a change from zero to non-zero or for a particular value, but not polling and comparing the relative magnitude of values. Users are cautioned that polling on one memory location and then accessing a different memory location has defined behavior only if the other rules given here and in this chapter are followed.



*Advice to users.* Some compiler optimizations can result in code that maintains the sequential semantics of the program, but violates this rule by introducing temporary values into locations in memory. Most compilers only apply such transformations under very high levels of optimization and users should be aware that such aggressive optimization may produce unexpected results. (*End of advice to users.*)

3. [Locally u]Updating a location in the window with a store operation that is also the target of a remote read (but not update) is valid (not erroneous) but the precise result will depend on the behavior of the implementation. [Updates from the local process]Store updates will appear in memory, but there are no atomicity or ordering guarantees if more than one byte is updated. Updates are stable in the sense that once data appears in memory, the data remains until replaced by another update. This permits [the local process] to update memory [in its local window] with store operations without requiring a lock/unlock or other RMA synchronization epoch. Users are cautioned that remote accesses to a window that is updated by the local process has defined behavior only if the other rules given here and in this chapter are followed.
4. A location in a window must not be accessed as a target of an RMA operation once an update to that location has started until the update completes at the target. There is one exception to this rule: in the case where the same variable is updated by two concurrent accumulates with the same predefined datatype on the same window. Additional restrictions on the operation apply; see the info key `accumulate_ops` in Section 11.2.1.
5. A put or accumulate must not access a target window once a [local update]store operation or a put or accumulate update to another (overlapping) target window has started on the same location in the target window until the update completes at the target window. Conversely, a [local update]store operation in-process memory to a location in a window must not start once a put or accumulate update to the same location in that target window has started until the put or accumulate update completes at the target.

Note that `MPI_WIN_FLUSH` and `MPI_WIN_FLUSH_ALL` may be used within a passive target epoch to complete RMA operations at the target process.

A program that violates these rules has undefined behavior.

*Advice to users.* A user can write correct programs by following the following rules:

**fence:** During each period between fence calls, each window is either updated by put or accumulate calls, or updated by [local] stores, but not both. Locations updated by put or accumulate calls should not be accessed during the same period (with the exception of concurrent updates to the same location by accumulate calls). Locations accessed by get calls should not be updated during the same period.

**post-start-complete-wait:** A window should not be updated [locally] with store operations while being posted, if it is being updated by put or accumulate calls. Locations updated by put or accumulate calls should not be accessed while the window is posted (with the exception of concurrent updates to the same location



by accumulate calls). Locations accessed by get calls should not be updated while the window is posted.

With the post-start synchronization, the target process can tell the origin process that its window is now ready for RMA access; with the complete-wait synchronization, the origin process can tell the target process that it has finished its RMA accesses to the window.

**lock:** Updates to the window are protected by exclusive locks if they may conflict. Nonconflicting accesses (such as read-only accesses or accumulate accesses) are protected by shared locks, both for `[local]load/store` accesses and for RMA accesses.

**changing window or synchronization mode:** One can change synchronization mode, or change the window used to access a location that belongs to two overlapping windows, when the process memory and the window copy are guaranteed to have the same values. This is true after a local call to `MPI_WIN_FENCE`, if RMA accesses to the window are synchronized with fences; after a local call to `MPI_WIN_WAIT`, if the accesses are synchronized with post-start-complete-wait; after the call at the origin (local or remote) to `MPI_WIN_UNLOCK` or `MPI_WIN_UNLOCK_ALL` if the accesses are synchronized with locks.

In addition, a process should not access the local buffer of a get operation until the operation is complete, and should not update the local buffer of a put or accumulate operation until that operation is complete.

The RMA synchronization operations define when updates are guaranteed to become visible in public and private windows. Updates may become visible earlier, but such behavior is implementation dependent. (*End of advice to users.*)

The semantics are illustrated by the following examples:

**Example 11.6** `[Rule 5:]` The following example demonstrates updating a memory location inside a window for the separate memory model, according to Rule 5. The `MPI_WIN_LOCK` and `MPI_WIN_UNLOCK` calls around the store to X in process B are necessary to ensure consistency between the public and private copies of the window.

Process A:	Process B:
	window location X
	<code>MPI_Win_lock(EXCLUSIVE,B)</code>
	store X /* local update to private copy of B */
	<code>MPI_Win_unlock(B)</code>
	/* now visible in public window copy */
<code>MPI_Barrier</code>	<code>MPI_Barrier</code>
<code>MPI_Win_lock(EXCLUSIVE,B)</code>	
<code>MPI_Get(X) /* ok, read from public window */</code>	
<code>MPI_Win_unlock(B)</code>	

**Example 11.7** In the RMA unified model, although the public and private copies of the windows are synchronized, caution must be used when combining [local] load/stores and multi-process synchronization. Although the following example appears correct, the compiler or hardware may delay the store to X after the barrier, possibly resulting in the MPI\_GET returning the incorrect value of X.

ticket284.

```

Process A:                Process B:
                             window location X

                             store X /* update to private&public copy of B */
MPI_Barrier                MPI_Barrier
MPI_Win_lock_all
MPI_Get(X) /* ok, read from window */
MPI_Win_flush_local(B)
/* read value in X */
MPI_Win_unlock_all

```

MPI\_BARRIER provides process synchronization, but not [local] memory synchronization. The example could potentially be made safe through the use of compiler and hardware specific notations to ensure the store to X occurs before process B enters the MPI\_BARRIER. The use of one-sided synchronization calls, as shown in Example 11.6, also ensures the correct result.

**Example 11.8** [Rule 6:]The following example demonstrates the reading of a memory location updated by a remote process (Rule 6) in the RMA separate memory model. Although the MPI\_WIN\_UNLOCK on process A and the MPI\_BARRIER ensure that the public copy on process B reflects the updated value of X, the call to MPI\_WIN\_LOCK by process B is necessary to synchronize the private copy with the public copy.

```

Process A:                Process B:
                             window location X

MPI_Win_lock(EXCLUSIVE,B)
MPI_Put(X) /* update to public window */
MPI_Win_unlock(B)

MPI_Barrier                MPI_Barrier

                             MPI_Win_lock(EXCLUSIVE,B)
                             /* now visible in private copy of B */
                             load X
                             MPI_Win_unlock(B)

```

Note that in this example, the barrier is not critical to the semantic correctness. The use of exclusive locks guarantees a remote process will not modify the public copy after MPI\_WIN\_LOCK synchronizes the private and public copies. A polling implementation looking for changes in X on process B would be semantically correct. The barrier is required to ensure that process A performs the put operation before process B performs the load of X.

**Example 11.9** Similar to Example 11.7, the following example is unsafe even in the unified model, because the load of X can not be guaranteed to occur after the MPI\_BARRIER. While Process B does not need to explicitly synchronize the public and private copies through MPI\_WIN\_LOCK as the MPI\_PUT will update both the public and private copies of the window, the scheduling of the load could result in old values of X being returned. Compiler and hardware specific notations could ensure the load occurs after the data is updated, or explicit one-sided synchronization calls can be used to ensure the proper result.

```

Process A:                Process B:
                           window location X

MPI_Win_lock_all
MPI_Put(X) /* update to window */
MPI_Win_flush(B)

MPI_Barrier                MPI_Barrier
                           load X

MPI_Win_unlock_all

```

**Example 11.10** [The rules do *not* guarantee that process A in the following sequence will see the value of X as updated by the local store by B before the lock.] The following example further clarifies Rule 5. MPI\_WIN\_LOCK and MPI\_WIN\_LOCK\_ALL do *not* update the public copy of a window with changes to the private copy. Therefore, there is no guarantee that process A in the following sequence will see the value of X as updated by the local store by process B before the lock.

```

Process A:                Process B:
                           window location X

                           store X /* update to private copy of B */
                           MPI_Win_lock(SHARED,B)
MPI_Barrier                MPI_Barrier

MPI_Win_lock(SHARED,B)
MPI_Get(X) /* X may be the X before the store */
MPI_Win_unlock(B)

                           MPI_Win_unlock(B)
                           /* update on X now visible in public window */

```

The addition of an MPI\_WIN\_SYNC before the call to MPI\_BARRIER by process B would guarantee process A would see the updated value of X, as the public copy of the window would be explicitly synchronized with the private copy.

**Example 11.11** [In the following sequence] Similar to the previous example, Rule 5 can have unexpected implications for general active target synchronization with the RMA separate memory model. It is *not* guaranteed that process B reads the value of X as per the local update by process A, because neither MPI\_WIN\_WAIT nor MPI\_WIN\_COMPLETE calls by process A ensure visibility in the public window copy.

```

1   Process A:                Process B:
2   window location X
3   window location Y
4
5   store Y
6   MPI_Win_post(A,B) /* Y visible in public window */
7   MPI_Win_start(A)      MPI_Win_start(A)
8
9   store X /* update to private window */
10
11  MPI_Win_complete      MPI_Win_complete
12  MPI_Win_wait
13  /* update on X may not yet visible in public window */
14
15  MPI_Barrier            MPI_Barrier
16
17                          MPI_Win_lock(EXCLUSIVE,A)
18                          MPI_Get(X) /* may return an obsolete value */
19                          MPI_Get(Y)
20                          MPI_Win_unlock(A)

```

[it is *not* guaranteed that process B reads the value of X as per the local update by process A, because neither `MPI_WIN_WAIT` nor `MPI_WIN_COMPLETE` calls by process A ensure visibility in the public window copy.] To allow process B to read the value of X stored by A the local store must be replaced by a local `MPI_PUT` that updates the public window copy. Note that by this replacement X may become visible in the private copy [in]of process [memory of ]A only after the `MPI_WIN_WAIT` call in process A. The update to Y made before the `MPI_WIN_POST` call is visible in the public window after the `MPI_WIN_POST` call and therefore [correctly gotten by process B]process B will read the proper value of Y. The `MPI_GET(Y)` call could be moved to the epoch started by the `MPI_WIN_START` operation, and process B would still get the value stored by process A.

**Example 11.12** [Finally, in the following sequence]The following example demonstrates the interaction of general active target synchronization with local read operations with the RMA separate memory model. Rules 5 and 6 do *not* guarantee that the private copy of X at process B has been updated before the load takes place.

```

37  Process A:                Process B:
38                              window location X
39
40  MPI_Win_lock(EXCLUSIVE,B)
41  MPI_Put(X) /* update to public window */
42  MPI_Win_unlock(B)
43
44  MPI_Barrier            MPI_Barrier
45
46                          MPI_Win_post(B)
47                          MPI_Win_start(B)
48

```

```

load X /* access to private window */
      /* may return an obsolete value */

MPI_Win_complete
MPI_Win_wait

```

[rules (5,6) do *not* guarantee that the private copy of X at B has been updated before the load takes place.] To ensure that the value put by process A is read, the local load must be replaced with a local MPI\_GET operation, or must be placed after the call to MPI\_WIN\_WAIT.

### 11.7.1 Atomicity

The outcome of concurrent `accumulate[ticket270.][s]` operations to the same location `[ticket270.][,]` with the same `[ticket270.][operation and]` predefined datatype `[ticket270.][,]` is as if the accumulates `[ticket270.][where]` were done at that location in some serial order. `[ticket270.][Additional restrictions on the operation apply, see the info key accumulate_ops in Section 11.2.1. [ticket270.][On the other hand, if two locations are both updated by two accumulate calls, then the updates may occur in reverse order at the two locations]` Concurrent accumulate operations with different origin and target pairs are not ordered. Thus, there is no guarantee that the entire call to `[ticket270.][MPI_ACCUMULATE]` an accumulate operation is executed atomically. The effect of this lack of atomicity is limited: The previous correctness conditions imply that a location updated by a call to `[ticket270.][MPI_ACCUMULATE]` an accumulate operation cannot be accessed by `[ticket270.][a` load or an RMA call other than `accumulate[ticket270.][,]` until the `[ticket270.][MPI_ACCUMULATE call]` accumulate operation has completed (at the target). Different interleavings can lead to different results only to the extent that computer arithmetics are not truly associative or commutative. `[ticket270.][The outcome of accumulate operations with overlapping types of different sizes or target displacements is undefined.`

### 11.7.2 Ordering

Accumulate calls enable element-wise atomic read and write to remote memory locations. MPI specifies ordering between accumulate operations from one process to the same (or overlapping) memory locations at another process on a per-datatype granularity. The default ordering is strict ordering, which guarantees that overlapping updates from the same source to a remote location are committed in program order and that reads (e.g., with MPI\_GET\_ACCUMULATE) and writes (e.g., with MPI\_ACCUMULATE) are executed and committed in program order. Ordering only applies to operations originating at the same origin that access overlapping target memory regions. MPI does not provide any guarantees for accesses or updates from different origins to overlapping target memory regions.

The default strict ordering may incur a significant performance penalty. MPI specifies the info key `accumulate_ordering` to allow relaxation of the ordering semantics when specified to any window creation function. The values for this key are as follows. If set to `none`, then no ordering will be guaranteed for accumulate calls. This was the behavior for RMA in MPI-2 but is *not* the default in MPI-3. The key can be set to a comma-separated list of required access orderings at the target. Allowed values in the comma-separated list are `rar`, `war`, `raw`, and `waw` for read-after-read, write-after-read, read-after-write, and write-after-write ordering, respectively. These indicate whether operations of the specified type complete in

the order they were issued. For example, `raw` means that any writes must complete at the target before any reads. These ordering requirements apply only to operations issued by the same origin process and targeting the same target process. [Note that `rar`, read-after-read, is included for completeness, as ordering is only important if an update (write) may be made.] The default value for `accumulate_ordering` is `rar,raw,war,waw`, which implies that writes complete at the target in the order in which they were issued, reads complete at the target before any writes that are issued after the reads, and writes complete at the target before any reads that are issued after the writes. Any subset of these four orderings can be specified. For example, if only read-after-read and write-after-write ordering is required, then the value of the `accumulate_ordering` key could be set to `rar,waw`. The order of values is not significant.

Note that the above ordering semantics apply only to accumulate operations, not puts and gets. Puts and gets within an epoch are unordered.

### 11.7.3 Progress

One-sided communication has the same progress requirements as point-to-point communication: once a communication is enabled[, then] it is guaranteed to complete. RMA calls must have local semantics, except when required for synchronization with other RMA calls.

There is some fuzziness in the definition of the time when a RMA communication becomes enabled. This fuzziness provides to the implementor more flexibility than with point-to-point communication. Access to a target window becomes enabled once the corresponding synchronization (such as `MPI_WIN_FENCE` or `MPI_WIN_POST`) has executed. On the origin process, an RMA communication may become enabled as soon as the corresponding put, get or accumulate call has executed, or as late as when the ensuing synchronization call is issued. Once the communication is enabled both at the origin and at the target, the communication must complete.

Consider the code fragment in Example 11.4, on page 457. Some of the calls may block if the target window is not posted. However, if the target window is posted, then the code fragment must complete. The data transfer may start as soon as the put call occurs, but may be delayed until the ensuing complete call occurs.

Consider the code fragment in Example 11.5, on page 463. Some of the calls may block if another process holds a conflicting lock. However, if no conflicting lock is held, then the code fragment must complete.

Consider the code illustrated in Figure 11.6. Each process updates the window of the other process using a put operation, then accesses its own window. The post calls are nonblocking, and should complete. Once the post calls occur, RMA access to the windows is enabled, so that each process should complete the sequence of calls start-put-complete. Once these are done, the wait calls should complete at both processes. Thus, this communication should not deadlock, irrespective of the amount of data transferred.

Assume, in the last example, that the order of the post and start calls is reversed, at each process. Then, the code may deadlock, as each process may block on the start call, waiting for the matching post to occur. Similarly, the program will deadlock, if the order of the complete and wait calls is reversed, at each process.

The following two examples illustrate the fact that the synchronization between complete and wait is not symmetric: the wait call blocks until the complete executes, but not vice[-] versa. Consider the code illustrated in Figure 11.7. This code will deadlock: the wait of process 1 blocks until process 0 calls complete, and the receive of process 0 blocks until

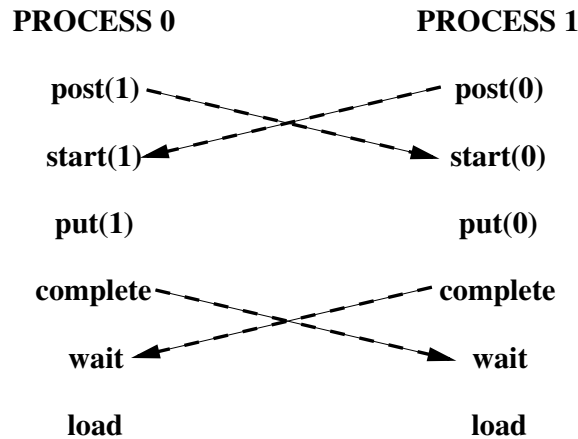


Figure 11.6: Symmetric communication

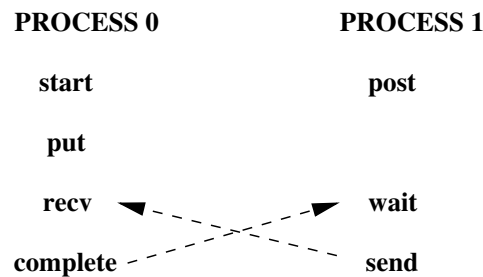


Figure 11.7: Deadlock situation

process 1 calls **send**. Consider, on the other hand, the code illustrated in Figure 11.8. This code will not deadlock. Once process 1 calls **post**, then the sequence **start**, **put**, **complete** on process 0 can proceed to completion. Process 0 will reach the **send** call, allowing the receive call of process 1 to complete.

*Rationale.* MPI implementations must guarantee that a process makes progress on all enabled communications it participates in, while blocked on an MPI call. This is true for send-receive communication and applies to RMA communication as well. Thus, in the example in Figure 11.8, the **put** and **complete** calls of process 0 should complete while process 1 is blocked on the receive call. This may require the involvement of process 1, e.g., to transfer the data put, while it is blocked on the receive call.

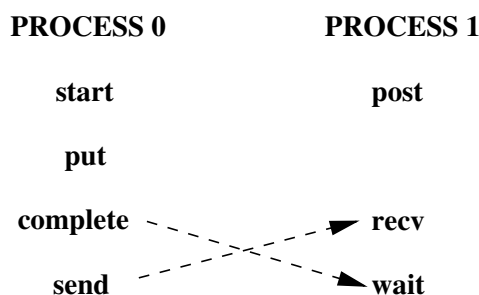


Figure 11.8: No deadlock

A similar issue is whether such progress must occur while a process is busy computing, or blocked in a non-MPI call. Suppose that in the last example the send-receive pair is replaced by a write-to-socket/read-from-socket pair. Then MPI does not specify whether deadlock is avoided. Suppose that the blocking receive of process 1 is replaced by a very long compute loop. Then, according to one interpretation of the MPI standard, process 0 must return from the complete call after a bounded delay, even if process 1 does not reach any MPI call in this period of time. According to another interpretation, the complete call may block until process 1 reaches the wait call, or reaches another MPI call. The qualitative behavior is the same, under both interpretations, unless a process is caught in an infinite compute loop, in which case the difference may not matter. However, the quantitative expectations are different. Different MPI implementations reflect these different interpretations. While this ambiguity is unfortunate, it does not seem to affect many real codes. The MPI [f]Forum decided not to decide which interpretation of the standard is the correct one, since the issue is very contentious, and a decision would have much impact on implementors but less impact on users. (*End of rationale.*)

#### 11.7.4 Registers and Compiler Optimizations

*Advice to users.* All the material in this section is an advice to users. (*End of advice to users.*)

A coherence problem exists between variables kept in registers and the memory value of these variables. An RMA call may access a variable in memory (or cache), while the up-to-date value of this variable is in register. A get will not return the latest variable value, and a put may be overwritten when the register is stored back in memory. **Note that these issues are unrelated to the RMA memory model; that is, these issues apply even if the memory model is MPI\_WIN\_UNIFIED.**

The problem is illustrated by the following code:

Source of Process 1	Source of Process 2	Executed in Process 2
bbbb = 777	buff = 999	reg_A:=999
call MPI_WIN_FENCE	call MPI_WIN_FENCE	
call MPI_PUT(bbbb		stop appl.thread
into buff of process 2)		buff:=777 in PUT handler
		continue appl.thread
call MPI_WIN_FENCE	call MPI_WIN_FENCE	
	ccc = buff	ccc:=reg_A

In this example, variable `buff` is allocated in the register `reg_A` and therefore `ccc` will have the old value of `buff` and not the new value 777.

This problem, which also afflicts in some cases send/receive communication, is discussed more at length in Section 16.2.16.

[MPI implementations will avoid this problem for standard conforming C programs.] **Programs written in C avoid this problem, because of the semantics of C.** Many Fortran compilers will avoid this problem, without disabling compiler optimizations. However, in order to avoid register coherence problems in a completely portable manner, users should restrict their use of RMA windows to variables stored in **[COMMON blocks, or to variables that were**



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declared `VOLATILE` (while `VOLATILE` is not a standard Fortran declaration, it is supported by many Fortran compilers) (but this attribute may inhibit optimization of any code containing the RMA window). [Details] Further details and an additional solution are discussed in Section 16.2.16, “A Problem with Register Optimization,” on page 677. See also[,] “Problems Due to Data Copying and Sequence Association,” on page 671, for additional Fortran [problems] issues. ] modules or `COMMON` blocks. To prevent problems with the argument copying and register optimization done by Fortran compilers, please note the hints in Sections 16.2.10-16.2.20, especially in Sections 16.2.12 and 16.2.13 on pages 671-674 about “Problems Due to Data Copying and Sequence Association with Subscript Triplets” and “Vector Subscripts”, and in Sections 16.2.16 to 16.2.19 on pages 677 to 686 about “Optimization Problems”, “Code Movements and Register Optimization”, “Temporary Data Movements” and “Permanent Data Movements”. Sections “Solutions” to “`VOLATILE`” on pages 680-684 discuss several solutions for the problem in this example.

## 11.8 Examples

[This section was moved from earlier in the chapter. Changes and additions to this section are marked in the same way as changes and additions in other parts of this chapter.]

**Example 11.13** The following example shows a generic loosely synchronous, iterative code, using fence synchronization. The window at each process consists of array `A`, which contains the origin and target buffers of the put calls.

```
...
while(!converged(A)){
    update(A);
    MPI_Win_fence(MPI_MODE_NOPRECEDE, win);
    for(i=0; i < toneighbors; i++)
        MPI_Put(&frombuf[i], 1, fromtype[i], toneighbor[i],
                todisp[i], 1, totype[i], win);
    MPI_Win_fence((MPI_MODE_NOSTORE | MPI_MODE_NOSUCCEED), win);
}
```

The same code could be written with `get`[,] rather than `put`. Note that, during the communication phase, each window is concurrently read (as origin buffer of puts) and written (as target buffer of puts). This is OK, provided that there is no overlap between the target buffer of a put and another communication buffer.

**Example 11.14** Same generic example, with more computation/communication overlap. We assume that the update phase is broken in two subphases: the first, where the “boundary,” which is involved in communication, is updated, and the second, where the “core,” which neither use nor provide communicated data, is updated.

```
...
while(!converged(A)){
    update_boundary(A);
    MPI_Win_fence((MPI_MODE_NOPUT | MPI_MODE_NOPRECEDE), win);
    for(i=0; i < fromneighbors; i++)
        MPI_Get(&tobuf[i], 1, totype[i], fromneighbor[i],
```

```

1         fromdisp[i], 1, fromtype[i], win);
2     update_core(A);
3     MPI_Win_fence(MPI_MODE_NOSUCCEED, win);
4 }

```

The get communication can be concurrent with the core update, since they do not access the same locations, and the local update of the origin buffer by the get call can be concurrent with the local update of the core by the `update_core` call. In order to get similar overlap with put communication we would need to use separate windows for the core and for the boundary. This is required because we do not allow local stores to be concurrent with puts on the same, or on overlapping, windows.

**Example 11.15** Same code as in Example 11.13, rewritten using post-start-complete-wait.

```

14 ...
15 while(!converged(A)){
16     update(A);
17     MPI_Win_post(fromgroup, 0, win);
18     MPI_Win_start(togroup, 0, win);
19     for(i=0; i < toneighbors; i++)
20         MPI_Put(&frombuf[i], 1, fromtype[i], toneighbor[i],
21               todisp[i], 1, totype[i], win);
22     MPI_Win_complete(win);
23     MPI_Win_wait(win);
24 }

```

**Example 11.16** Same example, with split phases, as in Example 11.14.

```

28 ...
29 while(!converged(A)){
30     update_boundary(A);
31     MPI_Win_post(togroup, MPI_MODE_NOPUT, win);
32     MPI_Win_start(fromgroup, 0, win);
33     for(i=0; i < fromneighbors; i++)
34         MPI_Get(&tobuf[i], 1, totype[i], fromneighbor[i],
35               fromdisp[i], 1, fromtype[i], win);
36     update_core(A);
37     MPI_Win_complete(win);
38     MPI_Win_wait(win);
39 }

```

**Example 11.17** A checkerboard, or double buffer communication pattern, that allows more computation/communication overlap. Array `A0` is updated using values of array `A1`, and vice versa. We assume that communication is symmetric: if process A gets data from process B, then process B gets data from process A. Window `wini` consists of array `Ai`.

```

...
if (!converged(A0,A1))
    MPI_Win_post(neighbors, (MPI_MODE_NOCHECK | MPI_MODE_NOPUT), win0);
MPI_Barrier(comm0);
/* the barrier is needed because the start call inside the
loop uses the nocheck option */
while(!converged(A0, A1)){
    /* communication on A0 and computation on A1 */
    update2(A1, A0); /* local update of A1 that depends on A0 (and A1) */
    MPI_Win_start(neighbors, MPI_MODE_NOCHECK, win0);
    for(i=0; i < neighbors; i++)
        MPI_Get(&tobuf0[i], 1, totype0[i], neighbor[i],
                fromdisp0[i], 1, fromtype0[i], win0);
    update1(A1); /* local update of A1 that is
                  concurrent with communication that updates A0 */
    MPI_Win_post(neighbors, (MPI_MODE_NOCHECK | MPI_MODE_NOPUT), win1);
    MPI_Win_complete(win0);
    MPI_Win_wait(win0);

    /* communication on A1 and computation on A0 */
    update2(A0, A1); /* local update of A0 that depends on A1 (and A0)*/
    MPI_Win_start(neighbors, MPI_MODE_NOCHECK, win1);
    for(i=0; i < neighbors; i++)
        MPI_Get(&tobuf1[i], 1, totype1[i], neighbor[i],
                fromdisp1[i], 1, fromtype1[i], win1);
    update1(A0); /* local update of A0 that depends on A0 only,
                  concurrent with communication that updates A1 */
    if (!converged(A0,A1))
        MPI_Win_post(neighbors, (MPI_MODE_NOCHECK | MPI_MODE_NOPUT), win0);
    MPI_Win_complete(win1);
    MPI_Win_wait(win1);
}

```

A process posts the local window associated with `win0` before it completes RMA accesses to the remote windows associated with `win1`. When the `wait(win1)` call returns, then all neighbors of the calling process have posted the windows associated with `win0`. Conversely, when the `wait(win0)` call returns, then all neighbors of the calling process have posted the windows associated with `win1`. Therefore, the `nocheck` option can be used with the calls to `MPI_WIN_START`.

Put calls can be used, instead of get calls, if the area of array `A0` (resp. `A1`) used by the `update(A1, A0)` (resp. `update(A0, A1)`) call is disjoint from the area modified by the RMA communication. On some systems, a put call may be more efficient than a get call, as it requires information exchange only in one direction.

In the next several examples, for conciseness, the expression

`z = MPI_Get_accumulate(...)`

means to perform an `MPI_GET_ACCUMULATE` with the result buffer (given by `result_addr` in the description of `MPI_GET_ACCUMULATE`) on the left side of the assignment; in this case, `z`. This format is also used with `MPI_COMPARE_AND_SWAP`.

**Example 11.18** The following example implements a naive, non-scalable counting semaphore. The example demonstrates the use of `MPI_WIN_SYNC` to manipulate the public copy of `X`, as well as `MPI_WIN_FLUSH` to complete operations without ending the access epoch opened with `MPI_WIN_LOCK_ALL`. To avoid the rules regarding synchronization of the public and private copies of windows, `MPI_ACCUMULATE` and `MPI_GET_ACCUMULATE` are used to write to or read from the local public copy.

Process A:	Process B:
<code>MPI_Win_lock_all</code>	<code>MPI_Win_lock_all</code>
window location <code>X</code>	
<code>X=2</code>	
<code>MPI_Win_sync</code>	
<code>MPI_Barrier</code>	<code>MPI_Barrier</code>
<code>MPI_Accumulate(X, MPI_SUM, -1)</code>	<code>MPI_Accumulate(X, MPI_SUM, -1)</code>
stack variable <code>z</code>	stack variable <code>z</code>
do	do
<code>z = MPI_Get_accumulate(X,</code>	<code>z = MPI_Get_accumulate(X,</code>
<code>MPI_NO_OP, 0)</code>	<code>MPI_NO_OP, 0)</code>
<code>MPI_Win_flush(A)</code>	<code>MPI_Win_flush(A)</code>
while( <code>z!=0</code> )	while( <code>z!=0</code> )
<code>MPI_Win_unlock_all</code>	<code>MPI_Win_unlock_all</code>

**Example 11.19** Implementing a critical region between two processes (Peterson's algorithm). Despite their appearance in the following example, `MPI_WIN_LOCK_ALL` and `MPI_WIN_UNLOCK_ALL` are not collective calls, but it is frequently useful to start shared access epochs to all processes from all other processes in a window. Once the access epochs are established, accumulate communication operations and flush and sync synchronization operations can be used to read from or write to the public copy of the window.

Process A:	Process B:
window location <code>X</code>	window location <code>Y</code>
window location <code>T</code>	
<code>MPI_Win_lock_all</code>	<code>MPI_Win_lock_all</code>
<code>X=1</code>	<code>Y=1</code>
<code>MPI_Win_sync</code>	<code>MPI_Win_sync</code>
<code>MPI_Barrier</code>	<code>MPI_Barrier</code>
<code>MPI_Accumulate(T, MPI_REPLACE, 1)</code>	<code>MPI_Accumulate(T, MPI_REPLACE, 0)</code>
stack variables <code>t,y</code>	stack variable <code>t,x</code>
<code>t=1</code>	<code>t=0</code>
<code>y=MPI_Get_accumulate(Y,</code>	<code>x=MPI_Get_accumulate(X,</code>
<code>MPI_NO_OP, 0)</code>	<code>MPI_NO_OP, 0)</code>
while( <code>y==1 &amp;&amp; t==1</code> ) do	while( <code>x==1 &amp;&amp; t==0</code> ) do
<code>y=MPI_Get_accumulate(Y,</code>	<code>x=MPI_Get_accumulate(X,</code>
<code>MPI_NO_OP, 0)</code>	<code>MPI_NO_OP, 0)</code>

```

t=MPI_Get_accumulate(T,          t=MPI_Get_accumulate(T,
    MPI_NO_OP, 0)                MPI_NO_OP, 0)
MPI_Win_flush_all                MPI_Win_flush(A)
done                             done
// critical region               // critical region
MPI_Accumulate(X, MPI_REPLACE, 0) MPI_Accumulate(Y, MPI_REPLACE, 0)
MPI_Win_unlock_all               MPI_Win_unlock_all

```

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**Example 11.20** Implementing a critical region between multiple processes with compare and swap. The call to `MPI_WIN_SYNC` is necessary on Process A after local initialization of A to guarantee the public copy has been updated with the initialization value found in the private copy. It would also be valid to call `MPI_ACCUMULATE` with `MPI_REPLACE` to directly initialize the public copy. A call to `MPI_WIN_FLUSH` would be necessary to assure A in the public copy of Process A had been updated before the barrier.

```

Process A:                      Process B...:
MPI_Win_lock_all                MPI_Win_lock_all
atomic location A
A=0
MPI_Win_sync
MPI_Barrier
stack variable r=1
while(r != 0) do
    r = MPI_Compare_and_swap(A, 0, 1)
    MPI_Win_flush(A)
done
// critical region
r = MPI_Compare_and_swap(A, 1, 0)
MPI_Win_unlock_all

MPI_Barrier
stack variable r=1
while(r != 0) do
    r = MPI_Compare_and_swap(A, 0, 1)
    MPI_Win_flush(A)
done
// critical region
r = MPI_Compare_and_swap(A, 1, 0)
MPI_Win_unlock_all

```

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**Example 11.21** The following example shows how request-based operations can be used to overlap communication with computation. Each process fetches, processes, and writes the result for *NSTEPS* chunks of data. Instead of a single buffer, *M* local buffers are used to allow up to *M* communication operations to overlap with computation.

```

int      i, j;
MPI_Win  win;
MPI_Request put_req[M] = { MPI_REQUEST_NULL };
MPI_Request get_req;
double   **baseptr;
double   data[M][N];

MPI_Win_allocate(NSTEPS*N*sizeof(double), sizeof(double), MPI_INFO_NULL,
    MPI_COMM_WORLD, baseptr, &win);

MPI_Win_lock_all(0, win);

for (i = 0; i < NSTEPS; i++) {

```

```

1   if (i<M)
2       j=i;
3   else
4       MPI_Waitany(M, put_req, &j, MPI_STATUS_IGNORE);
5
6   MPI_Rget(data[j], N, MPI_DOUBLE, target, i*N, N, MPI_DOUBLE, win,
7       &get_req);
8   MPI_Wait(&get_req, MPI_STATUS_IGNORE);
9   compute(i, data[j], ...);
10  MPI_Rput(data[j], N, MPI_DOUBLE, target, i*N, N, MPI_DOUBLE, win,
11      &put_req[j]);
12  }
13
14  MPI_Waitall(M, put_req, MPI_STATUSES_IGNORE);
15  MPI_Win_unlock_all(win);

```

### Example 11.22

The following example constructs a distributed shared linked list using dynamic windows. Initially process 0 creates the head of the list, attaches it to the window, and broadcasts the pointer to all processes. All processes then concurrently append  $N$  new elements to the list. When a process attempts to attach its element to the tail of the list it may discover that its tail pointer is stale and it must chase ahead to the new tail before the element can be attached. This example requires some modification to work in an environment where the length of a pointer is different on different processes.

```

26  ...
27  #define NUM_ELEMS 10
28
29  /* Linked list pointer */
30  typedef struct {
31      MPI_Aint disp;
32      int      rank;
33  } llist_ptr_t;
34
35  /* Linked list element */
36  typedef struct {
37      llist_ptr_t next;
38      int value;
39  } llist_elem_t;
40
41  const llist_ptr_t nil = { -1, (MPI_Aint) MPI_BOTTOM };
42
43  /* List of locally allocated list elements. */
44  static llist_elem_t **my_elems = NULL;
45  static int my_elems_size = 0;
46  static int my_elems_count = 0;

```

```

/* Allocate a new shared linked list element */
MPI_Aint alloc_elem(int value, MPI_Win win) {
    MPI_Aint disp;
    llist_elem_t *elem_ptr;

    /* Allocate the new element and register it with the window */
    MPI_Alloc_mem(sizeof(llist_elem_t), MPI_INFO_NULL, &elem_ptr);
    elem_ptr->value = value;
    elem_ptr->next = nil;
    MPI_Win_attach(win, elem_ptr, sizeof(llist_elem_t));

    /* Add the element to the list of local elements so we can free
       it later. */
    if (my_elems_size == my_elems_count) {
        my_elems_size += 100;
        my_elems = realloc(my_elems, my_elems_size);
    }
    my_elems[my_elems_count] = elem_ptr;
    my_elems_count++;

    MPI_Get_address(elem_ptr, &disp);
    return disp;
}

int main(int argc, char **argv) {
    int          procid, nproc, i;
    MPI_Win      llist_win;
    llist_ptr_t  head_ptr, tail_ptr;

    MPI_Init(&argc, &argv);

    MPI_Comm_rank(MPI_COMM_WORLD, &procid);
    MPI_Comm_size(MPI_COMM_WORLD, &nproc);

    MPI_Win_create_dynamic(MPI_INFO_NULL, MPI_COMM_WORLD, &llist_win);

    /* Process 0 creates the head node */
    if (procid == 0)
        head_ptr.disp = alloc_elem(-1, llist_win);

    /* Broadcast the head pointer to everyone */
    head_ptr.rank = 0;
    MPI_Bcast(&head_ptr.disp, 1, MPI_AINT, 0, MPI_COMM_WORLD);
    tail_ptr = head_ptr;

    /* Lock the window for shared access to all targets */
    MPI_Win_lock_all(0, llist_win);

```

```

1  /* All processes concurrently append NUM_ELEMS elements to the list */
2  for (i = 0; i < NUM_ELEMS; i++) {
3      llist_ptr_t new_elem_ptr;
4      int success;
5
6      /* Create a new list element and attach it to the window */
7      new_elem_ptr.rank = procid;
8      new_elem_ptr.disp = alloc_elem(procid, llist_win);
9
10     /* Append the new node to the list. This might take multiple
11        attempts if others have already appended and our tail pointer
12        is stale. */
13     do {
14         llist_ptr_t next_tail_ptr = nil;
15
16         MPI_Compare_and_swap((void*) &new_elem_ptr.rank, (void*) &nil.rank,
17                             (void*)&next_tail_ptr.rank, MPI_INT, tail_ptr.rank,
18                             (MPI_Aint) &(((llist_elem_t*)tail_ptr.disp)->next.rank),
19                             llist_win);
20
21         MPI_Win_flush(tail_ptr.rank, llist_win);
22         success = (next_tail_ptr.rank == nil.rank);
23
24         if (success) {
25             MPI_Accumulate(&new_elem_ptr.disp, 1, MPI_AINT, tail_ptr.rank,
26                           (MPI_Aint) &(((llist_elem_t*)tail_ptr.disp)->next.disp), 1,
27                           MPI_AINT, MPI_REPLACE, llist_win);
28
29             MPI_Win_flush(tail_ptr.rank, llist_win);
30             tail_ptr = new_elem_ptr;
31
32         } else {
33             /* Tail pointer is stale, fetch the displacement. May take
34                multiple tries if it is being updated. */
35             do {
36                 MPI_Get_accumulate( NULL, 0, MPI_AINT, &next_tail_ptr.disp,
37                                    1, MPI_AINT, tail_ptr.rank,
38                                    (MPI_Aint) &(((llist_elem_t*)tail_ptr.disp)->next.disp),
39                                    1, MPI_AINT, MPI_NO_OP, llist_win);
40
41                 MPI_Win_flush(tail_ptr.rank, llist_win);
42             } while (next_tail_ptr.disp == nil.disp);
43             tail_ptr = next_tail_ptr;
44         }
45     } while (!success);
46 }
47
48 MPI_Win_unlock_all(llist_win);

```



```
MPI_Barrier( MPI_COMM_WORLD );

/* Free all the elements in the list */
for ( ; my_elems_count > 0; my_elems_count-- ) {
    MPI_Win_detach(llist_win,my_elems[my_elems_count-1]);
    MPI_Free_mem(my_elems[my_elems_count-1]);
}
MPI_Win_free(&llist_win);
...
```

1  
2  
3  
4  
5  
6  
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# Chapter 12

## External Interfaces

### 12.1 Introduction

This chapter begins with calls used to create **generalized requests**, which allow users to create new nonblocking operations with an interface similar to what is present in MPI. This can be used to layer new functionality on top of MPI. Next, Section 12.3 deals with setting the information found in `status`. [This is] This functionality is needed for generalized requests. ticket0.

The chapter continues, in Section 12.4, with a discussion of how threads are to be handled in MPI. Although thread compliance is not required, the standard specifies how threads are to work if they are provided.

### 12.2 Generalized Requests

The goal of generalized requests is to allow users to define new nonblocking operations. Such an outstanding nonblocking operation is represented by a (generalized) request. A fundamental property of nonblocking operations is that progress toward the completion of this operation occurs asynchronously, i.e., concurrently with normal program execution. Typically, this requires execution of code concurrently with the execution of the user code, e.g., in a separate thread or in a signal handler. Operating systems provide a variety of mechanisms in support of concurrent execution. MPI does not attempt to standardize or replace these mechanisms: it is assumed programmers who wish to define new asynchronous operations will use the mechanisms provided by the underlying operating system. Thus, the calls in this section only provide a means for defining the effect of MPI calls such as `MPI_WAIT` or `MPI_CANCEL` when they apply to generalized requests, and for signaling to MPI the completion of a generalized operation.

*Rationale.* It is tempting to also define an MPI standard mechanism for achieving concurrent execution of user-defined nonblocking operations. However, it is very difficult to define such a mechanism without consideration of the specific mechanisms used in the operating system. The Forum feels that concurrency mechanisms are a proper part of the underlying operating system and should not be standardized by MPI; the MPI standard should only deal with the interaction of such mechanisms with MPI. (*End of rationale.*)

For a regular request, the operation associated with the request is performed by the MPI implementation, and the operation completes without intervention by the application. For a generalized request, the operation associated with the request is performed by the application; therefore, the application must notify MPI when the operation completes. This is done by making a call to `MPI_GREQUEST_COMPLETE`. MPI maintains the “completion” status of generalized requests. Any other request state has to be maintained by the user.

A new generalized request is started with

```
MPI_GREQUEST_START(query_fn, free_fn, cancel_fn, extra_state, request)
```

IN	query_fn	callback function invoked when request status is queried (function)
IN	free_fn	callback function invoked when request is freed (function)
IN	cancel_fn	callback function invoked when request is cancelled (function)
IN	extra_state	extra state
OUT	request	generalized request (handle)

```
int MPI_Grequest_start(MPI_Grequest_query_function *query_fn,
                      MPI_Grequest_free_function *free_fn,
                      MPI_Grequest_cancel_function *cancel_fn, void *extra_state,
                      MPI_Request *request)
```

```
MPI_Grequest_start(query_fn, free_fn, cancel_fn, extra_state, request,
                  ierror) BIND(C)
    PROCEDURE(MPI_Grequest_query_function) :: query_fn
    PROCEDURE(MPI_Grequest_free_function) :: free_fn
    PROCEDURE(MPI_Grequest_cancel_function) :: cancel_fn
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_GREQUEST_START(QUERY_FN, FREE_FN, CANCEL_FN, EXTRA_STATE, REQUEST,
                  IERROR)
    INTEGER REQUEST, IERROR
    EXTERNAL QUERY_FN, FREE_FN, CANCEL_FN
    INTEGER (KIND=MPI_ADDRESS_KIND) EXTRA_STATE
```

```
{static MPI::Grequest
    MPI::Grequest::Start(const MPI::Grequest::Query_function*
                        query_fn, const MPI::Grequest::Free_function* free_fn,
                        const MPI::Grequest::Cancel_function* cancel_fn,
                        void *extra_state) (binding deprecated, see Section 15.2) }
```

*Advice to users.* Note that a generalized request belongs, in C++, to the class `MPI::Grequest`, which is a derived class of `MPI::Request`. It is of the same type as regular requests, in C and Fortran. (*End of advice to users.*)

The call starts a generalized request and returns a handle to it in `request`.

The syntax and meaning of the callback functions are listed below. All callback functions are passed the `extra_state` argument that was associated with the request by the starting call `MPI_GREQUEST_START[. This can; extra_state can` be used to maintain user-defined state for the request.

In C, the query function is

```
typedef int MPI_Grequest_query_function(void *extra_state,
                                       MPI_Status *status);
```

in Fortran with the `mpi_f08` module

ABSTRACT INTERFACE

```
SUBROUTINE MPI_Grequest_query_function(extra_state, status, ierror)
BIND(C)
    TYPE(MPI_Status) :: status
    INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state
    INTEGER :: ierror
```

in Fortran with the `mpi` module and `mpif.h`

```
SUBROUTINE GREQUEST_QUERY_FUNCTION(EXTRA_STATE, STATUS, IERROR)
    INTEGER STATUS(MPI_STATUS_SIZE), IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE
```

and in C++

```
{typedef int MPI::Grequest::Query_function(void* extra_state,
      MPI::Status& status); (binding deprecated, see Section 15.2)}
```

`[query_fn]` The `query_fn` function computes the status that should be returned for the generalized request. The status also includes information about successful/unsuccessful cancellation of the request (result to be returned by `MPI_TEST_CANCELLED`).

`[query_fn]` The `query_fn` callback is invoked by the `MPI_{WAIT|TEST}{ANY|SOME|ALL}` call that completed the generalized request associated with this callback. The callback function is also invoked by calls to `MPI_REQUEST_GET_STATUS`, if the request is complete when the call occurs. In both cases, the callback is passed a reference to the corresponding status variable passed by the user to the MPI call; the status set by the callback function is returned by the MPI call. If the user provided `MPI_STATUS_IGNORE` or `MPI_STATUSES_IGNORE` to the MPI function that causes `query_fn` to be called, then MPI will pass a valid status object to `query_fn`, and this status will be ignored upon return of the callback function. Note that `query_fn` is invoked only after `MPI_GREQUEST_COMPLETE` is called on the request; it may be invoked several times for the same generalized request, e.g., if the user calls `MPI_REQUEST_GET_STATUS` several times for this request. Note also that a call to `MPI_{WAIT|TEST}{SOME|ALL}` may cause multiple invocations of `query_fn` callback functions, one for each generalized request that is completed by the MPI call. The order of these invocations is not specified by MPI.

In C, the free function is

```
typedef int MPI_Grequest_free_function(void *extra_state);
```

in Fortran with the `mpi_f08` module

1     **ABSTRACT INTERFACE**

2     **SUBROUTINE MPI\_Grequest\_free\_function(extra\_state, ierror) BIND(C)**  
 3         **INTEGER(KIND=MPI\_ADDRESS\_KIND) :: extra\_state**  
 4         **INTEGER :: ierror**

5     [and] in Fortran with the `mpi` module and `mpif.h`

6     **SUBROUTINE GREQUEST\_FREE\_FUNCTION(EXTRA\_STATE, IERROR)**  
 7         **INTEGER IERROR**  
 8         **INTEGER(KIND=MPI\_ADDRESS\_KIND) EXTRA\_STATE**

9     and in C++

10     {typedef int MPI::Grequest::Free\_function(void\* extra\_state); (*binding*  
 11         *deprecated, see Section 15.2*)}

12     [free\_fn] The `free_fn` function is invoked to clean up user-allocated resources when the  
 13     generalized request is freed.

14     [free\_fn] The `free_fn` callback is invoked by the `MPI_{WAIT|TEST}{ANY|SOME|ALL}`  
 15     call that completed the generalized request associated with this callback. `free_fn` is invoked  
 16     after the call to `query_fn` for the same request. However, if the MPI call completed multiple  
 17     generalized requests, the order in which `free_fn` callback functions are invoked is not specified  
 18     by MPI.

19     [free\_fn] The `free_fn` callback is also invoked for generalized requests that are freed by a  
 20     call to `MPI_REQUEST_FREE` (no call to `WAIT_{WAIT|TEST}{ANY|SOME|ALL}` will occur  
 21     for such a request). In this case, the callback function will be called either in the MPI call  
 22     `MPI_REQUEST_FREE(request)`, or in the MPI call `MPI_GREQUEST_COMPLETE(request)`,  
 23     whichever happens last, i.e., in this case the actual freeing code is executed as soon as both  
 24     calls `MPI_REQUEST_FREE` and `MPI_GREQUEST_COMPLETE` have occurred. The request  
 25     is not deallocated until after `free_fn` completes. Note that `free_fn` will be invoked only once  
 26     per request by a correct program.

27     *Advice to users.* Calling `MPI_REQUEST_FREE(request)` will cause the request handle  
 28     to be set to `MPI_REQUEST_NULL`. This handle to the generalized request is no longer  
 29     valid. However, user copies of this handle are valid until after `free_fn` completes since  
 30     MPI does not deallocate the object until then. Since `free_fn` is not called until after  
 31     `MPI_GREQUEST_COMPLETE`, the user copy of the handle can be used to make this  
 32     call. Users should note that MPI will deallocate the object after `free_fn` executes. At  
 33     this point, user copies of the request handle no longer point to a valid request. MPI  
 34     will not set user copies to `MPI_REQUEST_NULL` in this case, so it is up to the user to  
 35     avoid accessing this stale handle. This is a special case [where] in which MPI defers  
 36     deallocating the object until a later time that is known by the user. (*End of advice*  
 37     *to users.*)

38     In C, the cancel function is

39     typedef int MPI\_Grequest\_cancel\_function(void \*extra\_state, int complete);

40     in Fortran with the `mpi_f08` module

41     **ABSTRACT INTERFACE**

42     **SUBROUTINE MPI\_Grequest\_cancel\_function(extra\_state, complete, ierror)**  
 43         **BIND(C)**

```
INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state
LOGICAL :: complete
INTEGER :: ierror
```

in Fortran with the `mpi` module and `mpif.h`

```

SUBROUTINE GREQUEST_CANCEL_FUNCTION(EXTRA_STATE, COMPLETE, IERROR)
  INTEGER IERROR
  INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE
  LOGICAL COMPLETE

```

and in C++

```
{typedef int MPI::Grequest::Cancel_function(void* extra_state,
                                             bool complete); (binding deprecated, see Section 15.2)}
```

`[cancel_fn]`The `cancel_fn` function is invoked to start the cancelation of a generalized request. It is called by `MPI_CANCEL(request)`. MPI passes `[to the callback function complete=true]``complete=true` to the callback function if `MPI_GREQUEST_COMPLETE` was already called on the request, and `complete=false` otherwise.

All callback functions return an error code. The code is passed back and dealt with as appropriate for the error code by the MPI function that invoked the callback function. For example, if error codes are returned then the error code returned by the callback function will be returned by the MPI function that invoked the callback function. In the case of an `MPI_{WAIT|TEST}{ANY}` call that invokes both `query_fn` and `free_fn`, the MPI call will return the error code returned by the last callback, namely `free_fn`. If one or more of the requests in a call to `MPI_{WAIT|TEST}{SOME|ALL}` failed, then the MPI call will return `MPI_ERR_IN_STATUS`. In such a case, if the MPI call was passed an array of statuses, then MPI will return in each of the statuses that correspond to a completed generalized request the error code returned by the corresponding invocation of its `free_fn` callback function. However, if the MPI function was passed `MPI_STATUSES_IGNORE`, then the individual error codes returned by each callback functions will be lost.

*Advice to users.* `query_fn` must **not** set the error field of `status` since `query_fn` may be called by `MPI_WAIT` or `MPI_TEST`, in which case the error field of `status` should not change. The MPI library knows the “context” in which `query_fn` is invoked and can decide correctly when to put in the error field of `status` the returned error code. *(End of advice to users.)*

MPI\_GREQUEST\_COMPLETE(request)

INOUT	request	generalized request (handle)
-------	---------	------------------------------

```
int MPI_Grequest_complete(MPI_Request request)
```

```
MPI_Grequest_complete(request, ierror) BIND(C)
    TYPE(MPI_Request), INTENT(IN) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```

MPI GREQUEST COMPLETE(REQUEST, IERROR)

```

```

1      INTEGER REQUEST, IERROR
2
3      {void MPI::Grequest::Complete() (binding deprecated, see Section 15.2) }
```

The call informs MPI that the operations represented by the generalized request `request` are complete (see definitions in Section 2.4). A call to `MPI_WAIT(request, status)` will return and a call to `MPI_TEST(request, flag, status)` will return `flag=true` only after a call to `MPI_GREQUEST_COMPLETE` has declared that these operations are complete.

MPI imposes no restrictions on the code executed by the callback functions. However, new nonblocking operations should be defined so that the general semantic rules about MPI calls such as `MPI_TEST`, `MPI_REQUEST_FREE`, or `MPI_CANCEL` still hold. For example, all these calls are supposed to be local and nonblocking. Therefore, the callback functions `query_fn`, `free_fn`, or `cancel_fn` should invoke blocking MPI communication calls only if the context is such that these calls are guaranteed to return in finite time. Once `MPI_CANCEL` is invoked, the cancelled operation should complete in finite time, irrespective of the state of other processes (the operation has acquired “local” semantics). It should either succeed, or fail without side-effects. The user should guarantee these same properties for newly defined operations.

*Advice to implementors.* A call to `MPI_GREQUEST_COMPLETE` may unblock a blocked user process/thread. The MPI library should ensure that the blocked user computation will resume. (*End of advice to implementors.*)

### 12.2.1 Examples

**Example 12.1** This example shows the code for a user-defined reduce operation on an `int` using a binary tree: each non-root node receives two messages, sums them, and sends them up. We assume that no status is returned and that the operation cannot be cancelled.

```

29  typedef struct {
30      MPI_Comm comm;
31      int tag;
32      int root;
33      int valin;
34      int *valout;
35      MPI_Request request;
36  } ARGS;
37
38
39  int myreduce(MPI_Comm comm, int tag, int root,
40              int valin, int *valout, MPI_Request *request)
41  {
42      ARGS *args;
43      pthread_t thread;
44
45      /* start request */
46      MPI_Grequest_start(query_fn, free_fn, cancel_fn, NULL, request);
47
48      args = (ARGS*)malloc(sizeof(ARGS));
```



```

    args->comm = comm;
    args->tag = tag;
    args->root = root;
    args->valin = valin;
    args->valout = valout;
    args->request = *request;

    /* spawn thread to handle request */
    /* The availability of the pthread_create call is system dependent */
    pthread_create(&thread, NULL, reduce_thread, args);

    return MPI_SUCCESS;
}

/* thread code */
void* reduce_thread(void *ptr)
{
    int lchild, rchild, parent, lval, rval, val;
    MPI_Request req[2];
    ARGS *args;

    args = (ARGS*)ptr;

    /* compute left,right child and parent in tree; set
       to MPI_PROC_NULL if does not exist */
    /* code not shown */
    ...

    MPI_Irecv(&lval, 1, MPI_INT, lchild, args->tag, args->comm, &req[0]);
    MPI_Irecv(&rval, 1, MPI_INT, rchild, args->tag, args->comm, &req[1]);
    MPI_Waitall(2, req, MPI_STATUSES_IGNORE);
    val = lval + args->valin + rval;
    MPI_Send(&val, 1, MPI_INT, parent, args->tag, args->comm );
    if (parent == MPI_PROC_NULL) *(args->valout) = val;
    MPI_Grequest_complete((args->request));
    free(ptr);
    return(NULL);
}

int query_fn(void *extra_state, MPI_Status *status)
{
    /* always send just one int */
    MPI_Status_set_elements(status, MPI_INT, 1);
    /* can never cancel so always true */
    MPI_Status_set_cancelled(status, 0);
    /* choose not to return a value for this */
    status->MPI_SOURCE = MPI_UNDEFINED;
    /* tag has no meaning for this generalized request */

```

```

1      status->MPI_TAG = MPI_UNDEFINED;
2      /* this generalized request never fails */
3      return MPI_SUCCESS;
4  }
5
6
7  int free_fn(void *extra_state)
8  {
9      /* this generalized request does not need to do any freeing */
10     /* as a result it never fails here */
11     return MPI_SUCCESS;
12 }
13
14
15 int cancel_fn(void *extra_state, int complete)
16 {
17     /* This generalized request does not support cancelling.
18        Abort if not already done.  If done then treat as if cancel failed.*/
19     if (!complete) {
20         fprintf(stderr,
21             "Cannot cancel generalized request - aborting program\n");
22         MPI_Abort(MPI_COMM_WORLD, 99);
23     }
24     return MPI_SUCCESS;
25 }
26
27

```

## 12.3 Associating Information with Status

MPI supports several different types of requests besides those for point-to-point operations. These range from MPI calls for I/O to generalized requests. It is desirable to allow these calls *[use]to use* the same request *[mechanism. This]mechanism, which* allows one to wait or test on different types of requests. However, `MPI_{TEST|WAIT}{ANY|SOME|ALL}` returns a status with information about the request. With the generalization of requests, one needs to define what information will be returned in the status object.

Each MPI call fills in the appropriate fields in the status object. Any unused fields will have undefined values. A call to `MPI_{TEST|WAIT}{ANY|SOME|ALL}` can modify any of the fields in the status object. Specifically, it can modify fields that are undefined. The fields with meaningful *[value]values* for a given request are defined in the sections with the new request.

Generalized requests raise additional considerations. Here, the user provides the functions to deal with the request. Unlike other MPI calls, the user needs to provide the information to be returned in status. The status argument is provided directly to the callback function where the status needs to be set. Users can directly set the values in 3 of the 5 status values. The count and cancel fields are opaque. To overcome this, these calls are provided:

MPI\_STATUS\_SET\_ELEMENTS(status, datatype, count)

INOUT	status	status with which to associate count (Status)
IN	datatype	datatype associated with count (handle)
IN	count	number of elements to associate with status (integer)

```
int MPI_Status_set_elements(MPI_Status *status, MPI_Datatype datatype,
                           int count)
```

```
MPI_Status_set_elements(status, datatype, count, ierror) BIND(C)
  TYPE(MPI_Status), INTENT(INOUT) :: status
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  INTEGER, INTENT(IN) :: count
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_STATUS_SET_ELEMENTS(STATUS, DATATYPE, COUNT, IERROR)
  INTEGER STATUS(MPI_STATUS_SIZE), DATATYPE, COUNT, IERROR
```

```
{void MPI::Status::Set_elements(const MPI::Datatype& datatype, int
                               count) (binding deprecated, see Section 15.2) }
```

MPI\_STATUS\_SET\_ELEMENTS\_X(status, datatype, count)

INOUT	status	status with which to associate count (Status)
IN	datatype	datatype associated with count (handle)
IN	count	number of elements to associate with status (integer)

```
int MPI_Status_set_elements_x(MPI_Status *status, MPI_Datatype datatype,
                              MPI_Count count)
```

```
MPI_Status_set_elements_x(status, datatype, count, ierror) BIND(C)
  TYPE(MPI_Status), INTENT(INOUT) :: status
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  INTEGER(KIND = MPI_COUNT_KIND), INTENT(IN) :: count
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_STATUS_SET_ELEMENTS_X(STATUS, DATATYPE, COUNT, IERROR)
  INTEGER STATUS(MPI_STATUS_SIZE), DATATYPE, IERROR
  INTEGER (KIND=MPI_COUNT_KIND) COUNT
```

[This call modifies] These functions modify the opaque part of status so that a call to MPI\_GET\_ELEMENTS or MPI\_GET\_ELEMENTS\_X will return count. MPI\_GET\_COUNT will return a compatible value.

*Rationale.* The number of elements is set instead of the count because the former can deal with a nonintegral number of datatypes. (*End of rationale.*)

A subsequent call to MPI\_GET\_COUNT(status, datatype, count) [ or to ], MPI\_GET\_ELEMENTS(status, datatype, count) , or MPI\_GET\_ELEMENTS\_X(status, datatype,

`count`) must use a `datatype` argument that has the same type signature as the `datatype` argument that was used in the call to `MPI_STATUS_SET_ELEMENTS` or `MPI_STATUS_SET_ELEMENTS_X`.

*Rationale.* [This] The requirement of matching type signatures for these calls is similar to the restriction that holds when `count` is set by a receive operation: in that case, the calls to `MPI_GET_COUNT` [and], `MPI_GET_ELEMENTS`, and `MPI_GET_ELEMENTS_X` must use a `datatype` with the same signature as the `datatype` used in the receive call. (*End of rationale.*)

`MPI_STATUS_SET_CANCELLED(status, flag)`

INOUT	status	status with which to associate cancel flag (Status)
IN	flag	if true indicates request was cancelled (logical)

`int MPI_Status_set_cancelled(MPI_Status *status, int flag)`

`MPI_Status_set_cancelled(status, flag, ierror) BIND(C)`  
`TYPE(MPI_Status), INTENT(INOUT) :: status`  
`LOGICAL, INTENT(OUT) :: flag`  
`INTEGER, OPTIONAL, INTENT(OUT) :: ierror`

`MPI_STATUS_SET_CANCELLED(STATUS, FLAG, IERROR)`  
`INTEGER STATUS(MPI_STATUS_SIZE), IERROR`  
`LOGICAL FLAG`

`{void MPI::Status::Set_cancelled(bool flag) (binding deprecated, see Section 15.2)`  
`}`

If `flag` is set to true then a subsequent call to `MPI_TEST_CANCELLED(status, flag)` will also return `flag = true`, otherwise it will return false.

*Advice to users.* Users are advised not to reuse the status fields for values other than those for which they were intended. Doing so may lead to unexpected results when using the status object. For example, calling `MPI_GET_ELEMENTS` may cause an error if the value is out of range or it may be impossible to detect such an error. The `extra_state` argument provided with a generalized request can be used to return information that does not logically belong in status. Furthermore, modifying the values in a status set internally by MPI, e.g., `MPI_RECV`, may lead to unpredictable results and is strongly discouraged. (*End of advice to users.*)

## 12.4 MPI and Threads

This section specifies the interaction between MPI calls and threads. The section lists minimal requirements for **thread compliant** MPI implementations and defines functions that can be used for initializing the thread environment. MPI may be implemented in environments where threads are not supported or perform poorly. Therefore, it is not required that all MPI implementations fulfill all the requirements specified in this section.

This section generally assumes a thread package similar to POSIX threads [38], but the syntax and semantics of thread calls are not specified here — these are beyond the scope of this document.

### 12.4.1 General

In a thread-compliant implementation, an MPI process is a process that may be multi-threaded. Each thread can issue MPI calls; however, threads are not separately addressable: a rank in a send or receive call identifies a process, not a thread. A message sent to a process can be received by any thread in this process.

*Rationale.* This model corresponds to the POSIX model of interprocess communication: the fact that a process is multi-threaded, rather than single-threaded, does not affect the external interface of this process. MPI implementations [where]in which MPI ‘processes’ are POSIX threads inside a single POSIX process are not thread-compliant by this definition (indeed, their “processes” are single-threaded). (*End of rationale.*)

*Advice to users.* It is the user’s responsibility to prevent races when threads within the same application post conflicting communication calls. The user can make sure that two threads in the same process will not issue conflicting communication calls by using distinct communicators at each thread. (*End of advice to users.*)

The two main requirements for a thread-compliant implementation are listed below.

1. All MPI calls are *thread-safe*, i.e., two concurrently running threads may make MPI calls and the outcome will be as if the calls executed in some order, even if their execution is interleaved.
2. Blocking MPI calls will block the calling thread only, allowing another thread to execute, if available. The calling thread will be blocked until the event on which it is waiting occurs. Once the blocked communication is enabled and can proceed, then the call will complete and the thread will be marked runnable, within a finite time. A blocked thread will not prevent progress of other runnable threads on the same process, and will not prevent them from executing MPI calls.

**Example 12.2** Process 0 consists of two threads. The first thread executes a blocking send call `MPI_Send(buff1, count, type, 0, 0, comm)`, whereas the second thread executes a blocking receive call `MPI_Recv(buff2, count, type, 0, 0, comm, &status)`, i.e., the first thread sends a message that is received by the second thread. This communication should always succeed. According to the first requirement, the execution will correspond to some interleaving of the two calls. According to the second requirement, a call can only block the calling thread and cannot prevent progress of the other thread. If the send call went ahead of the receive call, then the sending thread may block, but this will not prevent the receiving thread from executing. Thus, the receive call will occur. Once both calls occur, the communication is enabled and both calls will complete. On the other hand, a single-threaded process that posts a send, followed by a matching receive, may deadlock. The progress requirement for multithreaded implementations is stronger, as a blocked call cannot prevent progress in other threads.

*Advice to implementors.* MPI calls can be made thread-safe by executing only one at a time, e.g., by protecting MPI code with one process-global lock. However, blocked operations cannot hold the lock, as this would prevent progress of other threads in the process. The lock is held only for the duration of an atomic, locally-completing suboperation such as posting a send or completing a send, and is released in between. Finer locks can provide more concurrency, at the expense of higher locking overheads. Concurrency can also be achieved by having some of the MPI protocol executed by separate server threads. (*End of advice to implementors.*)

## 12.4.2 Clarifications

**Initialization and Completion** The call to `MPI_FINALIZE` should occur on the same thread that initialized MPI. We call this thread the **main thread**. The call should occur only after all the process threads have completed their MPI calls, and have no pending communications or I/O operations.

*Rationale.* This constraint simplifies implementation. (*End of rationale.*)

**Multiple threads completing the same request.** A program where two threads block, waiting on the same request, is erroneous. Similarly, the same request cannot appear in the array of requests of two concurrent `MPI_{WAIT|TEST}_{ANY|SOME|ALL}` calls. In MPI, a request can only be completed once. Any combination of wait or test [which]that violates this rule is erroneous.

*Rationale.* [This]This restriction is consistent with the view that a multithreaded execution corresponds to an interleaving of the MPI calls. In a single threaded implementation, once a wait is posted on a request the request handle will be nullified before it is possible to post a second wait on the same handle. With threads, an `MPI_WAIT_{ANY|SOME|ALL}` may be blocked without having nullified its request(s) so it becomes the user's responsibility to avoid using the same request in an `MPI_WAIT` on another thread. This constraint also simplifies implementation, as only one thread will be blocked on any communication or I/O event. (*End of rationale.*)

**Probe** A receive call that uses source and tag values returned by a preceding call to `MPI_PROBE` or `MPI_IProbe` will receive the message matched by the probe call only if there was no other matching receive after the probe and before that receive. In a multi-threaded environment, it is up to the user to enforce this condition using suitable mutual exclusion logic. This can be enforced by making sure that each communicator is used by only one thread on each process.

**Collective calls** Matching of collective calls on a communicator, window, or file handle is done according to the order in which the calls are issued at each process. If concurrent threads issue such calls on the same communicator, window or file handle, it is up to the user to make sure the calls are correctly ordered, using interthread synchronization.

*Advice to users.* With three concurrent threads in each MPI process of a communicator `comm`, it is allowed that thread A in each MPI process calls a collective operation on `comm`, thread B calls a file operation on an existing filehandle that was formerly

opened on `comm`, and thread C invokes one-sided operations on an existing window handle that was also formerly created on `comm`. (*End of advice to users.*)

*Rationale.* As already specified in `MPI_FILE_OPEN` and `MPI_WIN_CREATE`, a file handle and a window handle inherit only the group of processes of the underlying communicator, but not the communicator itself. Accesses to communicators, window handles and file handles cannot affect one another. (*End of rationale.*)

*Advice to implementors.* [Advice to implementors.] If the implementation of file or window operations internally uses MPI communication then a duplicated communicator may be cached on the file or window object. (*End of advice to implementors.*)

**Exception handlers** An exception handler does not necessarily execute in the context of the thread that made the exception-raising MPI call; the exception handler may be executed by a thread that is distinct from the thread that will return the error code.

*Rationale.* The MPI implementation may be multithreaded, so that part of the communication protocol may execute on a thread that is distinct from the thread that made the MPI call. The design allows the exception handler to be executed on the thread where the exception occurred. (*End of rationale.*)

**Interaction with signals and cancellations** The outcome is undefined if a thread that executes an MPI call is cancelled (by another thread), or if a thread catches a signal while executing an MPI call. However, a thread of an MPI process may terminate, and may catch signals or be cancelled by another thread when not executing MPI calls.

*Rationale.* Few C library functions are signal safe, and many have cancellation points — points [where]at which the thread executing them may be cancelled. The above restriction simplifies implementation (no need for the MPI library to be “async-cancel-safe” or [“async-signal-safe.”]“async-signal-safe”). (*End of rationale.*)

*Advice to users.* Users can catch signals in separate, non-MPI threads (e.g., by masking signals on MPI calling threads, and unmasking them in one or more non-MPI threads). A good programming practice is to have a distinct thread blocked in a call to `sigwait` for each user expected signal that may occur. Users must not catch signals used by the MPI implementation; as each MPI implementation is required to document the signals used internally, users can avoid these signals. (*End of advice to users.*)

*Advice to implementors.* The MPI library should not invoke library calls that are not thread safe, if multiple threads execute. (*End of advice to implementors.*)

### 12.4.3 Initialization

The following function may be used to initialize MPI, and initialize the MPI thread environment, instead of `MPI_INIT`.

1 MPI\_INIT\_THREAD(required, provided)

2     IN         required                     desired level of thread support (integer)

3     OUT        provided                    provided level of thread support (integer)

4     int MPI\_Init\_thread(int \*argc, char \*((\*argv)[]), int required,  
5                             int \*provided)

6 MPI\_Init\_thread(required, provided, ierror) BIND(C)

7     INTEGER, INTENT(IN) :: required

8     INTEGER, INTENT(OUT) :: provided

9     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

10 MPI\_INIT\_THREAD(REQUIRED, PROVIDED, IERROR)

11     INTEGER REQUIRED, PROVIDED, IERROR

12 {int MPI::Init\_thread(int& argc, char\*\*& argv, int required) (*binding*  
13                             *deprecated, see Section 15.2*) }

14 {int MPI::Init\_thread(int required) (*binding deprecated, see Section 15.2*) }

15     *Advice to users.* In C and C++, the passing of argc and argv is [optional.]optional,  
16     as with MPI\_INIT as discussed in Section 8.7. In C, [this is accomplished by passing  
17     the appropriate null pointer.] null pointers may be passed in their place. In C++,  
18     [this is accomplished with two separate bindings to cover these two cases. This is as  
19     with MPI\_INIT as discussed in Section 8.7.]two separate bindings support this choice.  
20     (*End of advice to users.*)

21     This call initializes MPI in the same way that a call to MPI\_INIT would. In addition,  
22     it initializes the thread environment. The argument required is used to specify the desired  
23     level of thread support. The possible values are listed in increasing order of thread support.

24     **MPI\_THREAD\_SINGLE** Only one thread will execute.

25     **MPI\_THREAD\_FUNNELED** The process may be multi-threaded, but the application must  
26     ensure that only the main thread makes MPI calls (for the definition of main thread,  
27     see MPI\_IS\_THREAD\_MAIN on page 506).

28     **MPI\_THREAD\_SERIALIZED** The process may be multi-threaded, and multiple threads may  
29     make MPI calls, but only one at a time: MPI calls are not made concurrently from  
30     two distinct threads (all MPI calls are “serialized”).

31     **MPI\_THREAD\_MULTIPLE** Multiple threads may call MPI, with no restrictions.

32     These values are monotonic; i.e., MPI\_THREAD\_SINGLE < MPI\_THREAD\_FUNNELED <  
33     MPI\_THREAD\_SERIALIZED < MPI\_THREAD\_MULTIPLE.

34     Different processes in MPI\_COMM\_WORLD may require different levels of thread sup-  
35     port.

36     The call returns in provided information about the actual level of thread support that  
37     will be provided by MPI. It can be one of the four values listed above.

38     The level(s) of thread support that can be provided by MPI\_INIT\_THREAD will depend  
39     on the implementation, and may depend on information provided by the user before the



program started to execute (e.g., with arguments to `mpiexec`). If possible, the call will return `provided = required`. Failing this, the call will return the least supported level such that `provided > required` (thus providing a stronger level of support than required by the user). Finally, if the user requirement cannot be satisfied, then the call will return in `provided` the highest supported level.

A **thread compliant** MPI implementation will be able to return `provided = MPI_THREAD_MULTIPLE`. Such an implementation may always return `provided = MPI_THREAD_MULTIPLE`, irrespective of the value of `required`. At the other extreme, an MPI library that is not thread compliant may always return `provided = MPI_THREAD_SINGLE`, irrespective of the value of `required`.

A call to `MPI_INIT` has the same effect as a call to `MPI_INIT_THREAD` with a `required = MPI_THREAD_SINGLE`.

Vendors may provide (implementation dependent) means to specify the level(s) of thread support available when the MPI program is started, e.g., with arguments to `mpiexec`. This will affect the outcome of calls to `MPI_INIT` and `MPI_INIT_THREAD`. Suppose, for example, that an MPI program has been started so that only `MPI_THREAD_MULTIPLE` is available. Then `MPI_INIT_THREAD` will return `provided = MPI_THREAD_MULTIPLE`, irrespective of the value of `required`; a call to `MPI_INIT` will also initialize the MPI thread support level to `MPI_THREAD_MULTIPLE`. Suppose, on the other hand, that an MPI program has been started so that all four levels of thread support are available. Then, a call to `MPI_INIT_THREAD` will return `provided = required`; on the other hand, a call to `MPI_INIT` will initialize the MPI thread support level to `MPI_THREAD_SINGLE`.

*Rationale.* Various optimizations are possible when MPI code is executed single-threaded, or is executed on multiple threads, but not concurrently: mutual exclusion code may be omitted. Furthermore, if only one thread executes, then the MPI library can use library functions that are not thread safe, without risking conflicts with user threads. Also, the model of one communication thread, multiple computation threads fits many applications well, e.g., if the process code is a sequential Fortran/C/C++ program with MPI calls that has been parallelized by a compiler for execution on an SMP node, in a cluster of SMPs, then the process computation is multi-threaded, but MPI calls will likely execute on a single thread.

The design accommodates a static specification of the thread support level, for environments that require static binding of libraries, and for compatibility for current multi-threaded MPI codes. (*End of rationale.*)

*Advice to implementors.* If `provided` is not `MPI_THREAD_SINGLE` then the MPI library should not invoke C/ C++/Fortran library calls that are not thread safe, e.g., in an environment where `malloc` is not thread safe, then `malloc` should not be used by the MPI library.

Some implementors may want to use different MPI libraries for different levels of thread support. They can do so using dynamic linking and selecting which library will be linked when `MPI_INIT_THREAD` is invoked. If this is not possible, then optimizations for lower levels of thread support will occur only when the level of thread support required is specified at link time. (*End of advice to implementors.*)

The following function can be used to query the current level of thread support.

1 MPI\_QUERY\_THREAD(provided)

2 OUT provided provided level of thread support (integer)

4 int MPI\_Query\_thread(int \*provided)

6 MPI\_Query\_thread(provided, ierror) BIND(C)

7 INTEGER, INTENT(OUT) :: provided

8 INTEGER, OPTIONAL, INTENT(OUT) :: ierror

10 MPI\_QUERY\_THREAD(PROVIDED, IERROR)

11 INTEGER PROVIDED, IERROR

12 {int MPI::Query\_thread() (binding deprecated, see Section 15.2) }

14 The call returns in provided the current level of thread [support. This]support, which  
15 will be the value returned in provided by MPI\_INIT\_THREAD, if MPI was initialized by a  
16 call to MPI\_INIT\_THREAD().

18 MPI\_IS\_THREAD\_MAIN(flag)

20 OUT flag true if calling thread is main thread, false otherwise  
21 (logical)

23 int MPI\_Is\_thread\_main(int \*flag)

25 MPI\_Is\_thread\_main(flag, ierror) BIND(C)

26 LOGICAL, INTENT(OUT) :: flag

27 INTEGER, OPTIONAL, INTENT(OUT) :: ierror

28 MPI\_IS\_THREAD\_MAIN(FLAG, IERROR)

29 LOGICAL FLAG

30 INTEGER IERROR

31 {bool MPI::Is\_thread\_main() (binding deprecated, see Section 15.2) }

33 This function can be called by a thread to [find out whether]determine if it is the main  
34 thread (the thread that called MPI\_INIT or MPI\_INIT\_THREAD).

35 All routines listed in this section must be supported by all MPI implementations.

37 *Rationale.* MPI libraries are required to provide these calls even if they do not  
38 support threads, so that portable code that contains invocations to these functions  
39 [be able to]can link correctly. MPI\_INIT continues to be supported so as to provide  
40 compatibility with current MPI codes. (*End of rationale.*)

42 *Advice to users.* It is possible to spawn threads before MPI is initialized, but no  
43 MPI call other than MPI\_INITIALIZED should be executed by these threads, until  
44 MPI\_INIT\_THREAD is invoked by one thread (which, thereby, becomes the main  
45 thread). In particular, it is possible to enter the MPI execution with a multi-threaded  
46 process.

47 The level of thread support provided is a global property of the MPI process that can  
48 be specified only once, when MPI is initialized on that process (or before). Portable

third party libraries have to be written so as to accommodate any provided level of thread support. Otherwise, their usage will be restricted to specific level(s) of thread support. If such a library can run only with specific level(s) of thread support, e.g., only with `MPI_THREAD_MULTIPLE`, then `MPI_QUERY_THREAD` can be used to check whether the user initialized MPI to the correct level of thread support and, if not, raise an exception. (*End of advice to users.*)

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# Chapter 13

## I/O

### 13.1 Introduction

POSIX provides a model of a widely portable file system, but the portability and optimization needed for parallel I/O cannot be achieved with the POSIX interface.

The significant optimizations required for efficiency (e.g., grouping [47], collective buffering [7, 14, 48, 52, 59], and disk-directed I/O [43]) can only be implemented if the parallel I/O system provides a high-level interface supporting partitioning of file data among processes and a collective interface supporting complete transfers of global data structures between process memories and files. In addition, further efficiencies can be gained via support for asynchronous I/O, strided accesses, and control over physical file layout on storage devices (disks). The I/O environment described in this chapter provides these facilities.

Instead of defining I/O access modes to express the common patterns for accessing a shared file (broadcast, reduction, scatter, gather), we chose another approach in which data partitioning is expressed using derived datatypes. Compared to a limited set of predefined access patterns, this approach has the advantage of added flexibility and expressiveness.

#### 13.1.1 Definitions

**file** An MPI file is an ordered collection of typed data items. MPI supports random or sequential access to any integral set of these items. A file is opened collectively by a group of processes. All collective I/O calls on a file are collective over this group.

**displacement** A file *displacement* is an absolute byte position relative to the beginning of a file. The displacement defines the location where a *view* begins. Note that a “file displacement” is distinct from a “typemap displacement.”

**etype** An *etype* (*elementary* datatype) is the unit of data access and positioning. It can be any MPI predefined or derived datatype. Derived etypes can be constructed using any of the MPI datatype constructor routines, provided all resulting typemap displacements are non-negative and monotonically nondecreasing. Data access is performed in etype units, reading or writing whole data items of type etype. Offsets are expressed as a count of etypes; file pointers point to the beginning of etypes. Depending on context, the term “etype” is used to describe one of three aspects of an elementary datatype: a particular MPI type, a data item of that type, or the extent of that type.

**filetype** A *filetype* is the basis for partitioning a file among processes and defines a template for accessing the file. A filetype is either a single etype or a derived MPI datatype constructed from multiple instances of the same etype. In addition, the extent of any hole in the filetype must be a multiple of the etype's extent. The displacements in the typemap of the filetype are not required to be distinct, but they must be non-negative and monotonically nondecreasing.

**view** A *view* defines the current set of data visible and accessible from an open file as an ordered set of etypes. Each process has its own view of the file, defined by three quantities: a displacement, an etype, and a filetype. The pattern described by a filetype is repeated, beginning at the displacement, to define the view. The pattern of repetition is defined to be the same pattern that `MPI_TYPE_CONTIGUOUS` would produce if it were passed the filetype and an arbitrarily large count. Figure 13.1 shows how the tiling works; note that the filetype in this example must have explicit lower and upper bounds set in order for the initial and final holes to be repeated in the view. Views can be changed by the user during program execution. The default view is a linear byte stream (displacement is zero, etype and filetype equal to `MPI_BYTE`).

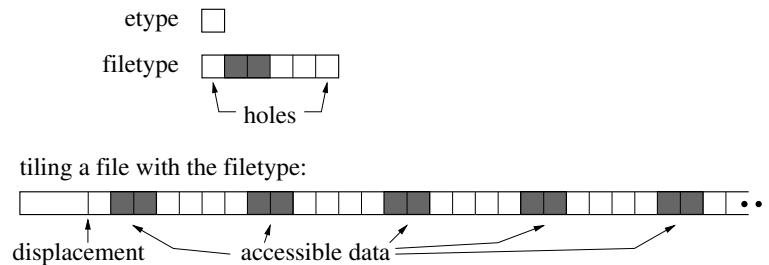


Figure 13.1: Etypes and filetypes

A group of processes can use complementary views to achieve a global data distribution such as a scatter/gather pattern (see Figure 13.2).

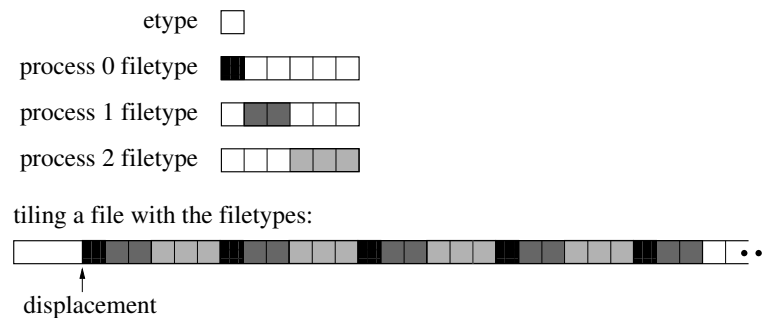


Figure 13.2: Partitioning a file among parallel processes

**offset** An *offset* is a position in the file relative to the current view, expressed as a count of etypes. Holes in the view's filetype are skipped when calculating this position. Offset 0 is the location of the first etype visible in the view (after skipping the displacement and any initial holes in the view). For example, an offset of 2 for process 1 in Figure 13.2 is the position of the 8th etype in the file after the displacement. An “explicit offset” is an offset that is used as a formal parameter in explicit data access routines.

**file size and end of file** The *size* of an MPI file is measured in bytes from the beginning of the file. A newly created file has a size of zero bytes. Using the size as an absolute displacement gives the position of the byte immediately following the last byte in the file. For any given view, the *end of file* is the offset of the first etype accessible in the current view starting after the last byte in the file.

**file pointer** A *file pointer* is an implicit offset maintained by MPI. “Individual file pointers” are file pointers that are local to each process that opened the file. A “shared file pointer” is a file pointer that is shared by the group of processes that opened the file.

**file handle** A *file handle* is an opaque object created by `MPI_FILE_OPEN` and freed by `MPI_FILE_CLOSE`. All operations on an open file reference the file through the file handle.

## 13.2 File Manipulation

### 13.2.1 Opening a File

`MPI_FILE_OPEN(comm, filename, amode, info, fh)`

IN	comm	communicator (handle)
IN	filename	name of file to open (string)
IN	amode	file access mode (integer)
IN	info	info object (handle)
OUT	fh	new file handle (handle)

```
int MPI_File_open(MPI_Comm comm, const char *filename, int amode,
                  MPI_Info info, MPI_File *fh)
```

```
MPI_File_open(comm, filename, amode, info, fh, ierror) BIND(C)
  TYPE(MPI_Comm), INTENT(IN) :: comm
  CHARACTER(LEN=*) INTENT(IN) :: filename
  INTEGER, INTENT(IN) :: amode
  TYPE(MPI_Info), INTENT(IN) :: info
  TYPE(MPI_File), INTENT(OUT) :: fh
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_OPEN(COMM, FILENAME, AMODE, INFO, FH, IERROR)
  CHARACTER*(*) FILENAME
  INTEGER COMM, AMODE, INFO, FH, IERROR
```

```
{static MPI::File MPI::File::Open(const MPI::Intracomm& comm,
    const char* filename, int amode, const MPI::Info& info) (binding
    deprecated, see Section 15.2) }
```

`MPI_FILE_OPEN` opens the file identified by the file name `filename` on all processes in the `comm` communicator group. `MPI_FILE_OPEN` is a collective routine: all processes must provide the same value for `amode`, and all processes must provide filenames that reference

the same file. (Values for `info` may vary.) `comm` must be an intracommunicator; it is erroneous to pass an intercommunicator to `MPI_FILE_OPEN`. Errors in `MPI_FILE_OPEN` are raised using the default file error handler (see Section 13.7, page 574). A process can open a file independently of other processes by using the `MPI_COMM_SELF` communicator. The file handle returned, `fh`, can be subsequently used to access the file until the file is closed using `MPI_FILE_CLOSE`. Before calling `MPI_FINALIZE`, the user is required to close (via `MPI_FILE_CLOSE`) all files that were opened with `MPI_FILE_OPEN`. Note that the communicator `comm` is unaffected by `MPI_FILE_OPEN` and continues to be usable in all MPI routines (e.g., `MPI_SEND`). Furthermore, the use of `comm` will not interfere with I/O behavior.

The format for specifying the file name in the `filename` argument is implementation dependent and must be documented by the implementation.

*Advice to implementors.* An implementation may require that `filename` include a string or strings specifying additional information about the file. Examples include the type of filesystem (e.g., a prefix of `ufs:`), a remote hostname (e.g., a prefix of `machine.univ.edu:`), or a file password (e.g., a suffix of `/PASSWORD=SECRET`). (*End of advice to implementors.*)

*Advice to users.* On some implementations of MPI, the file namespace may not be identical from all processes of all applications. For example, `"/tmp/foo"` may denote different files on different processes, or a single file may have many names, dependent on process location. The user is responsible for ensuring that a single file is referenced by the `filename` argument, as it may be impossible for an implementation to detect this type of namespace error. (*End of advice to users.*)

Initially, all processes view the file as a linear byte stream, and each process views data in its own native representation (no data representation conversion is performed). (POSIX files are linear byte streams in the native representation.) The file view can be changed via the `MPI_FILE_SET_VIEW` routine.

The following access modes are supported (specified in `amode`, a bit vector OR of the following integer constants):

- `MPI_MODE_RDONLY` — read only,
- `MPI_MODE_RDWR` — reading and writing,
- `MPI_MODE_WRONLY` — write only,
- `MPI_MODE_CREATE` — create the file if it does not exist,
- `MPI_MODE_EXCL` — error if creating file that already exists,
- `MPI_MODE_DELETE_ON_CLOSE` — delete file on close,
- `MPI_MODE_UNIQUE_OPEN` — file will not be concurrently opened elsewhere,
- `MPI_MODE_SEQUENTIAL` — file will only be accessed sequentially,
- `MPI_MODE_APPEND` — set initial position of all file pointers to end of file.



*Advice to users.* C/C++ users can use bit vector OR (|) to combine these constants; Fortran 90 users can use the bit vector IOR intrinsic. Fortran 77 users can use (non-portably) bit vector IOR on systems that support it. Alternatively, Fortran users can portably use integer addition to OR the constants (each constant should appear at most once in the addition.). (*End of advice to users.*)

*Advice to implementors.* The values of these constants must be defined such that the bitwise OR and the sum of any distinct set of these constants is equivalent. (*End of advice to implementors.*)

The modes MPI\_MODE\_RDONLY, MPI\_MODE\_RDWR, MPI\_MODE\_WRONLY, MPI\_MODE\_CREATE, and MPI\_MODE\_EXCL have identical semantics to their POSIX counterparts [38]. Exactly one of MPI\_MODE\_RDONLY, MPI\_MODE\_RDWR, or MPI\_MODE\_WRONLY, must be specified. It is erroneous to specify MPI\_MODE\_CREATE or MPI\_MODE\_EXCL in conjunction with MPI\_MODE\_RDONLY; it is erroneous to specify MPI\_MODE\_SEQUENTIAL together with MPI\_MODE\_RDWR.

The MPI\_MODE\_DELETE\_ON\_CLOSE mode causes the file to be deleted (equivalent to performing an MPI\_FILE\_DELETE) when the file is closed.

The MPI\_MODE\_UNIQUE\_OPEN mode allows an implementation to optimize access by eliminating the overhead of file locking. It is erroneous to open a file in this mode unless the file will not be concurrently opened elsewhere.

*Advice to users.* For MPI\_MODE\_UNIQUE\_OPEN, *not opened elsewhere* includes both inside and outside the MPI environment. In particular, one needs to be aware of potential external events which may open files (e.g., automated backup facilities). When MPI\_MODE\_UNIQUE\_OPEN is specified, the user is responsible for ensuring that no such external events take place. (*End of advice to users.*)

The MPI\_MODE\_SEQUENTIAL mode allows an implementation to optimize access to some sequential devices (tapes and network streams). It is erroneous to attempt nonsequential access to a file that has been opened in this mode.

Specifying MPI\_MODE\_APPEND only guarantees that all shared and individual file pointers are positioned at the initial end of file when MPI\_FILE\_OPEN returns. Subsequent positioning of file pointers is application dependent. In particular, the implementation does not ensure that all writes are appended.

Errors related to the access mode are raised in the class MPI\_ERR\_AMODE.

The info argument is used to provide information regarding file access patterns and file system specifics (see Section 13.2.8, page 519). The constant MPI\_INFO\_NULL can be used when no info needs to be specified.

*Advice to users.* Some file attributes are inherently implementation dependent (e.g., file permissions). These attributes must be set using either the info argument or facilities outside the scope of MPI. (*End of advice to users.*)

Files are opened by default using nonatomic mode file consistency semantics (see Section 13.6.1, page 565). The more stringent atomic mode consistency semantics, required for atomicity of conflicting accesses, can be set using MPI\_FILE\_SET\_ATOMICITY.

### 13.2.2 Closing a File

#### MPI\_FILE\_CLOSE(fh)

INOUT fh file handle (handle)

```
int MPI_File_close(MPI_File *fh)
```

```
MPI_File_close(fh, ierror) BIND(C)
    TYPE(MPI_File), INTENT(INOUT) :: fh
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_CLOSE(FH, IERROR)
    INTEGER FH, IERROR
```

```
{void MPI::File::Close() (binding deprecated, see Section 15.2) }
```

MPI\_FILE\_CLOSE first synchronizes file state (equivalent to performing an MPI\_FILE\_SYNC), then closes the file associated with fh. The file is deleted if it was opened with access mode MPI\_MODE\_DELETE\_ON\_CLOSE (equivalent to performing an MPI\_FILE\_DELETE). MPI\_FILE\_CLOSE is a collective routine.

*Advice to users.* If the file is deleted on close, and there are other processes currently accessing the file, the status of the file and the behavior of future accesses by these processes are implementation dependent. (*End of advice to users.*)

The user is responsible for ensuring that all outstanding nonblocking requests and split collective operations associated with fh made by a process have completed before that process calls MPI\_FILE\_CLOSE.

The MPI\_FILE\_CLOSE routine deallocates the file handle object and sets fh to MPI\_FILE\_NULL.

### 13.2.3 Deleting a File

#### MPI\_FILE\_DELETE(filename, info)

IN filename name of file to delete (string)

IN info info object (handle)

```
int MPI_File_delete(const char *filename, MPI_Info info)
```

```
MPI_File_delete(filename, info, ierror) BIND(C)
    CHARACTER(LEN=*) INTENT(IN) :: filename
    TYPE(MPI_Info), INTENT(IN) :: info
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_DELETE(FILENAME, INFO, IERROR)
    CHARACTER*(*) FILENAME
    INTEGER INFO, IERROR
```

```
{static void MPI::File::Delete(const char* filename,
                               const MPI::Info& info) (binding deprecated, see Section 15.2) }
```

MPI\_FILE\_DELETE deletes the file identified by the file name `filename`. If the file does not exist, MPI\_FILE\_DELETE raises an error in the class MPI\_ERR\_NO\_SUCH\_FILE.

The `info` argument can be used to provide information regarding file system specifics (see Section 13.2.8, page 519). The constant MPI\_INFO\_NULL refers to the null info, and can be used when no info needs to be specified.

If a process currently has the file open, the behavior of any access to the file (as well as the behavior of any outstanding accesses) is implementation dependent. In addition, whether an open file is deleted or not is also implementation dependent. If the file is not deleted, an error in the class MPI\_ERR\_FILE\_IN\_USE or MPI\_ERR\_ACCESS will be raised. Errors are raised using the default error handler (see Section 13.7, page 574).

#### 13.2.4 Resizing a File

MPI\_FILE\_SET\_SIZE(fh, size)

INOUT	fh	file handle (handle)
IN	size	size to truncate or expand file (integer)

```
int MPI_File_set_size(MPI_File fh, MPI_Offset size)
```

```
MPI_File_set_size(fh, size, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: size
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_SET_SIZE(FH, SIZE, IERROR)
    INTEGER FH, IERROR
    INTEGER(KIND=MPI_OFFSET_KIND) SIZE
```

```
{void MPI::File::Set_size(MPI::Offset size) (binding deprecated, see Section 15.2)
}
```

MPI\_FILE\_SET\_SIZE resizes the file associated with the file handle `fh`. `size` is measured in bytes from the beginning of the file. MPI\_FILE\_SET\_SIZE is collective; all processes in the group must pass identical values for `size`.

If `size` is smaller than the current file size, the file is truncated at the position defined by `size`. The implementation is free to deallocate file blocks located beyond this position.

If `size` is larger than the current file size, the file size becomes `size`. Regions of the file that have been previously written are unaffected. The values of data in the new regions in the file (those locations with displacements between old file size and `size`) are undefined. It is implementation dependent whether the MPI\_FILE\_SET\_SIZE routine allocates file space—use MPI\_FILE\_PREALLOCATE to force file space to be reserved.

MPI\_FILE\_SET\_SIZE does not affect the individual file pointers or the shared file pointer. If MPI\_MODE\_SEQUENTIAL mode was specified when the file was opened, it is erroneous to call this routine.

*Advice to users.* It is possible for the file pointers to point beyond the end of file after a `MPI_FILE_SET_SIZE` operation truncates a file. This is valid, and equivalent to seeking beyond the current end of file. (*End of advice to users.*)

All nonblocking requests and split collective operations on `fh` must be completed before calling `MPI_FILE_SET_SIZE`. Otherwise, calling `MPI_FILE_SET_SIZE` is erroneous. As far as consistency semantics are concerned, `MPI_FILE_SET_SIZE` is a write operation that conflicts with operations that access bytes at displacements between the old and new file sizes (see Section 13.6.1, page 565).

### 13.2.5 Preallocating Space for a File

`MPI_FILE_PREALLOCATE(fh, size)`

INOUT	<code>fh</code>	file handle (handle)
IN	<code>size</code>	size to preallocate file (integer)

`int MPI_File_preallocate(MPI_File fh, MPI_Offset size)`

```
MPI_File_preallocate(fh, size, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: size
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

`MPI_FILE_PREALLOCATE(FH, SIZE, IERROR)`

```
INTEGER FH, IERROR
INTEGER(KIND=MPI_OFFSET_KIND) SIZE
```

```
{void MPI::File::Preallocate(MPI::Offset size) (binding deprecated, see
    Section 15.2) }
```

`MPI_FILE_PREALLOCATE` ensures that storage space is allocated for the first `size` bytes of the file associated with `fh`. `MPI_FILE_PREALLOCATE` is collective; all processes in the group must pass identical values for `size`. Regions of the file that have previously been written are unaffected. For newly allocated regions of the file, `MPI_FILE_PREALLOCATE` has the same effect as writing undefined data. If `size` is larger than the current file size, the file size increases to `size`. If `size` is less than or equal to the current file size, the file size is unchanged.

The treatment of file pointers, pending nonblocking accesses, and file consistency is the same as with `MPI_FILE_SET_SIZE`. If `MPI_MODE_SEQUENTIAL` mode was specified when the file was opened, it is erroneous to call this routine.

*Advice to users.* In some implementations, file preallocation may be expensive. (*End of advice to users.*)

## 13.2.6 Querying the Size of a File

MPI\_FILE\_GET\_SIZE(fh, size)

IN	fh	file handle (handle)
OUT	size	size of the file in bytes (integer)

int MPI\_File\_get\_size(MPI\_File fh, MPI\_Offset \*size)

```

MPI_File_get_size(fh, size, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER(KIND=MPI_OFFSET_KIND), INTENT(OUT) :: size
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_FILE_GET_SIZE(FH, SIZE, IERROR)
    INTEGER FH, IERROR
    INTEGER(KIND=MPI_OFFSET_KIND) SIZE

```

{MPI::Offset MPI::File::Get\_size() const(*binding deprecated, see Section 15.2*) }

MPI\_FILE\_GET\_SIZE returns, in `size`, the current size in bytes of the file associated with the file handle `fh`. As far as consistency semantics are concerned, MPI\_FILE\_GET\_SIZE is a data access operation (see Section 13.6.1, page 565).

## 13.2.7 Querying File Parameters

MPI\_FILE\_GET\_GROUP(fh, group)

IN	fh	file handle (handle)
OUT	group	group which opened the file (handle)

int MPI\_File\_get\_group(MPI\_File fh, MPI\_Group \*group)

```

MPI_File_get_group(fh, group, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(MPI_Group), INTENT(OUT) :: group
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_FILE_GET_GROUP(FH, GROUP, IERROR)
    INTEGER FH, GROUP, IERROR

```

{MPI::Group MPI::File::Get\_group() const(*binding deprecated, see Section 15.2*) }

MPI\_FILE\_GET\_GROUP returns a duplicate of the group of the communicator used to open the file associated with `fh`. The group is returned in `group`. The user is responsible for freeing `group`.

1 MPI\_FILE\_GET\_AMODE(fh, amode)

2 IN fh file handle (handle)

3 OUT amode file access mode used to open the file (integer)

4  
5  
6 int MPI\_File\_get\_amode(MPI\_File fh, int \*amode)

7 MPI\_File\_get\_amode(fh, amode, ierror) BIND(C)

8 TYPE(MPI\_File), INTENT(IN) :: fh

9 INTEGER, INTENT(OUT) :: amode

10 INTEGER, OPTIONAL, INTENT(OUT) :: ierror

11  
12 MPI\_FILE\_GET\_AMODE(FH, AMODE, IERROR)

13 INTEGER FH, AMODE, IERROR

14 {int MPI::File::Get\_amode() const(*binding deprecated, see Section 15.2*) }

15  
16 MPI\_FILE\_GET\_AMODE returns, in amode, the access mode of the file associated with  
17 fh.

18  
19 **Example 13.1** In Fortran 77, decoding an amode bit vector will require a routine such as  
20 the following:

21  
22 SUBROUTINE BIT\_QUERY(TEST\_BIT, MAX\_BIT, AMODE, BIT\_FOUND)

23 !

24 ! TEST IF THE INPUT TEST\_BIT IS SET IN THE INPUT AMODE

25 ! IF SET, RETURN 1 IN BIT\_FOUND, 0 OTHERWISE

26 !

27 INTEGER TEST\_BIT, AMODE, BIT\_FOUND, CP\_AMODE, HIFOUND

28 BIT\_FOUND = 0

29 CP\_AMODE = AMODE

30 100 CONTINUE

31 LBIT = 0

32 HIFOUND = 0

33 DO 20 L = MAX\_BIT, 0, -1

34 MATCHER = 2\*\*L

35 IF (CP\_AMODE .GE. MATCHER .AND. HIFOUND .EQ. 0) THEN

36 HIFOUND = 1

37 LBIT = MATCHER

38 CP\_AMODE = CP\_AMODE - MATCHER

39 END IF

40 20 CONTINUE

41 IF (HIFOUND .EQ. 1 .AND. LBIT .EQ. TEST\_BIT) BIT\_FOUND = 1

42 IF (BIT\_FOUND .EQ. 0 .AND. HIFOUND .EQ. 1 .AND. &

43 CP\_AMODE .GT. 0) GO TO 100

44 END

45  
46 This routine could be called successively to decode amode, one bit at a time. For  
47 example, the following code fragment would check for MPI\_MODE\_RDONLY.

```

CALL BIT_QUERY(MPI_MODE_RDONLY, 30, AMODE, BIT_FOUND)
IF (BIT_FOUND .EQ. 1) THEN
    PRINT *, ' FOUND READ-ONLY BIT IN AMODE=', AMODE
ELSE
    PRINT *, ' READ-ONLY BIT NOT FOUND IN AMODE=', AMODE
END IF

```

### 13.2.8 File Info

Hints specified via info (see Section 9, page 381) allow a user to provide information such as file access patterns and file system specifics to direct optimization. Providing hints may enable an implementation to deliver increased I/O performance or minimize the use of system resources. However, hints do not change the semantics of any of the I/O interfaces. In other words, an implementation is free to ignore all hints. Hints are specified on a per file basis, in `MPI_FILE_OPEN`, `MPI_FILE_DELETE`, `MPI_FILE_SET_VIEW`, and `MPI_FILE_SET_INFO`, via the opaque info object. When an info object that specifies a subset of valid hints is passed to `MPI_FILE_SET_VIEW` or `MPI_FILE_SET_INFO`, there will be no effect on previously set or defaulted hints that the info does not specify.

*Advice to implementors.* It may happen that a program is coded with hints for one system, and later executes on another system that does not support these hints. In general, unsupported hints should simply be ignored. Needless to say, no hint can be mandatory. However, for each hint used by a specific implementation, a default value must be provided when the user does not specify a value for this hint. (*End of advice to implementors.*)

`MPI_FILE_SET_INFO(fh, info)`

INOUT	fh	file handle (handle)
IN	info	info object (handle)

`int MPI_File_set_info(MPI_File fh, MPI_Info info)`

```

MPI_File_set_info(fh, info, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(MPI_Info), INTENT(IN) :: info
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

`MPI_FILE_SET_INFO(FH, INFO, IERROR)`

`INTEGER FH, INFO, IERROR`

```

{void MPI::File::Set_info(const MPI::Info& info) (binding deprecated, see
    Section 15.2) }

```

`MPI_FILE_SET_INFO` sets new values for the hints of the file associated with `fh`. `MPI_FILE_SET_INFO` is a collective routine. The info object may be different on each process, but any info entries that an implementation requires to be the same on all processes must appear with the same value in each process's info object.

*Advice to users.* Many info items that an implementation can use when it creates or opens a file cannot easily be changed once the file has been created or opened. Thus, an implementation may ignore hints issued in this call that it would have accepted in an open call. (*End of advice to users.*)

`MPI_FILE_GET_INFO(fh, info_used)`

IN	fh	file handle (handle)
OUT	info_used	new info object (handle)

`int MPI_File_get_info(MPI_File fh, MPI_Info *info_used)`

```

MPI_File_get_info(fh, info_used, ierror) BIND(C)
  TYPE(MPI_File), INTENT(IN) :: fh
  TYPE(MPI_Info), INTENT(OUT) :: info_used
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

`MPI_FILE_GET_INFO(FH, INFO_USED, IERROR)`

`INTEGER FH, INFO_USED, IERROR`

`{MPI::Info MPI::File::Get_info() const(binding deprecated, see Section 15.2)}`

`MPI_FILE_GET_INFO` returns a new info object containing the hints of the file associated with `fh`. The current setting of all hints actually used by the system related to this open file is returned in `info_used`. If no such hints exist, a handle to a newly created info object is returned that contains no key/value pair. The user is responsible for freeing `info_used` via `MPI_INFO_FREE`.

*Advice to users.* The info object returned in `info_used` will contain all hints currently active for this file. This set of hints may be greater or smaller than the set of hints passed in to `MPI_FILE_OPEN`, `MPI_FILE_SET_VIEW`, and `MPI_FILE_SET_INFO`, as the system may not recognize some hints set by the user, and may recognize other hints that the user has not set. (*End of advice to users.*)

## Reserved File Hints

Some potentially useful hints (info key values) are outlined below. The following key values are reserved. An implementation is not required to interpret these key values, but if it does interpret the key value, it must provide the functionality described. (For more details on “info,” see Section 9, page 381.)

These hints mainly affect access patterns and the layout of data on parallel I/O devices. For each hint name introduced, we describe the purpose of the hint, and the type of the hint value. The “[**SAME**]” annotation specifies that the hint values provided by all participating processes must be identical; otherwise the program is erroneous. In addition, some hints are context dependent, and are only used by an implementation at specific times (e.g., `file_perm` is only useful during file creation).

**access\_style (comma separated list of strings):** This hint specifies the manner in which the file will be accessed until the file is closed or until the `access_style` key value is



altered. The hint value is a comma separated list of the following: `read_once`, `write_once`, `read_mostly`, `write_mostly`, `sequential`, `reverse_sequential`, and `random`.

**collective\_buffering (boolean) [SAME]:** This hint specifies whether the application may benefit from collective buffering. Collective buffering is an optimization performed on collective accesses. Accesses to the file are performed on behalf of all processes in the group by a number of target nodes. These target nodes coalesce small requests into large disk accesses. Valid values for this key are `true` and `false`. Collective buffering parameters are further directed via additional hints: `cb_block_size`, `cb_buffer_size`, and `cb_nodes`.

**cb\_block\_size (integer) [SAME]:** This hint specifies the block size to be used for collective buffering file access. *Target nodes* access data in chunks of this size. The chunks are distributed among target nodes in a round-robin (CYCLIC) pattern.

**cb\_buffer\_size (integer) [SAME]:** This hint specifies the total buffer space that can be used for collective buffering on each target node, usually a multiple of `cb_block_size`.

**cb\_nodes (integer) [SAME]:** This hint specifies the number of target nodes to be used for collective buffering.

**chunked (comma separated list of integers) [SAME]:** This hint specifies that the file consists of a multidimensional array that is often accessed by subarrays. The value for this hint is a comma separated list of array dimensions, starting from the most significant one (for an array stored in row-major order, as in C, the most significant dimension is the first one; for an array stored in column-major order, as in Fortran, the most significant dimension is the last one, and array dimensions should be reversed).

**chunked\_item (comma separated list of integers) [SAME]:** This hint specifies the size of each array entry, in bytes.

**chunked\_size (comma separated list of integers) [SAME]:** This hint specifies the dimensions of the subarrays. This is a comma separated list of array dimensions, starting from the most significant one.

**filename (string):** This hint specifies the file name used when the file was opened. If the implementation is capable of returning the file name of an open file, it will be returned using this key by `MPI_FILE_GET_INFO`. This key is ignored when passed to `MPI_FILE_OPEN`, `MPI_FILE_SET_VIEW`, `MPI_FILE_SET_INFO`, and `MPI_FILE_DELETE`.

**file\_perm (string) [SAME]:** This hint specifies the file permissions to use for file creation. Setting this hint is only useful when passed to `MPI_FILE_OPEN` with an `amode` that includes `MPI_MODE_CREATE`. The set of valid values for this key is implementation dependent.

**io\_node\_list (comma separated list of strings) [SAME]:** This hint specifies the list of I/O devices that should be used to store the file. This hint is most relevant when the file is created.

**nb\_proc (integer) [SAME]:** This hint specifies the number of parallel processes that will typically be assigned to run programs that access this file. This hint is most relevant when the file is created.

**num\_io\_nodes (integer) [SAME]:** This hint specifies the number of I/O devices in the system. This hint is most relevant when the file is created.

**striping\_factor (integer) [SAME]:** This hint specifies the number of I/O devices that the file should be striped across, and is relevant only when the file is created.

**striping\_unit (integer) [SAME]:** This hint specifies the suggested striping unit to be used for this file. The striping unit is the amount of consecutive data assigned to one I/O device before progressing to the next device, when striping across a number of devices. It is expressed in bytes. This hint is relevant only when the file is created.

### 13.3 File Views

**MPI\_FILE\_SET\_VIEW(fh, disp, etype, filetype, datarep, info)**

INOUT	fh	file handle (handle)
IN	disp	displacement (integer)
IN	etype	elementary datatype (handle)
IN	filetype	filetype (handle)
IN	datarep	data representation (string)
IN	info	info object (handle)

**int MPI\_File\_set\_view(MPI\_File fh, MPI\_Offset disp, MPI\_Datatype etype, MPI\_Datatype filetype, const char \*datarep, MPI\_Info info)**

```

MPI_File_set_view(fh, disp, etype, filetype, datarep, info, ierror) BIND(C)
  TYPE(MPI_File), INTENT(IN) :: fh
  INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: disp
  TYPE(MPI_Datatype), INTENT(IN) :: etype, filetype
  CHARACTER(LEN=*), INTENT(IN) :: datarep
  TYPE(MPI_Info), INTENT(IN) :: info
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_FILE_SET_VIEW(FH, DISP, ETYPE, FILETYPE, DATAREP, INFO, IERROR)
  INTEGER FH, ETYPE, FILETYPE, INFO, IERROR
  CHARACTER*(*) DATAREP
  INTEGER(KIND=MPI_OFFSET_KIND) DISP

```

```

{void MPI::File::Set_view(MPI::Offset disp, const MPI::Datatype& etype,
  const MPI::Datatype& filetype, const char* datarep,
  const MPI::Info& info) (binding deprecated, see Section 15.2) }

```

The **MPI\_FILE\_SET\_VIEW** routine changes the process's view of the data in the file. The start of the view is set to **disp**; the type of data is set to **etype**; the distribution of data to processes is set to **filetype**; and the representation of data in the file is set to **datarep**. In addition, **MPI\_FILE\_SET\_VIEW** resets the individual file pointers and the shared file pointer to zero. **MPI\_FILE\_SET\_VIEW** is collective; the values for **datarep** and the extents

of `etype` in the file data representation must be identical on all processes in the group; values for `disp`, `filetype`, and `info` may vary. The datatypes passed in `etype` and `filetype` must be committed.

The `etype` always specifies the data layout in the file. If `etype` is a portable datatype (see Section 2.4, page 11), the extent of `etype` is computed by scaling any displacements in the datatype to match the file data representation. If `etype` is not a portable datatype, no scaling is done when computing the extent of `etype`. The user must be careful when using nonportable `etypes` in heterogeneous environments; see Section 13.5.1, page 556 for further details.

If `MPI_MODE_SEQUENTIAL` mode was specified when the file was opened, the special displacement `MPI_DISPLACEMENT_CURRENT` must be passed in `disp`. This sets the displacement to the current position of the shared file pointer. `MPI_DISPLACEMENT_CURRENT` is invalid unless the amode for the file has `MPI_MODE_SEQUENTIAL` set.

*Rationale.* For some sequential files, such as those corresponding to magnetic tapes or streaming network connections, the *displacement* may not be meaningful. `MPI_DISPLACEMENT_CURRENT` allows the view to be changed for these types of files. (*End of rationale.*)

*Advice to implementors.* It is expected that a call to `MPI_FILE_SET_VIEW` will immediately follow `MPI_FILE_OPEN` in numerous instances. A high-quality implementation will ensure that this behavior is efficient. (*End of advice to implementors.*)

The `disp` displacement argument specifies the position (absolute offset in bytes from the beginning of the file) where the view begins.

*Advice to users.* `disp` can be used to skip headers or when the file includes a sequence of data segments that are to be accessed in different patterns (see Figure 13.3). Separate views, each using a different displacement and `filetype`, can be used to access each segment.

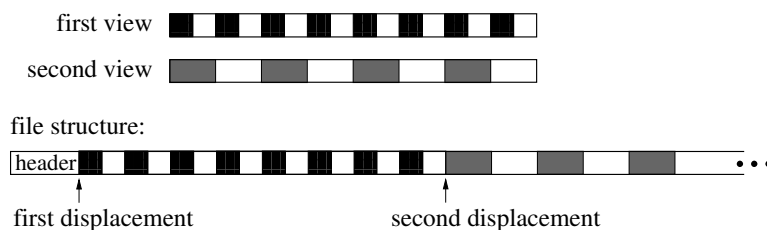


Figure 13.3: Displacements

(*End of advice to users.*)

An *etype* (*elementary* datatype) is the unit of data access and positioning. It can be any MPI predefined or derived datatype. Derived `etypes` can be constructed by using any of the MPI datatype constructor routines, provided all resulting typemap displacements are non-negative and monotonically nondecreasing. Data access is performed in `etype` units, reading or writing whole data items of type `etype`. Offsets are expressed as a count of `etypes`; file pointers point to the beginning of `etypes`.

*Advice to users.* In order to ensure interoperability in a heterogeneous environment, additional restrictions must be observed when constructing the **etype** (see Section 13.5, page 554). (*End of advice to users.*)

A filetype is either a single etype or a derived MPI datatype constructed from multiple instances of the same etype. In addition, the extent of any hole in the filetype must be a multiple of the etype’s extent. These displacements are not required to be distinct, but they cannot be negative, and they must be monotonically nondecreasing.

If the file is opened for writing, neither the **etype** nor the **filetype** is permitted to contain overlapping regions. This restriction is equivalent to the “datatype used in a receive cannot specify overlapping regions” restriction for communication. Note that filetypes from different processes may still overlap each other.

If filetype has holes in it, then the data in the holes is inaccessible to the calling process. However, the **disp**, **etype** and **filetype** arguments can be changed via future calls to **MPI\_FILE\_SET\_VIEW** to access a different part of the file.

It is erroneous to use absolute addresses in the construction of the **etype** and **filetype**.

The **info** argument is used to provide information regarding file access patterns and file system specifics to direct optimization (see Section 13.2.8, page 519). The constant **MPI\_INFO\_NULL** refers to the null info and can be used when no info needs to be specified.

The **datarep** argument is a string that specifies the representation of data in the file. See the file interoperability section (Section 13.5, page 554) for details and a discussion of valid values.

The user is responsible for ensuring that all nonblocking requests and split collective operations on **fh** have been completed before calling **MPI\_FILE\_SET\_VIEW**—otherwise, the call to **MPI\_FILE\_SET\_VIEW** is erroneous.

**MPI\_FILE\_GET\_VIEW**(**fh**, **disp**, **etype**, **filetype**, **datarep**)

IN	<b>fh</b>	file handle (handle)
OUT	<b>disp</b>	displacement (integer)
OUT	<b>etype</b>	elementary datatype (handle)
OUT	<b>filetype</b>	filetype (handle)
OUT	<b>datarep</b>	data representation (string)

```
int MPI_File_get_view(MPI_File fh, MPI_Offset *disp, MPI_Datatype *etype,
                     MPI_Datatype *filetype, char *datarep)
```

```
MPI_File_get_view(fh, disp, etype, filetype, datarep, ierror) BIND(C)
  TYPE(MPI_File), INTENT(IN) :: fh
  INTEGER(KIND=MPI_OFFSET_KIND), INTENT(OUT) :: disp
  TYPE(MPI_Datatype), INTENT(OUT) :: etype, filetype
  CHARACTER(LEN=*), INTENT(OUT) :: datarep
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_GET_VIEW(FH, DISP, ETYPE, FILETYPE, DATAREP, IERROR)
  INTEGER FH, ETYPE, FILETYPE, IERROR
  CHARACTER*(*) DATAREP
```

```

    INTEGER(KIND=MPI_OFFSET_KIND) DISP
{void MPI::File::Get_view(MPI::Offset& disp, MPI::Datatype& etype,
    MPI::Datatype& filetype, char* datarep) const(binding deprecated,
    see Section 15.2) }

```

MPI\_FILE\_GET\_VIEW returns the process's view of the data in the file. The current value of the displacement is returned in `disp`. The `etype` and `filetype` are new datatypes with typemaps equal to the typemaps of the current `etype` and `filetype`, respectively.

The data representation is returned in `datarep`. The user is responsible for ensuring that `datarep` is large enough to hold the returned data representation string. The length of a data representation string is limited to the value of `MPI_MAX_DATAREP_STRING`.

In addition, if a portable datatype was used to set the current view, then the corresponding datatype returned by `MPI_FILE_GET_VIEW` is also a portable datatype. If `etype` or `filetype` are derived datatypes, the user is responsible for freeing them. The `etype` and `filetype` returned are both in a committed state.

## 13.4 Data Access

### 13.4.1 Data Access Routines

Data is moved between files and processes by issuing read and write calls. There are three orthogonal aspects to data access: positioning (explicit offset *vs.* implicit file pointer), synchronism (blocking *vs.* nonblocking and split collective), and coordination (noncollective *vs.* collective). The following combinations of these data access routines, including two types of file pointers (individual and shared) are provided in Table 13.1.

positioning	synchronism	coordination	
		<i>noncollective</i>	<i>collective</i>
<i>explicit offsets</i>	<i>blocking</i>	MPI_FILE_READ_AT MPI_FILE_WRITE_AT	MPI_FILE_READ_AT_ALL MPI_FILE_WRITE_AT_ALL
	<i>nonblocking &amp; split collective</i>	MPI_FILE_IREAD_AT	MPI_FILE_READ_AT_ALL_BEGIN MPI_FILE_READ_AT_ALL_END
		MPI_FILE_IWRITE_AT	MPI_FILE_WRITE_AT_ALL_BEGIN MPI_FILE_WRITE_AT_ALL_END
<i>individual file pointers</i>	<i>blocking</i>	MPI_FILE_READ MPI_FILE_WRITE	MPI_FILE_READ_ALL MPI_FILE_WRITE_ALL
	<i>nonblocking &amp; split collective</i>	MPI_FILE_IREAD	MPI_FILE_READ_ALL_BEGIN MPI_FILE_READ_ALL_END
		MPI_FILE_IWRITE	MPI_FILE_WRITE_ALL_BEGIN MPI_FILE_WRITE_ALL_END
<i>shared file pointer</i>	<i>blocking</i>	MPI_FILE_READ_SHARED MPI_FILE_WRITE_SHARED	MPI_FILE_READ_ORDERED MPI_FILE_WRITE_ORDERED
	<i>nonblocking &amp; split collective</i>	MPI_FILE_IREAD_SHARED	MPI_FILE_READ_ORDERED_BEGIN MPI_FILE_READ_ORDERED_END
		MPI_FILE_IWRITE_SHARED	MPI_FILE_WRITE_ORDERED_BEGIN MPI_FILE_WRITE_ORDERED_END

Table 13.1: Data access routines

POSIX `read()`/`fread()` and `write()`/`fwrite()` are blocking, noncollective operations and

use individual file pointers. The MPI equivalents are `MPI_FILE_READ` and `MPI_FILE_WRITE`.

Implementations of data access routines may buffer data to improve performance. This does not affect reads, as the data is always available in the user’s buffer after a read operation completes. For writes, however, the `MPI_FILE_SYNC` routine provides the only guarantee that data has been transferred to the storage device.

## Positioning

MPI provides three types of positioning for data access routines: explicit offsets, individual file pointers, and shared file pointers. The different positioning methods may be mixed within the same program and do not affect each other.

The data access routines that accept explicit offsets contain `_AT` in their name (e.g., `MPI_FILE_WRITE_AT`). Explicit offset operations perform data access at the file position given directly as an argument—no file pointer is used nor updated. Note that this is not equivalent to an atomic seek-and-read or seek-and-write operation, as no “seek” is issued. Operations with explicit offsets are described in Section 13.4.2, page 528.

The names of the individual file pointer routines contain no positional qualifier (e.g., `MPI_FILE_WRITE`). Operations with individual file pointers are described in Section 13.4.3, page 532. The data access routines that use shared file pointers contain `_SHARED` or `_ORDERED` in their name (e.g., `MPI_FILE_WRITE_SHARED`). Operations with shared file pointers are described in Section 13.4.4, page 540.

The main semantic issues with MPI-maintained file pointers are how and when they are updated by I/O operations. In general, each I/O operation leaves the file pointer pointing to the next data item after the last one that is accessed by the operation. In a nonblocking or split collective operation, the pointer is updated by the call that initiates the I/O, possibly before the access completes.

More formally,

$$new\_file\_offset = old\_file\_offset + \frac{elements(datatype)}{elements(etype)} \times count$$

where *count* is the number of *datatype* items to be accessed, *elements(X)* is the number of predefined datatypes in the typemap of *X*, and *old\_file\_offset* is the value of the implicit offset before the call. The file position, *new\_file\_offset*, is in terms of a count of etypes relative to the current view.

## Synchronism

MPI supports blocking and nonblocking I/O routines.

A *blocking* I/O call will not return until the I/O request is completed.

A *nonblocking* I/O call initiates an I/O operation, but does not wait for it to complete. Given suitable hardware, this allows the transfer of data out/in the user’s buffer to proceed concurrently with computation. A separate *request complete* call (`MPI_WAIT`, `MPI_TEST`, or any of their variants) is needed to complete the I/O request, i.e., to confirm that the data has been read or written and that it is safe for the user to reuse the buffer. The nonblocking versions of the routines are named `MPI_FILE_IXXX`, where the *I* stands for immediate.

It is erroneous to access the local buffer of a nonblocking data access operation, or to use that buffer as the source or target of other communications, between the initiation and completion of the operation.

The split collective routines support a restricted form of “nonblocking” operations for collective data access (see Section 13.4.5, page 546).

### Coordination

Every noncollective data access routine `MPI_FILE_XXX` has a collective counterpart. For most routines, this counterpart is `MPI_FILE_XXX_ALL` or a pair of `MPI_FILE_XXX_BEGIN` and `MPI_FILE_XXX_END`. The counterparts to the `MPI_FILE_XXX_SHARED` routines are `MPI_FILE_XXX_ORDERED`.

The completion of a noncollective call only depends on the activity of the calling process. However, the completion of a collective call (which must be called by all members of the process group) may depend on the activity of the other processes participating in the collective call. See Section 13.6.4, page 568, for rules on semantics of collective calls.

Collective operations may perform much better than their noncollective counterparts, as global data accesses have significant potential for automatic optimization.

### Data Access Conventions

Data is moved between files and processes by calling read and write routines. Read routines move data from a file into memory. Write routines move data from memory into a file. The file is designated by a file handle, `fh`. The location of the file data is specified by an offset into the current view. The data in memory is specified by a triple: `buf`, `count`, and `datatype`. Upon completion, the amount of data accessed by the calling process is returned in a `status`.

An offset designates the starting position in the file for an access. The offset is always in `etype` units relative to the current view. Explicit offset routines pass `offset` as an argument (negative values are erroneous). The file pointer routines use implicit offsets maintained by MPI.

A data access routine attempts to transfer (read or write) `count` data items of type `datatype` between the user’s buffer `buf` and the file. The `datatype` passed to the routine must be a committed datatype. The layout of data in memory corresponding to `buf`, `count`, `datatype` is interpreted the same way as in MPI communication functions; see Section 3.2.2 on page 29 and Section 4.1.11 on page 119. The data is accessed from those parts of the file specified by the current view (Section 13.3, page 522). The type signature of `datatype` must match the type signature of some number of contiguous copies of the `etype` of the current view. As in a receive, it is erroneous to specify a `datatype` for reading that contains overlapping regions (areas of memory which would be stored into more than once).

The nonblocking data access routines indicate that MPI can start a data access and associate a request handle, `request`, with the I/O operation. Nonblocking operations are completed via `MPI_TEST`, `MPI_WAIT`, or any of their variants.

Data access operations, when completed, return the amount of data accessed in `status`.

*Advice to users.* To prevent problems with the argument copying and register optimization done by Fortran compilers, please note the hints in — [subsections “Problems Due to Data Copying and Sequence Association,” and “A Problem with Register Optimization” in Section 16.2.16, pages 671 and 677. subsections “Problems Due to] Sections 16.2.10-16.2.20, especially in Sections 16.2.12 and 16.2.13 on pages 671-674 about “Problems Due to Data Copying and Sequence Association with Subscript Triplets” and “Vector Subscripts”, and in Sections 16.2.16 to 16.2.19 on pages 677 to 686 about

ticket238-J.

ticket238-J.

ticket236-H.

ticket238-J.



“Optimization Problems”, “Code Movements and Register Optimization”, “Temporary Data Movements” and “Permanent Data Movements”. (*End of advice to users.*)

For blocking routines, `status` is returned directly. For nonblocking routines and split collective routines, `status` is returned when the operation is completed. The number of `datatype` entries and predefined elements accessed by the calling process can be extracted from `status` by using `MPI_GET_COUNT` and `MPI_GET_ELEMENTS`, respectively. The interpretation of the `MPI_ERROR` field is the same as for other operations — normally undefined, but meaningful if an MPI routine returns `MPI_ERR_IN_STATUS`. The user can pass (in C and Fortran) `MPI_STATUS_IGNORE` in the `status` argument if the return value of this argument is not needed. In C++, the `status` argument is optional. The `status` can be passed to `MPI_TEST_CANCELLED` to determine if the operation was cancelled. All other fields of `status` are undefined.

When reading, a program can detect the end of file by noting that the amount of data read is less than the amount requested. Writing past the end of file increases the file size. The amount of data accessed will be the amount requested, unless an error is raised (or a read reaches the end of file).

### 13.4.2 Data Access with Explicit Offsets

If `MPI_MODE_SEQUENTIAL` mode was specified when the file was opened, it is erroneous to call the routines in this section.

`MPI_FILE_READ_AT(fh, offset, buf, count, datatype, status)`

IN	fh	file handle (handle)
IN	offset	file offset (integer)
OUT	buf	initial address of buffer (choice)
IN	count	number of elements in buffer (integer)
IN	datatype	datatype of each buffer element (handle)
OUT	status	status object (Status)

```
int MPI_File_read_at(MPI_File fh, MPI_Offset offset, void *buf, int count,
                    MPI_Datatype datatype, MPI_Status *status)
```

```
MPI_File_read_at(fh, offset, buf, count, datatype, status, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
    TYPE(*), DIMENSION(..) :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_READ_AT(FH, OFFSET, BUF, COUNT, DATATYPE, STATUS, IERROR)
    <type> BUF(*)
    INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
```



```

    INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
1
2
3 {void MPI::File::Read_at(MPI::Offset offset, void* buf, int count,
4     const MPI::Datatype& datatype, MPI::Status& status) (binding
5     deprecated, see Section 15.2) }
6
7 {void MPI::File::Read_at(MPI::Offset offset, void* buf, int count,
8     const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
9     }
10
11

```

MPI\_FILE\_READ\_AT reads a file beginning at the position specified by offset.

MPI\_FILE\_READ\_AT\_ALL(fh, offset, buf, count, datatype, status)

IN	fh	file handle (handle)
IN	offset	file offset (integer)
OUT	buf	initial address of buffer (choice)
IN	count	number of elements in buffer (integer)
IN	datatype	datatype of each buffer element (handle)
OUT	status	status object (Status)

```

int MPI_File_read_at_all(MPI_File fh, MPI_Offset offset, void *buf,
    int count, MPI_Datatype datatype, MPI_Status *status)

```

MPI\_File\_read\_at\_all(fh, offset, buf, count, datatype, status, ierror)

```

    BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
    TYPE(*), DIMENSION(..) :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

MPI\_FILE\_READ\_AT\_ALL(FH, OFFSET, BUF, COUNT, DATATYPE, STATUS, IERROR)

```

    <type> BUF(*)
    INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
    INTEGER(KIND=MPI_OFFSET_KIND) OFFSET

```

```

{void MPI::File::Read_at_all(MPI::Offset offset, void* buf, int count,
    const MPI::Datatype& datatype, MPI::Status& status) (binding
    deprecated, see Section 15.2) }
{void MPI::File::Read_at_all(MPI::Offset offset, void* buf, int count,
    const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
    }

```

MPI\_FILE\_READ\_AT\_ALL is a collective version of the blocking MPI\_FILE\_READ\_AT interface.

ticket-248T.

```
1 MPI_FILE_WRITE_AT(fh, offset, buf, count, datatype, status)
```

```
2     INOUT    fh                file handle (handle)
```

```
4     IN      offset            file offset (integer)
```

```
5     IN      buf               initial address of buffer (choice)
```

```
6     IN      count             number of elements in buffer (integer)
```

```
8     IN      datatype          datatype of each buffer element (handle)
```

```
9     OUT     status            status object (Status)
```

```
11 int MPI_File_write_at(MPI_File fh, MPI_Offset offset, const void *buf,
12                      int count, MPI_Datatype datatype, MPI_Status *status)
```

```
13 MPI_File_write_at(fh, offset, buf, count, datatype, status, ierror) BIND(C)
14     TYPE(MPI_File), INTENT(IN) :: fh
15     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
16     TYPE(*), DIMENSION(..), INTENT(IN) :: buf
17     INTEGER, INTENT(IN) :: count
18     TYPE(MPI_Datatype), INTENT(IN) :: datatype
19     TYPE(MPI_Status) :: status
20     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
22 MPI_FILE_WRITE_AT(FH, OFFSET, BUF, COUNT, DATATYPE, STATUS, IERROR)
23     <type> BUF(*)
24     INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
25     INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
```

```
27 {void MPI::File::Write_at(MPI::Offset offset, const void* buf, int count,
28                          const MPI::Datatype& datatype, MPI::Status& status) (binding
29                          deprecated, see Section 15.2) }
30 {void MPI::File::Write_at(MPI::Offset offset, const void* buf, int count,
31                          const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
32 }
33 }
```

MPI\_FILE\_WRITE\_AT writes a file beginning at the position specified by offset.

```
37 MPI_FILE_WRITE_AT_ALL(fh, offset, buf, count, datatype, status)
```

```
38     INOUT    fh                file handle (handle)
```

```
39     IN      offset            file offset (integer)
```

```
40     IN      buf               initial address of buffer (choice)
```

```
41     IN      count             number of elements in buffer (integer)
```

```
42     IN      datatype          datatype of each buffer element (handle)
```

```
43     IN      status            status object (Status)
```

```
47 int MPI_File_write_at_all(MPI_File fh, MPI_Offset offset, const void *buf,
48                      int count, MPI_Datatype datatype, MPI_Status *status)
```

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ticket-248T.

```

1 MPI_File_write_at_all(fh, offset, buf, count, datatype, status, ierror)
2     BIND(C)
3     TYPE(MPI_File), INTENT(IN) :: fh
4     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
5     TYPE(*), DIMENSION(..), INTENT(IN) :: buf
6     INTEGER, INTENT(IN) :: count
7     TYPE(MPI_Datatype), INTENT(IN) :: datatype
8     TYPE(MPI_Status) :: status
9     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
10
11 MPI_FILE_WRITE_AT_ALL(FH, OFFSET, BUF, COUNT, DATATYPE, STATUS, IERROR)
12     <type> BUF(*)
13     INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
14     INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
15
16 {void MPI::File::Write_at_all(MPI::Offset offset, const void* buf,
17     int count, const MPI::Datatype& datatype,
18     MPI::Status& status) (binding deprecated, see Section 15.2) }
19
20 {void MPI::File::Write_at_all(MPI::Offset offset, const void* buf,
21     int count, const MPI::Datatype& datatype) (binding deprecated, see
22     Section 15.2) }
23
24 MPI_FILE_WRITE_AT_ALL is a collective version of the blocking
25 MPI_FILE_WRITE_AT interface.
26
27 MPI_FILE_IREAD_AT(fh, offset, buf, count, datatype, request)
28
29 IN      fh      file handle (handle)
30
31 IN      offset  file offset (integer)
32
33 OUT     buf      initial address of buffer (choice)
34
35 IN      count   number of elements in buffer (integer)
36
37 IN      datatype datatype of each buffer element (handle)
38
39 OUT     request  request object (handle)
40
41 int MPI_File_iread_at(MPI_File fh, MPI_Offset offset, void *buf, int count,
42     MPI_Datatype datatype, MPI_Request *request)
43
44 MPI_File_iread_at(fh, offset, buf, count, datatype, request, ierror)
45     BIND(C)
46     TYPE(MPI_File), INTENT(IN) :: fh
47     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
48     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
49     INTEGER, INTENT(IN) :: count
50     TYPE(MPI_Datatype), INTENT(IN) :: datatype
51     TYPE(MPI_Request), INTENT(OUT) :: request
52     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

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```

1 MPI_FILE_IREAD_AT(FH, OFFSET, BUF, COUNT, DATATYPE, REQUEST, IERROR)
2     <type> BUF(*)
3     INTEGER FH, COUNT, DATATYPE, REQUEST, IERROR
4     INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
5
6 {MPI::Request MPI::File::Iread_at(MPI::Offset offset, void* buf, int count,
7     const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
8     }

```

MPI\_FILE\_IREAD\_AT is a nonblocking version of the MPI\_FILE\_READ\_AT interface.

```

11 MPI_FILE_IWRITE_AT(fh, offset, buf, count, datatype, request)

```

13	INOUT	fh	file handle (handle)
14	IN	offset	file offset (integer)
15			
16	IN	buf	initial address of buffer (choice)
17	IN	count	number of elements in buffer (integer)
18			
19	IN	datatype	datatype of each buffer element (handle)
20	OUT	request	request object (handle)

```

21
ticket140. 22 int MPI_File_iwrite_at(MPI_File fh, MPI_Offset offset, const void *buf,
23     int count, MPI_Datatype datatype, MPI_Request *request)
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```

```

25 MPI_File_iwrite_at(fh, offset, buf, count, datatype, request, ierror)
26     BIND(C)
27     TYPE(MPI_File), INTENT(IN) :: fh
28     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
29     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
30     INTEGER, INTENT(IN) :: count
31     TYPE(MPI_Datatype), INTENT(IN) :: datatype
32     TYPE(MPI_Request), INTENT(OUT) :: request
33     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

34 MPI_FILE_IWRITE_AT(FH, OFFSET, BUF, COUNT, DATATYPE, REQUEST, IERROR)
35     <type> BUF(*)
36     INTEGER FH, COUNT, DATATYPE, REQUEST, IERROR
37     INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
38
39 {MPI::Request MPI::File::Iwrite_at(MPI::Offset offset, const void* buf,
40     int count, const MPI::Datatype& datatype) (binding deprecated, see
41     Section 15.2) }

```

MPI\_FILE\_IWRITE\_AT is a nonblocking version of the MPI\_FILE\_WRITE\_AT interface.

### 13.4.3 Data Access with Individual File Pointers

MPI maintains one individual file pointer per process per file handle. The current value of this pointer implicitly specifies the offset in the data access routines described in this

section. These routines only use and update the individual file pointers maintained by MPI. The shared file pointer is not used nor updated.

The individual file pointer routines have the same semantics as the data access with explicit offset routines described in Section 13.4.2, page 528, with the following modification:

- the offset is defined to be the current value of the MPI-maintained individual file pointer.

After an individual file pointer operation is initiated, the individual file pointer is updated to point to the next etype after the last one that will be accessed. The file pointer is updated relative to the current view of the file.

If MPI\_MODE\_SEQUENTIAL mode was specified when the file was opened, it is erroneous to call the routines in this section, with the exception of MPI\_FILE\_GET\_BYTE\_OFFSET.

MPI\_FILE\_READ(fh, buf, count, datatype, status)

INOUT	fh	file handle (handle)
OUT	buf	initial address of buffer (choice)
IN	count	number of elements in buffer (integer)
IN	datatype	datatype of each buffer element (handle)
OUT	status	status object (Status)

```
int MPI_File_read(MPI_File fh, void *buf, int count, MPI_Datatype datatype,
                 MPI_Status *status)
```

```
MPI_File_read(fh, buf, count, datatype, status, ierror) BIND(C)
```

```
TYPE(MPI_File), INTENT(IN) :: fh
TYPE(*), DIMENSION(..) :: buf
INTEGER, INTENT(IN) :: count
TYPE(MPI_Datatype), INTENT(IN) :: datatype
TYPE(MPI_Status) :: status
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_READ(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
```

```
<type> BUF(*)
INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
```

```
{void MPI::File::Read(void* buf, int count, const MPI::Datatype& datatype,
                     MPI::Status& status) (binding deprecated, see Section 15.2) }
```

```
{void MPI::File::Read(void* buf, int count,
                     const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
}
```

MPI\_FILE\_READ reads a file using the individual file pointer.

**Example 13.2** The following Fortran code fragment is an example of reading a file until the end of file is reached:

ticket-248T.

```

1  !   Read a preexisting input file until all data has been read.
2  !   Call routine "process_input" if all requested data is read.
3  !   The Fortran 90 "exit" statement exits the loop.
4
5      integer    bufsize, numread, totprocessed, status(MPI_STATUS_SIZE)
6      parameter (bufsize=100)
7      real       localbuffer(bufsize)
8      integer (kind=MPI_OFFSET_KIND) zero
9
10     zero = 0
11
12     call MPI_FILE_OPEN( MPI_COMM_WORLD, 'myoldfile', &
13                        MPI_MODE_RDONLY, MPI_INFO_NULL, myfh, ierr )
14     call MPI_FILE_SET_VIEW( myfh, zero, MPI_REAL, MPI_REAL, 'native', &
15                            MPI_INFO_NULL, ierr )
16     totprocessed = 0
17     do
18         call MPI_FILE_READ( myfh, localbuffer, bufsize, MPI_REAL, &
19                            status, ierr )
20         call MPI_GET_COUNT( status, MPI_REAL, numread, ierr )
21         call process_input( localbuffer, numread )
22         totprocessed = totprocessed + numread
23         if ( numread < bufsize ) exit
24     enddo
25
26     write(6,1001) numread, bufsize, totprocessed
27 1001 format( "No more data:  read", I3, "and expected", I3, &
28            "Processed total of", I6, "before terminating job." )
29
30     call MPI_FILE_CLOSE( myfh, ierr )
31
32
33

```

```

34 MPI_FILE_READ_ALL(fh, buf, count, datatype, status)

```

35	INOUT	fh	file handle (handle)
36	OUT	buf	initial address of buffer (choice)
37			
38	IN	count	number of elements in buffer (integer)
39	IN	datatype	datatype of each buffer element (handle)
40	OUT	status	status object (Status)
41			

```

42 int MPI_File_read_all(MPI_File fh, void *buf, int count,
43                      MPI_Datatype datatype, MPI_Status *status)

```

```

44 MPI_File_read_all(fh, buf, count, datatype, status, ierror) BIND(C)
45     TYPE(MPI_File), INTENT(IN) :: fh
46     TYPE(*), DIMENSION(..) :: buf
47     INTEGER, INTENT(IN) :: count

```

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```

    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_FILE_READ_ALL(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
    <type> BUF(*)
    INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR

{void MPI::File::Read_all(void* buf, int count,
    const MPI::Datatype& datatype, MPI::Status& status) (binding
    deprecated, see Section 15.2) }

{void MPI::File::Read_all(void* buf, int count,
    const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
    }

```

MPI\_FILE\_READ\_ALL is a collective version of the blocking MPI\_FILE\_READ interface.

MPI\_FILE\_WRITE(fh, buf, count, datatype, status)

INOUT	fh	file handle (handle)
IN	buf	initial address of buffer (choice)
IN	count	number of elements in buffer (integer)
IN	datatype	datatype of each buffer element (handle)
OUT	status	status object (Status)

```

int MPI_File_write(MPI_File fh, const void *buf, int count,
    MPI_Datatype datatype, MPI_Status *status)

```

```

MPI_File_write(fh, buf, count, datatype, status, ierror) BIND(C)

```

```

    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), INTENT(IN) :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_FILE_WRITE(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
    <type> BUF(*)
    INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR

```

```

{void MPI::File::Write(const void* buf, int count,
    const MPI::Datatype& datatype, MPI::Status& status) (binding
    deprecated, see Section 15.2) }

{void MPI::File::Write(const void* buf, int count,
    const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
    }

```

MPI\_FILE\_WRITE writes a file using the individual file pointer.

1 MPI\_FILE\_WRITE\_ALL(fh, buf, count, datatype, status)

2	INOUT	fh	file handle (handle)
3			
4	IN	buf	initial address of buffer (choice)
5	IN	count	number of elements in buffer (integer)
6	IN	datatype	datatype of each buffer element (handle)
7			
8	OUT	status	status object (Status)

9  
 ticket140. 10 int MPI\_File\_write\_all(MPI\_File fh, **const** void \*buf, int count,  
 11 MPI\_Datatype datatype, MPI\_Status \*status)

ticket-248T. 12 MPI\_File\_write\_all(fh, buf, count, datatype, status, ierror) BIND(C)  
 13 TYPE(MPI\_File), INTENT(IN) :: fh  
 14 TYPE(\*), DIMENSION(..), INTENT(IN) :: buf  
 15 INTEGER, INTENT(IN) :: count  
 16 TYPE(MPI\_Datatype), INTENT(IN) :: datatype  
 17 TYPE(MPI\_Status) :: status  
 18 INTEGER, OPTIONAL, INTENT(OUT) :: ierror  
 19  
 20 MPI\_FILE\_WRITE\_ALL(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)  
 21 <type> BUF(\*)  
 22 INTEGER FH, COUNT, DATATYPE, STATUS(MPI\_STATUS\_SIZE), IERROR  
 23  
 24 {void MPI::File::Write\_all(const void\* buf, int count,  
 25 const MPI::Datatype& datatype, MPI::Status& status) (*binding*  
 26 *deprecated, see Section 15.2*) }  
 27  
 28 {void MPI::File::Write\_all(const void\* buf, int count,  
 29 const MPI::Datatype& datatype) (*binding deprecated, see Section 15.2*)  
 30 }

31 MPI\_FILE\_WRITE\_ALL is a collective version of the blocking MPI\_FILE\_WRITE inter-  
 32 face.

33  
 34 MPI\_FILE\_IREAD(fh, buf, count, datatype, request)

35	INOUT	fh	file handle (handle)
36			
37	OUT	buf	initial address of buffer (choice)
38	IN	count	number of elements in buffer (integer)
39	IN	datatype	datatype of each buffer element (handle)
40			
41	OUT	request	request object (handle)

42  
 43 int MPI\_File\_iread(MPI\_File fh, void \*buf, int count,  
 44 MPI\_Datatype datatype, MPI\_Request \*request)

ticket-248T. 45 MPI\_File\_iread(fh, buf, count, datatype, request, ierror) BIND(C)  
 46 TYPE(MPI\_File), INTENT(IN) :: fh  
 47 TYPE(\*), DIMENSION(..), ASYNCHRONOUS :: buf  
 48



```

    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_FILE_IREAD(FH, BUF, COUNT, DATATYPE, REQUEST, IERROR)
    <type> BUF(*)
    INTEGER FH, COUNT, DATATYPE, REQUEST, IERROR

{MPI::Request MPI::File::Iread(void* buf, int count,
    const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
    }

```

MPI\_FILE\_IREAD is a nonblocking version of the MPI\_FILE\_READ interface.

**Example 13.3** The following Fortran code fragment illustrates file pointer update semantics:

```

!   Read the first twenty real words in a file into two local
!   buffers. Note that when the first MPI_FILE_IREAD returns,
!   the file pointer has been updated to point to the
!   eleventh real word in the file.

    integer    bufsize, req1, req2
    integer, dimension(MPI_STATUS_SIZE) :: status1, status2
    parameter (bufsize=10)
    real       buf1(bufsize), buf2(bufsize)
    integer (kind=MPI_OFFSET_KIND) zero

    zero = 0
    call MPI_FILE_OPEN( MPI_COMM_WORLD, 'myoldfile', &
                        MPI_MODE_RDONLY, MPI_INFO_NULL, myfh, ierr )
    call MPI_FILE_SET_VIEW( myfh, zero, MPI_REAL, MPI_REAL, 'native', &
                        MPI_INFO_NULL, ierr )
    call MPI_FILE_IREAD( myfh, buf1, bufsize, MPI_REAL, &
                        req1, ierr )
    call MPI_FILE_IREAD( myfh, buf2, bufsize, MPI_REAL, &
                        req2, ierr )

    call MPI_WAIT( req1, status1, ierr )
    call MPI_WAIT( req2, status2, ierr )

    call MPI_FILE_CLOSE( myfh, ierr )

```

1 MPI\_FILE\_IWRITE(fh, buf, count, datatype, request)

2	INOUT	fh	file handle (handle)
3			
4	IN	buf	initial address of buffer (choice)
5	IN	count	number of elements in buffer (integer)
6	IN	datatype	datatype of each buffer element (handle)
7			
8	OUT	request	request object (handle)

9  
 ticket140. 10 int MPI\_File\_ fwrite(MPI\_File fh, const void \*buf, int count,  
 11 MPI\_Datatype datatype, MPI\_Request \*request)

ticket-248T. 12 MPI\_File\_ fwrite(fh, buf, count, datatype, request, ierror) BIND(C)  
 13 TYPE(MPI\_File), INTENT(IN) :: fh  
 14 TYPE(\*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf  
 15 INTEGER, INTENT(IN) :: count  
 16 TYPE(MPI\_Datatype), INTENT(IN) :: datatype  
 17 TYPE(MPI\_Request), INTENT(OUT) :: request  
 18 INTEGER, OPTIONAL, INTENT(OUT) :: ierror  
 19  
 20 MPI\_FILE\_IWRITE(FH, BUF, COUNT, DATATYPE, REQUEST, IERROR)  
 21 <type> BUF(\*)  
 22 INTEGER FH, COUNT, DATATYPE, REQUEST, IERROR  
 23  
 24 {MPI::Request MPI::File::Iwrite(const void\* buf, int count,  
 25 const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)  
 26 }

27 MPI\_FILE\_IWRITE is a nonblocking version of the MPI\_FILE\_WRITE interface.

28  
 29  
 30 MPI\_FILE\_SEEK(fh, offset, whence)

31	INOUT	fh	file handle (handle)
32	IN	offset	file offset (integer)
33			
34	IN	whence	update mode (state)

35  
 ticket-248T. 36 int MPI\_File\_ seek(MPI\_File fh, MPI\_Offset offset, int whence)

37 MPI\_File\_ seek(fh, offset, whence, ierror) BIND(C)  
 38 TYPE(MPI\_File), INTENT(IN) :: fh  
 39 INTEGER(KIND=MPI\_OFFSET\_KIND), INTENT(IN) :: offset  
 40 INTEGER, INTENT(IN) :: whence  
 41 INTEGER, OPTIONAL, INTENT(OUT) :: ierror  
 42  
 43 MPI\_FILE\_SEEK(FH, OFFSET, WHENCE, IERROR)  
 44 INTEGER FH, WHENCE, IERROR  
 45 INTEGER(KIND=MPI\_OFFSET\_KIND) OFFSET  
 46  
 47 {void MPI::File::Seek(MPI::Offset offset, int whence) (binding deprecated, see  
 48 Section 15.2) }

MPI\_FILE\_SEEK updates the individual file pointer according to *whence*, which has the following possible values:

- MPI\_SEEK\_SET: the pointer is set to *offset*
- MPI\_SEEK\_CUR: the pointer is set to the current pointer position plus *offset*
- MPI\_SEEK\_END: the pointer is set to the end of file plus *offset*

The *offset* can be negative, which allows seeking backwards. It is erroneous to seek to a negative position in the view.

MPI\_FILE\_GET\_POSITION(*fh*, *offset*)

IN	<i>fh</i>	file handle (handle)
OUT	<i>offset</i>	offset of individual pointer (integer)

```
int MPI_File_get_position(MPI_File fh, MPI_Offset *offset)
```

```
MPI_File_get_position(fh, offset, ierror) BIND(C)
  TYPE(MPI_File), INTENT(IN) :: fh
  INTEGER(KIND=MPI_OFFSET_KIND), INTENT(OUT) :: offset
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_GET_POSITION(FH, OFFSET, IERROR)
```

```
  INTEGER FH, IERROR
  INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
```

```
{MPI::Offset MPI::File::Get_position() const(binding deprecated, see Section 15.2)
}
```

MPI\_FILE\_GET\_POSITION returns, in *offset*, the current position of the individual file pointer in etype units relative to the current view.

*Advice to users.* The *offset* can be used in a future call to MPI\_FILE\_SEEK using *whence* = MPI\_SEEK\_SET to return to the current position. To set the displacement to the current file pointer position, first convert *offset* into an absolute byte position using MPI\_FILE\_GET\_BYTE\_OFFSET, then call MPI\_FILE\_SET\_VIEW with the resulting displacement. (*End of advice to users.*)

MPI\_FILE\_GET\_BYTE\_OFFSET(*fh*, *offset*, *disp*)

IN	<i>fh</i>	file handle (handle)
IN	<i>offset</i>	offset (integer)
OUT	<i>disp</i>	absolute byte position of offset (integer)

```
int MPI_File_get_byte_offset(MPI_File fh, MPI_Offset offset,
                             MPI_Offset *disp)
```

```

1 MPI_File_get_byte_offset(fh, offset, disp, ierror) BIND(C)
2   TYPE(MPI_File), INTENT(IN) :: fh
3   INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
4   INTEGER(KIND=MPI_OFFSET_KIND), INTENT(OUT) :: disp
5   INTEGER, OPTIONAL, INTENT(OUT) :: ierror
6
7 MPI_FILE_GET_BYTE_OFFSET(FH, OFFSET, DISP, IERROR)
8   INTEGER FH, IERROR
9   INTEGER(KIND=MPI_OFFSET_KIND) OFFSET, DISP
10 {MPI::Offset MPI::File::Get_byte_offset(const MPI::Offset disp)
11     const(binding deprecated, see Section 15.2) }
12

```

MPI\_FILE\_GET\_BYTE\_OFFSET converts a view-relative offset into an absolute byte position. The absolute byte position (from the beginning of the file) of `offset` relative to the current view of `fh` is returned in `disp`.

#### 13.4.4 Data Access with Shared File Pointers

MPI maintains exactly one shared file pointer per collective MPI\_FILE\_OPEN (shared among processes in the communicator group). The current value of this pointer implicitly specifies the offset in the data access routines described in this section. These routines only use and update the shared file pointer maintained by MPI. The individual file pointers are not used nor updated.

The shared file pointer routines have the same semantics as the data access with explicit offset routines described in Section 13.4.2, page 528, with the following modifications:

- the `offset` is defined to be the current value of the MPI-maintained shared file pointer,
- the effect of multiple calls to shared file pointer routines is defined to behave as if the calls were serialized, and
- the use of shared file pointer routines is erroneous unless all processes use the same file view.

For the noncollective shared file pointer routines, the serialization ordering is not deterministic. The user needs to use other synchronization means to enforce a specific order.

After a shared file pointer operation is initiated, the shared file pointer is updated to point to the next etype after the last one that will be accessed. The file pointer is updated relative to the current view of the file.

## Noncollective Operations

MPI\_FILE\_READ\_SHARED(fh, buf, count, datatype, status)

INOUT	fh	file handle (handle)
OUT	buf	initial address of buffer (choice)
IN	count	number of elements in buffer (integer)
IN	datatype	datatype of each buffer element (handle)
OUT	status	status object (Status)

```
int MPI_File_read_shared(MPI_File fh, void *buf, int count,
                        MPI_Datatype datatype, MPI_Status *status)
```

MPI\_File\_read\_shared(fh, buf, count, datatype, status, ierror) BIND(C)

```
TYPE(MPI_File), INTENT(IN) :: fh
TYPE(*), DIMENSION(..) :: buf
INTEGER, INTENT(IN) :: count
TYPE(MPI_Datatype), INTENT(IN) :: datatype
TYPE(MPI_Status) :: status
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_READ_SHARED(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
<type> BUF(*)
INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
```

```
{void MPI::File::Read_shared(void* buf, int count,
                             const MPI::Datatype& datatype, MPI::Status& status) (binding
                             deprecated, see Section 15.2) }
```

```
{void MPI::File::Read_shared(void* buf, int count,
                             const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
}
```

MPI\_FILE\_READ\_SHARED reads a file using the shared file pointer.

MPI\_FILE\_WRITE\_SHARED(fh, buf, count, datatype, status)

INOUT	fh	file handle (handle)
IN	buf	initial address of buffer (choice)
IN	count	number of elements in buffer (integer)
IN	datatype	datatype of each buffer element (handle)
OUT	status	status object (Status)

```
int MPI_File_write_shared(MPI_File fh, const void *buf, int count,
                        MPI_Datatype datatype, MPI_Status *status)
```

MPI\_File\_write\_shared(fh, buf, count, datatype, status, ierror) BIND(C)

```

1      TYPE(MPI_File), INTENT(IN) :: fh
2      TYPE(*), DIMENSION(..), INTENT(IN) :: buf
3      INTEGER, INTENT(IN) :: count
4      TYPE(MPI_Datatype), INTENT(IN) :: datatype
5      TYPE(MPI_Status) :: status
6      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
7
8  MPI_FILE_WRITE_SHARED(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
9      <type> BUF(*)
10     INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
11
12 {void MPI::File::Write_shared(const void* buf, int count,
13     const MPI::Datatype& datatype, MPI::Status& status) (binding
14     deprecated, see Section 15.2) }
15
16 {void MPI::File::Write_shared(const void* buf, int count,
17     const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
18     }
19
20 MPI_FILE_WRITE_SHARED writes a file using the shared file pointer.
21
22 MPI_FILE_IREAD_SHARED(fh, buf, count, datatype, request)
23
24 INOUT    fh                file handle (handle)
25 OUT      buf                initial address of buffer (choice)
26 IN       count              number of elements in buffer (integer)
27 IN       datatype            datatype of each buffer element (handle)
28 OUT      request             request object (handle)
29
30 int MPI_File_iread_shared(MPI_File fh, void *buf, int count,
31     MPI_Datatype datatype, MPI_Request *request)
32
33 MPI_File_iread_shared(fh, buf, count, datatype, request, ierror) BIND(C)
34     TYPE(MPI_File), INTENT(IN) :: fh
35     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
36     INTEGER, INTENT(IN) :: count
37     TYPE(MPI_Datatype), INTENT(IN) :: datatype
38     TYPE(MPI_Request), INTENT(OUT) :: request
39     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
40
41 MPI_FILE_IREAD_SHARED(FH, BUF, COUNT, DATATYPE, REQUEST, IERROR)
42     <type> BUF(*)
43     INTEGER FH, COUNT, DATATYPE, REQUEST, IERROR
44
45 {MPI::Request MPI::File::Iread_shared(void* buf, int count,
46     const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
47     }
48
49 MPI_FILE_IREAD_SHARED is a nonblocking version of the MPI_FILE_READ_SHARED
50 interface.

```

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`MPI_FILE_IWRITE_SHARED(fh, buf, count, datatype, request)`

INOUT	fh	file handle (handle)
IN	buf	initial address of buffer (choice)
IN	count	number of elements in buffer (integer)
IN	datatype	datatype of each buffer element (handle)
OUT	request	request object (handle)

```
int MPI_File_iwrite_shared(MPI_File fh, const void *buf, int count,
                          MPI_Datatype datatype, MPI_Request *request)
```

```
MPI_File_iwrite_shared(fh, buf, count, datatype, request, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_IWRITE_SHARED(FH, BUF, COUNT, DATATYPE, REQUEST, IERROR)
    <type> BUF(*)
    INTEGER FH, COUNT, DATATYPE, REQUEST, IERROR
```

```
{MPI::Request MPI::File::Iwrite_shared(const void* buf, int count,
    const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
}
```

`MPI_FILE_IWRITE_SHARED` is a nonblocking version of the `MPI_FILE_WRITE_SHARED` interface.

### Collective Operations

The semantics of a collective access using a shared file pointer is that the accesses to the file will be in the order determined by the ranks of the processes within the group. For each process, the location in the file at which data is accessed is the position at which the shared file pointer would be after all processes whose ranks within the group less than that of this process had accessed their data. In addition, in order to prevent subsequent shared offset accesses by the same processes from interfering with this collective access, the call might return only after all the processes within the group have initiated their accesses. When the call returns, the shared file pointer points to the next etype accessible, according to the file view used by all processes, after the last etype requested.

*Advice to users.* There may be some programs in which all processes in the group need to access the file using the shared file pointer, but the program may not *require* that data be accessed in order of process rank. In such programs, using the shared ordered routines (e.g., `MPI_FILE_WRITE_ORDERED` rather than `MPI_FILE_WRITE_SHARED`) may enable an implementation to optimize access, improving performance. (*End of advice to users.*)

*Advice to implementors.* Accesses to the data requested by all processes do not have to be serialized. Once all processes have issued their requests, locations within the file for all accesses can be computed, and accesses can proceed independently from each other, possibly in parallel. (*End of advice to implementors.*)

**MPI\_FILE\_READ\_ORDERED(fh, buf, count, datatype, status)**

INOUT	fh	file handle (handle)
OUT	buf	initial address of buffer (choice)
IN	count	number of elements in buffer (integer)
IN	datatype	datatype of each buffer element (handle)
OUT	status	status object (Status)

```
int MPI_File_read_ordered(MPI_File fh, void *buf, int count,
    MPI_Datatype datatype, MPI_Status *status)
```

```
MPI_File_read_ordered(fh, buf, count, datatype, status, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..) :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_READ_ORDERED(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
    <type> BUF(*)
    INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
```

```
{void MPI::File::Read_ordered(void* buf, int count,
    const MPI::Datatype& datatype, MPI::Status& status) (binding
    deprecated, see Section 15.2) }
```

```
{void MPI::File::Read_ordered(void* buf, int count,
    const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
    }
```

MPI\_FILE\_READ\_ORDERED is a collective version of the MPI\_FILE\_READ\_SHARED interface.

**MPI\_FILE\_WRITE\_ORDERED(fh, buf, count, datatype, status)**

INOUT	fh	file handle (handle)
IN	buf	initial address of buffer (choice)
IN	count	number of elements in buffer (integer)
IN	datatype	datatype of each buffer element (handle)
OUT	status	status object (Status)



```

int MPI_File_write_ordered(MPI_File fh, const void *buf, int count,
    MPI_Datatype datatype, MPI_Status *status)

MPI_File_write_ordered(fh, buf, count, datatype, status, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), INTENT(IN) :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_FILE_WRITE_ORDERED(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
    <type> BUF(*)
    INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR

{void MPI::File::Write_ordered(const void* buf, int count,
    const MPI::Datatype& datatype, MPI::Status& status) (binding
    deprecated, see Section 15.2) }

{void MPI::File::Write_ordered(const void* buf, int count,
    const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
    }

    MPI_FILE_WRITE_ORDERED is a collective version of the MPI_FILE_WRITE_SHARED
    interface.

```

#### Seek

If MPI\_MODE\_SEQUENTIAL mode was specified when the file was opened, it is erroneous to call the following two routines (MPI\_FILE\_SEEK\_SHARED and MPI\_FILE\_GET\_POSITION\_SHARED).

```

MPI_FILE_SEEK_SHARED(fh, offset, whence)

    INOUT    fh                file handle (handle)
    IN       offset            file offset (integer)
    IN       whence            update mode (state)

int MPI_File_seek_shared(MPI_File fh, MPI_Offset offset, int whence)

MPI_File_seek_shared(fh, offset, whence, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
    INTEGER, INTENT(IN) :: whence
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_FILE_SEEK_SHARED(FH, OFFSET, WHENCE, IERROR)
    INTEGER FH, WHENCE, IERROR
    INTEGER(KIND=MPI_OFFSET_KIND) OFFSET

```

```
1 {void MPI::File::Seek_shared(MPI::Offset offset, int whence) (binding
2   deprecated, see Section 15.2) }
```

3 MPI\_FILE\_SEEK\_SHARED updates the shared file pointer according to whence, which  
4 has the following possible values:

- 6 • MPI\_SEEK\_SET: the pointer is set to offset
- 7
- 8 • MPI\_SEEK\_CUR: the pointer is set to the current pointer position plus offset
- 9
- 10 • MPI\_SEEK\_END: the pointer is set to the end of file plus offset

11 MPI\_FILE\_SEEK\_SHARED is collective; all the processes in the communicator group  
12 associated with the file handle fh must call MPI\_FILE\_SEEK\_SHARED with the same values  
13 for offset and whence.

14 The offset can be negative, which allows seeking backwards. It is erroneous to seek to  
15 a negative position in the view.

```
17
18 MPI_FILE_GET_POSITION_SHARED(fh, offset)
```

19	IN	fh	file handle (handle)
20			
21	OUT	offset	offset of shared pointer (integer)

```
22
23 int MPI_File_get_position_shared(MPI_File fh, MPI_Offset *offset)
```

```
24 MPI_File_get_position_shared(fh, offset, ierror) BIND(C)
25   TYPE(MPI_File), INTENT(IN) :: fh
26   INTEGER(KIND=MPI_OFFSET_KIND), INTENT(OUT) :: offset
27   INTEGER, OPTIONAL, INTENT(OUT) :: ierror
28
```

```
29 MPI_FILE_GET_POSITION_SHARED(FH, OFFSET, IERROR)
```

```
30   INTEGER FH, IERROR
31   INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
```

```
32 {MPI::Offset MPI::File::Get_position_shared() const (binding deprecated, see
33   Section 15.2) }
```

34

35 MPI\_FILE\_GET\_POSITION\_SHARED returns, in offset, the current position of the  
36 shared file pointer in etype units relative to the current view.

37

38 *Advice to users.* The offset can be used in a future call to MPI\_FILE\_SEEK\_SHARED  
39 using whence = MPI\_SEEK\_SET to return to the current position. To set the displace-  
40 ment to the current file pointer position, first convert offset into an absolute byte  
41 position using MPI\_FILE\_GET\_BYTE\_OFFSET, then call MPI\_FILE\_SET\_VIEW with  
42 the resulting displacement. (*End of advice to users.*)

#### 43 13.4.5 Split Collective Data Access Routines

44 MPI provides a restricted form of “nonblocking collective” I/O operations for all data ac-  
45 cesses using split collective data access routines. These routines are referred to as “split”  
46 collective routines because a single collective operation is split in two: a begin routine and  
47  
48

an end routine. The begin routine begins the operation, much like a nonblocking data access (e.g., `MPI_FILE_IREAD`). The end routine completes the operation, much like the matching test or wait (e.g., `MPI_WAIT`). As with nonblocking data access operations, the user must not use the buffer passed to a begin routine while the routine is outstanding; the operation must be completed with an end routine before it is safe to free buffers, etc.

Split collective data access operations on a file handle `fh` are subject to the semantic rules given below.

- On any MPI process, each file handle may have at most one active split collective operation at any time.
- Begin calls are collective over the group of processes that participated in the collective open and follow the ordering rules for collective calls.
- End calls are collective over the group of processes that participated in the collective open and follow the ordering rules for collective calls. Each end call matches the preceding begin call for the same collective operation. When an “end” call is made, exactly one unmatched “begin” call for the same operation must precede it.
- An implementation is free to implement any split collective data access routine using the corresponding blocking collective routine when either the begin call (e.g., `MPI_FILE_READ_ALL_BEGIN`) or the end call (e.g., `MPI_FILE_READ_ALL_END`) is issued. The begin and end calls are provided to allow the user and MPI implementation to optimize the collective operation.
- Split collective operations do not match the corresponding regular collective operation. For example, in a single collective read operation, an `MPI_FILE_READ_ALL` on one process does not match an `MPI_FILE_READ_ALL_BEGIN`/`MPI_FILE_READ_ALL_END` pair on another process.
- Split collective routines must specify a buffer in both the begin and end routines. By specifying the buffer that receives data in the end routine, we can avoid [many (though not all) of] the problems described in “A Problem with Code Movements and Register Optimization,” [Section 16.2.16, page 677. ]Section 16.2.17 on page 678, but not all of the problems described in Section 16.2.16 on page 677.
- No collective I/O operations are permitted on a file handle concurrently with a split collective access on that file handle (i.e., between the begin and end of the access). That is

```
MPI_File_read_all_begin(fh, ...);
...
MPI_File_read_all(fh, ...);
...
MPI_File_read_all_end(fh, ...);
```

is erroneous.

- In a multithreaded implementation, any split collective begin and end operation called by a process must be called from the same thread. This restriction is made to simplify

the implementation in the multithreaded case. (Note that we have already disallowed having two threads begin a split collective operation on the same file handle since only one split collective operation can be active on a file handle at any time.)

The arguments for these routines have the same meaning as for the equivalent collective versions (e.g., the argument definitions for `MPI_FILE_READ_ALL_BEGIN` and `MPI_FILE_READ_ALL_END` are equivalent to the arguments for `MPI_FILE_READ_ALL`). The begin routine (e.g., `MPI_FILE_READ_ALL_BEGIN`) begins a split collective operation that, when completed with the matching end routine (i.e., `MPI_FILE_READ_ALL_END`) produces the result as defined for the equivalent collective routine (i.e., `MPI_FILE_READ_ALL`).

For the purpose of consistency semantics (Section 13.6.1, page 565), a matched pair of split collective data access operations (e.g., `MPI_FILE_READ_ALL_BEGIN` and `MPI_FILE_READ_ALL_END`) compose a single data access.

`MPI_FILE_READ_AT_ALL_BEGIN(fh, offset, buf, count, datatype)`

IN	fh	file handle (handle)
IN	offset	file offset (integer)
OUT	buf	initial address of buffer (choice)
IN	count	number of elements in buffer (integer)
IN	datatype	datatype of each buffer element (handle)

```
int MPI_File_read_at_all_begin(MPI_File fh, MPI_Offset offset, void *buf,
    int count, MPI_Datatype datatype)
```

```
MPI_File_read_at_all_begin(fh, offset, buf, count, datatype, ierror)
```

```
    BIND(C)
```

```
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_READ_AT_ALL_BEGIN(FH, OFFSET, BUF, COUNT, DATATYPE, IERROR)
```

```
<type> BUF(*)
INTEGER FH, COUNT, DATATYPE, IERROR
INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
```

```
{void MPI::File::Read_at_all_begin(MPI::Offset offset, void* buf,
    int count, const MPI::Datatype& datatype) (binding deprecated, see
    Section 15.2) }
```

MPI\_FILE\_READ\_AT\_ALL\_END(fh, buf, status)

IN	fh	file handle (handle)
OUT	buf	initial address of buffer (choice)
OUT	status	status object (Status)

int MPI\_File\_read\_at\_all\_end(MPI\_File fh, void \*buf, MPI\_Status \*status)

MPI\_File\_read\_at\_all\_end(fh, buf, status, ierror) BIND(C)

```

TYPE(MPI_File), INTENT(IN) :: fh
TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
TYPE(MPI_Status) :: status
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

MPI\_FILE\_READ\_AT\_ALL\_END(FH, BUF, STATUS, IERROR)

```

<type> BUF(*)
INTEGER FH, STATUS(MPI_STATUS_SIZE), IERROR

```

{void MPI::File::Read\_at\_all\_end(void\* buf, MPI::Status& status) (*binding deprecated, see Section 15.2*) }

{void MPI::File::Read\_at\_all\_end(void\* buf) (*binding deprecated, see Section 15.2*) }

MPI\_FILE\_WRITE\_AT\_ALL\_BEGIN(fh, offset, buf, count, datatype)

INOUT	fh	file handle (handle)
IN	offset	file offset (integer)
IN	buf	initial address of buffer (choice)
IN	count	number of elements in buffer (integer)
IN	datatype	datatype of each buffer element (handle)

int MPI\_File\_write\_at\_all\_begin(MPI\_File fh, MPI\_Offset offset, const void \*buf, int count, MPI\_Datatype datatype)

MPI\_File\_write\_at\_all\_begin(fh, offset, buf, count, datatype, ierror) BIND(C)

```

TYPE(MPI_File), INTENT(IN) :: fh
INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
INTEGER, INTENT(IN) :: count
TYPE(MPI_Datatype), INTENT(IN) :: datatype
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

MPI\_FILE\_WRITE\_AT\_ALL\_BEGIN(FH, OFFSET, BUF, COUNT, DATATYPE, IERROR)

```

<type> BUF(*)
INTEGER FH, COUNT, DATATYPE, IERROR
INTEGER(KIND=MPI_OFFSET_KIND) OFFSET

```

```

1  {void MPI::File::Write_at_all_begin(MPI::Offset offset, const void* buf,
2      int count, const MPI::Datatype& datatype) (binding deprecated, see
3      Section 15.2) }
4
5
6  MPI_FILE_WRITE_AT_ALL_END(fh, buf, status)
7
8      INOUT    fh                file handle (handle)
9
10     IN       buf                initial address of buffer (choice)
11
12     OUT      status             status object (Status)
13
14
15  int MPI_File_write_at_all_end(MPI_File fh, const void *buf,
16      MPI_Status *status)
17
18  MPI_File_write_at_all_end(fh, buf, status, ierror) BIND(C)
19      TYPE(MPI_File), INTENT(IN) :: fh
20      TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
21      TYPE(MPI_Status) :: status
22      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
23
24  MPI_FILE_WRITE_AT_ALL_END(FH, BUF, STATUS, IERROR)
25      <type> BUF(*)
26      INTEGER FH, STATUS(MPI_STATUS_SIZE), IERROR
27
28  {void MPI::File::Write_at_all_end(const void* buf,
29      MPI::Status& status) (binding deprecated, see Section 15.2) }
30
31  {void MPI::File::Write_at_all_end(const void* buf) (binding deprecated, see
32      Section 15.2) }
33
34
35  MPI_FILE_READ_ALL_BEGIN(fh, buf, count, datatype)
36
37      INOUT    fh                file handle (handle)
38
39      OUT      buf                initial address of buffer (choice)
40
41      IN       count              number of elements in buffer (integer)
42
43      IN       datatype            datatype of each buffer element (handle)
44
45
46  int MPI_File_read_all_begin(MPI_File fh, void *buf, int count,
47      MPI_Datatype datatype)
48
49  MPI_File_read_all_begin(fh, buf, count, datatype, ierror) BIND(C)
50      TYPE(MPI_File), INTENT(IN) :: fh
51      TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
52      INTEGER, INTENT(IN) :: count
53      TYPE(MPI_Datatype), INTENT(IN) :: datatype
54      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
55
56  MPI_FILE_READ_ALL_BEGIN(FH, BUF, COUNT, DATATYPE, IERROR)

```

```

<type> BUF(*)
INTEGER FH, COUNT, DATATYPE, IERROR
{void MPI::File::Read_all_begin(void* buf, int count,
    const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
}

MPI_FILE_READ_ALL_END(fh, buf, status)

INOUT    fh                file handle (handle)
OUT      buf               initial address of buffer (choice)
OUT      status            status object (Status)

int MPI_File_read_all_end(MPI_File fh, void *buf, MPI_Status *status)

MPI_File_read_all_end(fh, buf, status, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_FILE_READ_ALL_END(FH, BUF, STATUS, IERROR)
    <type> BUF(*)
    INTEGER FH, STATUS(MPI_STATUS_SIZE), IERROR

{void MPI::File::Read_all_end(void* buf, MPI::Status& status) (binding
    deprecated, see Section 15.2) }

{void MPI::File::Read_all_end(void* buf) (binding deprecated, see Section 15.2) }

MPI_FILE_WRITE_ALL_BEGIN(fh, buf, count, datatype)

INOUT    fh                file handle (handle)
IN       buf               initial address of buffer (choice)
IN       count             number of elements in buffer (integer)
IN       datatype          datatype of each buffer element (handle)

int MPI_File_write_all_begin(MPI_File fh, const void *buf, int count,
    MPI_Datatype datatype)

MPI_File_write_all_begin(fh, buf, count, datatype, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_FILE_WRITE_ALL_BEGIN(FH, BUF, COUNT, DATATYPE, IERROR)

```

```

1      <type> BUF(*)
2      INTEGER FH, COUNT, DATATYPE, IERROR
3
4      {void MPI::File::Write_all_begin(const void* buf, int count,
5          const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
6          }
7
8
9      MPI_FILE_WRITE_ALL_END(fh, buf, status)
10
11      INOUT    fh                file handle (handle)
12
13      IN       buf              initial address of buffer (choice)
14
15      OUT      status           status object (Status)
16
17
18      int MPI_File_write_all_end(MPI_File fh, const void *buf,
19          MPI_Status *status)
20
21      MPI_File_write_all_end(fh, buf, status, ierror) BIND(C)
22      TYPE(MPI_File), INTENT(IN) :: fh
23      TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
24      TYPE(MPI_Status) :: status
25      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
26
27      MPI_FILE_WRITE_ALL_END(FH, BUF, STATUS, IERROR)
28
29      <type> BUF(*)
30      INTEGER FH, STATUS(MPI_STATUS_SIZE), IERROR
31
32      {void MPI::File::Write_all_end(const void* buf, MPI::Status& status) (binding
33          deprecated, see Section 15.2) }
34
35      {void MPI::File::Write_all_end(const void* buf) (binding deprecated, see
36          Section 15.2) }
37
38
39      MPI_FILE_READ_ORDERED_BEGIN(fh, buf, count, datatype)
40
41      INOUT    fh                file handle (handle)
42
43      OUT      buf              initial address of buffer (choice)
44
45      IN       count            number of elements in buffer (integer)
46
47      IN       datatype         datatype of each buffer element (handle)
48
49
50      int MPI_File_read_ordered_begin(MPI_File fh, void *buf, int count,
51          MPI_Datatype datatype)
52
53      MPI_File_read_ordered_begin(fh, buf, count, datatype, ierror) BIND(C)
54      TYPE(MPI_File), INTENT(IN) :: fh
55      TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
56      INTEGER, INTENT(IN) :: count
57      TYPE(MPI_Datatype), INTENT(IN) :: datatype

```



```

    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_FILE_READ_ORDERED_BEGIN(FH, BUF, COUNT, DATATYPE, IERROR)
    <type> BUF(*)
    INTEGER FH, COUNT, DATATYPE, IERROR
{void MPI::File::Read_ordered_begin(void* buf, int count,
    const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
}

MPI_FILE_READ_ORDERED_END(fh, buf, status)
    INOUT    fh                file handle (handle)
    OUT      buf                initial address of buffer (choice)
    OUT      status             status object (Status)

int MPI_File_read_ordered_end(MPI_File fh, void *buf, MPI_Status *status)
MPI_File_read_ordered_end(fh, buf, status, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_FILE_READ_ORDERED_END(FH, BUF, STATUS, IERROR)
    <type> BUF(*)
    INTEGER FH, STATUS(MPI_STATUS_SIZE), IERROR
{void MPI::File::Read_ordered_end(void* buf, MPI::Status& status) (binding
    deprecated, see Section 15.2) }
{void MPI::File::Read_ordered_end(void* buf) (binding deprecated, see Section 15.2)
}

MPI_FILE_WRITE_ORDERED_BEGIN(fh, buf, count, datatype)
    INOUT    fh                file handle (handle)
    IN      buf                initial address of buffer (choice)
    IN      count              number of elements in buffer (integer)
    IN      datatype            datatype of each buffer element (handle)

int MPI_File_write_ordered_begin(MPI_File fh, const void *buf, int count,
    MPI_Datatype datatype)
MPI_File_write_ordered_begin(fh, buf, count, datatype, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count

```

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```

1      TYPE(MPI_Datatype), INTENT(IN) :: datatype
2      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
3
4      MPI_FILE_WRITE_ORDERED_BEGIN(FH, BUF, COUNT, DATATYPE, IERROR)
5          <type> BUF(*)
6          INTEGER FH, COUNT, DATATYPE, IERROR
7
8      {void MPI::File::Write_ordered_begin(const void* buf, int count,
9          const MPI::Datatype& datatype) (binding deprecated, see Section 15.2)
10          }
11
12      MPI_FILE_WRITE_ORDERED_END(fh, buf, status)
13
14      INOUT    fh                file handle (handle)
15      IN       buf              initial address of buffer (choice)
16      OUT      status           status object (Status)
17
18      int MPI_File_write_ordered_end(MPI_File fh, const void *buf,
19          MPI_Status *status)
20
21      MPI_File_write_ordered_end(fh, buf, status, ierror) BIND(C)
22          TYPE(MPI_File), INTENT(IN) :: fh
23          TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
24          TYPE(MPI_Status) :: status
25          INTEGER, OPTIONAL, INTENT(OUT) :: ierror
26
27      MPI_FILE_WRITE_ORDERED_END(FH, BUF, STATUS, IERROR)
28          <type> BUF(*)
29          INTEGER FH, STATUS(MPI_STATUS_SIZE), IERROR
30
31      {void MPI::File::Write_ordered_end(const void* buf,
32          MPI::Status& status) (binding deprecated, see Section 15.2) }
33
34      {void MPI::File::Write_ordered_end(const void* buf) (binding deprecated, see
35          Section 15.2) }
36
37
38
39
40
41
42
43
44
45
46
47
48

```

## 13.5 File Interoperability

At the most basic level, file interoperability is the ability to read the information previously written to a file—not just the bits of data, but the actual information the bits represent. MPI guarantees full interoperability within a single MPI environment, and supports increased interoperability outside that environment through the external data representation (Section 13.5.2, page 558) as well as the data conversion functions (Section 13.5.3, page 559).

Interoperability within a single MPI environment (which could be considered “operability”) ensures that file data written by one MPI process can be read by any other MPI process, subject to the consistency constraints (see Section 13.6.1, page 565), provided that it would have been possible to start the two processes simultaneously and have them reside

in a single MPI\_COMM\_WORLD. Furthermore, both processes must see the same data values at every absolute byte offset in the file for which data was written.

This single environment file interoperability implies that file data is accessible regardless of the number of processes.

There are three aspects to file interoperability:

- transferring the bits,
- converting between different file structures, and
- converting between different machine representations.

The first two aspects of file interoperability are beyond the scope of this standard, as both are highly machine dependent. However, transferring the bits of a file into and out of the MPI environment (e.g., by writing a file to tape) is required to be supported by all MPI implementations. In particular, an implementation must specify how familiar operations similar to POSIX `cp`, `rm`, and `mv` can be performed on the file. Furthermore, it is expected that the facility provided maintains the correspondence between absolute byte offsets (e.g., after possible file structure conversion, the data bits at byte offset 102 in the MPI environment are at byte offset 102 outside the MPI environment). As an example, a simple off-line conversion utility that transfers and converts files between the native file system and the MPI environment would suffice, provided it maintained the offset coherence mentioned above. In a high-quality implementation of MPI, users will be able to manipulate MPI files using the same or similar tools that the native file system offers for manipulating its files.

The remaining aspect of file interoperability, converting between different machine representations, is supported by the typing information specified in the `etype` and `filetype`. This facility allows the information in files to be shared between any two applications, regardless of whether they use MPI, and regardless of the machine architectures on which they run.

MPI supports multiple data representations: “native,” “internal,” and “external32.” An implementation may support additional data representations. MPI also supports user-defined data representations (see Section 13.5.3, page 559). The “native” and “internal” data representations are implementation dependent, while the “external32” representation is common to all MPI implementations and facilitates file interoperability. The data representation is specified in the *datarep* argument to `MPI_FILE_SET_VIEW`.

*Advice to users.* MPI is not guaranteed to retain knowledge of what data representation was used when a file is written. Therefore, to correctly retrieve file data, an MPI application is responsible for specifying the same data representation as was used to create the file. (*End of advice to users.*)

**“native”** Data in this representation is stored in a file exactly as it is in memory. The advantage of this data representation is that data precision and I/O performance are not lost in type conversions with a purely homogeneous environment. The disadvantage is the loss of transparent interoperability within a heterogeneous MPI environment.

*Advice to users.* This data representation should only be used in a homogeneous MPI environment, or when the MPI application is capable of performing the data type conversions itself. (*End of advice to users.*)

*Advice to implementors.* When implementing read and write operations on top of MPI message-passing, the message data should be typed as `MPI_BYTE` to ensure that the message routines do not perform any type conversions on the data. (*End of advice to implementors.*)

**“internal”** This data representation can be used for I/O operations in a homogeneous or heterogeneous environment; the implementation will perform type conversions if necessary. The implementation is free to store data in any format of its choice, with the restriction that it will maintain constant extents for all predefined datatypes in any one file. The environment in which the resulting file can be reused is implementation-defined and must be documented by the implementation.

*Rationale.* This data representation allows the implementation to perform I/O efficiently in a heterogeneous environment, though with implementation-defined restrictions on how the file can be reused. (*End of rationale.*)

*Advice to implementors.* Since “external32” is a superset of the functionality provided by “internal,” an implementation may choose to implement “internal” as “external32.” (*End of advice to implementors.*)

**“external32”** This data representation states that read and write operations convert all data from and to the “external32” representation defined in Section 13.5.2, page 558. The data conversion rules for communication also apply to these conversions (see Section 3.3.2, page 25-27, of the MPI-1 document). The data on the storage medium is always in this canonical representation, and the data in memory is always in the local process’s native representation.

This data representation has several advantages. First, all processes reading the file in a heterogeneous MPI environment will automatically have the data converted to their respective native representations. Second, the file can be exported from one MPI environment and imported into any other MPI environment with the guarantee that the second environment will be able to read all the data in the file.

The disadvantage of this data representation is that data precision and I/O performance may be lost in data type conversions.

*Advice to implementors.* When implementing read and write operations on top of MPI message-passing, the message data should be converted to and from the “external32” representation in the client, and sent as type `MPI_BYTE`. This will avoid possible double data type conversions and the associated further loss of precision and performance. (*End of advice to implementors.*)

### 13.5.1 Datatypes for File Interoperability

If the file data representation is other than “native,” care must be taken in constructing etypes and filetypes. Any of the datatype constructor functions may be used; however, for those functions that accept displacements in bytes, the displacements must be specified in terms of their values in the file for the file data representation being used. MPI will interpret these byte displacements as is; no scaling will be done. The function `MPI_FILE_GET_TYPE_EXTENT` can be used to calculate the extents of datatypes in the

file. For etypes and filetypes that are portable datatypes (see Section 2.4, page 11), MPI will scale any displacements in the datatypes to match the file data representation. Datatypes passed as arguments to read/write routines specify the data layout in memory; therefore, they must always be constructed using displacements corresponding to displacements in memory.

*Advice to users.* One can logically think of the file as if it were stored in the memory of a file server. The **etype** and **filetype** are interpreted as if they were defined at this file server, by the same sequence of calls used to define them at the calling process. If the data representation is “native”, then this logical file server runs on the same architecture as the calling process, so that these types define the same data layout on the file as they would define in the memory of the calling process. If the **etype** and **filetype** are portable datatypes, then the data layout defined in the file is the same as would be defined in the calling process memory, up to a scaling factor. The routine `MPI_FILE_GET_FILE_EXTENT` can be used to calculate this scaling factor. Thus, two equivalent, portable datatypes will define the same data layout in the file, even in a heterogeneous environment with “internal”, “external32”, or user defined data representations. Otherwise, the **etype** and **filetype** must be constructed so that their typemap and extent are the same on any architecture. This can be achieved if they have an explicit upper bound and lower bound (defined either using `MPI_LB` and `MPI_UB` markers, or using `MPI_TYPE_CREATE_RESIZED`). This condition must also be fulfilled by any datatype that is used in the construction of the **etype** and **filetype**, if this datatype is replicated contiguously, either explicitly, by a call to `MPI_TYPE_CONTIGUOUS`, or implicitly, by a blocklength argument that is greater than one. If an **etype** or **filetype** is not portable, and has a typemap or extent that is architecture dependent, then the data layout specified by it on a file is implementation dependent.

File data representations other than “native” may be different from corresponding data representations in memory. Therefore, for these file data representations, it is important not to use hardwired byte offsets for file positioning, including the initial displacement that specifies the view. When a portable datatype (see Section 2.4, page 11) is used in a data access operation, any holes in the datatype are scaled to match the data representation. However, note that this technique only works when all the processes that created the file view build their etypes from the same predefined datatypes. For example, if one process uses an etype built from `MPI_INT` and another uses an etype built from `MPI_FLOAT`, the resulting views may be nonportable because the relative sizes of these types may differ from one data representation to another. (*End of advice to users.*)

`MPI_FILE_GET_TYPE_EXTENT(fh, datatype, extent)`

IN	fh	file handle (handle)
IN	datatype	datatype (handle)
OUT	extent	datatype extent (integer)

```

1  int MPI_File_get_type_extent(MPI_File fh, MPI_Datatype datatype,
2      MPI_Aint *extent)
3
4  MPI_File_get_type_extent(fh, datatype, extent, ierror) BIND(C)
5      TYPE(MPI_File), INTENT(IN) :: fh
6      TYPE(MPI_Datatype), INTENT(IN) :: datatype
7      INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: extent
8      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
9
10 MPI_FILE_GET_TYPE_EXTENT(FH, DATATYPE, EXTENT, IERROR)
11     INTEGER FH, DATATYPE, IERROR
12     INTEGER(KIND=MPI_ADDRESS_KIND) EXTENT
13
14 {MPI::Aint MPI::File::Get_type_extent(const MPI::Datatype& datatype)
15     const(binding deprecated, see Section 15.2) }
```

Returns the extent of `datatype` in the file `fh`. This extent will be the same for all processes accessing the file `fh`. If the current view uses a user-defined data representation (see Section 13.5.3, page 559), MPI uses the `dtype_file_extent_fn` callback to calculate the extent.

*Advice to implementors.* In the case of user-defined data representations, the extent of a derived datatype can be calculated by first determining the extents of the predefined datatypes in this derived datatype using `dtype_file_extent_fn` (see Section 13.5.3, page 559). (*End of advice to implementors.*)

### 13.5.2 External Data Representation: “external32”

All MPI implementations are required to support the data representation defined in this section. Support of optional datatypes (e.g., `MPI_INTEGER2`) is not required.

All floating point values are in big-endian IEEE format [36] of the appropriate size. Floating point values are represented by one of three IEEE formats. These are the IEEE “Single,” “Double,” and “Double Extended” formats, requiring 4, 8 and 16 bytes of storage, respectively. For the IEEE “Double Extended” formats, MPI specifies a Format Width of 16 bytes, with 15 exponent bits, bias = +16383, 112 fraction bits, and an encoding analogous to the “Double” format. All integral values are in two’s complement big-endian format. Big-endian means most significant byte at lowest address byte. For C `_Bool`, Fortran `LOGICAL` and C++ `bool`, 0 implies false and nonzero implies true. C `float` `_Complex`, `double` `_Complex` and long `double` `_Complex` as well as Fortran `COMPLEX` and `DOUBLE COMPLEX` are represented by a pair of floating point format values for the real and imaginary components. Characters are in ISO 8859-1 format [37]. Wide characters (of type `MPI_WCHAR`) are in Unicode format [60].

All signed numerals (e.g., `MPI_INT`, `MPI_REAL`) have the sign bit at the most significant bit. `MPI_COMPLEX` and `MPI_DOUBLE_COMPLEX` have the sign bit of the real and imaginary parts at the most significant bit of each part.

According to IEEE specifications [36], the “NaN” (not a number) is system dependent. It should not be interpreted within MPI as anything other than “NaN.”

*Advice to implementors.* The MPI treatment of “NaN” is similar to the approach used in XDR (see <ftp://ds.internic.net/rfc/rfc1832.txt>). (*End of advice to implementors.*)

All data is byte aligned, regardless of type. All data items are stored contiguously in the file (if the file view is contiguous).

*Advice to implementors.* All bytes of `LOGICAL` and `bool` must be checked to determine the value. (*End of advice to implementors.*)

*Advice to users.* The type `MPI_PACKED` is treated as bytes and is not converted. The user should be aware that `MPI_PACK` has the option of placing a header in the beginning of the pack buffer. (*End of advice to users.*)

The size of the predefined datatypes returned from `MPI_TYPE_CREATE_F90_REAL`, `MPI_TYPE_CREATE_F90_COMPLEX`, and `MPI_TYPE_CREATE_F90_INTEGER` are defined in Section 16.2.9, page 664.

*Advice to implementors.* When converting a larger size integer to a smaller size integer, only the less significant bytes are moved. Care must be taken to preserve the sign bit value. This allows no conversion errors if the data range is within the range of the smaller size integer. (*End of advice to implementors.*)

Table 13.2 specifies the sizes of predefined datatypes in “external32” format.

### 13.5.3 User-Defined Data Representations

There are two situations that cannot be handled by the required representations:

1. a user wants to write a file in a representation unknown to the implementation, and
2. a user wants to read a file written in a representation unknown to the implementation.

User-defined data representations allow the user to insert a third party converter into the I/O stream to do the data representation conversion.

```
MPI_REGISTER_DATAREP(datarep, read_conversion_fn, write_conversion_fn,
                     dtype_file_extent_fn, extra_state)
```

IN	<code>datarep</code>	data representation identifier (string)
IN	<code>read_conversion_fn</code>	function invoked to convert from file representation to native representation (function)
IN	<code>write_conversion_fn</code>	function invoked to convert from native representation to file representation (function)
IN	<code>dtype_file_extent_fn</code>	function invoked to get the extent of a datatype as represented in the file (function)
IN	<code>extra_state</code>	extra state

```
int MPI_Register_datarep(const char *datarep,
                        MPI_Datarep_conversion_function *read_conversion_fn,
                        MPI_Datarep_conversion_function *write_conversion_fn,
                        MPI_Datarep_extent_function *dtype_file_extent_fn,
                        void *extra_state)
```



Type	Length	Optional Type	Length
MPI_PACKED	1	MPI_INTEGER1	1
MPI_BYTE	1	MPI_INTEGER2	2
MPI_CHAR	1	MPI_INTEGER4	4
MPI_UNSIGNED_CHAR	1	MPI_INTEGER8	8
MPI_SIGNED_CHAR	1	MPI_INTEGER16	16
MPI_WCHAR	2		
MPI_SHORT	2	MPI_REAL2	2
MPI_UNSIGNED_SHORT	2	MPI_REAL4	4
MPI_INT	4	MPI_REAL8	8
MPI_UNSIGNED	4	MPI_REAL16	16
MPI_LONG	4		
MPI_UNSIGNED_LONG	4	MPI_COMPLEX4	2*2
MPI_LONG_LONG_INT	8	MPI_COMPLEX8	2*4
MPI_UNSIGNED_LONG_LONG	8	MPI_COMPLEX16	2*8
MPI_FLOAT	4	MPI_COMPLEX32	2*16
MPI_DOUBLE	8		
MPI_LONG_DOUBLE	16		
MPI_C_BOOL	[ticket171.] [4] 1		
MPI_INT8_T	1		
MPI_INT16_T	2		
MPI_INT32_T	4		
MPI_INT64_T	8		
MPI_UINT8_T	1		
MPI_UINT16_T	2		
MPI_UINT32_T	4		
MPI_UINT64_T	8		
MPI_AINT	8		
MPI_OFFSET	8		
MPI_C_COMPLEX	2*4		
MPI_C_FLOAT_COMPLEX	2*4		
MPI_C_DOUBLE_COMPLEX	2*8		
MPI_C_LONG_DOUBLE_COMPLEX	2*16		
MPI_CHARACTER	1		
MPI_LOGICAL	4		
MPI_INTEGER	4		
MPI_REAL	4		
MPI_DOUBLE_PRECISION	8		
MPI_COMPLEX	2*4		
MPI_DOUBLE_COMPLEX	2*8		

Table 13.2: “external32” sizes of predefined datatypes



```

MPI_Register_datarep(datarep, read_conversion_fn, write_conversion_fn,
    dtype_file_extent_fn, extra_state, ierror) BIND(C)
    CHARACTER(LEN=*), INTENT(IN) :: datarep
    PROCEDURE(MPI_Datarep_conversion_function) :: read_conversion_fn
    PROCEDURE(MPI_Datarep_conversion_function) :: write_conversion_fn
    PROCEDURE(MPI_Datarep_extent_function) :: dtype_file_extent_fn
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_REGISTER_DATAREP(DATAREP, READ_CONVERSION_FN, WRITE_CONVERSION_FN,
    DTYPE_FILE_EXTENT_FN, EXTRA_STATE, IERROR)
    CHARACTER*(*) DATAREP
    EXTERNAL READ_CONVERSION_FN, WRITE_CONVERSION_FN, DTYPE_FILE_EXTENT_FN
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE
    INTEGER IERROR

{void MPI::Register_datarep(const char* datarep,
    MPI::Datarep_conversion_function* read_conversion_fn,
    MPI::Datarep_conversion_function* write_conversion_fn,
    MPI::Datarep_extent_function* dtype_file_extent_fn,
    void* extra_state) (binding deprecated, see Section 15.2) }
```

The call associates `read_conversion_fn`, `write_conversion_fn`, and `dtype_file_extent_fn` with the data representation identifier `datarep`. `datarep` can then be used as an argument to `MPI_FILE_SET_VIEW`, causing subsequent data access operations to call the conversion functions to convert all data items accessed between file data representation and native representation. `MPI_REGISTER_DATAREP` is a local operation and only registers the data representation for the calling MPI process. If `datarep` is already defined, an error in the error class `MPI_ERR_DUP_DATAREP` is raised using the default file error handler (see Section 13.7, page 574). The length of a data representation string is limited to the value of `MPI_MAX_DATAREP_STRING`. `MPI_MAX_DATAREP_STRING` must have a value of at least 64. No routines are provided to delete data representations and free the associated resources; it is not expected that an application will generate them in significant numbers.

#### Extent Callback

```

typedef int MPI_Datarep_extent_function(MPI_Datatype datatype,
    MPI_Aint *file_extent, void *extra_state);

ABSTRACT INTERFACE
    SUBROUTINE MPI_Datarep_extent_function(datatype, extent, extra_state,
        ierror) BIND(C)
        TYPE(MPI_Datatype) :: datatype
        INTEGER(KIND=MPI_ADDRESS_KIND) :: extent, extra_state
        INTEGER :: ierror

SUBROUTINE DATAREP_EXTENT_FUNCTION(DATATYPE, EXTENT, EXTRA_STATE, IERROR)
    INTEGER DATATYPE, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTENT, EXTRA_STATE
```

```

1 {typedef void MPI::Datarep_extent_function(const MPI::Datatype& datatype,
2     MPI::Aint& file_extent, void* extra_state); (binding deprecated,
3     see Section 15.2)}

```

The function `dtype_file_extent_fn` must return, in `file_extent`, the number of bytes required to store `datatype` in the file representation. The function is passed, in `extra_state`, the argument that was passed to the `MPI_REGISTER_DATAREP` call. MPI will only call this routine with predefined datatypes employed by the user.

#### Datarep Conversion Functions

```

11 typedef int MPI_Datarep_conversion_function(void *userbuf,
12     MPI_Datatype datatype, int count, void *filebuf,
13     MPI_Offset position, void *extra_state);

```

#### ABSTRACT INTERFACE

```

16 SUBROUTINE MPI_Datarep_conversion_function(userbuf, datatype, count,
17     filebuf, position, extra_state, ierror) BIND(C)
18     USE, INTRINSIC :: ISO_C_BINDING, ONLY : C_PTR
19     TYPE(C_PTR), VALUE :: userbuf, filebuf
20     TYPE(MPI_Datatype) :: datatype
21     INTEGER :: count, ierror
22     INTEGER(KIND=MPI_OFFSET_KIND) :: position
23     INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state

```

```

24 SUBROUTINE DATAREP_CONVERSION_FUNCTION(USERBUF, DATATYPE, COUNT, FILEBUF,
25     POSITION, EXTRA_STATE, IERROR)
26     <TYPE> USERBUF(*), FILEBUF(*)
27     INTEGER COUNT, DATATYPE, IERROR
28     INTEGER(KIND=MPI_OFFSET_KIND) POSITION
29     INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE

```

```

31 {typedef void MPI::Datarep_conversion_function(void* userbuf,
32     MPI::Datatype& datatype, int count, void* filebuf,
33     MPI::Offset position, void* extra_state); (binding deprecated, see
34     Section 15.2)}

```

The function `read_conversion_fn` must convert from file data representation to native representation. Before calling this routine, MPI allocates and fills `filebuf` with `count` contiguous data items. The type of each data item matches the corresponding entry for the predefined datatype in the type signature of `datatype`. The function is passed, in `extra_state`, the argument that was passed to the `MPI_REGISTER_DATAREP` call. The function must copy all `count` data items from `filebuf` to `userbuf` in the distribution described by `datatype`, converting each data item from file representation to native representation. `datatype` will be equivalent to the datatype that the user passed to the read function. If the size of `datatype` is less than the size of the `count` data items, the conversion function must treat `datatype` as being contiguously tiled over the `userbuf`. The conversion function must begin storing converted data at the location in `userbuf` specified by `position` into the (tiled) `datatype`.

*Advice to users.* Although the conversion functions have similarities to `MPI_PACK`

and `MPI_UNPACK`, one should note the differences in the use of the arguments `count` and `position`. In the conversion functions, `count` is a count of data items (i.e., count of typemap entries of `datatype`), and `position` is an index into this typemap. In `MPI_PACK`, `incount` refers to the number of whole `datatypes`, and `position` is a number of bytes. (*End of advice to users.*)

*Advice to implementors.* A converted read operation could be implemented as follows:

1. Get file extent of all data items
2. Allocate a filebuf large enough to hold all count data items
3. Read data from file into filebuf
4. Call `read_conversion_fn` to convert data and place it into userbuf
5. Deallocate filebuf

(*End of advice to implementors.*)

If MPI cannot allocate a buffer large enough to hold all the data to be converted from a read operation, it may call the conversion function repeatedly using the same `datatype` and `userbuf`, and reading successive chunks of data to be converted in `filebuf`. For the first call (and in the case when all the data to be converted fits into `filebuf`), MPI will call the function with `position` set to zero. Data converted during this call will be stored in the `userbuf` according to the first `count` data items in `datatype`. Then in subsequent calls to the conversion function, MPI will increment the value in `position` by the `count` of items converted in the previous call, and the `userbuf` pointer will be unchanged.

*Rationale.* Passing the conversion function a position and one datatype for the transfer allows the conversion function to decode the datatype only once and cache an internal representation of it on the datatype. Then on subsequent calls, the conversion function can use the `position` to quickly find its place in the datatype and continue storing converted data where it left off at the end of the previous call. (*End of rationale.*)

*Advice to users.* Although the conversion function may usefully cache an internal representation on the datatype, it should not cache any state information specific to an ongoing conversion operation, since it is possible for the same datatype to be used concurrently in multiple conversion operations. (*End of advice to users.*)

The function `write_conversion_fn` must convert from native representation to file data representation. Before calling this routine, MPI allocates `filebuf` of a size large enough to hold `count` contiguous data items. The type of each data item matches the corresponding entry for the predefined datatype in the type signature of `datatype`. The function must copy `count` data items from `userbuf` in the distribution described by `datatype`, to a contiguous distribution in `filebuf`, converting each data item from native representation to file representation. If the size of `datatype` is less than the size of `count` data items, the conversion function must treat `datatype` as being contiguously tiled over the `userbuf`.

The function must begin copying at the location in `userbuf` specified by `position` into the (tiled) `datatype`. `datatype` will be equivalent to the datatype that the user passed to the

1 write function. The function is passed, in `extra_state`, the argument that was passed to the  
 2 `MPI_REGISTER_DATAREP` call.

3 The predefined constant `MPI_CONVERSION_FN_NULL` may be used as either  
 4 `write_conversion_fn` or `read_conversion_fn`. In that case, MPI will not attempt to invoke  
 5 `write_conversion_fn` or `read_conversion_fn`, respectively, but will perform the requested data  
 6 access using the native data representation.

7 An MPI implementation must ensure that all data accessed is converted, either by  
 8 using a filebuf large enough to hold all the requested data items or else by making repeated  
 9 calls to the conversion function with the same `datatype` argument and appropriate values  
 10 for `position`.

11 An implementation will only invoke the callback routines in this section  
 12 (`read_conversion_fn`, `write_conversion_fn`, and `dtype_file_extent_fn`) when one of the read or  
 13 write routines in Section 13.4, page 525, or `MPI_FILE_GET_TYPE_EXTENT` is called by  
 14 the user. `dtype_file_extent_fn` will only be passed predefined datatypes employed by the  
 15 user. The conversion functions will only be passed datatypes equivalent to those that the  
 16 user has passed to one of the routines noted above.

17 The conversion functions must be reentrant. User defined data representations are  
 18 restricted to use byte alignment for all types. Furthermore, it is erroneous for the conversion  
 19 functions to call any collective routines or to free `datatype`.

20 The conversion functions should return an error code. If the returned error code has  
 21 a value other than `MPI_SUCCESS`, the implementation will raise an error in the class  
 22 `MPI_ERR_CONVERSION`.

#### 24 13.5.4 Matching Data Representations

25 It is the user's responsibility to ensure that the data representation used to read data from  
 26 a file is *compatible* with the data representation that was used to write that data to the file.

27 In general, using the same data representation name when writing and reading a file  
 28 does not guarantee that the representation is compatible. Similarly, using different repre-  
 29 sentation names on two different implementations may yield compatible representations.

30 Compatibility can be obtained when “external32” representation is used, although  
 31 precision may be lost and the performance may be less than when “native” representation is  
 32 used. Compatibility is guaranteed using “external32” provided at least one of the following  
 33 conditions is met.

- 34 • The data access routines directly use types enumerated in Section 13.5.2, page 558,  
 35 that are supported by all implementations participating in the I/O. The predefined  
 36 type used to write a data item must also be used to read a data item.
- 37 • In the case of Fortran 90 programs, the programs participating in the data accesses  
 38 obtain compatible datatypes using MPI routines that specify precision and/or range  
 39 (Section 16.2.9, page 660).
- 40 • For any given data item, the programs participating in the data accesses use compat-  
 41 ible predefined types to write and read the data item.

42 User-defined data representations may be used to provide an implementation compat-  
 43 ibility with another implementation's “native” or “internal” representation.

*Advice to users.* Section 16.2.9, page 660, defines routines that support the use of matching datatypes in heterogeneous environments and contains examples illustrating their use. (*End of advice to users.*)

## 13.6 Consistency and Semantics

### 13.6.1 File Consistency

Consistency semantics define the outcome of multiple accesses to a single file. All file accesses in MPI are relative to a specific file handle created from a collective open. MPI provides three levels of consistency: sequential consistency among all accesses using a single file handle, sequential consistency among all accesses using file handles created from a single collective open with atomic mode enabled, and user-imposed consistency among accesses other than the above. Sequential consistency means the behavior of a set of operations will be as if the operations were performed in some serial order consistent with program order; each access appears atomic, although the exact ordering of accesses is unspecified. User-imposed consistency may be obtained using program order and calls to `MPI_FILE_SYNC`.

Let  $FH_1$  be the set of file handles created from one particular collective open of the file *FOO*, and  $FH_2$  be the set of file handles created from a different collective open of *FOO*. Note that nothing restrictive is said about  $FH_1$  and  $FH_2$ : the sizes of  $FH_1$  and  $FH_2$  may be different, the groups of processes used for each open may or may not intersect, the file handles in  $FH_1$  may be destroyed before those in  $FH_2$  are created, etc. Consider the following three cases: a single file handle (e.g.,  $fh_1 \in FH_1$ ), two file handles created from a single collective open (e.g.,  $fh_{1a} \in FH_1$  and  $fh_{1b} \in FH_1$ ), and two file handles from different collective opens (e.g.,  $fh_1 \in FH_1$  and  $fh_2 \in FH_2$ ).

For the purpose of consistency semantics, a matched pair (Section 13.4.5, page 546) of split collective data access operations (e.g., `MPI_FILE_READ_ALL_BEGIN` and `MPI_FILE_READ_ALL_END`) compose a single data access operation. Similarly, a non-blocking data access routine (e.g., `MPI_FILE_IREAD`) and the routine which completes the request (e.g., `MPI_WAIT`) also compose a single data access operation. For all cases below, these data access operations are subject to the same constraints as blocking data access operations.

*Advice to users.* For an `MPI_FILE_IREAD` and `MPI_WAIT` pair, the operation begins when `MPI_FILE_IREAD` is called and ends when `MPI_WAIT` returns. (*End of advice to users.*)

Assume that  $A_1$  and  $A_2$  are two data access operations. Let  $D_1$  ( $D_2$ ) be the set of absolute byte displacements of every byte accessed in  $A_1$  ( $A_2$ ). The two data accesses *overlap* if  $D_1 \cap D_2 \neq \emptyset$ . The two data accesses *conflict* if they overlap and at least one is a write access.

Let  $SEQ_{fh}$  be a sequence of file operations on a single file handle, bracketed by `MPI_FILE_SYNC`s on that file handle. (Both opening and closing a file implicitly perform an `MPI_FILE_SYNC`.)  $SEQ_{fh}$  is a “write sequence” if any of the data access operations in the sequence are writes or if any of the file manipulation operations in the sequence change the state of the file (e.g., `MPI_FILE_SET_SIZE` or `MPI_FILE_PREALLOCATE`). Given two sequences,  $SEQ_1$  and  $SEQ_2$ , we say they are not *concurrent* if one sequence is guaranteed to completely precede the other (temporally).

The requirements for guaranteeing sequential consistency among all accesses to a particular file are divided into the three cases given below. If any of these requirements are not met, then the value of all data in that file is implementation dependent.

**Case 1:**  $fh_1 \in FH_1$  All operations on  $fh_1$  are sequentially consistent if atomic mode is set. If nonatomic mode is set, then all operations on  $fh_1$  are sequentially consistent if they are either nonconcurrent, nonconflicting, or both.

**Case 2:**  $fh_{1a} \in FH_1$  and  $fh_{1b} \in FH_1$  Assume  $A_1$  is a data access operation using  $fh_{1a}$ , and  $A_2$  is a data access operation using  $fh_{1b}$ . If for any access  $A_1$ , there is no access  $A_2$  that conflicts with  $A_1$ , then MPI guarantees sequential consistency.

However, unlike POSIX semantics, the default MPI semantics for conflicting accesses do not guarantee sequential consistency. If  $A_1$  and  $A_2$  conflict, sequential consistency can be guaranteed by either enabling atomic mode via the `MPI_FILE_SET_ATOMICITY` routine, or meeting the condition described in Case 3 below.

**Case 3:**  $fh_1 \in FH_1$  and  $fh_2 \in FH_2$  Consider access to a single file using file handles from distinct collective opens. In order to guarantee sequential consistency, `MPI_FILE_SYNC` must be used (both opening and closing a file implicitly perform an `MPI_FILE_SYNC`).

Sequential consistency is guaranteed among accesses to a single file if for any write sequence  $SEQ_1$  to the file, there is no sequence  $SEQ_2$  to the file which is *concurrent* with  $SEQ_1$ . To guarantee sequential consistency when there are write sequences, `MPI_FILE_SYNC` must be used together with a mechanism that guarantees nonconcurrency of the sequences.

See the examples in Section 13.6.10, page 570, for further clarification of some of these consistency semantics.

`MPI_FILE_SET_ATOMICITY(fh, flag)`

INOUT	fh	file handle (handle)
IN	flag	true to set atomic mode, false to set nonatomic mode (logical)

`int MPI_File_set_atomicity(MPI_File fh, int flag)`

```
MPI_File_set_atomicity(fh, flag, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    LOGICAL, INTENT(IN) :: flag
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_SET_ATOMICITY(FH, FLAG, IERROR)
    INTEGER FH, IERROR
    LOGICAL FLAG
```

```
{void MPI::File::Set_atomicity(bool flag) (binding deprecated, see Section 15.2) }
```

Let  $FH$  be the set of file handles created by one collective open. The consistency semantics for data access operations using  $FH$  is set by collectively calling

MPI\_FILE\_SET\_ATOMICITY on *FH*. MPI\_FILE\_SET\_ATOMICITY is collective; all processes in the group must pass identical values for *fh* and *flag*. If *flag* is true, atomic mode is set; if *flag* is false, nonatomic mode is set.

Changing the consistency semantics for an open file only affects new data accesses. All completed data accesses are guaranteed to abide by the consistency semantics in effect during their execution. Nonblocking data accesses and split collective operations that have not completed (e.g., via MPI\_WAIT) are only guaranteed to abide by nonatomic mode consistency semantics.

*Advice to implementors.* Since the semantics guaranteed by atomic mode are stronger than those guaranteed by nonatomic mode, an implementation is free to adhere to the more stringent atomic mode semantics for outstanding requests. (*End of advice to implementors.*)

MPI\_FILE\_GET\_ATOMICITY(*fh*, *flag*)

IN	<i>fh</i>	file handle (handle)
OUT	<i>flag</i>	true if atomic mode, false if nonatomic mode (logical)

```
int MPI_File_get_atomicity(MPI_File fh, int *flag)
```

```
MPI_File_get_atomicity(fh, flag, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    LOGICAL, INTENT(OUT) :: flag
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_GET_ATOMICITY(FH, FLAG, IERROR)
    INTEGER FH, IERROR
    LOGICAL FLAG
```

```
{bool MPI::File::Get_atomicity() const(binding deprecated, see Section 15.2) }
```

MPI\_FILE\_GET\_ATOMICITY returns the current consistency semantics for data access operations on the set of file handles created by one collective open. If *flag* is true, atomic mode is enabled; if *flag* is false, nonatomic mode is enabled.

MPI\_FILE\_SYNC(*fh*)

INOUT	<i>fh</i>	file handle (handle)
-------	-----------	----------------------

```
int MPI_File_sync(MPI_File fh)
```

```
MPI_File_sync(fh, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_FILE_SYNC(FH, IERROR)
    INTEGER FH, IERROR
```

```
{void MPI::File::Sync() (binding deprecated, see Section 15.2) }
```



Calling `MPI_FILE_SYNC` with `fh` causes all previous writes to `fh` by the calling process to be transferred to the storage device. If other processes have made updates to the storage device, then all such updates become visible to subsequent reads of `fh` by the calling process. `MPI_FILE_SYNC` may be necessary to ensure sequential consistency in certain cases (see above).

`MPI_FILE_SYNC` is a collective operation.

The user is responsible for ensuring that all nonblocking requests and split collective operations on `fh` have been completed before calling `MPI_FILE_SYNC`—otherwise, the call to `MPI_FILE_SYNC` is erroneous.

### 13.6.2 Random Access vs. Sequential Files

MPI distinguishes ordinary random access files from sequential stream files, such as pipes and tape files. Sequential stream files must be opened with the `MPI_MODE_SEQUENTIAL` flag set in the `amode`. For these files, the only permitted data access operations are shared file pointer reads and writes. Filetypes and etypes with holes are erroneous. In addition, the notion of file pointer is not meaningful; therefore, calls to `MPI_FILE_SEEK_SHARED` and `MPI_FILE_GET_POSITION_SHARED` are erroneous, and the pointer update rules specified for the data access routines do not apply. The amount of data accessed by a data access operation will be the amount requested unless the end of file is reached or an error is raised.

*Rationale.* This implies that reading on a pipe will always wait until the requested amount of data is available or until the process writing to the pipe has issued an end of file. (*End of rationale.*)

Finally, for some sequential files, such as those corresponding to magnetic tapes or streaming network connections, writes to the file may be destructive. In other words, a write may act as a truncate (a `MPI_FILE_SET_SIZE` with `size` set to the current position) followed by the write.

### 13.6.3 Progress

The progress rules of MPI are both a promise to users and a set of constraints on implementors. In cases where the progress rules restrict possible implementation choices more than the interface specification alone, the progress rules take precedence.

All blocking routines must complete in finite time unless an exceptional condition (such as resource exhaustion) causes an error.

Nonblocking data access routines inherit the following progress rule from nonblocking point to point communication: a nonblocking write is equivalent to a nonblocking send for which a receive is eventually posted, and a nonblocking read is equivalent to a nonblocking receive for which a send is eventually posted.

Finally, an implementation is free to delay progress of collective routines until all processes in the group associated with the collective call have invoked the routine. Once all processes in the group have invoked the routine, the progress rule of the equivalent noncollective routine must be followed.

### 13.6.4 Collective File Operations

Collective file operations are subject to the same restrictions as collective communication operations. For a complete discussion, please refer to the semantics set forth in Section 5.13



on page 228.

Collective file operations are collective over a dup of the communicator used to open the file—this duplicate communicator is implicitly specified via the file handle argument. Different processes can pass different values for other arguments of a collective routine unless specified otherwise.

### 13.6.5 Type Matching

The type matching rules for I/O mimic the type matching rules for communication with one exception: if `etype` is `MPI_BYTE`, then this matches any `datatype` in a data access operation. In general, the `etype` of data items written must match the `etype` used to read the items, and for each data access operation, the current `etype` must also match the type declaration of the data access buffer.

*Advice to users.* In most cases, use of `MPI_BYTE` as a wild card will defeat the file interoperability features of MPI. File interoperability can only perform automatic conversion between heterogeneous data representations when the exact datatypes accessed are explicitly specified. (*End of advice to users.*)

### 13.6.6 Miscellaneous Clarifications

Once an I/O routine completes, it is safe to free any opaque objects passed as arguments to that routine. For example, the `comm` and `info` used in an `MPI_FILE_OPEN`, or the `etype` and `filetype` used in an `MPI_FILE_SET_VIEW`, can be freed without affecting access to the file. Note that for nonblocking routines and split collective operations, the operation must be completed before it is safe to reuse data buffers passed as arguments.

As in communication, datatypes must be committed before they can be used in file manipulation or data access operations. For example, the `etype` and `filetype` must be committed before calling `MPI_FILE_SET_VIEW`, and the `datatype` must be committed before calling `MPI_FILE_READ` or `MPI_FILE_WRITE`.

### 13.6.7 MPI\_Offset Type

`MPI_Offset` is an integer type of size sufficient to represent the size (in bytes) of the largest file supported by MPI. Displacements and offsets are always specified as values of type `MPI_Offset`.

In Fortran, the corresponding integer is an integer `[of kind]` with kind parameter `MPI_OFFSET_KIND`, `[defined in mpif.h and the mpi module]` which is defined in the `mpi_f08 module`, the `mpi module` and the `mpif.h include file`.

In Fortran 77 environments that do not support `KIND` parameters, `MPI_Offset` arguments should be declared as an `INTEGER` of suitable size. The language interoperability implications for `MPI_Offset` are similar to those for addresses (see Section 16.3, page 692).

### 13.6.8 Logical vs. Physical File Layout

MPI specifies how the data should be laid out in a virtual file structure (the view), not how that file structure is to be stored on one or more disks. Specification of the physical file structure was avoided because it is expected that the mapping of files to disks will be system specific, and any specific control over file layout would therefore restrict program

portability. However, there are still cases where some information may be necessary to optimize file layout. This information can be provided as *hints* specified via *info* when a file is created (see Section 13.2.8, page 519).

### 13.6.9 File Size

The size of a file may be increased by writing to the file after the current end of file. The size may also be changed by calling MPI *size changing* routines, such as `MPI_FILE_SET_SIZE`. A call to a size changing routine does not necessarily change the file size. For example, calling `MPI_FILE_PREALLOCATE` with a size less than the current size does not change the size.

Consider a set of bytes that has been written to a file since the most recent call to a size changing routine, or since `MPI_FILE_OPEN` if no such routine has been called. Let the *high byte* be the byte in that set with the largest displacement. The file size is the larger of

- One plus the displacement of the high byte.
- The size immediately after the size changing routine, or `MPI_FILE_OPEN`, returned.

When applying consistency semantics, calls to `MPI_FILE_SET_SIZE` and `MPI_FILE_PREALLOCATE` are considered writes to the file (which conflict with operations that access bytes at displacements between the old and new file sizes), and `MPI_FILE_GET_SIZE` is considered a read of the file (which overlaps with all accesses to the file).

*Advice to users.* Any sequence of operations containing the collective routines `MPI_FILE_SET_SIZE` and `MPI_FILE_PREALLOCATE` is a write sequence. As such, sequential consistency in nonatomic mode is not guaranteed unless the conditions in Section 13.6.1, page 565, are satisfied. (*End of advice to users.*)

File pointer update semantics (i.e., file pointers are updated by the amount accessed) are only guaranteed if file size changes are sequentially consistent.

*Advice to users.* Consider the following example. Given two operations made by separate processes to a file containing 100 bytes: an `MPI_FILE_READ` of 10 bytes and an `MPI_FILE_SET_SIZE` to 0 bytes. If the user does not enforce sequential consistency between these two operations, the file pointer may be updated by the amount requested (10 bytes) even if the amount accessed is zero bytes. (*End of advice to users.*)

### 13.6.10 Examples

The examples in this section illustrate the application of the MPI consistency and semantics guarantees. These address

- conflicting accesses on file handles obtained from a single collective open, and
- all accesses on file handles obtained from two separate collective opens.

The simplest way to achieve consistency for conflicting accesses is to obtain sequential consistency by setting atomic mode. For the code below, process 1 will read either 0 or 10 integers. If the latter, every element of `b` will be 5. If nonatomic mode is set, the results of the read are undefined.

```

/* Process 0 */
int i, a[10] ;
int TRUE = 1;

for ( i=0;i<10;i++)
    a[i] = 5 ;

MPI_File_open( MPI_COMM_WORLD, "workfile",
               MPI_MODE_RDWR | MPI_MODE_CREATE, MPI_INFO_NULL, &fh0 ) ;
MPI_File_set_view( fh0, 0, MPI_INT, MPI_INT, "native", MPI_INFO_NULL ) ;
MPI_File_set_atomicity( fh0, TRUE ) ;
MPI_File_write_at(fh0, 0, a, 10, MPI_INT, &status) ;
/* MPI_Barrier( MPI_COMM_WORLD ) ; */

/* Process 1 */
int b[10] ;
int TRUE = 1;
MPI_File_open( MPI_COMM_WORLD, "workfile",
               MPI_MODE_RDWR | MPI_MODE_CREATE, MPI_INFO_NULL, &fh1 ) ;
MPI_File_set_view( fh1, 0, MPI_INT, MPI_INT, "native", MPI_INFO_NULL ) ;
MPI_File_set_atomicity( fh1, TRUE ) ;
/* MPI_Barrier( MPI_COMM_WORLD ) ; */
MPI_File_read_at(fh1, 0, b, 10, MPI_INT, &status) ;

```

A user may guarantee that the write on process 0 precedes the read on process 1 by imposing temporal order with, for example, calls to `MPI_BARRIER`.

*Advice to users.* Routines other than `MPI_BARRIER` may be used to impose temporal order. In the example above, process 0 could use `MPI_SEND` to send a 0 byte message, received by process 1 using `MPI_RECV`. (*End of advice to users.*)

Alternatively, a user can impose consistency with nonatomic mode set:

```

/* Process 0 */
int i, a[10] ;
for ( i=0;i<10;i++)
    a[i] = 5 ;

MPI_File_open( MPI_COMM_WORLD, "workfile",
               MPI_MODE_RDWR | MPI_MODE_CREATE, MPI_INFO_NULL, &fh0 ) ;
MPI_File_set_view( fh0, 0, MPI_INT, MPI_INT, "native", MPI_INFO_NULL ) ;
MPI_File_write_at(fh0, 0, a, 10, MPI_INT, &status) ;
MPI_File_sync( fh0 ) ;
MPI_Barrier( MPI_COMM_WORLD ) ;
MPI_File_sync( fh0 ) ;

/* Process 1 */
int b[10] ;
MPI_File_open( MPI_COMM_WORLD, "workfile",

```

```

1      MPI_MODE_RDWR | MPI_MODE_CREATE, MPI_INFO_NULL, &fh1 ) ;
2  MPI_File_set_view( fh1, 0, MPI_INT, MPI_INT, "native", MPI_INFO_NULL ) ;
3  MPI_File_sync( fh1 ) ;
4  MPI_Barrier( MPI_COMM_WORLD ) ;
5  MPI_File_sync( fh1 ) ;
6  MPI_File_read_at(fh1, 0, b, 10, MPI_INT, &status ) ;

```

The “sync-barrier-sync” construct is required because:

- The barrier ensures that the write on process 0 occurs before the read on process 1.
- The first sync guarantees that the data written by all processes is transferred to the storage device.
- The second sync guarantees that all data which has been transferred to the storage device is visible to all processes. (This does not affect process 0 in this example.)

The following program represents an erroneous attempt to achieve consistency by eliminating the apparently superfluous second “sync” call for each process.

```

18  /* ----- THIS EXAMPLE IS ERRONEOUS ----- */
19  /* Process 0 */
20  int  i, a[10] ;
21  for ( i=0;i<10;i++)
22      a[i] = 5 ;
23
24  MPI_File_open( MPI_COMM_WORLD, "workfile",
25                MPI_MODE_RDWR | MPI_MODE_CREATE, MPI_INFO_NULL, &fh0 ) ;
26  MPI_File_set_view( fh0, 0, MPI_INT, MPI_INT, "native", MPI_INFO_NULL ) ;
27  MPI_File_write_at(fh0, 0, a, 10, MPI_INT, &status ) ;
28  MPI_File_sync( fh0 ) ;
29  MPI_Barrier( MPI_COMM_WORLD ) ;
30
31  /* Process 1 */
32  int  b[10] ;
33  MPI_File_open( MPI_COMM_WORLD, "workfile",
34                MPI_MODE_RDWR | MPI_MODE_CREATE, MPI_INFO_NULL, &fh1 ) ;
35  MPI_File_set_view( fh1, 0, MPI_INT, MPI_INT, "native", MPI_INFO_NULL ) ;
36  MPI_Barrier( MPI_COMM_WORLD ) ;
37  MPI_File_sync( fh1 ) ;
38  MPI_File_read_at(fh1, 0, b, 10, MPI_INT, &status ) ;
39
40  /* ----- THIS EXAMPLE IS ERRONEOUS ----- */

```

The above program also violates the MPI rule against out-of-order collective operations and will deadlock for implementations in which MPI\_FILE\_SYNC blocks.

*Advice to users.* Some implementations may choose to implement MPI\_FILE\_SYNC as a temporally synchronizing function. When using such an implementation, the “sync-barrier-sync” construct above can be replaced by a single “sync.” The results of using such code with an implementation for which MPI\_FILE\_SYNC is not temporally synchronizing is undefined. (*End of advice to users.*)

## Asynchronous I/O

The behavior of asynchronous I/O operations is determined by applying the rules specified above for synchronous I/O operations.

The following examples all access a preexisting file “myfile.” Word 10 in myfile initially contains the integer 2. Each example writes and reads word 10.

First consider the following code fragment:

```
int a = 4, b, TRUE=1;
MPI_File_open( MPI_COMM_WORLD, "myfile",
               MPI_MODE_RDWR, MPI_INFO_NULL, &fh ) ;
MPI_File_set_view( fh, 0, MPI_INT, MPI_INT, "native", MPI_INFO_NULL ) ;
/* MPI_File_set_atomicity( fh, TRUE ) ; Use this to set atomic mode. */
MPI_File_irewrite_at(fh, 10, &a, 1, MPI_INT, &reqs[0]) ;
MPI_File_iread_at(fh, 10, &b, 1, MPI_INT, &reqs[1]) ;
MPI_Waitall(2, reqs, statuses) ;
```

For asynchronous data access operations, MPI specifies that the access occurs at any time between the call to the asynchronous data access routine and the return from the corresponding request complete routine. Thus, executing either the read before the write, or the write before the read is consistent with program order. If atomic mode is set, then MPI guarantees sequential consistency, and the program will read either 2 or 4 into b. If atomic mode is not set, then sequential consistency is not guaranteed and the program may read something other than 2 or 4 due to the conflicting data access.

Similarly, the following code fragment does not order file accesses:

```
int a = 4, b;
MPI_File_open( MPI_COMM_WORLD, "myfile",
               MPI_MODE_RDWR, MPI_INFO_NULL, &fh ) ;
MPI_File_set_view( fh, 0, MPI_INT, MPI_INT, "native", MPI_INFO_NULL ) ;
/* MPI_File_set_atomicity( fh, TRUE ) ; Use this to set atomic mode. */
MPI_File_irewrite_at(fh, 10, &a, 1, MPI_INT, &reqs[0]) ;
MPI_File_iread_at(fh, 10, &b, 1, MPI_INT, &reqs[1]) ;
MPI_Wait(&reqs[0], &status) ;
MPI_Wait(&reqs[1], &status) ;
```

If atomic mode is set, either 2 or 4 will be read into b. Again, MPI does not guarantee sequential consistency in nonatomic mode.

On the other hand, the following code fragment:

```
int a = 4, b;
MPI_File_open( MPI_COMM_WORLD, "myfile",
               MPI_MODE_RDWR, MPI_INFO_NULL, &fh ) ;
MPI_File_set_view( fh, 0, MPI_INT, MPI_INT, "native", MPI_INFO_NULL ) ;
MPI_File_irewrite_at(fh, 10, &a, 1, MPI_INT, &reqs[0]) ;
MPI_Wait(&reqs[0], &status) ;
MPI_File_iread_at(fh, 10, &b, 1, MPI_INT, &reqs[1]) ;
MPI_Wait(&reqs[1], &status) ;
```

defines the same ordering as:

```

1  int a = 4, b;
2  MPI_File_open( MPI_COMM_WORLD, "myfile",
3                MPI_MODE_RDWR, MPI_INFO_NULL, &fh ) ;
4  MPI_File_set_view( fh, 0, MPI_INT, MPI_INT, "native", MPI_INFO_NULL ) ;
5  MPI_File_write_at(fh, 10, &a, 1, MPI_INT, &status ) ;
6  MPI_File_read_at(fh, 10, &b, 1, MPI_INT, &status ) ;

```

Since

- nonconcurrent operations on a single file handle are sequentially consistent, and
- the program fragments specify an order for the operations,

MPI guarantees that both program fragments will read the value 4 into **b**. There is no need to set atomic mode for this example.

Similar considerations apply to conflicting accesses of the form:

```

16 MPI_File_write_all_begin(fh,...) ;
17 MPI_File_iread(fh,...) ;
18 MPI_Wait(fh,...) ;
19 MPI_File_write_all_end(fh,...) ;

```

Recall that constraints governing consistency and semantics are not relevant to the following:

```

24 MPI_File_write_all_begin(fh,...) ;
25 MPI_File_read_all_begin(fh,...) ;
26 MPI_File_read_all_end(fh,...) ;
27 MPI_File_write_all_end(fh,...) ;

```

since split collective operations on the same file handle may not overlap (see Section 13.4.5, page 546).

## 13.7 I/O Error Handling

By default, communication errors are fatal—`MPI_ERRORS_ARE_FATAL` is the default error handler associated with `MPI_COMM_WORLD`. I/O errors are usually less catastrophic (e.g., “file not found”) than communication errors, and common practice is to catch these errors and continue executing. For this reason, MPI provides additional error facilities for I/O.

*Advice to users.* MPI does not specify the state of a computation after an erroneous MPI call has occurred. A high-quality implementation will support the I/O error handling facilities, allowing users to write programs using common practice for I/O. (*End of advice to users.*)

Like communicators, each file handle has an error handler associated with it. The MPI I/O error handling routines are defined in Section 8.3, page 355.

When MPI calls a user-defined error handler resulting from an error on a particular file handle, the first two arguments passed to the file error handler are the file handle and the error code. For I/O errors that are not associated with a valid file handle (e.g., in

MPI\_FILE\_OPEN or MPI\_FILE\_DELETE), the first argument passed to the error handler is MPI\_FILE\_NULL,

I/O error handling differs from communication error handling in another important aspect. By default, the predefined error handler for file handles is MPI\_ERRORS\_RETURN. The default file error handler has two purposes: when a new file handle is created (by MPI\_FILE\_OPEN), the error handler for the new file handle is initially set to the default error handler, and I/O routines that have no valid file handle on which to raise an error (e.g., MPI\_FILE\_OPEN or MPI\_FILE\_DELETE) use the default file error handler. The default file error handler can be changed by specifying MPI\_FILE\_NULL as the fh argument to MPI\_FILE\_SET\_ERRHANDLER. The current value of the default file error handler can be determined by passing MPI\_FILE\_NULL as the fh argument to MPI\_FILE\_GET\_ERRHANDLER.

*Rationale.* For communication, the default error handler is inherited from MPI\_COMM\_WORLD. In I/O, there is no analogous “root” file handle from which default properties can be inherited. Rather than invent a new global file handle, the default file error handler is manipulated as if it were attached to MPI\_FILE\_NULL. (*End of rationale.*)

## 13.8 I/O Error Classes

The implementation dependent error codes returned by the I/O routines can be converted into the error classes defined in Table 13.3.

In addition, calls to routines in this chapter may raise errors in other MPI classes, such as MPI\_ERR\_TYPE.

## 13.9 Examples

### 13.9.1 Double Buffering with Split Collective I/O

This example shows how to overlap computation and output. The computation is performed by the function `compute_buffer()`.

```
/*=====
*
* Function:          double_buffer
*
* Synopsis:
*   void double_buffer(
*       MPI_File fh,                ** IN
*       MPI_Datatype buftype,       ** IN
*       int bufcount                ** IN
*   )
*
* Description:
*   Performs the steps to overlap computation with a collective write
*   by using a double-buffering technique.
*
*/
```

MPI_ERR_FILE	Invalid file handle
MPI_ERR_NOT_SAME	Collective argument not identical on all processes, or collective routines called in a different order by different processes
MPI_ERR_AMODE	Error related to the <code>amode</code> passed to <code>MPI_FILE_OPEN</code>
MPI_ERR_UNSUPPORTED_DATAREP	Unsupported <code>datarep</code> passed to <code>MPI_FILE_SET_VIEW</code>
MPI_ERR_UNSUPPORTED_OPERATION	Unsupported operation, such as seeking on a file which supports sequential access only
MPI_ERR_NO_SUCH_FILE	File does not exist
MPI_ERR_FILE_EXISTS	File exists
MPI_ERR_BAD_FILE	Invalid file name (e.g., path name too long)
MPI_ERR_ACCESS	Permission denied
MPI_ERR_NO_SPACE	Not enough space
MPI_ERR_QUOTA	Quota exceeded
MPI_ERR_READ_ONLY	Read-only file or file system
MPI_ERR_FILE_IN_USE	File operation could not be completed, as the file is currently open by some process
MPI_ERR_DUP_DATAREP	Conversion functions could not be registered because a data representation identifier that was already defined was passed to <code>MPI_REGISTER_DATAREP</code>
MPI_ERR_CONVERSION	An error occurred in a user supplied data conversion function.
MPI_ERR_IO	Other I/O error

Table 13.3: I/O Error Classes



```

* Parameters:
*      fh              previously opened MPI file handle
*      buftype         MPI datatype for memory layout
*                      (Assumes a compatible view has been set on fh)
*      bufcount        # buftype elements to transfer
*-----*/
1
2
3
4
5
6
7
/* this macro switches which buffer "x" is pointing to */
8
#define TOGGLE_PTR(x) (((x)==(buffer1)) ? (x=buffer2) : (x=buffer1))
9
10
void double_buffer( MPI_File fh, MPI_Datatype buftype, int bufcount)
11
{
12
13
    MPI_Status status;          /* status for MPI calls */
14
    float *buffer1, *buffer2;   /* buffers to hold results */
15
    float *compute_buf_ptr;     /* destination buffer */
16
                                /* for computing */
17
    float *write_buf_ptr;       /* source for writing */
18
    int done;                   /* determines when to quit */
19
20
    /* buffer initialization */
21
    buffer1 = (float *)
22
                malloc(bufcount*sizeof(float)) ;
23
    buffer2 = (float *)
24
                malloc(bufcount*sizeof(float)) ;
25
    compute_buf_ptr = buffer1 ; /* initially point to buffer1 */
26
    write_buf_ptr   = buffer1 ; /* initially point to buffer1 */
27
28
29
    /* DOUBLE-BUFFER prolog:
30
    *   compute buffer1; then initiate writing buffer1 to disk
31
    */
32
    compute_buffer(compute_buf_ptr, bufcount, &done);
33
    MPI_File_write_all_begin(fh, write_buf_ptr, bufcount, buftype);
34
35
    /* DOUBLE-BUFFER steady state:
36
    *   Overlap writing old results from buffer pointed to by write_buf_ptr
37
    *   with computing new results into buffer pointed to by compute_buf_ptr.
38
    *
39
    *   There is always one write-buffer and one compute-buffer in use
40
    *   during steady state.
41
    */
42
    while (!done) {
43
        TOGGLE_PTR(compute_buf_ptr);
44
        compute_buffer(compute_buf_ptr, bufcount, &done);
45
        MPI_File_write_all_end(fh, write_buf_ptr, &status);
46
        TOGGLE_PTR(write_buf_ptr);
47
        MPI_File_write_all_begin(fh, write_buf_ptr, bufcount, buftype);
48

```

```

1      }
2
3      /* DOUBLE-BUFFER epilog:
4       *   wait for final write to complete.
5       */
6      MPI_File_write_all_end(fh, write_buf_ptr, &status);
7
8
9      /* buffer cleanup */
10     free(buffer1);
11     free(buffer2);
12 }

```

### 13.9.2 Subarray Filetype Constructor

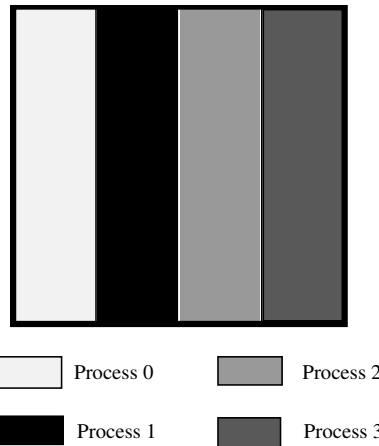


Figure 13.4: Example array file layout

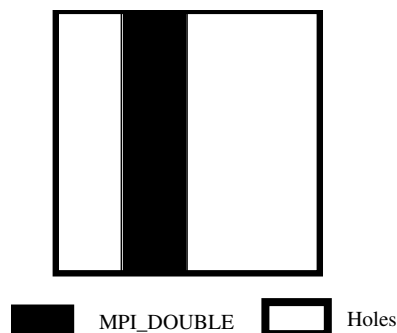


Figure 13.5: Example local array filetype for process 1

Assume we are writing out a 100x100 2D array of double precision floating point numbers that is distributed among 4 processes such that each process has a block of 25 columns (e.g., process 0 has columns 0-24, process 1 has columns 25-49, etc.; see Figure 13.4). To create the filetypes for each process one could use the following C program (see Section 4.1.3 on page 103):

```

double subarray[100][25];
MPI_Datatype filetype;
int sizes[2], subsizes[2], starts[2];
int rank;

MPI_Comm_rank(MPI_COMM_WORLD, &rank);
sizes[0]=100; sizes[1]=100;
subsizes[0]=100; subsizes[1]=25;
starts[0]=0; starts[1]=rank*subsizes[1];

MPI_Type_create_subarray(2, sizes, subsizes, starts, MPI_ORDER_C,
                        MPI_DOUBLE, &filetype);

```

Or, equivalently in Fortran:

```

double precision subarray(100,25)
integer filetype, rank, ierror
integer sizes(2), subsizes(2), starts(2)

call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierror)
sizes(1)=100
sizes(2)=100
subsizes(1)=100
subsizes(2)=25
starts(1)=0
starts(2)=rank*subsizes(2)

call MPI_TYPE_CREATE_SUBARRAY(2, sizes, subsizes, starts, &
                             MPI_ORDER_FORTRAN, MPI_DOUBLE_PRECISION, &
                             filetype, ierror)

```

The generated filetype will then describe the portion of the file contained within the process's subarray with holes for the space taken by the other processes. Figure 13.5 shows the filetype created for process 1.

1  
2  
3  
4  
5  
6  
7  
8  
9  
10  
11  
12  
13  
14  
15  
16  
17  
18  
19  
20  
21  
22  
23  
24  
25  
26  
27  
28  
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31  
32  
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43  
44  
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# Chapter 14

## Tool Support

### 14.1 Introduction

This chapter discusses a set of interfaces that allows debuggers, performance analyzers, and other tools to extract information about the operation of MPI processes. Specifically, this chapter defines both the MPI profiling interface (Section 14.2), which supports the transparent interception and inspection of MPI calls, and the MPI tool information interface (Section 14.3), which supports the inspection and manipulation of MPI control and performance variables. The interfaces described in this chapter are all defined in the context of an MPI process, i.e., are callable from the same code that invokes other MPI functions.

### 14.2 Profiling Interface

[ WAS: Chapter ]

#### 14.2.1 Requirements

[ WAS: Section ]

To meet [the]the requirements for the MPI profiling interface, an implementation of the MPI functions *must*

1. provide a mechanism through which all of the MPI defined [functions]functions, except those allowed as macros (See Section 2.6.5[]), may be accessed with a name shift. This requires, in C and Fortran, an alternate entry point name, with the prefix PMPI\_ for each MPI function in each provided language binding and language support method. The profiling interface in C++ is described in Section 16.1.10. For routines implemented as macros, it is still required that the PMPI\_ version be supplied and work as expected, but it is not possible to replace at link time the MPI\_ version with a user-defined version.

For Fortran, the different support methods cause several linker names. Therefore, several profiling routines (with these linker names) are needed for each Fortran MPI routine, as described in Section 16.2.5 on page 649.

2. ensure that those MPI functions that are not replaced may still be linked into an executable image without causing name clashes.

3. document the implementation of different language bindings of the MPI interface if they are layered on top of each other, so that the profiler developer knows whether she must implement the profile interface for each binding, or can [\[economise\]](#)[economize](#) by implementing it only for the lowest level routines.

4. where the implementation of different language bindings is done through a layered approach ([\[e.g.\]e.g.](#), the Fortran binding is a set of “wrapper” functions that call the C implementation), ensure that these wrapper functions are separable from the rest of the library.

This separability is necessary to allow a separate profiling library to be correctly implemented, since (at least with Unix linker semantics) the profiling library must contain these wrapper functions if it is to perform as expected. This requirement allows the person who builds the profiling library to extract these functions from the original MPI library and add them into the profiling library without bringing along any other unnecessary code.

5. provide a no-op routine `MPI_PCONTROL` in the MPI library.

### 14.2.2 Discussion

[\[ WAS: Section \]](#)

The objective of the MPI profiling interface is to ensure that it is relatively easy for authors of profiling (and other similar) tools to interface their codes to MPI implementations on different machines.

Since MPI is a machine independent standard with many different implementations, it is unreasonable to expect that the authors of profiling tools for MPI will have access to the source code that implements MPI on any particular machine. It is therefore necessary to provide a mechanism by which the implementors of such tools can collect whatever performance information they wish *without* access to the underlying implementation.

We believe that having such an interface is important if MPI is to be attractive to end users, since the availability of many different tools will be a significant factor in attracting users to the MPI standard.

The profiling interface is just that, an interface. It says *nothing* about the way in which it is used. There is therefore no attempt to lay down what information is collected through the interface, or how the collected information is saved, filtered, or displayed.

While the initial impetus for the development of this interface arose from the desire to permit the implementation of profiling tools, it is clear that an interface like that specified may also prove useful for other purposes, such as “internetworking” multiple MPI implementations. Since all that is defined is an interface, there is no objection to its being used wherever it is useful.

As the issues being addressed here are intimately tied up with the way in which executable images are built, which may differ greatly on different machines, the examples given below should be treated solely as one way of implementing the objective of the MPI profiling interface. The actual requirements made of an implementation are those detailed in the Requirements section above, the whole of the rest of this chapter is only present as justification and discussion of the logic for those requirements.

The examples below show one way in which an implementation could be constructed to meet the requirements on a Unix system (there are doubtless others that would be equally valid).

### 14.2.3 Logic of the Design

[ WAS: Section ]

Provided that an MPI implementation meets the requirements above, it is possible for the implementor of the profiling system to intercept all of the MPI calls that are made by the user program. She can then collect whatever information she requires before calling the underlying MPI implementation (through its name shifted entry points) to achieve the desired effects.

### 14.2.4 Miscellaneous Control of Profiling

[ WAS: Subsection, Now still a subsection to remove single subsubsection ]

There is a clear requirement for the user code to be able to control the profiler dynamically at run time. This is normally used for (at least) the purposes of

- Enabling and disabling profiling depending on the state of the calculation.
- Flushing trace buffers at non-critical points in the [calculation]calculation.
- Adding user events to a trace file.

These requirements are met by use of the MPI\_PCONTROL.

MPI\_PCONTROL(level, ...)

IN	level	Profiling level
----	-------	-----------------

int MPI\_Pcontrol(const int level, ...)

MPI\_Pcontrol(level) BIND(C)  
 INTEGER, INTENT(IN) :: level

MPI\_PCONTROL(LEVEL)  
 INTEGER LEVEL

{void MPI::Pcontrol(const int level, ...) (binding deprecated, see Section 15.2) }

MPI libraries themselves make no use of this routine, and simply return immediately to the user code. However the presence of calls to this routine allows a profiling package to be explicitly called by the user.

Since MPI has no control of the implementation of the profiling code, we are unable to specify precisely the semantics that will be provided by calls to MPI\_PCONTROL. This vagueness extends to the number of arguments to the function, and their datatypes.

However to provide some level of portability of user codes to different profiling libraries, we request the following meanings for certain values of level.

- level==0 Profiling is disabled.



- `level==1` Profiling is enabled at a normal default level of detail.
- `level==2` Profile buffers are `[flushed. (This may be a no-op in some profilers).]flushed,`  
`which may be a no-op in some profilers.`
- All other values of `level` have profile library defined effects and additional arguments.

We also request that the default state after `MPI_INIT` has been called is for profiling to be enabled at the normal default level. (i.e. as if `MPI_PCONTROL` had just been called with the argument 1). This allows users to link with a profiling library and obtain profile output without having to modify their source code at all.

The provision of `MPI_PCONTROL` as a no-op in the standard MPI library `[allows them to modify their source code to obtain]``supports the collection of` more detailed profiling information`[, but still be able to link exactly the]``with source [same code]``code that can still link` against the standard MPI library.

`[WAS: Subsection Examples]`

#### 14.2.5 Profiler Implementation `[Example]`

`[Suppose that the profiler wishes to]``A profiler can` accumulate the total amount of data sent by the `[MPI_SEND]``MPI_SEND` function, along with the total elapsed time spent in the `[function. This could trivially be achieved thus]``function, as follows:`

```
static int totalBytes = 0;
static double totalTime = 0.0;

int MPI_Send(void* buffer, int count, MPI_Datatype datatype,
             int dest, int tag, MPI_Comm comm)
{
    double tstart = MPI_Wtime();          /* Pass on all the arguments */
    int extent;
    int result    = PMPI_Send(buffer, count, datatype, dest, tag, comm);

    MPI_Type_size(datatype, &extent); /* Compute size */
    totalBytes += count*extent;

    totalTime += MPI_Wtime() - tstart;    /* and time */

    return result;
}
```

#### 14.2.6 MPI Library Implementation `[Example]`

`[On a Unix system, in which the MPI library is implemented in C, then]``If the MPI library is implemented in C on a Unix system, then there [there are various possible options, of which two of the most obvious]``are various options, including the two presented here, for supporting [are presented here. Which is better depends on whether the linker and]``the name-shift requirement. The choice between these two options [compiler support weak symbols.]depends partly on whether the linker and compiler support weak symbols.`

## Systems with Weak Symbols

If the compiler and linker support weak external symbols ([e.g.]e.g., Solaris 2.x, other system V.4 machines), then only a single library is required through the use of `#pragma weak` thus

```
#pragma weak MPI_Example = PMPI_Example

int PMPI_Example(/* appropriate args */)
{
    /* Useful content */
}
```

The effect of this `#pragma` is to define the external symbol `MPI_Example` as a weak definition. This means that the linker will not complain if there is another definition of the symbol (for instance in the profiling library), however if no other definition exists, then the linker will use the weak definition.

## Systems Without Weak Symbols

In the absence of weak symbols then one possible solution would be to use the C macro pre-processor thus

```
#ifdef PROFILELIB
#   ifdef __STDC__
#       define FUNCTION(name) P##name
#   else
#       define FUNCTION(name) P/**/name
#   endif
#else
#   define FUNCTION(name) name
#endif
```

Each of the user visible functions in the library would then be declared thus

```
int FUNCTION(MPI_Example)(/* appropriate args */)
{
    /* Useful content */
}
```

The same source file can then be compiled to produce both versions of the library, depending on the state of the `PROFILELIB` macro symbol.

It is required that the standard MPI library be built in such a way that the inclusion of MPI functions can be achieved one at a time. This is a somewhat unpleasant requirement, since it may mean that each external function has to be compiled from a separate file. However this is necessary so that the author of the profiling library need only define those MPI functions that she wishes to intercept, references to any others being fulfilled by the normal MPI library. Therefore the link step can look something like this

```
% cc ... -lmyprof -lpmpl -lmpi
```

Here `libmyprof.a` contains the profiler functions that intercept some of the MPI functions[,], `libpmpl.a` contains the “name shifted” MPI functions, and `libmpi.a` contains the normal definitions of the MPI functions.

### 14.2.7 Complications

#### Multiple Counting

Since parts of the MPI library may themselves be implemented using more basic MPI functions ([e.g.]e.g., a portable implementation of the collective operations implemented using point to point communications), there is potential for profiling functions to be called from within an MPI function that was called from a profiling function. This could lead to “double counting” of the time spent in the inner routine. Since this effect could actually be useful under some circumstances ([e.g.]e.g., it might allow one to answer the question “How much time is spent in the point to point routines when they’re called from collective functions?”), we have decided not to enforce any restrictions on the author of the MPI library that would overcome this. Therefore the author of the profiling library should be aware of this problem, and guard against it herself. In a single threaded world this is easily achieved through use of a static variable in the profiling code that remembers if you are already inside a profiling routine. It becomes more complex in a multi-threaded environment (as does the meaning of the times recorded[!]).

#### Linker Oddities

The Unix linker traditionally operates in one [pass :]pass: the effect of this is that functions from libraries are only included in the image if they are needed at the time the library is scanned. When combined with weak symbols, or multiple definitions of the same function, this can cause odd (and unexpected) effects.

Consider, for instance, an implementation of MPI in which the Fortran binding is achieved by using wrapper functions on top of the C implementation. The author of the profile library then assumes that it is reasonable only to provide profile functions for the C binding, since Fortran will eventually call these, and the cost of the wrappers is assumed to be small. However, if the wrapper functions are not in the profiling library, then none of the profiled entry points will be undefined when the profiling library is called. Therefore none of the profiling code will be included in the image. When the standard MPI library is scanned, the Fortran wrappers will be resolved, and will also pull in the base versions of the MPI functions. The overall effect is that the code will link successfully, but will not be profiled.

To overcome this we must ensure that the Fortran wrapper functions are included in the profiling version of the library. We ensure that this is possible by requiring that these be separable from the rest of the base MPI library. This allows them to be aared out of the base library and into the profiling one.

#### Fortran Support Methods

The different Fortran support methods and possible options for the support of subarrays (depending on whether the compiler can support TYPE(\*), DIMENSION(..) choice buffers) imply different linker names for the same Fortran MPI routine. The rules and implications for the profiling interface are described in Section 16.2.5 on page 649.

### 14.2.8 Multiple Levels of Interception

[ WAS: Section ] The scheme given here does not directly support the nesting of profiling functions, since it provides only a single alternative name for each MPI function. Consideration was given to an implementation that would allow multiple levels of call interception, however we were unable to construct an implementation of this that did not have the following disadvantages

- assuming a particular implementation language[.],
- imposing a run time cost even when no profiling was taking place.

Since one of the objectives of MPI is to permit efficient, low latency implementations, and it is not the business of a standard to require a particular implementation language, we decided to accept the scheme outlined above.

[Note, however, that it is possible to use the scheme above to implement a multi-level system, since the function called by the user may call many different profiling functions before calling the underlying MPI function.]

[Unfortunately such an implementation may require more cooperation between the different profiling libraries than is required for the single level implementation detailed above.]Note, however, that it is possible to use the scheme above to implement a multi-level system, since the function called by the user may call many different profiling functions before calling the underlying MPI function. This capability has been demonstrated in the P<sup>N</sup>MPI tool infrastructure [51].

### 14.3 The MPI Tool Information Interface

MPI implementations often use internal variables to control their operation and performance. Understanding and manipulating these variables can provide a more efficient execution environment or improve performance for many applications. This section describes the MPI tool information interface, which provides a mechanism for MPI implementors to expose a set of variables, each of which represents a particular property, setting, or performance measurement from within the MPI implementation. The interface is split into two parts: the first part provides information about and supports the setting of control variables through which the MPI implementation tunes its configuration. The second part provides access to performance variables that can provide insight into internal performance information of the MPI implementation.

To avoid restrictions on the MPI implementation, the MPI tool information interface allows the implementation to specify which control and performance variables exist. Additionally, the user of the MPI tool information interface can obtain metadata about each available variable, such as its datatype, and a textual description. The MPI tool information interface provides the necessary routines to find all variables that exist in a particular MPI implementation, to query their properties, to retrieve descriptions about their meaning, and to access and, if appropriate, to alter their values.

The MPI tool information interface can be used independently from the MPI communication functionality. In particular, the routines of this interface can be called before MPI\_INIT (or equivalent) and after MPI\_FINALIZE. In order to support this behavior cleanly, the MPI tool information interface uses separate initialization and finalization routines. All identifiers used in the MPI tool information interface have the prefix MPI\_T\_.

On success, all MPI tool information interface routines return MPI\_SUCCESS, otherwise they return an appropriate and unique return code indicating the reason why the call was

not successfully completed. Details on return codes can be found in Section 14.3.9. However, unsuccessful calls to the MPI tool information interface are not fatal and do not impact the execution of subsequent MPI routines.

Since the MPI tool information interface primarily focuses on tools and support libraries, MPI implementations are only required to provide C bindings for functions introduced in this Section 14.3. Except where otherwise noted, all conventions and principles governing the C bindings of the MPI API also apply to the MPI tool information interface, which is available by including the `mpi.h` header file. All routines in this interface have local semantics.

*Advice to users.* The number and type of control variables and performance variables can vary between MPI implementations, platforms and different builds of the same implementation on the same platform as well as between runs. Hence, any application relying on a particular variable will not be portable. Further, there is no guarantee that number of variables, variable indices, and variable names are the same across processes.

This interface is primarily intended for performance monitoring tools, support tools, and libraries controlling the application's environment. When maximum portability is desired, application programmers should either avoid using the MPI tool information interface or avoid being dependent on the existence of a particular control or performance variable. (*End of advice to users.*)

### 14.3.1 Verbosity Levels

The MPI tool information interface provides access to internal configuration and performance information through a set of control and performance variables defined by the MPI implementation. Since some implementations may export a large number of variables, variables are classified by a verbosity level that categorizes both their intended audience (end users, performance tuners or MPI implementors) and a relative measure of level of detail (basic, detailed or all). These verbosity levels are described by a single integer. Table 14.1 lists the constants for all possible verbosity levels. The values of the constants are monotonic in the order listed in the table; i.e., `MPI_T_VERBOSITY_USER_BASIC` < `MPI_T_VERBOSITY_USER_DETAIL` < ... < `MPI_T_VERBOSITY_MPIDEV_ALL`.

<code>MPI_T_VERBOSITY_USER_BASIC</code>	Basic information of interest to users
<code>MPI_T_VERBOSITY_USER_DETAIL</code>	Detailed information of interest to users
<code>MPI_T_VERBOSITY_USER_ALL</code>	All information of interest to users
<code>MPI_T_VERBOSITY_TUNER_BASIC</code>	Basic information required for tuning
<code>MPI_T_VERBOSITY_TUNER_DETAIL</code>	Detailed information required for tuning
<code>MPI_T_VERBOSITY_TUNER_ALL</code>	All information required for tuning
<code>MPI_T_VERBOSITY_MPIDEV_BASIC</code>	Basic information for MPI implementors
<code>MPI_T_VERBOSITY_MPIDEV_DETAIL</code>	Detailed information for MPI implementors
<code>MPI_T_VERBOSITY_MPIDEV_ALL</code>	All information for MPI implementors

Table 14.1: MPI tool information interface verbosity levels.

### 14.3.2 Binding MPI Tool Information Interface Variables to MPI Objects

Each MPI tool information interface variable provides access to a particular control setting or performance property of the MPI implementation. A variable may refer to a specific MPI object such as a communicator, datatype, or one-sided communication window, or the variable may refer more generally to the MPI environment of the process. Except for the last case, the variable must be bound to exactly one MPI object before it can be used. Table 14.2 lists all MPI object types to which an MPI tool information interface variable can be bound, together with the matching constant that MPI tool information interface routines return to identify the object type.

Constant	MPI object
MPI_T_BIND_NO_OBJECT	N/A; applies globally to entire MPI process
MPI_T_BIND_MPI_COMM	MPI communicators
MPI_T_BIND_MPI_DATATYPE	MPI datatypes
MPI_T_BIND_MPI_ERRHANDLER	MPI error handlers
MPI_T_BIND_MPI_FILE	MPI file handles
MPI_T_BIND_MPI_GROUP	MPI groups
MPI_T_BIND_MPI_OP	MPI reduction operators
MPI_T_BIND_MPI_REQUEST	MPI requests
MPI_T_BIND_MPI_WIN	MPI windows for one-sided communication
MPI_T_BIND_MPI_MESSAGE	MPI message object
MPI_T_BIND_MPI_INFO	MPI info object

Table 14.2: Constants to identify associations of variables.

*Rationale.* Some variables have meanings tied to a specific MPI object. Examples include the number of send or receive operations using a particular datatype, the number of times a particular error handler has been called, or the communication protocol and “eager limit” used for a particular communicator. Creating a new MPI tool information interface variable for each MPI object would cause the number of variables to grow without bounds, since they cannot be reused to avoid naming conflicts. By associating MPI tool information interface variables with a specific MPI object, the MPI implementation only must specify and maintain a single variable, which can then be applied to as many MPI objects of the respective type as created during the program’s execution. (*End of rationale.*)

### 14.3.3 Convention for Returning Strings

Several MPI tool information interface functions return one or more strings. These functions have two arguments for each string to be returned: an OUT parameter that identifies a pointer to the buffer in which the string will be returned, and an IN/OUT parameter to pass the length of the buffer. The user is responsible for the memory allocation of the buffer and must pass the size of the buffer ( $n$ ) as the length argument. Let  $n$  be the length value specified to the function. On return, the function writes at most  $n - 1$  of the string’s characters into the buffer, followed by a null terminator. If the returned string’s length is greater than or equal to  $n$ , the string will be truncated to  $n - 1$  characters. In this case, the length of the string plus one (for the terminating null character) is returned in the length

argument. If the user passes the null pointer as the buffer argument or passes 0 as the length argument, the function does not return the string and only returns the length of the string plus one in the length argument. If the user passes the null pointer as the length argument, the buffer argument is ignored and nothing is returned.

#### 14.3.4 Initialization and Finalization

The MPI tool information interface requires a separate set of initialization and finalization routines.

**MPI\_T\_INIT\_THREAD(required, provided)**

IN	required	desired level of thread support (integer)
OUT	provided	provided level of thread support (integer)

**int MPI\_T\_init\_thread(int required, int \*provided)**

All programs or tools that use the MPI tool information interface must initialize the MPI tool information interface in the processes that will use the interface before calling any other of its routines. A user can initialize the MPI tool information interface by calling **MPI\_T\_INIT\_THREAD**, which can be called multiple times. In addition, this routine initializes the thread environment for all routines in the MPI tool information interface. Calling this routine when the MPI tool information interface is already initialized has no effect beyond increasing the reference count of how often the interface has been initialized. The argument **required** is used to specify the desired level of thread support. The possible values and their semantics are identical to the ones that can be used with **MPI\_INIT\_THREAD** listed in Section 12.4. The call returns in **provided** information about the actual level of thread support that will be provided by the MPI implementation for calls to MPI tool information interface routines. It can be one of the four values listed in Section 12.4.

The MPI specification does not require all MPI processes to exist before the call to **MPI\_INIT**. If the MPI tool information interface is used before **MPI\_INIT** has been called, **MPI\_T\_INIT\_THREAD** must be called on each process that will use the MPI tool information interface. Processes created by the MPI implementation during **MPI\_INIT** inherit the status of the MPI tool information interface (whether it is initialized or not as well as all active sessions and handles) from the process from which they are created.

Processes created at runtime as a result of calls MPI's dynamic process management require their own initialization before they can use the MPI tool information interface.

*Advice to users.* If **MPI\_T\_INIT\_THREAD** is called before **MPI\_INIT\_THREAD**, the requested and granted thread level for **MPI\_T\_INIT\_THREAD** may influence the behavior and return value of **MPI\_INIT\_THREAD**. The same is true for the reverse order. (*End of advice to users.*)

*Advice to implementors.* MPI implementations should strive to make as many control or performance variables available before **MPI\_INIT** (instead of adding them within **MPI\_INIT**) to allow tools the most flexibility. In particular, control variables should be available before **MPI\_INIT** if their value cannot be changed after **MPI\_INIT**. (*End of advice to implementors.*)



```
1 MPI_T_FINALIZE( )
```

```
2
3 int MPI_T_finalize(void)
```

4 This routine finalizes the use of the MPI tool information interface and may be called  
5 as often as the corresponding MPI\_T\_INIT\_THREAD routine up to the current point of  
6 execution. Calling it more times returns a corresponding error code. As long as the number  
7 of calls to MPI\_T\_FINALIZE is smaller than the number of calls to MPI\_T\_INIT\_THREAD  
8 up to the current point of execution, the MPI tool information interface remains initialized  
9 and calls to its routines are permissible. Further, additional calls to MPI\_T\_INIT\_THREAD  
10 after one or more calls to MPI\_T\_FINALIZE are permissible.

11 Once MPI\_T\_FINALIZE is called the same number of times as the routine  
12 MPI\_T\_INIT\_THREAD up to the current point of execution, the MPI tool information in-  
13 terface is no longer initialized. The interface can be reinitialized by subsequent calls to  
14 MPI\_T\_INIT\_THREAD.

15 At the end of the program execution, unless MPI\_ABORT is called, an application must  
16 have called MPI\_T\_INIT\_THREAD and MPI\_T\_FINALIZE an equal number of times.  
17

### 18 14.3.5 Datatype System

19 All variables managed through the MPI tool information interface represent their values  
20 through typed buffers of a given length and type using an MPI datatype (similar to regular  
21 send/receive buffers). Since the initialization of the MPI tool information interface is sep-  
22 arate from the initialization of MPI, MPI tool information interface routines can be called  
23 before MPI\_INIT. Consequently, these routines can also use MPI datatypes before MPI\_INIT.  
24 Therefore, within the context of the MPI tool information interface, it is permissible to use  
25 a subset of MPI datatypes as specified below before a call to MPI\_INIT (or equivalent).  
26

```
27
28 MPI_INT
29 MPI_UNSIGNED
30 MPI_UNSIGNED_LONG
31 MPI_UNSIGNED_LONG_LONG
32 MPI_COUNT
33 MPI_CHAR
34 MPI_DOUBLE
```

35  
36 Table 14.3: MPI datatypes that can be used by the MPI tool information interface.  
37

38  
39 *Rationale.* The MPI tool information interface relies mainly on unsigned datatypes  
40 for integer values since most variables are expected to represent counters or resource  
41 sizes. MPI\_INT is provided for additional flexibility and is expected to be used mainly  
42 for control variables and enumeration types (see below).  
43

44 Providing all basic datatypes, in particular providing all signed and unsigned variants  
45 of integer types, would lead to a larger number of types, which tools need to interpret.  
46 This would cause unnecessary complexity in the implementation of tools based on the  
47 MPI tool information interface. (*End of rationale.*)  
48



The MPI tool information interface only relies on a subset of the basic MPI datatypes and does not use any derived MPI datatypes. Table 14.3 lists all MPI datatypes that can be returned by the MPI tool information interface to represent its variables.

*Rationale.* The MPI tool information interface requires a significantly simpler type system than MPI itself. Therefore, only its required subset must be present before MPI\_INIT (or equivalent) and MPI implementations do not need to initialize the complete MPI datatype system. (*End of rationale.*)

For variables of type MPI\_INT, an MPI implementation can provide additional information by associating names with a fixed number of values. We refer to this information in the following as an enumeration. In this case, the respective calls that provide additional metadata for each control or performance variable, i.e., MPI\_T\_CVAR\_GET\_INFO (Section 14.3.6) and MPI\_T\_PVAR\_GET\_INFO (Section 14.3.7), return a handle of type MPI\_T\_enum that can be passed to the following functions to extract additional information. Thus, the MPI implementation can describe variables with a fixed set of values that each represents a particular state. Each enumeration type can have  $N$  different values, with a fixed  $N$  that can be queried using MPI\_T\_ENUM\_GET\_INFO.

MPI\_T\_ENUM\_GET\_INFO(enumtype, num, name, name\_len)

IN	enumtype	enumeration to be queried (handle)
OUT	num	number of discrete values represented by this enumeration (integer)
OUT	name	buffer to return the string containing the name of the enumeration (string)
INOUT	name_len	length of the string and/or buffer for name (integer)

```
int MPI_T_enum_get_info(MPI_T_enum enumtype, int *num, char *name, int
                        *name_len)
```

If enumtype is a valid enumeration, this routine returns the number of items represented by this enumeration type. range and the name of the enumeration.  $N$  must be greater than 0, i.e., the enumeration must represent at least one value.

The arguments name and name\_len are used to return the name of the enumerations as described in Section 14.3.3.

The routine is required to return a name of at least length one. This name must be unique with respect to all other names for enumerations that the MPI implementation uses.

Names associated with individual values in each enumeration enumtype can be queried using MPI\_T\_ENUM\_GET\_ITEM.

```
1 MPI_T_ENUM_GET_ITEM(enumtype, index, value, name, name_len)
```

2	IN	enumtype	enumeration to be queried (handle)
3	IN	index	number of the value to be queried in this enumeration
4			(integer)
5			
6	OUT	value	variable value (integer)
7	OUT	name	buffer to return the string containing the name of the
8			enumeration item (string)
9			
10	INOUT	name_len	length of the string and/or buffer for name (integer)

```
11
12 int MPI_T_enum_get_item(MPI_T_enum enumtype, int index, int value, char
13                        *name, int *name_len)
```

14 The arguments `name` and `name_len` are used to return the name of the enumeration  
15 item as described in Section 14.3.3.

16 If completed successfully, the routine returns the name/value pair describing the enu-  
17 meration at the specified index. The call is further required to return a name of at least  
18 length one. This name must be unique with respect to all other names of items for the same  
19 enumeration.  
20

### 21 14.3.6 Control Variables

22 The routines described in this section of the MPI tool information interface specification  
23 focus on the ability to list, query, and possibly set control variables exposed by the MPI  
24 implementation. These variables can typically be used by the user to fine tune properties  
25 and configuration settings of the MPI implementation. On many systems, such variables  
26 can be set using environment variables, although other configuration mechanisms may be  
27 available, such as configuration files or central configuration registries. A typical example  
28 that is available in several existing MPI implementations is the ability to specify an “eager  
29 limit”, i.e., an upper bound on the size of messages sent or received using an eager protocol.  
30  
31

#### 32 Control Variable Query Functions

33 An MPI implementation exports a set of  $N$  control variables through the MPI tool infor-  
34 mation interface. If  $N$  is zero, then the MPI implementation does not export any control  
35 variables, otherwise the provided control variables are indexed from 0 to  $N - 1$ . This index  
36 number is used in subsequent calls to identify the individual variables.  
37

38 An MPI implementation is allowed to increase the number of control variables during  
39 the execution of an MPI application when new variables become available through dynamic  
40 loading. However, MPI implementations are not allowed to change the index of a control  
41 variable or delete a variable once it has been added to the set. When variables become  
42 inactive, e.g., through dynamic unloading, accessing its value should return a corresponding  
43 error code.  
44

45 *Advice to users.* While the MPI tool information interface guarantees that indices or  
46 variable properties do not change during a particular run of an MPI program, it does  
47 not provide a similar guarantee between runs. (*End of advice to users.*)  
48

The following function can be used to query the number of control variables, *num\_cvar*:

`MPI_T_CVAR_GET_NUM(num_cvar)`

OUT      *num\_cvar*                      returns number of control variables (integer)

`int MPI_T_cvar_get_num(int *num_cvar)`

The function `MPI_T_CVAR_GET_INFO` provides access to additional information for each variable.

`MPI_T_CVAR_GET_INFO(cvar_index, name, name_len, verbosity, datatype, enumtype, desc, desc_len, bind, scope)`

IN	<i>cvar_index</i>	index of the control variable to be queried, value between 0 and <i>num_cvar</i> - 1 (integer)
OUT	<i>name</i>	buffer to return the string containing the name of the control variable (string)
INOUT	<i>name_len</i>	length of the string and/or buffer for <i>name</i> (integer)
OUT	<i>verbosity</i>	verbosity level of this variable (integer)
OUT	<i>datatype</i>	MPI datatype of the information stored in the control variable (handle)
OUT	<i>enumtype</i>	optional descriptor for enumeration information (handle)
OUT	<i>desc</i>	buffer to return the string containing a description of the control variable (string)
INOUT	<i>desc_len</i>	length of the string and/or buffer for <i>desc</i> (integer)
OUT	<i>bind</i>	type of MPI object to which this variable must be bound (integer)
OUT	<i>scope</i>	scope of when changes to this variable are possible (integer)

`int MPI_T_cvar_get_info(int cvar_index, char *name, int *name_len, int *verbosity, MPI_Datatype *datatype, MPI_T_enum *enumtype, char *desc, int *desc_len, int *bind, int *scope)`

After a successful call to `MPI_T_CVAR_GET_INFO` for a particular variable, subsequent calls to this routine querying information about the same variable must return the same information. An MPI implementation is not allowed to alter any of the returned values.

The arguments *name* and *name\_len* are used to return the name of the control variable as described in Section 14.3.3.

If completed successfully, the routine is required to return a name of at least length one. The name must be unique with respect to all other names for control variables used by the MPI implementation.

The argument *verbosity* returns the verbosity level of the variable (see Section 14.3.1).

The argument `datatype` returns the MPI datatype that is used to represent the control variable.

If the variable is of type `MPI_INT`, MPI can optionally specify an enumeration for the values represented by this variable and return it in `enumtype`. In this case, MPI returns an enumeration identifier, which can then be used to gather more information as described in Section 14.3.5. If the datatype is not `MPI_INT` or the argument `enumtype` is the constant `MPI_T_ENUM_NULL`, no enumeration type is returned.

The arguments `desc` and `desc_len` are used to return a description of the control variable as described in Section 14.3.3.

Returning a description is optional. If an MPI implementation decides not to return a description, the first character for `desc` must be set to the null character and `desc_len` must be set to one at the return of this call.

The parameter `bind` returns the type of the MPI object to which the variable must be bound or the value `MPI_T_BIND_NO_OBJECT` (see Section 14.3.2).

The scope of a variable determines whether changing a variable's value is either local to the process or must be done by the user across multiple processes. The latter is further split into variables that require changes in a group of processes and those that require collective changes among all connected processes. Both cases can require all processes to either be set to consistent (but potentially different) values or to equal values on every participating process. The description provided with the variable must contain an explanation about the requirements and/or restrictions for setting the particular variable.

On successful return from `MPI_T_CVAR_GET_INFO`, the argument `scope` will be set to one of the constants listed in Table 14.4.

Scope Constant	Description
<code>MPI_T_SCOPE_READONLY</code>	read-only, cannot be written
<code>MPI_T_SCOPE_LOCAL</code>	may be writeable, writing is a local operation
<code>MPI_T_SCOPE_GROUP</code>	may be writeable, must be done to a group of processes, all processes in a group must be set to consistent values
<code>MPI_T_SCOPE_GROUP_EQ</code>	may be writeable, must be done to a group of processes, all processes in a group must be set to the same value
<code>MPI_T_SCOPE_ALL</code>	may be writeable, must be done to all processes, all connected processes must be set to consistent values
<code>MPI_T_SCOPE_ALL_EQ</code>	may be writeable, must be done to all processes, all connected processes must be set to the same value

Table 14.4: Scopes for control variables.

*Advice to users.* The `scope` of a variable only indicates if a variable might be changeable; it is not a guarantee that it can be changed at any time. (*End of advice to users.*)

#### Example: Printing All Control Variables

#### Example 14.1

The following example shows how the MPI tool information interface can be used to query and print the names of all available control variables.

```

#include <stdio.h>
#include <stdlibh.h>
#include <mpi.h>

int main(int argc, char **argv) {
    int i, err, num, namelen, bind, verbose, scope;
    int threadsupport;
    char name[100];
    MPI_Datatype datatype;

    err=MPI_T_init_thread(MPI_THREAD_SINGLE,&threadsupport);
    if (err!=MPI_SUCCESS)
        return err;

    err=MPI_T_cvar_get_num(&num);
    if (err!=MPI_SUCCESS)
        return err;

    for (i=0; i<num; i++) {
        namelen=100;
        err=MPI_T_cvar_get_info(i, name, &namelen,
                                &verbose, &datatype, MPI_T_ENUM_NULL,
                                NULL, NULL, /*no description */
                                &bind, &scope);
        if (err!=MPI_SUCCESS) return err;
        printf("Var %i: %s\n", i, name);
    }

    err=MPI_T_finalize();
    if (err!=MPI_SUCCESS)
        return 1;
    else
        return 0;
}

```

#### Handle Allocation and Deallocation

Before reading or writing the value of a variable, a user must first allocate a handle of type `MPI_T_cvar_handle` for the variable by binding it to an MPI object (see also Section 14.3.2).

*Rationale.* Handles used in the MPI tool information interface are distinct from handles used in the remaining parts of the MPI standard because they must be usable before `MPI_INIT` and after `MPI_FINALIZE`. Further, accessing handles, in particular for performance variables, can be time critical and having a separate handle space enables optimizations. (*End of rationale.*)

```
1 MPI_T_CVAR_HANDLE_ALLOC(cvar_index, object, handle, count)
```

2	IN	cvar_index	index of control variable for which handle is to be allocated (index)
3			
4			
5	IN	obj_handle	reference to a handle of the MPI object to which this variable is supposed to be bound (pointer)
6			
7	OUT	handle	allocated handle (handle)
8	OUT	count	number of elements used to represent this variable (integer)
9			
10			

```
11
12 int MPI_T_cvar_handle_alloc(int cvar_index, void *obj_handle,
13 MPI_T_cvar_handle *handle, int *count)
```

14 This routine binds the control variable specified by the argument `index` to an MPI  
 15 object. The object is passed in the argument `obj_handle` as an address to a local variable  
 16 that stores the object's handle. The handle allocated to reference the variable is returned  
 17 in the argument `handle`. Upon successful return, `count` contains the number of elements (of  
 18 the datatype returned by a previous `MPI_T_CVAR_GET_INFO` call) used to represent this  
 19 variable.

20  
 21 *Advice to users.* The `count` can be different based on the MPI object to which it  
 22 was bound. For example, variables bound to communicators could have a count that  
 23 matches the size of the communicator.

24 It is not portable to pass references to predefined MPI object handles, such as  
 25 `MPI_COMM_WORLD` to this routine, since their implementation depends on the MPI  
 26 library. Instead, such object handles should be stored in a local variable and the  
 27 address of this local variables should be passed into `MPI_T_CVAR_HANDLE_ALLOC`.  
 28 (*End of advice to users.*)

29  
 30 The value of `cvar_index` should be in the range 0 to `num_cvar - 1`, where `num_cvar`  
 31 is the number of available control variables as determined from a prior call to  
 32 `MPI_T_CVAR_GET_NUM`. The type of the MPI object it references must be consistent  
 33 with the type returned in the `bind` argument in a prior call to `MPI_T_CVAR_GET_INFO`.

34 In the case the `bind` argument returned by `MPI_T_CVAR_GET_INFO` equals  
 35 `MPI_T_BIND_NO_OBJECT`, the argument `obj_handle` is ignored.

```
36
37
38 MPI_T_CVAR_HANDLE_FREE(handle)
```

39	INOUT	handle	handle to be freed (handle)
40			

```
41
42 int MPI_T_cvar_handle_free(MPI_T_cvar_handle *handle)
```

43 When a handle is no longer needed, a user of the MPI tool information interface should  
 44 call `MPI_T_CVAR_HANDLE_FREE` to free the handle and the associated resources in the  
 45 MPI implementation. On a successful return, MPI sets the handle to  
 46 `MPI_T_CVAR_HANDLE_NULL`.

## Control Variable Access Functions

## MPI\_T\_CVAR\_READ(handle, buf)

IN	handle	handle to the control variable to be read (handle)
OUT	buf	initial address of storage location for variable value (choice)

```
int MPI_T_cvar_read(MPI_T_cvar_handle handle, void* buf)
```

This routine queries the value of the control variable identified by the argument `handle` and stores the result in the buffer identified by the parameter `buf`. The user must ensure that the buffer is of the appropriate size to hold the entire value of the control variable (based on the returned datatype and count from prior corresponding calls to `MPI_T_CVAR_GET_INFO` and `MPI_T_CVAR_HANDLE_ALLOC`, respectively).

## MPI\_T\_CVAR\_WRITE(handle, buf)

IN	handle	handle to the control variable to be written (handle)
IN	buf	initial address of storage location for variable value (choice)

```
int MPI_T_cvar_write(MPI_T_cvar_handle handle, const void* buf)
```

This routine sets the value of the control variable identified by the argument `handle` to the data stored in the buffer identified by the parameter `buf`. The user must ensure that the buffer is of the appropriate size to hold the entire value of the control variable (based on the returned datatype and count from prior corresponding calls to `MPI_T_CVAR_GET_INFO` and `MPI_T_CVAR_HANDLE_ALLOC`, respectively).

If the variable has a global scope (as returned by a prior corresponding `MPI_T_CVAR_GET_INFO` call) any write call to this variable must be issued by the user in all connected (as defined in Section 10.5.4) MPI processes. If the variable has a group scope, any write call to this variable must be issued by the user in all MPI processes in the group, which must be described by the MPI implementation in the description by the `MPI_T_CVAR_GET_INFO`.

In both cases, the user must ensure that the writes in all processes are consistent. If the scope is either `MPI_T_SCOPE_GLOBAL_EQ` or `MPI_T_SCOPE_GROUP_EQ` this means that the variable in all processes must be set to the same value.

If it is not possible to change the variable at the time the call is made, the function returns either `MPI_T_ERR_CVAR_SETNOW`, if there may be a later time at which the variable could be set, or `MPI_T_ERR_CVAR_SETNEVER`, if the variable cannot be set for the remainder of the application's execution.

## Example: Reading the Value of a Control Variable

## Example 14.2

The following example shows a routine that can be used to query the value with a control variable with a given index. The example assumes that the variable is intended to be bound to an MPI communicator.

```

1  int getValue_int_comm(int index, MPI_Comm comm, int *val) {
2      int err,count;
3      MPI_T_cvar_handle handle;
4
5      /* This is example assumes that the variable index */
6      /* can be bound to a communicator */
7
8
9      err=MPI_T_cvar_handle_alloc(index,&comm,&handle,&count);
10     if (err!=MPI_SUCCESS) return err;
11
12     /* The following assumes that the variable is */
13     /* represented by a single integer */
14
15     err=MPI_T_cvar_read(handle,val);
16     if (err!=MPI_SUCCESS) return err;
17
18     err=MPI_T_cvar_handle_free(&handle);
19     return err;
20 }
21
22
23
24

```

### 14.3.7 Performance Variables

The following section focuses on the ability to list and query performance variables provided by the MPI implementation. Performance variables provide insight into MPI implementation specific internals and can represent information such as the state of the MPI implementation (e.g., waiting blocked, receiving, not active), aggregated timing data for submodules, or queue sizes and lengths.

*Rationale.* The interface for performance variables is separate from the interface for control variables, since performance variables have different requirements and parameters. By keeping them separate, the interface provides cleaner semantics and allows for more performance optimization opportunities. (*End of rationale.*)

#### Performance Variable Classes

Each performance variable is associated with a class that describes its basic semantics, possible datatypes, basic behavior, its starting value, whether it can overflow, and when and how an MPI implementation can change the variable's value. The starting value is the value the variable assumes when it is used for the first time or whenever it is reset.

*Advice to users.* If a performance variable belongs to a class that can overflow, it is up to the user to appropriately protect against this, e.g., by frequently reading and resetting the variable value. (*End of advice to users.*)



*Advice to implementors.* MPI implementations should use large enough datatypes for each performance variable to avoid overflows under normal circumstances. (*End of advice to implementors.*)

The classes are defined by the following constants:

- **MPI\_T\_PVAR\_CLASS\_STATE**

A performance variable in this class represents a set of discrete states. Variables of this class are represented by MPI\_INT and can be set by the MPI implementation at any time. Variables of this type should be described further using an enumeration, as discussed in Section 14.3.5. The starting value is the current state of the implementation at the time the starting value is set. MPI implementations must ensure that variables of this class cannot overflow.

- **MPI\_T\_PVAR\_CLASS\_LEVEL**

A performance variable in this class represents a value that describes the utilization level of a resource. The value of a variable of this class can change at any time to match the current utilization level of the resource. Values returned from variables in this class are non-negative and represented by one of the following datatypes: MPI\_UNSIGNED, MPI\_UNSIGNED\_LONG, MPI\_UNSIGNED\_LONG\_LONG, MPI\_DOUBLE. The starting value is the current utilization level of the resource at the time the starting value is set. MPI implementations must ensure that variables of this class cannot overflow.

- **MPI\_T\_PVAR\_CLASS\_SIZE**

A performance variable in this class represents a value that is the fixed size of a resource. Values returned from variables in this class are non-negative and represented by one of the following datatypes: MPI\_UNSIGNED, MPI\_UNSIGNED\_LONG, MPI\_UNSIGNED\_LONG\_LONG, MPI\_DOUBLE. The starting value is the current utilization level of the resource at the time the starting value is set. MPI implementations must ensure that variables of this class cannot overflow.

- **MPI\_T\_PVAR\_CLASS\_PERCENTAGE**

The value of a performance variable in this class represents the percentage utilization of a finite resource. The value of a variable of this class can change at any time to match the current utilization level of the resource. It will be returned as an MPI\_DOUBLE datatype. The value must always be between 0.0 (resource not used at all) and 1.0 (resource completely used). The starting value is the current percentage utilization level of the resource at the time the starting value is set. MPI implementations must ensure that variables of this class cannot overflow.

- **MPI\_T\_PVAR\_CLASS\_HIGHWATERMARK**

A performance variable in this class represents a value that describes the high watermark utilization of a resource. The value of a variable of this class is non-negative and grows monotonically from the initialization or reset of the variable. It can be represented by one of the following datatypes: MPI\_UNSIGNED, MPI\_UNSIGNED\_LONG, MPI\_UNSIGNED\_LONG\_LONG, MPI\_DOUBLE. The starting value is the current utilization level of the resource at the time the starting value is set. MPI implementations must ensure that variables of this class cannot overflow.

- **MPI\_T\_PVAR\_CLASS\_LOWWATERMARK**

A performance variable in this class represents a value that describes the low watermark utilization of a resource. The value of a variable of this class is non-negative and decreases monotonically from the initialization or reset of the variable. It can be represented by one of the following datatypes: `MPI_UNSIGNED`, `MPI_UNSIGNED_LONG`, `MPI_UNSIGNED_LONG_LONG`, `MPI_DOUBLE`. The starting value is the current utilization level of the resource at the time the starting value is set. MPI implementations must ensure that variables of this class cannot overflow.

- **MPI\_T\_PVAR\_CLASS\_COUNTER**

A performance variable in this class counts the number of occurrences of a specific event (e.g., the number of memory allocations within an MPI library). The value of a variable of this class increases monotonically from the initialization or reset of the performance variable by one for each specific event that is observed. Values must be non-negative and represented by one of the following datatypes: `MPI_UNSIGNED`, `MPI_UNSIGNED_LONG`, `MPI_UNSIGNED_LONG_LONG`. The starting value for variables of this class is 0. Variables of this class can overflow.

- **MPI\_T\_PVAR\_CLASS\_AGGREGATE**

The value of a performance variable in this class is an aggregated value that represents a sum of arguments processed during a specific event (e.g., the amount of memory allocated by all memory allocations). This class is similar to the counter class, but instead of counting individual events, the value can be incremented by arbitrary amounts. The value of a variable of this class increases monotonically from the initialization or reset of the performance variable. It must be non-negative and represented by one of the following datatypes: `MPI_UNSIGNED`, `MPI_UNSIGNED_LONG`, `MPI_UNSIGNED_LONG_LONG`, `MPI_DOUBLE`. The starting value for variables of this class is 0. Variables of this class can overflow.

- **MPI\_T\_PVAR\_CLASS\_TIMER**

The value of a performance variable in this class represents the aggregated time that the MPI implementation spends executing a particular event, type of event, or section of the MPI library. This class has the same basic semantics as `MPI_T_PVAR_CLASS_AGGREGATE`, but explicitly records a timing value. The value of a variable of this class increases monotonically from the initialization or reset of the performance variable. It must be non-negative and represented by one of the following datatypes: `MPI_UNSIGNED`, `MPI_UNSIGNED_LONG`, `MPI_UNSIGNED_LONG_LONG`, `MPI_DOUBLE`. The starting value for variables of this class is 0. If the type `MPI_DOUBLE` is used, the units representing time in this datatype must match the units used by `MPI_WTIME`. Otherwise, the time units should be documented, e.g., in the description returned by `MPI_T_PVAR_GET_INFO`. Variables of this class can overflow.

- **MPI\_T\_PVAR\_CLASS\_GENERIC**

This class can be used to describe a variable that does not fit into any of the other classes. For variables in this class, the starting value is variable specific and implementation defined.

## Performance Variable Query Functions

An MPI implementation exports a set of  $N$  performance variables through the MPI tool information interface. If  $N$  is zero, then the MPI implementation does not export any performance variables, otherwise the provided performance variables are indexed from 0 to  $N - 1$ . This index number is used in subsequent calls to identify the individual variables.

An MPI implementation is allowed to increase the number of performance variables during the execution of an MPI application when new variables become available through dynamic loading. However, MPI implementations are not allowed to change the index of a performance variable or delete a variable once it has been added to the set. When variables become inactive, e.g., through dynamic unloading, accessing its value should return a corresponding error code.

The following function can be used to query the number of performance variables,  $N$ :

`MPI_T_PVAR_GET_NUM(num_pvar)`

OUT      `num_pvar`      returns number of performance variables (integer)

`int MPI_T_pvar_get_num(int *num_pvar)`

The function `MPI_T_PVAR_GET_INFO` provides access to additional information for each variable.

```

1 MPI_T_PVAR_GET_INFO(pvar_index, name, name_len, verbosity, varclass, datatype, enum-
2   type, desc, desc_len, bind, readonly, continuous, atomic)
3
4   IN      pvar_index      index of the performance variable to be queried be-
5                               tween 0 and num_pvar - 1 (integer)
6
7   OUT     name            buffer to return the string containing the name of the
8                               performance variable (string)
9
10  INOUT   name_len        length of the string and/or buffer for name (integer)
11
12  OUT     verbosity        verbosity level of this variable (integer)
13
14  OUT     var_class        class of performance variable (integer)
15
16  OUT     datatype         MPI datatype of the information stored in the perfor-
17                               mance variable (handle)
18
19  OUT     enumtype         optional descriptor for enumeration information (han-
20                               dle)
21
22  OUT     desc             buffer to return the string containing a description of
23                               the performance variable (string)
24
25  INOUT   desc_len        length of the string and/or buffer for desc (integer)
26
27  OUT     bind             type of MPI object to which this variable must be
28                               bound (integer)
29
30  OUT     readonly         flag indicating whether the variable can be written/reset
31                               (integer)
32
33  OUT     continuous       flag indicating whether the variable can be started and
34                               stopped or is continuously active (integer)
35
36  OUT     atomic           flag indicating whether the variable can be atomically
37                               read and reset (integer)
38
39  int MPI_T_pvar_get_info(int pvar_index, char *name, int *name_len, int
40      *verbosity, int *var_class, MPI_Datatype *datatype, MPI_T_enum
41      *enumtype, char *desc, int *desc_len, int *bind, int
42      *readonly, int *continuous, int *atomic)

```

After a successful call to `MPI_T_PVAR_GET_INFO` for a particular variable, subsequent calls to this routine querying information about the same variable must return the same information. An MPI implementation is not allowed to alter any of the returned values.

The arguments `name` and `name_len` are used to return the name of the performance variable as described in Section 14.3.3. If completed successfully, the routine is required to return a name of at least length one.

The argument `verbosity` returns the verbosity level of the variable (see Section 14.3.1).

The class of the performance variable is returned in the parameter `var_class`. The class must be one of the constants defined in Section 14.3.7.

The combination of the name and the class of the performance variable must be unique with respect to all other names for performance variables used by the MPI implementation.

*Advice to implementors.* Groups of variables that belong closely together, but have different classes, can have the same name. This is, e.g., useful to refer to multiple

variables that describe a single resource (like the level, the total size, as well as high and low watermarks). (*End of advice to implementors.*)

The argument `datatype` returns the MPI datatype that is used to represent the performance variable.

If the variable is of type `MPI_INT`, MPI can optionally specify an enumeration for the values represented by this variable and return it in `enumtype`. In this case, MPI returns an enumeration identifier, which can then be used as described in Section 14.3.5 to gather more information. If the datatype is not `MPI_INT` or the argument `enumtype` is the constant `MPI_T_ENUM_NULL`, no enumeration type is returned.

Returning a description is optional. If an MPI implementation decides not to return a description, the first character for `desc` must be set to the null character and `desc_len` must be set to one at the return from this function.

The parameter `bind` returns the type of the MPI object to which the variable must be bound or the value `MPI_T_BIND_NO_OBJECT` (see Section 14.3.2).

Upon return, the argument `readonly` is set to zero if the variable can be written or reset by the user. It is set to one if the variable can only be read.

Upon return, the argument `continuous` is set to zero if the variable can be started and stopped by the user, i.e., it is possible for the user to control if and when the value of a variable is updated. It is set to one if the variable is always active and cannot be controlled by the user.

Upon return, the argument `atomic` is set to zero if the variable cannot be atomically read and reset. Only variables for which the call sets `atomic` to one, can be used in a call to `MPI_T_PVAR_READRESET`.

### Performance Experiment Sessions

Within a single program, multiple components can use the MPI tool information interface. To avoid collisions with respect to accesses to performance variables, users of the MPI tool information interface must first create a session. Subsequent calls accessing performance variables can then be made within the context of this session. Any call executed in a session must not influence the results in any other session.

#### `MPI_T_PVAR_SESSION_CREATE(session)`

OUT      `session`      identifier of performance session (handle)

```
int MPI_T_pvar_session_create(MPI_T_pvar_session *session)
```

This call creates a new session for accessing performance variables and returns a handle for this session in the argument `session` of type `MPI_T_pvar_session`.

#### `MPI_T_PVAR_SESSION_FREE(session)`

INOUT    `session`      identifier of performance experiment session (handle)

```
int MPI_T_pvar_session_free(MPI_T_pvar_session *session)
```

This call frees an existing session. Calls to the MPI tool information interface can no longer be made within the context of a session after it is freed. On a successful return, MPI sets the session identifier to `MPI_T_PVAR_SESSION_NULL`.

#### Handle Allocation and Deallocation

Before using a performance variable, a user must first allocate a handle of type `MPI_T_pvar_handle` for the variable by binding it to an MPI object (see also Section 14.3.2).

`MPI_T_PVAR_HANDLE_ALLOC(session, pvar_index, obj_handle, handle, count)`

IN	session	identifier of performance experiment session (handle)
IN	pvar_index	index of performance variable for which handle is to be allocated (integer)
IN	obj_handle	reference to a handle of the MPI object to which this variable is supposed to be bound (pointer)
OUT	handle	allocated handle (handle)
OUT	count	number of elements used to represent this variable (integer)

```
int MPI_T_pvar_handle_alloc(MPI_T_pvar_session session, int pvar_index,
                           void *obj_handle, MPI_T_pvar_handle *handle, int *count)
```

This routine binds the performance variable specified by the argument `index` to an MPI object in the session identified by the parameter `session`. The object is passed in the argument `obj_handle` as an address to a local variable that stores the object's handle. The handle allocated to reference the variable is returned in the argument `handle`. Upon successful return, `count` contains the number of elements (of the datatype returned by a previous `MPI_T_PVAR_GET_INFO` call) used to represent this variable.

*Advice to users.* The count can be different based on the MPI object, to which is was bound. For example, variables bound to communicators could have a count that matches the size of the communicator.

It is not portable to pass references to predefined MPI object handles, such as `MPI_COMM_WORLD`, to this routine, since their implementation depends on the MPI library. Instead, such object handles should be stored in a local variable and the address of this local variables should be passed into `MPI_T_PVAR_HANDLE_ALLOC`. (*End of advice to users.*)

The value of `index` should be in the range 0 to `num_pvar - 1`, where `num_pvar` is the number of available control variables as determined from a prior call to `MPI_T_PVAR_GET_NUM`. The type of the MPI object it references must be consistent with the type returned in the `bind` argument in a prior call to `MPI_T_PVAR_GET_INFO`.

In the case the `bind` argument equals `MPI_T_BIND_NO_OBJECT`, the argument `obj_handle` is ignored.

**MPI\_T\_PVAR\_HANDLE\_FREE**(session, handle)

IN	session	identifier of performance experiment session (handle)
INOUT	handle	handle to be freed (handle)

```
int MPI_T_pvar_handle_free(MPI_T_pvar_session session, MPI_T_pvar_handle
                          *handle)
```

When a handle is no longer needed, a user of the MPI tool information interface should call **MPI\_T\_PVAR\_HANDLE\_FREE** to free the handle in the session identified by the parameter **session** and the associated resources in the MPI implementation. On a successful return, MPI sets the handle to **MPI\_T\_PVAR\_HANDLE\_NULL**.

### Starting and Stopping of Performance Variables

Performance variables that have the continuous flag set during the query operation are continuously operating once a handle has been allocated. Such variables may be queried at any time, but they cannot be started or stopped by the user. All other variables are in a stopped state after their handle has been allocated; their values are not updated until they have been started by the user.

**MPI\_T\_PVAR\_START**(session, handle)

IN	session	identifier of performance experiment session (handle)
IN	handle	handle of a performance variable (handle)

```
int MPI_T_pvar_start(MPI_T_pvar_session session, MPI_T_pvar_handle handle)
```

This function starts the performance variable with the handle identified by the parameter **handle** in the session identified by the parameter **session**.

If the constant **MPI\_T\_PVAR\_ALL\_HANDLES** is passed in **handle**, the MPI implementation attempts to start all variables within the session identified by the parameter **session** for which handles have been allocated. In this case, the routine returns **MPI\_SUCCESS** if all variables are started successfully, otherwise **MPI\_T\_ERR\_PVAR\_NOSTARTSTOP** is returned. Continuous variables and variables that are already started are ignored when **MPI\_T\_PVAR\_ALL\_HANDLES** is specified.

**MPI\_T\_PVAR\_STOP**(session, handle)

IN	session	identifier of performance experiment session (handle)
IN	handle	handle of a performance variable (handle)

```
int MPI_T_pvar_stop(MPI_T_pvar_session session, MPI_T_pvar_handle handle)
```

This function stops the performance variable with the handle identified by the parameter **handle** in the session identified by the parameter **session**.

If the constant **MPI\_T\_PVAR\_ALL\_HANDLES** is passed in **handle**, the MPI implementation attempts to stop all variables within the session identified by the parameter **session** for

which handles have been allocated. In this case, the routine returns `MPI_SUCCESS` if all variables are stopped successfully, otherwise `MPI_T_ERR_PVAR_NOSTARTSTOP` is returned. Continuous variables and variables that are already stopped are ignored when `MPI_T_PVAR_ALL_HANDLES` is specified.

## Performance Variable Access Functions

### `MPI_T_PVAR_READ(session, handle, buf)`

IN	session	identifier of performance experiment session (handle)
IN	handle	handle of a performance variable (handle)
OUT	buf	initial address of storage location for variable value (choice)

```
int MPI_T_pvar_read(MPI_T_pvar_session session, MPI_T_pvar_handle handle,
                    void* buf)
```

The `MPI_T_PVAR_READ` call queries the value of the performance variable with the handle `handle` in the session identified by the parameter `session` and stores the result in the buffer identified by the parameter `buf`. The user is responsible to ensure that the buffer is of the appropriate size to hold the entire value of the performance variable (based on the datatype and count returned by the corresponding previous calls to `MPI_T_PVAR_GET_INFO` and `MPI_T_PVAR_HANDLE_ALLOC`, respectively).

The constant `MPI_T_PVAR_ALL_HANDLES` cannot be used as an argument for the function `MPI_T_PVAR_READ`.

### `MPI_T_PVAR_WRITE(session, handle, buf)`

IN	session	identifier of performance experiment session (handle)
IN	handle	handle of a performance variable (handle)
IN	buf	initial address of storage location for variable value (choice)

```
int MPI_T_pvar_write(MPI_T_pvar_session session, MPI_T_pvar_handle handle,
                     const void* buf)
```

The `MPI_T_PVAR_WRITE` call attempts to write the value of the performance variable with the handle identified by the parameter `handle` in the session identified by the parameter `session`. The value to be written is passed in the buffer identified by the parameter `buf`. The user must ensure that the buffer is of the appropriate size to hold the entire value of the performance variable (based on the datatype and count returned by the corresponding previous calls to `MPI_T_PVAR_GET_INFO` and `MPI_T_PVAR_HANDLE_ALLOC`, respectively).

If it is not possible to change the variable, the function returns `MPI_T_ERR_PVAR_NOWRITE`.

The constant `MPI_T_PVAR_ALL_HANDLES` cannot be used as an argument for the function `MPI_T_PVAR_WRITE`.



`MPI_T_PVAR_RESET(session, handle)`

IN	session	identifier of performance experiment session (handle)
IN	handle	handle of a performance variable (handle)

`int MPI_T_pvar_reset(MPI_T_pvar_session session, MPI_T_pvar_handle handle)`

The `MPI_T_PVAR_RESET` call sets the performance variable with the handle identified by the parameter `handle` to its starting value specified in Section 14.3.7. If it is not possible to change the variable, the function returns `MPI_T_ERR_PVAR_NOWRITE`.

If the constant `MPI_T_PVAR_ALL_HANDLES` is passed in `handle`, the MPI implementation attempts to reset all variables within the session identified by the parameter `session` for which handles have been allocated. In this case, the routine returns `MPI_SUCCESS` if all variables are reset successfully, otherwise `MPI_T_ERR_PVAR_NOWRITE` is returned. Read-only variables are ignored when `MPI_T_PVAR_ALL_HANDLES` is specified.

`MPI_T_PVAR_READRESET(session, handle, buf)`

IN	session	identifier of performance experiment session (handle)
IN	handle	handle of a performance variable (handle)
OUT	buf	initial address of storage location for variable value (choice)

`int MPI_T_pvar_readreset(MPI_T_pvar_session session, MPI_T_pvar_handle handle, void* buf)`

This call atomically combines the functionality of `MPI_T_PVAR_READ` and `MPI_T_PVAR_RESET` with the same semantics as if these two calls were called separately. If atomic operations on this variable are not supported, this routine returns `MPI_ERR_NOATOMIC`.

The constant `MPI_T_PVAR_ALL_HANDLES` can not be used as an argument for the function `MPI_T_PVAR_READRESET`.

*Advice to implementors.* Sampling based tools rely on the ability to call the MPI tool information interface, in particular routines to start, stop, read, write and reset performance variables, from any program context, including asynchronous contexts such as signal handlers. MPI implementations should strive, if possible in their particular environment, to enable these usage scenarios for all or a subset of the routines mentioned above. If implementing only a subset, the read, write, and reset routines are typically the most critical for sampling based tools. An MPI implementation should clearly document any restrictions on the program contexts in which the MPI tool information interface can be used. Restrictions might include guaranteeing usage outside of all signals or outside a specific set of signals. Any restrictions could be documented, for example, through the description returned by `MPI_T_PVAR_GET_INFO`. *(End of advice to implementors.)*

*Rationale.* All routines to read, write or reset performance variables require the session argument. This keeps the interface consistent and allows the use

MPI\_T\_PVAR\_ALL\_HANDLES where appropriate. Further, this opens up additional performance optimizations for the implementation of handles. (*End of rationale.*)

Example: Tool to Detect Receives with Long Unexpected Message Queues

### Example 14.3

The following example shows a sample tool to identify receive operations that occur during times with long message queues. This examples assumes that the MPI implementation exports a variable with the name "MPI\_T\_UMQ\_LENGTH" to represent the current length of the unexpected message queue. The tool is implemented as a PMPI tool using the MPI profiling interface.

The tool consists of three parts: (1) the initialization (by intercepting the call to MPI\_INIT), (2) the test for long unexpected message queues (by intercepting calls to MPI\_RECV), and (3) the clean up phase (by intercepting the call to MPI\_FINALIZE. To capture all receives, the example would have to be extended to have similar wrappers for all receive operations.

**Part 1— Initialization:** During initialization, the tool searches for the variable and, once the right index is found, allocates a session and a handle for the variable with the found index, and starts the performance variable.

```
#include <stdio.h>
#include <stdlib.h>
#include <assert.h>
#include <mpi.h>

/* Global variables for the tool */
static MPI_T_pvar_session session;
static MPI_T_pvar_handle handle;

int MPI_Init(int *argc, char ***argv) {
    int err, num, i, index, namelen, verbosity;
    int var_class, bind, threadsup;
    int readonly, continuous, atomic, count;
    char name[17];
    MPI_Comm comm;
    MPI_Datatype datatype;
    MPI_T_enum enumtype;

    err=PMPI_Init(argc,argv);
    if (err!=MPI_SUCCESS) return err;

    err=PMPI_T_init_thread(MPI_THREAD_SINGLE,&threaddup);
    if (err!=MPI_SUCCESS) return err;

    err=PMPI_T_pvar_get_num(&num);
    if (err!=MPI_SUCCESS) return err;
```

```

index=-1;
i=0;
while ((i<num) && (index<0)) {
    namelen=17;
    err=PMPI_T_pvar_get_info(i, name, namelen, &verbosity,
        &var_class, &datatype, &enumtype, &bind,
        &readonly, &continuous, &atomic);
    if (strcmp(name,"MPI_T_UMQ_LENGTH")==0) index=i;
    i++; }

/* this could be handled in a more flexible way for a generic tool */
assert(index>=0);
assert(var_class==MPI_T_PVAR_CLASS_LEVEL);
assert(datatype==MPI_INT);
assert(bind==MPI_T_BIND_MPI_COMM);

/* Create a session */
err=PMPI_T_pvar_session_create(&session);
if (err!=MPI_SUCCESS) return err;

/* Get a handle and bind to MPI_COMM_WORLD */
comm=MPI_COMM_WORLD;
err=PMPI_T_pvar_handle_alloc(session, index, &comm, &handle, &count);
if (err!=MPI_SUCCESS) return err;

/* this could be handled in a more flexible way for a generic tool */
assert(count==1);

/* Start variable */
err=PMPI_T_pvar_start(session, handle);
if (err!=MPI_SUCCESS) return err;

return MPI_SUCCESS;
}

```

Part 2 — Testing the Queue Lengths During Receives: During every receive operation, the tool reads the unexpected queue length through the matching performance variable and compares it against a predefined threshold.

```

#define THRESHOLD 5

int MPI_Recv(void *buf, int count, MPI_Datatype datatype, int source, int tag,
    MPI_Comm comm, MPI_Status *status)
{
    int value, err;

    if (comm==MPI_COMM_WORLD) {
        err=PMPI_T_pvar_read(session, handle, &value);

```

```

1         if ((err==MPI_SUCCESS) && (value>THRESHOLD))
2         {
3             /* tool identified receive called with long UMQ */
4             /* execute tool functionality, */
5             /* e.g., gather and print call stack */
6         }
7     }
8
9     return PMPI_Recv(buf, count, datatype, source, tag, comm, status);
10 }
11

```

Part 3 — Termination: In the wrapper for MPI\_FINALIZE, the MPI tool information interface is finalized.

```

15 int MPI_Finalize()
16 {
17     int err;
18     err=PMPI_T_handle_free(&session, &handle);
19     err=PMPI_T_session_free(&session);
20     err=PMPI_T_finalize();
21     return PMPI_Finalize();
22 }
23

```

#### 14.3.8 Variable Categorization

MPI implementations can optionally group performance and control variables into categories to express logical relationships between various variables. For example, an MPI implementation could group all control and performance variables that refer to message transfers in the MPI implementation and thereby distinguish them from variables that refer to local resources such as memory allocations or other interactions with the operating system.

Categories can also contain other categories to form a hierarchical grouping. Categories can never include themselves, either directly or transitively within other included categories. Expanding on the example above, this allows MPI to refine the grouping of variables referring to message transfers into variables to control and monitor message queues, message matching activities and communication protocols. Each of these groups of variables would be represented by a separate category and these categories would then be listed in a single category representing variables for message transfers.

The category information may be queried in a fashion similar to the mechanism for querying variable information. The MPI implementation exports a set of  $N$  categories via the MPI tool information interface. If  $N = 0$ , then the MPI implementation does not export any categories, otherwise the provided categories are indexed from 0 to  $N - 1$ . This index number is used in subsequent calls to functions of the MPI tool information interface to identify the individual categories.

An MPI implementation is permitted to increase the number of categories during the execution of an MPI program when new categories become available through dynamic loading. However, MPI implementations are not allowed to change the index of a category or delete it once it has been added to the set.

Similarly, MPI implementations are allowed to add variables to categories, but they are not allowed to remove variables from categories or change the order in which they are returned.

The following function can be used to query the number of control variables,  $N$ .

**MPI\_T\_CATEGORY\_GET\_NUM**(num\_cat)

OUT      num\_cat                      current number of categories (integer)

int MPI\_T\_category\_get\_num(int \*num\_cat)

Individual category information can then be queried by calling the following function:

**MPI\_T\_CATEGORY\_GET\_INFO**(cat\_index, name, name\_len, desc, desc\_len, num\_cvars, num\_pvars, num\_categories)

IN	cat_index	index of the category to be queried (integer)
OUT	name	buffer to return the string containing the name of the category (string)
INOUT	name_len	length of the string and/or buffer for name (integer)
OUT	desc	buffer to return the string containing the description of the category (string)
INOUT	desc_len	length of the string and/or buffer for desc (integer)
OUT	num_cvars	number of control variables in the category (integer)
OUT	num_pvars	number of performance variables in the category (integer)
OUT	num_categories	number of categories contained in the category (integer)

int MPI\_T\_category\_get\_info(int cat\_index, char \*name, int \*name\_len, char \*desc, int \*desc\_len, int \*num\_cvars, int \*num\_pvars, int \*num\_categories)

The arguments name and name\_len are used to return the name of the category as described in Section 14.3.3.

The routine is required to return a name of at least length one. This name must be unique with respect to all other names for categories used by the MPI implementation.

The arguments desc and desc\_len are used to return the description of the category as described in Section 14.3.3.

Returning a description is optional. If an MPI implementation decides not to return a description, the first character for desc must be set to the null character and desc\_len must be set to one at the return of this call.

The function returns the number of control variables, performance variables and other categories contained in the queried category in the arguments num\_cvars, num\_pvars, and num\_categories, respectively.

```
1 MPI_T_CATEGORY_GET_CVARS(cat_index, len, indices)
```

```
2     IN      cat_index      index of the category to be queried, in the range [0, N-
3                               1] (integer)
4
5     IN      len            the length of the indices array (integer)
6
7     OUT     indices        an integer array of size len, indicating control variable
8                               indices (array of integers)
```

```
9 int MPI_T_category_get_cvars(int cat_index, int len, int indices[])
10
```

11 MPI\_T\_CATEGORY\_GET\_CVARS can be used to query which control variables are  
12 contained in a particular category. A category contains zero or more control variables.

```
14 MPI_T_CATEGORY_GET_PVARS(cat_index,len,indices)
```

```
16     IN      cat_index      index of the category to be queried, in the range [0, N-
17                               1] (integer)
18
19     IN      len            the length of the indices array (integer)
20
21     OUT     indices        an integer array of size len, indicating performance
22                               variable indices (array of integers)
```

```
23 int MPI_T_category_get_pvars(int cat_index, int len, int indices[])
24
```

25 MPI\_T\_CATEGORY\_GET\_PVARS can be used to query which performance variables  
26 are contained in a particular category. A category contains zero or more performance  
27 variables.

```
29 MPI_T_CATEGORY_GET_CATEGORIES(cat_index,len,indices)
```

```
30     IN      cat_index      index of the category to be queried, in the range [0, N-
31                               1] (integer)
32
33     IN      len            the length of the indices array (integer)
34
35     OUT     indices        an integer array of size len, indicating category indices
36                               (array of integers)
```

```
37 int MPI_T_category_get_categories(int cat_index, int len, int indices[])
38
```

39 MPI\_T\_CATEGORY\_GET\_CATEGORIES can be used to query which other categories  
40 are contained in a particular category. A category contains zero or more other categories.

41 As mentioned above, MPI implementations can grow the number of categories as well  
42 as the number of variables or other categories within a category. In order to allow users  
43 of the MPI tool information interface to quickly check whether new categories have been  
44 added or new variables or categories have been added to a category, MPI maintains a  
45 virtual timestamp. This timestamp is monotonically increasing during the execution and is  
46 returned by the following function:

**MPI\_T\_CATEGORY\_CHANGED(stamp)**

OUT      stamp                      a virtual time stamp to indicate the last change to the categories (integer)

```
int MPI_T_category_changed(int *stamp)
```

If two subsequent calls to this routine return the same timestamp, it is guaranteed that the category information has not changed between the two calls. If the timestamp retrieved from the second call is higher, then some categories have been added or expanded.

*Advice to users.* The timestamp value is purely virtual and only intended to check for changes in the category information. It should not be used for any other purpose. (*End of advice to users.*)

The index values returned in indices by `MPI_T_CATEGORY_GET_CVARS`, `MPI_T_CATEGORY_GET_PVARS` and `MPI_T_CATEGORY_GET_CATEGORIES` can be used as input to `MPI_T_CVAR_GET_INFO`, `MPI_T_PVAR_GET_INFO` and `MPI_T_CATEGORY_GET_INFO`, respectively.

The user is responsible for allocating the arrays passed into the functions `MPI_T_CATEGORY_GET_CVARS`, `MPI_T_CATEGORY_GET_PVARS` and `MPI_T_CATEGORY_GET_CATEGORIES`. Starting from array index 0, each function writes up to `len` elements into the array. If the category contains more than `len` elements, the function returns an arbitrary subset of size `len`. Otherwise, the entire set of elements is returned in the beginning entries of the array, and any remaining array entries are not modified.

**14.3.9 Return Codes for the MPI tool information interface**

All functions defined as part of the MPI tool information interface return an integer return code (see Table 14.5) to indicate whether the function has completed successfully or aborted its execution. In the latter case the return code indicates the reason for not completing the routine. None of the return codes returned by an routine impact the execution of the MPI process and do not invoke MPI error handlers. The execution of the MPI process continues as if the call would have completed. However, the MPI implementation is not required to check all user provided parameters; if a user passes invalid parameter values to any routine the behavior of the implementation is undefined.

All return codes with the prefix `MPI_T_` must be unique values and cannot overlap with any other return values returned by the MPI implementation.

**14.3.10 Profiling Interface**

All requirements for the profiling interfaces, as described in Section 14.2, also apply to the MPI tool information interface. All rules, guidelines, and recommendations from Section 14.2 apply equally to calls defined as part of the MPI tool information interface.

Return Code	Description
Return Codes for all Functions in the MPI tool information interface	
MPI_SUCCESS	Call completed successfully
MPI_T_ERR_MEMORY	Out of memory
MPI_T_ERR_NOTINITIALIZED	Interface not initialized
MPI_T_ERR_CANTINIT	Interface not in the state to be initialized
Return Codes for Datatype Functions: MPI_T_ENUM_*	
MPI_T_ERR_INVALIDINDEX	The enumeration index is invalid or has been deleted.
MPI_T_ERR_INVALIDITEM	The item index queried is out of range (for MPI_T_ENUMITEM only)
Return Codes for variable and category query functions: MPI_T_*.GET_INFO	
MPI_T_ERR_INVALIDINDEX	The variable or category index is invalid
Return Codes for Handle Functions: MPI_T_*.ALLOCATE,FREE	
MPI_T_ERR_INVALIDINDEX	The variable index is invalid or has been deleted
MPI_T_ERR_INVALIDHANDLE	The handle is invalid
MPI_T_ERR_OUTOFHANDLES	No more handles available
Return Codes for Session Functions: MPI_T_PVAR_SESSION_*	
MPI_T_ERR_OUTOFSESSIONS	No more sessions available
MPI_T_ERR_INVALIDSESSION	Session argument is not a valid session
Return Codes for Control Variable Access Functions:	
MPI_T_CVAR_READ, WRITE	
MPI_T_ERR_CVAR_SETNOTNOW	Variable cannot be set at this moment
MPI_T_ERR_CVAR_SETNEVER	Variable cannot be set until end of execution
MPI_T_ERR_INVALIDHANDLE	The handle is invalid
Return Codes for Performance Variable Access and Control:	
MPI_T_PVAR_START, STOP, READ, WRITE, RESET, READRESET	
MPI_T_ERR_INVALIDHANDLE	The handle is invalid
MPI_T_ERR_INVALIDSESSION	Session argument is not a valid session
MPI_T_ERR_PVAR_NOSTARTSTOP	Variable can not be started or stopped (for MPI_T_PVAR_START and MPI_T_PVAR_STOP)
MPI_T_ERR_PVAR_NOWRITE	Variable can not be written or reset (for MPI_T_PVAR_WRITE and MPI_T_PVAR_RESET)
MPI_T_NOATOMIC	Variable cannot be read and written atomically (for MPI_T_PVAR_READRESET)
Return Codes for Category Functions: MPI_T_CATEGORY_*	
MPI_T_ERR_INVALIDINDEX	The category index is invalid

Table 14.5: Return codes used in functions of the MPI tool information interface.



## Chapter 15

# Deprecated Functions

### 15.1 Deprecated since MPI-2.0

The following function is deprecated and is superseded by `MPI_TYPE_CREATE_HVECTOR` in MPI-2.0. The language independent definition and the C binding of the deprecated function is the same as of the new function, except of the function name. Only the Fortran language binding is different.

`MPI_TYPE_HVECTOR( count, blocklength, stride, oldtype, newtype)`

IN	count	number of blocks (non-negative integer)
IN	blocklength	number of elements in each block (non-negative integer)
IN	stride	number of bytes between start of each block (integer)
IN	oldtype	old datatype (handle)
OUT	newtype	new datatype (handle)

```
int MPI_Type_hvector(int count, int blocklength, MPI_Aint stride,  
                    MPI_Datatype oldtype, MPI_Datatype *newtype)
```

For this routine, an interface within the `mpi_f08` module was never defined.

```
MPI_TYPE_HVECTOR(COUNT, BLOCKLENGTH, STRIDE, OLDTYPE, NEWTYPE, IERROR)  
INTEGER COUNT, BLOCKLENGTH, STRIDE, OLDTYPE, NEWTYPE, IERROR
```

The following function is deprecated and is superseded by `MPI_TYPE_CREATE_HINDEXED` in MPI-2.0. The language independent definition and the C binding of the deprecated function is the same as of the new function, except of the function name. Only the Fortran language binding is different.

```

1 MPI_TYPE_HINDEXED( count, array_of_blocklengths, array_of_displacements, oldtype, new-
2     type)
3
4     IN          count          number of blocks – also number of entries in
5                               array_of_displacements and array_of_blocklengths (non-
6                               negative integer)
7
8     IN          array_of_blocklengths  number of elements in each block (array of non-negative
9                               integers)
10
11    IN          array_of_displacements  byte displacement of each block (array of integer)
12
13    IN          oldtype              old datatype (handle)
14
15    OUT         newtype              new datatype (handle)

```

```

16 int MPI_Type_hindexed(int count, int *array_of_blocklengths,
17     MPI_Aint *array_of_displacements, MPI_Datatype oldtype,
18     MPI_Datatype *newtype)

```

For this routine, an interface within the `mpi_f08` module was never defined.

```

19 MPI_TYPE_HINDEXED(COUNT, ARRAY_OF_BLOCKLENGTHS, ARRAY_OF_DISPLACEMENTS,
20     OLDTYPE, NEWTYPE, IERROR)
21     INTEGER COUNT, ARRAY_OF_BLOCKLENGTHS(*), ARRAY_OF_DISPLACEMENTS(*),
22     OLDTYPE, NEWTYPE, IERROR

```

The following function is deprecated and is superseded by `MPI_TYPE_CREATE_STRUCT` in MPI-2.0. The language independent definition and the C binding of the deprecated function is the same as of the new function, except of the function name. Only the Fortran language binding is different.

```

29 MPI_TYPE_STRUCT(count, array_of_blocklengths, array_of_displacements, array_of_types,
30     newtype)
31
32    IN          count          number of blocks (integer) (non-negative integer) –
33                               also number of entries in arrays array_of_types,
34                               array_of_displacements and array_of_blocklengths
35
36    IN          array_of_blocklength  number of elements in each block (array of non-negative
37                               integer)
38
39    IN          array_of_displacements  byte displacement of each block (array of integer)
40
41    IN          array_of_types        type of elements in each block (array of handles to
42                               datatype objects)
43
44    OUT         newtype              new datatype (handle)

```

```

45 int MPI_Type_struct(int count, int *array_of_blocklengths,
46     MPI_Aint *array_of_displacements,
47     MPI_Datatype *array_of_types, MPI_Datatype *newtype)

```

For this routine, an interface within the `mpi_f08` module was never defined.

```

MPI_TYPE_STRUCT(COUNT, ARRAY_OF_BLOCKLENGTHS, ARRAY_OF_DISPLACEMENTS,
                ARRAY_OF_TYPES, NEWTYPE, IERROR)
    INTEGER COUNT, ARRAY_OF_BLOCKLENGTHS(*), ARRAY_OF_DISPLACEMENTS(*),
    ARRAY_OF_TYPES(*), NEWTYPE, IERROR

```

The following function is deprecated and is superseded by `MPI_GET_ADDRESS` in MPI-2.0. The language independent definition and the C binding of the deprecated function is the same as of the new function, except of the function name. Only the Fortran language binding is different.

```

MPI_ADDRESS(location, address)

```

IN	location	location in caller memory (choice)
OUT	address	address of location (integer)

```

int MPI_Address(void* location, MPI_Aint *address)

```

For this routine, an interface within the `mpi_f08` module was never defined.

```

MPI_ADDRESS(LOCATION, ADDRESS, IERROR)
    <type> LOCATION(*)
    INTEGER ADDRESS, IERROR

```

The following functions are deprecated and are superseded by `MPI_TYPE_GET_EXTENT` in MPI-2.0.

```

MPI_TYPE_EXTENT(datatype, extent)

```

IN	datatype	datatype (handle)
OUT	extent	datatype extent (integer)

```

int MPI_Type_extent(MPI_Datatype datatype, MPI_Aint *extent)

```

For this routine, an interface within the `mpi_f08` module was never defined.

```

MPI_TYPE_EXTENT(DATATYPE, EXTENT, IERROR)
    INTEGER DATATYPE, EXTENT, IERROR

```

Returns the extent of a datatype, where extent is as defined on page 113.

The two functions below can be used for finding the lower bound and the upper bound of a datatype.

```

MPI_TYPE_LB( datatype, displacement)

```

IN	datatype	datatype (handle)
OUT	displacement	displacement of lower bound from origin, in bytes (integer)

```

int MPI_Type_lb(MPI_Datatype datatype, MPI_Aint* displacement)

```

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For this routine, an interface within the `mpi_f08` module was never defined.

```
MPI_TYPE_LB( DATATYPE, DISPLACEMENT, IERROR)
    INTEGER DATATYPE, DISPLACEMENT, IERROR
```

```
MPI_TYPE_UB( datatype, displacement)
```

IN	datatype	datatype (handle)
OUT	displacement	displacement of upper bound from origin, in bytes (integer)

```
int MPI_Type_ub(MPI_Datatype datatype, MPI_Aint* displacement)
```

For this routine, an interface within the `mpi_f08` module was never defined.

```
MPI_TYPE_UB( DATATYPE, DISPLACEMENT, IERROR)
    INTEGER DATATYPE, DISPLACEMENT, IERROR
```

The following function is deprecated and is superseded by `MPI_COMM_CREATE_KEYVAL` in MPI-2.0. The language independent definition of the deprecated function is the same as that of the new function, except for the function name and a different behavior in the C/Fortran language interoperability, see Section 16.3.7 on page 702. The language bindings are modified.

```
MPI_KEYVAL_CREATE(copy_fn, delete_fn, keyval, extra_state)
```

IN	copy_fn	Copy callback function for keyval
IN	delete_fn	Delete callback function for keyval
OUT	keyval	key value for future access (integer)
IN	extra_state	Extra state for callback functions

```
int MPI_Keyval_create(MPI_Copy_function *copy_fn, MPI_Delete_function
    *delete_fn, int *keyval, void* extra_state)
```

For this routine, an interface within the `mpi_f08` module was never defined.

```
MPI_KEYVAL_CREATE(COPY_FN, DELETE_FN, KEYVAL, EXTRA_STATE, IERROR)
    EXTERNAL COPY_FN, DELETE_FN
    INTEGER KEYVAL, EXTRA_STATE, IERROR
```

The `copy_fn` function is invoked when a communicator is duplicated by `MPI_COMM_DUP`. `copy_fn` should be of type `MPI_Copy_function`, which is defined as follows:

```
typedef int MPI_Copy_function(MPI_Comm oldcomm, int keyval,
    void *extra_state, void *attribute_val_in,
    void *attribute_val_out, int *flag)
```

ticket229.1.	A Fortran declaration for such a function is as follows:	1
	For this routine, an interface within the <code>mpi_f08</code> module was never defined.	2
		3
	SUBROUTINE COPY_FUNCTION(OLDCOMM, KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,	4
	ATTRIBUTE_VAL_OUT, FLAG, IERR)	5
	INTEGER OLDCOMM, KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,	6
	ATTRIBUTE_VAL_OUT, IERR	7
	LOGICAL FLAG	8
		9
	copy_fn may be specified as MPI_NULL_COPY_FN or MPI_DUP_FN from either C or	10
	FORTTRAN; MPI_NULL_COPY_FN is a function that does nothing other than returning	11
	flag = 0 and MPI_SUCCESS. MPI_DUP_FN is a simple-minded copy function that sets flag =	12
	1, returns the value of attribute_val_in in attribute_val_out, and returns MPI_SUCCESS. Note	13
	that MPI_NULL_COPY_FN and MPI_DUP_FN are also deprecated.	14
	Analogous to copy_fn is a callback deletion function, defined as follows. The delete_fn	15
	function is invoked when a communicator is deleted by MPI_COMM_FREE or when a call	16
	is made explicitly to MPI_ATTR_DELETE. delete_fn should be of type MPI_Delete_function,	17
	which is defined as follows:	18
		19
	typedef int MPI_Delete_function(MPI_Comm comm, int keyval,	20
	void *attribute_val, void *extra_state);	21
		22
	A Fortran declaration for such a function is as follows:	23
	For this routine, an interface within the <code>mpi_f08</code> module was never defined.	24
		25
	SUBROUTINE DELETE_FUNCTION(COMM, KEYVAL, ATTRIBUTE_VAL, EXTRA_STATE, IERR)	26
	INTEGER COMM, KEYVAL, ATTRIBUTE_VAL, EXTRA_STATE, IERR	27
		28
	delete_fn may be specified as MPI_NULL_DELETE_FN from either C or FORTRAN;	29
	MPI_NULL_DELETE_FN is a function that does nothing, other than returning	30
	MPI_SUCCESS. Note that MPI_NULL_DELETE_FN is also deprecated.	31
	The following function is deprecated and is superseded by MPI_COMM_FREE_KEYVAL	32
	in MPI-2.0. The language independent definition of the deprecated function is the same as	33
	of the new function, except of the function name. The language bindings are modified.	34
		35
	MPI_KEYVAL_FREE(keyval)	36
	INOUT keyval Frees the integer key value (integer)	37
		38
	int MPI_Keyval_free(int *keyval)	39
	For this routine, an interface within the <code>mpi_f08</code> module was never defined.	40
		41
	MPI_KEYVAL_FREE(KEYVAL, IERROR)	42
	INTEGER KEYVAL, IERROR	43
		44
	The following function is deprecated and is superseded by MPI_COMM_SET_ATTR in	45
	MPI-2.0. The language independent definition of the deprecated function is the same as of	46
	the new function, except of the function name. The language bindings are modified.	47
		48

```
1 MPI_ATTR_PUT(comm, keyval, attribute_val)
```

```
2     INOUT    comm                communicator to which attribute will be attached (han-
3                                     dle)
```

```
5     IN      keyval              key value, as returned by
6                                     MPI_KEYVAL_CREATE (integer)
```

```
7     IN      attribute_val       attribute value
```

```
9 int MPI_Attr_put(MPI_Comm comm, int keyval, void* attribute_val)
```

```
11 For this routine, an interface within the mpi_f08 module was never defined.
```

```
12 MPI_ATTR_PUT(COMM, KEYVAL, ATTRIBUTE_VAL, IERROR)
```

```
13     INTEGER COMM, KEYVAL, ATTRIBUTE_VAL, IERROR
```

```
15     The following function is deprecated and is superseded by MPI_COMM_GET_ATTR in
16     MPI-2.0. The language independent definition of the deprecated function is the same as of
17     the new function, except of the function name. The language bindings are modified.
```

```
19 MPI_ATTR_GET(comm, keyval, attribute_val, flag)
```

```
21     IN      comm                communicator to which attribute is attached (handle)
```

```
22     IN      keyval              key value (integer)
```

```
23     OUT     attribute_val       attribute value, unless flag = false
```

```
25     OUT     flag                true if an attribute value was extracted; false if no
26                                     attribute is associated with the key
```

```
28 int MPI_Attr_get(MPI_Comm comm, int keyval, void *attribute_val, int *flag)
```

```
30 For this routine, an interface within the mpi_f08 module was never defined.
```

```
31 MPI_ATTR_GET(COMM, KEYVAL, ATTRIBUTE_VAL, FLAG, IERROR)
```

```
32     INTEGER COMM, KEYVAL, ATTRIBUTE_VAL, IERROR
```

```
33     LOGICAL FLAG
```

```
35     The following function is deprecated and is superseded by MPI_COMM_DELETE_ATTR
36     in MPI-2.0. The language independent definition of the deprecated function is the same as
37     of the new function, except of the function name. The language bindings are modified.
```

```
39 MPI_ATTR_DELETE(comm, keyval)
```

```
41     INOUT    comm                communicator to which attribute is attached (handle)
```

```
42     IN      keyval              The key value of the deleted attribute (integer)
```

```
44 int MPI_Attr_delete(MPI_Comm comm, int keyval)
```

```
46 For this routine, an interface within the mpi_f08 module was never defined.
```

```
47 MPI_ATTR_DELETE(COMM, KEYVAL, IERROR)
```

INTEGER COMM, KEYVAL, IERROR

The following function is deprecated and is superseded by MPI\_COMM\_CREATE\_ERRHANDLER in MPI-2.0. The language independent definition of the deprecated function is the same as of the new function, except of the function name. The language bindings are modified.

MPI\_ERRHANDLER\_CREATE( [function]handler\_fn, errhandler )

IN [ticket252-W.] [function]handler\_fn user defined error handling procedure  
OUT errhandler MPI error handler (handle)

int MPI\_Errhandler\_create(MPI\_Handler\_function \* [function]handler\_fn,  
MPI\_Errhandler \*errhandler)

For this routine, an interface within the mpi\_f08 module was never defined.

MPI\_ERRHANDLER\_CREATE( [FUNCTION]HANDLER\_FN, ERRHANDLER, IERROR)  
EXTERNAL [FUNCTION]HANDLER\_FN  
INTEGER ERRHANDLER, IERROR

Register the user routine [function]handler\_fn for use as an MPI exception handler. Returns in errhandler a handle to the registered exception handler.

In the C language, the user routine should be a C function of type MPI\_Handler\_function, which is defined as:

typedef void (MPI\_Handler\_function)(MPI\_Comm \*, int \*, ...);

The first argument is the communicator in use, the second is the error code to be returned.

In the Fortran language, the user routine should be of the form:

SUBROUTINE HANDLER\_FUNCTION(COMM, ERROR\_CODE)  
INTEGER COMM, ERROR\_CODE

The following function is deprecated and is superseded by MPI\_COMM\_SET\_ERRHANDLER in MPI-2.0. The language independent definition of the deprecated function is the same as of the new function, except of the function name. The language bindings are modified.

MPI\_ERRHANDLER\_SET( comm, errhandler )

INOUT comm communicator to set the error handler for (handle)  
IN errhandler new MPI error handler for communicator (handle)

int MPI\_Errhandler\_set(MPI\_Comm comm, MPI\_Errhandler errhandler)

For this routine, an interface within the mpi\_f08 module was never defined.

MPI\_ERRHANDLER\_SET(COMM, ERRHANDLER, IERROR)

1       INTEGER COMM, ERRHANDLER, IERROR

2       Associates the new error handler `errorhandler` with communicator `comm` at the calling  
3 process. Note that an error handler is always associated with the communicator.

4       The following function is deprecated and is superseded by  
5 `MPI_COMM_GET_ERRHANDLER` in MPI-2.0. The language independent definition of the  
6 deprecated function is the same as of the new function, except of the function name. The  
7 language bindings are modified.  
8

9  
10    MPI\_ERRHANDLER\_GET( comm, errhandler )

11       IN       comm                           communicator to get the error handler from (handle)  
12  
13       OUT       errhandler                   MPI error handler currently associated with commu-  
14 nicator (handle)

15  
16    int MPI\_Errhandler\_get(MPI\_Comm comm, MPI\_Errhandler \*errhandler)

17    For this routine, an interface within the `mpi_f08` module was never defined.

18  
19    MPI\_ERRHANDLER\_GET(COMM, ERRHANDLER, IERROR)

20       INTEGER COMM, ERRHANDLER, IERROR

21  
22       Returns in `errhandler` (a handle to) the error handler that is currently associated with  
23 communicator `comm`.  
24

## 25   15.2   Deprecated since MPI-2.2

26  
27    The entire set of C++ language bindings have been deprecated.

28  
29       *Rationale.* The C++ bindings add minimal functionality over the C bindings while  
30 incurring a significant amount of maintenance to the MPI specification. Since the  
31 C++ bindings are effectively a one-to-one mapping of the C bindings, it should be  
32 relatively easy to convert existing C++ MPI applications to use the MPI C bindings.  
33 Additionally, there are third party packages available that provide C++ class library  
34 functionality (i.e., C++-specific functionality layered on top of the MPI C bindings)  
35 that are likely more expressive and/or natural to C++ programmers and are not  
36 suitable for standardization in this specification. (*End of rationale.*)  
37

38       The following function typedefs have been deprecated and are superseded by new  
39 names. Other than the typedef names, the function signatures are exactly the same; the  
40 names were updated to match conventions of other function typedef names.

Deprecated Name	New Name
MPI_Comm_errhandler_fn	MPI_Comm_errhandler_function
MPI::Comm::Errhandler_fn	MPI::Comm::Errhandler_function
MPI_File_errhandler_fn	MPI_File_errhandler_function
MPI::File::Errhandler_fn	MPI::File::Errhandler_function
MPI_Win_errhandler_fn	MPI_Win_errhandler_function
MPI::Win::Errhandler_fn	MPI::Win::Errhandler_function



## 15.3 Deprecated since MPI-3.0

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3 ticket229.1.  
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# Chapter 16

## Language Bindings

### 16.1 C++

#### 16.1.1 Overview

The C++ language bindings have been deprecated. **A compliant MPI implementation providing C++ language bindings must provide the entire set defined in this document.** [ticket279](#).

There are some issues specific to C++ that must be considered in the design of an interface that go beyond the simple description of language bindings. In particular, in C++, we must be concerned with the design of objects and their interfaces, rather than just the design of a language-specific functional interface to MPI. Fortunately, the design of MPI was based on the notion of objects, so a natural set of classes is already part of MPI.

MPI-2 includes C++ bindings as part of its function specifications. In some cases, MPI-2 provides new names for the C bindings of MPI-1 functions. In this case, the C++ binding matches the new C name — there is no binding for the deprecated name.

#### 16.1.2 Design

The C++ language interface for MPI is designed according to the following criteria:

1. The C++ language interface consists of a small set of classes with a lightweight functional interface to MPI. The classes are based upon the fundamental MPI object types (e.g., communicator, group, etc.).
2. The MPI C++ language bindings provide a semantically correct interface to MPI.
3. To the greatest extent possible, the C++ bindings for MPI functions are member functions of MPI classes.

*Rationale.* Providing a lightweight set of MPI objects that correspond to the basic MPI types is the best fit to MPI's implicit object-based design; methods can be supplied for these objects to realize MPI functionality. The existing C bindings can be used in C++ programs, but much of the expressive power of the C++ language is forfeited. On the other hand, while a comprehensive class library would make user programming more elegant, such a library it is not suitable as a language binding for MPI since a binding must provide a direct and unambiguous mapping to the specified functionality of MPI. (*End of rationale.*)

### 16.1.3 C++ Classes for MPI

All MPI classes, constants, and functions are declared within the scope of an MPI `namespace`. Thus, instead of the `MPI_` prefix that is used in C and Fortran, MPI functions essentially have an `MPI::` prefix.

The members of the MPI namespace are those classes corresponding to objects implicitly used by MPI. An abbreviated definition of the MPI namespace and its member classes is as follows:

```
namespace MPI {
    class Comm { ... };
    class Intracomm : public Comm { ... };
    class Graphcomm : public Intracomm { ... };
    class Distgraphcomm : public Intracomm { ... };
    class Cartcomm : public Intracomm { ... };
    class Intercomm : public Comm { ... };
    class Datatype { ... };
    class Errhandler { ... };
    class Exception { ... };
    class File { ... };
    class Group { ... };
    class Info { ... };
    class Op { ... };
    class Request { ... };
    class Prerequest : public Request { ... };
    class Grequest : public Request { ... };
    class Status { ... };
    class Win { ... };
};
```

Note that there are a small number of derived classes, and that virtual inheritance is *not* used.

### 16.1.4 Class Member Functions for MPI

Besides the member functions which constitute the C++ language bindings for MPI, the C++ language interface has additional functions (as required by the C++ language). In particular, the C++ language interface must provide a constructor and destructor, an assignment operator, and comparison operators.

The complete set of C++ language bindings for MPI is presented in Annex A.5. The bindings take advantage of some important C++ features, such as references and `const`. Declarations (which apply to all MPI member classes) for construction, destruction, copying, assignment, comparison, and mixed-language operability are also provided.

Except where indicated, all non-static member functions (except for constructors and the assignment operator) of MPI member classes are virtual functions.

*Rationale.* Providing virtual member functions is an important part of design for inheritance. Virtual functions can be bound at run-time, which allows users of libraries to re-define the behavior of objects already contained in a library. There is a small

performance penalty that must be paid (the virtual function must be looked up before it can be called). However, users concerned about this performance penalty can force compile-time function binding. (*End of rationale.*)

**Example 16.1** Example showing a derived MPI class.

```
class foo_comm : public MPI::Intracomm {
public:
    void Send(const void* buf, int count, const MPI::Datatype& type,
              int dest, int tag) const
    {
        // Class library functionality
        MPI::Intracomm::Send(buf, count, type, dest, tag);
        // More class library functionality
    }
};
```

*Advice to implementors.* Implementors must be careful to avoid unintended side effects from class libraries that use inheritance, especially in layered implementations. For example, if MPI\_BCAST is implemented by repeated calls to MPI\_SEND or MPI\_RECV, the behavior of MPI\_BCAST cannot be changed by derived communicator classes that might redefine MPI\_SEND or MPI\_RECV. The implementation of MPI\_BCAST must explicitly use the MPI\_SEND (or MPI\_RECV) of the base MPI::Comm class. (*End of advice to implementors.*)

### 16.1.5 Semantics

The semantics of the member functions constituting the C++ language binding for MPI are specified by the MPI function description itself. Here, we specify the semantics for those portions of the C++ language interface that are not part of the language binding. In this subsection, functions are prototyped using the type MPI::<CLASS> rather than listing each function for every MPI class; the word <CLASS> can be replaced with any valid MPI class name (e.g., Group), except as noted.

**Construction / Destruction** The default constructor and destructor are prototyped as follows:

```
{ MPI::<CLASS>() (binding deprecated, see Section 15.2) }
{ ~MPI::<CLASS>() (binding deprecated, see Section 15.2) }
```

In terms of construction and destruction, opaque MPI user level objects behave like handles. Default constructors for all MPI objects except MPI::Status create corresponding MPI::\*\_NULL handles. That is, when an MPI object is instantiated, comparing it with its corresponding MPI::\*\_NULL object will return true. The default constructors do not create new MPI opaque objects. Some classes have a member function Create() for this purpose.

**Example 16.2** In the following code fragment, the test will return true and the message will be sent to cout.

```

1 void foo()
2 {
3     MPI::Intracomm bar;
4
5     if (bar == MPI::COMM_NULL)
6         cout << "bar is MPI::COMM_NULL" << endl;
7 }

```

The destructor for each MPI user level object does *not* invoke the corresponding `MPI_*_FREE` function (if it exists).

*Rationale.* `MPI_*_FREE` functions are not automatically invoked for the following reasons:

1. Automatic destruction contradicts the shallow-copy semantics of the MPI classes.
2. The model put forth in MPI makes memory allocation and deallocation the responsibility of the user, not the implementation.
3. Calling `MPI_*_FREE` upon destruction could have unintended side effects, including triggering collective operations (this also affects the copy, assignment, and construction semantics). In the following example, we would want neither `foo_comm` nor `bar_comm` to automatically invoke `MPI_*_FREE` upon exit from the function.

```

23 void example_function()
24 {
25     MPI::Intracomm foo_comm(MPI::COMM_WORLD), bar_comm;
26     bar_comm = MPI::COMM_WORLD.Dup();
27     // rest of function
28 }

```

*(End of rationale.)*

**Copy / Assignment** The copy constructor and assignment operator are prototyped as follows:

```

34 { MPI::<CLASS>(const MPI::<CLASS>& data) (binding deprecated, see Section 15.2) }
35
36 { MPI::<CLASS>& MPI::<CLASS>::operator=(const MPI::<CLASS>& data) (binding
37     deprecated, see Section 15.2) }

```

In terms of copying and assignment, opaque MPI user level objects behave like handles. Copy constructors perform handle-based (shallow) copies. `MPI::Status` objects are exceptions to this rule. These objects perform deep copies for assignment and copy construction.

*Advice to implementors.* Each MPI user level object is likely to contain, by value or by reference, implementation-dependent state information. The assignment and copying of MPI object handles may simply copy this value (or reference). *(End of advice to implementors.)*

**Example 16.3** Example using assignment operator. In this example, `MPI::Intracomm::Dup()` is *not* called for `foo_comm`. The object `foo_comm` is simply an alias for `MPI::COMM_WORLD`. But `bar_comm` is created with a call to `MPI::Intracomm::Dup()` and is therefore a different communicator than `foo_comm` (and thus different from `MPI::COMM_WORLD`). `baz_comm` becomes an alias for `bar_comm`. If one of `bar_comm` or `baz_comm` is freed with `MPI_COMM_FREE` it will be set to `MPI::COMM_NULL`. The state of the other handle will be undefined — it will be invalid, but not necessarily set to `MPI::COMM_NULL`.

```
MPI::Intracomm foo_comm, bar_comm, baz_comm;

foo_comm = MPI::COMM_WORLD;
bar_comm = MPI::COMM_WORLD.Dup();
baz_comm = bar_comm;
```

**Comparison** The comparison operators are prototyped as follows:

```
{bool MPI::<CLASS>::operator==(const MPI::<CLASS>& data) const(binding
    deprecated, see Section 15.2) }

{bool MPI::<CLASS>::operator!=(const MPI::<CLASS>& data) const(binding
    deprecated, see Section 15.2) }
```

The member function `operator==()` returns `true` only when the handles reference the same internal MPI object, `false` otherwise. `operator!=()` returns the boolean complement of `operator==()`. However, since the `Status` class is not a handle to an underlying MPI object, it does not make sense to compare `Status` instances. Therefore, the `operator==()` and `operator!=()` functions are not defined on the `Status` class.

**Constants** Constants are singleton objects and are declared `const`. Note that not all globally defined MPI objects are constant. For example, `MPI::COMM_WORLD` and `MPI::COMM_SELF` are not `const`.

### 16.1.6 C++ Datatypes

Table 16.1 lists all of the C++ predefined MPI datatypes and their corresponding C and C++ datatypes, Table 16.2 lists all of the Fortran predefined MPI datatypes and their corresponding Fortran 77 datatypes. Table 16.3 lists the C++ names for all other MPI datatypes.

`MPI::BYTE` and `MPI::PACKED` conform to the same restrictions as `MPI_BYTE` and `MPI_PACKED`, listed in Sections 3.2.2 on page 29 and Sections 4.2 on page 139, respectively.

The following table defines groups of MPI predefined datatypes:

C integer:	<code>MPI::INT</code> , <code>MPI::LONG</code> , <code>MPI::SHORT</code> , <code>MPI::UNSIGNED_SHORT</code> , <code>MPI::UNSIGNED</code> , <code>MPI::UNSIGNED_LONG</code> , <code>MPI::_LONG_LONG</code> , <code>MPI::UNSIGNED_LONG_LONG</code> , <code>MPI::SIGNED_CHAR</code> , <code>MPI::UNSIGNED_CHAR</code>
Fortran integer:	<code>MPI::INTEGER</code> and handles returned from

MPI datatype	C datatype	C++ datatype
MPI::CHAR	char	char
MPI::SHORT	signed short	signed short
MPI::INT	signed int	signed int
MPI::LONG	signed long	signed long
MPI::LONG_LONG	signed long long	signed long long
MPI::SIGNED_CHAR	signed char	signed char
MPI::UNSIGNED_CHAR	unsigned char	unsigned char
MPI::UNSIGNED_SHORT	unsigned short	unsigned short
MPI::UNSIGNED	unsigned int	unsigned int
MPI::UNSIGNED_LONG	unsigned long	unsigned long int
MPI::UNSIGNED_LONG_LONG	unsigned long long	unsigned long long
MPI::FLOAT	float	float
MPI::DOUBLE	double	double
MPI::LONG_DOUBLE	long double	long double
MPI::BOOL		bool
MPI::COMPLEX		Complex<float>
MPI::DOUBLE_COMPLEX		Complex<double>
MPI::LONG_DOUBLE_COMPLEX		Complex<long double>
MPI::WCHAR	wchar_t	wchar_t
MPI::BYTE		
MPI::PACKED		

Table 16.1: C++ names for the MPI C and C++ predefined datatypes, and their corresponding C/C++ datatypes.

MPI datatype	Fortran datatype
MPI::INTEGER	INTEGER
MPI::REAL	REAL
MPI::DOUBLE_PRECISION	DOUBLE PRECISION
MPI::F_COMPLEX	COMPLEX
MPI::LOGICAL	LOGICAL
MPI::CHARACTER	CHARACTER(1)
MPI::BYTE	
MPI::PACKED	

Table 16.2: C++ names for the MPI Fortran predefined datatypes, and their corresponding Fortran 77 datatypes.



MPI datatype	Description
MPI::FLOAT_INT	C/C++ reduction type
MPI::DOUBLE_INT	C/C++ reduction type
MPI::LONG_INT	C/C++ reduction type
MPI::TWOINT	C/C++ reduction type
MPI::SHORT_INT	C/C++ reduction type
MPI::LONG_DOUBLE_INT	C/C++ reduction type
MPI::TWOREAL	Fortran reduction type
MPI::TWODOUBLE_PRECISION	Fortran reduction type
MPI::TWOINTEGER	Fortran reduction type
MPI::F_DOUBLE_COMPLEX	Optional Fortran type
MPI::INTEGER1	Explicit size type
MPI::INTEGER2	Explicit size type
MPI::INTEGER4	Explicit size type
MPI::INTEGER8	Explicit size type
MPI::INTEGER16	Explicit size type
MPI::REAL2	Explicit size type
MPI::REAL4	Explicit size type
MPI::REAL8	Explicit size type
MPI::REAL16	Explicit size type
MPI::F_COMPLEX4	Explicit size type
MPI::F_COMPLEX8	Explicit size type
MPI::F_COMPLEX16	Explicit size type
MPI::F_COMPLEX32	Explicit size type

Table 16.3: C++ names for other MPI datatypes. Implementations may also define other optional types (e.g., `MPI::INTEGER8`).

```

1      MPI::Datatype::Create_f90_integer,
2      and if available: MPI::INTEGER1,
3      MPI::INTEGER2, MPI::INTEGER4,
4      MPI::INTEGER8, MPI::INTEGER16
5      Floating point: MPI::FLOAT, MPI::DOUBLE, MPI::REAL,
6                      MPI::DOUBLE_PRECISION,
7                      MPI::LONG_DOUBLE
8      and handles returned from
9      MPI::Datatype::Create_f90_real,
10     and if available: MPI::REAL2,
11     MPI::REAL4, MPI::REAL8, MPI::REAL16
12     Logical: MPI::LOGICAL, MPI::BOOL
13     Complex: MPI::F_COMPLEX, MPI::COMPLEX,
14             MPI::F_DOUBLE_COMPLEX,
15             MPI::DOUBLE_COMPLEX,
16             MPI::LONG_DOUBLE_COMPLEX
17     and handles returned from
18     MPI::Datatype::Create_f90_complex,
19     and if available: MPI::F_DOUBLE_COMPLEX,
20     MPI::F_COMPLEX4, MPI::F_COMPLEX8,
21     MPI::F_COMPLEX16, MPI::F_COMPLEX32
22     Byte: MPI::BYTE

```

Valid datatypes for each reduction operation are specified below in terms of the groups defined above.

Op	Allowed Types
MPI::MAX, MPI::MIN	C integer, Fortran integer, Floating point
MPI::SUM, MPI::PROD	C integer, Fortran integer, Floating point, Complex
MPI::LAND, MPI::LOR, MPI::LXOR	C integer, Logical
MPI::BAND, MPI::BOR, MPI::BXOR	C integer, Fortran integer, Byte

MPI::MINLOC and MPI::MAXLOC perform just as their C and Fortran counterparts; see Section 5.9.4 on page 191.

### 16.1.7 Communicators

The `MPI::Comm` class hierarchy makes explicit the different kinds of communicators implicitly defined by MPI and allows them to be strongly typed. Since the original design of MPI defined only one type of handle for all types of communicators, the following clarifications are provided for the C++ design.

**Types of communicators** There are six different types of communicators: `MPI::Comm`, `MPI::Intercomm`, `MPI::Intracomm`, `MPI::Cartcomm`, `MPI::Graphcomm`, and `MPI::Distgraphcomm`. `MPI::Comm` is the abstract base communicator class, encapsulating the functionality common to all MPI communicators. `MPI::Intercomm` and `MPI::Intracomm` are derived from `MPI::Comm`. `MPI::Cartcomm`, `MPI::Graphcomm`, and `MPI::Distgraphcomm` are derived from `MPI::Intracomm`.

*Advice to users.* Initializing a derived class with an instance of a base class is not [legal]valid in C++. For instance, it is not [legal]valid to initialize a Cartcomm from an Intracomm. Moreover, because MPI::Comm is an abstract base class, it is non-instantiable, so that it is not possible to have an object of class MPI::Comm. However, it is possible to have a reference or a pointer to an MPI::Comm.

**Example 16.4** The following code is erroneous.

```
Intracomm intra = MPI::COMM_WORLD.Dup();
Cartcomm cart(intra);           // This is erroneous
```

*(End of advice to users.)*

MPI::COMM\_NULL The specific type of MPI::COMM\_NULL is implementation dependent. MPI::COMM\_NULL must be able to be used in comparisons and initializations with all types of communicators. MPI::COMM\_NULL must also be able to be passed to a function that expects a communicator argument in the parameter list (provided that MPI::COMM\_NULL is an allowed value for the communicator argument).

*Rationale.* There are several possibilities for implementation of MPI::COMM\_NULL. Specifying its required behavior, rather than its realization, provides maximum flexibility to implementors. *(End of rationale.)*

**Example 16.5** The following example demonstrates the behavior of assignment and comparison using MPI::COMM\_NULL.

```
MPI::Intercomm comm;
comm = MPI::COMM_NULL;           // assign with COMM_NULL
if (comm == MPI::COMM_NULL)      // true
    cout << "comm is NULL" << endl;
if (MPI::COMM_NULL == comm)      // note -- a different function!
    cout << "comm is still NULL" << endl;
```

Dup() is not defined as a member function of MPI::Comm, but it is defined for the derived classes of MPI::Comm. Dup() is not virtual and it returns its OUT parameter by value.

MPI::Comm::Clone() The C++ language interface for MPI includes a new function Clone(). MPI::Comm::Clone() is a pure virtual function. For the derived communicator classes, Clone() behaves like Dup() except that it returns a new object by reference. The Clone() functions are prototyped as follows:

```
Comm& Comm::Clone() const = 0

Intracomm& Intracomm::Clone() const

Intercomm& Intercomm::Clone() const

Cartcomm& Cartcomm::Clone() const

Graphcomm& Graphcomm::Clone() const
```

```
1 Distgraphcomm& Distgraphcomm::Clone() const
```

3       *Rationale.* `Clone()` provides the “virtual dup” functionality that is expected by C++  
4 programmers and library writers. Since `Clone()` returns a new object by reference,  
5 users are responsible for eventually deleting the object. A new name is introduced  
6 rather than changing the functionality of `Dup()`. (*End of rationale.*)

8       *Advice to implementors.* Within their class declarations, prototypes for `Clone()` and  
9 `Dup()` would look like the following:

```
11 namespace MPI {
12     class Comm {
13         virtual Comm& Clone() const = 0;
14     };
15     class Intracomm : public Comm {
16         Intracomm Dup() const { ... };
17         virtual Intracomm& Clone() const { ... };
18     };
19     class Intercomm : public Comm {
20         Intercomm Dup() const { ... };
21         virtual Intercomm& Clone() const { ... };
22     };
23     // Cartcomm, Graphcomm,
24     // and Distgraphcomm are similarly defined
25 };
```

26       (*End of advice to implementors.*)

### 28 16.1.8 Exceptions

30 The C++ language interface for MPI includes the predefined error handler  
31 `MPI::ERRORS_THROW_EXCEPTIONS` for use with the `Set_errhandler()` member functions.  
32 `MPI::ERRORS_THROW_EXCEPTIONS` can only be set or retrieved by C++ functions. If a  
33 non-C++ program causes an error that invokes the `MPI::ERRORS_THROW_EXCEPTIONS` error  
34 handler, the exception will pass up the calling stack until C++ code can catch it. If there  
35 is no C++ code to catch it, the behavior is undefined. In a multi-threaded environment  
36 or if a nonblocking MPI call throws an exception while making progress in the background,  
37 the behavior is implementation dependent.

38 The error handler `MPI::ERRORS_THROW_EXCEPTIONS` causes an `MPI::Exception` to be  
39 thrown for any MPI result code other than `MPI::SUCCESS`. The public interface to  
40 `MPI::Exception` class is defined as follows:

```
41
42 namespace MPI {
43     class Exception {
44     public:
45
46         Exception(int error_code);
47
48         int Get_error_code() const;
```

```

int Get_error_class() const;
const char *Get_error_string() const;
};
};

```

*Advice to implementors.*

The exception will be thrown within the body of `MPI::ERRORS_THROW_EXCEPTIONS`. It is expected that control will be returned to the user when the exception is thrown. Some MPI functions specify certain return information in their parameters in the case of an error and `MPI_ERRORS_RETURN` is specified. The same type of return information must be provided when exceptions are thrown.

For example, `MPI_WAITALL` puts an error code for each request in the corresponding entry in the status array and returns `MPI_ERR_IN_STATUS`. When using `MPI::ERRORS_THROW_EXCEPTIONS`, it is expected that the error codes in the status array will be set appropriately before the exception is thrown.

*(End of advice to implementors.)*

### 16.1.9 Mixed-Language Operability

The C++ language interface provides functions listed below for mixed-language operability. These functions provide for a seamless transition between C and C++. For the case where the C++ class corresponding to `<CLASS>` has derived classes, functions are also provided for converting between the derived classes and the C `MPI_<CLASS>`.

```
MPI::<CLASS>& MPI::<CLASS>::operator=(const MPI_<CLASS>& data)
```

```
MPI::<CLASS>(const MPI_<CLASS>& data)
```

```
MPI::<CLASS>::operator MPI_<CLASS>() const
```

These functions are discussed in Section [16.3.4](#).

### 16.1.10 Profiling

This section specifies the requirements of a C++ profiling interface to MPI.

*Advice to implementors.* Since the main goal of profiling is to intercept function calls from user code, it is the implementor's decision how to layer the underlying implementation to allow function calls to be intercepted and profiled. If an implementation of the MPI C++ bindings is layered on top of MPI bindings in another language (such as C), or if the C++ bindings are layered on top of a profiling interface in another language, no extra profiling interface is necessary because the underlying MPI implementation already meets the MPI profiling interface requirements.

Native C++ MPI implementations that do not have access to other profiling interfaces must implement an interface that meets the requirements outlined in this section.

High-quality implementations can implement the interface outlined in this section in order to promote portable C++ profiling libraries. Implementors may wish to provide an option whether to build the C++ profiling interface or not; C++ implementations that are already layered on top of bindings in another language or another profiling

interface will have to insert a third layer to implement the C++ profiling interface.  
(*End of advice to implementors.*)

To meet the requirements of the C++ MPI profiling interface, an implementation of the MPI functions *must*:

1. Provide a mechanism through which all of the MPI defined functions may be accessed with a name shift. Thus all of the MPI functions (which normally start with the prefix “MPI:.”) should also be accessible with the prefix “PMPI:.”
2. Ensure that those MPI functions which are not replaced may still be linked into an executable image without causing name clashes.
3. Document the implementation of different language bindings of the MPI interface if they are layered on top of each other, so that profiler developer knows whether they must implement the profile interface for each binding, or can economize by implementing it only for the lowest level routines.
4. Where the implementation of different language bindings is done through a layered approach (e.g., the C++ binding is a set of “wrapper” functions which call the C implementation), ensure that these wrapper functions are separable from the rest of the library.

This is necessary to allow a separate profiling library to be correctly implemented, since (at least with Unix linker semantics) the profiling library must contain these wrapper functions if it is to perform as expected. This requirement allows the author of the profiling library to extract these functions from the original MPI library and add them into the profiling library without bringing along any other unnecessary code.

5. Provide a no-op routine MPI::Pcontrol in the MPI library.

*Advice to implementors.* There are (at least) two apparent options for implementing the C++ profiling interface: inheritance or caching. An inheritance-based approach may not be attractive because it may require a virtual inheritance implementation of the communicator classes. Thus, it is most likely that implementors will cache PMPI objects on their corresponding MPI objects. The caching scheme is outlined below.

The “real” entry points to each routine can be provided within a namespace PMPI. The non-profiling version can then be provided within a namespace MPI.

Caching instances of PMPI objects in the MPI handles provides the “has a” relationship that is necessary to implement the profiling scheme.

Each instance of an MPI object simply “wraps up” an instance of a PMPI object. MPI objects can then perform profiling actions before invoking the corresponding function in their internal PMPI object.

The key to making the profiling work by simply re-linking programs is by having a header file that *declares* all the MPI functions. The functions must be *defined* elsewhere, and compiled into a library. MPI constants should be declared *extern* in the MPI namespace. For example, the following is an excerpt from a sample `mpi.h` file:

**Example 16.6** Sample `mpi.h` file.

```

namespace PMPI {
    class Comm {
    public:
        int Get_size() const;
    };
    // etc.
};

namespace MPI {
public:
    class Comm {
    public:
        int Get_size() const;

    private:
        PMPI::Comm pmpi_comm;
    };
};

```

Note that all constructors, the assignment operator, and the destructor in the `MPI` class will need to initialize/destroy the internal `PMPI` object as appropriate.

The definitions of the functions must be in separate object files; the `PMPI` class member functions and the non-profiling versions of the `MPI` class member functions can be compiled into `libmpi.a`, while the profiling versions can be compiled into `libpmpi.a`. Note that the `PMPI` class member functions and the `MPI` constants must be in different object files than the non-profiling `MPI` class member functions in the `libmpi.a` library to prevent multiple definitions of `MPI` class member function names when linking both `libmpi.a` and `libpmpi.a`. For example:

**Example 16.7** `pmapi.cc`, to be compiled into `libmpi.a`.

```

int PMPI::Comm::Get_size() const
{
    // Implementation of MPI_COMM_SIZE
}

```

**Example 16.8** `constants.cc`, to be compiled into `libmpi.a`.

```

const MPI::Intracomm MPI::COMM_WORLD;

```

**Example 16.9** `mpi_no_profile.cc`, to be compiled into `libpmpi.a`.

```

int MPI::Comm::Get_size() const
{
    return pmpi_comm.Get_size();
}

```

**Example 16.10** `mpi_profile.cc`, to be compiled into `libmpi.a`.

```
int MPI::Comm::Get_size() const
{
    // Do profiling stuff
    int ret = pmpi_comm.Get_size();
    // More profiling stuff
    return ret;
}
```

*(End of advice to implementors.)*

## 16.2 Fortran Support

### 16.2.1 Overview

The Fortran [MPI-2]MPI language bindings have been designed to be compatible with the Fortran 90 standard [(and later)]with additional features from Fortran 2003 and Fortran 2008 [39] + TR 29113 [41]. [These bindings are in most cases compatible with Fortran 77, implicit-style interfaces.]

*Rationale.* Fortran 90 contains numerous features designed to make it a more “modern” language than Fortran 77. It seems natural that [MPI]MPI should be able to take advantage of these new features with a set of bindings tailored to Fortran 90. [MPI does not (yet) use many of these features because of a number of technical difficulties.]In Fortran 2008 + TR 29113, the major new language features used are the **ASYNCHRONOUS** attribute to protect nonblocking MPI operations, and assumed-type and assumed-rank dummy arguments for choice buffer arguments. Further requirements for compiler support are listed in Section 16.2.7 on page 657. *(End of rationale.)*

MPI defines [two levels]three methods of Fortran support[, described in Sections 16.2.4 and 16.2.3. In the rest of this section, “Fortran” and “Fortran 90” shall refer to “Fortran 90” and its successors, unless qualified.]:

1. **USE mpi\_f08:** This method is described in Section 16.2.2 and requires compile-time argument checking with unique MPI handle types and provides techniques to fully solve the optimization problems with nonblocking calls. This is the only Fortran support method that is consistent with the Fortran standard (Fortran 2008 + TR 29113 and later). This method is highly recommended for all MPI applications.
2. **[Extended Fortran Support]USE mpi:** [An implementation with this level of Fortran support provides Basic Fortran Support plus additional features that specifically support Fortran 90, as]This method is described in Section 16.2.3 and requires compile-time argument checking. Handles are defined as **INTEGER**. This Fortran support method is inconsistent with the Fortran standard, and its use is therefore not recommended. It exists only for backwards compatibility.
3. **[Basic Fortran Support]INCLUDE 'mpif.h':** [An implementation with this level of Fortran support provides the original Fortran bindings specified in MPI-1, with small



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additional requirements specified] This method is described in Section 16.2.4. The use of the include file `mpif.h` is strongly discouraged starting with MPI-3.0, because this method neither guarantees compile-time argument checking nor provides sufficient techniques to solve the optimization problems with nonblocking calls, and is therefore inconsistent with the Fortran standard. It exists only for backwards compatibility with legacy MPI applications.

[A compliant MPI-2] Compliant MPI-3 implementations providing a Fortran interface must provide [Extended Fortran Support unless the target compiler does not support modules or KIND-parameterized types] [all three Fortran support methods. ] provide one or both of the following:

- The `USE mpi_f08` Fortran support method.
- The `USE mpi` and `INCLUDE 'mpif.h'` Fortran support methods.

Section 16.2.6 on page 654 describes restrictions if the compiler does not support all the needed features.

Application[s] subroutines and functions may use either [the `mpi`] one of the modules or the `mpif.h` include file. An implementation may require the use of one of the modules to prevent type mismatch errors [ (see below)].

*Advice to users.* [It is recommended to use the `mpi` module even if it is not necessary to use it to avoid type mismatch errors ] Users are advised to utilize one of the MPI modules even if `mpif.h` enforces type checking on a particular system. Using a module provides several potential advantages over using an include file; the `mpi_f08` module offers the most [advantages] robust and complete Fortran support. (*End of advice to users.*)

[It] In a single application, it must be possible to link together routines [some of which `USE mpi` and others of which `INCLUDE 'mpif.h'`. ] which `USE mpi_f08`, `USE mpi`, and `INCLUDE mpif.h`.

The `INTEGER` compile-time constant `MPI_SUBARRAYS_SUPPORTED` is set to `.TRUE.` if all buffer choice arguments are defined in explicit interfaces with assumed-type and assumed-rank [41]; otherwise it is set to `.FALSE.`. The `INTEGER` compile-time constant `MPI_ASYNC_PROTECTS_NONBLOCKING` is set to `.TRUE.` if the `ASYNCHRONOUS` attribute was added to the choice buffer arguments of all nonblocking interfaces **and** the underlying Fortran compiler supports the `ASYNCHRONOUS` attribute for MPI communication (as part of TR 29113), otherwise it is set to `.FALSE.`. These constants exist with each Fortran support method, but not in the C/C++ header files. The values may be different for each Fortran support method. All other constants and the integer values of handles must be the same for each Fortran support method.

Section 16.2.2 through 16.2.4 define the Fortran support methods. The Fortran interfaces of each MPI routine are shorthands. Section 16.2.5 defines the corresponding full interface specification together with the used linker names and implications for the profiling interface. Section 16.2.6 the implementation of the MPI routines for different versions of the Fortran standard. Section 16.2.7 summarizes major requirements for valid MPI-3.0 implementations with Fortran support. Section 16.2.8 and Section 16.2.9 describe additional functionality that is part of the Fortran support. `MPI_F_SYNC_REG` is needed for one of the methods to prevent register optimization problems. A [new] set

of functions[ to] provides additional support for Fortran intrinsic numeric types, including parameterized types: `MPI_SIZEOF`, `MPI_TYPE_MATCH_SIZE`, `MPI_TYPE_CREATE_F90_INTEGER`, `MPI_TYPE_CREATE_F90_REAL` and `MPI_TYPE_CREATE_F90_COMPLEX`. [Parameterized]In the context of MPI, parameterized types are Fortran intrinsic types which are specified using `KIND` type parameters. [These routines are described in detail in Section 16.2.9.]Sections 16.2.10 through 16.2.19 give an overview and details on known problems when using Fortran together with MPI; Section 16.2.20 compares the Fortran problems with those in C.

## 16.2.2 Fortran Support Through the `mpi_f08` Module

An MPI implementation providing a Fortran interface must provide a module named `mpi_f08` that can be used in a Fortran program. Section 16.2.6 on page 654 describes restrictions if the compiler does not support all the needed features. Within all MPI function specifications, the first of the set of two Fortran routine interface specifications is provided by this module. This module must:

- Define all named MPI constants.
- Declare MPI functions that return a value.
- Provide explicit interfaces according to the Fortran routine interface specifications. This module therefore guarantees compile-time argument checking for all arguments which are not `TYPE(*)`, with the following exception:

Only one Fortran interface is defined for functions that are deprecated as of MPI-3.0. This interface must be provided as an explicit interface according to the rules defined for the `mpi` module, see Section 16.2.3 on page 644.

*Advice to users.* It is strongly recommended that developers substitute calls to deprecated routines when upgrading from `mpif.h` or the `mpi` module to the `mpi_f08` module. (*End of advice to users.*)

- Define all MPI handles with uniquely named handle types (instead of `INTEGER` handles, as in the `mpi` module). This is reflected in the first Fortran binding in each MPI function definition throughout this document (except for the deprecated routines).
- Use the `ASYNCHRONOUS` attribute to protect the buffers of nonblocking operations, and set the `INTEGER` compile-time constant `MPI_ASYNC_PROTECTS_NONBLOCKING` to `.TRUE.` if the underlying Fortran compiler supports the `ASYNCHRONOUS` attribute for MPI communication (as part of TR 29113). See Section 16.2.6 on page 654 for older compiler versions.
- Set the `INTEGER` compile-time constant `MPI_SUBARRAYS_SUPPORTED` to `.TRUE.` and declare choice buffers using the Fortran 2008 TR 29113 feature assumed-type and assumed-rank, i.e., `TYPE(*)`, `DIMENSION(..)` in all nonblocking, split collective and persistent communication routines, if the underlying Fortran compiler supports it. With this, non-contiguous sub-arrays can be used as buffers in nonblocking routines.

*Rationale.* In all blocking routines, i.e., if the choice-buffer is not declared as `ASYNCHRONOUS`, the TR 29113 feature is not needed for the support of non-contiguous buffers because the compiler can pass the buffer by in-and-out-copy through a contiguous scratch array. (*End of rationale.*)

- Set the `MPI_SUBARRAYS_SUPPORTED` compile-time constant to `.FALSE.` and declare choice buffers with a compiler-dependent mechanism that overrides type checking if the underlying Fortran compiler does not support the Fortran 2008 TR 29113 assumed-type and assumed-rank notation. In this case, the use of non-contiguous sub-arrays as buffers in nonblocking calls may be invalid. See Section 16.2.6 on page 654 for details.

- Declare each argument with an `INTENT` of `IN`, `OUT`, or `INOUT` as defined in this standard.

*Rationale.* For these definitions in the `mpi_f08` bindings, in most cases, `INTENT(IN)` is used if the C interface uses call-by-value. For all buffer arguments and for `OUT` and `INOUT` dummy arguments that allow one of the non-ordinary Fortran constants (see `MPI_BOTTOM`, etc. in Section 2.5.4 on page 15) as input, an `INTENT` is not specified. (*End of rationale.*)

*Advice to users.* If a dummy argument is declared with `INTENT(OUT)`, then the Fortran standard stipulates that the actual argument becomes undefined upon invocation of the MPI routine, i.e., it may be overwritten by some other values, e.g. zeros; according to [39], 12.5.2.4 Ordinary dummy variables, Paragraph 17: “If a dummy argument has `INTENT(OUT)`, the actual argument becomes undefined at the time the association is established, except [...]”. For example, if the dummy argument is an assumed-size array and the actual argument is a strided array, the call may be implemented with copy-in and copy-out of the argument. In the case of `INTENT(OUT)` the copy-in may be suppressed by the optimization and the routine starts execution using an array of undefined values. If the routine stores fewer elements into the dummy argument than is provided in the actual argument, then the remaining locations are overwritten with these undefined values. See also both advices to implementors in Section 16.2.3 on page 644. (*End of advice to users.*)

- Declare all `ierror` output arguments as `OPTIONAL`, except for user-defined callback functions (e.g., `COMM_COPY_ATTR_FUNCTION`) and predefined callbacks (e.g., `MPI_NULL_COPY_FN`).

*Rationale.* For user-defined callback functions (e.g., `COMM_COPY_ATTR_FUNCTION`) and their predefined callbacks (e.g., `MPI_NULL_COPY_FN`), the `ierror` argument is not optional. The MPI library must always call these routines with an actual `ierror` argument. Therefore, these user-defined functions need not check whether the MPI library calls these routines with or without an actual `ierror` output argument. (*End of rationale.*)

The MPI Fortran bindings in the `mpi_f08` module are designed based on the Fortran 2008 standard [39] together with the Technical Report “TR 29113 Further Interoperability with C” [41] of the ISO/IEC JTC1/SC22/WG5 (Fortran) working group.

*Rationale.* The features in TR 29113 on further interoperability with C were decided on by ISO/IEC JTC1/SC22/WG5 and designed by PL22.3 (formerly J3) to support a higher level of integration between Fortran-specific features and C than was provided in the Fortran 2008 standard; part of this design is based on requirements from the MPI Forum to support MPI-3.0. According to [40] page iv, last paragraph, “it is the intention of ISO/IEC JTC1/SC22/WG5 that the semantics and syntax specified by this technical report be included in the next revision of the Fortran International Standard without change unless experience in the implementation and use of this feature identifies errors that need to be corrected, or changes are needed to achieve proper integration, in which case every reasonable effort will be made to minimize the impact of such changes on existing implementations”.

The TR 29113 contains the following language features that are needed for the MPI bindings in the `mpi_f08` module: assumed-type and assumed-rank. It is important that any possible actual argument can be used for such dummy arguments, e.g., scalars, arrays, assumed-shape arrays, assumed-size arrays, allocatable arrays, and with any element type, e.g., `REAL`, `CHARACTER*5`, `CHARACTER(*)`, sequence derived types, or `BIND(C)` derived types. Especially for backward compatibility reasons, it is important that any possible actual argument in an implicit interface implementation of a choice buffer dummy argument (e.g., with `mpif.h` without argument-checking) can be used in an implementation with assumed-type and assumed-rank argument in an explicit interface (e.g., with the `mpi_f08` module).

The `INTERFACE` construct in combination with `BIND(C)` allows the implementation of the Fortran `mpi_f08` interface with a single set of portable wrapper routines written in C, which supports all desired features in the `mpi_f08` interface. TR 29113 also has a provision for `OPTIONAL` arguments in `BIND(C)` interfaces.

A further feature useful for MPI is the extension of the semantics of the `ASYNCHRONOUS` attribute: In F2003 and F2008, this attribute could be used only to protect buffers of Fortran asynchronous I/O. With TR29113, this attribute now also covers asynchronous communication occurring within library routines written in C.

The MPI Forum hereby wishes to acknowledge this important effort by the Fortran PL22.3 and WG5 committee. (*End of rationale.*)

### 16.2.3 [Extended] Fortran Support Through the `mpi` Module

[Implementations with Extended Fortran support must provide:

1. An `mpi` module
2. (List item was moved to another location, and modified) A new set of functions to provide additional support for Fortran intrinsic numeric types, including parameterized types: `MPI_SIZEOF`, `MPI_TYPE_MATCH_SIZE`, `MPI_TYPE_CREATE_F90_INTEGER`, `MPI_TYPE_CREATE_F90_REAL` and

`MPI_TYPE_CREATE_F90_COMPLEX`. Parameterized types are Fortran intrinsic types which are specified using `KIND` type parameters. These routines are described in detail in Section 16.2.9.

Additionally, high-quality implementations should provide a mechanism to prevent fatal type mismatch errors for MPI routines with choice arguments.

### The `mpi` Module

]

An MPI implementation **providing a Fortran interface** must provide a module named `mpi` that can be used in a Fortran [90] program. Within all MPI function specifications, the second of the set of two Fortran routine interface specifications is provided by this module. This module must:

- Define all named MPI constants
- Declare MPI functions that return a value.
- Provide explicit interfaces according to the Fortran routine interface specifications. This module therefore guarantees compile-time argument checking and allows positional and keyword-based argument lists.
- Define all MPI handles as type `INTEGER`.
- Define all named handle types and the derived type `MPI_Status` that are used in the `mpi_f08` module.

*Rationale.* They are needed only when the application converts old-style `INTEGER` handles into new-style handles with a named type. (*End of rationale.*)

- A high quality MPI implementation may enhance the interface by using the `ASYNCHRONOUS` attribute in the same way as in the `mpi_f08` module if it is supported by the underlying compiler.
- Set the `INTEGER` compile-time constant `MPI_ASYNC_PROTECTS_NONBLOCKING` to `.TRUE.` if the `ASYNCHRONOUS` attribute is used in all nonblocking interfaces **and** the underlying Fortran compiler supports the `ASYNCHRONOUS` attribute for MPI communication (as part of TR 29113), otherwise to `.FALSE.`

*Advice to users.* For an MPI implementation that fully supports nonblocking calls with the `ASYNCHRONOUS` attribute for choice buffers, an existing MPI-2.2 application may fail to compile even if it compiled and executed with expected results with an MPI-2.2 implementation. One reason may be that the application uses ‘contiguous’ but not ‘simply contiguous’ `ASYNCHRONOUS` arrays as actual arguments for choice buffers of nonblocking routines, e.g., by using subscript triplets with stride one or specifying `(1:n)` for a whole dimension instead of using `(:)`. This should be fixed to fulfill the Fortran constraints for `ASYNCHRONOUS` dummy arguments. This is not considered a violation of backward compatibility because existing applications can not use the

ASYNCHRONOUS attribute to protect nonblocking calls. Onother reason may be that the application does not conform either to MPI-2.2, or to MPI-3.0, or to the Fortran standard, typically because the program forces the compiler to perform copyin/out for a choice buffer argument in a nonblocking MPI call. This is also not a violation of backward compatibility because the application itself is non-conforming. See Section 16.2.12 on page 671 for more details. (*End of advice to users.*)

- A high quality MPI implementation may enhance the interface by using `TYPE(*)`, `DIMENSION(..)` choice buffer dummy arguments instead of using non-standardized extensions such as `!$PRAGMA IGNORE_TKR` or a set of overloaded functions as described by M. Hennecke in [28], if the compiler supports this TR 29113 language feature. See Section 16.2.6 on page 654 for further details.
- Set the `INTEGER` compile-time constant `MPI_SUBARRAYS_SUPPORTED` to `.TRUE.` if all choice buffer arguments in all nonblocking, split collective and persistent communication routines are declared with `TYPE(*)`, `DIMENSION(..)`, otherwise set it to `.FALSE..` With `MPI_SUBARRAYS_SUPPORTED==.TRUE.`, non-contiguous sub-arrays can be used as buffers in nonblocking routines.
- Set the `MPI_SUBARRAYS_SUPPORTED` compile-time constant to `.FALSE.` and declare choice buffers with a compiler-dependent mechanism that overrides type checking if the underlying Fortran compiler does not support the TR 29113 assumed-type and assumed-rank features. In this case, the use of non-contiguous sub-arrays in nonblocking calls may be disallowed. See Section 16.2.6 on page 654 for details.

An MPI implementation may provide other features in the `mpi` module [other features] that enhance the usability of MPI while maintaining adherence to the standard. For example, it may[:] provide `INTENT` information in these interface blocks. [

- Provide interfaces for all or for a subset of MPI routines.
- Provide `INTENT` information in these interface blocks.

]

*Advice to implementors.* The appropriate `INTENT` may be different from what is given in the MPI [generic interface]language-neutral bindings. Implementations must choose `INTENT` so that the function adheres to the MPI standard, e.g., by defining the `INTENT` as provided in the `mpi_f08` bindings. (*End of advice to implementors.*)

*Rationale.* The intent given by the MPI generic interface is not precisely defined and does not in all cases correspond to the correct Fortran `INTENT`. For instance, receiving into a buffer specified by a datatype with absolute addresses may require associating `MPI_BOTTOM` with a dummy `OUT` argument. Moreover, “constants” such as `MPI_BOTTOM` and `MPI_STATUS_IGNORE` are not constants as defined by Fortran, but “special addresses” used in a nonstandard way. Finally, the MPI-1 generic intent [is]was changed in several places [by]in MPI-2. For instance, `MPI_IN_PLACE` changes the [sense]intent of an `OUT` argument to be `INOUT`. (*End of rationale.*)



*Advice to implementors.* The Fortran 2008 standard illustrates in its Note 5.17 that “*INTENT(OUT)* means that the value of the argument after invoking the procedure is entirely the result of executing that procedure. If an argument should retain its value rather than being redefined, *INTENT(INOUT)* should be used rather than *INTENT(OUT)*, even if there is no explicit reference to the value of the dummy argument. Furthermore, *INTENT(INOUT)* is not equivalent to omitting the *INTENT* attribute, because *INTENT(INOUT)* always requires that the associated actual argument is *definable*”. Applications that include `mpif.h` may not expect that *INTENT(OUT)* is used. In particular, output array arguments are expected to keep their content as long as the MPI routine does not modify them. To keep this behavior, it is recommended that implementations not use *INTENT(OUT)* in the `mpi` module and the `mpif.h` include file, even though *INTENT(OUT)* is specified in an interface description of the `mpi_f08` module. (*End of advice to implementors.*)

[(Paragraph was moved to another location, and modified) Applications may use either the `mpi` module or the `mpif.h` include file. An implementation may require use of the module to prevent type mismatch errors (see below).

*Advice to users.* It is recommended to use the `mpi` module even if it is not necessary to use it to avoid type mismatch errors on a particular system. Using a module provides several potential advantages over using an include file. (*End of advice to users.*)

It must be possible to link together routines some of which `USE mpi` and others of which `INCLUDE mpif.h`. ]

[

#### No Type Mismatch Problems for Subroutines with Choice Arguments

A high-quality MPI implementation should provide a mechanism to ensure that MPI choice arguments do not cause fatal compile-time or run-time errors due to type mismatch. An MPI implementation may require applications to use the `mpi` module, or require that it be compiled with a particular compiler flag, in order to avoid type mismatch problems.

*Advice to implementors.* In the case where the compiler does not generate errors, nothing needs to be done to the existing interface. In the case where the compiler may generate errors, a set of overloaded functions may be used. See the paper of M. Hennecke [28]. Even if the compiler does not generate errors, explicit interfaces for all routines would be useful for detecting errors in the argument list. Also, explicit interfaces which give *INTENT* information can reduce the amount of copying for `BUF(*)` arguments. (*End of advice to implementors.*)

]

#### 16.2.4 [Basic ]Fortran Support Through the `mpif.h` Include File

The use of the `mpif.h` include file is strongly discouraged and may be deprecated in a future version of MPI.

[Because Fortran 90 is (for all practical purposes) a superset of Fortran 77, Fortran 90 (and future) programs can use the original Fortran interface. ][The following additional requirements are added:] An MPI implementation providing a Fortran interface must pro-

vide an include file named `mpif.h` that can be used in a Fortran program. Within all MPI function specifications, the second of the set of two Fortran routine interface specifications is supported by this include file. This include file must:

[

1. Implementations are required to provide the file `mpif.h`, as described in the original MPI-1 specification.
2. `mpif.h` must be valid and equivalent for both fixed- and free- source form.

]

- Define all named MPI constants.
- Declare MPI functions that return a value.
- Define all handles as `INTEGER`.
- Be valid and equivalent for both fixed and free source form.

For each MPI routine, an implementation can choose to use an implicit or explicit interface for the second Fortran binding (in deprecated routines, the first one may be omitted).

- Set the `INTEGER` compile-time constants `MPI_SUBARRAYS_SUPPORTED` and `MPI_ASYNC_PROTECTS_NONBLOCKING` according to the same rules as for the `mpi` module. In the case of implicit interfaces for choice buffer or nonblocking routines, the constants must be set to `.FALSE..`

*Advice to users.* Instead of using `mpif.h`, the use of the `mpi_f08` or `mpi` module is strongly encouraged for the following reasons:

- Most `mpif.h` implementations do not include compile-time argument checking.
- Therefore, too many bugs in MPI applications remain undetected at compile-time, such as:
  - Missing `ierror` as last argument in most Fortran bindings.
  - Declaration of a `status` as an `INTEGER` variable instead of an `INTEGER` array with size `MPI_STATUS_SIZE`.
  - Wrong argument positions; e.g., interchanging the `count` and `datatype` arguments.
  - Passing wrong MPI handles; e.g., passing a `datatype` instead of a `communicator`.
- The migration from `mpif.h` to the `mpi` module should be relatively straightforward (i.e., substituting `include 'mpif.h'` after an `implicit` statement by `use mpi` before such `implicit` statement) as long as the application syntax is correct.
- Migrating portable and correctly written applications to the `mpi` module is not expected to be difficult. No compile or runtime problems should occur because an `mpif.h` include file was always allowed to provide explicit Fortran interfaces.

*(End of advice to users.)*



*Rationale.* With MPI-3.0, the `mpif.h` include file was not deprecated in order to retain strong backward compatibility. Internally, `mpif.h` and the `mpi` module may be implemented so that the same (or similar) library implementation of the MPI routines can be used. (*End of rationale.*)

*Advice to implementors.* To make `mpif.h` compatible with both fixed- and free-source forms, to allow automatic inclusion by preprocessors, and to allow extended fixed-form line length, it is recommended that [requirement two]the requirement of usability in free and fixed source form applications be met by constructing `mpif.h` without any continuation lines. This should be possible because `mpif.h` [contains]may contain only declarations, and because common block declarations can be split among several lines. The argument names may need to be shortened to keep the SUBROUTINE statement within the allowed 72-6=66 characters, e.g.,

```
INTERFACE
SUBROUTINE PMPI_DIST_GRAPH_CREATE_ADJACENT(a,b,c,d,e,f,g,h,i,j,k)
... ! dummy argument declarations
```

This line has [65]65 characters and is the longest in MPI-3.0. [

**TODO: This is only checked for MPI-2.2. We have to check all new MPI-3.0 interfaces that they stay within these 66 characters. Otherwise the routine name should be shortened before the name is standardized. ]**

[If `mpif.h` contains also explicit interfaces with `BIND(C,NAME='...')` for providing `MPI_SUBARRAYS_SUPPORTED` and `MPI_ASYNC_PROTECTS_NONBLOCKING` equals `.TRUE.`, the linker routine name may need to be shortened. ]As long as the MPI standard contains routines with choice buffers and a name length and argument count that implies that a `BIND(C)` implementation would need to shorten their linker names in `mpif.h`, the `mpif.h` cannot set `MPI_SUBARRAYS_SUPPORTED` and `MPI_ASYNC_PROTECTS_NONBLOCKING` equals `.TRUE.`, because such shortening is invalid. For example, `MPI_FILE_WRITE_AT_ALL_BEGIN` with 6 arguments, may be defined:

```
INTERFACE MPI_FILE_WRITE_AT_ALL_BEGIN
SUBROUTINE MPI_X(a,b,c,d,e,f)BIND(C,NAME='MPI_File_write_at_all_begin_f')
... ! dummy argument declarations
```

This would need a line length of 73 characters, i.e., the C routine name [must]would need to be shortened by 7 characters to stay within the available 66 characters. **[TODO: Do we want to define these shortened routine names for `mpif.h`; this would help the tools people. ]**Note that the name `MPI_X` has no meaning for the compilation, and that this problem occurs only with routines with choice buffers implemented with the assumed-type and assumed-rank facility of TR 29113. To support Fortran 77 as well as Fortran 90 and later, it may be necessary to eliminate all comments from `mpif.h`. (*End of advice to implementors.*)

### 16.2.5 Interface Specifications, Linker Names and the Profiling Interface

The Fortran interface specifications of each MPI routine specifies the routine name that must be called by the application program, and the names and types of the dummy arguments

together with additional attributes. The rules for the linker names and its implications for the profiling interface are specified within this section. The linker name of a Fortran routine is defined as the name that a C routine would have if both routines would have the same name visible for the linker. A typical linker name of the Fortran routine `FOOfoo` is `foofoo__`. In the case of `BIND(C,NAME='...')`, the linker name is directly defined through the given string.

The following rules for linker names apply:

- With the Fortran `mpi_f08` module, if `MPI_SUBARRAYS_SUPPORTED` equals `.TRUE.:`  
The Fortran binding must use `BIND(C)` interfaces with an interface name identical to the language independent name, e.g., `MPI_SEND`. The linker name is a combination of the C name and an `_f08` suffix, e.g., `MPI_Send_f08`. Prototype example:

```
INTERFACE
SUBROUTINE MPI_Send(...) BIND(C,NAME='MPI_Send_f08')
```

- With the Fortran `mpi_f08` module, if `MPI_SUBARRAYS_SUPPORTED` equals `.FALSE.` (i.e., with a preliminary implementation of this module without TR 29113):  
The linker name of each routine is defined through the linker name mapping of the Fortran compiler for the name defined when subarrays are supported. For example, `MPI_Send_f08` may be mapped to `mpi_send_f08__`. Example:

```
INTERFACE MPI_Send
SUBROUTINE MPI_Send_f08(...)
```

- With the Fortran `mpi` module or `mpif.h` include file, if `MPI_SUBARRAYS_SUPPORTED` equals `.FALSE.:`  
The linker name of each routine is defined through the linker-name mapping of the Fortran compiler. For example, `MPI_SEND` may be mapped to `mpi_send__`. Example:

```
INTERFACE
SUBROUTINE MPI_SEND(...)
```

- With the Fortran `mpi` module or `mpif.h` include file, if `MPI_SUBARRAYS_SUPPORTED` equals `.TRUE.:`  
The Fortran binding must use `BIND(C)` interfaces with an interface name identical to the language independent name, e.g., `MPI_SEND`. The linker name is a combination of the C name and an `_f` suffix, e.g., `MPI_Send_f`. Prototype example:

```
INTERFACE
SUBROUTINE MPI_SEND(...) BIND(C,NAME='MPI_Send_f')
```

If the support of subarrays is different for the `mpi` module and the `mpif.h` include file, then both linker-name methods can be used in the same application. If the application also uses the `mpi_f08` module and was compiled with this module partially before and after the subarrays were supported, then all four interfaces are used within the same application.

*Rationale.* After a compiler provides the facilities from TR29113, i.e., `TYPE(*)`, `DIMENSION(...)`, it is possible to change the bindings within a Fortran support method to support subarrays and without recompiling the complete application. Of course, only recompiled routines can benefit from the added facilities. There is no binary compatibility conflict because each interface uses its own linker names and all interfaces use the same constants and type definitions. (*End of rationale.*)

A user-written or middleware profiling routine that is written according to the same binding rules will have the same linker name, and therefore, can interpose itself as the MPI library routine. The profiling routine can internally call the matching PMPI routine with any of its existing bindings, except for routines that have callback routine dummy arguments. In this case, the profiling software must use the same Fortran support method as used in the calling application program, because the C, `mpi_f08` and `mpi` callback prototypes are different.

*Advice to users.* This advice is mainly for tool writers. Even if an MPI library supports subarrays in all three Fortran support methods, a portable profiling layer should also provide the two interfaces for `MPI_SUBARRAYS_SUPPORTED==.FALSE.` to support older binary user routines that were compiled before TR29113 level support was achieved.

If a user application calls `MPI_SEND`, then the chosen Fortran support method together with the MPI implement decision about `MPI_SUBARRAYS_SUPPORTED` imply, to which linker name the compiler will translate this call, i.e., whether the application calls `mpi_send__`, or `MPI_Send_f`, or `mpi_send_f08__`, or `MPI_Send_f08`. If the profiling layer wants to be independent of the decision of the user program and MPI implementation, then it should provide all four routines. For example:

```
SUBROUTINE MPI_SEND(...) BIND(C,NAME='MPI_Send_f')
  USE mpi
  CALL PMPI_SEND(...)
END SUBROUTINE
```

The MPI library must provide the `PMPI_SEND` routine according to the same rules as for providing the `MPI_SEND` routine. (*End of advice to users.*)

*Advice to implementors.* If an implementation provides in a first step two sets of routines, one for the `mpi` module and `mpif.h`, and the other for the `mpi_f08` module, and both sets without TR 29113, i.e., `MPI_SUBARRAYS_SUPPORTED` equals `.FALSE.`, [If]and the implementor wants to add a TR 29113 based set of routines, then it is not necessary to add two full sets of routines. For full quality, it is enough to implement in each set only those routines that have a choice buffer argument. (*End of advice to implementors.*)

In the case that a Fortran binding consists of multiple routines through function overloading, the base names of overloaded routines are appended by a suffix notifying the difference in the argument list. For example, `MPI_ALLOC_MEM` (in the `mpi` module and `mpif.h`) has an `INTEGER(KIND=...) baseptr` argument without a suffix. This routine is overloaded by a routine with `TYPE(C_PTR) baseptr` and the suffix `_CPTR`. The implied linker name base is `MPI_ALLOC_MEM_CPTR`. It is mapped to the linker names `MPI_Alloc_mem_cptr_f`, and, e.g., `mpi_alloc_mem_cptr__`. Note that these routines are always called via the interface name `MPI_ALLOC_MEM` by the application within all Fortran support methods.

For routines without `ASYNCHRONOUS` choice buffers and that are not predefined callback routines, the implementor can freely choose to implement the routines according to the rules for `MPI_SUBARRAYS_SUPPORTED` equals `.TRUE.` or `.FALSE.`, provided that the following rule about routine grouping is fulfilled. The implementation of routines with `ASYNCHRONOUS` choice buffers depends on the rules for the provided Fortran support method and language

level of the underlying compiler. Predefined callback routines for the `mpi_f08` module must be implemented with `BIND(C)` interfaces, and for the `mpi` module and `mpif.h` without `BIND(C)`.

Similar MPI routines are grouped together for linker symbol scheme classification. If the peer routine of a group is available within an MPI library with one of its possible linker names then all of the routines in this group must be provided according to the same linker name scheme. If the peer routine is not available through a linker name scheme then all other routines in the group must not be available through this scheme.

Peer routines and their groups:

<code>MPI_ALLOC_MEM</code> .....	<code>MPI_ALLOC_MEM</code> and <code>MPI_WIN_ALLOCATE</code> .
<code>MPI_FREE_MEM</code> .....	Only this routine is in this group.
<code>MPI_GET_ADDRESS</code> .....	<code>MPI_GET_ADDRESS</code> and <code>MPI_ADDRESS</code> .
<code>MPI_SEND</code> .....	All routines with choice buffer arguments that are not declared as <code>ASYNCHRONOUS</code> within the <code>mpi_f08</code> module.
<code>MPI_ISEND</code> .....	All routines with choice buffer arguments that are declared as <code>ASYNCHRONOUS</code> within the <code>mpi_f08</code> module.
<code>MPI_OP_CREATE</code> .....	Only this routine is in this group.
<code>MPI_REGISTER_DATAREP</code> ....	Only this routine is in this group.
<code>MPI_COMM_KEYVAL_CREATE</code>	All other routines with callback function arguments.
<code>MPI_COMM_DUP_FN</code> .....	All predefined callback routines.
<code>MPI_COMM_RANK</code> .....	All other MPI routines.

Additionally, four C preprocessor macros are available in `mpi.h` for each routine group. The name of the macros are the peer routine name written as in the list above and appended with one of the following suffixes and meanings:

<code>_mpi_f08_BIND_C</code>	The macro is set to 1 if the <code>BIND(C)</code> linker name with the linker suffix <code>_f08</code> is available for all routines within this group (e.g., <code>MPI_Send_f08</code> ), otherwise it is set to 0.
<code>_mpi_f08_BIND_F</code>	The macro is set to 1 if the Fortran linker name with the linker suffix <code>_f08</code> is available for all routines within this group (e.g., <code>mpi_send_f08__</code> ), otherwise it is set to 0.
<code>_mpi_BIND_C</code> ...	The macro is set to 1 if the <code>BIND(C)</code> linker name with the linker suffix <code>_f</code> is available for all routines within this group (e.g., <code>MPI_Send_f</code> ), otherwise it is set to 0.
<code>_mpi_BIND_F</code> ...	The macro is set to 1 if the Fortran linker name without a linker suffix is available for all routines within this group (e.g., <code>mpi_send__</code> ), otherwise it is set to 0.

For example

```
...
#define MPI_SEND_mpi_f08_BIND_C 0
#define MPI_SEND_mpi_f08_BIND_F 1
#define MPI_SEND_mpi_BIND_C 0
#define MPI_SEND_mpi_BIND_F 1
```

```

#define MPI_ISEND_mpi_f08_BIND_C 1
#define MPI_ISEND_mpi_f08_BIND_F 1
#define MPI_ISEND_mpi_BIND_C      1
#define MPI_ISEND_mpi_BIND_F      1
...
#define MPI_COMM_DUP_FN_mpi_f08_BIND_C 1
#define MPI_COMM_DUP_FN_mpi_f08_BIND_F 0
#define MPI_COMM_DUP_FN_mpi_BIND_C      0
#define MPI_COMM_DUP_FN_mpi_BIND_F      1
...

```

shows, that

- the routines in the `MPI_SEND` group are only available through their Fortran linker names (e.g., `mpi_send_f08__`, `mpi_send__`, `mpi_recv_f08__`, `mpi_recv__`, ...),
- the routines in the `MPI_ISEND` group are available with all four interfaces: the `MPI` library, the `mpi_f08` and `mpi` modules (that provide the TR 29113 quality), and this `MPI` library supports application routines that are compiled with an older `MPI` library version with `_BIND_C` set to 0 and `_BIND_F` set to 1.

For the predefined callbacks, there is no choice, because the interfaces must fit to the callback function prototypes which are `BIND(C)` based for `mpi_f08` and without `BIND(C)` for the `mpi` module and `mpif.h`.

*Advice to implementors.* If all following conditions are fulfilled (which is the case for most compilers):

- the handles in the `mpi_f08` module occupy one Fortran numerical storage unit (same as an `INTEGER` handle),
- the internal argument passing used to pass an actual `ierror` argument to a non optional `ierror` dummy argument is binary compatible to passing an actual `ierror` argument to an `ierror` dummy argument that is declared as `OPTIONAL`,
- the internal argument passing for `ASYNCHRONOUS` and non-`ASYNCHRONOUS` arguments is the same,
- the internal routine call mechanism is the same for the Fortran and the C compiler,
- the compiler does not provide TR 29113,

then for most groups, the implementor may use the same internal routine implementations for all Fortran support methods but with several different linker names. For TR 29113 quality, new routines are needed only for the routine group of `MPI_ISEND`. Typical settings for `_mpi_f08_BIND_C` / `_mpi_f08_BIND_F` / `_mpi_BIND_C` / `_mpi_BIND_F` may be:

	Without TR 29113	Upgrade to TR 29113	Upgrade for strided data optimization	New impl. with TR 29113
MPI_ALLOC_MEM_....	0/1/0/1	0/1/0/1	0/1/0/1	1/0/1/0
MPI_FREE_MEM_....	0/1/0/1	0/1/0/1	0/1/0/1	1/0/1/0
MPI_GET_ADDRESS_....	0/1/0/1	0/1/0/1	0/1/0/1	1/0/1/0
MPI_SEND_....	0/1/0/1	0/1/0/1	1/1/1/1	1/0/1/0
MPI_ISEND_....	0/1/0/1	1/1/1/1	1/1/1/1	1/0/1/0
MPI_OP_CREATE_....	0/1/0/1	0/1/0/1	0/1/0/1	1/0/1/0
MPI_REGISTER_DATAREP_....	0/1/0/1	0/1/0/1	0/1/0/1	1/0/1/0
MPI_COMM_KEYVAL_CREATE_....	0/1/0/1	0/1/0/1	0/1/0/1	1/0/1/0
MPI_COMM_DUP_FN_....	1/0/0/1	1/0/0/1	1/0/0/1	1/0/0/1
MPI_COMM_RANK_....	0/1/0/1	0/1/0/1	0/1/0/1	1/0/1/0

(End of advice to implementors.)

## 16.2.6 MPI for Different Fortran Standard Versions

This section describes which Fortran interface functionality can be provided for different versions of the Fortran standard.

- For Fortran 77 with some extensions:

- MPI identifiers are limited to thirty or more, not six, significant characters.
- MPI identifiers may contain underscores after the first character.
- An MPI subroutine with a choice argument may be called with different argument types.
- Although not required by the MPI standard, the `INCLUDE` statement should be available for including `mpif.h` into the user application source code.

Only MPI-1.1, MPI-1.2, and MPI-1.3 can be implemented. The use of absolute addresses from `MPI_ADDRESS` and `MPI_BOTTOM` may cause problems if an address does not fit into the memory space provided by an `INTEGER`. (In MPI-2.0 this problem is solved with `MPI_GET_ADDRESS`, but not for Fortran 77.)

- For Fortran 90:

The major additional features that are needed from Fortran 90 are:

- The `MODULE` and `INTERFACE` concept.
- The `KIND=` and `SELECTED_..._KIND` concept.
- Fortran derived `TYPE`s and the `SEQUENCE` attribute.
- The `OPTIONAL` attribute for dummy arguments.
- Cray pointers, which are a non-standard compiler extension, are needed for the use of `MPI_ALLOC_MEM`.

With these features, MPI-1.1 - MPI-2.2 can be implemented without restrictions. MPI-3.0 can be implemented with some restrictions. The Fortran support methods are abbreviated with `S1` = the `mpi_f08` module, `S2` = the `mpi` module, and `S3` = the `mpif.f` include file. If not stated otherwise, restrictions exist for each method which prevent implementing the complete semantics of MPI-3.0.

- MPI\_SUBARRAYS\_SUPPORTED equals `.FALSE.`, i.e., subscript triplets and non-contiguous subarrays cannot be used as buffers in nonblocking routines, RMA, or split-collective I/O.
- S1, S2, and S3 can be implemented, but for S1, only a preliminary implementation is possible.
- In this preliminary interface of S1, the following changes are necessary:
  - \* The routines are not `BIND(C)`.
  - \* `TYPE(*)`, `DIMENSION(...)` is substituted by non-standardized extensions like `!$PRAGMA IGNORE_TKR`.
  - \* The `ASYNCHRONOUS` attribute is omitted.
  - \* `PROCEDURE(...)` callback declarations are substituted by `EXTERNAL`.
- The linker names are specified in Section 16.2.5 on page 649.
- Due to the rules specified in Section 16.2.5 on page 649, choice buffer declarations should be implemented only with non-standardized extensions like `!$PRAGMA IGNORE_TKR` (as long as F2008+TR29113 is not available).

In S2 and S3: Without such extensions, routines with choice buffers should be provided with an implicit interface, instead of overloading with a different MPI function for each possible buffer type (as mentioned in Section 16.2.11 on page 670). Such overloading would also imply restrictions for passing Fortran derived types as choice buffer, see also Section 16.2.15 on page 675.

Only in S1: The implicit interfaces for routines with choice buffer arguments imply that the `ierror` argument cannot be defined as `OPTIONAL`. For this reason, it is recommended not to provide the `mpi_f08` module if such an extension is not available.

- The `ASYNCHRONOUS` attribute can **not** be used in applications to protect buffers in nonblocking MPI calls (S1-S3).
- The `TYPE(C_PTR)` binding of the `MPI_ALLOC_MEM` and `MPI_WIN_ALLOCATE` routines is not available.
- In S1 and S2, the definition of the handle types (e.g., `TYPE(MPI_Comm)` and the status type `TYPE(MPI_Status)` must be modified: The `SEQUENCE` attribute must be used instead of `BIND(C)` (which is not available in Fortran 90/95). This restriction implies that the application must be fully recompiled if one switches to an MPI library for Fortran 2003 and later because the internal memory size of the handles may have changed. For this reason, an implementor may choose not to provide the `mpi_f08` module for Fortran 90 compilers. In this case, the `mpi_f08` handle types and all routines, constants and types related to `TYPE(MPI_Status)` (see Section 16.3.5 on page 696) are also not available in the `mpi` module and `mpif.h`.

- For Fortran 95:

The quality of the MPI interface and the restrictions are the same as with Fortran 90.

- For Fortran 2003:

The major features that are needed from Fortran 2003 are:

- Interoperability with C, i.e.,

- \* `BIND(C, NAME='...')` interfaces.
- \* `BIND(C)` derived types.
- \* The `ISO_C_BINDING` intrinsic type `C_PTR` and routine `C_F_POINTER`.
- The ability to define an `ABSTRACT INTERFACE` and to use it for `PROCEDURE` dummy arguments.
- The `ASYNCHRONOUS` attribute is available to protect Fortran asynchronous I/O. This feature is not yet used by MPI, but it is the basis for the enhancement for MPI communication in the TR 29113.

With these features (but still without the features of TR29113), MPI-1.1 - MPI-2.2 can be implemented without restrictions, but with one enhancement:

- The user application can use `TYPE(C_PTR)` together with `MPI_ALLOC_MEM` as long as `MPI_ALLOC_MEM` is defined with an implicit interface because a `C_PTR` and an `INTEGER(KIND=MPI_ADDRESS_KIND)` argument must both map to a `void *` argument.

MPI-3.0 can be implemented with the following restrictions:

- `MPI_SUBARRAYS_SUPPORTED` equals `.FALSE..`
  - For `S1`, only a preliminary implementation is possible. The following changes are necessary:
    - \* The routines are not `BIND(C)`.
    - \* `TYPE(*)`, `DIMENSION(..)` is substituted by non-standardized extensions like `!$PRAGMA IGNORE_TKR`.
  - The linker names are specified in Section 16.2.5 on page 649.
  - With `S1`, the `ASYNCHRONOUS` is required as specified in the second Fortran interfaces. With `S2` and `S3` the implementation can also add this attribute if explicit interfaces are used.
  - The `ASYNCHRONOUS` Fortran attribute can be used in applications to **try to** protect buffers in nonblocking MPI calls, but the protection can work only if the compiler is able to protect asynchronous Fortran I/O and makes no difference between such asynchronous Fortran I/O and MPI communication.
  - The `TYPE(C_PTR)` binding of the `MPI_ALLOC_MEM` and `MPI_WIN_ALLOCATE` routines can be used only for Fortran types that are C compatible.
  - The same restriction as for Fortran 90 applies if non-standardized extensions like `!$PRAGMA IGNORE_TKR` are not available.
- For Fortran 2008 + TR 29113 and later and  
For Fortran 2003 + TR 29113:  
 The major feature that are needed from TR29113 are:
    - `TYPE(*)`, `DIMENSION(..)` is available.
    - The `ASYNCHRONOUS` attribute is extended to protect also nonblocking MPI communication.
    - `OPTIONAL` dummy arguments are allowed in combination with `BIND(C)` interfaces.



- CHARACTER(LEN=\*) dummy arguments are allowed in combination with BIND(C) interfaces.
- The array dummy argument of the ISO\_C\_BINDING intrinsic C\_F\_POINTER is not restricted to Fortran types for which a corresponding type in C exists.

Using these features, MPI-3.0 can be implemented without any restrictions.

- With S1, MPI\_SUBARRAYS\_SUPPORTED equals .TRUE.. The ASYNCHRONOUS attribute can be used to protect buffers in nonblocking MPI calls. The TYPE(C\_PTR) binding of the MPI\_ALLOC\_MEM and MPI\_WIN\_ALLOCATE routines can be used for any Fortran type.
- With S2 and S3, the value of MPI\_SUBARRAYS\_SUPPORTED is implementation dependent. A high quality implementation will also provide MPI\_SUBARRAYS\_SUPPORTED==.TRUE. and will use the ASYNCHRONOUS attribute in the same way as in S1.
- If non-standardized extensions like !\$PRAGMA IGNORE\_TKR are not available then S2 must be implemented with TYPE(\*), DIMENSION(..).

*Advice to implementors.* If MPI\_SUBARRAYS\_SUPPORTED==.FALSE., the choice argument may be implemented with an explicit interface using compiler directives, for example:

```

INTERFACE
  SUBROUTINE MPI...(buf, ...)
    !DEC$ ATTRIBUTES NO_ARG_CHECK :: buf
    !$PRAGMA IGNORE_TKR buf
    !DIR$ IGNORE_TKR buf
    !IBM* IGNORE_TKR buf
    REAL, DIMENSION(*) :: buf
    ... ! declarations of the other arguments
  END SUBROUTINE
END INTERFACE

```

*(End of advice to implementors.)*

### 16.2.7 Requirements on Fortran Compilers

MPI-3.0 (and later) compliant Fortran bindings are not only a property of the MPI library itself, but rather a property of an MPI library together with the Fortran compiler suite for which it is compiled.

*Advice to users.* Users must take appropriate steps to ensure that proper options are specified to compilers. MPI libraries must document these options. Some MPI libraries are shipped together with special compilation scripts (e.g., mpif90, mpicc) that set these options automatically. *(End of advice to users.)*

An MPI library together with the Fortran compiler suite is only compliant with MPI-3.0 (and later), as referred by MPI\_GET\_VERSION, if all the solutions described in Sections 16.2.11 through 16.2.19 work correctly. Based on this rule, major requirements for all three Fortran support methods (i.e., the mpi\_f08 and mpi modules, and mpif.h) are:

- The language features assumed-type and assumed-rank from Fortran 2008 TR 29113 [41] are available. This is required only for `mpi_f08`. As long as this requirement is not supported by the compiler, it is valid to build [a preliminary MPI-3.0 (and not later)]an MPI library that implements the `mpi_f08` module with `MPI_SUBARRAYS_SUPPORTED` set to `.FALSE..`
- ‘Simply contiguous’ arrays and scalars must be passed to choice buffer dummy arguments of nonblocking routines with call by reference. This is needed only if one of the support methods does not use the `ASYNCHRONOUS` attribute. See Section 16.2.12 on page 671 for more details.
- `SEQUENCE` and `BIND(C)` derived types are valid as actual arguments passed to choice buffer dummy arguments, and, in the case of `MPI_SUBARRAYS_SUPPORTED==.FALSE.`, they are passed with call by reference, and passed by descriptor in the case of `.TRUE..`
- All actual arguments that are allowed for a dummy argument in an implicitly defined and separately compiled Fortran routine with the given compiler (e.g., `CHARACTER(LEN=*)` strings and array of strings) must also be valid for choice buffer dummy arguments with all Fortran support methods.
- The array dummy argument of the `ISO_C_BINDING` intrinsic module procedure `C_F_POINTER` is not restricted to Fortran types for which a corresponding type in C exists.
- The Fortran compiler shall not provide `TYPE(*)` unless the `ASYNCHRONOUS` attribute protects MPI communication as described in TR 29113. Specifically, the TR 29113 must be implemented as a whole.

The following rules are required at least as long as the compiler does not provide the extension of the `ASYNCHRONOUS` attribute as part of TR 29113 and there is still one Fortran support method with `MPI_ASYNC_PROTECTS_NONBLOCKING==.FALSE..` It is helpful when these rules are observed, especially for backward compatibility of existing applications that use the `mpi` module or the `mpif.h` include file. The rules are as follows:

- Separately compiled empty Fortran routines with implicit interfaces and separately compiled empty C routines with `BIND(C)` Fortran interfaces (e.g., `MPI_F_SYNC_REG` on page 682 and Section 16.2.8 on page 659, and `DD` on page 683) solve the problems described in Section 16.2.17 on page 678.
- The problems with temporary data movement (described in detail in Section 16.2.18 on page 685) are solved as long as the application uses different sets of variables for the nonblocking communication (or nonblocking or split collective IO) and the computation when overlapping communication and computation.
- Problems caused by automatic and permanent data movement (e.g., within a garbage collection, see Section 16.2.19 on page 686) are resolved **without** any further requirements on the application program, neither on the usage of the buffers, nor on the declaration of application routines that are involved in calling MPI operations.

All of these rules are valid independently of whether the MPI routine interfaces in the `mpi_f08` and `mpi` modules are internally defined with an `INTERFACE` or `CONTAINS` construct, and with or without `BIND(C)`, and also when `mpif.h` uses explicit interfaces.

*Advice to implementors.* Some of these rules are already part of the Fortran 2003 standard if the MPI interfaces are defined without BIND(C). Additional compiler support may be necessary if BIND(C) is used. Some of these additional requirements are defined in the Fortran 2008 TR 29113 [41]. Some of these requirements for MPI-3.0 are beyond the scope of TR 29113. (*End of advice to implementors.*)

Further requirements apply when the MPI library internally uses BIND(C) routine interfaces (i.e, for a full implementation of `mpi_f08`):

- Non-buffer arguments are `INTEGER`, `INTEGER(KIND=...)`, `CHARACTER(LEN=*)`, `LOGICAL`, and `BIND(C)` derived types, (handles and status in `mpi_f08`) variables and arrays; function results are `DOUBLE PRECISION`. All these types must be valid as dummy arguments in the `BIND(C)` MPI routine interfaces. When compiling an MPI application, the compiler should not issue warnings indicating that these types may not be interoperable with an existing type in C. Some of these types are already valid in `BIND(C)` interfaces since Fortran 2003, some may be valid based on TR 29113 (e.g., `CHARACTER*(*)`).
- `OPTIONAL` dummy arguments are also valid within `BIND(C)` interfaces. This requirement is fulfilled if TR 29113 is fully supported by the compiler.

### 16.2.8 Additional Support for Fortran Register-Memory-Synchronization

As described in Section 16.2.17 on page 678, a dummy call may be necessary to tell the compiler that registers are to be flushed for a given buffer or that accesses to a buffer may not be moved across a given point in the execution sequence. Only a Fortran binding exists for this call.

MPI\_F\_SYNC\_REG(buf)

**INOUT**    buf                          initial address of buffer (choice)

```
MPI_F_sync_reg(buf) BIND(C)
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
```

```
MPI_F_SYNC_REG(buf)
    <type> buf(*)
```

This routine is a no-operation. It must be compiled in the MPI library in such a manner that a Fortran compiler cannot detect in the module that the routine has an empty body. It is used only to force the compiler to flush a cached register value of a variable or buffer back to memory (when necessary), or to invalidate the register value.

*Rationale.* This function is not available in other languages because it would not be useful. This routine has no `ierror` return argument because there is no operation that can fail. (*End of rationale.*)

*Advice to implementors.* This routine can be bound to a C routine to minimize the risk that the Fortran compiler can learn that this routine is empty (and that the call to this routine can be removed as part of an optimization). However, it is

explicitly allowed to implement this routine within the `mpi_f08` module according to the definition for the `mpi` module or `mpif.h` to circumvent the overhead of building the internal dope vector to handle the assumed-type, assumed-rank argument. (*End of advice to implementors.*)

*Rationale.* This routine is not defined with `TYPE(*)`, `DIMENSION(*)`, i.e., assumed size instead of assumed rank, because this would restrict the usability to ‘simply contiguous’ arrays and would require overloading with another interface for scalar arguments. (*End of rationale.*)

*Advice to users.* If only a part of an array (e.g., defined by a subscript triplet) is used in a nonblocking routine, it is recommended to pass the whole array to `MPI_F_SYNC_REG` anyway to minimize the overhead of this no-operation call. Note that this routine need not to be called if `MPI_ASYNC_PROTECTS_NONBLOCKING` is `.TRUE.` and the application fully uses the facilities of `ASYNCHRONOUS` arrays. (*End of advice to users.*)

## 16.2.9 Additional Support for Fortran Numeric Intrinsic Types

[The routines in this section are part of Extended Fortran Support described in Section 16.2.3. ]

MPI provides a small number of named datatypes that correspond to named intrinsic types supported by C and Fortran. These include `MPI_INTEGER`, `MPI_REAL`, `MPI_INT`, `MPI_DOUBLE`, etc., as well as the optional types `MPI_REAL4`, `MPI_REAL8`, etc. There is a one-to-one correspondence between language declarations and MPI types.

Fortran (starting with Fortran 90) provides so-called `KIND`-parameterized types. These types are declared using an intrinsic type (one of `INTEGER`, `REAL`, `COMPLEX`, `LOGICAL` and `CHARACTER`) with an optional integer `KIND` parameter that selects from among one or more variants. The specific meaning of different `KIND` values themselves are implementation dependent and not specified by the language. Fortran provides the `KIND` selection functions `selected_real_kind` for `REAL` and `COMPLEX` types, and `selected_int_kind` for `INTEGER` types that allow users to declare variables with a minimum precision or number of digits. These functions provide a portable way to declare `KIND`-parameterized `REAL`, `COMPLEX` and `INTEGER` variables in Fortran. This scheme is backward compatible with Fortran 77. `REAL` and `INTEGER` Fortran variables have a default `KIND` if none is specified. Fortran `DOUBLE PRECISION` variables are of intrinsic type `REAL` with a non-default `KIND`. The following two declarations are equivalent:

```
double precision x
real(KIND(0.0d0)) x
```

MPI provides two orthogonal methods to communicate using numeric intrinsic types. The first method (see the following section) can be used when variables have been declared in a portable way — using default `KIND` or using `KIND` parameters obtained with the `selected_int_kind` or `selected_real_kind` functions. With this method, MPI automatically selects the correct data size (e.g., 4 or 8 bytes) and provides representation conversion in heterogeneous environments. The second method (see Support for size-specific MPI Datatypes on page 665) gives the user complete control over communication by exposing machine representations.

## Parameterized Datatypes with Specified Precision and Exponent Range

MPI provides named datatypes corresponding to standard Fortran 77 numeric types — `MPI_INTEGER`, `MPI_COMPLEX`, `MPI_REAL`, `MPI_DOUBLE_PRECISION` and `MPI_DOUBLE_COMPLEX`. MPI automatically selects the correct data size and provides representation conversion in heterogeneous environments. The mechanism described in this section extends this model to support portable parameterized numeric types.

The model for supporting portable parameterized types is as follows. Real variables are declared (perhaps indirectly) using `selected_real_kind(p, r)` to determine the `KIND` parameter, where `p` is decimal digits of precision and `r` is an exponent range. Implicitly MPI maintains a two-dimensional array of predefined MPI datatypes `D(p, r)`. `D(p, r)` is defined for each value of `(p, r)` supported by the compiler, including pairs for which one value is unspecified. Attempting to access an element of the array with an index `(p, r)` not supported by the compiler is erroneous. MPI implicitly maintains a similar array of `COMPLEX` datatypes. For integers, there is a similar implicit array related to `selected_int_kind` and indexed by the requested number of digits `r`. Note that the predefined datatypes contained in these implicit arrays are not the same as the named MPI datatypes `MPI_REAL`, etc., but a new set.

*Advice to implementors.* The above description is for explanatory purposes only. It is not expected that implementations will have such internal arrays. (*End of advice to implementors.*)

*Advice to users.* `selected_real_kind()` maps a large number of `(p,r)` pairs to a much smaller number of `KIND` parameters supported by the compiler. `KIND` parameters are not specified by the language and are not portable. From the language point of view intrinsic types of the same base type and `KIND` parameter are of the same type. In order to allow interoperability in a heterogeneous environment, MPI is more stringent. The corresponding MPI datatypes match if and only if they have the same `(p,r)` value (`REAL` and `COMPLEX`) or `r` value (`INTEGER`). Thus MPI has many more datatypes than there are fundamental language types. (*End of advice to users.*)

`MPI_TYPE_CREATE_F90_REAL(p, r, newtype)`

IN	<code>p</code>	precision, in decimal digits (integer)
IN	<code>r</code>	decimal exponent range (integer)
OUT	<code>newtype</code>	the requested MPI datatype (handle)

```
int MPI_Type_create_f90_real(int p, int r, MPI_Datatype *newtype)
```

```
MPI_Type_create_f90_real(p, r, newtype, ierror) BIND(C)
```

```
    INTEGER, INTENT(IN) :: p, r
```

```
    TYPE(MPI_Datatype), INTENT(OUT) :: newtype
```

```
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_TYPE_CREATE_F90_REAL(P, R, NEWTYPE, IERROR)
```

```
    INTEGER P, R, NEWTYPE, IERROR
```

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```

1 {static MPI::Datatype MPI::Datatype::Create_f90_real(int p, int r) (binding
2   deprecated, see Section 15.2) }
3

```

This function returns a predefined MPI datatype that matches a REAL variable of KIND `selected_real_kind(p, r)`. In the model described above it returns a handle for the element `D(p, r)`. Either `p` or `r` may be omitted from calls to `selected_real_kind(p, r)` (but not both). Analogously, either `p` or `r` may be set to `MPI_UNDEFINED`. In communication, an MPI datatype `A` returned by `MPI_TYPE_CREATE_F90_REAL` matches a datatype `B` if and only if `B` was returned by `MPI_TYPE_CREATE_F90_REAL` called with the same values for `p` and `r` or `B` is a duplicate of such a datatype. Restrictions on using the returned datatype with the “external32” data representation are given on page 664.

It is erroneous to supply values for `p` and `r` not supported by the compiler.

```

14 MPI_TYPE_CREATE_F90_COMPLEX(p, r, newtype)
15

```

IN	<code>p</code>	precision, in decimal digits (integer)
IN	<code>r</code>	decimal exponent range (integer)
OUT	<code>newtype</code>	the requested MPI datatype (handle)

```

20
21 int MPI_Type_create_f90_complex(int p, int r, MPI_Datatype *newtype)

```

```

22 MPI_Type_create_f90_complex(p, r, newtype, ierror) BIND(C)
23   INTEGER, INTENT(IN) :: p, r
24   TYPE(MPI_Datatype), INTENT(OUT) :: newtype
25   INTEGER, OPTIONAL, INTENT(OUT) :: ierror
26

```

```

27 MPI_TYPE_CREATE_F90_COMPLEX(P, R, NEWTYPE, IERROR)
28   INTEGER P, R, NEWTYPE, IERROR
29

```

```

30 {static MPI::Datatype MPI::Datatype::Create_f90_complex(int p,
31   int r) (binding deprecated, see Section 15.2) }
32

```

This function returns a predefined MPI datatype that matches a COMPLEX variable of KIND `selected_real_kind(p, r)`. Either `p` or `r` may be omitted from calls to `selected_real_kind(p, r)` (but not both). Analogously, either `p` or `r` may be set to `MPI_UNDEFINED`. Matching rules for datatypes created by this function are analogous to the matching rules for datatypes created by `MPI_TYPE_CREATE_F90_REAL`. Restrictions on using the returned datatype with the “external32” data representation are given on page 664.

It is erroneous to supply values for `p` and `r` not supported by the compiler.

```

41 MPI_TYPE_CREATE_F90_INTEGER(r, newtype)
42

```

IN	<code>r</code>	decimal exponent range, i.e., number of decimal digits (integer)
OUT	<code>newtype</code>	the requested MPI datatype (handle)

```

47 int MPI_Type_create_f90_integer(int r, MPI_Datatype *newtype)
48

```

```

MPI_Type_create_f90_integer(r, newtype, ierror) BIND(C)
  INTEGER, INTENT(IN) :: r
  TYPE(MPI_Datatype), INTENT(OUT) :: newtype
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_TYPE_CREATE_F90_INTEGER(R, NEWTYPE, IERROR)
  INTEGER R, NEWTYPE, IERROR

{static MPI::Datatype MPI::Datatype::Create_f90_integer(int r) (binding
  deprecated, see Section 15.2) }

```

This function returns a predefined MPI datatype that matches a `INTEGER` variable of `KIND selected_int_kind(r)`. Matching rules for datatypes created by this function are analogous to the matching rules for datatypes created by `MPI_TYPE_CREATE_F90_REAL`. Restrictions on using the returned datatype with the “external32” data representation are given on page 664.

It is erroneous to supply a value for `r` that is not supported by the compiler.

Example:

```

integer      longtype, quadtype
integer, parameter :: long = selected_int_kind(15)
integer(long) ii(10)
real(selected_real_kind(30)) x(10)
call MPI_TYPE_CREATE_F90_INTEGER(15, longtype, ierror)
call MPI_TYPE_CREATE_F90_REAL(30, MPI_UNDEFINED, quadtype, ierror)
...

call MPI_SEND(ii, 10, longtype, ...)
call MPI_SEND(x, 10, quadtype, ...)

```

*Advice to users.* The datatypes returned by the above functions are predefined datatypes. They cannot be freed; they do not need to be committed; they can be used with predefined reduction operations. There are two situations in which they behave differently syntactically, but not semantically, from the MPI named predefined datatypes.

1. `MPI_TYPE_GET_ENVELOPE` returns special combinators that allow a program to retrieve the values of `p` and `r`.
2. Because the datatypes are not named, they cannot be used as compile-time initializers or otherwise accessed before a call to one of the `MPI_TYPE_CREATE_F90_XXXX` routines.

If a variable was declared specifying a non-default `KIND` value that was not obtained with `selected_real_kind()` or `selected_int_kind()`, the only way to obtain a matching MPI datatype is to use the size-based mechanism described in the next section.

*(End of advice to users.)*

*Advice to implementors.* An application may often repeat a call to `MPI_TYPE_CREATE_F90_XXXX` with the same combination of `(xxxx,p,r)`. The application is not allowed to free the returned predefined, unnamed datatype handles. To



prevent the creation of a potentially huge amount of handles, a high quality MPI implementation should return the same datatype handle for the same (REAL/COMPLEX/INTEGER,p,r) combination. Checking for the combination (p,r) in the preceding call to MPI\_TYPE\_CREATE\_F90\_xxxx and using a hash[-] table to find formerly generated handles should limit the overhead of finding a previously generated datatype with same combination of (xxxx,p,r). (*End of advice to implementors.*)

*Rationale.* The MPI\_TYPE\_CREATE\_F90\_REAL/COMPLEX/INTEGER interface needs as input the original range and precision values to be able to define useful and compiler-independent external (Section 13.5.2 on page 558) or user-defined (Section 13.5.3 on page 559) data representations, and in order to be able to perform automatic and efficient data conversions in a heterogeneous environment. (*End of rationale.*)

We now specify how the datatypes described in this section behave when used with the “external32” external data representation described in Section 13.5.2 on page 558.

The external32 representation specifies data formats for integer and floating point values. Integer values are represented in two’s complement big-endian format. Floating point values are represented by one of three IEEE formats. These are the IEEE “Single,” “Double” and “Double Extended” formats, requiring 4, 8 and 16 bytes of storage, respectively. For the IEEE “Double Extended” formats, MPI specifies a Format Width of 16 bytes, with 15 exponent bits, bias = +10383, 112 fraction bits, and an encoding analogous to the “Double” format.

The external32 representations of the datatypes returned by MPI\_TYPE\_CREATE\_F90\_REAL/COMPLEX/INTEGER are given by the following rules.

For MPI\_TYPE\_CREATE\_F90\_REAL:

```

if      (p > 33) or (r > 4931) then  external32 representation
                                     is undefined
else if (p > 15) or (r > 307) then  external32_size = 16
else if (p > 6) or (r > 37) then   external32_size = 8
else                                external32_size = 4

```

For MPI\_TYPE\_CREATE\_F90\_COMPLEX: twice the size as for MPI\_TYPE\_CREATE\_F90\_REAL.

For MPI\_TYPE\_CREATE\_F90\_INTEGER:

```

if      (r > 38) then  external32 representation is undefined
else if (r > 18) then  external32_size = 16
else if (r > 9) then   external32_size = 8
else if (r > 4) then   external32_size = 4
else if (r > 2) then   external32_size = 2
else                  external32_size = 1

```

If the external32 representation of a datatype is undefined, the result of using the datatype directly or indirectly (i.e., as part of another datatype or through a duplicated datatype) in operations that require the external32 representation is undefined. These operations include MPI\_PACK\_EXTERNAL, MPI\_UNPACK\_EXTERNAL and many MPI\_FILE functions, when the “external32” data representation is used. The ranges for which the external32 representation is undefined are reserved for future standardization.



## Support for Size-specific MPI Datatypes

MPI provides named datatypes corresponding to optional Fortran 77 numeric types that contain explicit byte lengths — `MPI_REAL4`, `MPI_INTEGER8`, etc. This section describes a mechanism that generalizes this model to support all Fortran numeric intrinsic types.

We assume that for each **typeclass** (integer, real, complex) and each word size there is a unique machine representation. For every pair (**typeclass**, **n**) supported by a compiler, MPI must provide a named size-specific datatype. The name of this datatype is of the form `MPI_<TYPE>n` in C and Fortran and of the form `MPI::<TYPE>n` in C++ where `<TYPE>` is one of `REAL`, `INTEGER` and `COMPLEX`, and **n** is the length in bytes of the machine representation. This datatype locally matches all variables of type (**typeclass**, **n**). The list of names for such types includes:

```
MPI_REAL4
MPI_REAL8
MPI_REAL16
MPI_COMPLEX8
MPI_COMPLEX16
MPI_COMPLEX32
MPI_INTEGER1
MPI_INTEGER2
MPI_INTEGER4
MPI_INTEGER8
MPI_INTEGER16
```

One datatype is required for each representation supported by the compiler. To be backward compatible with the interpretation of these types in MPI-1, we assume that the nonstandard declarations `REAL*n`, `INTEGER*n`, always create a variable whose representation is of size **n**. These datatypes may also be used for variables declared with `KIND=INT8/16/32/64` or `KIND=REAL32/64/128`, which are defined in the `ISO_FORTRAN_ENV` intrinsic module. Note that the MPI datatypes and the `REAL*n`, `INTEGER*n` declarations count bytes whereas the Fortran `KIND` values count bits. All these datatypes are predefined.

The following functions allow a user to obtain a size-specific MPI datatype for any intrinsic Fortran type.

```
MPI_SIZEOF(x, size)
```

IN	x	a Fortran variable of numeric intrinsic type (choice)
OUT	size	size of machine representation of that type (integer)

```
MPI_Sizeof(x, size, ierror) BIND(C)
  TYPE(*), DIMENSION(..) :: x
  INTEGER, INTENT(OUT) :: size
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_SIZEOF(X, SIZE, IERROR)
  <type> X
  INTEGER SIZE, IERROR
```

This function returns the size in bytes of the machine representation of the given variable. It is a generic Fortran routine and has a Fortran binding only.

*Advice to users.* This function is similar to the C and C++ `sizeof` operator but behaves slightly differently. If given an array argument, it returns the size of the base element, not the size of the whole array. (*End of advice to users.*)

*Rationale.* This function is not available in other languages because it would not be useful. (*End of rationale.*)

`MPI_TYPE_MATCH_SIZE(typeclass, size, datatype)`

IN	typeclass	generic type specifier (integer)
IN	size	size, in bytes, of representation (integer)
OUT	[ticket252-W.]datatype	datatype with correct type, size (handle)

`int MPI_Type_match_size(int typeclass, int size, MPI_Datatype *datatype)`

```
MPI_Type_match_size(typeclass, size, datatype, ierror) BIND(C)
  INTEGER, INTENT(IN) :: typeclass, size
  TYPE(MPI_Datatype), INTENT(OUT) :: datatype
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

`MPI_TYPE_MATCH_SIZE(TYPECLASS, SIZE, DATATYPE, IERROR)`

INTEGER TYPECLASS, SIZE, DATATYPE, IERROR

```
{static MPI::Datatype MPI::Datatype::Match_size(int typeclass,
  int size) (binding deprecated, see Section 15.2) }
```

typeclass is one of `MPI_TYPECLASS_REAL`, `MPI_TYPECLASS_INTEGER` and `MPI_TYPECLASS_COMPLEX`, corresponding to the desired **typeclass**. The function returns an MPI datatype matching a local variable of type (**typeclass**, **size**).

This function returns a reference (handle) to one of the predefined named datatypes, not a duplicate. This type cannot be freed. `MPI_TYPE_MATCH_SIZE` can be used to obtain a size-specific type that matches a Fortran numeric intrinsic type by first calling `MPI_SIZEOF` in order to compute the variable size, and then calling `MPI_TYPE_MATCH_SIZE` to find a suitable datatype. In C and C++, one can use the C function `sizeof()`, instead of `MPI_SIZEOF`. In addition, for variables of default kind the variable's size can be computed by a call to `MPI_TYPE_GET_EXTENT`, if the **typeclass** is known. It is erroneous to specify a size not supported by the compiler.

*Rationale.* This is a convenience function. Without it, it can be tedious to find the correct named type. See note to implementors below. (*End of rationale.*)

*Advice to implementors.* This function could be implemented as a series of tests.

```
int MPI_Type_match_size(int typeclass, int size, MPI_Datatype *rtype)
{
  switch(typeclass) {
```

```

    case MPI_TYPECLASS_REAL: switch(size) {
        case 4: *rtype = MPI_REAL4; return MPI_SUCCESS;
        case 8: *rtype = MPI_REAL8; return MPI_SUCCESS;
        default: error(...);
    }
    case MPI_TYPECLASS_INTEGER: switch(size) {
        case 4: *rtype = MPI_INTEGER4; return MPI_SUCCESS;
        case 8: *rtype = MPI_INTEGER8; return MPI_SUCCESS;
        default: error(...);
    }
    ... etc. ...
}

```

*(End of advice to implementors.)*

### Communication With Size-specific Types

The usual type matching rules apply to size-specific datatypes: a value sent with datatype `MPI_<TYPE>n` can be received with this same datatype on another process. Most modern computers use 2's complement for integers and IEEE format for floating point. Thus, communication using these size-specific datatypes will not entail loss of precision or truncation errors.

*Advice to users.* Care is required when communicating in a heterogeneous environment. Consider the following code:

```

real(selected_real_kind(5)) x(100)
call MPI_SIZEOF(x, size, ierror)
call MPI_TYPE_MATCH_SIZE(MPI_TYPECLASS_REAL, size, xtype, ierror)
if (myrank .eq. 0) then
    ... initialize x ...
    call MPI_SEND(x, xtype, 100, 1, ...)
else if (myrank .eq. 1) then
    call MPI_RECV(x, xtype, 100, 0, ...)
endif

```

This may not work in a heterogeneous environment if the value of `size` is not the same on process 1 and process 0. There should be no problem in a homogeneous environment. To communicate in a heterogeneous environment, there are at least four options. The first is to declare variables of default type and use the MPI datatypes for these types, e.g., declare a variable of type `REAL` and use `MPI_REAL`. The second is to use `selected_real_kind` or `selected_int_kind` and with the functions of the previous section. The third is to declare a variable that is known to be the same size on all architectures (e.g., `selected_real_kind(12)` on almost all compilers will result in an 8-byte representation). The fourth is to carefully check representation size before communication. This may require explicit conversion to a variable of size that can be communicated and handshaking between sender and receiver to agree on a size.

Note finally that using the “external32” representation for I/O requires explicit attention to the representation sizes. Consider the following code:

```

1      real(selected_real_kind(5)) x(100)
2      call MPI_SIZEOF(x, size, ierror)
3      call MPI_TYPE_MATCH_SIZE(MPI_TYPECLASS_REAL, size, xtype, ierror)
4
5      if (myrank .eq. 0) then
6          call MPI_FILE_OPEN(MPI_COMM_SELF, 'foo',
7                              &
8                              MPI_MODE_CREATE+MPI_MODE_WRONLY, &
9                              MPI_INFO_NULL, fh, ierror)
10         call MPI_FILE_SET_VIEW(fh, zero, xtype, xtype, 'external32', &
11                                 MPI_INFO_NULL, ierror)
12         call MPI_FILE_WRITE(fh, x, 100, xtype, status, ierror)
13         call MPI_FILE_CLOSE(fh, ierror)
14     endif
15
16     call MPI_BARRIER(MPI_COMM_WORLD, ierror)
17
18     if (myrank .eq. 1) then
19         call MPI_FILE_OPEN(MPI_COMM_SELF, 'foo', MPI_MODE_RDONLY, &
20                             MPI_INFO_NULL, fh, ierror)
21         call MPI_FILE_SET_VIEW(fh, zero, xtype, xtype, 'external32', &
22                                 MPI_INFO_NULL, ierror)
23         call MPI_FILE_WRITE(fh, x, 100, xtype, status, ierror)
24         call MPI_FILE_CLOSE(fh, ierror)
25     endif
26
27
28
29
30
31
32
33
34
35
36
37
38
39
40
41
42
43
44
45
46
47
48

```

If processes 0 and 1 are on different machines, this code may not work as expected if the size is different on the two machines. (*End of advice to users.*)

### 16.2.10 Problems With Fortran Bindings for MPI

This section discusses a number of problems that may arise when using MPI in a Fortran program. It is intended as advice to users, and clarifies how MPI interacts with Fortran. It does not add to the standard, but is intended to clarify the standard.

As noted in the original MPI specification, the interface violates the Fortran standard in several ways. While these **may** cause few problems for Fortran 77 programs, they become more significant for Fortran 90 programs, so that users must exercise care when using new Fortran 90 features. **With Fortran 2008 and the new semantics defined in TR 29113, most violations are resolved, and this is hinted at in an addendum to each item.** The violations were originally adopted and have been retained because they are important for the usability of MPI. The rest of this section describes the potential problems in detail. **[It supersedes and replaces the discussion of Fortran bindings in the original MPI specification (for Fortran 90, not Fortran 77). ]**

The following MPI features are inconsistent with Fortran 90 **and Fortran 77**.

1. An MPI subroutine with a choice argument may be called with different argument types. When using the `mpi_f08` module together with a compiler that supports Fortran 2008 + TR 29113, this problem is resolved.
2. An MPI subroutine with an assumed-size dummy argument may be passed an actual scalar argument. This is only solved for choice buffers through the use of `DIMENSION(...)`.
3. [Many]Nonblocking and split-collective MPI routines assume that actual arguments are passed by address or descriptor and that arguments and the associated data are not copied on entrance to or exit from the subroutine. This problem is solved with the use of the `ASYNCHRONOUS` attribute.
4. An MPI implementation may read or modify user data (e.g., communication buffers used by nonblocking communications) concurrently with a user program that is executing outside of MPI calls. This problem is resolved by relying on the extended semantics of the `ASYNCHRONOUS` attribute as specified in TR 29113.
5. Several named “constants,” such as `MPI_BOTTOM`, `MPI_IN_PLACE`, `MPI_STATUS_IGNORE`, `MPI_STATUSES_IGNORE`, `MPI_ERRCODES_IGNORE`, `MPI_UNWEIGHTED`, `MPI_ARGV_NULL`, and `MPI_ARGVS_NULL` are not ordinary Fortran constants and require a special implementation. See Section 2.5.4 on page 15 for more information.
6. The memory allocation routine `MPI_ALLOC_MEM` can’t be usefully used in Fortran 77/90/95 without a language extension (for example, Cray pointers) that allows the allocated memory to be associated with a Fortran variable. Therefore, address sized integers were used in MPI-2.0 - MPI-2.2. In Fortran 2003, `TYPE(C_PTR)` entities were added, which allow a standard-conforming implementation of the semantics of `MPI_ALLOC_MEM`. In MPI-3.0 and later, `MPI_ALLOC_MEM` has an additional, overloaded interface to support this language feature. The use of Cray pointers is deprecated. The `mpi_f08` module only supports `TYPE(C_PTR)` pointers.

Additionally, MPI is inconsistent with Fortran 77 in a number of ways, as noted below.

- MPI identifiers exceed 6 characters.
- MPI identifiers may contain underscores after the first character.
- MPI requires an include file, `mpif.h`. On systems that do not support include files, the implementation should specify the values of named constants.
- Many routines in MPI have `KIND`-parameterized integers (e.g., `MPI_ADDRESS_KIND` and `MPI_OFFSET_KIND`) that hold address information. On systems that do not support Fortran 90-style parameterized types, `INTEGER*8` or `INTEGER` should be used instead.

MPI-1 contained several routines that take address-sized information as input or return address-sized information as output. In C such arguments were of type `MPI_Aint` and in Fortran of type `INTEGER`. On machines where integers are smaller than addresses, these routines can lose information. In MPI-2 the use of these functions has been deprecated and they have been replaced by routines taking `INTEGER` arguments of `KIND=MPI_ADDRESS_KIND`.

A number of new MPI-2 functions also take INTEGER arguments of non-default KIND. See Section 2.6 on page 17 and Section 4.1.1 on page 93 for more information.

ticket230-B.

Sections 16.2.11 through 16.2.19 describe several problems in detail which concern the interaction of MPI and Fortran as well as their solutions. Some of these solutions require special capabilities from the compilers. Major requirements are summarized in Section 16.2.7 on page 657.

### 16.2.11 Problems Due to Strong Typing

#### [ Problems Due to Strong Typing ]

All MPI functions with choice arguments associate actual arguments of different Fortran datatypes with the same dummy argument. This is not allowed by Fortran 77, and in Fortran 90, it is technically only allowed if the function is overloaded with a different function for each type (see also Section 16.2.6 on page 654). In C, the use of `void*` formal arguments avoids these problems. Similar to C, with Fortran 2008 + TR 29113 (and later) together with the `mpi_f08` module, the problem is avoided by declaring choice arguments with `TYPE(*)`, `DIMENSION(..)`, i.e., as assumed-type and assumed-rank dummy arguments.

[The] Using `INCLUDE mpif.h`, the following code fragment is technically [illegal]invalid and may generate a compile-time error.

```
integer i(5)
real    x(5)
...
call mpi_send(x, 5, MPI_REAL, ...)
call mpi_send(i, 5, MPI_INTEGER, ...)
```

In practice, it is rare for compilers to do more than issue a warning[, though there is concern that Fortran 90 compilers are more likely to return errors]. Using the `mpi_f08` or `mpi` module, the problem is usually resolved through the assumed-type and assumed-rank declarations of the dummy arguments, or with a compiler-dependent mechanism that overrides type checking for choice arguments.

It is also technically [illegal]invalid in Fortran to pass a scalar actual argument to an array dummy argument that is not a choice buffer argument. Thus, when using the `mpi_f08` or `mpi` module, the following code fragment[ may generate] usually generates an error since the[ buf argument] `dims` and `periods` arguments to `[MPI_SEND is declared as an assumed-size array <type> buf(*)]` `MPI_CART_CREATE` are declared as assumed size arrays `INTEGER :: DIMS(*)` and `LOGICAL :: PERIODS(*)`.

```
[integer a
call mpi_send(a, 1, MPI_INTEGER, ...) ]

USE mpi_f08      ! or  USE mpi
INTEGER size
CALL MPI_Cart_create( comm_old,1,size,.TRUE.,.TRUE.,comm_cart,ierror )
```

Although this is a non-conforming MPI call, compiler warnings are not expected (but may occur) when using `INCLUDE 'mpif.h'` and this include file does not use Fortran explicit interfaces.

*Advice to users.* In the event that you run into one of the problems related to type checking, you may be able to work around it by using a compiler flag, by compiling separately, or by using an MPI implementation with Extended Fortran Support as described in Section 16.2.3. An alternative that will usually work with variables local to a routine but not with arguments to a function or subroutine is to use the EQUIVALENCE statement to create another variable with a type accepted by the compiler. (*End of advice to users.*)

#### 16.2.12 Problems Due to Data Copying and Sequence Association with Subscript Triplets

##### Problems Due to Data Copying and Sequence Association

Arrays with subscript **triplets** describe Fortran subarrays with or without strides, e.g.,

```
REAL a(100,100,100)
CALL MPI_Send( a(11:17, 12:99:3, 1:100), 7*30*100, MPI_REAL, ...)
```

The handling of subscript triplets depends on the value of the constant `MPI_SUBARRAYS_SUPPORTED`:

- If `MPI_SUBARRAYS_SUPPORTED` equals `.TRUE.:`

Choice buffer arguments are declared as `TYPE(*), DIMENSION(...)`. For example, consider the following code fragment:

```
REAL s(100), r(100)
CALL MPI_Isend(s(1:100:5), 3, MPI_REAL, ..., rq, ierror)
CALL MPI_Wait(rq, status, ierror)
CALL MPI_Irecv(r(1:100:5), 3, MPI_REAL, ..., rq, ierror)
CALL MPI_Wait(rq, status, ierror)
```

In this case, the individual elements `s(1)`, `s(6)`, and `s(11)` are sent between the start of `MPI_ISEND` and the end of `MPI_WAIT` even though the compiled code will not copy `s(1:100:5)` to a real contiguous temporary scratch buffer. Instead, the compiled code will pass a descriptor to `MPI_ISEND` that allows MPI to operate directly on `s(1)`, `s(6)`, `s(11)`, ..., `s(96)`. The called `MPI_ISEND` routine will take only the first three of these elements due to the type signature “3, `MPI_REAL`”.

All nonblocking MPI functions (e.g., `MPI_ISEND`, `MPI_PUT`, `MPI_FILE_WRITE_ALL_BEGIN`) behave as if *the user-specified elements of choice buffers are copied to a contiguous scratch buffer in the MPI runtime environment*. All datatype descriptions (in the example above, “3, `MPI_REAL`”) read and store data from and to this virtual contiguous scratch buffer. Displacements in MPI derived datatypes are relative to the beginning of this virtual contiguous scratch buffer.

Upon completion of a nonblocking receive operation (e.g., when `MPI_WAIT` on a corresponding `MPI_Request` returns), it is as if the received data has been copied from the virtual contiguous scratch buffer back to the non-contiguous application buffer. In the example above, `r(1)`, `r(6)`, and `r(11)` are guaranteed to be defined with the received data when `MPI_WAIT` returns.

*Advice to implementors.* The Fortran descriptor for `TYPE(*)`, `DIMENSION(...)` arguments contains enough information that, if desired, the MPI library can make a real contiguous copy of non-contiguous user buffers when the nonblocking operation is started, and released this buffer not before the nonblocking communication has completed (e.g., in an MPI wait routine). Efficient implementations may avoid such additional memory-to-memory data copying. (*End of advice to implementors.*)

*Rationale.* If `MPI_SUBARRAYS_SUPPORTED` equals `.TRUE.`, non-contiguous buffers are handled inside of the MPI library instead of by the compiler through argument association conventions. Therefore, the scope of MPI library scratch buffers can be from the beginning of a nonblocking operation until the completion of the operation although beginning and completion are implemented in different routines. (*End of rationale.*)

- If `MPI_SUBARRAYS_SUPPORTED` equals `.FALSE.`:

Implicit in MPI is the idea of a contiguous chunk of memory accessible through a linear address space. MPI copies data to and from this memory. An MPI program specifies the location of data by providing memory addresses and offsets. In the C language, sequence association rules plus pointers provide all the necessary low-level structure.

In Fortran[ 90], `[user]array` data is not necessarily stored contiguously. For example, the array section `A(1:N:2)` involves only the elements of `A` with indices 1, 3, 5, ... . The same is true for a pointer array whose target is such a section. Most compilers ensure that an array that is a dummy argument is held in contiguous memory if it is declared with an explicit shape (e.g., `B(N)`) or is of assumed size (e.g., `B(*)`). If necessary, they do this by making a copy of the array into contiguous memory.[Both Fortran 77 and Fortran 90 are carefully worded to allow such copying to occur, but few Fortran 77 compilers do it.]<sup>1</sup>

Because MPI dummy buffer arguments are assumed-size arrays if `MPI_SUBARRAYS_SUPPORTED` equals `.FALSE.`, this leads to a serious problem for a nonblocking call: the compiler copies the temporary array back on return but MPI continues to copy data to the memory that held it. For example, consider the following code fragment:

```
real a(100)
call MPI_IRECV(a(1:100:2), MPI_REAL, 50, ...)
```

Since the first dummy argument to `MPI_IRECV` is an assumed-size array (`<type> buf(*)`), the array section `a(1:100:2)` is copied to a temporary before being passed

<sup>1</sup>Technically, the Fortran standard is worded to allow non-contiguous storage of any array data, unless the dummy argument has the `CONTIGUOUS` attribute.



to MPI\_IRecv, so that it is contiguous in memory. MPI\_IRecv returns immediately, and data is copied from the temporary back into the array `a`. Sometime later, MPI may write to the address of the deallocated temporary. Copying is also a problem for MPI\_Isend since the temporary array may be deallocated before the data has all been sent from it.

Most Fortran 90 compilers do not make a copy if the actual argument is the whole of an explicit-shape or assumed-size array or is a ‘[simple]simply contiguous’ section such as `A(1:N)` of such an array. ([We define ‘simple’ more fully]‘Simply contiguous’ is defined in the next paragraph.) Also, many compilers treat allocatable arrays the same as they treat explicit-shape arrays in this regard (though we know of one that does not). However, the same is not true for assumed-shape and pointer arrays; since they may be discontinuous, copying is often done. It is this copying that causes problems for MPI as described in the previous paragraph.

[Our formal definition of a ‘simple’]According to the Fortran 2008 Standard, Section 6.5.4, a ‘simply contiguous’ array section is

```
name ( [ :, ] ... [ <subscript> : [ <subscript> ] [, <subscript> ] ... )
```

That is, there are zero or more dimensions that are selected in full, then one dimension selected without a stride, then zero or more dimensions that are selected with a simple subscript. The compiler can detect from analyzing the source code that the array is contiguous. Examples are

```
A(1:N), A(:,N), A(:,1:N,1), A(1:6,N), A(:, :, 1:N)
```

Because of Fortran’s column-major ordering, where the first index varies fastest, a ‘[simple]simply contiguous’ section of a contiguous array will also be contiguous. [

*Footnote removed :*

To keep the definition of ‘simply contiguous’ simple, we have chosen to require all but one of the section subscripts to be without bounds. A colon without bounds makes it obvious both to the compiler and to the reader that the whole of the dimension is selected. It would have been possible to allow cases where the whole dimension is selected with one or two bounds, but this means for the reader that the array declaration or most recent allocation has to be consulted and for the compiler that a run-time check may be required.]

The same problem can occur with a scalar argument. [Some]A compiler[s, even for Fortran 77,] may make a copy of[ some] scalar dummy arguments within a called procedure when passed as an actual argument to a choice buffer routine. That this can cause a problem is illustrated by the example

```
[ticket236-H.]real :: a
call user1(a,rq)
call MPI_WAIT(rq,status,ierr)
write (*,*) a

subroutine user1(buf,request)
call MPI_IRecv(buf,...,request,...)
end
```

If `a` is copied, `MPI_IRECV` will alter the copy when it completes the communication and will not alter `a` itself.

Note that copying will almost certainly occur for an argument that is a non-trivial expression (one with at least one operator or function call), a section that does not select a contiguous part of its parent (e.g., `A(1:n:2)`), a pointer whose target is such a section, or an assumed-shape array that is (directly or indirectly) associated with such a section.

If `[there is]` a compiler option `exists` that inhibits copying of arguments, in either the calling or called procedure, this `[should]must` be employed.

If a compiler makes copies in the calling procedure of arguments that are explicit-shape or assumed-size arrays, `'[simple]simply contiguous'` array sections of such arrays, or scalars, and if `[there is no compiler option to inhibit this]no compiler option exists to inhibit such copying`, then the compiler cannot be used for applications that use `MPI_GET_ADDRESS`, or any nonblocking MPI routine. If a compiler copies scalar arguments in the called procedure and there is no compiler option to inhibit this, then this compiler cannot be used for applications that use memory references across subroutine calls as in the example above.

### 16.2.13 Problems Due to Data Copying and Sequence Association with Vector Subscripts

Fortran arrays with `vector` subscripts describe subarrays containing a possibly irregular set of elements

```
REAL a(100)
CALL MPI_Send( A((/7,9,23,81,82/)), 5, MPI_REAL, ...)
```

Arrays with a vector subscript must not be used as actual choice buffer arguments in any nonblocking or split collective MPI operations. They may, however, be used in blocking MPI operations.

### 16.2.14 Special Constants

`[`  
`Special Constants`

`]`  
MPI requires a number of special “constants” that cannot be implemented as normal Fortran constants, e.g., `MPI_BOTTOM`. The complete list can be found in Section 2.5.4 on page 15. In C, these are implemented as constant pointers, usually as `NULL` and are used where the function prototype calls for a pointer to a variable, not the variable itself.

In Fortran, `[the implementation of these special constants may require the use of language constructs that are outside the Fortran standard. Using ]using` special values for the constants (e.g., by defining them through `parameter` statements) is not possible because an implementation cannot distinguish these values from `[legal]valid` data. Typically these constants are implemented as predefined static variables (e.g., a variable in an MPI-declared `COMMON` block), relying on the fact that the target compiler passes data by address. Inside the subroutine, `[this address can be extracted by some mechanism outside the Fortran standard (e.g., by Fortran extensions or by implementing the function in C)]the address of the`

actual choice buffer argument can be compared with the address of such a predefined static variable.

These special constants also cause an exception with the usage of Fortran `INTENT`: with `USE mpi_f08`, the attributes `INTENT(IN)`, `INTENT(OUT)`, and `INTENT(INOUT)` are used in the Fortran interface. In most cases, `INTENT(IN)` is used if the C interface uses call-by-value. For all buffer arguments and for dummy arguments that may be modified and allow one of these special constants as input, an `INTENT` is not specified.

### 16.2.15 Fortran Derived Types

[  
Fortran 90 Derived Types  
]

MPI[ does not explicitly] supports passing Fortran[ 90] entities of `BIND(C)` and `SEQUENCE` derived types to choice dummy arguments, provided no type component has the `ALLOCATABLE` or `POINTER` attribute. [Indeed, for MPI implementations that provide explicit interfaces through the `mpi` module a compiler will reject derived type actual arguments at compile time. Even when no explicit interfaces are given, users should be aware that Fortran 90 provides no guarantee of sequence association for derived types or arrays of derived types. For instance, an array of a derived type consisting of two elements may be implemented as an array of the first elements followed by an array of the second. Use of the `SEQUENCE` attribute may help here, somewhat. ]

The following code fragment shows [one possible way to send a]some possible ways to send scalars or arrays of interoperable derived type in Fortran. The example assumes that all data is passed by address.

```

type[ticket237-I.], BIND(C) :: mytype
  integer [ticket229.2.]:: i
  real [ticket229.2.]:: x
  double precision [ticket229.2.]:: d
  [ticket229.2.]logical :: l
end type mytype

type(mytype) [ticket250-V.]:: foo[ticket237-I.], fooarr(5)
integer [ticket250-V.]:: blocklen(4), type(4)
integer(KIND=MPI_ADDRESS_KIND) [ticket250-V.]:: disp(4), base[ticket237-I.]
lb, extent

call MPI_GET_ADDRESS(foo%i, disp(1), ierr)
call MPI_GET_ADDRESS(foo%x, disp(2), ierr)
call MPI_GET_ADDRESS(foo%d, disp(3), ierr)
[ticket229.2.]call MPI_GET_ADDRESS(foo%l, disp(4), ierr)

base = disp(1)
disp(1) = disp(1) - base
disp(2) = disp(2) - base
disp(3) = disp(3) - base
[ticket229.2.]disp(4) = disp(4) - base

```

```

1
2     blocklen(1) = 1
3     blocklen(2) = 1
4     blocklen(3) = 1
5     [ticket229.2.]blocklen(4) = 1
6
7     type(1) = MPI_INTEGER
8     type(2) = MPI_REAL
9     type(3) = MPI_DOUBLE_PRECISION
10    [ticket229.2.]type(4) = MPI_LOGICAL
11
12    call MPI_TYPE_CREATE_STRUCT(4, blocklen, disp, type, newtype, ierr)
13    call MPI_TYPE_COMMIT(newtype, ierr)
14
15    [ticket237-I.] [! unpleasant to send foo%i instead of foo, but it works for scalar]
16    [ticket237-I.] [! entities of type mytype]
17    call MPI_SEND(foo%i, 1, newtype, ...)
18    [ticket237-I.] ! or
19    [ticket237-I.] call MPI_SEND(foo, 1, newtype, ...)
20    [ticket237-I.] ! expects that base == address(foo%i) == address(foo)
21
22    [ticket237-I.] call MPI_GET_ADDRESS(fooarr(1), disp(1), ierr)
23    [ticket237-I.] call MPI_GET_ADDRESS(fooarr(2), disp(2), ierr)
24    [ticket237-I.] extent = disp(2) - disp(1)
25    [ticket237-I.] lb = 0
26    [ticket237-I.] call MPI_TYPE_CREATE_RESIZED(newtype, lb, extent, newarrtype, ierr)
27    [ticket237-I.] call MPI_TYPE_COMMIT(newarrtype, ierr)
28    [ticket237-I.]
29    [ticket237-I.] call MPI_SEND(fooarr, 5, newarrtype, ...)

```

Using the derived type variable `foo` instead of its first basic type element `foo%i` may be impossible if the MPI library implements choice buffer arguments through overloading instead of using `TYPE(*)`, `DIMENSION(..)`, or through a non-standardized extensions such as `!$PRAGMA IGNORE_TKR`; see Section 16.2.6 on page 654. [\[1\]](#)

To use a derived type in an array requires a correct extent of the datatype handle to take care of the alignment rules applied by the compiler. These alignment rules may imply that there are gaps between the elements of a derived type, and also between the array elements. [\[2\]](#) The extent of an iteroperable derived type (i.e., defined with `BIND(C)`) and a `SEQUENCE` derived type with the same content may be different because C and Fortran may apply different alignment rules. [\[mytype is a SEQUENCE derived type.\]](#) As recommended in the advice to users in Section 4.1 on page 91, one should add an additional fifth structure element with one numerical storage unit at the end of this structure to force in most cases that the array of structures is contiguous. Even with such an additional element, one should keep this resizing due to the special alignment rules that can be used by the compiler for structures, as also mentioned in this advice.

Using the extended semantics defined in TR 29113, it is also possible to use entities or derived types without either the `BIND(C)` or the `SEQUENCE` attribute as choice buffer arguments; some additional constraints must be observed e.g., no `ALLOCATABLE` or `POINTER`

type components may exist. In this case, the `base` address in the example must be changed to become the address of `foo` instead of `foo%i`, because the Fortran compiler may rearrange type components or add padding as it may fit for such types. Sending the structure `foo` should then also be performed by providing it (and not `foo%i`) as actual argument for `MPI_Send`.

### 16.2.16 Optimization Problems, an Overview

[  
A Problem with Register Optimization  
]

MPI provides operations that may be hidden from the user code and run concurrently with it, accessing the same memory as user code. Examples include the data transfer for an `MPI_IRECV`. The optimizer of a compiler will assume that it can recognize periods when a copy of a variable can be kept in a register without reloading from or storing to memory. When the user code is working with a register copy of some variable while the hidden operation reads or writes the memory copy, problems occur. [ This section discusses register optimization pitfalls.] These problems are independent of the Fortran support method; i.e., they occur with the `mpi_f08` module, the `mpi` module, and the `mpif.h` include file.

This section shows four problematic usage areas (the abbreviations in parentheses are used in the table below):

- Use of nonblocking routines or persistent requests (*Nonbl.*).
- Use of one-sided routines (*1-sided*).
- Use of MPI parallel file I/O split collective operations (*Split*).
- Use of `MPI_BOTTOM` together with absolute displacements in MPI datatypes, or relative displacements between two variables in such datatypes (*Bottom*).

The following compiler optimization strategies (valid for serial code) may cause problems in MPI applications:

- Code movement and register optimization problems; see Section 16.2.17 on page 678.
- Temporary data movement and temporary memory modifications; see Section 16.2.18 on page 685.
- Permanent data movement (e.g., through garbage collection); see Section 16.2.19 on page 686.

Table 16.4 shows in which usage areas the optimization problems may only occur.

The solutions in the following sections are based on compromises:

- to minimize the burden for the application programmer, e.g., as shown in Sections “Solutions” to “VOLATILE” on pages 680-681,
- to minimize the drawbacks on compiler based optimization, and
- to minimize the requirements defined in Section 16.2.7 on page 657.

Optimization ...	... may cause a problem in following usage areas			
	Nonbl.	1-sided	Split	Bottom
Code movement and register optimization	yes	yes	no	yes
Temporary data movement	yes	yes	yes	no
Permanent data movement	yes	yes	yes	yes

Table 16.4: Occurrence of Fortran optimization problems in several usage areas

16.2.17 Problems with Code Movement and Register Optimization

Nonblocking operations

If a variable is local to a Fortran subroutine (i.e., not in a module or a COMMON block), the compiler will assume that it cannot be modified by a called subroutine unless it is an actual argument of the call. In the most common linkage convention, the subroutine is expected to save and restore certain registers. Thus, the optimizer will assume that a register which held a valid copy of such a variable before the call will still hold a valid copy on return.

**Example 16.11** Fortran 90 register optimization – extreme.

Source	compiled as	or compiled as
[ticket238-J.] REAL :: buf, b1 call MPI_Irecv(buf,..req)  call MPI_WAIT(req,..) b1 = buf	 REAL :: buf, b1 call MPI_Irecv(buf,..req) register = buf call MPI_WAIT(req,..) b1 = register	 REAL :: buf, b1 call MPI_Irecv(buf,..req) b1 = buf call MPI_WAIT(req,..)

Example 16.11 shows extreme, but allowed, possibilities. MPI\_WAIT on a concurrent thread modifies buf between the invocation of MPI\_Irecv and the finish of MPI\_WAIT. But the compiler cannot see any possibility that buf can be changed after MPI\_Irecv has returned, and may schedule the load of buf earlier than typed in the source. The compiler has no reason to avoid using a register to hold buf across the call to MPI\_WAIT. It also may reorder the instructions as [in the case on the right]illustrated in the rightmost column.

[ticket238-J.]

**Example 16.12** Similar example with MPI\_Isend

Source	compiled as	with a possible MPI-internal execution sequence
REAL :: buf, copy buf = val call MPI_Isend(buf,..req) copy = buf  call MPI_WAIT(req,..) buf = val_overwrite	 REAL :: buf, copy buf = val call MPI_Isend(buf,..req) copy= buf buf = val_overwrite call MPI_WAIT(req,..)	 REAL :: buf, copy buf = val addr = &buf copy= buf buf = val_overwrite send(*addr) ! within MPI_WAIT

Due to valid compiler code movement optimizations in Example 16.12, the content of `buf` may already be overwritten by the compiler when the content of `buf` is sent. The code movement is permitted because the compiler cannot detect a possible access to `buf` in `MPI_WAIT` (or in a second thread between the start of `MPI_ISEND` and the end of `MPI_WAIT`).

Such register optimization is based on moving code; here, the access to `buf` was moved from after `MPI_WAIT` to before `MPI_WAIT`. Note that code movement may also occur across subroutine boundaries when subroutines or functions are inlined.

This register optimization / code movement problem for nonblocking operations does not occur with MPI parallel file I/O split collective operations, because in the `..._BEGIN` and `..._END` calls, the same buffer has to be provided as an actual argument. The register optimization / code movement problem for `MPI_BOTTOM` and derived MPI datatypes may occur in each blocking and nonblocking communication or parallel file I/O operation.

### One-sided communication

An example with instruction reordering due to register optimization can be found in Section 11.7.4 on page 480.

### MPI\_BOTTOM and combining independent variables in datatypes

[Normally users are not afflicted with this. But the user should pay attention to this section if in his/her program ]This section is only relevant if the MPI program uses a buffer argument to an `MPI_SEND`, `MPI_RECV` etc., [ uses a name ] which hides the actual variables involved. `MPI_BOTTOM` with an `MPI_Datatype` containing absolute addresses is one example. Creating a datatype which uses one variable as an anchor and brings along others by using `MPI_GET_ADDRESS` to determine their offsets from the anchor is another. The anchor variable would be the only one [mentioned]referenced in the call. Also attention must be paid if MPI operations are used that run in parallel with the user's application.

Example 16.13 shows what Fortran compilers are allowed to do.

#### Example 16.13 Fortran 90 register optimization.

This source ...	can be compiled as:
<code>call MPI_GET_ADDRESS(buf,bufaddr,</code>	<code>call MPI_GET_ADDRESS(buf,...)</code>
<code>                  ierror)</code>	
<code>call MPI_TYPE_CREATE_STRUCT(1,1,</code>	<code>call MPI_TYPE_CREATE_STRUCT(...)</code>
<code>                  bufaddr,</code>	
<code>                  MPI_REAL,type,ierror)</code>	
<code>call MPI_TYPE_COMMIT(type,ierror)</code>	<code>call MPI_TYPE_COMMIT(...)</code>
<code>val_old = buf</code>	<code>register = buf</code>
	<code>val_old = register</code>
<code>call MPI_RECV(MPI_BOTTOM,1,type,...)</code>	<code>call MPI_RECV(MPI_BOTTOM,...)</code>
<code>val_new = buf</code>	<code>val_new = register</code>

[The]In Example 16.13, the compiler does not invalidate the register because it cannot see that `MPI_RECV` changes the value of `buf`. The access [of]to `buf` is hidden by the use of `MPI_GET_ADDRESS` and `MPI_BOTTOM`.



[ticket238-J.]

### Example 16.14 Similar example with MPI\_SEND

<p>This source ...</p> <pre> ! buf contains val_old buf = val_new call MPI_SEND(MPI_BOTTOM,1,type,...) ! with buf as a displacement in type  buf = val_overwrite </pre>	<p>can be compiled as:</p> <pre> ! buf contains val_old call MPI_SEND(...) ! i.e. val_old is sent ! ! buf=val_new is moved to here ! and detected as dead code ! and therefore removed ! buf = val_overwrite </pre>
-------------------------------------------------------------------------------------------------------------------------------------------------------------------------	---------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------

In Example 16.14, several successive assignments to the same variable `buf` can be combined in a way such that only the last assignment is executed. “Successive” means that no interfering load access to this variable occurs between the assignments. The compiler cannot detect that the call to `MPI_SEND` statement is interfering because the load access to `buf` is hidden by the usage of `MPI_BOTTOM`.

### Solutions

The following sections show in detail how the problems with code movement and register optimization can be solved in a portable way. Application writers can partially or fully avoid these compiler optimization problems by using one or more of the special Fortran declarations with the send and receive buffers used in nonblocking operations, or in operations in which `MPI_BOTTOM` is used, or datatype handles that combine several variables are used:

- Use of the Fortran `ASYNCHRONOUS` attribute.
- Use of the helper routine `MPI_F_SYNC_REG`, or an equivalent user-written dummy routine.
- Declare the buffer as a Fortran module variable or within a Fortran common block.
- Use of the Fortran `VOLATILE` attribute.

Each of these methods solves the problems of code movement and register optimization, but may involve different degrees of performance impact, and may not be usable in every application context. These methods may not be guaranteed by the Fortran standard, but they must be guaranteed by a MPI-3.0 compliant (and later) MPI library and their compiler according to the requirements listed in Section 16.2.7 on page 657. The methods may have different impact on performance. `MPI_F_SYNC_REG` may have low impact, module data and the `ASYNCHRONOUS` attribute low through medium, and the `VOLATILE` attribute may have the most negative impact on performance. Note that there is one attribute that cannot be used for this purpose: the Fortran `TARGET` attribute does not solve code movement problems in MPI applications.



### The Fortran ASYNCHRONOUS attribute

Declaring an actual buffer argument with the **ASYNCHRONOUS** Fortran attribute in a scoping unit (or **BLOCK**) tells the compiler that any statement in the scoping unit may be executed while the buffer is affected by a pending asynchronous Fortran input/output operation (since Fortran 2003) or by an asynchronous communication (TR 29113 extension). Without the extensions specified in TR 29113, a Fortran compiler may totally ignore this attribute if the Fortran compiler implements asynchronous Fortran input/output operations with blocking I/O. The **ASYNCHRONOUS** attribute protects the buffer accesses from optimizations through code movements across routine calls, and the buffer itself from temporary and permanent data movements. If the choice buffer dummy argument of a nonblocking MPI routine is declared with **ASYNCHRONOUS** (which is mandatory for the `mpi_f08` module, with allowable exceptions listed in Section 16.2.6 on page 654), then the compiler has to guarantee call by reference and should report a compile-time error if call by reference is impossible, e.g., if vector subscripts are used. The `MPI_ASYNC_PROTECTS_NONBLOCKING` is set to `.TRUE.` if both the protection of the actual buffer argument through **ASYNCHRONOUS** according to the TR 29113 extension and the declaration of the dummy argument with **ASYNCHRONOUS** in the Fortran support method is guaranteed for all nonblocking routines, otherwise it is set to `.FALSE.`

The **ASYNCHRONOUS** attribute has some restrictions. The TR 29113 defines (in the PDTR N1869):

“Asynchronous communication for a Fortran variable occurs through the action of procedures defined by means other than Fortran. It is initiated by execution of an asynchronous communication initiation procedure and completed by execution of an asynchronous communication completion procedure. Between the execution of the initiation and completion procedures, any variable of which any part is associated with any part of the asynchronous communication variable is a pending communication affector. Whether a procedure is an asynchronous communication initiation or completion procedure is processor dependent. Asynchronous communication is either input communication or output communication. For input communication, a pending communication affector shall not be referenced, become defined, become undefined, become associated with a dummy argument that has the **VALUE** attribute, or have its pointer association status changed. For output communication, a pending communication affector shall not be redefined, become undefined, or have its pointer association status changed.”

In Example 16.15 Case (a) on page 688, the read accesses to **b** within `function(b(i-1), b(i), b(i+1))` cannot be moved by compiler optimizations to before the wait call because **b** was declared as **ASYNCHRONOUS**. Note that only the elements 0, 1, 100, and 101 of **b** are involved in asynchronous communication but by definition, the total variable **b** is the pending communication affector and is usable for input and output asynchronous communication between the `MPI_I...` routines and `MPI_Waitall`. Case (a) works fine because the read accesses to **b** occur after the communication completed.

In Case (b), the read accesses to `b(1:100)` in the loop `i=2,99` are read accesses to a pending communication affector while input communication (i.e., the two `MPI_Irecv` calls) is pending. This is a contradiction to the rule that *for input communication, a pending communication affector shall not be referenced*. The problem can be solved by using separate

variables for the halos and the inner array, or by splitting a common array into disjunct subarrays which are passed through different dummy arguments into a subroutine, as shown in Example 16.19 on page 690.

If one does not overlap communication and computation on the same variable, then all optimization problems can be solved through the `ASYNCHRONOUS` attribute.

The problems with `MPI_BOTTOM`, as shown in Example 16.13 and Example 16.14, can also be solved by declaring the buffer `buf` with the `ASYNCHRONOUS` attribute.

In some MPI routines, a buffer dummy argument is defined as `ASYNCHRONOUS` to guarantee passing by reference, provided that the actual argument is also defined as `ASYNCHRONOUS`.

[(Example 16.11 and its following paragraph were moved to an earlier position)]

[To prevent instruction reordering or the allocation of a buffer in a register there are two possibilities in portable Fortran code: ]

### Calling `MPI_F_SYNC_REG`

The compiler may be prevented from moving a reference to a buffer across a call to an MPI subroutine by surrounding the call by calls to an external subroutine with the buffer as an actual argument. The MPI library provides the `MPI_F_SYNC_REG` routine for this purpose; see Section 16.2.8 on page 659.

□

- The problems illustrated by the Examples 16.11 and 16.12 can be solved by calling `MPI_F_SYNC_REG(buf)` once immediately after `MPI_WAIT`.

#### Example 16.11

can be solved with

```
call MPI_Irecv(buf,...req)
```

```
call MPI_WAIT(req,...)
```

```
call MPI_F_SYNC_REG(buf)
```

```
b1 = buf
```

#### Example 16.12

can be solved with

```
buf = val
```

```
call MPI_Isend(buf,...req)
```

```
copy = buf
```

```
call MPI_WAIT(req,...)
```

```
call MPI_F_SYNC_REG(buf)
```

```
buf = val_overwrite
```

The call to `MPI_F_SYNC_REG(buf)` prevents moving the last line before the `MPI_WAIT` call. Further calls to `MPI_F_SYNC_REG(buf)` are not needed because it is still correct if the additional read access `copy=buf` is moved below `MPI_WAIT` and before `buf=val_overwrite`.

- The problems illustrated by the Examples 16.13 and 16.14 can be solved with two additional `MPI_F_SYNC_REG(buf)` statements; one directly before `MPI_RECV/`  
`MPI_SEND`, and one directly after this communication operation.

#### Example 16.13

can be solved with

```
call MPI_F_SYNC_REG(buf)
```

```
call MPI_RECV(MPI_BOTTOM,...)
```

```
call MPI_F_SYNC_REG(buf)
```

#### Example 16.14

can be solved with

```
call MPI_F_SYNC_REG(buf)
```

```
call MPI_SEND(MPI_BOTTOM,...)
```

```
call MPI_F_SYNC_REG(buf)
```

The first call to `MPI_F_SYNC_REG(buf)` is needed to finish all load and store references to `buf` prior to `MPI_RECV/MPI_SEND`; the second call is needed to assure that the subsequent access to `buf` are not moved before `MPI_RECV/SEND`.

- In the example in Section 11.7.4 on page 480, two asynchronous accesses must be protected: in Process 1, the access to `bbbb` must be protected similar to Example 16.11, i.e., a call to `MPI_F_SYNC_REG(bbbb)` is needed after the second `MPI_WIN_FENCE` to guarantee that further accesses to `bbbb` are not moved ahead of the call to `MPI_WIN_FENCE`. In Process 2, both calls to `MPI_WIN_FENCE` together act as a communication call with `MPI_BOTTOM` as the buffer. That is, before the first fence and after the second fence, a call to `MPI_F_SYNC_REG(buff)` is needed to guarantee that accesses to `buff` are not moved after or ahead of the calls to `MPI_WIN_FENCE`. Using `MPI_GET` instead of `MPI_PUT`, the same calls to `MPI_F_SYNC_REG` are necessary.

#### Source of Process 1

```
bbbb = 777

call MPI_WIN_FENCE
call MPI_PUT(bbbb
into buff of process 2)
```

```
call MPI_WIN_FENCE
call MPI_F_SYNC_REG(bbbb)
```

#### Source of Process 2

```
buff = 999
call MPI_F_SYNC_REG(buff)
call MPI_WIN_FENCE

call MPI_WIN_FENCE
call MPI_F_SYNC_REG(buff)
ccc = buff
```

- The temporary memory modification problem, i.e., Example 16.16 on page 689, can **not** be solved with this method.

□

#### A user defined routine instead of `MPI_F_SYNC_REG`

Instead of `MPI_F_SYNC_REG`, one can also use a user defined external subroutine, which is separately compiled: [with the separately compiled ]

```
subroutine DD(buf)
  integer buf
end
```

Note that if the intent is declared in an explicit interface for the external subroutine, it must be `OUT` or `INOUT`. The subroutine itself may have an empty body, but the compiler does not know this and has to assume that the buffer may be altered. For example, [the above] a call [of] to `MPI_RECV` with `MPI_BOTTOM` as buffer might be replaced by

```
call DD(buf)
call MPI_RECV(MPI_BOTTOM,...)
call DD(buf)
```

ticket238-J.

Such a user-defined routine was introduced in MPI-2.0 and is still included here to document such usage in existing application programs although new applications should prefer `MPI_F_SYNC_REG` or one of the other possibilities. In an existing application, calls to such a user-written routine should be substituted by a call to `MPI_F_SYNC_REG` because the user-written routine may not be implemented according to the rules specified in Section 16.2.7 on page 657.

[(assuming that `buf` has type `INTEGER`). The compiler may be similarly prevented from moving a reference to a variable across a call to an MPI subroutine.

In the case of a nonblocking call, as in the above call of `MPI_WAIT`, no reference to the buffer is permitted until it has been verified that the transfer has been completed. Therefore, in this case, the extra call ahead of the MPI call is not necessary, i.e., the call of `MPI_WAIT` in the example might be replaced by

```
/tt call MPI_WAIT(req,..)
```

```
call DD(buf) ]
```

```
[]
```

### Module variables and COMMON blocks

An alternative to the already mentioned methods is to put the buffer or variable into a module or a common block and access it through a `USE` or `COMMON` statement in each scope where it is referenced, defined or appears as an actual argument in a call to an MPI routine. The compiler will then have to assume that the MPI procedure [(MPI\_RECV in the above example) ] may alter the buffer or variable, provided that the compiler cannot [analyze]infer that the MPI procedure does not reference the module or common block.

- This method solves problems of instruction reordering, code movement, and register optimization related to nonblocking and one-sided communication, or related to the usage of `MPI_BOTTOM` and derived datatype handles.
- Unfortunately, this method does **not** solve problems caused by asynchronous accesses between the start and end of a nonblocking or one-sided communication. Specifically, problems caused by temporary memory modifications are not solved.

```
[]
```

### The (poorly performing) Fortran VOLATILE attribute

The `VOLATILE` attribute[, available in later versions of Fortran,] gives the buffer or variable the properties needed, but it may inhibit optimization of any code containing references or definitions of the buffer or variable.

### The Fortran TARGET attribute

The `TARGET` attribute does not solve the code movement problem because it is not specified for the choice buffer dummy arguments of nonblocking routines. If the compiler detects that the application program specifies the `TARGET` attribute for an actual buffer argument used in the call to a nonblocking routine, the compiler may ignore this attribute if no pointer reference to this buffer exists.

*Rationale.* The Fortran standardization body decided to extend the `ASYNCHRONOUS` attribute within the TR 29113 to protect buffers in nonblocking calls from all kinds of optimization, instead of extending the `TARGET` attribute. (*End of rationale.*)

### 16.2.18 Temporary Data Movement and Temporary Memory Modification

The compiler is allowed to temporarily modify data in memory. Normally, this problem may occur only when overlapping communication and computation, as in Example 16.15, Case (b) on page 688. Example 16.16 on page 689 shows a possibility that could be problematic.

In the compiler-generated, possible optimization in Example 16.17, `buf(100,100)` from Example 16.16 is equivalenced with the 1-dimensional array `buf_1dim(10000)`. The nonblocking receive may asynchronously receive the data in the boundary `buf(1,1:100)` while the fused loop is temporarily using this part of the buffer. When the `tmp` data is written back to `buf`, the previous data of `buf(1,1:100)` is restored and the received data is lost. The principle behind this optimization is that the receive buffer data `buf(1,1:100)` was temporarily moved to `tmp`.

Example 16.18 shows a second possible optimization. The whole array is temporarily moved to `local_buf`. When storing `local_buf` back to the original location `buf`, then this includes also an overwriting of the receive buffer part `buf(1,1:100)`, i.e., this storing back may overwrite the asynchronously received data.

Note, that this problem may also occur:

- With the local buffer at the origin process, between an RMA communication call and the ensuing synchronization call; see Chapter 11 on page 419.
- With the window buffer at the target process between two ensuing RMA synchronization calls.
- With the local buffer in MPI parallel file I/O split collective operations with between the `..._BEGIN` and `..._END` calls; see Section 13.4.5 on page 546.

As already mentioned in subsection *The Fortran ASYNCHRONOUS attribute* on page 681 in Section 16.2.17 on page 678, the `ASYNCHRONOUS` attribute can prevent compiler optimization with temporary data movement, but only if the receive buffer and the numerical read accesses are separated into different variables, as shown in Example 16.19 on page 690 and in Example 16.20 on page 691.

Note also that the methods

- calling `MPI_F_SYNC_REG` (or such a user-defined routine),
- using module variables and `COMMON` blocks, and
- the `TARGET` attribute

cannot be used to prevent such temporary data movement. These methods influence compiler optimization when library routines are called. They cannot prevent the optimizations of the numerical code shown in Example 16.16 and 16.17.

Note also that compiler optimization with temporary data movement should **not** be prevented by declaring `buf` as `VOLATILE` because the `VOLATILE` implies that all accesses to

any storage unit (word) of `buf` must be directly done in the main memory exactly in the sequence defined by the application program. The `VOLATILE` attribute prevents all register and cache optimizations. Therefore, `VOLATILE` may cause a huge performance degradation.

Instead of solving the problem, it is needed to **prevent** the problem. When overlapping communication and computation, the nonblocking communication (or nonblocking or split collective IO) and the computation should be executed **on different sets of variables**. In this case, the temporary memory modifications are done only on the variables used in the computation and cannot have any side effect on the data used in the nonblocking MPI operations.

*Rationale.* This is a strong restriction for application programs. To weaken this restriction, a new or modified asynchronous feature in the Fortran language would be necessary: an asynchronous attribute that can be used on parts of an array and together with asynchronous operations outside the scope of Fortran. If such a feature is available in a later version of the Fortran standard, then this restriction also may be weakened in a later version of the MPI standard. (*End of rationale.*)

In Example 16.19 on page 690 (which is a solution for the problem shown in Example 16.15 on page 688) and in Example 16.20 on page 691 (which is a solution for the problem shown in Example 16.18 on page 689), the array is split into inner and halo part and both disjunct parts are passed to a subroutine `separated_sections`. This routine overlaps the receiving of the halo data and the calculations on the inner part of the array. In a second step, the whole array is used to do the calculation on the elements where inner+halo is needed. Note that the halo and the inner area are strided arrays. Those can be used in non-blocking communication only with a TR 29113 based MPI library.

### 16.2.19 Permanent Data Movement

A Fortran compiler may implement permanent data movement during the execution of a Fortran program. This would require that pointers to such data are appropriately updated. Automatic garbage collection implementation is one use case. Such permanent data movement is in conflict with MPI in several areas:

- MPI datatype handles with absolute addresses in combination with `MPI_BOTTOM`.
- Nonblocking MPI operations (communication, one-sided, I/O) if the internally used pointers to the buffers are not updated by the Fortran runtime, or if within an MPI process, the data movement is executed in parallel with the MPI operation.

This problem can be also solved by using the `ASYNCHRONOUS` attribute for such buffers. This MPI standard requires that the problems with permanent data movement do not occur by imposing suitable restrictions on the MPI library together with the compiler used; see Section 16.2.7 on page 657.

### 16.2.20 Comparison with C

In C, subroutines which modify variables that are not in the argument list will not cause register optimization problems. This is because taking pointers to storage objects by using the `&` operator and later referencing the objects by way of the pointer is an integral part of the language. A C compiler understands the implications, so that the problem should not

occur, in general. However, some compilers do offer optional aggressive optimization levels which may not be safe. Problems due to temporary memory modifications can also occur in C. As above, the best advice is to avoid the problem: use different variables for buffers in nonblocking MPI operations and computation that is executed while the nonblocking operations are pending.

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[ticket238-J.]

**Example 16.15** Protecting nonblocking communication with the `ASYNCHRONOUS` attribute.

```

USE mpi_f08
REAL, ASYNCHRONOUS :: b(0:101) ! elements 0 and 101 are halo cells
REAL :: bnew(0:101)           ! elements 1 and 100 are newly computed
TYPE(MPI_Request) :: req(4)
INTEGER :: left, right, i
CALL MPI_Cart_shift(...,left,right,...)
CALL MPI_Irecv(b( 0), ..., left, ..., req(1), ...)
CALL MPI_Irecv(b(101), ..., right, ..., req(2), ...)
CALL MPI_Isend(b( 1), ..., left, ..., req(3), ...)
CALL MPI_Isend(b(100), ..., right, ..., req(4), ...)

#ifdef WITHOUT_OVERLAPPING_COMMUNICATION_AND_COMPUTATION
! Case (a)
CALL MPI_Waitall(4,req,...)
DO i=1,100 ! compute all new local data
  bnew(i) = function(b(i-1), b(i), b(i+1))
END DO
#endif

#ifdef WITH_OVERLAPPING_COMMUNICATION_AND_COMPUTATION
! Case (b)
DO i=2,99 ! compute only elements for which halo data is not needed
  bnew(i) = function(b(i-1), b(i), b(i+1))
END DO
CALL MPI_Waitall(4,req,...)
i=1 ! compute leftmost element
  bnew(i) = function(b(i-1), b(i), b(i+1))
i=100 ! compute rightmost element
  bnew(i) = function(b(i-1), b(i), b(i+1))
#endif

```



[ticket238-J.]

**Example 16.16** Overlapping Communication and Computation.

```
USE mpi_f08
REAL :: buf(100,100)
CALL MPI_Irecv(buf(1,1:100),...req,...)
DO j=1,100
  DO i=2,100
    buf(i,j)=....
  END DO
END DO
CALL MPI_Wait(req,...)
```

[ticket238-J.]

**Example 16.17** The compiler may substitute the nested loops through loop fusion.

```
REAL :: buf(100,100), buf_1dim(10000)
EQUIVALENCE (buf(1,1), buf_1dim(1))
CALL MPI_Irecv(buf(1,1:100),...req,...)
tmp(1:100) = buf(1,1:100)
DO j=1,10000
  buf_1dim(h)=...
END DO
buf(1,1:100) = tmp(1:100)
CALL MPI_Wait(req,...)
```

[ticket238-J.]

**Example 16.18** Another optimization is based on the usage of a separate memory storage area, e.g., in a GPU.

```
REAL :: buf(100,100), local_buf(100,100)
CALL MPI_Irecv(buf(1,1:100),...req,...)
local_buf = buf
DO j=1,100
  DO i=2,100
    local_buf(i,j)=....
  END DO
END DO
buf = local_buf ! may overwrite asynchronously received
                ! data in buf(1,1:100)
CALL MPI_Wait(req,...)
```

[ticket238-J.]

**Example 16.19** Using separated variables for overlapping communication and computation to allow the protection of nonblocking communication with the `ASYNCHRONOUS` attribute.

```

USE mpi_f08
REAL :: b(0:101)      ! elements 0 and 101 are halo cells
REAL :: bnew(0:101)   ! elements 1 and 100 are newly computed
INTEGER :: i
CALL separated_sections(b(0), b(1:100), b(101), bnew(0:101))
i=1 ! compute leftmost element
  bnew(i) = function(b(i-1), b(i), b(i+1))
i=100 ! compute rightmost element
  bnew(i) = function(b(i-1), b(i), b(i+1))
END

SUBROUTINE separated_sections(b_lefthalo, b_inner, b_righthalo, bnew)
USE mpi_f08
REAL, ASYNCHRONOUS :: b_lefthalo(0:0), b_inner(1:100), b_righthalo(101:101)
REAL :: bnew(0:101) ! elements 1 and 100 are newly computed
TYPE(MPI_Request) :: req(4)
INTEGER :: left, right, i
CALL MPI_Cart_shift(...,left,right,...)
CALL MPI_Irecv(b_lefthalo ( 0), ..., left, ..., req(1), ...)
CALL MPI_Irecv(b_righthalo(101), ..., right, ..., req(2), ...)
! b_lefthalo and b_righthalo is written asynchronously.
! There is no other concurrent access to b_lefthalo and b_righthalo.
CALL MPI_Isend(b_inner( 1), ..., left, ..., req(3), ...)
CALL MPI_Isend(b_inner(100), ..., right, ..., req(4), ...)

DO i=2,99 ! compute only elements for which halo data is not needed
  bnew(i) = function(b_inner(i-1), b_inner(i), b_inner(i+1))
  ! b_inner is read and send at the same time.
  ! This is allowed based on the rules for ASYNCHRONOUS.
END DO
CALL MPI_Waitall(4,req,...)
END SUBROUTINE

```

[ticket238-J.]

**Example 16.20** Protecting GPU optimizations with the ASYNCHRONOUS attribute.

```
USE mpi_f08
REAL :: buf(100,100)
CALL separated_sections(buf(1:1,1:100), buf(2:100,1:100))
END

SUBROUTINE separated_sections(buf_halo, buf_inner)
REAL, ASYNCHRONOUS :: buf_halo(1:1,1:100)
REAL :: buf_inner(2:100,1:100)
REAL :: local_buf(2:100,100)

CALL MPI_Irecv(buf_halo(1,1:100),...req,...)
local_buf = buf_inner
DO j=1,100
  DO i=2,100
    local_buf(i,j)=....
  END DO
END DO
buf_inner = local_buf ! buf_halo is not touched!!!

CALL MPI_Wait(req,...)
```

## 16.3 Language Interoperability

### 16.3.1 Introduction

It is not uncommon for library developers to use one language to develop an applications library that may be called by an application program written in a different language. MPI currently supports ISO (previously ANSI) C, C++, and Fortran bindings. It should be possible for applications in any of the supported languages to call MPI-related functions in another language.

Moreover, MPI allows the development of client-server code, with MPI communication used between a parallel client and a parallel server. It should be possible to code the server in one language and the clients in another language. To do so, communications should be possible between applications written in different languages.

There are several issues that need to be addressed in order to achieve interoperability.

**Initialization** We need to specify how the MPI environment is initialized for all languages.

**Interlanguage passing of MPI opaque objects** We need to specify how MPI object handles are passed between languages. We also need to specify what happens when an MPI object is accessed in one language, to retrieve information (e.g., attributes) set in another language.

**Interlanguage communication** We need to specify how messages sent in one language can be received in another language.

It is highly desirable that the solution for interlanguage interoperability be extensible to new languages, should MPI bindings be defined for such languages.

### 16.3.2 Assumptions

We assume that conventions exist for programs written in one language to call routines written in another language. These conventions specify how to link routines in different languages into one program, how to call functions in a different language, how to pass arguments between languages, and the correspondence between basic data types in different languages. In general, these conventions will be implementation dependent. Furthermore, not every basic datatype may have a matching type in other languages. For example, C/C++ character strings may not be compatible with Fortran `CHARACTER` variables. However, we assume that a Fortran `INTEGER`, as well as a (sequence associated) Fortran array of `INTEGER`s, can be passed to a C or C++ program. We also assume that Fortran, C, and C++ have address-sized integers. This does not mean that the default-size integers are the same size as default-sized pointers, but only that there is some way to hold (and pass) a C address in a Fortran integer. It is also assumed that `INTEGER(KIND=MPI_OFFSET_KIND)` can be passed from Fortran to C as `MPI_Offset`.

### 16.3.3 Initialization

A call to `MPI_INIT` or `MPI_INIT_THREAD`, from any language, initializes MPI for execution in all languages.

*Advice to users.* Certain implementations use the (inout) `argc`, `argv` arguments of the C/C++ version of `MPI_INIT` in order to propagate values for `argc` and `argv` to

all executing processes. Use of the Fortran version of MPI\_INIT to initialize MPI may result in a loss of this ability. (*End of advice to users.*)

The function MPI\_INITIALIZED returns the same answer in all languages.

The function MPI\_FINALIZE finalizes the MPI environments for all languages.

The function MPI\_FINALIZED returns the same answer in all languages.

The function MPI\_ABORT kills processes, irrespective of the language used by the caller or by the processes killed.

The MPI environment is initialized in the same manner for all languages by MPI\_INIT. E.g., MPI\_COMM\_WORLD carries the same information regardless of language: same processes, same environmental attributes, same error handlers.

Information can be added to info objects in one language and retrieved in another.

*Advice to users.* The use of several languages in one MPI program may require the use of special options at compile and/or link time. (*End of advice to users.*)

*Advice to implementors.* Implementations may selectively link language specific MPI libraries only to codes that need them, so as not to increase the size of binaries for codes that use only one language. The MPI initialization code need perform initialization for a language only if that language library is loaded. (*End of advice to implementors.*)

#### 16.3.4 Transfer of Handles

Handles are passed between Fortran and C or C++ by using an explicit C wrapper to convert Fortran handles to C handles. There is no direct access to C or C++ handles in Fortran. Handles are passed between C and C++ using overloaded C++ operators called from C++ code. There is no direct access to C++ objects from C.

The type definition MPI\_Fint is provided in C/C++ for an integer of the size that matches a Fortran INTEGER; usually, MPI\_Fint will be equivalent to int. **With the Fortran mpi module or the mpif.h include file, a Fortran handle is a Fortran INTEGER value that can be used in the following conversion functions. With the Fortran mpi\_f08 module, a Fortran handle is a BIND(C) derived type that contains an INTEGER field named MPI\_VAL. This INTEGER value can be used in the following conversion functions.**

The following functions are provided in C to convert from a Fortran communicator handle (which is an integer) to a C communicator handle, and vice versa. See also Section 2.6.5 on page 23.

```
MPI_Comm MPI_Comm_f2c(MPI_Fint comm)
```

If comm is a valid Fortran handle to a communicator, then MPI\_Comm\_f2c returns a valid C handle to that same communicator; if comm = MPI\_COMM\_NULL (Fortran value), then MPI\_Comm\_f2c returns a null C handle; if comm is an invalid Fortran handle, then MPI\_Comm\_f2c returns an invalid C handle.

```
MPI_Fint MPI_Comm_c2f(MPI_Comm comm)
```

The function MPI\_Comm\_c2f translates a C communicator handle into a Fortran handle to the same communicator; it maps a null handle into a null handle and an invalid handle into an invalid handle.

Similar functions are provided for the other types of opaque objects.

```

1  MPI_Datatype MPI_Type_f2c(MPI_Fint datatype)
2
3  MPI_Fint MPI_Type_c2f(MPI_Datatype datatype)
4
5  MPI_Group MPI_Group_f2c(MPI_Fint group)
6  MPI_Fint MPI_Group_c2f(MPI_Group group)
7
8  MPI_Request MPI_Request_f2c(MPI_Fint request)
9  MPI_Fint MPI_Request_c2f(MPI_Request request)
10
11 MPI_File MPI_File_f2c(MPI_Fint file)
12 MPI_Fint MPI_File_c2f(MPI_File file)
13
14 MPI_Win MPI_Win_f2c(MPI_Fint win)
15 MPI_Fint MPI_Win_c2f(MPI_Win win)
16
17 MPI_Op MPI_Op_f2c(MPI_Fint op)
18 MPI_Fint MPI_Op_c2f(MPI_Op op)
19
20 MPI_Info MPI_Info_f2c(MPI_Fint info)
21 MPI_Fint MPI_Info_c2f(MPI_Info info)
22
23 MPI_Errhandler MPI_Errhandler_f2c(MPI_Fint errhandler)
24 MPI_Fint MPI_Errhandler_c2f(MPI_Errhandler errhandler)
25
26 MPI_Message MPI_Message_f2c(MPI_Fint message)
27 MPI_Fint MPI_Message_c2f(MPI_Message message)
28

```

**Example 16.21** The example below illustrates how the Fortran MPI function `MPI_TYPE_COMMIT` can be implemented by wrapping the C MPI function `MPI_Type_commit` with a C wrapper to do handle conversions. In this example a Fortran-C interface is assumed where a Fortran function is all upper case when referred to from C and arguments are passed by addresses.

```

35  ! FORTRAN PROCEDURE
36  SUBROUTINE MPI_TYPE_COMMIT( DATATYPE, IERR)
37  INTEGER [ticket250-V.]:: DATATYPE, IERR
38  CALL MPI_X_TYPE_COMMIT(DATATYPE, IERR)
39  RETURN
40  END
41
42  /* C wrapper */
43
44  void MPI_X_TYPE_COMMIT( MPI_Fint *f_handle, MPI_Fint *ierr)
45  {
46      MPI_Datatype datatype;
47
48      datatype = MPI_Type_f2c( *f_handle);

```

```

    *ierr = (MPI_Fint)MPI_Type_commit( &datatype);
    *f_handle = MPI_Type_c2f(datatype);
    return;
}

```

The same approach can be used for all other MPI functions. The call to MPI\_XXX\_f2c (resp. MPI\_XXX\_c2f) can be omitted when the handle is an OUT (resp. IN) argument, rather than INOUT.

*Rationale.* The design here provides a convenient solution for the prevalent case, where a C wrapper is used to allow Fortran code to call a C library, or C code to call a Fortran library. The use of C wrappers is much more likely than the use of Fortran wrappers, because it is much more likely that a variable of type INTEGER can be passed to C, than a C handle can be passed to Fortran.

Returning the converted value as a function value rather than through the argument list allows the generation of efficient inlined code when these functions are simple (e.g., the identity). The conversion function in the wrapper does not catch an invalid handle argument. Instead, an invalid handle is passed below to the library function, which, presumably, checks its input arguments. (*End of rationale.*)

**C and C++** The C++ language interface provides the functions listed below for mixed-language interoperability. The token <CLASS> is used below to indicate any valid MPI opaque handle name (e.g., Group), except where noted. For the case where the C++ class corresponding to <CLASS> has derived classes, functions are also provided for converting between the derived classes and the C MPI\_<CLASS>.

The following function allows assignment from a C MPI handle to a C++ MPI handle.

```
MPI::<CLASS>& MPI::<CLASS>::operator=(const MPI_<CLASS>& data)
```

The constructor below creates a C++ MPI object from a C MPI handle. This allows the automatic promotion of a C MPI handle to a C++ MPI handle.

```
MPI::<CLASS>::<CLASS>(const MPI_<CLASS>& data)
```

**Example 16.22** In order for a C program to use a C++ library, the C++ library must export a C interface that provides appropriate conversions before invoking the underlying C++ library call. This example shows a C interface function that invokes a C++ library call with a C communicator; the communicator is automatically promoted to a C++ handle when the underlying C++ function is invoked.

```

// C++ library function prototype
void cpp_lib_call(MPI::Intracomm cpp_comm);

// Exported C function prototype
extern "C" {
    void c_interface(MPI_Comm c_comm);
}

void c_interface(MPI_Comm c_comm)

```

```

1  {
2      // the MPI_Comm (c_comm) is automatically promoted to MPI::Intracomm
3      cpp_lib_call(c_comm);
4  }

```

The following function allows conversion from C++ objects to C MPI handles. In this case, the casting operator is overloaded to provide the functionality.

```

8  MPI::<CLASS>::operator MPI_<CLASS>() const

```

**Example 16.23** A C library routine is called from a C++ program. The C library routine is prototyped to take an MPI\_Comm as an argument.

```

13 // C function prototype
14 extern "C" {
15     void c_lib_call(MPI_Comm c_comm);
16 }
17
18 void cpp_function()
19 {
20     // Create a C++ communicator, and initialize it with a dup of
21     // MPI::COMM_WORLD
22     MPI::Intracomm cpp_comm(MPI::COMM_WORLD.Dup());
23     c_lib_call(cpp_comm);
24 }

```

*Rationale.* Providing conversion from C to C++ via constructors and from C++ to C via casting allows the compiler to make automatic conversions. Calling C from C++ becomes trivial, as does the provision of a C or Fortran interface to a C++ library. (*End of rationale.*)

*Advice to users.* Note that the casting and promotion operators return new handles by value. Using these new handles as INOUT parameters will affect the internal MPI object, but will *not* affect the original handle from which it was cast. (*End of advice to users.*)

It is important to note that all C++ objects with corresponding C handles can be used interchangeably by an application. For example, an application can cache an attribute on MPI\_COMM\_WORLD and later retrieve it from MPI::COMM\_WORLD.

### 16.3.5 Status

The following two procedures are provided in C to convert from a Fortran (with the mpi module or mpif.h) status (which is an array of integers) to a C status (which is a structure), and vice versa. The conversion occurs on all the information in status, including that which is hidden. That is, no status information is lost in the conversion.

```

45 int MPI_Status_f2c(const MPI_Fint *f_status, MPI_Status *c_status)

```

If f\_status is a valid Fortran status, but not the Fortran value of MPI\_STATUS\_IGNORE or MPI\_STATUSES\_IGNORE, then MPI\_Status\_f2c returns in c\_status a valid C status with



the same content. If `f_status` is the Fortran value of `MPI_STATUS_IGNORE` or `MPI_STATUSES_IGNORE`, or if `f_status` is not a valid Fortran status, then the call is erroneous.

The C status has the same source, tag and error code values as the Fortran status, and returns the same answers when queried for count, elements, and cancellation. The conversion function may be called with a Fortran status argument that has an undefined error field, in which case the value of the error field in the C status argument is undefined.

Two global variables of type `MPI_Fint*`, `MPI_F_STATUS_IGNORE` and `MPI_F_STATUSES_IGNORE` are declared in `mpi.h`. They can be used to test, in C, whether `f_status` is the Fortran value of `MPI_STATUS_IGNORE` or `MPI_STATUSES_IGNORE`, respectively, defined in the `mpi` module or `mpif.h`. These are global variables, not C constant expressions and cannot be used in places where C requires constant expressions. Their value is defined only between the calls to `MPI_INIT` and `MPI_FINALIZE` and should not be changed by user code.

To do the conversion in the other direction, we have the following:

```
int MPI_Status_c2f(const MPI_Status *c_status, MPI_Fint *f_status)
```

This call converts a C status into a Fortran status, and has a behavior similar to `MPI_Status_f2c`. That is, the value of `c_status` must not be either `MPI_STATUS_IGNORE` or `MPI_STATUSES_IGNORE`.

*Advice to users.* There is not a separate conversion function for arrays of statuses, since one can simply loop through the array, converting each status with the routines in Fig. 16.1 on page 698. (End of advice to users.)

*Rationale.* The handling of `MPI_STATUS_IGNORE` is required in order to layer libraries with only a C wrapper: if the Fortran call has passed `MPI_STATUS_IGNORE`, then the C wrapper must handle this correctly. Note that this constant need not have the same value in Fortran and C. If `MPI_Status_f2c` were to handle `MPI_STATUS_IGNORE`, then the type of its result would have to be `MPI_Status**`, which was considered an inferior solution. (End of rationale.)

Using the `mpi_f08` Fortran module, a status is declared as `TYPE(MPI_Status)`. The C type `MPI_F08_status` can be used to pass a Fortran `TYPE(MPI_Status)` argument into a C routine. Figure 16.1 illustrates all status conversion routines. Some are only available in C, some in both C and Fortran.

```
int MPI_Status_f082c(const MPI_F08_status *f08_status, MPI_Status
                    *c_status)
```

This C routine converts a Fortran `mpi_f08 TYPE(MPI_Status)` into a C `MPI_Status`.

```
int MPI_Status_c2f08(const MPI_Status *c_status, MPI_F08_status
                    *f08_status)
```

This C routine converts a C `MPI_Status` into a Fortran `mpi_f08 TYPE(MPI_Status)`. Two global variables of type `MPI_F08_status*`, `MPI_F08_STATUS_IGNORE` and `MPI_F08_STATUSES_IGNORE` are declared in `mpi.h`. They can be used to test, in C, whether `f_status` is the Fortran value of `MPI_STATUS_IGNORE` or `MPI_STATUSES_IGNORE` defined in the `mpi_f08` module. These are global variables, not C constant expressions and cannot be used in places where C requires constant expressions. Their value is defined only between the calls to `MPI_INIT` and `MPI_FINALIZE` and should not be changed by user code.

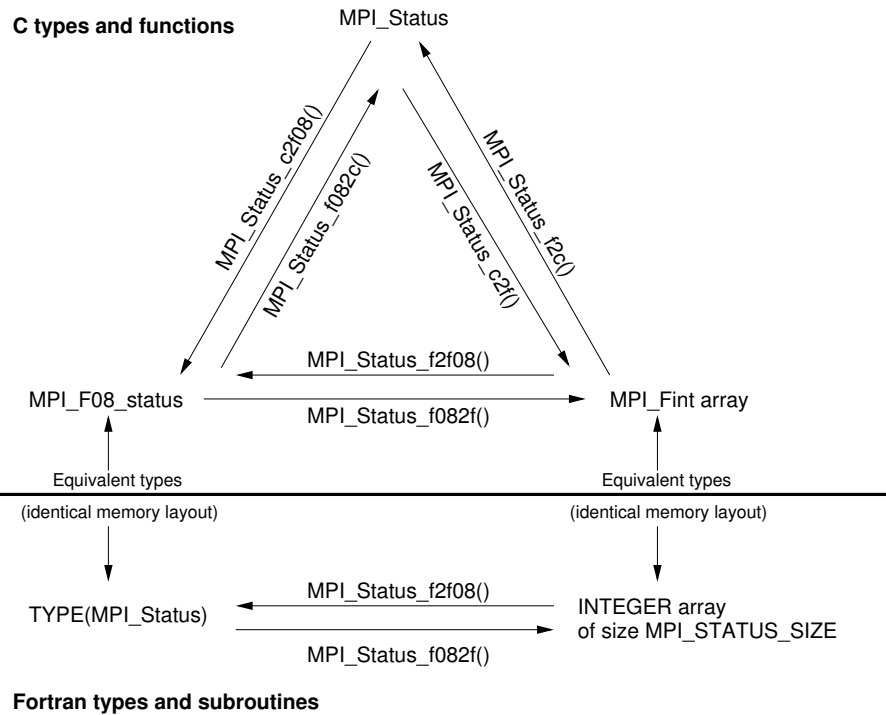


Figure 16.1: Status conversion routines

Conversion between the two Fortran versions of a status can be done with:

```
MPI_STATUS_F2F08(f_status, f08_status)
```

```
IN      f_status      status object declared as array
OUT     f08_status     status object declared as named type
```

```
int MPI_Status_f2f08(MPI_Fint *f_status, MPI_F08_status *f08_status)
```

```
MPI_Status_f2f08(f_status, f08_status, ierror) BIND(C)
```

```
  INTEGER, INTENT(IN) :: f_status(MPI_STATUS_SIZE)
```

```
  TYPE(MPI_Status), INTENT(OUT) :: f08_status
```

```
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_STATUS_F2F08(F_STATUS, F08_STATUS, IERROR)
```

```
  INTEGER :: F_STATUS(MPI_STATUS_SIZE)
```

```
  TYPE(MPI_Status) :: F08_STATUS
```

```
  INTEGER IERROR
```

This routine converts a Fortran `INTEGER, DIMENSION(MPI_STATUS_SIZE)` status array into a Fortran `mpi_f08 TYPE(MPI_Status)`.

```
MPI_STATUS_F082F(f08_status, f_status)
```

```
IN      f08_status      status object declared as named type
OUT     f_status        status object declared as array
```

```
int MPI_Status_f082f(MPI_F08_status *f08_status, MPI_Fint *f_status)
```

```
MPI_Status_f082f(f08_status, f_status, ierror) BIND(C)
  TYPE(MPI_Status), INTENT(IN) :: f08_status
  INTEGER, INTENT(OUT) :: f_status(MPI_STATUS_SIZE)
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror
```

```
MPI_STATUS_F082F(F08_STATUS, F_STATUS, IERROR)
  TYPE(MPI_Status) :: F08_STATUS
  INTEGER :: F_STATUS(MPI_STATUS_SIZE)
  INTEGER IERROR
```

This routine converts a Fortran `mpi_f08` `TYPE(MPI_Status)` into a Fortran `INTEGER`, `DIMENSION(MPI_STATUS_SIZE)` status array.

### 16.3.6 MPI Opaque Objects

Unless said otherwise, opaque objects are “the same” in all languages: they carry the same information, and have the same meaning in both languages. The mechanism described in the previous section can be used to pass references to MPI objects from language to language. An object created in one language can be accessed, modified or freed in another language.

We examine below in more detail, issues that arise for each type of MPI object.

#### Datatypes

Datatypes encode the same information in all languages. E.g., a datatype accessor like `MPI_TYPE_GET_EXTENT` will return the same information in all languages. If a datatype defined in one language is used for a communication call in another language, then the message sent will be identical to the message that would be sent from the first language: the same communication buffer is accessed, and the same representation conversion is performed, if needed. All predefined datatypes can be used in datatype constructors in any language. If a datatype is committed, it can be used for communication in any language.

The function `MPI_GET_ADDRESS` returns the same value in all languages. Note that we do not require that the constant `MPI_BOTTOM` have the same value in all languages (see [16.3.9](#), page 706).

#### Example 16.24

```
! FORTRAN CODE
REAL [ticket250-V.]:: R(5)
INTEGER [ticket250-V.]:: TYPE, IERR, AOBLLEN(1), AOTYPE(1)
INTEGER (KIND=MPI_ADDRESS_KIND) [ticket250-V.]:: AODISP(1)

! create an absolute datatype for array R
```

```

1  AOBLEN(1) = 5
2  CALL MPI_GET_ADDRESS( R, AODISP(1), IERR)
3  AOTYPE(1) = MPI_REAL
4  CALL MPI_TYPE_CREATE_STRUCT(1, AOBLEN,AODISP,AOTYPE, TYPE, IERR)
5  CALL C_ROUTINE(TYPE)
6
7  /* C code */
8
9  void C_ROUTINE(MPI_Fint *ftype)
10 {
11     int count = 5;
12     int lens[2] = {1,1};
13     MPI_Aint displs[2];
14     MPI_Datatype types[2], newtype;
15
16     /* create an absolute datatype for buffer that consists */
17     /* of count, followed by R(5) */
18
19     MPI_Get_address(&count, &displs[0]);
20     displs[1] = 0;
21     types[0] = MPI_INT;
22     types[1] = MPI_Type_f2c(*ftype);
23     MPI_Type_create_struct(2, lens, displs, types, &newtype);
24     MPI_Type_commit(&newtype);
25
26     MPI_Send(MPI_BOTTOM, 1, newtype, 1, 0, MPI_COMM_WORLD);
27     /* the message sent contains an int count of 5, followed */
28     /* by the 5 REAL entries of the Fortran array R. */
29 }
30

```

*Advice to implementors.* The following implementation can be used: MPI addresses, as returned by `MPI_GET_ADDRESS`, will have the same value in all languages. One obvious choice is that MPI addresses be identical to regular addresses. The address is stored in the datatype, when datatypes with absolute addresses are constructed. When a send or receive operation is performed, then addresses stored in a datatype are interpreted as displacements that are all augmented by a base address. This base address is (the address of) `buf`, or zero, if `buf = MPI_BOTTOM`. Thus, if `MPI_BOTTOM` is zero then a send or receive call with `buf = MPI_BOTTOM` is implemented exactly as a call with a regular buffer argument: in both cases the base address is `buf`. On the other hand, if `MPI_BOTTOM` is not zero, then the implementation has to be slightly different. A test is performed to check whether `buf = MPI_BOTTOM`. If true, then the base address is zero, otherwise it is `buf`. In particular, if `MPI_BOTTOM` does not have the same value in Fortran and C/C++, then an additional test for `buf = MPI_BOTTOM` is needed in at least one of the languages.

It may be desirable to use a value other than zero for `MPI_BOTTOM` even in C/C++, so as to distinguish it from a NULL pointer. If `MPI_BOTTOM = c` then one can still avoid the test `buf = MPI_BOTTOM`, by using the displacement from `MPI_BOTTOM`,

i.e., the regular address - c, as the MPI address returned by `MPI_GET_ADDRESS` and stored in absolute datatypes. (*End of advice to implementors.*)

### Callback Functions

MPI calls may associate callback functions with MPI objects: error handlers are associated with communicators and files, attribute copy and delete functions are associated with attribute keys, reduce operations are associated with operation objects, etc. In a multilanguage environment, a function passed in an MPI call in one language may be invoked by an MPI call in another language. MPI implementations must make sure that such invocation will use the calling convention of the language the function is bound to.

*Advice to implementors.* Callback functions need to have a language tag. This tag is set when the callback function is passed in by the library function (which is presumably different for each language **and language support method**), and is used to generate the right calling sequence when the callback function is invoked. (*End of advice to implementors.*)

*Advice to users.* If a subroutine written in one language or Fortran support method wants to pass a callback routine including the predefined Fortran functions (e.g., `MPI_COMM_NULL_COPY_FN`) to another application routine written in another language or Fortran support method, then it must be guaranteed that both routines use the callback interface definition that is defined for the argument when passing the callback to an MPI routine (e.g., `MPI_COMM_CREATE_KEYVAL`); see also the advice to users on page 282. (*End of advice to users.*)

### Error Handlers

*Advice to implementors.* Error handlers, have, in C and C++, a “`stdargs`” argument list. It might be useful to provide to the handler information on the language environment where the error occurred. (*End of advice to implementors.*)

### Reduce Operations

*Advice to users.* Reduce operations receive as one of their arguments the datatype of the operands. Thus, one can define “polymorphic” reduce operations that work for C, C++, and Fortran datatypes. (*End of advice to users.*)

### Addresses

Some of the datatype accessors and constructors have arguments of type `MPI_Aint` (in C) or `MPI::Aint` in C++, to hold addresses. The corresponding arguments, in Fortran, have type `INTEGER`. This causes Fortran and C/C++ to be incompatible, in an environment where addresses have 64 bits, but Fortran `INTEGER`s have 32 bits.

This is a problem, irrespective of interlanguage issues. Suppose that a Fortran process has an address space of  $\geq 4$  GB. What should be the value returned in Fortran by `MPI_ADDRESS`, for a variable with an address above  $2^{32}$ ? The design described here addresses this issue, while maintaining compatibility with current Fortran codes.

The constant `MPI_ADDRESS_KIND` is defined so that, in Fortran 90, `INTEGER(KIND=MPI_ADDRESS_KIND)` is an address sized integer type (typically, but not necessarily, the size of an `INTEGER(KIND=MPI_ADDRESS_KIND)` is 4 on 32 bit address machines and 8 on 64 bit address machines). Similarly, the constant `MPI_INTEGER_KIND` is defined so that `INTEGER(KIND=MPI_INTEGER_KIND)` is a default size `INTEGER`.

There are seven functions that have address arguments: `MPI_TYPE_HVECTOR`, `MPI_TYPE_HINDEXED`, `MPI_TYPE_STRUCT`, `MPI_ADDRESS`, `MPI_TYPE_EXTENT`, `MPI_TYPE_LB` and `MPI_TYPE_UB`.

Four new functions are provided to supplement the first four functions in this list. These functions are described in Section 4.1.1 on page 93. The remaining three functions are supplemented by the new function `MPI_TYPE_GET_EXTENT`, described in that same section. The new functions have the same functionality as the old functions in C/C++, or on Fortran systems where default `INTEGER`s are address sized. In Fortran, they accept arguments of type `INTEGER(KIND=MPI_ADDRESS_KIND)`, wherever arguments of type `MPI_Aint` and `MPI::Aint` are used in C and C++. On Fortran 77 systems that do not support the Fortran 90 `KIND` notation, and where addresses are 64 bits whereas default `INTEGER`s are 32 bits, these arguments will be of an appropriate integer type. The old functions will continue to be provided, for backward compatibility. However, users are encouraged to switch to the new functions, in Fortran, so as to avoid problems on systems with an address range  $> 2^{32}$ , and to provide compatibility across languages.

### 16.3.7 Attributes

Attribute keys can be allocated in one language and freed in another. Similarly, attribute values can be set in one language and accessed in another. To achieve this, attribute keys will be allocated in an integer range that is valid all languages. The same holds true for system-defined attribute values (such as `MPI_TAG_UB`, `MPI_WTIME_IS_GLOBAL`, etc.)

Attribute keys declared in one language are associated with copy and delete functions in that language (the functions provided by the `MPI_{TYPE,COMM,WIN}_CREATE_KEYVAL` call). When a communicator is duplicated, for each attribute, the corresponding copy function is called, using the right calling convention for the language of that function; and similarly, for the delete callback function.

*Advice to implementors.* This requires that attributes be tagged either as “C,” “C++” or “Fortran,” and that the language tag be checked in order to use the right calling convention for the callback function. (*End of advice to implementors.*)

The attribute manipulation functions described in Section 6.7 on page 276 define attributes arguments to be of type `void*` in C, and of type `INTEGER`, in Fortran. On some systems, `INTEGER`s will have 32 bits, while C/C++ pointers will have 64 bits. This is a problem if communicator attributes are used to move information from a Fortran caller to a C/C++ callee, or vice-versa.

MPI behaves as if it stores, internally, address sized attributes. If Fortran `INTEGER`s are smaller, then the Fortran function `MPI_ATTR_GET` will return the least significant part of the attribute word; the Fortran function `MPI_ATTR_PUT` will set the least significant part of the attribute word, which will be sign extended to the entire word. (These two functions may be invoked explicitly by user code, or implicitly, by attribute copying callback functions.)

As for addresses, new functions are provided that manipulate Fortran address sized attributes, and have the same functionality as the old functions in C/C++. These functions are described in Section 6.7, page 276. Users are encouraged to use these new functions.

MPI supports two types of attributes: address-valued (pointer) attributes, and integer valued attributes. C and C++ attribute functions put and get address valued attributes. Fortran attribute functions put and get integer valued attributes. When an integer valued attribute is accessed from C or C++, then `MPI_xxx_get_attr` will return the address of (a pointer to) the integer valued attribute, which is a pointer to `MPI_Aint` if the attribute was stored with Fortran `MPI_xxx_SET_ATTR`, and a pointer to `int` if it was stored with the deprecated Fortran `MPI_ATTR_PUT`. When an address valued attribute is accessed from Fortran, then `MPI_xxx_GET_ATTR` will convert the address into an integer and return the result of this conversion. This conversion is lossless if new style attribute functions are used, and an integer of kind `MPI_ADDRESS_KIND` is returned. The conversion may cause truncation if deprecated attribute functions are used. In C, the deprecated routines `MPI_Attr_put` and `MPI_Attr_get` behave identical to `MPI_Comm_set_attr` and `MPI_Comm_get_attr`.

### Example 16.25

A. Setting an attribute value in C

```
int set_val = 3;
struct foo set_struct;

/* Set a value that is a pointer to an int */

MPI_Comm_set_attr(MPI_COMM_WORLD, keyval1, &set_val);
/* Set a value that is a pointer to a struct */
MPI_Comm_set_attr(MPI_COMM_WORLD, keyval2, &set_struct);
/* Set an integer value */
MPI_Comm_set_attr(MPI_COMM_WORLD, keyval3, (void *) 17);
```

B. Reading the attribute value in C

```
int flag, *get_val;
struct foo *get_struct;

/* Upon successful return, get_val == &set_val
   (and therefore *get_val == 3) */
MPI_Comm_get_attr(MPI_COMM_WORLD, keyval1, &get_val, &flag);
/* Upon successful return, get_struct == &set_struct */
MPI_Comm_get_attr(MPI_COMM_WORLD, keyval2, &get_struct, &flag);
/* Upon successful return, get_val == (void*) 17 */
/*      i.e., (MPI_Aint) get_val == 17 */
MPI_Comm_get_attr(MPI_COMM_WORLD, keyval3, &get_val, &flag);
```

C. Reading the attribute value with (deprecated) Fortran MPI-1 calls

```

1  LOGICAL FLAG
2  INTEGER IERR, GET_VAL, GET_STRUCT
3
4  ! Upon successful return, GET_VAL == &set_val, possibly truncated
5  CALL MPI_ATTR_GET(MPI_COMM_WORLD, KEYVAL1, GET_VAL, FLAG, IERR)
6  ! Upon successful return, GET_STRUCT == &set_struct, possibly truncated
7  CALL MPI_ATTR_GET(MPI_COMM_WORLD, KEYVAL2, GET_STRUCT, FLAG, IERR)
8  ! Upon successful return, GET_VAL == 17
9  CALL MPI_ATTR_GET(MPI_COMM_WORLD, KEYVAL3, GET_VAL, FLAG, IERR)

```

D. Reading the attribute value with Fortran MPI-2 calls

```

12 LOGICAL FLAG
13 INTEGER IERR
14 INTEGER (KIND=MPI_ADDRESS_KIND) GET_VAL, GET_STRUCT
15
16 ! Upon successful return, GET_VAL == &set_val
17 CALL MPI_COMM_GET_ATTR(MPI_COMM_WORLD, KEYVAL1, GET_VAL, FLAG, IERR)
18 ! Upon successful return, GET_STRUCT == &set_struct
19 CALL MPI_COMM_GET_ATTR(MPI_COMM_WORLD, KEYVAL2, GET_STRUCT, FLAG, IERR)
20 ! Upon successful return, GET_VAL == 17
21 CALL MPI_COMM_GET_ATTR(MPI_COMM_WORLD, KEYVAL3, GET_VAL, FLAG, IERR)

```

### Example 16.26

A. Setting an attribute value with the (deprecated) Fortran MPI-1 call

```

27 INTEGER IERR, VAL
28 VAL = 7
29 CALL MPI_ATTR_PUT(MPI_COMM_WORLD, KEYVAL, VAL, IERR)

```

B. Reading the attribute value in C

```

33 int flag;
34 int *value;
35
36 /* Upon successful return, value points to internal MPI storage and
37    *value == (int) 7 */
38 MPI_Comm_get_attr(MPI_COMM_WORLD, keyval, &value, &flag);

```

C. Reading the attribute value with (deprecated) Fortran MPI-1 calls

```

42 LOGICAL FLAG
43 INTEGER IERR, VALUE
44
45 ! Upon successful return, VALUE == 7
46 CALL MPI_ATTR_GET(MPI_COMM_WORLD, KEYVAL, VALUE, FLAG, IERR)

```

D. Reading the attribute value with Fortran MPI-2 calls



```

LOGICAL FLAG
INTEGER IERR
INTEGER (KIND=MPI_ADDRESS_KIND) VALUE

! Upon successful return, VALUE == 7 (sign extended)
CALL MPI_COMM_GET_ATTR(MPI_COMM_WORLD, KEYVAL, VALUE, FLAG, IERR)

```

**Example 16.27** A. Setting an attribute value via a Fortran MPI-2 call

```

INTEGER IERR
INTEGER(KIND=MPI_ADDRESS_KIND) VALUE1
INTEGER(KIND=MPI_ADDRESS_KIND) VALUE2
VALUE1 = 42
VALUE2 = INT(2, KIND=MPI_ADDRESS_KIND) ** 40

CALL MPI_COMM_SET_ATTR(MPI_COMM_WORLD, KEYVAL1, VALUE1, IERR)
CALL MPI_COMM_SET_ATTR(MPI_COMM_WORLD, KEYVAL2, VALUE2, IERR)

```

B. Reading the attribute value in C

```

int flag;
MPI_Aint *value1, *value2;

/* Upon successful return, value1 points to internal MPI storage and
   *value1 == 42 */
MPI_Comm_get_attr(MPI_COMM_WORLD, keyval1, &value1, &flag);
/* Upon successful return, value2 points to internal MPI storage and
   *value2 == 2^40 */
MPI_Comm_get_attr(MPI_COMM_WORLD, keyval2, &value2, &flag);

```

C. Reading the attribute value with (deprecated) Fortran MPI-1 calls

```

LOGICAL FLAG
INTEGER IERR, VALUE1, VALUE2

! Upon successful return, VALUE1 == 42
CALL MPI_ATTR_GET(MPI_COMM_WORLD, KEYVAL1, VALUE1, FLAG, IERR)
! Upon successful return, VALUE2 == 2^40, or 0 if truncation
! needed (i.e., the least significant part of the attribute word)
CALL MPI_ATTR_GET(MPI_COMM_WORLD, KEYVAL2, VALUE2, FLAG, IERR)

```

D. Reading the attribute value with Fortran MPI-2 calls

```

LOGICAL FLAG
INTEGER IERR
INTEGER (KIND=MPI_ADDRESS_KIND) VALUE1, VALUE2

! Upon successful return, VALUE1 == 42
CALL MPI_COMM_GET_ATTR(MPI_COMM_WORLD, KEYVAL1, VALUE1, FLAG, IERR)
! Upon successful return, VALUE2 == 2^40
CALL MPI_COMM_GET_ATTR(MPI_COMM_WORLD, KEYVAL2, VALUE2, FLAG, IERR)

```

The predefined MPI attributes can be integer valued or address valued. Predefined integer valued attributes, such as `MPI_TAG_UB`, behave as if they were put by a call to the deprecated Fortran routine `MPI_ATTR_PUT`, i.e., in Fortran, `MPI_COMM_GET_ATTR(MPI_COMM_WORLD, MPI_TAG_UB, val, flag, ierr)` will return in `val` the upper bound for tag value; in C, `MPI_Comm_get_attr(MPI_COMM_WORLD, MPI_TAG_UB, &p, &flag)` will return in `p` a pointer to an int containing the upper bound for tag value.

Address valued predefined attributes, such as `MPI_WIN_BASE` behave as if they were put by a C call, i.e., in Fortran, `MPI_WIN_GET_ATTR(win, MPI_WIN_BASE, val, flag, ierr)` will return in `val` the base address of the window, converted to an integer. In C, `MPI_Win_get_attr(win, MPI_WIN_BASE, &p, &flag)` will return in `p` a pointer to the window base, cast to `(void *)`.

*Rationale.* The design is consistent with the behavior specified for predefined attributes, and ensures that no information is lost when attributes are passed from language to language. Because the language interoperability for predefined attributes was defined based on `MPI_ATTR_PUT`, this definition is kept for compatibility reasons although the routine itself is now deprecated. (*End of rationale.*)

*Advice to implementors.* Implementations should tag attributes either as (1) address attributes, (2) as `INTEGER(KIND=MPI_ADDRESS_KIND)` attributes or (3) as `INTEGER` attributes, according to whether they were set in (1) C (with `MPI_Attr_put` or `MPI_Xxx_set_attr`), (2) in Fortran with `MPI_XXX_SET_ATTR` or (3) with the deprecated Fortran routine `MPI_ATTR_PUT`. Thus, the right choice can be made when the attribute is retrieved. (*End of advice to implementors.*)

### 16.3.8 Extra State

Extra-state should not be modified by the copy or delete callback functions. (This is obvious from the C binding, but not obvious from the Fortran binding). However, these functions may update state that is indirectly accessed via extra-state. E.g., in C, extra-state can be a pointer to a data structure that is modified by the copy or callback functions; in Fortran, extra-state can be an index into an entry in a `COMMON` array that is modified by the copy or callback functions. In a multithreaded environment, users should be aware that distinct threads may invoke the same callback function concurrently: if this function modifies state associated with extra-state, then mutual exclusion code must be used to protect updates and accesses to the shared state.

### 16.3.9 Constants

MPI constants have the same value in all languages, unless specified otherwise. This does not apply to constant handles (`MPI_INT`, `MPI_COMM_WORLD`, `MPI_ERRORS_RETURN`, `MPI_SUM`, etc.) These handles need to be converted, as explained in Section 16.3.4. Constants that specify maximum lengths of strings (see Section A.1.1 for a listing) have a value one less in Fortran than C/C++ since in C/C++ the length includes the null terminating character. Thus, these constants represent the amount of space which must be allocated to hold the largest possible such string, rather than the maximum number of printable characters the string could contain.

*Advice to users.* This definition means that it is safe in C/C++ to allocate a buffer to receive a string using a declaration like

```
char name [MPI_MAX_OBJECT_NAME];
```

*(End of advice to users.)*

Also constant “addresses,” i.e., special values for reference arguments that are not handles, such as MPI\_BOTTOM or MPI\_STATUS\_IGNORE may have different values in different languages.

*Rationale.* The current MPI standard specifies that MPI\_BOTTOM can be used in initialization expressions in C, but not in Fortran. Since Fortran does not normally support call by value, then MPI\_BOTTOM must be in Fortran the name of a predefined static variable, e.g., a variable in an MPI declared COMMON block. On the other hand, in C, it is natural to take MPI\_BOTTOM = 0 (Caveat: Defining MPI\_BOTTOM = 0 implies that NULL pointer cannot be distinguished from MPI\_BOTTOM; it may be that MPI\_BOTTOM = 1 is better ...) Requiring that the Fortran and C values be the same will complicate the initialization process. *(End of rationale.)*

### 16.3.10 Interlanguage Communication

The type matching rules for communication[s] in MPI are not changed: the datatype specification for each item sent should match, in type signature, the datatype specification used to receive this item (unless one of the types is MPI\_PACKED). Also, the type of a message item should match the type declaration for the corresponding communication buffer location, unless the type is MPI\_BYTE or MPI\_PACKED. Interlanguage communication is allowed if it complies with these rules.

**Example 16.28** In the example below, a Fortran array is sent from Fortran and received in C.

```
! FORTRAN CODE
USE mpi_f08
REAL [ticket250-V.]:: R(5)
INTEGER [ticket250-V.]:: IERR, MYRANK, AOBLN(1), AOTYPE(1)
[ticket250-V.]TYPE(MPI_Type) :: TYPE
INTEGER (KIND=MPI_ADDRESS_KIND) [ticket250-V.]:: AODISP(1)

! create an absolute datatype for array R
AOBLN(1) = 5
CALL MPI_GET_ADDRESS( R, AODISP(1), IERR)
AOTYPE(1) = MPI_REAL
CALL MPI_TYPE_CREATE_STRUCT(1, AOBLN,AODISP,AOTYPE, TYPE, IERR)
CALL MPI_TYPE_COMMIT(TYPE, IERR)

CALL MPI_COMM_RANK( MPI_COMM_WORLD, MYRANK, IERR)
IF (MYRANK.EQ.0) THEN
    CALL MPI_SEND( MPI_BOTTOM, 1, TYPE, 1, 0, MPI_COMM_WORLD, IERR)
```

ticket250-V.

```
1  ELSE
2      CALL C_ROUTINE(TYPE[ticket250-V.]%MPI_VAL)
3  END IF
4
5
6  /* C code */
7
8  void C_ROUTINE(MPI_Fint *fhandle)
9  {
10     MPI_Datatype type;
11     MPI_Status status;
12
13     type = MPI_Type_f2c(*fhandle);
14
15     MPI_Recv( MPI_BOTTOM, 1, type, 0, 0, MPI_COMM_WORLD, &status);
16 }
17
```

MPI implementors may weaken these type matching rules, and allow messages to be sent with Fortran types and received with C types, and vice versa, when those types match. I.e., if the Fortran type `INTEGER` is identical to the C type `int`, then an MPI implementation may allow data to be sent with datatype `MPI_INTEGER` and be received with datatype `MPI_INT`. However, such code is not portable.

# Annex A

## Language Bindings Summary

In this section we summarize the specific bindings for C, Fortran, and C++. First we present the constants, type definitions, info values and keys. Then we present the routine prototypes separately for each binding. Listings are alphabetical within chapter.

### A.1 Defined Values and Handles

#### A.1.1 Defined Constants

The C and Fortran name is listed in the left column and the C++ name is listed in the middle or right column. Constants with the type `const int` may also be implemented as literal integer constants substituted by the preprocessor.

Return Codes	
C type: <code>const int</code> (or unnamed <code>enum</code> )	C++ type: <code>const int</code>
Fortran type: <code>INTEGER</code>	(or unnamed <code>enum</code> )
<code>MPI_SUCCESS</code>	<code>MPI::SUCCESS</code>
<code>MPI_ERR_BUFFER</code>	<code>MPI::ERR_BUFFER</code>
<code>MPI_ERR_COUNT</code>	<code>MPI::ERR_COUNT</code>
<code>MPI_ERR_TYPE</code>	<code>MPI::ERR_TYPE</code>
<code>MPI_ERR_TAG</code>	<code>MPI::ERR_TAG</code>
<code>MPI_ERR_COMM</code>	<code>MPI::ERR_COMM</code>
<code>MPI_ERR_RANK</code>	<code>MPI::ERR_RANK</code>
<code>MPI_ERR_REQUEST</code>	<code>MPI::ERR_REQUEST</code>
<code>MPI_ERR_ROOT</code>	<code>MPI::ERR_ROOT</code>
<code>MPI_ERR_GROUP</code>	<code>MPI::ERR_GROUP</code>
<code>MPI_ERR_OP</code>	<code>MPI::ERR_OP</code>
<code>MPI_ERR_TOPOLOGY</code>	<code>MPI::ERR_TOPOLOGY</code>
<code>MPI_ERR_DIMS</code>	<code>MPI::ERR_DIMS</code>
<code>MPI_ERR_ARG</code>	<code>MPI::ERR_ARG</code>
<code>MPI_ERR_UNKNOWN</code>	<code>MPI::ERR_UNKNOWN</code>
<code>MPI_ERR_TRUNCATE</code>	<code>MPI::ERR_TRUNCATE</code>
<code>MPI_ERR_OTHER</code>	<code>MPI::ERR_OTHER</code>
<code>MPI_ERR_INTERN</code>	<code>MPI::ERR_INTERN</code>
<code>MPI_ERR_PENDING</code>	<code>MPI::ERR_PENDING</code>

(Continued on next page)

## Return Codes (continued)

MPI_ERR_IN_STATUS	MPI::ERR_IN_STATUS
MPI_ERR_ACCESS	MPI::ERR_ACCESS
MPI_ERR_AMODE	MPI::ERR_AMODE
MPI_ERR_ASSERT	MPI::ERR_ASSERT
MPI_ERR_BAD_FILE	MPI::ERR_BAD_FILE
MPI_ERR_BASE	MPI::ERR_BASE
MPI_ERR_CONVERSION	MPI::ERR_CONVERSION
MPI_ERR_DISP	MPI::ERR_DISP
MPI_ERR_DUP_DATAREP	MPI::ERR_DUP_DATAREP
MPI_ERR_FILE_EXISTS	MPI::ERR_FILE_EXISTS
MPI_ERR_FILE_IN_USE	MPI::ERR_FILE_IN_USE
MPI_ERR_FILE	MPI::ERR_FILE
MPI_ERR_INFO_KEY	MPI::ERR_INFO_VALUE
MPI_ERR_INFO_NOKEY	MPI::ERR_INFO_NOKEY
MPI_ERR_INFO_VALUE	MPI::ERR_INFO_KEY
MPI_ERR_INFO	MPI::ERR_INFO
MPI_ERR_IO	MPI::ERR_IO
MPI_ERR_KEYVAL	MPI::ERR_KEYVAL
MPI_ERR_LOCKTYPE	MPI::ERR_LOCKTYPE
MPI_ERR_NAME	MPI::ERR_NAME
MPI_ERR_NO_MEM	MPI::ERR_NO_MEM
MPI_ERR_NOT_SAME	MPI::ERR_NOT_SAME
MPI_ERR_NO_SPACE	MPI::ERR_NO_SPACE
MPI_ERR_NO_SUCH_FILE	MPI::ERR_NO_SUCH_FILE
MPI_ERR_PORT	MPI::ERR_PORT
MPI_ERR_QUOTA	MPI::ERR_QUOTA
MPI_ERR_READ_ONLY	MPI::ERR_READ_ONLY
MPI_ERR_RMA_CONFLICT	MPI::ERR_RMA_CONFLICT
MPI_ERR_RMA_SYNC	MPI::ERR_RMA_SYNC
MPI_ERR_SERVICE	MPI::ERR_SERVICE
MPI_ERR_SIZE	MPI::ERR_SIZE
MPI_ERR_SPAWN	MPI::ERR_SPAWN
MPI_ERR_UNSUPPORTED_DATAREP	MPI::ERR_UNSUPPORTED_DATAREP
MPI_ERR_UNSUPPORTED_OPERATION	MPI::ERR_UNSUPPORTED_OPERATION
MPI_ERR_WIN	MPI::ERR_WIN
MPI_ERR_LASTCODE	MPI::ERR_LASTCODE
[ticket270.]MPI_ERR_RMA_RANGE	[ticket270.]Not defined for C++
[ticket270.]MPI_ERR_RMA_ATTACH	[ticket270.]Not defined for C++

ticket266.

**Return Codes for the MPI tool information interface**


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MPI\_T\_ERR\_CANTINIT  
 MPI\_T\_ERR\_NOTINITIALIZED  
 MPI\_T\_ERR\_MEMORY  
 MPI\_T\_ERR\_INVALIDINDEX  
 MPI\_T\_ERR\_INVALIDITEM  
 MPI\_T\_ERR\_INVALIDSESSION  
 MPI\_T\_ERR\_INVALIDHANDLE  
 MPI\_T\_ERR\_OUTOFHANDLES  
 MPI\_T\_ERR\_OUTOFSESSIONS  
 MPI\_T\_ERR\_CVAR\_SETNOTNOW  
 MPI\_T\_ERR\_CVAR\_SETNEVER  
 MPI\_T\_ERR\_PVAR\_NOWRITE  
 MPI\_T\_ERR\_PVAR\_NOSTARTSTOP  
 MPI\_T\_ERR\_PVAR\_NOATOMIC

---

**Buffer Address Constants**

C type: void * const	C++ type:
Fortran type: (predefined memory location)	void * const
MPI_BOTTOM	MPI::BOTTOM
MPI_IN_PLACE	MPI::IN_PLACE

---

**Assorted Constants**

C type: const int (or unnamed enum)	C++ type:
Fortran type: INTEGER	const int (or unnamed enum)
MPI_PROC_NULL	MPI::PROC_NULL
MPI_ANY_SOURCE	MPI::ANY_SOURCE
MPI_ANY_TAG	MPI::ANY_TAG
MPI_UNDEFINED	MPI::UNDEFINED
MPI_BSEND_OVERHEAD	MPI::BSEND_OVERHEAD
MPI_KEYVAL_INVALID	MPI::KEYVAL_INVALID
MPI_LOCK_EXCLUSIVE	MPI::LOCK_EXCLUSIVE
MPI_LOCK_SHARED	MPI::LOCK_SHARED
MPI_ROOT	MPI::ROOT

---

**[ticket247-S.] Fortran Support Method Specific Constants**


---

Fortran type: LOGICAL

---

[ticket234-F.] **MPI\_SUBARRAYS\_SUPPORTED** (Fortran only)

---

[ticket238-J.] [ticket229.1.] **MPI\_ASYNC\_PROTECTS\_NONBLOCKING** (Fortran only)

---

**Status size and reserved index values (Fortran only)**

Fortran type: INTEGER

MPI_STATUS_SIZE	Not defined for C++
MPI_SOURCE	Not defined for C++
MPI_TAG	Not defined for C++
MPI_ERROR	Not defined for C++

**Variable Address Size (Fortran only)**

Fortran type: INTEGER

MPI_ADDRESS_KIND	Not defined for C++
[ticket265.]MPI_COUNT_KIND	[ticket265.]Not defined for C++
MPI_INTEGER_KIND	Not defined for C++
MPI_OFFSET_KIND	Not defined for C++

**Error-handling specifiers**

C type: MPI_Errhandler	C++ type: MPI::Errhandler
Fortran type: INTEGER	
[ticket231-C.]or TYPE(MPI_Errhandler)	
MPI_ERRORS_ARE_FATAL	MPI::ERRORS_ARE_FATAL
MPI_ERRORS_RETURN	MPI::ERRORS_RETURN
	MPI::ERRORS_THROW_EXCEPTIONS

**Maximum Sizes for Strings**

C type: const int (or unnamed enum)	C++ type:
Fortran type: INTEGER	const int (or unnamed enum)
MPI_MAX_PROCESSOR_NAME	MPI::MAX_PROCESSOR_NAME
[ticket204.]MPI_MAX_LIBRARY_VERSION_STRING	
MPI_MAX_ERROR_STRING	MPI::MAX_ERROR_STRING
MPI_MAX_DATAREP_STRING	MPI::MAX_DATAREP_STRING
MPI_MAX_INFO_KEY	MPI::MAX_INFO_KEY
MPI_MAX_INFO_VAL	MPI::MAX_INFO_VAL
MPI_MAX_OBJECT_NAME	MPI::MAX_OBJECT_NAME
MPI_MAX_PORT_NAME	MPI::MAX_PORT_NAME



Named Predefined Datatypes		C/C++ types	1
C type: MPI_Datatype	C++ type: MPI::Datatype		2
Fortran type: INTEGER			3
[ticket231-C.]or TYPE(MPI_Datatype)			4
MPI_CHAR	MPI::CHAR	char	5
		(treated as printable character)	6
MPI_SHORT	MPI::SHORT	signed short int	8
MPI_INT	MPI::INT	signed int	10
MPI_LONG	MPI::LONG	signed long	11
MPI_LONG_LONG_INT	MPI::LONG_LONG_INT	signed long long	
MPI_LONG_LONG	MPI::LONG_LONG	long long (synonym)	
MPI_SIGNED_CHAR	MPI::SIGNED_CHAR	signed char	14
		(treated as integral value)	
MPI_UNSIGNED_CHAR	MPI::UNSIGNED_CHAR	unsigned char	16
		(treated as integral value)	
MPI_UNSIGNED_SHORT	MPI::UNSIGNED_SHORT	unsigned short	18
MPI_UNSIGNED	MPI::UNSIGNED	unsigned int	19
MPI_UNSIGNED_LONG	MPI::UNSIGNED_LONG	unsigned long	20
MPI_UNSIGNED_LONG_LONG	MPI::UNSIGNED_LONG_LONG	unsigned long long	
MPI_FLOAT	MPI::FLOAT	float	22
MPI_DOUBLE	MPI::DOUBLE	double	23
MPI_LONG_DOUBLE	MPI::LONG_DOUBLE	long double	24
MPI_WCHAR	MPI::WCHAR	wchar_t	25
		(defined in <stddef.h>)	
		(treated as printable character)	28
MPI_C_BOOL	(use C datatype handle)	_Bool	29
MPI_INT8_T	(use C datatype handle)	int8_t	30
MPI_INT16_T	(use C datatype handle)	int16_t	31
MPI_INT32_T	(use C datatype handle)	int32_t	32
MPI_INT64_T	(use C datatype handle)	int64_t	33
MPI_UINT8_T	(use C datatype handle)	uint8_t	34
MPI_UINT16_T	(use C datatype handle)	uint16_t	35
MPI_UINT32_T	(use C datatype handle)	uint32_t	36
MPI_UINT64_T	(use C datatype handle)	uint64_t	37
MPI_AINT	(use C datatype handle)	MPI_Aint	38
[ticket265.]MPI_COUNT	[ticket265.](use C datatype handle)	[ticket265.]MPI_Count	39
MPI_OFFSET	(use C datatype handle)	MPI_Offset	40
MPI_C_COMPLEX	(use C datatype handle)	float _Complex	41
MPI_C_FLOAT_COMPLEX	(use C datatype handle)	float _Complex	42
MPI_C_DOUBLE_COMPLEX	(use C datatype handle)	double _Complex	43
MPI_C_LONG_DOUBLE_COMPLEX	(use C datatype handle)	long double _Complex	
MPI_BYTE	MPI::BYTE	(any C/C++ type)	
MPI_PACKED	MPI::PACKED	(any C/C++ type)	

47

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Named Predefined Datatypes		Fortran types
C type: MPI_Datatype	C++ type: MPI::Datatype	
Fortran type: INTEGER		
[ticket231-C.]or TYPE(MPI_Datatype)		
MPI_INTEGER	MPI::INTEGER	INTEGER
MPI_REAL	MPI::REAL	REAL
MPI_DOUBLE_PRECISION	MPI::DOUBLE_PRECISION	DOUBLE PRECISION
MPI_COMPLEX	MPI::F_COMPLEX	COMPLEX
MPI_LOGICAL	MPI::LOGICAL	LOGICAL
MPI_CHARACTER	MPI::CHARACTER	CHARACTER(1)
MPI_AINT	(use C datatype handle)	INTEGER (KIND=MPI_ADDRESS_KIND)
[ticket265.]MPI_COUNT	(use C datatype handle)	[ticket265.]INTEGER (KIND=MPI_COUNT_KIND)
MPI_OFFSET	(use C datatype handle)	INTEGER (KIND=MPI_OFFSET_KIND)
MPI_BYTE	MPI::BYTE	(any Fortran type)
MPI_PACKED	MPI::PACKED	(any Fortran type)

C++-Only Named Predefined Datatypes	C++ types
C++ type: MPI::Datatype	
MPI::BOOL	bool
MPI::COMPLEX	Complex<float>
MPI::DOUBLE_COMPLEX	Complex<double>
MPI::LONG_DOUBLE_COMPLEX	Complex<long double>

Optional datatypes (Fortran)		Fortran types
C type: MPI_Datatype	C++ type: MPI::Datatype	
Fortran type: INTEGER		
[ticket231-C.]or TYPE(MPI_Datatype)		
MPI_DOUBLE_COMPLEX	MPI::F_DOUBLE_COMPLEX	DOUBLE COMPLEX
MPI_INTEGER1	MPI::INTEGER1	INTEGER*1
MPI_INTEGER2	MPI::INTEGER2	INTEGER*[ticket231-C.] [8] 2
MPI_INTEGER4	MPI::INTEGER4	INTEGER*4
MPI_INTEGER8	MPI::INTEGER8	INTEGER*8
MPI_INTEGER16		INTEGER*16
MPI_REAL2	MPI::REAL2	REAL*2
MPI_REAL4	MPI::REAL4	REAL*4
MPI_REAL8	MPI::REAL8	REAL*8
MPI_REAL16		REAL*16
MPI_COMPLEX4		COMPLEX*4
MPI_COMPLEX8		COMPLEX*8
MPI_COMPLEX16		COMPLEX*16
MPI_COMPLEX32		COMPLEX*32

**Datatypes for reduction functions (C and C++)**

C type: MPI_Datatype	C++ type: MPI::Datatype
Fortran type: INTEGER	
[ticket231-C.]or TYPE(MPI_Datatype)	
MPI_FLOAT_INT	MPI::FLOAT_INT
MPI_DOUBLE_INT	MPI::DOUBLE_INT
MPI_LONG_INT	MPI::LONG_INT
MPI_2INT	MPI::TWOINT
MPI_SHORT_INT	MPI::SHORT_INT
MPI_LONG_DOUBLE_INT	MPI::LONG_DOUBLE_INT

**Datatypes for reduction functions (Fortran)**

C type: MPI_Datatype	C++ type: MPI::Datatype
Fortran type: INTEGER	
[ticket231-C.]or TYPE(MPI_Datatype)	
MPI_2REAL	MPI::TWOREAL
MPI_2DOUBLE_PRECISION	MPI::TWODOUBLE_PRECISION
MPI_2INTEGER	MPI::TWOINTEGER

**Special datatypes for constructing derived datatypes**

C type: MPI_Datatype	C++ type: MPI::Datatype
Fortran type: INTEGER	
[ticket231-C.]or TYPE(MPI_Datatype)	
MPI_UB	MPI::UB
MPI_LB	MPI::LB

**Reserved communicators**

C type: MPI_Comm	C++ type: MPI::Intracomm
Fortran type: INTEGER	
[ticket231-C.]or TYPE(MPI_Comm)	
MPI_COMM_WORLD	MPI::COMM_WORLD
MPI_COMM_SELF	MPI::COMM_SELF

**Communicator split type constants**C type: `const int` (or unnamed `enum`)Fortran type: `INTEGER``MPI_COMM_TYPE_SHARED`

ticket287.

**Results of communicator and group comparisons**

C type: <code>const int</code> (or unnamed <code>enum</code> )	C++ type: <code>const int</code>
Fortran type: <code>INTEGER</code>	(or unnamed <code>enum</code> )
<code>MPI_IDENT</code>	<code>MPI::IDENT</code>
<code>MPI_CONGRUENT</code>	<code>MPI::CONGRUENT</code>
<code>MPI_SIMILAR</code>	<code>MPI::SIMILAR</code>
<code>MPI_UNEQUAL</code>	<code>MPI::UNEQUAL</code>

**Environmental inquiry keys**

C type: <code>const int</code> (or unnamed <code>enum</code> )	C++ type: <code>const int</code>
Fortran type: <code>INTEGER</code>	(or unnamed <code>enum</code> )
<code>MPI_TAG_UB</code>	<code>MPI::TAG_UB</code>
<code>MPI_IO</code>	<code>MPI::IO</code>
<code>MPI_HOST</code>	<code>MPI::HOST</code>
<code>MPI_WTIME_IS_GLOBAL</code>	<code>MPI::WTIME_IS_GLOBAL</code>

**Collective Operations**

C type: <code>MPI_Op</code>	C++ type: <code>const MPI::Op</code>
Fortran type: <code>INTEGER</code>	
[ticket231-C.]or <code>TYPE(MPI_Op)</code>	
<code>MPI_MAX</code>	<code>MPI::MAX</code>
<code>MPI_MIN</code>	<code>MPI::MIN</code>
<code>MPI_SUM</code>	<code>MPI::SUM</code>
<code>MPI_PROD</code>	<code>MPI::PROD</code>
<code>MPI_MAXLOC</code>	<code>MPI::MAXLOC</code>
<code>MPI_MINLOC</code>	<code>MPI::MINLOC</code>
<code>MPI_BAND</code>	<code>MPI::BAND</code>
<code>MPI_BOR</code>	<code>MPI::BOR</code>
<code>MPI_BXOR</code>	<code>MPI::BXOR</code>
<code>MPI_LAND</code>	<code>MPI::LAND</code>
<code>MPI_LOR</code>	<code>MPI::LOR</code>
<code>MPI_LXOR</code>	<code>MPI::LXOR</code>
<code>MPI_REPLACE</code>	<code>MPI::REPLACE</code>
[ticket270.] <code>MPI_NO_OP</code>	i[ticket270.]Not defined for C++

Null Handles	
C/Fortran name	C++ name
C type / Fortran type	C++ type
MPI_GROUP_NULL	MPI::GROUP_NULL
MPI_Group / INTEGER	const MPI::Group
[ticket231-C.] or TYPE(MPI_Group)	
MPI_COMM_NULL	MPI::COMM_NULL
MPI_Comm / INTEGER	<sup>1)</sup>
[ticket231-C.] or TYPE(MPI_Comm)	
MPI_DATATYPE_NULL	MPI::DATATYPE_NULL
MPI_Datatype / INTEGER	const MPI::Datatype
[ticket231-C.] or TYPE(MPI_Datatype)	
MPI_REQUEST_NULL	MPI::REQUEST_NULL
MPI_Request / INTEGER	const MPI::Request
[ticket231-C.] or TYPE(MPI_Request)	
MPI_OP_NULL	MPI::OP_NULL
MPI_Op / INTEGER	const MPI::Op
[ticket231-C.] or TYPE(MPI_Op)	
MPI_ERRHANDLER_NULL	MPI::ERRHANDLER_NULL
MPI_Errhandler / INTEGER	const MPI::Errhandler
[ticket231-C.] or TYPE(MPI_Errhandler)	
MPI_FILE_NULL	MPI::FILE_NULL
MPI_File / INTEGER	
[ticket231-C.] or TYPE(MPI_File)	
MPI_INFO_NULL	MPI::INFO_NULL
MPI_Info / INTEGER	const MPI::Info
[ticket231-C.] or TYPE(MPI_Info)	
MPI_WIN_NULL	MPI::WIN_NULL
MPI_Win / INTEGER	
[ticket231-C.] or TYPE(MPI_Win)	
[ticket274.] MPI_MESSAGE_NULL	[ticket274.] Not defined for C++
[ticket274.] MPI_Message / INTEGER	
[ticket231-C.] or TYPE(MPI_Message)	
<sup>1)</sup> C++ type: See Section 16.1.7 on page 634 regarding class hierarchy and the specific type of MPI::COMM_NULL	

Empty group	
C type: MPI_Group	C++ type: const MPI::Group
Fortran type: INTEGER	
[ticket231-C.] or TYPE(MPI_Group)	
MPI_GROUP_EMPTY	MPI::GROUP_EMPTY

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Topologies	
C type: <code>const int</code> (or unnamed <code>enum</code> )	C++ type: <code>const int</code>
Fortran type: <code>INTEGER</code>	(or unnamed <code>enum</code> )
MPI_GRAPH	MPI::GRAPH
MPI_CART	MPI::CART
MPI_DIST_GRAPH	MPI::DIST_GRAPH

Predefined functions		1
C/Fortran name	C++ name	2
C type / Fortran type [ticket230-B.]with <b>mpi</b> module	C++ type	3
[ticket230-B.]/ <b>Fortran</b> type with <b>mpi_f08</b> module		4
		5
MPI_COMM_NULL_COPY_FN	MPI_COMM_NULL_COPY_FN	6
MPI_Comm_copy_attr_function	same as in C <sup>1)</sup>	7
/ COMM_COPY_ATTR_[ticket250-V.] <b>[FN] FUNCTION</b>		8
/ [ticket230-B.] <b>PROCEDURE(MPI_Comm_copy_attr_function) <sup>2)</sup></b>		9
MPI_COMM_DUP_FN	MPI_COMM_DUP_FN	10
MPI_Comm_copy_attr_function	same as in C <sup>1)</sup>	11
/ COMM_COPY_ATTR_[ticket250-V.] <b>[FN] FUNCTION</b>		12
/ [ticket230-B.] <b>PROCEDURE(MPI_Comm_copy_attr_function) <sup>2)</sup></b>		13
MPI_COMM_NULL_DELETE_FN	MPI_COMM_NULL_DELETE_FN	14
MPI_Comm_delete_attr_function	same as in C <sup>1)</sup>	15
/ COMM_DELETE_ATTR_[ticket250-V.] <b>[FN] FUNCTION</b>		16
/ [ticket230-B.] <b>PROCEDURE(MPI_Comm_delete_attr_function) <sup>2)</sup></b>		17
MPI_WIN_NULL_COPY_FN	MPI_WIN_NULL_COPY_FN	18
MPI_Win_copy_attr_function	same as in C <sup>1)</sup>	19
/ WIN_COPY_ATTR_[ticket250-V.] <b>[FN] FUNCTION</b>		20
/ [ticket230-B.] <b>PROCEDURE(MPI_Win_copy_attr_function) <sup>2)</sup></b>		21
MPI_WIN_DUP_FN	MPI_WIN_DUP_FN	22
MPI_Win_copy_attr_function	same as in C <sup>1)</sup>	23
/ WIN_COPY_ATTR_[ticket250-V.] <b>[FN] FUNCTION</b>		24
/ [ticket230-B.] <b>PROCEDURE(MPI_Win_copy_attr_function) <sup>2)</sup></b>		25
MPI_WIN_NULL_DELETE_FN	MPI_WIN_NULL_DELETE_FN	26
MPI_Win_delete_attr_function	same as in C <sup>1)</sup>	27
/ WIN_DELETE_ATTR_[ticket250-V.] <b>[FN] FUNCTION</b>		28
/ [ticket230-B.] <b>PROCEDURE(MPI_Win_delete_attr_function) <sup>2)</sup></b>		29
MPI_TYPE_NULL_COPY_FN	MPI_TYPE_NULL_COPY_FN	30
MPI_Type_copy_attr_function	same as in C <sup>1)</sup>	31
/ TYPE_COPY_ATTR_[ticket250-V.] <b>[FN] FUNCTION</b>		32
/ [ticket230-B.] <b>PROCEDURE(MPI_Type_copy_attr_function) <sup>2)</sup></b>		33
MPI_TYPE_DUP_FN	MPI_TYPE_DUP_FN	34
MPI_Type_copy_attr_function	same as in C <sup>1)</sup>	35
/ TYPE_COPY_ATTR_[ticket250-V.] <b>[FN] FUNCTION</b>		36
/ [ticket230-B.] <b>PROCEDURE(MPI_Type_copy_attr_function) <sup>2)</sup></b>		37
MPI_TYPE_NULL_DELETE_FN	MPI_TYPE_NULL_DELETE_FN	38
MPI_Type_delete_attr_function	same as in C <sup>1)</sup>	39
/ TYPE_DELETE_ATTR_[ticket250-V.] <b>[FN] FUNCTION</b>		40
/ [ticket230-B.] <b>PROCEDURE(MPI_Type_delete_attr_function) <sup>2)</sup></b>		41
<sup>1</sup> See the advice to implementors [ticket230-B.](on page 281) and advice to users (on page 281)		42
on [ticket230-B.]the predefined C functions MPI_COMM_NULL_COPY_FN, ... in		43
Section 6.7.2 on page 278		44
[ticket230-B.] <sup>2</sup> See the advice to implementors (on page 281) and advice to users (on page 281)		45
[ticket230-B.] on the predefined Fortran functions MPI_COMM_NULL_COPY_FN, ... in		46
[ticket230-B.] Section 6.7.2 on page 278		47
		48

**Deprecated predefined functions**

C/Fortran name	C++ name
C type / Fortran type	C++ type
MPI_NULL_COPY_FN	MPI::NULL_COPY_FN
MPI_Copy_function / COPY_FUNCTION	MPI::Copy_function
MPI_DUP_FN	MPI::DUP_FN
MPI_Copy_function / COPY_FUNCTION	MPI::Copy_function
MPI_NULL_DELETE_FN	MPI::NULL_DELETE_FN
MPI_Delete_function / DELETE_FUNCTION	MPI::Delete_function

**Predefined Attribute Keys**

C type: <code>const int</code> (or unnamed <code>enum</code> )	C++ type:
Fortran type: <code>INTEGER</code>	<code>const int</code> (or unnamed <code>enum</code> )
MPI_APPNUM	MPI::APPNUM
MPI_LASTUSED_CODE	MPI::LASTUSED_CODE
MPI_UNIVERSE_SIZE	MPI::UNIVERSE_SIZE
MPI_WIN_BASE	MPI::WIN_BASE
MPI_WIN_DISP_UNIT	MPI::WIN_DISP_UNIT
MPI_WIN_SIZE	MPI::WIN_SIZE
[ticket270.]MPI_WIN_CREATE_FLAVOR	[ticket270.]Not defined for C++
[ticket270.]MPI_WIN_MODEL	[ticket270.]Not defined for C++

**MPI Window Create Flavors**C type: `const int` (or unnamed `enum`)Fortran type: `INTEGER`

MPI\_WIN\_FLAVOR\_CREATE

MPI\_WIN\_FLAVOR\_ALLOCATE

MPI\_WIN\_FLAVOR\_DYNAMIC

[ticket284.]MPI\_WIN\_FLAVOR\_SHARED

**MPI Window Models**C type: `const int` (or unnamed `enum`)Fortran type: `INTEGER`

MPI\_WIN\_SEPARATE

MPI\_WIN\_UNIFIED



**Mode Constants**

C type: <code>const int</code> (or unnamed <code>enum</code> )	C++ type:
Fortran type: <code>INTEGER</code>	<code>const int</code> (or unnamed <code>enum</code> )
<code>MPI_MODE_APPEND</code>	<code>MPI::MODE_APPEND</code>
<code>MPI_MODE_CREATE</code>	<code>MPI::MODE_CREATE</code>
<code>MPI_MODE_DELETE_ON_CLOSE</code>	<code>MPI::MODE_DELETE_ON_CLOSE</code>
<code>MPI_MODE_EXCL</code>	<code>MPI::MODE_EXCL</code>
<code>MPI_MODE_NOCHECK</code>	<code>MPI::MODE_NOCHECK</code>
<code>MPI_MODE_NOPRECEDE</code>	<code>MPI::MODE_NOPRECEDE</code>
<code>MPI_MODE_NOPUT</code>	<code>MPI::MODE_NOPUT</code>
<code>MPI_MODE_NOSTORE</code>	<code>MPI::MODE_NOSTORE</code>
<code>MPI_MODE_NOSUCCEED</code>	<code>MPI::MODE_NOSUCCEED</code>
<code>MPI_MODE_RDONLY</code>	<code>MPI::MODE_RDONLY</code>
<code>MPI_MODE_RDWR</code>	<code>MPI::MODE_RDWR</code>
<code>MPI_MODE_SEQUENTIAL</code>	<code>MPI::MODE_SEQUENTIAL</code>
<code>MPI_MODE_UNIQUE_OPEN</code>	<code>MPI::MODE_UNIQUE_OPEN</code>
<code>MPI_MODE_WRONLY</code>	<code>MPI::MODE_WRONLY</code>

**Datatype Decoding Constants**

C type: <code>const int</code> (or unnamed <code>enum</code> )	C++ type:
Fortran type: <code>INTEGER</code>	<code>const int</code> (or unnamed <code>enum</code> )
<code>MPI_COMBINER_CONTIGUOUS</code>	<code>MPI::COMBINER_CONTIGUOUS</code>
<code>MPI_COMBINER_DARRAY</code>	<code>MPI::COMBINER_DARRAY</code>
<code>MPI_COMBINER_DUP</code>	<code>MPI::COMBINER_DUP</code>
<code>MPI_COMBINER_F90_COMPLEX</code>	<code>MPI::COMBINER_F90_COMPLEX</code>
<code>MPI_COMBINER_F90_INTEGER</code>	<code>MPI::COMBINER_F90_INTEGER</code>
<code>MPI_COMBINER_F90_REAL</code>	<code>MPI::COMBINER_F90_REAL</code>
<code>MPI_COMBINER_HINDEXED_INTEGER</code>	<code>MPI::COMBINER_HINDEXED_INTEGER</code>
<code>MPI_COMBINER_HINDEXED</code>	<code>MPI::COMBINER_HINDEXED</code>
<code>MPI_COMBINER_HVECTOR_INTEGER</code>	<code>MPI::COMBINER_HVECTOR_INTEGER</code>
<code>MPI_COMBINER_HVECTOR</code>	<code>MPI::COMBINER_HVECTOR</code>
<code>MPI_COMBINER_INDEXED_BLOCK</code>	<code>MPI::COMBINER_INDEXED_BLOCK</code>
<code>[ticket280.]MPI_COMBINER_HINDEXED_BLOCK</code>	
<code>MPI_COMBINER_INDEXED</code>	<code>MPI::COMBINER_INDEXED</code>
<code>MPI_COMBINER_NAMED</code>	<code>MPI::COMBINER_NAMED</code>
<code>MPI_COMBINER_RESIZED</code>	<code>MPI::COMBINER_RESIZED</code>
<code>MPI_COMBINER_STRUCT_INTEGER</code>	<code>MPI::COMBINER_STRUCT_INTEGER</code>
<code>MPI_COMBINER_STRUCT</code>	<code>MPI::COMBINER_STRUCT</code>
<code>MPI_COMBINER_SUBARRAY</code>	<code>MPI::COMBINER_SUBARRAY</code>
<code>MPI_COMBINER_VECTOR</code>	<code>MPI::COMBINER_VECTOR</code>

**Threads Constants**

C type: <code>const int</code> (or unnamed <code>enum</code> )	C++ type:
Fortran type: <code>INTEGER</code>	<code>const int</code> (or unnamed <code>enum</code> )
<code>MPI_THREAD_FUNNELED</code>	<code>MPI::THREAD_FUNNELED</code>
<code>MPI_THREAD_MULTIPLE</code>	<code>MPI::THREAD_MULTIPLE</code>
<code>MPI_THREAD_SERIALIZED</code>	<code>MPI::THREAD_SERIALIZED</code>
<code>MPI_THREAD_SINGLE</code>	<code>MPI::THREAD_SINGLE</code>

**File Operation Constants, Part 1**

C type: <code>const MPI_Offset</code> (or unnamed <code>enum</code> )	C++ type:
Fortran type: <code>INTEGER (KIND=MPI_OFFSET_KIND)</code>	<code>const MPI::Offset</code> (or unnamed <code>enum</code> )
<code>MPI_DISPLACEMENT_CURRENT</code>	<code>MPI::DISPLACEMENT_CURRENT</code>

**File Operation Constants, Part 2**

C type: <code>const int</code> (or unnamed <code>enum</code> )	C++ type:
Fortran type: <code>INTEGER</code>	<code>const int</code> (or unnamed <code>enum</code> )
<code>MPI_DISTRIBUTE_BLOCK</code>	<code>MPI::DISTRIBUTE_BLOCK</code>
<code>MPI_DISTRIBUTE_CYCLIC</code>	<code>MPI::DISTRIBUTE_CYCLIC</code>
<code>MPI_DISTRIBUTE_DFLT_DARG</code>	<code>MPI::DISTRIBUTE_DFLT_DARG</code>
<code>MPI_DISTRIBUTE_NONE</code>	<code>MPI::DISTRIBUTE_NONE</code>
<code>MPI_ORDER_C</code>	<code>MPI::ORDER_C</code>
<code>MPI_ORDER_FORTRAN</code>	<code>MPI::ORDER_FORTRAN</code>
<code>MPI_SEEK_CUR</code>	<code>MPI::SEEK_CUR</code>
<code>MPI_SEEK_END</code>	<code>MPI::SEEK_END</code>
<code>MPI_SEEK_SET</code>	<code>MPI::SEEK_SET</code>

**F90 Datatype Matching Constants**

C type: <code>const int</code> (or unnamed <code>enum</code> )	C++ type:
Fortran type: <code>INTEGER</code>	<code>const int</code> (or unnamed <code>enum</code> )
<code>MPI_TYPECLASS_COMPLEX</code>	<code>MPI::TYPECLASS_COMPLEX</code>
<code>MPI_TYPECLASS_INTEGER</code>	<code>MPI::TYPECLASS_INTEGER</code>
<code>MPI_TYPECLASS_REAL</code>	<code>MPI::TYPECLASS_REAL</code>

Constants Specifying Empty or Ignored Input			1
C/Fortran name		C++ name	2
C type / Fortran type		C++ type	3
MPI_ARGVS_NULL		MPI::ARGVS_NULL	4
char*** / 2-dim. array of CHARACTER*(*)		const char ***	5
MPI_ARGV_NULL		MPI::ARGV_NULL	6
char** / array of CHARACTER*(*)		const char **	7
MPI_ERRCODES_IGNORE		Not defined for C++	8
int* / INTEGER array			9
MPI_STATUSES_IGNORE		Not defined for C++	10
MPI_Status* / INTEGER, DIMENSION(MPI_STATUS_SIZE,*)			11
[ticket231-C.] or TYPE(MPI_Status), DIMENSION(*)			12
MPI_STATUS_IGNORE		Not defined for C++	13
MPI_Status* / INTEGER, DIMENSION(MPI_STATUS_SIZE)			14
[ticket231-C.] or TYPE(MPI_Status)			15
MPI_UNWEIGHTED		Not defined for C++	16
int* / INTEGER array			17
			18
			19
C Constants Specifying Ignored Input (no C++ or Fortran)			20
C type: MPI_Fint*	[ticket243-O.]	equivalent to Fortran	21
MPI_F_STATUSES_IGNORE	[ticket243-O.]	MPI_STATUSES_IGNORE in mpi / mpif.h	22
MPI_F_STATUS_IGNORE	[ticket243-O.]	MPI_STATUS_IGNORE in mpi / mpif.h	23
[ticket243-O.] C type: MPI_F08_status*	[ticket243-O.]	equivalent to Fortran	24
[ticket243-O.] MPI_F08_STATUSES_IGNORE	[ticket243-O.]	MPI_STATUSES_IGNORE in mpi_f08	25
[ticket243-O.] MPI_F08_STATUS_IGNORE	[ticket243-O.]	MPI_STATUS_IGNORE in mpi_f08	26
			27
			28
C and C++ preprocessor Constants and Fortran Parameters			29
C/C++ type: const int (or unnamed enum)			30
Fortran type: INTEGER			31
MPI_SUBVERSION			32
MPI_VERSION			33
			34
			35
			36
Null handles used in the MPI tool information interface			37
MPI_T_ENUM_NULL			38
MPI_T_CVAR_HANDLE_NULL			39
MPI_T_PVAR_HANDLE_NULL			40
MPI_T_PVAR_SESSION_NULL			41
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ticket266.

### Verbosity Levels in the MPI tool information interface

---

MPI\_T\_VERBOSITY\_USER\_BASIC  
 MPI\_T\_VERBOSITY\_USER\_DETAIL  
 MPI\_T\_VERBOSITY\_USER\_ALL  
 MPI\_T\_VERBOSITY\_TUNER\_BASIC  
 MPI\_T\_VERBOSITY\_TUNER\_DETAIL  
 MPI\_T\_VERBOSITY\_TUNER\_ALL  
 MPI\_T\_VERBOSITY\_MPIDEV\_BASIC  
 MPI\_T\_VERBOSITY\_MPIDEV\_DETAIL  
 MPI\_T\_VERBOSITY\_MPIDEV\_ALL

---

### Constants to identify associations of variables in the MPI tool information interface

---

MPI\_T\_BIND\_NO\_OBJECT  
 MPI\_T\_BIND\_MPI\_COMM  
 MPI\_T\_BIND\_MPI\_DATATYPE  
 MPI\_T\_BIND\_MPI\_ERRHANDLER  
 MPI\_T\_BIND\_MPI\_FILE  
 MPI\_T\_BIND\_MPI\_GROUP  
 MPI\_T\_BIND\_MPI\_OP  
 MPI\_T\_BIND\_MPI\_REQUEST  
 MPI\_T\_BIND\_MPI\_WIN  
 MPI\_T\_BIND\_MPI\_MESSAGE  
 MPI\_T\_BIND\_MPI\_INFO

---

### Constants describing the scope of a control variable in the MPI tool information interface

---

MPI\_T\_SCOPE\_READONLY  
 MPI\_T\_SCOPE\_LOCAL  
 MPI\_T\_SCOPE\_GROUP  
 MPI\_T\_SCOPE\_GROUP\_EQ  
 MPI\_T\_SCOPE\_GLOBAL  
 MPI\_T\_SCOPE\_GLOBAL\_EQ

---

### Additional constants used by the MPI tool information interface

---

MPI\_T\_PVAR\_ALL\_HANDLES

---

### Performance variables classes used by the MPI tool information interface

---

```

MPI_T_PVAR_CLASS_STATE
MPI_T_PVAR_CLASS_LEVEL
MPI_T_PVAR_CLASS_SIZE
MPI_T_PVAR_CLASS_PERCENTAGE
MPI_T_PVAR_CLASS_HIGHWATERMARK
MPI_T_PVAR_CLASS_LOWWATERMARK
MPI_T_PVAR_CLASS_COUNTER
MPI_T_PVAR_CLASS_AGGREGATE
MPI_T_PVAR_CLASS_TIMER
MPI_T_PVAR_CLASS_GENERIC

```

---

#### A.1.2 Types

The following are defined C type definitions, included in the file `mpi.h`.

```

/* C opaque types */
MPI_Aint
MPI_Count
MPI_Fint
MPI_Offset
MPI_Status
MPI_F08_status

/* C handles to assorted structures */
MPI_Comm
MPI_Datatype
MPI_Errhandler
MPI_File
MPI_Group
MPI_Info
MPI_Message
MPI_Op
MPI_Request
MPI_Win

/* Types for the MPI_T interface */
MPI_T_enum
MPI_T_cvar_handle
MPI_T_pvar_handle
MPI_T_pvar_session

// C++ opaque types (all within the MPI namespace)
MPI::Aint
MPI::Offset
MPI::Status

```

```

1
2 // C++ handles to assorted structures (classes,
3 // all within the MPI namespace)
4 MPI::Comm
5 MPI::Intracomm
6 MPI::Graphcomm
7 MPI::Distgraphcomm
8 MPI::Cartcomm
9 MPI::Intercomm
10 MPI::Datatype
11 MPI::Errhandler
12 MPI::Exception
13 MPI::File
14 MPI::Group
15 MPI::Info
16 MPI::Op
17 MPI::Request
18 MPI::Prequest
19 MPI::Grequest
20 MPI::Win

```

ticket243-O. The following are defined Fortran type definitions, included in the `mpi_f08` and `mpi` module.

```

21
22
23
24 ! Fortran opaque types in the mpi_f08 and mpi module
25 TYPE(MPI_Status)
26

```

ticket231-C. ! Fortran handles in the `mpi_f08` and `mpi` module

```

27
28 TYPE(MPI_Comm)
29 TYPE(MPI_Datatype)
30 TYPE(MPI_Errhandler)
31 TYPE(MPI_File)
32 TYPE(MPI_Group)
33 TYPE(MPI_Info)
34 TYPE(MPI_Op)
35 TYPE(MPI_Request)
36 TYPE(MPI_Win)
37

```

ticket0. A.1.3 Prototype [d]Definitions

ticket230-B. C Bindings

The following are defined C typedefs for user-defined functions, also included in the file `mpi.h`.

```

44 /* prototypes for user-defined functions */
45 typedef void MPI_User_function(void *invec, void *inoutvec, int *len,
46                               MPI_Datatype *datatype);
47
48

```

```

typedef int MPI_Comm_copy_attr_function(MPI_Comm oldcomm,
    int comm_keyval, void *extra_state, void *attribute_val_in,
    void *attribute_val_out, int*flag);
typedef int MPI_Comm_delete_attr_function(MPI_Comm comm,
    int comm_keyval, void *attribute_val, void *extra_state);

typedef int MPI_Win_copy_attr_function(MPI_Win oldwin, int win_keyval,
    void *extra_state, void *attribute_val_in,
    void *attribute_val_out, int *flag);
typedef int MPI_Win_delete_attr_function(MPI_Win win, int win_keyval,
    void *attribute_val, void *extra_state);

typedef int MPI_Type_copy_attr_function(MPI_Datatype oldtype,
    int type_keyval, void *extra_state,
    void *attribute_val_in, void *attribute_val_out, int *flag);
typedef int MPI_Type_delete_attr_function(MPI_Datatype [ticket252-W.]datatype,
    int type_keyval, void *attribute_val, void *extra_state);

typedef void MPI_Comm_errhandler_function(MPI_Comm *, int *, ...);
typedef void MPI_Win_errhandler_function(MPI_Win *, int *, ...);
typedef void MPI_File_errhandler_function(MPI_File *, int *, ...);

typedef int MPI_Grequest_query_function(void *extra_state,
    MPI_Status *status);
typedef int MPI_Grequest_free_function(void *extra_state);
typedef int MPI_Grequest_cancel_function(void *extra_state, int complete);

typedef int MPI_Datarep_extent_function(MPI_Datatype datatype,
    MPI_Aint *file_extent, void *extra_state);
typedef int MPI_Datarep_conversion_function(void *userbuf,
    MPI_Datatype datatype, int count, void *filebuf,
    MPI_Offset position, void *extra_state);

```

ticket230-B.

### Fortran 2008 Bindings with the mpi\_f08 Module

With the Fortran `mpi_f08` module, the callback prototypes are:

The user-function argument to `MPI_Op_create` should be declared according to:

#### ABSTRACT INTERFACE

```

SUBROUTINE MPI_User_function(invec, inoutvec, len, datatype) BIND(C)
    USE, INTRINSIC :: ISO_C_BINDING, ONLY : C_PTR
    TYPE(C_PTR), VALUE :: invec, inoutvec
    INTEGER :: len
    TYPE(MPI_Datatype) :: datatype

```

The copy and delete function arguments to `MPI_Comm_create_keyval` should be declared according to:

#### ABSTRACT INTERFACE

ticket230-B.

ticket-248T.

ticket230-B.

ticket-248T.

```

1      SUBROUTINE MPI_Comm_copy_attr_function(oldcomm, comm_keyval, extra_state,
2      attribute_val_in, attribute_val_out, flag, ierror) BIND(C)
3          TYPE(MPI_Comm) :: oldcomm
4          INTEGER :: comm_keyval, ierror
5          INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state, attribute_val_in,
6          attribute_val_out
7          LOGICAL :: flag

```

ticket-248T. 8

#### ABSTRACT INTERFACE

```

9      SUBROUTINE MPI_Comm_delete_attr_function(comm, comm_keyval,
10     attribute_val, extra_state, ierror) BIND(C)
11          TYPE(MPI_Comm) :: comm
12          INTEGER :: comm_keyval, ierror
13          INTEGER(KIND=MPI_ADDRESS_KIND) :: attribute_val, extra_state

```

ticket230-B. 14

15 The copy and delete function arguments to MPI\_Win\_create\_keyval should be declared  
 according to:

ticket-248T. 16

#### ABSTRACT INTERFACE

```

17     SUBROUTINE MPI_Win_copy_attr_function(oldwin, win_keyval, extra_state,
18     attribute_val_in, attribute_val_out, flag, ierror) BIND(C)
19         TYPE(MPI_Win) :: oldwin
20         INTEGER :: win_keyval, ierror
21         INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state, attribute_val_in,
22         attribute_val_out
23         LOGICAL :: flag

```

ticket-248T. 25

#### ABSTRACT INTERFACE

```

26     SUBROUTINE MPI_Win_delete_attr_function(win, win_keyval, attribute_val,
27     extra_state, ierror) BIND(C)
28         TYPE(MPI_Win) :: win
29         INTEGER :: win_keyval, ierror
30         INTEGER(KIND=MPI_ADDRESS_KIND) :: attribute_val, extra_state

```

ticket230-B. 31

32 The copy and delete function arguments to MPI\_Type\_create\_keyval should be declared  
 according to:

ticket-248T. 33

#### ABSTRACT INTERFACE

```

34     SUBROUTINE MPI_Type_copy_attr_function(oldtype, type_keyval, extra_state,
35     attribute_val_in, attribute_val_out, flag, ierror) BIND(C)
36         TYPE(MPI_Datatype) :: oldtype
37         INTEGER :: type_keyval, ierror
38         INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state, attribute_val_in,
39         attribute_val_out
40         LOGICAL :: flag

```

ticket-248T. 42

#### ABSTRACT INTERFACE

```

43     SUBROUTINE MPI_Type_delete_attr_function(datatype, type_keyval,
44     attribute_val, extra_state, ierror) BIND(C)
45         TYPE(MPI_Datatype) :: datatype
46         INTEGER :: type_keyval, ierror
47         INTEGER(KIND=MPI_ADDRESS_KIND) :: attribute_val, extra_state

```



The handler-function argument to `MPI_Comm_create_errhandler` should be declared like this:

ABSTRACT INTERFACE

```
SUBROUTINE MPI_Comm_errhandler_function(comm, error_code) BIND(C)
    TYPE(MPI_Comm) :: comm
    INTEGER :: error_code
```

The handler-function argument to `MPI_Win_create_errhandler` should be declared like this:

ABSTRACT INTERFACE

```
SUBROUTINE MPI_Win_errhandler_function(win, error_code) BIND(C)
    TYPE(MPI_Win) :: win
    INTEGER :: error_code
```

The handler-function argument to `MPI_File_create_errhandler` should be declared like this:

ABSTRACT INTERFACE

```
SUBROUTINE MPI_File_errhandler_function(file, error_code) BIND(C)
    TYPE(MPI_File) :: file
    INTEGER :: error_code
```

The query, free, and cancel function arguments to `MPI_Grequest_start` should be declared according to:

ABSTRACT INTERFACE

```
SUBROUTINE MPI_Grequest_query_function(extra_state, status, ierror)
    BIND(C)
    TYPE(MPI_Status) :: status
    INTEGER :: ierror
    INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state
```

ABSTRACT INTERFACE

```
SUBROUTINE MPI_Grequest_free_function(extra_state, ierror) BIND(C)
    INTEGER :: ierror
    INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state
```

ABSTRACT INTERFACE

```
SUBROUTINE MPI_Grequest_cancel_function(extra_state, complete, ierror)
    BIND(C)
    INTEGER :: ierror
    INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state
    LOGICAL :: complete
```

The extend and conversion function arguments to `MPI_Register_datarep` should be declared according to:

ABSTRACT INTERFACE

```
SUBROUTINE MPI_Datarep_extent_function(datatype, extent, extra_state,
    ierror) BIND(C)
    TYPE(MPI_Datatype) :: datatype
    INTEGER(KIND=MPI_ADDRESS_KIND) :: extent, extra_state
```

```

1      INTEGER :: ierror
2
3  ABSTRACT INTERFACE
4      SUBROUTINE MPI_Datarep_conversion_function(userbuf, datatype, count,
5      filebuf, position, extra_state, ierror) BIND(C)
6          USE, INTRINSIC :: ISO_C_BINDING, ONLY : C_PTR
7          TYPE(C_PTR), VALUE :: userbuf, filebuf
8          TYPE(MPI_Datatype) :: datatype
9          INTEGER :: count, ierror
10         INTEGER(KIND=MPI_OFFSET_KIND) :: position
11         INTEGER(KIND=MPI_ADDRESS_KIND) :: extra_state

```

### Fortran Bindings with mpif.h or the mpi Module

[For Fortran] With the Fortran `mpi` module or `mpif.h`, here are examples of how each of the user-defined subroutines should be declared.

The user-function argument to `MPI_OP_CREATE` should be declared like this:

```

17  SUBROUTINE USER_FUNCTION(INVEC, INOUTVEC, LEN, [ticket252-W.]DATATYPE)
18      <type> INVEC(LEN), INOUTVEC(LEN)
19      INTEGER LEN, [ticket252-W.]DATATYPE

```

The copy and delete function arguments to `MPI_COMM_CREATE_KEYVAL` should be declared like these:

```

23  SUBROUTINE COMM_COPY_ATTR_[ticket250-V.] [FN]FUNCTION(OLDCOMM, COMM_KEYVAL, EXTRA_STATE,
24      ATTRIBUTE_VAL_IN, ATTRIBUTE_VAL_OUT, FLAG, IERROR)
25      INTEGER OLDCOMM, COMM_KEYVAL, IERROR
26      INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE, ATTRIBUTE_VAL_IN,
27      ATTRIBUTE_VAL_OUT
28      LOGICAL FLAG
29
30  SUBROUTINE COMM_DELETE_ATTR_[ticket250-V.] [FN]FUNCTION(COMM, COMM_KEYVAL, ATTRIBUTE_VAL,
31      EXTRA_STATE, IERROR)
32      INTEGER COMM, COMM_KEYVAL, IERROR
33      INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL, EXTRA_STATE
34

```

The copy and delete function arguments to `MPI_WIN_CREATE_KEYVAL` should be declared like these:

```

38  SUBROUTINE WIN_COPY_ATTR_[ticket250-V.] [FN]FUNCTION(OLDWIN, WIN_KEYVAL, EXTRA_STATE,
39      ATTRIBUTE_VAL_IN, ATTRIBUTE_VAL_OUT, FLAG, IERROR)
40      INTEGER OLDWIN, WIN_KEYVAL, IERROR
41      INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE, ATTRIBUTE_VAL_IN,
42      ATTRIBUTE_VAL_OUT
43      LOGICAL FLAG
44
45  SUBROUTINE WIN_DELETE_ATTR_[ticket250-V.] [FN]FUNCTION(WIN, WIN_KEYVAL, ATTRIBUTE_VAL,
46      EXTRA_STATE, IERROR)
47      INTEGER WIN, WIN_KEYVAL, IERROR
48      INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL, EXTRA_STATE

```

The copy and delete function arguments to MPI\_TYPE\_CREATE\_KEYVAL should be declared like these:

```

SUBROUTINE TYPE_COPY_ATTR_[ticket250-V.] [FN]FUNCTION(OLDTYPE, TYPE_KEYVAL, EXTRA_STATE,
    ATTRIBUTE_VAL_IN, ATTRIBUTE_VAL_OUT, FLAG, IERROR)
    INTEGER OLDTYPE, TYPE_KEYVAL, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE,
    ATTRIBUTE_VAL_IN, ATTRIBUTE_VAL_OUT
    LOGICAL FLAG
SUBROUTINE TYPE_DELETE_ATTR_[ticket250-V.] [FN]FUNCTION([ticket252-W.]DATATYPE, TYPE_KEYVAL,
    EXTRA_STATE, IERROR)
    INTEGER [ticket252-W.]DATATYPE, TYPE_KEYVAL, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL, EXTRA_STATE

```

The handler-function argument to MPI\_COMM\_CREATE\_ERRHANDLER should be declared like this:

```

SUBROUTINE COMM_ERRHANDLER_FUNCTION(COMM, ERROR_CODE)
    INTEGER COMM, ERROR_CODE

```

The handler-function argument to MPI\_WIN\_CREATE\_ERRHANDLER should be declared like this:

```

SUBROUTINE WIN_ERRHANDLER_FUNCTION(WIN, ERROR_CODE)
    INTEGER WIN, ERROR_CODE

```

The handler-function argument to MPI\_FILE\_CREATE\_ERRHANDLER should be declared like this:

```

SUBROUTINE FILE_ERRHANDLER_FUNCTION(FILE, ERROR_CODE)
    INTEGER FILE, ERROR_CODE

```

The query, free, and cancel function arguments to MPI\_GREQUEST\_START should be declared like these:

```

SUBROUTINE GREQUEST_QUERY_FUNCTION(EXTRA_STATE, STATUS, IERROR)
    INTEGER STATUS(MPI_STATUS_SIZE), IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE
SUBROUTINE GREQUEST_FREE_FUNCTION(EXTRA_STATE, IERROR)
    INTEGER IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE
SUBROUTINE GREQUEST_CANCEL_FUNCTION(EXTRA_STATE, COMPLETE, IERROR)
    INTEGER IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE
    LOGICAL COMPLETE

```

The extend and conversion function arguments to MPI\_REGISTER\_DATAREP should be declared like these:

```

SUBROUTINE DATAREP_EXTENT_FUNCTION(DATATYPE, EXTENT, EXTRA_STATE, IERROR)
  INTEGER DATATYPE, IERROR
  INTEGER(KIND=MPI_ADDRESS_KIND) EXTENT, EXTRA_STATE

SUBROUTINE DATAREP_CONVERSION_FUNCTION(USERBUF, DATATYPE, COUNT, FILEBUF,
  POSITION, EXTRA_STATE, IERROR)
  <TYPE> USERBUF(*), FILEBUF(*)
  INTEGER COUNT, DATATYPE, IERROR
  INTEGER(KIND=MPI_OFFSET_KIND) POSITION
  INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE

```

### C++ Bindings (deprecated)

The following are **deprecated** defined C++ typedefs, also included in the file `mpi.h`.

```

namespace MPI {
  typedef void User_function(const void* invec, void *inoutvec,
    int len, const Datatype& datatype);

  typedef int Comm::Copy_attr_function(const Comm& oldcomm,
    int comm_keyval, void* extra_state, void* attribute_val_in,
    void* attribute_val_out, bool& flag);
  typedef int Comm::Delete_attr_function(Comm& comm, int
    comm_keyval, void* attribute_val, void* extra_state);

  typedef int Win::Copy_attr_function(const Win& oldwin,
    int win_keyval, void* extra_state, void* attribute_val_in,
    void* attribute_val_out, bool& flag);
  typedef int Win::Delete_attr_function(Win& win, int
    win_keyval, void* attribute_val, void* extra_state);

  typedef int Datatype::Copy_attr_function(const Datatype& oldtype,
    int type_keyval, void* extra_state,
    const void* attribute_val_in, void* attribute_val_out,
    bool& flag);
  typedef int Datatype::Delete_attr_function(Datatype& [ticket3.]0){252-W}{data}type,
    int type_keyval, void* attribute_val, void* extra_state);

  typedef void Comm::Errhandler_function(Comm &, int *, ...);
  typedef void Win::Errhandler_function(Win &, int *, ...);
  typedef void File::Errhandler_function(File &, int *, ...);

  typedef int Grequest::Query_function(void* extra_state, Status& status);
  typedef int Grequest::Free_function(void* extra_state);
  typedef int Grequest::Cancel_function(void* extra_state, bool complete);

```

```

typedef void Datarep_extent_function(const Datatype& datatype,
                                     Aint& file_extent, void* extra_state);
typedef void Datarep_conversion_function(void* userbuf,
                                         Datatype& datatype, int count, void* filebuf,
                                         Offset position, void* extra_state);
}

```

#### A.1.4 Deprecated [p]Prototype [d]Definitions

The following are defined C typedefs for deprecated user-defined functions, also included in the file `mpi.h`.

```

/* prototypes for user-defined functions */
typedef int MPI_Copy_function(MPI_Comm oldcomm, int keyval,
                             void *extra_state, void *attribute_val_in,
                             void *attribute_val_out, int *flag);
typedef int MPI_Delete_function(MPI_Comm comm, int keyval,
                               void *attribute_val, void *extra_state);
typedef void MPI_Handler_function(MPI_Comm *, int *, ...);

```

The following are deprecated Fortran user-defined callback subroutine prototypes. The deprecated copy and delete function arguments to `MPI_KEYVAL_CREATE` should be declared like these:

```

SUBROUTINE COPY_FUNCTION(OLDCOMM, KEYVAL, EXTRA_STATE,
                        ATTRIBUTE_VAL_IN, ATTRIBUTE_VAL_OUT, FLAG, IERR)
    INTEGER OLDCOMM, KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,
           ATTRIBUTE_VAL_OUT, IERR
    LOGICAL FLAG

SUBROUTINE DELETE_FUNCTION(COMM, KEYVAL, ATTRIBUTE_VAL, EXTRA_STATE, IERR)
    INTEGER COMM, KEYVAL, ATTRIBUTE_VAL, EXTRA_STATE, IERR

```

The deprecated handler-function for error handlers should be declared like this:

```

SUBROUTINE HANDLER_FUNCTION(COMM, ERROR_CODE)
    INTEGER COMM, ERROR_CODE

```

#### A.1.5 Info Keys

```

access_style
appnum
arch
cb_block_size
cb_buffer_size
cb_nodes
chunked_item
chunked_size
chunked

```

1 collective\_buffering  
2 file\_perm  
3 filename  
4 file  
5 host  
6 io\_node\_list  
7 ip\_address  
8 ip\_port  
9 nb\_proc  
10 no\_locks  
11 num\_io\_nodes  
12 path  
13 soft  
14 striping\_factor  
15 striping\_unit  
16 wdir  
17  
18

#### 19 A.1.6 Info Values

20  
21 false  
22 random  
23 read\_mostly  
24 read\_once  
25 reverse\_sequential  
26 sequential  
27 true  
28 write\_mostly  
29 write\_once  
30  
31  
32  
33  
34  
35  
36  
37  
38  
39  
40  
41  
42  
43  
44  
45  
46  
47  
48

## A.2 C Bindings

### A.2.1 Point-to-Point Communication C Bindings

```

int MPI_Bsend_init(const void* buf, int count, MPI_Datatype datatype,
                  int dest, int tag, MPI_Comm comm, MPI_Request *request)
int MPI_Bsend(const void* buf, int count, MPI_Datatype datatype, int dest,
              int tag, MPI_Comm comm)
int MPI_Buffer_attach(void* buffer, int size)
int MPI_Buffer_detach(void* buffer_addr, int* size)
int MPI_Cancel(MPI_Request *request)
int MPI_Get_count(const MPI_Status *status, MPI_Datatype datatype,
                  int *count)
int MPI_Ibsend(const void* buf, int count, MPI_Datatype datatype, int dest,
              int tag, MPI_Comm comm, MPI_Request *request)
int MPI_Improbe(int source, int tag, MPI_Comm comm, int *flag,
                MPI_Message *message, MPI_Status *status)
int MPI_Imrecv(void* buf, int count, MPI_Datatype datatype,
               MPI_Message *message, MPI_Request *request)
int MPI_Iprobe(int source, int tag, MPI_Comm comm, int *flag,
               MPI_Status *status)
int MPI_Irecv(void* buf, int count, MPI_Datatype datatype, int source,
              int tag, MPI_Comm comm, MPI_Request *request)
int MPI_Irsend(const void* buf, int count, MPI_Datatype datatype, int dest,
               int tag, MPI_Comm comm, MPI_Request *request)
int MPI_Isend(const void* buf, int count, MPI_Datatype datatype, int dest,
              int tag, MPI_Comm comm, MPI_Request *request)
int MPI_Issend(const void* buf, int count, MPI_Datatype datatype, int dest,
               int tag, MPI_Comm comm, MPI_Request *request)
int MPI_Mprobe(int source, int tag, MPI_Comm comm, MPI_Message *message,
               MPI_Status *status)
int MPI_Mrecv(void* buf, int count, MPI_Datatype datatype,
               MPI_Message *message, MPI_Status *status)
int MPI_Probe(int source, int tag, MPI_Comm comm, MPI_Status *status)
int MPI_Recv_init(void* buf, int count, MPI_Datatype datatype, int source,
                  int tag, MPI_Comm comm, MPI_Request *request)
int MPI_Recv(void* buf, int count, MPI_Datatype datatype, int source,
              int tag, MPI_Comm comm, MPI_Status *status)

```

```

1  int MPI_Request_free(MPI_Request *request)
2
3  int MPI_Request_get_status(MPI_Request request, int *flag,
4                             MPI_Status *status)
ticket140. 5  int MPI_Rsend_init(const void* buf, int count, MPI_Datatype datatype,
6                  int dest, int tag, MPI_Comm comm, MPI_Request *request)
7
ticket140. 8  int MPI_Rsend(const void* buf, int count, MPI_Datatype datatype, int dest,
9                  int tag, MPI_Comm comm)
10
ticket140. 11 int MPI_Send_init(const void* buf, int count, MPI_Datatype datatype,
12                  int dest, int tag, MPI_Comm comm, MPI_Request *request)
13
ticket140. 14 int MPI_Send(const void* buf, int count, MPI_Datatype datatype, int dest,
15                  int tag, MPI_Comm comm)
16
ticket140. 17 int MPI_Sendrecv(const void *sendbuf, int sendcount, MPI_Datatype sendtype,
18                  int dest, int sendtag, void *recvbuf, int recvcount,
19                  MPI_Datatype recvtype, int source, int recvtag, MPI_Comm comm,
20                  MPI_Status *status)
21
22 int MPI_Sendrecv_replace(void* buf, int count, MPI_Datatype datatype,
23                          int dest, int sendtag, int source, int recvtag, MPI_Comm comm,
24                          MPI_Status *status)
25
ticket140. 26 int MPI_Ssend_init(const void* buf, int count, MPI_Datatype datatype,
27                  int dest, int tag, MPI_Comm comm, MPI_Request *request)
28
ticket140. 29 int MPI_Ssend(const void* buf, int count, MPI_Datatype datatype, int dest,
30                  int tag, MPI_Comm comm)
31
ticket125. 32 int MPI_Startall(int count, MPI_Request [*]array_of_requests[])
33
ticket125. 34 int MPI_Start(MPI_Request *request)
35
ticket125. 36 int MPI_Testall(int count, MPI_Request [*]array_of_requests[], int *flag,
37                  MPI_Status [*]array_of_statuses[])
38
ticket125. 39 int MPI_Testany(int count, MPI_Request [*]array_of_requests[], int *index,
40                  int *flag, MPI_Status *status)
41
ticket125. 42 int MPI_Test_cancelled(const MPI_Status *status, int *flag)
43
ticket140. 44 int MPI_Test(MPI_Request *request, int *flag, MPI_Status *status)
45
ticket125. 46 int MPI_Testsome(int incount, MPI_Request [*]array_of_requests[],
47                  int *outcount, int [*]array_of_indices[],
48                  MPI_Status [*]array_of_statuses[])
49
ticket125. 50 int MPI_Waitall(int count, MPI_Request [*]array_of_requests[],
51                  MPI_Status [*]array_of_statuses[])
52
ticket125. 53 int MPI_Waitany(int count, MPI_Request [*]array_of_requests[], int *index,
54                  MPI_Status *status)
55
ticket125.

```



```

int MPI_Wait(MPI_Request *request, MPI_Status *status)
int MPI_Waitsome(int incout, MPI_Request [*]array_of_requests[],
                 int *outcount, int [*]array_of_indices[],
                 MPI_Status [*]array_of_statuses[])

```

### A.2.2 Datatypes C Bindings

```

int MPI_Get_address(void *location, MPI_Aint *address)
int MPI_Get_elements(MPI_Status *status, MPI_Datatype datatype, int *count)
int MPI_Pack_external(char *datarep, void *inbuf, int incout,
                     MPI_Datatype datatype, void *outbuf, MPI_Aint outsize,
                     MPI_Aint *position)
int MPI_Pack_external_size(char *datarep, int incout,
                           MPI_Datatype datatype, MPI_Aint *size)
int MPI_Pack_size(int incout, MPI_Datatype datatype, MPI_Comm comm,
                  int *size)
int MPI_Pack(void* inbuf, int incout, MPI_Datatype datatype, void *outbuf,
             int outsize, int *position, MPI_Comm comm)
int MPI_Type_commit(MPI_Datatype *datatype)
int MPI_Type_contiguous(int count, MPI_Datatype oldtype,
                        MPI_Datatype *newtype)
int MPI_Type_create_darray(int size, int rank, int ndims,
                           int array_of_gsizes[], int array_of_distribs[], int
                           array_of_dargs[], int array_of_psizes[], int order,
                           MPI_Datatype oldtype, MPI_Datatype *newtype)
int MPI_Type_create_hindexed(int count, int array_of_blocklengths[],
                             MPI_Aint array_of_displacements[], MPI_Datatype oldtype,
                             MPI_Datatype *newtype)
int MPI_Type_create_hvector(int count, int blocklength, MPI_Aint stride,
                             MPI_Datatype oldtype, MPI_Datatype *newtype)
int MPI_Type_create_indexed_block(int count, int blocklength,
                                  int array_of_displacements[], MPI_Datatype oldtype,
                                  MPI_Datatype *newtype)
int MPI_Type_create_resized(MPI_Datatype oldtype, MPI_Aint lb, MPI_Aint
                             extent, MPI_Datatype *newtype)
int MPI_Type_create_struct(int count, int array_of_blocklengths[],
                           MPI_Aint array_of_displacements[],
                           MPI_Datatype array_of_types[], MPI_Datatype *newtype)
int MPI_Type_create_subarray(int ndims, int array_of_sizes[],

```

```

1         int array_of_subsizes[], int array_of_starts[], int order,
2         MPI_Datatype oldtype, MPI_Datatype *newtype)
3
4 int MPI_Type_dup(MPI_Datatype oldtype, MPI_Datatype *newtype)
5
6 int MPI_Type_free(MPI_Datatype *datatype)
7
8 int MPI_Type_get_contents(MPI_Datatype datatype, int max_integers,
9         int max_addresses, int max_datatypes, int array_of_integers[],
10        MPI_Aint array_of_addresses[],
11        MPI_Datatype array_of_datatypes[])
12
13 int MPI_Type_get_envelope(MPI_Datatype datatype, int *num_integers,
14        int *num_addresses, int *num_datatypes, int *combiner)
15
16 int MPI_Type_get_extent(MPI_Datatype datatype, MPI_Aint *lb,
17        MPI_Aint *extent)
18
19 int MPI_Type_get_true_extent(MPI_Datatype datatype, MPI_Aint *true_lb,
20        MPI_Aint *true_extent)
21
22 int MPI_Type_indexed(int count, int *array_of_blocklengths,
23        int *array_of_displacements, MPI_Datatype oldtype,
24        MPI_Datatype *newtype)
25
26 int MPI_Type_size(MPI_Datatype datatype, int *size)
27
28 int MPI_Type_vector(int count, int blocklength, int stride,
29        MPI_Datatype oldtype, MPI_Datatype *newtype)
30
31 int MPI_Unpack_external(char *datarep, void *inbuf, MPI_Aint insize,
32        MPI_Aint *position, void *outbuf, int outcount,
33        MPI_Datatype datatype)
34
35 int MPI_Unpack(void* inbuf, int insize, int *position, void *outbuf,
36        int outcount, MPI_Datatype datatype, MPI_Comm comm)

```

### A.2.3 Collective Communication C Bindings

```

35 int MPI_Allgather(const void* sendbuf, int sendcount,
36        MPI_Datatype sendtype, void* recvbuf, int recvcount,
37        MPI_Datatype recvtype, MPI_Comm comm)
38
39 int MPI_Allgatherv(const void* sendbuf, int sendcount,
40        MPI_Datatype sendtype, void* recvbuf, const int recvcounts[],
41        const int displs[], MPI_Datatype recvtype, MPI_Comm comm)
42
43 int MPI_Allreduce(const void* sendbuf, void* recvbuf, int count,
44        MPI_Datatype datatype, MPI_Op op, MPI_Comm comm)
45
46 int MPI_Alltoall(const void* sendbuf, int sendcount, MPI_Datatype sendtype,
47        void* recvbuf, int recvcount, MPI_Datatype recvtype,
48        MPI_Comm comm)

```

```

int MPI_Alltoallv(const void* sendbuf, const int sendcounts[], const
    int sdispls[], MPI_Datatype sendtype, void* recvbuf, const
    int recvcnts[], const int rdispls[], MPI_Datatype recvtype,
    MPI_Comm comm)
int MPI_Alltoallw(const void* sendbuf, const int sendcounts[], const
    int sdispls[], const MPI_Datatype sendtypes[], void* recvbuf,
    const int recvcnts[], const int rdispls[], const
    MPI_Datatype recvtypes[], MPI_Comm comm)
int MPI_Barrier(MPI_Comm comm)
int MPI_Bcast(void* buffer, int count, MPI_Datatype datatype, int root,
    MPI_Comm comm)
int MPI_Exscan(const void* sendbuf, void* recvbuf, int count,
    MPI_Datatype datatype, MPI_Op op, MPI_Comm comm)
int MPI_Gather(const void* sendbuf, int sendcount, MPI_Datatype sendtype,
    void* recvbuf, int recvcnt, MPI_Datatype recvtype, int root,
    MPI_Comm comm)
int MPI_Gatherv(const void* sendbuf, int sendcount, MPI_Datatype sendtype,
    void* recvbuf, const int recvcnts[], const int displs[],
    MPI_Datatype recvtype, int root, MPI_Comm comm)
int MPI_Iallgather(const void* sendbuf, int sendcount,
    MPI_Datatype sendtype, void* recvbuf, int recvcnt,
    MPI_Datatype recvtype, MPI_Comm comm, MPI_Request *request)
int MPI_Iallgatherv(const void* sendbuf, int sendcount,
    MPI_Datatype sendtype, void* recvbuf, const int recvcnts[],
    const int displs[], MPI_Datatype recvtype, MPI_Comm comm,
    MPI_Request* request)
int MPI_Iallreduce(const void* sendbuf, void* recvbuf, int count,
    MPI_Datatype datatype, MPI_Op op, MPI_Comm comm,
    MPI_Request *request)
int MPI_Ialltoall(const void* sendbuf, int sendcount,
    MPI_Datatype sendtype, void* recvbuf, int recvcnt,
    MPI_Datatype recvtype, MPI_Comm comm, MPI_Request *request)
int MPI_Ialltoallv(const void* sendbuf, const int sendcounts[], const
    int sdispls[], MPI_Datatype sendtype, void* recvbuf, const
    int recvcnts[], const int rdispls[], MPI_Datatype recvtype,
    MPI_Comm comm, MPI_Request *request)
int MPI_Ialltoallw(const void* sendbuf, const int sendcounts[], const
    int sdispls[], const MPI_Datatype sendtypes[], void* recvbuf,
    const int recvcnts[], const int rdispls[], const
    MPI_Datatype recvtypes[], MPI_Comm comm, MPI_Request *request)
int MPI_Ibarrier(MPI_Comm comm, MPI_Request *request)

```

```

ticket140. 1  int MPI_Ibcast(const void* buffer, int count, MPI_Datatype datatype,
2              int root, MPI_Comm comm, MPI_Request *request)
3
ticket140. 4  int MPI_Iexscan(const void* sendbuf, void* recvbuf, int count,
5              MPI_Datatype datatype, MPI_Op op, MPI_Comm comm,
6              MPI_Request *request)
7
ticket140. 8  int MPI_Igather(const void* sendbuf, int sendcount, MPI_Datatype sendtype,
9              void* recvbuf, int recvcount, MPI_Datatype recvtype, int root,
10             MPI_Comm comm, MPI_Request *request)
11
ticket140. 12 int MPI_Igatherv(const void* sendbuf, int sendcount, MPI_Datatype sendtype,
13             void* recvbuf, const int recvcounts[], const int displs[],
14             MPI_Datatype recvtype, int root, MPI_Comm comm,
15             MPI_Request *request)
16
ticket140. 17 int MPI_Ireduce(const void* sendbuf, void* recvbuf, int count,
18             MPI_Datatype datatype, MPI_Op op, int root, MPI_Comm comm,
19             MPI_Request *request)
20
ticket140. 21 int MPI_Ireduce_scatter_block(const void* sendbuf, void* recvbuf,
22             int recvcount, MPI_Datatype datatype, MPI_Op op,
23             MPI_Comm comm, MPI_Request *request)
24
ticket140. 25 int MPI_Ireduce_scatter(const void* sendbuf, void* recvbuf, const
26             int recvcounts[], MPI_Datatype datatype, MPI_Op op,
27             MPI_Comm comm, MPI_Request *request)
28
ticket140. 29 int MPI_Iscan(const void* sendbuf, void* recvbuf, int count,
30             MPI_Datatype datatype, MPI_Op op, MPI_Comm comm,
31             MPI_Request *request)
32
ticket140. 33 int MPI_Iscatter(const void* sendbuf, int sendcount, MPI_Datatype sendtype,
34             void* recvbuf, int recvcount, MPI_Datatype recvtype, int root,
35             MPI_Comm comm, MPI_Request *request)
36
ticket140. 37 int MPI_Iscatterv(const void* sendbuf, const int sendcounts[], const
38             int displs[], MPI_Datatype sendtype, void* recvbuf,
39             int recvcount, MPI_Datatype recvtype, int root, MPI_Comm comm,
40             MPI_Request *request)
41
42 int MPI_Op_commutative(MPI_Op op, int *commute)
43
ticket252-W. 39 int MPI_Op_create(MPI_User_function* [function]user_fn, int commute,
44             MPI_Op* op)
45
46 int MPI_Op_free(MPI_Op *op)
47
ticket140. 48 int MPI_Reduce_local(const void* inbuf, void* inoutbuf, int count,
49             MPI_Datatype datatype, MPI_Op op)
50
ticket140. 51 int MPI_Reduce(const void* sendbuf, void* recvbuf, int count,
52             MPI_Datatype datatype, MPI_Op op, int root, MPI_Comm comm)
53
ticket140. 54 int MPI_Reduce_scatter_block(const void* sendbuf, void* recvbuf,

```

```

        int recvcount, MPI_Datatype datatype, MPI_Op op,
        MPI_Comm comm)
int MPI_Reduce_scatter(const void* sendbuf, void* recvbuf, const
        int recvcounts[], MPI_Datatype datatype, MPI_Op op,
        MPI_Comm comm)
int MPI_Scan(const void* sendbuf, void* recvbuf, int count,
        MPI_Datatype datatype, MPI_Op op, MPI_Comm comm)
int MPI_Scatter(const void* sendbuf, int sendcount, MPI_Datatype sendtype,
        void* recvbuf, int recvcount, MPI_Datatype recvtype, int root,
        MPI_Comm comm)
int MPI_Scatterv(const void* sendbuf, const int sendcounts[], const
        int displs[], MPI_Datatype sendtype, void* recvbuf,
        int recvcount, MPI_Datatype recvtype, int root, MPI_Comm comm)

```

#### A.2.4 Groups, Contexts, Communicators, and Caching C Bindings

```

int MPI_Comm_compare(MPI_Comm comm1, MPI_Comm comm2, int *result)
int MPI_Comm_create_keyval(MPI_Comm_copy_attr_function *comm_copy_attr_fn,
        MPI_Comm_delete_attr_function *comm_delete_attr_fn,
        int *comm_keyval, void *extra_state)
int MPI_Comm_create(MPI_Comm comm, MPI_Group group, MPI_Comm *newcomm)
int MPI_Comm_delete_attr(MPI_Comm comm, int comm_keyval)
int MPI_COMM_DUP_FN(MPI_Comm oldcomm, int comm_keyval, void *extra_state,
        void *attribute_val_in, void *attribute_val_out, int *flag)
int MPI_Comm_dup(MPI_Comm comm, MPI_Comm *newcomm)
int MPI_Comm_free_keyval(int *comm_keyval)
int MPI_Comm_free(MPI_Comm *comm)
int MPI_Comm_get_attr(MPI_Comm comm, int comm_keyval, void *attribute_val,
        int *flag)
int MPI_Comm_get_name(MPI_Comm comm, char *comm_name, int *resultlen)
int MPI_Comm_group(MPI_Comm comm, MPI_Group *group)
int MPI_COMM_NULL_COPY_FN(MPI_Comm oldcomm, int comm_keyval,
        void *extra_state, void *attribute_val_in,
        void *attribute_val_out, int *flag)
int MPI_COMM_NULL_DELETE_FN(MPI_Comm comm, int comm_keyval, void
        *attribute_val, void *extra_state)
int MPI_Comm_rank(MPI_Comm comm, int *rank)
int MPI_Comm_remote_group(MPI_Comm comm, MPI_Group *group)

```

```

1  int MPI_Comm_remote_size(MPI_Comm comm, int *size)
2
3  int MPI_Comm_set_attr(MPI_Comm comm, int comm_keyval, void *attribute_val)
ticket140. 4  int MPI_Comm_set_name(MPI_Comm comm, const char *comm_name)
5
6  int MPI_Comm_size(MPI_Comm comm, int *size)
7
8  int MPI_Comm_split(MPI_Comm comm, int color, int key, MPI_Comm *newcomm)
9
10 int MPI_Comm_test_inter(MPI_Comm comm, int *flag)
11
12 int MPI_Group_compare(MPI_Group group1, MPI_Group group2, int *result)
13
14 int MPI_Group_difference(MPI_Group group1, MPI_Group group2,
ticket140. 15 MPI_Group *newgroup)
16
17 int MPI_Group_free(MPI_Group *group)
18
19 int MPI_Group_incl(MPI_Group group, int n, const int *ranks,
ticket140. 20 MPI_Group *newgroup)
21
22 int MPI_Group_intersection(MPI_Group group1, MPI_Group group2,
23 MPI_Group *newgroup)
24
25 int MPI_Group_range_excl(MPI_Group group, int n, int ranges[][3],
26 MPI_Group *newgroup)
27
28 int MPI_Group_range_incl(MPI_Group group, int n, int ranges[][3],
29 MPI_Group *newgroup)
30
31 int MPI_Group_rank(MPI_Group group, int *rank)
32
33 int MPI_Group_size(MPI_Group group, int *size)
34
35 int MPI_Group_translate_ranks (MPI_Group group1, int n, const int *ranks1,
ticket140. 36 MPI_Group group2, int *ranks2)
37
38 int MPI_Group_union(MPI_Group group1, MPI_Group group2,
39 MPI_Group *newgroup)
40
41 int MPI_Intercomm_create(MPI_Comm local_comm, int local_leader,
42 MPI_Comm peer_comm, int remote_leader, int tag,
43 MPI_Comm *newintercomm)
44
45 int MPI_Intercomm_merge(MPI_Comm intercomm, int high,
46 MPI_Comm *newintracomm)
47
48 int MPI_Type_create_keyval(MPI_Type_copy_attr_function *type_copy_attr_fn,
49 MPI_Type_delete_attr_function *type_delete_attr_fn,
50 int *type_keyval, void *extra_state)
51
52 int MPI_Type_delete_attr(MPI_Datatype datatype, int type_keyval)
ticket252-W. 53
54 int MPI_TYPE_DUP_FN(MPI_Datatype oldtype, int type_keyval,

```

```

        void *extra_state, void *attribute_val_in,
        void *attribute_val_out, int *flag)
int MPI_Type_free_keyval(int *type_keyval)
int MPI_Type_get_attr(MPI_Datatype datatype, int type_keyval, void
    *attribute_val, int *flag)
int MPI_Type_get_name(MPI_Datatype datatype, char *type_name, int
    *resultlen)
int MPI_TYPE_NULL_COPY_FN(MPI_Datatype oldtype, int type_keyval,
    void *extra_state, void *attribute_val_in,
    void *attribute_val_out, int *flag)
int MPI_TYPE_NULL_DELETE_FN(MPI_Datatype datatype, int type_keyval, void
    *attribute_val, void *extra_state)
int MPI_Type_set_attr(MPI_Datatype datatype, int type_keyval,
    void *attribute_val)
int MPI_Type_set_name(MPI_Datatype datatype, const char *type_name)
int MPI_Win_create_keyval(MPI_Win_copy_attr_function *win_copy_attr_fn,
    MPI_Win_delete_attr_function *win_delete_attr_fn,
    int *win_keyval, void *extra_state)
int MPI_Win_delete_attr(MPI_Win win, int win_keyval)
int MPI_WIN_DUP_FN(MPI_Win oldwin, int win_keyval, void *extra_state,
    void *attribute_val_in, void *attribute_val_out, int *flag)
int MPI_Win_free_keyval(int *win_keyval)
int MPI_Win_get_attr(MPI_Win win, int win_keyval, void *attribute_val,
    int *flag)
int MPI_Win_get_name(MPI_Win win, char *win_name, int *resultlen)
int MPI_WIN_NULL_COPY_FN(MPI_Win oldwin, int win_keyval, void *extra_state,
    void *attribute_val_in, void *attribute_val_out, int *flag)
int MPI_WIN_NULL_DELETE_FN(MPI_Win win, int win_keyval, void
    *attribute_val, void *extra_state)
int MPI_Win_set_attr(MPI_Win win, int win_keyval, void *attribute_val)
int MPI_Win_set_name(MPI_Win win, const char *win_name)

```

### A.2.5 Process Topologies C Bindings

```

int MPI_Cart_coords(MPI_Comm comm, int rank, int maxdims,
    int [*coords]coords[])
int MPI_Cart_create(MPI_Comm comm_old, int ndims, const int [*dims]dims[],
    const int [*periods]periods[], int reorder,

```



```

1           MPI_Comm *comm_cart)
2
3   int MPI_Cartdim_get(MPI_Comm comm, int *ndims)
4
5   int MPI_Cart_get(MPI_Comm comm, int maxdims, int [*dims]dims[],
6                   int [*periods]periods[], int [*coords]coords[])
7
8   int MPI_Cart_map(MPI_Comm comm, int ndims, const int [*dims]dims[], const
9                   int *periods, int *newrank)
10
11  int MPI_Cart_rank(MPI_Comm comm, const int [*coords]coords[], int *rank)
12
13  int MPI_Cart_shift(MPI_Comm comm, int direction, int disp,
14                    int *rank_source, int *rank_dest)
15
16  int MPI_Cart_sub(MPI_Comm comm, const int [*remain_dims]remain_dims[],
17                  MPI_Comm *newcomm)
18
19  int MPI_Dims_create(int nnodes, int ndims, int *dims)
20
21  int MPI_Dist_graph_create_adjacent(MPI_Comm comm_old, int indegree, const
22                                   int sources[], const int sourceweights[], int outdegree, const
23                                   int destinations[], const int destweights[], MPI_Info info,
24                                   int reorder, MPI_Comm *comm_dist_graph)
25
26  int MPI_Dist_graph_create(MPI_Comm comm_old, int n, const int sources[],
27                            const int degrees[], const int destinations[], const
28                            int weights[], MPI_Info info, int reorder,
29                            MPI_Comm *comm_dist_graph)
30
31  int MPI_Dist_graph_neighbors_count(MPI_Comm comm, int *indegree,
32                                    int *outdegree, int *weighted)
33
34  int MPI_Dist_graph_neighbors(MPI_Comm comm, int maxindegree, int sources[],
35                                int sourceweights[], int maxoutdegree, int destinations[],
36                                int destweights[])
37
38  int MPI_Graph_create(MPI_Comm comm_old, int nnodes, const
39                      int [*index]index[], const int [*edges]edges[], int reorder,
40                      MPI_Comm *comm_graph)
41
42  int MPI_Graphdims_get(MPI_Comm comm, int *nnodes, int *nedges)
43
44  int MPI_Graph_get(MPI_Comm comm, int maxindex, int maxedges,
45                   int [*index]index[], int [*edges]edges[])
46
47  int MPI_Graph_map(MPI_Comm comm, int nnodes, const int [*index]index[],
48                   const int [*edges]edges[], int *newrank)
49
50  int MPI_Graph_neighbors_count(MPI_Comm comm, int rank, int *nneighbors)
51
52  int MPI_Graph_neighbors(MPI_Comm comm, int rank, int maxneighbors,
53                          int [*neighbors]neighbors[])
54
55  int MPI_Inighbor_allgather(const void* sendbuf, int sendcount,

```



```

        MPI_Datatype sendtype, void* recvbuf, int recvcnt,
        MPI_Datatype recvtype, MPI_Comm comm, MPI_Request *request)
1
2
3
int MPI_Ineighbor_allgatherv(const void* sendbuf, int sendcount,
4 ticket140.
    MPI_Datatype sendtype, void* recvbuf, const int recvcnts[],
5 ticket140.
    const int displs[], MPI_Datatype recvtype, MPI_Comm comm,
6 ticket140.
    MPI_Request *request)
7
8
int MPI_Ineighbor_alltoall(const void* sendbuf, int sendcount, MPI_Datatype
9 ticket140.
    sendtype, void* recvbuf, int recvcnt, MPI_Datatype recvtype,
10
    MPI_Comm comm, MPI_Request *request)
11
12
int MPI_Ineighbor_alltoallv(const void* sendbuf, const int sendcounts[],
12 ticket140.
    const int sdispls[], MPI_Datatype sendtype, void* recvbuf,
13 ticket140.
    const int recvcnts[], const int rdispls[], MPI_Datatype
14 ticket140.
    recvtype, MPI_Comm comm, MPI_Request *request)
15 ticket140.
16
int MPI_Ineighbor_alltoallw(const void* sendbuf, const int sendcounts[],
16 ticket140.
    const [int]MPI_Aint sdispls[], const MPI_Datatype sendtypes[],
17 ticket140.
    void* recvbuf, const int recvcnts[], const [int]MPI_Aint
18 ticket140.
    rdispls[], const MPI_Datatype recvtypes[], MPI_Comm comm,
19 ticket299.
    MPI_Request *request)
20 ticket140.
21
int MPI_Neighbor_allgather(const void* sendbuf, int sendcount, MPI_Datatype
22 ticket140.
    sendtype, void* recvbuf, int recvcnt, MPI_Datatype recvtype,
23 ticket299.
    MPI_Comm comm)
24 ticket140.
25
int MPI_Neighbor_allgatherv(const void* sendbuf, int sendcount,
25 ticket140.
    MPI_Datatype sendtype, void* recvbuf, const int recvcnts[],
26 ticket140.
    const int displs[], MPI_Datatype recvtype, MPI_Comm comm)
27 ticket140.
28
int MPI_Neighbor_alltoall(const void* sendbuf, int sendcount, MPI_Datatype
29 ticket140.
    sendtype, void* recvbuf, int recvcnt, MPI_Datatype recvtype,
30
    MPI_Comm comm)
31
32
int MPI_Neighbor_alltoallv(const void* sendbuf, const int sendcounts[],
32 ticket140.
    const int sdispls[], MPI_Datatype sendtype, void* recvbuf,
33 ticket140.
    const int recvcnts[], const int rdispls[], MPI_Datatype
34 ticket140.
    recvtype, MPI_Comm comm)
35 ticket140.
36
int MPI_Neighbor_alltoallw(const void* sendbuf, const int sendcounts[],
37 ticket140.
    const [int]MPI_Aint sdispls[], const MPI_Datatype sendtypes[],
38 ticket140.
    void* recvbuf, const int recvcnts[], const [int]MPI_Aint
39 ticket140.
    rdispls[], const MPI_Datatype recvtypes[], MPI_Comm comm)
40 ticket299.
41
int MPI_Topo_test(MPI_Comm comm, int *status)
41 ticket140.
42
43
44
45
46
47
48

```

## A.2.6 MPI Environmental Management C Bindings

```
int MPI_Abort(MPI_Comm comm, int errorcode)
```

```
int MPI_Add_error_class(int *errorclass)
```

```

1  int MPI_Add_error_code(int errorclass, int *errorcode)
2
3  int MPI_Add_error_string(int errorcode, char *string)
4
5  int MPI_Alloc_mem(MPI_Aint size, MPI_Info info, void *baseptr)
6
7  int MPI_Comm_call_errhandler(MPI_Comm comm, int errorcode)
8  int MPI_Comm_create_errhandler(MPI_Comm_errhandler_function
ticket252-W. 9      *[function]comm_errhandler_fn, MPI_Errhandler *errhandler)
10
11 int MPI_Comm_get_errhandler(MPI_Comm comm, MPI_Errhandler *errhandler)
12
13 int MPI_Comm_set_errhandler(MPI_Comm comm, MPI_Errhandler errhandler)
14
15 int MPI_Error_class(int errorcode, int *errorclass)
16
17 int MPI_Error_string(int errorcode, char *string, int *resultlen)
18
19 int MPI_File_call_errhandler(MPI_File fh, int errorcode)
20 int MPI_File_create_errhandler(MPI_File_errhandler_function
ticket252-W. 21      *[function]file_errhandler_fn, MPI_Errhandler *errhandler)
22
23 int MPI_File_get_errhandler(MPI_File file, MPI_Errhandler *errhandler)
24
25 int MPI_File_set_errhandler(MPI_File file, MPI_Errhandler errhandler)
26
27 int MPI_Finalized(int *flag)
28
29 int MPI_Finalize(void)
30
31 int MPI_Free_mem(void *base)
32
33 int MPI_Get_processor_name(char *name, int *resultlen)
34
35 int MPI_Get_version(int *version, int *subversion)
36
37 int MPI_Initialized(int *flag)
38
39 int MPI_Init(int *argc, char ***argv)
40
41 int MPI_Win_call_errhandler(MPI_Win win, int errorcode)
42
43 int MPI_Win_create_errhandler(MPI_Win_errhandler_function
ticket252-W. 44      *[function]win_errhandler_fn, MPI_Errhandler *errhandler)
45
46 int MPI_Win_get_errhandler(MPI_Win win, MPI_Errhandler *errhandler)
47
48 int MPI_Win_set_errhandler(MPI_Win win, MPI_Errhandler errhandler)
49
50 double MPI_Wtick(void)
51
52 double MPI_Wtime(void)
53
54
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99

```

## A.2.7 The Info Object C Bindings

```

int MPI_Info_create(MPI_Info *info)

```

```

int MPI_Info_delete(MPI_Info info, const char *key)
int MPI_Info_dup(MPI_Info info, MPI_Info *newinfo)
int MPI_Info_free(MPI_Info *info)
int MPI_Info_get(MPI_Info info, const char *key, int valuelen, char *value,
                  int *flag)
int MPI_Info_get_nkeys(MPI_Info info, int *nkeys)
int MPI_Info_get_nthkey(MPI_Info info, int n, char *key)
int MPI_Info_get_valuelen(MPI_Info info, const char *key, int *valuelen,
                           int *flag)
int MPI_Info_set(MPI_Info info, const char *key, const char *value)

```

### A.2.8 Process Creation and Management C Bindings

```

int MPI_Close_port(char *port_name)
int MPI_Comm_accept(char *port_name, MPI_Info info, int root,
                    MPI_Comm comm, MPI_Comm *newcomm)
int MPI_Comm_connect(char *port_name, MPI_Info info, int root,
                     MPI_Comm comm, MPI_Comm *newcomm)
int MPI_Comm_disconnect(MPI_Comm *comm)
int MPI_Comm_get_parent(MPI_Comm *parent)
int MPI_Comm_join(int fd, MPI_Comm *intercomm)
int MPI_Comm_spawn(char *command, char *argv[], int maxprocs, MPI_Info
                   info, int root, MPI_Comm comm, MPI_Comm *intercomm,
                   int array_of_errcodes[])
int MPI_Comm_spawn_multiple(int count, char *array_of_commands[],
                             char **array_of_argv[], int array_of_maxprocs[],
                             MPI_Info array_of_info[], int root, MPI_Comm comm,
                             MPI_Comm *intercomm, int array_of_errcodes[])
int MPI_Lookup_name(char *service_name, MPI_Info info, char *port_name)
int MPI_Open_port(MPI_Info info, char *port_name)
int MPI_Publish_name(char *service_name, MPI_Info info, char *port_name)
int MPI_Unpublish_name(char *service_name, MPI_Info info, char *port_name)

```

### A.2.9 One-Sided Communications C Bindings

```

int MPI_Accumulate(const void *origin_addr, int origin_count,
                  MPI_Datatype origin_datatype, int target_rank,

```

```

1          MPI_Aint target_disp, int target_count,
2          MPI_Datatype target_datatype, MPI_Op op, MPI_Win win)
3
ticket140a. 4 int MPI_Compare_and_swap(const void *origin_addr, const void *compare_addr,
ticket140a. 5          void *result_addr, MPI_Datatype datatype, int target_rank,
6          MPI_Aint target_disp, MPI_Win win)
7
ticket140a. 8 int MPI_Fetch_and_op(const void *origin_addr, void *result_addr,
9          MPI_Datatype datatype, int target_rank, MPI_Aint target_disp,
10         MPI_Op op, MPI_Win win)
11
ticket140a. 12 int MPI_Get_accumulate(const void *origin_addr, int origin_count,
13         MPI_Datatype origin_datatype, void *result_addr,
14         int result_count, MPI_Datatype result_datatype,
15         int target_rank, MPI_Aint target_disp, int target_count,
16         MPI_Datatype target_datatype, MPI_Op op, MPI_Win win)
17
18 int MPI_Get(void *origin_addr, int origin_count,
19         MPI_Datatype origin_datatype, int target_rank,
20         MPI_Aint target_disp, int target_count,
21         MPI_Datatype target_datatype, MPI_Win win)
22
ticket140. 23 int MPI_Put(const void *origin_addr, int origin_count, MPI_Datatype
24         origin_datatype, int target_rank, MPI_Aint target_disp, int
25         target_count, MPI_Datatype target_datatype, MPI_Win win)
26
ticket140a. 27 int MPI_Raccumulate(const void *origin_addr, int origin_count,
28         MPI_Datatype origin_datatype, int target_rank,
29         MPI_Aint target_disp, int target_count,
30         MPI_Datatype target_datatype, MPI_Op op, MPI_Win win,
31         MPI_Request *request)
32
ticket140a. 33 int MPI_Rget_accumulate(const void *origin_addr, int origin_count,
34         MPI_Datatype origin_datatype, void *result_addr,
35         int result_count, MPI_Datatype result_datatype,
36         int target_rank, MPI_Aint target_disp, int target_count,
37         MPI_Datatype target_datatype, MPI_Op op, MPI_Win win,
38         MPI_Request *request)
39
40 int MPI_Rget(void *origin_addr, int origin_count,
41         MPI_Datatype origin_datatype, int target_rank,
42         MPI_Aint target_disp, int target_count,
43         MPI_Datatype target_datatype, MPI_Win win,
44         MPI_Request *request)
45
ticket140a. 46 int MPI_Rput(const void *origin_addr, int origin_count,
47         MPI_Datatype origin_datatype, int target_rank,
48         MPI_Aint target_disp, int target_count,
49         MPI_Datatype target_datatype, MPI_Win win,
50         MPI_Request *request)
51
52 int MPI_Win_allocate(MPI_Aint size, int disp_unit, MPI_Info info,

```

```

        MPI_Comm comm, void *baseptr, MPI_Win *win)
int MPI_Win_allocate_shared(MPI_Aint size, MPI_Info info, MPI_Comm comm,
        void *baseptr, MPI_Win *win)
int MPI_Win_attach(MPI_Win win, void *base, MPI_Aint size)
int MPI_Win_complete(MPI_Win win)
int MPI_Win_create_dynamic(MPI_Info info, MPI_Comm comm, MPI_Win *win)
int MPI_Win_create(void *base, MPI_Aint size, int disp_unit, MPI_Info info,
        MPI_Comm comm, MPI_Win *win)
int MPI_Win_detach(MPI_Win win, const void *base)
int MPI_Win_fence(int assert, MPI_Win win)
int MPI_Win_flush_all(MPI_Win win)
int MPI_Win_flush(int rank, MPI_Win win)
int MPI_Win_flush_local_all(MPI_Win win)
int MPI_Win_flush_local(int rank, MPI_Win win)
int MPI_Win_free(MPI_Win *win)
int MPI_Win_get_group(MPI_Win win, MPI_Group *group)
int MPI_Win_lock_all(int assert, MPI_Win win)
int MPI_Win_lock(int lock_type, int rank, int assert, MPI_Win win)
int MPI_Win_post(MPI_Group group, int assert, MPI_Win win)
int MPI_Win_shared_query(MPI_Win win, int rank, MPI_Aint *size,
        void *baseptr)
int MPI_Win_start(MPI_Group group, int assert, MPI_Win win)
int MPI_Win_sync(MPI_Win win)
int MPI_Win_test(MPI_Win win, int *flag)
int MPI_Win_unlock_all(MPI_Win win)
int MPI_Win_unlock(int rank, MPI_Win win)
int MPI_Win_wait(MPI_Win win)

```

#### A.2.10 External Interfaces C Bindings

```

int MPI_Grequest_complete(MPI_Request request)
int MPI_Grequest_start(MPI_Grequest_query_function *query_fn,
        MPI_Grequest_free_function *free_fn,
        MPI_Grequest_cancel_function *cancel_fn, void *extra_state,
        MPI_Request *request)

```

```

1  int MPI_Init_thread(int *argc, char *((*argv)[]), int required,
2      int *provided)
3
4  int MPI_Is_thread_main(int *flag)
5
6  int MPI_Query_thread(int *provided)
7
8  int MPI_Status_set_cancelled(MPI_Status *status, int flag)
9
10 int MPI_Status_set_elements(MPI_Status *status, MPI_Datatype datatype,
11     int count)
12
13 int MPI_Status_set_elements_x(MPI_Status *status, MPI_Datatype datatype,
14     MPI_Count count)

```

#### A.2.11 I/O C Bindings

```

15 int MPI_File_close(MPI_File *fh)
16
17 int MPI_File_delete(const char *filename, MPI_Info info)
18
19 int MPI_File_get_amode(MPI_File fh, int *amode)
20
21 int MPI_File_get_atomicity(MPI_File fh, int *flag)
22
23 int MPI_File_get_byte_offset(MPI_File fh, MPI_Offset offset,
24     MPI_Offset *disp)
25
26 int MPI_File_get_group(MPI_File fh, MPI_Group *group)
27
28 int MPI_File_get_info(MPI_File fh, MPI_Info *info_used)
29
30 int MPI_File_get_position(MPI_File fh, MPI_Offset *offset)
31
32 int MPI_File_get_position_shared(MPI_File fh, MPI_Offset *offset)
33
34 int MPI_File_get_size(MPI_File fh, MPI_Offset *size)
35
36 int MPI_File_get_type_extent(MPI_File fh, MPI_Datatype datatype,
37     MPI_Aint *extent)
38
39 int MPI_File_get_view(MPI_File fh, MPI_Offset *disp, MPI_Datatype *etype,
40     MPI_Datatype *filetype, char *datarep)
41
42 int MPI_File_iread_at(MPI_File fh, MPI_Offset offset, void *buf, int count,
43     MPI_Datatype datatype, MPI_Request *request)
44
45 int MPI_File_iread(MPI_File fh, void *buf, int count,
46     MPI_Datatype datatype, MPI_Request *request)
47
48 int MPI_File_iread_shared(MPI_File fh, void *buf, int count,
49     MPI_Datatype datatype, MPI_Request *request)
50
51 int MPI_File_iwrite_at(MPI_File fh, MPI_Offset offset, const void *buf,
52     int count, MPI_Datatype datatype, MPI_Request *request)
53
54 int MPI_File_iwrite(MPI_File fh, const void *buf, int count,

```

```

        MPI_Datatype datatype, MPI_Request *request)
1
2
int MPI_File_iwrite_shared(MPI_File fh, const void *buf, int count,
3 ticket140.
        MPI_Datatype datatype, MPI_Request *request)
4
int MPI_File_open(MPI_Comm comm, const char *filename, int amode,
5 ticket140.
        MPI_Info info, MPI_File *fh)
6
int MPI_File_preallocate(MPI_File fh, MPI_Offset size)
7
int MPI_File_read_all_begin(MPI_File fh, void *buf, int count,
8
        MPI_Datatype datatype)
9
int MPI_File_read_all_end(MPI_File fh, void *buf, MPI_Status *status)
10
11
int MPI_File_read_all(MPI_File fh, void *buf, int count,
12
        MPI_Datatype datatype, MPI_Status *status)
13
14
int MPI_File_read_at_all_begin(MPI_File fh, MPI_Offset offset, void *buf,
15
        int count, MPI_Datatype datatype)
16
17
int MPI_File_read_at_all_end(MPI_File fh, void *buf, MPI_Status *status)
18
19
int MPI_File_read_at_all(MPI_File fh, MPI_Offset offset, void *buf,
20
        int count, MPI_Datatype datatype, MPI_Status *status)
21
22
int MPI_File_read_at(MPI_File fh, MPI_Offset offset, void *buf, int count,
23
        MPI_Datatype datatype, MPI_Status *status)
24
int MPI_File_read(MPI_File fh, void *buf, int count, MPI_Datatype datatype,
25
        MPI_Status *status)
26
int MPI_File_read_ordered_begin(MPI_File fh, void *buf, int count,
27
        MPI_Datatype datatype)
28
29
int MPI_File_read_ordered_end(MPI_File fh, void *buf, MPI_Status *status)
30
31
int MPI_File_read_ordered(MPI_File fh, void *buf, int count,
32
        MPI_Datatype datatype, MPI_Status *status)
33
int MPI_File_read_shared(MPI_File fh, void *buf, int count,
34
        MPI_Datatype datatype, MPI_Status *status)
35
36
int MPI_File_seek(MPI_File fh, MPI_Offset offset, int whence)
37
int MPI_File_seek_shared(MPI_File fh, MPI_Offset offset, int whence)
38
39
int MPI_File_set_atomicity(MPI_File fh, int flag)
40
int MPI_File_set_info(MPI_File fh, MPI_Info info)
41
42
int MPI_File_set_size(MPI_File fh, MPI_Offset size)
43
int MPI_File_set_view(MPI_File fh, MPI_Offset disp, MPI_Datatype etype,
44
        MPI_Datatype filetype, const char *datarep, MPI_Info info)
45 ticket140.
46
int MPI_File_sync(MPI_File fh)
47
48

```

```

ticket140. 1  int MPI_File_write_all_begin(MPI_File fh, const void *buf, int count,
2             MPI_Datatype datatype)
3
ticket140. 4  int MPI_File_write_all_end(MPI_File fh, const void *buf,
5             MPI_Status *status)
6
ticket140. 7  int MPI_File_write_all(MPI_File fh, const void *buf, int count,
8             MPI_Datatype datatype, MPI_Status *status)
9
ticket140. 10 int MPI_File_write_at_all_begin(MPI_File fh, MPI_Offset offset, const
11          void *buf, int count, MPI_Datatype datatype)
12
ticket140. 13 int MPI_File_write_at_all_end(MPI_File fh, const void *buf,
14          MPI_Status *status)
15
ticket140. 16 int MPI_File_write_at_all(MPI_File fh, MPI_Offset offset, const void *buf,
17          int count, MPI_Datatype datatype, MPI_Status *status)
18
ticket140. 19 int MPI_File_write(MPI_File fh, const void *buf, int count,
20          MPI_Datatype datatype, MPI_Status *status)
21
ticket140. 22 int MPI_File_write_ordered_begin(MPI_File fh, const void *buf, int count,
23          MPI_Datatype datatype)
24
ticket140. 25 int MPI_File_write_ordered_end(MPI_File fh, const void *buf,
26          MPI_Status *status)
27
ticket140. 28 int MPI_File_write_ordered(MPI_File fh, const void *buf, int count,
29          MPI_Datatype datatype, MPI_Status *status)
30
ticket140. 31 int MPI_File_write_shared(MPI_File fh, const void *buf, int count,
32          MPI_Datatype datatype, MPI_Status *status)
33
ticket140. 34 int MPI_Register_datarep(const char *datarep,
35          MPI_Datarep_conversion_function *read_conversion_fn,
36          MPI_Datarep_conversion_function *write_conversion_fn,
37          MPI_Datarep_extent_function *dtype_file_extent_fn,
38          void *extra_state)
39
A.2.12 Language Bindings C Bindings
40
int MPI_Status_f082f(MPI_F08_status *f08_status, MPI_Fint *f_status)
41
int MPI_Status_f2f08(MPI_Fint *f_status, MPI_F08_status *f08_status)
42
int MPI_Type_create_f90_complex(int p, int r, MPI_Datatype *newtype)
43
int MPI_Type_create_f90_integer(int r, MPI_Datatype *newtype)
44
int MPI_Type_create_f90_real(int p, int r, MPI_Datatype *newtype)
45
ticket252-W. 47 int MPI_Type_match_size(int typeclass, int size, MPI_Datatype *datatype)
48

```



```

MPI_Fint MPI_Comm_c2f(MPI_Comm comm) 1
MPI_Comm MPI_Comm_f2c(MPI_Fint comm) 2
MPI_Fint MPI_Errhandler_c2f(MPI_Errhandler errhandler) 3
MPI_Errhandler MPI_Errhandler_f2c(MPI_Fint errhandler) 4
MPI_Fint MPI_File_c2f(MPI_File file) 5
MPI_File MPI_File_f2c(MPI_Fint file) 6
MPI_Fint MPI_Group_c2f(MPI_Group group) 7
MPI_Group MPI_Group_f2c(MPI_Fint group) 8
MPI_Fint MPI_Info_c2f(MPI_Info info) 9
MPI_Info MPI_Info_f2c(MPI_Fint info) 10
MPI_Fint MPI_Message_c2f(MPI_Message message) 11
MPI_Message MPI_Message_f2c(MPI_Fint message) 12
MPI_Fint MPI_Op_c2f(MPI_Op op) 13
MPI_Op MPI_Op_f2c(MPI_Fint op) 14
MPI_Fint MPI_Request_c2f(MPI_Request request) 15
MPI_Request MPI_Request_f2c(MPI_Fint request) 16
int MPI_Status_c2f08(const MPI_Status *c_status, MPI_F08_status 17
    *f08_status) 18 ticket140.
int MPI_Status_c2f(const MPI_Status *c_status, MPI_Fint *f_status) 19 ticket140.
int MPI_Status_f082c(const MPI_F08_status *f08_status, MPI_Status 20 ticket140.
    *c_status) 21
int MPI_Status_f2c(const MPI_Fint *f_status, MPI_Status *c_status) 22 ticket140.
MPI_Fint MPI_Type_c2f(MPI_Datatype datatype) 23
MPI_Datatype MPI_Type_f2c(MPI_Fint datatype) 24
MPI_Fint MPI_Win_c2f(MPI_Win win) 25
MPI_Win MPI_Win_f2c(MPI_Fint win) 26

```

### A.2.13 Profiling Interface C Bindings

```
int MPI_Pcontrol(const int level, ...)
```

### A.2.14 Deprecated C Bindings

```
int MPI_Address(void* location, MPI_Aint *address)
```

```

1  int MPI_Attr_delete(MPI_Comm comm, int keyval)
2
3  int MPI_Attr_get(MPI_Comm comm, int keyval, void *attribute_val, int *flag)
4
5  int MPI_Attr_put(MPI_Comm comm, int keyval, void* attribute_val)
6
7  int MPI_DUP_FN(MPI_Comm oldcomm, int keyval, void *extra_state,
8                void *attribute_val_in, void *attribute_val_out, int *flag)
9
10 int MPI_Errhandler_create(MPI_Handler_function *[function]handler_fn,
11                           MPI_Errhandler *errhandler)
12
13 int MPI_Errhandler_get(MPI_Comm comm, MPI_Errhandler *errhandler)
14
15 int MPI_Errhandler_set(MPI_Comm comm, MPI_Errhandler errhandler)
16
17 int MPI_Keyval_create(MPI_Copy_function *copy_fn, MPI_Delete_function
18                      *delete_fn, int *keyval, void* extra_state)
19
20 int MPI_Keyval_free(int *keyval)
21
22 int MPI_NULL_COPY_FN(MPI_Comm oldcomm, int keyval, void *extra_state,
23                      void *attribute_val_in, void *attribute_val_out, int *flag)
24
25 int MPI_NULL_DELETE_FN(MPI_Comm comm, int keyval, void *attribute_val,
26                        void *extra_state)
27
28 int MPI_Type_extent(MPI_Datatype datatype, MPI_Aint *extent)
29
30 int MPI_Type_hindexed(int count, int *array_of_blocklengths,
31                      MPI_Aint *array_of_displacements, MPI_Datatype oldtype,
32                      MPI_Datatype *newtype)
33
34 int MPI_Type_hvector(int count, int blocklength, MPI_Aint stride,
35                     MPI_Datatype oldtype, MPI_Datatype *newtype)
36
37 int MPI_Type_lb(MPI_Datatype datatype, MPI_Aint* displacement)
38
39 int MPI_Type_struct(int count, int *array_of_blocklengths,
40                    MPI_Aint *array_of_displacements,
41                    MPI_Datatype *array_of_types, MPI_Datatype *newtype)
42
43 int MPI_Type_ub(MPI_Datatype datatype, MPI_Aint* displacement)
44
45
46
47
48

```

## A.3 Fortran 2008 Bindings with the mpi\_f08 Module

### A.3.1 Point-to-Point Communication Fortran 2008 Bindings

```

MPI_Bsend(buf, count, datatype, dest, tag, comm, ierror) BIND(C)
  TYPE(*), DIMENSION(..), INTENT(IN) :: buf
  INTEGER, INTENT(IN) :: count, dest, tag
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  TYPE(MPI_Comm), INTENT(IN) :: comm
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Bsend_init(buf, count, datatype, dest, tag, comm, request, ierror)
  BIND(C)
  TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
  INTEGER, INTENT(IN) :: count, dest, tag
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  TYPE(MPI_Comm), INTENT(IN) :: comm
  TYPE(MPI_Request), INTENT(OUT) :: request
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Buffer_attach(buffer, size, ierror) BIND(C)
  TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buffer
  INTEGER, INTENT(IN) :: size
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Buffer_detach(buffer_addr, size, ierror) BIND(C)
  USE, INTRINSIC :: ISO_C_BINDING, ONLY : C_PTR
  TYPE(C_PTR), INTENT(OUT) :: buffer_addr
  INTEGER, INTENT(OUT) :: size
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Cancel(request, ierror) BIND(C)
  TYPE(MPI_Request), INTENT(IN) :: request
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Get_count(status, datatype, count, ierror) BIND(C)
  TYPE(MPI_Status), INTENT(IN) :: status
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  INTEGER, INTENT(OUT) :: count
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Ibsend(buf, count, datatype, dest, tag, comm, request, ierror) BIND(C)
  TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
  INTEGER, INTENT(IN) :: count, dest, tag
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  TYPE(MPI_Comm), INTENT(IN) :: comm
  TYPE(MPI_Request), INTENT(OUT) :: request
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

1  MPI_Improbe(source, tag, comm, flag, message, status, ierror) BIND(C)
2      INTEGER, INTENT(IN) :: source, tag
3      TYPE(MPI_Comm), INTENT(IN) :: comm
4      INTEGER, INTENT(OUT) :: flag
5      TYPE(MPI_Message), INTENT(OUT) :: message
6      TYPE(MPI_Status) :: status
7      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T.
9  MPI_Imrecv(buf, count, datatype, message, request, ierror) BIND(C)
10     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
11     INTEGER, INTENT(IN) :: count
12     TYPE(MPI_Datatype), INTENT(IN) :: datatype
13     TYPE(MPI_Message), INTENT(INOUT) :: message
14     TYPE(MPI_Request), INTENT(OUT) :: request
15     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T.
16  MPI_Iprobe(source, tag, comm, flag, status, ierror) BIND(C)
17     INTEGER, INTENT(IN) :: source, tag
18     TYPE(MPI_Comm), INTENT(IN) :: comm
19     LOGICAL, INTENT(OUT) :: flag
20     TYPE(MPI_Status) :: status
21     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T.
22  MPI_Irecv(buf, count, datatype, source, tag, comm, request, ierror) BIND(C)
23     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
24     INTEGER, INTENT(IN) :: count, source, tag
25     TYPE(MPI_Datatype), INTENT(IN) :: datatype
26     TYPE(MPI_Comm), INTENT(IN) :: comm
27     TYPE(MPI_Request), INTENT(OUT) :: request
28     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T.
29  MPI_Irsend(buf, count, datatype, dest, tag, comm, request, ierror) BIND(C)
30     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
31     INTEGER, INTENT(IN) :: count, dest, tag
32     TYPE(MPI_Datatype), INTENT(IN) :: datatype
33     TYPE(MPI_Comm), INTENT(IN) :: comm
34     TYPE(MPI_Request), INTENT(OUT) :: request
35     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T.
36  MPI_Isend(buf, count, datatype, dest, tag, comm, request, ierror) BIND(C)
37     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
38     INTEGER, INTENT(IN) :: count, dest, tag
39     TYPE(MPI_Datatype), INTENT(IN) :: datatype
40     TYPE(MPI_Comm), INTENT(IN) :: comm
41     TYPE(MPI_Request), INTENT(OUT) :: request
42     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T.
43  MPI_Issend(buf, count, datatype, dest, tag, comm, request, ierror) BIND(C)
44     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
45     INTEGER, INTENT(IN) :: count, dest, tag
46
47
48

```

```

TYPE(MPI_Datatype), INTENT(IN) :: datatype
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Request), INTENT(OUT) :: request
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Mprobe(source, tag, comm, message, status, ierror) BIND(C)
  INTEGER, INTENT(IN) :: source, tag
  TYPE(MPI_Comm), INTENT(IN) :: comm
  TYPE(MPI_Message), INTENT(OUT) :: message
  TYPE(MPI_Status) :: status
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Mrecv(buf, count, datatype, message, status, ierror) BIND(C)
  TYPE(*), DIMENSION(..) :: buf
  INTEGER, INTENT(IN) :: count
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  TYPE(MPI_Message), INTENT(INOUT) :: message
  TYPE(MPI_Status) :: status
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Probe(source, tag, comm, status, ierror) BIND(C)
  INTEGER, INTENT(IN) :: source, tag
  TYPE(MPI_Comm), INTENT(IN) :: comm
  TYPE(MPI_Status) :: status
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Recv(buf, count, datatype, source, tag, comm, status, ierror) BIND(C)
  TYPE(*), DIMENSION(..) :: buf
  INTEGER, INTENT(IN) :: count, source, tag
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  TYPE(MPI_Comm), INTENT(IN) :: comm
  TYPE(MPI_Status) :: status
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Recv_init(buf, count, datatype, source, tag, comm, request, ierror)
  BIND(C)
  TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
  INTEGER, INTENT(IN) :: count, source, tag
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  TYPE(MPI_Comm), INTENT(IN) :: comm
  TYPE(MPI_Request), INTENT(OUT) :: request
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Request_free(request, ierror) BIND(C)
  TYPE(MPI_Request), INTENT(INOUT) :: request
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Request_get_status(request, flag, status, ierror) BIND(C)
  TYPE(MPI_Request), INTENT(IN) :: request
  LOGICAL, INTENT(OUT) :: flag
  TYPE(MPI_Status) :: status

```

```

1      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 2
3      MPI_Rsend(buf, count, datatype, dest, tag, comm, ierror) BIND(C)
4      TYPE(*), DIMENSION(..), INTENT(IN) :: buf
5      INTEGER, INTENT(IN) :: count, dest, tag
6      TYPE(MPI_Datatype), INTENT(IN) :: datatype
7      TYPE(MPI_Comm), INTENT(IN) :: comm
8      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 9
10     MPI_Rsend_init(buf, count, datatype, dest, tag, comm, request, ierror)
11     BIND(C)
12     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
13     INTEGER, INTENT(IN) :: count, dest, tag
14     TYPE(MPI_Datatype), INTENT(IN) :: datatype
15     TYPE(MPI_Comm), INTENT(IN) :: comm
16     TYPE(MPI_Request), INTENT(OUT) :: request
17     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 18
19     MPI_Send(buf, count, datatype, dest, tag, comm, ierror) BIND(C)
20     TYPE(*), DIMENSION(..), INTENT(IN) :: buf
21     INTEGER, INTENT(IN) :: count, dest, tag
22     TYPE(MPI_Datatype), INTENT(IN) :: datatype
23     TYPE(MPI_Comm), INTENT(IN) :: comm
24     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 25
26     MPI_Send_init(buf, count, datatype, dest, tag, comm, request, ierror)
27     BIND(C)
28     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
29     INTEGER, INTENT(IN) :: count, dest, tag
30     TYPE(MPI_Datatype), INTENT(IN) :: datatype
31     TYPE(MPI_Comm), INTENT(IN) :: comm
32     TYPE(MPI_Request), INTENT(OUT) :: request
33     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 34
35     MPI_Sendrecv_replace(buf, count, datatype, dest, sendtag, source, recvtag,
36     comm, status, ierror) BIND(C)
37     TYPE(*), DIMENSION(..) :: buf
38     INTEGER, INTENT(IN) :: count, dest, sendtag, source, recvtag
39     TYPE(MPI_Datatype), INTENT(IN) :: datatype
40     TYPE(MPI_Comm), INTENT(IN) :: comm
41     TYPE(MPI_Status) :: status
42     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 43
44     MPI_Sendrecv(sendbuf, sendcount, sendtype, dest, sendtag, recvbuf,
45     recvcount, recvtype, source, recvtag, comm, status, ierror)
46     BIND(C)
47     TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
48     TYPE(*), DIMENSION(..) :: recvbuf
49     INTEGER, INTENT(IN) :: sendcount, dest, sendtag, recvcount, source,
50     recvtag

```

```

TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Status) :: status
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Ssend(buf, count, datatype, dest, tag, comm, ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN) :: buf
INTEGER, INTENT(IN) :: count, dest, tag
TYPE(MPI_Datatype), INTENT(IN) :: datatype
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Ssend_init(buf, count, datatype, dest, tag, comm, request, ierror)
BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
INTEGER, INTENT(IN) :: count, dest, tag
TYPE(MPI_Datatype), INTENT(IN) :: datatype
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Request), INTENT(OUT) :: request
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Startall(count, array_of_requests, ierror) BIND(C)
INTEGER, INTENT(IN) :: count
TYPE(MPI_Request), INTENT(INOUT) :: array_of_requests(count)
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Start(request, ierror) BIND(C)
TYPE(MPI_Request), INTENT(INOUT) :: request
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Testall(count, array_of_requests, flag, array_of_statuses, ierror)
BIND(C)
INTEGER, INTENT(IN) :: count
TYPE(MPI_Request), INTENT(INOUT) :: array_of_requests(count)
LOGICAL, INTENT(OUT) :: flag
TYPE(MPI_Status) :: array_of_statuses(*)
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Testany(count, array_of_requests, index, flag, status, ierror) BIND(C)
INTEGER, INTENT(IN) :: count
TYPE(MPI_Request), INTENT(INOUT) :: array_of_requests(count)
INTEGER, INTENT(OUT) :: index
LOGICAL, INTENT(OUT) :: flag
TYPE(MPI_Status) :: status
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Test_cancelled(status, flag, ierror) BIND(C)
TYPE(MPI_Status), INTENT(IN) :: status
LOGICAL, INTENT(OUT) :: flag
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

1  MPI_Test(request, flag, status, ierror) BIND(C)
2      TYPE(MPI_Request), INTENT(INOUT) :: request
3      LOGICAL, INTENT(OUT) :: flag
4      TYPE(MPI_Status) :: status
5      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 6
7  MPI_Testsome(incount, array_of_requests, outcount, array_of_indices,
8              array_of_statuses, ierror) BIND(C)
9      INTEGER, INTENT(IN) :: incount
10     TYPE(MPI_Request), INTENT(INOUT) :: array_of_requests(incount)
11     INTEGER, INTENT(OUT) :: outcount, array_of_indices(*)
12     TYPE(MPI_Status) :: array_of_statuses(*)
13     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 14
15 MPI_Waitall(count, array_of_requests, array_of_statuses, ierror) BIND(C)
16     INTEGER, INTENT(IN) :: count
17     TYPE(MPI_Request), INTENT(INOUT) :: array_of_requests(count)
18     TYPE(MPI_Status) :: array_of_statuses(*)
19     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 20
21 MPI_Waitany(count, array_of_requests, index, status, ierror) BIND(C)
22     INTEGER, INTENT(IN) :: count
23     TYPE(MPI_Request), INTENT(INOUT) :: array_of_requests(count)
24     INTEGER, INTENT(OUT) :: index
25     TYPE(MPI_Status) :: status
26     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 27
28 MPI_Wait(request, status, ierror) BIND(C)
29     TYPE(MPI_Request), INTENT(INOUT) :: request
30     TYPE(MPI_Status) :: status
31     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 32
33 MPI_Waitsome(incount, array_of_requests, outcount, array_of_indices,
34             array_of_statuses, ierror) BIND(C)
35     INTEGER, INTENT(IN) :: incount
36     TYPE(MPI_Request), INTENT(INOUT) :: array_of_requests(incount)
37     INTEGER, INTENT(OUT) :: outcount, array_of_indices(*)
38     TYPE(MPI_Status) :: array_of_statuses(*)
39     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

### A.3.2 Datatypes Fortran 2008 Bindings

```

40 MPI_Get_address(location, address, ierror) BIND(C)
41     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: location
42     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: address
43     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 44
45 MPI_Get_elements(status, datatype, count, ierror) BIND(C)
46     TYPE(MPI_Status), INTENT(IN) :: status
47     TYPE(MPI_Datatype), INTENT(IN) :: datatype
48

```



```

INTEGER, INTENT(OUT) :: count
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Pack_external(datarep, inbuf, incount, datatype, outbuf, outsize,
                  position, ierror) BIND(C)
CHARACTER(LEN=*), INTENT(IN) :: datarep
TYPE(*), DIMENSION(..), INTENT(IN) :: inbuf
TYPE(*), DIMENSION(..) :: outbuf
INTEGER, INTENT(IN) :: incount
TYPE(MPI_Datatype), INTENT(IN) :: datatype
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: outsize
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(INOUT) :: position
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Pack_external_size(datarep, incount, datatype, size, ierror) BIND(C)
TYPE(MPI_Datatype), INTENT(IN) :: datatype
INTEGER, INTENT(IN) :: incount
CHARACTER(LEN=*), INTENT(IN) :: datarep
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: size
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Pack(inbuf, incount, datatype, outbuf, outsize, position, comm, ierror)
  BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN) :: inbuf
TYPE(*), DIMENSION(..) :: outbuf
INTEGER, INTENT(IN) :: incount, outsize
TYPE(MPI_Datatype), INTENT(IN) :: datatype
INTEGER, INTENT(INOUT) :: position
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Pack_size(incount, datatype, comm, size, ierror) BIND(C)
INTEGER, INTENT(IN) :: incount
TYPE(MPI_Datatype), INTENT(IN) :: datatype
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, INTENT(OUT) :: size
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Type_commit(datatype, ierror) BIND(C)
TYPE(MPI_Datatype), INTENT(INOUT) :: datatype
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Type_contiguous(count, oldtype, newtype, ierror) BIND(C)
INTEGER, INTENT(IN) :: count
TYPE(MPI_Datatype), INTENT(IN) :: oldtype
TYPE(MPI_Datatype), INTENT(OUT) :: newtype
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Type_create_darray(size, rank, ndims, array_of_gsizes,
                      array_of_distribs, array_of_dargs, array_of_psize, order,
                      oldtype, newtype, ierror) BIND(C)

```

```

1      INTEGER, INTENT(IN) :: size, rank, ndims, array_of_gsizes(ndims),
2      array_of_distrib(ndims), array_of_dargs(ndims),
3      array_of_psize(ndims), order
4      TYPE(MPI_Datatype), INTENT(IN) :: oldtype
5      TYPE(MPI_Datatype), INTENT(OUT) :: newtype
6      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 7
8      MPI_Type_create_hindexed(count, array_of_blocklengths,
9      array_of_displacements, oldtype, newtype, ierror) BIND(C)
10     INTEGER, INTENT(IN) :: count, array_of_blocklengths(count)
11     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) ::
12     array_of_displacements(count)
13     TYPE(MPI_Datatype), INTENT(IN) :: oldtype
14     TYPE(MPI_Datatype), INTENT(OUT) :: newtype
15     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 16
17     MPI_Type_create_hvector(count, blocklength, stride, oldtype, newtype,
18     ierror) BIND(C)
19     INTEGER, INTENT(IN) :: count, blocklength
20     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: stride
21     TYPE(MPI_Datatype), INTENT(IN) :: oldtype
22     TYPE(MPI_Datatype), INTENT(OUT) :: newtype
23     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 24
25     MPI_Type_create_indexed_block(count, blocklength, array_of_displacements,
26     oldtype, newtype, ierror) BIND(C)
27     INTEGER, INTENT(IN) :: count, blocklength,
28     array_of_displacements(count)
29     TYPE(MPI_Datatype), INTENT(IN) :: oldtype
30     TYPE(MPI_Datatype), INTENT(OUT) :: newtype
31     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 32
33     MPI_Type_create_resized(oldtype, lb, extent, newtype, ierror) BIND(C)
34     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: lb, extent
35     TYPE(MPI_Datatype), INTENT(IN) :: oldtype
36     TYPE(MPI_Datatype), INTENT(OUT) :: newtype
37     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 38
39     MPI_Type_create_struct(count, array_of_blocklengths,
40     array_of_displacements, array_of_types, newtype, ierror)
41     BIND(C)
42     INTEGER, INTENT(IN) :: count, array_of_blocklengths(count)
43     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) ::
44     array_of_displacements(count)
45     TYPE(MPI_Datatype), INTENT(IN) :: array_of_types(count)
46     TYPE(MPI_Datatype), INTENT(OUT) :: newtype
47     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 48
49     MPI_Type_create_subarray(ndims, array_of_sizes, array_of_subsizes,
50     array_of_starts, order, oldtype, newtype, ierror) BIND(C)

```

```

INTEGER, INTENT(IN) :: ndims, array_of_sizes(ndims),
array_of_subsizes(ndims), array_of_starts(ndims), order
TYPE(MPI_Datatype), INTENT(IN) :: oldtype
TYPE(MPI_Datatype), INTENT(OUT) :: newtype
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Type_dup(oldtype, newtype, ierror) BIND(C)
TYPE(MPI_Datatype), INTENT(IN) :: oldtype
TYPE(MPI_Datatype), INTENT(OUT) :: newtype
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Type_free(datatype, ierror) BIND(C)
TYPE(MPI_Datatype), INTENT(INOUT) :: datatype
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Type_get_contents(datatype, max_integers, max_addresses, max_datatypes,
array_of_integers, array_of_addresses, array_of_datatypes,
ierror) BIND(C)
TYPE(MPI_Datatype), INTENT(IN) :: datatype
INTEGER, INTENT(IN) :: max_integers, max_addresses, max_datatypes
INTEGER, INTENT(OUT) :: array_of_integers(max_integers)
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) ::
array_of_addresses(max_addresses)
TYPE(MPI_Datatype), INTENT(OUT) :: array_of_datatypes(max_datatypes)
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Type_get_envelope(datatype, num_integers, num_addresses, num_datatypes,
combiner, ierror) BIND(C)
TYPE(MPI_Datatype), INTENT(IN) :: datatype
INTEGER, INTENT(OUT) :: num_integers, num_addresses, num_datatypes,
combiner
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Type_get_extent(datatype, lb, extent, ierror) BIND(C)
TYPE(MPI_Datatype), INTENT(IN) :: datatype
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: lb, extent
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Type_get_true_extent(datatype, true_lb, true_extent, ierror) BIND(C)
TYPE(MPI_Datatype), INTENT(IN) :: datatype
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: true_lb, true_extent
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Type_indexed(count, array_of_blocklengths, array_of_displacements,
oldtype, newtype, ierror) BIND(C)
INTEGER, INTENT(IN) :: count, array_of_blocklengths(count),
array_of_displacements(count)
TYPE(MPI_Datatype), INTENT(IN) :: oldtype
TYPE(MPI_Datatype), INTENT(OUT) :: newtype
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

1  MPI_Type_size(datatype, size, ierror) BIND(C)
2      TYPE(MPI_Datatype), INTENT(IN) :: datatype
3      INTEGER, INTENT(OUT) :: size
4      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 5
6  MPI_Type_vector(count, blocklength, stride, oldtype, newtype, ierror)
7      BIND(C)
8      INTEGER, INTENT(IN) :: count, blocklength, stride
9      TYPE(MPI_Datatype), INTENT(IN) :: oldtype
10     TYPE(MPI_Datatype), INTENT(OUT) :: newtype
11     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 12
13 MPI_Unpack_external(datarep, inbuf, insize, position, outbuf, outcount,
14     datatype, ierror) BIND(C)
15     CHARACTER(LEN=*), INTENT(IN) :: datarep
16     TYPE(*), DIMENSION(..), INTENT(IN) :: inbuf
17     TYPE(*), DIMENSION(..) :: outbuf
18     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: insize
19     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(INOUT) :: position
20     INTEGER, INTENT(IN) :: outcount
21     TYPE(MPI_Datatype), INTENT(IN) :: datatype
22     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 23
24 MPI_Unpack(inbuf, insize, position, outbuf, outcount, datatype, comm,
25     ierror) BIND(C)
26     TYPE(*), DIMENSION(..), INTENT(IN) :: inbuf
27     TYPE(*), DIMENSION(..) :: outbuf
28     INTEGER, INTENT(IN) :: insize, outcount
29     INTEGER, INTENT(INOUT) :: position
30     TYPE(MPI_Datatype), INTENT(IN) :: datatype
31     TYPE(MPI_Comm), INTENT(IN) :: comm
32     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
33

```

### A.3.3 Collective Communication Fortran 2008 Bindings

```

34 MPI_Allgather(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype,
35     comm, ierror) BIND(C)
36     TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
37     TYPE(*), DIMENSION(..) :: recvbuf
38     INTEGER, INTENT(IN) :: sendcount, recvcount
39     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
40     TYPE(MPI_Comm), INTENT(IN) :: comm
41     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 42
43 MPI_Allgatherv(sendbuf, sendcount, sendtype, recvbuf, recvcounts, displs,
44     recvtype, comm, ierror) BIND(C)
45     TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
46     TYPE(*), DIMENSION(..) :: recvbuf
47     INTEGER, INTENT(IN) :: sendcount, recvcounts(*), displs(*)
48

```

```

TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Allreduce(sendbuf, recvbuf, count, datatype, op, comm, ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
TYPE(*), DIMENSION(..) :: recvbuf
INTEGER, INTENT(IN) :: count
TYPE(MPI_Datatype), INTENT(IN) :: datatype
TYPE(MPI_Op), INTENT(IN) :: op
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Alltoall(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype,
             comm, ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
TYPE(*), DIMENSION(..) :: recvbuf
INTEGER, INTENT(IN) :: sendcount, recvcount
TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Alltoallv(sendbuf, sendcounts, sdispls, sendtype, recvbuf, recvcounts,
              rdispls, recvtype, comm, ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
TYPE(*), DIMENSION(..) :: recvbuf
INTEGER, INTENT(IN) :: sendcounts(*), sdispls(*), recvcounts(*),
rdispls(*)
TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Alltoallw(sendbuf, sendcounts, sdispls, sendtypes, recvbuf, recvcounts,
              rdispls, recvtypes, comm, ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
TYPE(*), DIMENSION(..) :: recvbuf
INTEGER, INTENT(IN) :: sendcounts(*), sdispls(*), recvcounts(*),
rdispls(*)
TYPE(MPI_Datatype), INTENT(IN) :: sendtypes(*)
TYPE(MPI_Datatype), INTENT(IN) :: recvtypes(*)
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Barrier(comm, ierror) BIND(C)
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Bcast(buffer, count, datatype, root, comm, ierror) BIND(C)
TYPE(*), DIMENSION(..) :: buffer
INTEGER, INTENT(IN) :: count, root

```

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41 ticket-248T.  
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44 ticket-248T.  
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```

1      TYPE(MPI_Datatype), INTENT(IN) :: datatype
2      TYPE(MPI_Comm), INTENT(IN) :: comm
3      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 4
5      MPI_Exscan(sendbuf, recvbuf, count, datatype, op, comm, ierror) BIND(C)
6      TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
7      TYPE(*), DIMENSION(..) :: recvbuf
8      INTEGER, INTENT(IN) :: count
9      TYPE(MPI_Datatype), INTENT(IN) :: datatype
10     TYPE(MPI_Op), INTENT(IN) :: op
11     TYPE(MPI_Comm), INTENT(IN) :: comm
ticket-248T. 12     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
13
14     MPI_Gather(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvttype,
15               root, comm, ierror) BIND(C)
16     TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
17     TYPE(*), DIMENSION(..) :: recvbuf
18     INTEGER, INTENT(IN) :: sendcount, recvcount, root
19     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvttype
20     TYPE(MPI_Comm), INTENT(IN) :: comm
ticket-248T. 21     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
22
23     MPI_Gatherv(sendbuf, sendcount, sendtype, recvbuf, recvcounts, displs,
24               recvttype, root, comm, ierror) BIND(C)
25     TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
26     TYPE(*), DIMENSION(..) :: recvbuf
27     INTEGER, INTENT(IN) :: sendcount, recvcounts(*), displs(*), root
28     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvttype
29     TYPE(MPI_Comm), INTENT(IN) :: comm
ticket-248T. 30     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
31
32     MPI_Iallgather(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvttype,
33                   comm, request, ierror) BIND(C)
34     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
35     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
36     INTEGER, INTENT(IN) :: sendcount, recvcount
37     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvttype
38     TYPE(MPI_Comm), INTENT(IN) :: comm
39     TYPE(MPI_Request), INTENT(OUT) :: request
ticket-248T. 40     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
41
42     MPI_Iallgatherv(sendbuf, sendcount, sendtype, recvbuf, recvcounts, displs,
43                   recvttype, comm, request, ierror) BIND(C)
44     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
45     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
46     INTEGER, INTENT(IN) :: sendcount
47     INTEGER, INTENT(IN), ASYNCHRONOUS :: recvcounts(*), displs(*)
48     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvttype
49     TYPE(MPI_Comm), INTENT(IN) :: comm
50     TYPE(MPI_Request), INTENT(OUT) :: request

```

```

    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Iallreduce(sendbuf, recvbuf, count, datatype, op, comm, request,
               ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Op), INTENT(IN) :: op
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Ialltoall(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvttype,
               comm, request, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN) :: sendcount, recvcount
    TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvttype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Ialltoallv(sendbuf, sendcounts, sdispls, sendtype, recvbuf, recvcounts,
                rdispls, recvttype, comm, request, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN), ASYNCHRONOUS :: sendcounts(*), sdispls(*),
    recvcounts(*), rdispls(*)
    TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvttype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Ialltoallw(sendbuf, sendcounts, sdispls, sendtypes, recvbuf,
                recvcounts, rdispls, recvtypes, comm, request, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN), ASYNCHRONOUS :: sendcounts(*), sdispls(*),
    recvcounts(*), rdispls(*)
    TYPE(MPI_Datatype), INTENT(IN), ASYNCHRONOUS :: sendtypes(*),
    recvtypes(*)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Ibarrier(comm, request, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

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44 ticket-248T.  
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ticket-248T. 2 MPI_Ibcast(buffer, count, datatype, root, comm, request, ierror) BIND(C)
3     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buffer
4     INTEGER, INTENT(IN) :: count, root
5     TYPE(MPI_Datatype), INTENT(IN) :: datatype
6     TYPE(MPI_Comm), INTENT(IN) :: comm
7     TYPE(MPI_Request), INTENT(OUT) :: request
8     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 9
10 MPI_Iexscan(sendbuf, recvbuf, count, datatype, op, comm, request, ierror)
11     BIND(C)
12     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
13     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
14     INTEGER, INTENT(IN) :: count
15     TYPE(MPI_Datatype), INTENT(IN) :: datatype
16     TYPE(MPI_Op), INTENT(IN) :: op
17     TYPE(MPI_Comm), INTENT(IN) :: comm
18     TYPE(MPI_Request), INTENT(OUT) :: request
19     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 20
21 MPI_Igather(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype,
22     root, comm, request, ierror) BIND(C)
23     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
24     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
25     INTEGER, INTENT(IN) :: sendcount, recvcount, root
26     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
27     TYPE(MPI_Comm), INTENT(IN) :: comm
28     TYPE(MPI_Request), INTENT(OUT) :: request
29     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 30
31 MPI_Igatherv(sendbuf, sendcount, sendtype, recvbuf, recvcounts, displs,
32     recvtype, root, comm, request, ierror) BIND(C)
33     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
34     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
35     INTEGER, INTENT(IN) :: sendcount, root
36     INTEGER, INTENT(IN), ASYNCHRONOUS :: recvcounts(*), displs(*)
37     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
38     TYPE(MPI_Comm), INTENT(IN) :: comm
39     TYPE(MPI_Request), INTENT(OUT) :: request
40     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 41
42 MPI_Ireduce_scatter_block(sendbuf, recvbuf, recvcount, datatype, op, comm,
43     request, ierror) BIND(C)
44     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
45     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
46     INTEGER, INTENT(IN) :: recvcount
47     TYPE(MPI_Datatype), INTENT(IN) :: datatype
48     TYPE(MPI_Op), INTENT(IN) :: op
49     TYPE(MPI_Comm), INTENT(IN) :: comm
50     TYPE(MPI_Request), INTENT(OUT) :: request

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    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Ireduce_scatter(sendbuf, recvbuf, recvcnt, datatype, op, comm,
    request, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN), ASYNCHRONOUS :: recvcnt(*)
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Op), INTENT(IN) :: op
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Ireduce(sendbuf, recvbuf, count, datatype, op, root, comm, request,
    ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN) :: count, root
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Op), INTENT(IN) :: op
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Iscan(sendbuf, recvbuf, count, datatype, op, comm, request, ierror)
    BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Op), INTENT(IN) :: op
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Iscatter(sendbuf, sendcount, sendtype, recvbuf, recvcnt, recvtype,
    root, comm, request, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN) :: sendcount, recvcnt, root
    TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Iscatterv(sendbuf, sendcounts, displs, sendtype, recvbuf, recvcnt,
    recvtype, root, comm, request, ierror) BIND(C)
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: sendbuf
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: recvbuf
    INTEGER, INTENT(IN), ASYNCHRONOUS :: sendcounts(*), displs(*)

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43 ticket-248T.  
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1      INTEGER, INTENT(IN) :: recvcount, root
2      TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
3      TYPE(MPI_Comm), INTENT(IN) :: comm
4      TYPE(MPI_Request), INTENT(OUT) :: request
5      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 6
7      MPI_Op_commutative(op, commute, ierror) BIND(C)
8      TYPE(MPI_Op), INTENT(IN) :: op
9      LOGICAL, INTENT(OUT) :: commute
ticket-248T. 10
11     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
12
13     MPI_Op_create(user_fn, commute, op, ierror) BIND(C)
14     PROCEDURE(MPI_User_function) :: user_fn
15     LOGICAL, INTENT(IN) :: commute
16     TYPE(MPI_Op), INTENT(OUT) :: op
17     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 18
19     MPI_Op_free(op, ierror) BIND(C)
20     TYPE(MPI_Op), INTENT(INOUT) :: op
21     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 22
23     MPI_Reduce_local(inbuf, inoutbuf, count, datatype, op, ierror) BIND(C)
24     TYPE(*), DIMENSION(..), INTENT(IN) :: inbuf
25     TYPE(*), DIMENSION(..) :: inoutbuf
26     INTEGER, INTENT(IN) :: count
27     TYPE(MPI_Datatype), INTENT(IN) :: datatype
28     TYPE(MPI_Op), INTENT(IN) :: op
29     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 30
31     MPI_Reduce_scatter_block(sendbuf, recvbuf, recvcount, datatype, op, comm,
32                               ierror) BIND(C)
33     TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
34     TYPE(*), DIMENSION(..) :: recvbuf
35     INTEGER, INTENT(IN) :: recvcount
36     TYPE(MPI_Datatype), INTENT(IN) :: datatype
37     TYPE(MPI_Op), INTENT(IN) :: op
38     TYPE(MPI_Comm), INTENT(IN) :: comm
39     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 40
41     MPI_Reduce_scatter(sendbuf, recvbuf, recvcounts, datatype, op, comm,
42                               ierror) BIND(C)
43     TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
44     TYPE(*), DIMENSION(..) :: recvbuf
45     INTEGER, INTENT(IN) :: recvcounts(*)
46     TYPE(MPI_Datatype), INTENT(IN) :: datatype
47     TYPE(MPI_Op), INTENT(IN) :: op
48     TYPE(MPI_Comm), INTENT(IN) :: comm
49     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 50
51     MPI_Reduce(sendbuf, recvbuf, count, datatype, op, root, comm, ierror)
52     BIND(C)

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TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
TYPE(*), DIMENSION(..) :: recvbuf
INTEGER, INTENT(IN) :: count, root
TYPE(MPI_Datatype), INTENT(IN) :: datatype
TYPE(MPI_Op), INTENT(IN) :: op
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Scan(sendbuf, recvbuf, count, datatype, op, comm, ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
TYPE(*), DIMENSION(..) :: recvbuf
INTEGER, INTENT(IN) :: count
TYPE(MPI_Datatype), INTENT(IN) :: datatype
TYPE(MPI_Op), INTENT(IN) :: op
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Scatter(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype,
            root, comm, ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
TYPE(*), DIMENSION(..) :: recvbuf
INTEGER, INTENT(IN) :: sendcount, recvcount, root
TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Scatterv(sendbuf, sendcounts, displs, sendtype, recvbuf, recvcount,
            recvtype, root, comm, ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
TYPE(*), DIMENSION(..) :: recvbuf
INTEGER, INTENT(IN) :: sendcounts(*), displs(*), recvcount, root
TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

A.3.4 Groups, Contexts, Communicators, and Caching Fortran 2008 Bindings

MPI_Comm_compare(comm1, comm2, result, ierror) BIND(C)
TYPE(MPI_Comm), INTENT(IN) :: comm1, comm2
INTEGER, INTENT(OUT) :: result
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Comm_create(comm, group, newcomm, ierror) BIND(C)
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Group), INTENT(IN) :: group
TYPE(MPI_Comm), INTENT(OUT) :: newcomm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Comm_create_keyval(comm_copy_attr_fn, comm_delete_attr_fn, comm_keyval,
                    extra_state, ierror) BIND(C)

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ticket-248T.

ticket-248T.

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ticket-248T.

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1      PROCEDURE(MPI_Comm_copy_attr_function) :: comm_copy_attr_fn
2      PROCEDURE(MPI_Comm_delete_attr_function) :: comm_delete_attr_fn
3      INTEGER, INTENT(OUT) :: comm_keyval
4      INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state
5      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 6
7      MPI_Comm_delete_attr(comm, comm_keyval, ierror) BIND(C)
8      TYPE(MPI_Comm), INTENT(IN) :: comm
9      INTEGER, INTENT(IN) :: comm_keyval
10     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 11
12     MPI_Comm_dup(comm, newcomm, ierror) BIND(C)
13     TYPE(MPI_Comm), INTENT(IN) :: comm
14     TYPE(MPI_Comm), INTENT(OUT) :: newcomm
15     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 16
17     MPI_COMM_DUP_FN(oldcomm, comm_keyval, extra_state, attribute_val_in,
18         attribute_val_out, flag, ierror) BIND(C)
19     TYPE(MPI_Comm), INTENT(IN) :: oldcomm
20     INTEGER, INTENT(IN) :: comm_keyval
21     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state,
22     attribute_val_in
23     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val_out
24     LOGICAL, INTENT(OUT) :: flag
25     INTEGER, INTENT(OUT) :: ierror
ticket-248T. 26
27     MPI_Comm_free(comm, ierror) BIND(C)
28     TYPE(MPI_Comm), INTENT(INOUT) :: comm
29     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 30
31     MPI_Comm_free_keyval(comm_keyval, ierror) BIND(C)
32     INTEGER, INTENT(INOUT) :: comm_keyval
33     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 34
35     MPI_Comm_get_attr(comm, comm_keyval, attribute_val, flag, ierror) BIND(C)
36     TYPE(MPI_Comm), INTENT(IN) :: comm
37     INTEGER, INTENT(IN) :: comm_keyval
38     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val
39     LOGICAL, INTENT(OUT) :: flag
40     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 41
42     MPI_Comm_get_name(comm, comm_name, resultlen, ierror) BIND(C)
43     TYPE(MPI_Comm), INTENT(IN) :: comm
44     CHARACTER(LEN=MPI_MAX_OBJECT_NAME), INTENT(OUT) :: comm_name
45     INTEGER, INTENT(OUT) :: resultlen
46     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 47
48     MPI_Comm_group(comm, group, ierror) BIND(C)
49     TYPE(MPI_Comm), INTENT(IN) :: comm
50     TYPE(MPI_Group), INTENT(OUT) :: group
51     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

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MPI_COMM_NULL_COPY_FN(oldcomm, comm_keyval, extra_state, attribute_val_in,
    attribute_val_out, flag, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: oldcomm
    INTEGER, INTENT(IN) :: comm_keyval
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state,
    attribute_val_in
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val_out
    LOGICAL, INTENT(OUT) :: flag
    INTEGER, INTENT(OUT) :: ierror
MPI_COMM_NULL_DELETE_FN(comm, comm_keyval, attribute_val, extra_state,
    ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, INTENT(IN) :: comm_keyval
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: attribute_val,
    extra_state
    INTEGER, INTENT(OUT) :: ierror
MPI_Comm_rank(comm, rank, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, INTENT(OUT) :: rank
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Comm_remote_group(comm, group, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Group), INTENT(OUT) :: group
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Comm_remote_size(comm, size, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, INTENT(OUT) :: size
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Comm_set_attr(comm, comm_keyval, attribute_val, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, INTENT(IN) :: comm_keyval
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: attribute_val
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Comm_set_name(comm, comm_name, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    CHARACTER(LEN=*), INTENT(IN) :: comm_name
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Comm_size(comm, size, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, INTENT(OUT) :: size
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Comm_split(comm, color, key, newcomm, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm

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1      INTEGER, INTENT(IN) :: color, key
2      TYPE(MPI_Comm), INTENT(OUT) :: newcomm
3      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 4
5      MPI_Comm_test_inter(comm, flag, ierror) BIND(C)
6      TYPE(MPI_Comm), INTENT(IN) :: comm
7      LOGICAL, INTENT(OUT) :: flag
8      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 9
10     MPI_Group_compare(group1, group2, result, ierror) BIND(C)
11     TYPE(MPI_Group), INTENT(IN) :: group1, group2
12     INTEGER, INTENT(OUT) :: result
13     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 14
15     MPI_Group_difference(group1, group2, newgroup, ierror) BIND(C)
16     TYPE(MPI_Group), INTENT(IN) :: group1, group2
17     TYPE(MPI_Group), INTENT(OUT) :: newgroup
18     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 19
20     MPI_Group_excl(group, n, ranks, newgroup, ierror) BIND(C)
21     TYPE(MPI_Group), INTENT(IN) :: group
22     INTEGER, INTENT(IN) :: n, ranks(n)
23     TYPE(MPI_Group), INTENT(OUT) :: newgroup
24     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 25
26     MPI_Group_free(group, ierror) BIND(C)
27     TYPE(MPI_Group), INTENT(INOUT) :: group
28     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 29
30     MPI_Group_incl(group, n, ranks, newgroup, ierror) BIND(C)
31     TYPE(MPI_Group), INTENT(IN) :: group
32     INTEGER, INTENT(IN) :: n, ranks(n)
33     TYPE(MPI_Group), INTENT(OUT) :: newgroup
34     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 35
36     MPI_Group_intersection(group1, group2, newgroup, ierror) BIND(C)
37     TYPE(MPI_Group), INTENT(IN) :: group1, group2
38     TYPE(MPI_Group), INTENT(OUT) :: newgroup
39     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 40
41     MPI_Group_range_excl(group, n, ranges, newgroup, ierror) BIND(C)
42     TYPE(MPI_Group), INTENT(IN) :: group
43     INTEGER, INTENT(IN) :: n, ranges(3,n)
44     TYPE(MPI_Group), INTENT(OUT) :: newgroup
45     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 46
47     MPI_Group_range_incl(group, n, ranges, newgroup, ierror) BIND(C)
48     TYPE(MPI_Group), INTENT(IN) :: group
49     INTEGER, INTENT(IN) :: n, ranges(3,n)
50     TYPE(MPI_Group), INTENT(OUT) :: newgroup
51     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

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MPI_Group_rank(group, rank, ierror) BIND(C)
  TYPE(MPI_Group), INTENT(IN) :: group
  INTEGER, INTENT(OUT) :: rank
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Group_size(group, size, ierror) BIND(C)
  TYPE(MPI_Group), INTENT(IN) :: group
  INTEGER, INTENT(OUT) :: size
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Group_translate_ranks(group1, n, ranks1, group2, ranks2, ierror)
  BIND(C)
  TYPE(MPI_Group), INTENT(IN) :: group1, group2
  INTEGER, INTENT(IN) :: n, ranks1(n)
  INTEGER, INTENT(OUT) :: ranks2(n)
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Group_union(group1, group2, newgroup, ierror) BIND(C)
  TYPE(MPI_Group), INTENT(IN) :: group1, group2
  TYPE(MPI_Group), INTENT(OUT) :: newgroup
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Intercomm_create(local_comm, local_leader, peer_comm, remote_leader,
  tag, newintercomm, ierror) BIND(C)
  TYPE(MPI_Comm), INTENT(IN) :: local_comm, peer_comm
  INTEGER, INTENT(IN) :: local_leader, remote_leader, tag
  TYPE(MPI_Comm), INTENT(OUT) :: newintercomm
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Intercomm_merge(intercomm, high, newintracomm, ierror) BIND(C)
  TYPE(MPI_Comm), INTENT(IN) :: intercomm
  LOGICAL, INTENT(IN) :: high
  TYPE(MPI_Comm), INTENT(OUT) :: newintracomm
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Type_create_keyval(type_copy_attr_fn, type_delete_attr_fn, type_keyval,
  extra_state, ierror) BIND(C)
  PROCEDURE(MPI_Type_copy_attr_function) :: type_copy_attr_fn
  PROCEDURE(MPI_Type_delete_attr_function) :: type_delete_attr_fn
  INTEGER, INTENT(OUT) :: type_keyval
  INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Type_delete_attr(datatype, type_keyval, ierror) BIND(C)
  TYPE(MPI_Datatype), INTENT(IN) :: datatype
  INTEGER, INTENT(IN) :: type_keyval
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Type_Dup_Fn(oldtype, type_keyval, extra_state, attribute_val_in,
  attribute_val_out, flag, ierror) BIND(C)
  TYPE(MPI_Datatype), INTENT(IN) :: oldtype

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1      INTEGER, INTENT(IN) :: type_keyval
2      INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state,
3      attribute_val_in
4      INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val_out
5      LOGICAL, INTENT(OUT) :: flag
6      INTEGER, INTENT(OUT) :: ierror
ticket-248T. 7
8      MPI_Type_free_keyval(type_keyval, ierror) BIND(C)
9      INTEGER, INTENT(INOUT) :: type_keyval
10     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 11
12     MPI_Type_get_attr(datatype, type_keyval, attribute_val, flag, ierror)
13     BIND(C)
14     TYPE(MPI_Datatype), INTENT(IN) :: datatype
15     INTEGER, INTENT(IN) :: type_keyval
16     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val
17     LOGICAL, INTENT(OUT) :: flag
18     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 19
20     MPI_Type_get_name(datatype, type_name, resultlen, ierror) BIND(C)
21     TYPE(MPI_Datatype), INTENT(IN) :: datatype
22     CHARACTER(LEN=MPI_MAX_OBJECT_NAME), INTENT(OUT) :: type_name
23     INTEGER, INTENT(OUT) :: resultlen
24     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 25
26     MPI_TYPE_NULL_COPY_FN(oldtype, type_keyval, extra_state, attribute_val_in,
27     attribute_val_out, flag, ierror) BIND(C)
28     TYPE(MPI_Datatype), INTENT(IN) :: oldtype
29     INTEGER, INTENT(IN) :: type_keyval
30     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state,
31     attribute_val_in
32     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val_out
33     LOGICAL, INTENT(OUT) :: flag
34     INTEGER, INTENT(OUT) :: ierror
ticket-248T. 35
36     MPI_TYPE_NULL_DELETE_FN(datatype, type_keyval, attribute_val, extra_state,
37     ierror) BIND(C)
38     TYPE(MPI_Datatype), INTENT(IN) :: datatype
39     INTEGER, INTENT(IN) :: type_keyval
40     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: attribute_val,
41     extra_state
42     INTEGER, INTENT(OUT) :: ierror
ticket-248T. 43
44     MPI_Type_set_attr(datatype, type_keyval, attribute_val, ierror) BIND(C)
45     TYPE(MPI_Datatype), INTENT(IN) :: datatype
46     INTEGER, INTENT(IN) :: type_keyval
47     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: attribute_val
48     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 49
50     MPI_Type_set_name(datatype, type_name, ierror) BIND(C)
51     TYPE(MPI_Datatype), INTENT(IN) :: datatype

```



```

    CHARACTER(LEN=*), INTENT(IN) :: type_name
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Win_create_keyval(win_copy_attr_fn, win_delete_attr_fn, win_keyval,
    extra_state, ierror) BIND(C)
    PROCEDURE(MPI_Win_copy_attr_function) :: win_copy_attr_fn
    PROCEDURE(MPI_Win_delete_attr_function) :: win_delete_attr_fn
    INTEGER, INTENT(OUT) :: win_keyval
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Win_delete_attr(win, win_keyval, ierror) BIND(C)
    TYPE(MPI_Win), INTENT(IN) :: win
    INTEGER, INTENT(IN) :: win_keyval
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_WIN_DUP_FN(oldwin, win_keyval, extra_state, attribute_val_in,
    attribute_val_out, flag, ierror) BIND(C)
    INTEGER, INTENT(IN) :: oldwin, win_keyval
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state,
    attribute_val_in
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val_out
    LOGICAL, INTENT(OUT) :: flag
    INTEGER, INTENT(OUT) :: ierror

MPI_Win_free_keyval(win_keyval, ierror) BIND(C)
    INTEGER, INTENT(INOUT) :: win_keyval
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Win_get_attr(win, win_keyval, attribute_val, flag, ierror) BIND(C)
    TYPE(MPI_Win), INTENT(IN) :: win
    INTEGER, INTENT(IN) :: win_keyval
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val
    LOGICAL, INTENT(OUT) :: flag
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Win_get_name(win, win_name, resultlen, ierror) BIND(C)
    TYPE(MPI_Win), INTENT(IN) :: win
    CHARACTER(LEN=MPI_MAX_OBJECT_NAME), INTENT(OUT) :: win_name
    INTEGER, INTENT(OUT) :: resultlen
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_WIN_NULL_COPY_FN(oldwin, win_keyval, extra_state, attribute_val_in,
    attribute_val_out, flag, ierror) BIND(C)
    INTEGER, INTENT(IN) :: oldwin, win_keyval
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state,
    attribute_val_in
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: attribute_val_out
    LOGICAL, INTENT(OUT) :: flag
    INTEGER, INTENT(OUT) :: ierror

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23 ticket-248T.  
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27 ticket-248T.  
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33 ticket-248T.  
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39 ticket-248T.  
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46  
47 ticket-248T.  
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```

1  MPI_WIN_NULL_DELETE_FN(win, win_keyval, attribute_val, extra_state, ierror)
2      BIND(C)
3      TYPE(MPI_Win), INTENT(IN) :: win
4      INTEGER, INTENT(IN) :: win_keyval
5      INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: attribute_val,
6      extra_state
7      INTEGER, INTENT(OUT) :: ierror

```

```

8  MPI_Win_set_attr(win, win_keyval, attribute_val, ierror) BIND(C)
9      TYPE(MPI_Win), INTENT(IN) :: win
10     INTEGER, INTENT(IN) :: win_keyval
11     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: attribute_val
12     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

13 MPI_Win_set_name(win, win_name, ierror) BIND(C)
14     TYPE(MPI_Win), INTENT(IN) :: win
15     CHARACTER(LEN=*), INTENT(IN) :: win_name
16     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

### A.3.5 Process Topologies Fortran 2008 Bindings

```

21 MPI_Cart_coords(comm, rank, maxdims, coords, ierror) BIND(C)
22     TYPE(MPI_Comm), INTENT(IN) :: comm
23     INTEGER, INTENT(IN) :: rank, maxdims
24     INTEGER, INTENT(OUT) :: coords(maxdims)
25     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

26 MPI_Cart_create(comm_old, ndims, dims, periods, reorder, comm_cart, ierror)
27     BIND(C)
28     TYPE(MPI_Comm), INTENT(IN) :: comm_old
29     INTEGER, INTENT(IN) :: ndims, dims(ndims)
30     LOGICAL, INTENT(IN) :: periods(ndims), reorder
31     TYPE(MPI_Comm), INTENT(OUT) :: comm_cart
32     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

33 MPI_Cartdim_get(comm, ndims, ierror) BIND(C)
34     TYPE(MPI_Comm), INTENT(IN) :: comm
35     INTEGER, INTENT(OUT) :: ndims
36     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

37 MPI_Cart_get(comm, maxdims, dims, periods, coords, ierror) BIND(C)
38     TYPE(MPI_Comm), INTENT(IN) :: comm
39     INTEGER, INTENT(IN) :: maxdims
40     INTEGER, INTENT(OUT) :: dims(maxdims), coords(maxdims)
41     LOGICAL, INTENT(OUT) :: periods(maxdims)
42     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

43 MPI_Cart_map(comm, ndims, dims, periods, newrank, ierror) BIND(C)
44     TYPE(MPI_Comm), INTENT(IN) :: comm
45     INTEGER, INTENT(IN) :: ndims, dims(ndims)

```

```

LOGICAL, INTENT(IN) :: periods(ndims)
INTEGER, INTENT(OUT) :: newrank
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Cart_rank(comm, coords, rank, ierror) BIND(C)
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, INTENT(IN) :: coords(*)
INTEGER, INTENT(OUT) :: rank
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Cart_shift(comm, direction, disp, rank_source, rank_dest, ierror)
    BIND(C)
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, INTENT(IN) :: direction, disp
INTEGER, INTENT(OUT) :: rank_source, rank_dest
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Cart_sub(comm, remain_dims, newcomm, ierror) BIND(C)
TYPE(MPI_Comm), INTENT(IN) :: comm
LOGICAL, INTENT(IN) :: remain_dims(*)
TYPE(MPI_Comm), INTENT(OUT) :: newcomm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Dims_create(nnodes, ndims, dims, ierror) BIND(C)
INTEGER, INTENT(IN) :: nnodes, ndims
INTEGER, INTENT(INOUT) :: dims(ndims)
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Dist_graph_create_adjacent(comm_old, indegree, sources, sourceweights,
    outdegree, destinations, destweights, info, reorder,
    comm_dist_graph, ierror) BIND(C)
TYPE(MPI_Comm), INTENT(IN) :: comm_old
INTEGER, INTENT(IN) :: indegree, sources(indegree), outdegree,
    destinations(outdegree)
INTEGER, INTENT(IN) :: sourceweights(*), destweights(*)
TYPE(MPI_Info), INTENT(IN) :: info
LOGICAL, INTENT(IN) :: reorder
TYPE(MPI_Comm), INTENT(OUT) :: comm_dist_graph
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Dist_graph_create(comm_old, n, sources, degrees, destinations, weights,
    info, reorder, comm_dist_graph, ierror) BIND(C)
TYPE(MPI_Comm), INTENT(IN) :: comm_old
INTEGER, INTENT(IN) :: n, sources(n), degrees(n), destinations(*)
INTEGER, INTENT(IN) :: weights(*)
TYPE(MPI_Info), INTENT(IN) :: info
LOGICAL, INTENT(IN) :: reorder
TYPE(MPI_Comm), INTENT(OUT) :: comm_dist_graph
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Dist_graph_neighbors(comm, maxindegree, sources, sourceweights,

```

```

1          maxoutdegree, destinations, destweights, ierror) BIND(C)
2      TYPE(MPI_Comm), INTENT(IN) :: comm
3      INTEGER, INTENT(IN) :: maxindegree, maxoutdegree
4      INTEGER, INTENT(OUT) :: sources(maxindegree),
5      destinations(maxoutdegree)
6      INTEGER :: sourceweights(*), destweights(*)
7      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 8
9      MPI_Dist_graph_neighbors_count(comm, indegree, outdegree, weighted, ierror)
10         BIND(C)
11      TYPE(MPI_Comm), INTENT(IN) :: comm
12      INTEGER, INTENT(OUT) :: indegree, outdegree
13      LOGICAL, INTENT(OUT) :: weighted
ticket-248T. 14
15      MPI_Graph_create(comm_old, nnodes, index, edges, reorder, comm_graph,
16         ierror) BIND(C)
17      TYPE(MPI_Comm), INTENT(IN) :: comm_old
18      INTEGER, INTENT(IN) :: nnodes, index(nnodes), edges(*)
19      LOGICAL, INTENT(IN) :: reorder
20      TYPE(MPI_Comm), INTENT(OUT) :: comm_graph
21      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 22
23      MPI_Graphdims_get(comm, nnodes, nedges, ierror) BIND(C)
24      TYPE(MPI_Comm), INTENT(IN) :: comm
25      INTEGER, INTENT(OUT) :: nnodes, nedges
ticket-248T. 26
27      MPI_Graph_get(comm, maxindex, maxedges, index, edges, ierror) BIND(C)
28      TYPE(MPI_Comm), INTENT(IN) :: comm
29      INTEGER, INTENT(IN) :: maxindex, maxedges
30      INTEGER, INTENT(OUT) :: index(maxindex), edges(maxedges)
31      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 32
33      MPI_Graph_map(comm, nnodes, index, edges, newrank, ierror) BIND(C)
34      TYPE(MPI_Comm), INTENT(IN) :: comm
35      INTEGER, INTENT(IN) :: nnodes, index(nnodes), edges(*)
36      INTEGER, INTENT(OUT) :: newrank
37      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 38
39      MPI_Graph_neighbors(comm, rank, maxneighbors, neighbors, ierror) BIND(C)
40      TYPE(MPI_Comm), INTENT(IN) :: comm
41      INTEGER, INTENT(IN) :: rank, maxneighbors
42      INTEGER, INTENT(OUT) :: neighbors(maxneighbors)
43      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 44
45      MPI_Graph_neighbors_count(comm, rank, nneighbors, ierror) BIND(C)
46      TYPE(MPI_Comm), INTENT(IN) :: comm
47      INTEGER, INTENT(IN) :: rank
48      INTEGER, INTENT(OUT) :: nneighbors
49      INTEGER, OPTIONAL, INTENT(OUT) :: ierror

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```

1      TYPE(MPI_Datatype), INTENT(IN), ASYNCHRONOUS :: sendtypes(*),
2      recvtypes(*)
3      TYPE(MPI_Comm), INTENT(IN) :: comm
4      TYPE(MPI_Request), INTENT(OUT) :: request
5      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 6
7      MPI_Neighbor_allgather(sendbuf, sendcount, sendtype, recvbuf, recvcount,
8          recvtype, comm, ierror) BIND(C)
9      TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
10     TYPE(*), DIMENSION(..) :: recvbuf
11     INTEGER, INTENT(IN) :: sendcount, recvcount
12     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
13     TYPE(MPI_Comm), INTENT(IN) :: comm
ticket-248T. 14     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
15     MPI_Neighbor_allgatherv(sendbuf, sendcount, sendtype, recvbuf, recvcounts,
16         displs, recvtype, comm, ierror) BIND(C)
17     TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
18     TYPE(*), DIMENSION(..) :: recvbuf
19     INTEGER, INTENT(IN) :: sendcount, recvcounts(*), displs(*)
20     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
21     TYPE(MPI_Comm), INTENT(IN) :: comm
22     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 23
24     MPI_Neighbor_alltoall(sendbuf, sendcount, sendtype, recvbuf, recvcount,
25         recvtype, comm, ierror) BIND(C)
26     TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
27     TYPE(*), DIMENSION(..) :: recvbuf
28     INTEGER, INTENT(IN) :: sendcount, recvcount
29     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
30     TYPE(MPI_Comm), INTENT(IN) :: comm
ticket-248T. 31     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
32     MPI_Neighbor_alltoallv(sendbuf, sendcounts, sdispls, sendtype, recvbuf,
33         recvcounts, rdispls, recvtype, comm, ierror) BIND(C)
34     TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
35     TYPE(*), DIMENSION(..) :: recvbuf
36     INTEGER, INTENT(IN) :: sendcounts(*), sdispls(*), recvcounts(*),
37         rdispls(*)
38     TYPE(MPI_Datatype), INTENT(IN) :: sendtype, recvtype
39     TYPE(MPI_Comm), INTENT(IN) :: comm
40     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 41
42     MPI_Neighbor_alltoallw(sendbuf, sendcounts, sdispls, sendtypes, recvbuf,
43         recvcounts, rdispls, recvtypes, comm, ierror) BIND(C)
44     TYPE(*), DIMENSION(..), INTENT(IN) :: sendbuf
45     TYPE(*), DIMENSION(..) :: recvbuf
46     INTEGER, INTENT(IN) :: sendcounts(*), recvcounts(*)
47     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: sdispls(*), rdispls(*)
48     TYPE(MPI_Datatype), INTENT(IN) :: sendtypes(*), recvtypes(*)

```

```

TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Topo_test(comm, status, ierror) BIND(C)
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, INTENT(OUT) :: status
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

### A.3.6 MPI Environmental Management Fortran 2008 Bindings

```

DOUBLE PRECISION MPI_Wtick() BIND(C)

DOUBLE PRECISION MPI_Wtime() BIND(C)

MPI_Abort(comm, errorcode, ierror) BIND(C)
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, INTENT(IN) :: errorcode
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Add_error_class(errorclass, ierror) BIND(C)
INTEGER, INTENT(OUT) :: errorclass
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Add_error_code(errorclass, errorcode, ierror) BIND(C)
INTEGER, INTENT(IN) :: errorclass
INTEGER, INTENT(OUT) :: errorcode
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Add_error_string(errorcode, string, ierror) BIND(C)
INTEGER, INTENT(IN) :: errorcode
CHARACTER(LEN=*), INTENT(IN) :: string
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Alloc_mem(size, info, baseptr, ierror) BIND(C)
USE, INTRINSIC :: ISO_C_BINDING, ONLY : C_PTR
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: size
TYPE(MPI_Info), INTENT(IN) :: info
TYPE(C_PTR), INTENT(OUT) :: baseptr
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Comm_call_errhandler(comm, errorcode, ierror) BIND(C)
TYPE(MPI_Comm), INTENT(IN) :: comm
INTEGER, INTENT(IN) :: errorcode
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Comm_create_errhandler(comm_errhandler_fn, errhandler, ierror) BIND(C)
PROCEDURE(MPI_Comm_errhandler_function) :: comm_errhandler_fn
TYPE(MPI_Errhandler), INTENT(OUT) :: errhandler
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Comm_get_errhandler(comm, errhandler, ierror) BIND(C)
TYPE(MPI_Comm), INTENT(IN) :: comm

```

```

1      TYPE(MPI_Errhandler), INTENT(OUT) :: errhandler
2      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 3
4      MPI_Comm_set_errhandler(comm, errhandler, ierror) BIND(C)
5      TYPE(MPI_Comm), INTENT(IN) :: comm
6      TYPE(MPI_Errhandler), INTENT(IN) :: errhandler
7      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 8
9      MPI_Errhandler_free(errhandler, ierror) BIND(C)
10     TYPE(MPI_Errhandler), INTENT(INOUT) :: errhandler
11     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 12
13     MPI_Error_class(errorcode, errorclass, ierror) BIND(C)
14     INTEGER, INTENT(IN) :: errorcode
15     INTEGER, INTENT(OUT) :: errorclass
16     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 17
18     MPI_Error_string(errorcode, string, resultlen, ierror) BIND(C)
19     INTEGER, INTENT(IN) :: errorcode
20     CHARACTER(LEN=MPI_MAX_ERROR_STRING), INTENT(OUT) :: string
21     INTEGER, INTENT(OUT) :: resultlen
22     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 23
24     MPI_File_call_errhandler(fh, errorcode, ierror) BIND(C)
25     TYPE(MPI_File), INTENT(IN) :: fh
26     INTEGER, INTENT(IN) :: errorcode
27     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 28
29     MPI_File_create_errhandler(file_errhandler_fn, errhandler, ierror) BIND(C)
30     PROCEDURE(MPI_File_errhandler_function) :: file_errhandler_fn
31     TYPE(MPI_Errhandler), INTENT(OUT) :: errhandler
32     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 33
34     MPI_File_get_errhandler(file, errhandler, ierror) BIND(C)
35     TYPE(MPI_File), INTENT(IN) :: file
36     TYPE(MPI_Errhandler), INTENT(OUT) :: errhandler
37     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 38
39     MPI_File_set_errhandler(file, errhandler, ierror) BIND(C)
40     TYPE(MPI_File), INTENT(IN) :: file
41     TYPE(MPI_Errhandler), INTENT(IN) :: errhandler
42     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 43
44     MPI_Finalized(flag, ierror) BIND(C)
45     LOGICAL, INTENT(OUT) :: flag
46     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 47
48     MPI_Finalize(ierr) BIND(C)
49     INTEGER, OPTIONAL, INTENT(OUT) :: ierr
ticket-248T. 50
51     MPI_Free_mem(base, ierror) BIND(C)
52     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: base
53     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```



```

MPI_Get_processor_name(name, resultlen, ierror) BIND(C)
  CHARACTER(LEN=MPI_MAX_PROCESSOR_NAME), INTENT(OUT) :: name
  INTEGER, INTENT(OUT) :: resultlen
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Get_version(version, subversion, ierror) BIND(C)
  INTEGER, INTENT(OUT) :: version, subversion
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Initialized(flag, ierror) BIND(C)
  LOGICAL, INTENT(OUT) :: flag
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Init(ierror) BIND(C)
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Win_call_errhandler(win, errorcode, ierror) BIND(C)
  TYPE(MPI_Win), INTENT(IN) :: win
  INTEGER, INTENT(IN) :: errorcode
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Win_create_errhandler(win_errhandler_fn, errhandler, ierror) BIND(C)
  PROCEDURE(MPI_Win_errhandler_function) :: win_errhandler_fn
  TYPE(MPI_Errhandler), INTENT(OUT) :: errhandler
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Win_get_errhandler(win, errhandler, ierror) BIND(C)
  TYPE(MPI_Win), INTENT(IN) :: win
  TYPE(MPI_Errhandler), INTENT(OUT) :: errhandler
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Win_set_errhandler(win, errhandler, ierror) BIND(C)
  TYPE(MPI_Win), INTENT(IN) :: win
  TYPE(MPI_Errhandler), INTENT(IN) :: errhandler
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

### A.3.7 The Info Object Fortran 2008 Bindings

```

MPI_Info_create(info, ierror) BIND(C)
  TYPE(MPI_Info), INTENT(OUT) :: info
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Info_delete(info, key, ierror) BIND(C)
  TYPE(MPI_Info), INTENT(IN) :: info
  CHARACTER(LEN=*), INTENT(IN) :: key
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Info_dup(info, newinfo, ierror) BIND(C)
  TYPE(MPI_Info), INTENT(IN) :: info
  TYPE(MPI_Info), INTENT(OUT) :: newinfo
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

1
ticket-248T. 2 MPI_Info_free(info, ierror) BIND(C)
3 TYPE(MPI_Info), INTENT(INOUT) :: info
4 INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 5
6 MPI_Info_get(info, key, valuelen, value, flag, ierror) BIND(C)
7 TYPE(MPI_Info), INTENT(IN) :: info
8 CHARACTER(LEN=*), INTENT(IN) :: key
9 INTEGER, INTENT(IN) :: valuelen
10 CHARACTER(LEN=valuelen), INTENT(OUT) :: value
11 LOGICAL, INTENT(OUT) :: flag
12 INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 13
14 MPI_Info_get_nkeys(info, nkeys, ierror) BIND(C)
15 TYPE(MPI_Info), INTENT(IN) :: info
16 INTEGER, INTENT(OUT) :: nkeys
17 INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 18
19 MPI_Info_get_nthkey(info, n, key, ierror) BIND(C)
20 TYPE(MPI_Info), INTENT(IN) :: info
21 INTEGER, INTENT(IN) :: n
22 CHARACTER(LEN=*), INTENT(OUT) :: key
23 INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 24
25 MPI_Info_get_valuelen(info, key, valuelen, flag, ierror) BIND(C)
26 TYPE(MPI_Info), INTENT(IN) :: info
27 CHARACTER(LEN=*), INTENT(IN) :: key
28 INTEGER, INTENT(OUT) :: valuelen
29 LOGICAL, INTENT(OUT) :: flag
30 INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 31
32 MPI_Info_set(info, key, value, ierror) BIND(C)
33 TYPE(MPI_Info), INTENT(IN) :: info
34 CHARACTER(LEN=*), INTENT(IN) :: key, value
35 INTEGER, OPTIONAL, INTENT(OUT) :: ierror
36
37 A.3.8 Process Creation and Management Fortran 2008 Bindings
ticket-248T. 38
39 MPI_Close_port(port_name, ierror) BIND(C)
40 CHARACTER(LEN=*), INTENT(IN) :: port_name
41 INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 42
43 MPI_Comm_accept(port_name, info, root, comm, newcomm, ierror) BIND(C)
44 CHARACTER(LEN=*), INTENT(IN) :: port_name
45 TYPE(MPI_Info), INTENT(IN) :: info
46 INTEGER, INTENT(IN) :: root
47 TYPE(MPI_Comm), INTENT(IN) :: comm
48 TYPE(MPI_Comm), INTENT(OUT) :: newcomm
49 INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 50
51 MPI_Comm_connect(port_name, info, root, comm, newcomm, ierror) BIND(C)

```

```

CHARACTER(LEN=*), INTENT(IN) :: port_name
TYPE(MPI_Info), INTENT(IN) :: info
INTEGER, INTENT(IN) :: root
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Comm), INTENT(OUT) :: newcomm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Comm_disconnect(comm, ierror) BIND(C)
TYPE(MPI_Comm), INTENT(INOUT) :: comm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Comm_get_parent(parent, ierror) BIND(C)
TYPE(MPI_Comm), INTENT(OUT) :: parent
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Comm_join(fd, intercomm, ierror) BIND(C)
INTEGER, INTENT(IN) :: fd
TYPE(MPI_Comm), INTENT(OUT) :: intercomm
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Comm_spawn(command, argv, maxprocs, info, root, comm, intercomm,
               array_of_errcodes, ierror) BIND(C)
CHARACTER(LEN=*), INTENT(IN) :: command, argv(*)
INTEGER, INTENT(IN) :: maxprocs, root
TYPE(MPI_Info), INTENT(IN) :: info
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Comm), INTENT(OUT) :: intercomm
INTEGER :: array_of_errcodes(*)
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Comm_spawn_multiple(count, array_of_commands, array_of_argv,
                       array_of_maxprocs, array_of_info, root, comm, intercomm,
                       array_of_errcodes, ierror) BIND(C)
INTEGER, INTENT(IN) :: count, array_of_maxprocs(*), root
CHARACTER(LEN=*), INTENT(IN) :: array_of_commands(*),
array_of_argv(count, *)
TYPE(MPI_Info), INTENT(IN) :: array_of_info(*)
TYPE(MPI_Comm), INTENT(IN) :: comm
TYPE(MPI_Comm), INTENT(OUT) :: intercomm
INTEGER :: array_of_errcodes(*)
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Lookup_name(service_name, info, port_name, ierror) BIND(C)
CHARACTER(LEN=*), INTENT(IN) :: service_name
TYPE(MPI_Info), INTENT(IN) :: info
CHARACTER(LEN=MPI_MAX_PORT_NAME), INTENT(OUT) :: port_name
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Open_port(info, port_name, ierror) BIND(C)
TYPE(MPI_Info), INTENT(IN) :: info
CHARACTER(LEN=MPI_MAX_PORT_NAME), INTENT(OUT) :: port_name

```

```

1      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 2
3      MPI_Publish_name(service_name, info, port_name, ierror) BIND(C)
4      TYPE(MPI_Info), INTENT(IN) :: info
5      CHARACTER(LEN=*), INTENT(IN) :: service_name, port_name
6      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 7
8      MPI_Unpublish_name(service_name, info, port_name, ierror) BIND(C)
9      CHARACTER(LEN=*), INTENT(IN) :: service_name, port_name
10     TYPE(MPI_Info), INTENT(IN) :: info
11     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
12
13     A.3.9 One-Sided Communications Fortran 2008 Bindings
ticket-248T. 14
15     MPI_Accumulate(origin_addr, origin_count, origin_datatype, target_rank,
16                   target_disp, target_count, target_datatype, op, win, ierror)
17                   BIND(C)
18     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr
19     INTEGER, INTENT(IN) :: origin_count, target_rank, target_count
20     TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype
21     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
22     TYPE(MPI_Op), INTENT(IN) :: op
23     TYPE(MPI_Win), INTENT(IN) :: win
24     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 25
26     MPI_Compare_and_swap(origin_addr, compare_addr, result_addr, datatype,
27                          target_rank, target_disp, win, ierror) BIND(C)
28     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr,
29     compare_addr
30     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: result_addr
31     TYPE(MPI_Datatype), INTENT(IN) :: datatype
32     INTEGER, INTENT(IN) :: target_rank
33     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
34     TYPE(MPI_Win), INTENT(IN) :: win
35     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 36
37     MPI_Fetch_and_op(origin_addr, result_addr, datatype, target_rank,
38                     target_disp, op, win, ierror) BIND(C)
39     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr
40     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: result_addr
41     TYPE(MPI_Datatype), INTENT(IN) :: datatype
42     INTEGER, INTENT(IN) :: target_rank
43     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
44     TYPE(MPI_Op), INTENT(IN) :: op
45     TYPE(MPI_Win), INTENT(IN) :: win
46     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 47
48     MPI_Get_accumulate(origin_addr, origin_count, origin_datatype, result_addr,
49                       result_count, result_datatype, target_rank, target_disp,
50                       target_count, target_datatype, op, win, ierror) BIND(C)

```

```

TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr
TYPE(*), DIMENSION(..), ASYNCHRONOUS :: result_addr
INTEGER, INTENT(IN) :: origin_count, result_count, target_rank,
target_count
TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype,
result_datatype
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
TYPE(MPI_Op), INTENT(IN) :: op
TYPE(MPI_Win), INTENT(IN) :: win
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Get(origin_addr, origin_count, origin_datatype, target_rank,
        target_disp, target_count, target_datatype, win, ierror)
BIND(C)
TYPE(*), DIMENSION(..), ASYNCHRONOUS :: origin_addr
INTEGER, INTENT(IN) :: origin_count, target_rank, target_count
TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
TYPE(MPI_Win), INTENT(IN) :: win
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Put(origin_addr, origin_count, origin_datatype, target_rank,
        target_disp, target_count, target_datatype, win, ierror)
BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr
INTEGER, INTENT(IN) :: origin_count, target_rank, target_count
TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
TYPE(MPI_Win), INTENT(IN) :: win
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Raccumulate(origin_addr, origin_count, origin_datatype, target_rank,
        target_disp, target_count, target_datatype, op, win, request,
        ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr
INTEGER, INTENT(IN) :: origin_count, target_rank, target_count
TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype
INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
TYPE(MPI_Op), INTENT(IN) :: op
TYPE(MPI_Win), INTENT(IN) :: win
TYPE(MPI_Request), INTENT(OUT) :: request
INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Rget_accumulate(origin_addr, origin_count, origin_datatype,
        result_addr, result_count, result_datatype, target_rank,
        target_disp, target_count, target_datatype, op, win, request,
        ierror) BIND(C)
TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr
TYPE(*), DIMENSION(..), ASYNCHRONOUS :: result_addr

```

```

1      INTEGER, INTENT(IN) :: origin_count, result_count, target_rank,
2      target_count
3      TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype,
4      result_datatype
5      INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
6      TYPE(MPI_Op), INTENT(IN) :: op
7      TYPE(MPI_Win), INTENT(IN) :: win
8      TYPE(MPI_Request), INTENT(OUT) :: request
9      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 10
11     MPI_Rget(origin_addr, origin_count, origin_datatype, target_rank,
12             target_disp, target_count, target_datatype, win, request,
13             ierror) BIND(C)
14     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: origin_addr
15     INTEGER, INTENT(IN) :: origin_count, target_rank, target_count
16     TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype
17     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
18     TYPE(MPI_Win), INTENT(IN) :: win
19     TYPE(MPI_Request), INTENT(OUT) :: request
ticket-248T. 20     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
21
22     MPI_Rput(origin_addr, origin_count, origin_datatype, target_rank,
23             target_disp, target_count, target_datatype, win, request,
24             ierror) BIND(C)
25     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: origin_addr
26     INTEGER, INTENT(IN) :: origin_count, target_rank, target_count
27     TYPE(MPI_Datatype), INTENT(IN) :: origin_datatype, target_datatype
28     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: target_disp
29     TYPE(MPI_Win), INTENT(IN) :: win
30     TYPE(MPI_Request), INTENT(OUT) :: request
ticket-248T. 31     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
32
33     MPI_Win_allocate(size, disp_unit, info, comm, baseptr, win, ierror) BIND(C)
34     USE, INTRINSIC :: ISO_C_BINDING, ONLY : C_PTR
35     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: size
36     INTEGER, INTENT(IN) :: disp_unit
37     TYPE(MPI_Info), INTENT(IN) :: info
38     TYPE(MPI_Comm), INTENT(IN) :: comm
39     TYPE(C_PTR), INTENT(OUT) :: baseptr
40     TYPE(MPI_Win), INTENT(OUT) :: win
ticket-248T. 41     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
42
43     MPI_Win_attach(win, base, size, ierror) BIND(C)
44     TYPE(MPI_Win), INTENT(IN) :: win
45     TYPE(*), ASYNCHRONOUS :: base
46     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: size
47     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 48
49     MPI_Win_complete(win, ierror) BIND(C)
50     TYPE(MPI_Win), INTENT(IN) :: win

```

```

    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Win_create(base, size, disp_unit, info, comm, win, ierror) BIND(C)
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: base
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: size
    INTEGER, INTENT(IN) :: disp_unit
    TYPE(MPI_Info), INTENT(IN) :: info
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Win), INTENT(OUT) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Win_create_dynamic(info, comm, win, ierror) BIND(C)
    TYPE(MPI_Info), INTENT(IN) :: info
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Win), INTENT(OUT) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Win_detach(win, base, ierror) BIND(C)
    TYPE(MPI_Win), INTENT(IN) :: win
    TYPE(*), ASYNCHRONOUS :: base
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Win_fence(assert, win, ierror) BIND(C)
    INTEGER, INTENT(IN) :: assert
    TYPE(MPI_Win), INTENT(IN) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Win_flush_all(win, ierror) BIND(C)
    TYPE(MPI_Win), INTENT(IN) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Win_flush_local_all(win, ierror) BIND(C)
    TYPE(MPI_Win), INTENT(IN) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Win_flush_local(rank, win, ierror) BIND(C)
    INTEGER, INTENT(IN) :: rank
    TYPE(MPI_Win), INTENT(IN) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Win_flush(rank, win, ierror) BIND(C)
    INTEGER, INTENT(IN) :: rank
    TYPE(MPI_Win), INTENT(IN) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Win_free(win, ierror) BIND(C)
    TYPE(MPI_Win), INTENT(INOUT) :: win
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_Win_get_group(win, group, ierror) BIND(C)
    TYPE(MPI_Win), INTENT(IN) :: win
    TYPE(MPI_Group), INTENT(OUT) :: group

```

```

1      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 2
3      MPI_Win_lock_all(assert, win, ierror) BIND(C)
4      INTEGER, INTENT(IN) :: assert
5      TYPE(MPI_Win), INTENT(IN) :: win
6      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 7
8      MPI_Win_lock(lock_type, rank, assert, win, ierror) BIND(C)
9      INTEGER, INTENT(IN) :: lock_type, rank, assert
10     TYPE(MPI_Win), INTENT(IN) :: win
11     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 12
13     MPI_Win_post(group, assert, win, ierror) BIND(C)
14     TYPE(MPI_Group), INTENT(IN) :: group
15     INTEGER, INTENT(IN) :: assert
16     TYPE(MPI_Win), INTENT(IN) :: win
17     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 18
19     MPI_Win_start(group, assert, win, ierror) BIND(C)
20     TYPE(MPI_Group), INTENT(IN) :: group
21     INTEGER, INTENT(IN) :: assert
22     TYPE(MPI_Win), INTENT(IN) :: win
23     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 24
25     MPI_Win_sync(win, ierror) BIND(C)
26     TYPE(MPI_Win), INTENT(IN) :: win
27     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 28
29     MPI_Win_test(win, flag, ierror) BIND(C)
30     TYPE(MPI_Win), INTENT(IN) :: win
31     LOGICAL, INTENT(OUT) :: flag
32     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 33
34     MPI_Win_unlock_all(win, ierror) BIND(C)
35     TYPE(MPI_Win), INTENT(IN) :: win
36     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 37
38     MPI_Win_unlock(rank, win, ierror) BIND(C)
39     INTEGER, INTENT(IN) :: rank
40     TYPE(MPI_Win), INTENT(IN) :: win
41     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 42
43     MPI_Win_wait(win, ierror) BIND(C)
44     TYPE(MPI_Win), INTENT(IN) :: win
45     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
46
47
48

```

#### A.3.10 External Interfaces Fortran 2008 Bindings

```

MPI_Grequest_complete(request, ierror) BIND(C)
TYPE(MPI_Request), INTENT(IN) :: request
INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```



```

MPI_Grequest_start(query_fn, free_fn, cancel_fn, extra_state, request,
                  ierror) BIND(C)
    PROCEDURE(MPI_Grequest_query_function) :: query_fn
    PROCEDURE(MPI_Grequest_free_function) :: free_fn
    PROCEDURE(MPI_Grequest_cancel_function) :: cancel_fn
    INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Init_thread(required, provided, ierror) BIND(C)
    INTEGER, INTENT(IN) :: required
    INTEGER, INTENT(OUT) :: provided
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Is_thread_main(flag, ierror) BIND(C)
    LOGICAL, INTENT(OUT) :: flag
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Query_thread(provided, ierror) BIND(C)
    INTEGER, INTENT(OUT) :: provided
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Status_set_cancelled(status, flag, ierror) BIND(C)
    TYPE(MPI_Status), INTENT(INOUT) :: status
    LOGICAL, INTENT(OUT) :: flag
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Status_set_elements(status, datatype, count, ierror) BIND(C)
    TYPE(MPI_Status), INTENT(INOUT) :: status
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    INTEGER, INTENT(IN) :: count
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Status_set_elements_x(status, datatype, count, ierror) BIND(C)
    TYPE(MPI_Status), INTENT(INOUT) :: status
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    INTEGER(KIND = MPI_COUNT_KIND), INTENT(IN) :: count
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

#### A.3.11 I/O Fortran 2008 Bindings

```

MPI_File_close(fh, ierror) BIND(C)
    TYPE(MPI_File), INTENT(INOUT) :: fh
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_File_delete(filename, info, ierror) BIND(C)
    CHARACTER(LEN=*), INTENT(IN) :: filename
    TYPE(MPI_Info), INTENT(IN) :: info
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_File_get_amode(fh, amode, ierror) BIND(C)

```

```

1      TYPE(MPI_File), INTENT(IN) :: fh
2      INTEGER, INTENT(OUT) :: amode
3      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 4
5      MPI_File_get_atomicity(fh, flag, ierror) BIND(C)
6      TYPE(MPI_File), INTENT(IN) :: fh
7      LOGICAL, INTENT(OUT) :: flag
8      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 9
10     MPI_File_get_byte_offset(fh, offset, disp, ierror) BIND(C)
11     TYPE(MPI_File), INTENT(IN) :: fh
12     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
13     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(OUT) :: disp
14     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 15
16     MPI_File_get_group(fh, group, ierror) BIND(C)
17     TYPE(MPI_File), INTENT(IN) :: fh
18     TYPE(MPI_Group), INTENT(OUT) :: group
19     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 20
21     MPI_File_get_info(fh, info_used, ierror) BIND(C)
22     TYPE(MPI_File), INTENT(IN) :: fh
23     TYPE(MPI_Info), INTENT(OUT) :: info_used
24     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 25
26     MPI_File_get_position(fh, offset, ierror) BIND(C)
27     TYPE(MPI_File), INTENT(IN) :: fh
28     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(OUT) :: offset
29     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 30
31     MPI_File_get_position_shared(fh, offset, ierror) BIND(C)
32     TYPE(MPI_File), INTENT(IN) :: fh
33     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(OUT) :: offset
34     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 35
36     MPI_File_get_size(fh, size, ierror) BIND(C)
37     TYPE(MPI_File), INTENT(IN) :: fh
38     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(OUT) :: size
39     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 40
41     MPI_File_get_type_extent(fh, datatype, extent, ierror) BIND(C)
42     TYPE(MPI_File), INTENT(IN) :: fh
43     TYPE(MPI_Datatype), INTENT(IN) :: datatype
44     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(OUT) :: extent
45     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 46
47     MPI_File_get_view(fh, disp, etype, filetype, datarep, ierror) BIND(C)
48     TYPE(MPI_File), INTENT(IN) :: fh
49     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(OUT) :: disp
50     TYPE(MPI_Datatype), INTENT(OUT) :: etype, filetype
51     CHARACTER(LEN=*), INTENT(OUT) :: datarep
52     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_File_iread_at(fh, offset, buf, count, datatype, request, ierror)
    BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_File_iread(fh, buf, count, datatype, request, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_File_iread_shared(fh, buf, count, datatype, request, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_File_iwrite_at(fh, offset, buf, count, datatype, request, ierror)
    BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_File_iwrite(fh, buf, count, datatype, request, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Request), INTENT(OUT) :: request
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_File_iwrite_shared(fh, buf, count, datatype, request, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype

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43 ticket-248T.  
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1      TYPE(MPI_Request), INTENT(OUT) :: request
2      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 3
4      MPI_File_open(comm, filename, amode, info, fh, ierror) BIND(C)
5      TYPE(MPI_Comm), INTENT(IN) :: comm
6      CHARACTER(LEN=*), INTENT(IN) :: filename
7      INTEGER, INTENT(IN) :: amode
8      TYPE(MPI_Info), INTENT(IN) :: info
9      TYPE(MPI_File), INTENT(OUT) :: fh
10     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 10
11     MPI_File_preallocate(fh, size, ierror) BIND(C)
12     TYPE(MPI_File), INTENT(IN) :: fh
13     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: size
14     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 15
16     MPI_File_read_all_begin(fh, buf, count, datatype, ierror) BIND(C)
17     TYPE(MPI_File), INTENT(IN) :: fh
18     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
19     INTEGER, INTENT(IN) :: count
20     TYPE(MPI_Datatype), INTENT(IN) :: datatype
21     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 21
22     MPI_File_read_all_end(fh, buf, status, ierror) BIND(C)
23     TYPE(MPI_File), INTENT(IN) :: fh
24     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
25     TYPE(MPI_Status) :: status
26     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 27
28     MPI_File_read_all(fh, buf, count, datatype, status, ierror) BIND(C)
29     TYPE(MPI_File), INTENT(IN) :: fh
30     TYPE(*), DIMENSION(..) :: buf
31     INTEGER, INTENT(IN) :: count
32     TYPE(MPI_Datatype), INTENT(IN) :: datatype
33     TYPE(MPI_Status) :: status
34     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 34
35     MPI_File_read_at_all_begin(fh, offset, buf, count, datatype, ierror)
36     BIND(C)
37     TYPE(MPI_File), INTENT(IN) :: fh
38     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
39     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
40     INTEGER, INTENT(IN) :: count
41     TYPE(MPI_Datatype), INTENT(IN) :: datatype
42     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 43
44     MPI_File_read_at_all_end(fh, buf, status, ierror) BIND(C)
45     TYPE(MPI_File), INTENT(IN) :: fh
46     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
47     TYPE(MPI_Status) :: status
48     INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

```

MPI_File_read_at_all(fh, offset, buf, count, datatype, status, ierror)
    BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
    TYPE(*), DIMENSION(..) :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_File_read_at(fh, offset, buf, count, datatype, status, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
    TYPE(*), DIMENSION(..) :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_File_read(fh, buf, count, datatype, status, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..) :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_File_read_ordered_begin(fh, buf, count, datatype, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_File_read_ordered_end(fh, buf, status, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_File_read_ordered(fh, buf, count, datatype, status, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..) :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_File_read_shared(fh, buf, count, datatype, status, ierror) BIND(C)

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27 ticket-248T.  
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33 ticket-248T.  
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39 ticket-248T.  
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46 ticket-248T.  
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```

1      TYPE(MPI_File), INTENT(IN) :: fh
2      TYPE(*), DIMENSION(..) :: buf
3      INTEGER, INTENT(IN) :: count
4      TYPE(MPI_Datatype), INTENT(IN) :: datatype
5      TYPE(MPI_Status) :: status
6      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 7
8      MPI_File_seek(fh, offset, whence, ierror) BIND(C)
9      TYPE(MPI_File), INTENT(IN) :: fh
10     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
11     INTEGER, INTENT(IN) :: whence
12     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 13
14     MPI_File_seek_shared(fh, offset, whence, ierror) BIND(C)
15     TYPE(MPI_File), INTENT(IN) :: fh
16     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
17     INTEGER, INTENT(IN) :: whence
18     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 19
20     MPI_File_set_atomicity(fh, flag, ierror) BIND(C)
21     TYPE(MPI_File), INTENT(IN) :: fh
22     LOGICAL, INTENT(IN) :: flag
23     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 24
25     MPI_File_set_info(fh, info, ierror) BIND(C)
26     TYPE(MPI_File), INTENT(IN) :: fh
27     TYPE(MPI_Info), INTENT(IN) :: info
28     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 29
30     MPI_File_set_size(fh, size, ierror) BIND(C)
31     TYPE(MPI_File), INTENT(IN) :: fh
32     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: size
33     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 34
35     MPI_File_set_view(fh, disp, etype, filetype, datarep, info, ierror) BIND(C)
36     TYPE(MPI_File), INTENT(IN) :: fh
37     INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: disp
38     TYPE(MPI_Datatype), INTENT(IN) :: etype, filetype
39     CHARACTER(LEN=*), INTENT(IN) :: datarep
40     TYPE(MPI_Info), INTENT(IN) :: info
41     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 42
43     MPI_File_sync(fh, ierror) BIND(C)
44     TYPE(MPI_File), INTENT(IN) :: fh
45     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 46
47     MPI_File_write_all_begin(fh, buf, count, datatype, ierror) BIND(C)
48     TYPE(MPI_File), INTENT(IN) :: fh
49     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
50     INTEGER, INTENT(IN) :: count
51     TYPE(MPI_Datatype), INTENT(IN) :: datatype

```

```

    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_File_write_all_end(fh, buf, status, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_File_write_all(fh, buf, count, datatype, status, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), INTENT(IN) :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_File_write_at_all_begin(fh, offset, buf, count, datatype, ierror)
    BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_File_write_at_all_end(fh, buf, status, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_File_write_at_all(fh, offset, buf, count, datatype, status, ierror)
    BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
    TYPE(*), DIMENSION(..), INTENT(IN) :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_File_write_at(fh, offset, buf, count, datatype, status, ierror) BIND(C)
    TYPE(MPI_File), INTENT(IN) :: fh
    INTEGER(KIND=MPI_OFFSET_KIND), INTENT(IN) :: offset
    TYPE(*), DIMENSION(..), INTENT(IN) :: buf
    INTEGER, INTENT(IN) :: count
    TYPE(MPI_Datatype), INTENT(IN) :: datatype
    TYPE(MPI_Status) :: status
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_File_write(fh, buf, count, datatype, status, ierror) BIND(C)

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47 ticket-248T.  
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```

1      TYPE(MPI_File), INTENT(IN) :: fh
2      TYPE(*), DIMENSION(..), INTENT(IN) :: buf
3      INTEGER, INTENT(IN) :: count
4      TYPE(MPI_Datatype), INTENT(IN) :: datatype
5      TYPE(MPI_Status) :: status
6      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 7
8      MPI_File_write_ordered_begin(fh, buf, count, datatype, ierror) BIND(C)
9      TYPE(MPI_File), INTENT(IN) :: fh
10     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
11     INTEGER, INTENT(IN) :: count
12     TYPE(MPI_Datatype), INTENT(IN) :: datatype
13     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 14
15     MPI_File_write_ordered_end(fh, buf, status, ierror) BIND(C)
16     TYPE(MPI_File), INTENT(IN) :: fh
17     TYPE(*), DIMENSION(..), INTENT(IN), ASYNCHRONOUS :: buf
18     TYPE(MPI_Status) :: status
19     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 20
21     MPI_File_write_ordered(fh, buf, count, datatype, status, ierror) BIND(C)
22     TYPE(MPI_File), INTENT(IN) :: fh
23     TYPE(*), DIMENSION(..), INTENT(IN) :: buf
24     INTEGER, INTENT(IN) :: count
25     TYPE(MPI_Datatype), INTENT(IN) :: datatype
26     TYPE(MPI_Status) :: status
27     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 28
29     MPI_File_write_shared(fh, buf, count, datatype, status, ierror) BIND(C)
30     TYPE(MPI_File), INTENT(IN) :: fh
31     TYPE(*), DIMENSION(..), INTENT(IN) :: buf
32     INTEGER, INTENT(IN) :: count
33     TYPE(MPI_Datatype), INTENT(IN) :: datatype
34     TYPE(MPI_Status) :: status
35     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
ticket-248T. 36
37     MPI_Register_datarep(datarep, read_conversion_fn, write_conversion_fn,
38                          dtype_file_extent_fn, extra_state, ierror) BIND(C)
39     CHARACTER(LEN=*), INTENT(IN) :: datarep
40     PROCEDURE(MPI_Datarep_conversion_function) :: read_conversion_fn
41     PROCEDURE(MPI_Datarep_conversion_function) :: write_conversion_fn
42     PROCEDURE(MPI_Datarep_extent_function) :: dtype_file_extent_fn
43     INTEGER(KIND=MPI_ADDRESS_KIND), INTENT(IN) :: extra_state
44     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
45
46     A.3.12 Language Bindings Fortran 2008 Bindings
47     MPI_F_sync_reg(buf) BIND(C)
48     TYPE(*), DIMENSION(..), ASYNCHRONOUS :: buf

```



```

MPI_Sizeof(x, size, ierror) BIND(C)
  TYPE(*), DIMENSION(..) :: x
  INTEGER, INTENT(OUT) :: size
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Status_f082f(f08_status, f_status, ierror) BIND(C)
  TYPE(MPI_Status), INTENT(IN) :: f08_status
  INTEGER, INTENT(OUT) :: f_status(MPI_STATUS_SIZE)
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Status_f2f08(f_status, f08_status, ierror) BIND(C)
  INTEGER, INTENT(IN) :: f_status(MPI_STATUS_SIZE)
  TYPE(MPI_Status), INTENT(OUT) :: f08_status
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Type_create_f90_complex(p, r, newtype, ierror) BIND(C)
  INTEGER, INTENT(IN) :: p, r
  TYPE(MPI_Datatype), INTENT(OUT) :: newtype
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Type_create_f90_integer(r, newtype, ierror) BIND(C)
  INTEGER, INTENT(IN) :: r
  TYPE(MPI_Datatype), INTENT(OUT) :: newtype
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Type_create_f90_real(p, r, newtype, ierror) BIND(C)
  INTEGER, INTENT(IN) :: p, r
  TYPE(MPI_Datatype), INTENT(OUT) :: newtype
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

MPI_Type_match_size(typeclass, size, datatype, ierror) BIND(C)
  INTEGER, INTENT(IN) :: typeclass, size
  TYPE(MPI_Datatype), INTENT(OUT) :: datatype
  INTEGER, OPTIONAL, INTENT(OUT) :: ierror

```

### A.3.13 Profiling Interface Fortran 2008 Bindings

### A.3.14 Deprecated Fortran 2008 Bindings

## A.4 Fortran Bindings with mpif.h or the mpi Module

### A.4.1 Point-to-Point Communication Fortran Bindings

```

MPI_BSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, IERROR)
    <type> BUF(*)
    INTEGER COUNT, DATATYPE, DEST, TAG, COMM, IERROR

MPI_BSEND_INIT(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)
    <type> BUF(*)
    INTEGER [REQUEST, ]COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR

MPI_BUFFER_ATTACH(BUFFER, SIZE, IERROR)
    <type> BUFFER(*)
    INTEGER SIZE, IERROR

MPI_BUFFER_DETACH(BUFFER_ADDR, SIZE, IERROR)
    <type> BUFFER_ADDR(*)
    INTEGER SIZE, IERROR

MPI_CANCEL(REQUEST, IERROR)
    INTEGER REQUEST, IERROR

MPI_GET_COUNT(STATUS, DATATYPE, COUNT, IERROR)
    INTEGER STATUS(MPI_STATUS_SIZE), DATATYPE, COUNT, IERROR

MPI_IBSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)
    <type> BUF(*)
    INTEGER COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR

MPI_IMPROBE(SOURCE, TAG, COMM, FLAG, MESSAGE, STATUS, IERROR)
    INTEGER SOURCE, TAG, COMM, FLAG, MESSAGE, STATUS(MPI_STATUS_SIZE),
    IERROR

MPI_IMRECV(BUF, COUNT, DATATYPE, MESSAGE, REQUEST, IERROR)
    <type> BUF(*)
    INTEGER COUNT, DATATYPE, MESSAGE, REQUEST, IERROR

MPI_IPROBE(SOURCE, TAG, COMM, FLAG, STATUS, IERROR)
    LOGICAL FLAG
    INTEGER SOURCE, TAG, COMM, STATUS(MPI_STATUS_SIZE), IERROR

MPI_IRECV(BUF, COUNT, DATATYPE, SOURCE, TAG, COMM, REQUEST, IERROR)
    <type> BUF(*)
    INTEGER COUNT, DATATYPE, SOURCE, TAG, COMM, REQUEST, IERROR

MPI_IRSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)
    <type> BUF(*)
    INTEGER COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR

MPI_ISEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)
    <type> BUF(*)
    INTEGER COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR

```

```

MPI_ISSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)      1
    <type> BUF(*)                                                         2
    INTEGER COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR            3
                                                                           4
MPI_MPROBE(SOURCE, TAG, COMM, MESSAGE, STATUS, IERROR)                  5
    INTEGER SOURCE, TAG, COMM, MESSAGE, STATUS(MPI_STATUS_SIZE), IERROR  6
                                                                           7
MPI_MRECV(BUF, COUNT, DATATYPE, MESSAGE, STATUS, IERROR)                8
    <type> BUF(*)                                                         9
    INTEGER COUNT, DATATYPE, MESSAGE, STATUS(MPI_STATUS_SIZE), IERROR   10
                                                                           11
MPI_PROBE(SOURCE, TAG, COMM, STATUS, IERROR)                            12
    INTEGER SOURCE, TAG, COMM, STATUS(MPI_STATUS_SIZE), IERROR          13
                                                                           14
MPI_RECV(BUF, COUNT, DATATYPE, SOURCE, TAG, COMM, STATUS, IERROR)       15
    <type> BUF(*)                                                         16
    INTEGER COUNT, DATATYPE, SOURCE, TAG, COMM, STATUS(MPI_STATUS_SIZE), 17
    IERROR                                                                18
                                                                           19
MPI_RECV_INIT(BUF, COUNT, DATATYPE, SOURCE, TAG, COMM, REQUEST, IERROR) 20
    <type> BUF(*)                                                         21
    INTEGER COUNT, DATATYPE, SOURCE, TAG, COMM, REQUEST, IERROR         22
                                                                           23
MPI_REQUEST_FREE(REQUEST, IERROR)                                       24
    INTEGER REQUEST, IERROR                                              25
                                                                           26
MPI_REQUEST_GET_STATUS( REQUEST, FLAG, STATUS, IERROR)                  27
    INTEGER REQUEST, STATUS(MPI_STATUS_SIZE), IERROR                    28
    LOGICAL FLAG                                                         29
                                                                           30
MPI_RSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, IERROR)                31
    <type> BUF(*)                                                         32
    INTEGER COUNT, DATATYPE, DEST, TAG, COMM, IERROR                    33
                                                                           34
MPI_RSEND_INIT(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)  35
    <type> BUF(*)                                                         36
    INTEGER COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR           37
                                                                           38
MPI_SEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, IERROR)                 39
    <type> BUF(*)                                                         40
    INTEGER COUNT, DATATYPE, DEST, TAG, COMM, IERROR                    41
                                                                           42
MPI_SEND_INIT(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)   43
    <type> BUF(*)                                                         44
    INTEGER [REQUEST, ]COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR 45
                                                                           46
MPI_SENDRECV_REPLACE(BUF, COUNT, DATATYPE, DEST, SENDTAG, SOURCE, RECVTAG, 47
    COMM, STATUS, IERROR)                                                48
    <type> BUF(*)                                                         49
    INTEGER COUNT, DATATYPE, DEST, SENDTAG, SOURCE, RECVTAG, COMM,       50
    STATUS(MPI_STATUS_SIZE), IERROR                                     51
                                                                           52
MPI_SENDRECV(SENDBUF, SENDCOUNT, SENDTYPE, DEST, SENDTAG, RECVBUF,    53
    RECVCOUNT, RECVTYPE, SOURCE, RECVTAG, COMM, STATUS, IERROR)         54

```

```

1      <type> SENDBUF(*), RECVBUF(*)
2      INTEGER SENDCOUNT, SENDTYPE, DEST, SENDTAG, RECVCOUNT, RECVMYPE,
3      SOURCE, RECMTAG, COMM, STATUS(MPI_STATUS_SIZE), IERROR
4
5      MPI_SSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, IERROR)
6      <type> BUF(*)
7      INTEGER COUNT, DATATYPE, DEST, TAG, COMM, IERROR
8
9      MPI_SSEND_INIT(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)
10     <type> BUF(*)
11     INTEGER COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR
12
13     MPI_STARTALL(COUNT, ARRAY_OF_REQUESTS, IERROR)
14     INTEGER COUNT, ARRAY_OF_REQUESTS(*), IERROR
15
16     MPI_START(REQUEST, IERROR)
17     INTEGER REQUEST, IERROR
18
19     MPI_TESTALL(COUNT, ARRAY_OF_REQUESTS, FLAG, ARRAY_OF_STATUSES, IERROR)
20     LOGICAL FLAG
21     INTEGER COUNT, ARRAY_OF_REQUESTS(*),
22     ARRAY_OF_STATUSES(MPI_STATUS_SIZE,*), IERROR
23
24     MPI_TESTANY(COUNT, ARRAY_OF_REQUESTS, INDEX, FLAG, STATUS, IERROR)
25     LOGICAL FLAG
26     INTEGER COUNT, ARRAY_OF_REQUESTS(*), INDEX, STATUS(MPI_STATUS_SIZE),
27     IERROR
28
29     MPI_TEST_CANCELLED(STATUS, FLAG, IERROR)
30     LOGICAL FLAG
31     INTEGER STATUS(MPI_STATUS_SIZE), IERROR
32
33     MPI_TEST(REQUEST, FLAG, STATUS, IERROR)
34     LOGICAL FLAG
35     INTEGER REQUEST, STATUS(MPI_STATUS_SIZE), IERROR
36
37     MPI_TESTSOME(INCOUNT, ARRAY_OF_REQUESTS, OUTCOUNT, ARRAY_OF_INDICES,
38     ARRAY_OF_STATUSES, IERROR)
39     INTEGER INCOUNT, ARRAY_OF_REQUESTS(*), OUTCOUNT, ARRAY_OF_INDICES(*),
40     ARRAY_OF_STATUSES(MPI_STATUS_SIZE,*), IERROR
41
42     MPI_WAITALL(COUNT, ARRAY_OF_REQUESTS, ARRAY_OF_STATUSES, IERROR)
43     INTEGER COUNT, ARRAY_OF_REQUESTS(*)
44     INTEGER ARRAY_OF_STATUSES(MPI_STATUS_SIZE,*), IERROR
45
46     MPI_WAITANY(COUNT, ARRAY_OF_REQUESTS, INDEX, STATUS, IERROR)
47     INTEGER COUNT, ARRAY_OF_REQUESTS(*), INDEX, STATUS(MPI_STATUS_SIZE),
48     IERROR
49
50     MPI_WAIT(REQUEST, STATUS, IERROR)
51     INTEGER REQUEST, STATUS(MPI_STATUS_SIZE), IERROR
52
53     MPI_WAITSOME(INCOUNT, ARRAY_OF_REQUESTS, OUTCOUNT, ARRAY_OF_INDICES,
54     ARRAY_OF_STATUSES, IERROR)

```

```

    INTEGER INCOUNT, ARRAY_OF_REQUESTS(*), OUTCOUNT, ARRAY_OF_INDICES(*),
    ARRAY_OF_STATUSES(MPI_STATUS_SIZE,*), IERROR

```

#### A.4.2 Datatypes Fortran Bindings

```

MPI_GET_ADDRESS(LOCATION, ADDRESS, IERROR)
    <type> LOCATION(*)
    INTEGER IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) ADDRESS

MPI_GET_ELEMENTS(STATUS, DATATYPE, COUNT, IERROR)
    INTEGER STATUS(MPI_STATUS_SIZE), DATATYPE, COUNT, IERROR

MPI_PACK_EXTERNAL(DATAREP, INBUF, INCOUNT, DATATYPE, OUTBUF, OUTSIZE,
    POSITION, IERROR)
    INTEGER INCOUNT, DATATYPE, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) OUTSIZE, POSITION
    CHARACTER*(*) DATAREP
    <type> INBUF(*), OUTBUF(*)

MPI_PACK_EXTERNAL_SIZE(DATAREP, INCOUNT, DATATYPE, SIZE, IERROR)
    INTEGER INCOUNT, DATATYPE, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) SIZE
    CHARACTER*(*) DATAREP

MPI_PACK(INBUF, INCOUNT, DATATYPE, OUTBUF, OUTSIZE, POSITION, COMM, IERROR)
    <type> INBUF(*), OUTBUF(*)
    INTEGER INCOUNT, DATATYPE, OUTSIZE, POSITION, COMM, IERROR

MPI_PACK_SIZE(INCOUNT, DATATYPE, COMM, SIZE, IERROR)
    INTEGER INCOUNT, DATATYPE, COMM, SIZE, IERROR

MPI_TYPE_COMMIT(DATATYPE, IERROR)
    INTEGER DATATYPE, IERROR

MPI_TYPE_CONTIGUOUS(COUNT, OLDTYPE, NEWTYPE, IERROR)
    INTEGER COUNT, OLDTYPE, NEWTYPE, IERROR

MPI_TYPE_CREATE_DARRAY(SIZE, RANK, NDIMS, ARRAY_OF_GSIZES,
    ARRAY_OF_DISTRIBS, ARRAY_OF_DARGS, ARRAY_OF_PSIZEs, ORDER,
    OLDTYPE, NEWTYPE, IERROR)
    INTEGER SIZE, RANK, NDIMS, ARRAY_OF_GSIZES(*), ARRAY_OF_DISTRIBS(*),
    ARRAY_OF_DARGS(*), ARRAY_OF_PSIZEs(*), ORDER, OLDTYPE, NEWTYPE, IERROR

MPI_TYPE_CREATE_HINDEXED(COUNT, ARRAY_OF_BLOCKLENGTHS,
    ARRAY_OF_DISPLACEMENTS, OLDTYPE, NEWTYPE, IERROR)
    INTEGER COUNT, ARRAY_OF_BLOCKLENGTHS(*), OLDTYPE, NEWTYPE, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) ARRAY_OF_DISPLACEMENTS(*)

MPI_TYPE_CREATE_HVECTOR(COUNT, BLOCKLENGTH, STRIDE, OLDTYPE, NEWTYPE,
    IERROR)
    INTEGER COUNT, BLOCKLENGTH, OLDTYPE, NEWTYPE, IERROR

```

```

1      INTEGER(KIND=MPI_ADDRESS_KIND) STRIDE
2
3      MPI_TYPE_CREATE_INDEXED_BLOCK(COUNT, BLOCKLENGTH, ARRAY_OF_DISPLACEMENTS,
4          OLDTYPE, NEWTYPE, IERROR)
5          INTEGER COUNT, BLOCKLENGTH, ARRAY_OF_DISPLACEMENTS(*), OLDTYPE,
6          NEWTYPE, IERROR
7
8      MPI_TYPE_CREATE_RESIZED(OLDTYPE, LB, EXTENT, NEWTYPE, IERROR)
9          INTEGER OLDTYPE, NEWTYPE, IERROR
10         INTEGER(KIND=MPI_ADDRESS_KIND) LB, EXTENT
11
12      MPI_TYPE_CREATE_STRUCT(COUNT, ARRAY_OF_BLOCKLENGTHS,
13          ARRAY_OF_DISPLACEMENTS, ARRAY_OF_TYPES, NEWTYPE, IERROR)
14          INTEGER COUNT, ARRAY_OF_BLOCKLENGTHS(*), ARRAY_OF_TYPES(*), NEWTYPE,
15          IERROR
16          INTEGER(KIND=MPI_ADDRESS_KIND) ARRAY_OF_DISPLACEMENTS(*)
17
18      MPI_TYPE_CREATE_SUBARRAY(NDIMS, ARRAY_OF_SIZES, ARRAY_OF_SUBSIZES,
19          ARRAY_OF_STARTS, ORDER, OLDTYPE, NEWTYPE, IERROR)
20          INTEGER NDIMS, ARRAY_OF_SIZES(*), ARRAY_OF_SUBSIZES(*),
21          ARRAY_OF_STARTS(*), ORDER, OLDTYPE, NEWTYPE, IERROR
22
23      MPI_TYPE_DUP(OLDTYPE, NEWTYPE, IERROR)
24          INTEGER OLDTYPE, NEWTYPE, IERROR
25
26      MPI_TYPE_FREE(DATATYPE, IERROR)
27          INTEGER DATATYPE, IERROR
28
29      MPI_TYPE_GET_CONTENTS(DATATYPE, MAX_INTEGERS, MAX_ADDRESSES, MAX_DATATYPES,
30          ARRAY_OF_INTEGERS, ARRAY_OF_ADDRESSES, ARRAY_OF_DATATYPES,
31          IERROR)
32          INTEGER DATATYPE, MAX_INTEGERS, MAX_ADDRESSES, MAX_DATATYPES,
33          ARRAY_OF_INTEGERS(*), ARRAY_OF_DATATYPES(*), IERROR
34          INTEGER(KIND=MPI_ADDRESS_KIND) ARRAY_OF_ADDRESSES(*)
35
36      MPI_TYPE_GET_ENVELOPE(DATATYPE, NUM_INTEGERS, NUM_ADDRESSES, NUM_DATATYPES,
37          COMBINER, IERROR)
38          INTEGER DATATYPE, NUM_INTEGERS, NUM_ADDRESSES, NUM_DATATYPES, COMBINER,
39          IERROR
40
41      MPI_TYPE_GET_EXTENT(DATATYPE, LB, EXTENT, IERROR)
42          INTEGER DATATYPE, IERROR
43          INTEGER(KIND = MPI_ADDRESS_KIND) LB, EXTENT
44
45      MPI_TYPE_GET_TRUE_EXTENT(DATATYPE, TRUE_LB, TRUE_EXTENT, IERROR)
46          INTEGER DATATYPE, IERROR
47          INTEGER(KIND = MPI_ADDRESS_KIND) TRUE_LB, TRUE_EXTENT
48
49      MPI_TYPE_INDEXED(COUNT, ARRAY_OF_BLOCKLENGTHS, ARRAY_OF_DISPLACEMENTS,
50          OLDTYPE, NEWTYPE, IERROR)
51          INTEGER COUNT, ARRAY_OF_BLOCKLENGTHS(*), ARRAY_OF_DISPLACEMENTS(*),
52          OLDTYPE, NEWTYPE, IERROR

```

```

MPI_TYPE_SIZE(DATATYPE, SIZE, IERROR)
    INTEGER DATATYPE, SIZE, IERROR
MPI_TYPE_VECTOR(COUNT, BLOCKLENGTH, STRIDE, OLDTYPE, NEWTYPE, IERROR)
    INTEGER COUNT, BLOCKLENGTH, STRIDE, OLDTYPE, NEWTYPE, IERROR
MPI_UNPACK_EXTERNAL(DATAREP, INBUF, INSIZE, POSITION, OUTBUF, OUTCOUNT,
    DATATYPE, IERROR)
    INTEGER OUTCOUNT, DATATYPE, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) INSIZE, POSITION
    CHARACTER*(*) DATAREP
    <type> INBUF(*), OUTBUF(*)
MPI_UNPACK(INBUF, INSIZE, POSITION, OUTBUF, OUTCOUNT, DATATYPE, COMM,
    IERROR)
    <type> INBUF(*), OUTBUF(*)
    INTEGER INSIZE, POSITION, OUTCOUNT, DATATYPE, COMM, IERROR

```

#### A.4.3 Collective Communication Fortran Bindings

```

MPI_ALLGATHER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, REVCOUNT, RECVTYPE,
    COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNT, SENDTYPE, REVCOUNT, RECVTYPE, COMM, IERROR
MPI_ALLGATHERV(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, REVCOUNTS, DISPLS,
    RECVTYPE, COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNT, SENDTYPE, REVCOUNTS(*), DISPLS(*), RECVTYPE, COMM,
    IERROR
MPI_ALLREDUCE(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER COUNT, DATATYPE, OP, COMM, IERROR
MPI_ALLTOALL(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, REVCOUNT, RECVTYPE,
    COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNT, SENDTYPE, REVCOUNT, RECVTYPE, COMM, IERROR
MPI_ALLTOALLV(SENDBUF, SENDCOUNTS, SDISPLS, SENDTYPE, RECVBUF, REVCOUNTS,
    RDISPLS, RECVTYPE, COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNTS(*), SDISPLS(*), SENDTYPE, REVCOUNTS(*), RDISPLS(*),
    RECVTYPE, COMM, IERROR
MPI_ALLTOALLW(SENDBUF, SENDCOUNTS, SDISPLS, SENDTYPES, RECVBUF, REVCOUNTS,
    RDISPLS, RECVTYPES, COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNTS(*), SDISPLS(*), SENDTYPES(*), REVCOUNTS(*),
    RDISPLS(*), RECVTYPES(*), COMM, IERROR

```

```

1  MPI_BARRIER(COMM, IERROR)
2      INTEGER COMM, IERROR
3
4  MPI_BCAST(BUFFER, COUNT, DATATYPE, ROOT, COMM, IERROR)
5      <type> BUFFER(*)
6      INTEGER COUNT, DATATYPE, ROOT, COMM, IERROR
7
8  MPI_EXSCAN(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, COMM, IERROR)
9      <type> SENDBUF(*), RECVBUF(*)
10     INTEGER COUNT, DATATYPE, OP, COMM, IERROR
11
12 MPI_GATHER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, REVCOUNT, RECVTYPE,
13           ROOT, COMM, IERROR)
14     <type> SENDBUF(*), RECVBUF(*)
15     INTEGER SENDCOUNT, SENDTYPE, REVCOUNT, RECVTYPE, ROOT, COMM, IERROR
16
17 MPI_GATHERV(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, REVCOUNTS, DISPLS,
18           RECVTYPE, ROOT, COMM, IERROR)
19     <type> SENDBUF(*), RECVBUF(*)
20     INTEGER SENDCOUNT, SENDTYPE, REVCOUNTS(*), DISPLS(*), RECVTYPE, ROOT,
21     COMM, IERROR
22
23 MPI_IALLGATHER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, REVCOUNT, RECVTYPE,
24           COMM, REQUEST, IERROR)
25     <type> SENDBUF(*), RECVBUF(*)
26     INTEGER SENDCOUNT, SENDTYPE, REVCOUNT, RECVTYPE, COMM, REQUEST, IERROR
27
28 MPI_IALLGATHERV(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, REVCOUNTS, DISPLS,
29           RECVTYPE, COMM, REQUEST, IERROR)
30     <type> SENDBUF(*), RECVBUF(*)
31     INTEGER SENDCOUNT, SENDTYPE, REVCOUNTS(*), DISPLS(*), RECVTYPE, COMM,
32     REQUEST, IERROR
33
34 MPI_IALLREDUCE(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, COMM, REQUEST,
35           IERROR)
36     <type> SENDBUF(*), RECVBUF(*)
37     INTEGER COUNT, DATATYPE, OP, COMM, REQUEST, IERROR
38
39 MPI_IALLTOALL(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, REVCOUNT, RECVTYPE,
40           COMM, REQUEST, IERROR)
41     <type> SENDBUF(*), RECVBUF(*)
42     INTEGER SENDCOUNT, SENDTYPE, REVCOUNT, RECVTYPE, COMM, REQUEST, IERROR
43
44 MPI_IALLTOALLV(SENDBUF, SENDCOUNTS, SDISPLS, SENDTYPE, RECVBUF, REVCOUNTS,
45           RDISPLS, RECVTYPE, COMM, REQUEST, IERROR)
46     <type> SENDBUF(*), RECVBUF(*)
47     INTEGER SENDCOUNTS(*), SDISPLS(*), SENDTYPE, REVCOUNTS(*), RDISPLS(*),
48     RECVTYPE, COMM, REQUEST, IERROR

```



```

    INTEGER SENDCOUNTS(*), SDISPLS(*), SENDTYPES(*), RECVCOUNTS(*),
    RDISPLS(*), RECVTYPES(*), COMM, REQUEST, IERROR
1
2
3
MPI_IBARRIER(COMM, REQUEST, IERROR)
4
    INTEGER COMM, REQUEST, IERROR
5
MPI_IBCAST(BUFFER, COUNT, DATATYPE, ROOT, COMM, REQUEST, IERROR)
6
    <type> BUFFER(*)
7
    INTEGER COUNT, DATATYPE, ROOT, COMM, REQUEST, IERROR
8
9
MPI_IEXSCAN(SENDBUF, RECVBUFF, COUNT, DATATYPE, OP, COMM, REQUEST, IERROR)
10
    <type> SENDBUF(*), RECVBUFF(*)
11
    INTEGER COUNT, DATATYPE, OP, COMM, REQUEST, IERROR
12
13
MPI_IGATHER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUFF, RECVCOUNT, RECVTYPE,
14
    ROOT, COMM, REQUEST, IERROR)
15
    <type> SENDBUF(*), RECVBUFF(*)
16
    INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, ROOT, COMM, REQUEST,
17
    IERROR
18
19
MPI_IGATHERV(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUFF, RECVCOUNTS, DISPLS,
20
    RECVTYPE, ROOT, COMM, REQUEST, IERROR)
21
    <type> SENDBUF(*), RECVBUFF(*)
22
    INTEGER SENDCOUNT, SENDTYPE, RECVCOUNTS(*), DISPLS(*), RECVTYPE, ROOT,
23
    COMM, REQUEST, IERROR
24
25
MPI_IREDUCE_SCATTER_BLOCK(SENDBUF, RECVBUFF, RECVCOUNT, DATATYPE, OP, COMM,
26
    REQUEST, IERROR)
27
    <type> SENDBUF(*), RECVBUFF(*)
28
    INTEGER RECVCOUNT, DATATYPE, OP, COMM, REQUEST, IERROR
29
30
MPI_IREDUCE_SCATTER(SENDBUF, RECVBUFF, RECVCOUNTS, DATATYPE, OP, COMM,
31
    REQUEST, IERROR)
32
    <type> SENDBUF(*), RECVBUFF(*)
33
    INTEGER RECVCOUNTS(*), DATATYPE, OP, COMM, REQUEST, IERROR
34
35
MPI_IREDUCE(SENDBUF, RECVBUFF, COUNT, DATATYPE, OP, ROOT, COMM, REQUEST,
36
    IERROR)
37
    <type> SENDBUF(*), RECVBUFF(*)
38
    INTEGER COUNT, DATATYPE, OP, ROOT, COMM, REQUEST, IERROR
39
40
MPI_ISCAN(SENDBUF, RECVBUFF, COUNT, DATATYPE, OP, COMM, REQUEST, IERROR)
41
    <type> SENDBUF(*), RECVBUFF(*)
42
    INTEGER COUNT, DATATYPE, OP, COMM, REQUEST, IERROR
43
44
MPI_ISCATTER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUFF, RECVCOUNT, RECVTYPE,
45
    ROOT, COMM, REQUEST, IERROR)
46
    <type> SENDBUF(*), RECVBUFF(*)
47
    INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, ROOT, COMM, REQUEST,
48
    IERROR

```

```

1      <type> SENDBUF(*), RECVBUF(*)
2      INTEGER SENDCOUNTS(*), DISPLS(*), SENDTYPE, RECVCOUNT, RECVTYPE, ROOT,
3      COMM, REQUEST, IERROR
4
5      MPI_OP_COMMUTATIVE(OP, COMMUTE, IERROR)
6      LOGICAL COMMUTE
7      INTEGER OP, IERROR
8
9      MPI_OP_CREATE( [FUNCTION]USER_FN, COMMUTE, OP, IERROR)
10     EXTERNAL [FUNCTION]USER_FN
11     LOGICAL COMMUTE
12     INTEGER OP, IERROR
13
14     MPI_OP_FREE(OP, IERROR)
15     INTEGER OP, IERROR
16
17     MPI_REDUCE_LOCAL(INBUF, INOUTBUF, COUNT, DATATYPE, OP, IERROR)
18     <type> INBUF(*), INOUTBUF(*)
19     INTEGER COUNT, DATATYPE, OP, IERROR
20
21     MPI_REDUCE_SCATTER_BLOCK(SENDBUF, RECVBUF, RECVCOUNT, DATATYPE, OP, COMM,
22     IERROR)
23     <type> SENDBUF(*), RECVBUF(*)
24     INTEGER RECVCOUNT, DATATYPE, OP, COMM, IERROR
25
26     MPI_REDUCE_SCATTER(SENDBUF, RECVBUF, REVCOUNTS, DATATYPE, OP, COMM,
27     IERROR)
28     <type> SENDBUF(*), RECVBUF(*)
29     INTEGER REVCOUNTS(*), DATATYPE, OP, COMM, IERROR
30
31     MPI_REDUCE(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, ROOT, COMM, IERROR)
32     <type> SENDBUF(*), RECVBUF(*)
33     INTEGER COUNT, DATATYPE, OP, ROOT, COMM, IERROR
34
35     MPI_SCAN(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, COMM, IERROR)
36     <type> SENDBUF(*), RECVBUF(*)
37     INTEGER COUNT, DATATYPE, OP, COMM, IERROR
38
39     MPI_SCATTER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT, RECVTYPE,
40     ROOT, COMM, IERROR)
41     <type> SENDBUF(*), RECVBUF(*)
42     INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, ROOT, COMM, IERROR
43
44     MPI_SCATTERV(SENDBUF, SENDCOUNTS, DISPLS, SENDTYPE, RECVBUF, RECVCOUNT,
45     RECVTYPE, ROOT, COMM, IERROR)
46     <type> SENDBUF(*), RECVBUF(*)
47     INTEGER SENDCOUNTS(*), DISPLS(*), SENDTYPE, RECVCOUNT, RECVTYPE, ROOT,
48     COMM, IERROR

```

A.4.4 Groups, Contexts, Communicators, and Caching Fortran Bindings

```

MPI_COMM_COMPARE(COMM1, COMM2, RESULT, IERROR)

```

```

    INTEGER COMM1, COMM2, RESULT, IERROR
1
MPI_COMM_CREATE(COMM, GROUP, NEWCOMM, IERROR)
2
    INTEGER COMM, GROUP, NEWCOMM, IERROR
3
MPI_COMM_CREATE_KEYVAL(COMM_COPY_ATTR_FN, COMM_DELETE_ATTR_FN, COMM_KEYVAL,
4
    EXTRA_STATE, IERROR)
5
    EXTERNAL COMM_COPY_ATTR_FN, COMM_DELETE_ATTR_FN
6
    INTEGER COMM_KEYVAL, IERROR
7
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE
8
MPI_COMM_DELETE_ATTR(COMM, COMM_KEYVAL, IERROR)
9
    INTEGER COMM, COMM_KEYVAL, IERROR
10
MPI_COMM_DUP(COMM, NEWCOMM, IERROR)
11
    INTEGER COMM, NEWCOMM, IERROR
12
MPI_COMM_DUP_FN(OLDCOMM, COMM_KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,
13
    ATTRIBUTE_VAL_OUT, FLAG, IERROR)
14
    INTEGER OLDCOMM, COMM_KEYVAL, IERROR
15
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE, ATTRIBUTE_VAL_IN,
16
    ATTRIBUTE_VAL_OUT
17
    LOGICAL FLAG
18
MPI_COMM_FREE(COMM, IERROR)
19
    INTEGER COMM, IERROR
20
MPI_COMM_FREE_KEYVAL(COMM_KEYVAL, IERROR)
21
    INTEGER COMM_KEYVAL, IERROR
22
MPI_COMM_GET_ATTR(COMM, COMM_KEYVAL, ATTRIBUTE_VAL, FLAG, IERROR)
23
    INTEGER COMM, COMM_KEYVAL, IERROR
24
    INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL
25
    LOGICAL FLAG
26
MPI_COMM_GET_NAME(COMM, COMM_NAME, RESULTLEN, IERROR)
27
    INTEGER COMM, RESULTLEN, IERROR
28
    CHARACTER*(*) COMM_NAME
29
MPI_COMM_GROUP(COMM, GROUP, IERROR)
30
    INTEGER COMM, GROUP, IERROR
31
MPI_COMM_NULL_COPY_FN(OLDCOMM, COMM_KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,
32
    ATTRIBUTE_VAL_OUT, FLAG, IERROR)
33
    INTEGER OLDCOMM, COMM_KEYVAL, IERROR
34
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE, ATTRIBUTE_VAL_IN,
35
    ATTRIBUTE_VAL_OUT
36
    LOGICAL FLAG
37
MPI_COMM_NULL_DELETE_FN(COMM, COMM_KEYVAL, ATTRIBUTE_VAL, EXTRA_STATE,
38
    IERROR)
39
    INTEGER COMM, COMM_KEYVAL, IERROR
40
    INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL, EXTRA_STATE
41
42
43
44
45
46
47
48

```

```

1  MPI_COMM_RANK(COMM, RANK, IERROR)
2      INTEGER COMM, RANK, IERROR
3
4  MPI_COMM_REMOTE_GROUP(COMM, GROUP, IERROR)
5      INTEGER COMM, GROUP, IERROR
6
7  MPI_COMM_REMOTE_SIZE(COMM, SIZE, IERROR)
8      INTEGER COMM, SIZE, IERROR
9
10 MPI_COMM_SET_ATTR(COMM, COMM_KEYVAL, ATTRIBUTE_VAL, IERROR)
11     INTEGER COMM, COMM_KEYVAL, IERROR
12     INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL
13
14 MPI_COMM_SET_NAME(COMM, COMM_NAME, IERROR)
15     INTEGER COMM, IERROR
16     CHARACTER*(*) COMM_NAME
17
18 MPI_COMM_SIZE(COMM, SIZE, IERROR)
19     INTEGER COMM, SIZE, IERROR
20
21 MPI_COMM_SPLIT(COMM, COLOR, KEY, NEWCOMM, IERROR)
22     INTEGER COMM, COLOR, KEY, NEWCOMM, IERROR
23
24 MPI_COMM_TEST_INTER(COMM, FLAG, IERROR)
25     INTEGER COMM, IERROR
26     LOGICAL FLAG
27
28 MPI_GROUP_COMPARE(GROUP1, GROUP2, RESULT, IERROR)
29     INTEGER GROUP1, GROUP2, RESULT, IERROR
30
31 MPI_GROUP_DIFFERENCE(GROUP1, GROUP2, NEWGROUP, IERROR)
32     INTEGER GROUP1, GROUP2, NEWGROUP, IERROR
33
34 MPI_GROUP_EXCL(GROUP, N, RANKS, NEWGROUP, IERROR)
35     INTEGER GROUP, N, RANKS(*), NEWGROUP, IERROR
36
37 MPI_GROUP_FREE(GROUP, IERROR)
38     INTEGER GROUP, IERROR
39
40 MPI_GROUP_INCL(GROUP, N, RANKS, NEWGROUP, IERROR)
41     INTEGER GROUP, N, RANKS(*), NEWGROUP, IERROR
42
43 MPI_GROUP_INTERSECTION(GROUP1, GROUP2, NEWGROUP, IERROR)
44     INTEGER GROUP1, GROUP2, NEWGROUP, IERROR
45
46 MPI_GROUP_RANGE_EXCL(GROUP, N, RANGES, NEWGROUP, IERROR)
47     INTEGER GROUP, N, RANGES(3,*), NEWGROUP, IERROR
48
49 MPI_GROUP_RANGE_INCL(GROUP, N, RANGES, NEWGROUP, IERROR)
50     INTEGER GROUP, N, RANGES(3,*), NEWGROUP, IERROR
51
52 MPI_GROUP_RANK(GROUP, RANK, IERROR)
53     INTEGER GROUP, RANK, IERROR
54
55 MPI_GROUP_SIZE(GROUP, SIZE, IERROR)

```

```

    INTEGER GROUP, SIZE, IERROR
1
MPI_GROUP_TRANSLATE_RANKS(GROUP1, N, RANKS1, GROUP2, RANKS2, IERROR)
2
    INTEGER GROUP1, N, RANKS1(*), GROUP2, RANKS2(*), IERROR
3
4
MPI_GROUP_UNION(GROUP1, GROUP2, NEWGROUP, IERROR)
5
    INTEGER GROUP1, GROUP2, NEWGROUP, IERROR
6
7
MPI_INTERCOMM_CREATE(LOCAL_COMM, LOCAL_LEADER, PEER_COMM, REMOTE_LEADER,
8
    TAG, NEWINTERCOMM, IERROR)
9
    INTEGER LOCAL_COMM, LOCAL_LEADER, PEER_COMM, REMOTE_LEADER, TAG,
10
    NEWINTERCOMM, IERROR
11
MPI_INTERCOMM_MERGE(INTERCOMM, HIGH, NEWINTRACOMM, IERROR)
12 ticket250-V.
    INTEGER INTERCOMM, NEWINTRACOMM, IERROR
13 ticket250-V.
    LOGICAL HIGH
14
15
MPI_TYPE_CREATE_KEYVAL(TYPE_COPY_ATTR_FN, TYPE_DELETE_ATTR_FN, TYPE_KEYVAL,
16
    EXTRA_STATE, IERROR)
17
    EXTERNAL TYPE_COPY_ATTR_FN, TYPE_DELETE_ATTR_FN
18
    INTEGER TYPE_KEYVAL, IERROR
19
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE
20
21
MPI_TYPE_DELETE_ATTR(DATATYPE, TYPE_KEYVAL, IERROR)
21 ticket252-W.
    INTEGER DATATYPE, TYPE_KEYVAL, IERROR
22 ticket252-W.
23
24
MPI_TYPE_DUP_FN(OLDTYPE, TYPE_KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,
24
    ATTRIBUTE_VAL_OUT, FLAG, IERROR)
25
    INTEGER OLDTYPE, TYPE_KEYVAL, IERROR
26
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE, ATTRIBUTE_VAL_IN,
27
    ATTRIBUTE_VAL_OUT
28
    LOGICAL FLAG
29
30
MPI_TYPE_FREE_KEYVAL(TYPE_KEYVAL, IERROR)
30
    INTEGER TYPE_KEYVAL, IERROR
31
32
MPI_TYPE_GET_ATTR(DATATYPE, TYPE_KEYVAL, ATTRIBUTE_VAL, FLAG, IERROR)
33 ticket252-W.
    INTEGER DATATYPE, TYPE_KEYVAL, IERROR
34 ticket252-W.
    INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL
35
    LOGICAL FLAG
36
37
MPI_TYPE_GET_NAME(DATATYPE, TYPE_NAME, RESULTLEN, IERROR)
37 ticket252-W.
    INTEGER DATATYPE, RESULTLEN, IERROR
38 ticket252-W.
    CHARACTER*(*) TYPE_NAME
39
40
MPI_TYPE_NULL_COPY_FN(OLDTYPE, TYPE_KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,
41
    ATTRIBUTE_VAL_OUT, FLAG, IERROR)
42
    INTEGER OLDTYPE, TYPE_KEYVAL, IERROR
43
    INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE, ATTRIBUTE_VAL_IN,
44
    ATTRIBUTE_VAL_OUT
45
    LOGICAL FLAG
46
47
MPI_TYPE_NULL_DELETE_FN(DATATYPE, TYPE_KEYVAL, ATTRIBUTE_VAL, EXTRA_STATE,
47 ticket252-W.
48

```

```

1          IERROR)
ticket252-W. 2      INTEGER DATATYPE, TYPE_KEYVAL, IERROR
3      INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL, EXTRA_STATE
4
ticket252-W. 5      MPI_TYPE_SET_ATTR(DATATYPE, TYPE_KEYVAL, ATTRIBUTE_VAL, IERROR)
ticket252-W. 6      INTEGER DATATYPE, TYPE_KEYVAL, IERROR
7      INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL
8
ticket252-W. 8      MPI_TYPE_SET_NAME(DATATYPE, TYPE_NAME, IERROR)
ticket252-W. 9      INTEGER DATATYPE, IERROR
10     CHARACTER*(*) TYPE_NAME
11
12     MPI_WIN_CREATE_KEYVAL(WIN_COPY_ATTR_FN, WIN_DELETE_ATTR_FN, WIN_KEYVAL,
13         EXTRA_STATE, IERROR)
14     EXTERNAL WIN_COPY_ATTR_FN, WIN_DELETE_ATTR_FN
15     INTEGER WIN_KEYVAL, IERROR
16     INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE
17
18     MPI_WIN_DELETE_ATTR(WIN, WIN_KEYVAL, IERROR)
19     INTEGER WIN, WIN_KEYVAL, IERROR
20
21     MPI_WIN_DUP_FN(OLDWIN, WIN_KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,
22         ATTRIBUTE_VAL_OUT, FLAG, IERROR)
23     INTEGER OLDWIN, WIN_KEYVAL, IERROR
24     INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE, ATTRIBUTE_VAL_IN,
25         ATTRIBUTE_VAL_OUT
26     LOGICAL FLAG
27
28     MPI_WIN_FREE_KEYVAL(WIN_KEYVAL, IERROR)
29     INTEGER WIN_KEYVAL, IERROR
30
31     MPI_WIN_GET_ATTR(WIN, WIN_KEYVAL, ATTRIBUTE_VAL, FLAG, IERROR)
32     INTEGER WIN, WIN_KEYVAL, IERROR
33     INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL
34     LOGICAL FLAG
35
36     MPI_WIN_GET_NAME(WIN, WIN_NAME, RESULTLEN, IERROR)
37     INTEGER WIN, RESULTLEN, IERROR
38     CHARACTER*(*) WIN_NAME
39
40     MPI_WIN_NULL_COPY_FN(OLDWIN, WIN_KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,
41         ATTRIBUTE_VAL_OUT, FLAG, IERROR)
42     INTEGER OLDWIN, WIN_KEYVAL, IERROR
43     INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE, ATTRIBUTE_VAL_IN,
44         ATTRIBUTE_VAL_OUT
45     LOGICAL FLAG
46
47     MPI_WIN_NULL_DELETE_FN(WIN, WIN_KEYVAL, ATTRIBUTE_VAL, EXTRA_STATE, IERROR)
48     INTEGER WIN, WIN_KEYVAL, IERROR
49     INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL, EXTRA_STATE
50
51     MPI_WIN_SET_ATTR(WIN, WIN_KEYVAL, ATTRIBUTE_VAL, IERROR)
52     INTEGER WIN, WIN_KEYVAL, IERROR

```

```

    INTEGER(KIND=MPI_ADDRESS_KIND) ATTRIBUTE_VAL
1
MPI_WIN_SET_NAME(WIN, WIN_NAME, IERROR)
2
    INTEGER WIN, IERROR
3
    CHARACTER*(*) WIN_NAME
4
5
6
A.4.5 Process Topologies Fortran Bindings
7
MPI_CART_COORDS(COMM, RANK, MAXDIMS, COORDS, IERROR)
8
    INTEGER COMM, RANK, MAXDIMS, COORDS(*), IERROR
9
10
MPI_CART_CREATE(COMM_OLD, NDIMS, DIMS, PERIODS, REORDER, COMM_CART, IERROR)
11
    INTEGER COMM_OLD, NDIMS, DIMS(*), COMM_CART, IERROR
12
    LOGICAL PERIODS(*), REORDER
13
14
MPI_CARTDIM_GET(COMM, NDIMS, IERROR)
15
    INTEGER COMM, NDIMS, IERROR
16
17
MPI_CART_GET(COMM, MAXDIMS, DIMS, PERIODS, COORDS, IERROR)
18
    INTEGER COMM, MAXDIMS, DIMS(*), COORDS(*), IERROR
19
    LOGICAL PERIODS(*)
20
21
MPI_CART_MAP(COMM, NDIMS, DIMS, PERIODS, NEWRANK, IERROR)
22
    INTEGER COMM, NDIMS, DIMS(*), NEWRANK, IERROR
23
    LOGICAL PERIODS(*)
24
25
MPI_CART_RANK(COMM, COORDS, RANK, IERROR)
26
    INTEGER COMM, COORDS(*), RANK, IERROR
27
28
MPI_CART_SHIFT(COMM, DIRECTION, DISP, RANK_SOURCE, RANK_DEST, IERROR)
29
    INTEGER COMM, DIRECTION, DISP, RANK_SOURCE, RANK_DEST, IERROR
30
31
MPI_CART_SUB(COMM, REMAIN_DIMS, NEWCOMM, IERROR)
32
    INTEGER COMM, NEWCOMM, IERROR
33
    LOGICAL REMAIN_DIMS(*)
34
35
MPI_DIMS_CREATE(NNODES, NDIMS, DIMS, IERROR)
36
    INTEGER NNODES, NDIMS, DIMS(*), IERROR
37
38
MPI_DIST_GRAPH_CREATE_ADJACENT(COMM_OLD, INDEGREE, SOURCES, SOURCEWEIGHTS,
39
    OUTDEGREE, DESTINATIONS, DESTWEIGHTS, INFO, REORDER,
40
    COMM_DIST_GRAPH, IERROR)
41
    INTEGER COMM_OLD, INDEGREE, SOURCES(*), SOURCEWEIGHTS(*), OUTDEGREE,
42
    DESTINATIONS(*), DESTWEIGHTS(*), INFO, COMM_DIST_GRAPH, IERROR
43
    LOGICAL REORDER
44
45
MPI_DIST_GRAPH_CREATE(COMM_OLD, N, SOURCES, DEGREES, DESTINATIONS, WEIGHTS,
46
    INFO, REORDER, COMM_DIST_GRAPH, IERROR)
47
    INTEGER COMM_OLD, N, SOURCES(*), DEGREES(*), DESTINATIONS(*),
48
    WEIGHTS(*), INFO, COMM_DIST_GRAPH, IERROR
49
    LOGICAL REORDER
50

```

```

1  MPI_DIST_GRAPH_NEIGHBORS(COMM, MAXINDEGREE, SOURCES, SOURCEWEIGHTS,
2      MAXOUTDEGREE, DESTINATIONS, DESTWEIGHTS, IERROR)
3      INTEGER COMM, MAXINDEGREE, SOURCES(*), SOURCEWEIGHTS(*), MAXOUTDEGREE,
4      DESTINATIONS(*), DESTWEIGHTS(*), IERROR
5
6  MPI_DIST_GRAPH_NEIGHBORS_COUNT(COMM, INDEGREE, OUTDEGREE, WEIGHTED, IERROR)
7      INTEGER COMM, INDEGREE, OUTDEGREE, IERROR
8      LOGICAL WEIGHTED
9
10 MPI_GRAPH_CREATE(COMM_OLD, NNODES, INDEX, EDGES, REORDER, COMM_GRAPH,
11      IERROR)
12     INTEGER COMM_OLD, NNODES, INDEX(*), EDGES(*), COMM_GRAPH, IERROR
13     LOGICAL REORDER
14
15 MPI_GRAPHDIMS_GET(COMM, NNODES, NEDGES, IERROR)
16     INTEGER COMM, NNODES, NEDGES, IERROR
17
18 MPI_GRAPH_GET(COMM, MAXINDEX, MAXEDGES, INDEX, EDGES, IERROR)
19     INTEGER COMM, MAXINDEX, MAXEDGES, INDEX(*), EDGES(*), IERROR
20
21 MPI_GRAPH_MAP(COMM, NNODES, INDEX, EDGES, NEWRANK, IERROR)
22     INTEGER COMM, NNODES, INDEX(*), EDGES(*), NEWRANK, IERROR
23
24 MPI_GRAPH_NEIGHBORS(COMM, RANK, MAXNEIGHBORS, NEIGHBORS, IERROR)
25     INTEGER COMM, RANK, MAXNEIGHBORS, NEIGHBORS(*), IERROR
26
27 MPI_GRAPH_NEIGHBORS_COUNT(COMM, RANK, NNEIGHBORS, IERROR)
28     INTEGER COMM, RANK, NNEIGHBORS, IERROR
29
30 MPI_INEIGHBOR_ALLGATHER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT,
31     RECVTYPE, COMM, REQUEST, IERROR)
32     <type> SENDBUF(*), RECVBUF(*)
33     INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, COMM, REQUEST, IERROR
34
35 MPI_INEIGHBOR_ALLGATHERV(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNTS,
36     DISPLS, RECVTYPE, COMM, REQUEST, IERROR)
37     <type> SENDBUF(*), RECVBUF(*)
38     INTEGER SENDCOUNT, SENDTYPE, RECVCOUNTS(*), DISPLS(*), RECVTYPE, COMM,
39     REQUEST, IERROR
40
41 MPI_INEIGHBOR_ALLTOALL(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT,
42     RECVTYPE, COMM, REQUEST, IERROR)
43     <type> SENDBUF(*), RECVBUF(*)
44     INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, COMM, REQUEST, IERROR
45
46 MPI_INEIGHBOR_ALLTOALLV(SENDBUF, SENDCOUNTS, SDISPLS, SENDTYPE, RECVBUF,
47     RECVCOUNTS, RDISPLS, RECVTYPE, COMM, REQUEST, IERROR)
48     <type> SENDBUF(*), RECVBUF(*)
49     INTEGER SENDCOUNTS(*), SDISPLS(*), SENDTYPE, RECVCOUNTS(*), RDISPLS(*),
50     RECVTYPE, COMM, REQUEST, IERROR

```



```

    <type> SENDBUF(*), RECVBUF(*)
    INTEGER(KIND=MPI_ADDRESS_KIND) SDISPLS(*), RDISPLS(*)
    INTEGER SENDCOUNTS(*), SENDTYPES(*), RECVCOUNTS(*), RECVTYPES(*), COMM,
    REQUEST, IERROR
MPI_NEIGHBOR_ALLGATHER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT,
    RECVTYPE, COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, COMM, IERROR
MPI_NEIGHBOR_ALLGATHERV(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNTS,
    DISPLS, RECVTYPE, COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNT, SENDTYPE, RECVCOUNTS(*), DISPLS(*), RECVTYPE, COMM,
    IERROR
MPI_NEIGHBOR_ALLTOALL(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT,
    RECVTYPE, COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, COMM, IERROR
MPI_NEIGHBOR_ALLTOALLV(SENDBUF, SENDCOUNTS, SDISPLS, SENDTYPE, RECVBUF,
    RECVCOUNTS, RDISPLS, RECVTYPE, COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER SENDCOUNTS(*), SDISPLS(*), SENDTYPE, RECVCOUNTS(*), RDISPLS(*),
    RECVTYPE, COMM, IERROR
MPI_NEIGHBOR_ALLTOALLW(SENDBUF, SENDCOUNTS, SDISPLS, SENDTYPES, RECVBUF,
    RECVCOUNTS, RDISPLS, RECVTYPES, COMM, IERROR)
    <type> SENDBUF(*), RECVBUF(*)
    INTEGER(KIND=MPI_ADDRESS_KIND) SDISPLS(*), RDISPLS(*)
    INTEGER SENDCOUNTS(*), SENDTYPES(*), RECVCOUNTS(*), RECVTYPES(*), COMM,
    IERROR
MPI_TOPO_TEST(COMM, STATUS, IERROR)
    INTEGER COMM, STATUS, IERROR

```

#### A.4.6 MPI Environmental Management Fortran Bindings

```

DOUBLE PRECISION MPI_WTICK()
DOUBLE PRECISION MPI_WTIME()
MPI_ABORT(COMM, ERRORCODE, IERROR)
    INTEGER COMM, ERRORCODE, IERROR
MPI_ADD_ERROR_CLASS(ERRORCLASS, IERROR)
    INTEGER ERRORCLASS, IERROR
MPI_ADD_ERROR_CODE(ERRORCLASS, ERRORCODE, IERROR)
    INTEGER ERRORCLASS, ERRORCODE, IERROR
MPI_ADD_ERROR_STRING(ERRORCODE, STRING, IERROR)

```

```

1      INTEGER ERRORCODE, IERROR
2      CHARACTER*(*) STRING
3
4      MPI_ALLOC_MEM(SIZE, INFO, BASEPTR, IERROR)
5          INTEGER INFO, IERROR
6          INTEGER(KIND=MPI_ADDRESS_KIND) SIZE, BASEPTR
7
8      MPI_COMM_CALL_ERRHANDLER(COMM, ERRORCODE, IERROR)
9          INTEGER COMM, ERRORCODE, IERROR
10
11      MPI_COMM_CREATE_ERRHANDLER([FUNCTION] COMM_ERRHANDLER_FN, ERRHANDLER,
12          IERROR)
13
14      MPI_COMM_GET_ERRHANDLER(COMM, ERRHANDLER, IERROR)
15          INTEGER COMM, ERRHANDLER, IERROR
16
17      MPI_COMM_SET_ERRHANDLER(COMM, ERRHANDLER, IERROR)
18          INTEGER COMM, ERRHANDLER, IERROR
19
20      MPI_ERRHANDLER_FREE(ERRHANDLER, IERROR)
21          INTEGER ERRHANDLER, IERROR
22
23      MPI_ERROR_CLASS(ERRORCODE, ERRORCLASS, IERROR)
24          INTEGER ERRORCODE, ERRORCLASS, IERROR
25
26      MPI_ERROR_STRING(ERRORCODE, STRING, RESULTLEN, IERROR)
27          INTEGER ERRORCODE, RESULTLEN, IERROR
28          CHARACTER*(*) STRING
29
30      MPI_FILE_CALL_ERRHANDLER(FH, ERRORCODE, IERROR)
31          INTEGER FH, ERRORCODE, IERROR
32
33      MPI_FILE_CREATE_ERRHANDLER([FUNCTION] FILE_ERRHANDLER_FN, ERRHANDLER,
34          IERROR)
35
36      MPI_FILE_GET_ERRHANDLER(FILE, ERRHANDLER, IERROR)
37          INTEGER FILE, ERRHANDLER, IERROR
38
39      MPI_FILE_SET_ERRHANDLER(FILE, ERRHANDLER, IERROR)
40          INTEGER FILE, ERRHANDLER, IERROR
41
42      MPI_FINALIZED(FLAG, IERROR)
43          LOGICAL FLAG
44          INTEGER IERROR
45
46      MPI_FINALIZE(IERROR)
47          INTEGER IERROR
48
49      MPI_FREE_MEM(BASE, IERROR)
50          <type> BASE(*)
51          INTEGER IERROR

```

MPI_GET_PROCESSOR_NAME( NAME, RESULTLEN, IERROR)	1
CHARACTER*(*) NAME	2
INTEGER RESULTLEN, IERROR	3
	4
MPI_GET_VERSION(VERSION, SUBVERSION, IERROR)	5
INTEGER VERSION, SUBVERSION, IERROR	6
	7
MPI_INITIALIZED(FLAG, IERROR)	8
LOGICAL FLAG	9
INTEGER IERROR	10
	11
MPI_INIT(IERROR)	12
INTEGER IERROR	13
	14
MPI_WIN_CALL_ERRHANDLER(WIN, ERRORCODE, IERROR)	15
INTEGER WIN, ERRORCODE, IERROR	16
	17
MPI_WIN_CREATE_ERRHANDLER([FUNCTION]WIN_ERRHANDLER_FN, ERRHANDLER, IERROR)	18
EXTERNAL [FUNCTION]WIN_ERRHANDLER_FN	19
INTEGER ERRHANDLER, IERROR	20
	21
MPI_WIN_GET_ERRHANDLER(WIN, ERRHANDLER, IERROR)	22
INTEGER WIN, ERRHANDLER, IERROR	23
	24
MPI_WIN_SET_ERRHANDLER(WIN, ERRHANDLER, IERROR)	25
INTEGER WIN, ERRHANDLER, IERROR	26
	27
	28
	29
A.4.7 The Info Object Fortran Bindings	30
	31
	32
MPI_INFO_CREATE(INFO, IERROR)	33
INTEGER INFO, IERROR	34
	35
MPI_INFO_DELETE(INFO, KEY, IERROR)	36
INTEGER INFO, IERROR	37
CHARACTER*(*) KEY	38
	39
MPI_INFO_DUP(INFO, NEWINFO, IERROR)	40
INTEGER INFO, NEWINFO, IERROR	41
	42
MPI_INFO_FREE(INFO, IERROR)	43
INTEGER INFO, IERROR	44
	45
MPI_INFO_GET(INFO, KEY, VALUELEN, VALUE, FLAG, IERROR)	46
INTEGER INFO, VALUELEN, IERROR	47
CHARACTER*(*) KEY, VALUE	48
LOGICAL FLAG	
MPI_INFO_GET_NKEYS(INFO, NKEYS, IERROR)	
INTEGER INFO, NKEYS, IERROR	
MPI_INFO_GET_NTHKEY(INFO, N, KEY, IERROR)	
INTEGER INFO, N, IERROR	
CHARACTER*(*) KEY	

```

1  MPI_INFO_GET_VALUELEN(INFO, KEY, VALUELEN, FLAG, IERROR)
2      INTEGER INFO, VALUELEN, IERROR
3      LOGICAL FLAG
4      CHARACTER*(*) KEY

```

```

5  MPI_INFO_SET(INFO, KEY, VALUE, IERROR)
6      INTEGER INFO, IERROR
7      CHARACTER*(*) KEY, VALUE

```

#### A.4.8 Process Creation and Management Fortran Bindings

```

12 MPI_CLOSE_PORT(PORT_NAME, IERROR)
13     CHARACTER*(*) PORT_NAME
14     INTEGER IERROR

```

```

15 MPI_COMM_ACCEPT(PORT_NAME, INFO, ROOT, COMM, NEWCOMM, IERROR)
16     CHARACTER*(*) PORT_NAME
17     INTEGER INFO, ROOT, COMM, NEWCOMM, IERROR

```

```

19 MPI_COMM_CONNECT(PORT_NAME, INFO, ROOT, COMM, NEWCOMM, IERROR)
20     CHARACTER*(*) PORT_NAME
21     INTEGER INFO, ROOT, COMM, NEWCOMM, IERROR

```

```

22 MPI_COMM_DISCONNECT(COMM, IERROR)
23     INTEGER COMM, IERROR

```

```

25 MPI_COMM_GET_PARENT(PARENT, IERROR)
26     INTEGER PARENT, IERROR

```

```

27 MPI_COMM_JOIN(FD, INTERCOMM, IERROR)
28     INTEGER FD, INTERCOMM, IERROR

```

```

30 MPI_COMM_SPAWN(COMMAND, ARGV, MAXPROCS, INFO, ROOT, COMM, INTERCOMM,
31     ARRAY_OF_ERRCODES, IERROR)
32     CHARACTER*(*) COMMAND, ARGV(*)
33     INTEGER INFO, MAXPROCS, ROOT, COMM, INTERCOMM, ARRAY_OF_ERRCODES(*),
34     IERROR

```

```

35 MPI_COMM_SPAWN_MULTIPLE(COUNT, ARRAY_OF_COMMANDS, ARRAY_OF_ARGV,
36     ARRAY_OF_MAXPROCS, ARRAY_OF_INFO, ROOT, COMM, INTERCOMM,
37     ARRAY_OF_ERRCODES, IERROR)
38     INTEGER COUNT, ARRAY_OF_INFO(*), ARRAY_OF_MAXPROCS(*), ROOT, COMM,
39     INTERCOMM, ARRAY_OF_ERRCODES(*), IERROR
40     CHARACTER*(*) ARRAY_OF_COMMANDS(*), ARRAY_OF_ARGV(COUNT, *)

```

```

42 MPI_LOOKUP_NAME(SERVICE_NAME, INFO, PORT_NAME, IERROR)
43     CHARACTER*(*) SERVICE_NAME, PORT_NAME
44     INTEGER INFO, IERROR

```

```

45 MPI_OPEN_PORT(INFO, PORT_NAME, IERROR)
46     CHARACTER*(*) PORT_NAME
47     INTEGER INFO, IERROR

```

```
MPI_PUBLISH_NAME(SERVICE_NAME, INFO, PORT_NAME, IERROR)
```

```
    INTEGER INFO, IERROR
```

```
    CHARACTER*(*) SERVICE_NAME, PORT_NAME
```

```
MPI_UNPUBLISH_NAME(SERVICE_NAME, INFO, PORT_NAME, IERROR)
```

```
    INTEGER INFO, IERROR
```

```
    CHARACTER*(*) SERVICE_NAME, PORT_NAME
```

#### A.4.9 One-Sided Communications Fortran Bindings

```
MPI_ACCUMULATE(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK,  
               TARGET_DISP, TARGET_COUNT, TARGET_DATATYPE, OP, WIN, IERROR)
```

```
    <type> ORIGIN_ADDR(*)
```

```
    INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
```

```
    INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK, TARGET_COUNT,
```

```
    TARGET_DATATYPE, OP, WIN, IERROR
```

```
MPI_COMPARE_AND_SWAP(ORIGIN_ADDR, COMPARE_ADDR, RESULT_ADDR, DATATYPE,  
                     TARGET_RANK, TARGET_DISP, WIN, IERROR)
```

```
    <type> ORIGIN_ADDR(*), COMPARE_ADDR(*), RESULT_ADDR(*)
```

```
    INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
```

```
    INTEGER DATATYPE, TARGET_RANK, WIN, IERROR
```

```
MPI_FETCH_AND_OP(ORIGIN_ADDR, RESULT_ADDR, DATATYPE, TARGET_RANK,  
                 TARGET_DISP, OP, WIN, IERROR)
```

```
    <type> ORIGIN_ADDR(*), RESULT_ADDR(*)
```

```
    INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
```

```
    INTEGER DATATYPE, TARGET_RANK, OP, WIN, IERROR
```

```
MPI_GET_ACCUMULATE(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE, RESULT_ADDR,  
                   RESULT_COUNT, RESULT_DATATYPE, TARGET_RANK, TARGET_DISP,  
                   TARGET_COUNT, TARGET_DATATYPE, OP, WIN, IERROR)
```

```
    <type> ORIGIN_ADDR(*), RESULT_ADDR(*)
```

```
    INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
```

```
    INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, RESULT_COUNT, RESULT_DATATYPE,
```

```
    TARGET_RANK, TARGET_COUNT, TARGET_DATATYPE, OP, WIN, IERROR
```

```
MPI_GET(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK,  
        TARGET_DISP, TARGET_COUNT, TARGET_DATATYPE, WIN, IERROR)
```

```
    <type> ORIGIN_ADDR(*)
```

```
    INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
```

```
    INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK, TARGET_COUNT,
```

```
    TARGET_DATATYPE, WIN, IERROR
```

```
MPI_PUT(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK,  
        TARGET_DISP, TARGET_COUNT, TARGET_DATATYPE, WIN, IERROR)
```

```
    <type> ORIGIN_ADDR(*)
```

```
    INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
```

```
    INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK, TARGET_COUNT,
```

```
    TARGET_DATATYPE, WIN, IERROR
```

```

1  MPI_RACCUMULATE(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK,
2                  TARGET_DISP, TARGET_COUNT, TARGET_DATATYPE, OP, WIN, REQUEST,
3                  IERROR)
4      <type> ORIGIN_ADDR(*)
5      INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
6      INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK, TARGET_COUNT,
7      TARGET_DATATYPE, OP, WIN, REQUEST, IERROR
8
9  MPI_RGET_ACCUMULATE(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE,
10                     RESULT_ADDR, RESULT_COUNT, RESULT_DATATYPE, TARGET_RANK,
11                     TARGET_DISP, TARGET_COUNT, TARGET_DATATYPE, OP, WIN, REQUEST,
12                     IERROR)
13      <type> ORIGIN_ADDR(*), RESULT_ADDR(*)
14      INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
15      INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, RESULT_COUNT, RESULT_DATATYPE,
16      TARGET_RANK, TARGET_COUNT, TARGET_DATATYPE, OP, WIN, REQUEST, IERROR
17
18  MPI_RGET(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK,
19           TARGET_DISP, TARGET_COUNT, TARGET_DATATYPE, WIN, REQUEST,
20           IERROR)
21      <type> ORIGIN_ADDR(*)
22      INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
23      INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK, TARGET_COUNT,
24      TARGET_DATATYPE, WIN, REQUEST, IERROR
25
26  MPI_RPUT(ORIGIN_ADDR, ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK,
27           TARGET_DISP, TARGET_COUNT, TARGET_DATATYPE, WIN, REQUEST,
28           IERROR)
29      <type> ORIGIN_ADDR(*)
30      INTEGER(KIND=MPI_ADDRESS_KIND) TARGET_DISP
31      INTEGER ORIGIN_COUNT, ORIGIN_DATATYPE, TARGET_RANK, TARGET_COUNT,
32      TARGET_DATATYPE, WIN, REQUEST, IERROR
33
34  MPI_WIN_ALLOCATE_SHARED(SIZE, INFO, COMM, BASEPTR, WIN, IERROR)
35      INTEGER SIZE, INFO, COMM, WIN, IERROR
36      INTEGER(KIND=MPI_ADDRESS_KIND) SIZE, BASEPTR
37
38  MPI_WIN_ALLOCATE(SIZE, DISP_UNIT, INFO, COMM, BASEPTR, WIN, IERROR)
39      INTEGER DISP_UNIT, INFO, COMM, WIN, IERROR
40      INTEGER(KIND=MPI_ADDRESS_KIND) SIZE, BASEPTR
41
42  MPI_WIN_ATTACH(WIN, BASE, SIZE, IERROR)
43      INTEGER WIN, IERROR
44      <type> [base]BASE(*)
45      INTEGER (KIND=MPI_ADDRESS_[SIZE]KIND) [size]SIZE
46
47  MPI_WIN_COMPLETE(WIN, IERROR)
48      INTEGER WIN, IERROR
49
50  MPI_WIN_CREATE(BASE, SIZE, DISP_UNIT, INFO, COMM, WIN, IERROR)
51      <type> BASE(*)

```

```

    INTEGER(KIND=MPI_ADDRESS_KIND) SIZE
    INTEGER DISP_UNIT, INFO, COMM, WIN, IERROR
MPI_WIN_CREATE_DYNAMIC(INFO, COMM, WIN, IERROR)
    INTEGER INFO, COMM, WIN, IERROR
MPI_WIN_DETACH(WIN, BASE, IERROR)
    INTEGER WIN, IERROR
    <type> [base]BASE(*)
MPI_WIN_FENCE(ASSERT, WIN, IERROR)
    INTEGER ASSERT, WIN, IERROR
MPI_WIN_FLUSH_ALL(WIN, IERROR)
    INTEGER WIN, IERROR
MPI_WIN_FLUSH_LOCAL_ALL(WIN, IERROR)
    INTEGER WIN, IERROR
MPI_WIN_FLUSH_LOCAL(RANK, WIN, IERROR)
    INTEGER RANK, WIN, IERROR
MPI_WIN_FLUSH(RANK, WIN, IERROR)
    INTEGER RANK, WIN, IERROR
MPI_WIN_FREE(WIN, IERROR)
    INTEGER WIN, IERROR
MPI_WIN_GET_GROUP(WIN, GROUP, IERROR)
    INTEGER WIN, GROUP, IERROR
MPI_WIN_LOCK_ALL(ASSERT, WIN, IERROR)
    INTEGER ASSERT, WIN, IERROR
MPI_WIN_LOCK(LOCK_TYPE, RANK, ASSERT, WIN, IERROR)
    INTEGER LOCK_TYPE, RANK, ASSERT, WIN, IERROR
MPI_WIN_POST(GROUP, ASSERT, WIN, IERROR)
    INTEGER GROUP, ASSERT, WIN, IERROR
MPI_WIN_SHARED_QUERY(WIN, RANK, SIZE, BASEPTR, IERROR)
    INTEGER WIN, RANK, IERROR
    INTEGER (KIND=MPI_ADDRESS_KIND) SIZE, BASEPTR
MPI_WIN_START(GROUP, ASSERT, WIN, IERROR)
    INTEGER GROUP, ASSERT, WIN, IERROR
MPI_WIN_SYNC(WIN, IERROR)
    INTEGER WIN, IERROR
MPI_WIN_TEST(WIN, FLAG, IERROR)
    INTEGER WIN, IERROR
    LOGICAL FLAG
MPI_WIN_UNLOCK_ALL(WIN, IERROR)

```

1       INTEGER WIN, IERROR

2       MPI\_WIN\_UNLOCK(RANK, WIN, IERROR)

3       INTEGER RANK, WIN, IERROR

4       MPI\_WIN\_WAIT(WIN, IERROR)

5       INTEGER WIN, IERROR

#### 6       A.4.10 External Interfaces Fortran Bindings

7       MPI\_GREQUEST\_COMPLETE(REQUEST, IERROR)

8       INTEGER REQUEST, IERROR

9       MPI\_GREQUEST\_START(QUERY\_FN, FREE\_FN, CANCEL\_FN, EXTRA\_STATE, REQUEST,  
10       IERROR)

11       INTEGER REQUEST, IERROR

12       EXTERNAL QUERY\_FN, FREE\_FN, CANCEL\_FN

13       INTEGER (KIND=MPI\_ADDRESS\_KIND) EXTRA\_STATE

14       MPI\_INIT\_THREAD(REQUIRED, PROVIDED, IERROR)

15       INTEGER REQUIRED, PROVIDED, IERROR

16       MPI\_IS\_THREAD\_MAIN(FLAG, IERROR)

17       LOGICAL FLAG

18       INTEGER IERROR

19       MPI\_QUERY\_THREAD(PROVIDED, IERROR)

20       INTEGER PROVIDED, IERROR

21       MPI\_STATUS\_SET\_CANCELLED(STATUS, FLAG, IERROR)

22       INTEGER STATUS(MPI\_STATUS\_SIZE), IERROR

23       LOGICAL FLAG

24       MPI\_STATUS\_SET\_ELEMENTS(STATUS, DATATYPE, COUNT, IERROR)

25       INTEGER STATUS(MPI\_STATUS\_SIZE), DATATYPE, COUNT, IERROR

26       MPI\_STATUS\_SET\_ELEMENTS\_X(STATUS, DATATYPE, COUNT, IERROR)

27       INTEGER STATUS(MPI\_STATUS\_SIZE), DATATYPE, IERROR

28       INTEGER (KIND=MPI\_COUNT\_KIND) COUNT

#### 29       A.4.11 I/O Fortran Bindings

30       MPI\_FILE\_CLOSE(FH, IERROR)

31       INTEGER FH, IERROR

32       MPI\_FILE\_DELETE(FILENAME, INFO, IERROR)

33       CHARACTER\*(\*) FILENAME

34       INTEGER INFO, IERROR

35       MPI\_FILE\_GET\_AMODE(FH, AMODE, IERROR)

36       INTEGER FH, AMODE, IERROR

37       MPI\_FILE\_GET\_ATOMICITY(FH, FLAG, IERROR)



INTEGER FH, IERROR	1
LOGICAL FLAG	2
	3
MPI_FILE_GET_BYTE_OFFSET(FH, OFFSET, DISP, IERROR)	4
INTEGER FH, IERROR	5
INTEGER(KIND=MPI_OFFSET_KIND) OFFSET, DISP	6
MPI_FILE_GET_GROUP(FH, GROUP, IERROR)	7
INTEGER FH, GROUP, IERROR	8
	9
MPI_FILE_GET_INFO(FH, INFO_USED, IERROR)	10
INTEGER FH, INFO_USED, IERROR	11
	12
MPI_FILE_GET_POSITION(FH, OFFSET, IERROR)	13
INTEGER FH, IERROR	14
INTEGER(KIND=MPI_OFFSET_KIND) OFFSET	15
MPI_FILE_GET_POSITION_SHARED(FH, OFFSET, IERROR)	16
INTEGER FH, IERROR	17
INTEGER(KIND=MPI_OFFSET_KIND) OFFSET	18
	19
MPI_FILE_GET_SIZE(FH, SIZE, IERROR)	20
INTEGER FH, IERROR	21
INTEGER(KIND=MPI_OFFSET_KIND) SIZE	22
MPI_FILE_GET_TYPE_EXTENT(FH, DATATYPE, EXTENT, IERROR)	23
INTEGER FH, DATATYPE, IERROR	24
INTEGER(KIND=MPI_ADDRESS_KIND) EXTENT	25
	26
MPI_FILE_GET_VIEW(FH, DISP, ETYPE, FILETYPE, DATAREP, IERROR)	27
INTEGER FH, ETYPE, FILETYPE, IERROR	28
CHARACTER*(*) DATAREP	29
INTEGER(KIND=MPI_OFFSET_KIND) DISP	30
MPI_FILE_IREAD_AT(FH, OFFSET, BUF, COUNT, DATATYPE, REQUEST, IERROR)	31
<type> BUF(*)	32
INTEGER FH, COUNT, DATATYPE, REQUEST, IERROR	33
INTEGER(KIND=MPI_OFFSET_KIND) OFFSET	34
	35
MPI_FILE_IREAD(FH, BUF, COUNT, DATATYPE, REQUEST, IERROR)	36
<type> BUF(*)	37
INTEGER FH, COUNT, DATATYPE, REQUEST, IERROR	38
	39
MPI_FILE_IREAD_SHARED(FH, BUF, COUNT, DATATYPE, REQUEST, IERROR)	40
<type> BUF(*)	41
INTEGER FH, COUNT, DATATYPE, REQUEST, IERROR	42
	43
MPI_FILE_IWRITE_AT(FH, OFFSET, BUF, COUNT, DATATYPE, REQUEST, IERROR)	44
<type> BUF(*)	45
INTEGER FH, COUNT, DATATYPE, REQUEST, IERROR	46
INTEGER(KIND=MPI_OFFSET_KIND) OFFSET	47
	48
MPI_FILE_IWRITE(FH, BUF, COUNT, DATATYPE, REQUEST, IERROR)	
<type> BUF(*)	

```

1      INTEGER FH, COUNT, DATATYPE, REQUEST, IERROR
2
3      MPI_FILE_IWRITE_SHARED(FH, BUF, COUNT, DATATYPE, REQUEST, IERROR)
4          <type> BUF(*)
5      INTEGER FH, COUNT, DATATYPE, REQUEST, IERROR
6
7      MPI_FILE_OPEN(COMM, FILENAME, AMODE, INFO, FH, IERROR)
8          CHARACTER*(*) FILENAME
9      INTEGER COMM, AMODE, INFO, FH, IERROR
10
11     MPI_FILE_PREALLOCATE(FH, SIZE, IERROR)
12         INTEGER FH, IERROR
13         INTEGER(KIND=MPI_OFFSET_KIND) SIZE
14
15     MPI_FILE_READ_ALL_BEGIN(FH, BUF, COUNT, DATATYPE, IERROR)
16         <type> BUF(*)
17         INTEGER FH, COUNT, DATATYPE, IERROR
18
19     MPI_FILE_READ_ALL_END(FH, BUF, STATUS, IERROR)
20         <type> BUF(*)
21         INTEGER FH, STATUS(MPI_STATUS_SIZE), IERROR
22
23     MPI_FILE_READ_ALL(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
24         <type> BUF(*)
25         INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
26
27     MPI_FILE_READ_AT_ALL_BEGIN(FH, OFFSET, BUF, COUNT, DATATYPE, IERROR)
28         <type> BUF(*)
29         INTEGER FH, COUNT, DATATYPE, IERROR
30         INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
31
32     MPI_FILE_READ_AT_ALL_END(FH, BUF, STATUS, IERROR)
33         <type> BUF(*)
34         INTEGER FH, STATUS(MPI_STATUS_SIZE), IERROR
35         INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
36
37     MPI_FILE_READ_AT(FH, OFFSET, BUF, COUNT, DATATYPE, STATUS, IERROR)
38         <type> BUF(*)
39         INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
40         INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
41
42     MPI_FILE_READ(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
43         <type> BUF(*)
44         INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
45
46     MPI_FILE_READ_ORDERED_BEGIN(FH, BUF, COUNT, DATATYPE, IERROR)
47         <type> BUF(*)
48         INTEGER FH, COUNT, DATATYPE, IERROR
49
50     MPI_FILE_READ_ORDERED_END(FH, BUF, STATUS, IERROR)

```

```

    <type> BUF(*)
    INTEGER FH, STATUS(MPI_STATUS_SIZE), IERROR
MPI_FILE_READ_ORDERED(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
    <type> BUF(*)
    INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
MPI_FILE_READ_SHARED(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
    <type> BUF(*)
    INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
MPI_FILE_SEEK(FH, OFFSET, WHENCE, IERROR)
    INTEGER FH, WHENCE, IERROR
    INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
MPI_FILE_SEEK_SHARED(FH, OFFSET, WHENCE, IERROR)
    INTEGER FH, WHENCE, IERROR
    INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
MPI_FILE_SET_ATOMICITY(FH, FLAG, IERROR)
    INTEGER FH, IERROR
    LOGICAL FLAG
MPI_FILE_SET_INFO(FH, INFO, IERROR)
    INTEGER FH, INFO, IERROR
MPI_FILE_SET_SIZE(FH, SIZE, IERROR)
    INTEGER FH, IERROR
    INTEGER(KIND=MPI_OFFSET_KIND) SIZE
MPI_FILE_SET_VIEW(FH, DISP, ETYPE, FILETYPE, DATAREP, INFO, IERROR)
    INTEGER FH, ETYPE, FILETYPE, INFO, IERROR
    CHARACTER*(*) DATAREP
    INTEGER(KIND=MPI_OFFSET_KIND) DISP
MPI_FILE_SYNC(FH, IERROR)
    INTEGER FH, IERROR
MPI_FILE_WRITE_ALL_BEGIN(FH, BUF, COUNT, DATATYPE, IERROR)
    <type> BUF(*)
    INTEGER FH, COUNT, DATATYPE, IERROR
MPI_FILE_WRITE_ALL_END(FH, BUF, STATUS, IERROR)
    <type> BUF(*)
    INTEGER FH, STATUS(MPI_STATUS_SIZE), IERROR
MPI_FILE_WRITE_ALL(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
    <type> BUF(*)
    INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
MPI_FILE_WRITE_AT_ALL_BEGIN(FH, OFFSET, BUF, COUNT, DATATYPE, IERROR)
    <type> BUF(*)
    INTEGER FH, COUNT, DATATYPE, IERROR
    INTEGER(KIND=MPI_OFFSET_KIND) OFFSET

```

```

1  MPI_FILE_WRITE_AT_ALL_END(FH, BUF, STATUS, IERROR)
2      <type> BUF(*)
3      INTEGER FH, STATUS(MPI_STATUS_SIZE), IERROR
4
5  MPI_FILE_WRITE_AT_ALL(FH, OFFSET, BUF, COUNT, DATATYPE, STATUS, IERROR)
6      <type> BUF(*)
7      INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
8      INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
9
10 MPI_FILE_WRITE_AT(FH, OFFSET, BUF, COUNT, DATATYPE, STATUS, IERROR)
11     <type> BUF(*)
12     INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
13     INTEGER(KIND=MPI_OFFSET_KIND) OFFSET
14
15 MPI_FILE_WRITE(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
16     <type> BUF(*)
17     INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
18
19 MPI_FILE_WRITE_ORDERED_BEGIN(FH, BUF, COUNT, DATATYPE, IERROR)
20     <type> BUF(*)
21     INTEGER FH, COUNT, DATATYPE, IERROR
22
23 MPI_FILE_WRITE_ORDERED_END(FH, BUF, STATUS, IERROR)
24     <type> BUF(*)
25     INTEGER FH, STATUS(MPI_STATUS_SIZE), IERROR
26
27 MPI_FILE_WRITE_ORDERED(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
28     <type> BUF(*)
29     INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
30
31 MPI_FILE_WRITE_SHARED(FH, BUF, COUNT, DATATYPE, STATUS, IERROR)
32     <type> BUF(*)
33     INTEGER FH, COUNT, DATATYPE, STATUS(MPI_STATUS_SIZE), IERROR
34
35 MPI_REGISTER_DATAREP(DATAREP, READ_CONVERSION_FN, WRITE_CONVERSION_FN,
36     DTYPE_FILE_EXTENT_FN, EXTRA_STATE, IERROR)
37     CHARACTER*(*) DATAREP
38     EXTERNAL READ_CONVERSION_FN, WRITE_CONVERSION_FN, DTYPE_FILE_EXTENT_FN
39     INTEGER(KIND=MPI_ADDRESS_KIND) EXTRA_STATE
40     INTEGER IERROR

```

#### A.4.12 Language Bindings Fortran Bindings

```

41 MPI_F_SYNC_REG(buf)
42     <type> buf(*)
43
44 MPI_SIZEOF(X, SIZE, IERROR)
45     <type> X
46     INTEGER SIZE, IERROR
47
48 MPI_STATUS_F082F(F08_STATUS, F_STATUS, IERROR)
49     TYPE(MPI_Status) :: F08_STATUS

```

```

    INTEGER :: F_STATUS(MPI_STATUS_SIZE)
    INTEGER IERROR

MPI_STATUS_F2F08(F_STATUS, F08_STATUS, IERROR)
    INTEGER :: F_STATUS(MPI_STATUS_SIZE)
    TYPE(MPI_Status) :: F08_STATUS
    INTEGER IERROR

MPI_TYPE_CREATE_F90_COMPLEX(P, R, NEWTYPE, IERROR)
    INTEGER P, R, NEWTYPE, IERROR

MPI_TYPE_CREATE_F90_INTEGER(R, NEWTYPE, IERROR)
    INTEGER R, NEWTYPE, IERROR

MPI_TYPE_CREATE_F90_REAL(P, R, NEWTYPE, IERROR)
    INTEGER P, R, NEWTYPE, IERROR

MPI_TYPE_MATCH_SIZE(TYPECLASS, SIZE, DATATYPE, IERROR)
    INTEGER TYPECLASS, SIZE, DATATYPE, IERROR

```

#### A.4.13 Profiling Interface Fortran Bindings

```

MPI_PCONTROL(LEVEL)
    INTEGER LEVEL

```

#### A.4.14 Deprecated Fortran Bindings

```

MPI_ADDRESS(LOCATION, ADDRESS, IERROR)
    <type> LOCATION(*)
    INTEGER ADDRESS, IERROR

MPI_ATTR_DELETE(COMM, KEYVAL, IERROR)
    INTEGER COMM, KEYVAL, IERROR

MPI_ATTR_GET(COMM, KEYVAL, ATTRIBUTE_VAL, FLAG, IERROR)
    INTEGER COMM, KEYVAL, ATTRIBUTE_VAL, IERROR
    LOGICAL FLAG

MPI_ATTR_PUT(COMM, KEYVAL, ATTRIBUTE_VAL, IERROR)
    INTEGER COMM, KEYVAL, ATTRIBUTE_VAL, IERROR

MPI_DUP_FN(OLDCOMM, KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,
           ATTRIBUTE_VAL_OUT, FLAG, IERR)
    INTEGER OLDCOMM, KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,
    ATTRIBUTE_VAL_OUT, IERR
    LOGICAL FLAG

MPI_ERRHANDLER_CREATE([FUNCTION]HANDLER_FN, ERRHANDLER, IERROR)
    EXTERNAL [FUNCTION]HANDLER_FN
    INTEGER ERRHANDLER, IERROR

MPI_ERRHANDLER_GET(COMM, ERRHANDLER, IERROR)

```

```

1      INTEGER COMM, ERRHANDLER, IERROR
2
3      MPI_ERRHANDLER_SET(COMM, ERRHANDLER, IERROR)
4      INTEGER COMM, ERRHANDLER, IERROR
5
6      MPI_KEYVAL_CREATE(COPY_FN, DELETE_FN, KEYVAL, EXTRA_STATE, IERROR)
7      EXTERNAL COPY_FN, DELETE_FN
8      INTEGER KEYVAL, EXTRA_STATE, IERROR
9
10     MPI_KEYVAL_FREE(KEYVAL, IERROR)
11     INTEGER KEYVAL, IERROR
12
13     MPI_NULL_COPY_FN(OLDCOMM, KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,
14                     ATTRIBUTE_VAL_OUT, FLAG, IERR)
15     INTEGER OLDCOMM, KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,
16     ATTRIBUTE_VAL_OUT, IERR
17     LOGICAL FLAG
18
19     MPI_NULL_DELETE_FN(COMM, KEYVAL, ATTRIBUTE_VAL, EXTRA_STATE, IERROR)
20     INTEGER COMM, KEYVAL, ATTRIBUTE_VAL, EXTRA_STATE, IERROR
21
22     MPI_TYPE_EXTENT(DATATYPE, EXTENT, IERROR)
23     INTEGER DATATYPE, EXTENT, IERROR
24
25     MPI_TYPE_HINDEXED(COUNT, ARRAY_OF_BLOCKLENGTHS, ARRAY_OF_DISPLACEMENTS,
26                     OLDTYPE, NEWTYPE, IERROR)
27     INTEGER COUNT, ARRAY_OF_BLOCKLENGTHS(*), ARRAY_OF_DISPLACEMENTS(*),
28     OLDTYPE, NEWTYPE, IERROR
29
30     MPI_TYPE_HVECTOR(COUNT, BLOCKLENGTH, STRIDE, OLDTYPE, NEWTYPE, IERROR)
31     INTEGER COUNT, BLOCKLENGTH, STRIDE, OLDTYPE, NEWTYPE, IERROR
32
33     MPI_TYPE_LB( DATATYPE, DISPLACEMENT, IERROR)
34     INTEGER DATATYPE, DISPLACEMENT, IERROR
35
36     MPI_TYPE_STRUCT(COUNT, ARRAY_OF_BLOCKLENGTHS, ARRAY_OF_DISPLACEMENTS,
37                     ARRAY_OF_TYPES, NEWTYPE, IERROR)
38     INTEGER COUNT, ARRAY_OF_BLOCKLENGTHS(*), ARRAY_OF_DISPLACEMENTS(*),
39     ARRAY_OF_TYPES(*), NEWTYPE, IERROR
40
41     MPI_TYPE_UB( DATATYPE, DISPLACEMENT, IERROR)
42     INTEGER DATATYPE, DISPLACEMENT, IERROR
43
44     SUBROUTINE COPY_FUNCTION(OLDCOMM, KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,
45                             ATTRIBUTE_VAL_OUT, FLAG, IERR)
46     INTEGER OLDCOMM, KEYVAL, EXTRA_STATE, ATTRIBUTE_VAL_IN,
47     ATTRIBUTE_VAL_OUT, IERR
48     LOGICAL FLAG
49
50     SUBROUTINE DELETE_FUNCTION(COMM, KEYVAL, ATTRIBUTE_VAL, EXTRA_STATE, IERR)
51     INTEGER COMM, KEYVAL, ATTRIBUTE_VAL, EXTRA_STATE, IERR

```

## A.5 C++ Bindings (deprecated)

### A.5.1 Point-to-Point Communication C++ Bindings

```

namespace MPI {

    {void Attach_buffer(void* buffer, int size) (binding deprecated, see
        Section 15.2) }

    {void Comm::Bsend(const void* buf, int count, const Datatype& datatype,
        int dest, int tag) const (binding deprecated, see Section 15.2) }

    {Prequest Comm::Bsend_init(const void* buf, int count, const
        Datatype& datatype, int dest, int tag) const (binding deprecated,
        see Section 15.2) }

    {void Request::Cancel() const (binding deprecated, see Section 15.2) }

    {int Detach_buffer(void*& buffer) (binding deprecated, see Section 15.2) }

    {void Request::Free() (binding deprecated, see Section 15.2) }

    {int Status::Get_count(const Datatype& datatype) const (binding deprecated,
        see Section 15.2) }

    {int Status::Get_error() const (binding deprecated, see Section 15.2) }

    {int Status::Get_source() const (binding deprecated, see Section 15.2) }

    {bool Request::Get_status() const (binding deprecated, see Section 15.2) }

    {bool Request::Get_status(Status& status) const (binding deprecated, see
        Section 15.2) }

    {int Status::Get_tag() const (binding deprecated, see Section 15.2) }

    {Request Comm::Ibsend(const void* buf, int count, const
        Datatype& datatype, int dest, int tag) const (binding deprecated,
        see Section 15.2) }

    {bool Comm::Iprobe(int source, int tag) const (binding deprecated, see
        Section 15.2) }

    {bool Comm::Iprobe(int source, int tag, Status& status) const (binding
        deprecated, see Section 15.2) }

    {Request Comm::Irecv(void* buf, int count, const Datatype& datatype,
        int source, int tag) const (binding deprecated, see Section 15.2) }

    {Request Comm::Irsend(const void* buf, int count, const
        Datatype& datatype, int dest, int tag) const (binding deprecated,
        see Section 15.2) }

    {bool Status::Is_cancelled() const (binding deprecated, see Section 15.2) }

    {Request Comm::Isend(const void* buf, int count, const
        Datatype& datatype, int dest, int tag) const (binding deprecated,

```

```

1         see Section 15.2) }
2
3     {Request Comm::Issend(const void* buf, int count, const
4         Datatype& datatype, int dest, int tag) const(binding deprecated,
5         see Section 15.2) }
6
7     {void Comm::Probe(int source, int tag) const(binding deprecated, see
8         Section 15.2) }
9
10    {void Comm::Probe(int source, int tag, Status& status) const(binding
11        deprecated, see Section 15.2) }
12
13    {Prequest Comm::Recv_init(void* buf, int count, const Datatype& datatype,
14        int source, int tag) const(binding deprecated, see Section 15.2) }
15
16    {void Comm::Recv(void* buf, int count, const Datatype& datatype,
17        int source, int tag) const(binding deprecated, see Section 15.2) }
18
19    {void Comm::Recv(void* buf, int count, const Datatype& datatype,
20        int source, int tag, Status& status) const(binding deprecated, see
21        Section 15.2) }
22
23    {void Comm::Rsend(const void* buf, int count, const Datatype& datatype,
24        int dest, int tag) const(binding deprecated, see Section 15.2) }
25
26    {Prequest Comm::Rsend_init(const void* buf, int count, const
27        Datatype& datatype, int dest, int tag) const(binding deprecated,
28        see Section 15.2) }
29
30    {void Comm::Send(const void* buf, int count, const Datatype& datatype,
31        int dest, int tag) const(binding deprecated, see Section 15.2) }
32
33    {Prequest Comm::Send_init(const void* buf, int count, const
34        Datatype& datatype, int dest, int tag) const(binding deprecated,
35        see Section 15.2) }
36
37    {void Comm::Sendrecv(const void *sendbuf, int sendcount, const
38        Datatype& sendtype, int dest, int sendtag, void *recvbuf,
39        int recvcount, const Datatype& recvtype, int source,
40        int recvttag) const(binding deprecated, see Section 15.2) }
41
42    {void Comm::Sendrecv(const void *sendbuf, int sendcount, const
43        Datatype& sendtype, int dest, int sendtag, void *recvbuf,
44        int recvcount, const Datatype& recvtype, int source,
45        int recvttag, Status& status) const(binding deprecated, see
46        Section 15.2) }
47
48    {void Comm::Sendrecv_replace(void* buf, int count, const
49        Datatype& datatype, int dest, int sendtag, int source,
50        int recvttag) const(binding deprecated, see Section 15.2) }
51
52    {void Comm::Sendrecv_replace(void* buf, int count, const
53        Datatype& datatype, int dest, int sendtag, int source,
54        int recvttag, Status& status) const(binding deprecated, see
55        Section 15.2) }

```



```

    Section 15.2) }
1
2
3 {void Status::Set_error(int error) (binding deprecated, see Section 15.2) }
4 {void Status::Set_source(int source) (binding deprecated, see Section 15.2) }
5 {void Status::Set_tag(int tag) (binding deprecated, see Section 15.2) }
6 {void Comm::Ssend(const void* buf, int count, const Datatype& datatype,
7     int dest, int tag) const (binding deprecated, see Section 15.2) }
8 {Prequest Comm::Ssend_init(const void* buf, int count, const
9     Datatype& datatype, int dest, int tag) const (binding deprecated,
10     see Section 15.2) }
11 {static void Prequest::Startall(int count,
12     Prequest array_of_requests[]) (binding deprecated, see Section 15.2) }
13 {void Prequest::Start() (binding deprecated, see Section 15.2) }
14 {static bool Request::Testall(int count, Request array_of_requests[],
15     Status array_of_statuses[]) (binding deprecated, see Section 15.2) }
16 {static bool Request::Testall(int count,
17     Request array_of_requests[]) (binding deprecated, see Section 15.2) }
18 {static bool Request::Testany(int count, Request array_of_requests[],
19     int& index, Status& status) (binding deprecated, see Section 15.2) }
20 {static bool Request::Testany(int count, Request array_of_requests[],
21     int& index) (binding deprecated, see Section 15.2) }
22 {static bool Request::Testany(int count, Request array_of_requests[],
23     int& index) (binding deprecated, see Section 15.2) }
24 {static bool Request::Testany(int count, Request array_of_requests[],
25     int& index) (binding deprecated, see Section 15.2) }
26 {bool Request::Test() (binding deprecated, see Section 15.2) }
27 {bool Request::Test(Status& status) (binding deprecated, see Section 15.2) }
28 {static int Request::Testsome(int incount, Request array_of_requests[],
29     int array_of_indices[], Status array_of_statuses[]) (binding
30     deprecated, see Section 15.2) }
31 {static int Request::Testsome(int incount, Request array_of_requests[],
32     int array_of_indices[]) (binding deprecated, see Section 15.2) }
33 {static void Request::Waitall(int count, Request array_of_requests[],
34     Status array_of_statuses[]) (binding deprecated, see Section 15.2) }
35 {static void Request::Waitall(int count,
36     Request array_of_requests[]) (binding deprecated, see Section 15.2) }
37 {static int Request::Waitany(int count, Request array_of_requests[],
38     Status& status) (binding deprecated, see Section 15.2) }
39 {static int Request::Waitany(int count,
40     Request array_of_requests[]) (binding deprecated, see Section 15.2) }
41 {static int Request::Waitany(int count,
42     Request array_of_requests[]) (binding deprecated, see Section 15.2) }
43 {void Request::Wait(Status& status) (binding deprecated, see Section 15.2) }
44
45
46
47
48

```

```

1  {static int Request::Waitsome(int incount, Request array_of_requests[],
2      int array_of_indices[], Status array_of_statuses[]) (binding
3      deprecated, see Section 15.2) }
4
5  {static int Request::Waitsome(int incount, Request array_of_requests[],
6      int array_of_indices[]) (binding deprecated, see Section 15.2) }
7
8  {void Request::Wait() (binding deprecated, see Section 15.2) }
9
10 };

```

## A.5.2 Datatypes C++ Bindings

```

13 namespace MPI {
14
15     {void Datatype::Commit() (binding deprecated, see Section 15.2) }
16
17     {Datatype Datatype::Create_contiguous(int count) const (binding deprecated,
18         see Section 15.2) }
19
20     {Datatype Datatype::Create_darray(int size, int rank, int ndims,
21         const int array_of_gsizes[], const int array_of_distrib[],
22         const int array_of_dargs[], const int array_of_psize[],
23         int order) const (binding deprecated, see Section 15.2) }
24
25     {Datatype Datatype::Create_hindexed(int count,
26         const int array_of_blocklengths[],
27         const Aint array_of_displacements[]) const (binding deprecated, see
28         Section 15.2) }
29
30     {Datatype Datatype::Create_hvector(int count, int blocklength, Aint
31         stride) const (binding deprecated, see Section 15.2) }
32
33     {Datatype Datatype::Create_indexed_block(int count, int blocklength,
34         const int array_of_displacements[]) const (binding deprecated, see
35         Section 15.2) }
36
37     {Datatype Datatype::Create_indexed(int count,
38         const int array_of_blocklengths[],
39         const int array_of_displacements[]) const (binding deprecated, see
40         Section 15.2) }
41
42     {Datatype Datatype::Create_resized(const Aint lb, const Aint extent)
43         const (binding deprecated, see Section 15.2) }
44
45     {static Datatype Datatype::Create_struct(int count,
46         const int array_of_blocklengths[], const Aint
47         array_of_displacements[],
48         const Datatype array_of_types[]) (binding deprecated, see
49         Section 15.2) }
50
51     {Datatype Datatype::Create_subarray(int ndims,
52         const int array_of_sizes[], const int array_of_subsizes[],

```

```

        const int array_of_starts[], int order) const(binding deprecated,
        see Section 15.2) }
1
2
3
{Datatype Datatype::Create_vector(int count, int blocklength, int stride)
    const(binding deprecated, see Section 15.2) }
4
5
{Datatype Datatype::Dup() const(binding deprecated, see Section 15.2) }
6
7
{void Datatype::Free() (binding deprecated, see Section 15.2) }
8
9
{Aint Get_address(void* location) (binding deprecated, see Section 15.2) }
10
11
{void Datatype::Get_contents(int max_integers, int max_addresses,
    int max_datatypes, int array_of_integers[],
    Aint array_of_addresses[], Datatype array_of_datatypes[])
    const(binding deprecated, see Section 15.2) }
12
13
14
{int Status::Get_elements(const Datatype& datatype) const(binding deprecated,
    see Section 15.2) }
15
16
17
{void Datatype::Get_envelope(int& num_integers, int& num_addresses,
    int& num_datatypes, int& combiner) const(binding deprecated, see
    Section 15.2) }
18
19
20
{void Datatype::Get_extent(Aint& lb, Aint& extent) const(binding deprecated,
    see Section 15.2) }
21
22
23
{int Datatype::Get_size() const(binding deprecated, see Section 15.2) }
24
25
{void Datatype::Get_true_extent(Aint& true_lb, Aint& true_extent)
    const(binding deprecated, see Section 15.2) }
26
27
{void Datatype::Pack(const void* inbuf, int incount, void *outbuf,
    int outsize, int& position, const Comm &comm) const(binding
    deprecated, see Section 15.2) }
28
29
30
{void Datatype::Pack_external(const char* datarep, const void* inbuf,
    int incount, void* outbuf, Aint outsize, Aint& position)
    const(binding deprecated, see Section 15.2) }
31
32
33
34
{Aint Datatype::Pack_external_size(const char* datarep, int incount)
    const(binding deprecated, see Section 15.2) }
35
36
37
{int Datatype::Pack_size(int incount, const Comm& comm) const(binding
    deprecated, see Section 15.2) }
38
39
40
{void Datatype::Unpack(const void* inbuf, int insize, void *outbuf,
    int outcount, int& position, const Comm& comm) const(binding
    deprecated, see Section 15.2) }
41
42
43
{void Datatype::Unpack_external(const char* datarep, const void* inbuf,
    Aint insize, Aint& position, void* outbuf, int outcount) const
    (binding deprecated, see Section 15.2) }
44
45
46
};
47
48

```

## A.5.3 Collective Communication C++ Bindings

```

namespace MPI {

    {void Comm::Allgather(const void* sendbuf, int sendcount, const
        Datatype& sendtype, void* recvbuf, int recvcount,
        const Datatype& recvtpe) const = 0(binding deprecated, see
        Section 15.2) }

    {void Comm::Allgatherv(const void* sendbuf, int sendcount, const
        Datatype& sendtype, void* recvbuf, const int recvcounts[],
        const int displs[], const Datatype& recvtpe) const = 0(binding
        deprecated, see Section 15.2) }

    {void Comm::Allreduce(const void* sendbuf, void* recvbuf, int count,
        const Datatype& datatype, const Op& op) const = 0(binding
        deprecated, see Section 15.2) }

    {void Comm::Alltoall(const void* sendbuf, int sendcount, const
        Datatype& sendtype, void* recvbuf, int recvcount,
        const Datatype& recvtpe) const = 0(binding deprecated, see
        Section 15.2) }

    {void Comm::Alltoallv(const void* sendbuf, const int sendcounts[],
        const int sdispls[], const Datatype& sendtype, void* recvbuf,
        const int recvcounts[], const int rdispls[],
        const Datatype& recvtpe) const = 0(binding deprecated, see
        Section 15.2) }

    {void Comm::Alltoallw(const void* sendbuf, const int sendcounts[], const
        int sdispls[], const Datatype sendtypes[], void* recvbuf,
        const int recvcounts[], const int rdispls[], const Datatype
        recvtypes[]) const = 0(binding deprecated, see Section 15.2) }

    {void Comm::Barrier() const = 0(binding deprecated, see Section 15.2) }

    {void Comm::Bcast(void* buffer, int count, const Datatype& datatype,
        int root) const = 0(binding deprecated, see Section 15.2) }

    {void Intracomm::Exscan(const void* sendbuf, void* recvbuf, int count,
        const Datatype& datatype, const Op& op) const(binding deprecated,
        see Section 15.2) }

    {void Op::Free() (binding deprecated, see Section 15.2) }

    {void Comm::Gather(const void* sendbuf, int sendcount, const
        Datatype& sendtype, void* recvbuf, int recvcount,
        const Datatype& recvtpe, int root) const = 0(binding deprecated,
        see Section 15.2) }

    {void Comm::Gatherv(const void* sendbuf, int sendcount, const
        Datatype& sendtype, void* recvbuf, const int recvcounts[],
        const int displs[], const Datatype& recvtpe, int root)
        const = 0(binding deprecated, see Section 15.2) }

```

```

{Request Comm::Iallgather(const void* sendbuf, int sendcount, const
    Datatype& sendtype, void* recvbuf, int recvcount,
    const Datatype& recvtype) const = 0(binding deprecated, see
    Section 15.2) }

{Request Comm::Iallgatherv(const void* sendbuf, int sendcount, const
    Datatype& sendtype, void* recvbuf, const int recvcounts[],
    const int displs[], const Datatype& recvtype) const = 0(binding
    deprecated, see Section 15.2) }

{Request Comm::Iallreduce(const void* sendbuf, void* recvbuf, int count,
    const Datatype& datatype, const Op& op) const = 0(binding
    deprecated, see Section 15.2) }

{Request Comm::Ialltoall(const void* sendbuf, int sendcount, const
    Datatype& sendtype, void* recvbuf, int recvcount,
    const Datatype& recvtype) const = 0(binding deprecated, see
    Section 15.2) }

{Request Comm::Ialltoallv(const void* sendbuf, const int sendcounts[],
    const int sdispls[], const Datatype& sendtype, void* recvbuf,
    const int recvcounts[], const int rdispls[],
    const Datatype& recvtype) const = 0(binding deprecated, see
    Section 15.2) }

{Request Comm::Ialltoallw(const void* sendbuf, const int sendcounts[],
    const int sdispls[], const Datatype sendtypes[], void*
    recvbuf, const int recvcounts[], const int rdispls[], const
    Datatype recvtypes[]) const = 0(binding deprecated, see
    Section 15.2) }

{Request Comm::Ibarrier() const = 0(binding deprecated, see Section 15.2) }

{Request Comm::Ibcast(void* buffer, int count, const Datatype& datatype,
    int root) const = 0(binding deprecated, see Section 15.2) }

{Request Intracomm::Iexscan(const void* sendbuf, void* recvbuf, int
    count, const Datatype& datatype, const Op& op) const(binding
    deprecated, see Section 15.2) }

{Request Comm::Igather(const void* sendbuf, int sendcount, const
    Datatype& sendtype, void* recvbuf, int recvcount,
    const Datatype& recvtype, int root) const = 0(binding deprecated,
    see Section 15.2) }

{Request Comm::Igatherv(const void* sendbuf, int sendcount, const
    Datatype& sendtype, void* recvbuf, const int recvcounts[],
    const int displs[], const Datatype& recvtype, int root)
    const = 0(binding deprecated, see Section 15.2) }

{void Op::Init(User_function* [function]user_fn, bool commute)(binding
    deprecated, see Section 15.2) }

```

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```

1  {Request Comm::Ireduce(const void* sendbuf, void* recvbuf, int count,
2      const Datatype& datatype, const Op& op, int root)
3      const = 0(binding deprecated, see Section 15.2) }
4
5  {Request Comm::Ireduce_scatter_block(const void* sendbuf, void* recvbuf,
6      int recvcnt, const Datatype& datatype, const Op& op)
7      const = 0(binding deprecated, see Section 15.2) }
8
9  {Request Comm::Ireduce_scatter(const void* sendbuf, void* recvbuf,
10     int recvcnts[], const Datatype& datatype, const Op& op)
11     const = 0(binding deprecated, see Section 15.2) }
12
13 {Request Intracomm::Iscan(const void* sendbuf, void* recvbuf, int count,
14     const Datatype& datatype, const Op& op) const(binding deprecated,
15     see Section 15.2) }
16
17 {Request Comm::Iscatter(const void* sendbuf, int sendcount, const
18     Datatype& sendtype, void* recvbuf, int recvcnt,
19     const Datatype& recvtype, int root) const = 0(binding deprecated,
20     see Section 15.2) }
21
22 {Request Comm::Iscatterv(const void* sendbuf, const int sendcounts[],
23     const int displs[], const Datatype& sendtype, void* recvbuf,
24     int recvcnt, const Datatype& recvtype, int root)
25     const = 0(binding deprecated, see Section 15.2) }
26
27 {bool Op::Is_commutative() const(binding deprecated, see Section 15.2) }
28
29 {void Comm::Reduce(const void* sendbuf, void* recvbuf, int count,
30     const Datatype& datatype, const Op& op, int root)
31     const = 0(binding deprecated, see Section 15.2) }
32
33 {void Op::Reduce_local(const void* inbuf, void* inoutbuf, int count,
34     const Datatype& datatype) const(binding deprecated, see
35     Section 15.2) }
36
37 {void Comm::Reduce_scatter_block(const void* sendbuf, void* recvbuf,
38     int recvcnt, const Datatype& datatype, const Op& op)
39     const = 0(binding deprecated, see Section 15.2) }
40
41 {void Comm::Reduce_scatter(const void* sendbuf, void* recvbuf,
42     int recvcnts[], const Datatype& datatype, const Op& op)
43     const = 0(binding deprecated, see Section 15.2) }
44
45 {void Intracomm::Scan(const void* sendbuf, void* recvbuf, int count,
46     const Datatype& datatype, const Op& op) const(binding deprecated,
47     see Section 15.2) }
48
49 {void Comm::Scatter(const void* sendbuf, int sendcount, const
50     Datatype& sendtype, void* recvbuf, int recvcnt,
51     const Datatype& recvtype, int root) const = 0(binding deprecated,
52     see Section 15.2) }

```

```

{void Comm::Scatterv(const void* sendbuf, const int sendcounts[],
    const int displs[], const Datatype& sendtype, void* recvbuf,
    int recvcnt, const Datatype& recvttype, int root)
    const = 0(binding deprecated, see Section 15.2) }

};

```

#### A.5.4 Groups, Contexts, Communicators, and Caching C++ Bindings

```

namespace MPI {

    {Comm& Comm::Clone() const = 0(binding deprecated, see Section 15.2) }

    {Cartcomm& Cartcomm::Clone() const(binding deprecated, see Section 15.2) }

    {Distgraphcomm& Distgraphcomm::Clone() const(binding deprecated, see
        Section 15.2) }

    {Graphcomm& Graphcomm::Clone() const(binding deprecated, see Section 15.2) }

    {Intercomm& Intercomm::Clone() const(binding deprecated, see Section 15.2) }

    {Intracomm& Intracomm::Clone() const(binding deprecated, see Section 15.2) }

    {static int Comm::Compare(const Comm& comm1, const Comm& comm2)(binding
        deprecated, see Section 15.2) }

    {static int Group::Compare(const Group& group1,
        const Group& group2)(binding deprecated, see Section 15.2) }

    {Intercomm Intercomm::Create(const Group& group) const(binding deprecated,
        see Section 15.2) }

    {Intracomm Intracomm::Create(const Group& group) const(binding deprecated,
        see Section 15.2) }

    {Intercomm Intracomm::Create_intercomm(int local_leader, const
        Comm& peer_comm, int remote_leader, int tag) const(binding
        deprecated, see Section 15.2) }

    {static int Comm::Create_keyval(Comm::Copy_attr_function*
        comm_copy_attr_fn,
        Comm::Delete_attr_function* comm_delete_attr_fn,
        void* extra_state)(binding deprecated, see Section 15.2) }

    {static int Datatype::Create_keyval(Datatype::Copy_attr_function*
        type_copy_attr_fn, Datatype::Delete_attr_function*
        type_delete_attr_fn, void* extra_state)(binding deprecated, see
        Section 15.2) }

    {static int Win::Create_keyval(Win::Copy_attr_function* win_copy_attr_fn,
        Win::Delete_attr_function* win_delete_attr_fn,
        void* extra_state)(binding deprecated, see Section 15.2) }

    {void Comm::Delete_attr(int comm_keyval)(binding deprecated, see Section 15.2) }
}

```



```

1  {void Datatype::Delete_attr(int type_keyval) (binding deprecated, see
2      Section 15.2) }
3
4  {void Win::Delete_attr(int win_keyval) (binding deprecated, see Section 15.2) }
5
6  {static Group Group::Difference(const Group& group1,
7      const Group& group2) (binding deprecated, see Section 15.2) }
8
9  {Cartcomm Cartcomm::Dup() const (binding deprecated, see Section 15.2) }
10
11 {Distgraphcomm Distgraphcomm::Dup() const (binding deprecated, see Section 15.2) }
12
13 {Graphcomm Graphcomm::Dup() const (binding deprecated, see Section 15.2) }
14
15 {Intercomm Intercomm::Dup() const (binding deprecated, see Section 15.2) }
16
17 {Intracomm Intracomm::Dup() const (binding deprecated, see Section 15.2) }
18
19 {Group Group::Excl(int n, const int ranks[]) const (binding deprecated, see
20     Section 15.2) }
21
22 {static void Comm::Free_keyval(int& comm_keyval) (binding deprecated, see
23     Section 15.2) }
24
25 {static void Datatype::Free_keyval(int& type_keyval) (binding deprecated, see
26     Section 15.2) }
27
28 {static void Win::Free_keyval(int& win_keyval) (binding deprecated, see
29     Section 15.2) }
30
31 {void Comm::Free() (binding deprecated, see Section 15.2) }
32
33 {void Group::Free() (binding deprecated, see Section 15.2) }
34
35 {bool Comm::Get_attr(int comm_keyval, void* attribute_val) const (binding
36     deprecated, see Section 15.2) }
37
38 {bool Datatype::Get_attr(int type_keyval, void* attribute_val)
39     const (binding deprecated, see Section 15.2) }
40
41 {bool Win::Get_attr(int win_keyval, void* attribute_val) const (binding
42     deprecated, see Section 15.2) }
43
44 {Group Comm::Get_group() const (binding deprecated, see Section 15.2) }
45
46 {void Comm::Get_name(char* comm_name, int& resultlen) const (binding
47     deprecated, see Section 15.2) }
48
49 {void Datatype::Get_name(char* type_name, int& resultlen) const (binding
50     deprecated, see Section 15.2) }
51
52 {void Win::Get_name(char* win_name, int& resultlen) const (binding deprecated,
53     see Section 15.2) }
54
55 {int Comm::Get_rank() const (binding deprecated, see Section 15.2) }
56
57 {int Group::Get_rank() const (binding deprecated, see Section 15.2) }

```



```

{Group Intercomm::Get_remote_group() const(binding deprecated, see Section 15.2)
    }
{int Intercomm::Get_remote_size() const(binding deprecated, see Section 15.2) }
{int Comm::Get_size() const(binding deprecated, see Section 15.2) }
{int Group::Get_size() const(binding deprecated, see Section 15.2) }
{Group Group::Incl(int n, const int ranks[]) const(binding deprecated, see
    Section 15.2) }
{static Group Group::Intersect(const Group& group1,
    const Group& group2)(binding deprecated, see Section 15.2) }
{bool Comm::Is_inter() const(binding deprecated, see Section 15.2) }
{Intracomm Intercomm::Merge(bool high) const(binding deprecated, see
    Section 15.2) }
{Group Group::Range_excl(int n, const int ranges[][3]) const(binding
    deprecated, see Section 15.2) }
{Group Group::Range_incl(int n, const int ranges[][3]) const(binding
    deprecated, see Section 15.2) }
{void Comm::Set_attr(int comm_keyval, const void* attribute_val)
    const(binding deprecated, see Section 15.2) }
{void Datatype::Set_attr(int type_keyval, const void*
    attribute_val)(binding deprecated, see Section 15.2) }
{void Win::Set_attr(int win_keyval, const void* attribute_val)(binding
    deprecated, see Section 15.2) }
{void Comm::Set_name(const char* comm_name)(binding deprecated, see
    Section 15.2) }
{void Datatype::Set_name(const char* type_name)(binding deprecated, see
    Section 15.2) }
{void Win::Set_name(const char* win_name)(binding deprecated, see Section 15.2) }
{Intercomm Intercomm::Split(int color, int key) const(binding deprecated, see
    Section 15.2) }
{Intracomm Intracomm::Split(int color, int key) const(binding deprecated, see
    Section 15.2) }
{static void Group::Translate_ranks (const Group& group1, int n,
    const int ranks1[], const Group& group2, int ranks2[])(binding
    deprecated, see Section 15.2) }
{static Group Group::Union(const Group& group1,
    const Group& group2)(binding deprecated, see Section 15.2) }

```

```
1 };
2
```

### 3 A.5.5 Process Topologies C++ Bindings

```
4 namespace MPI {
5
6     {void Compute_dims(int nnodes, int ndims, int dims[]) (binding deprecated, see
7         Section 15.2) }
8
9     {Cartcomm Intracomm::Create_cart(int ndims, const int dims[],
10         const bool periods[], bool reorder) const (binding deprecated, see
11         Section 15.2) }
12
13     {Graphcomm Intracomm::Create_graph(int nnodes, const int index[],
14         const int edges[], bool reorder) const (binding deprecated, see
15         Section 15.2) }
16
17     {Distgraphcomm Intracomm::Dist_graph_create_adjacent(int indegree,
18         const int sources[], const int sourceweights[], int outdegree,
19         const int destinations[], const int destweights[],
20         const Info& info, bool reorder) const (binding deprecated, see
21         Section 15.2) }
22
23     {Distgraphcomm Intracomm::Dist_graph_create_adjacent(int indegree,
24         const int sources[], int outdegree, const int destinations[],
25         const Info& info, bool reorder) const (binding deprecated, see
26         Section 15.2) }
27
28     {Distgraphcomm Intracomm::Dist_graph_create(int n, const int sources[],
29         const int degrees[], const int destinations[],
30         const int weights[], const Info& info, bool reorder)
31         const (binding deprecated, see Section 15.2) }
32
33     {Distgraphcomm Intracomm::Dist_graph_create(int n, const int sources[],
34         const int degrees[], const int destinations[],
35         const Info& info, bool reorder) const (binding deprecated, see
36         Section 15.2) }
37
38     {int Cartcomm::Get_cart_rank(const int coords[]) const (binding deprecated,
39         see Section 15.2) }
40
41     {void Cartcomm::Get_coords(int rank, int maxdims, int coords[])
42         const (binding deprecated, see Section 15.2) }
43
44     {int Cartcomm::Get_dim() const (binding deprecated, see Section 15.2) }
45
46     {void Graphcomm::Get_dims(int nnodes[], int nedges[]) const (binding
47         deprecated, see Section 15.2) }
48
49     {void Distgraphcomm::Get_dist_neighbors_count(int rank, int indegree[],
50         int outdegree[], bool& weighted) const (binding deprecated, see
51         Section 15.2) }
```

```

{void Distgraphcomm::Get_dist_neighbors(int maxindegree, int sources[],
    int sourceweights[], int maxoutdegree, int destinations[],
    int destweights[]) (binding deprecated, see Section 15.2) }
{int Graphcomm::Get_neighbors_count(int rank) const (binding deprecated, see
    Section 15.2) }
{void Graphcomm::Get_neighbors(int rank, int maxneighbors, int
    neighbors[]) const (binding deprecated, see Section 15.2) }
{void Cartcomm::Get_topo(int maxdims, int dims[], bool periods[],
    int coords[]) const (binding deprecated, see Section 15.2) }
{void Graphcomm::Get_topo(int maxindex, int maxedges, int index[],
    int edges[]) const (binding deprecated, see Section 15.2) }
{int Comm::Get_topology() const (binding deprecated, see Section 15.2) }
{int Cartcomm::Map(int ndims, const int dims[], const bool periods[])
    const (binding deprecated, see Section 15.2) }
{int Graphcomm::Map(int nnodes, const int index[], const int edges[])
    const (binding deprecated, see Section 15.2) }
{void Cartcomm::Shift(int direction, int disp, int& rank_source,
    int& rank_dest) const (binding deprecated, see Section 15.2) }
{Cartcomm Cartcomm::Sub(const bool remain_dims[]) const (binding deprecated,
    see Section 15.2) }
};

```

#### A.5.6 MPI Environmental Management C++ Bindings

```

namespace MPI {
{void Comm::Abort(int errorcode) (binding deprecated, see Section 15.2) }
{int Add_error_class() (binding deprecated, see Section 15.2) }
{int Add_error_code(int errorclass) (binding deprecated, see Section 15.2) }
{void Add_error_string(int errorcode, const char* string) (binding deprecated,
    see Section 15.2) }
{void* Alloc_mem(Aint size, const Info& info) (binding deprecated, see
    Section 15.2) }
{void Comm::Call_errhandler(int errorcode) const (binding deprecated, see
    Section 15.2) }
{void File::Call_errhandler(int errorcode) const (binding deprecated, see
    Section 15.2) }
{void Win::Call_errhandler(int errorcode) const (binding deprecated, see
    Section 15.2) }
}

```

```

1      {static Errhandler Comm::Create_errhandler(Comm::Errhandler_function*
ticket252-W. 2          [function]comm_errhandler_fn)(binding deprecated, see Section 15.2) }
3
4      {static Errhandler File::Create_errhandler(File::Errhandler_function*
ticket252-W. 5          [function]file_errhandler_fn)(binding deprecated, see Section 15.2) }
6
7      {static Errhandler Win::Create_errhandler(Win::Errhandler_function*
ticket252-W. 8          [function]win_errhandler_fn)(binding deprecated, see Section 15.2) }
9
10     {void Finalize()(binding deprecated, see Section 15.2) }
11
12     {void Free_mem(void *base)(binding deprecated, see Section 15.2) }
13
14     {void Errhandler::Free()(binding deprecated, see Section 15.2) }
15
16     {Errhandler Comm::Get_errhandler() const(binding deprecated, see Section 15.2) }
17
18     {Errhandler File::Get_errhandler() const(binding deprecated, see Section 15.2) }
19
20     {Errhandler Win::Get_errhandler() const(binding deprecated, see Section 15.2) }
21
22     {int Get_error_class(int errorcode)(binding deprecated, see Section 15.2) }
23
24     {void Get_error_string(int errorcode, char* name, int& resultlen)(binding
25         deprecated, see Section 15.2) }
26
27     {void Get_processor_name(char* name, int& resultlen)(binding deprecated, see
28         Section 15.2) }
29
30     {void Get_version(int& version, int& subversion)(binding deprecated, see
31         Section 15.2) }
32
33     {void Init(int& argc, char**& argv)(binding deprecated, see Section 15.2) }
34
35     {void Init()(binding deprecated, see Section 15.2) }
36
37     {bool Is_finalized()(binding deprecated, see Section 15.2) }
38
39     {bool Is_initialized()(binding deprecated, see Section 15.2) }
40
41     {void Comm::Set_errhandler(const Errhandler& errhandler)(binding deprecated,
42         see Section 15.2) }
43
44     {void File::Set_errhandler(const Errhandler& errhandler)(binding deprecated,
45         see Section 15.2) }
46
47     {void Win::Set_errhandler(const Errhandler& errhandler)(binding deprecated,
48         see Section 15.2) }
49
50     {double Wtick()(binding deprecated, see Section 15.2) }
51
52     {double Wtime()(binding deprecated, see Section 15.2) }
53
54 };

```

## A.5.7 The Info Object C++ Bindings

```
namespace MPI {
```

```

{static Info Info::Create() (binding deprecated, see Section 15.2) }
{void Info::Delete(const char* key) (binding deprecated, see Section 15.2) }
{Info Info::Dup() const (binding deprecated, see Section 15.2) }
{void Info::Free() (binding deprecated, see Section 15.2) }
{bool Info::Get(const char* key, int valuelen, char* value) const (binding
    deprecated, see Section 15.2) }
{int Info::Get_nkeys() const (binding deprecated, see Section 15.2) }
{void Info::Get_nthkey(int n, char* key) const (binding deprecated, see
    Section 15.2) }
{bool Info::Get_valuelen(const char* key, int& valuelen) const (binding
    deprecated, see Section 15.2) }
{void Info::Set(const char* key, const char* value) (binding deprecated, see
    Section 15.2) }
};

```

#### A.5.8 Process Creation and Management C++ Bindings

```

namespace MPI {
    {Intercomm Intracomm::Accept(const char* port_name, const Info& info,
        int root) const (binding deprecated, see Section 15.2) }
    {void Close_port(const char* port_name) (binding deprecated, see Section 15.2) }
    {Intercomm Intracomm::Connect(const char* port_name, const Info& info,
        int root) const (binding deprecated, see Section 15.2) }
    {void Comm::Disconnect() (binding deprecated, see Section 15.2) }
    {static Intercomm Comm::Get_parent() (binding deprecated, see Section 15.2) }
    {static Intercomm Comm::Join(const int fd) (binding deprecated, see Section 15.2)
        }
    {void Lookup_name(const char* service_name, const Info& info,
        char* port_name) (binding deprecated, see Section 15.2) }
    {void Open_port(const Info& info, char* port_name) (binding deprecated, see
        Section 15.2) }
    {void Publish_name(const char* service_name, const Info& info,
        const char* port_name) (binding deprecated, see Section 15.2) }
    {Intercomm Intracomm::Spawn(const char* command, const char* argv[],
        int maxprocs, const Info& info, int root) const (binding
        deprecated, see Section 15.2) }
}

```

```

1  {Intercomm Intracomm::Spawn(const char* command, const char* argv[],
2      int maxprocs, const Info& info, int root,
3      int array_of_errcodes[]) const(binding deprecated, see Section 15.2)
4      }
5
6  {Intercomm Intracomm::Spawn_multiple(int count,
7      const char* array_of_commands[], const char** array_of_argv[],
8      const int array_of_maxprocs[], const Info array_of_info[],
9      int root, int array_of_errcodes[]) (binding deprecated, see
10     Section 15.2) }
11
12 {Intercomm Intracomm::Spawn_multiple(int count,
13     const char* array_of_commands[], const char** array_of_argv[],
14     const int array_of_maxprocs[], const Info array_of_info[],
15     int root) (binding deprecated, see Section 15.2) }
16
17 {void Unpublish_name(const char* service_name, const Info& info,
18     const char* port_name) (binding deprecated, see Section 15.2) }
19 };
20

```

#### A.5.9 One-Sided Communications C++ Bindings

```

21 namespace MPI {
22
23     {void Win::Accumulate(const void* origin_addr, int origin_count, const
24         Datatype& origin_datatype, int target_rank, Aint target_disp,
25         int target_count, const Datatype& target_datatype, const Op&
26         op) const(binding deprecated, see Section 15.2) }
27
28     {void Win::Complete() const(binding deprecated, see Section 15.2) }
29
30     {static Win Win::Create(const void* base, Aint size, int disp_unit, const
31         Info& info, const Intracomm& comm) (binding deprecated, see
32         Section 15.2) }
33
34     {void Win::Fence(int assert) const(binding deprecated, see Section 15.2) }
35
36     {void Win::Free() (binding deprecated, see Section 15.2) }
37
38     {Group Win::Get_group() const(binding deprecated, see Section 15.2) }
39
40     {void Win::Get(void *origin_addr, int origin_count,
41         const Datatype& origin_datatype, int target_rank,
42         Aint target_disp, int target_count,
43         const Datatype& target_datatype) const(binding deprecated, see
44         Section 15.2) }
45
46     {void Win::Lock(int lock_type, int rank, int assert) const(binding
47         deprecated, see Section 15.2) }
48
49     {void Win::Post(const Group& group, int assert) const(binding deprecated, see
50         Section 15.2) }
51

```

```

{void Win::Put(const void* origin_addr, int origin_count,
               const Datatype& origin_datatype, int target_rank,
               Aint target_disp, int target_count,
               const Datatype& target_datatype) const(binding deprecated, see
               Section 15.2) }
{void Win::Start(const Group& group, int assert) const(binding deprecated,
               see Section 15.2) }
{bool Win::Test() const(binding deprecated, see Section 15.2) }
{void Win::Unlock(int rank) const(binding deprecated, see Section 15.2) }
{void Win::Wait() const(binding deprecated, see Section 15.2) }
};

```

#### A.5.10 External Interfaces C++ Bindings

```

namespace MPI {
    {void Grequest::Complete() (binding deprecated, see Section 15.2) }
    {int Init_thread(int& argc, char**& argv, int required) (binding deprecated,
        see Section 15.2) }
    {int Init_thread(int required) (binding deprecated, see Section 15.2) }
    {bool Is_thread_main() (binding deprecated, see Section 15.2) }
    {int Query_thread() (binding deprecated, see Section 15.2) }
    {void Status::Set_cancelled(bool flag) (binding deprecated, see Section 15.2) }
    {void Status::Set_elements(const Datatype& datatype, int count) (binding
        deprecated, see Section 15.2) }
    {static Grequest Grequest::Start(const Grequest::Query_function*
        query_fn, const Grequest::Free_function* free_fn,
        const Grequest::Cancel_function* cancel_fn,
        void *extra_state) (binding deprecated, see Section 15.2) }
};

```

#### A.5.11 I/O C++ Bindings

```

namespace MPI {
    {void File::Close() (binding deprecated, see Section 15.2) }
    {static void File::Delete(const char* filename, const Info& info) (binding
        deprecated, see Section 15.2) }
    {int File::Get_amode() const(binding deprecated, see Section 15.2) }
}

```



```

1      {bool File::Get_atomicity() const(binding deprecated, see Section 15.2) }
2
3      {Offset File::Get_byte_offset(const Offset disp) const(binding deprecated,
4          see Section 15.2) }
5
6      {Group File::Get_group() const(binding deprecated, see Section 15.2) }
7
8      {Info File::Get_info() const(binding deprecated, see Section 15.2) }
9
10     {Offset File::Get_position() const(binding deprecated, see Section 15.2) }
11
12     {Offset File::Get_position_shared() const(binding deprecated, see Section 15.2) }
13
14     {Offset File::Get_size() const(binding deprecated, see Section 15.2) }
15
16     {Aint File::Get_type_extent(const Datatype& datatype) const(binding
17         deprecated, see Section 15.2) }
18
19     {void File::Get_view(Offset& disp, Datatype& etype, Datatype& filetype,
20         char* datarep) const(binding deprecated, see Section 15.2) }
21
22     {Request File::Iread_at(Offset offset, void* buf, int count,
23         const Datatype& datatype)(binding deprecated, see Section 15.2) }
24
25     {Request File::Iread_shared(void* buf, int count,
26         const Datatype& datatype)(binding deprecated, see Section 15.2) }
27
28     {Request File::Iread(void* buf, int count,
29         const Datatype& datatype)(binding deprecated, see Section 15.2) }
30
31     {Request File::Iwrite_at(Offset offset, const void* buf, int count,
32         const Datatype& datatype)(binding deprecated, see Section 15.2) }
33
34     {Request File::Iwrite(const void* buf, int count,
35         const Datatype& datatype)(binding deprecated, see Section 15.2) }
36
37     {Request File::Iwrite_shared(const void* buf, int count,
38         const Datatype& datatype)(binding deprecated, see Section 15.2) }
39
40     {static File File::Open(const Intracomm& comm, const char* filename,
41         int amode, const Info& info)(binding deprecated, see Section 15.2) }
42
43     {void File::Preallocate(Offset size)(binding deprecated, see Section 15.2) }
44
45     {void File::Read_all_begin(void* buf, int count,
46         const Datatype& datatype)(binding deprecated, see Section 15.2) }
47
48     {void File::Read_all_end(void* buf, Status& status)(binding deprecated, see
49         Section 15.2) }
50
51     {void File::Read_all_end(void* buf)(binding deprecated, see Section 15.2) }
52
53     {void File::Read_all(void* buf, int count, const Datatype& datatype,
54         Status& status)(binding deprecated, see Section 15.2) }
55
56     {void File::Read_all(void* buf, int count,
57         const Datatype& datatype)(binding deprecated, see Section 15.2) }

```



```

{void File::Read_at_all_begin(Offset offset, void* buf, int count,
    const Datatype& datatype) (binding deprecated, see Section 15.2) }
{void File::Read_at_all_end(void* buf, Status& status) (binding deprecated,
    see Section 15.2) }
{void File::Read_at_all_end(void* buf) (binding deprecated, see Section 15.2) }
{void File::Read_at_all(Offset offset, void* buf, int count,
    const Datatype& datatype, Status& status) (binding deprecated, see
    Section 15.2) }
{void File::Read_at_all(Offset offset, void* buf, int count,
    const Datatype& datatype) (binding deprecated, see Section 15.2) }
{void File::Read_at(Offset offset, void* buf, int count,
    const Datatype& datatype, Status& status) (binding deprecated, see
    Section 15.2) }
{void File::Read_at(Offset offset, void* buf, int count,
    const Datatype& datatype) (binding deprecated, see Section 15.2) }
{void File::Read_ordered_begin(void* buf, int count,
    const Datatype& datatype) (binding deprecated, see Section 15.2) }
{void File::Read_ordered_end(void* buf, Status& status) (binding deprecated,
    see Section 15.2) }
{void File::Read_ordered_end(void* buf) (binding deprecated, see Section 15.2) }
{void File::Read_ordered(void* buf, int count, const Datatype& datatype,
    Status& status) (binding deprecated, see Section 15.2) }
{void File::Read_ordered(void* buf, int count,
    const Datatype& datatype) (binding deprecated, see Section 15.2) }
{void File::Read_shared(void* buf, int count, const Datatype& datatype,
    Status& status) (binding deprecated, see Section 15.2) }
{void File::Read_shared(void* buf, int count,
    const Datatype& datatype) (binding deprecated, see Section 15.2) }
{void File::Read(void* buf, int count, const Datatype& datatype, Status&
    status) (binding deprecated, see Section 15.2) }
{void File::Read(void* buf, int count, const Datatype& datatype) (binding
    deprecated, see Section 15.2) }
{void Register_datarep(const char* datarep,
    Datarep_conversion_function* read_conversion_fn,
    Datarep_conversion_function* write_conversion_fn,
    Datarep_extent_function* dtype_file_extent_fn,
    void* extra_state) (binding deprecated, see Section 15.2) }

```

```

1  {void File::Seek(Offset offset, int whence) (binding deprecated, see
2      Section 15.2) }
3
4  {void File::Seek_shared(Offset offset, int whence) (binding deprecated, see
5      Section 15.2) }
6
7  {void File::Set_atomicity(bool flag) (binding deprecated, see Section 15.2) }
8  {void File::Set_info(const Info& info) (binding deprecated, see Section 15.2) }
9  {void File::Set_size(Offset size) (binding deprecated, see Section 15.2) }
10
11 {void File::Set_view(Offset disp, const Datatype& etype,
12     const Datatype& filetype, const char* datarep,
13     const Info& info) (binding deprecated, see Section 15.2) }
14
15 {void File::Sync() (binding deprecated, see Section 15.2) }
16
17 {void File::Write_all_begin(const void* buf, int count,
18     const Datatype& datatype) (binding deprecated, see Section 15.2) }
19
20 {void File::Write_all(const void* buf, int count,
21     const Datatype& datatype, Status& status) (binding deprecated, see
22     Section 15.2) }
23
24 {void File::Write_all(const void* buf, int count,
25     const Datatype& datatype) (binding deprecated, see Section 15.2) }
26
27 {void File::Write_all_end(const void* buf, Status& status) (binding
28     deprecated, see Section 15.2) }
29
30 {void File::Write_at_all_begin(Offset offset, const void* buf, int count,
31     const Datatype& datatype) (binding deprecated, see Section 15.2) }
32
33 {void File::Write_at_all_end(const void* buf, Status& status) (binding
34     deprecated, see Section 15.2) }
35
36 {void File::Write_at_all_end(const void* buf) (binding deprecated, see
37     Section 15.2) }
38
39 {void File::Write_at_all(Offset offset, const void* buf, int count,
40     const Datatype& datatype, Status& status) (binding deprecated, see
41     Section 15.2) }
42
43 {void File::Write_at_all(Offset offset, const void* buf, int count,
44     const Datatype& datatype) (binding deprecated, see Section 15.2) }
45
46 {void File::Write_at(Offset offset, const void* buf, int count,
47     const Datatype& datatype, Status& status) (binding deprecated, see
48     Section 15.2) }

```

```

{void File::Write(const void* buf, int count, const Datatype& datatype,
                  Status& status) (binding deprecated, see Section 15.2) }
{void File::Write(const void* buf, int count,
                  const Datatype& datatype) (binding deprecated, see Section 15.2) }
{void File::Write_ordered_begin(const void* buf, int count,
                                const Datatype& datatype) (binding deprecated, see Section 15.2) }
{void File::Write_ordered(const void* buf, int count,
                           const Datatype& datatype, Status& status) (binding deprecated, see
                           Section 15.2) }
{void File::Write_ordered(const void* buf, int count,
                           const Datatype& datatype) (binding deprecated, see Section 15.2) }
{void File::Write_ordered_end(const void* buf, Status& status) (binding
                              deprecated, see Section 15.2) }
{void File::Write_ordered_end(const void* buf) (binding deprecated, see
                              Section 15.2) }
{void File::Write_shared(const void* buf, int count,
                          const Datatype& datatype, Status& status) (binding deprecated, see
                          Section 15.2) }
{void File::Write_shared(const void* buf, int count,
                          const Datatype& datatype) (binding deprecated, see Section 15.2) }
};

```

#### A.5.12 Language Bindings C++ Bindings

```

namespace MPI {
    {static Datatype Datatype::Create_f90_complex(int p, int r) (binding
        deprecated, see Section 15.2) }
    {static Datatype Datatype::Create_f90_integer(int r) (binding deprecated, see
        Section 15.2) }
    {static Datatype Datatype::Create_f90_real(int p, int r) (binding deprecated,
        see Section 15.2) }
    Exception::Exception(int error_code)
    {int Exception::Get_error_class() const (binding deprecated, see Section 15.2) }
    {int Exception::Get_error_code() const (binding deprecated, see Section 15.2) }
    {const char* Exception::Get_error_string() const (binding deprecated, see
        Section 15.2) }
    {static Datatype Datatype::Match_size(int typeclass, int size) (binding
        deprecated, see Section 15.2) }
}

```

```
1 };
```

### 2 A.5.13 Profiling Interface C++ Bindings

```
3 namespace MPI {
4
5     {void Pcontrol(const int level, ...) (binding deprecated, see Section 15.2) }
6
7
8 };
```

ticket11.

[\[C++ Deprecated Functions section\]](#)

### 11 A.5.14 C++ Bindings on all MPI Classes

12 The C++ language requires all classes to have four special functions: a default constructor, a copy constructor, a destructor, and an assignment operator. The bindings for these functions are listed below; their semantics are discussed in Section 16.1.5. The two constructors are *not virtual*. The bindings prototype functions are using the type `<CLASS>` rather than listing each function for every MPI class. The token `<CLASS>` can be replaced with valid MPI-2 class names, such as `Group`, `Datatype`, etc., except when noted. In addition, bindings are provided for comparison and inter-language operability from Sections 16.1.5 and 16.1.9.

### 21 A.5.15 Construction / Destruction

```
22 namespace MPI {
23
24     <CLASS>::<CLASS>()
25
26     <CLASS>::~~<CLASS>()
27
28 };
```

### 30 A.5.16 Copy / Assignment

```
31 namespace MPI {
32
33     <CLASS>::<CLASS>(const <CLASS>& data)
34
35     <CLASS>& <CLASS>::operator=(const <CLASS>& data)
36
37 };
```

### 38 A.5.17 Comparison

39 Since `Status` instances are not handles to underlying MPI objects, the `operator==( )` and `operator!=( )` functions are not defined on the `Status` class.

```
40 namespace MPI {
41
42
43     bool <CLASS>::operator==(const <CLASS>& data) const
44
45     bool <CLASS>::operator!=(const <CLASS>& data) const
46
47
48 };
```

## A.5.18 Inter-language Operability

Since there are no C++ `MPI::STATUS_IGNORE` and `MPI::STATUSES_IGNORE` objects, the result of promoting the C or Fortran handles (`MPI_STATUS_IGNORE` and `MPI_STATUSES_IGNORE`) to C++ is undefined.

```
namespace MPI {
    <CLASS>& <CLASS>::operator=(const MPI_<CLASS>& data)
    <CLASS>::<CLASS>(const MPI_<CLASS>& data)
    <CLASS>::operator MPI_<CLASS>() const
};
```

1  
2  
3  
4  
5  
6  
7  
8  
9  
10  
11  
12  
13  
14  
15  
16  
17  
18  
19  
20  
21  
22  
23  
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48

# Annex B

## Change-Log

This annex summarizes changes from the previous version of the MPI standard to the version presented by this document. Only significant changes (i.e., clarifications and new features) that might either require implementation effort in the MPI libraries or change the understanding of MPI from a user’s perspective are presented. Editorial modifications, formatting, typo corrections and minor clarifications are not shown.

### B.1 Changes from Version 2.2 to Version 3.0

1. Section 16.1.6 on page 631, and MPI-2.2 Section 16.1.16 on page 471, line 45.  
This is an MPI-2.2 errata: The constant `MPI::_LONG_LONG` should be `MPI::LONG_LONG`.  
TICKET NOT YET PASSED. NEW CHANGE-LOG TEXT.
2. Section 13.5.2, Table 13.2 on page 560, and MPI-2.2, Section 13.5.3, Table 13.2 on page 433.  
This is an MPI-2.2 errata: The `MPI_C_BOOL` “external32” representation is corrected to a 1-byte size.  
TICKET PASSED. NEW CHANGE-LOG TEXT.
3. Section 7.5.5 on page 315, and MPI-2.2, Section 7.5.5 on page 257, C++ interface on page 264, line 3.  
This is an MPI-2.2 errata: In the C++ interface of `MPI_DIST_GRAPH_NEIGHBORS_COUNT`, the argument `rank` is removed.  
TICKET NOT YET PASSED (Had 1st vote). NEW CHANGE-LOG TEXT.
4. Annex A.1.1 on page 709, Table “Optional datatypes (Fortran)”, and MPI-2.2, Annex A.1.1, Table on page 517, lines 34, and 37-41.  
This is an MPI-2.2 errata: The C++ datatype handles `MPI::INTEGER16`, `MPI::REAL16`, `MPI::F_COMPLEX4`, `MPI::F_COMPLEX8`, `MPI::F_COMPLEX16`, `MPI::F_COMPLEX32` where added to the table.  
TICKET NOT YET PASSED. NEW CHANGE-LOG TEXT.
5. Sections 3.8.2, 3.8.3, 16.3.4, A.1.1 on pages 75, 76, 693, 709.  
Like `MPI_PROBE` and `MPI_IProbe`, the new `MPI_MProbe` and `MPI_IMProbe` operations allow incoming messages to be queried without actually receiving them, except that `MPI_MProbe` and `MPI_IMProbe` provide a mechanism

to receive the specific message with the new routine `MPI_MRECV` regardless of other intervening probe or receive operations. The opaque object `MPI_Message`, the null handle `MPI_MESSAGE_NULL`, and the conversion functions `MPI_Message_c2f` and `MPI_Message_f2c` are defined.

TICKETS 38+274 PASSED. NEW CHANGE-LOG TEXT.

6. Chapter 5 on page 151 and Section 5.12 on page 209.  
Added nonblocking interfaces to all collective operations.

TEXT AS PASSED.

7. Section 2.3 on page 10.  
Clarified parameter usage for IN parameters. C bindings are now const-correct where backward compatibility is preserved.

TICKET AS PASSED. Reference updated from Section 2 to Section 2.3

8. Chapter 3 on page 27 until Chapter 16 on page 627.  
In the C language bindings, the array-arguments' interfaces are modified to consistently to always use `[]` instead of `*`, and the `'const'` keyword has been added to many functions.

TICKET PASSED. NEW CHANGE-LOG TEXT.

9. Section 7.5.8 on page 326.  
`MPI_CART_MAP` can also be used for a zero-dimensional topologies.

TEXT AS PASSED.

10. Section 6.4.2 on page 252.  
Added `MPI_COMM_IDUP`.

TEXT AS PASSED.

11. Section 2.5.4 on page 15 and Section 8.1.1 on page 349.  
Added new routine `MPI_GET_LIBRARY_VERSION` to query library specific versions, and the constant `MPI_MAX_LIBRARY_VERSION_STRING`.

TICKET PASSED. MODIFIED CHANGE-LOG TEXT.

12. Section 6.8 on page 294.  
Section 6.8 on page 238. The constant `MPI_MAX_OBJECT_NAME` also applies for type and window names.

13. Section ?? on page ??.

I ASKED GEORGE TO SET THE MISSING LABEL AT "12.4.3 Initialization"

IT MUST BE DECIDED, WHICH OPTION WE USE ABOUT SAME  
required ARGUMENT WHEN CALLING `MPI_INIT_THREAD`.

TICKET NOT YET PASSED. CHANGE-LOG TEXT MUST BE ALSO DEFINED.

14. Section 3.8 on page 71 and Section 3.11 on page 89.  
The use of `MPI_PROC_NULL` in probe and matching probe operations was clarified. A special predefined message `MPI_MESSAGE_NO_PROC` is defined for the use of matching probe with `MPI_PROC_NULL`.

TICKET 256 NOT YET PASSED (Had 1st vote). CHANGE-LOG TEXT AS DEFINED IN TICKET.

TICKET 328 NOT YET PASSED (Had 1st vote). NEW CHANGE-LOG TEXT.



ticket299.  
ticket258.

15. Section 7.6 on page 328 and Section 7.7 on page 337.  
The neighborhood collective communication routines are added to support sparse communication on virtual topology grids: MPI\_NEIGHBOR\_ALLGATHER, MPI\_NEIGHBOR\_ALLGATHERV, MPI\_NEIGHBOR\_ALLTOALL, MPI\_NEIGHBOR\_ALLTOALLV, MPI\_NEIGHBOR\_ALLTOALLW and the nonblocking variants MPI\_INEIGHBOR\_ALLGATHER, MPI\_INEIGHBOR\_ALLGATHERV, MPI\_INEIGHBOR\_ALLTOALL, MPI\_INEIGHBOR\_ALLTOALLV, and MPI\_INEIGHBOR\_ALLTOALLW. The displacement arguments in MPI\_NEIGHBOR\_ALLTOALLW and MPI\_INEIGHBOR\_ALLTOALLW are defined as address size integers. In MPI\_DIST\_GRAPH\_NEIGHBORS, an ordering rule is added for communicators created with MPI\_DIST\_GRAPH\_CREATE\_ADJACENT.

TICKET PASSED. MODIFIED CHANGE-LOG TEXT.

16. Sections 2.5.8, 3.2.2, 3.3, 5.9.2, on pages 17, 29, 31, 188, Sections ??, ??, ??, ??, 12.3 on pages ??, ??, ??, ??, 499, and Annex A.1.1 on page 709.  
New inquiry functions, MPI\_TYPE\_SIZE\_X, MPI\_TYPE\_GET\_EXTENT\_X, MPI\_TYPE\_GET\_TRUE\_EXTENT\_X, and MPI\_GET\_ELEMENTS\_X, return their results as an MPI\_Count value, which is a new type large enough to represent element counts in memory, file views, etc. A new function, MPI\_STATUS\_SET\_ELEMENTS\_X, modifies the opaque part of MPI\_STATUS so that a call to MPI\_GET\_ELEMENTS\_X returns the provided MPI\_Count value (in Fortran, INTEGER (KIND=MPI\_COUNT\_KIND)). The corresponding predefined datatype is MPI\_COUNT.

TICKET PASSED. CHANGE-LOG TEXT EXTENDED BY Fortran AND predefined datatype.

17. Sections ??, ??, ??, ?? on pages ??, ??, ??, ??.  
The functions MPI\_GET\_COUNT, MPI\_TYPE\_SIZE, and MPI\_GET\_ELEMENTS are now defined to set the count parameter to MPI\_UNDEFINED when that parameter would overflow. The function MPI\_PACK\_SIZE is now defined to set the size parameter to MPI\_UNDEFINED when that parameter would overflow. In all other MPI-2.2 routines, the type and semantics of the count arguments are kept unchanged, i.e., int or INTEGER.

TEXT AS PASSED.

18. Section 8.7 on page 372.  
Allow calls to MPI\_T routines before MPI\_INIT and after MPI\_FINALIZE.

TICKET PASSED. MODIFIED CHANGE-LOG TEXT.

19. Section 14.3 on page 588.  
A new MPI Tool Information Interface is added.

TICKET PASSED. MODIFIED CHANGE-LOG TEXT.

20. Chapter 11 on page 419.  
Substantial revision of the entire One-sided chapter, with new routines for window creation, additional synchronization methods in passive target, new one-sided communication routines, a new memory model, and other changes.

Ticket 270, TEXT AS PASSED. Ticket 284, NO ADDITIONAL CHANGE-LOG.

21. Sections 6.4.2, ??, ??, on pages 252, ??, ??.

I ASKED ADAM FOR THE MISSING LABELS of “6.4.4 Communicator Info” and “11.2.3 Window Info”

The new routines `MPI_COMM_DUP_INFO`, `MPI_COMM_SET_INFO`, `MPI_COMM_GET_INFO`, `MPI_WIN_SET_INFO`, and `MPI_WIN_GET_INFO` are added. The routine `MPI_COMM_DUP` must also duplicate topology information and info hints.

TICKET NOT YET PASSED (Had 1st vote). NEW CHANGE-LOG TEXT.

22. Chapter 16.1.1 on page 627.

Added a sentence making the C++ bindings optional.

TEXT AS PASSED.

CAUTION: MAY BE OBSOLETE WITH TICKET 281!

23. Section 4.1.2 on page 93 and Section 4.1.13 on page 123.

The routine `MPI_TYPE_CREATE_HINDEXED_BLOCK` and constant `MPI_COMBINER_HINDEXED_BLOCK` are added.

TICKET PASSED. MODIFIED CHANGE-LOG TEXT.

24. Section ?? on page ?? and all other chapters.

THIS LABEL IS ONLY AVAILABLE AFTER TICKET 281 IS INCLUDED.

The C++ bindings are removed from the standard. See MPI-2.2 errata at the beginning of this list for latest changes to the MPI C++ binding defined in MPI-2.2.

TICKET NOT YET PASSED (Had 1st vote). NEW CHANGE-LOG TEXT.

25. Section 6.4.2 on page 252.

New communicator construction routine `MPI_COMM_CREATE_GROUP`, which is invoked only by the processes in the group of the new communicator being constructed.

TICKET PASSED. NEW CHANGE-LOG TEXT.

26. Section 6.4.2 on page 252.

Added `MPI_COMM_SPLIT_TYPE` function and the communicator split type constant `MPI_COMM_TYPE_SHARED`.

TICKET PASSED. MODIFIED CHANGE-LOG TEXT.

27. Section 2.5.4 on page 15 and Section 7.5.4 on page 308.

The recommended C implementation value for `MPI_UNWEIGHTED` was changed from `NULL` to `non-NULL`. An additional weight array constant (`MPI_WEIGHTS_EMPTY`) was introduced.

TICKET NOT YET PASSED (Had 1st vote). CHANGE-LOG TEXT AS DEFINED IN TICKET.

28. Section 15.1 on page 617 and Section ?? on page ??.

The deprecated functions `MPI_TYPE_HVECTOR`, `MPI_TYPE_HINDEXED`, `MPI_TYPE_STRUCT`, `MPI_ADDRESS`, `MPI_TYPE_EXTENT`, `MPI_TYPE_LB`, `MPI_TYPE_UB`, `MPI_ERRHANDLER_CREATE` (and its callback function prototype `MPI_Handler_function`), `MPI_ERRHANDLER_SET`, `MPI_ERRHANDLER_GET`, the deprecated special datatype handles `MPI_LB`, `MPI_UB`, and the constants `MPI_COMBINER_HINDEXED_INTEGER`, `MPI_COMBINER_HVECTOR_INTEGER`, `MPI_COMBINER_STRUCT_INTEGER` are removed from the standard.

TICKET NOT YET PASSED (Had 1st vote). NEW CHANGE-LOG TEXT.

29. Section 6.6.2 on page 271.

The scope of the `tag` argument in `MPI_INTERCOMM_CREATE` is shrunk to the use

in this routine.

TICKET PASSED. NEW CHANGE-LOG TEXT.

30. Section 8.7 on page 372 and Section ?? on page ??.

I ASKED GEORGE TO SET THE MISSING LABEL AT “12.4.3 Initialization”

The use of MPI\_INIT, MPI\_INIT\_THREAD and MPI\_FINALIZE is clarified. After MPI is initialized, the application can access information about the execution environment by querying the new predefined info object MPI\_INFO\_GET\_ENV.

TICKET NOT YET PASSED (Had 1st vote). NEW CHANGE-LOG TEXT.

CAUTION: There are three locations of MPI\_INFO\_KEY. I expect that this is a typo and should mean MPI\_INFO\_GET\_ENV.

31. Sections 5.9.2, 13.5.2 Table 13.2, and Annex A.1.1 on pages 188, 560, and 709.

New named optional predefined datatypes MPI\_QUAD, MPI\_C\_QUAD\_COMPLEX, MPI\_FLOAT128, and MPI\_C\_FLOAT128\_COMPLEX for the C types \_Quad, \_\_float128, \_Quad\_Complex, and \_\_float128\_Complex.

TICKET NOT YET PASSED. CHANGE-LOG TEXT AS ON TICKET.

32. Section 6.7.2 on page 278.

Section 6.7.2 on page 226. It was clarified that in Fortran, the flag values of a comm\_copy\_attr\_fn callback and of MPI\_COMM\_NULL\_COPY\_FN and MPI\_COMM\_DUP\_FN are .FALSE. and .TRUE.; see MPI\_COMM\_CREATE\_KEYVAL.

TICKET PASSED. MODIFIED CHANGE-LOG TEXT.

33. Sections 5.9.2, 13.5.2 Table 13.2, and Annex A.1.1 on pages 188, 560, and 709, and MPI-2.2 Sections 5.9.2, 13.5.2 Table 13.2, and Annex A.1.1 on pages 164, 433, and 513.

MPI-2.2 errata: New named predefined datatypes MPI\_CXX\_BOOL, MPI\_CXX\_FLOAT\_COMPLEX, MPI\_CXX\_DOUBLE\_COMPLEX, and MPI\_CXX\_LONG\_DOUBLE\_COMPLEX in C and Fortran for the C++ types bool, Complex<float>, Complex<double>, and Complex<long double>, corresponding to the deprecated C++ predefined datatypes MPI::BOOL, MPI::COMPLEX, MPI::DOUBLE\_COMPLEX, and MPI::LONG\_DOUBLE\_COMPLEX, which are removed in MPI-3.0.

TICKET NOT YET PASSED. CHANGE-LOG TEXT AS ON TICKET.

34. Sections 5.9.2 on pages 188.

MPI\_C\_COMPLEX is added to the “Complex” reduction group.

TICKET NOT YET PASSED. CHANGE-LOG TEXT AS ON TICKET.

35. Section 2.3 on page 10, and Sections 16.2.1, 16.2.2, 16.2.7 on pages 640, 642, and 657. The new mpi\_08 Fortran module is introduced.

36. Section 2.5.1 on page 12, Section 16.2.3 on page 644, Section 16.2.2 on page 642, and Section 16.2.7 on page 657.

Handles to opaque objects are defined as named types within the mpi\_08 Fortran module. The handle types are also available through the mpi Fortran module.

37. Sections 2.5.4, 2.5.5 on pages 15, 16, Sections 16.2.1, 16.2.10, 16.2.11, 16.2.12, 16.2.13 on pages 640, 668, 670, 671, 674, and Sections 16.2.3, 16.2.2, 16.2.7 on pages 644, 642, 657.

Within the `mpi_08` Fortran module, choice buffers are defined as assumed-type and assumed-rank according to Fortran 2008 TR 29113 [41], and the compile-time constant `MPI_SUBARRAYS_SUPPORTED` is set to `.TRUE..` With this, Fortran subscript triplets can be used in nonblocking MPI operations; vector subscripts are not supported in nonblocking operations. If the compiler does not support this Fortran TR 29113 feature, the constant is set to `.FALSE..`

38. Section 2.6.2 on page 18, Section 16.2.2 on page 642, and Section 16.2.7 on page 657. The `ierror` dummy arguments are `OPTIONAL` within the `mpi_08` Fortran module.

39. Section 3.2.5 on page 34, Section 16.2.3 on page 644, Section 16.2.2 on page 642, Section 16.2.7 on page 657, and Section 16.3.5 on page 696.

Within the `mpi_08` Fortran module, the status is defined as `TYPE(MPI_Status)`. New conversion routines are added: `MPI_STATUS_F2F08`, `MPI_STATUS_F082F`, `MPI_Status_c2f08`, and `MPI_Status_f082c`. In `mpi.h`, the new type `MPI_F08_status`, and the external variables `MPI_F08_STATUS_IGNORE` and `MPI_F08_STATUSES_IGNORE` are added.

40. Section 3.2.6 on page 36, and Section 3.8 on page 71.

`MPI_STATUS_IGNORE` can be also used in `MPI_IProbe`, `MPI_Probe`, `MPI_Improbe`, and `MPI_Mprobe`.

41. Section 3.6 on page 49.

In Fortran with the `mpi` module or `mpif.h`, the type of the `buffer_addr` argument of `MPI_BUFFER_DETACH` is wrongly defined and the argument is therefore unused.

42. Section 4.1 on page 91, Section 4.1.6 on page 112, and Section 16.2.15 on page 675.

The Fortran alignments of basic datatypes are implementation dependent. It is recommended that they are computed according to `BIND(C)` derived types. If an array of structures (in C/C++) or derived types (in Fortran) should be communicated, it is recommended that the user creates a portable datatype handle and applies additionally `MPI_TYPE_CREATE_RESIZED` to this datatype handle.

43. Sections 4.1.10, 5.9.5, 5.9.7, 6.7.4, 6.8, 8.3.1, 8.3.2, 8.3.3, 15.1, 16.2.9 on pages 118,

195, 201, 288, 294, 357, 359, 361, 617, and 660. In some routines, the dummy argument names were changed because they were identical to the Fortran keywords `TYPE` and `FUNCTION`. The new dummy argument names must be used because the `mpi` and `mpi_08` modules guarantee keyword-based actual argument lists. The argument name `type` was changed into `oldtype` in `MPI_TYPE_DUP`, and into `datatype` in the Fortran `USER_FUNCTION` of `MPI_OP_CREATE`, and in `MPI_TYPE_SET_ATTR`, `MPI_TYPE_GET_ATTR`, `MPI_TYPE_DELETE_ATTR`, `MPI_TYPE_SET_NAME`, `MPI_TYPE_GET_NAME`, `MPI_TYPE_MATCH_SIZE`, in the callback prototype definition `MPI_Type_delete_attr_function`, and the predefined callback function `MPI_TYPE_NULL_DELETE_FN`; function was changed into `user_fn` in `MPI_OP_CREATE`, into `comm_errhandler_fn` in `MPI_COMM_CREATE_ERRHANDLER`, into `win_errhandler_fn` in `MPI_WIN_CREATE_ERRHANDLER`, into `file_errhandler_fn` in `MPI_FILE_CREATE_ERRHANDLER`, into `handler_fn` in `MPI_ERRHANDLER_CREATE`.

For consistency reasons, `INOUBUF` was changed into `INOUTBUF` in

`MPI_REDUCE_LOCAL`, and `intracomm` into `newintracomm` in `MPI_INTERCOMM_MERGE`.

ticket245-Q.

44. Section 8.2 on page 352.  
In Fortran with the `mpi` and `mpi_f08` modules, `MPI_ALLOC_MEM` now also supports `TYPE(C_PTR)` C-pointer instead of only returning an address-sized integer that may be usable together with non-standard Cray-pointer. The Fortran interfaces with `INTEGER(KIND=MPI_ADDRESS_KIND)` `BASEPTR` in the `mpi` module and the `mpif.h` include file are deprecated since MPI-3.0.
45. Section 16.2.15 on page 675, and Section 16.2.7 on page 657.  
Fortran `SEQUENCE` and `BIND(C)` derived application types can be used as buffers in MPI operations.
46. Section 16.2.16 on page 677 to Section 16.2.19 on page 686, Section 16.2.7 on page 657, and Section 16.2.8 on page 659.  
The sections about Fortran optimization problems and their solution is partially rewritten and new methods are added, e.g., the use of the `ASYNCHRONOUS` attribute. The constant `MPI_ASYNC_PROTECTS_NONBLOCKING` tells whether the meaning of the `ASYNCHRONOUS` attribute is extended to protect nonblocking operations. The Fortran routine `MPI_F_SYNC_REG` is added. To achieve a secure and portable programming interfaces, in Section 16.2.7, several requirements are defined for the combination of an MPI library and a Fortran compiler to be MPI-3.0 compliant.
47. Section 16.2.4 on page 647.  
The use of the `mpif.h` Fortran include file is strongly discouraged.
48. Section 16.2.3 on page 644, and Section 16.2.7 on page 657.  
The existing `mpi` Fortran module must implement compile-time argument checking.
49. Section 16.2.2 on page 642.  
Within the `mpi_08` Fortran module, dummy arguments are declared with `INTENT=IN`, `OUT`, or `INOUT` as defined in the `mpi_08` interfaces.
50. Section 16.2.7 on page 657.  
This new section summarizes requirements that an MPI library together with a Fortran compiler is compliant to the MPI standard.
51. Section A.1.1, Table “*Predefined functions*” on page 718, Section A.1.3 on page 726, and Section A.3.4 on page 771.  
Within the new `mpi_f08` module, all callback prototype definitions are defined with explicit interfaces `PROCEDURE(MPI_...)` with `BIND(C)` attribute.
52. Section A.1.3 on page 726.  
In some routines, the Fortran callback prototype names were changed from `..._FN` to `..._FUNCTION` to be consistent with the other language bindings.

## B.2 TEST for Tickets 271, 168, 204, 280, 286, 287

This section is not part of the MPI standard and will be removed after the next meeting, July 2012.

The correctness of these interfaces must be verified!

```

1  MPI_Comm_dup_info(comm, info, newcomm, ierror) BIND(C)
2      TYPE(MPI_Comm), INTENT(IN) :: comm
3      TYPE(MPI_Info), INTENT(IN) :: info
4      TYPE(MPI_Comm), INTENT(OUT) :: newcomm
5      INTEGER, OPTIONAL, INTENT(OUT) :: ierror
6
7  MPI_COMM_DUP_INFO(COMM, INFO, NEWCOMM, IERROR)
8      INTEGER COMM, INFO, NEWCOMM, IERROR
ticket-248T.
9  MPI_Comm_set_info(comm, info, ierror) BIND(C)
10     TYPE(MPI_Comm), INTENT(IN) :: comm
11     TYPE(MPI_Info), INTENT(IN) :: info
12     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
13
14  MPI_COMM_SET_INFO(COMM, INFO, IERROR)
15     INTEGER COMM, INFO, IERROR
ticket-248T.
16  MPI_Comm_get_info(comm, info_used, ierror) BIND(C)
17     TYPE(MPI_Comm), INTENT(IN) :: comm
18     TYPE(MPI_Info), INTENT(OUT) :: info_used
19     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
20
21  MPI_COMM_GET_INFO(COMM, INFO_USED, IERROR)
22     INTEGER COMM, INFO_USED, IERROR
ticket-248T.
23  MPI_Win_set_info(win, info, ierror) BIND(C)
24     TYPE(MPI_Win), INTENT(IN) :: win
25     TYPE(MPI_Info), INTENT(IN) :: info
26     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
27
28  MPI_WIN_SET_INFO(WIN, INFO, IERROR)
29     INTEGER WIN, INFO, IERROR
ticket-248T.
30  MPI_Win_get_info(win, info_used, ierror) BIND(C)
31     TYPE(MPI_Win), INTENT(IN) :: win
32     TYPE(MPI_Info), INTENT(OUT) :: info_used
33     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
34
35  MPI_WIN_GET_INFO(WIN, INFO_USED, IERROR)
36     INTEGER WIN, INFO_USED, IERROR
ticket-248T.
37  MPI_Comm_idup(comm, newcomm, request, ierror) BIND(C)
38     TYPE(MPI_Comm), INTENT(IN) :: comm
39     TYPE(MPI_Comm), ASYNCHRONOUS :: newcomm
40     TYPE(MPI_Request), INTENT(OUT) :: request
41     INTEGER, OPTIONAL, INTENT(OUT) :: ierror
42
43  MPI_COMM_IDUP(COMM, NEWCOMM, REQUEST, IERROR)
44     INTEGER COMM, NEWCOMM, REQUEST, IERROR
ticket-248T.
45  MPI_Get_library_version(version, resultlen, ierror) BIND(C)
46     CHARACTER(LEN=MPI_MAX_LIBRARY_VERSION_STRING), INTENT(OUT) :: version
47     INTEGER, INTENT(OUT) :: resultlen
48

```

```

    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_GET_LIBRARY_VERSION(VERSION, RESULTEN, IERROR)
    CHARACTER*(*) VERSION
    INTEGER RESULTLEN, IERROR
MPI_Type_create_hindexed_block(count, blocklength, array_of_displacements,
    oldtype, newtype, ierror) BIND(C)
    INTEGER, INTENT(IN) :: count, blocklength
    INTEGER(kind=MPI_Address_kind), INTENT(IN) ::
    array_of_displacements(count)
    TYPE(MPI_Datatype), INTENT(IN) :: oldtype
    TYPE(MPI_Datatype), INTENT(OUT) :: newtype
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_TYPE_CREATE_HINDEXED_BLOCK(COUNT, BLOCKLENGTH, ARRAY_OF_DISPLACEMENTS,
    OLDTYPE, NEWTYPE, IERROR)
    INTEGER COUNT, BLOCKLENGTH, OLDTYPE, NEWTYPE, IERROR
    INTEGER(KIND=MPI_ADDRESS_KIND) ARRAY_OF_DISPLACEMENTS(*)
MPI_Comm_create_group(comm, group, tag, newcomm, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    TYPE(MPI_Group), INTENT(IN) :: group
    INTEGER, INTENT(IN) :: tag
    TYPE(MPI_Comm), INTENT(OUT) :: newcomm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_COMM_CREATE_GROUP(COMM, GROUP, TAG, NEWCOMM, IERROR)
    INTEGER COMM, GROUP, TAG, NEWCOMM, IERROR
MPI_Comm_split_type(comm, split_type, key, info, newcomm, ierror) BIND(C)
    TYPE(MPI_Comm), INTENT(IN) :: comm
    INTEGER, INTENT(IN) :: split_type, key
    TYPE(MPI_Info), INTENT(IN) :: info
    TYPE(MPI_Comm), INTENT(OUT) :: newcomm
    INTEGER, OPTIONAL, INTENT(OUT) :: ierror
MPI_COMM_SPLIT_TYPE(COMM, SPLIT_TYPE, KEY, INFO, NEWCOMM, IERROR)
    INTEGER COMM, SPLIT_TYPE, KEY, INFO, NEWCOMM, IERROR

```

ticket-248T.

ticket-248T.

ticket-248T.

### B.3 Changes from Version 2.1 to Version 2.2

1. Section 2.5.4 on page 15.  
It is now guaranteed that predefined named constant handles (as other constants) can be used in initialization expressions or assignments, i.e., also before the call to MPI\_INIT.
2. Section 2.6 on page 17, Section 2.6.4 on page 20, and Section 16.1 on page 627.  
The C++ language bindings have been deprecated and may be removed in a future version of the MPI specification.



- 1       3. Section 3.2.2 on page 29.  
2       MPI\_CHAR for printable characters is now defined for C type char (instead of signed  
3       char). This change should not have any impact on applications nor on MPI libraries  
4       (except some comment lines), because printable characters could and can be stored in  
5       any of the C types char, signed char, and unsigned char, and MPI\_CHAR is not allowed  
6       for predefined reduction operations.
- 7  
8       4. Section 3.2.2 on page 29.  
9       MPI\_(U)INT{8,16,32,64}\_T, MPI\_AINT, MPI\_OFFSET, MPI\_C\_BOOL,  
10       MPI\_C\_COMPLEX, MPI\_C\_FLOAT\_COMPLEX, MPI\_C\_DOUBLE\_COMPLEX, and  
11       MPI\_C\_LONG\_DOUBLE\_COMPLEX are now valid predefined MPI datatypes.
- 12  
13       5. Section 3.4 on page 41, Section 3.7.2 on page 54, Section 3.9 on page 80, and Section 5.1  
14       on page 151.  
15       The read access restriction on the send buffer for blocking, non blocking and collective  
16       API has been lifted. It is permitted to access for read the send buffer while the  
17       operation is in progress.
- 18  
19       6. Section 3.7 on page 52.  
20       The Advice to users for IBSEND and IRSEND was slightly changed.
- 21  
22       7. Section 3.7.3 on page 58.  
23       The advice to free an active request was removed in the Advice to users for  
24       MPI\_REQUEST\_FREE.
- 25  
26       8. Section 3.7.6 on page 70.  
27       MPI\_REQUEST\_GET\_STATUS changed to permit inactive or null requests as input.
- 28  
29       9. Section 5.8 on page 180.  
30       “*In place*” option is added to MPI\_ALLTOALL, MPI\_ALLTOALLV, and  
31       MPI\_ALLTOALLW for intracommunicators.
- 32  
33       10. Section 5.9.2 on page 188.  
34       Predefined parameterized datatypes (e.g., returned by  
35       MPI\_TYPE\_CREATE\_F90\_REAL) and optional named predefined datatypes (e.g.  
36       MPI\_REAL8) have been added to the list of valid datatypes in reduction operations.
- 37  
38       11. Section 5.9.2 on page 188.  
39       MPI\_(U)INT{8,16,32,64}\_T are all considered C integer types for the purposes of the  
40       predefined reduction operators. MPI\_AINT and MPI\_OFFSET are considered Fortran  
41       integer types. MPI\_C\_BOOL is considered a Logical type.  
42       MPI\_C\_COMPLEX, MPI\_C\_FLOAT\_COMPLEX, MPI\_C\_DOUBLE\_COMPLEX, and  
43       MPI\_C\_LONG\_DOUBLE\_COMPLEX are considered Complex types.
- 44  
45       12. Section 5.9.7 on page 201.  
46       The local routines MPI\_REDUCE\_LOCAL and MPI\_OP\_COMMUTATIVE have been  
47       added.
- 48  
49       13. Section 5.10.1 on page 203.  
50       The collective function MPI\_REDUCE\_SCATTER\_BLOCK is added to the MPI stan-  
51       dard.



14. Section 5.11.2 on page 207. 1  
 Added in place argument to MPI\_EXSCAN. 2  
3
15. Section 6.4.2 on page 252, and Section 6.6 on page 268. 4  
 Implementations that did not implement MPI\_COMM\_CREATE on intercommuni- 5  
 cators will need to add that functionality. As the standard described the behav- 6  
 ior of this operation on intercommunicators, it is believed that most implementa- 7  
 tions already provide this functionality. Note also that the C++ binding for both 8  
 MPI\_COMM\_CREATE and MPI\_COMM\_SPLIT explicitly allow Intercomms. 9  
10
16. Section 6.4.2 on page 252. 11  
 MPI\_COMM\_CREATE is extended to allow several disjoint subgroups as input if comm 12  
 is an intracommunicator. If comm is an intercommunicator it was clarified that all 13  
 processes in the same local group of comm must specify the same value for group. 14
17. Section 7.5.4 on page 308. 15  
 New functions for a scalable distributed graph topology interface has been added. 16  
 In this section, the functions MPI\_DIST\_GRAPH\_CREATE\_ADJACENT and 17  
 MPI\_DIST\_GRAPH\_CREATE, the constants MPI\_UNWEIGHTED, and the derived C++ 18  
 class Distgraphcomm were added. 19  
20
18. Section 7.5.5 on page 315. 21  
 For the scalable distributed graph topology interface, the functions [ 22 ticket201.  
 MPI\_DIST\_NEIGHBORS\_COUNT]MPI\_DIST\_GRAPH\_NEIGHBORS\_COUNT and [ 23 ticket201.  
 MPI\_DIST\_NEIGHBORS]MPI\_DIST\_GRAPH\_NEIGHBORS and the constant 24  
 MPI\_DIST\_GRAPH were added. 25  
26
19. Section 7.5.5 on page 315. 27  
 Remove ambiguity regarding duplicated neighbors with MPI\_GRAPH\_NEIGHBORS 28  
 and MPI\_GRAPH\_NEIGHBORS\_COUNT. 29
20. Section 8.1.1 on page 349. 30  
 The subversion number changed from 1 to 2. 31  
32
21. Section 8.3 on page 355, Section 15.2 on page 624, and Annex A.1.3 on page 726. 33  
 Changed function pointer typedef names MPI\_{Comm,File,Win}\_errhandler\_fn to 34  
 MPI\_{Comm,File,Win}\_errhandler\_function. Deprecated old “\_fn” names. 35  
36
22. Section 8.7.1 on page 377. 37  
 Attribute deletion callbacks on MPI\_COMM\_SELF are now called in LIFO order. Imple- 38  
 mentors must now also register all implementation-internal attribute deletion callbacks 39  
 on MPI\_COMM\_SELF before returning from MPI\_INIT/MPI\_INIT\_THREAD. 40
23. Section 11.3.4 on page 438. 41  
 The restriction added in MPI 2.1 that the operation MPI\_REPLACE in 42  
 MPI\_ACCUMULATE can be used only with predefined datatypes has been removed. 43  
 MPI\_REPLACE can now be used even with derived datatypes, as it was in MPI 2.0. 44  
 Also, a clarification has been made that MPI\_REPLACE can be used only in 45  
 MPI\_ACCUMULATE, not in collective operations that do reductions, such as 46  
 MPI\_REDUCE and others. 47  
48

24. Section 12.2 on page 491.  
Add “\*” to the `query_fn`, `free_fn`, and `cancel_fn` arguments to the C++ binding for `MPI::Grequest::Start()` for consistency with the rest of MPI functions that take function pointer arguments.
25. Section 13.5.2 on page 558, and Table 13.2 on page 560.  
`MPI_(U)INT{8,16,32,64}_T`, `MPI_AINT`, `MPI_OFFSET`, `MPI_C_COMPLEX`, `MPI_C_FLOAT_COMPLEX`, `MPI_C_DOUBLE_COMPLEX`, `MPI_C_LONG_DOUBLE_COMPLEX`, and `MPI_C_BOOL` are added as predefined datatypes in the external32 representation.
26. Section 16.3.7 on page 702.  
The description was modified that it only describes how an MPI implementation behaves, but not how MPI stores attributes internally. The erroneous MPI-2.1 Example 16.17 was replaced with three new examples 16.25, 16.26, and 16.27 on pages 703-705 explicitly detailing cross-language attribute behavior. Implementations that matched the behavior of the old example will need to be updated.
27. Annex A.1.1 on page 709.  
Removed type `MPI::Fint` (compare `MPI_Fint` in Section A.1.2 on page 725).
28. Annex A.1.1 on page 709. Table *Named Predefined Datatypes*.  
Added `MPI_(U)INT{8,16,32,64}_T`, `MPI_AINT`, `MPI_OFFSET`, `MPI_C_BOOL`, `MPI_C_FLOAT_COMPLEX`, `MPI_C_COMPLEX`, `MPI_C_DOUBLE_COMPLEX`, and `MPI_C_LONG_DOUBLE_COMPLEX` are added as predefined datatypes.

## B.4 Changes from Version 2.0 to Version 2.1

1. Section 3.2.2 on page 29, Section 16.1.6 on page 631, and Annex A.1 on page 709.  
In addition, the `MPI_LONG_LONG` should be added as an optional type; it is a synonym for `MPI_LONG_LONG_INT`.
2. Section 3.2.2 on page 29, Section 16.1.6 on page 631, and Annex A.1 on page 709.  
`MPI_LONG_LONG_INT`, `MPI_LONG_LONG` (as synonym), `MPI_UNSIGNED_LONG_LONG`, `MPI_SIGNED_CHAR`, and `MPI_WCHAR` are moved from optional to official and they are therefore defined for all three language bindings.
3. Section 3.2.5 on page 34.  
`MPI_GET_COUNT` with zero-length datatypes: The value returned as the count argument of `MPI_GET_COUNT` for a datatype of length zero where zero bytes have been transferred is zero. If the number of bytes transferred is greater than zero, `MPI_UNDEFINED` is returned.
4. Section 4.1 on page 91.  
General rule about derived datatypes: Most datatype constructors have replication count or block length arguments. Allowed values are non-negative integers. If the value is zero, no elements are generated in the type map and there is no effect on datatype bounds or extent.

5. Section 4.3 on page 146.  
MPI\_BYTE should be used to send and receive data that is packed using MPI\_PACK\_EXTERNAL.
6. Section 5.9.6 on page 199.  
If `comm` is an intercommunicator in MPI\_ALLREDUCE, then both groups should provide `count` and `datatype` arguments that specify the same type signature (i.e., it is not necessary that both groups provide the same `count` value).
7. Section 6.3.1 on page 242.  
MPI\_GROUP\_TRANSLATE\_RANKS and MPI\_PROC\_NULL: MPI\_PROC\_NULL is a valid rank for input to MPI\_GROUP\_TRANSLATE\_RANKS, which returns MPI\_PROC\_NULL as the translated rank.
8. Section 6.7 on page 276.  
About the attribute caching functions:  
  
*Advice to implementors.* High-quality implementations should raise an error when a keyval that was created by a call to MPI\_XXX\_CREATE\_KEYVAL is used with an object of the wrong type with a call to MPI\_YYY\_GET\_ATTR, MPI\_YYY\_SET\_ATTR, MPI\_YYY\_DELETE\_ATTR, or MPI\_YYY\_FREE\_KEYVAL. To do so, it is necessary to maintain, with each keyval, information on the type of the associated user function. (*End of advice to implementors.*)
9. Section 6.8 on page 294.  
In MPI\_COMM\_GET\_NAME: In C, a null character is additionally stored at `name[resultlen]`. `resultlen` cannot be larger than MPI\_MAX\_OBJECT\_NAME-1. In Fortran, `name` is padded on the right with blank characters. `resultlen` cannot be larger than MPI\_MAX\_OBJECT\_NAME.
10. Section 7.4 on page 302.  
About MPI\_GRAPH\_CREATE and MPI\_CART\_CREATE: All input arguments must have identical values on all processes of the group of `comm_old`.
11. Section 7.5.1 on page 304.  
In MPI\_CART\_CREATE: If `ndims` is zero then a zero-dimensional Cartesian topology is created. The call is erroneous if it specifies a grid that is larger than the group size or if `ndims` is negative.
12. Section 7.5.3 on page 306.  
In MPI\_GRAPH\_CREATE: If the graph is empty, i.e., `nnodes == 0`, then MPI\_COMM\_NULL is returned in all processes.
13. Section 7.5.3 on page 306.  
In MPI\_GRAPH\_CREATE: A single process is allowed to be defined multiple times in the list of neighbors of a process (i.e., there may be multiple edges between two processes). A process is also allowed to be a neighbor to itself (i.e., a self loop in the graph). The adjacency matrix is allowed to be non-symmetric.

- Advice to users.* Performance implications of using multiple edges or a non-symmetric adjacency matrix are not defined. The definition of a node-neighbor edge does not imply a direction of the communication. (*End of advice to users.*)
14. Section 7.5.5 on page 315.  
In `MPI_CARTDIM_GET` and `MPI_CART_GET`: If `comm` is associated with a zero-dimensional Cartesian topology, `MPI_CARTDIM_GET` returns `ndims=0` and `MPI_CART_GET` will keep all output arguments unchanged.
  15. Section 7.5.5 on page 315.  
In `MPI_CART_RANK`: If `comm` is associated with a zero-dimensional Cartesian topology, `coord` is not significant and 0 is returned in `rank`.
  16. Section 7.5.5 on page 315.  
In `MPI_CART_COORDS`: If `comm` is associated with a zero-dimensional Cartesian topology, `coords` will be unchanged.
  17. Section 7.5.6 on page 323.  
In `MPI_CART_SHIFT`: It is erroneous to call `MPI_CART_SHIFT` with a direction that is either negative or greater than or equal to the number of dimensions in the Cartesian communicator. This implies that it is erroneous to call `MPI_CART_SHIFT` with a `comm` that is associated with a zero-dimensional Cartesian topology.
  18. Section 7.5.7 on page 325.  
In `MPI_CART_SUB`: If all entries in `remain_dims` are false or `comm` is already associated with a zero-dimensional Cartesian topology then `newcomm` is associated with a zero-dimensional Cartesian topology.
  - 18.1. Section 8.1.1 on page 349.  
The subversion number changed from 0 to 1.
  19. Section 8.1.2 on page 350.  
In `MPI_GET_PROCESSOR_NAME`: In C, a null character is additionally stored at `name[resultlen]`. `resultlen` cannot be larger than `MPI_MAX_PROCESSOR_NAME-1`. In Fortran, `name` is padded on the right with blank characters. `resultlen` cannot be larger than `MPI_MAX_PROCESSOR_NAME`.
  20. Section 8.3 on page 355.  
`MPI_{COMM,WIN,FILE}_GET_ERRHANDLER` behave as if a new error handler object is created. That is, once the error handler is no longer needed, `MPI_ERRHANDLER_FREE` should be called with the error handler returned from `MPI_ERRHANDLER_GET` or `MPI_{COMM,WIN,FILE}_GET_ERRHANDLER` to mark the error handler for deallocation. This provides behavior similar to that of `MPI_COMM_GROUP` and `MPI_GROUP_FREE`.
  21. Section 8.7 on page 372, see explanations to `MPI_FINALIZE`.  
`MPI_FINALIZE` is collective over all connected processes. If no processes were spawned, accepted or connected then this means over `MPI_COMM_WORLD`; otherwise it is collective over the union of all processes that have been and continue to be connected, as explained in Section 10.5.4 on page 414.

22. Section 8.7 on page 372.  
About MPI\_ABORT:  
  
*Advice to users.* Whether the errorcode is returned from the executable or from the MPI process startup mechanism (e.g., mpiexec), is an aspect of quality of the MPI library but not mandatory. (*End of advice to users.*)  
  
*Advice to implementors.* Where possible, a high-quality implementation will try to return the errorcode from the MPI process startup mechanism (e.g. mpiexec or singleton init). (*End of advice to implementors.*)
23. Section 9 on page 381.  
An implementation must support info objects as caches for arbitrary (key, value) pairs, regardless of whether it recognizes the key. Each function that takes hints in the form of an MPI\_Info must be prepared to ignore any key it does not recognize. This description of info objects does not attempt to define how a particular function should react if it recognizes a key but not the associated value. MPI\_INFO\_GET\_NKEYS, MPI\_INFO\_GET\_NTHKEY, MPI\_INFO\_GET\_VALUELEN, and MPI\_INFO\_GET must retain all (key,value) pairs so that layered functionality can also use the Info object.
24. Section 11.3 on page 431.  
MPI\_PROC\_NULL is a valid target rank in the MPI RMA calls MPI\_ACCUMULATE, MPI\_GET, and MPI\_PUT. The effect is the same as for MPI\_PROC\_NULL in MPI point-to-point communication. See also item 25 in this list.
25. Section 11.3 on page 431.  
After any RMA operation with rank MPI\_PROC\_NULL, it is still necessary to finish the RMA epoch with the synchronization method that started the epoch. See also item 24 in this list.
26. Section 11.3.4 on page 438.  
MPI\_REPLACE in MPI\_ACCUMULATE, like the other predefined operations, is defined only for the predefined MPI datatypes.
27. Section 13.2.8 on page 519.  
About MPI\_FILE\_SET\_VIEW and MPI\_FILE\_SET\_INFO: When an info object that specifies a subset of valid hints is passed to MPI\_FILE\_SET\_VIEW or MPI\_FILE\_SET\_INFO, there will be no effect on previously set or defaulted hints that the info does not specify.
28. Section 13.2.8 on page 519.  
About MPI\_FILE\_GET\_INFO: If no hint exists for the file associated with fh, a handle to a newly created info object is returned that contains no key/value pair.
29. Section 13.3 on page 522.  
If a file does not have the mode MPI\_MODE\_SEQUENTIAL, then MPI\_DISPLACEMENT\_CURRENT is invalid as disp in MPI\_FILE\_SET\_VIEW.
30. Section 13.5.2 on page 558.  
The bias of 16 byte doubles was defined with 10383. The correct value is 16383.

31. Section [16.1.4](#) on page [628](#).

In the example in this section, the buffer should be declared as `const void* buf`.

32. Section [16.2.9](#) on page [660](#).

About `MPI_TYPE_CREATE_F90_XXXX`:

*Advice to implementors.* An application may often repeat a call to `MPI_TYPE_CREATE_F90_XXXX` with the same combination of `(XXXX,p,r)`. The application is not allowed to free the returned predefined, unnamed datatype handles. To prevent the creation of a potentially huge amount of handles, the MPI implementation should return the same datatype handle for the same `(REAL/COMPLEX/INTEGER,p,r)` combination. Checking for the combination `(p,r)` in the preceding call to `MPI_TYPE_CREATE_F90_XXXX` and using a hash-table to find formerly generated handles should limit the overhead of finding a previously generated datatype with same combination of `(XXXX,p,r)`. (*End of advice to implementors.*)

33. Section [A.1.1](#) on page [709](#).

`MPI_BOTTOM` is defined as `void * const MPI::BOTTOM`.

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# MPI Callback Function Prototype Index

This index lists the C typedef names for callback routines, such as those used with attribute caching or user-defined reduction operations. C++ names for these typedefs and Fortran example prototypes are given near the text of the C name.

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