# Fixing Probe for Multi-Threaded MPI Applications (Revision 1)

Douglas Gregor, Torsten Hoefler, and Andrew Lumsdaine {dgregor,htor,lums}@osl.iu.edu

June 5, 2008

### 1 Introduction

MPI's message-probing operations, MPI\_Probe and MPI\_Iprobe, are useful in MPI applications that do not know *a priori* what messages they will receive or how much data those messages will contain. Such applications often have irregular, data-driven communication patterns or deal with data structures that require serialization for transmission.

Unfortunately, MPI's message-probing operations are unusable in multi-threaded applications where they could be most useful. The fundamental problem with MPI\_Probe and MPI\_lprobe functions is that a message found by a probe can still be matched and received by a receive operation in a different thread. Thus, despite the fact that a probe operation returns a source and tag that can be used to receive a message, there is no guarantee that the message will still be available when that receive operation is invoked. For example, the following code can not be executed concurrently by two threads in an MPI process, because a message could be found by the MPI\_Probe in both threads, while only one of the threads could successfully receive the message (the other will block):

The lack of a usable threaded MPI\_Probe or MPI\_Iprobe causes serious problems for the construction of language bindings for high-level object-oriented and generic languages, where users would like to be able to transmit objects that require serialization, including C++ [?,?], Java [?], Python [?], and C# [?]. MPI provides a mechanism to "pack" (serialize) an object into a buffer that can be transmitted via MPI, then "unpack" (de-serialize) that object at the receiver's end. However, while MPI provides good support for serialization and sending serialized data, it does not provide adequate support for receiving serialized data. The problem is that, in general, the receiver cannot know the length of the serialized data before it posts the receive. MPI\_Probe and MPI\_Iprobe provide this functionality, but they are

unusable in a multi-threaded environment. Thus, bindings for these languages must resort to elaborate and inefficient workarounds [?].

There is no known workaround that addresses all of the problems with MPI\_Probe and MPI\_Iprobe in multi-threaded MPI applications. Therefore, we propose extensions for MPI-3 that introduce a new kind of probe—a "matched" probe—and a set of corresponding receive operations. The new probe matches a message and returns a handle to that specific message, which cannot be found by any other probe operation or matched by any other receive. The new receive operations allow the receipt of a message based on the message handle returned from this probe. These extensions allow the use of probe in a multi-threaded context, ensuring that the message found by probe is the message received.

In the following example, we illustrate how the new probe operation, MPI\_Mprobe, can be used to receive a message of unknown length. Note that this code can be concurrently executed in several threads, each of which will receive different messages.

```
MPI_Message msg;
MPI_Status status;
/* Match a message */
MPI_Mprobe(MPI_ANY_SOURCE, MPI_ANY_TAG, MPI_COMM_WORLD, &msg, &status);
/* Allocate memory to receive the message */
int count;
MPI_get_count(&status, MPI_BYTE, &count);
char* buffer = malloc(count);

/* Receive this message. */
MPI_Mrecv(buffer, count, MPI_BYTE, &msg, MPI_STATUS_IGNORE);
```

# 2 Proposed Extensions

## 2.1 Message handles

A message handle returned by a matching probe (MPI\_Mprobe or MPI\_Improbe) has type MPI\_Message.

```
int MPI_Message_cancel(MPI_Message *message)

MPI_MESSAGE_CANCEL(MESSAGE, IERROR)
INTEGER MESSAGE, IERROR

void Message::Cancel()

INOUT message the message to be cancelled (Message)
```

A call to MPI\_MESSAGE\_CANCEL cancels the receipt of a message matched by a matching probe. A cancelled message cannot be received.

Advice to implementers. Because no receive buffers have been posted for a receive, cancellation always succeeds even if the underlying interconnect does not permit the cancellation of transmissions after they have been matched. A valid implementation of MPI\_Message\_cancel that supports such interconnects is:

```
int count; /* set to the size of the message */
char* buffer = malloc(count);
MPI_Mrecv(buffer, count, MPI_BYTE, &msg, MPI_STATUS_IGNORE);
free(buffer); \Box
```

## 2.2 Matching Probe

The MPI\_MPROBE and MPI\_IMPROBE operations allow incoming messages to be queried without actually receiving them. The user can then decide how to receive them, based on the information returned by the probe. In particular, the user may allocate memory for the receive buffer, according to the length of the probed message.

```
MPI_IMPROBE(SOURCE, TAG, COMM, FLAG, MESSAGE, STATUS, IERROR)
LOGICAL FLAG
INTEGER SOURCE, TAG, COMM, MESSAGE, STATUS(MPI_STATUS_SIZE), IERROR
```

```
bool Comm::Improbe(int source, int tag, Message& message) const
bool Comm::Improbe(int source, int tag, Message& message, Status& status) const
```

```
IN source source rank or MPI_ANY_SOURCE (integer)
IN tag tag value or MPI_ANY_TAG (integer)
IN comm communicator (handle)
OUT flag (logical)
OUT message message handle (Message)
OUT status status object (Status)
```

MPI\_IMPROBE(source, tag, comm, flag, message, status) returns flag = true if there is a message that can be received and that matches the pattern specified by the arguments source, tag, and comm. The call matches the same message that would have been received by a call to MPI\_RECV(..., source, tag, comm, status) executed at the same point in the program, and returns in message a handle to that message and in status the same value that would have been returned by MPI\_RECV(). Otherwise, the call returns flag = false, and leaves message undefined.

A matched receive executed with the message handle will receive the message that was matched by the probe. Unlike MPI\_IPROBE, no other probe or receive operation may match the message returned by MPI\_IMPROBE. Each message returned by MPI\_IMPROBE must either be completed with a matched receive or cancelled with MPI\_MESSAGE\_CANCEL.

The source argument of MPI\_IMPROBE can be MPI\_ANY\_SOURCE, and the tag argument can be MPI\_ANY\_TAG, so that one can probe for messages from an arbitrary source and/or with an arbitrary tag. However, a specific communication context must be provided with the comm argument.

MPI\_MPROBE(SOURCE, TAG, COMM, MESSAGE, STATUS, IERROR)
INTEGER SOURCE, TAG, COMM, MESSAGE, STATUS(MPI\_STATUS\_SIZE), IERROR

void Comm::Mprobe(int source, int tag, Message& message) const
void Comm::Mprobe(int source, int tag, Message& message, Status& status) const

IN source source rank or MPI\_ANY\_SOURCE (integer)

IN tag tag value or MPI\_ANY\_TAG (integer)

IN comm communicator (handle)

OUT message message handle (Message)

**OUT status** status object (Status)

MPI\_MPROBE behaves like MPI\_IMPROBE except that it is a blocking call that returns only after a matching message has been found.

Advice to users. Unlike the (deprecated) MPI\_PROBE and MPI\_IPROBE, MPI\_MPROBE and MPI\_IMPROBE can be safely used in a multi-threaded MPI program. A message returned by MPI\_MPROBE or MPI\_IMPROBE has already been matched, and can only be received with a matched receive (section 2.3) executed with the corresponding message handle. □

The MPI implementation of MPI\_MPROBE and MPI\_IMPROBE needs to guarantee progress: if a call to MPI\_MPROBE has been issued by a process, and a send that matches the probe has been initiated by some process, then the call to MPI\_MPROBE will return, unless the message is matched by a concurrent matching probe operation or received by another concurrent receive operation (that is executed by another thread at the probing process). Similarly, if a process busy waits with MPI\_IMPROBE and a matching message has been issued, then the call to MPI\_IMPROBE will eventually return flag = true unless the message is matched by a concurrent matching probe operation or received by another concurrent receive operation.

Editorial note: the definitions of MPI\_IPROBE and MPI\_Probe should remain the same as they are now, but we deprecate them by adding the following text:

MPI\_PROBE and MPI\_IPROBE are deprecated.

Rationale. MPI\_PROBE and MPI\_IPROBE find messages, but do not match them, which makes MPI\_PROBE and MPI\_IPROBE unusable in multi-threaded MPI programs. MPI\_MPROBE and MPI\_IMPROBE provide better semantics than MPI\_PROBE and MPI\_IPROBE for multi-threaded MPI programs.  $\Box$ 

#### 2.3 Matched receives

Messages that have been matched by a matching probe (section 2.2) can be received by a matched receive.

This call receives a message found by a matching probe operation (section 2.2).

The receive buffer consists of the storage containing **count** consecutive elements of the type specified by **datatype**, starting at address **buf**. The length of the received message must be less than or equal to the length of the receive buffer. An overflow error occurs if all incoming data does not fit, without truncation, into the receive buffer.

If a message that is shorter than the receive buffer arrives, then only those locations corresponding to the (shorter) message are modified.

**Example** The following example uses a matching probe and a ready receive to receive any message of any size. This code can be executed in multiple threads concurrently.

```
MPI_Message message;
MPI_Status status;
/* Match a message */
MPI_Mprobe(MPI_ANY_SOURCE, MPI_ANY_TAG, MPI_COMM_WORLD, &message, &status);
/* Allocate memory to receive the message */
int count;
MPI_Get_count(&status, MPI_BYTE, &count);
char* buffer = malloc(count);

/* Receive this message. */
MPI_Mrecv(buffer, count, MPI_BYTE, &message, MPI_STATUS_IGNORE); 

Description:
```

Rationale. MPI\_MRECV does not have a communicator parameter because the communicator was part of the matching probe operation. Requiring the communicator to also be passed into MPI\_MRECV would involve addition user code and additional error checking in the MPI implementation, with no clear benefit.  $\Box$ 

```
\label{eq:model} \begin{array}{ll} \textbf{int} \  \, \mathsf{MPI\_Imrecv}(\textbf{void}* \ \mathsf{buf}, \ \textbf{int} \  \, \mathsf{count}, \  \, \mathsf{MPI\_Datatype} \  \, \mathsf{datatype}, \  \, \mathsf{MPI\_Message}* \  \, \mathsf{message}, \\ & \mathsf{MPI\_Request} \  \, *\mathsf{request}) \end{array}
```

MPI\_IMRECV(BUF, COUNT, DATATYPE, MESSAGE, REQUEST, IERROR)

```
<type> BUF(*)
INTEGER COUNT, DATATYPE, MESSAGE, REQUEST, IERROR
```

void Message::Imrecv(void\* buf, int count, const Datatype& datatype, Request& request)

OUT buf initial address of receive buffer (choice)
IN count number of elements in receive buffer (integer)
IN datatype datatype of each receive buffer element (handle)
INOUT message message to be received (Message)
OUT request request object (Status)

Start a ready, non-blocking receive of a message found by a matching probe operation (section 2.2).

# 3 Revision History

#### Revision 1

- Renamed MPI\_Rprobe, MPI\_Irprobe, MPI\_Rrecv and MPI\_Irrecv to MPI\_Mprobe, MPI\_Improbe, MPI\_Mrecv, and MPI\_Imrecv, respectively.
- Made MPI\_Message an opaque handle; MPI\_Mprobe and MPI\_Improbe now return an MPI\_Status result to provide information about the message (source, tag, count). Removed MPI\_Message\_get\_count.
- Corrected numerous small errors in the C++ bindings.